
1

234567

$\mathbb{F}\mathbb{F}_{\omega}^8$

<https://5ht.github.io/bertrand/>

https://llis.nasa.gov/llis_lib/pdf/1009464main1_0641-mr.pdf

<http://www-users.math.umn.edu/~arnold/disasters/ariane5rep.html>

<http://mdbailey.ece.illinois.edu/publications/imc14-heartbleed.pdf>

<https://arxiv.org/pdf/1809.03981.pdf>

<https://www.cs.umd.edu/~aseem/solidetherplas.pdf>

<https://www.dcs.ed.ac.uk/home/mlj/>



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|-------------|
| $\Pi\Sigma$ |
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$$\star : \star \star \leadsto \star \omega, \,), \, .$$

$$\begin{array}{ccc} \lambda \omega & \longrightarrow & \lambda P_\omega \\ \lambda_2 \xrightarrow{\quad} \lambda P_2 & & \downarrow \\ \downarrow \lambda \omega \cdots & \downarrow & \lambda P_\omega \\ \lambda \rightarrow & \longrightarrow & \lambda P_\star \end{array}$$

$$[0,1]$$



1112

1314

<https://groupoid.github.io/languages>
<https://arxiv.org/pdf/1804.07608.pdf>
<https://n2o.dev/ua>
<https://anders.groupoid.space/lib>

12

34

567

<http://mrg.doc.ic.ac.uk/kohei/>

89

1011

$\pi\pi$

$O_{CPS\pi}$
 O_{CPS}^{1213}

<https://5ht.co/ctt.pdf>
<https://5ht.co/cctt.pdf>
<https://web.sfc.keio.ac.jp/~hagino/thesis.pdf>
<https://github.com/devaspot/charity>
<http://nickbenton.name/coqasm.pdf>
<https://www.cl.cam.ac.uk/~mom22/tphols09-lisp.pdf>

$$O_\lambda \lambda O_\pi \pi O_\mu O_\Pi O_\Sigma O_= O_W O_I$$

$$\Pi\Sigma=+\perp\top N\mathcal{U}_i\mathcal{E}\mathcal{I}$$

—
—

—

$\Pi x : \text{op}(o) : \mathcal{U}p : \text{concept}(o)$

$\text{Ob} : \mathcal{U}\text{Hom} : \text{Ob} \rightarrow \text{Ob} \rightarrow \mathcal{U}$

id_o

$A : \text{Ob}(c) \text{isContr}(\Sigma(B : \text{Ob}(C)), A = B)$

$$p:E\rightarrow Bf:Y\rightarrow Xp(f)Bp(f)=id_B$$

$$p:E\rightarrow Bf:Y\rightarrow XEt=p(f)u:Z\rightarrow Yp(f)=p(u)a:Z\rightarrow Yf\circ a=u$$

$$p:E\rightarrow Bf:Y\rightarrow XEt=p(f)u:K\rightarrow JBv:Z\rightarrow XEp(v)=t\circ uw:Z\rightarrow Y\\ Ev=f\circ wp(w)=u$$

$$B^{\rightarrow}Ba:Ob_{\overrightarrow{B}}=Hom_B(x,y)Hom_{\overrightarrow{B}}=[f:Hom_B,g:Hom_B]B$$

$$\frac{f(x)}{g(y)}a$$

$$p:E\rightarrow BBu:J\rightarrow IBX\in p(I)Bf:Y\rightarrow XuXu$$

$$p^{-1}B^{op}\rightarrow CatPsh(B)=[B^{op},Cat]$$

$$p:X\rightarrow Bq:Y\rightarrow BBF:X\rightarrow Yq\circ F=P_X\chi pF(\chi)q$$

$$p:E\rightarrow Bq:D\rightarrow ABFib_B(p,q)Cat/Bp\rightarrow q(H:E\rightarrow D,K:B\rightarrow A)fp\\ H(f)q$$

$$p:E\rightarrow B^{\rightarrow}cod\circ p:E\rightarrow Bf\in EpfB$$

$$Psh(B)Fib(B)$$

$$\int: Psh(B) \stackrel{\cong}{\rightarrow} Fib(B)$$

$\Pi\Sigma$

$$\text{COb}_C \Pi(A, B) \text{Hom}_C(\Pi(A, B), \Pi(A', B')) [f : A \rightarrow A', g(x : A) : B(x) \rightarrow B'(f(x))]$$

```
def CwF : U :=  $\Sigma$  (C: precategory) (T: catfunctor C Fam)
  (context: isContext C) (terminal: isTerminal C), isComprehension C T
```

$$CF : C^{op} \rightarrow \mathbf{SetCSet}$$

$$Ct \in CTy, Tm : C^{op} \rightarrow \mathbf{Set}p : Tm \rightarrow Ty$$

```
def naturalModel : U := Σ (C : precategory) ( _ : isCategory C)
  (t : terminal C) (Tm : carrier C) (Ty : carrier C)
  (p : hom C VT V), Π (f : homTo C V), hasPullback C (Tm, f, Ty, p)
```

$$C-C-$$

$$P,Q : C^{op} \rightarrow \alpha : Q \rightarrow P \alpha Ob(C) x : Ob(C) p_x : D \rightarrow Cy : Q(D)$$

$$\frac{y(D) \rightarrow Q}{y(C) \rightarrow P}$$

$$F : C \rightarrow D \phi_{Ty} : F_!Ty_C \rightarrow Ty_D \phi_{Tm} : F_!Tm_C \rightarrow Tm_D$$

$$\frac{F_!Tm_C \rightarrow Tm_D}{F_!Ty_C \rightarrow Ty_D}$$

$$F_! : C^{op} \rightarrow D^{op}$$

| | |
|----------------------|-----------|
| O_λ | λ |
| O_π | π |
| O_μ | |
| O_Π | |
| O_Σ | |
| $O_=_$ | |
| O_W | |
| O_I | |
| O_{\triangleright} | π |
| $O_{/}$ | |
| O_H | |
| O_{\lrcorner} | |

TmTy

βη

$O_\Pi O_\Sigma O_=_{PTS} O_{MLTT-80} O_{HTS}$

$$O_\Pi \rightarrow O_\Sigma \rightarrow O_=_ \rightarrow O_W \rightarrow O_I.$$

$$\begin{aligned} O_{PTS}(O_\Pi) &\rightarrow O_{MLTT-72}(O_\Pi, O_\Sigma) \rightarrow O_{MLTT-75}(..., O_\Sigma, O_=_) \rightarrow \\ &\rightarrow O_{MLTT-80}(..., O_=_, O_W) \rightarrow O_{HTS}(..., O_W, O_I). \end{aligned}$$

$$\begin{aligned} O_{PTS} &: O_\Pi \rightarrow \mathbb{U} \\ O_{MLTT-72} &: O_\Pi \rightarrow O_\Sigma \rightarrow \mathbb{U} \\ O_{MLTT-75} &: O_\Pi \rightarrow O_\Sigma \rightarrow O_=_ \rightarrow \mathbb{U} \\ O_{MLTT-80} &: O_\Pi \rightarrow O_\Sigma \rightarrow O_=_ \rightarrow O_W \rightarrow \mathbb{U} \\ O_{HTS} &: O_\Pi \rightarrow O_\Sigma \rightarrow O_=_ \rightarrow O_W \rightarrow O_I \rightarrow \mathbb{U} \end{aligned}$$

$$O_{HTS} = O_{\Pi\Sigma=W_I}$$

$$O_{HTS} = O_{\Pi\Sigma=W_I} : O_\Pi \rightarrow O_\Sigma \rightarrow O_=_ \rightarrow O_W \rightarrow O_I \rightarrow \mathbb{U}.$$

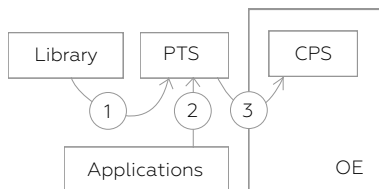
$$O_{MLTT-80}$$

βηβη

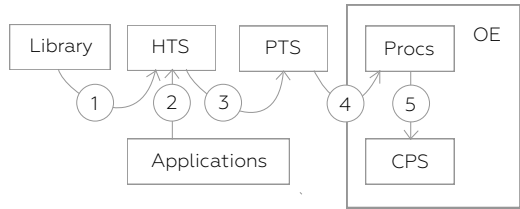
$$O_\infty : O_{CPS} \rightarrow O_{PTS} \rightarrow O_{MLTT-80} \rightarrow O_{HTS} \rightarrow \dots$$

$$O_{I*} \rightarrow O_{\Pi=} O_{\Pi} \rightarrow O_{\Pi\Sigma} O_{\Pi} \rightarrow O_{\Pi\Sigma} O_{\Pi*} \rightarrow O_{\Pi}$$

$$\text{PTS}_{\text{CPS}} = \begin{cases} \text{Ob} : \{\text{O}_{\text{CPS}}, \text{O}_{\text{PTS}}\} \\ \text{Hom} : \{1, 2 : \mathbb{1} \rightarrow \text{O}_{\text{PTS}}, 3 : \text{O}_{\text{PTS}} \rightarrow \text{O}_{\text{CPS}}\} \end{cases}$$



$$\text{Total} = \begin{cases} \text{Ob} : \{O_{\text{CPS}}, O_{\text{PTS}}, O_{\text{MLTT-75}}, O_{\text{MLTT-80}}, O_{\text{HTS}}\} \\ \text{Hom} : \begin{cases} 1, 2 : \mathbb{1} \rightarrow O_{\text{HTS}}, 3 : O_{\text{MLTT-75}} \rightarrow O_{\text{MLTT-80}} \\ 4 : O_{\text{HTS}} \rightarrow O_{\text{MLTT-80}}, 5 : O_{\text{MLTT-80}} \rightarrow O_{\text{PTS}}, 6 : O_{\text{PTS}} \rightarrow O_{\text{CPS}} \end{cases} \end{cases}$$



O_{CPS}

$$O_{CPS} = \begin{cases} Ob : \{\text{maybeCPS}\} \\ Hom : \{\text{eval} : Ob \rightarrow Ob\} \end{cases}$$

$O_\lambda O_\pi O_\mu$

O_{CPS}

```
data CPS = lambda (_: church CPS)
          | process (_: milner CPS)
          | tensor (_: futhark CPS)
```

$\lambda O_{CPS} \text{ maybe}$

$O_{CPS} : O_\lambda \rightarrow O_\pi \rightarrow O_\mu \rightarrow U$

O_λ

```
data church = var (x: nat)
             | lam (l: nat) (d: cps)
             | app (f a: cps)
```

$O_\pi O_*$

```
data milner (lang: U)
  = process (protocol: lang)
  | spawn (cursors: lang) (core: nat) (program: lang)
  | snd (cursor: lang) (data: lang)
  | rcv (cursor: lang)
  | pub (size: nat)
  | sub (cursor: lang)
```

O_μ

```
data futhark (lang: U)
  = iota (cursor: lang)
  | map (cursor: lang)
  | fold (size: nat)
  | scan (cursor: lang)
  | for (cursor: lang)
  | while (cursor: lang)
  | concat (cursor: lang)
  | zip (cursor: lang)
  | transpose (cursor: lang)
```

Π

Σ

$N \rightarrow u$

O_{PTS}

$$O_{PTS} = \begin{cases} \text{Ob} : \{X : \text{maybePTS}, \text{target} : \text{maybeCPS}\} \\ \text{Hom} : \begin{cases} \text{type}, \text{norm} : X \rightarrow X, \text{extract} : X \rightarrow \text{target} \\ \text{certify} : X \rightarrow \text{target} = \text{type} \circ \text{norm} \circ \text{extract} \end{cases} \end{cases}$$

$O_{PTS} O_{PTS} \Pi$

```
data PTS = forall (_: Forall PTS)
```

O_{Π}

```
data Forall (lang: U)
  = fibrant (n: nat)
  | variable (x: name) (l: nat)
  | pi (x: name) (l: nat) (f: lang)
  | lambda (x: name) (l: nat) (f: lang)
  | application (f a: lang)
```

$$\Pi\Sigma\text{MLTT} - 75\Pi\Sigma^2\Pi\forall^3$$

$$O_{\text{MLTT}-75}$$

$$O_{\text{MLTT}-75} = \begin{cases} \text{Ob} : \{\text{maybeMLTT} - 75\} \\ \text{Hom} : \begin{cases} \text{type, norm} : \text{Ob} \rightarrow \text{Ob} \\ \text{certify} : \text{Ob} \rightarrow \text{Ob} = \text{type} \circ \text{norm} \end{cases} \end{cases}$$

$$O_{\text{MLTT}-75}O_{\text{MLTT}-75}O_{\Pi}O_{\Sigma}O_{=}$$

```
data MLTT
  = forall (_: Forall MLTT)
  | sigma (_: Sigma MLTT)
  | id (_: Id MLTT)
```

$$O_{\Sigma}\Sigma O_{\text{MLTT}-72}O_{\Pi\Sigma}$$

```
data Sigma (lang: U)
  = sigma (n: name) (a b: lang)
  | pair (a b: lang)
  | fst (p: lang)
  | snd (p: lang)
```

$$O_{=}O_{\text{MLTT}-75}O_{\Pi\Sigma}=$$

```
data Id (lang: U)
  = identity (t a b: lang)
  | id_intro (a b: lang)
  | id_elim (a b c d e: lang)
  | id_compute (a b c d e: lang)
```

$$O_{=}\eta$$

$O_{\text{MLTT-80}}$

$$O_{\text{MLTT-80}} = \begin{cases} \text{Ob} : \{X : \text{maybePM}, \text{target} : \text{maybeCPS}\} \\ \text{Hom} : \begin{cases} \text{type}, \text{norm}, \text{induction} : X \rightarrow X, \text{extract} : X \rightarrow \text{target} \\ \text{certify} : X \rightarrow \text{target} \\ \text{cerfity} = \text{type} \circ \text{norm} \circ \text{induction} \circ \text{extract} \end{cases} \end{cases}$$

$O = O_{\Sigma} O_{\Pi} O_0 O_1 O_2 O_W$

$O_{\text{MLTT-80}}$

```
data MLTT-80
  = forall (_: Forall MLTT-80)
  | sigma (_: Sigma MLTT-80)
  | id (_: Id MLTT-80)
  | 0 (_: Empty MLTT-80)
  | 1 (_: Unit MLTT-80)
  | 2 (_: Bool MLTT-80)
  | W (_: W MLTT-80)

data W (lang: U) =
  | W_Form (n: name) (a b: lang)
  | W_Sup (a b: lang)
  | W_Ind (a b c: lang)

data 2 (lang: U) =
  | 2_bool | 2_true | 2_false
  | 2_Ind (a: lang)

data 1 (lang: U) =
  | 1_unit | 1_star
  | 1_Ind (a: lang)

data 0 (lang: U) =
  | 0_Ind (a: lang)
```


O_{PM}

$$O_{PM} = \begin{cases} \text{Ob} : \{X : \text{maybePM}, \text{target} : \text{maybeCPS}\} \\ \text{Hom} : \begin{cases} \text{type}, \text{norm}, \text{induction} : X \rightarrow X, \text{extract} : X \rightarrow \text{target} \\ \text{certify} : X \rightarrow \text{target} \\ \text{cerfity} = \text{type} \circ \text{norm} \circ \text{induction} \circ \text{extract} \end{cases} \end{cases}$$

$O = O_{\Sigma} O_{\Pi}$

O_{PM}

```
data PM = forall (_: Forall PM)
  | sigma (_: Sigma PM)
  | id (_: Id PM)
  | inductive (_: InductiveSchemes PM)
```

```
data tele (A: U) = emp | tel (n: name) (b: A) (t: tele A)
data branch (A: U) = br (n: name) (args: list name) (term: A)
data label (A: U) = lab (n: name) (t: tele A)
  | com (n: name) (t: tele A) (dim: list name)
    (s: list (prod (prod name bool) A))
```

O_*O_*

```
data InductiveSchemes (lang: U)
  = data (n: name) (t: tele lang) (labels: list (label lang))
  | case (n: name) (t: lang) (branches: list (branch lang))
  | constructor (n: name) (args: list lang)
```

$$O_{\text{HTS}} = \begin{cases} \text{Ob} : \{\text{maybeHTS}\} \\ \text{Hom} : \begin{cases} \text{type}, \text{norm} : \text{Ob} \rightarrow \text{Ob} \\ \text{certify} : \text{Ob} \rightarrow \text{Ob} = \text{type} \circ \text{norm} \end{cases} \end{cases}$$

$$O_{\text{HTS}} O_I O_W O = O_\Sigma O_\Pi$$

```
data HTS = forall (_: Forall HTS)
  | sigma (_: Sigma HTS)
  | id (_: Id HTS)
  | homotopy (_: Homotopy HTS)
```

$$O_{\text{MLTT-80}}$$

$$O_I$$

```
data Homotopy (lang: U)
  = pretype (n: nat)
  | path (A x y: lang)
  | path_lambda (name: name) (a: lang)
  | path_app (f a: lang)
  | interval | zero | one
  | meet (a b: lang) | join (a b: lang) | neg (e: lang)
  | transp (a b c: lang) | hcomp (a b: lang)
  | glue (a b c: lang) | Glue (a b: lang) | unglue (a b: lang)
```

$$O_{\text{HTS}} O = O_I$$

O_{CPS}

```
data cps = var (x: nat)
         | lam (l: nat) (d: cps)
         | app (f a: cps)
```

CPS


```
objdump ./target/release/o -d | grep mulpd
223f1: c5 f5 59 0c d3    vmulpd (%rbx,%rdx,8),%ymm1,%ymm1
223f6: c5 dd 59 64 d3 20 vmulpd 0x20(%rbx,%rdx,8),%ymm4,%ymm4
22416: c5 f5 59 4c d3 40 vmulpd 0x40(%rbx,%rdx,8),%ymm1,%ymm1
2241c: c5 dd 59 64 d3 60 vmulpd 0x60(%rbx,%rdx,8),%ymm4,%ymm4
2264d: c5 f5 59 0c d3    vmulpd (%rbx,%rdx,8),%ymm1,%ymm1
22652: c5 e5 59 5c d3 20 vmulpd 0x20(%rbx,%rdx,8),%ymm3,%ymm3
```

O_{CPS}

```
data Lazy = Defer (otree: NodeId) (a: AST) (cont: Cont)
  | Continuation (otree: NodeId) (a: AST) (cont: Cont)
  | Return (a: AST)
  | Start

data Cont = Expressions (a: AST) (v: Option (Iter AST)) (c: Cont)
  | Assign (ast: AST) (cont: Cont)
  | Cond (c,d: AST) (cont: Cont)
  | Func (a,b,c: AST) (cont: Cont)
  | List (acc: Vec AST) (vec: Iter AST) (i: Nat) (c: Cont)
  | Call (a: AST) (i: Nat) (cont: Cont)
  | Return
  | Intercore (m: Message) (cont: Cont)
  | Yield (cont: Cont)
```

O_{CPS}

```
E: V | A | C
NC: ";" = [] | ";" m:NL = m
FC: ";" = [] | ";" m:FL = m
EC: ";" = [] | ";" m:EL = m
NL: NAME | o:NAME m:NC = Cons o m
FL: E | o:E | m:FC = Cons o m
EL: E | EC | o:E m:EC = Cons o m
C: N | c:N a:C = Call c a
N: NAME | S | HEX | L | F
L: "(" ")" = [] | "(" [" c:NL "]" m:FL ")" = Table c m
  | "(" "1:EL ")" = List 1
F: "{" "}" = Lambda [] [] []
  | "[" [" c:NL "]" m:EL "]" = Lambda [] c m
  | "{" m:EL "}" = Lambda [] [] m
```

```

data AST      = Atom (a: Scalar)
               | Vector (a: Vec AST)

data Value    = Nil
               | SymbolInt (a: u16)
               | SequenceInt (a: u16)
               | Number (a: i64)
               | Float (a: f64)
               | VecNumber (Vec i64)
               | VecFloat (Vec f64)

data Scalar   = Nil
               | Any
               | List (a: AST)
               | Dict (a: AST)
               | Call (a b: AST)
               | Assign (a b: AST)
               | Cond (a b c: AST)
               | Lambda (otree: Option NodeId) (a b: AST)
               | Yield (c: Context)
               | Value (v: Value)
               | Name (s: String)

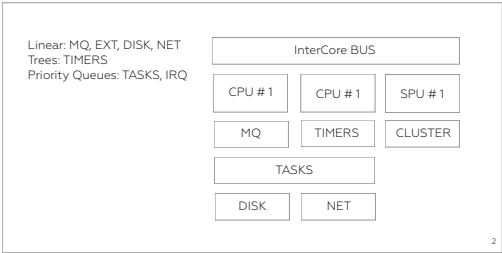
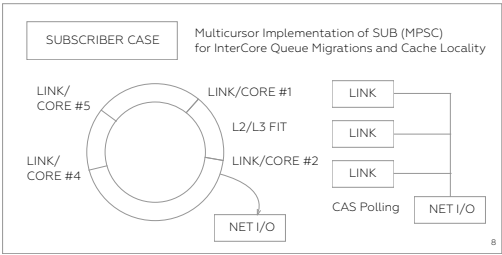
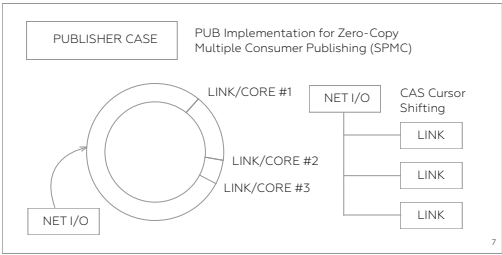
```

1

```

reactor[aux;0;mod[console;network]];
reactor[timercore;1;mod[timer]];
reactor[core1;2;mod[task]];
reactor[core2;3;mod[task]];

```




```
pub [ capacity ]
```

```
sub [ publisher ]
```

```
spawn [ core ; program ; cursors ]
```

```
snd [ writer ; data ]
```

```
rcv [ reader ]
```

O_{CPS}

```
pub struct Publisher<T> {  
    ring: Arc<RingBuffer<T>>,  
    next: Cell<Sequence>,  
    cursors: UncheckedUnsafeArc<Vec<Cursor>>,  
}
```

```
pub struct Subscriber<T> {  
    ring: Arc<RingBuffer<T>>,  
    token: usize,  
    next: Cell<Sequence>,  
    cursors: UncheckedUnsafeArc<Vec<Cursor>>,  
}
```

```

pub struct Channel {
    publisher: Publisher<Message>,
    subscribers: Vec<Subscriber<Message>>,
}

pub struct Memory<'a> {
    publishers: Vec<Publisher<Value<'a>>>,
    subscribers: Vec<Subscriber<Value<'a>>>,
}

pub struct Scheduler<'a> {
    pub tasks: Vec<T3<Job<'a>>>,
    pub bus: Channel,
    pub queues: Memory<'a>,
    pub io: IO,
}

pub enum Message {
    Pub(Pub),
    Sub(Sub),
    Print(String),
    Spawn(Spawn),
    AckSub(AckSub),
    AckPub(AckPub),
    AckSpawn(AckSpawn),
    Exec(usize, String),
    Select(String, u16),
    QoS(u8, u8, u8),
    Halt,
    Nop,
}

```

$$e_i : O_{n+1} \rightarrow O_n O_{CPS} = O_0$$

$$\mathcal{O}_{\text{PTS}}^2 \mathcal{O}_{\text{CCHM}}^3$$

1
2
3
4

<https://github.com/o1>
<https://github.com/o3>
<https://ws.erp.uno>
<https://github.com/o89>



5

6



${}^7F_{\omega}$

```

axiom all : IO (List TableId)
axiom browse : IO String
axiom delete :  $\Pi$  (k: U), TableId -> k -> IO Unit
axiom first :  $\Pi$  (t: TableId) (k: U), IO (Maybe k)
axiom last :  $\Pi$  (k: U), TableId -> IO (Maybe k)
axiom foldl :  $\Pi$  (v acc: U), (v -> acc -> acc) -> acc -> TableId -> IO acc
axiom foldr :  $\Pi$  (v acc: U), (v -> acc -> acc) -> acc -> TableId -> IO acc
axiom insert :  $\Pi$  (v: U), TableId -> v -> IO Boolean
axiom lookup :  $\Pi$  (k v: U), TableId -> k -> IO (List v)
axiom member :  $\Pi$  (k: U), TableId -> k -> IO Boolean
axiom new : Atom -> TableOptions -> IO TableId
axiom next :  $\Pi$  (k: U), TableId -> k -> IO (Maybe k)
axiom prev :  $\Pi$  (k: U), TableId -> k -> IO (Maybe k)
axiom rename : TableId -> Atom -> IO Atom
axiom take :  $\Pi$  (k v: U), TableId -> k -> IO (List v)
axiom match :  $\Pi$  (a v: U), TableId -> a -> IO (List v)
axiom slot :  $\Pi$  (v: U), TableId -> Integer -> IO (Maybe (List v))

```

```

axiom pickle : Binary -> Binary
axiom depickle : Binary -> Binary
axiom encode :  $\Pi$  (k: U), k -> Binary
axiom decode :  $\Pi$  (k: U), Binary -> IO k
axiom reg:  $\Pi$  (k: U), k -> IO k
axiom unreg :  $\Pi$  (k: U), k -> IO k
axiom send :  $\Pi$  (k v z: U), k -> v -> IO z
axiom getSession :  $\Pi$  (k v: U), k -> IO v
axiom putSession :  $\Pi$  (k v: U), k -> v -> IO v
axiom getCache :  $\Pi$  (k v: U), Atom -> k -> IO v
axiom putCache :  $\Pi$  (k v: U), Atom -> k -> v -> IO v

```

```

data PI = PI String Atom Atom Atom Integer RestartType
data Sup = Ok Pid String | Error String

```

```

axiom start : PI -> IO Sup

```



```

axiom get :  $\Pi$  (f k v: U), f -> k -> IO (Maybe v)
axiom put :  $\Pi$  (r: U), r -> IO StoreResult
axiom delete :  $\Pi$  (f k: U), f -> k -> StoreResult
axiom index :  $\Pi$  (f p v r: U), f -> Atom -> v -> List r

```

```

data Reader = Reader Integer Binary ETF String Integer
data Writer = Writer Integer Binary ETF String Integer
data StoreResult = Ok Integer String Binary
                  | Error Integer String Binary

```

```

axiom next : Reader -> IO Reader
axiom prev : Reader -> IO Reader
axiom take : Reader -> IO Reader
axiom drop : Reader -> IO Reader
axiom save : Reader -> IO Reader
axiom append :  $\Pi$  (f r: U), f -> r -> IO StoreResult
axiom remove :  $\Pi$  (f r: U), f -> r -> IO StoreResult

```

```

axiom start : Proc -> IO Sup
axiom stop : String -> IO Sup
axiom next : ProcId -> IO ProcRes
axiom load : ProcId -> IO ProcRes
axiom proc : ProcId -> IO ProcRes
axiom assign : ProcId -> IO ProcRes
axiom persist : ProcId -> Proc -> IO ProcRes
axiom amend :  $\Pi$  (k: U), ProcId -> k -> IO ProcRes
axiom discard :  $\Pi$  (k: U), ProcId -> k -> IO ProcRes
axiom modify :  $\Pi$  (k: U), ProcId -> k -> Atom -> IO ProcRes
axiom event : ProcId -> String -> IO ProcRes
axiom head : ProcId -> IO Hist
axiom hist : ProcId -> IO (List Hist)
axiom step : ProcId -> Atom -> IO (List Hist)
axiom docs : ProcId -> IO (List Tuple)
axiom events : ProcId -> IO (List Tuple)
axiom tasks : ProcId -> IO (List Tuple)
axiom doc : Tuple -> ProcId -> IO (List Tuple)

```

```

data ProcId = String
data Proc = Proc ProcId String
data ProcRes = Ok Integer String Binary
              | Error Integer String Binary

```

```

axiom q :  $\Pi$  (k: U), Atom -> k
axiom qc :  $\Pi$  (k: U), Atom -> k

```

```

axiom jse : Maybe Binary -> Binary
axiom hte : Maybe Binary -> Binary
axiom wire (actions: List Action) : IO (List Action)
axiom render (content: Either Action Element) : Binary
axiom insert_top (dom: Atom) (content: List Element) : IO (List Action)
axiom insert_bottom (dom: Atom) (content: List Element) : IO (List Action)
axiom update (dom: Atom) (content: List Element) : IO (List Action)
axiom clear (dom: Atom) : IO Unit
axiom remove (dom: Atom) : IO Unit

```

```

#textbox { id=userName, body= <<"Anonymous">> },
#panel { id=chatHistory, class=chat_history }

```

```

<input value="Anonymous" id="userName" type="text"/>
<div id="chatHistory" class="chat_history"></div>

```

```

nitro:update(loginButton,
  #button{id = loginButton,
    body = "Login",
    postback = login,
    source = [user, pass]});

```

```

qi('loginButton').outerHTML=<'button id=\"loginButton\"
  type=\"button\">Login</button>'; { var x=qi('loginButton');
  x && x.addEventListener('click',function (event){
    event.preventDefault(); { if (validateSources(['user','pass'])) {
      ws.send(enc(tuple(atom('pickle'),bin('loginButton'),
        bin('b840bca20b3295619d1157105e355880f850bf0223f2f081549dc
        8934ecbcd3653f617bd96cc9b36b2e7a19d2d47fb8f9fbe32d3c768866
        cb9d6d85700416edf47b9b90742b0632c750a4240a62dfc56789e0f5d8
        590f9afdfb31f35fbab1563ec54fcb17a8e3bad463218d6402f1304'),
        [tuple(tuple(string('loginButton'),bin('detail')),[]),
        tuple(atom('user'),querySource('user')),
        tuple(atom('pass'),querySource('pass'))]])); }
      else console.log('Validation Error'); }));};

```

$$\begin{array}{c} 123\Pi\Sigma^\infty \\ 4 \end{array}$$

$$\begin{array}{c} \omega \\ \text{morte}^5\text{cubical}^6\text{caramel}^7 \end{array}$$

$$\left\{ \begin{array}{l} \text{Sorts} = \mathcal{U}.\{i\}, i : \text{Nat} \\ \text{Axioms} = \mathcal{U}.\{i\} : \mathcal{U}.\{\text{inci}\} \\ \text{Rules} = \mathcal{U}.\{i\} \sim \mathcal{U}.\{j\} : \mathcal{U}.\{\text{maxij}\} \end{array} \right.$$

$$\Pi\lambda\beta\eta$$

$$\frac{\pi\lambda\mu}{\infty}$$

$$\frac{8}{\omega}$$

$$\omega$$

```

I := #list #nat
U := * + * . #nat
0 := U + I + ( 0 ) + 0 0 + 0 → 0
    + λ ( I : 0 ) → 0
    + ∀ ( I : 0 ) → 0

```

$$\rightarrow\lambda\forall O_{\text{PTS}}O_{\Pi}$$

```

data pts (lang: U)
= star      (n: nat)
| var      (x: name) (l: nat)
| pi      (x: name) (l: nat) (d c: lang)
| remote (n: name) (n: nat)
| lambda (x: name) (l: nat) (d c: lang)
| app      (f a: lang)

```

$$\infty_{\omega}{}^9\mathsf{Fixpointremote}$$

$$u_0:u_1:u_2:u_3:\ldots$$

$$u_0u_1u_2u_3$$

$$\frac{o:\mathsf{Nat}}{u_o}$$

$$\frac{}{u:u}$$

$$\frac{i : \text{Nat}, j : \text{Nat}, i < j}{U_i : U_j}$$

$$\frac{i : \text{Nat}, j : \text{Nat}}{U_i \rightarrow U_j : U_{\max(i,j)}}$$

$$\frac{i : \text{Nat}}{U_i : U_{i+1}}$$

$$\frac{i : \text{Nat}, j : \text{Nat}}{U_i \rightarrow U_j : U_j}$$

list Sigma

$$\overline{\Gamma : \text{Ctx}}$$

$$\frac{\Gamma : \text{Ctx}}{\emptyset : \Gamma}$$

$$\frac{A : U_i, x : A, \Gamma : \text{Ctx}}{(x : A) \vdash \Gamma : \text{Ctx}}$$

$\beta\eta\text{O}_{\text{PTS}}\beta\eta$

$$\frac{A : U_i, x : A \vdash B : U_j}{\Pi(x : A) \rightarrow B : U_{p(i,j)}}$$

$$\frac{x : A \vdash b : B}{\lambda(x : A) \rightarrow b : \Pi(x : A) \rightarrow B}$$

$$\frac{f : (\Pi(x : A) \rightarrow B) a : A}{fa : B[a/x]}$$

$$\frac{x : A \vdash b : B a : A}{(\lambda(x : A) \rightarrow b) a = b[a/x] : B[a/x]}$$

$$\frac{\pi_1 : Au : A \vdash \pi_2 : B}{[\pi_1 / u] \pi_2 : B}$$

```

PTS (A: U): U
= (Form: (A -> U) -> U)
* (Ctor: (B: A -> U) -> ((a: A) -> B a) -> (Pi A B))
* (Elim: (B: A -> U)(a: A) -> (Pi A B) -> B a)
* (Beta: (B: A -> U)(a: A)(f: Pi A B)-> Equ (B a)(Elim B a (Ctor B f))(f a))
* (Eta: (B: A -> U)(a: A)(f: Pi A B)-> Equ (Pi A B) f (\(x:A) -> f x)) * 1

```

$O = O_I$

remoteremote

```

type (:star,N)      D -> (:star,N+1)
  (:var,N,I)        D -> :true = proplists:is_defined N B, om:keyget N D I
  (:remote,N)       D -> om:cache (type N D)
  (:pi,N,0,I,0)     D -> (:star,h(star(type I D)),star(type 0 [(N,norm I)|D]))
  (:fn,N,0,I,0)     D -> let star (type I D), NI = norm I
                        in (:pi,N,0,NI,type(0,[(N,NI)|D]))
  (:app,F,A)        D -> let T = type(F,D),
                        (:pi,N,0,I,0) = T, :true = eq I (type A D)
                        in norm (subst 0 N A)

```

```

sh (:star,X)      N P -> (:star,X)
  (:var,N,I)      N P -> (:var,N,I+1) when I >= P
                  -> (:var,N,I)
  (:remote,X)     N P -> (:remote,X)
  (:pi,N,0,I,0)   N P -> (:pi,N,0,sh I N P,sh 0 N P+1)
  (:fn,N,0,I,0)   N P -> (:fn,N,0,sh I N P,sh 0 N P+1)
  (:app,L,R)      N P -> (:app,L,R)

```

```

sub (:star,X)      N V L -> (:star,X)
  (:var,N,L)       N V L -> V
  (:var,N,I)       N V L -> (:var,N,I-1) when I > L
  (:remote,X)      N V L -> (:remote,X)
  (:pi,N,0,I,0)    N V L -> (:pi,N,0,sub I N V L,sub 0 N (sh V N 0) L+1)
  (:pi,F,X,I,0)    N V L -> (:pi,F,X,sub I N V L,sub 0 N (sh V F 0) L)
  (:fn,N,0,I,0)    N V L -> (:fn,N,0,sub I N V L,sub 0 N (sh V N 0) L+1)
  (:fn,F,X,I,0)    N V L -> (:fn,F,X,sub I N V L,sub 0 N (sh V F 0) L)
  (:app,F,A)       N V L -> (:app, sub F N V L,sub A N V L)

```

```

norm (:star,X)      → (:star,X)
  (:var,X)          → (:var,X)
  (:remote,N)       → cache (norm N [])
  (:pi,N,0,I,0)     → (:pi,N,0,norm I,norm 0)
  (:fn,N,0,I,0)     → (:fn,N,0,norm I,norm 0)
  (:app,F,A)        → case norm F of
                        (:fn,N,0,I,0) → norm (subst 0 N A)
                        NF           → (:app,NF,norm A) end

```

```

eq (:star,N)      (:star,N)      → true
  (:var,N,I)      (:var,(N,I))   → true
  (:remote,N)     (:remote,N)    → true
  (:pi,N1,0,I1,01) (:pi,N2,0,I2,02) →
    let :true = eq I1 I2
    in eq 01 (subst (shift 02 N1 0) N2 (:var,N1,0) 0)
  (:fn,N1,0,I1,01) (:fn,N2,0,I2,02) →
    let :true = eq I1 I2
    in eq 01 (subst (shift 02 N1 0) N2 (:var,N1,0) 0)
  (:app,F1,A1)    (:app,F2,A2)   → let :true = eq F1 F2 in eq A1 A2
  (A,B)           → (:error,(:eq,A,B))

```

```

> ./om help me
[{a,[expr],"to parse. Returns {_,_} or {error,_}.",
 {type,[term],"typechecks and returns type."},
 {erase,[term],"to untyped term. Returns {_,_}.",
 {norm,[term],"normalize term. Returns term's normal form."},
 {file,[name],"load file as binary."},
 {str,[binary],"lexical tokenizer."},
 {parse,[tokens],"parse given tokens into {_,_} term."},
 {fst,[{x,y}], "returns first element of a pair."},
 {snd,[{x,y}], "returns second element of a pair."},
 {debug,[bool],"enable/disable debug output."},
 {mode,[name],"select metaverse folder."},
 {modes,[],"list all metaverses."}]

> ./om print fst erase norm a "#List/Cons"
  \ Head
-> \ Tail
-> \ Cons
-> \ Nil
-> Cons Head (Tail Cons Nil)
ok

```

PTS_∞λλ

```

ext (:var,X,N,F)      → (:var,X)
  (:app,A,B,N,F)      → (:call,N,ext(F,A,N),[ext(F,B,N)])
  (:fn,S,_,I,O,N,F) → (:fun,N,(:clauses,[{:clause,N,
                                         [(:var,N,S)],[],[ext(F,O,N)]}]))
  _ → []

∞

```

O_{PTS}


```

def := data id tele = sum + id tele : exp = exp +
      id tele : exp where def
exp := cotele*exp + cotele → exp + exp → exp + (exp) + app + id +
      (exp,exp) + \ cotele → exp + split cobrs + exp .1 + exp .2

0 := #empty      imp      := [ import id ]
brs := 0 + cobrs  tele     := 0 + cotele
app := exp exp    cotele := ( exp : exp ) tele
id  := [ #nat ]   sum      := 0 + id tele + id tele | sum
ids := [ id ]     br       := ids → exp
cod := def dec    mod      := module id where imp def
dec := 0 + codec  cobrs    := | br brs

data tele (A: U) = emp | tel (n: name) (b: A) (t: tele A)
data branch (A: U) = br (n: name) (args: list name) (term: A)
data label (A: U) = lab (n: name) (t: tele A)
                  | com (n: name) (t: tele A) (dim: list name)
                    (s: list (prod (prod name bool) A))

data ind (lang: U)
  = datum (n: name) (t: tele lang) (labels: list (label lang))
  | case (n: name) (t: lang) (branches: list (branch lang))
  | ctor (n: name) (args: list lang)

```

$$\begin{aligned}
F_A(X) &= 1 + A \times XF_A(X) = A + X \times X1AF_A(X) = A \times X \\
\mu X &\rightarrow 1 + X \\
\mu X &\rightarrow 1 + A \times X \\
\mu X &\rightarrow 1 + X \times X + X \\
\nu X &\rightarrow A \times X \\
\nu X &\rightarrow 1 + A \times X \\
\mu X &\rightarrow \mu Y \rightarrow 1 + X \times Y = \mu X = \text{List}X
\end{aligned}$$

⋯⋯⋯FixpointfunExthomotopy

$$\begin{aligned}
(\mu L_A, \text{in})L_A(X) &= 1 + (A \times X)\mu L_A = \text{List}(A)\text{nil} : 1 \rightarrow \text{List}(A)\text{cons} : \\
A \times \text{List}(A) &\rightarrow \text{List}(A)\text{nil} = \text{in} \circ \text{inl} \text{cons} = \text{in} \circ \text{inr} \text{in} = [\text{nil}, \text{cons}] \\
c : 1 \rightarrow Ch : A \times C &\rightarrow Cf = \llbracket [c, h] \rrbracket : \text{List}(A) \rightarrow C
\end{aligned}$$

$$\begin{cases} f \circ \text{nil} = c \\ f \circ \text{cons} = h \circ (\text{id} \times f) \end{cases}$$

$f = \text{foldr}(c, h)\mu(1 + A \times X)[1 \rightarrow \text{List}(A), A \times \text{List}(A) \rightarrow \text{List}(A)]$
 O_{PTS}

$$\begin{cases} \text{foldr} = \llbracket [f \circ \text{nil}, h] \rrbracket, f \circ \text{cons} = h \circ (\text{id} \times f) \\ \text{len} = \llbracket [\text{zero}, \lambda n \rightarrow \text{succ}n] \rrbracket \\ (++) = \lambda x \text{sys} \rightarrow \llbracket [\lambda(x) \rightarrow \text{ys}, \text{cons}] \rrbracket (xs) \\ \text{map} = \lambda f \rightarrow \llbracket [\text{nil}, \text{cons} \circ (f \times \text{id})] \rrbracket \end{cases}$$

`data list (A: U) = cons (x: A) (cs: list A) | nil`

$$\begin{cases} \text{list} = \lambda \text{ctor} \rightarrow \lambda \text{cons} \rightarrow \lambda \text{nil} \rightarrow \text{ctor} \\ \text{cons} = \lambda x \rightarrow \lambda xs \rightarrow \lambda \text{list} \rightarrow \lambda \text{cons} \rightarrow \lambda \text{nil} \rightarrow \text{cons}x(\text{xslstconsnil}) \\ \text{nil} = \lambda \text{list} \rightarrow \lambda \text{cons} \rightarrow \lambda \text{nil} \rightarrow \text{nil} \end{cases}$$

`module list where`
`map (A B: U) (f: A -> B) : list A -> list B`
`length (A: U): list A -> nat`
`append (A: U): list A -> list A -> list A`
`foldl (A B: U) (f: B -> A -> B) (Z: B): list A -> B`
`filter (A: U) (p: A -> bool) : list A -> list A`

$$\begin{cases} \text{len} = \text{foldr}(\lambda x n \rightarrow \text{succ}n)0 \\ (++) = \lambda \text{ys} \rightarrow \text{foldrconsys} \\ \text{map} = \lambda f \rightarrow \text{foldr}(\lambda xxs \rightarrow \text{cons}(fx)xs)\text{nil} \\ \text{filter} = \lambda p \rightarrow \text{foldr}(\lambda xxs \rightarrow \text{if } p x \text{ then } \text{cons}xxx \text{ else } \text{sex}x)\text{nil} \\ \text{foldl} = \lambda f v xs = \text{foldr}(\lambda x g \rightarrow (\lambda \rightarrow g(fax)))\text{id}xsv \end{cases}$$

W012

$$\frac{A : \text{Type} \quad x : A \vdash B(x) : \text{Type}}{W(x : A) \rightarrow B(x) : \text{Type}} \quad W$$

$$\frac{a : A \vdash B(a) \rightarrow W}{\text{sup}(a, t) : W} \quad W$$

$$\frac{\begin{array}{l} w : W \vdash C(w) : \text{Type} \\ x : A, u : B(x) \rightarrow W, \\ v : \Pi(y : B(x)) \rightarrow C(u(y)) \vdash c(x, u, v) : C(\text{sup}(x, u)) \end{array}}{w : W \vdash \text{wrec}(w, c) : C(w)}$$

$$\begin{array}{l}
w : W \vdash C(w) : \text{Type} \\
x : A, u : B(x) \rightarrow W \\
v : \Pi(y : B(x)) \rightarrow C(u(y)) \vdash c(x, u, v) : C(\text{sup}(x, u)) \\
\hline
x : A, u : B(x) \rightarrow W \vdash \text{wrec}(\text{sup}(x, u), c) \\
= c(x, u, \lambda(y : B(x)), \text{wrec}(u(y), c)) : C(\text{sup}(x, u))
\end{array}$$

$$I : \Box_n^{\text{op}} \rightarrow \text{Set}$$

```

sys := [ sides ]           side := (id=0)→exp+(id=1)→exp
form := form\f1+f1+f2     sides := #empty+cos+side
cos := side,side+side,cos  mod := module id where impls dec
f1 := f1\f2                f2 := -f2+id+0+1
imp := import id           brs := #empty+cobrs
app := exp exp             tel := #empty+cotel
impls := #list imp         cotel := (exp:exp) tel
id := #list #nat           dec := #empty+codec
u2 := glue+unglue+Glue     u1 := fill+comp
ids := #list id            br := ids→exp+ids@ids→exp
codec := def dec
cobrs := | br brs
sum := #empty+id tel+id tel|sum+id tel<ids>sys
def := data id tel=sum+id tel:exp=exp+id tel:exp where def
exp := cotel*exp+cotel→exp+exp→exp+(exp)+id
      (exp,exp)+\cotele→exp+split cobrs+exp.1+exp.2+
      (ids)exp+exp@form+app+u2 exp exp sys+u1 exp sys

|⟨⟩\→moduleimportdatasplitwherecompfillGlueglueunglue
.1.2,
```

cubical

```
data hts (lang: U)
= pre (n: nat)
| path (A x y: lang)
| plam (name: name) (a: lang)
| papp (f a: lang)
| interval
| zero
| one
| meet (a b: lang)
| join (a b: lang)
| neg (e: lang)
| comp (a b: lang)
| fill (a b c: lang)
| glue (a b c: lang)
| glue-1 (a b: lang)
| unglue-1 (a b: lang)
```

$\mathsf{Henk} \mathcal{U}_i \Pi \tau = \mathsf{typePerAnders} \Sigma \tau \Pi$

ω

$\mathbb{N}\mathsf{Anders}\omega = \{\mathsf{V}_i, \mathcal{U}_i\}$

$\alpha\beta$

$= \mathsf{V}_i \mathcal{U}_i \Pi$

τ

$\tau = \{\mathsf{infer}, \mathsf{app}, \mathsf{check}, \mathsf{act}, \mathsf{conv}, \mathsf{eval}\}$

Σ

$\mathcal{I} = \{\Pi, \Sigma, =, \mathsf{W}, 0, 1, 2, \mathsf{Path}, \mathsf{Glue}\}$

$e_i : \mathcal{O}_{n+1} \rightarrow \mathcal{O}_n \mathcal{O}_{\mathsf{CPS}} = \mathcal{O}_0$
 $\mathcal{O}_{\mathsf{PTS}}^{10} \mathcal{O}_{\mathsf{CCHM}}^{11}$

$\Pi\Sigma=$

123

⁴_∞

PerAndersAndersAnders

\perp
 \top
 $A \vee B$
 $A \wedge B$
 $A \Rightarrow B$
 $x) \quad \exists_{x:A} B(x)$
 $x) \quad \forall_{x:A} B(x)$
 $=_A \quad A^I$

Anders

PerAnders

Anders

$$\Pi\Sigma\mathrm{Id}$$

$$\Pi\Sigma$$

$$\Pi\Sigma$$

$$\Pi$$

$$\Pi\Pi$$

$$\Pi\Pi f : \Pi(x : A), B(x) \rightarrow A \rightarrow \mathcal{U}_i$$

$$\Pi : \mathcal{U} =_{\text{def}} \prod_{x : \mathcal{A}} B(x).$$

$$\text{def Pi } (A : \mathcal{U}) \ (B : A \rightarrow \mathcal{U}) : \mathcal{U} := \Pi(x : A), B(x)$$

$$\Pi \lambda x. b(x) x \mapsto b(x)$$

$$\backslash(x : A) \rightarrow b(x) =_{\text{def}} \prod_{A : \mathcal{U}} \prod_{B : A \rightarrow \mathcal{U}} \prod_{b : \prod_{a : A} B(a)} \lambda x. b(x) : \prod_{y : \mathcal{A}} B(y).$$

$$\begin{aligned} &\text{lambda } (A\ B : \mathcal{U}) \ (b : B) : A \rightarrow B = \backslash(x : A) \rightarrow b \\ &\text{lam } (A : \mathcal{U}) \ (B : A \rightarrow \mathcal{U}) \ (b : (a : A) \rightarrow B\ a) : \text{Pi } A\ B = \backslash(x : A) \rightarrow b\ x \end{aligned}$$

\prod

$$f a =_{\text{def}} \prod_{A:\mathcal{U}} \prod_{B:A \rightarrow \mathcal{U}} \prod_{a:A} \prod_{f:\prod_{x:A} B(x)} f(a) : B(a).$$

```

apply (A B:  $\mathcal{U}$ ) (f: A  $\rightarrow$  B) (a: A) : B = f a
app (A:  $\mathcal{U}$ ) (B: A  $\rightarrow$   $\mathcal{U}$ ) (a: A) (f: Pi A B): B a = f a

```

Π

$$f(a) =_{B(a)} (\lambda(x:A) \rightarrow f(a))(a).$$

```
Beta (A:U) (B:A->U) (a: A) (f: Pi A B)
      : Path (B a) (app A B a (lam A B f)) (f a)
```

Π

$$f =_{(x:A) \rightarrow B(a)} (\lambda(y:A) \rightarrow f(y)).$$

```
Eta (A:U) (B:A->U) (a:A) (f: Pi A B)
      : Path (Pi A B) f \ (x:A) -> f x
```

ΠΣ

$$g : B \rightarrow A \Pi_g : C_{/B} \rightarrow C_{/A}$$

$${\mathbf H}(\infty,1)E \rightarrow B : {\mathbf H}_{/B}{\mathbf H}\Gamma_\Sigma(E)$$

$$\Gamma_\Sigma(E) = \Pi_\Sigma(E) \in {\mathbf H}.$$

```
setFun (A B : U) ( _: isSet B) : isSet (A -> B)
```

```
piIsContr (A: U) (B: A -> U) (u: isContr A)
      (q: (x: A) -> isContr (B x)) : isContr (Pi A B)
```

$$f:A\rightarrow Bg:B\rightarrow Af\circ g:B\stackrel{g}{\rightarrow}A\stackrel{f}{\rightarrow}B$$

ΠΠΠ

$$p:E\rightarrow By:Bx:Ep(x)=y$$

$$F\rightarrow E\stackrel{p}{\rightarrow}BEFB(F,E,p,B)p:E\rightarrow By:BU_bf:p^{-1}(\mathcal{U}_b)\rightarrow\mathcal{U}_b\times F$$

$$p^{-1}(\mathcal{U}_b)\mathop{\times}\limits_{\mathcal{U}_b\leftarrow pr_1}\mathcal{U}_b\times F$$

$$F_{\mathrm{app}}:F\times B\rightarrow E$$

$$F\times B\stackrel{\mathrm{app}}{\longrightarrow}E\stackrel{pr_1}{\rightarrow}B$$

$$pr_1pr_1appSet_{/B}FAA\times B\rightarrow ESet_{/B}A\rightarrow F$$

$$E\Sigma(B,F)p=pr_1(F,\Sigma(B,F),pr_1,B)$$

$$f:(x:A)\rightarrow B(x)ap_f:x=_A y\rightarrow f(x)=_{B(x)}f(y)fcong$$

$$(F, B * F, pr_1, B)y : BF(y)$$

```
FiberPi (B: U) (F: B -> U) (y: B)
      : Path U (fiber (Sigma B F) B (pi1 B F) y) (F y)
```

$$f,g:(x:A)\rightarrow B(x)$$

```
setPi (A: U) (B: A -> U) (h: (x: A) -> isSet (B x)) (f g: Pi A B)
      (p q: Path (Pi A B) f g) : Path (Path (Pi A B) f g) p q
```

$$\Sigma$$

$$\Sigma\Sigma$$

$$\Sigma$$

```
Sigma (A : U) (B : A -> U) : U = (x : A) * B x
```

$$\Sigma$$

```
dpair (A: U) (B: A -> U) (a: A) (b: B a) : Sigma A B = (a,b)
```

$$\Sigma$$

```
pr1 (A: U) (B: A -> U)
      (x: Sigma A B): A = x.1
```

```
pr2 (A: U) (B: A -> U)
      (x: Sigma A B): B (pr1 A B x) = x.2
```

```
sigInd (A: U) (B: A -> U) (C: Sigma A B -> U)
      (g: (a: A) (b: B a) -> C (a, b))
      (p: Sigma A B) : C p = g p.1 p.2
```

$$\Sigma$$

```
Beta1 (A: U) (B: A -> U)
      (a:A) (b: B a)
      : Equ A a (pr1 A B (a,b))
```

```
Beta2 (A: U) (B: A -> U)
      (a: A) (b: B a)
      : Equ (B a) b (pr2 A B (a,b))
```

$$\Sigma$$

```
Eta2 (A: U) (B: A -> U) (p: Sigma A B)
      : Equ (Sigma A B) p (pr1 A B p,pr2 A B p)
```

$$f:A\rightarrow BC\Sigma_f:C/_A\rightarrow C/_B$$

$$x : Ay : BR(x,y)f : A \rightarrow Bx : AR(x,f(x))$$

```
ac (A B: U) (R: A -> B -> U)
: (p: (x:A) -> (y:B)*(R x y)) -> (f:A->B) * ((x:A)->R(x)(f x))
```

```
total (A:U) (B C: A -> U)
(f: (x:A) -> B x -> C x) (w: Sigma A B)
: Sigma A C = (w.1,f (w.1) (w.2))
```

$$\Sigma\Sigma$$

```
setSig (A:U) (B: A -> U) (sA: isSet A)
(sB : (x:A) -> isSet (B x)) : isSet (Sigma A B)
```

$$t,u:\Sigma(A,B)p:t_1=_A u_1)(t_2=_B(p@i) u_2)$$

```
pathSig (A:U) (B : A -> U) (t u : Sigma A B)
: Path U (Path (Sigma A B) t u)
((p: Path A t.1 u.1) * PathP (<i>B(p@i)) t.2 u.2)
```

$$[0,1]^{56789}$$

```
Hetero (A B: U) (a: A) (b: B) (P: Path U A B) : U = PathP P a b
Path (A: U) (a b: A) : U = PathP (<i>A) a b
```

$$[0,1]<i>a\lambda(i:I)\rightarrow a$$

```
refl (A: U) (a: A) : Path A a a
```

```
app1 (A: U) (a b: A) (p: Path A a b): A = p @ 0
app2 (A: U) (a b: A) (p: Path A a b): A = p @ 1
```

$$\lambda(i:I)\rightarrow \overset{\text{comp}}{\underset{\text{def}}{a_i}}$$

<https://5ht.co/cubicaltt.pdf>
<https://5ht.co/ccctt.pdf>
<http://www.cse.chalmers.se/~coquand/mod1.pdf>
<https://www.cs.cmu.edu/~cangiuli/papers/ccctt.pdf>
<https://arxiv.org/pdf/1712.04864.pdf>

```
composition (A: U) (a b c: A) (p: Path A a b) (q: Path A b c)
  : Path A a c = comp (<i>Path A a (q@i)) p []
```

```
inv (A: U) (a b: A) (p: Path A a b): Path A b a = <i> p @ -i
```

```
λ(i,j:I) → p@min(i,j)λ(i,j:I) → p@max(i,j)
```

$$\lambda(i:I) \lambda(j:I) \lambda(p: \text{Path } A \ a \ b) \rightarrow a \quad \lambda(i:I) \rightarrow b$$

```
connection1 (A: U) (a b: A) (p: Path A a b)
  : PathP (<x> Path A (p@x) b) p (<i>b)
  = <y x> p @ (x \ / y)
```

```
connection2 (A: U) (a b: A) (p: Path A a b)
  : PathP (<x> Path A a (p@x)) (<i>a) p
  = <x y> p @ (x /\ y)
```

```
[0,1]λ→
```

```
ap (A B: U) (f: A -> B)
  (a b: A) (p: Path A a b)
  : Path B (f a) (f b)
```

```
apd (A: U) (a x:A) (B: A -> U) (f: A -> B a)
  (b: B a) (p: Path A a x)
  : Path (B a) (f a) (f x)
```

p

```
trans (A B: U) (p: Path U A B) (a: A) : B
```

```
singl (A: U) (a: A): U = (x: A) * Path A a x
```

```
eta (A: U) (a: A): singl A a = (a,refl A a)
```

```
contr (A: U) (a b: A) (p: Path A a b)
  : Path (singl A a) (eta A a) (b,p)
  = <i> (p @ i,<j> p @ i/\j)
```

```

D (A: U) : U = (x y: A) -> Path A x y -> U
J (A: U) (x y: A) (C: D A)
  (d: C x x (refl A x))
  (p: Path A x y) : C x y p
= subst (singl A x) T (eta A x) (y, p) (contr A x y p) d where
  T (z: singl A x) : U = C x (z.1) (z.2)

```

```

J (A: U) (a b: A)
  (P: singl A a -> U)
  (u: P (a,refl A a))
  (p: Path A a b) : P (b,p)

```

```

J (A: U) (a b: A)
  (C: (x: A) -> Path A a x -> U)
  (d: C a (refl A a))
  (p: Path A a b) : C b p

```

```

trans_comp (A: U) (a: A)
  : Path A a (trans A A (<_> A) a)
  = fill (<i> A) a []
subst_comp (A: U) (P: A -> U) (a: A) (e: P a)
  : Path (P a) e (subst A P a a (refl A a) e)
  = trans_comp (P a) e
J_comp (A: U) (a: A) (C: (x: A) -> Path A a x -> U) (d: C a (refl A a))
  : Path (C a (refl A a)) d (J A a C d a (refl A a))
  = subst_comp (singl A a) T (eta A a) d where T (z: singl A a)
  : U = C a (z.1) (z.2)

```

$$U_{n \in \mathbb{N}} U_0$$

$$U_i : U_j, i, j \in \mathbb{N}, j$$

$$U_i \rightarrow U_j : U_{\lambda(i,j), i, j \in \mathbb{N}} \lambda : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$$

$$\lambda_{\max} : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$$

$$\lambda_{\text{snd}} : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$$

$U_i, i \in \mathbb{N} \in$

$U_{n \in \mathbb{N}} U_i : U_j, i, j \in \mathbb{N} U_i \rightarrow U_j : U_{\lambda(i,j), i, j \in \mathbb{N}} \lambda \max \text{snd}$

$\{\{\star, \square\}, \{\star : \square\}, \{i \rightarrow j : j; i, j \in \{\star, \square\}\} \star \square \lambda = \text{snd}$

$\text{PTS}^\infty\{U_{i \in \mathbb{N}}, U_i : U_{j; i < j; i, j \in \mathbb{N}}, U_i \rightarrow U_j : U_{\lambda(i,j); i, j \in \mathbb{N}}\} U_i \text{iiPTS}^{\infty 10}$

$\gamma_0 : \Gamma =_{\text{def}} \star.$

$\Gamma; A =_{\text{def}} \sum_{\gamma: \Gamma} A(\gamma).$

$\Gamma \vdash A =_{\text{def}} \prod_{\gamma: \Gamma} A(\gamma).$

$\Pi\Sigma$

$\Pi\Sigma\Sigma$

```
def MLTT (A: U) : U := Σ
  (Π-form : Π (B: A → U), U)
  (Π-ctor1 : Π (B: A → U), Pi A B → Pi A B)
  (Π-elim1 : Π (B: A → U), Pi A B → Pi A B)
  (Π-comp1 : Π (B: A → U) (a: A) (f: Pi A B),
    Equ (B a) (Π-elim1 B (Π-ctor1 B f) a) (f a))
  (Π-comp2 : Π (B: A → U) (a: A) (f: Pi A B),
    Equ (Pi A B) f (λ (x : A), f x))
  (Σ-form : Π (B: A → U), U)
  (Σ-ctor1 : Π (B: A → U) (a : A) (b : B a), Sigma A B)
  (Σ-elim1 : Π (B: A → U) (p : Sigma A B), A)
  (Σ-elim2 : Π (B: A → U) (p : Sigma A B), B (pr1 A B p))
  (Σ-comp1 : Π (B: A → U) (a : A) (b: B a),
    Equ A a (Σ-elim1 B (Σ-ctor1 B a b)))
  (Σ-comp2 : Π (B: A → U) (a : A) (b: B a),
    Equ (B a) b (Σ-elim2 B (a, b)))
  (Σ-comp3 : Π (B: A → U) (p : Sigma A B),
    Equ (Sigma A B) p (pr1 A B p, pr2 A B p))
  (=-form : Π (a: A), A → U)
  (=-ctor1 : Π (a: A), Equ A a a)
  (=-elim1 : Π (a: A) (C: D A) (d: C a a (=-ctor1 a))
    (y: A) (p: Equ A a y), C a y p)
  (=-comp1 : Π (a: A) (C: D A) (d: C a a (=-ctor1 a)),
    Equ (C a a (=-ctor1 a)) d (=-elim1 a C d a (=-ctor1 a))), U
```

```

theorem instance (A : U) : MLTT A :=
  (Pi A, lambda A, app A, comp1 A, comp2 A,
   Sigma A, pair A, pr1 A, pr2 A, comp3 A, comp4 A, comp5 A,
   Equ A, refl A, J A, comp6 A, A)

```

```
mltt.cttcubicaltt
```

```
$ rllwrap ./anders.native check ./experiments/mltt.anders
```

```
File loaded.
```

```
> :n instance
```

```
TYPE:  $\prod (A : U), \sum (\Pi\text{-form} : \prod (B : (A \rightarrow U)), U), \sum (\Pi\text{-ctor}_1 : \prod (B : (A \rightarrow U)), (\prod (x : A), (B x) \rightarrow \prod (x : A), (B x))), \sum (\Pi\text{-elim}_1 : \prod (B : (A \rightarrow U)), (\prod (x : A), (B x) \rightarrow \prod (x : A), (B x))), \sum (\Pi\text{-comp}_1 : \prod (B : (A \rightarrow U)), \prod (a : A), \prod (f : \prod (x : A), (B x)), \prod (P : ((B a) \rightarrow U)), ((P ((\Pi\text{-elim}_1 B) ((\Pi\text{-ctor}_1 B) f)) a)) \rightarrow (P (f a))))), \sum (\Pi\text{-comp}_2 : \prod (B : (A \rightarrow U)), \prod (a : A), \prod (f : \prod (x : A), (B x)), \prod (P : (\prod (x : A), (B x) \rightarrow U)), ((P f) \rightarrow (P \lambda (x : A), (f x))))), \sum (\Sigma\text{-form} : \prod (B : (A \rightarrow U)), U), \sum (\Sigma\text{-ctor}_1 : \prod (B : (A \rightarrow U)), \prod (a : A), \prod (b : (B a))), \sum (x : A), (B x)), \sum (\Sigma\text{-elim}_1 : \prod (B : (A \rightarrow U)), \prod (p : \sum (x : A), (B x))), A), \sum (\Sigma\text{-elim}_2 : \prod (B : (A \rightarrow U)), \prod (p : \sum (x : A), (B x)), (B p.1)), \sum (\Sigma\text{-comp}_1 : \prod (B : (A \rightarrow U)), \prod (a : A), \prod (b : (B a)), \prod (P : (A \rightarrow U)), ((P a) \rightarrow (P ((\Sigma\text{-elim}_1 B) ((\Sigma\text{-ctor}_1 B) a) b))))), \sum (\Sigma\text{-comp}_2 : \prod (B : (A \rightarrow U)), \prod (a : A), \prod (b : (B a)), \prod (P : ((B a) \rightarrow U)), ((P b) \rightarrow (P ((\Sigma\text{-elim}_2 B) (a, b))))) , \sum (\Sigma\text{-comp}_3 : \prod (B : (A \rightarrow U)), \prod (p : \sum (x : A), (B x)), \prod (P : (\sum (x : A), (B x) \rightarrow U)), ((P p) \rightarrow (P (p.1, p.2))))), \sum (\Sigma\text{-form} : \prod (a : A), (A \rightarrow U)), \sum (\Sigma\text{-ctor}_1 : \prod (a : A), \prod (P : (A \rightarrow U)), ((P a) \rightarrow (P a))), \sum (\Sigma\text{-elim}_1 : \prod (a : A), \prod (C : \prod (x : A), \prod (y : A), (\prod (P : (A \rightarrow U)), ((P x) \rightarrow (P y)) \rightarrow U)), \prod (d : (((C a) a) (\Sigma\text{-ctor}_1 a))), \prod (y : A), \prod (p : \prod (P : (A \rightarrow U)), ((P a) \rightarrow (P y))), (((C a) y) p)), \sum (\Sigma\text{-comp}_1 : \prod (a : A), \prod (C : \prod (x : A), \prod (y : A), (\prod (P : (A \rightarrow U)), ((P x) \rightarrow (P y)) \rightarrow U)), \prod (d : (((C a) a) (\Sigma\text{-ctor}_1 a))), \prod (P : (((C a) a) (\Sigma\text{-ctor}_1 a)) \rightarrow U)), ((P d) \rightarrow (P (((\Sigma\text{-elim}_1 a) C) d) a) (\Sigma\text{-ctor}_1 a))))), U$ 
```

```
NORMEVAL:  $\lambda (A : U), (\lambda (B : (A \rightarrow U)), \prod (x : A), (B x)), (\lambda (B : (A \rightarrow U)), \lambda (b : \prod (x : A), (B x)), \lambda (x : A), (b x)), (\lambda (B : (A \rightarrow U)), \lambda (f : \prod (x : A), (B x)), \lambda (a : A), (f a)), (\lambda (B : (A \rightarrow U)), \lambda (a : A), \lambda (f : \prod (x : A), (B x)), \lambda (P : ((B a) \rightarrow U)), \lambda (u : (P (f a))), u, (\lambda (B : (A \rightarrow U)), \lambda (a : A), \lambda (f : \prod (x : A), (B x)), \lambda (P : (\prod (x : A), (B x) \rightarrow U)), \lambda (u : (P f)), u, (\lambda (B : (A \rightarrow U)), \sum (x : A), (B x)), (\lambda (B : (A \rightarrow U)), \lambda (a : A), \lambda (b : (B a)), (a, b)), (\lambda (B : (A \rightarrow U)), \lambda (x : \sum (x : A), (B x)), x.1, (\lambda (B : (A \rightarrow U)), \lambda (x : \sum (x : A), (B x)), x.2, (\lambda (B : (A \rightarrow U)), \lambda (a : A), \lambda (b : (B a)), \lambda (P : (A \rightarrow U)), \lambda (u : (P a))), u, (\lambda (B : (A \rightarrow U)), \lambda (a : A), \lambda (b : (B a)), \lambda (P : ((B a) \rightarrow U)), \lambda (u : (P b))), u, (\lambda (B : (A \rightarrow U)), \lambda (p : \sum (x : A), (B x)), \lambda (P : (\sum (x : A), (B x) \rightarrow U)), \lambda (u : (P p))), u, (\lambda (x : A), \lambda (y : A), \prod (P : (A \rightarrow U)), ((P x) \rightarrow (P y))), (\lambda (x : A), \lambda (P : (A \rightarrow U)), \lambda (u : (P x)), u, ((J A), ((comp_6 A), A))))))))))$ 
```

```
data empty =
```

```
emptyRec (C: U): empty -> C = split {}
```

```
emptyInd (C: empty -> U): (z: empty) -> C z = split {}
```

```

data unit = star
unitRec (C: U) (x: C): unit -> C = split tt -> x
unitInd (C: unit -> U) (x: C tt): (z: unit) -> C z = split tt -> x

```

```

data bool = false | true
b1: U = bool -> bool
b2: U = bool -> bool -> bool
negation: b1 = split { false -> true; true -> false }
or: b2 = split { false -> idfun bool; true -> lambda bool bool true }
and: b2 = split { false -> lambda bool bool false; true -> idfun boo }
boolEq: b2 = lamb bool (bool -> bool) negation
boolRec (C: U) (f t: C): bool -> C = split { false -> f ; true -> t }
boolInd (C: bool -> U) (f: A false) (t: A true): (n:bool) -> A n
    = split { false -> f ; true -> t }

```

$$M_A(X) = 1 + A$$

```

data maybe (A: U) = nothing | just (x: A)
maybeRec (A P: U) (n: P) (j: A -> P): maybe A -> P
    = split { nothing -> n; just a -> j a }

maybeInd (A: U) (P: maybe A -> U) (n: P nothing)
    (j: (a: A) -> P (just a)): (a: maybe A) -> P a
    = split { nothing -> n ; just x -> j x }

```

```

data either (A B: U) = left (x: A) | right (y: B)
eitherRec (A B C: U) (b: A -> C) (c: B -> C): either A B -> C
    = split { inl x -> b(x) ; inr y -> c(y) }

```

```

eitherInd (A B: U) (C: either A B -> U)
    (x: (a: A) -> C (inl a))
    (y: (b: B) -> C (inr b))
: (x: either A B) -> C x
    = split { inl i -> x i ; inr j -> y j }

```

```

data tuple (A B: U) = pair (x: A) (y: B)
prod (A B: U) (x: A) (y: B): ( _: A ) * B = (x,y)
tupleRec (A B C: U) (c: (x:A) (y:B) -> C): (x: tuple A B) -> C
    = split pair a b -> c a b
tupleInd (A B: U) (C: tuple A B -> U)
    (c: (x:A)(y:B) -> C (pair x y))
: (x: tuple A B) -> C x
    = split pair a b -> c a b

```

```

data nat = zero | succ (n: nat)
natEq: nat -> nat -> bool
natCase (C:U) (a b: C): nat -> C
natRec (C:U) (z: C) (s: nat->C->C) : (n:nat) -> C

natElim (C:nat->U) (z: C zero)
  (s: (n:nat)->C(succ n)): (n:nat) -> C(n)
natInd (C:nat->U) (z: C zero)
  (s: (n:nat)->C(n)->C(succ n)): (n:nat) -> C(n)

```

$(\mu L_A, \text{in})L_A(X) = 1 + (AX)\mu L_A = \text{List}(A)\text{nil} : 1 \rightarrow \text{List}(A)\text{cons} :$
 $A \times \text{List}(A) \rightarrow \text{List}(A)\text{nil} = \text{in} \circ \text{inlcons} = \text{in} \circ \text{inrin} = [\text{nil}, \text{cons}]$

```

data list (A: U) = nil | cons (x:A) (xs: list A)
listCase (A C:U) (a b: C): list A -> C
listRec (A C:U) (z: C) (s: A->list A->C->C): (n:list A) -> C
listElim (A: U) (C:list A->U) (z: C nil)
  (s: (x:A)(xs:list A)->C(cons x xs)): (n:list A) -> C(n)
listInd (A: U) (C:list A->U) (z: C nil)
  (s: (x:A)(xs:list A)->C(xs)->C(cons x xs)): (n:list A) -> C(n)

```

```

null (A:U): list A -> bool
head (A:U): list A -> maybe A
tail (A:U): list A -> maybe (list A)
nth (A:U): nat -> list A -> maybe A
append (A: U): list A -> list A -> list A
reverse (A: U): list A -> list A
map (A B: U): (A -> B) -> list A -> list B
zip (AB: U): list A -> list B -> list (tuple A B)
foldr (AB: U): (A -> B -> B) -> B -> list A -> B
foldl (AB: U): (B -> A -> B) -> B -> list A -> B
switch (A: U): (Unit -> list A) -> bool -> list A
filter (A: U): (A -> bool) -> list A -> list A
length (A: U): list A -> nat
listEq (A: eq): list A.1 -> list A.1 -> bool

```

```

data stream (A: U) = cons (x: A) (xs: stream A)

```

```

data fin (n: nat)
  = fzero | fsucc (_, fin (pred n))

fz (n: nat): fin (succ n) = fzero
fs (n: nat): fin n -> fin (succ n) = \ (x: fin n) -> fsucc x

```

```

data vector (A: U) (n: nat)
  = nil | cons (_, A) (_, vector A (pred n))

data seq (A: U) (B: A -> A -> U) (X Y: A)
  = seqNil (_, A)
  | seqCons (X Y Z: A) (_, B X Y) (_, Seq A B Y Z)

nat = (X:U) -> (X -> X) -> X -> X
      (X -> X) succXzero

list (A: U) = (X:U) -> X -> (A -> X) -> X

NAT (A: U) = (X:U) -> isSet X -> X -> (A -> X) -> X

TRUN (A:U) type = (X: U) -> isProp X -> (A -> X) -> X
S1 = (X:U) -> isGroupoid X -> ((x:X) -> Path X x x) -> X
MONOPL (A:U) = (X:U) -> isSet X -> (A -> X) -> X
NAT = (X:U) -> isSet X -> X -> (A -> X) -> X

upPath      (X Y:U) (f:X->Y) (a:X->X): X -> Y = o X X Y f a
downPath    (X Y:U) (f:X->Y) (b:Y->Y): X -> Y = o X Y Y b f
naturality  (X Y:U) (f:X->Y) (a:X->X) (b:Y->Y): U
  = Path (X->Y) (upPath X Y f a) (downPath X Y f b)

unitEnc': U = (X: U) -> isSet X -> X -> X
isUnitEnc (one: unitEnc'): U
  = (X Y:U) (x:isSet X) (y:isSet Y) (f:X->Y) ->
    naturality X Y f (one X x) (one Y y)

unitEnc: U = (x: unitEnc') * isUnitEnc x
unitEncStar: unitEnc = (\(X:U) (_, isSet X) ->
  idfun X, \ (X Y: U) (_, isSet X) (_, isSet Y) -> refl (X->Y))
unitEncRec (C: U) (s: isSet C) (c: C): unitEnc -> C
  = \ (z: unitEnc) -> z.1 C s c
unitEncBeta (C: U) (s: isSet C) (c: C)
  : Path C (unitEncRec C s c unitEncStar) c = refl C c
unitEncEta (z: unitEnc): Path unitEnc unitEncStar z = undefined
unitEncInd (P: unitEnc -> U) (a: unitEnc): P unitEncStar -> P a
  = subst unitEnc P unitEncStar a (unitEncEta a)
unitEncCondition (n: unitEnc'): isProp (isUnitEnc n)
  = \ (f g: isUnitEnc n) ->
    <h> \ (x y: U) -> \ (X: isSet x) -> \ (Y: isSet y)
    -> \ (F: x -> y) -> <i> \ (R: x) -> Y (F (n x X R)) (n y Y (F R))
    (<j> f x y X Y F @ j R) (<j> g x y X Y F @ j R) @ h @ i

```

$$\frac{\infty}{\quad}$$

$$\mathbb{R}^{n11}$$

$$I=[0,1]$$

```
data I = i0
      | i1
      | seg <i> [(i=0) -> i0,
                (i=1) -> i1]
```

$$i_0,i_1 : x,y : A$$

```
I
pathToHtpy (A: U) (x y: A) (p: Path A x y): I -> A
= split { i0 -> x; i1 -> y; seg @ i -> p @ i }
```

$$f,g:X\rightarrow YH:X\times\rightarrow Y\qquad\qquad\qquad\begin{cases} H(x,0)=f(x),\\ H(x,1)=g(x). \end{cases}$$

```
homotopy (X Y: U) (f g: X -> Y)
  (p: (x: X) -> Path Y (f x) (g x))
  (x: X): I -> Y = pathToHtpy Y (f x) (g x) (p x)
```

$$1213$$

```
cat: U = (A: U) * (A -> A -> U)
groupoid: U = (X: cat) * isCatGroupoid X
PathCat (X: U): cat = (X,\(x y:X)->Path X x y)
```

```

isCatGroupoid (C: cat): U
= (id: (x: C.1) -> C.2 x x)
* (c: (x y z:C.1) -> C.2 x y -> C.2 y z -> C.2 x z)
* (inv: (x y: C.1) -> C.2 x y -> C.2 y x)
* (inv_left: (x y: C.1) (p: C.2 x y) ->
  Path (C.2 x x) (c x y x p (inv x y p)) (id x))
* (inv_right: (x y: C.1) (p: C.2 x y) ->
  Path (C.2 y y) (c y x y (inv x y p) p) (id y))
* (left: (x y: C.1) (f: C.2 x y) ->
  Path (C.2 x y) (c x x y (id x) f) f)
* (right: (x y: C.1) (f: C.2 x y) ->
  Path (C.2 x y) (c x y y f (id y)) f)
* ((x y z w:C.1)(f:C.2 x y)(g:C.2 y z)(h:C.2 z w)->
  Path (C.2 x w) (c x z w (c x y z f g) h)
    (c x y w f (c y z w g h)))

PathGrpd (X: U)
: groupoid
= ((Ob,Hom),id,c,sym X,compPathInv X,compInvPath X,L,R,Q) where
Ob: U = X
Hom (A B: Ob): U = Path X A B
id (A: Ob): Path X A A = refl X A
c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C
= comp (<i> Path X A (g@i)) f []

```

$\Pi\Sigma$

```

funext_form (A B: U) (f g: A -> B): U
= Path (A -> B) f g

funext (A B: U) (f g: A -> B) (p: (x:A) -> Path B (f x) (g x))
: funext_form A B f g
= <i> \ (a: A) -> p a @ i

happly (A B: U) (f g: A -> B) (p: funext_form A B f g) (x: A)
: Path B (f x) (g x)
= cong (A -> B) B (\ (h: A -> B) -> apply A B h x) f g p

funext_Beta (A B: U) (f g: A -> B) (p: (x:A) -> Path B (f x) (g x))
: (x:A) -> Path B (f x) (g x)
= \ (x:A) -> happly A B f g (funext A B f g p) x

funext_Eta (A B: U) (f g: A -> B) (p: Path (A -> B) f g)
: Path (Path (A -> B) f g) (funext A B f g (happly A B f g p)) p
= refl (Path (A -> B) f g) p

```



```

pullback (A B C:U) (f: A -> C) (g: B -> C): U
  = (a: A)
    * (b: B)
      * Path C (f a) (g b)

```

```

pb1 (A B C: U) (f: A -> C) (g: B -> C)
  : pullback A B C f g -> A
  = \ (x: pullback A B C f g) -> x.1

```

```

pb2 (A B C: U) (f: A -> C) (g: B -> C)
  : pullback A B C f g -> B
  = \ (x: pullback A B C f g) -> x.2.1

```

```

pb3 (A B C: U) (f: A -> C) (g: B -> C)
  : (x: pullback A B C f g) -> Path C (f x.1) (g x.2.1)
  = \ (x: pullback A B C f g) -> x.2.2

```

```

kernel (A B: U) (f: A -> B): U
  = pullback A A B f f

```

```

hofiber (A B: U) (f: A -> B) (y: B): U
  = pullback A unit B f (\ (x: unit) -> y)

```

```

pullbackSq (Z A B C: U) (f: A -> C) (g: B -> C) (z1: Z -> A) (z2: Z -> B): U
  = (h: (z:Z) -> Path C ((o Z A C f z1) z) (((o Z B C g z2)) z))
    * isEquiv Z (pullback A B C f g) (induced Z A B C f g z1 z2 h)

```

```

completePullback (A B C: U) (f: A -> C) (g: B -> C)
  : pullbackSq (pullback A B C f g) A B C f g (pb1 A B C f g) (pb2 A B C f g)
  )

```

```

data pushout (A B C: U) (f: C -> A) (g: C -> B)
  = po1 (_: A)
    | po2 (_: B)
    | po3 (c: C) <i> [ (i = 0) -> po1 (f c) ,
                      (i = 1) -> po2 (g c) ]

```

```
isFBundle1 (B: U) (p: B → U) (F: U): U
  = (λ (b: B) → isContr (Path U (p b) F))
  * (λ (x: Sigma B p) → B)
```

```
isFBundle2 (B: U) (p: B → U) (F: U): U
  = (V: U)
  * (v: surjective V B)
  * (λ (x: V) → Path U (p (v.1 x)) F)
```

```
im1 (A B: U) (f: A → B): U = (b: B) * pTrunc ((λ (a: A) * Path B (f a) b)
BAut (F: U): U = im1 unit U (λ (x: unit) → F)
unitIm1 (A B: U) (f: A → B): im1 A B f → B = λ (x: im1 A B f) → x.1
unitBAut (F: U): BAut F → U = unitIm1 unit U (λ (x: unit) → F)
```

```
isFBundle3 (E B: U) (p: E → B) (F: U): U
  = (X: B → BAut F)
  * (classify B (BAut F) (λ (b: B) → fiber E B p b) (unitBAut F) X) where
  classify (A' A: U) (E': A' → U) (E: A → U) (f: A' → A): U
    = (x: A') → Path U (E'(x)) (E(f(x)))
```

```
isFBundle4 (E B: U) (p: E → B) (F: U): U
  = (V: U)
  * (v: surjective V B)
  * (v': prod V F → E)
  * pullbackSq (prod V F) E V B p v.1 v' (λ (x: prod V F) → x.1)
```

```
fiber (A B: U) (f: A → B) (y: B): U = (x: A) * Path B y (f x)
isSingleton (X: U): U = (c: X) * ((x: X) → Path X c x)
isEquiv (A B: U) (f: A → B): U = (y: B) → isContr (fiber A B f y)
equiv (A B: U): U = (f: A → B) * isEquiv A B f
```

```
isSurjective (A B: U) (f: A → B): U
  = (b: B) * pTrunc (fiber A B f b)
```

```
surjective (A B: U): U
  = (f: A → B)
  * isSurjective A B f
```

```
isInjective' (A B: U) (f: A → B): U
  = (b: B) → isProp (fiber A B f b)
```

```
injective (A B: U): U
  = (f: A → B)
  * isInjective A B f
```

```

isEmbedding (A B: U) (f: A -> B) : U
  = (x y: A) -> isEquiv (Path A x y) (Path B (f x) (f y)) (cong A B f x y)

```

```

embedding (A B: U): U
  = (f: A -> B)
  * isEmbedding A B f

```

```

isHae (A B: U) (f: A -> B): U
  = (g: B -> A)
  * (eta_: Path (id A) (o A B A g f) (idfun A))
  * (eps_: Path (id B) (o B A B f g) (idfun B))
  * ((x: A) -> Path B (f ((eta_ @ 0) x)) ((eps_ @ 0) (f x)))

```

```

hae (A B: U): U
  = (f: A -> B)
  * isHae A B f

```

```

iso_Form (A B: U): U = isIso A B -> Path U A B

```

```

iso_Intro (A B: U): iso_Form A B

```

```

iso_Elim (A B: U): Path U A B -> isIso A B

```

```

iso_Comp (A B : U) (p : Path U A B)
  : Path (Path U A B) (iso_Intro A B (iso_Elim A B p)) p

```

```

iso_Uniq (A B : U) (p: isIso A B)
  : Path (isIso A B) (iso_Elim A B (iso_Intro A B p)) p

```

```

univ_Formation (A B: U): U = equiv A B -> Path U A B

```

```

equivToPath (A B: U): univ_Formation A B
  = \ (p: equiv A B) -> <i> Glue B [(i=0) -> (A,p),
    (i=1) -> (B, subst U (equiv B) B B (<_>B) (idEquiv B))] ]

```

```

pathToEquiv (A B : U) (p : Path U A B) : equiv A B
= subst U (equiv A) A B p (idEquiv A)

```

```

eqToEq (A B : U) (p : Path U A B)
: Path (Path U A B) (equivToPath A B (pathToEquiv A B p)) p
= <j i> let Ai: U = p@i in Glue B
  [ (i=0) -> (A,pathToEquiv A B p),
    (i=1) -> (B,pathToEquiv B B (<k> B)),
    (j=1) -> (p@i,pathToEquiv Ai B (<k> p @ (i \ / k))) ]

```

```

transPathFun (A B : U) (w: equiv A B)
: Path (A -> B) w.1 (pathToEquiv A B (equivToPath A B w)).1

```

```

data I = i0
      | i1
      | seg <i> [(i=0) -> i0,
                (i=1) -> i1]

```

$I_{i0}, I_{i1} : Ix, y : A$

```

data S1
= base
| loop <i> [ (i = 0) -> base,
            (i = 1) -> base ]

```

```

data S2
= point
| surf <i j> [ (i = 0) -> point, (i = 1) -> point,
              (j = 0) -> point, (j = 1) -> point ]
              (j = 0) -> point, (j = 1) -> point ]

```

```

data susp (A: U)
= north
| south
| merid (a: A) <i> [ (i = 0) -> north ,
                    (i = 1) -> south ]

data pTrunc (A: U) -- (-1)-trunc, mere proposition truncation
= pinc (a: A)
| pline (x y: pTrunc A) <i>
  [ (i = 0) -> x,
    (i = 1) -> y ]

data sTrunc (A: U) -- (0)-trunc, set truncation
= sinc (a: A)
| sline (a b: sTrunc A)
  (p q: Path (sTrunc A) a b) &lt;i>i j&gt;;
  [ (i = 0) -> p @ j,
    (i = 1) -> q @ j,
    (j = 0) -> a,
    (j = 1) -> b ]

data gTrunc (A: U) -- (1)-trunc, groupoid truncation
= ginc (a: A)
| gline (a b: gTrunc A)
  (p q: Path (gTrunc A) a b)
  (r s: Path (Path (gTrunc A) a b) p q) &lt;i>i j k&gt;;
  [ (i = 0) -> r @ j @ k,
    (i = 1) -> s @ j @ k,
    (j = 0) -> p @ k,
    (j = 1) -> q @ k,
    (k = 0) -> a,
    (k = 1) -> b ]

data quot (A: U) (R: A -> A -> U)
= inj (a: A)
| quoteq (a b: A) (r: R a b) &lt;i>i&gt;;
  [ (i = 0) -> inj a,
    (i = 1) -> inj b ]

data setquot (A: U) (R: A -> A -> U)
= quotient (a: A)
| identification (a b: A) (r: R a b) &lt;i>i&gt;;
  [ (i = 0) -> quotient a,
    (i = 1) -> quotient b ]

```

```

| setTruncation (a b: setquot A R)
  (p q: Path (setquot A R) a b) &lt;i j>;
[ (i = 0) -> p @ j,
  (i = 1) -> q @ j,
  (j = 0) -> a,
  (j = 1) -> b ]

```

```
storage: U -> U = list
```

Σ

```

process : U
= (protocol state: U)
* (current: prod protocol state)
* (act: id (prod protocol state))
* (storage (prod protocol state))

spawn (protocol state: U) (init: prod protocol state)
  (action: id (prod protocol state)) : process
= (protocol,state,init,action,nil)

```

```

protocol (p: process): U = p.1
state (p: process): U = p.2.1
signature (p: process): U = prod p.1 p.2.1
current (p: process): signature p = p.2.2.1
action (p: process): id (signature p) = p.2.2.2.1
trace (p: process): storage (signature p) = p.2.2.2.2

```

$$P \times S \rightarrow SP \times S \rightarrow P \times S$$

```
receive (p: process) : protocol p = axiom
```

```
send (p: process) (message: protocol p) : unit = axiom
```

```

execute (p: process) (message: protocol p) : process
= let step: signature p = (action p) (message, (current p).2)
  in (protocol p, state p, step, action p, cons step (trace p))

```

14
15
16

$$\Sigma_{A:\mathcal{U}} A \rightarrow A \rightarrow \mathcal{U}\mathcal{U}\text{pr}_1\text{Obpr}_2\text{Hom}(a,b)a,b:\text{Ob}$$

$$\text{cat}: \mathcal{U} = (A: \mathcal{U}) * (A \rightarrow A \rightarrow \mathcal{U})$$

$$C\text{Hom}_C(a,b)a,b:\text{Ob}_C\text{idHom}_C(x,x)$$

$$\begin{aligned} C\text{Ob}_C a,b &: \text{Ob}_C \text{Hom}_C(a,b) a : \text{Ob}_C 1_a : \text{Hom}_C(a,a) a,b,c : \\ \text{Ob}_C \text{Hom}_C(b,c) &\rightarrow \text{Hom}_C(a,b) \rightarrow \text{Hom}_C(a,c) g \circ f a,b : \text{Ob}_C f : \\ \text{Hom}_C(a,b) f &= 1_b \circ f f = f \circ 1_a a,b,c,d : A f : \text{Hom}_C(a,b) g : \text{Hom}_C(b,c) \\ h : \text{Hom}_C(c,d) h \circ (g \circ f) &= (h \circ g) \circ f \end{aligned}$$

$$a,b:\text{ObHom}_C(a,b)$$

$$\begin{aligned} \text{isPrecategory } (C: \text{cat}): \mathcal{U} \\ &= (\text{id}: (x: C.1) \rightarrow C.2 \ x \ x) \\ &* (c: (x \ y \ z: C.1) \rightarrow C.2 \ x \ y \rightarrow C.2 \ y \ z \rightarrow C.2 \ x \ z) \\ &* (\text{homSet}: (x \ y: C.1) \rightarrow \text{isSet } (C.2 \ x \ y)) \\ &* (\text{left}: (x \ y: C.1) \rightarrow (f: C.2 \ x \ y) \\ &\rightarrow \text{Path } (C.2 \ x \ y) \ (c \ x \ x \ y \ (\text{id } x) \ f) \ f) \\ &* (\text{right}: (x \ y: C.1) \rightarrow (f: C.2 \ x \ y) \\ &\rightarrow \text{Path } (C.2 \ x \ y) \ (c \ x \ y \ y \ f \ (\text{id } y)) \ f) \\ &* ((x \ y \ z \ w: C.1) (f: C.2 \ x \ y) (g: C.2 \ y \ z) \\ &\quad (h: C.2 \ z \ w) \rightarrow \text{Path } (C.2 \ x \ w) \\ &\quad (c \ x \ z \ w \ (c \ x \ y \ z \ f \ g) \ h) \ (c \ x \ y \ w \ f \ (c \ y \ z \ w \ g \ h))) \end{aligned}$$

ObHom

```

carrier (C: precategory): U = C.1.1
hom      (C: precategory) (a b: carrier C): U = C.1.2 a b
path     (C: precategory) (x: carrier C): hom C x x = C.2.1 x
compose  (C: precategory) (x y z: carrier C)
          (f: hom C x y) (g: hom C y z): hom C x z = C.2.2.1 x y z f g

```

$\text{Ob}_C \Pi_{x,y:\text{Ob}_C} \text{isContr}(\text{Hom}_C(x, y))$

$\text{Ob}_C \Pi_{x,y:\text{Ob}_C} \text{isContr}(\text{Hom}_C(y, x))$

```

isInitial (C: precategory) (x: carrier C): U
  = (y: carrier C) -> isContr (hom C x y)
isTerminal (C: precategory) (y: carrier C): U
  = (x: carrier C) -> isContr (hom C x y)
initial (C: precategory): U
  = (x: carrier C) * isInitial C x
terminal (C: precategory): U
  = (y: carrier C) * isTerminal C y

```

$\text{ABF} : A \rightarrow \text{BF}_{\text{Ob}} : \text{Ob}_h A \rightarrow \text{Ob}_B a, b : \text{Ob}_A \text{FHom} : \text{Hom}_A(a, b) \rightarrow$
 $\text{Hom}_B(\text{F}_{\text{Ob}}(a), \text{F}_{\text{Ob}}(b))a : \text{Ob}_A \text{F}_{\text{Ob}}(1_a) = 1_{\text{F}_{\text{Ob}}(a)} a, b, c : \text{Ob}_A f :$
 $\text{Hom}_A(a, b)g : \text{Hom}_A(b, c)\text{F}(g \circ f) = \text{F}_{\text{Hom}}(g) \circ \text{F}_{\text{Hom}}(f)$

```

catfunctor (A B: precategory): U
  = (ob: carrier A -> carrier B)
  * (mor: (x y: carrier A) -> hom A x y -> hom B (ob x) (ob y))
  * (id: (x: carrier A) -> Path (hom B (ob x) (ob x))
      (mor x x (path A x)) (path B (ob x)))
  * ((x y z: carrier A) -> (f: hom A x y) -> (g: hom A y z) ->
      Path (hom B (ob x) (ob z)) (mor x z (compose A x y z f g))
      (compose B (ob x) (ob y) (ob z) (mor x y f) (mor y z g)))

```

$F, G : C \rightarrow D \gamma : F \rightarrow Gx : C \gamma_a : \text{Hom}_D(F(x), G(x))x, y : Cf : \text{Hom}_C(x, y)$
 $G(f) \circ \gamma_x = \gamma_y \circ F(g)$

```

isNaturalTrans (C D: precategory)
  (F G: catfunctor C D)
  (eta: (x: carrier C) -> hom D (F.1 x) (G.1 x)): U
  = (x y: carrier C) (h: hom C x y) ->
    Path (hom D (F.1 x) (G.1 y))
      (compose D (F.1 x) (F.1 y) (G.1 y) (F.2.1 x y h) (eta y))
      (compose D (F.1 x) (G.1 x) (G.1 y) (eta x) (G.2.1 x y h))

```

```

ntrans (C D: precategory) (F G: catfunctor C D): U
  = (eta: (x: carrier C) -> hom D (F.1 x) (G.1 x))
  * (isNaturalTrans C D F G eta)

```

```

extension (C C' D: precategory)
  (K: catfunctor C C') (G: catfunctor C D) : U
= (F: catfunctor C' D)
* (ntrans C D (compFunctor C C' D K F) G)

```

$f : \text{Hom}_A(a, b) g : \text{Hom}_A(b, a) l_a =_{\eta} g \circ \text{ff} \circ g =_{\epsilon} l_b = g a, b : A$
 $a = b \rightarrow \text{iso}_A(a, b)$

```

iso (C: precategory) (A B: carrier C): U
= (f: hom C A B)
* (g: hom C B A)
* (eta: Path (hom C A A) (compose C A B A f g) (path C A))
* (Path (hom C B B) (compose C B A B g f) (path C B))

```

$a : \text{Ob}_C \Pi_{A: \text{Ob}_C} \text{isContr} \Sigma_{B: \text{Ob}_C} \text{iso}_C(A, B)$

```

isCategory (C: precategory): U
= (A: carrier C) -> isContr ((B: carrier C) * iso C A B)
category: U = (C: precategory) * isCategory C

```

```

Product    (X Y: precategory) : precategory
Coproduct  (X Y: precategory) : precategory

```

CC^{op}

```

opCat (P: precategory): precategory

```

```

sliceCat (C D: precategory)
  (a: carrier (opCat C))
  (F: catfunctor D (opCat C))
: precategory
= cosliceCat (opCat C) D a F

```

```
cosliceCat (C D: precategory)
  (a: carrier C)
  (F: catfunctor D C) : precategory
```

```
initArr (C D: precategory)
  (a: carrier C)
  (F: catfunctor D C): U = initial (cosliceCat C D a F)
```

```
termArr (C D: precategory)
  (a: carrier (opCat C))
  (F: catfunctor D (opCat C)): U = terminal (sliceCat C D a F)
```

$$\text{Ob} = \top \text{Hom} = \top$$

```
unitCat: precategory
```

```
Set: precategory = ((Ob,Hom),id,c,HomSet,L,R,Q) where
  Ob: U = SET
  Hom (A B: Ob): U = A.1 -> B.1
  id (A: Ob): Hom A A = idfun A.1
  c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C
    = o A.1 B.1 C.1 g f
  HomSet (A B: Ob): isSet (Hom A B) = setFun A.1 B.1 B.2
  L (A B: Ob) (f: Hom A B): Path (Hom A B) (c A A B (id A) f) f
    = refl (Hom A B) f
  R (A B: Ob) (f: Hom A B): Path (Hom A B) (c A B B f (id B)) f
    = refl (Hom A B) f
  Q (A B C D: Ob) (f: Hom A B) (g: Hom B C) (h: Hom C D)
    : Path (Hom A D) (c A C D (c A B C f g) h) (c A B D f (c B C D g h))
    = refl (Hom A D) (c A B D f (c B C D g h))
```

```
Functions (X Y: U) (Z: isSet Y): precategory
  = ((Ob,Hom),id,c,HomSet,L,R,Q) where
  Ob: U = X -> Y
  Hom (A B: Ob): U = id (X -> Y)
  id (A: Ob): Hom A A = idfun (X -> Y)
  c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C = idfun (X -> Y)
```

```

HomSet (A B: Ob): isSet (Hom A B) = setFun Ob Ob (setFun X Y Z)
L (A B: Ob) (f: Hom A B): Path (Hom A B) (c A A B (id A) f) f = axiom
R (A B: Ob) (f: Hom A B): Path (Hom A B) (c A B B f (id B)) f = axiom
Q (A B C D: Ob) (f: Hom A B) (g: Hom B C) (h: Hom C D)
  : Path (Hom A D) (c A C D (c A B C f g) h)
    (c A B D f (c B C D g h)) = axiom

Cat: precategory = ((Ob,Hom),id,c,HomSet,L,R,Q) where
  Ob: U = precategory
  Hom (A B: Ob): U = catfunctor A B
  id (A: Ob): catfunctor A A = idFunctor A
  c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C
    = compFunctor A B C f g
  HomSet (A B: Ob): isSet (Hom A B) = axiom
  L (A B: Ob) (f: Hom A B): Path (Hom A B) (c A A B (id A) f) f = axiom
  R (A B: Ob) (f: Hom A B): Path (Hom A B) (c A B B f (id B)) f = axiom
  Q (A B C D: Ob) (f: Hom A B) (g: Hom B C) (h: Hom C D)
    : Path (Hom A D) (c A C D (c A B C f g) h)
      (c A B D f (c B C D g h)) = axiom

Func (X Y: precategory): precategory
  = ((Ob,Hom),id,c,HomSet,L,R,Q) where
  Ob: U = catfunctor X Y
  Hom (A B: Ob): U = ntrans X Y A B
  id (A: Ob): ntrans X Y A A = axiom
  c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C = axiom
  HomSet (A B: Ob): isSet (Hom A B) = axiom
  L (A B: Ob) (f: Hom A B): Path (Hom A B) (c A A B (id A) f) f = axiom
  R (A B: Ob) (f: Hom A B): Path (Hom A B) (c A B B f (id B)) f = axiom
  Q (A B C D: Ob) (f: Hom A B) (g: Hom B C) (h: Hom C D)
    : Path (Hom A D) (c A C D (c A B C f g) h)
      (c A B D f (c B C D g h)) = axiom

```

$k(k-1)01$

```

equiv: U
functor (C D: cat): U
ntrans (C D: cat) (F G: functor C D): U
modification (C D: cat) (F G: functor C D) (I J: ntrans C D F G): U

```

```

Cat2 : U
= (Ob: U)
* (Hom: (A B: Ob) -> U)
* (Hom2: (A B: Ob) -> (C F: Hom A B) -> U)
* (id: (A: Ob) -> Hom A A)
* (id2: (A: Ob) -> (B: Hom A A) -> Hom2 A A B B)
* (c: (A B C: Ob) (f: Hom A B) (g: Hom B C) -> Hom A C)
* (c2: (A B: Ob) (X Y Z: Hom A B)
  (f: Hom2 A B X Y) (g: Hom2 A B Y Z) -> Hom2 A B X Z)

```

$$A \xrightarrow{f} C \xleftarrow{g} BA \times_C B \text{pb}_1 : \times_C \rightarrow \text{Apb}_2 : \times_C \rightarrow B$$

$$A \times_C B \xrightarrow{\text{pb}_2} B$$

$$A \xrightarrow{\text{pb}_1} C$$

$$(\times_C, \text{pb}_1, \text{pb}_2)(D, q_1, q_2)u : D \rightarrow \times_C \text{pb}_1 \circ u = q_1 \text{pb}_2 \circ q_2$$

```

homTo (C: precategory) (X: carrier C): U
= (Y: carrier C) * hom C Y X
cospan (C: precategory): U
= (X: carrier C) * ( _: homTo C X) * homTo C X
cospanCone (C: precategory) (D: cospan C): U
= (W: carrier C) * hasCospanCone C D W
cospanConeHom (C: precategory) (D: cospan C)
  (E1 E2: cospanCone C D) : U
= (h: hom C E1.1 E2.1) * isCospanConeHom C D E1 E2 h
isPullback (C: precategory) (D: cospan C) (E: cospanCone C D) : U
= (h: cospanCone C D) -> isContr (cospanConeHom C D h E)
hasPullback (C: precategory) (D: cospan C) : U
= (E: cospanCone C D) * isPullback C D E

```

$$ABF : A \rightarrow BF_{Ob} : Ob_h A \rightarrow Ob_B a, b : Ob_A F_{Hom} : Hom_A(a, b) \rightarrow Hom_B(F_{Ob}(a), F_{Ob}(b))a : Ob_A F_{Ob}(1_a) = 1_{F_{Ob}}(a)a, b, c : Ob_A f : Hom_A(a, b)g : Hom_A(b, c)F(g \circ f) = F_{Hom}(g) \circ F_{Hom}(f)$$

```

catfunctor (A B: precategory): U
= (ob: carrier A -> carrier B)
* (mor: (x y:carrier A)->hom A x y->hom B(ob x)(ob y))
* (id: (x: carrier A) -> Path (hom B (ob x) (ob x))
  (mor x x (path A x)) (path B (ob x)))

```

```

* ((x y z: carrier A) -> (f: hom A x y) -> (g: hom A y z) ->
  Path (hom B (ob x) (ob z)) (mor x z (compose A x y z f g))
    (compose B (ob x) (ob y) (ob z) (mor x y f) (mor y z g)))

```

Ob_C

$$\prod_{x,y:\text{Ob}_C} \text{isContr}(\text{Hom}_C(y,x)).$$

```

isTerminal (C: precategory) (y: carrier C): U
= (x: carrier C) -> isContr (hom C x y)
terminal (C: precategory): U
= (y: carrier C) * isTerminal C y

```

∞

$\text{PROP}_{x,y} : P_x = y$

$$\text{isProp}(P) = \prod_{x,y:P} (x = y).$$

$x,y : \mathcal{A}p, q : x =_{\mathcal{A}} yp = q$

$x,y : \mathcal{A}p, q : x =_{\mathcal{A}} yr, s : p =_{=\mathcal{A}} qr = s$

$\text{SETSET}_{x,y} : \mathcal{A}p, q : x = yp = q$

$$\text{isSet}(A) = \prod_{x,y:A} \prod_{p,q:x=y} (p = q).$$

```

data N = Z | S (n: N)

```

```

n_grpd (A: U) (n: N): U = (a b: A) -> rec A a b n where
  rec (A: U) (a b: A) : (k: N) -> U
    = split { Z -> Path A a b ; S n -> n_grpd (Path A a b) n }

```

```

isContr (A: U): U = (x: A) * ((y: A) -> Path A x y)
isProp  (A: U): U = n_grpd A Z
isSet   (A: U): U = n_grpd A (S Z)
PROP    : U = (X:U) * isProp X
SET     : U = (X:U) * isSet X

```

\prod

```

setPi (A: U) (B: A -> U) (h: (x: A) -> isSet (B x)) (f g: Pi A B)
  (p q: Path (Pi A B) f g)
  : Path (Path (Pi A B) f g) p q

```

$\Sigma\Sigma$

```

setSig (A:U) (B: A -> U) (base: isSet A)
  (fiber: (x:A) -> isSet (B x)) : isSet (Sigma A B)

```

```

data unit = tt
unitRec (C: U) (x: C): unit -> C = split tt -> x
unitInd (C: unit -> U) (x: C tt): (z:unit) -> C z
    = split tt -> x

```

SetHomΠ

```

Set: precategory = ((Ob,Hom),id,c,HomSet,L,R,Q) where
  Ob: U = SET
  Hom (A B: Ob): U = A.1 -> B.1
  id (A: Ob): Hom A A = idfun A.1
  c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C
    = o A.1 B.1 C.1 g f
  HomSet (A B: Ob): isSet (Hom A B) = setFun A.1 B.1 B.2
  L (A B:Ob) (f:Hom A B): Path (Hom A B)(c A A B (id A)f)f
    = refl (Hom A B) f
  R (A B:Ob) (f:Hom A B): Path (Hom A B)(c A B B f(id B))f
    = refl (Hom A B) f
  Q (A B C D: Ob) (f:Hom A B) (g:Hom B C) (h:Hom C D)
    : Path (Hom A D) (c A C D (c A B C f g) h)
      (c A B D f (c B C D g h))
    = refl (Hom A D) (c A B D f (c B C D g h))

```

AS ∈ ASSSSSS

```

Structure topology (A : Type) := {
  open :> (A -> Prop) -> Prop;
  empty_open: open (empty _);
  full_open: open (full _);
  inter_open: forall u,
    open u -> forall v, open v
      -> open (inter A u v) ;
  union_open: forall s, (subset _ s open)
    -> open (union A s) }.

```


$$R \subset \operatorname{Hom}_{\mathbf{C}}(, \mathbf{U}), \mathbf{U} \in \mathbf{C},$$

$$R \subset \operatorname{Hom}_{\mathbf{C}}(, \mathbf{U}) \phi : V \rightarrow \mathbf{U} \mathbf{C}$$

$$\phi^{-1}(R) = \{\gamma : W \rightarrow V \parallel \phi \cdot \gamma \in R\}$$

$$VR, R' \subset \operatorname{Hom}_{\mathbf{C}}(, \mathbf{U}) \phi^{-1}(R') \phi : V \rightarrow \mathbf{U} R R' \operatorname{Hom}_{\mathbf{C}}(, \mathbf{U}) \mathbf{U} \in \mathbf{C}$$

$$\operatorname{Ob}_{\mathbf{C}}\{f_i : \mathbf{U}_i \rightarrow \mathbf{U}\}_{i \in \mathbf{I}g} : V \rightarrow \mathbf{U}\{h : V_j \rightarrow V\}_{j \in \mathbf{J}h_j \circ \operatorname{gf}_i} \frac{V_j \mathbf{U}_i}{V \mathbf{U}_i}$$

$$\begin{array}{l} \text{Co } (C\text{: } \textcolor{blue}{\text{precategory}}) \text{ (cod: carrier } C) : \mathbf{U} \\ \quad = (\text{dom: carrier } C) \\ \quad * (\text{hom } C \text{ dom cod}) \end{array}$$

$$\begin{array}{l} \text{Delta } (C\text{: } \textcolor{blue}{\text{precategory}}) \text{ (d: carrier } C) : \mathbf{U} \\ \quad = (\text{index: } \mathbf{U}) \\ \quad * (\text{index} \rightarrow \text{Co } C \text{ d}) \end{array}$$

$$\begin{array}{l} \text{Coverage } (C\text{: } \textcolor{blue}{\text{precategory}}) : \mathbf{U} \\ \quad = (\text{cod: carrier } C) \\ \quad * (\text{fam: Delta } C \text{ cod}) \\ \quad * (\text{coverings: carrier } C \rightarrow \text{Delta } C \text{ cod} \rightarrow \mathbf{U}) \\ \quad * (\text{coverings cod fam}) \end{array}$$

$$\mathbf{CCC}$$

$$\{\phi_\alpha : \mathbf{U}_\alpha \rightarrow \mathbf{U}\}, \mathbf{U} \in \mathbf{C},$$

$$\phi_\alpha : \mathbf{U}_\alpha \rightarrow \mathbf{U} \psi : V \rightarrow \mathbf{U} \mathbf{C} V \times_{\mathbf{U}} \mathbf{U}_\alpha \rightarrow \mathbf{V} \mathbf{V} \{\phi_\alpha : \mathbf{U}_\alpha \rightarrow \mathbf{U}\}_{\{\gamma_{\alpha,\beta} : W_{\alpha,\beta} \rightarrow \mathbf{U}_\alpha\} \alpha}$$

$$W_{\alpha,\beta} \xrightarrow{\gamma_{\alpha,\beta}} \mathbf{U}_\alpha \xrightarrow{\phi_\alpha} \mathbf{U}$$

$$\{1 : \mathbf{U} \rightarrow \mathbf{U}\} \mathbf{U} \in \mathbf{C}$$

$$\begin{array}{l} \text{site } (C\text{: } \textcolor{blue}{\text{precategory}}) : \mathbf{U} \\ \quad = (C\text{: } \textcolor{blue}{\text{precategory}}) * \text{Coverage } C \end{array}$$

$$\mathbf{CC}^{\operatorname{op}} \rightarrow \mathbf{Set}$$

$$\begin{array}{l} \text{presheaf } (C\text{: } \textcolor{blue}{\text{precategory}}) : \mathbf{U} \\ \quad = \textcolor{blue}{\text{catfunctor}} \text{ (opCat } C) \mathbf{Set} \end{array}$$

$$\mathbf{C}$$

$$F : \mathbf{C}^{\operatorname{op}} \rightarrow \mathbf{Set}$$

$$F(\mathbf{U}) \rightarrow \varprojlim_{V \rightarrow \mathbf{U} \in \mathbf{R}} F(V)$$

$$\begin{array}{l} R \subset \operatorname{Hom}_{\mathbf{C}}(, \mathbf{U}) \\ \operatorname{Hom}_{\mathbf{C}}(\operatorname{Hom}_{\mathbf{C}}(, \mathbf{U}), F) \rightarrow \operatorname{Hom}_{\mathbf{C}}(R, F) \end{array}$$

```

sheaf (C: precategory): U
= (S: site C)
* presheaf S.1

```

$$(C,J)CJ$$

$$CR \rightarrow ER \rightarrow E \rightarrow QC$$

$$\mathcal{C}D$$

$$f:\mathbf{Sh}(C)\rightarrow \mathbf{Sh}(D)$$

$$f_*:\mathbf{Sh}(C)\rightarrow \mathbf{Sh}(D)f^*:\mathbf{Sh}(D)\rightarrow \mathbf{Sh}(C)f^*f_*f^*f_*f^*$$

$$ES(p^*,p_*) : E \rightarrow Sp^! \vdash p_*p^! \dashv p_*p^*p^!p_!$$

$$\int \dashv \vdash \dashv \#$$

$$f:Y\rightarrow ZXg_1,g_2:X\rightarrow Y$$

$$f\circ g_1=f\circ g_2\rightarrow g_1=g_2.$$

$$X\mathrm{Hom}(X,)\mathrm{Hom}(X,Y)\rightarrow \mathrm{Hom}(X,Z)$$

```

mono (P: precategory) (Y Z: carrier P) (f: hom P Y Z): U
= (X: carrier P) (g1 g2: hom P X Y)
-> Path (hom P X Z) (compose P X Y Z g1 f)
      (compose P X Y Z g2 f)
-> Path (hom P X Y) g1 g2

```

$$Ctrue:1\rightarrow \Omega 1U\rightarrow X_{XU}:X\rightarrow \Omega_{XU}^{1,k_1}$$

```

subobjectClassifier (C: precategory): U
= (omega: carrier C)
* (end: terminal C)
* (trueHom: hom C end.1 omega)
* (chi: (V X: carrier C) (j: hom C V X) -> hom C X omega)
* (square: (V X: carrier C) (j: hom C V X) -> mono C V X j
-> hasPullback C (omega,(end.1,trueHom),(X,chi V X j)))
* ((V X: carrier C) (j: hom C V X) (k: hom C X omega)
-> mono C V X j
-> hasPullback C (omega,(end.1,trueHom),(X,k))
-> Path (hom C X omega) (chi V X j) k)

```

$$\mathcal{CMLTT}$$

```

isCCC (C: precategory): U
= (Exp: (A B: carrier C) -> carrier C)
* (Prod: (A B: carrier C) -> carrier C)
* (Apply: (A B: carrier C) -> hom C (Prod (Exp A B) A) B)
* (P1: (A B: carrier C) -> hom C (Prod A B) A)
* (P2: (A B: carrier C) -> hom C (Prod A B) B)
* (Term: terminal C)
* unit

```

$\Pi\Sigma\text{SETisSetpropPiMLTT}$

```

cartesianClosure : isCCC Set
= (expo,prod,appli,proj1,proj2,term,tt) where
  exp (A B: SET): SET = (A.1 -> B.1, setFun A.1 B.1 B.2)
  pro (A B: SET): SET = (prod A.1 B.1, setSig A.1 (\(_ : A.1)
    -> B.1) A.2 (\(_ : A.1) -> B.2))
  expo: (A B: SET) -> SET = \ (A B: SET) -> exp A B
  prod: (A B: SET) -> SET = \ (A B: SET) -> pro A B
  appli: (A B: SET) -> hom Set (pro (exp A B) A) B
    = \ (A B: SET) -> \ (x: (pro (exp A B) A).1) -> x.1 x.2
  proj1: (A B: SET) -> hom Set (pro A B) A
    = \ (A B: SET) (x: (pro A B).1) -> x.1
  proj2: (A B: SET) -> hom Set (pro A B) B
    = \ (A B: SET) (x: (pro A B).1) -> x.2
  unitContr (x: SET) (f: x.1 -> unit) : isContr (x.1 -> unit)
    = (f, \ (z: x.1 -> unit) -> propPi x.1 (\ (_:x.1) -> unit)
      (\ (x:x.1) -> propUnit) f z)
  term: terminal Set = ((unit,setUnit),
    \ (x: SET) -> unitContr x (\ (z: x.1) -> tt))

```

```

Topos (cat: precategory) : U
= (cartesianClosure: isCCC cat)
* subobjectClassifier cat

```

```

internal : Topos Set
= (cartesianClosure,hasSubobject)

```

$$\text{Cb} : \text{Ob}_C \downarrow \text{bf} : a \rightarrow \text{bf}^* : C \downarrow b \rightarrow c \downarrow a \sum_f \prod_f$$

$$S^{n-1} \hookrightarrow D^n D^n S^{n-1}$$

$$Xf\colon S^{n-1}\rightarrow X$$

$$\frac{S^{n-1}\times X}{D^n\cup_f D^n}$$

$$X\cup_f D^n f$$

$$-1\varnothing\leqslant nXXn-1$$

$$Xcolimit(X_i)X_{-1}=\varnothing\hookrightarrow X_0\hookrightarrow X_1\hookrightarrow X_2\hookrightarrow \ldots XX_i\leqslant nX_{i+1}X_i$$

$$\varnothing\hookrightarrow X_0\hookrightarrow X_1\hookrightarrow X_2\hookrightarrow \ldots X$$

$$X_i\leqslant i$$

$$(A,a)A:Ua:A$$

```
pointed: U = (A: U) * A
point (A: pointed): A.1 = A.2
space (A: pointed): U = A.1
```

$$\Omega(A,a)=_{\mathrm{def}}((a=_A a),\mathrm{refl}_A(a)).$$

```
omega1 (A: pointed) : pointed
= (Path (space A) (point A) (point A), refl A.1 (point A))
```

$$\begin{cases} \Omega^0(A,a)=_{\mathrm{def}}(A,a) \\ \Omega^{n+1}(A,a)=_{\mathrm{def}}\Omega^n(\Omega(A,a)) \end{cases}$$

```
omega : nat -> pointed -> pointed = split
zero -> idfun pointed
succ n -> \(A: pointed) -> omega n (omega1 A)
```

$$\pi_n S^m = ||\Omega^n(S^m)||_0.$$

```

piS (n: nat): (m: nat) -> U = split
  zero  -> sTrunc (space (omega n (bool,false)))
  succ x -> sTrunc (space (omega n (Sn (succ x),north)))

```

$$\Omega(S^1)=\mathbb{Z}$$

```

data S1 = base
  | loop <i> [ (i=0) -> base ,
              (i=1) -> base ]

loopS1 : U = Path S1 base base

encode (x:S1) (p:Path S1 base x)
  : helix x
  = subst S1 helix base x p zeroZ

decode : (x:S1) -> helix x -> Path S1 base x = split
  base -> loopIt
  loop @ i -> rem @ i where
    p : Path U (Z -> loopS1) (Z -> loopS1)
    = <j> helix (loop1@j) -> Path S1 base (loop1@j)
  rem : PathP p loopIt loopIt
    = corFib1 S1 helix (\(x:S1)->Path S1 base x) base
      loopIt loopIt loop1 (\(n:Z) ->
        comp (<i> Path loopS1 (oneTurn (loopIt n))
              (loopIt (testIsoPath Z Z sucZ predZ
                sucpredZ predsucZ n @ i)))
              (<i>(lem1It n)@-i) [])

```

```

loopS1eqZ : Path U Z loopS1
  = isoPath Z loopS1 (decode base) (encode base)
  sectionZ retractZ

```

$$\frac{S^3S^1S^2S^3R^3S^0S^1S^3S^7}{S^3S^2}$$

$$S^3$$

$$S^3\mathbb{R}^4$$

$$S^3 = \{(x_0,x_1,x_2,x_3) \in \mathbb{R}^4 : \sum_{i=0}^3 x_i^2 = 1\};$$

$$\mathbb{H}$$

$$S^3 = \{x \in \mathbb{H} : \|x\| = 1\}.$$

$$S^3(\eta, \theta_1, \theta_2)$$

$$\begin{cases} x_0 = \cos(\theta_1)\sin(\eta), \\ x_1 = \sin(\theta_1)\sin(\eta), \\ x_2 = \cos(\theta_2)\cos(\eta), \\ x_3 = \sin(\theta_2)\cos(\eta). \end{cases}$$

$$\eta \in [0, \frac{\pi}{2}] \theta_{1,2} \in [0, 2\pi]$$

$$S^2S^2\theta_2$$

$$\begin{cases} x = \sin(2\eta)\cos(\theta_1), \\ y = \sin(2\eta)\sin(\theta_1), \\ z = \cos(2\eta). \end{cases}$$

```
var fiber = new THREE.Curve(),
    color = sphericalCoords.color;

fiber.getPoint = function(t) {
    var eta = sphericalCoords.eta,
        phi = sphericalCoords.phi,
        theta = 2 * Math.PI * t;
    var x1 = Math.cos(phi+theta) * Math.sin(eta/2),
        x2 = Math.sin(phi+theta) * Math.sin(eta/2),
        x3 = Math.cos(phi-theta) * Math.cos(eta/2),
        x4 = Math.sin(phi-theta) * Math.cos(eta/2);
    var m = mag([x1,x2,x3]),
        r = Math.sqrt((1-x4)/(1+x4));
    return new THREE.Vector3(r*x1/m,r*x2/m, r*x3/m);
};
```

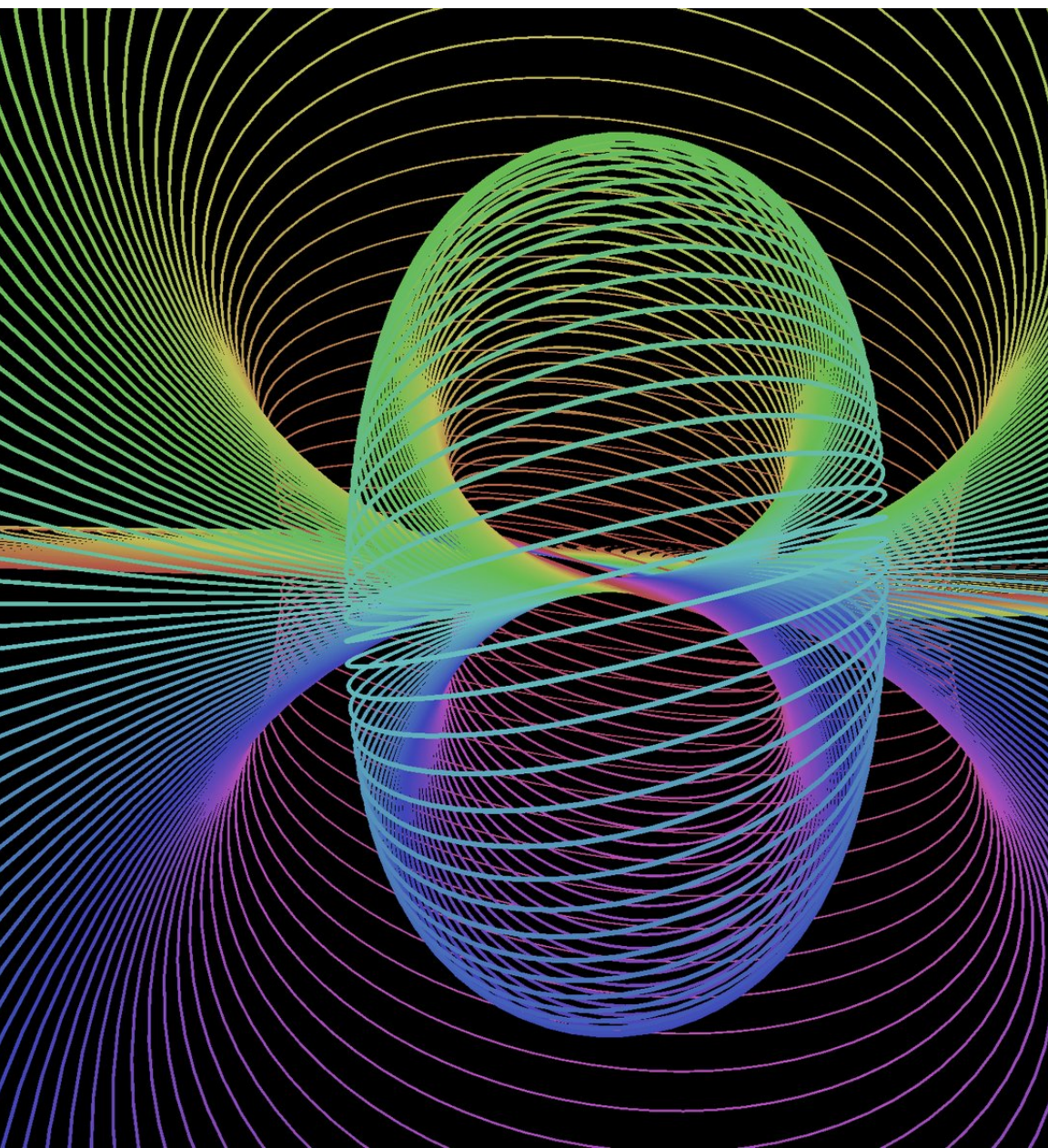
$$S^3$$

```
rot: (x : S1) -> Path S1 x x = split
  base -> loop1
  loop @ i -> constSquare S1 base loop1 @ i

mu : S1 -> equiv S1 S1 = split
  base -> idEquiv S1
  loop @ i -> equivPath S1 S1 (idEquiv S1)
    (idEquiv S1) (<j> \ (x : S1) -> rot x @ j) @ i

H : S2 -> U = split
  north -> S1
  south -> S1
  merid x @ i -> ua S1 S1 (mu x) @ i

total : U = (c : S2) * H c
```



$$\mathsf{A}$$

$$\mathsf{H}_\mathsf{A} = \begin{cases} \mathsf{A} : \mathsf{U} \\ e : \mathsf{A} \\ \mu : \mathsf{A} \rightarrow \mathsf{A} \rightarrow \mathsf{A} \\ \beta : (\mathsf{a} : \mathsf{A}) \rightarrow \Sigma(\mu(e,\mathsf{a}) = \mathsf{a})(\mu(\mathsf{a},e) = \mathsf{a}) \end{cases}$$

$$(S^0,S^1,p,S^1)(S^1,S^3,p,S^2)(S^3,S^7,p,S^4)(S^7,S^{15},p,S^8)$$

$$\phi:S^{2n-1}\rightarrow S^n\phi\mathrm{cofib}(\phi)=S^n\cup_\phi\mathbb{D}^{2n}$$

$$\mathrm{H}^k(\mathrm{cofib}(\phi),\mathbb{Z})=\begin{cases}\mathbb{Z}\mathrm{fork}=n,2n\\0\mathrm{otherwise}\end{cases}$$

$$\alpha,\beta n2n\mathrm{h}(\phi)\alpha\beta\alpha\sqcup\alpha=\mathrm{h}(\phi)\cdot\beta\mathrm{h}(\phi)\phi$$

$$\mathbf{1}$$

O_{HTS}

```
module buddhism where
import path

concept (o: U): U
  = o -> U

nondual (o: U) (p: concept o): U
  = (x y: o) -> Path U (p x) (p y)

allpaths (o: U): U
  = (x y: o) -> Path o x y

  u(py)x =o y.(coerce)(cong).

encode (o:U): ((p: concept o) -> nondual o p) -> allpaths o
  = \ (nd: (p: concept o) -> nondual o p) (a b: o)
  -> coerce(Path o a a)(Path o a b)(nd(\(z:o)->Path o a z)a b)(refl o a)

decode (o:U): allpaths o -> ((p: concept o) -> nondual o p)
  = \ (all: allpaths o)(p: concept o)(x y: o) -> cong o U p x y (all x y)

1
```

```

CoInductive Co (E : Effect.t) : Type -> Type :=
| Bind : forall (A B : Type), Co E A -> (A -> Co E B) -> Co E B
| Split : forall (A : Type), Co E A -> Co E A -> Co E A
| Join : forall (A B : Type), Co E A -> Co E B -> Co E (A * B).
| Ret : forall (A : Type) (x : A), Co E A
| Call : forall (command : Effect.command E),
    Co E (Effect.answer E command)

Definition run (argv : list LString.t) : Co effect unit :=
  ido! log (LString.s "What is your name?") in
  ile! name := read_line in
  match name with
  | None => ret tt
  | Some name => log (LString.s "Hello " ++ name ++ LString.s "!!")
  end.

Parameter infinity : nat.
Definition eval {A} (x : Co effect A) : Lwt.t A := eval_aux infinity x.

Fixpoint eval_aux {A} (s : nat) (x : Co effect A) : Lwt.t A :=
  match s with
  | 0 => error tt
  | S s =>
    match x with
    | Bind _ _ x f => Lwt.bind (eval_aux s x) (fun v_x => eval_aux s (f
      v_x))
    | Split _ x y => Lwt.choose (eval_aux s x) (eval_aux s y)
    | Join _ _ x y => Lwt.join (eval_aux s x) (eval_aux s y)
    | Ret _ v => Lwt.ret v
    | Call c => eval_command c
    end
  end.

CoFixpoint handle_commands : Co effect unit :=
  ile! name := read_line in
  match name with
  | None => ret tt
  | Some command =>
    ile! result := log (LString.s "Input: "
      ++ command ++ LString.s ".")
  in handle_commands
  end.

Definition launch (m : list LString.t -> Co effect unit) : unit :=
  let argv := List.map String.to_lstring Sys.argv in
  Lwt.launch (eval (m argv)).

Definition corun (argv : list LString.t) : Co effect unit :=
  handle_commands.

Definition main := launch corun.

```

```

String: Type = List Nat
data IO: Type =
  (getLine: (String -> IO) -> IO)
  (putLine: String -> IO)
  (pure: () -> IO)

-- IOI/@: (r: U) [x: U] [[s: U] -> s -> [s -> #IOI/F r s] -> x] x
  \ (r : *)
-> \ / (x : *)
-> (\ / (s : *)
  -> s
  -> (s -> #IOI/F r s)
  -> x)
-> x

-- IOI/F
  \ (a : *)
-> \ (State : *)
-> \ / (IOF : *)
-> \ / (PutLine_ : #IOI/data -> State -> IOF)
-> \ / (GetLine_ : (#IOI/data -> State) -> IOF)
-> \ / (Pure_ : a -> IOF)
-> IOF

-- IOI/MkIO
  \ (r : *)
-> \ (s : *)
-> \ (seed : s)
-> \ (step : s -> #IOI/F r s)
-> \ (x : *)
-> \ (k : forall (s : *) -> s -> (s -> #IOI/F r s) -> x)
-> k s seed step

-- Morte/corecursive
( \ (r: *1)
-> ( (((#IOI/MkIO r) (#Maybe/@ #IOI/data)) (#Maybe/Nothing #IOI/data))
  ( \ (m: (#Maybe/@ #IOI/data))
    -> (((((#Maybe/maybe #IOI/data) m) ((#IOI/F r) (#Maybe/@ #IOI/data)))
      ( \ (str: #IOI/data)
        -> ((((#IOI/putLine r) (#Maybe/@ #IOI/data)) str)
          (#Maybe/Nothing #IOI/data))))
      (((#IOI/getLine r) (#Maybe/@ #IOI/data))
        (#Maybe/Just #IOI/data))))))

```

```

copure() ->
  fun (_) -> fun (IO) -> IO end end.

cogetLine() ->
  fun (IO) -> fun (_) ->
    L = ch:list(io:get_line("> ")),
    ch:ap(IO,[L]) end end.

coputLine() ->
  fun (S) -> fun (IO) ->
    X = ch:unlist(S),
    io:put_chars(": "++X),
    case X of "0\n" -> list([]);
    _ -> corec() end end end.

corec() ->
  ap('Morte':corecursive(),
    [copure(),cogetLine(),coputLine(),copure(),list([])]).

```

```

> om_extract:extract("priv/normal/IOI").
ok
> Active: module loaded: {reloaded,'IOI'}

```

```

> om:corec().
> 1
: 1
> 0
: 0
#Fun<List.3.113171260>

```

```

-- IO/0
\ (a : *)
-> \ / (IO : *)
-> \ / (GetLine_ : (#IO/data -> IO) -> IO)
-> \ / (PutLine_ : #IO/data -> IO -> IO)
-> \ / (Pure_ : a -> IO)
-> IO

-- IO/replicateM
\ (n: #Nat/0)
-> \ (io: #IO/0 #Unit/0)
-> #Nat/fold n (#IO/0 #Unit/0)
    (#IO/[>>] io)
    (#IO/pure #Unit/0 #Unit/Make)

```

```

-- Morte/recursive
((#IO/replicateM #Nat/Five)
  (((#IO/[>>=] #IO/data) #Unit/0) #IO/getLine) #IO/putLine))

```

```

pure() ->
    fun(IO) -> IO end.

getLine() ->
    fun(IO) -> fun(_) ->
        L = ch:list(io:get_line("> ")),
        ch:ap(IO,[L]) end end.

putLine() ->
    fun(S) -> fun(IO) ->
        io:put_chars(": "++ch:unlist(S)),
        ch:ap(IO,[S]) end end.

rec() ->
    ap('Morte':recursive(),
        [getLine(),putLine(),pure(),list([])]).

> om:rec().
> 1
: 1
> 2
: 2
> 3
: 3
> 4
: 4
> 5
: 5
#Fun<List.28.113171260>

```