

## Module 5–6

Name: NAME HERE

Questions requiring written answers.

1. You are given an array  $A[1..n]$  of  $n$  distinct elements such that each element in  $A$  is at most  $k$  positions away from its position in sorted order ( $2 \leq k < n$ ). Give an  $O(n \log k)$  time algorithm to sort the array. Prove your algorithm correct and analyze its running time.
2. You are given an array  $A[1..n]$  of  $n$  integers. We say the array is *well-positioned* with respect to a parameter  $k \in [1..n]$  iff for any  $k$  consecutive indices in the array, there are no two indices such that the value of one index is at least twice the value as the other. Give an  $O(n \log k)$  time algorithm to determine if  $A$  is well-positioned. Prove your algorithm correct and analyze its running time.
3. You are given three arrays of integers  $A$ ,  $B$ , and  $C$  with  $n$  elements each as well as a target integer  $t$ . Design an algorithm to determine if there exists three numbers  $a \in A$ ,  $b \in B$ , and  $c \in C$  such that  $a + b + c = t$ . Prove your algorithm correct and analyze its running time. Your algorithm should run in worst-case  $O(n^2 \log n)$  or  $O(n^2)$  expected time.
4. You are given a directed graph  $G = (V, E)$  with  $n$  vertices labeled 1 through  $n$ . For each vertex  $u$ , let  $H(u)$  be the highest-valued vertex that is reachable from  $u$ . Design an efficient algorithm to compute  $H(u)$  for each vertex in the graph.
5. Given a *tree* there is a unique simple path between any two vertices. Give an efficient algorithm to find the length of the longest such path.
6. You are given a directed graph  $G = (V, E)$ . Let  $S$  be the subset of vertices in  $G$  that are able to reach some cycle. Design an efficient algorithm to compute  $S$ .
7. You are given a directed acyclic graph  $G = (V, E)$ , where each vertex  $v$  that has indegree 0 has a value  $value(v)$  associated with it. For each other vertex  $u$ , define  $Pred(u)$  to be the set of vertices that have incoming edges to  $u$ . We now define  $value(u) = \sum_{v \in Pred(u)} value(v)$ . Design an efficient algorithm to compute  $value(u)$  for all vertices  $u$ . Analyze the running time of your algorithm and prove it correct.