

Today: Advanced Topics

- ① Coordinate Descent
- ② Equality-constrained Newton
- ③ Interior-Point methods (out of scope)

Coordinate Descent

- Sequential version:

$$F(\vec{x}) = g(\vec{x}) + \sum_{i=1}^n h_i(x_i)$$

g : convex, differentiable

h_i : convex

LASSO: $F(\vec{x}) = \frac{1}{2} \|A\vec{x} - \vec{b}\|_2^2 + \lambda \|\vec{x}\|_1$

non-differentiable,
but separable
i.e., $\sum |x_i|$

Coordinate Descent:

Initial guess: $\vec{x}^{(0)} = [x_1^{(0)} \ x_2^{(0)} \ \dots \ x_n^{(0)}]$

k^{th} guess for x_i :

$$x_i^{(k)} = \underset{x_i}{\operatorname{argmin}} F(x_i, x_2^{(k-1)}, \dots, x_n^{(k-1)})$$

x_i ← minimization variable

Fixed terms

doing a minimization over scalar x_i

↳ LASSO is a type of problem where this sequential version of coordinate descent works really well

eg 1st guess for x_i :

$$x_i^{(0)} = \underset{x_i}{\operatorname{argmin}} F(x_i, x_2^{(0)}, x_3^{(0)}, \dots, x_n^{(0)})$$

↳ get this from previous step

Algorithm:

① $x_1^{(k)} = \underset{y}{\operatorname{argmin}} F(y, x_2^{(k-1)}, \dots, x_n^{(k-1)})$

② $x_2^{(k)} = \underset{y}{\operatorname{argmin}} F(x_1^{(k)}, y, x_3^{(k-1)}, \dots, x_n^{(k-1)})$

⋮

③ $x_n^{(k)} = \underset{y}{\operatorname{argmin}} F(x_1^{(k)}, x_2^{(k)}, \dots, x_{n-1}^{(k)}, y)$

LASSO outline w/ coordinate descent:

$$f(\vec{x}) = \underbrace{\frac{1}{2} \|\mathbf{A}\vec{x} - \vec{b}\|_2^2}_{\text{quadratic Fcn of } x_i} + \underbrace{\lambda \|\vec{x}\|_1}_{\lambda |\vec{x}|}$$

$\nabla_{x_i} f$

- \hookrightarrow compute derivative
- \hookrightarrow do soft-thresholding coordinate by coordinate
- \hookrightarrow converges as $K \rightarrow \infty$ to optimum point

\hookrightarrow casts $x_1 > 0$
 \hookrightarrow (compute derivative) $x_2 < 0$
 $x_2 = 0$

- Note: a lot of times, convergence is determined by looking at the duality gap.

\rightarrow ie, solve primal \hat{z} corresponding dual \hat{y} see gap
 b/w dual feasible \hat{y} point we just found

Newton's Method for Equality-constrained QPs

- Equality-constrained QPs \rightarrow unconstrained QPs

$$\min \frac{1}{2} \vec{x}^T \mathbf{H} \vec{x} + \vec{c}^T \vec{x} + d \quad \mathbf{H} \text{ PD} \rightarrow \text{strong duality}$$

s.t. $\mathbf{A}\vec{x} = \vec{b}$

$$L(\vec{x}, \vec{y}) = \frac{1}{2} \vec{x}^T \mathbf{H} \vec{x} + \vec{c}^T \vec{x} + d + \vec{y}^T (\mathbf{A}\vec{x} - \vec{b})$$

KKT conditions: (necessary \hat{z} sufficient bc convex \hat{z} strong duality)
 $\mathbf{A}\vec{x}^* = \vec{b}$
 $\mathbf{H}\vec{x}^* + \vec{c} + \mathbf{A}^T \vec{y}^* = 0$
 \hookrightarrow no λ^* cond bc no inequalities

\hookrightarrow linear:

$$\begin{bmatrix} \mathbf{H} & \mathbf{A}^T \\ \mathbf{A} & \mathbf{0} \end{bmatrix} \begin{bmatrix} \vec{x}^* \\ \vec{y}^* \end{bmatrix} = \begin{bmatrix} -\vec{c} \\ \vec{b} \end{bmatrix}$$

IF $\begin{bmatrix} \mathbf{H} & \mathbf{A}^T \\ \mathbf{A} & \mathbf{0} \end{bmatrix}$ IS non-singular \wedge invertible, we have a unique solution

$$\begin{aligned} \min \quad & F(\vec{x}) \\ \text{s.t.} \quad & A\vec{x} = \vec{b} \end{aligned}$$

Say I start with some \vec{x}_0 s.t.

$$A\vec{x}_0 = \vec{b}$$

$$\begin{aligned} \min \quad & \hat{F}(\vec{x}_0 + \vec{v}) = \underbrace{F(\vec{x}_0)}_{\text{const}} + \underbrace{\nabla F(\vec{x}_0)^T \vec{v}}_{\text{linear}} + \underbrace{\frac{1}{2} \vec{v}^T \nabla^2 F(\vec{x}_0) \vec{v}}_{\text{quadratic}} \\ \text{s.t.} \quad & A(\vec{x}_0 + \vec{v}) = \vec{b} \\ & A\vec{v} = \vec{0} \end{aligned}$$

→ want to find that \vec{v} that minimizes the quadratic

$$\text{solution: } \begin{bmatrix} \nabla^2 F(x) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} \vec{v} \\ \vec{\lambda} \end{bmatrix} = \begin{bmatrix} -\nabla F(x_0)^T \\ 0 \end{bmatrix}$$

bc you've started in the feasible set \Rightarrow you're requiring $A\vec{v} = \vec{0}$, we always stay in the feasible set

→ the step remains feasible every time