Efficient Algorithms for Geometric Partial Matching

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Geometric (bipartite) matching







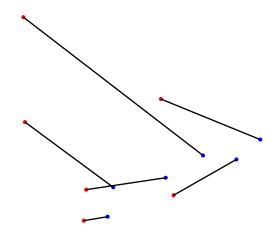






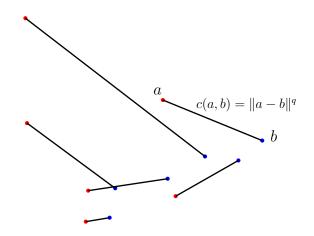
Geometric (bipartite) matching





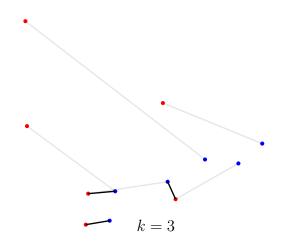
Geometric (bipartite) matching





Geometric (bipartite) partial matching





Prior work

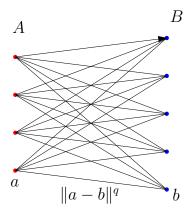


	approx.	time	
Hungarian algorithm (Kuhn)	exact	$O(km + k^2 \log n)$	$q \ge 1$
		$O(kn \operatorname{polylog} n)$	
Ramshaw, Tarjan 2012	exact ¹	$O(m\sqrt{k}\log(kC))$	$q \ge 1$
	$(1+\varepsilon)$	$O(n\sqrt{k}\operatorname{polylog} n\log(1/\varepsilon))$	
Sharathkumar, Agarwal 2012	$(1+\varepsilon)$	$O(n\operatorname{poly}(\log n, 1/\varepsilon)$	q = 1
new (Hungarian)	1	$O((n+k^2)\operatorname{polylog} n)$	$q \ge 1$
new (cost-scaling)	$(1+\varepsilon)$	$O((n+k\sqrt{k}) \operatorname{polylog} n \log(1/\varepsilon))$	$q \ge 1$

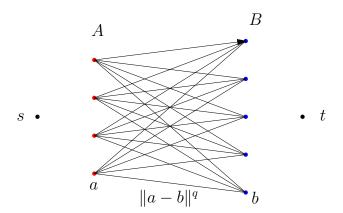


¹Assuming integer costs $\leq C$.

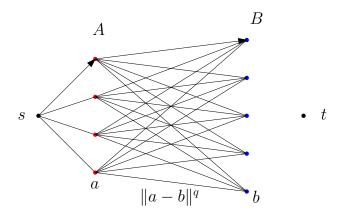




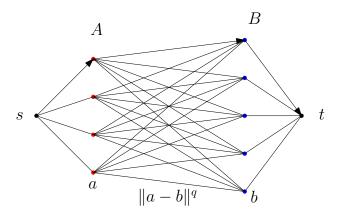




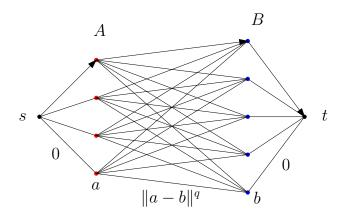




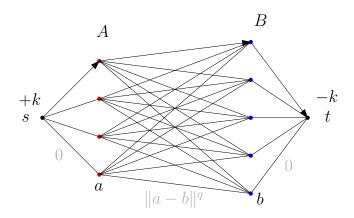




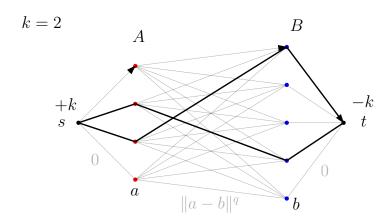




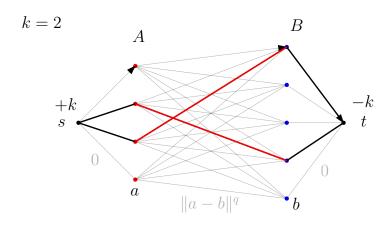






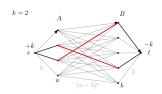






Primal-dual formulation and terminology





- primal variables f: 0-1 for each arc
- **dual** variables π : real values for each node
- Pseudoflow may have imbalance, circulations are balanced
- reduced cost: $c_{\pi}(v \rightarrow w) \coloneqq c(v \rightarrow w) \pi(v) + \pi(w)$
- ▶ dual feasibility: $c_{\pi}(v \rightarrow w) \ge 0$ on all residual arcs
- ▶ θ -optimality: $c_{\pi}(v \rightarrow w) \geq -\theta$ on all residual arcs
- ▶ admissible residual arcs: $c_{\pi}(v \rightarrow w) \leq 0$



Cost-scaling (Ramshaw-Tarjan)



- θ -optimality: $c_{\pi}(v \rightarrow w) \geq -\theta$ on all residual arcs
- ▶ admissible residual arcs: $c_{\pi}(v \rightarrow w) \leq 0$
- ightharpoonup heta-optimal circulation is $+n\theta$ approx.
- ▶ Find θ -optimal circulations for geometrically decreasing values of θ :
 - 1. Reduce $\theta \leftarrow \theta/2$, while creating O(k) excess.
 - 2. Refine this pseudoflow into a circulation, while preserving θ -optimality

Cost-scaling for geometric partial matching



• $(1+\varepsilon)$ -approx. geometric partial matching reduces into executing $O(\log(n^q/\varepsilon))$ cost scales.

Refinement by blocking flows (Ramshaw-Tarjan)



- Inside Refine:
 - Hungarian search: raise potentials until an excess-deficit admissible path exists.
 - 2. Augment by an admissible blocking flow.
- ► $O(\sqrt{k})$ blocking flows before 0 excess.

High-level goal per scale



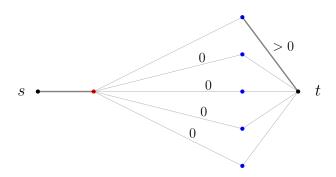
- Inside Refine:
 - 1. Hungarian search: raise potentials until an excess-deficit admissible path exists.
 - 2. Augment by an admissible blocking flow.
- After $O(n \operatorname{polylog} n)$ -time preprocessing, perform Hungarian search and find each blocking flow in $O(k \operatorname{polylog} n)$ time.

Hungarian search with BCP (Agarwal-Efrat-Sharir)

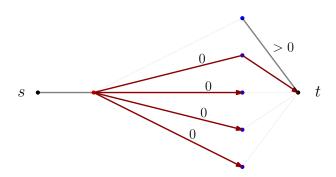


▶ Dynamic 2D BCP with $O(\operatorname{polylog} n)$ update time, $O(\log^2 n)$ query time (Kaplan *et al.* SODA'17)

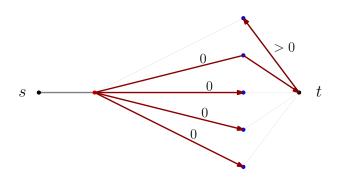




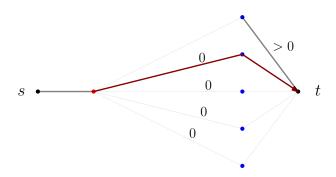






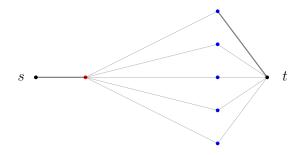








- ▶ Alive nodes: nonzero excess/deficit, or adjoining flow support arcs.
- ▶ Dead nodes: ones which aren't alive.

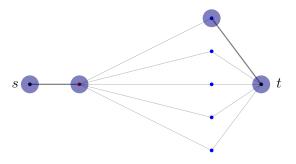


- ► Alive path: residual path between two alive nodes with no other alive nodes in between.
- Don't need to track potential of dead nodes.





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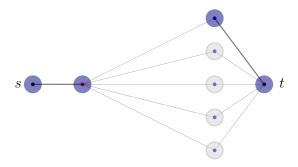


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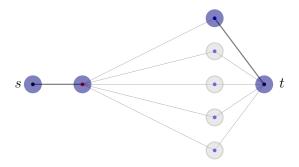


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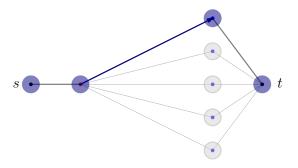


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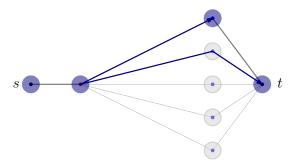


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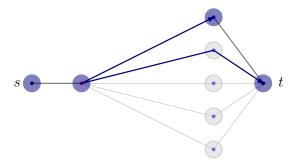


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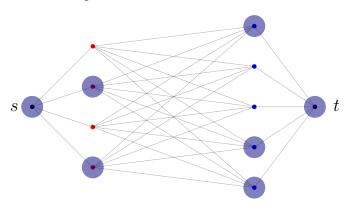
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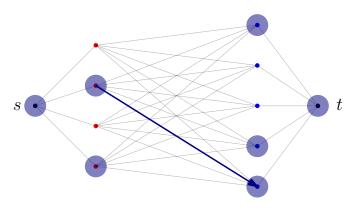




- ► Telescoping: $c_{\pi}(s \rightarrow a \rightarrow b) = c(a \rightarrow b) \pi(s) + \pi(b)$
- Can still use BCP.



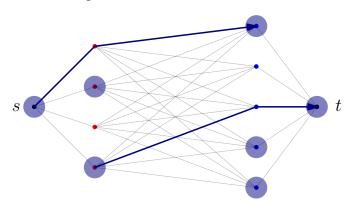




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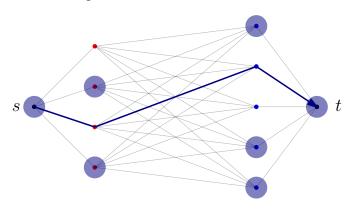




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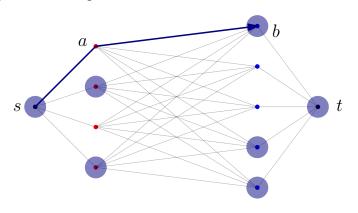




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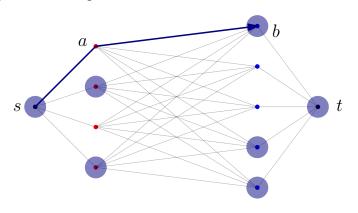




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Problem: BCP initialization



Initial BCP by rewinding



- $ightharpoonup X_t$ and X_{t+1} differ by the newly-matched A points.
- ► Persistence?

The End



Thank you.