

MST - DCS - Sem IV

Q1.] a.) Here, we are given

$$V_1 = (1000) \quad V_2 = (1010)$$

Their linear combinations are $a_1, a_2 = \{0, 1\}$

Now, by the linear combinations, we'll find out the value of $a_1 V_1 + a_2 V_2$

$$\begin{aligned} &\therefore 0 \cdot (1000) + 0 \cdot (1010) \\ &= (0000) + (0000) \\ &= 0000 \end{aligned}$$

$$[a_1 = 0, a_2 = 0]$$

$$\begin{aligned} &0 \cdot (1000) + 1 \cdot (1010) \\ &= (0000) + (1010) \\ &= 1010 \end{aligned}$$

$$[a_1 = 0, a_2 = 1]$$

$$\begin{aligned} &1 \cdot (1000) + 0 \cdot (1010) \\ &= 1000 + 0000 \\ &= 1000 \end{aligned}$$

$$[a_1 = 1, a_2 = 0]$$

$$\begin{aligned} &1 \cdot (1000) + 1 \cdot (1010) \\ &= (1000) + (1010) \\ &= (0010) \end{aligned}$$

$$[a_1 = a_2 = 1]$$

Here, we can see that $a_1 V_1 + a_2 V_2 \neq 0$

Only when $a_1 = a_2 = 0$, $a_1 V_1 + a_2 V_2 = 0$

Hence, they are linearly independent.

Affham

Q1.] b.) $P(x) = x^3 + x^2 + 1$

$\therefore m = 3$

$\therefore n = 2^m - 1 = 7$

\therefore It should divide a polynomial given with degree 7.

$$\begin{array}{r}
 x^4 + x^3 + x^2 + 1 \\
 \hline
 \therefore x^3 + x^2 + 1 \overline{) x^7} \\
 \underline{x^7 + x^6 + x^4} \\
 (-) x^6 + x^4 + 1 \\
 \underline{x^6 + x^5 + x^3} \\
 (-) x^4 + x^5 + x^3 + 1 \\
 \underline{x^4 + x^5} \\
 (-) x^3 + x^2 + 1 \\
 \underline{x^3 + x^2 + 1} \\
 (-) 0
 \end{array}$$

\therefore The remainder = 0,

$\therefore P(x) = x^3 + x^2 + 1$ is primitive.

Also, for $GF(2)$ let us consider $\{0, 1\}$

$\therefore P(0) = 0^3 + 0^2 + 1 \neq 0$

$P(1) = 1 + 1 + 1 = 3 \neq 0$

Hence, the polynomial is irreducible as well.

Wahid

Q1.] c.) $GF(8) = GF(2^3) = GF(2^m)$

$\therefore m = 3$

$\therefore GF(8) = \{0, 1, \alpha, \alpha^2, \alpha^3, \alpha^4, \dots, \alpha^{2^m-2}\}$
 $\therefore GF(8) = \{0, 1, \alpha, \alpha^2, \alpha^3, \alpha^4, \alpha^5, \alpha^6\}$

Given primitive polynomial is $p(x) = x^3 + x + 1$
 Putting $x = \alpha$ ($\because \alpha$ is a primitive element)

$\therefore \alpha^3 + \alpha + 1 = 0$
 $\therefore \alpha^3 = \alpha + 1$

$\therefore \alpha^4 = \alpha^3 \cdot \alpha = \alpha(\alpha + 1) = \alpha^2 + \alpha$

$\therefore \alpha^5 = \alpha^4 \cdot \alpha = \alpha(\alpha^2 + \alpha) = \alpha^3 + \alpha^2 = 1 + \alpha + \alpha^2$

$\therefore \alpha^6 = \alpha^5 \cdot \alpha = \alpha(1 + \alpha + \alpha^2) = \alpha + \alpha^2 + \alpha^3$
 $= \alpha + \alpha^2 + \alpha + 1$
 $= 1 + \alpha^2$

$\therefore GF(8) = \{0, 1, \alpha, \alpha^2, 1 + \alpha, \alpha^2 + \alpha, 1 + \alpha + \alpha^2, 1 + \alpha^2\}$

\therefore Multiplicative inverse of α will be α^6 .

$\alpha^6 = 1 + \alpha^2$

\therefore In binary form, $\alpha^6 = [101]$

Althant

Q2.]

a.) The conjugacy class for α^3 can be found by the following method:-

$$\text{Let } \beta = \alpha^3$$

$\therefore (\beta)^{2^l}$ for $l = 1, 2, 3, \dots$ is to be found out.

$$\therefore (\alpha^3)^{2^1} = \alpha^6$$

$$(\alpha^3)^{2^2} = \alpha^{12}$$

$$(\alpha^3)^{2^3} = \alpha^{24}$$

$$(\alpha^3)^{2^4} = \alpha^{48} = \alpha^{31} \cdot \alpha^{17} = \alpha^{17}$$

$$(\alpha^3)^{2^5} = \alpha^{96} = (\alpha^{31})^3 \cdot \alpha^3 = \alpha^3 \Rightarrow \text{Repetition starts}$$

\therefore Conjugacy class for

$$\alpha^3 = \{ \alpha^3, \alpha^6, \alpha^{12}, \alpha^{17}, \alpha^{24} \}$$

Alphale

Q2.] b.) $F = \{ 0, 1, 2, 3, 4 \}$

Additive Table :-

+	0	1	2	3	4
0	0	1	2	3	4
1	0	2	3	4	0
2	2	3	4	0	1
3	3	4	0	1	2
4	4	0	1	2	3

\Rightarrow Commutative

i.e., $a + b = b + a$

$$1 + 2 = 2 + 1 = 3$$

$$1 + 4 = 4 + 1 = 0$$

\therefore Commutative

Multiplicative Table :-

\times	0	1	2	3	4
0	0	0	0	0	0
1	0	1	2	3	4
2	0	2	4	1	3
3	0	3	1	4	2
4	0	4	3	2	1

\Rightarrow Commutative

i.e., $a \times b = b \times a$

$$1 \times 2 = 2 \times 1 = 2$$

$$4 \times 3 = 3 \times 4 = 2$$

A) Here, we can see that closure property is followed, i.e., the result of addition or multiplication of two elements in set lies in the set.

B) Additive identity \rightarrow let it be e .

$$\therefore a + e = a$$

$$\text{Here, } e = 0$$

Multiplicative identity \rightarrow let it be e

$$\therefore a \cdot e = a$$

$$\text{where, } e = 1$$

Wahab

Q2.] c) Order m for an extension field $GF(n)$ is given by,

$$m = \log_2 n$$

Here, $n = 32$

$$\therefore m = \log_2 32 = 5$$

\therefore order of this extension field $GF(32) = 5$

Answer

Q3] a) We know,

$$G = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 1 \end{bmatrix} \rightarrow \text{given}$$

$$m = [101] \rightarrow \text{given}$$

$$\begin{aligned} \text{Also, } c &= m \cdot G \\ &= [101] \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 1 \end{bmatrix} \end{aligned}$$

$$= [1 \cdot 1 + 0 \cdot 0 + 1 \cdot 0, 1 \cdot 0 + 0 \cdot 1 + 1 \cdot 0, 1 \cdot 0 + 0 \cdot 0 + 1 \cdot 1, 1 \cdot 1 + 0 \cdot 0 + 1 \cdot 1, 1 \cdot 1 + 0 \cdot 1 + 1 \cdot 1, 1 \cdot 0 + 0 \cdot 1 + 1 \cdot 1]$$

$$= [1, 0, 1, 1, 1, 1]$$

$$1 \cdot 0 + 0 \cdot 0 + 1 \cdot 1 = 1$$

$$1 \cdot 1 + 0 \cdot 0 + 1 \cdot 1 = 1$$

$$1 \cdot 1 + 0 \cdot 1 + 1 \cdot 1 = 1$$

$$1 \cdot 0 + 0 \cdot 1 + 1 \cdot 1 = 1$$

$$\therefore c = [101111]$$

Q3] b) Now, we know, $G = [I_k \ P_{n-k}]$
From given Generator matrix,

$$P = \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix}$$

$$\text{Also, } H = [P^T \ I_{n-k}]$$

$$\therefore H = \begin{bmatrix} 1 & 1 & 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 \end{bmatrix}$$

To check if the codeword is valid, $S = r \cdot H^T = 0$

Attract

$$\therefore S = [0 \ 1 \ 1 \ 1 \ 1 \ 1] \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$\therefore S = \begin{bmatrix} 0 \cdot 1 + 1 \cdot 0 + 1 \cdot 1 + 1 \cdot 1 + 1 \cdot 0 + 1 \cdot 0 \\ 0 \cdot 1 + 1 \cdot 1 + 1 \cdot 1 + 1 \cdot 0 + 1 \cdot 1 + 1 \cdot 0 \\ 0 \cdot 0 + 1 \cdot 1 + 1 \cdot 1 + 1 \cdot 0 + 1 \cdot 0 + 1 \cdot 1 \end{bmatrix}$$

$$\therefore S = [0 \ 1 \ 1]$$

\therefore The syndrome is not $[0 \ 0 \ 0]$,

\therefore The received codeword $[1 \ 0 \ 1 \ 1 \ 1 \ 1]$ is not a valid codeword of the code set and has an error at one bit. (~~first bit~~).

Q3.] c.) Now, we have been given

$$x^{15} + 1 = (x+1)(x^4 + x^3 + x^2 + x + 1)(x^4 + x^3 + 1) \cdot (x^4 + x + 1)(x^2 + x + 1)$$

For a (n, k) systematic cyclic code, degree of generator polynomial $G(x)$ is $n-k$.

$$\therefore \text{Degree of } G(x) = 15 - 9 = 6$$

$$\begin{aligned} \therefore G(x) &= (x^4 + x + 1)(x^2 + x + 1) \\ &= x^6 + x^5 + x^4 + x^3 + x^2 + x + 1 \\ &= x^6 + x^5 + x^4 + x^3 + 1 \end{aligned}$$

$$\therefore G(x) = x^6 + x^5 + x^4 + x^3 + 1$$

Alham