

 $\begin{array}{ll} \text{Alphabet:} & \Sigma \!=\! \{0,1,2,3\} \!\leftrightarrow\! \{\mathsf{A},\mathsf{T},\mathsf{G},\mathsf{C}\} \\ \text{Code length:} & \nu & \text{symbol (even)} \\ \text{Block length:} & N \!=\! \nu n & \text{symbol (block} = n & \text{inner codewords)} \\ \text{Constraints (binary):} & \text{run length:} & \mathsf{RL-}\rho & \\ \text{local balance:} & \mathsf{LB-}(\ell,\delta) & |w(b_i^{i+\ell-1}) \!-\! \frac{\ell}{2}| \! \leq \! \delta \\ \text{indicator func:} & \mathbb{1}_{\rho,\ell,\delta}(\boldsymbol{b}) \\ \end{array}$ 

IDS channel:

error prob:  $p_{\mathrm{i}}, p_{\mathrm{d}}, p_{\mathrm{s}}$ 

drift vector:

 $oldsymbol{d}^j \!=\! (d_0^j,\ldots,d_N^j)$  (symbol level)

$$ilde{m{d}}^j \!=\! ( ilde{d}_0^j, \dots, ilde{d}_n^j)$$
 (word level)  $ilde{d}_i^j \!=\! d_{i
u}^j$ 

## Binary constraint vector set [SITA2024]

$$\begin{split} \mathcal{B} \subset \mathbb{B}^{\nu} & \quad (|\mathcal{B}| = 2^{k_{\mathrm{a}}}) \\ \forall \pmb{b} \in \mathcal{B}, w(\pmb{b}) = \nu/2 & \quad \text{(LB: inside CW)} \\ \forall \pmb{b} \in \mathcal{B}, \forall i \in [\nu]_{\mathrm{even}}, \left| \tilde{w}(\pmb{b}_i^{\nu-1}) \right| \leq 2\delta & \quad \text{(LB: CW boundary)} \\ \forall \pmb{b} \in \mathcal{B}, \forall i \in [\nu]_{\mathrm{odd}}, \left| \tilde{w}(\pmb{b}_i^{\nu-1}) \right| \leq 2\delta - 1 & \quad \text{(LB: CW boundary)} \\ \forall \pmb{b} \in \mathcal{B}, \lambda(\pmb{b}) \leq \rho & \quad \text{(RL: inside CW)} \\ \forall \pmb{b} \in \mathcal{B}, \lambda(\pmb{b}, \nu-1) \leq \rho - 1 & \quad \text{(RL: CW boundary)} \\ \pmb{b} \in \mathcal{B} \to \overline{\pmb{b}} \in \mathcal{B} & \quad \text{(re-balance)} \\ (01)^{\frac{\nu}{2}}, (10)^{\frac{\nu}{2}} \notin \mathcal{B} & \quad \text{(erasure symbol)} \\ \mathcal{\tilde{B}} = \mathcal{B} \cup \left\{ (01)^{\frac{\nu}{2}}, (10)^{\frac{\nu}{2}} \right\} & \quad \text{(re-balance)} \end{split}$$

# Composite symbol sets

$$\begin{split} &\Phi_0 \subseteq \left\{ (\sigma_0, \sigma_1, \sigma_2, \sigma_3) \in \mathbb{Z}_{k+1}^4 \ \middle| \ \sigma_0 + \sigma_1 = k, \sigma_2 = \sigma_3 = 0 \right\} \text{ AT } \\ &\Phi_1 \subseteq \left\{ (\sigma_0, \sigma_1, \sigma_2, \sigma_3) \in \mathbb{Z}_{k+1}^4 \ \middle| \ \sigma_0 = \sigma_1 = 0, \sigma_2 + \sigma_3 = k \right\} \text{ GC } \\ &|\Phi_0| = |\Phi_1| = 2^{k_{\rm b}} \le k + 1 \qquad \Phi = \Phi_0 \cup \Phi_1 \end{split}$$

#### Inner codebook

 $(\sigma_0, \sigma_1, \sigma_2, \sigma_3) = f_{b^j} (C_{i,t}^\intercal)$ 

$$C = \{(c_0, \dots, c_{\nu-1}) \mid c_j \in \Phi_{b_j}, (b_0, \dots, b_{\nu-1}) \in \mathcal{B}\}$$

$$\begin{array}{lll} \textbf{Mapping} & f_{\mathbf{a}}: \mathbb{B}^{k_{\mathbf{a}}} \to \mathcal{B} & \textbf{(bijection)} \\ & f_i: \mathbb{B}^{k_{\mathbf{b}}} \to \Phi_i \ (i \! \in \! \mathbb{B}) & \textbf{(bijection)} \\ & f_{\mathbf{d}}: \Sigma \to \mathbb{B} & f_{\mathbf{d}}(x) \! = \! \mathbb{1}[x \! \in \! \{2,3\}] \\ \end{array}$$

## **Encoding**

input: 
$$egin{aligned} c = (c_0, \dots, c_{n-1}) & c_i \in \mathbb{B}^{k_{\mathrm{a}}} \\ C = (C_0, \dots, C_{n-1}) & C_i \in \mathbb{B}^{k_{\mathrm{b}} \times \nu} \end{aligned}$$
 output:  $egin{aligned} s = (s_0, \dots, s_{n-1}) & s_i \in \Phi^{\nu} \end{aligned}$  (1) encode  $c$  to  $b = (b_0, \dots, b_{n-1}) \in \tilde{\mathcal{B}}^n$  
$$\hline c_0 & \cdots & c_{i-1} & c_i & \cdots & c_{n-1} \\ \hline k_{\mathrm{a}} & & & & & & & & \\ \hline b_0 & \cdots & & & & & & \\ \hline b_{i-1} & b_i & \cdots & & & & \\ \hline b_{n-1} & & & & & & \\ \hline \end{aligned}$$

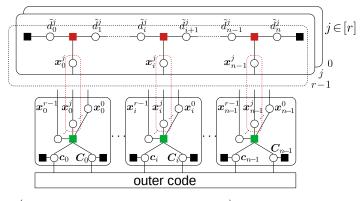
$$\boldsymbol{b}_i = \begin{cases} \boldsymbol{b}_i' & (\mathbbm{1}_{\rho,\ell,\delta}(\tilde{\boldsymbol{b}},\boldsymbol{b}_i') \!=\! 1) & \text{(unchanged)} \\ \overline{\boldsymbol{b}}_i' & (\mathbbm{1}_{\rho,\ell,\delta}(\tilde{\boldsymbol{b}},\overline{\boldsymbol{b}}_i') \!=\! 1) & \text{(invert)} \\ (\overline{b}b)^{\frac{\nu}{2}} & (\mathbbm{1}_{\rho,\ell,\delta}(\tilde{\boldsymbol{b}},(\overline{b}b)^{\frac{\nu}{2}}) \!=\! 1) & \text{(rebalance)} \\ (b\overline{b})^{\frac{\nu}{2}} & \text{(otherwise)} & \text{(rebalance)} \end{cases}$$

(2) encode  $(C_i, b_i)$  to  $s_i$   $(i \in [n])$ 

$$\begin{split} s_{i,t} \!=\! f_{b_{i,t}}(\boldsymbol{C}_{i,t}^{\intercal}) & \quad \boldsymbol{s}_i \!=\! (s_{i,0}, \dots, s_{i,\nu-1}) \!\in\! \boldsymbol{\Phi}^{\nu} \\ (t \!\in\! [\nu]) & \quad \boldsymbol{b}_i \!=\! (b_{i,0}, \dots, b_{i,\nu-1}) \!\in\! \tilde{\mathcal{B}} \\ & \quad \boldsymbol{C}_{i,t} : t \text{ th column of } \boldsymbol{C}_i \end{split}$$

$$\begin{array}{c} (\text{ex}) \ \nu = 6, \ |\mathcal{B}| = 16, k = 7, k_{\text{a}} = 4, k_{\text{b}} = 3 \\ \mathcal{B} : 0 : 001011 \ 8 : 110100 \ \Phi_0 : 0 : 0700 \ \Phi_1 : 0 : 0007 \\ 1 : 001101 \ 9 : 110010 \ 1 : 1600 \ 1 : 0016 \\ 2 : 001110 \ A : 110001 \ 2 : 2500 \ 2 : 0025 \\ 3 : 010011 \ B : 101100 \ 3 : 3400 \ 3 : 0034 \\ - : 010101 \ - : 101010 \ 4 : 4300 \ 4 : 0043 \\ 4 : 010110 \ C : 101001 \ 5 : 5200 \ 5 : 0052 \\ 5 : 011001 \ D : 100110 \ 6 : 6100 \ 6 : 0061 \\ 6 : 011010 \ E : 100101 \ 7 : 7000 \ 7 : 0070 \\ 7 : 011100 \ F : 100011 \ (Lee \ distance?) \\ \hline c_i : \ 0110(6) = 011010 \\ \hline c_i^* : \ 0110(6) = 011010 \\ \hline C_i^{\mathsf{T}} : \ 011(3) \ (1,4) = 0043 \\ 010(2) \ 011(3) \ (1,2) = 0025 \\ 011(3) \ (1,6) = 0061 \\ 100(4) \ (0,4) = 4300 \\ \end{array}$$

# Decoding



$$\begin{split} p\left(\boldsymbol{y}^{0},\dots,\boldsymbol{y}^{r-1},\boldsymbol{x}^{0},\dots,\boldsymbol{x}^{r-1},\tilde{\boldsymbol{d}}^{0},\dots,\tilde{\boldsymbol{d}}^{r-1},\boldsymbol{c},\boldsymbol{C}\right) \\ &= p\left(\boldsymbol{y}^{0},\dots,\boldsymbol{y}^{r-1}\Big|\boldsymbol{x}^{0},\dots,\boldsymbol{x}^{r-1},\tilde{\boldsymbol{d}}^{0},\dots,\tilde{\boldsymbol{d}}^{r-1},\boldsymbol{c},\boldsymbol{C}\right) \times \\ &\quad p\left(\tilde{\boldsymbol{d}}^{0},\dots,\tilde{\boldsymbol{d}}^{r-1}\right)p\left(\boldsymbol{x}^{0},\dots,\boldsymbol{x}^{r-1}\Big|\boldsymbol{c},\boldsymbol{C}\right)p\left(\boldsymbol{c},\boldsymbol{C}\right) \\ &= p\left(\boldsymbol{c}\right)p\left(\boldsymbol{C}\right)\prod_{j=0}^{r-1}p\left(\boldsymbol{y}^{j}\Big|\boldsymbol{x}^{j},\tilde{\boldsymbol{d}}^{j}\right)p\left(\tilde{\boldsymbol{d}}^{j}\right)p\left(\boldsymbol{x}^{j}\Big|\boldsymbol{c},\boldsymbol{C}\right) \\ &\simeq \left(\prod_{i=0}^{n-1}p\left(\boldsymbol{c}_{i}\right)p\left(\boldsymbol{C}_{i}\right)\right) \times \\ &\prod_{j=0}^{r-1}\left(p(\tilde{\boldsymbol{d}}_{0}^{j})\prod_{i=0}^{n-1}p\left(\left[\boldsymbol{y}^{j}\right]_{i\nu+d_{i}^{j}}^{i\nu+\nu-1+\tilde{\boldsymbol{d}}_{i+1}^{j}}\Big|\boldsymbol{x}_{i}^{j},\tilde{\boldsymbol{d}}_{i}^{j},\tilde{\boldsymbol{d}}_{i}^{j},\tilde{\boldsymbol{d}}_{i+1}^{j},\right)p\left(\tilde{\boldsymbol{d}}_{i+1}^{j}\Big|\tilde{\boldsymbol{d}}_{i}^{j}\right)p\left(\boldsymbol{x}_{i}^{j}\Big|\boldsymbol{c}_{i},\boldsymbol{C}_{i}\right)\right) \\ &\prod_{j=0}^{r-1}p\left(\boldsymbol{x}_{i}^{j}\Big|\boldsymbol{c}_{i},\boldsymbol{C}_{i}\right) = \prod_{j=0}^{r-1}p\left(\boldsymbol{b}_{i}^{j}\Big|\boldsymbol{c}_{i}\right)p\left(\boldsymbol{x}_{i}^{j}\Big|\boldsymbol{b}_{i}^{j},\boldsymbol{C}_{i}\right) = \prod_{j=0}^{r-1}p\left(\boldsymbol{b}_{i}^{j}\Big|\boldsymbol{c}_{i}\right)\prod_{t=0}^{\nu-1}p\left(\boldsymbol{x}_{i,t}^{j}\Big|\boldsymbol{b}_{i,t}^{j},\boldsymbol{C}_{i,t}\right) \\ &\boldsymbol{b}_{i}^{j} = \left(f_{\mathrm{d}}(\boldsymbol{x}_{i,0}^{j}),\dots,f_{\mathrm{d}}(\boldsymbol{x}_{i,\nu-1}^{j})\right) \\ &\boldsymbol{p}\left(\boldsymbol{x}_{i,t}^{j}\Big|\boldsymbol{b}_{i,t}^{j},\boldsymbol{C}_{i,t}\right) = \sigma_{\boldsymbol{x}_{i}^{j},\boldsymbol{k}}^{j}/\boldsymbol{k} \quad \left(\boldsymbol{x}_{i,t}^{j} \in \boldsymbol{\Sigma},\boldsymbol{b}_{i,t}^{j} \in \boldsymbol{\mathbb{B}},\boldsymbol{C}_{i,t}^{\mathsf{T}} \in \boldsymbol{\mathbb{B}}^{k_{\mathrm{b}}}\right) \end{split}$$