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注意

优先队列是大的在前面 如果要小的 要重载 long long 二分答案的时候...精度 也有可能 爆 int (?

哈希 自然溢出 yyds 双哈希 输出限制..

匈牙利的复杂度常数非常小(...

```
递归爆栈 re
for i 进行计算的时候 (i 开 long long)
边界问题各种01的特判
模 多模一点 都可以模 (
char 数组开小了也可能报错 tle 和 wa (
图是否连通 是否重边 是否自环
读题!!与或(
重点 重边
当保证 n 的总和不会很大,但数据组数可能很多的时候,注意初始化造成的 tle 问
题(
pow()的精度问题
unique erase 先排序
图论初始化!!
没开 longlong,中间有个判定条件爆了
他卡快排,由于答案不超过10000,可以计数排序
re: 没有开 longlong, (以为是 dfs 爆栈
定义
() (a,b)=1 最大公约数即 a, b 互质
| 整除 a|b b%a==0
STL
优先队列重载
priority_queue<int, vector<int>, cmp>s;
struct cmp{
    bool operator()(const int &a,const int &b){
         return a>b;
    }
};
```

```
set 重载
#include <bits/stdc++.h>
using namespace std;
#define 1 first
#define r second
struct cmp{
      bool operator() (const pair <int, int> &a, const pair<int, int> &
b) const{
            int lena = a.r - a.l + 1;
            int lenb = b.r - b.l + 1;
            if(lena == lenb) return a.l < b.l;</pre>
            return lena > lenb;
      }
};
int main(){
      ios :: sync_with_stdio(0); cin.tie(0); cout.tie(0);
      int T:
      cin >> T;
      while(T -- ){
            int n;
            cin >> n;
            set<pair<int, int>, cmp> segs;
            segs.insert({0, n - 1});
            vector<int> a(n);
            for(int i = 1; i <= n; ++ i){</pre>
                  pair<int, int> cur = *segs.begin();
                  segs.erase(segs.begin());
                  int id = (cur.1 + cur.r) / 2;
                  a[id] = i;
                  if(cur.l < id) segs.insert({cur.l, id - 1});</pre>
                  if(id < cur.r) segs.insert({id + 1, cur.r});</pre>
            }
            for(auto it : a) cout << it << " ";</pre>
            cout << endl;</pre>
      }
}
动态开数组
int a[15], n, m;
cin >> n >> m;
int (*b)[m] = (int (*)[m])a;
new / delete
#define M 10U
#define N 20
第一种,可以直接[][]访问。但是内存不连续,不是很推荐使用,除非 M \ N 都不确定
```

```
//定义的时候
int** pNum;//以int 为例
pNum = new int*[M];
for(int i = 0;i < M;i ++){</pre>
     pNum[i]=new int[N];
}
//删除的时候是
for(int j = 0; j < M; j ++){}
     delete []pNum[i];
delete []pNum;
malloc / free
#include<stdio.h>
#include<stdlib.h>
int main() {
   int **a; //用二级指针动态申请二维数组
   int i,j;
   int m,n;
   printf("请输入行数\n");
   scanf("%d",&m);
   printf("请输入列数\n");
   scanf("%d",&n);
   a=(int**)malloc(sizeof(int*)*m);
   for(i=0;i<m;i++)</pre>
   a[i]=(int*)malloc(sizeof(int)*n);
   for(i=0;i<m;i++) {</pre>
       for(j=0;j<n;j++) {
           printf("%p\n",&a[i][j]);
                                       //输出每个元素地址,每行的列与列
之间的地址时连续的, 行与行之间的地址不连续
       }
   for(i=0;i<m;i++)</pre>
   free(a[i]);
   free(a);
   return 0;
}
#include<stdio.h>
#include<stdlib.h>
int main()
{
```

```
int i,j;
   //申请一个3 行2 列的整型数组
   int (*a)[2]=(int(*)[2])malloc(sizeof(int)*3*2);
   for(i=0;i<3;i++) {</pre>
       for(j=0;j<2;j++) {</pre>
          printf("%p\n",&a[i][j]); //输出数组每个元素地址,每个元素的地
址是连续的
       }
   }
   free(a);
   return 0;
}
vector
//二维 vector
vector<vector <int> > ivec(m ,vector<int>(n)); //m*n 的二维vector
//动态创建 m*n 的二维 vector
//方法一:
vector<vector <int> > ivec;
ivec.resize(m);
for(int i=0;i<m;i++)</pre>
 ivec[i].resize(n);
//方法二:
vector<vector <int> > ivec;
ivec.resize(m, vector<int>(n));
set
begin(),返回 set 容器的第一个元素
end(),返回 set 容器的最后一个元素
clear(),删除 set 容器中的所有的元素
empty(),判断 set 容器是否为空
max_size(),返回 set 容器可能包含的元素最大个数
size(),返回当前 set 容器中的元素个数
rbegin,返回的值和 end()相同
rend(),返回的值和 rbegin()相同
count() 用来查找 set 中某个某个键值出现的次数。
```

equal_range(),返回一对定位器,分别表示第一个大于或等于给定关键值的元素和第一个大于给定关键值的元素,这个返回值是一个 pair 类型,如果这一对定位器中哪个返回失败,就会等于 end()的值。

erase(iterator),删除定位器 iterator 指向的值

erase(first,second),删除定位器 first 和 second 之间的值

erase(key_value),删除键值 key_value 的值

find(),返回给定值值得定位器,如果没找到则返回 end()。

insert(key_value); 将 keyvalue 插入到 set 中,返回值是 pair<set::iterator,bool>, bool 标志着插入是否成功,而 iterator 代表插入的位置,若 keyvalue 已经在 set 中,则 iterator 表示的 key_value 在 set 中的位置。

inset(first,second);将定位器 first 到 second 之间的元素插入到 set 中,返回值是void.

lowerbound(keyvalue),返回第一个大于等于 key_value 的定位器 upperbound(keyvalue),返回最后一个大于等于 key_value 的定位器

map

插入操作

使用[]进行单个插入

```
map<int, string> ID_Name;
// 如果已经存在键值 2015,则会作赋值修改操作,如果没有则插入
ID Name[2015] = "Tom";1234
```

使用 insert 进行单个和多个插入 (insert 共有 4 个重载函数:

```
// 插入单个键值对,并返回插入位置和成功标志,插入位置已经存在值时,插入失败 pair<iterator,bool> insert (const value_type& val);
//在指定位置插入,在不同位置插入效率是不一样的,因为涉及到重排
iterator insert (const_iterator position, const value_type& val);
// 插入多个
void insert (InputIterator first, InputIterator last);
//c++11 开始支持,使用列表插入多个
void insert (initializer_list<value_type> il);
```

取值

Map 中元素取值主要有 at 和[]两种操作, at 会作下标检查, 而[]不会。

```
map<int, string> ID Name;
//ID Name 中没有关键字2016,使用[]取值会导致插入
//因此,下面语句不会报错,但打印结果为空
cout<<ID_Name[2016].c_str()<<endl;</pre>
//使用at 会进行关键字检查,因此下面语句会报错
ID_Name.at(2016) = "Bob";
容量查询
// 查询 map 是否为空
bool empty();
// 查询map 中键值对的数量
size t size();
// 查询 map 所能包含的最大键值对数量,和系统和应用库有关。
// 此外,这并不意味着用户一定可以存这么多,很可能还没达到就已经开辟内存失败了
size t max size();
// 查询关键字为 key 的元素的个数,在 map 里结果非 0 即 1
size_t count( const Key& key ) const; //
迭代器
共有八个获取迭代器的函数: begin, end, rbegin, rend 以及对应的 cbegin, cend,
crbegin, crend o
二者的区别在于,后者一定返回 constiterator, 而前者则根据 map 的类型返回
iterator 或者 constiterator。const 情况下,不允许对值进行修改。如下面代码所
示:
map<int,int>::iterator it;
map<int,int> mmap;
const map<int,int> const_mmap;
it = mmap.begin(); //iterator
mmap.cbegin(); //const iterator
const mmap.begin(); //const iterator
const mmap.cbegin(); //const iterator123456789
返回的迭代器可以进行加减操作,此外,如果 map 为空,则 begin = end。
删除
// 删除迭代器指向位置的键值对,并返回一个指向下一元素的迭代器
iterator erase( iterator pos )
// 删除一定范围内的元素,并返回一个指向下一元素的迭代器
iterator erase( const_iterator first, const_iterator last );
// 根据Key 来进行删除, 返回删除的元素数量,在map 里结果非 0 即 1
size_t erase( const key_type& key );
// 清空map,清空后的size 为0
void clear();
```

交换

```
// 就是两个map 的内容互换
void swap( map& other );
顺序比较
// 比较两个关键字在 map 中位置的先后
key_compare key_comp() const;
查找
// 关键字查询,找到则返回指向该关键字的迭代器,否则返回指向 end 的迭代器
// 根据map 的类型,返回的迭代器为 iterator 或者 const_iterator
iterator find (const key type& k);
const_iterator find (const key_type& k) const;
操作符
operator: == != < <= > >=
注意 对于==运算符,只有键值对以及顺序完全相等才算成立。
unordered_map
查找元素是否存在
若有 unordered_map <int, int> mp;查找 x 是否在 map 中
方法 1: 若存在 mp.find(x)!=mp.end()
方法 2: 若存在 mp.count(x)!=0123
插入数据
mp.insert(Map::value_type(1, "Raoul"));1
遍历 map
 unordered map<key,T>::iterator it;
   (*it).first; //the key value
   (*it).second
               //the mapped value
   for(unordered_map<key,T>::iterator iter=mp.begin();iter!=mp.end();i
ter++)
         cout<<"key value is"<<iter->first<<" the mapped value is "<<</pre>
iter->second;
   //也可以这样
   for(auto& v : mp)
       print v.first and v.second
```

bitset

C++的 bitset 在 bitset 头文件中,它是一种类似数组的结构,它的每一个元素只能是 0 或 1 ,每个元素仅用 1 bit 空间。

bitset 数组与 vector 数组区别

bitset 声明数组:bitset<100> number[10] vector 声明数组:vector number[10];

bitset<每个 bitset 元素的长度(没有占满前面全部自动补 0)> 元素 bitset 内置转化函数: 可将 bitset 转化为 string,unsigned long,unsigned long long。

构造

```
bitset<4> bitset1; //无参构造,长度为4,默认每一位为0
   bitset<8> bitset2(12);
                        //长度为8,二进制保存,前面用0补充
   string s = "100101";
   bitset<10> bitset3(s); //长度为10,前面用0补充
   char s2[] = "10101";
   bitset<13> bitset4(s2); //长度为13,前面用 0 补充
   cout << bitset1 << endl;</pre>
                          //0000
   cout << bitset2 << endl;</pre>
                          //00001100
   cout << bitset3 << endl;</pre>
                          //0000100101
   cout << bitset4 << endl;</pre>
                          //0000000010101
函数
     bitset<8> foo ("10011011");
   cout << foo.count() << endl; //5 (count 函数用来求bitset 中1的
位数,foo 中共有5个1
   cout << foo.size() << endl; //8 (size 函数用来求bitset 的大
小,一共有8位
   cout << foo.test(0) << endl; //true</pre>
                                       (test 函数用来查下标处的元
素是 0 还是 1 ,并返回 false 或 true,此处 foo[0]为 1 ,返回 true
   cout << foo.test(2) << endl; //false (同理, foo[2]为0,返回fa
Lse
   cout << foo.any() << endl; //true (any 函数检查bitset 中是否有
   cout << foo.none() << endl; //false (none 函数检查bitset 中是否
没有1
   cout << foo.all() << endl; //false (all 函数检查bitset 中是全部
为1
```

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```
#include <bits/stdc++.h>
#define 11 long long
using namespace std;
int t,n,m;
char str[1010];
bitset<500> number[30];
int main() {
      ios::sync_with_stdio(false); cin.tie(0); cout.tie(0);
    //freopen("test.in","r",stdin);
    //freopen("test.out", "w", stdout);
      scanf("%d",&t);
      while(t--)
      {
            scanf("%d %d",&n,&m);
            for(int i=0;i<m;i++)</pre>
                  scanf("%s",str);
                  number[i]=bitset<500>(str);
            int len=1<<m,ans=m+1;</pre>
            for(int i=1;i<len;i++)</pre>
                  int t=i,s=0;
                  bitset<500> num(0);
                  for(int j=0;j<m&&t>0;j++)
                         if(t&1)
                         {
                               num=num|number[j];
                               S++;
                         t>>=1;
                   if(num.count()==n) ans=min(ans,s);
            if(ans==m+1) printf("-1\n");
            else printf("%d\n",ans);
      return 0;
}
计算几何
zyx 的计算几何
const double eps = 1e-9;
const double PI = acos(-1.0);
struct Line;
```

Н

```
struct Point {
    double x, y;
    Point() { x = y = 0; }
    Point(const Line &a);
    Point(const double &a, const double &b) : x(a), y(b) {}
    Point operator+(const Point &a) const {
        return \{x + a.x, y + a.y\};
    Point operator-(const Point &a) const {
        return {x - a.x, y - a.y};
    Point operator*(const double &a) const {
        return {x * a, y * a};
    Point operator/(const double &d) const {
        return {x / d, y / d};
    bool operator==(const Point &a) const {
        return abs(x - a.x) + abs(y - a.y) < eps;
    }
    void standardize() {
        *this = *this / sqrt(x * x + y * y);
    }
};
Point normal(const Point &a) { return Point(-a.y, a.x); }
double dist(const Point &a, const Point &b) {
    return sqrt((a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y));
}
double dist2(const Point &a, const Point &b) {
    return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y);
}
struct Line {
    Point s, t;
    Line() {}
    Line(const Point &a, const Point &b) : s(a), t(b) {}
};
struct circle {
    Point o;
    double r;
    circle() {}
    circle(Point P, double R = 0) { o = P, r = R; }
};
double length(const Point &p) {
    return sqrt(p.x * p.x + p.y * p.y);
double length(const Line &1) {
    Point p(1);
    return length(p);
Point::Point(const Line &a) { *this = a.t - a.s; }
```

```
istream &operator>>(istream &in, Point &a) {
    in >> a.x >> a.y;
    return in;
}
double dot(const Point &a, const Point &b) {
    return a.x * b.x + a.y * b.y;
}
double det(const Point &a, const Point &b) {
    return a.x * b.y - a.y * b.x;
int sgn(const\ double\ &x) \{ return\ fabs(x) < eps ? 0 : (x > 0 ? 1 : -1);
double sqr(const double &x) { return x * x; }
Point rotate(const Point &a, const double &ang) {
    double x = cos(ang) * a.x - sin(ang) * a.y;
    double y = sin(ang) * a.x + cos(ang) * a.y;
    return {x, y};
}
//点在线段上 <=0 包含端点
bool sp_on(const Line &seg, const Point &p) {
    Point a = seg.s, b = seg.t;
    return !sgn(det(p - a, b - a)) && sgn(dot(p - a, p - b)) <= 0;
bool lp on(const Line &line, const Point &p) {
    Point a = line.s, b = line.t;
    return !sgn(det(p - a, b - a));
//等于不包含共线
int andrew(Point *point, Point *convex, int n) {
    sort(point, point + n, [](Point a, Point b) {
        if (a.x != b.x) return a.x < b.x;
        return a.y < b.y;</pre>
    });
    int top = 0;
    for (int i = 0; i < n; i++) {</pre>
       while ((top > 1) && det(convex[top - 1] - convex[top - 2], poin
t[i] - convex[top - 1]) <= 0
            top--;
       convex[top++] = point[i];
    }
    int tmp = top;
    for (int i = n - 2; i >= 0; i--) {
       while ((top > tmp) && det(convex[top - 1] - convex[top - 2], po
int[i] - convex[top - 1]) <= 0
            top--;
       convex[top++] = point[i];
    }
    if (n > 1) top--;
    return top;
}
```

```
double slope(const Point &a, const Point &b) {
    return (a.y - b.y) / (a.x - b.x);
double slope(const Line &a) {
    return slope(a.s, a.t);
Point 11 intersection(const Line &a, const Line &b) {
    double s1 = det(Point(a), b.s - a.s), s2 = det(Point(a), b.t - a.s);
    return (b.s * s2 - b.t * s1) / (s2 - s1);
int ss_cross(const Line &a, const Line &b, Point &p) {
    int d1 = sgn(det(a.t - a.s, b.s - a.s));
    int d2 = sgn(det(a.t - a.s, b.t - a.s));
    int d3 = sgn(det(b.t - b.s, a.s - b.s));
    int d4 = sgn(det(b.t - b.s, a.t - b.s));
    if ((d1 ^ d2) == -2 \&\& (d3 ^ d4) == -2) {
        p = ll_intersection(a, b);
        return 1;
    if (!d1 && sp_on(a, b.s)) {
        p = b.s;
        return 2;
    if (!d2 && sp on(a, b.t)) {
       p = b.t;
        return 2;
    if (!d3 && sp_on(b, a.s)) {
        p = a.s;
        return 2;
    if (!d4 && sp_on(b, a.t)) {
        p = a.t;
        return 2;
    return 0;
Point project(const Line &l, const Point &p) {
    Point base(1);
    double r = dot(base, p - 1.s) / sqr(length(base));
    return 1.s + (base * r);
double sp dist(const Line &l, const Point &p) {
    if (l.s == l.t) return dist(l.s, p);
    Point x = p - 1.s, y = p - 1.t, z = 1.t - 1.s;
    if (sgn(dot(x, z)) < 0)return length(x);//P 距离A 更近
    if (sgn(dot(y, z)) > 0)return length(y);//P 距离 B 更近
    return abs(det(x, z) / length(z));//面积除以底边长
}
```

```
double lp dist(const Line &1, const Point &p) {
    Point x = p - 1.s, y = p - 1.t, z = 1.t - 1.s;
    return abs(det(x, z) / length(z));//面积除以底边长
}
int lc cross(const Line &1, const Point &a, const double &r, pair<Point,</pre>
 Point> &ans) {
    int num = 0;
    Point pr = project(1, a);
    double dis = dist(pr, a);
    double tmp = r * r - dis * dis;
    if (sgn(tmp) == 1) num = 2;
    else if (sgn(tmp) == 0) num = 1;
    else return 0;
    double base = sqrt(r * r - dis * dis);
    Point e(1);
    e.standardize();
    e = e * base;
    ans = make pair(pr + e, pr - e);
    return num;
}
int cc cross(const Point &c1, const double &r1, const Point &c2, const
double &r2, pair<Point, Point> &ans) {
    double x1 = c1.x, x2 = c2.x, y1 = c1.y, y2 = c2.y;
    double d = length(c1 - c2);
    if (sgn(fabs(r1 - r2) - d) > 0) return -1; //内含
    if (sgn(r1 + r2 - d) < 0) return 0; //相离
    double a = r1 * (x1 - x2) * 2, b = r1 * (y1 - y2) * 2, c = r2 * r2
- r1 * r1 - d * d:
    double p = a * a + b * b, q = -a * c * 2, r = c * c - b * b;
    double cosa, sina, cosb, sinb;
    //One Intersection
    if (sgn(d - (r1 + r2)) == 0 | | sgn(d - fabs(r1 - r2)) == 0) {
        cosa = -q / p / 2;
        sina = sqrt(1 - sqr(cosa));
        Point p0(x1 + r1 * cosa, y1 + r1 * sina);
        if (sgn(dist(p0, c2) - r2)) p0.y = y1 - r1 * sina;
        ans = pair<Point, Point>(p0, p0);
        return 1;
    }
    //Two Intersections
    double delta = sqrt(q * q - p * r * 4);
    cosa = (delta - q) / p / 2;
    cosb = (-delta - q) / p / 2;
    sina = sqrt(1 - sqr(cosa));
    sinb = sqrt(1 - sqr(cosb));
    Point p1(x1 + r1 * cosa, y1 + r1 * sina);
    Point p2(x1 + r1 * cosb, y1 + r1 * sinb);
    if (sgn(dist(p1, c2) - r2)) p1.y = y1 - r1 * sina;
    if (sgn(dist(p2, c2) - r2)) p2.y = y1 - r1 * sinb;
```

```
if (p1 == p2) p1.y = y1 - r1 * sina;
   ans = pair<Point, Point>(p1, p2);
   return 2;
Point lp_sym(const Line &1, const Point &p) {
   return p + (project(1, p) - p) * 2;
}
double alpha(const Point &t1, const Point &t2) {
   double theta;
   theta = atan2((double) t2.y, (double) t2.x) - atan2((double) t1.y,
(double) t1.x);
   if (sgn(theta) < 0)</pre>
       theta += 2.0 * PI;
   return theta;
int pip(const Point *P, const int &n, const Point &a) {//【射线法】判断
点A 是否在任意多边形 Poly 以内
   int cnt = 0;
   int tmp;
   for (int i = 1; i <= n; ++i) {</pre>
       int j = i < n ? i + 1 : 1;
       if (sp on(Line(P[i], P[j]), a))return 2;//点在多边形上
       if (a.y >= min(P[i].y, P[j].y) && a.y < max(P[i].y, P[j].y))//</pre>
纵坐标在该线段两端点之间
           tmp = P[i].x + (a.y - P[i].y) / (P[j].y - P[i].y) * (P[j].x
 - P[i].x), cnt += sgn(tmp - a.x) > 0;//交点在A 右方
   return cnt & 1;//穿过奇数次则在多边形以内
bool pip_convex_jud(const Point &a, const Point &L, const Point &R) {//
判断AL 是否在AR 右边
   return sgn(det(L - a, R - a)) > 0;//必须严格以内
bool pip_convex(const Point *P, const int &n, const Point &a) {// 【二分
法】判断点A 是否在凸多边形Polv 以内
   //点按逆时针给出
   if (pip_convex_jud(P[0], a, P[1]) || pip_convex_jud(P[0], P[n - 1],
 a)) return 0;//在P[0_1]或P[0_n-1]外
   if (sp_on(Line(P[0], P[1]), a) \mid | sp_on(Line(P[0], P[n - 1]), a)) r
eturn 2;//在P[0 1]或P[0 n-1]上
   int l = 1, r = n - 2;
   while (1 < r) {//二分找到一个位置 pos 使得 P[0]_A 在 P[0_pos], P[0_(pos+
1) ]之间
       int mid = (1 + r + 1) >> 1;
       if (pip_convex_jud(P[0], P[mid], a))l = mid;
       else r = mid - 1;
   if (pip convex jud(P[1], a, P[1 + 1]))return 0;//在P[pos (pos+1)]外
```

```
if (sp_on(Line(P[1], P[1 + 1]), a))return 2;//在P[pos_(pos+1)]上
   return 1;
}
// 多边形是否包含线段
// 因此我们可以先求出所有和线段相交的多边形的顶点,然后按照X-Y 坐标排序(X 坐标
小的排在前面,对于X 坐标相同的点,Y 坐标小的排在前面,
// 这种排序准则也是为了保证水平和垂直情况的判断正确),这样相邻的两个点就是在线
段上相邻的两交点,如果任意相邻两点的中点也在多边形内,
// 则该线段一定在多边形内。
int pp judge(Point *A, int n, Point *B, int m) {// 【判断多边形 A 与多边形
B 是否相离】
   for (int i1 = 1; i1 <= n; ++i1) {</pre>
       int j1 = i1 < n ? i1 + 1 : 1;
       for (int i2 = 1; i2 <= m; ++i2) {</pre>
           int j2 = i2 < m ? i2 + 1 : 1;
           Point tmp;
           if (ss_cross(Line(A[i1], A[j1]), Line(B[i2], B[j2]), tmp))
return 0;//两线段相交
           if (pip(B, m, A[i1]) | | pip(A, n, B[i2]))return 0;//点包含在
内
       }
   }
   return 1;
double area(Point *P, int n) {//【任意多边形 P 的面积】
   double S = 0;
   for (int i = 1; i <= n; i++) S += det(P[i], P[i < n ? i + 1 : 1]);</pre>
   return S / 2.0;
Line O[N];
int judge(Line L, Point a) { return sgn(det(a - L.s, L.t - L.s)) > 0; }
//判断点 a 是否在直线 L 的右边
int halfcut(Line *L, int n, Point *P) {//【半平面交】
   sort(L, L + n, [](const Line &a, const Line &b) {
       double d = atan2((a.t - a.s).y, (a.t - a.s).x) - atan2((b.t - b.
s).y, (b.t - b.s).x);
       return sgn(d) ? sgn(d) < 0 : judge(a, b.s);</pre>
   });
   int m = n;
   n = 0:
   for (int i = 0; i < m; ++i)
       if (i == 0 || sgn(atan2(Point(L[i]).y, Point(L[i]).x) - atan2(P
oint(L[i-1]).y, Point(L[i-1]).x)))
           L[n++] = L[i];
   int h = 1, t = 0;
   for (int i = 0; i < n; ++i) {</pre>
       while (h < t && judge(L[i], ll_intersection(Q[t], Q[t - 1]))) -</pre>
-t;//当队尾两个直线交点不是在直线L[i]上或者左边时就出队
```

```
while (h < t \&\& judge(L[i], ll intersection(Q[h], Q[h + 1]))) +
+h;//当队头两个直线交点不是在直线L[i]上或者左边时就出队
       Q[++t] = L[i];
   while (h < t \&\& judge(Q[h], ll_intersection(Q[t], Q[t - 1]))) --t;
   while (h < t && judge(O[t], ll intersection(O[h], O[h + 1]))) ++h;
   n = 0;
   for (int i = h; i <= t; ++i) {</pre>
       P[n++] = ll\_intersection(Q[i], Q[i < t ? i + 1 : h]);
    }
   return n;
}
Point V1[N], V2[N];
int mincowski(Point *P1, int n, Point *P2, int m, Point *V) {//【闵可夫
斯基和】求两个凸包{P1},{P2}的向量集合{V}={P1+P2}构成的凸包
   for (int i = 0; i < n; ++i) V1[i] = P1[(i + 1) \% n] - P1[i];
   for (int i = 0; i < m; ++i) V2[i] = P2[(i + 1) \% m] - P2[i];
   int t = 0, i = 0, j = 0;
   V[t++] = P1[0] + P2[0];
   while (i < n \& j < m) V[t] = V[t - 1] + (sgn(det(V1[i], V2[j])) >
0 ? V1[i++] : V2[j++]), t++;
   while (i < n) V[t] = V[t - 1] + V1[i++], t++;
   while (j < m) V[t] = V[t - 1] + V2[j++], t++;
   return t;
circle getcircle(const Point &A, const Point &B, const Point &C) {//
【三点确定一圆】向量垂心法
   Point P1 = (A + B) * 0.5, P2 = (A + C) * 0.5;
   Line R1 = Line(P1, P1 + normal(B - A));
   Line R2 = Line(P2, P2 + normal(C - A));
   circle 0;
   0.o = ll_intersection(R1, R2);
   0.r = length(A - 0.0);
   return 0;
struct ConvexHull {
   int op;
   struct cmp {
       bool operator()(const Point &a, const Point &b) const {
            return sgn(a.x - b.x) < 0 \mid \mid sgn(a.x - b.x) == 0 && sgn(a.y)
 - b.y) < 0;
       }
   };
   set<Point, cmp> s;
   ConvexHull(int o) {
       op = o;
       s.clear();
    inline int PIP(Point P) {
```

```
set<Point>::iterator it = s.lower bound(Point(P.x, -dinf));//按
到第一个横坐标大于P 的点
       if (it == s.end())return 0;
        if (sgn(it\rightarrow x - P.x) == 0) return sgn((P.y - it\rightarrow y) * op) <= 0;
//比较纵坐标大小
        if (it == s.begin())return 0;
        set<Point>::iterator j = it, k = it;
        --j;
        return sgn(det(P - *j, *k - *j) * op) >= 0;//看叉姬1
   inline int judge(set<Point>::iterator it) {
        set<Point>::iterator j = it, k = it;
        if (j == s.begin())return 0;
        --j;
        if (++k == s.end())return 0;
        return sgn(det(*it - *j, *k - *j) * op) >= 0;//看叉姬
   inline void insert(Point P) {
        if (PIP(P))return;//如果点P已经在凸壳上或凸包里就不插入了
        set<Point>::iterator tmp = s.lower_bound(Point(P.x, -inf));
        if (tmp != s.end() && sgn(tmp->x - P.x) == \emptyset)s.erase(tmp);//特
判横坐标相等的点要去掉
       s.insert(P);
        set<Point>::iterator it = s.find(P), p = it;
        if (p != s.begin()) {
            --p;
           while (judge(p)) {
                set<Point>::iterator temp = p--;
                s.erase(temp);
            }
        if ((p = ++it) != s.end()) {
           while (judge(p)) {
                set<Point>::iterator temp = p++;
                s.erase(temp);
            }
        }
} up(1), down(-1);
int PIC(circle C, Point a) { return sgn(length(a - C.o) - C.r) <= 0; }/</pre>
/判断点A是否在圆C内
void Random(Point *P, int n) { for (int i = 0; i < n; ++i)swap(P[i], P</pre>
[(rand() + 1) % n]); }//随机一个排列
circle min circle(Point *P, int n) {// 【求点集 P 的最小覆盖圆】 O(n)
// random shuffle(P,P+n);
   Random(P, n);
   circle C = circle(P[0], 0);
   for (int i = 1; i < n; ++i)
        if (!PIC(C, P[i])) {
```

计算几何全家桶 #include <bits/stdc++.h> using namespace std; typedef long long 11; const $11 N = 1 \ll 20$; const $11 \mod = 1e9 + 7$; const double dinf = 1e99; const int inf = 0x3f3f3f3f3f; const 11 linf = 0x3f3f3f3f3f3f3f3f3f; const double eps = 1e-9; const double PI = acos(-1.0); struct Line; struct Point { double x, y; Point() { x = y = 0; } Point(const Line &a); Point(const double &a, const double &b) : x(a), y(b) {} Point operator+(const Point &a) const { **return** {x + a.x, y + a.y}; Point operator-(const Point &a) const { **return** {x - a.x, y - a.y}; Point operator*(const double &a) const { **return** {x * a, y * a}; Point operator/(const double &d) const { **return** {x / d, y / d}; bool operator==(const Point &a) const { return abs(x - a.x) + abs(y - a.y) < eps;void standardize() { *this = *this / sqrt(x * x + y * y); } **}**; Point normal(const Point &a) { return Point(-a.y, a.x); }

```
double dist(const Point &a, const Point &b) {
    return sqrt((a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y));
double dist2(const Point &a, const Point &b) {
    return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y);
struct Line {
    Point s, t;
    Line() {}
    Line(const Point &a, const Point &b) : s(a), t(b) {}
};
struct circle {
    Point o;
    double r;
    circle() {}
    circle(Point P, double R = 0) { o = P, r = R; }
};
double length(const Point &p) {
    return sqrt(p.x * p.x + p.y * p.y);
}
double length(const Line &1) {
    Point p(1);
    return length(p);
Point::Point(const Line &a) { *this = a.t - a.s; }
istream &operator>>(istream &in, Point &a) {
    in >> a.x >> a.y;
    return in;
double dot(const Point &a, const Point &b) {
    return a.x * b.x + a.y * b.y;
}
double det(const Point &a, const Point &b) {
    return a.x * b.y - a.y * b.x;
int sgn(const double &x) \{ return fabs(x) < eps ? 0 : (x > 0 ? 1 : -1); 
double sqr(const double &x) { return x * x; }
Point rotate(const Point &a, const double &ang) {
    double x = cos(ang) * a.x - sin(ang) * a.y;
    double y = sin(ang) * a.x + cos(ang) * a.y;
    return {x, y};
}
//点在线段上 <=0 包含端点
bool sp_on(const Line &seg, const Point &p) {
    Point a = seg.s, b = seg.t;
    return !sgn(det(p - a, b - a)) && sgn(dot(p - a, p - b)) <= 0;
bool lp_on(const Line &line, const Point &p) {
    Point a = line.s, b = line.t;
```

```
return !sgn(det(p - a, b - a));
}
//等于不包含共线
int andrew(Point *point, Point *convex, int n) {
    sort(point, point + n, [](Point a, Point b) {
        if (a.x != b.x) return a.x < b.x;
        return a.y < b.y;</pre>
    });
    int top = 0;
    for (int i = 0; i < n; i++) {</pre>
        while ((top > 1) && det(convex[top - 1] - convex[top - 2], poin
t[i] - convex[top - 1]) \leftarrow 0
            top--;
        convex[top++] = point[i];
    }
    int tmp = top;
    for (int i = n - 2; i >= 0; i--) {
        while ((top > tmp) && det(convex[top - 1] - convex[top - 2], po
int[i] - convex[top - 1]) <= 0)
            top--;
        convex[top++] = point[i];
    if (n > 1) top--;
    return top;
}
double slope(const Point &a, const Point &b) {
    return (a.y - b.y) / (a.x - b.x);
}
double slope(const Line &a) {
    return slope(a.s, a.t);
Point 11 intersection(const Line &a, const Line &b) {
    double s1 = det(Point(a), b.s - a.s), s2 = det(Point(a), b.t - a.s);
    return (b.s * s2 - b.t * s1) / (s2 - s1);
int ss_cross(const Line &a, const Line &b, Point &p) {
    int d1 = sgn(det(a.t - a.s, b.s - a.s));
    int d2 = sgn(det(a.t - a.s, b.t - a.s));
    int d3 = sgn(det(b.t - b.s, a.s - b.s));
    int d4 = sgn(det(b.t - b.s, a.t - b.s));
    if ((d1 ^ d2) == -2 && (d3 ^ d4) == -2) {
        p = 11 intersection(a, b);
        return 1;
    if (!d1 && sp_on(a, b.s)) {
        p = b.s;
        return 2;
    if (!d2 && sp_on(a, b.t)) {
        p = b.t;
```

```
return 2;
    }
    if (!d3 && sp_on(b, a.s)) {
        p = a.s;
        return 2;
    if (!d4 && sp on(b, a.t)) {
        p = a.t;
        return 2;
    return 0;
Point project(const Line &1, const Point &p) {
    Point base(1);
    double r = dot(base, p - 1.s) / sqr(length(base));
    return 1.s + (base * r);
}
double sp dist(const Line &1, const Point &p) {
    if (l.s == l.t) return dist(l.s, p);
    Point x = p - 1.s, y = p - 1.t, z = 1.t - 1.s;
    if (sgn(dot(x, z)) < 0)return length(x);//P 距离A 更近
    if (sgn(dot(y, z)) > 0)return length(y);//P 距离 B 更近
    return abs(det(x, z) / length(z));//面积除以底边长
double lp dist(const Line &1, const Point &p) {
    Point x = p - 1.s, y = p - 1.t, z = 1.t - 1.s;
    return abs(det(x, z) / length(z));//面积除以底边长
int lc_cross(const Line &1, const Point &a, const double &r, pair<Point,</pre>
 Point> &ans) {
    int num = 0;
    Point pr = project(1, a);
    double dis = dist(pr, a);
    double tmp = r * r - dis * dis;
    if (sgn(tmp) == 1) num = 2;
    else if (sgn(tmp) == 0) num = 1;
    else return 0;
    double base = sqrt(r * r - dis * dis);
    Point e(1);
    e.standardize();
    e = e * base;
    ans = make_pair(pr + e, pr - e);
    return num;
int cc cross(const Point &c1, const double &r1, const Point &c2, const
double &r2, pair<Point, Point> &ans) {
    double x1 = c1.x, x2 = c2.x, y1 = c1.y, y2 = c2.y;
    double d = length(c1 - c2);
    if (sgn(fabs(r1 - r2) - d) > 0) return -1; //内含
```

```
if (sgn(r1 + r2 - d) < 0) return 0; //相离
   double a = r1 * (x1 - x2) * 2, b = r1 * (y1 - y2) * 2, c = r2 * r2
- r1 * r1 - d * d;
   double p = a * a + b * b, q = -a * c * 2, r = c * c - b * b;
    double cosa, sina, cosb, sinb;
   //One Intersection
   if (sgn(d - (r1 + r2)) == 0 \mid | sgn(d - fabs(r1 - r2)) == 0) {
        cosa = -q / p / 2;
        sina = sqrt(1 - sqr(cosa));
       Point p0(x1 + r1 * cosa, y1 + r1 * sina);
       if (sgn(dist(p0, c2) - r2)) p0.y = y1 - r1 * sina;
       ans = pair<Point, Point>(p0, p0);
       return 1;
   //Two Intersections
   double delta = sqrt(q * q - p * r * 4);
   cosa = (delta - q) / p / 2;
   cosb = (-delta - q) / p / 2;
   sina = sqrt(1 - sqr(cosa));
    sinb = sqrt(1 - sqr(cosb));
   Point p1(x1 + r1 * cosa, y1 + r1 * sina);
   Point p2(x1 + r1 * cosb, y1 + r1 * sinb);
   if (sgn(dist(p1, c2) - r2)) p1.y = y1 - r1 * sina;
   if (sgn(dist(p2, c2) - r2)) p2.y = y1 - r1 * sinb;
   if (p1 == p2) p1.y = y1 - r1 * sina;
   ans = pair<Point, Point>(p1, p2);
   return 2;
Point lp sym(const Line &1, const Point &p) {
   return p + (project(1, p) - p) * 2;
}
double alpha(const Point &t1, const Point &t2) {
   double theta;
   theta = atan2((double) t2.y, (double) t2.x) - atan2((double) t1.y,
(double) t1.x);
   if (sgn(theta) < 0)</pre>
        theta += 2.0 * PI;
   return theta;
}
int pip(const Point *P, const int &n, const Point &a) {// 【射线法】判断
点 A 是否在任意多边形 Poly 以内
   int cnt = 0;
   int tmp;
   for (int i = 1; i <= n; ++i) {</pre>
        int j = i < n ? i + 1 : 1;
       if (sp_on(Line(P[i], P[j]), a))return 2;//点在多边形上
       if (a.y >= min(P[i].y, P[j].y) && a.y < max(P[i].y, P[j].y))//
纵坐标在该线段两端点之间
           tmp = P[i].x + (a.y - P[i].y) / (P[j].y - P[i].y) * (P[j].x
```

```
- P[i].x), cnt += sgn(tmp - a.x) > 0;//交点在A 右方
   return cnt & 1;//穿过奇数次则在多边形以内
bool pip convex jud(const Point &a, const Point &L, const Point &R) {//
判断AL 是否在AR 右边
   return sgn(det(L - a, R - a)) > 0;//必须严格以内
}
bool pip convex(const Point *P, const int &n, const Point &a) {// 【二分
法】判断点A 是否在凸多边形 Poly 以内
   //点按逆时针给出
   if (pip_convex_jud(P[0], a, P[1]) || pip_convex_jud(P[0], P[n - 1],
 a)) return 0;//在P[0 1]或P[0 n-1]外
   if (sp_on(Line(P[0], P[1]), a) || sp_on(Line(P[0], P[n - 1]), a)) r
eturn 2;//在P[0 1]或P[0 n-1]上
   int l = 1, r = n - 2;
   while (1 < r) {//二分找到一个位置 pos 使得 P[0] A 在 P[0 pos], P[0 (pos+
1) ]之间
       int mid = (1 + r + 1) >> 1;
       if (pip_convex_jud(P[0], P[mid], a))1 = mid;
       else r = mid - 1;
   }
   if (pip_convex_jud(P[1], a, P[1 + 1]))return 0;//\pmP[pos_(pos+1)]外
   if (sp on(Line(P[1], P[1 + 1]), a))return 2;//在P[pos (pos+1)]上
   return 1;
}
// 多边形是否包含线段
// 因此我们可以先求出所有和线段相交的多边形的顶点, 然后按照X-Y 坐标排序(X 坐标
小的排在前面,对于X 坐标相同的点,Y 坐标小的排在前面,
// 这种排序准则也是为了保证水平和垂直情况的判断正确),这样相邻的两个点就是在线
段上相邻的两交点,如果任意相邻两点的中点也在多边形内,
// 则该线段一定在多边形内。
int pp judge(Point *A, int n, Point *B, int m) {// 【判断多边形 A 与多边形
B 是否相离】
   for (int i1 = 1; i1 <= n; ++i1) {</pre>
       int j1 = i1 < n ? i1 + 1 : 1;
       for (int i2 = 1; i2 <= m; ++i2) {</pre>
           int j2 = i2 < m ? i2 + 1 : 1;
          Point tmp;
          if (ss_cross(Line(A[i1], A[j1]), Line(B[i2], B[j2]), tmp))
return 0;//两线段相交
          if (pip(B, m, A[i1]) || pip(A, n, B[i2]))return 0;//点包含在
内
       }
   }
   return 1;
}
```

```
double area(Point *P, int n) {//【任意多边形 P 的面积】
   double S = 0;
   for (int i = 1; i <= n; i++) S += det(P[i], P[i < n ? i + 1 : 1]);</pre>
   return S / 2.0:
Line Q[N];
int judge(Line L, Point a) { return sgn(det(a - L.s, L.t - L.s)) > 0; }
//判断点 a 是否在直线 L 的右边
int halfcut(Line *L, int n, Point *P) {//【半平面交】
    sort(L, L + n, [](const Line &a, const Line &b) {
        double d = atan2((a.t - a.s).y, (a.t - a.s).x) - atan2((b.t - b.
s).y, (b.t - b.s).x);
       return sgn(d) ? sgn(d) < 0 : judge(a, b.s);</pre>
    });
    int m = n;
   n = 0;
   for (int i = 0; i < m; ++i)</pre>
        if (i == 0 || sgn(atan2(Point(L[i]).y, Point(L[i]).x) - atan2(P
oint(L[i-1]).y, Point(L[i-1]).x)))
           L[n++] = L[i];
    int h = 1, t = 0;
   for (int i = 0; i < n; ++i) {</pre>
       while (h < t && judge(L[i], ll_intersection(Q[t], Q[t - 1]))) -
-t;//当队尾两个直线交点不是在直线L[i]上或者左边时就出队
       while (h < t \&\& judge(L[i], ll_intersection(Q[h], Q[h + 1]))) +
+h;//当队头两个直线交点不是在直线L[i]上或者左边时就出队
       O[++t] = L[i];
   while (h < t && judge(Q[h], ll intersection(Q[t], Q[t-1]))) --t;
   while (h < t \&\& judge(Q[t], ll_intersection(Q[h], Q[h + 1]))) ++h;
   n = 0;
   for (int i = h; i <= t; ++i) {</pre>
       P[n++] = ll\_intersection(Q[i], Q[i < t ? i + 1 : h]);
   return n;
Point V1[N], V2[N];
int mincowski(Point *P1, int n, Point *P2, int m, Point *V) {// 【闵可夫
斯基和】求两个凸包{P1},{P2}的向量集合{V}={P1+P2}构成的凸包
   for (int i = 0; i < n; ++i) V1[i] = P1[(i + 1) % n] - P1[i];</pre>
   for (int i = 0; i < m; ++i) V2[i] = P2[(i + 1) \% m] - P2[i];
    int t = 0, i = 0, j = 0;
   V[t++] = P1[0] + P2[0];
   while (i < n && j < m) V[t] = V[t - 1] + (sgn(det(V1[i], V2[j])) >
0 ? V1[i++] : V2[j++]), t++;
   while (i < n) V[t] = V[t - 1] + V1[i++], t++;
   while (j < m) V[t] = V[t - 1] + V2[j++], t++;
   return t;
}
```

```
circle getcircle(const Point &A, const Point &B, const Point &C) {//
【三点确定一圆】向量垂心法
   Point P1 = (A + B) * 0.5, P2 = (A + C) * 0.5;
   Line R1 = Line(P1, P1 + normal(B - A));
   Line R2 = Line(P2, P2 + normal(C - A));
   circle 0;
   0.o = 11 intersection(R1, R2);
   0.r = length(A - 0.o);
   return 0;
}
struct ConvexHull {
   int op;
   struct cmp {
       bool operator()(const Point &a, const Point &b) const {
           return sgn(a.x - b.x) < 0 \mid | sgn(a.x - b.x) == 0 && sgn(a.y)
 - b.y) < 0;
       }
   };
   set<Point, cmp> s;
   ConvexHull(int o) {
       op = o;
       s.clear();
   inline int PIP(Point P) {
       set<Point>::iterator it = s.lower_bound(Point(P.x, -dinf));//救
到第一个横坐标大干P 的点
       if (it == s.end())return 0;
       if (sgn(it->x - P.x) == 0) return sgn((P.y - it->y) * op) <= 0;
//比较纵坐标大小
       if (it == s.begin())return 0;
       set<Point>::iterator j = it, k = it;
       return sgn(det(P - *j, *k - *j) * op) >= 0;//看叉姬1
    inline int judge(set<Point>::iterator it) {
       set<Point>::iterator j = it, k = it;
       if (j == s.begin())return 0;
       if (++k == s.end())return 0;
       return sgn(det(*it - *j, *k - *j) * op) >= 0;//看叉姬
   inline void insert(Point P) {
       if (PIP(P))return;//如果点P已经在凸壳上或凸包里就不插入了
       set<Point>::iterator tmp = s.lower_bound(Point(P.x, -inf));
       if (tmp != s.end() && sgn(tmp->x - P.x) == 0)s.erase(tmp);//特
判横坐标相等的点要去掉
       s.insert(P):
       set<Point>::iterator it = s.find(P), p = it;
       if (p != s.begin()) {
```

```
--p;
            while (judge(p)) {
                set<Point>::iterator temp = p--;
                s.erase(temp);
            }
        if ((p = ++it) != s.end()) {
            while (judge(p)) {
                set<Point>::iterator temp = p++;
                s.erase(temp);
            }
        }
} up(1), down(-1);
int PIC(circle C, Point a) { return sgn(length(a - C.o) - C.r) <= 0; }/</pre>
/判断点A 是否在圆C 内
void Random(Point *P, int n) { for (int i = 0; i < n; ++i)swap(P[i], P</pre>
[(rand() + 1) % n]); }//随机一个排列
circle min circle(Point *P, int n) {// 【求点集 P 的最小覆盖圆】 O(n)
// random shuffle(P,P+n);
    Random(P, n);
    circle C = circle(P[0], 0);
    for (int i = 1; i < n; ++i)
        if (!PIC(C, P[i])) {
            C = circle(P[i], 0);
            for (int j = 0; j < i; ++j)</pre>
                if (!PIC(C, P[j])) {
                    C.o = (P[i] + P[j]) * 0.5, C.r = length(P[j] - C.o);
                    for (int k = 0; k < j; ++k) if (!PIC(C, P[k])) C =</pre>
getcircle(P[i], P[j], P[k]);
    return C;
}
自适应辛普森
double f(double x) {
double simpson(double 1, double r) {
    double mid = (1 + r) / 2;
    return (r - 1) * (f(1) + 4 * f(mid) + f(r)) / 6; // 辛普森公式
}
double asr(double 1, double r, double EPS, double ans) {
    double mid = (1 + r) / 2;
    double fl = simpson(l, mid), fr = simpson(mid, r);
    if (abs(fl + fr - ans) <= 15 * EPS)
        return fl + fr + (fl + fr - ans) / 15; // 足够相似的话就直接返回
```

```
return asr(1, mid, EPS / 2, fl) +
           asr(mid, r, EPS / 2, fr); // 否则分割成两段递归求解
}
数据结构
kruskal 重构树
int pa[N];
void init(int n) {
    for (int i = 0; i <= n; i++) {</pre>
        pa[i] = i;
    }
}
int find(int a) {
    return pa[a] == a ? a : pa[a] = find(pa[a]);
int kruskal() {
    int kcnt = n;
    init(n);
    sort(e + 1, e + 1 + m, [](edge a, edge b) { return a.l < b.l; });</pre>
    for (int i = 1; i <= m; i++) {</pre>
        int u = find(e[i].from);
        int v = find(e[i].to);
        if (u == v) continue;
        w[++kcnt] = e[i].1;
        pa[kcnt] = pa[u] = pa[v] = kcnt;
        g[u].push_back(kcnt);
        g[v].push_back(kcnt);
        g[kcnt].push_back(u);
        g[kcnt].push back(v);
    }
   return kcnt;
}
срр
普通莫队
#include <bits/stdc++.h>
using namespace std;
const int N = 1e6 + 10, M = 1e6 + 10;
int a[N];
struct node {
      int id, l, r;
} mp[M];
int len;
int ans[M], cnt[1000010];
int getNum(int 1) {
```

```
return 1 / len;
}
//左指针的分块,右指针的大小
bool cmp (const node &a, const node & b) {
      if(getNum(a.1) == getNum(b.1)) return a.r < b.r;</pre>
      return a.l < b.l;</pre>
}
/* 奇偶优化
struct node {
 int l, r, id;
 bool operator<(const node &x) const {</pre>
   if (l / unit != x.l / unit) return l < x.l;</pre>
   if ((l / unit) & 1)
      return r < x.r; // 注意这里和下面一行不能写小于 (大于) 等于
   return r > x.r;
};
*/
void add(int x, int& res) {
      if(cnt[x] == 0) res++;
      cnt[x] ++;
void del(int x, int& res) {
      cnt[x] --;
      if(cnt[x] == 0) res --;
int main() {
      ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
      int n;
      cin >> n;
      for(int i = 1; i <= n; ++ i) {</pre>
            cin >> a[i];
      }
      int m;
      cin >> m;
      len = sqrt((double)n * n / m);
      for(int i = 1; i <= m; ++ i) {</pre>
            mp[i].id = i;
            cin >> mp[i].l >> mp[i].r;
      sort(mp + 1, mp + m + 1, cmp);
      //离线处理询问
      int res = 0, i = 0, j = 0;
      for(int k = 1; k <= m; ++ k) {</pre>
            int id = mp[k].id, l = mp[k].l, r = mp[k].r;
            while(j < r) add(a[++j], res);</pre>
            while(j > r) del(a[j--], res);
```

```
while(i < 1) del(a[i++], res);
    while(i > 1) add(a[--i], res);
    ans[id] = res;
}

for(int i = 1; i <= m; ++ i) {
    cout << ans[i] << endl;
}
    return 0;
}</pre>
```

带修莫队 #include <bits/stdc++.h> using namespace std; const int N = 10010; int a[N], cnt[1000010], ans[N]; int len, mq, mc; struct Query { int id, 1, r, t; } q[N]; struct Modify { int p, c; } c[N]; int getNum(int x) { return x / len; } // L 所在块的编号,r 所在块的编号,t 升序 bool cmp(const Query& a, const Query& b) { if(getNum(a.1) == getNum(b.1) && getNum(a.r) == getNum(b.r)) { return a.t < b.t;</pre> if(getNum(a.1) == getNum(b.1)) return a.r < b.r;</pre> return a.l < b.l;</pre> } void add(int x, int& res) { **if** (!cnt[x]) res ++ ; cnt[x] ++ ; void del(int x, int& res) { cnt[x] --; **if** (!cnt[x]) res -- ; } int main() { ios::sync_with_stdio(0); cin.tie(0); cout.tie(0); int n, m; cin >> n >> m;

```
char op;
    int x, y;
    for(int i = 1; i <= n; ++ i) {</pre>
          cin >> a[i];
    for(int i = 1; i <= m; ++ i) {</pre>
          cin >> op >> x >> y;
      if (op == 'Q') q[++ mq] = \{mq, x, y, mc\};
      else c[ ++ mc] = \{x, y\};
    }
///
    len = cbrt((double)n * mc) + 1;
sort(q + 1, q + mq + 1, cmp);
    int i = 1, j = 0, t = 0, res = 0;
    for(int k = 1; k <= mq; ++ k) {</pre>
           int id = q[k].id, 1 = q[k].1, r = q[k].r, tm = q[k].t;
           while(j < r) add(a[++ j], res);</pre>
           while(j > r) del(a[j --], res);
           while(i < 1) del(a[i ++], res);</pre>
           while(i > 1) add(a[-- i], res);
           while(t < tm) {</pre>
                 ++ t;
                 if(c[t].p >= i && c[t].p <= j) {
                        del(a[c[t].p], res);
                        add(c[t].c, res);
                 swap(a[c[t].p], c[t].c);
           while(t > tm) {
                 if(c[t].p >= i && c[t].p <= j) {
                        del(a[c[t].p], res);
                        add(c[t].c, res);
                 }
                 swap(a[c[t].p], c[t].c);
                 -- t;
           }
           ans[id] = res;
    }
    for(int i = 1; i <= mq; ++ i) {</pre>
           cout << ans[i] << endl;</pre>
    }
```

}

线段树合并分裂

```
11 nodetot, recycnt, bac[N << 5], ch[N << 5][2], rt[N];</pre>
ll val[N << 5];
11 newnod() { return (recycnt ? bac[recycnt--] : ++nodetot); }
void recyc(ll p) {
    bac[++recycnt] = p, ch[p][0] = ch[p][1] = val[p] = 0;
    return;
}
void pushdown(ll p) {
void pushup(ll p) {
    val[p] = 0;
    if (ch[p][0]) val[p] += val[ch[p][0]];
    if (ch[p][1]) val[p] += val[ch[p][1]];
}
void modify(ll &p, ll l, ll r, ll pos, ll v) {
    if (!p) { p = newnod(); }
    if (1 == r) {
        val[p] += v;
        return:
    11 \text{ mid} = (1 + r) >> 1;
     pushdown(p);
    if (pos \leftarrow mid) { modify(ch[p][0], 1, mid, pos, v); }
    else { modify(ch[p][1], mid + 1, r, pos, v); }
    pushup(p);
    return;
}
11 query(11 p, 11 l, 11 r, 11 xl, 11 xr) {
    if (xr < 1 || r < x1) { return 0; }</pre>
    if (x1 <= 1 && r <= xr) { return val[p]; }</pre>
    11 \text{ mid} = (1 + r) >> 1;
// pushdown(p);
    return query(ch[p][0], 1, mid, x1, xr) + query(ch[p][1], mid + 1, r,
x1, xr);
11 kth(ll p, ll l, ll r, ll k) {
    if (l == r) { return l; }
    11 \text{ mid} = (1 + r) >> 1;
    pushdown(p);
    if (val[ch[p][0]] >= k) { return kth(ch[p][0], 1, mid, k); }
    else { return kth(ch[p][1], mid + 1, r, k - val[ch[p][0]]); }
11 merge(ll x, ll y, ll l, ll r) {
    if (!x || !y) {
        return x + y;
      // 只有一边有点,不用合并
    ll p = newnod(); // 创建一个新结点 p
                                    // 边界(某些时候可以省略,见下面一个代
    if (1 == r) {
```

```
码)
        val[p] = val[x] + val[y];
        return p;
    }
     pushdown(x), pushdown(y);
    11 \text{ mid} = (1 + r) >> 1;
    ch[p][0] = merge(ch[x][0], ch[y][0], 1, mid);
    ch[p][1] = merge(ch[x][1], ch[y][1], mid + 1, r);
    recyc(x), recyc(y);
                                   // 垃圾回收
    pushup(p);
                                     // pushup
    return p;
}
void split(ll x, ll &y, ll k) {
    if (x == 0) return;
    y = newnod();
    ll v = val[ch[x][\emptyset]];
   pushdown(x);
    if (k > v) { split(ch[x][1], ch[y][1], k - v); }
    else { swap(ch[x][1], ch[y][1]); }
    if (k < v) { split(ch[x][0], ch[y][0], k); }</pre>
    val[y] = val[x] - k;
    val[x] = k;
    return;
}
```

主席树

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 N = 1 << 20;
11 ch[N << 5][2], rt[N], tot;</pre>
ll val[N << 5];
11 update(11 a, 11 b) {
    return a + b;
}
ll build(ll l, ll r) { // 建树
    11 p = ++tot;
    if (1 == r) {
        //初始化
        val[p] = 0;
        return p;
    11 \text{ mid} = (1 + r) >> 1;
    ch[p][0] = build(1, mid);
    ch[p][1] = build(mid + 1, r);
    val[p] = update(val[ch[p][0]], val[ch[p][1]]);
```

```
return p; // 返回该子树的根节点
}
ll modify(ll pre, ll l, ll r, ll pos, ll v) { // 插入操作
   11 \text{ now} = ++\text{tot};
   ch[now][0] = ch[pre][0], ch[now][1] = ch[pre][1];
   if (1 == r) {
       val[now] = val[pre] + v;
       return now;
   11 \text{ mid} = (1 + r) >> 1;
   if (pos <= mid)</pre>
        ch[now][0] = modify(ch[now][0], 1, mid, pos, v);
   else
        ch[now][1] = modify(ch[now][1], mid + 1, r, pos, v);
   val[now] = update(val[ch[now][0]], val[ch[now][1]]);
   return now;
ll kth(ll pre, ll now, ll l, ll r, ll k) { // 查询操作
   11 \text{ mid} = (1 + r) >> 1;
   ll x = val[ch[now][0]] - val[ch[pre][0]]; // 通过区间减法得到左儿子
的信息
   if (1 == r) return 1;
   if (k <= x) // 说明在左儿子中
        return kth(ch[pre][0], ch[now][0], 1, mid, k);
   else // 说明在右儿子中
       return kth(ch[pre][1], ch[now][1], mid + 1, r, k - x);
}
ll query(ll pre, ll now, ll l, ll r, ll ql, ll qr) { // 查询操作
   if (ql <= 1 && r <= qr) {
       return val[now] - val[pre];
   if (qr < 1 || r < q1) {
       return 0;
    }
   11 \text{ mid} = (1 + r) >> 1;
   11 lv = query(ch[pre][0], ch[now][0], l, mid, ql, qr);
   11 rv = query(ch[pre][1], ch[now][1], mid + 1, r, ql, qr);
   return update(lv, rv);
}
LCT
11 ch[N][2], f[N], sum[N], val[N], tag[N], siz[N], siz2[N];
inline void pushup(ll p) {
```

```
sum[p] = sum[ch[p][0]] ^ sum[ch[p][1]] ^ val[p];
    siz[p] = siz[ch[p][0]] + siz[ch[p][1]] + 1 + siz2[p];
}
inline void pushdown(ll p) {
    if (tag[p]) {
       if (ch[p][0]) swap(ch[ch[p][0]][0], ch[ch[p][0]][1]), tag[ch[p]
[0] ^= 1;
       if (ch[p][1]) swap(ch[ch[p][1]][0], ch[ch[p][1]][1]), tag[ch[p]
[1]] ^= 1;
       tag[p] = 0;
    }
}
11 getch(11 x) { return ch[f[x]][1] == x; }
bool isroot(ll x) { return ch[f[x]][0] != x && ch[f[x]][1] != x; }
inline void rotate(ll x) {
    11 y = f[x], z = f[y], k = getch(x);
    if (!isroot(y)) ch[z][ch[z][1] == y] = x;
   // 上面这句一定要写在前面,普通的 Splay 是不用的,因为 is Root (后面会讲)
    ch[y][k] = ch[x][!k], f[ch[x][!k]] = y;
    ch[x][!k] = y, f[y] = x, f[x] = z;
    pushup(y), pushup(x);
```

```
}
// 从上到下一层一层 pushDown 即可
void update(ll p) {
   if (!isroot(p)) update(f[p]);
   pushdown(p);
}
inline void splay(ll x) {
   update(x); // 马上就能看到啦。 在
   // Splay 之前要把旋转会经过的路径上的点都 PushDown
   for (ll fa; fa = f[x], !isroot(x); rotate(x)) {
       if (!isroot(fa)) rotate(getch(fa) == getch(x) ? fa : x);
   }
}
// 回顾一下代码
inline void access(ll x) {
   for (11 p = 0; x; p = x, x = f[x]) {
       splay(x), siz2[x] += siz[ch[x][1]] - siz[p], ch[x][1] = p, push
up(x);
   }
}
```

```
inline void makeroot(ll p) {
    access(p);
   splay(p);
    swap(ch[p][0], ch[p][1]);
   tag[p] ^= 1;
}
inline void split(ll a, ll b) {
   makeroot(a);
    access(b);
   splay(b);
}
inline 11 find(11 p) {
    access(p), splay(p);
   while (ch[p][0]) pushdown(p), p = ch[p][0];
    splay(p);
    return p;
}
inline void link(ll x, ll y) {
   makeroot(y);
   makeroot(x);
```

```
if (find(y) != x) {
       f[x] = y;
       siz2[y] += siz[x];
    }
}
inline void cut(ll x, ll y) {
   makeroot(x);
    if (find(y) == x \&\& f[y] == x) {
       ch[x][1] = f[y] = 0;
       pushup(x);
    }
}
void init(int n) {
   for (int i = 1; i <= n; i++) siz[i] = 1;</pre>
}
Splay1
#include <bits/stdc++.h>
using namespace std;
struct Splay {
    static const int N = 100005;
    int rt, tot, fa[N], ch[N][2], val[N], cnt[N], sz[N];
   // rt=根编号, tot=总节点, fa=父节点编号, ch=左/右儿子编号, val=节点的
值, cnt=权值出现次数, sz=子树大小
    void maintain(int x) { //更新 x 节点字数大小
        sz[x] = sz[ch[x][0]] + sz[ch[x][1]] + cnt[x];
```

```
bool get(int x) {
    return x == ch[fa[x]][1];
} //返回节点是父亲的0/1-左/右儿子
void clear(int x) { //销毁节点 x
    ch[x][0] = ch[x][1] = fa[x] = val[x] = sz[x] = cnt[x] = 0;
}
void rotate(int x) { //旋转
    int y = fa[x], z = fa[y], chk = get(x);
    ch[y][chk] = ch[x][chk ^ 1];
    fa[ch[x][chk ^ 1]] = y;
    ch[x][chk ^ 1] = y;
    fa[y] = x;
    fa[x] = z;
    if (z) ch[z][y == ch[z][1]] = x;
    maintain(x);
    maintain(y);
}
void splay(int x) { //将x 节点移动到根
    for (int f = fa[x]; f = fa[x], f; rotate(x))
        if (fa[f]) rotate(get(x) == get(f) ? f : x);
    rt = x;
}
void ins(int k) { //插入
    if (!rt) {
        val[++tot] = k;
        cnt[tot]++;
        rt = tot;
        maintain(rt);
        return;
    }
    int cnr = rt, f = 0;
   while (1) {
        if (val[cnr] == k) {
            cnt[cnr]++;
            maintain(cnr);
            maintain(f);
            splay(cnr);
            break;
        }
        f = cnr;
        cnr = ch[cnr][val[cnr] < k];</pre>
        if (!cnr) {
            val[++tot] = k;
            cnt[tot]++;
            fa[tot] = f;
            ch[f][val[f] < k] = tot;
            maintain(tot);
            maintain(f);
```

```
splay(tot);
                break;
           }
       }
   }
   int rk(int k) { // k 权值的排名
        int res = 0, cnr = rt;
       while (1) {
            if (k < val[cnr]) {</pre>
                cnr = ch[cnr][0];
            } else {
                res += sz[ch[cnr][0]];
                if (k == val[cnr]) {
                    splay(cnr);
                    return res + 1;
                }
                res += cnt[cnr];
                cnr = ch[cnr][1];
            }
       }
   }
   int kth(int k) { //第 k 名的权值
       int cnr = rt;
       while (1) {
            if (ch[cnr][0] && k <= sz[ch[cnr][0]]) {</pre>
                cnr = ch[cnr][0];
            } else {
                k -= cnt[cnr] + sz[ch[cnr][0]];
                if (k <= 0) {
                    splay(cnr);
                    return val[cnr];
                cnr = ch[cnr][1];
            }
       }
   }
   int pre() { //前驱节点编号
        int cnr = ch[rt][0];
       while (ch[cnr][1]) cnr = ch[cnr][1];
       splay(cnr);
       return cnr;
   } // 若需要得到前驱 tree.ins(x), printf("%d\n", tree.val[tree.pre
()]),
      // tree.del(x);
   int nxt() { //后驱节点编号
        int cnr = ch[rt][1];
       while (ch[cnr][0]) cnr = ch[cnr][0];
        splay(cnr);
        return cnr;
```

```
} // 若需要得到后驱 tree.ins(x), printf("%d\n", tree.val[tree.pre
()]),
       // tree.del(x);
    void del(int k) { //删除k值
        rk(k);
        if (cnt[rt] > 1) {
            cnt[rt]--;
            maintain(rt);
            return;
        }
        if (!ch[rt][0] && !ch[rt][1]) {
            clear(rt);
            rt = 0;
            return;
        }
        if (!ch[rt][0]) {
            int cnr = rt;
            rt = ch[rt][1];
            fa[rt] = 0;
            clear(cnr);
            return;
        if (!ch[rt][1]) {
            int cnr = rt;
            rt = ch[rt][0];
            fa[rt] = 0;
            clear(cnr);
            return;
        }
        int cnr = rt;
        int x = pre();
        splay(x);
        fa[ch[cnr][1]] = x;
        ch[x][1] = ch[cnr][1];
        clear(cnr);
       maintain(rt);
} tree;
splay2
11 ch[N][2], f[N], sum[N], val[N], tag[N], siz[N];
inline void pushup(ll p) {
    sum[p] = sum[ch[p][0]] ^ sum[ch[p][1]] ^ val[p];
    siz[p] = siz[ch[p][0]] + siz[ch[p][1]] + 1;
inline void pushdown(ll p) {
    if (tag[p]) {
        if (ch[p][0]) swap(ch[ch[p][0]][0], ch[ch[p][0]][1]), tag[ch[p]
```

```
[0] ^= 1;
       if (ch[p][1]) swap(ch[ch[p][1]][0], ch[ch[p][1]][1]), tag[ch[p]
[1]] ^= 1;
       tag[p] = 0;
   }
ll getch(ll x) { return ch[f[x]][1] == x; }
bool isroot(ll x) { return ch[f[x]][0] != x && ch[f[x]][1] != x; }
inline void rotate(ll x) {
   11 y = f[x], z = f[y], k = getch(x);
   if (!isroot(y)) ch[z][ch[z][1] == y] = x;
   // 上面这句一定要写在前面,普通的Splay 是不用的,因为 isRoot (后面会讲)
   ch[y][k] = ch[x][!k], f[ch[x][!k]] = y;
   ch[x][!k] = y, f[y] = x, f[x] = z;
   pushup(y), pushup(x);
}
// 从上到下一层一层 pushDown 即可
void update(ll p) {
   if (!isroot(p)) update(f[p]);
   pushdown(p);
}
inline void splay(ll x) {
   update(x); // 马上就能看到啦。 在
   // Splay 之前要把旋转会经过的路径上的点都 PushDown
   for (11 fa; fa = f[x], !isroot(x); rotate(x)) {
       if (!isroot(fa)) rotate(getch(fa) == getch(x) ? fa : x);
   }
}
// 回顾一下代码
inline void access(ll x) {
   for (11 p = 0; x; p = x, x = f[x]) {
       splay(x), ch[x][1] = p, pushup(x);
   }
inline void makeroot(ll p) {
   access(p);
   splay(p);
   swap(ch[p][0], ch[p][1]);
   tag[p] ^= 1;
inline void split(ll a, ll b) {
   makeroot(a);
   access(b);
   splay(b);
}
inline ll find(ll p) {
   access(p), splay(p);
   while (ch[p][0]) pushdown(p), p = ch[p][0];
```

```
splay(p);
    return p;
inline void link(ll x, ll y) {
    makeroot(x);
    if (find(y) != x) f[x] = y;
inline void cut(ll x, ll y) {
    makeroot(x);
    if (find(y) == x && f[y] == x) {
         ch[x][1] = f[y] = 0;
         pushup(x);
    }
}
Treap
#include <bits/stdc++.h>
using namespace std;
struct node {
    node* ch[2];
    int r;
    int v;
    int cmp(int const& a) const {
         if (v == a) return -a;
         return a > v ? 1 : 0;
    }
};
void rotate(node*& a, int d) {
    node* k = a \rightarrow ch[d ^ 1];
    a \rightarrow ch[d ^ 1] = k \rightarrow ch[d];
    k \rightarrow ch[d] = a;
    a = k;
void insert(node*& a, int x) {
    if (a == NULL) {
         a = new node;
         a \rightarrow ch[0] = a \rightarrow ch[1] = NULL;
         a \rightarrow v = x;
         a \rightarrow r = rand();
    } else {
         int d = a \rightarrow cmp(x);
         insert(a->ch[d], x);
         if (a->ch[d]->r > a->r) rotate(a, d ^ 1);
    }
void remove(node*& a, int x) {
```

```
int d = a \rightarrow cmp(x);
    if (d == -1) {
         if (a->ch[0] == NULL)
              a = a \rightarrow ch[1];
         else if (a->ch[1] == NULL)
              a = a \rightarrow ch[0];
         else {
              int d2 = a->ch[1]->r > a->ch[0]->r ? 0 : 1;
              rotate(a, d2);
              remove(a \rightarrow ch[d2], x);
    } else {
         remove(a->ch[d], x);
    }
}
int find(node*& a, int x) {
    if (a == NULL)
         return 0;
    else if (a->v == x)
         return 1;
    else {
         int d = a \rightarrow cmp(x);
         return find(a->ch[d], x);
    }
}
int main() {
    node* a = NULL;
    int k, 1;
    while (cin >> k >> 1) {
         if (k == 1)
              insert(a, 1);
         else if (k == 2)
              remove(a, 1);
              cout << find(a, 1) << endl;</pre>
         }
    }
}
```

```
舞蹈链 (多重覆盖)
#include <bits/stdc++.h>
using namespace std;
struct DLX {
    static const int maxn = 1000; //列的上限
    static const int maxr = 1000; //解的上限
    static const int maxnode = 5000; //总结点数上限
    static const int INF = 10000000000;
    int n, sz;
```

```
int S[maxn];
    int row[maxnode], col[maxnode];
   int L[maxnode], R[maxnode], U[maxnode], D[maxnode];
   int ansd, ans[maxr];
   int vis[maxnode];
   void init(int n) {
       this->n = n;
       //虚拟节点
       for (int i = 0; i <= n; i++) {</pre>
           U[i] = i;
           D[i] = i;
           L[i] = i - 1;
           R[i] = i + 1;
       R[n] = 0;
       L[0] = n;
       sz = n + 1;
       memset(S, 0, sizeof(S));
   void addRow(int r, vector<int> columns) {
        int first = sz;
       for (int i = 0; i < columns.size(); i++) {</pre>
            int c = columns[i];
           L[sz] = sz - 1;
           R[sz] = sz + 1;
           D[sz] = c;
           U[sz] = U[c];
           D[U[c]] = sz;
           U[c] = sz;
           row[sz] = r;
            col[sz] = c;
           S[c]++;
           SZ++;
        }
       R[sz - 1] = first;
        L[first] = sz - 1;
#define FOR(i, A, s) for (int i = A[s]; i != s; i = A[i])
   void remove(int c) {
       FOR(i, D, c) { L[R[i]] = L[i], R[L[i]] = R[i]; }
   void restore(int c) {
       FOR(i, U, c) { L[R[i]] = i, R[L[i]] = i; }
    }
   int f_check() //精确覆盖区估算剪枝
   {
       /*
        强剪枝。这个
        剪枝利用的思想是A*搜索中的估价函数。即,对于当前的递归深度K 下的矩
```

前列的所有行全部选中,去掉这些行能够覆盖到的列,将这个操作作为一层深 度。重复此操作直到所有列全部出解的深度是多少。如果当前深度加上这个估价函数返 回值,其和已然不能更优(也就是已经超过当前最优解),则直接返回,不必 再搜。 */ int ret = 0; FOR(c, R, 0) vis[c] = true; FOR(c, R, ∅) **if** (vis[c]) { ret++; vis[c] = false; FOR(i, D, c)FOR(j, R, i) vis[col[j]] = false; return ret; // d 为递归深度 void dfs(int d, vector<int>& v) { if (d + f_check() >= ansd) return; **if** (R[0] == 0) { if (d < ansd) {</pre> ansd = d; v.clear(); for (int i = 0; i < ansd; i++) {</pre> v.push_back(ans[i]); } //找到解 } return; //记录解的长度 } //找到S 最小的列 c int c = R[0]; FOR(i, R, ∅) **if** (S[i] < S[c]) //第一个未删除的列 c = i;//删除第c列 FOR(i, D, c) { //用结点i 所在的行能覆盖的所有其他列 ans[d] = row[i];remove(i); FOR(j, R, i) remove(j); //删除结点i 所在的能覆的所有其他列 dfs(d + 1, v);FOR(j, L, i) restore(j); restore(i); //恢复结点i所在的行能覆盖的所有其他列 } //恢复第c列 bool solve(vector<int>& v) { v.clear(); ansd = INF;

阵,估计其最好情况下(即最少还需要多少步)才能出解。也就是,如果将能够覆盖当

```
dfs(0, v);
       return !v.empty();
   }
};
//使用时 init 初始化,vector 中存入 r 行结点列表用 addRow 加行,solve(ans)后答
案按行的选择在 ans 中
DLX dlx;
int main() {
   int n, m;
   cin >> n >> m;
   dlx.init(m);
   for (int i = 1; i <= n; i++) {</pre>
       vector<int> v;
       for (int j = 1; j <= m; j++) {</pre>
           int a;
            cin >> a;
           if (a == 1) v.push_back(j);
       dlx.addRow(i, v);
   }
   vector<int> ans;
   dlx.solve(ans);
   for (int i = 0; i < ans.size(); i++) cout << ans[i];</pre>
}
舞蹈链 (精确覆盖)
#include <bits/stdc++.h>
using namespace std;
struct DLX {
                                  //列的上限
   static const int maxn = 1000;
   static const int maxr = 1000;
                                    //解的上限
   static const int maxnode = 5000; //总结点数上限
   int n, sz;
   int S[maxn];
   int row[maxnode], col[maxnode];
   int L[maxnode], R[maxnode], U[maxnode], D[maxnode];
   int ansd, ans[maxr];
   void init(int n) {
       this->n = n;
       //虚拟节点
       for (int i = 0; i <= n; i++) {</pre>
           U[i] = i;
           D[i] = i;
           L[i] = i - 1;
           R[i] = i + 1;
       R[n] = 0;
```

```
L[0] = n;
        sz = n + 1;
        memset(S, 0, sizeof(S));
    }
    void addRow(int r, vector<int> columns) {
        int first = sz;
        for (int i = 0; i < columns.size(); i++) {</pre>
            int c = columns[i];
            L[sz] = sz - 1;
            R[sz] = sz + 1;
            D[sz] = c;
            U[sz] = U[c];
            D[U[c]] = sz;
            U[c] = sz;
            row[sz] = r;
            col[sz] = c;
            S[c]++;
            SZ++;
        R[sz - 1] = first;
        L[first] = sz - 1;
#define FOR(i, A, s) for (int i = A[s]; i != s; i = A[i])
    void remove(int c) {
        L[R[c]] = L[c];
        R[L[c]] = R[c];
        FOR(i, D, c)
        FOR(j, R, i) {
            U[D[j]] = U[j];
            D[U[j]] = D[j];
            --S[col[j]];
        }
    }
    void restore(int c) {
        FOR(i, U, c)
        FOR(j, L, i) {
            ++S[col[j]];
            U[D[j]] = j;
            D[U[j]] = j;
        L[R[c]] = c;
        R[L[c]] = c;
    }
    // d 为递归深度
    bool dfs(int d) {
        if (R[0] == 0) {
                          //找到解
            ansd = d;
            return true; //记录解的长度
        }
```

```
//找到S 最小的列c
       int c = R[0];
      FOR(i, R, 0) if (S[i] < S[c]) c = i; //第一个未删除的列
       remove(c);
                    //删除第c列
       FOR(i, D, c) { //用结点i 所在的行能覆盖的所有其他列
          ans[d] = row[i];
          FOR(j, R, i) remove(col[j]); //删除结点i 所在的能覆的所有其他
列
          if (dfs(d + 1)) return true;
          FOR(j, L, i) restore(col[j]); //恢复结点i 所在的行能覆盖的所
有其他列
       restore(c); //恢复第c列
      return false;
   bool solve(vector<int>& v) {
      v.clear();
      if (!dfs(0)) return false;
      for (int i = 0; i < ansd; i++) v.push_back(ans[i]);</pre>
      return true;
   }
};
//使用时 init 初始化,vector 中存入 r 行结点列表用 addRow 加行,solve(ans)后答
案按行的选择在 ans 中
```

数论

```
lucas 求组合数
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
11 p;
const int maxn = 1e5 + 10;
ll \ qpow(ll \ x, ll \ n)
      11 \text{ res} = 1;
      while(n){
            if(n & 1) res = (res * x) % p;
            x = (x * x) % p;
            n >>= 1;
      }
      return res;
11 C(11 up, 11 down){
      if(up > down) return 0;
```

```
11 \text{ res} = 1;
//
      for(int i = up + 1; i <= down; ++ i){</pre>
//
            res = (res * i) % p;
//
//
      for(int i = 1; i <= down - up; ++ i){
//
            res = (res * qpow(i, p - 2)) % p;
//
      for(int i = 1, j = down; i <= up; ++ i, -- j){</pre>
            res = (res * j) % p;
            res = (res * qpow(i, p - 2)) % p;
      }
      return res;
}
11 lucas(ll up, ll down){
      if(up 
      return C(up % p, down % p) * lucas(up / p, down / p) % p;
}
int main(){
      ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
      int T;
      cin >> T;
      while (T --){
            11 down, up;
            cin >> down >> up >> p;
            cout << lucas(up, down) % p << endl;</pre>
      }
      return 0;
}
```

扩展欧几里得求逆元

```
typedef long long ll;
void extgcd(ll a,ll b,ll& d,ll& x,ll& y){
    if(!b){ d=a; x=1; y=0;}
    else{ extgcd(b,a%b,d,y,x); y-=x*(a/b); }
}
ll inverse(ll a,ll n){
    ll d,x,y;
    extgcd(a,n,d,x,y);
    return d==1?(x+n)%n:-1;
}
```

```
逆元线性递推 inv 阶乘组合数
11 fac[maxn];// n!
ll invfac[maxn]; // n!的inv
ll invn[maxn]; //n 的inv
int init(){
      int len=(int)(1e5+5);
     fac[0]=fac[1]=invfac[0]=invfac[1]=invn[0]=invn[1]=1;
     for(int i=2;i<=len;++i){</pre>
           fac[i]=fac[i-1]*i%mod;
           invn[i]=(mod-mod/i)*invn[mod%i]%mod;
           invfac[i]=invfac[i-1]*invn[i]%mod;
      }
ll C(ll n,ll m){
     if(n>m) return 0;
     if(n<0 || m<0) return 0;</pre>
     11 res=fac[m];
     res=res*invfac[m-n]%mod;
     res=res*invfac[n]%mod;
     return res;
}
// 先 init (), C(n,m) n 在上面..
//改init 里面的Len (为maxn 的长度 注意不要数组越界)
数学
一些范围
1~n的质数个数
1~2e9中拥有最多约数个数的数拥有的约数个数
约 1600
n 个不同的点可以构成n^{n-2}棵不同的树
BSGS
求a^t \equiv b \pmod{p} (a,p) = 1 的最小的 t
t = x \times k - y, x \in [1, k], y \in [0, k - 1]
t \in [1, k^2]
```

 $a^k x \equiv b \times a^y \pmod{p}$

对 $b \times a^y$ 建立 hash 表, 枚举 x 看是否有解

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
unordered_map<int , int> mp;
int bsgs(int a, int p, int b) {
      if (1 % p == b % p) return 0; // 特判0 是不是解
      mp.clear();
      int k = sqrt(p) + 1;
      for(int i = 0, j = b % p; i < k; ++ i, j = (ll)j * a % p) {</pre>
            mp[j] = i;
      }
      int ak = 1;
      for(int i = 0; i < k; ++i) {</pre>
            ak = (11)ak * a % p;
      }
      for(int i = 1, j = ak % p; i <= k; ++ i, j = (11)j * ak % p) {
            if(mp.count(j)) return (ll)i * k - mp[j];
      }
      return -1;
int main() {
      ios::sync_with_stdio(∅);
      cin.tie(0); cout.tie(0);
      int a, p, b;
      while(cin >> a >> p >> b, a | p | b) {
            int res;
            res = bsgs(a, p, b);
            if(res == -1) {
                  cout << "No Solution\n";</pre>
            }
            else {
                  cout << res << endl;</pre>
            }
      }
      return 0;
}
```

扩展 BSGS

```
求a^t \equiv b \pmod{p} 的最小的 t
当(a, p)! = 1
(a, p) = d d \mid b 无解
a^t \equiv b \pmod{p}, a^t + kp = b 两边同时除以 d, \frac{a}{d} a^{t-1} + k \frac{p}{d} = \frac{b}{d}
a^{t-1} \equiv \frac{b}{d} (\frac{a}{d})^{-1}
t' = t - 1, p' = \frac{p}{d}, b' = \frac{b}{a} (\frac{a}{d})^{-1}
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
unordered_map<11, 11> mp;
11 bsgs(ll a, ll p, ll b) {
       if(1 % p == b % p) return 0; // 特判 0 是不是解
       mp.clear();
       ll k = sqrt(p) + 1;
       for(ll i = 0, j = b % p; i < k; ++i, j = (ll)j * a % p) {</pre>
              mp[j] = i;
       }
       11 ak = 1;
       for(ll i = 0; i < k; ++i) {</pre>
              ak = (11) ak * a % p;
       }
       for(ll i = 1, j = ak % p; i <= k; ++i, j = (ll)j * ak % p) {
              if(mp.count(j)) return (ll) i * k - mp[j];
       }
       return -1;
11 gcd(l1 x, l1 y) {
       return x % y == 0 ? y : gcd(y, x % y);
void extgcd(ll a,ll b,ll& d,ll& x,ll& y){
     if(!b){
         d = a; x = 1; y = 0;
     else{
```

```
extgcd(b, a%b, d, y, x);
        y -= x * (a / b);
    }
11 inverse(ll a,ll n){
    11 d,x,y;
    extgcd(a,n,d,x,y);
    return d == 1 ? (x + n) % n : -1;
int main() {
      11 a, p, b;
      while(cin >> a >> p >> b, a | p | b) {
            11 d = gcd(a, p);
            if(d == 1) {
                   11 res = bsgs(a, p, b);
                   if(res == -1) {
                         cout << "No Solution\n";</pre>
                   }
                   else {
                         cout << res << endl;</pre>
                   }
            }
            else {
                   if(b % d != 0) {
                         cout << "No Solution\n";</pre>
                         continue;
                   }
                   else {
                          p = p / d;
                         b = (b / d) * inverse(a / d, p);
                         11 res = bsgs(a, p, b);
                         if(res == -1) {
                                cout << "No Solution\n";</pre>
                          }
                         else {
                                cout << res + 1 << endl;</pre>
                          }
                   }
            }
      }
      return 0;
}
二次剩余
```

解的数量

对于 $x^2 \equiv n \pmod{p}$ 能满足 n 是 mod p 的二次剩余的 n 一共有 $\frac{p-1}{2}$ 个(不包括 0),非二次剩余为 $\frac{p-1}{2}$ 个

勒让德符号

欧拉判别准则

$$(\frac{n}{p}) \equiv n^{\frac{p-1}{2}} (\bmod p)$$

若 n 是二次剩余,当且仅当 $n^{\frac{p-1}{2}} \equiv 1 \pmod{p}$

若 n 是非二次剩余,当且仅当 $n^{\frac{p-1}{2}} \equiv -1 \pmod{p}$

Cipolla

找到一个数 a 满足 a^2-n 是 **非二次剩余**,至于为什么要找满足非二次剩余的数,在下文会给出解释。 这里通过生成随机数再检验的方法来实现,由于非二次剩余的数量为 $\frac{p-1}{2}$,接近 $\frac{p}{2}$,所以期望约 2 次就可以找到这个数。

建立一个 "复数域",并不是实际意义上的复数域,而是根据复数域的概念建立的一个类似的域。 在复数中 $i^2 = -1$,这里定义 $i^2 = a^2 - n$,于是就可以将所有的数表达为A + Bi的形式,这里的 和 都是模意义下的数,类似复数中的实部和虚部。

在有了 i 和 a 后可以直接得到答案, $x^2 \equiv n \pmod{p}$ 的解为 $(a+i)^{\frac{p+1}{2}}$ 。

```
ll binpow_real(ll a, ll b, ll p) { //实部快速幂
      ll ans = 1;
     while (b) {
            if (b & 1) ans = ans * a % p;
            a = a * a % p;
            b >>= 1;
      }
      return ans % p;
ll binpow_imag(num a, ll b, ll p) { //虚部快速幂
      num ans = \{1, 0\};
     while (b) {
            if (b & 1) ans = mul(ans, a, p);
            a = mul(a, a, p);
            b >>= 1;
      }
      return ans.x % p;
ll cipolla(ll n, ll p) {
      n \% = p;
      if (p == 2) return n;
      if (binpow_real(n, (p - 1) / 2, p) == p - 1) return -1;
      11 a;
     while (1) { //生成随机数再检验找到满足非二次剩余的 a
            a = rand() \% p;
            w = ((a * a % p - n) % p + p) % p;
            if (binpow_real(w, (p - 1) / 2, p) == p - 1) break;
      }
      num x = \{a, 1\};
      return binpow_imag(x, (p + 1) / 2, p);
}
```

卡特兰数

卡特兰数 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012,...

$$C_n = \frac{1}{n+1} C_{2n}^n = C_{2n}^n - C_{2n}^{n-1}$$

$$C_n = \frac{1}{n+1} \sum_{i=0}^n (C_n^i)^2$$

$$C_n = \frac{4n-2}{n+1} C_{n-1} (C_0 = 1)$$

$$C_{n+1} = \sum_{i=0}^n C_i C_{n-i} (C_0 = 1)$$

超级卡特兰数 1, 1, 3, 11, 45, 197, 903, 4279, 20793, 103049,... (从第 0 项开始)

```
F_n * (n + 1) = (6 * n - 3) * F_{n-1} - (n - 2) * F_{n-2}
```

大施罗德数(OEIS A006318)1, 2, 6, 22, 90, 394, 1806, 8558, 41586, 206098,...

超级卡特兰数的两倍(除第一项)

快速幂

```
11 qpow(11 a, 11 b) {
    11 ans = 1;
    while (b) {
        if (b & 1) ans = (ans * a) % mod;
        a = (a * a) % mod;
        b >>= 1;
    }
    return ans;
}
```

龟速乘快速幂(快速幂爆 longlong

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
11 qmul(ll a, ll b, ll p) {
      11 \text{ res} = 0;
      while(b) {
             if(b \& 1) res = (res + a) \% p;
             a = (a + a) \% p;
             b >>= 1;
      return res;
11 qpow(ll x, ll n, ll p) {
      11 \text{ res} = 1;
      while(n) {
             if(n & 1) res = qmul(res, x, p);
             x = qmul(x, x, p);
             n >>= 1;
      return res % p; // 1 0 1
int main() {
      11 b, p, k;
      cin >> b >> p >> k;
      11 \text{ ans} = \text{qpow}(b, p, k);
```

```
printf("%lld^%lld mod %lld=%lld", b, p, k, ans);
    return 0;
}
```

莫比乌斯反演

莫比乌斯函数

对n进行因数分解:
$$n = P_1^{\alpha_1} P_2^{\alpha_2} ... P_k^{\alpha_k}$$
, 则 $\mu(n) = \begin{cases} 1, n = 1 \\ 0, \forall \alpha_i \geq 2 \\ \pm 1, (-1)^k \end{cases}$

n 的所有约数的莫比乌斯的和

$$S(n) = \sum_{d|n} \mu(d) = \begin{cases} 1, & n = 1 \\ 0, & else \end{cases}$$

反演

$$(一般不用)$$
1. 若 $F(n) = \sum_{d|n} f(d)$, 则 $f(n) = \sum_{d|n} \mu(d)F(\frac{n}{d})$

$$(\sqrt{2}. \, \overline{Z}F(n) = \sum_{n/d} f(d), \, \mathcal{D}f(n) = \sum_{n/d} \mu(\frac{d}{n})F(d)$$

构造F(n)和f(n)使 f(n)为目标,F(n)好求

1

求满足 $a \le x \le b, c \le y \le d$ 且 gcd(x, y) = k的 xy的对数

$$F(n) = gcd(x, y) = n$$
的倍数的xy的对数

$$f(n) = gcd(x, y) = n 的xy 的对数$$

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
const int N = 50010;
```

```
11 primes[N], mu[N], sum[N], cnt;
bool st[N];
void init() {
      mu[1] = 1;
      for(int i = 2; i < N; ++ i) {</pre>
            if(!st[i]) {
                   primes[cnt ++] = i;
                   mu[i] = -1;
            }
            for(int j = 0; primes[j] * i < N; ++ j) {</pre>
                   st[primes[j] * i] = 1;
                   if(i % primes[j] == 0) break;
                   mu[primes[j] * i] = -mu[i];
            }
      }
      for(int i = 1; i < N; ++ i) {</pre>
            sum[i] = sum[i - 1] + mu[i];
      }
ll g(ll n, ll x) {
      return n / (n / x);
11 f (int a, int b, int k) {
      a = a / k, b = b / k;
      11 \text{ res} = 0;
      11 n = min(a, b);
      for(11 1 = 1, r; 1 \le n; 1 = r + 1) {
            r = min(n, min(g(a, 1), g(b, 1)));
            res += (sum[r] - sum[l - 1]) * (a / 1) * (b / 1);
      }
      return res;
}
int main() {
      ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
      init();
      int T;
      cin >> T;
      while(T --) {
            int a, b, c, d, k;
            cin >> a >> b >> c >> d >> k;
```

```
cout << f(b, d, k) - f(a - 1, d, k) - f(b, c - 1, k)
                                     + f(a - 1, c - 1, k) << endl;
          }
          return 0;
}
 2
求\sum_{i=1}^{N}\sum_{j=1}^{M}d(ij)
// d(ij) = \sum_{x|i} \sum_{y|i} [(x, y) = 1]
F(n) = \sum_{i=1}^{N} \sum_{j=1}^{M} \sum_{x \mid i} \sum_{y \mid i} [n/(x, y)]
f(n) = \sum_{i=1}^{N} \sum_{j=1}^{M} \sum_{x \mid i} \sum_{x \mid j} [(x, y) = n]
F(n) = \sum_{i=1}^{N} \sum_{j=1}^{M} \sum_{x|i} \sum_{y|i} [n/(x,y)] = \sum_{x=1}^{N} \sum_{y=1}^{M} \lfloor \frac{N}{x} \rfloor \lfloor \frac{M}{y} \rfloor [n/(x,y)] = \sum_{i}^{N} \sum_{y=1}^{M} \lfloor \frac{N}{x'n} \rfloor \lfloor \frac{M}{y'n} \rfloor
 两次整数分块
#include <bits/stdc++.h>
 using namespace std;
 typedef long long 11;
 const int N = 50010;
int primes[N], cnt, mu[N], sum[N], h[N];
 bool st[N];
inline int g(int n, int x) {
         return n / (n / x);
void init() {
         mu[1] = 1;
          for(int i = 2; i < N; ++i) {</pre>
                   if(!st[i]){
                            primes[cnt++] = i;
                            mu[i] = -1;
                   for(int j = 0; primes[j] * i < N; ++j) {</pre>
                            st[primes[j] * i] = 1;
                            if(i % primes[j] == 0) break;
                            mu[primes[j] * i] = -mu[i];
```

```
}
      }
      for(int i = 1; i < N; ++ i) {</pre>
            sum[i] = sum[i - 1] + mu[i];
      }
      for(int i = 1; i < N; ++i) {</pre>
            for(int l = 1, r; l \leftarrow i; l = r + 1) {
                   r = min(i, g(i, 1));
                   h[i] += (r - l + 1) * (i / l);
            }
      }
}
int main() {
      //ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
      init();
      int T;
      scanf("%d", &T);
      while(T--) {
            int n, m;
            scanf("%d %d", &n, &m);
            11 \text{ res} = 0;
            int k = min(n, m);
            for(int l = 1, r; l <= k; l = r + 1) {
                   r = min(k, min(g(n, 1), g(m, 1)));
                   res += (11)(sum[r] - sum[1 - 1]) * h[n / 1] * h[m /
1];
          printf("%lld\n", res);
      }
      return 0;
}
```

博弈

SG 定理:

mex(minimal excludant)运算,表示最小的不属于这个集合的非负整数。例如 $mex\{0,1,2,4\}=3$ 、 $mex\{2,3,5\}=0$ 、 $mex\{\}=0$ 。

Sprague-Grundy 定理(SG 定理):游戏和的 SG 函数等于各个游戏 SG 函数的

Nim 和。这样就可以将每一个子游戏分而治之,从而简化了问题。而 Bouton 定理 就是 Sprague-Grundy 定理在 Nim 游戏中的直接应用,因为单堆的 Nim 游戏 SG 函数满足 SG(x) = x。

Nimk:

普通的 NIM 游戏是在 n 堆石子中每次选一堆,取任意个石子,而 NIMK 游戏是在 n 堆石子中每次选择 k 堆, 1 <= k <= n,从这 k 堆中每堆里都取出任意数目的石子,取的石子数可以不同,其他规则相同。

对于普通的 NIM 游戏,我们采取的是对每堆的 SG 值进行异或,异或其实就是对每一个 SG 值二进制位上的数求和然后模 2,比如说 3^5 就是 011+101=112,然后对每一位都模 2 就变成了 110,所以 3^5=6。而 NIMK 游戏和 NIM 游戏的区别就在于模的不是 2,如果是取 k 堆,就模 k+1,所以取 1 堆的普通 NIM 游戏是模 2。当 k=2 时,3^5 \rightarrow 011+101=112,对每一位都模 3 之后三位二进制位上对应的数仍然是 1,1,2。那么当且仅当每一位二进制位上的数都是 0 的时候,先手必败,否则先手必胜。

anti_nim

描述

和最普通的 Nim 游戏相同,不过是取走最后一个石子的人输。

先手必胜条件

以下两个条件满足其一即可:

- 1. 所有堆的石子个数=1,且异或和=0(其实这里就是有偶数堆的意思)。
- 2. 至少存在一堆石子个数>1, 且异或和≠0。

高精度 GCD

```
#include <bits/stdc++.h>
using namespace std;
string add(string a, string b) {
    const int L = 1e5;
    string ans;
    int na[L] = {0}, nb[L] = {0};
    int la = a.size(), lb = b.size();
    for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0';
    for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0';
    int lmax = la > lb ? la : lb;
    for (int i = 0; i < lmax; i++)
        na[i] += nb[i], na[i + 1] += na[i] / 10, na[i] %= 10;
    if (na[lmax]) lmax++;
    for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
```

```
return ans;
}
string mul(string a, string b) {
   const int L = 1e5;
   string s;
   int na[L], nb[L], nc[L],
       La = a.size(), Lb = b.size(); // na 存储被乘数, nb 存储乘数, nc 存
储积
   fill(na, na + L, \emptyset);
   fill(nb, nb + L, \emptyset);
   fill(nc, nc + L, 0); //将 na, nb, nc 都置为0
   for (int i = La - 1; i >= 0; i--)
       na[La - i] =
           a[i] - '0'; //将字符串表示的大整形数转成 i 整形数组表示的大整形
数
   for (int i = Lb - 1; i >= 0; i--) nb[Lb - i] = b[i] - '0';
   for (int i = 1; i <= La; i++)</pre>
       for (int j = 1; j <= Lb; j++)</pre>
           nc[i + j - 1] +=
               na[i] *
               nb[j]; // a 的第 i 位乘以 b 的第 j 位为积的第 i+j-1 位 ( 先不考
虑进位)
   for (int i = 1; i <= La + Lb; i++)
       nc[i + 1] += nc[i] / 10, nc[i] %= 10; //统一处理进位
   if (nc[La + Lb]) s += nc[La + Lb] + '0'; //判断第 i+j 位上的数字是不
是0
   for (int i = La + Lb - 1; i >= 1; i--)
       s += nc[i] + '0'; //将整形数组转成字符串
   return s;
int sub(int *a, int *b, int La, int Lb) {
   if (La < Lb) return -1; //如果 a 小于 b,则返回-1
   if (La == Lb) {
       for (int i = La - 1; i >= 0; i--)
           if (a[i] > b[i])
               break;
           else if (a[i] < b[i])
               return -1; //如果 a 小于 b ,则返回-1
   for (int i = 0; i < La; i++) //高精度减法
   {
       a[i] -= b[i];
       if (a[i] < 0) a[i] += 10, a[i + 1]--;</pre>
   for (int i = La - 1; i >= 0; i--)
       if (a[i]) return i + 1; //返回差的位数
                               //返回差的位数
   return 0;
}
```

```
string div(string n1, string n2,
          int nn) // n1,n2 是字符串表示的被除数,除数,nn 是选择返回商还是
余数
{
   const int L = 1e5;
   string s, v; // s 存商, v 存余数
   int a[L], b[L], r[L],
       La = n1.size(), Lb = n2.size(), i,
       tp = La; // a, b 是整形数组表示被除数,除数,tp 保存被除数的长度
   fill(a, a + L, 0);
   fill(b, b + L, 0);
   fill(r, r + L, 0); //数组元素都置为0
   for (i = La - 1; i >= 0; i--) a[La - 1 - i] = n1[i] - '0';
   for (i = Lb - 1; i >= 0; i--) b[Lb - 1 - i] = n2[i] - '0';
   if (La < Lb || (La == Lb && n1 < n2)) {</pre>
       // cout<<0<<endl;</pre>
       return n1;
                    //如果 a<b, 则商为 0, 余数为被除数
   int t = La - Lb; //除被数和除数的位数之差
   for (int i = La - 1; i >= 0; i--) //将除数扩大10^t 倍
       if (i >= t)
           b[i] = b[i - t];
       else
           b[i] = 0;
   Lb = La;
   for (int j = 0; j <= t; j++) {</pre>
       int temp;
       while ((temp = sub(a, b + j, La, Lb - j)) >=
              0) //如果被除数比除数大继续减
       {
           La = temp;
           r[t - j]++;
       }
   for (i = 0; i < L - 10; i++)
       r[i + 1] += r[i] / 10, r[i] %= 10; //统一处理进位
   while (!r[i]) i--; //将整形数组表示的商转化成字符串表示的
   while (i >= 0) s += r[i--] + '0';
   // cout<<s<<endl;</pre>
   i = tp;
   while (!a[i]) i--; //将整形数组表示的余数转化成字符串表示的</span>
   while (i >= 0) v += a[i--] + '0';
   if (v.empty()) v = "0";
   // cout<<v<<endl;</pre>
   if (nn == 1) return s;
   if (nn == 2) return v;
bool judge(string s) //判断 s 是否为全 0 串
```

```
{
    for (int i = 0; i < s.size(); i++)</pre>
        if (s[i] != '0') return false;
    return true;
string gcd(string a, string b) //求最大公约数
    string t;
   while (!judge(b)) //如果余数不为0,继续除
                           //保存被除数的值
       t = a;
       a = b;
                           //用除数替换被除数
       b = div(t, b, 2); //用余数替换除数
    return a;
//o(无法估计)
高精度乘法(FFT)
#include <bits/stdc++.h>
using namespace std;
#define L(x) (1 << (x))
const double PI = acos(-1.0);
const int Maxn = 133015;
double ax[Maxn], ay[Maxn], bx[Maxn], by[Maxn];
char sa[Maxn / 2], sb[Maxn / 2];
int sum[Maxn];
int x1[Maxn], x2[Maxn];
int revv(int x, int bits) {
    int ret = 0;
    for (int i = 0; i < bits; i++) {</pre>
        ret <<= 1;
        ret |= x \& 1;
       x >>= 1;
    }
    return ret;
}
void fft(double* a, double* b, int n, bool rev) {
    int bits = 0;
   while (1 << bits < n) ++bits;
    for (int i = 0; i < n; i++) {</pre>
        int j = revv(i, bits);
        if (i < j) swap(a[i], a[j]), swap(b[i], b[j]);</pre>
    for (int len = 2; len <= n; len <<= 1) {</pre>
        int half = len >> 1;
        double wmx = cos(2 * PI / len), wmy = sin(2 * PI / len);
        if (rev) wmy = -wmy;
        for (int i = 0; i < n; i += len) {</pre>
```

```
double wx = 1, wy = 0;
            for (int j = 0; j < half; j++) {</pre>
                double cx = a[i + j], cy = b[i + j];
                double dx = a[i + j + half], dy = b[i + j + half];
                double ex = dx * wx - dy * wy, ey = dx * wy + dy * wx;
                a[i + j] = cx + ex, b[i + j] = cy + ey;
                a[i + j + half] = cx - ex, b[i + j + half] = cy - ey;
                double wnx = wx * wmx - wy * wmy, wny = wx * wmy + wy *
wmx;
                wx = wnx, wy = wny;
            }
        }
    if (rev) {
        for (int i = 0; i < n; i++) a[i] /= n, b[i] /= n;</pre>
    }
}
int solve(int a[], int na, int b[], int nb, int ans[]) {
    int len = max(na, nb), ln;
    for (ln = 0; L(ln) < len; ++ln)
    len = L(++ln);
    for (int i = 0; i < len; ++i) {</pre>
        if (i >= na)
            ax[i] = 0, ay[i] = 0;
        else
            ax[i] = a[i], ay[i] = 0;
    fft(ax, ay, len, 0);
    for (int i = 0; i < len; ++i) {</pre>
        if (i >= nb)
            bx[i] = 0, by[i] = 0;
        else
            bx[i] = b[i], by[i] = 0;
    fft(bx, by, len, ∅);
    for (int i = 0; i < len; ++i) {</pre>
        double cx = ax[i] * bx[i] - ay[i] * by[i];
        double cy = ax[i] * by[i] + ay[i] * bx[i];
        ax[i] = cx, ay[i] = cy;
    }
    fft(ax, ay, len, 1);
    for (int i = 0; i < len; ++i) ans[i] = (int)(ax[i] + 0.5);
    return len;
string mul(string sa, string sb) {
    int 11, 12, 1;
    int i;
    string ans;
    memset(sum, 0, sizeof(sum));
```

```
11 = sa.size();
   12 = sb.size();
   for (i = 0; i < l1; i++) x1[i] = sa[l1 - i - 1] - '0';</pre>
   for (i = 0; i < 12; i++) \times 2[i] = sb[12 - i - 1] - '0';
   1 = solve(x1, 11, x2, 12, sum);
   for (i = 0; i < l || sum[i] >= 10; i++) // 进位
       sum[i + 1] += sum[i] / 10;
       sum[i] %= 10;
   }
   l = i;
   return ans:
int main() {
   cin.sync_with_stdio(false);
   string a, b;
   while (cin >> a >> b) cout << mul(a, b) << endl;</pre>
   return 0;
//o(nlogn)
高精度乘法(乘单精度
#include <bits/stdc++.h>
using namespace std;
string mul(string a, int b) //高精度 a 乘单精度 b
   const int L = 100005;
   int na[L];
   string ans;
   int La = a.size();
   fill(na, na + L, ∅);
   for (int i = La - 1; i >= 0; i--) na[La - i - 1] = a[i] - '0';
   int w = 0;
   for (int i = 0; i < La; i++)</pre>
       na[i] = na[i] * b + w, w = na[i] / 10, na[i] = na[i] % 10;
   while (w) na[La++] = w % 10, w /= 10;
   La--;
   while (La >= 0) ans += na[La--] + '0';
   return ans;
}
//o(n)
高精度乘法(朴素)
#include <bits/stdc++.h>
using namespace std;
string mul(string a, string b) //高精度乘法a,b,均为非负整数
{
```

```
const int L = 1e5;
   string s;
   int na[L], nb[L], nc[L],
       La = a.size(), Lb = b.size(); // na 存储被乘数, nb 存储乘数, nc 存
储积
   fill(na, na + L, \emptyset);
   fill(nb, nb + L, \emptyset);
   fill(nc, nc + L, 0); //将 na, nb, nc 都置为0
   for (int i = La - 1; i >= 0; i--)
       na[La - i] =
           a[i] - '0'; //将字符串表示的大整形数转成 i 整形数组表示的大整形
数
   for (int i = Lb - 1; i \ge 0; i--) nb[Lb - i] = b[i] - '0';
   for (int i = 1; i <= La; i++)</pre>
       for (int j = 1; j <= Lb; j++)</pre>
           nc[i + j - 1] +=
               na[i] *
               nb[j]; // a 的第 i 位乘以 b 的第 j 位为积的第 i+j-1 位 (先不考
虑进位)
   for (int i = 1; i <= La + Lb; i++)
       nc[i + 1] += nc[i] / 10, nc[i] %= 10; //统一处理进位
   if (nc[La + Lb]) s += nc[La + Lb] + '0'; //判断第i+j 位上的数字是不
是0
   for (int i = La + Lb - 1; i >= 1; i--)
       s += nc[i] + '0'; //将整形数组转成字符串
   return s;
}
//o(n^2)
高精度除法 (除单精度)
#include <bits/stdc++.h>
using namespace std;
string div(string a, int b) //高精度a 除以单精度b
   string r, ans;
   int d = 0;
   if (a == "0") return a; //特判
   for (int i = 0; i < a.size(); i++) {</pre>
       r += (d * 10 + a[i] - '0') / b + '0'; //求出商
       d = (d * 10 + (a[i] - '0')) \% b;
                                             //求出余数
   }
   int p = 0;
   for (int i = 0; i < r.size(); i++)</pre>
       if (r[i] != '0') {
           p = i;
           break;
   return r.substr(p);
```

```
}
//o(n)
```

```
高精度除法 (除高精度)
#include <bits/stdc++.h>
using namespace std;
int sub(int *a, int *b, int La, int Lb) {
   if (La < Lb) return -1; //如果 a 小于 b ,则返回-1
   if (La == Lb) {
       for (int i = La - 1; i >= 0; i--)
           if (a[i] > b[i])
               break;
           else if (a[i] < b[i])
               return -1; //如果 a 小于 b ,则返回-1
   }
   for (int i = 0; i < La; i++) //高精度减法
   {
       a[i] -= b[i];
       if (a[i] < 0) a[i] += 10, a[i + 1]--;</pre>
   for (int i = La - 1; i >= 0; i--)
       if (a[i]) return i + 1; //返回差的位数
                               //返回差的位数
   return 0;
string div(string n1, string n2, int nn)
// n1,n2 是字符串表示的被除数,除数,nn 是选择返回商还是余数
{
   const int L = 1e5;
   string s, v; // s 存商, v 存余数
   int a[L], b[L], r[L], La = n1.size(), Lb = n2.size(), i, tp = La;
   // a,b 是整形数组表示被除数,除数,tp 保存被除数的长度
   fill(a, a + L, 0);
   fill(b, b + L, 0);
   fill(r, r + L, 0); //数组元素都置为0
   for (i = La - 1; i >= 0; i--) a[La - 1 - i] = n1[i] - '0';
   for (i = Lb - 1; i >= 0; i--) b[Lb - 1 - i] = n2[i] - '0';
   if (La < Lb | (La == Lb && n1 < n2)) {
       // cout<<0<<endl;</pre>
       return n1;
                    //如果 a < b , 则商为 0 , 余数为被除数
   int t = La - Lb; //除被数和除数的位数之差
   for (int i = La - 1; i >= 0; i--) //将除数扩大10<sup>^</sup>t 倍
       if (i >= t)
           b[i] = b[i - t];
       else
           b[i] = 0;
   Lb = La;
```

```
for (int j = 0; j <= t; j++) {</pre>
       int temp;
       while ((temp = sub(a, b + j, La, Lb - j)) >=
              0) //如果被除数比除数大继续减
       {
           La = temp;
           r[t - j]++;
       }
   for (i = 0; i < L - 10; i++)
       r[i + 1] += r[i] / 10, r[i] %= 10; //统一处理进位
   while (!r[i]) i--; //将整形数组表示的商转化成字符串表示的
   while (i >= 0) s += r[i--] + '0';
   // cout<<s<<endl:</pre>
   i = tp;
   while (!a[i]) i--; //将整形数组表示的余数转化成字符串表示的</span>
   while (i >= 0) v += a[i--] + '0';
   if (v.empty()) v = "0";
   // cout<<v<<endl;</pre>
   if (nn == 1) return s; //返回商
   if (nn == 2) return v; //返回余数
}
//o(n^2)
```

高精度加法

```
#include <bits/stdc++.h>
using namespace std;
string add(string a, string b) // 只限两个非负整数相加
{
   const int L = 1e5;
   string ans;
   int na[L] = {0}, nb[L] = {0};
   int la = a.size(), lb = b.size();
   for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0';</pre>
   for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0';
   int lmax = la > lb ? la : lb;
   for (int i = 0; i < lmax; i++)</pre>
        na[i] += nb[i], na[i + 1] += na[i] / 10, na[i] %= 10;
   if (na[lmax]) lmax++;
   for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
   return ans;
//o(n)
```

```
高精度减法
```

```
#include <bits/stdc++.h>
using namespace std;
string sub(string a, string b) // 只限大的非负整数减小的非负整数
{
    const int L = 1e5;
    string ans;
    int na[L] = \{0\}, nb[L] = \{0\};
    int la = a.size(), lb = b.size();
    for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0';</pre>
    for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0';
    int lmax = la > lb ? la : lb;
    for (int i = 0; i < lmax; i++) {</pre>
        na[i] -= nb[i];
        if (na[i] < 0) na[i] += 10, na[i + 1]--;</pre>
   while (!na[--lmax] && lmax > 0)
       ;
    lmax++;
    for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
    return ans;
//o(n)
高精度阶乘
#include <bits/stdc++.h>
using namespace std;
string fac(int n) {
    const int L = 100005;
    int a[L];
    string ans;
    if (n == 0) return "1";
    fill(a, a + L, 0);
    int s = 0, m = n;
   while (m) a[++s] = m \% 10, m /= 10;
    for (int i = n - 1; i >= 2; i--) {
        int w = 0;
        for (int j = 1; j <= s; j++)
            a[j] = a[j] * i + w, w = a[j] / 10, a[j] = a[j] % 10;
        while (w) a[++s] = w \% 10, w /= 10;
   while (!a[s]) s--;
   while (s >= 1) ans += a[s--] + '0';
    return ans;
//o(n^2)
```

```
高精度进制转换
```

```
#include <bits/stdc++.h>
using namespace std;
//将字符串表示的10 进制大整数转换为m 进制的大整数
//并返回 m 进制大整数的字符串
bool judge(string s) //判断串是否为全零串
   for (int i = 0; i < s.size(); i++)</pre>
       if (s[i] != '0') return 1;
   return 0;
}
string solve(
   string s, int n,
   int m) // n 进制转 m 进制只限 0-9 进制, 若涉及带字母的进制, 稍作修改即可
{
   string r, ans;
   int d = 0;
   if (!judge(s)) return "0"; //特判
                           //被除数不为0则继续
   while (judge(s))
   {
       for (int i = 0; i < s.size(); i++) {</pre>
          r += (d * n + s[i] - '0') / m + '0'; //求出商
          d = (d * n + (s[i] - '0')) % m;
                                         //求出余数
       }
                     //把商赋给下一次的被除数
       s = r;
                     //把商清空
       r = "";
      ans += d + '0'; //加上进制转换后数字
                     //清空余数
       d = 0;
   reverse(ans.begin(), ans.end()); //倒置下
   return ans;
//o(n^2)
```

高精度幂

```
#include <bits/stdc++.h>
#define L(x) (1 << (x))
using namespace std;
const double PI = acos(-1.0);
const int Maxn = 133015;
double ax[Maxn], ay[Maxn], bx[Maxn], by[Maxn];
char sa[Maxn / 2], sb[Maxn / 2];
int sum[Maxn];
int x1[Maxn], x2[Maxn];
int revv(int x, int bits) {
   int ret = 0;</pre>
```

```
for (int i = 0; i < bits; i++) {</pre>
        ret <<= 1;
        ret |= x \& 1;
        x >>= 1;
    return ret;
}
void fft(double* a, double* b, int n, bool rev) {
    int bits = 0;
    while (1 << bits < n) ++bits;</pre>
    for (int i = 0; i < n; i++) {
        int j = revv(i, bits);
        if (i < j) swap(a[i], a[j]), swap(b[i], b[j]);</pre>
    }
    for (int len = 2; len <= n; len <<= 1) {</pre>
        int half = len >> 1;
        double wmx = cos(2 * PI / len), wmy = sin(2 * PI / len);
        if (rev) wmy = -wmy;
        for (int i = 0; i < n; i += len) {</pre>
            double wx = 1, wy = 0;
            for (int j = 0; j < half; j++) {</pre>
                 double cx = a[i + j], cy = b[i + j];
                 double dx = a[i + j + half], dy = b[i + j + half];
                 double ex = dx * wx - dy * wy, ey = dx * wy + dy * wx;
                 a[i + j] = cx + ex, b[i + j] = cy + ey;
                 a[i + j + half] = cx - ex, b[i + j + half] = cy - ey;
                 double wnx = wx * wmx - wy * wmy, wny = wx * wmy + wy *
wmx;
                wx = wnx, wy = wny;
            }
        }
    if (rev) {
        for (int i = 0; i < n; i++) a[i] /= n, b[i] /= n;</pre>
    }
int solve(int a[], int na, int b[], int nb, int ans[]) {
    int len = max(na, nb), ln;
    for (ln = 0; L(ln) < len; ++ln)</pre>
    len = L(++ln);
    for (int i = 0; i < len; ++i) {</pre>
        if (i >= na)
            ax[i] = 0, ay[i] = 0;
        else
            ax[i] = a[i], ay[i] = 0;
    fft(ax, ay, len, ∅);
    for (int i = 0; i < len; ++i) {</pre>
        if (i >= nb)
```

```
bx[i] = 0, by[i] = 0;
        else
            bx[i] = b[i], by[i] = 0;
   fft(bx, by, len, ∅);
    for (int i = 0; i < len; ++i) {</pre>
        double cx = ax[i] * bx[i] - ay[i] * by[i];
        double cy = ax[i] * by[i] + ay[i] * bx[i];
        ax[i] = cx, ay[i] = cy;
   fft(ax, ay, len, 1);
    for (int i = 0; i < len; ++i) ans[i] = (int)(ax[i] + 0.5);
    return len;
string mul(string sa, string sb) {
    int 11, 12, 1;
    int i;
    string ans;
    memset(sum, 0, sizeof(sum));
    11 = sa.size();
    12 = sb.size();
   for (i = 0; i < 11; i++) \times 1[i] = sa[11 - i - 1] - '0';
    for (i = 0; i < 12; i++) \times 2[i] = sb[12 - i - 1] - '0';
    1 = solve(x1, 11, x2, 12, sum);
    for (i = 0; i < l || sum[i] >= 10; i++) // 进位
        sum[i + 1] += sum[i] / 10;
       sum[i] %= 10;
    }
    l = i;
                                                  // 检索最高位
   while (sum[1] <= 0 && 1 > 0) 1--;
    for (i = l; i >= 0; i--) ans += sum[i] + '0'; // 倒序输出
    return ans;
}
string Pow(string a, int n) {
    if (n == 1) return a;
    if (n & 1) return mul(Pow(a, n - 1), a);
    string ans = Pow(a, n / 2);
    return mul(ans, ans);
//o(nlognlogm)
高精度平方根
#include <bits/stdc++.h>
using namespace std;
const int L = 2015;
string add(string a, string b) // 只限两个非负整数相加
{
    string ans;
```

```
int na[L] = \{0\}, nb[L] = \{0\};
    int la = a.size(), lb = b.size();
    for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0';</pre>
    for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0';
    int lmax = la > lb ? la : lb;
    for (int i = 0; i < lmax; i++)</pre>
        na[i] += nb[i], na[i + 1] += na[i] / 10, na[i] %= 10;
    if (na[lmax]) lmax++;
    for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
    return ans;
}
string sub(string a, string b) // 只限大的非负整数减小的非负整数
{
    string ans;
    int na[L] = \{0\}, nb[L] = \{0\};
    int la = a.size(), lb = b.size();
    for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0';
    for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0';
    int lmax = la > lb ? la : lb;
    for (int i = 0; i < lmax; i++) {</pre>
        na[i] -= nb[i];
        if (na[i] < 0) na[i] += 10, na[i + 1]--;</pre>
   while (!na[--lmax] && lmax > 0)
    lmax++;
    for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
    return ans;
}
string mul(string a, string b) //高精度乘法a,b,均为非负整数
{
    string s;
    int na[L], nb[L], nc[L],
        La = a.size(), Lb = b.size(); // na 存储被乘数, nb 存储乘数, nc 存
储积
   fill(na, na + L, 0);
    fill(nb, nb + L, \emptyset);
    fill(nc, nc + L, 0); //将 na, nb, nc 都置为0
    for (int i = La - 1; i >= 0; i--)
        na[La - i] =
            a[i] - '0': //将字符串表示的大整形数转成 i 整形数组表示的大整形
数
    for (int i = Lb - 1; i >= 0; i--) nb[Lb - i] = b[i] - '0';
    for (int i = 1; i <= La; i++)
        for (int j = 1; j <= Lb; j++)</pre>
            nc[i + j - 1] +=
                na[i] *
                nb[j]; // a 的第 i 位乘以 b 的第 j 位为积的第 i+j-1 位 ( 先不考
虑进位)
```

```
for (int i = 1; i <= La + Lb; i++)</pre>
       nc[i + 1] += nc[i] / 10, nc[i] %= 10; //统一处理进位
   if (nc[La + Lb]) s += nc[La + Lb] + '0'; //判断第i+j 位上的数字是不
是0
   for (int i = La + Lb - 1; i >= 1; i--)
       s += nc[i] + '0'; //将整形数组转成字符串
   return s;
int sub(int *a, int *b, int La, int Lb) {
   if (La < Lb) return -1; //如果 a 小于 b ,则返回-1
   if (La == Lb) {
       for (int i = La - 1; i >= 0; i--)
           if (a[i] > b[i])
               break;
           else if (a[i] < b[i])
               return -1; //如果 a 小于 b ,则返回-1
   for (int i = 0; i < La; i++) //高精度减法
   {
       a[i] -= b[i];
       if (a[i] < 0) a[i] += 10, a[i + 1]--;</pre>
   for (int i = La - 1; i >= 0; i--)
       if (a[i]) return i + 1; //返回差的位数
                               //返回差的位数
   return 0;
}
string div(string n1, string n2,
          int nn) // n1, n2 是字符串表示的被除数,除数, nn 是选择返回商还是
余数
{
   string s, v; // s 存商,v 存余数
   int a[L], b[L], r[L],
       La = n1.size(), Lb = n2.size(), i,
       tp = La; // a, b 是整形数组表示被除数,除数,tp 保存被除数的长度
   fill(a, a + L, 0);
   fill(b, b + L, \emptyset);
   fill(r, r + L, 0); //数组元素都置为0
   for (i = La - 1; i >= 0; i--) a[La - 1 - i] = n1[i] - '0';
   for (i = Lb - 1; i >= 0; i--) b[Lb - 1 - i] = n2[i] - '0';
   if (La < Lb || (La == Lb && n1 < n2)) {
       // cout<<0<<endl;</pre>
       return n1;
                    //如果 a<b,则商为 0,余数为被除数
   int t = La - Lb; //除被数和除数的位数之差
   for (int i = La - 1; i >= 0; i--) //将除数扩大10<sup>^</sup>t 倍
       if (i >= t)
           b[i] = b[i - t];
       else
```

```
b[i] = 0;
   Lb = La;
   for (int j = 0; j <= t; j++) {</pre>
       int temp;
       while ((temp = sub(a, b + j, La, Lb - j)) >=
              0) //如果被除数比除数大继续减
        {
           La = temp;
           r[t - j]++;
        }
   for (i = 0; i < L - 10; i++)
        r[i + 1] += r[i] / 10, r[i] %= 10; //统一处理进位
   while (!r[i]) i--; //将整形数组表示的商转化成字符串表示的
   while (i >= 0) s += r[i--] + '0';
   // cout<<s<<endl;</pre>
   i = tp;
   while (!a[i]) i--; //将整形数组表示的余数转化成字符串表示的</span>
   while (i >= 0) v += a[i--] + '0';
   if (v.empty()) v = "0";
   // cout<<v<<endl;</pre>
   if (nn == 1) return s;
   if (nn == 2) return v;
}
bool cmp(string a, string b) {
   if (a.size() < b.size()) return 1; // a 小于等于 b 返回真
   if (a.size() == b.size() && a <= b) return 1;</pre>
   return 0;
}
string DeletePreZero(string s) {
   int i;
   for (i = 0; i < s.size(); i++)</pre>
       if (s[i] != '0') break;
   return s.substr(i);
string BigInterSqrt(string n) {
   n = DeletePreZero(n);
   string l = "1", r = n, mid, ans;
   while (cmp(1, r)) {
       mid = div(add(1, r), "2", 1);
        if (cmp(mul(mid, mid), n))
           ans = mid, l = add(mid, "1");
       else
            r = sub(mid, "1");
   return ans;
// o(n^3)
```

```
高精度取模 (对单精度)
#include <bits/stdc++.h>
using namespace std;
int mod(string a, int b)//高精度 a 除以单精度 b
{
    int d=0;
    for(int i=0;i<a.size();i++) d=(d*10+(a[i]-'0'))%b;//求出余数
    return d;
}
//o(n)
欧拉筛
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const int N = 1000005;
int phi[N], prime[N], cnt;
bool st[N];
void get eulers() {
    phi[1] = 1;
    for (int i = 2; i < N; i++) {</pre>
        if (!st[i]) {
            prime[cnt++] = i;
            phi[i] = i - 1;
        for (int j = 0; prime[j] * i < N; j++) {</pre>
            st[prime[j] * i] = 1;
            if (i % prime[j] == 0) {
                phi[prime[j] * i] = phi[i] * prime[j];
                break;
            }
            phi[prime[j] * i] = phi[i] * (prime[j] - 1);
        }
    }
int main() {
    get_eulers();
    11 n;
    cin >> n;
    ll ans = 0;
    for (int i = 1; i <= n; i++) ans += phi[i];</pre>
    cout << ans;</pre>
}
组合数(逆元线性递推
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 \mod = 1e9 + 7;
```

```
const 11 \text{ maxn} = 3e4 + 5;
11 inv[maxn], fac[maxn];
11 qpow(ll a, ll b) {
    ll ans = 1;
    while (b) {
        if (b & 1) ans = (ans * a) % mod;
        a = (a * a) % mod;
        b >>= 1;
    }
    return ans;
11 c(ll n, ll m) {
    if (n < 0 || m < 0 || n < m) return 0;
    return fac[n] * inv[n - m] % mod * inv[m] % mod;
}
void init() {
    fac[0] = 1;
    for (int i = 1; i < maxn; i++) {</pre>
        fac[i] = fac[i - 1] * i % mod;
    }
    inv[maxn - 1] = qpow(fac[maxn - 1], mod - 2);
    for (ll i = maxn - 2; i >= 0; i--) {
        inv[i] = (inv[i + 1] * (i + 1)) % mod;
    }
}
中国剩余定理
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const int maxn = 20;
11 A[maxn], B[maxn];
11 exgcd(ll a, ll b, ll & x, ll & y) {
      if(b == 0) {
            x = 1, y = 0;
            return a;
      }
      11 d = exgcd(b, a % b, y, x);
      y -= (a / b) * x;
      return d;
int main() {
      int n;
      cin >> n;
      11 M = 111;
```

```
for(int i = 0; i < n; ++ i) {</pre>
            cin >> A[i] >> B[i];
            M = M * A[i];
      }
     11 ans = 0;
     11 x, y;
      for(int i = 0; i < n; ++ i) {</pre>
            11 Mi = M / A[i];
            exgcd(Mi, A[i], x, y);
            ans += B[i] * Mi * x;
      }
      cout << (ans % M + M) % M;
}
图论
有源汇上下界最大小流
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
struct Edge {
    11 from, to, cap, flow, mn;
    Edge(ll a, ll b, ll c, ll d, ll e) : from(a), to(b), cap(c), flow
(d), mn(e) {}
};
11 n, m;
struct Dinic {
    static const ll maxn = 50010; // 点的大小,记得改
    static const 11 inf = 0x3f3f3f3f3f3f3f3f3f;
    11 N, M, S, T;
    vector<Edge> edges;
    vector<ll> G[maxn];
    bool vis[maxn];
    11 d[maxn];
    11 cur[maxn];
    void AddEdge(ll from, ll to, ll cap, ll c) {
        edges.push back(Edge(from, to, cap, 0, c));
        edges.push_back(Edge(to, from, 0, 0, c));
       M = edges.size();
        G[from].push_back(M - 2);
        G[to].push_back(M - 1);
    }
```

```
bool BFS() {
    memset(vis, 0, sizeof(vis));
    queue<11> Q;
    Q.push(S);
    d[S] = 0;
    vis[S] = 1;
    while (!Q.empty()) {
        11 \times = Q.front();
        Q.pop();
        for (11 i = 0; i < G[x].size(); i++) {</pre>
            Edge& e = edges[G[x][i]];
            if (!vis[e.to] && e.cap > e.flow) {
                vis[e.to] = 1;
                d[e.to] = d[x] + 1;
                Q.push(e.to);
            }
        }
    return vis[T];
11 DFS(11 x, 11 a) {
    if (x == T || a == 0) return a;
    11 flow = 0, f;
    for (l1& i = cur[x]; i < G[x].size(); i++) {</pre>
        Edge& e = edges[G[x][i]];
        if (d[x] + 1 == d[e.to] &&
            (f = DFS(e.to, min(a, e.cap - e.flow))) > 0) {
            e.flow += f;
            edges[G[x][i] ^ 1].flow -= f;
            flow += f;
            a -= f;
            if (a == 0) break;
        }
    }
    return flow;
void deleteEdge(ll u, ll v) {
    11 siz = edges.size();
    for(ll i = 0; i < siz; ++ i) {</pre>
        if(edges[i].from == u && edges[i].to == v) {
            edges[i].cap = edges[i].flow = 0;
            edges[i ^1].cap = edges[i ^1].flow = 0;
            break;
        }
    }
11 getValue() {
    return edges[2 * m].flow;
11 Maxflow(11 S, 11 T) {
```

```
this->S = S, this->T = T;
        11 flow = 0;
       while (BFS()) {
            memset(cur, 0, sizeof(cur));
            flow += DFS(S, inf);
        return flow;
    }
} MF;
int main() {
    11 s, t;
    cin >> n >> m >> s >> t;
 // n 个点, m 条边, 给的源点汇点
    ll mp[50010] = {0}; // 点的大小,记得改
    for(ll i = 1; i <= m; ++ i) {
        11 a, b, c, d; // 从 a 到 b 有一条下界 c 上界 d 的边
        cin >> a >> b >> c >> d;
       mp[b] += c;
       mp[a] -= c;
       MF.AddEdge(a, b, d - c, c);
    }
   MF.AddEdge(t, s, 1e18, 0); //
    11 tot = 0;
    for(ll i = 1; i <= n; ++ i) {</pre>
        if(mp[i] > 0) {
            tot += mp[i];
            MF.AddEdge(0, i , mp[i], 0);
        }
        else {
            MF.AddEdge(i, n + 1, -mp[i], \emptyset);
        }
    if( MF.Maxflow(0, n + 1) != tot) {
        cout << "No Solution" << endl;</pre>
    }
    else {
        ll res = MF.getValue(); // 从t 到s 边的流量
       MF.deleteEdge(t, s);
     //cout << res + MF.Maxflow(s, t) << endl; // 最大流
        cout << res - MF.Maxflow(t, s) << endl; // 最小流
    return 0;
}
```

树链剖分

```
11 fa[N], son[N], dep[N], siz[N], dfn[N], rnk[N], top[N];
11 dfscnt;
```

```
vector<ll> g[N];
ll tree[N << 1];
11 lazy[N << 1];</pre>
void dfs1(ll u, ll f, ll d) {
    son[u] = -1;
    siz[u] = 1;
    fa[u] = f;
    dep[u] = d;
    for (auto v:g[u]) {
        if (v == f) continue;
        dfs1(v, u, d + 1);
        siz[u] += siz[v];
        if (son[u] == -1 \mid | siz[v] > siz[son[u]]) son[u] = v;
    }
}
void dfs2(11 u, 11 t) {
    dfn[u] = ++dfscnt;
    rnk[dfscnt] = u;
    top[u] = t;
    if (son[u] == -1) return;
    dfs2(son[u], t);
    for (auto v:g[u]) {
        if (v == son[u] || v == fa[u]) continue;
        dfs2(v, v);
    }
}
11 lca(ll a, ll b) {
    while (top[a] != top[b]) {
        if (dep[top[a]] < dep[top[b]]) swap(a, b);</pre>
        a = fa[top[a]];
    return dep[a] < dep[b] ? a : b;</pre>
}
void init() {
    for (ll i = 0; i < N; i++) g[i].clear();</pre>
    for (11 i = 0; i < (N << 1); i++) {
        tree[i] = 0;
        lazy[i] = 0;
    }
    dfscnt = 0;
}
void pushdown(ll k, ll l, ll r) {
    if (k >= N || lazy[k] == 0) return;
    11 len = (r - 1 + 1) / 2;
    tree[k << 1] = tree[k << 1] + len * lazy[k];</pre>
    tree[k << 1 | 1] = tree[k << 1 | 1] + len * lazy[k];
    lazy[k << 1] = lazy[k << 1] + lazy[k];
    lazy[k << 1 | 1] = lazy[k << 1 | 1] + lazy[k];
    lazy[k] = 0;
```

```
11 merge range(ll a, ll b) {
    11 \text{ ans} = a + b;
    return ans;
void change_range(ll k, ll l, ll r, ll ql, ll qr, ll x) {
    if (r < ql || qr < l)return;
    if (ql <= 1 && r <= qr) {</pre>
        tree[k] = tree[k] + x * (r - 1 + 1);
        lazy[k] = lazy[k] + x;
        return;
    }
    pushdown(k, 1, r);
    11 \text{ mid} = (1 + r) >> 1;
    change_range(k << 1, 1, mid, ql, qr, x);</pre>
    change_range(k << 1 | 1, mid + 1, r, ql, qr, x);
    tree[k] = merge_range(tree[k << 1], tree[k << 1 | 1]);
11 query range(11 k, 11 l, 11 r, 11 q1, 11 qr) {
    if (r < ql || qr < l)return 0;
    if (ql <= 1 && r <= qr) {
        return tree[k];
    }
    pushdown(k, 1, r);
    11 \text{ mid} = (1 + r) >> 1;
    11 lq = query_range(k << 1, 1, mid, ql, qr);</pre>
    ll rq = query range(k \ll 1 \mid 1, mid + 1, r, ql, qr);
    return merge_range(lq, rq);
11 query path(ll a, ll b) {
    11 sum = 0;
    while (top[a] != top[b]) {
        if (dep[top[a]] < dep[top[b]]) swap(a, b);</pre>
        sum = sum + query_range(1, 1, N, dfn[top[a]], dfn[a]);
        //dfn[top[a]]~dfn[a]
        a = fa[top[a]];
    if (dep[a] > dep[b]) swap(a, b);
    //点权
    sum = sum + query_range(1, 1, N, dfn[a], dfn[b]);
    //边权
    //if (a != b) sum = sum + query range(1, 1, N, dfn[a] + 1, dfn[b]);
    //dfn[a]\sim dfn[b],x
    return sum;
void change_path(ll a, ll b, ll x) {
    while (top[a] != top[b]) {
        if (dep[top[a]] < dep[top[b]]) swap(a, b);</pre>
        change_range(1, 1, N, dfn[top[a]], dfn[a], x);
        //dfn[top[a]]~dfn[a]
```

```
a = fa[top[a]];
    }
    if (dep[a] > dep[b]) swap(a, b);
    //点权
    change_range(1, 1, N, dfn[a], dfn[b], x);
   //边权
    //if (a != b) change_range(1, 1, N, dfn[a] + 1, dfn[b], x);
   //dfn[a]\sim dfn[b].x
}
虚树
11 fa[N], son[N], dep[N], siz[N], dfn[N], rnk[N], top[N];
11 dfscnt;
vector<ll> g[N];
11 mmin[N];
void dfs1(ll u, ll f, ll d) {
    son[u] = -1;
    siz[u] = 1;
    fa[u] = f;
    dep[u] = d;
    for (auto v:g[u]) {
        if (v == f) continue;
        dfs1(v, u, d + 1);
        siz[u] += siz[v];
        if (son[u] == -1 \mid | siz[v] > siz[son[u]]) son[u] = v;
    }
}
void dfs2(ll u, ll t) {
    dfn[u] = ++dfscnt;
    rnk[dfscnt] = u;
    top[u] = t;
    if (son[u] == -1) return;
    dfs2(son[u], t);
    for (auto v:g[u]) {
        if (v == son[u] || v == fa[u]) continue;
        dfs2(v, v);
    }
11 lca(ll a, ll b) {
    while (top[a] != top[b]) {
        if (dep[top[a]] < dep[top[b]]) swap(a, b);</pre>
        a = fa[top[a]];
    return dep[a] < dep[b] ? a : b;</pre>
struct edge {
    11 s, t, v;
```

```
};
edge e[N];
vector<int> vg[N];
int sta[N], tot;
int h[N];
void build(int *H, int num) {
   sort(H + 1, H + 1 + num, [](int a, int b) { return dfn[a] < dfn[b];</pre>
});
   sta[tot = 1] = 1, vg[1].clear();// 1 号节点入栈, 清空 1 号节点对应的邻
接表,设置邻接表边数为1
   for (int i = 1, 1; i <= num; ++i) {</pre>
       if (H[i] == 1) continue; //如果 1 号节点是关键节点就不要重复添加
       1 = lca(H[i], sta[tot]); //计算当前节点与栈顶节点的 LCA
      if (1 != sta[tot]) { //如果 LCA 和栈顶元素不同,则说明当前节点不再
当前栈所存的链上
          while (dfn[l] < dfn[sta[tot - 1]]) {//当次大节点的 Dfs 序大于
LCA 的 Dfs 序
             vg[sta[tot - 1]].push_back(sta[tot]);
             vg[sta[tot]].push_back(sta[tot - 1]);
             tot--:
          } //把与当前节点所在的链不重合的链连接掉并且弹出
          if (dfn[1] > dfn[sta[tot - 1]]) { //如果 LCA 不等于次大节点
(这里的大于其实和不等于没有区别)
             vg[1].clear();
             vg[1].push back(sta[tot]);
             vg[sta[tot]].push_back(1);
              sta[tot] = 1;//说明 LCA 是第一次入栈,清空其邻接表,连边后
弹出栈顶元素,并将 LCA 入栈
          } else {
              vg[1].push back(sta[tot]);
             vg[sta[tot]].push_back(1);
             tot--; //说明 LCA 就是次大节点,直接弹出栈顶元素
          }
       }
      vg[H[i]].clear();
       sta[++tot] = H[i];
      //当前节点必然是第一次入栈,清空邻接表并入栈
   for (int i = 1; i < tot; ++i) {</pre>
      vg[sta[i]].push_back(sta[i + 1]);
      vg[sta[i + 1]].push_back(sta[i]);
   } //剩余的最后一条链连接一下
   return;
}
```

```
spfa 最短路及负环
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 N = 1e6 + 10;
struct edge {
    11 to, v;
    edge() {}
    edge(ll a, ll b) : to(a), v(b) {}
};
edge e[N];
11 h[N], vis[N], nxt[N], inque[N], tot, d[N];
void init(ll n, ll m) {
    tot = 0;
    for (int i = 0; i <= n; i++) {</pre>
        h[i] = vis[i] = inque[i] = 0;
    }
    for (int i = 0; i <= m; i++) {</pre>
        nxt[i] = 0;
    }
}
void addedge(ll a, ll b, ll v) {
    nxt[++tot] = h[a];
    e[tot] = edge(b, v);
    h[a] = tot;
bool spfa(ll s, ll n) {
    for (int i = 0; i <= n; i++) d[i] = INT_MAX;</pre>
    d[s] = 0;
    queue<int> que;
    que.push(s);
    inque[s] = 1;
    while (!que.empty()) {
        int p = que.front();
        que.pop();
        vis[p] = 0;
        for (int k = h[p]; k; k = nxt[k]) {
            if (d[e[k].to] > d[p] + e[k].v) {
                d[e[k].to] = d[p] + e[k].v;
                if (!vis[e[k].to]) {
                     inque[e[k].to]++;
                    if (inque[e[k].to] > n) return 0;
                    vis[e[k].to] = 1;
                    que.push(e[k].to);
                }
            }
        }
    return 1;
}
```

```
int main() {
   ios::sync with stdio(0), cin.tie(0), cout.tie(0);
   int t;
   cin >> t;
   while (t--) {
       int n, m;
       cin >> n >> m;
       init(n, m * 2);
       for (int i = 1; i <= m; i++) {</pre>
           int a, b, c;
           cin >> a >> b >> c;
           addedge(a, b, c);
           if (c >= 0) addedge(b, a, c);
       if (spfa(1, n))
           cout << "NO" << endl;</pre>
       else
           cout << "YES" << endl;</pre>
   }
}
二分图匹配(匈牙利)
//大量使用了memset,但常数貌似很小?HDU6808 跑了998ms (限制5000ms),然而
这个代int main()不是HDU6808 的
#include<bits/stdc++.h>
using namespace std;
const int maxn=505;// 最大点数
const int inf=0x3f3f3f3f;// 距离初始值
struct HK_Hungary{//这个板子从1开始,0点不能用,nx 为左边点数,ny 为右边点数
   int nx,ny;//左右顶点数量
   vector<int>bmap[maxn];
   int cx[maxn];//cx[i]表示左集合 i 顶点所匹配的右集合的顶点序号
   int cy[maxn]; //cy[i]表示右集合i 顶点所匹配的左集合的顶点序号
   int dx[maxn];
   int dy[maxn];
   int dis;
   bool bmask[maxn];
   void init(int a,int b){
       nx=a,ny=b;
       for(int i=0;i<=nx;i++){</pre>
           bmap[i].clear();
       }
   void add edge(int u,int v){
       bmap[u].push_back(v);
   bool searchpath(){//寻找 增广路径
```

```
queue<int>0;
       dis=inf;
       memset(dx,-1,sizeof(dx));
       memset(dy,-1,sizeof(dy));
       for(int i=1;i<=nx;i++){//cx[i]表示左集合i 顶点所匹配的右集合的顶点
序号
           if(cx[i]==-1){//将未遍历的节点 入队 并初始化次节点距离为0
              Q.push(i);
              dx[i]=0;
           }
       }//广度搜索增广路径
       while(!Q.empty()){
           int u=Q.front();
           Q.pop();
           if(dx[u]>dis) break;//取右侧节点
           for(int i=0;i<bmap[u].size();i++){</pre>
               int v=bmap[u][i];//右侧节点的增广路径的距离
               if(dy[v]==-1){
                  dy[v]=dx[u]+1;//v 对应的距离 为 u 对应距离加1
                  if(cy[v]==-1)dis=dy[v];
                  else{
                      dx[cy[v]]=dy[v]+1;
                      Q.push(cy[v]);
                  }
               }
           }
       }
       return dis!=inf;
   int findpath(int u){//寻找路径 深度搜索
       for(int i=0;i<bmap[u].size();i++){</pre>
           int v=bmap[u][i];//如果该点没有被遍历过 并且距离为上一节点+1
           if(!bmask[v]&&dy[v]==dx[u]+1){//对该点染色
               bmask[v]=1;
               if(cy[v]!=-1&&dy[v]==dis)continue;
               if(cy[v]==-1||findpath(cy[v])){
                  cy[v]=u;cx[u]=v;
                  return 1;
               }
           }
       }
       return 0;
   int MaxMatch(){//得到最大匹配的数目
       int res=0;
       memset(cx,-1,sizeof(cx));
       memset(cy,-1,sizeof(cy));
       while(searchpath()){
           memset(bmask,0,sizeof(bmask));
```

```
for(int i=1;i<=nx;i++){</pre>
                 if(cx[i]==-1){
                     res+=findpath(i);
                 }
            }
        }
        return res;
    }
}HK;
int main(){
    int nn,n,m;
    cin>>nn;
    while(nn--){
        scanf("%d%d",&n,&m);
        HK.init(n,m);//左端点和右端点数量
        for(int i=1;i<=n;i++){</pre>
            int snum;
            cin>>snum;
            int v;
            for(int j=1;j<=snum;j++){</pre>
                 cin>>v;
                HK.add_edge(i,v);//连边
            }
        }
        cout<<HK.MaxMatch()<<endl;//求最大匹配
    }
    return 0;
}
```

```
强连通(kosaraju
#include <bits/stdc++.h>
using namespace std;
struct SCC {
    static const int MAXV = 100000;
    int V;
    vector<int> g[MAXV], rg[MAXV], vs;
    bool used[MAXV];
    int cmp[MAXV];
    void add_edge(int from, int to) {
        g[from].push_back(to);
        rg[to].push_back(from);
    }
    void dfs(int v) {
        used[v] = 1;
        for (int i = 0; i < g[v].size(); i++) {</pre>
            if (!used[g[v][i]]) dfs(g[v][i]);
        }
        vs.push_back(v);
```

```
void rdfs(int v, int k) {
        used[v] = 1;
        cmp[v] = k;
        for (int i = 0; i < rg[v].size(); i++) {</pre>
            if (!used[rg[v][i]]) rdfs(rg[v][i], k);
        }
    }
    int solve() {
        memset(used, 0, sizeof(used));
        vs.clear();
        for (int v = 1; v \le V; v++) {
            if (!used[v]) dfs(v);
        }
        memset(used, 0, sizeof(used));
        int k = 0;
        for (int i = (int)vs.size() - 1; i >= 0; i--) {
            if (!used[vs[i]]) rdfs(vs[i], ++k);
        return k;
    void init(int n) {
        V = n;
        vs.clear();
        for (int i = 0; i < MAXV; i++) {</pre>
            g[i].clear();
            rg[i].clear();
            used[i] = 0;
            cmp[i] = 0;
        }
    }
} scc;
//记得调用 init()
```

```
强连通(tarjan
#include <bits/stdc++.h>
using namespace std;
struct SCC {
    static const int MAXN = 100000;
    vector<int> g[MAXN];
    int dfn[MAXN], lowlink[MAXN], sccno[MAXN], dfs_clock, scc_cnt;
    stack<int> S;
    void dfs(int u) {
        dfn[u] = lowlink[u] = ++dfs_clock;
        S.push(u);
        for (int i = 0; i < g[u].size(); i++) {
            int v = g[u][i];
            if (!dfn[v]) {
```

```
dfs(v);
                lowlink[u] = min(lowlink[u], lowlink[v]);
            } else if (!sccno[v]) {
                lowlink[u] = min(lowlink[u], dfn[v]);
            }
        if (lowlink[u] == dfn[u]) {
            ++scc_cnt;
            for (;;) {
                int x = S.top();
                S.pop();
                sccno[x] = scc cnt;
                if (x == u) break;
            }
        }
    }
    void solve(int n) {
        dfs clock = scc cnt = 0;
        memset(sccno, 0, sizeof(sccno));
        memset(dfn, 0, sizeof(dfn));
        memset(lowlink, 0, sizeof(lowlink));
        for (int i = 1; i <= n; i++) {</pre>
            if (!dfn[i]) dfs(i);
        }
    }
} scc;
// scc_cnt 为SCC 计数器, sccno[i]为i 所在SCC 的编号
// vector<int> g[MAXN] 中加边
//之后再补充 init()
```

```
强连通(tarjan 无 vector
#include <bits/stdc++.h>
using namespace std;
struct SCC {
    static const int MAXN = 5000;
    static const int MAXM = 2000000;
    int dfs_clock, edge_cnt = 1, scc_cnt;
    int head[MAXN];
    int dfn[MAXN], lowlink[MAXN];
    int sccno[MAXN];
    stack<int> s;
    struct edge {
        int v, next;
    } e[MAXM];
    void add_edge(int u, int v) {
        e[edge_cnt].v = v;
        e[edge cnt].next = head[u];
        head[u] = edge_cnt++;
```

```
void tarjan(int u) {
       int v;
       dfn[u] = lowlink[u] = ++dfs clock; //每次dfs, u 的次序号增加1
                                           //将u入栈
       s.push(u);
       for (int i = head[u]; i != -1; i = e[i].next) //访问从 u 出发的
边
       {
           v = e[i].v;
           if (!dfn[v]) //如果 v 没被处理过
               tarjan(v); // dfs(v)
               lowlink[u] = min(lowlink[u], lowlink[v]);
            } else if (!sccno[v])
               lowlink[u] = min(lowlink[u], dfn[v]);
       if (dfn[u] == lowlink[u]) {
            scc_cnt++;
           do {
               v = s.top();
               s.pop();
               sccno[v] = scc_cnt;
           } while (u != v);
       }
   int find_scc(int n) {
       for (int i = 1; i <= n; i++)</pre>
           if (!dfn[i]) tarjan(i);
       return scc cnt;
    }
   void init() {
       scc_cnt = dfs_clock = 0;
       edge_cnt = 1; //不用初始化 e 数组,省时间
       while (!s.empty()) s.pop();
       memset(head, -1, sizeof(head));
       memset(sccno, 0, sizeof(sccno));
       memset(dfn, 0, sizeof(dfn));
       memset(lowlink, 0, sizeof(lowlink));
} scc;
```

最大流

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
struct Edge {
    ll from, to, cap, flow;
```

```
Edge(ll a, ll b, ll c, ll d) : from(a), to(b), cap(c), flow(d) {}
};
struct Dinic {
    static const ll maxn = 10000;
    11 N, M, S, T;
    vector<Edge> edges;
    vector<ll> G[maxn];
    bool vis[maxn];
    11 d[maxn];
    11 cur[maxn];
    void AddEdge(ll from, ll to, ll cap) {
        edges.push_back(Edge(from, to, cap, 0));
        edges.push_back(Edge(to, from, 0, 0));
       M = edges.size();
        G[from].push_back(M - 2);
        G[to].push_back(M - 1);
    }
    bool BFS() {
        memset(vis, 0, sizeof(vis));
        queue<11> Q;
        Q.push(S);
        d[S] = 0;
        vis[S] = 1;
        while (!Q.empty()) {
            11 \times = Q.front();
            Q.pop();
            for (11 i = 0; i < G[x].size(); i++) {</pre>
                Edge& e = edges[G[x][i]];
                if (!vis[e.to] && e.cap > e.flow) {
                    vis[e.to] = 1;
                    d[e.to] = d[x] + 1;
                    Q.push(e.to);
                }
            }
        }
        return vis[T];
    ll DFS(ll x, ll a) {
        if (x == T | | a == 0) return a;
        11 flow = 0, f;
        for (l1& i = cur[x]; i < G[x].size(); i++) {</pre>
            Edge& e = edges[G[x][i]];
            if (d[x] + 1 == d[e.to] &&
                (f = DFS(e.to, min(a, e.cap - e.flow))) > 0) {
                e.flow += f;
                edges[G[x][i] ^ 1].flow -= f;
                flow += f;
                a -= f;
                if (a == 0) break;
```

```
}
        }
        return flow;
    11 Maxflow(11 S, 11 T) {
        this->S = S, this->T = T;
        11 flow = 0;
        while (BFS()) {
            memset(cur, 0, sizeof(cur));
            flow += DFS(S, inf);
        return flow;
} MF;
最大流 (double)
#include <iostream>
#include <cstring>
#include <algorithm>
using namespace std;
struct Dinic {
      static constexpr int N = 10010, M = 100010, INF = 1e8;
      static constexpr double eps = 1e-8;
//
      int n, m, S, T;
      int S, T;
      int h[N], e[M], ne[M], idx;
      double f[M];
      int q[N], d[N], cur[N]; // d 表示从源点开始走到该点的路径上所有边的容
量的最小值
      void AddEdge(int a, int b, double c)
      {
          e[idx] = b, f[idx] = c, ne[idx] = h[a], h[a] = idx ++ ;
          e[idx] = a, f[idx] = 0, ne[idx] = h[b], h[b] = idx ++ ;
      }
      bool bfs()
          int hh = 0, tt = 0;
          memset(d, -1, sizeof d);
          q[0] = S, d[S] = 0, cur[S] = h[S];
          while (hh <= tt)</pre>
              int t = q[hh ++];
              for (int i = h[t]; ~i; i = ne[i])
                  int ver = e[i];
                  if (d[ver] == -1 && f[i] > 0)
```

```
{
                      d[ver] = d[t] + 1;
                      cur[ver] = h[ver];
                      if (ver == T) return true;
                      q[ ++ tt] = ver;
                  }
              }
          }
          return false;
      }
      double find(int u, double limit)
      {
          if (u == T) return limit;
          double flow = 0;
          for (int i = cur[u]; ~i && flow < limit; i = ne[i])</pre>
              cur[u] = i;
              int ver = e[i];
              if (d[ver] == d[u] + 1 && f[i] > 0)
                  double t = find(ver, min(f[i], limit - flow));
                  if (t < eps) d[ver] = -1;
                  f[i] -= t, f[i ^ 1] += t, flow += t;
              }
          }
          return flow;
      }
      double Maxflow(int S, int T)
      {
            this->S = S, this->T = T;
          double r = 0, flow;
          while (bfs()) while (flow = find(S, INF)) r += flow;
          return r;
      void init() ///////
            memset(h, -1, sizeof h);
            idx = 0;
} MF;
// ?èinit
```

最小费用最大流

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
struct Edge {
    ll from, to, cap, flow, cost;
    Edge(ll u, ll v, ll c, ll f, ll w):from(u), to(v), cap(c), flow(f),
cost(w) {}
};
struct MCMF {
    static const 11 maxn = 6000;
    static const 11 INF = 0x3f3f3f3f3f3f3f3f;
    11 n, m;
    vector<Edge> edges;
    vector<11> G[maxn];
    11 inq[maxn];
    11 d[maxn];
    11 p[maxn];
    11 a[maxn];
    void init(ll n) {
        this->n = n;
        for (ll i = 1; i <= n; i++) G[i].clear();</pre>
        edges.clear();
    }
    void add_edge(ll from, ll to, ll cap, ll cost) {
        from++, to++;//原板子无法使用 0 点,故修改
        edges.push_back(Edge(from, to, cap, 0, cost));
        edges.push back(Edge(to, from, 0, 0, -cost));
        m = edges.size();
        G[from].push_back(m - 2);
        G[to].push back(m - 1);
    bool BellmanFord(ll s, ll t, ll& flow, ll& cost) {
        for (ll i = 1; i <= n; ++i) d[i] = INF;</pre>
        memset(inq, 0, sizeof(inq));
        d[s] = 0, inq[s] = 1, p[s] = 0, a[s] = INF;
        queue<11> Q;
        Q.push(s);
        while (!Q.empty()) {
            11 u = Q.front();
            Q.pop();
            inq[u] = 0;
            for (ll i = 0; i < G[u].size(); ++i) {</pre>
                Edge& e = edges[G[u][i]];
                if (e.cap > e.flow && d[e.to] > d[u] + e.cost) {
                    d[e.to] = d[u] + e.cost;
                    p[e.to] = G[u][i];
                    a[e.to] = min(a[u], e.cap - e.flow);
                    if (!inq[e.to]) {
```

```
Q.push(e.to);
                      inq[e.to] = 1;
                  }
              }
           }
       if (d[t] == INF) return false;
       flow += a[t];
       cost += (11)d[t] * (11)a[t];
       for (ll u = t; u != s; u = edges[p[u]].from) {
           edges[p[u]].flow += a[t];
           edges[p[u] ^ 1].flow -= a[t];
       return true;
   }
   //需要保证初始网络中没有负权圈
   11 MincostMaxflow(ll s, ll t, ll& cost) {
       S++, t++;//原板子无法使用 0 点,故修改
       11 flow = 0;
       cost = 0;
       while (BellmanFord(s, t, flow, cost));
       return flow;
} mcmf; // 若固定流量k,增广时在flow+a>=k的时候只增广k-flow单位的流量,
然后终止程序
```

```
树分治
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 10005;
const int INF = 1000000000;
struct edge {
    int to, length;
    edge() {}
    edge(int a, int b) : to(a), length(b) {}
};
int N, K;
vector<edge> g[MAXN];
bool centroid[MAXN];
int subtree_size[MAXN];
int ans;
//计算子树大小
int compute_subtree_size(int v, int p) {
    int c = 1;
    for (int i = 0; i < g[v].size(); i++) {</pre>
        int w = g[v][i].to;
        if (w == p || centroid[w]) continue;
```

```
c += compute subtree size(w, v);
    }
    subtree_size[v] = c;
    return c;
}
//查找重心,t 为连通分量大小
// pair (最大子树顶点数, 顶点编号)
pair<int, int> search centroid(int v, int p, int t) {
    pair<int, int> res = pair<int, int>(INF, -1);
    int s = 1, m = 0;
    for (int i = 0; i < g[v].size(); i++) {</pre>
        int w = g[v][i].to;
        if (w == p || centroid[w]) continue;
        res = min(res, search centroid(w, v, t));
        m = max(m, subtree_size[w]);
        s += subtree_size[w];
    }
    m = max(m, t - s);
    res = min(res, pair<int, int>(m, v));
    return res;
}
void init() {
    memset(centroid, 0, sizeof(centroid));
    memset(subtree_size, 0, sizeof(subtree_size));
    for (int i = 0; i <= N; i++) g[i].clear();</pre>
    ans = 0;
int solve(int u) {
    compute_subtree_size(u, -1);
    int s = search_centroid(u, -1, subtree_size[u]).second;
    centroid[s] = 1;
    int ans:
    for (int i = 0; i < g[s].size(); i++) {</pre>
        int v = g[s][i].to;
        if (centroid[v]) continue;
        /*solve()*/
    /*do something*/
    centroid[s] = 0;
    return ans;
}
```

```
拓扑排序
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 100000;
int c[MAXN];
int topo[MAXN], t, V;
```

```
vector<int> g[MAXN];
bool dfs(int u) {
    c[u] = -1;
    for (int i = 0; i < g[u].size(); i++) {</pre>
        int v = g[u][i];
        if (c[v] < 0)
            return false;
        else if (!c[v] && !dfs(v))
            return false;
    }
    c[u] = 1;
    topo[t--] = u;
    return true;
bool toposort(int n) {
   V = n;
    t = n;
    memset(c, 0, sizeof(c));
    for (int u = 1; u <= V; u++)</pre>
        if (!c[u] && !dfs(u)) return false;
    return true;
}
```

```
最近公共祖先(倍增)
#include <algorithm>
#include <cstdio>
#include <cstring>
#include <iostream>
using namespace std;
const int MAX = 600000;
struct edge {
    int t, nex;
} e[MAX << 1];</pre>
int head[MAX], tot;
int depth[MAX], fa[MAX][22], lg[MAX];
void add_edge(int x, int y) {
    e[++tot].t = y;
    e[tot].nex = head[x];
    head[x] = tot;
    e[++tot].t = x;
    e[tot].nex = head[y];
    head[y] = tot;
void dfs(int now, int fath) {
    fa[now][0] = fath;
    depth[now] = depth[fath] + 1;
    for (int i = 1; i <= lg[depth[now]]; ++i)</pre>
        fa[now][i] = fa[fa[now][i - 1]][i - 1];
```

```
for (int i = head[now]; i; i = e[i].nex)
        if (e[i].t != fath) dfs(e[i].t, now);
int lca(int x, int y) {
    if (depth[x] < depth[y]) swap(x, y);</pre>
    while (depth[x] > depth[y]) x = fa[x][lg[depth[x] - depth[y]] - 1];
    if (x == y) return x;
    for (int k = \lg[depth[x]] - 1; k \ge 0; --k)
        if (fa[x][k] != fa[y][k]) x = fa[x][k], y = fa[y][k];
    return fa[x][0];
}
void init(int n, int root) {
    for (int i = 1; i \le n; ++i) lg[i] = lg[i - 1] + (1 << lg[i - 1] ==
 i);
    dfs(root, ∅);
}
最近公共祖先(线段树)
#include <bits/stdc++.h>
using namespace std;
int n, m, root;
const int MAX_N = 500005;
const int MAX = 1 << 20;
vector<int> g[MAX_N];
vector<int> vs;
pair<int, int> tree[MAX * 2 + 10];
int fir[MAX N];
int fa[MAX N];
int dep[MAX_N];
void dfs(int k, int p, int d) {
    fa[k] = p;
    dep[k] = d;
    vs.push back(k);
    for (int i = 0; i < g[k].size(); i++) {</pre>
        if (g[k][i] != p) {
            dfs(g[k][i], k, d + 1);
            vs.push_back(k);
        }
    }
}
void build(int k) {
    if (k >= MAX) return;
    build(k \ll 1);
    build(k \ll 1 | 1);
    tree[k] = min(tree[k << 1], tree[k << 1 | 1]);
pair<int, int> query(int k, int s, int e, int l, int r) {
    if (e < 1 | | r < s) return pair<int, int>(INT MAX, 0);
    if (1 <= s && e <= r) return tree[k];
    return min(query(k \langle\langle 1, s, (s + e) \rangle\rangle 1, l, r),
```

```
query(k \ll 1 \mid 1, ((s + e) >> 1) + 1, e, l, r));
}
void init() {
    dfs(root, root, ∅);
    for (int i = 0; i < MAX * 2 + 10; i++) tree[i] = pair<int, int>(INT
_MAX, 0);
    for (int i = MAX; i < MAX + vs.size(); i++)</pre>
        tree[i] = pair<int, int>(dep[vs[i - MAX]], vs[i - MAX]);
    for (int i = 0; i < vs.size(); i++) {</pre>
        if (fir[vs[i]] == 0) fir[vs[i]] = i + 1;
    build(1);
int lca(int a, int b) {
    return query(1, 1, MAX, min(fir[a], fir[b]), max(fir[a], fir[b])).s
econd;
}
int main() {
    scanf("%d%d%d", &n, &m, &root);
    for (int i = 1; i < n; i++) {</pre>
        int a, b;
        scanf("%d%d", &a, &b);
        g[a].push_back(b);
        g[b].push_back(a);
    }
    init();
    for (int i = 1; i <= m; i++) {</pre>
        int a, b;
        scanf("%d%d", &a, &b);
        printf("%d\n", lca(a, b));
    }
}
```

线性代数

```
高斯消元
#include <iostream>
#include <vector>
using namespace std;
const double eps = 1e-8;
void sway(vector<double>& a, vector<double>& b) {
    vector<double> s;
    for (int i = 0; i < a.size(); i++) {
        s.push_back(a[i]);
    }
    a.clear();
```

```
for (int i = 0; i < b.size(); i++) {</pre>
        a.push back(b[i]);
    b.clear();
    for (int i = 0; i < s.size(); i++) {</pre>
        b.push_back(s[i]);
    }
}
vector<double> gauss_jordan(const vector<vector<double> >& A,
                              const vector<double>& b) {
    int n = A.size();
    vector<vector<double> > B(n, vector<double>(n + 1));
    for (int i = 0; i < n; i++)</pre>
        for (int j = 0; j < n; j++) B[i][j] = A[i][j];</pre>
    for (int i = 0; i < n; i++) B[i][n] = b[i];</pre>
    for (int i = 0; i < n; i++) {
        int pivot = i;
        for (int j = i; j < n; j++) {</pre>
             if (abs(B[j][i]) > abs(B[pivot][i])) pivot = j;
        }
        swap(B[i], B[pivot]);
        if (abs(B[i][i]) < eps) return vector<double>();
        for (int j = i + 1; j <= n; j++) B[i][j] /= B[i][i];</pre>
        for (int j = 0; j < n; j++) {</pre>
             if (i != j) {
                 for (int k = i + 1; k \le n; k++) B[j][k] -= B[j][i] * B
[i][k];
             }
        }
    }
    vector<double> x(n);
    for (int i = 0; i < n; i++) x[i] = B[i][n];</pre>
    return x;
int main() {
    int n, m;
    cin >> n >> m;
    vector<vector<double> > mat(n, vector<double>(m));
    for (int i = 0; i < n; i++) {</pre>
        for (int j = 0; j < m; j++) {
             cin >> mat[i][j];
        }
    }
    vector<double> val(n);
    for (int i = 0; i < n; i++) cin >> val[i];
    vector<double> ans = gauss jordan(mat, val);
    for (int i = 0; i < ans.size(); i++) cout << ans[i] << ' ';</pre>
}
```

```
矩阵行列式
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 \mod = 1e9 + 7;
struct Matrix {
    static const 11 MAXN = 300;
    11 a[MAXN][MAXN];
    void init() { memset(a, 0, sizeof(a)); }
    11 det(11 n) {
        for (int i = 0; i < n; i++)</pre>
             for (int j = 0; j < n; j++) a[i][j] = (a[i][j] + mod) % mod;</pre>
        11 \text{ res} = 1;
        for (int i = 0; i < n; i++) {</pre>
             if (!a[i][i]) {
                 bool flag = false;
                 for (int j = i + 1; j < n; j++) {</pre>
                     if (a[j][i]) {
                          flag = true;
                          for (int k = i; k < n; k++) {</pre>
                              swap(a[i][k], a[j][k]);
                          }
                          res = -res;
                          break;
                     }
                 if (!flag) return 0;
             }
            for (int j = i + 1; j < n; j++) {</pre>
                 while (a[j][i]) {
                     11 t = a[i][i] / a[j][i];
                     for (int k = i; k < n; k++) {</pre>
                          a[i][k] = (a[i][k] - t * a[j][k]) % mod;
                          swap(a[i][k], a[j][k]);
                     }
                     res = -res;
                 }
             }
            res *= a[i][i];
            res %= mod;
        return (res + mod) % mod;
} mat;
```

```
线性基
//
const int maxbit = 62; //maxbit 不能太大
```

```
struct L B{
     11 lba[maxbit];
     L_B(){
       memset(lba, 0, sizeof(lba));
   }
     void Insert(ll val){
                               //插入
       for(int i = maxbit - 1; i >= 0; -- i) // 从高位向低位扫
          if(val & (1ll << i)){ //</pre>
              if(!lba[i]){
                  lba[i] = val;
                  break;
              }
              val ^= lba[i];
          }
   }
};
//对原集合的每个数 val 转为 2 进制,从高位向低位扫,对于当前位为 1 的,若 Lba[i]
不存在就令 Lba[i]=x, 否则令 val=val`xor`Lba[i]
//使用: 直接insert
// -----线性基模板
```

线性基 2

线性基 能表示的线性空间与原向量 能表示的线性空间等价

用高斯消元得到线性基

先输入数组 a[] 中

```
}
    ++k;
    if(k == n) break;
}
```

这里注意最后的线性基是 a[]中从 0 到 k-1 个,在前的是高位

矩阵(加减乘快速幂

```
//矩阵类模板
struct Matrix{
    int n,m;
    int a[maxn][maxm];
    void clear(){
        n=m=0;
        memset(a,0,sizeof(a));
    }
    Matrix operator +(const Matrix &b) const {
        Matrix tmp;
        tmp.n=n;tmp.m=m;
        for (int i=0;i<n;++i)</pre>
            for(int j=0;j<m;++j)</pre>
                 tmp.a[i][j]=a[i][j]+b.a[i][j];
        return tmp;
    Matrix operator -(const Matrix &b)const{
        Matrix tmp;
        tmp.n=n;tmp.m=m;
        for (int i=0;i< n;++i){
            for(int j=0;j<m;++j)</pre>
                 tmp.a[i][j]=a[i][j]-b.a[i][j];
            }
        return tmp;
    Matrix operator * (const Matrix &b) const{
        Matrix tmp;
        tmp.clear();
        tmp.n=n;tmp.m=b.m;
        for (int i=0;i<n;++i)</pre>
            for(int j=0;j<b.m;++j)</pre>
                 for (int k=0; k<m; ++k){
                     tmp.a[i][j]+=a[i][k]*b.a[k][j];
                     tmp.a[i][j]%=mod;
                 }
```

```
return tmp;
    }
   Matrix get(int x){//幂运算
        Matrix E;
        E.clear();
        E.n=E.m=n;
        for(int i=0;i<n;++i)</pre>
            E.a[i][i]=1;
        if(x==0) return E;
        else if(x==1) return *this;
        Matrix tmp=get(x/2);
        tmp=tmp*tmp;
        if(x%2) tmp=tmp*(*this);
        return tmp;
    }
};
//矩阵模板结束
```

稀疏矩阵乘法

```
struct Matrix{
    int n,m;
    int a[maxn][maxn];////
    void clear(){
        n=m=0;
        memset(a,0,sizeof(a));
   Matrix operator * (const Matrix &b) const{
        Matrix tmp;
        tmp.clear();
        tmp.n=n;tmp.m=b.m;
        for (int k=0;k<m;++k){</pre>
            for (int i=0;i<n;++i){</pre>
                   if(a[i][k]==0) continue;
                   for(int j=0;j<b.m;++j){</pre>
                         if(b.a[k][j]==0) continue;
                         tmp.a[i][j]+=a[i][k]*b.a[k][j];
                     tmp.a[i][j]%=mod;
                         }
                   }
        return tmp;
    }
};
//稀疏矩阵乘法
```

杂项

快读

```
inline int read(){
   int X=0,w=0;char ch=0;
   while(!isdigit(ch)){w|=ch=='-';ch=getchar();}
   while(isdigit(ch))X=(X<<3)+(X<<1)+(ch^48),ch=getchar();
   return w?-X:X;
}</pre>
```

fread 快读

```
#include <bits/stdc++.h>
using namespace std;
char next_char() {
      static char buf[1 << 20], *first, *last;</pre>
      if(first == last) {
            last = buf + fread(buf, 1, 1 << 20, stdin);</pre>
            first = buf;
      return first == last ? EOF : *first ++;
inline int read(){
      int x = 0, w = 0; char ch = 0;
      while(!isdigit(ch)) {w |= ch == '-'; ch = next_char(); }
      while(isdigit(ch)) \{x = (x << 3) + (x << 1) + (ch ^ 48), ch = nex\}
t char(); }
      return w ? -x : x;
int main(){
      freopen("1.txt", "r", stdin); // 交代码的时候一定要去掉 aaa
      int T;
      cin >> T;
      while(T --){
            int x = read();
            cout << x << endl;</pre>
      }
}
```

模拟退火

"优化的随机算法"

连续函数找区间最优

// 找一个点, 与平面中的 n 个点的距离和最近

//进行多次模拟退火避免局部最大值

```
#include <bits/stdc++.h>
#include <ctime>
using namespace std;
const int maxn = 110;
int n;
#define x first
#define y second
typedef pair<double, double> PDD;
PDD q[maxn];
double ans = 1e8;
double rand(double 1, double r) {
   return (double) rand() / RAND_MAX * (r - 1) + 1;
}
double getDist(PDD a, PDD b) {
   double dx = a.x - b.x;
   double dy = a.y - b.y;
   return sqrt(dx * dx + dy * dy);
}
double calc(PDD p) {
   double res = 0;
   for(int i = 0; i < n; ++ i) {
       res += getDist(q[i], p);
   ans = min(ans, res);
   return res;
}
double simulate anneal() {
   PDD cur(rand(0, 10000), rand(0, 10000)); // 随机一个起点
   for(double T = 1e4; T > 1e-4; T = T * 0.99) { // 初始温度, 末态温度,
衰减系数,一般调整衰减系数0.999 0.95
       PDD np(rand(cur.x - T, cur.x + T), rand(cur.y - T, cur.y + T));
// 随机新点
       double delta = calc(np) - calc(cur);
       if(exp(-delta / T) > rand(0, 1)) cur = np; //如果新点比现在的点更
优,必过去,不然有一定概率过去
int main() {
   cin >> n;
   for(int i = 0; i < n; ++ i) {
       cin >> q[i].x >> q[i].y;
   while((double) clock() / CLOCKS_PER_SEC < 0.8) { // 卡时 // 或for
 (100)
       simulate anneal();
   cout << (int)(ans + 0.5) << endl;
```

```
return 0;
}
//n 个点带权费马点 // 平衡点||吊打 XXX
//n 个二维坐标点, 带重物重量, 找平衡点
//进行一次模拟退火,但是在局部最大值周围多次跳动(以提高精度
#include <cmath>
#include <cstdio>
#include <cstdlib>
#include <ctime>
const int N = 10005;
int n, x[N], y[N], w[N];
double ansx, ansy, dis;
double Rand() { return (double)rand() / RAND_MAX; }
double calc(double xx, double yy) {
  double res = 0;
  for (int i = 1; i <= n; ++i) {</pre>
    double dx = x[i] - xx, dy = y[i] - yy;
   res += sqrt(dx * dx + dy * dy) * w[i];
  if (res < dis) dis = res, ansx = xx, ansy = yy;</pre>
  return res;
void simulateAnneal() {
  double t = 100000;
  double nowx = ansx, nowy = ansy;
 while (t > 0.001) {
    double nxtx = nowx + t * (Rand() * 2 - 1);
    double nxty = nowy + t * (Rand() * 2 - 1);
    double delta = calc(nxtx, nxty) - calc(nowx, nowy);
    if (exp(-delta / t) > Rand()) nowx = nxtx, nowy = nxty;
   t *= 0.97;
  for (int i = 1; i <= 1000; ++i) {
    double nxtx = ansx + t * (Rand() * 2 - 1);
    double nxty = ansy + t * (Rand() * 2 - 1);
    calc(nxtx, nxty);
  }
}
int main() {
  srand(time(∅));
  scanf("%d", &n);
 for (int i = 1; i <= n; ++i) {</pre>
    scanf("%d%d%d", &x[i], &y[i], &w[i]);
    ansx += x[i], ansy += y[i];
  }
```

```
ansx /= n, ansy /= n, dis = calc(ansx, ansy);
  simulateAnneal();
  printf("%.31f %.31f\n", ansx, ansy);
  return 0;
}
整体二分
11 bit[N];
void add_bit(ll k, ll a) {
   while (k < N) {
        bit[k] = bit[k] + a;
        k += k \& -k;
    }
11 query_bit(ll k) {
    11 \text{ ans} = 0;
   while (k) {
        ans = ans + bit[k];
        k -= k \& -k;
    }
    return ans;
}
struct node {
    ll x, y, k, id, type;
};
node q[N], q1[N], q2[N];
11 ans[N], now[N], tot, totx;
void solve(11 1, 11 r, 11 q1, 11 qr) {
    if (ql > qr) return;
    if (1 == r) {
        for (ll i = ql; i <= qr; i++) {</pre>
            if (q[i].type == 2) {
                ans[q[i].id] = 1;
            }
        }
        return;
    11 \text{ mid} = (1 + r) >> 1;
    11 cq1 = 0, cq2 = 0;
    for (ll i = ql; i <= qr; i++) {</pre>
        if (q[i].type == 1) {
            if (q[i].y <= mid) {
                 add_bit(q[i].x, q[i].k);
                q1[++cq1] = q[i];
            } else {
                q2[++cq2] = q[i];
        } else {
            11 sum = query_bit(q[i].y) - query_bit(q[i].x - 1);
            if (sum >= q[i].k) {
```

```
q1[++cq1] = q[i];
            } else {
                q2[++cq2] = q[i];
                q2[cq2].k -= sum;
            }
        }
    }
    for (ll i = 1; i <= cq1; i++) if (q1[i].type == 1) add_bit(q1[i].x,</pre>
    for (ll i = 1; i \le cq1; i++) q[ql + i - 1] = q1[i];
    for (ll i = 1; i <= cq2; i++) q[ql + cq1 + i - 1] = q2[i];
    solve(1, mid, ql, ql + cq1 - 1);
    solve(mid + 1, r, ql + cq1, qr);
void init() {
   totx = 0;
    tot = 0;
    memset(bit, 0, sizeof bit);
}
```

字符串

```
马拉车
```

```
#include <bits/stdc++.h>
using namespace std;
const int maxn = 100005;
char s[maxn];
char s_new[maxn * 2];
int p[maxn * 2];
int Manacher(char* a, int 1) {
    s_new[0] = '$';
    s new[1] = '#';
    int len = 2;
    for (int i = 0; i < 1; i++) {
        s_new[len++] = a[i];
        s_new[len++] = '#';
    }
    s new[len] = '\0';
    int id;
    int mx = 0;
    int mmax = 0;
    for (int i = 1; i < len; i++) {</pre>
        p[i] = i < mx ? min(p[2 * id - i], mx - i) : 1;
        while (s_new[i + p[i]] == s_new[i - p[i]]) p[i]++;
```

```
if (mx < i + p[i]) {</pre>
            id = i;
            mx = i + p[i];
        mmax = max(mmax, p[i] - 1);
    return mmax;
}
int main() {
    cin >> s;
    cout << Manacher(s, strlen(s));</pre>
}
AC 自动机
#include <bits/stdc++.h>
using namespace std;
struct AC {
    static const int maxnode = 200005;
    static const int sigma size = 26;
    char T[maxnode];
    int ch[maxnode][sigma size];
    int val[maxnode], fail[maxnode], last[maxnode];
    int sz;
    vector<pair<int, int> > ans;
    void init() {
        sz = 1;
        memset(ch[0], 0, sizeof(ch[0]));
        ans.clear();
    int idx(const char &c) { return c - 'a'; }
    void insert(string s, int v) {
        int u = 0, n = s.length();
        for (int i = 0; i < n; i++) {</pre>
            int c = idx(s[i]);
            if (!ch[u][c]) {
                memset(ch[sz], 0, sizeof(ch[sz]));
                val[sz] = 0;
                ch[u][c] = sz++;
            u = ch[u][c];
        val[u] = v;
    void get_fail() {
        queue<int> que;
        fail[0] = 0;
        for (int c = 0; c < sigma_size; c++) {</pre>
            int u = ch[0][c];
            if (u) {
                fail[u] = 0;
```

```
que.push(u);
                last[u] = 0;
            }
        }
        while (!que.empty()) {
            int r = que.front();
            que.pop();
            for (int c = 0; c < sigma_size; c++) {</pre>
                int u = ch[r][c];
                if (!u) continue;
                que.push(u);
                int v = fail[r];
                while (v \&\& !ch[v][c]) v = fail[v];
                fail[u] = ch[v][c];
                last[u] = val[fail[u]] ? fail[u] : last[fail[u]];
            }
        }
    }
    void print(int j) {
        if (j) {
            ans.push_back(pair<int, int>(j, val[j]));
            print(last[j]);
        }
    void find() {
        int n = strlen(T);
        int j = 0;
        for (int i = 0; i < n; i++) {</pre>
            int c = idx(T[i]);
            while (j && !ch[j][c]) j = fail[j];
            j = ch[j][c];
            if (val[j])
                print(j);
            else if (last[j])
                print(last[j]);
        }
        //字符串下标从 0 开始
} ac;
```

KMP

```
//next 数组等价于前缀函数
#include<bits/stdc++.h>
using namespace std;
typedef long long ll;
int kmp(char *s1,int *p1,char *s2=0,int *p2=0){//必须先求s1 的 next 数组,即 kmp(s1,p1);再 kmp(s1,p1,s2,p2);
    int n=strlen(s1);
    if(p2==0){
```

```
p1[0]=0;
        for(int i=1;s1[i]!='\0';i++){
            int j=p1[i-1];
            while(j>0&&s1[i]!=s1[j])j=p1[j-1];
            if(s1[i]==s1[j])j++;
            p1[i]=j;
        }
    }
    else{
        for(int i=0;s2[i]!='\0';i++){
            int j=i==0?0:p2[i-1];
            while(j>0&&s2[i]!=s1[j])j=p1[j-1];
            if(s2[i]==s1[j])j++;
            p2[i]=j;
            if(j==n)return i-n+2;//返回位置
        }
    }
    return 0;
int main(){
    char s1[15],s2[105];
    int p1[15],p2[105];
    cin>>s1>>s2;
    kmp(s1,p1);
    cout<<kmp(s1,p1,s2,p2)<<endl;</pre>
    return 0;
}
KMP<sub>2</sub>
#include <bits/stdc++.h>
using namespace std;
struct KMP {
    static const int MAXN = 1000010;
    char T[MAXN], P[MAXN];
    int fail[MAXN];
    vector<int> ans;
    void init() { ans.clear(); }
    void get_fail() {
        int m = strlen(P);
        fail[0] = fail[1] = 0;
        for (int i = 1; i < m; i++) {</pre>
            int j = fail[i];
```

while (j && P[i] != P[j]) j = fail[j];
fail[i + 1] = (P[i] == P[j] ? j + 1 : 0);

int n = strlen(T), m = strlen(P);

}

void find() {

}

```
get_fail();
int j = 0;
for (int i = 0; i < n; i++) {
    while (j && P[j] != T[i]) j = fail[j];
    if (P[j] == T[i]) j++;
    if (j == m) ans.push_back(i - m + 1);
    }
} kmp; //P 为模式串,下标从 0 开始,输入后直接调用 find()</pre>
```

```
Tire
#include <bits/stdc++.h>
using namespace std;
struct Trie {
    static const int maxnode = 200005;
    static const int sigma_size = 26;
    int ch[maxnode][sigma_size];
    int val[maxnode];
    int sz;
    Trie() {
        sz = 1;
        memset(ch[0], 0, sizeof(ch[0]));
    int idx(const char &c) { return c - 'a'; }
    void insert(string s, int v) {
        int u = 0, n = s.length();
        for (int i = 0; i < n; i++) {</pre>
            int c = idx(s[i]);
            if (!ch[u][c]) {
                memset(ch[sz], 0, sizeof(ch[sz]));
                val[sz] = 0;
                ch[u][c] = sz++;
            }
            u = ch[u][c];
        val[u] = v;
    int find(string s) {
        int u = 0, n = s.length();
        for (int i = 0; i < n; i++) {</pre>
            int c = idx(s[i]);
            if (!ch[u][c]) return 0;
            u = ch[u][c];
        return val[u];
} trie;
```

```
后缀数组
#include <bits/stdc++.h>
using namespace std;
struct SuffixArray {
    static const int MAXN = 1100000;
    char s[MAXN];
    int sa[MAXN], t[MAXN], t1[MAXN], c[MAXN], ra[MAXN], height[MAXN], m;
    inline void init() { memset(this, 0, sizeof(SuffixArray)); }
    inline void get_sa(int n) {
        m = 256;
        int *x = t, *y = t1;
        for (int i = 1; i <= m; i++) c[i] = 0;
        for (int i = 1; i <= n; i++) c[x[i] = s[i]]++;
        for (int i = 1; i <= m; i++) c[i] += c[i - 1];</pre>
        for (int i = n; i >= 1; i--) sa[c[x[i]]--] = i;
        for (int k = 1; k <= n; k <<= 1) {
            int p = 0;
            for (int i = n - k + 1; i <= n; i++) y[++p] = i;
            for (int i = 1; i <= n; i++)</pre>
                if (sa[i] > k) y[++p] = sa[i] - k;
            for (int i = 1; i <= m; i++) c[i] = 0;</pre>
            for (int i = 1; i <= n; i++) c[x[y[i]]]++;
            for (int i = 1; i <= m; i++) c[i] += c[i - 1];</pre>
            for (int i = n; i >= 1; i--) sa[c[x[y[i]]]--] = y[i];
            std::swap(x, y);
            p = x[sa[1]] = 1;
            for (int i = 2; i <= n; i++) {
                x[sa[i]] = (y[sa[i - 1]] == y[sa[i]] &&
                             y[sa[i - 1] + k] == y[sa[i] + k])
                                ? p
                                : ++p;
            if (p >= n) break;
            m = p;
        }
    }
    inline void get height(int n) {
        int i, j, k = 0;
        for (int i = 1; i <= n; i++) ra[sa[i]] = i;</pre>
        for (int i = 1; i <= n; i++) {</pre>
            if (k) k--;
            int j = sa[ra[i] - 1];
            while (s[i + k] == s[j + k]) k++;
            height[ra[i]] = k;
        }
```

//字符串下标从一开始

} SA;

```
可持久化字典树
struct Trie01 {
    static const int maxnode = 2000005;
    static const int sigma size = 2;
    int ch[maxnode << 5][sigma_size], val[maxnode << 5];</pre>
    int rt[maxnode];
    int sz;
    Trie01() {
        sz = 0;
        memset(ch[0], 0, sizeof(ch[0]));
    void insert(int &now, int pre, int v) {
        now = ++sz;
        for (int i = 30; i >= 0; i--) {
            int k = ((v >> i) & 1);
            ch[now][k] = ++sz;
            ch[now][k ^ 1] = ch[pre][k ^ 1];
            val[ch[now][k]] = val[ch[pre][k]] + 1;
            now = ch[now][k];
            pre = ch[pre][k];
        }
} trie;
```

对拍

windows 环境下 bat 对拍

\$\color{red}{A.exe}\$ < data.txt > test.txt

```
fc ac.txt test.txt
if not errorlevel 1 goto loop
pause
goto loop
```

文件以.bat 作为后缀

将三个程序(数据生成文件(dataa),标程或暴力代码(biaocheng), 要看的代码(A))放在同一目录下,

记得加 freopen

随机数记得加 srand((int)time(0));

随机数生成 code

```
#include <iostream>
#include <cstdlib>
#include <ctime>
using namespace std;
int main(){
    freopen("data.txt", "w", stdout);

    srand((int)time(0));
    int T = rand() % 100000;
    cout << T << endl;

    for (int i = 0; i < T; i++){
        cout << rand() % 100;
    }
}</pre>
```

rand()似乎只有三万多,需要更大的数的话要乘一下