

Code length: ν (even)

Binary constraint vector set: $\mathcal{B} \subset \mathbb{B}^{\nu}$ $(|\mathcal{B}| \geq 2^{k_{\mathrm{a}}})$ Composite symbol set:

$$\begin{split} &\Phi_0 \subseteq \left\{ (\sigma_0, \sigma_1, \sigma_2, \sigma_3) \in \mathbb{Z}_{k+1}^4 \ \middle| \ \sigma_0 + \sigma_1 = k, \sigma_2 = \sigma_3 = 0 \right\} \text{ AT} \\ &\Phi_1 \subseteq \left\{ (\sigma_0, \sigma_1, \sigma_2, \sigma_3) \in \mathbb{Z}_{k+1}^4 \ \middle| \ \sigma_0 = \sigma_1 = 0, \sigma_2 + \sigma_3 = k \right\} \text{ GC} \\ &|\Phi_0| = |\Phi_1| = 2^{k_b} \le k + 1 \qquad \Phi = \Phi_0 \cup \Phi_1 \end{split}$$

Inner codebook:

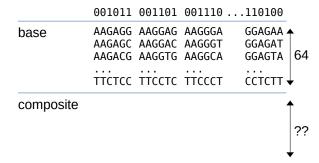
$$\mathcal{C} \subseteq \left\{ (c_0, \dots, c_{\nu-1}) \mid c_j \in \Phi_{b_j}, (b_0, \dots, b_{\nu-1}) \in \mathcal{B} \right\}$$

Encode:

$$\begin{split} f_{\mathrm{e}}^{\mathrm{a}} : \mathbb{B}^{k_{\mathrm{a}}} &\to \mathcal{B} \\ f_{\mathrm{e}}^{i} : \mathbb{B}^{k_{\mathrm{b}}} &\to \Phi_{i} \ (i \!\in\! \mathbb{B}) \\ \end{bmatrix} \text{bijection} \\ f_{\mathrm{e}}(\boldsymbol{c}_{i}^{\mathrm{a}}, \boldsymbol{c}_{i}^{\mathrm{b}}) &= \boldsymbol{s}_{i} = (s_{i,0}, \dots, s_{i,\nu-1}) : \mathbb{B}^{k_{\mathrm{a}}} \!\times\! \mathbb{B}^{k_{\mathrm{b}}\nu} &\to \Phi^{\nu} \\ s_{i,j} &= f_{\mathrm{e}}^{b_{i}}(\boldsymbol{c}_{i,j}^{\mathrm{b}}) \\ (b_{0}, \dots, b_{\nu-1}) &= f_{\mathrm{e}}^{\mathrm{a}}(\boldsymbol{c}_{i}^{\mathrm{a}}) \\ \boldsymbol{c}_{i,j}^{\mathrm{b}} &= [\boldsymbol{c}_{i}^{\mathrm{b}}]_{j\nu}^{j\nu+\nu-1} \\ \text{flip (RL, LB)} \end{split}$$

(example) $\nu = 6$, $|\mathcal{B}| = 16$, k = 7, $k_a = 4$, $k_b = 3$, Lee dist? 011010 100101

Decode:



multi-base composite (Manchester coding)

 $\nu = 1 : (A,C,G,T)$

 $\nu\!=\!2$: (AA,AC,AG,AT, CA,CC,CG,CT, GA,GC,GG,GT, TA,TC,TG,TT)

 $\nu = 3$: (AAA, AAC, ..., TTT)

001011 001101 001110 ...110100

AAGAGG AAGGAG AAGGGA GGAGAA AAGAGC AAGGAC AAGGGT GGAGAC

AAGACG AAGGTG AAGGCA GGAGCA 64 ... CCTCTT CCTTCT CCTTTC

composite

base