

Question 1:

- (a) Show $f(n) = 5n^3 + 4n^2 + 10$ is $O(n^4)$.

There should be appropriate c and n_0 that satisfy the equation:

$$5n^3 + 4n^2 + 10 \leq cn^4 \text{ for all } n \geq n_0$$

If we pick $c = 19$ and $n_0 = 1$, then this equation is satisfied. Then, by the definition of Big-O, this statement is true.

- (b) Trace the sorting of the array in ascending order: [24, 8, 51, 28, 20, 29, 21, 17, 38, 27]

1- Insertion sort

Divide the array into two parts: sorted and unsorted. Take [24] as the sorted part, the rest is unsorted. Then, loop from $i=1$ to $i=9$.

For $i=1$, the current element is 8. It is smaller than 24, so we insert 8 before 24. The logic is to pick an element from the unsorted array and find the right location in the sorted array to insert this element, by comparing the elements of the sorted array with this element.

Current Array: [8, 24, 51, 28, 20, 29, 21, 17, 38, 27]

For $i=2$, the current element = 51 and it is greater than all elements in the sorted array, so it will remain at the same place.

Current Array: [8, 24, 51, 28, 20, 29, 21, 17, 38, 27]

For $i=3$, the current element = 28 and it is greater 24 and smaller than 51. It will be inserted in between these elements.

Current Array: [8, 24, 28, 51, 20, 29, 21, 17, 38, 27]

For $i=4$, the current element = 20 and it is greater 8 and smaller than 24. It will be inserted in between these elements.

Current Array: [8, 20, 24, 28, 51, 29, 21, 17, 38, 27]

For $i=5$, the current element = 29 and it is greater 28 and smaller than 51. It will be inserted in between these elements.

Current Array: [8, 20, 24, 28, 29, 51, 21, 17, 38, 27]

For $i=6$, the current element = 21 and it is greater 20 and smaller than 24. It will be inserted in between these elements.

Current Array: [8, 20, 21, 24, 28, 29, 51, 17, 38, 27]

For $i=7$, the current element = 17 and it is greater 8 and smaller than 20. It will be inserted in between these elements.

Current Array: [8, 17, 20, 21, 24, 28, 29, 51, 38, 27]

For $i=8$, the current element = 38 and it is greater 29 and smaller than 51. It will be inserted in between these elements.

Current Array: [8, 17, 20, 21, 24, 28, 29, 38, 51, 27]

For $i=9$, the current element = 27 and it is greater 24 and smaller than 28. It will be inserted in between these elements.

Current Array: [8, 17, 20, 21, 24, 27, 28, 29, 38, 51]

Now, the unsorted part of the array has no elements left and we obtained the sorted array in ascending order.

2- Bubble sort

We again start by dividing the array into two parts: sorted and unsorted. Then, we will swap the adjacent elements starting from indexes 0 and 1 by comparing these elements. We will loop from $i=0$ to $i=9$ (unsorted array size-1). We will go through the array from the beginning until there is no need for swapping, in other words, until the array is sorted. There will be (array size-1) number of passes and at the end of each pass, the greatest element in the unsorted array will be put at the last index of the unsorted array.

Current array: [24, 8, 51, 28, 20, 29, 21, 17, 38, 27]

Pass 1:

For $i=0$, the pair is (24,8) and we will swap between them, since $24 > 8$.

Current array: [8, 24, 51, 28, 20, 29, 21, 17, 38, 27]

For $i=1$, the pair is (24,51) and we will not swap them, because they are in the right order.

Current array: [8, 24, 51, 28, 20, 29, 21, 17, 38, 27]

For $i=2$, the pair is (51,28) and we will swap them.

Current array: [8, 24, 28, 51, 20, 29, 21, 17, 38, 27]

Since 51 is greatest element in the whole array, it will be at the last index of this array when $i=9$. The sorted part of the array only includes 51 now. We need to perform 8 more passes to obtain the full sorted array. I will not show the details of each pass, since I showed it in the first one.

Current array: [8, 24, 28, 20, 29, 21, 17, 38, 27, 51]

Pass 2: The greatest element in the unsorted array is 38. There will be switches in the unsorted array, because there are pairs in between that should be switched. (For example, the pair (24,20) is switched, but neither of the elements are the greatest element of the unsorted array.)

Current array: [8, 24, 20, 28, 21, 17, 29, 27, 38, 51]

Pass 3: The greatest element in the unsorted array is 29.

Current array: [8, 20, 24, 21, 17, 28, 27, 29, 38, 51]

Pass 4: The greatest element in the unsorted array is 28.

Current array: [8, 20, 21, 17, 24, 27, 28, 29, 38, 51]

Pass 5: The greatest element in the unsorted array is 27.

Current array: [8, 20, 17, 21, 24, 27, 28, 29, 38, 51]

Pass 6: The greatest element in the unsorted array is 24. At this point, the array is fully sorted. Bubble sort will continue checking the pairs, but will not perform any swap operations. So, the array will remain the same for the remaining passes.

Current array: [8, 17, 20, 21, 24, 27, 28, 29, 38, 51]

Pass 7:

Current array: [8, 17, 20, 21, 24, 27, 28, 29, 38, 51]

Pass 8:

Current array: [8, 17, 20, 21, 24, 27, 28, 29, 38, 51]

Pass 9:

Current array: [8, 17, 20, 21, 24, 27, 28, 29, 38, 51]

The algorithm stops, we have obtained the sorted array.

Question 2:

```
[ayda.yurtoglu@dijkstra cs202hw1]$ ./hw1
Selection sort:
Number of comparisons: 120
Number of data moves: 45
{ 3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21 }

Merge sort:
Number of comparisons: 46
Number of data moves: 128
{ 3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21 }

Quick:
Number of comparisons: 45
Number of data moves: 102
{ 3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21 }

Radix sort:
{ 3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21 }

[ayda.yurtoglu@dijkstra cs202hw1]$
```

RANDOM ARRAYS

Analysis of Selection Sort

Array Size	Elapsed Time	compCount	moveCount
6000	80 ms	17997000	17997
10000	220 ms	49995000	29997
14000	440 ms	97993000	41997
18000	720 ms	161991000	53997
22000	1080 ms	241989000	65997
26000	1510 ms	337987000	77997
30000	2000 ms	449985000	89997

Analysis of Merge Sort

Array Size	Elapsed Time	compCount	moveCount
6000	0 ms	67944	151616
10000	10 ms	120481	267232
14000	0 ms	175324	387232
18000	10 ms	232021	510464
22000	10 ms	290153	638464
26000	10 ms	348883	766464
30000	10 ms	408571	894464

Analysis of Quick Sort

Array Size	Elapsed Time	compCount	moveCount
6000	0 ms	93621	158664
10000	0 ms	156927	262737
14000	10 ms	220809	326196
18000	0 ms	302232	503347
22000	0 ms	379335	653469
26000	0 ms	438508	699248
30000	10 ms	509046	788361

Analysis of Radix Sort

Array Size	Elapsed Time
6000	0 ms
10000	0 ms
14000	0 ms
18000	10 ms
22000	10 ms
26000	10 ms
30000	10 ms

ASCENDING ARRAYS

Analysis of Selection Sort

Array Size	Elapsed Time	compCount	moveCount
6000	140 ms	17997000	17997
10000	500 ms	49995000	29997
14000	1050 ms	97993000	41997
18000	1860 ms	161991000	53997
22000	2840 ms	241989000	65997
26000	3970 ms	337987000	77997
30000	5540 ms	449985000	89997

Analysis of Merge Sort

Array Size	Elapsed Time	compCount	moveCount
6000	0 ms	39984	151616
10000	0 ms	70497	267232
14000	0 ms	101699	387232
18000	0 ms	133904	510464
22000	0 ms	168080	638464
26000	0 ms	200771	766464
30000	10 ms	233274	894464

Analysis of Quick Sort

Array Size	Elapsed Time	compCount	moveCount
6000	80 ms	17997000	23996
10000	220 ms	49995000	39996
14000	410 ms	97993000	55996
18000	690 ms	161991000	71996
22000	1020 ms	241989000	87996
26000	1430 ms	337987000	103996
30000	1900 ms	449985000	119996

Analysis of Radix Sort

Array Size	Elapsed Time
6000	0 ms
10000	0 ms
14000	10 ms
18000	0 ms
22000	10 ms
26000	10 ms
30000	10 ms

DESCENDING ARRAYS

Analysis of Selection Sort

Array Size	Elapsed Time	compCount	moveCount
6000	90 ms	17997000	17997
10000	300 ms	49995000	29997
14000	650 ms	97993000	41997
18000	1130 ms	161991000	53997
22000	1750 ms	241989000	65997
26000	2560 ms	337987000	77997
30000	3420 ms	449985000	89997

Analysis of Merge Sort

Array Size	Elapsed Time	compCount	moveCount
6000	0 ms	36656	151616
10000	0 ms	64608	267232
14000	0 ms	94256	387232
18000	0 ms	124640	510464
22000	10 ms	154208	638464
26000	0 ms	186160	766464
30000	10 ms	219504	894464

Analysis of Quick Sort

Array Size	Elapsed Time	compCount	moveCount
6000	120 ms	14088243	21153226
10000	340 ms	39054266	58616386
14000	670 ms	76522927	114833336
18000	1090 ms	124669785	187066919
22000	1630 ms	186425116	279713558
26000	2300 ms	263110172	394755781
30000	3060 ms	350085794	525233545

Analysis of Radix Sort

Array Size	Elapsed Time
6000	10 ms
10000	10 ms
14000	0 ms
18000	10 ms
22000	0 ms
26000	10 ms
30000	10 ms

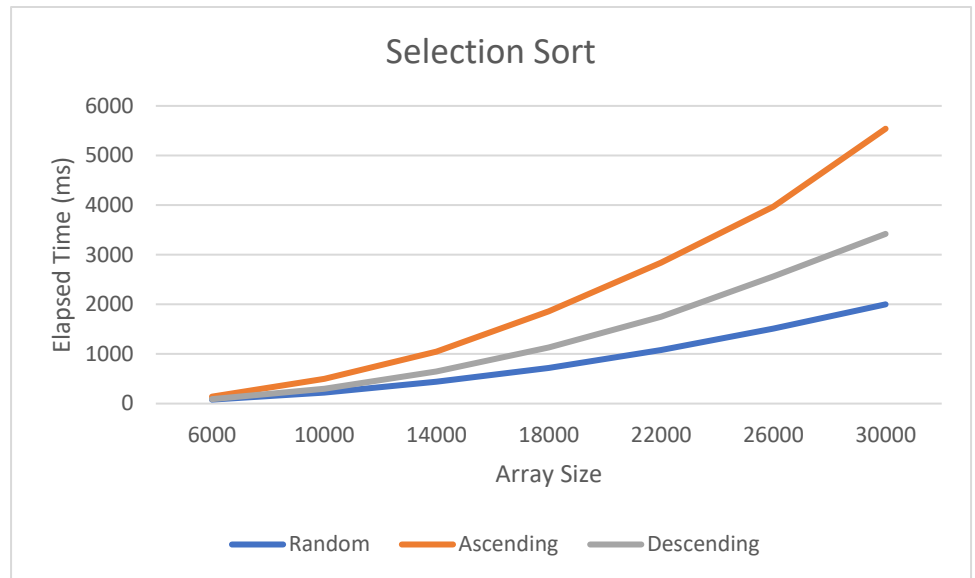
Question 3:

Selection Sort:

Time complexity: $O(n^2)$

As the array size increased, the time to sort also increased with time complexity $O(n^2)$ for all cases. However, the elapsed time was the greatest in the ascending scenario, when the array was already sorted. This was different from what I expected, because the best, worst and average cases

should be the same in theory. My empirical results suggest that when the array is random the time it took to sort is the smallest. I think this is just a coincidence and it doesn't matter whether the array is actually sorted or not. Also, the time we measure is very small and although the graph suggests that there is a huge difference, the results are coincidental and might be the same when I run the program again.

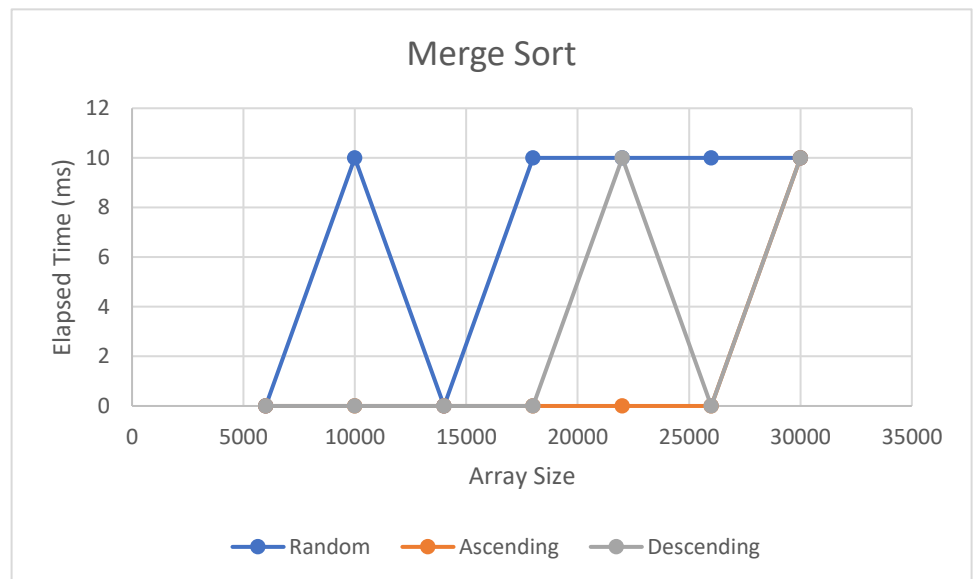


Merge Sort:

Time complexity is $O(n \cdot \log n)$ for all cases of merge sort.

Because of Dijkstra's rounding, the results were interesting in merge sort. If it didn't round up, I would have gathered more precise information about elapsed time. In merge sort, theoretically the worst case and the average case should be the same in terms of time complexity. This is what I observed in my results,

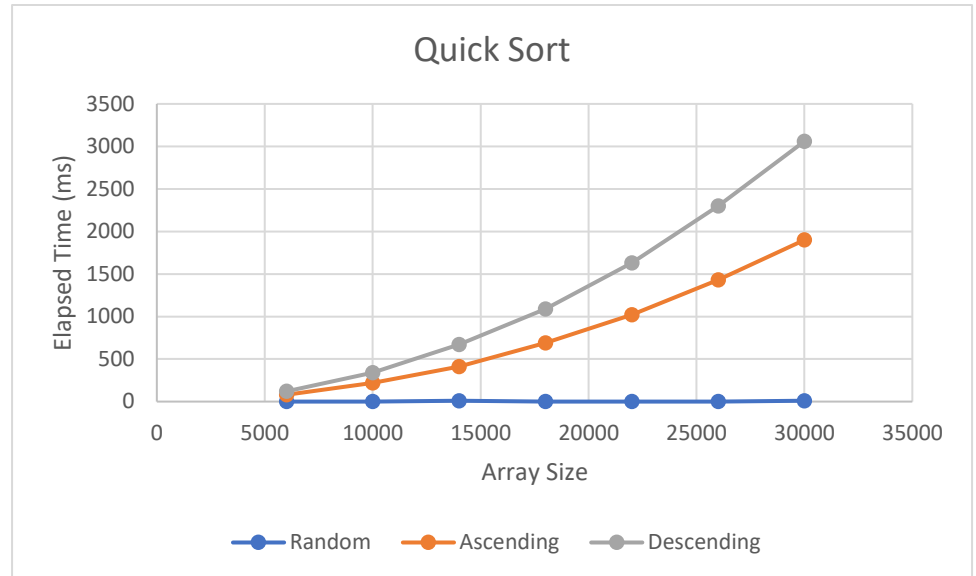
because the elapsed time were almost the same for all three cases. So, it didn't matter whether the array was sorted at first or not. Since the elapsed time didn't change drastically as it did in selection sort, we can say that merge sort is a very efficient sorting function.



Quick Sort:

For quick sort's average case, the time complexity is $O(n \cdot \log n)$. For the worst case, it is $O(n^2)$. When we select the first element as pivot, worst case happens when the array is already sorted. This is because the function will keep dividing the array into two sub arrays of sizes 0 and (size-1) instead of size/2 and size/2 (like in the average case). My results suggest that when the array was sorted in descending

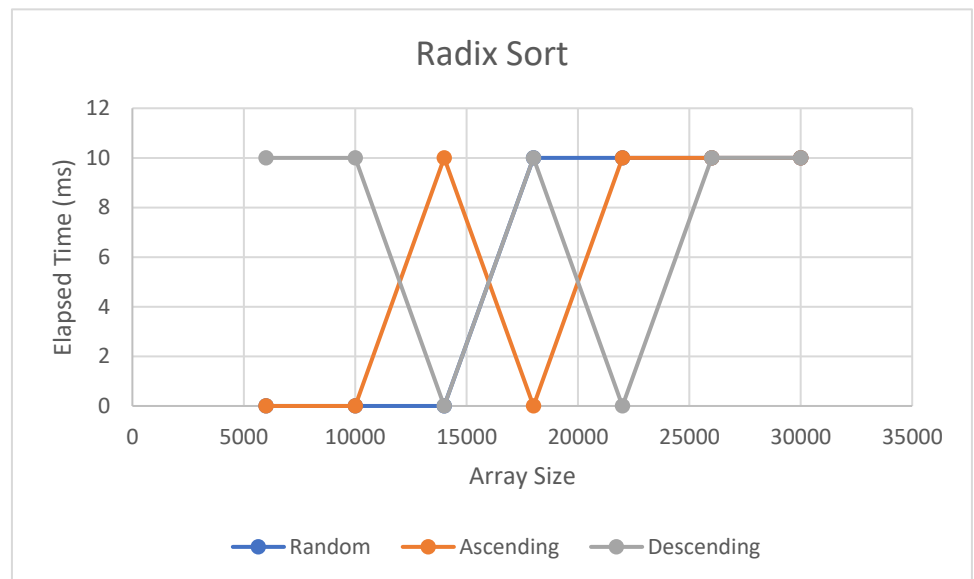
order, the time the function required was the greatest. When the array was random, it took very small amount of time to sort. We can also observe that for both ascending and descending arrays, the time complexity is $O(n^2)$ from the graph. This coincides with the theory for quick sort. So, the average case would be when the array is random.



Radix Sort:

For all cases of radix sort, the time complexity is $O(n)$. In theory, it shouldn't matter whether the arrays are sorted or not and my results suggest this. Since radix sort is very efficient, Dijkstra rounded up the elapsed time just like it did with merge sort, so results are not as precise as the selection sort. They are either 0 or 10. Still, we can understand that the scenarios don't matter

when it comes to time complexity. The duration is random and is not dependent on the scenario. If we used larger array sizes or Dijkstra didn't round up, we would also see that as the array size increased the time also increased with $O(n)$ complexity.



Comparing all functions:

I used MS Excel to compare the results of all sorting functions and their required time to sort random arrays. Since the time for selection sort was too large than the others, we can't see the red and the grey lines, because they are behind the yellow line. So, I created another graph that only compares merge sort, quick sort and radix sort.

Starting with the first graph, we can say that selection sort is the least efficient one. This is also the case when we compare the time complexities of these functions.

$O(n^2) > O(n \cdot \log n) > O(n)$

If we compare the other functions, we see that Dijkstra again rounded the results because they were so small. So, all of these functions are very efficient and the results we see here

are random. Still, in theory radix sort should be more efficient than merge sort and quick sort due to its time complexity. We might have seen this result here if we have used greater array sizes to see the difference. We are able to see the difference of selection sort and the others because $O(n^2)$ grows very rapidly. This is not the case for other functions.

