

Tutorial 3

Ans1 \Rightarrow pseudofunction for linear search \Rightarrow

```
 $\Rightarrow$  int linearS ( int * arr[], int n, int key) {  
    for (int i=0; i<n; i++) {  
        if (arr[i] == key)  
            return i;  
    }  
    return -1;  
}
```

Ans2 \Rightarrow Pseudo code of Insertion Sort \Rightarrow

```
 $\Rightarrow$  void insertion ( int arr[], int n) {  
    for (int i=1; i<n; i++) {  
        int key = arr[i];  
        int j = i-1;  
        while (j >= 0 && arr[j] > key) {  
            arr[j+1] = arr[j];  
            j--;  
        }  
        arr[j+1] = key;  
    }  
}
```

}

* Insertion sort is called online sorting as, as soon as an element comes in an array \Rightarrow it is automatically inserted at its correct position.

* Other algos. discussed in lectures are not online.

Ans 3 \Rightarrow Avg. case complexities of Sorting Algos \Rightarrow

* Bubble = $O(n^2)$

* Quick $\Rightarrow O(n \log n)$

* Insertion $\Rightarrow O(n^2)$

* Heap $\Rightarrow O(n \log n)$

* Selection $\Rightarrow O(n^2)$

* Merge $\Rightarrow O(n \log n)$

<u>Ans 4</u> \Rightarrow	Stable	Inplace
Bubble \Rightarrow	✓	✓
Selection \Rightarrow	✗	✓
Insertion \Rightarrow	✓	✓
Merge \Rightarrow	✓	✗
Quick \Rightarrow	✗	✗
Heap \Rightarrow	✗	✓

Ans 5 \Rightarrow pseudocode for Binary Search \Rightarrow

\Rightarrow int start = 0

⊙ Time $\Rightarrow O(\log n)$

int end = size - 1

⊙ space $\Rightarrow O(1)$

while (start \leq end) {

int mid = start + (end - start) / 2;

if (key == arr[mid])

return mid;

else if (key < arr[mid])

end = mid - 1;

else

start = mid + 1;

}

return -1;

⊙ Time (of linear search) $\Rightarrow O(n)$

Ans 6 \Rightarrow Recurrence

⊙ space $\Rightarrow O(1)$

relation of Binary

Search $\Rightarrow T(n) = T(n/2) + 1$

Ans 8 \Rightarrow Quick Sort is best sorting algo. in practical

uses as it follows the locality of reference and

also best case time comp. is $O(n \log n)$.

Ans 9 \Rightarrow Number of inversions \Rightarrow It tells us how far is array from being sorted.

\hookrightarrow if $a[i] > a[j]$ and $i < j$.

\Rightarrow 7 21 31 8 10 1 20 6 4 5

\hookrightarrow no. of inversions $\Rightarrow 4 + 7 + 7 + 4 + 4 + 3 + 2$
 $\Rightarrow 31$

Ans 10 \Rightarrow Quick Sort will give \Rightarrow

* Best case complexity \Rightarrow when array is totally unsorted

* worst comp. \Rightarrow when arr. is sorted or reverse sorted.

Ans 11 \Rightarrow Recurrence relations of \Rightarrow

\Rightarrow merge Sort

\Rightarrow quickSort

Best \Rightarrow $2T(n/2) + \theta(n)$

$T(n) = T(k) + T(n-k-1) + \theta(n)$

Worst \Rightarrow

$T(n) = T(n-1) + \theta(n)$

* **Similarity** \Rightarrow Both are types of Divide and Conquer Algorithms.

* Differences \Rightarrow Worst case complexity of Merge sort is $O(n \log n)$ whereas of Quick sort is $O(n^2)$.

Ans 13 \Rightarrow Optimised Bubble Sort \Rightarrow

```
 $\Rightarrow$  for (int i=0; i<n; i++) {  
    swap = false;  
    for (j=0; j<n-i-1; j++) {  
        if (arr[j] > arr[j+1]) {  
            swap(arr[j], arr[j+1]);  
            swapped = true;  
        }  
    }  
}
```

Ans 14 \Rightarrow In such case, Merge sort would be efficient as it is an External sorting algo \Rightarrow data is divided into chunks and then sorted using Merge Sort.

\Rightarrow Sorted data is dumped into files.

① Internal Sorting \Rightarrow It is type of sort in which whole sorting takes place in main memory of computer.