

When to Commit to an Action in Online Planning

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Plan

- Introduction.
- Problem setting.
- Proposed method.
- Committing VS not committing .
- Summary.

Introduction

- When a planner commits to an action, it re-roots its search tree at the node that was a result of that action.
- A time presser results in real time search, when the planner have to commit to a new action before the previously one ended. This action can be a no-operation action, that keeps the state unchanged.
- Time presser can be beneficial to commit early to actions, so that more look ahead search are performed on upcoming states.
- The paper proposes a method for making commitment decision.

Problem setting

- The system must be controlled all the time (actions must be executing at any time).
- The planner can commit to actions earlier in order to re-root the tree at a deeper node → focus the search later on the future.
- Assumptions:
 1. Actions are serials.
 2. The world is observable and deterministic.
- Objective of the system:

Achieve the goal as soon as possible.

Proposed method

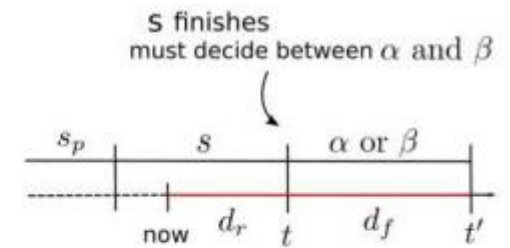
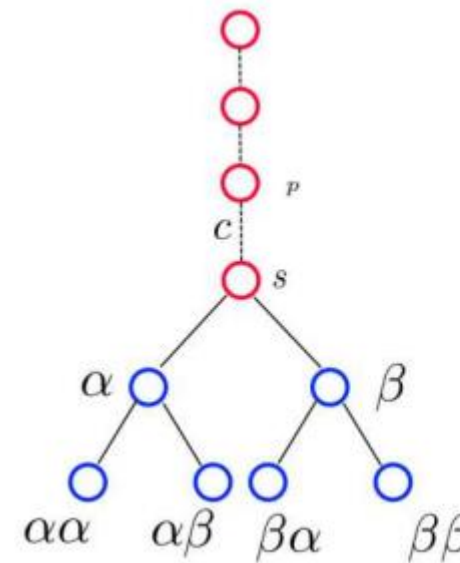
- A meta-reasoning scheme for action commitment called Flexible Action Commitment Search (FACS).
- The paper used additional assumption to estimate the utility of the committing or not committing to an action and to simplify the analysis such as:
 1. Each node has exactly two children.
 2. The time required to fully execute an action is identical for all actions.
 3. The order of decision is a fixed search tree.
 4. No re-planning after action commitment.
 5. The search may re-started if necessary.

Committing VS not committing

- During the search, after each expansion or periodically after a set of expansions a meta reasoning process decides whether:
- Commit to the current top-level action **now and** re-root the search from that node, **or**
- postpone the commitment and continue the search.

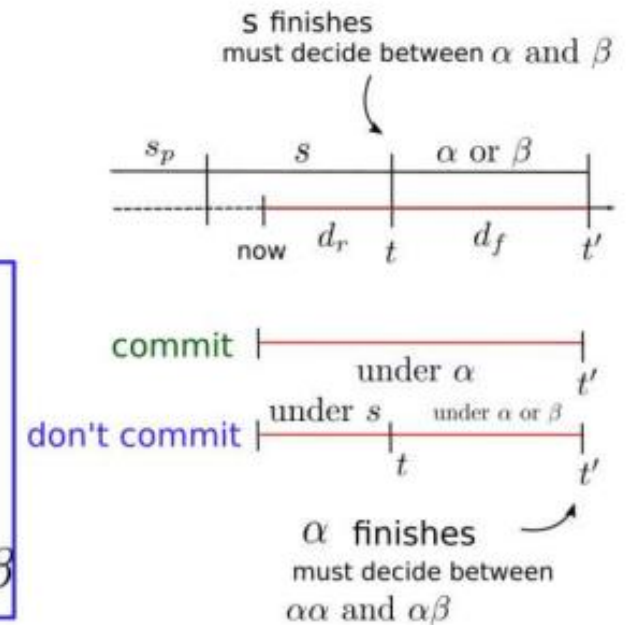
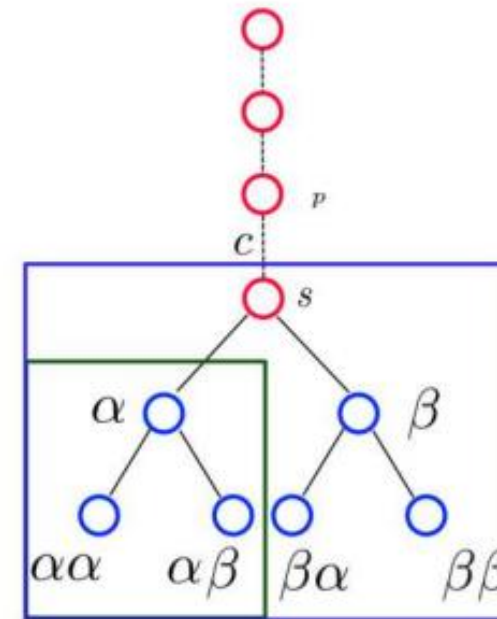
Committing VS not committing

- Current search tree **rooted** at node **s**.
- Available search time denoted by d consist of d_r and d_f
- d_r : is the remaining time induced by previous commitments
- d_f : The time required to execute a full action.



Committing VS not committing

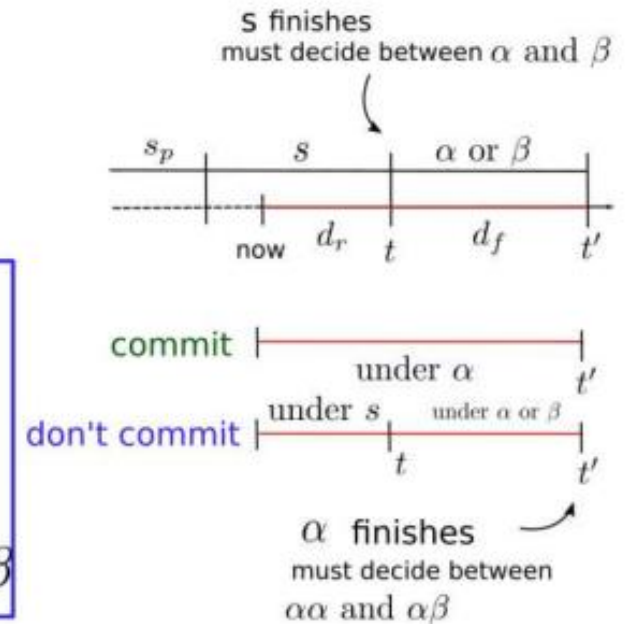
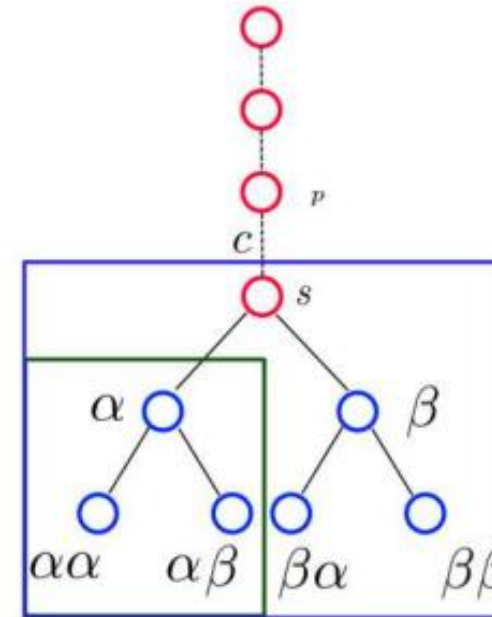
- By committing to α , all the available search time is used to search under children of α ($\alpha\alpha$ and $\alpha\beta$). each one receives a search budget of $d = (d_r + d_f) / 2$.
- The utility of committing to α is the expectation of the minimum \hat{f} -value of $\alpha\alpha$ and $\alpha\beta$, after searching d time units under each of them: $U_{\text{commit}} = E [\min(X_d \alpha\alpha, X_d \alpha\beta)]$



Committing VS not committing

- If choosing to not commit yet. Half of d_r ($d_r/2$) is used to search under each child of the root (s).
- The rest of the time d_f is used to search under whichever child of the root is judged most promising at that time.
- The search duration under each grand child is

$$d' = ((d_r/2) + d_f) / 2.$$



Summary

- Committing means re-rooting the search tree.
- Not committing means to not re-root the search tree.
- Algorithms committing to one action at a time and re-root the search tree at every step cannot benefit from gaining future searching time.
- The method proposed in this paper focuses on real time search to get more time to search ahead in the search tree.
- The evaluation of the method outperform previously proposed fixed strategies such as LSS-LRTA*(all), LSS-LRTA*(one) and Dynamic f.