

Alphabet: $\Sigma = \{0, 1, 2, 3\} \leftrightarrow \{A, T, G, C\}$

Code length: ν (even) Block length: $N = \nu n$ symbol

(n codewords/block)

[ex]
$$\nu = 6$$
, $|\mathcal{B}| = 16$, $k = 7$, $k_a = 4$, $k_b = 3$

15:		Φ_0 :	Ψ_1 :
0:001011	8:110100	0:0700	0:0007
1:001101	9:110010	1:1600	1:0016
2:001110	A:110001	2:2500	2:0025
3:010011	B:101100	3:3400	3:0034
-:010101	-:101010	4:4300	4:0043
4:010110	C:101001	5:5200	5:0052
5:011001	D:100110	6:6100	6:0061
6:011010	E:100101	7:7000	7:0070
7:011100	F:100011	(Lee	distance?)

[Binary constraint vector set]

$$\mathcal{B} \subset \mathbb{B}^{\nu} \quad (|\mathcal{B}| \ge 2^{k_{a}}, \mathcal{B} \cap \{(01)^{\frac{\nu}{2}}, (10)^{\frac{\nu}{2}}\} = \phi)$$

$$\tilde{\mathcal{B}} = \mathcal{B} \cup \{(01)^{\frac{\nu}{2}}, (10)^{\frac{\nu}{2}}\}$$
 (re-balance)

[Composite symbol sets]

$$\begin{split} \Phi_0 &\subseteq \left\{ (\sigma_0, \sigma_1, \sigma_2, \sigma_3) \in \mathbb{Z}_{k+1}^4 \ \middle| \ \sigma_0 + \sigma_1 = k, \sigma_2 = \sigma_3 = 0 \right\} \text{ AT } \\ \Phi_1 &\subseteq \left\{ (\sigma_0, \sigma_1, \sigma_2, \sigma_3) \in \mathbb{Z}_{k+1}^4 \ \middle| \ \sigma_0 = \sigma_1 = 0, \sigma_2 + \sigma_3 = k \right\} \text{ GC } \\ &|\Phi_0| = |\Phi_1| = 2^{k_b} \le k + 1 \qquad \Phi = \Phi_0 \cup \Phi_1 \end{split}$$

[Inner codebook]

$$C \subseteq \{(c_0,\ldots,c_{\nu-1}) \mid c_j \in \Phi_{b_j}, (b_0,\ldots,b_{\nu-1}) \in \mathcal{B}\}$$

[Encode]

mapping:
$$f_{\mathrm{e}}^{\mathrm{a}}: \mathbb{B}^{k_{\mathrm{a}}n} \to \mathcal{B}^n$$
 (lossy: SITA2024)

$$f_e^i: \mathbb{B}^{k_{\mathrm{b}}} \to \Phi_i \ (i \in \mathbb{B})$$
 (bijection)

input:
$$\boldsymbol{c}^{\mathrm{a}} = (\boldsymbol{c}_0^{\mathrm{a}}, \dots, \boldsymbol{c}_{n-1}^{\mathrm{a}}) \quad \boldsymbol{c}_i^{\mathrm{a}} \in \mathbb{B}^{\kappa_{\mathrm{a}}}$$

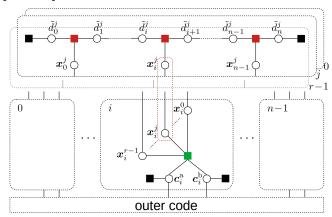
$$egin{aligned} f_{\mathrm{e}}^i : \mathbb{B}^{k_{\mathrm{b}}} &
ightarrow \Phi_i \ (i\!\in\!\mathbb{B}) \end{aligned} egin{aligned} ext{(bijection)} \ c^{\mathrm{a}} &= (c_0^{\mathrm{a}}, \dots, c_{n-1}^{\mathrm{a}}) & c_i^{\mathrm{a}} &\in \mathbb{B}^{k_{\mathrm{a}}} \ c^{\mathrm{b}} &= (c_0^{\mathrm{b}}, \dots, c_{n-1}^{\mathrm{b}}) & c_i^{\mathrm{b}} &\in \mathbb{B}^{k_{\mathrm{b}}
u}, c_{i,j}^{\mathrm{b}} &= [c_i^{\mathrm{b}}]_{j
u}^{j
u+
u-1} \end{aligned}$$

output:
$$s = (s_0, \ldots, s_{n-1})$$
 $s_i = (s_{i,0}, \ldots, s_{i,\nu-1}) \in \Phi^{\nu}$

(1)
$$b = (b_0, \dots, b_{n-1}) = f_e^a(c^a)$$
 $b_i = (b_{i,0}, \dots, b_{i,\nu-1}) \in \tilde{\mathcal{B}}$

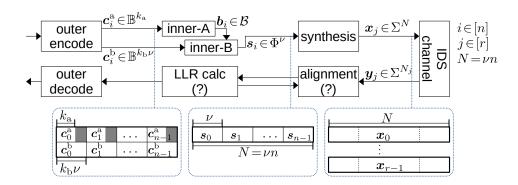
(2)
$$s_{i,j} = f_{e}^{b_i}(\boldsymbol{c}_{i,j}^{b}) \quad (i \in [n], j \in [\nu])$$

[Decode]



$$\tilde{d}_i^j = d_{i\nu}^j$$
 (symbol level) $j \in [r]$

$$\begin{split} &p\left(\boldsymbol{y}^{0},\ldots,\boldsymbol{y}^{r-1},\boldsymbol{x}^{0},\ldots,\boldsymbol{x}^{r-1},\tilde{\boldsymbol{d}}^{0},\ldots,\tilde{\boldsymbol{d}}^{r-1},\boldsymbol{c}^{\mathbf{a}},\boldsymbol{c}^{\mathbf{b}}\right) \\ &= p\left(\boldsymbol{y}^{0},\ldots,\boldsymbol{y}^{r-1}\Big|\boldsymbol{x}^{0},\ldots,\boldsymbol{x}^{r-1},\tilde{\boldsymbol{d}}^{0},\ldots,\tilde{\boldsymbol{d}}^{r-1},\boldsymbol{c}^{\mathbf{a}},\boldsymbol{c}^{\mathbf{b}}\right)p\left(\tilde{\boldsymbol{d}}^{0},\ldots,\tilde{\boldsymbol{d}}^{r-1}\right)p\left(\boldsymbol{x}^{0},\ldots,\boldsymbol{x}^{r-1}\Big|\boldsymbol{c}^{\mathbf{a}},\boldsymbol{c}^{\mathbf{b}}\right)p\left(\boldsymbol{c}^{\mathbf{a}},\boldsymbol{c}^{\mathbf{b}}\right)\\ &= p\left(\boldsymbol{c}^{\mathbf{a}}\right)p\left(\boldsymbol{c}^{\mathbf{b}}\right)\prod_{j=0}^{r-1}p\left(\boldsymbol{y}^{j}\Big|\boldsymbol{x}^{j},\tilde{\boldsymbol{d}}^{j}\right)p\left(\tilde{\boldsymbol{d}}^{j}\right)p\left(\boldsymbol{x}^{j}\Big|\boldsymbol{c}^{\mathbf{a}},\boldsymbol{c}^{\mathbf{b}}\right)\\ &= \left(\prod_{i=0}^{n-1}p\left(\boldsymbol{c}^{\mathbf{a}}_{i}\right)p\left(\boldsymbol{c}^{\mathbf{b}}\right)\right)\prod_{j=0}^{r-1}\left(p\left(\tilde{\boldsymbol{d}}^{j}_{0}\right)\prod_{i=0}^{n-1}p\left(\boldsymbol{y}^{j}\Big|_{i\nu+\tilde{\boldsymbol{d}}^{j}_{i}}^{i\nu+\nu-1+\tilde{\boldsymbol{d}}^{j}_{i+1}}\Big|\boldsymbol{x}^{j}_{i},\tilde{\boldsymbol{d}}^{j}_{i},\tilde{\boldsymbol{d}}^{j}_{i},\tilde{\boldsymbol{d}}^{j}_{i+1},\right)p\left(\tilde{\boldsymbol{d}}^{j}_{i+1}\Big|\tilde{\boldsymbol{d}}^{j}_{i}\right)p\left(\boldsymbol{x}^{j}_{i}\Big|\boldsymbol{c}^{\mathbf{a}},\boldsymbol{c}^{\mathbf{b}}\right)\right) \end{split}$$



Code length: ν (even)

Binary constraint vector set: $\mathcal{B} \subset \mathbb{B}^{\nu}$ $(|\mathcal{B}| \geq 2^{k_{\mathrm{a}}})$ Composite symbol set:

$$\begin{split} &\Phi_0\subseteq\left\{(\sigma_0,\sigma_1,\sigma_2,\sigma_3)\in\mathbb{Z}_{k+1}^4\ \middle|\ \sigma_0+\sigma_1=k,\sigma_2=\sigma_3=0\right\} \text{ AT }\\ &\Phi_1\subseteq\left\{(\sigma_0,\sigma_1,\sigma_2,\sigma_3)\in\mathbb{Z}_{k+1}^4\ \middle|\ \sigma_0=\sigma_1=0,\sigma_2+\sigma_3=k\right\} \text{ GC }\\ &|\Phi_0|=|\Phi_1|=2^{k_{\mathrm{b}}}\!\leq\!k\!+\!1\qquad \Phi=\!\Phi_0\!\cup\!\Phi_1 \end{split}$$

Inner codebook:

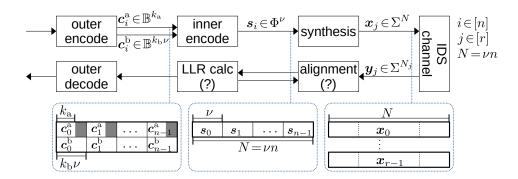
$$C \subseteq \{(c_0, \ldots, c_{\nu-1}) \mid c_j \in \Phi_{b_j}, (b_0, \ldots, b_{\nu-1}) \in \mathcal{B}\}$$

Encode:

$$\begin{split} f_{\mathrm{e}}^{\mathrm{a}} : \mathbb{B}^{k_{\mathrm{a}}} &\rightarrow \mathcal{B} \qquad \text{(lossy)} \\ f_{\mathrm{e}}^{i} : \mathbb{B}^{k_{\mathrm{b}}} &\rightarrow \Phi_{i} \ (i \in \mathbb{B}) \quad \text{(bijection)} \\ f_{\mathrm{e}}(\boldsymbol{c}_{i}^{\mathrm{a}}, \boldsymbol{c}_{i}^{\mathrm{b}}) = \boldsymbol{s}_{i} = (s_{i,0}, \dots, s_{i,\nu-1}) : \mathbb{B}^{k_{\mathrm{a}}} \times \mathbb{B}^{k_{\mathrm{b}}\nu} &\rightarrow \Phi^{\nu} \\ s_{i,j} = f_{\mathrm{e}}^{b_{i}}(\boldsymbol{c}_{i,j}^{\mathrm{b}}) & (b_{0}, \dots, b_{\nu-1}) = f_{\mathrm{e}}^{\mathrm{a}}(\boldsymbol{c}_{i}^{\mathrm{a}}) \\ c_{i,j}^{b} = [\boldsymbol{c}_{i}^{\mathrm{b}}]_{j\nu}^{j\nu+\nu-1} \end{split}$$

```
(example) \nu = 6, |\mathcal{B}| = 16, k = 7, k_a = 4, k_b = 3,
B: 0:001011 8:110100
                         \Phi_0:0:0700 \Phi_1:0:0007
    1:001101 9:110010
                                         1:0016
                             1:1600
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                             2:2500
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   3:010011 B:101100
                             3:3400
                                         3:0034
    -:010101
                             4:4300
                                         4:0043
              -:101010
    4:010110 C:101001
                             5:5200
                                          5:0052
   5:011001 D:100110
                             6:6100
                                         6:0061
                                         7:0070
    6:011010 E:100101
                             7:7000
    7:011100 F:100011
                                    Lee dist?
```

Decode:



Code length: ν (even)

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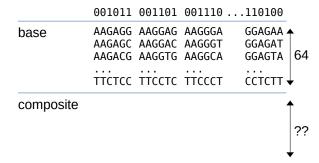
$$\begin{split} &\Phi_0\subseteq\left\{(\sigma_0,\sigma_1,\sigma_2,\sigma_3)\in\mathbb{Z}_{k+1}^4\;\big|\;\sigma_0+\sigma_1=k,\sigma_2=\sigma_3=0\right\}\;\text{AT}\\ &\Phi_1\subseteq\left\{(\sigma_0,\sigma_1,\sigma_2,\sigma_3)\in\mathbb{Z}_{k+1}^4\;\big|\;\sigma_0=\sigma_1=0,\sigma_2+\sigma_3=k\right\}\;\text{GC}\\ &|\Phi_0|=|\Phi_1|=2^{k_{\mathrm{b}}}\!\leq\!k\!+\!1\qquad\Phi=\Phi_0\cup\Phi_1 \end{split}$$

Inner codebook:

$$\mathcal{C} \subseteq \left\{ (c_0, \dots, c_{\nu-1}) \mid c_j \in \Phi_{b_j}, (b_0, \dots, b_{\nu-1}) \in \mathcal{B} \right\}$$

Encode:

Decode:



multi-base composite (Manchester coding)

 $\nu = 1 : (A,C,G,T)$

 $\nu\!=\!2$: (AA,AC,AG,AT, CA,CC,CG,CT, GA,GC,GG,GT, TA,TC,TG,TT)

 $\nu = 3$: (AAA, AAC, ..., TTT)

001011 001101 001110 ...110100 AAGAGG AAGGAG AAGGGA GGAGAA ♠

AAGAGC AAGGAC AAGGGT GGAGAC AAGACG AAGGTG AAGGCA GGAGCA

composite

base