6CCS3CFL-CW3

```
Q1
```

```
Grammer:
Stmts ::= Stmt; Stmts \mid Stmt
Block := \{Stmts\}|Stmt
AExp := Te + AExp \mid Te - AExp \mid Te
Te ::= Fa * Te | Fa / Te | Fa \% Te | Fa
Fa := (AExp) \mid num \mid id
Stmt := skip
        id := Aexp
        write (id) |
        write (string) |
        write id |
        write string |
        read (id) |
        read id |
        if BExp then Block else Block |
        while BExp do Block
Type equation here.
BExp := AExp == AExp
        AExp! = AExp
        AExp < AExp
        AExp \leq AExp
        AExp > AExp
        AExp \ge AExp
        (BExp) \&\& BExp
        (BExp) \mid\mid BExp\mid
        true|
        false|
        (BExp)
```

For making my parser recognise tokens, I had to change the implementation of a few essential functions.

```
I introduced a few new implicit definitions which correspond to 
Parser List Token . Token
```

And also introduced some atomic parsers for some specific tokens.

Boolen Expressions

```
// boolean expressions with some simple nesting

lazy val BExp: Parser[List[Token], BExp] =

(AExp ~ T_OP("=") ~ AExp).map[BExp]{ case x ~ _ ~ z => Bop("==", x, z) } ||

(AExp ~ T_OP("!=") ~ AExp).map[BExp]{ case x ~ _ ~ z => Bop("!=", x, z) } ||

(AExp ~ T_OP("\) ~ AExp).map[BExp]{ case x ~ _ ~ z => Bop("\], x, z) } ||

(AExp ~ T_OP("\) ~ AExp).map[BExp]{ case x ~ _ ~ z => Bop("\], x, z) } ||

(AExp ~ T_OP("\) ~ AExp).map[BExp]{ case x ~ _ ~ z => Bop("\), x, z) } ||

(AExp ~ T_OP("\) ~ AExp).map[BExp]{ case x ~ _ ~ z => Bop("\), x, z) } ||

(AExp ~ T_OP("\) ~ AExp).map[BExp]{ case x ~ _ ~ z => Bop("\), x, z) } ||

(T_LPAREN_N ~ BExp ~ T_RPAREN_N ~ T_OP("\), w BExp).map[BExp]{ case _ ~ y ~ _ ~ _ ~ v => And(y, v) } ||

(T_LPAREN_N ~ BExp ~ T_RPAREN_N ~ T_OP("\)|" ~ BExp).map[BExp]{ case _ ~ y ~ _ ~ _ ~ v => Or(y, v) } ||

(T_KMD("\true").map[BExp]{ _ => True }) ||

(T_KMD("\true").map[BExp]{ _ => False }) ||

(T_LPAREN_N ~ BExp ~ T_RPAREN_N).map[BExp]{ case _ ~ x ~ _ => x }
```

Arithmetic Expressions

```
// arithmetic expressions
lazy val AExp: Parser[List[Token], AExp] =
    (Te ~ T_OP("+") ~ AExp).map[AExp]{    case x ~ _ ~ z => Aop("+", x, z) } ||
    (Te ~ T_OP("-") ~ AExp).map[AExp]{    case x ~ _ ~ z => Aop("-", x, z) } ||
    Te

lazy val Te: Parser[List[Token], AExp] =
    (Fa ~ T_OP("*") ~ Te).map[AExp]{    case x ~ _ ~ z => Aop("*", x, z) } ||
    (Fa ~ T_OP("/") ~ Te).map[AExp]{    case x ~ _ ~ z => Aop("/", x, z) } ||
    (Fa ~ T_OP("%") ~ Te).map[AExp]{    case x ~ _ ~ z => Aop("%", x, z) } ||
    Fa

lazy val Fa: Parser[List[Token], AExp] =
    (T_LPAREN_N ~ AExp ~ T_RPAREN_N).map{    case _ ~ y ~ _ => y } ||
    IdParserToken.map(Var) ||
    NumParserToken.map(Num)
```

Single Statements

Had to introduce extra cases of write since there were various different ways it was implemented in test programs.

```
// a single statement

//* we need 4 types of writes here
write id
write "string"
write (id)
write ("string")

*/

/ lazy val Stmt: Parser[List[Token], Stmt] =

((T_KWD("skip").map[Stmt]{_ => Skip }) ||
    (IdParserToken ~ T_OP(":=") ~ AExp).map[Stmt]{ case x ~ _ ~ z => Assign(x, z) } ||
    (I_KWD("write") ~ I_LPAREN_N ~ IdParserToken ~ I_RPAREN_N).map[Stmt]{ case _ ~ y ~ _ => WriteVar(y) } ||
    (I_KWD("write") ~ StrParserToken).map[Stmt]{ case _ ~ y => WriteStr(y) } ||
    (I_KWD("write") ~ IdParserToken).map[Stmt]{ case _ ~ y => WriteVar(y)} ||
    (I_KWD("write") ~ I_LPAREN_N ~ StrParserToken ~ I_RPAREN_N).map[Stmt]{ case _ ~ y ~ _ => WriteStr(y) } ||
    (I_KWD("write") ~ IdParserToken).map[Stmt]{ case _ ~ y ~> Read(y) } ||
    (I_KWD("read") ~ IdParserToken).map[Stmt]{ case _ ~ y ~> Read(y) } ||
    (I_KWD("if") ~ BExp ~ I_KWD("then") ~ Block ~ I_KWD("else") ~ Block)
    .map[Stmt]{ case _ ~ y ~ _ ~ u ~ _ ~ w => If(y, u, w) } ||
    (I_KWD("while") ~ BExp ~ I_KWD("do") ~ Block).map[Stmt]{ case _ ~ y ~ _ ~ w => While(y, w) })
```

Compound Statements

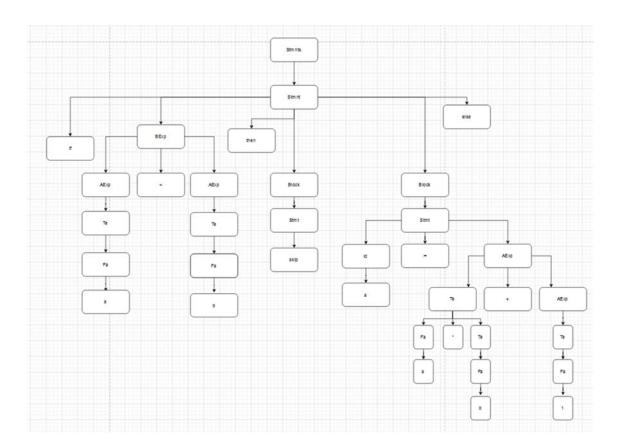
```
// statements
v lazy val Stmts: Parser[List[Token], Block] =
   (Stmt ~ T_SEMI ~ Stmts).map[Block]{ case x ~ _ ~ z => x :: z } ||
   (Stmt.map[Block]{ s => List(s) })
```

Blocks

```
// blocks (enclosed in curly braces)
lazy val Block: Parser[List[Token], Block] =
  ((T_LPAREN_C ~ Stmts ~ T_RPAREN_C).map{ case _ ~ y ~ _ => y } ||
  (Stmt.map(s => List(s))))
```

Q2

```
Output of running the respective programs:
       Fib
HashSet(List(WriteStr("Fib"), Read(n), Assign(minus1,Num(0)),
Assign(minus2,Num(1)),
While(Bop(>,Var(n),Num(0)),List(Assign(temp,Var(minus2)),
Assign(minus2,Aop(+,Var(minus1),Var(minus2))),                               Assign(minus1,Var(temp)),
Assign(n,Aop(-,Var(n),Num(1))))),    WriteStr("Result"),    WriteVar(minus2)))
        Loops
While(Bop(<,Num(0),Var(x)),List(While(Bop(<,Num(0),Var(y)),List(While(Bop(<,Nu
m(0),Var(z)),List(Assign(z,Aop(-,Var(z),Num(1))))),                          Assign(z,Var(start)),
Var(x), Num(1))))))
         Primes
Set(List(Assign(end,Num(100)), Assign(n,Num(2)),
While(Bop(<,Var(n),Var(end)),List(Assign(f,Num(2)),                          Assign(tmp,Num(0)),
While(And(Bop(<,Var(f),Aop(+,Aop(/,Var(n),Num(2)),Num(1))),Bop(==,Var(tmp),Num
(0))),List(If(Bop(==,Aop(*,Aop(/,Var(n),Var(f)),Var(f)),Var(n)),List(Assign(tm
If(Bop(==,Var(tmp),Num(0)),List(WriteVar(n)),List(Skip)),
Assign(n,Aop(+,Var(n),Num(1))))))
         Collatz
Set(List(Assign(bnd,Num(1)),
While(Bop(<,Var(bnd),Num(101)),List(WriteVar(bnd), WriteStr(": "),
Assign(n,Var(bnd)), Assign(cnt,Num(0)),
While(Bop(>,Var(n),Num(1)),List(WriteVar(n), WriteStr(","),
If(Bop(==,Aop(%,Var(n),Num(2)),Num(0)),List(Assign(n,Aop(/,Var(n),Num(2)))),Li
st(Assign(n,Aop(+,Aop(*,Num(3),Var(n)),Num(1))))),
Assign(cnt,Aop(+,Var(cnt),Num(1))))),    WriteStr(" => "),    WriteVar(cnt),
WriteStr("\n"), Assign(bnd,Aop(+,Var(bnd),Num(1))))))
The tokenised output has \n and quotes since my implementation of the program manipulates
strings in the eval stmt function.
Output for (if (a < b) then skip else a := a * b + 1)
Set(List(If(Bop(<,Var(a),Var(b)),List(Skip),List(Assign(a,Aop(+,Aop(*,Var(a),Var(b)),Num(1))))))
Parse Tree for (if (a < b) then skip else a := a * b + 1) is denoted below.
```



Q3

The Time measurements are outlined below.

1. Fibonacci

```
ibWaiting for User Input....
Result8Code Run Time: 8.0255445 s
```

The formatting of the code is not ideal since there are no newlines in the actual code of the program.

The write function doesn't use **println** instead it uses **print**.

2. Three Nested Loops

```
Loop Program - start: 100
Code Run Time: 0.1543181 s

Loop Program - start: 500
Code Run Time: 17.3906449 s

Loop Program - start: 800
Code Run Time: 68.5276533 s
```

Started off with 100 which took less than half a second.

Then 500 took 18 seconds to execute and lastly having a start value of 800 took just a little more than 1 minute.

3. Factors

```
2
3
5
7
11
13
17
19
23
29
31
37
41
43
47
53
59
61
67
71
73
79
83
89
97
Code Run Time: 4.2165406 s
```

The formatting of this code is better since when the screenshot was taken, the write function was using **println** to print output. But since it is now using **print**, the new output is denoted below:

```
2357111317192329313741434753596167717379838997Code Run Time: 1.0296003 s
```

4. Collatz

```
50,425,1276,638,319,958,479,1438,719,2158,1079,3238,16
, => 118
98: 98,49,148,74,37,112,56,28,14,7,22,11,34,17,52,26,1
99: 99,298,149,448,224,112,56,28,14,7,22,11,34,17,52,2
100: 100,50,25,76,38,19,58,29,88,44,22,11,34,17,52,26,
Code Run Time: 0.3878173 s
```