

Inner code

constrained coding

- * run length
- * GC-balance
- * prior: $p(x)$ for IDS

channel coding

- * 4-ary IDS/asymmetric
- * (multi-read)

inner code performance:

- * run length distribution
- * GC-weight distribution (sliding window)
- * synchronization: cross entropy of APP(?)

alphabet: $\Sigma = \{A, T, G, C\}$
 segment size: β [symbols]
 segment: $\mathbf{u} = (u_0, \dots, u_{\beta-1}) \in \Sigma^\beta$

outer code: $\tilde{\mathcal{C}} \subset \mathbb{F}_2^n$ ($b < 2\beta$)

inner code: $\mathcal{C}_i \subset \Sigma^\beta$ ($i \in [\nu]$)

(num CW)

$$|\mathcal{C}_i| = 2^b$$

(rate)

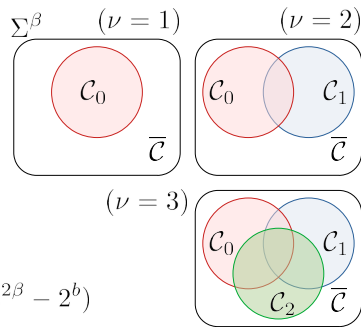
$$R = b/2\beta$$

(encoding)

$$\phi_i : \mathbb{F}_2^b \rightarrow \mathcal{C}_i$$

(forbidden set)

$$\bar{\mathcal{C}} \subseteq \Sigma^\beta \setminus \bigcup_{i \in [\nu]} \mathcal{C}_i \quad (|\bar{\mathcal{C}}| \leq 2^{2\beta} - 2^b)$$



vector over Σ : $\mathbf{u} = (u_0, \dots, u_{n-1}) \in \Sigma^n$

weight: $w(\mathbf{u})_x = |\{i \in [n] \mid u_i = x\}|$
 $(x \in \Sigma)$

balance: $f_B(\mathbf{u}) = w_G(\mathbf{u}) + w_C(\mathbf{u}) - w_A(\mathbf{u}) - w_T(\mathbf{u})$

max run length: $f_R(\mathbf{u})$

[example 1]

forbidden set $\bar{\mathcal{C}}$:

ϕ (empty)

$\bar{\mathcal{C}}_{\omega,*} = \{\mathbf{u} \in \Sigma^\beta \mid |f_B(\mathbf{u})| \geq \omega\}$ (GC-balance)

$\bar{\mathcal{C}}_{*,\lambda} = \{\mathbf{u} \in \Sigma^\beta \mid f_R(\mathbf{u}) \geq \lambda\}$ (run length)

$\bar{\mathcal{C}}_{\omega,\lambda} = \bar{\mathcal{C}}_{\omega,*} \cup \bar{\mathcal{C}}_{*,\lambda}$ (both)

[example 2] $\beta = 4, b = 7$

$$\mathcal{C}_i = \{(\phi(x_0 y_0), \dots, \phi(x_3 y_3)) \mid (x_0, \dots, x_3) \in \mathcal{B}_i, (y_0, \dots, y_3) \in \mathbb{B}^4\}$$

$\mathcal{B}_0 : 0011, 0101, 0110, 1001, 1010, 1100, \mathbf{0010, 0100}$

$\mathcal{B}_1 : 0011, 0101, 0110, 1001, 1010, 1100, \mathbf{1101, 1011}$

$$|\mathcal{C}_i| = 8 \times 16 = 128 = 2^b$$

\mathbf{v}	$\phi(\mathbf{v})$
00	A
01	T
10	G
11	C

[run-length of random sequence]

prob. of run-length ℓ :

$$p(\ell) = \left(\frac{1}{q}\right)^{\ell-1} \left(\frac{q-1}{q}\right) = \frac{q-1}{q^\ell}$$

expectation:

$$E(\ell) = \sum_{\ell=1}^n \ell p(\ell) = \sum_{\ell=1}^n \frac{(q-1)\ell}{q^\ell}$$

$$= (q-1) \sum_{\ell=1}^n \frac{\ell}{q^\ell} = \frac{q}{q-1} \left(1 - \frac{1}{q^n}\right) - \frac{n}{q^n}$$

[GC-balance of random sequence]

window size: ℓ prob. binary vector having weight w :

$$p_\ell(w) = \binom{\ell}{w} / 2^\ell$$

absolute GC-balance: $\sum_{w=0}^\ell p_\ell(w) |\ell - 2w|$

NB-IDS channel

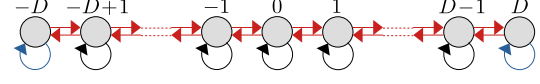
block length: n
 alphabet: Σ ($|\Sigma| = q$)
 error prob.: p_{id} ($< \frac{1}{2}$) (ins/del)
 p_s (sub)
 input: $\mathbf{x} = (x_0, \dots, x_{n-1}) = \Sigma^n$
 output: $\mathbf{y} = (y_0, \dots, y_{n'-1}) = \Sigma^{n'}$
 $(n-D \leq n' \leq n+D)$

[transmission]

1) drift vector: $\mathbf{d} = (d_0, \dots, d_{n-1}, d_n) \in \mathcal{D}^{n+1}$
 2) intermediate vector: $\mathbf{z} = (z_0, \dots, z_{n'-1}) \in \Sigma^{n'}$
 $z_j = x_i$ ($j \in [i + d_i, i + d_{i+1}]$)
 $i \in [n], n' = n + d_n$
 3) output vector: $p(y_i|z_i) = \begin{cases} 1 - p_s & (y_i = z_i) \\ \frac{p_s}{q-1} & (y_i \neq z_i) \end{cases}$

max drift: D
 set of drift values: $\mathcal{D} = \{-D, \dots, 0, \dots, D\}$
 drift vector: $\mathbf{d} = (d_0=0, \dots, d_{n-1}, d_n) \in \mathcal{D}^{n+1}$

$$p(d_{i+1}|d_i) = \begin{cases} 1 - 2p_{id} & (d_{i+1} = d_i, |d_i| < D) \rightarrow \\ 1 - p_{id} & (d_{i+1} = d_i, |d_i| = D) \rightarrow \\ p_{id} & (|d_{i+1} - d_i| = 1, |d_{i+1}| \leq D) \rightarrow \\ 0 & (\text{otherwise}) \end{cases}$$



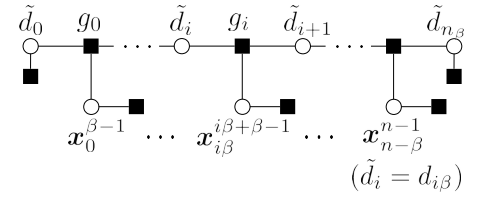
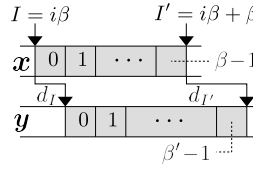
APP by FBA

$$p(\mathbf{x}, \mathbf{y}, \mathbf{d}) = p(\mathbf{y}|\mathbf{x}, \mathbf{d})p(\mathbf{x}, \mathbf{d}) = p(\mathbf{y}|\mathbf{x}, \mathbf{d})p(\mathbf{x})p(\mathbf{d})$$

$$= p(d_0) \prod_{i=0}^{n-1} p\left(\mathbf{y}_{i+d_i}^{i+d_{i+1}} \middle| x_i, d_i, d_{i+1}\right) p(d_{i+1}|d_i) p(x_i)$$

segment level synchronization: $n = n_\beta \beta$

$$= p(d_0) \prod_{i=0}^{n_\beta-1} \frac{p\left(\mathbf{y}_{I+d_I}^{I'-1+d_{I'}} \middle| \mathbf{x}_I^{I'-1}, d_I, d_{I'}\right) p(d_{I'}|d_I) p\left(\mathbf{x}_I^{I'-1}\right)}{(I = i\beta, I' = i\beta + \beta)}$$

[LUT] (for a large D)

$$G\left(\beta', \mathbf{y}_0^{\beta'-1}\right) \text{ list of } \left[\mathbf{x}_0^{\beta-1}, p_0\left(\mathbf{y}_0^{\beta'-1} \middle| \mathbf{x}_0^{\beta-1}, 0, d'\right) p_0(d'|0) (> p_{th}) \right]$$

$$p_\alpha\left(\mathbf{y}_{\alpha'}^{\beta'-1} \middle| \mathbf{x}_\alpha^{\beta-1}, d, d'\right) p_\alpha(d'|d) = \begin{matrix} (d' = \beta' - \beta) \\ (d = \alpha' - \alpha) \end{matrix}$$

$$\sum_{\delta \in \{-1, 0, 1\}} p\left(\mathbf{y}_{\alpha'+\delta}^{\alpha'+\delta} \middle| \mathbf{x}_\alpha, d, d + \delta\right) p(d + \delta|d)$$

$$\times p_{\alpha+1}\left(\mathbf{y}_{\alpha'+\delta+1}^{\beta'} \middle| \mathbf{x}_{\alpha+1}^{\beta-1}, d + \delta, d'\right) p_{\alpha+1}(d'|d + \delta)$$

Gtable:
 int beta
 double Pid
 double Ps
 double Pth
 LIST[Len][Y]

LIST:
 int num
 XP[num]

XP:
 int x
 float p

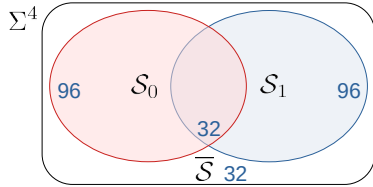
constrained coding

- * run length
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channel coding

- * IDS/asymmetric
- * multi-read

4-ary



$$|\bar{S}| = 32$$

$$|S| = 224$$

$$|S_0| = |S_1| = 128$$

$$|S_0 \cap S_1| = 32$$

$$\bar{S} = \{A, T\}^4 \cup \{G, C\}^4$$

$$S = S_0 \cup S_1$$

$$S \cup \bar{S} = \Sigma^4$$

0000
0001,0010,0011,0100,0110,0111
0101,1010
1110,1101,1100,1011,1001,1000
1111

alphabet:

$$\Sigma = \{A, T, G, C\}$$

segment size:

$$\beta \quad [\text{symbols}]$$

symbol size:

$$b = 2\beta - 1 \quad [\text{bits}]$$

inner code rate:

$$R = (2\beta - 1)/2\beta$$

set of binary vectors:

$$\begin{aligned} \mathcal{B}_0 &= \{(0, b_1, \dots, b_{\beta-1}) \mid b_i \in \mathbb{B}\} & |\mathcal{B}_0| &= |\mathcal{B}_1| = 2^{\beta-1} \\ \mathcal{B}_1 &= \{(1, b_1, \dots, b_{\beta-1}) \mid b_i \in \mathbb{B}\} & |\mathcal{B}_a| &= |\mathcal{B}_c| = 2 \\ \mathcal{B}_a &= \{(b, \bar{b}, b, \bar{b}, \dots) \in \mathbb{B}^\beta \mid b \in \mathbb{B}\} \\ \mathcal{B}_c &= \{(b, b, b, b, \dots) \in \mathbb{B}^\beta \mid b \in \mathbb{B}\} \\ \tilde{\mathcal{B}}_i &= (\mathcal{B}_i \setminus \mathcal{B}_c) \cup \mathcal{B}_a \quad (i \in \mathbb{B}) \end{aligned}$$

segment:

$$s = (s_0, s_1, \dots, s_{\beta-1}) \in \Sigma^\beta$$

symbol mapping:

$$\begin{aligned} \phi: \mathbb{B} \times \mathbb{B} &\rightarrow \Sigma \\ \phi: \mathbb{B}^\beta \times \mathbb{B}^\beta &\rightarrow \Sigma^\beta \\ \phi(x, y) &= (\phi(x_0, y_0), \dots, \phi(x_{\beta-1}, y_{\beta-1})) \\ x &= (x_0, \dots, x_{\beta-1}) \in \mathbb{B}^\beta \\ y &= (y_0, \dots, y_{\beta-1}) \in \mathbb{B}^\beta \end{aligned}$$

v	$\phi(v)$
00	A
01	T
10	G
11	C

set of segments:

$$\bar{S} = \{\phi(x, y) \mid x \in \mathcal{B}_c, y \in \mathbb{B}^\beta\} = \{A, T\}^\beta \cup \{G, C\}^\beta$$

$$S_i = \{\phi(x, y) \mid x \in \tilde{\mathcal{B}}_i, y \in \mathbb{B}^\beta\} \quad (i \in \mathbb{B})$$

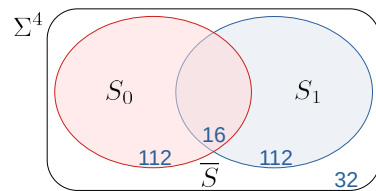
encoding func.:

$$\begin{aligned} \psi_i: \mathbb{B}^{\beta-1} \times \mathbb{B}^\beta &\rightarrow S_i \quad (i \in \mathbb{B}) \\ \psi_i(\tilde{x}, y) &= \phi(x, y) \\ x &= \begin{cases} (i, \tilde{x}) & ((i, \tilde{x}) \notin \mathcal{B}_c) \\ (\bar{i}, i, \bar{i}, i, \dots) & ((i, \tilde{x}) \in \mathcal{B}_c) \end{cases} \end{aligned}$$

$$\begin{aligned} (S_0 \cup S_1) &= (\Sigma^\beta \setminus \bar{S}) \\ |S_0| &= |S_1| = 2^b \\ |S_0 \cap S_1| &= 2^{b+1} \\ |\bar{S}| &= 2^{b+1} \end{aligned}$$

(base caller)
(NB marker)

set partitioning:



$$|\bar{S}| = 32$$

$$|S| = 224$$

$$|S_0| = |S_1| = 128$$

$$|S_0 \cap S_1| = 16$$

$$\bar{S} = \{A, T\}^4 \cup \{G, C\}^4$$

$$S = S_0 \cup S_1$$

$$S \cup \bar{S} = \Sigma^4$$

$$S_0 \cap S_1 =$$

$$\{x_0x_1x_2x_3 \mid x_0x_2 \in \{A, T\}^2, x_1x_3 \in \{G, C\}^2\}$$