EARTHCAT 2021年11月模板

目录

汪怠	4
定义	4
STL	5
优先队列重载	5
set 重载	
动态开数组	
vector	
set	
map	
插入操作	9
取值	
容量查询	9
删除	
交换	
顺序比较	
查找	
型X操作符	
採旧句 unordered_map unordered_	
插入数据	
遍历 map	
她// mapbitset	
计算几何	
几何的一些定理	
多面体欧拉定理	
zyx 的计算几何	
计算几何全家桶	
自适应辛普森	
球体积交和并	
数据结构	
仙人掌	
CDQ	
kruskal 重构树	
普通莫队	
带修莫队	58
回滚莫队	
线段树合并分裂	64
主席树	66
LCT	68
Splay1	70
splay2	
Treap	
舞蹈链(多重覆盖)	
舞蹈链(精确覆盖)	
数 论	84

lucas 求组合数	84
扩展欧几里得求逆元	85
逆元线性递推 inv 阶乘组合数	86
数学	86
一些范围	86
- 2e9 中拥有最多约数个数的数拥有的约数个数	86
勾股数/圆上格点数	86
勾股数	
完全公式	87
高斯整数/高斯素数	87
exgcd	88
Pollard_Rho+Miller-Robin	88
FFT	
BSGS	
扩展 BSGS	
二次剩余	
解的数量	
勒让德符号	
欧拉判别准则	
Cipolla	
卡特兰数	
快速幂	
龟速乘快速幂(快速幂爆 longlong	
莫比乌斯反演	
莫比乌斯函数	
n 的所有约数的莫比乌斯的和	
反演	
博弈	
SG 定理:	
Nimk:	
anti_nim	
描述	
先手必胜条件	
高精度 GCD	
高精度乘法(FFT)	
高精度乘法(乘单精度	
高精度乘法(朴素)	
高精度除法(除单精度)	
高精度除法(除高精度)	
高精度加法	
高精度减法	
高精度阶乘	
高精度进制转换	
高精度幂	
高精度平方根	
高精度取模(对单精度)	
欧拉筛	
组合数(逆元线性递推	
中国剩余定理	
图论	

	有源汇上下界最大小流	122
	树链剖分	125
	虚树	128
	spfa 最短路及负环	130
	<u> </u>	131
	强连通(kosaraju	
	强连通(tarjan	
	强连通(tarjan 无 vector	
	最大流	
	最大流(double)	
	最小费用最大流	
	树分治	
	拓扑排序	
	最近公共祖先(倍增)	
	最近公共祖先(线段树)	
	KM	
	prufer 序列	
	朱刘算法	
	欧拉回路	
	点分树	
烂 州:	代数	
汉江	高斯消元	
	<u> </u>	
	发性基	
	· · · ·	
	线性基 2	
	矩阵(加减乘快速幂	
力, 元五	稀疏矩阵乘法	
	140027	
	mt19937	
	洗牌算法	
	快读	
	fread 快读	
	朝鲜大哥快读	
	模拟退火	
	整体二分	
字符	串	
	马拉车	
	AC 自动机	
	KMP	
	KMP 2	
	Tire.	
	后缀数组	
_,,,,	可持久化字典树	
对拍.	and the same and t	
	windows 环境下 bat 对拍	194

注意

优先队列是大的在前面 如果要小的 要重载 long long 二分答案的时候..精度 也有可能 爆 int (?

哈希 自然溢出 yyds 双哈希 输出限制..

匈牙利的复杂度常数非常小(...

递归爆栈 re

for i 进行计算的时候 (i 开 long long)

边界问题 各种 01 的特判

模 多模一点 都可以模 (

char 数组开小了也可能报错 tle 和 wa (

图是否连通 是否重边 是否自环

读题!!与或(

重点 重边

当保证 \mathbf{n} 的总和不会很大,但数据组数可能很多的时候,注意初始化造成的 \mathbf{tle} 问题(

pow() 的精度问题

unique erase 先排序

图论初始化!!

没开 longlong,中间有个判定条件爆了

他卡快排,由于答案不超过10000,可以计数排序

re: 没有开 longlong, (以为是 dfs 爆栈

定义

() (a,b)=1 最大公约数 即 a, b 互质

| 整除 a|b b%a==0

```
STL
```

```
优先队列重载
priority_queue<int, vector<int>, cmp>s;
struct cmp{
      bool operator()(const int &a,const int &b){
            return a>b;
      }
};
set 重载
#include <bits/stdc++.h>
using namespace std;
#define 1 first
#define r second
struct cmp{
      bool operator() (const pair <int, int> &a, const pair <int, int> &b)
const{
            int lena = a.r - a.l + 1;
            int lenb = b.r - b.l + 1;
            if(lena == lenb) return a.l < b.l;</pre>
            return lena > lenb;
      }
};
int main(){
      ios :: sync_with_stdio(∅); cin.tie(∅); cout.tie(∅);
      int T;
      cin >> T;
      while(T -- ){
            int n;
            cin >> n;
            set<pair<int, int>, cmp> segs;
            segs.insert(\{0, n - 1\});
            vector<int> a(n);
            for(int i = 1; i <= n; ++ i){
                  pair<int, int> cur = *segs.begin();
                   segs.erase(segs.begin());
                  int id = (cur.l + cur.r) / 2;
                  a[id] = i;
                   if(cur.l < id) segs.insert({cur.l, id - 1});</pre>
                   if(id < cur.r) segs.insert({id + 1, cur.r});</pre>
            }
```

```
for(auto it : a) cout << it << " ";</pre>
           cout << endl;</pre>
      }
}
动态开数组
int a[15], n, m;
cin >> n >> m;
int (*b)[m] = (int (*)[m])a;
new / delete
#define M 10U
#define N 20
第一种,可以直接[][]访问。但是内存不连续,不是很推荐使用,除非 M \ N 都不确定
//定义的时候
int** pNum;//以int 为例
pNum = new int*[M];
for(int i = 0;i < M;i ++){</pre>
      pNum[i]=new int[N];
}
//删除的时候是
for(int j = 0; j < M; j ++){</pre>
     delete []pNum[i];
}
delete []pNum;
malloc / free
#include<stdio.h>
#include<stdlib.h>
int main() {
   int **a; //用二级指针动态申请二维数组
   int i,j;
   int m,n;
   printf("请输入行数\n");
```

```
scanf("%d",&m);
   printf("请输入列数\n");
   scanf("%d",&n);
   a=(int**)malloc(sizeof(int*)*m);
   for(i=0;i<m;i++)</pre>
   a[i]=(int*)malloc(sizeof(int)*n);
   for(i=0;i<m;i++) {</pre>
       for(j=0;j<n;j++) {</pre>
                                       //输出每个元素地址,每行的列与列之
          printf("%p\n",&a[i][j]);
间的地址时连续的, 行与行之间的地址不连续
       }
   }
   for(i=0;i<m;i++)</pre>
   free(a[i]);
   free(a);
   return 0;
}
#include<stdio.h>
#include<stdlib.h>
int main()
{
   int i,j;
   //申请一个3 行2 列的整型数组
   int (*a)[2]=(int(*)[2])malloc(sizeof(int)*3*2);
   for(i=0;i<3;i++) {</pre>
       for(j=0;j<2;j++) {</pre>
          printf("%p\n",&a[i][j]); //输出数组每个元素地址,每个元素的地
址是连续的
       }
   }
   free(a);
   return 0;
}
vector
//二维 vector
vector<vector <int> > ivec(m ,vector<int>(n)); //m*n 的二维vector
// 动态创建 m*n 的二维 vector
//方法一:
vector<vector <int> > ivec;
ivec.resize(m);
for(int i=0;i<m;i++)</pre>
 ivec[i].resize(n);
```

//方法二:

vector<vector <int> > ivec; ivec.resize(m,vector<int>(n));

set

begin(),返回 set 容器的第一个元素

end(),返回 set 容器的最后一个元素

clear(),删除 set 容器中的所有的元素

empty(),判断 set 容器是否为空

max_size(),返回 set 容器可能包含的元素最大个数

size(),返回当前 set 容器中的元素个数

rbegin,返回的值和 end()相同

rend(),返回的值和 rbegin()相同

count() 用来查找 set 中某个某个键值出现的次数。

equal_range() ,返回一对定位器,分别表示第一个大于或等于给定关键值的元素和 第一个大于给定关键值的元素,这个返回值是一个 pair 类型,如果这一对定位器中哪个返回失败,就会等于 end()的值。

erase(iterator),删除定位器 iterator 指向的值

erase(first,second),删除定位器 first 和 second 之间的值

erase(key_value),删除键值 key_value 的值

find(), 返回给定值值得定位器,如果没找到则返回 end()。

insert(key_value);将 keyvalue 插入到 set 中 ,返回值是 pair<set::iterator,bool>,bool 标志着插入是否成功,而 iterator 代表插入的位置,若 keyvalue 已经在 set 中,则 iterator 表示的 key value 在 set 中的位置。

inset(first,second);将定位器 first 到 second 之间的元素插入到 set 中,返回值是 void.

lowerbound(keyvalue) ,返回第一个大于等于 key_value 的定位器

upperbound(keyvalue),返回最后一个大于等于 key_value 的定位器

map

```
插入操作
使用[]进行单个插入
map<int, string> ID Name;
// 如果已经存在键值 2015,则会作赋值修改操作,如果没有则插入
ID_Name[2015] = "Tom";1234
使用 insert 进行单个和多个插入 (insert 共有 4 个重载函数:
// 插入单个键值对,并返回插入位置和成功标志,插入位置已经存在值时,插入失败
pair<iterator,bool> insert (const value type& val);
//在指定位置插入,在不同位置插入效率是不一样的,因为涉及到重排
iterator insert (const_iterator position, const value_type& val);
// 插入多个
void insert (InputIterator first, InputIterator last);
//c++11 开始支持,使用列表插入多个
void insert (initializer_list<value_type> il);
取值
Map 中元素取值主要有 at 和[]两种操作, at 会作下标检查,而[]不会。
map<int, string> ID_Name;
//ID Name 中没有关键字 2016, 使用[]取值会导致插入
//因此,下面语句不会报错,但打印结果为空
cout<<ID_Name[2016].c_str()<<endl;</pre>
//使用at 会进行关键字检查,因此下面语句会报错
ID Name.at(2016) = "Bob";
容量查询
// 查询 map 是否为空
bool empty();
// 查询map 中键值对的数量
size_t size();
// 查询map 所能包含的最大键值对数量,和系统和应用库有关。
// 此外,这并不意味着用户一定可以存这么多,很可能还没达到就已经开辟内存失败了
size t max size();
```

```
// 查询关键字为 key 的元素的个数,在 map 里结果非 0 即 1
size t count( const Key& key ) const; //
迭代器
共有八个获取迭代器的函数: begin, end, rbegin, rend 以及对应的 cbegin, cend,
crbegin, crend.
二者的区别在于,后者一定返回 constiterator, 而前者则根据 map 的类型返回
iterator 或者 constiterator。const 情况下,不允许对值进行修改。如下面代码所
示:
map<int,int>::iterator it;
map<int,int> mmap;
const map<int,int> const mmap;
it = mmap.begin(); //iterator
mmap.cbegin(); //const iterator
const mmap.begin(); //const iterator
const_mmap.cbegin(); //const_iterator123456789
返回的迭代器可以进行加减操作,此外,如果 map 为空,则 begin = end。
删除
// 删除迭代器指向位置的键值对,并返回一个指向下一元素的迭代器
iterator erase( iterator pos )
// 删除一定范围内的元素,并返回一个指向下一元素的迭代器
iterator erase( const_iterator first, const_iterator last );
// 根据Key 来进行删除, 返回删除的元素数量,在 map 里结果非 0 即 1
size t erase( const key type& key );
// 清空map,清空后的size 为0
void clear();
交换
// 就是两个map 的内容互换
void swap( map& other );
顺序比较
// 比较两个关键字在 map 中位置的先后
key_compare key_comp() const;
```

查找

```
// 关键字查询,找到则返回指向该关键字的迭代器,否则返回指向 end 的迭代器
// 根据map 的类型,返回的迭代器为 iterator 或者 const iterator
iterator find (const key_type& k);
const iterator find (const key type& k) const;
操作符
operator: == != < <= > >=
注意 对于==运算符, 只有键值对以及顺序完全相等才算成立。
unordered map
查找元素是否存在
若有 unordered map <int, int> mp;查找 x 是否在 map 中
方法 1: 若存在 mp.find(x)!=mp.end()
方法 2: 若存在 mp.count(x)!=0123
插入数据
mp.insert(Map::value_type(1, "Raoul"));1
遍历 map
unordered_map<key,T>::iterator it;
   (*it).first; //the key value
   (*it).second //the mapped value
   for(unordered map<key,T>::iterator iter=mp.begin();iter!=mp.end();i
ter++)
        cout<<"key value is"<<iter->first<<" the mapped value is "<< it
er->second;
   //也可以这样
   for(auto& v : mp)
      print v.first and v.second
bitset
   C++的 bitset 在 bitset 头文件中,它是一种类似数组的结构,它的每一
   个元素只能是0或1,每个元素仅用1 bit 空间。
   bitset 数组与 vector 数组区别
   bitset 声明数组:bitset<100> number[10]
   vector 声明数组:vector number[10];
```

bitset<每个 bitset 元素的长度(没有占满前面全部自动补 0)> 元素

bitset 内置转化函数: 可将 bitset 转化为 string,unsigned long,unsigned long long。

```
构造
     bitset<4> bitset1; //无参构造,长度为4,默认每一位为0
   bitset<8> bitset2(12); //长度为8,二进制保存,前面用0补充
   string s = "100101";
                         //长度为10,前面用0补充
   bitset<10> bitset3(s);
   char s2[] = "10101";
   bitset<13> bitset4(s2); //长度为13,前面用0补充
   cout << bitset1 << endl; //0000</pre>
   cout << bitset2 << endl;</pre>
                          //00001100
   cout << bitset3 << endl;</pre>
                           //0000100101
   cout << bitset4 << endl;</pre>
                           //0000000010101
函数
     bitset<8> foo ("10011011");
   cout << foo.count() << endl; //5 (count 函数用来求bitset 中1的
位数,foo 中共有5个1
   cout << foo.size() << endl; //8 (size 函数用来求bitset 的大小,
一共有8位
   cout << foo.test(0) << endl; //true (test 函数用来查下标处的元素
是 0 还是 1 ,并返回 false 或 true,此处 foo[0]为 1 ,返回 true
   cout << foo.test(2) << endl; //false (同理, foo[2]为0,返回
false
   cout << foo.any() << endl; //true (any 函数检查 bitset 中是否有
                                       (none 函数检查bitset 中是否
   cout << foo.none() << endl; //false</pre>
没有1
   cout << foo.all() << endl; //false (all 函数检查 bitset 中是全部
为1
2019-2020 ICPC Asia Taipei-Hsinchu Regional Contest (H
```

```
Н
#include <bits/stdc++.h>
#define 11 long long
using namespace std;
int t,n,m;
char str[1010];
bitset<500> number[30];
int main() {
      ios::sync with stdio(false); cin.tie(0); cout.tie(0);
   //freopen("test.in","r",stdin);
   //freopen("test.out", "w", stdout);
      scanf("%d",&t);
      while(t--)
      {
            scanf("%d %d",&n,&m);
            for(int i=0;i<m;i++)</pre>
                  scanf("%s",str);
                  number[i]=bitset<500>(str);
            int len=1<<m,ans=m+1;</pre>
            for(int i=1;i<len;i++)</pre>
                  int t=i,s=0;
                  bitset<500> num(0);
                  for(int j=0;j<m&&t>0;j++)
                        if(t&1)
                              num=num|number[j];
                              S++;
                        t>>=1;
                  if(num.count()==n) ans=min(ans,s);
            if(ans==m+1) printf("-1\n");
            else printf("%d\n",ans);
      return 0;
}
计算几何
几何的一些定理
多面体欧拉定理
```

多面体欧拉定理是指对于简单多面体,其各维对象数总满足一定的数学关系,在三维空间中多面体欧拉定理可表示为:

"顶点数-棱长数+表面数=2"。

}

简单多面体即表面经过连续变形可以变为球面的多面体。

```
zyx 的计算几何
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const int N = 1e6 + 10;
const double eps = 1e-9;
const double PI = acos(-1.0);
const double dinf = 1e99;
const 11 inf = 0x3f3f3f3f3f3f3f3f3f;
struct Line;
struct Point {
   double x, y;
   Point() { x = y = 0; }
   Point(const Line &a);
   Point(const double &a, const double &b) : x(a), y(b) {}
   Point operator+(const Point &a) const {
       return \{x + a.x, y + a.y\};
```

```
Point operator-(const Point &a) const {
       return {x - a.x, y - a.y};
   }
   Point operator*(const double &a) const {
       return {x * a, y * a};
   }
   Point operator/(const double &d) const {
       return \{x / d, y / d\};
   }
   bool operator==(const Point &a) const {
       return abs(x - a.x) + abs(y - a.y) < eps;
   }
   void standardize() {
       *this = *this / sqrt(x * x + y * y);
   }
double norm(const Point &p) { return p.x * p.x + p.y * p.y; }
Point orth(const Point &a) { return Point(-a.y, a.x); }
double dist(const Point &a, const Point &b) {
   return sqrt((a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y));
```

};

```
}
double dist2(const Point &a, const Point &b) {
   return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y);
}
struct Line {
   Point s, t;
   Line() {}
   Line(const Point &a, const Point &b) : s(a), t(b) {}
};
struct Circle {
   Point o;
   double r;
   Circle() {}
   Circle(Point P, double R = 0) { o = P, r = R; }
};
double length(const Point &p) {
   return sqrt(p.x * p.x + p.y * p.y);
}
```

```
double length(const Line &1) {
   Point p(1);
   return length(p);
}
Point::Point(const Line &a) { *this = a.t - a.s; }
istream &operator>>(istream &in, Point &a) {
   in >> a.x >> a.y;
   return in;
}
ostream &operator<<(ostream &out, Point &a) {</pre>
   out << fixed << setprecision(10) << a.x << ' ' << a.y;</pre>
   return out;
}
double dot(const Point &a, const Point &b) { return a.x * b.x + a.y * b.
y; }
double det(const Point &a, const Point &b) { return a.x * b.y - a.y * b.
x; }
int sgn(const double &x) { return fabs(x) < eps ? 0 : (x > 0 ? 1 : -1); }
double sqr(const double &x) { return x * x; }
Point rotate(const Point &a, const double &ang) {
```

```
double x = cos(ang) * a.x - sin(ang) * a.y;
   double y = sin(ang) * a.x + cos(ang) * a.y;
   return {x, y};
}
//点在线段上 <=0 包含端点
bool sp_on(const Line &seg, const Point &p) {
   Point a = seg.s, b = seg.t;
   return !sgn(det(p - a, b - a)) && sgn(dot(p - a, p - b)) <= 0;
}
bool lp_on(const Line &line, const Point &p) {
   Point a = line.s, b = line.t;
   return !sgn(det(p - a, b - a));
}
//等于不包含共线
int andrew(Point *point, Point *convex, int n) {
   sort(point, point + n, [](Point a, Point b) {
       if (a.x != b.x) return a.x < b.x;
       return a.y < b.y;</pre>
   });
   int top = 0;
   for (int i = 0; i < n; i++) {</pre>
       while ((top > 1) && det(convex[top - 1] - convex[top - 2], point
[i] - convex[top - 1]) <= ∅)
           top--;
       convex[top++] = point[i];
   }
```

```
int tmp = top;
   for (int i = n - 2; i >= 0; i--) {
       while ((top > tmp) && det(convex[top - 1] - convex[top - 2], poin
t[i] - convex[top - 1]) <= 0
           top--;
       convex[top++] = point[i];
   }
   if (n > 1) top--;
   return top;
}
double slope(const Point &a, const Point &b) { return (a.y - b.y) / (a.x
 - b.x); }
double slope(const Line &a) { return slope(a.s, a.t); }
Point 11 intersection(const Line &a, const Line &b) {
   double s1 = det(Point(a), b.s - a.s), s2 = det(Point(a), b.t - a.s);
   if (sgn(s1) == 0 && sgn(s2) == 0) return a.s;
   return (b.s * s2 - b.t * s1) / (s2 - s1);
}
int ss_cross(const Line &a, const Line &b, Point &p) {
   int d1 = sgn(det(a.t - a.s, b.s - a.s));
   int d2 = sgn(det(a.t - a.s, b.t - a.s));
   int d3 = sgn(det(b.t - b.s, a.s - b.s));
   int d4 = sgn(det(b.t - b.s, a.t - b.s));
   if ((d1 ^ d2) == -2 && (d3 ^ d4) == -2) {
       p = ll_intersection(a, b);
```

```
return 1;
   }
   if (!d1 && sp_on(a, b.s)) {
       p = b.s;
       return 2;
   }
   if (!d2 && sp_on(a, b.t)) {
       p = b.t;
       return 2;
   }
   if (!d3 && sp_on(b, a.s)) {
       p = a.s;
       return 2;
   }
   if (!d4 && sp_on(b, a.t)) {
       p = a.t;
       return 2;
   }
   return 0;
}
int ccw(const Point &a, Point b, Point c) {
   b = b - a, c = c - a;
   if (sgn(det(b, c)) > 0) return +1; // "COUNTER_CLOCKWISE"
   if (sgn(det(b, c)) < 0) return -1; // "CLOCKWISE"</pre>
   if (sgn(dot(b, c)) < 0) return +2;  // "ONLINE_BACK"</pre>
   if (sgn(norm(b) - norm(c)) < 0) return -2; // "ONLINE_FRONT"</pre>
```

```
return 0;
                                   // "ON SEGMENT"
}
Point project(const Line &1, const Point &p) {
   Point base(1);
   double r = dot(base, p - 1.s) / sqr(length(base));
   return 1.s + (base * r);
}
double sp_dist(const Line &l, const Point &p) {
   if (l.s == l.t) return dist(l.s, p);
   Point x = p - 1.s, y = p - 1.t, z = 1.t - 1.s;
   if (sgn(dot(x, z)) < 0)return length(x);//P 距离A 更近
   if (sgn(dot(y, z)) > 0)return length(y);//P 距离 B 更近
   return abs(det(x, z) / length(z));//面积除以底边长
}
double lp_dist(const Line &l, const Point &p) {
   Point x = p - 1.s, y = p - 1.t, z = 1.t - 1.s;
   return abs(det(x, z) / length(z));//面积除以底边长
}
int cl_cross(const Circle &c, const Line &l, pair<Point, Point> &ans) {
   Point a = c.o;
   double r = c.r;
   Point pr = project(1, a);
   double dis = dist(pr, a);
```

```
double tmp = r * r - dis * dis;
   if (sgn(tmp) == 1) {
       double base = sqrt(max(0.0, r * r - dis * dis));
       Point e(1);
       e.standardize();
       e = e * base;
       ans = make_pair(pr + e, pr - e);
       return 2;
   } else if (sgn(tmp) == 0) {
       ans = make_pair(pr, pr);
       return 1;
   } else return 0;
}
int intersectCS(Circle c, Line 1) {//交点个数,下面 cs_cross 用到
   if (sgn(norm(project(1, c.o) - c.o) - c.r * c.r) > 0) return 0;
   double d1 = length(c.o - l.s), d2 = length(c.o - l.t);
   if (sgn(d1 - c.r) <= 0 && sgn(d2 - c.r) <= 0) return 0;
   if ((sgn(d1 - c.r) < 0 \&\& sgn(d2 - c.r) > 0) || (sgn(d1 - c.r) > 0 \&\&
 sgn(d2 - c.r) < 0)) return 1;
   Point h = project(1, c.o);
   if (dot(1.s - h, 1.t - h) < 0) return 2;
   return 0;
}
int cs_cross(Circle c, Line s, pair<Point, Point> &ans) {//圆和线段交点
   Line l(s);
   int num = cl_cross(c, 1, ans);
   int res = intersectCS(c, s);
```

```
if (res == 2) return 2;
   if (num > 1) {
       if (dot(l.s - ans.first, l.t - ans.first) > 0) swap(ans.first, an
s.second);
       ans.second = ans.first;
   }
   return res;
}
int cc_cross(const Circle &cir1, const Circle &cir2, pair<Point, Point>
&ans) {
   const Point &c1 = cir1.o, &c2 = cir2.o;
   const double &r1 = cir1.r, &r2 = cir2.r;
   double x1 = c1.x, x2 = c2.x, y1 = c1.y, y2 = c2.y;
   double d = length(c1 - c2);
   if (sgn(fabs(r1 - r2) - d) > 0) return 0; //内含
   if (sgn(r1 + r2 - d) < 0) return 4; //相离
   double a = r1 * (x1 - x2) * 2, b = r1 * (y1 - y2) * 2, c = r2 * r2 - r
1 * r1 - d * d;
   double p = a * a + b * b, q = -a * c * 2, r = c * c - b * b;
   double cosa, sina, cosb, sinb;
   //One Intersection
   if (sgn(d - (r1 + r2)) == 0 \mid | sgn(d - fabs(r1 - r2)) == 0) {
       cosa = -q / p / 2;
       sina = sqrt(1 - sqr(cosa));
       Point p0(x1 + r1 * cosa, y1 + r1 * sina);
       if (sgn(dist(p0, c2) - r2)) p0.y = y1 - r1 * sina;
```

```
ans = pair<Point, Point>(p0, p0);
       if (sgn(r1 + r2 - d) == 0) return 3;
                                             //外切
       else return 1; //内切
   }
   //Two Intersections
   double delta = sqrt(q * q - p * r * 4);
   cosa = (delta - q) / p / 2;
   cosb = (-delta - q) / p / 2;
   sina = sqrt(1 - sqr(cosa));
   sinb = sqrt(1 - sqr(cosb));
   Point p1(x1 + r1 * cosa, y1 + r1 * sina);
   Point p2(x1 + r1 * cosb, y1 + r1 * sinb);
   if (sgn(dist(p1, c2) - r2)) p1.y = y1 - r1 * sina;
   if (sgn(dist(p2, c2) - r2)) p2.y = y1 - r1 * sinb;
   if (p1 == p2) p1.y = y1 - r1 * sina;
   ans = pair<Point, Point>(p1, p2);
   return 2; // 相交
Point lp sym(const Line &1, const Point &p) {
   return p + (project(1, p) - p) * 2;
double alpha(const Point &t1, const Point &t2) {
   double theta;
   theta = atan2((double) t2.y, (double) t2.x) - atan2((double) t1.y,
(double) t1.x);
   if (sgn(theta) < 0)</pre>
       theta += 2.0 * PI;
```

}

}

```
return theta;
}
 int pip(const Point *P, const int &n, const Point &a) {// 【射线法】判断点
A 是否在任意多边形 Poly 以内
            int cnt = 0;
           double tmp;
           for (int i = 1; i <= n; ++i) {</pre>
                      int j = i < n ? i + 1 : 1;
                      if (sp_on(Line(P[i], P[j]), a))return 2;//点在多边形上
                       if (a.y >= min(P[i].y, P[j].y) && a.y < max(P[i].y, P[j].y))//纵
 坐标在该线段两端点之间
                                 tmp = P[i].x + (a.y - P[i].y) / (P[j].y - P[i].y) * (P[j].x - P[i].y) * (P[j].x - P[i].y) * (P[j].x - P[i].y) * (P[j].y) * (P[j].y
   P[i].x), cnt += sgn(tmp - a.x) > 0;//交点在A 右方
            }
           return cnt & 1;//穿过奇数次则在多边形以内
 }
 bool pip convex jud(const Point &a, const Point &L, const Point &R) {//
 判断AL 是否在AR 右边
           return sgn(det(L - a, R - a)) > 0;//必须严格以内
 }
 bool pip convex(const Point *P, const int &n, const Point &a) {//【二分
 法】判断点A是否在凸多边形Poly以内
           //点按逆时针给出
            if (pip_convex_jud(P[0], a, P[1]) || pip_convex_jud(P[0], P[n - 1],
a)) return 0;//在P[0_1]或P[0_n-1]外
            if (sp_on(Line(P[0], P[1]), a) || sp_on(Line(P[0], P[n - 1]), a)) re
turn 2;//在P[0_1]或P[0_n-1]上
```

```
int l = 1, r = n - 2;
   while (1 < r) {//二分找到一个位置 pos 使得 P[0] A 在 P[0 pos], P[0 (pos+
1) 1之间
      int mid = (1 + r + 1) >> 1;
      if (pip_convex_jud(P[0], P[mid], a))l = mid;
      else r = mid - 1;
   }
   if (pip convex jud(P[1], a, P[1 + 1]))return 0;//在P[pos (pos+1)]外
   if (sp_on(Line(P[1], P[1 + 1]), a)) return 2; // #EP[pos_os+1)] \bot
   return 1;
}
// 多边形是否包含线段
// 因此我们可以先求出所有和线段相交的多边形的顶点,然后按照X-Y 坐标排序(X 坐标
小的排在前面,对于X坐标相同的点,Y坐标小的排在前面,
// 这种排序准则也是为了保证水平和垂直情况的判断正确),这样相邻的两个点就是在线
段上相邻的两交点,如果任意相邻两点的中点也在多边形内,
// 则该线段一定在多边形内。
int pp judge(Point *A, int n, Point *B, int m) {// 【判断多边形 A 与多边形 B
是否相离】
   for (int i1 = 1; i1 <= n; ++i1) {</pre>
      int j1 = i1 < n ? i1 + 1 : 1;
      for (int i2 = 1; i2 <= m; ++i2) {</pre>
          int j2 = i2 < m ? i2 + 1 : 1;
          Point tmp;
          if (ss_cross(Line(A[i1], A[j1]), Line(B[i2], B[j2]), tmp)) re
turn 0://两线段相交
          if (pip(B, m, A[i1]) || pip(A, n, B[i2]))return 0;//点包含在内
      }
```

```
}
   return 1;
}
double area(Point *P, int n) {//【任意多边形 P 的面积】
   double S = 0;
   for (int i = 0; i < n; i++) S += det(P[i], P[(i + 1) % n]);</pre>
   return S * 0.5;
}
double pc area(Point *p, int n, const Circle &c) {
   if (n < 3) return 0;
   function<double(Circle, Point, Point)> dfs = [&](Circle c, Point a,
Point b) {
       Point va = c.o - a, vb = c.o - b;
       double f = det(va, vb), res = 0;
       if (sgn(f) == 0) return res;
       if (sgn(max(length(va), length(vb)) - c.r) <= 0) return f;</pre>
       Point d(dot(va, vb), det(va, vb));
       if (sgn(sp_dist(Line(a, b), c.o) - c.r) >= 0) return c.r * c.r *
atan2(d.y, d.x);
       pair<Point, Point> u;
       int cnt = cs_cross(c, Line(a, b), u);
       if (cnt == 0) return res;
       if (cnt > 1 && sgn(dot(u.second - u.first, a - u.first)) > 0) swa
p(u.first, u.second);
       res += dfs(c, a, u.first);
       if (cnt == 2) res += dfs(c, u.first, u.second) + dfs(c, u.second,
b);
       else if (cnt == 1) res += dfs(c, u.first, b);
```

```
return res;
   };
   double res = 0;
   for (int i = 0; i < n; i++) {</pre>
       res += dfs(c, p[i], p[(i + 1) % n]);
   }
   return res * 0.5;
}
Line Q[N];
int judge(Line L, Point a) { return sgn(det(a - L.s, L.t - L.s)) > 0; }/
/判断点 a 是否在直线 L 的右边
int halfcut(Line *L, int n, Point *P) {//【半平面交】
   sort(L, L + n, [](const Line &a, const Line &b) {
       double d = atan2((a.t - a.s).y, (a.t - a.s).x) - atan2((b.t - b.
s).y, (b.t - b.s).x);
       return sgn(d) ? sgn(d) < 0 : judge(a, b.s);</pre>
   });
   int m = n;
   n = 0;
   for (int i = 0; i < m; ++i)</pre>
       if (i == 0 || sgn(atan2(Point(L[i]).y, Point(L[i]).x) - atan2(Po
int(L[i-1]).y, Point(L[i-1]).x)))
           L[n++] = L[i];
   int h = 1, t = 0;
   for (int i = 0; i < n; ++i) {</pre>
       while (h < t \&\& judge(L[i], ll_intersection(Q[t], Q[t - 1]))) --t;
// 当队尾两个直线交点不是在直线 L[i] 上或者左边时就出队
```

```
while (h < t \&\& judge(L[i], ll_intersection(Q[h], Q[h + 1]))) ++h;
// 当队头两个直线交点不是在直线 L[i] 上或者左边时就出队
       Q[++t] = L[i];
   }
   while (h < t \&\& judge(Q[h], ll_intersection(Q[t], Q[t - 1]))) --t;
   while (h < t \&\& judge(Q[t], ll_intersection(Q[h], Q[h + 1]))) ++h;
   n = 0;
   for (int i = h; i <= t; ++i) {</pre>
       P[n++] = 11 intersection(Q[i], Q[i < t ? i + 1 : h]);
   }
   return n;
}
Point V1[N], V2[N];
int mincowski(Point *P1, int n, Point *P2, int m, Point *V) {//【闵可夫斯
基和】求两个凸包{P1},{P2}的向量集合{V}={P1+P2}构成的凸包
   for (int i = 0; i < n; ++i) V1[i] = P1[(i + 1) \% n] - P1[i];
   for (int i = 0; i < m; ++i) V2[i] = P2[(i + 1) \% m] - P2[i];
   int t = 0, i = 0, j = 0;
   V[t++] = P1[0] + P2[0];
   while (i < n \&\& j < m) V[t] = V[t - 1] + (sgn(det(V1[i], V2[j])) > 0?
V1[i++] : V2[j++]), t++;
   while (i < n) V[t] = V[t - 1] + V1[i++], t++;
   while (j < m) V[t] = V[t - 1] + V2[j++], t++;
   return t;
}
```

```
Circle external circle(const Point &A, const Point &B, const Point &C)
{//【三点确定一圆】向量垂心法
             Point P1 = (A + B) * 0.5, P2 = (A + C) * 0.5;
             Line R1 = Line(P1, P1 + orth(B - A));
             Line R2 = Line(P2, P2 + orth(C - A));
             Circle 0;
             0.o = 11_intersection(R1, R2);
             0.r = length(A - 0.o);
             return 0;
}
Circle internal circle(const Point &A, const Point &B, const Point &C)
             double a = dist(B, C), b = dist(A, C), c = dist(A, B);
             double s = (a + b + c) / 2;
             double S = sqrt(max(0.0, s * (s - a) * (s - b) * (s - c)));
             double r = S / s;
             return Circle((A * a + B * b + C * c) / (a + b + c), r);
}
struct ConvexHull {
             int op;
             struct cmp {
                          bool operator()(const Point &a, const Point &b) const {
                                        return sgn(a.x - b.x) < 0 \mid | sgn(a.x - b.x) == 0 && sgn(a.y - b.x) == 0 & sgn(a.y - b
   b.y) < 0;
```

```
}
   };
   set<Point, cmp> s;
   ConvexHull(int o) {
       op = o;
       s.clear();
   }
   inline int PIP(Point P) {
       set<Point>::iterator it = s.lower bound(Point(P.x, -dinf));//救
到第一个横坐标大于P 的点
       if (it == s.end())return 0;
       if (sgn(it\rightarrow x - P.x) == 0) return sgn((P.y - it\rightarrow y) * op) <= 0;//
比较纵坐标大小
       if (it == s.begin())return 0;
       set<Point>::iterator j = it, k = it;
       --j;
       return sgn(det(P - *j, *k - *j) * op) >= 0;//看叉姬1
   }
   inline int judge(set<Point>::iterator it) {
       set<Point>::iterator j = it, k = it;
       if (j == s.begin())return 0;
       --j;
       if (++k == s.end())return 0;
       return sgn(det(*it - *j, *k - *j) * op) >= 0;//看叉姬
   }
```

```
inline void insert(Point P) {
       if (PIP(P))return;//如果点P已经在凸壳上或凸包里就不插入了
       set<Point>::iterator tmp = s.lower_bound(Point(P.x, -dinf));
       if (tmp != s.end() && sgn(tmp->x - P.x) == 0)s.erase(tmp);//特判
横坐标相等的点要去掉
       s.insert(P);
       set<Point>::iterator it = s.find(P), p = it;
       if (p != s.begin()) {
          --p;
          while (judge(p)) {
              set<Point>::iterator temp = p--;
              s.erase(temp);
          }
       }
       if ((p = ++it) != s.end()) {
          while (judge(p)) {
              set<Point>::iterator temp = p++;
              s.erase(temp);
          }
       }
   }
} up(1), down(-1);
int PIC(Circle C, Point a) { return sgn(length(a - C.o) - C.r) <= 0; }//</pre>
判断点A 是否在圆C 内
void Random(Point *P, int n) { for (int i = 0; i < n; ++i)swap(P[i], P</pre>
[(rand() + 1) % n]); }//随机一个排列
Circle min_circle(Point *P, int n) {//【求点集 P 的最小覆盖圆】 O(n)
```

```
// random shuffle(P,P+n);
    Random(P, n);
   Circle C = Circle(P[0], 0);
   for (int i = 1; i < n; ++i)</pre>
       if (!PIC(C, P[i])) {
           C = Circle(P[i], 0);
           for (int j = 0; j < i; ++j)</pre>
               if (!PIC(C, P[j])) {
                   C.o = (P[i] + P[j]) * 0.5, C.r = length(P[j] - C.o);
                   for (int k = 0; k < j; ++k) if (!PIC(C, P[k])) C = ext
ernal_circle(P[i], P[j], P[k]);
               }
        }
    return C;
}
int temp[N];
double closest_point(Point *p, int n) {
    function<double(int, int)> merge = [&](int 1, int r) {
       double d = dinf;
       if (1 == r) return d;
       if (1 + 1 == r) return dist(p[1], p[r]);
       int mid = (1 + r) >> 1;
       double d1 = merge(1, mid);
       double d2 = merge(mid + 1, r);
       d = min(d1, d2);
       int i, j, k = 0;
```

```
for (i = 1; i <= r; i++) {</pre>
           if (sgn(abs(p[mid].x - p[i].x) - d) <= 0)
               temp[k++] = i;
       }
       sort(temp, temp + k, [&](const int &a, const int &b) {
           return sgn(p[a].y - p[b].y) < 0;</pre>
       });
       for (i = 0; i < k; i++) {
           for (j = i + 1; j < k \&\& sgn(p[temp[j]].y - p[temp[i]].y - d)
<= 0; j++) {
               double d3 = dist(p[temp[i]], p[temp[j]]);
               d = min(d, d3);
           }
       }
       return d;
   };
   sort(p, p + n, [&](const Point &a, const Point &b) {
       if (sgn(a.x - b.x) == 0) return sgn(a.y - b.y) < 0;
       else return sgn(a.x - b.x) < 0;
   });
   return merge(0, n - 1);
}
int tangent(const Circle &c1, const Point &p2, pair<Point, Point> &ans)
{ //圆和点切线
   Point tmp = c1.o - p2;
   int sta;
   if (sgn(norm(tmp) - c1.r * c1.r) < 0) return 0;</pre>
```

```
else if (sgn(norm(tmp) - c1.r * c1.r) == 0) sta = 1;
   else sta = 2;
   Circle c2 = Circle(p2, sqrt(max(0.0, norm(tmp) - c1.r * c1.r)));
   cc_cross(c1, c2, ans);
   return sta;
}
int tangent(Circle c1, Circle c2, vector<Line> &ans) { //圆和点切线
   ans.clear();
   if (sgn(c1.r - c2.r) < 0) swap(c1, c2);</pre>
   double g = norm(c1.o - c2.o);
   if (sgn(g) == 0) return 0;
   Point u = (c2.o - c1.o) / sqrt(g);
   Point v = orth(u);
   for (int s = 1; s >= -1; s -= 2) {
       double h = (c1.r + s * c2.r) / sqrt(g);
       if (sgn(1 - h * h) == 0) {
           ans.push_back(Line(c1.o + u * c1.r, c1.o + (u + v) * c1.r));
       } else if (sgn(1 - h * h) >= 0) {
           Point uu = u * h, vv = v * sqrt(1 - h * h);
           ans.push_back(Line(c1.o + (uu + vv) * c1.r, c2.o - (uu + vv)
* c2.r * s));
           ans.push_back(Line(c1.o + (uu - vv) * c1.r, c2.o - (uu - vv)
* c2.r * s));
       }
   }
   return ans.size();
}
```

```
double areaofCC(Circle c1, Circle c2) { //两圆面积交
   if (c1.r > c2.r) swap(c1, c2);
   double nor = norm(c1.o - c2.o);
   double dist = sqrt(max(0.0, nor));
   if (sgn(c1.r + c2.r - dist) <= 0) return 0;</pre>
   if (sgn(dist + c1.r - c2.r) <= 0) return c1.r * c1.r * PI;</pre>
   double val;
   val = (nor + c1.r * c1.r - c2.r * c2.r) / (2 * c1.r * dist);
   val = max(val, -1.0), val = min(val, 1.0);
   double theta1 = acos(val);
   val = (nor + c2.r * c2.r - c1.r * c1.r) / (2 * c2.r * dist);
   val = max(val, -1.0), val = min(val, 1.0);
   double theta2 = acos(val);
   return (theta1 - sin(theta1 + theta1) * 0.5) * c1.r * c1.r + (theta2
- sin(theta2 + theta2) * 0.5) * c2.r * c2.r;
}
int convexCut(Point *p, Point *ans, int n, Line 1) {
   int top = 0;
   for (int i = 0; i < n; i++) {</pre>
       Point a = p[i], b = p[(i + 1) \% n];
       if (ccw(l.s, l.t, a) != -1) ans[top++] = a;
       if (ccw(l.s, l.t, a) * ccw(l.s, l.t, b) < 0)
           ans[top++] = ll_intersection(Line(a, b), l);
```

```
}
   return top;
}
double SphereCross(double d, double r1, double r2) {
   if (r1 < r2) swap(r1, r2);
   if (sgn(d - r1 - r2) >= 0) return 0;
   if (sgn(d + r2 - r1) <= 0) return 4.0 / 3 * PI * r2 * r2 * r2;</pre>
   double co = (r1 * r1 + d * d - r2 * r2) / (2.0 * d * r1);
   double h = r1 * (1 - co);
   double ans = (1.0 / 3) * PI * (3.0 * r1 - h) * h * h;
   co = (r2 * r2 + d * d - r1 * r1) / (2.0 * d * r2);
   h = r2 * (1 - co);
   ans += (1.0 / 3) * PI * (3.0 * r2 - h) * h * h;
   return ans;
}
计算几何全家桶
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 N = 1 << 20;
const 11 \mod = 1e9 + 7;
const double dinf = 1e99;
const int inf = 0x3f3f3f3f;
const 11 linf = 0x3f3f3f3f3f3f3f3f3;
const double eps = 1e-9;
const double PI = acos(-1.0);
struct Line;
```

struct Point {

double x, y;

```
Point() { x = y = 0; }
   Point(const Line &a);
   Point(const double &a, const double &b) : x(a), y(b) {}
   Point operator+(const Point &a) const {
       return {x + a.x, y + a.y};
   Point operator-(const Point &a) const {
       return {x - a.x, y - a.y};
   }
   Point operator*(const double &a) const {
       return {x * a, y * a};
   Point operator/(const double &d) const {
       return \{x / d, y / d\};
   }
   bool operator==(const Point &a) const {
       return abs(x - a.x) + abs(y - a.y) < eps;
   }
   void standardize() {
       *this = *this / sqrt(x * x + y * y);
};
Point normal(const Point &a) { return Point(-a.y, a.x); }
double dist(const Point &a, const Point &b) {
   return sqrt((a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y));
}
double dist2(const Point &a, const Point &b) {
   return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y);
}
struct Line {
   Point s, t;
   Line() {}
   Line(const Point &a, const Point &b) : s(a), t(b) {}
```

```
};
struct circle {
   Point o;
   double r;
   circle() {}
   circle(Point P, double R = 0) { o = P, r = R; }
};
double length(const Point &p) {
   return sqrt(p.x * p.x + p.y * p.y);
}
double length(const Line &1) {
   Point p(1);
   return length(p);
}
Point::Point(const Line &a) { *this = a.t - a.s; }
istream &operator>>(istream &in, Point &a) {
   in >> a.x >> a.y;
   return in;
}
double dot(const Point &a, const Point &b) {
   return a.x * b.x + a.y * b.y;
}
double det(const Point &a, const Point &b) {
   return a.x * b.y - a.y * b.x;
}
int sgn(const\ double\ \&x) \{ return\ fabs(x) < eps ? 0 : (x > 0 ? 1 : -1); \}
double sqr(const double &x) { return x * x; }
Point rotate(const Point &a, const double &ang) {
   double x = cos(ang) * a.x - sin(ang) * a.y;
   double y = sin(ang) * a.x + cos(ang) * a.y;
   return {x, y};
}
//点在线段上 <=0 包含端点
bool sp on(const Line &seg, const Point &p) {
   Point a = seg.s, b = seg.t;
   return !sgn(det(p - a, b - a)) && sgn(dot(p - a, p - b)) <= 0;</pre>
```

```
}
bool lp on(const Line &line, const Point &p) {
   Point a = line.s, b = line.t;
   return !sgn(det(p - a, b - a));
}
//等于不包含共线
int andrew(Point *point, Point *convex, int n) {
   sort(point, point + n, [](Point a, Point b) {
       if (a.x != b.x) return a.x < b.x;
       return a.y < b.y;</pre>
   });
   int top = 0;
   for (int i = 0; i < n; i++) {</pre>
       while ((top > 1) && det(convex[top - 1] - convex[top - 2], point
[i] - convex[top - 1] <= 0)
           top--;
       convex[top++] = point[i];
   int tmp = top;
   for (int i = n - 2; i >= 0; i--) {
       while ((top > tmp) && det(convex[top - 1] - convex[top - 2], poin
t[i] - convex[top - 1]) <= 0
           top--;
       convex[top++] = point[i];
   if (n > 1) top--;
   return top;
}
double slope(const Point &a, const Point &b) {
   return (a.y - b.y) / (a.x - b.x);
}
double slope(const Line &a) {
   return slope(a.s, a.t);
}
Point ll_intersection(const Line &a, const Line &b) {
   double s1 = det(Point(a), b.s - a.s), s2 = det(Point(a), b.t - a.s);
   return (b.s * s2 - b.t * s1) / (s2 - s1);
}
int ss_cross(const Line &a, const Line &b, Point &p) {
   int d1 = sgn(det(a.t - a.s, b.s - a.s));
   int d2 = sgn(det(a.t - a.s, b.t - a.s));
   int d3 = sgn(det(b.t - b.s, a.s - b.s));
   int d4 = sgn(det(b.t - b.s, a.t - b.s));
```

```
if ((d1 ^ d2) == -2 && (d3 ^ d4) == -2) {
       p = 11 intersection(a, b);
       return 1;
   if (!d1 && sp_on(a, b.s)) {
       p = b.s;
       return 2;
   if (!d2 && sp_on(a, b.t)) {
       p = b.t;
       return 2;
   if (!d3 && sp_on(b, a.s)) {
       p = a.s;
       return 2;
   if (!d4 && sp_on(b, a.t)) {
       p = a.t;
       return 2;
   return 0;
}
Point project(const Line &l, const Point &p) {
   Point base(1);
   double r = dot(base, p - 1.s) / sqr(length(base));
   return 1.s + (base * r);
}
double sp_dist(const Line &1, const Point &p) {
   if (l.s == l.t) return dist(l.s, p);
   Point x = p - 1.s, y = p - 1.t, z = 1.t - 1.s;
   if (sgn(dot(x, z)) < 0)return length(x);//P 距离 A 更近
   if (sgn(dot(y, z)) > 0)return length(y);//P 距离B 更近
   return abs(det(x, z) / length(z));//面积除以底边长
}
double lp dist(const Line &1, const Point &p) {
   Point x = p - 1.s, y = p - 1.t, z = 1.t - 1.s;
   return abs(det(x, z) / length(z));//面积除以底边长
}
int 1c cross(const Line &1, const Point &a, const double &r, pair<Point,
Point> &ans) {
   int num = 0;
   Point pr = project(1, a);
   double dis = dist(pr, a);
   double tmp = r * r - dis * dis;
   if (sgn(tmp) == 1) num = 2;
```

```
else if (sgn(tmp) == 0) num = 1;
   else return ∅;
   double base = sqrt(r * r - dis * dis);
   Point e(1);
   e.standardize();
   e = e * base;
   ans = make_pair(pr + e, pr - e);
   return num;
}
int cc_cross(const Point &c1, const double &r1, const Point &c2, const d
ouble &r2, pair<Point, Point> &ans) {
   double x1 = c1.x, x2 = c2.x, y1 = c1.y, y2 = c2.y;
   double d = length(c1 - c2);
   if (sgn(fabs(r1 - r2) - d) > 0) return -1; //内含
   if (sgn(r1 + r2 - d) < 0) return 0; //相离
   double a = r1 * (x1 - x2) * 2, b = r1 * (y1 - y2) * 2, c = r2 * r2 - r
1 * r1 - d * d;
   double p = a * a + b * b, q = -a * c * 2, r = c * c - b * b;
   double cosa, sina, cosb, sinb;
   //One Intersection
   if (sgn(d - (r1 + r2)) == 0 \mid | sgn(d - fabs(r1 - r2)) == 0) {
       cosa = -q / p / 2;
       sina = sqrt(1 - sqr(cosa));
       Point p0(x1 + r1 * cosa, y1 + r1 * sina);
       if (sgn(dist(p0, c2) - r2)) p0.y = y1 - r1 * sina;
       ans = pair<Point, Point>(p0, p0);
       return 1;
   //Two Intersections
   double delta = sqrt(q * q - p * r * 4);
   cosa = (delta - q) / p / 2;
   cosb = (-delta - q) / p / 2;
   sina = sqrt(1 - sqr(cosa));
   sinb = sqrt(1 - sqr(cosb));
   Point p1(x1 + r1 * cosa, y1 + r1 * sina);
   Point p2(x1 + r1 * cosb, y1 + r1 * sinb);
   if (sgn(dist(p1, c2) - r2)) p1.y = y1 - r1 * sina;
   if (sgn(dist(p2, c2) - r2)) p2.y = y1 - r1 * sinb;
   if (p1 == p2) p1.y = y1 - r1 * sina;
   ans = pair<Point, Point>(p1, p2);
   return 2;
}
Point lp_sym(const Line &1, const Point &p) {
   return p + (project(1, p) - p) * 2;
}
```

```
double alpha(const Point &t1, const Point &t2) {
   double theta;
   theta = atan2((double) t2.y, (double) t2.x) - atan2((double) t1.y,
(double) t1.x);
   if (sgn(theta) < 0)</pre>
       theta += 2.0 * PI;
   return theta;
}
int pip(const Point *P, const int &n, const Point &a) {// 【射线法】判断点
A 是否在任意多边形 Polv 以内
   int cnt = 0;
   int tmp;
   for (int i = 1; i <= n; ++i) {</pre>
       int j = i < n ? i + 1 : 1;
       if (sp_on(Line(P[i], P[j]), a))return 2;//点在多边形上
       if (a.y >= min(P[i].y, P[j].y) && a.y < max(P[i].y, P[j].y))//纵
坐标在该线段两端点之间
          tmp = P[i].x + (a.y - P[i].y) / (P[j].y - P[i].y) * (P[j].x -
P[i].x), cnt += sgn(tmp - a.x) > 0;//交点在A 右方
   return cnt & 1;//穿过奇数次则在多边形以内
}
bool pip convex jud(const Point &a, const Point &L, const Point &R) {//
判断AL 是否在AR 右边
   return sgn(det(L - a, R - a)) > 0;//必须严格以内
}
bool pip convex(const Point *P, const int &n, const Point &a) {// 【二分
法】判断点A 是否在凸多边形 Poly 以内
   //点按逆时针给出
   if (pip_convex_jud(P[0], a, P[1]) || pip_convex_jud(P[0], P[n - 1],
a)) return 0;//在P[0 1]或P[0 n-1]外
   if (sp_on(Line(P[0], P[1]), a) || sp_on(Line(P[0], P[n - 1]), a)) re
turn 2;//在P[0 1]或P[0 n-1]上
   int l = 1, r = n - 2;
   while (1 < r) {//二分找到一个位置 pos 使得 P[0]_A 在 P[0_pos], P[0_(pos+
1)]之间
       int mid = (1 + r + 1) >> 1;
       if (pip_convex_jud(P[0], P[mid], a))1 = mid;
       else r = mid - 1;
   if (pip_convex_jud(P[1], a, P[1 + 1]))return 0;//在P[pos_(pos+1)]外
   if (sp on(Line(P[1], P[1 + 1]), a))return 2;//EP[pos (pos+1)]
   return 1;
}
```

```
// 多边形是否包含线段
// 因此我们可以先求出所有和线段相交的多边形的顶点,然后按照X-Y 坐标排序(X 坐标
小的排在前面,对于X 坐标相同的点,Y 坐标小的排在前面,
// 这种排序准则也是为了保证水平和垂直情况的判断正确),这样相邻的两个点就是在线
段上相邻的两交点,如果任意相邻两点的中点也在多边形内,
// 则该线段一定在多边形内。
int pp judge(Point *A, int n, Point *B, int m) {// 【判断多边形 A 与多边形 B
是否相离】
   for (int i1 = 1; i1 <= n; ++i1) {</pre>
       int j1 = i1 < n ? i1 + 1 : 1;</pre>
       for (int i2 = 1; i2 <= m; ++i2) {
          int j2 = i2 < m ? i2 + 1 : 1;
          Point tmp;
          if (ss cross(Line(A[i1], A[j1]), Line(B[i2], B[j2]), tmp)) re
turn 0;//两线段相交
          if (pip(B, m, A[i1]) || pip(A, n, B[i2]))return 0;//点包含在内
       }
   }
   return 1;
}
double area(Point *P, int n) {//【任意多边形 P 的面积】
   double S = 0:
   for (int i = 1; i <= n; i++) S += det(P[i], P[i < n ? <math>i + 1 : 1]);
   return S / 2.0;
}
Line Q[N];
int judge(Line L, Point a) { return sgn(det(a - L.s, L.t - L.s)) > 0; }/
/判断点 a 是否在直线 L 的右边
int halfcut(Line *L, int n, Point *P) {//【半平面交】
   sort(L, L + n, [](const Line &a, const Line &b) {
      double d = atan2((a.t - a.s).y, (a.t - a.s).x) - atan2((b.t - b.
s).y, (b.t - b.s).x);
       return sgn(d) ? sgn(d) < 0 : judge(a, b.s);</pre>
   });
   int m = n;
   n = 0;
   for (int i = 0; i < m; ++i)</pre>
      if (i == 0 || sgn(atan2(Point(L[i]).y, Point(L[i]).x) - atan2(Po
int(L[i-1]).y, Point(L[i-1]).x)))
          L[n++] = L[i];
   int h = 1, t = 0;
   for (int i = 0; i < n; ++i) {</pre>
      while (h < t && judge(L[i], ll_intersection(Q[t], Q[t - 1]))) --t;</pre>
```

```
//当队尾两个直线交点不是在直线L[i]上或者左边时就出队
                 while (h < t \&\& judge(L[i], ll_intersection(Q[h], Q[h + 1]))) ++h;
//当队头两个直线交点不是在直线 L[i]上或者左边时就出队
                 Q[++t] = L[i];
        while (h < t \&\& judge(Q[h], ll_intersection(Q[t], Q[t - 1]))) --t;
        while (h < t \&\& judge(Q[t], ll_intersection(Q[h], Q[h + 1]))) ++h;
        for (int i = h; i <= t; ++i) {</pre>
                 P[n++] = 11_{intersection}(Q[i], Q[i < t ? i + 1 : h]);
        return n;
}
Point V1[N], V2[N];
int mincowski(Point *P1, int n, Point *P2, int m, Point *V) {//【闵可夫斯
基和】求两个凸包\{P1\}, \{P2\}的向量集合\{V\}=\{P1+P2\}构成的凸包
         for (int i = 0; i < n; ++i) V1[i] = P1[(i + 1) \% n] - P1[i];
         for (int i = 0; i < m; ++i) V2[i] = P2[(i + 1) % m] - P2[i];</pre>
         int t = 0, i = 0, j = 0;
        V[t++] = P1[0] + P2[0];
        while (i < n \&\& j < m) V[t] = V[t - 1] + (sgn(det(V1[i], V2[j])) > 0?
  V1[i++] : V2[j++]), t++;
        while (i < n) V[t] = V[t - 1] + V1[i++], t++;
        while (j < m) V[t] = V[t - 1] + V2[j++], t++;
         return t;
}
circle getcircle(const Point &A, const Point &B, const Point &C) {// [=
点确定一圆】向量垂心法
         Point P1 = (A + B) * 0.5, P2 = (A + C) * 0.5;
         Line R1 = Line(P1, P1 + normal(B - A));
         Line R2 = Line(P2, P2 + normal(C - A));
        circle 0;
        0.o = ll_intersection(R1, R2);
        0.r = length(A - 0.0);
        return 0;
}
struct ConvexHull {
         int op;
        struct cmp {
                 bool operator()(const Point &a, const Point &b) const {
                           return sgn(a.x - b.x) < 0 \mid | sgn(a.x - b.x) == 0 && sgn(a.y - b.x) == 0 & sgn(a.y - b
  b.y) < 0;
```

```
}
   };
   set<Point, cmp> s;
   ConvexHull(int o) {
       op = o;
       s.clear();
   }
   inline int PIP(Point P) {
       set<Point>::iterator it = s.lower_bound(Point(P.x, -dinf));//救
到第一个横坐标大于P 的点
       if (it == s.end())return 0;
       if (sgn(it->x - P.x) == 0) return sgn((P.y - it->y) * op) <= 0;//
比较纵坐标大小
       if (it == s.begin())return 0;
       set<Point>::iterator j = it, k = it;
       return sgn(det(P - *j, *k - *j) * op) >= 0;//看叉姬1
   }
   inline int judge(set<Point>::iterator it) {
       set<Point>::iterator j = it, k = it;
       if (j == s.begin())return 0;
       --j;
       if (++k == s.end())return 0;
       return sgn(det(*it - *j, *k - *j) * op) >= 0;//看叉姬
   }
   inline void insert(Point P) {
       if (PIP(P))return;//如果点P已经在凸壳上或凸包里就不插入了
       set<Point>::iterator tmp = s.lower_bound(Point(P.x, -inf));
       if (tmp != s.end() && sgn(tmp->x - P.x) == 0)s.erase(tmp);//特判
横坐标相等的点要去掉
       s.insert(P);
       set<Point>::iterator it = s.find(P), p = it;
       if (p != s.begin()) {
          --p;
          while (judge(p)) {
              set<Point>::iterator temp = p--;
              s.erase(temp);
          }
       if ((p = ++it) != s.end()) {
          while (judge(p)) {
              set<Point>::iterator temp = p++;
              s.erase(temp);
```

```
}
       }
int PIC(circle C, Point a) { return sgn(length(a - C.o) - C.r) <= 0; }//</pre>
判断点A 是否在圆C 内
void Random(Point *P, int n) { for (int i = 0; i < n; ++i)swap(P[i], P</pre>
[(rand() + 1) % n]); }//随机一个排列
circle min circle(Point *P, int n) {//【求点集 P 的最小覆盖圆】 O(n)
// random shuffle(P,P+n);
   Random(P, n);
   circle C = circle(P[0], 0);
   for (int i = 1; i < n; ++i)</pre>
       if (!PIC(C, P[i])) {
          C = circle(P[i], 0);
          for (int j = 0; j < i; ++j)</pre>
              if (!PIC(C, P[j])) {
                 C.o = (P[i] + P[j]) * 0.5, C.r = length(P[j] - C.o);
                  for (int k = 0; k < j; ++k) if (!PIC(C, P[k])) C = get
circle(P[i], P[j], P[k]);
   return C;
}
自适应辛普森
double f(double x) {
}
double simpson(double 1, double r) {
   double mid = (1 + r) / 2;
   return (r - 1) * (f(1) + 4 * f(mid) + f(r)) / 6; // 辛普森公式
}
double asr(double 1, double r, double EPS, double ans) {
   double mid = (1 + r) / 2;
   double fl = simpson(l, mid), fr = simpson(mid, r);
   if (abs(fl + fr - ans) <= 15 * EPS)
       return fl + fr + (fl + fr - ans) / 15; // 足够相似的话就直接返回
   return asr(1, mid, EPS / 2, fl) +
         asr(mid, r, EPS / 2, fr); // 否则分割成两段递归求解
}
```

```
球体积交和并
#include<bits/stdc++.h>
#define fi first
#define sf scanf
#define se second
#define pf printf
#define pb push_back
#define mp make_pair
#define sz(x) ((int)(x).size())
#define all(x) (x).begin(),(x).end()
#define mem(x,y) memset((x),(y),sizeof(x))
#define fup(i,x,y) for(int i=(x);i<=(y);++i)
#define fdn(i,x,y) for(int i=(x);i>=(y);--i)
typedef long long 11;
typedef long double ld;
typedef unsigned long long ull;
typedef std::pair<int,int> pii;
using namespace std;
const ld pi=acos(-1);
ld pow2(ld x){return x*x;}
ld pow3(ld x){return x*x*x;}
ld dis(ld x1,ld y1,ld z1,ld x2,ld y2,ld z2)
{
   return pow2(x1-x2)+pow2(y1-y2)+pow2(z1-z2);
```

```
}
ld cos(ld a,ld b,ld c){return (b*b+c*c-a*a)/(2*b*c);}
ld cap(ld r,ld h){return pi*(r*3-h)*h*h/3;} // 球缺体积公式,h 为球缺的高
//2 球体积交
ld sphere_intersect(ld x1,ld y1,ld z1,ld r1,ld x2,ld y2,ld z2,ld r2)
{
   ld d=dis(x1,y1,z1,x2,y2,z2);
   //相离
   if(d>=pow2(r1+r2))return 0;
   //包含
   if(d<=pow2(r1-r2))return pow3(min(r1,r2))*4*pi/3;</pre>
   //相交
   ld h1=r1-r1*cos(r2,r1,sqrt(d)),h2=r2-r2*cos(r1,r2,sqrt(d));
   return cap(r1,h1)+cap(r2,h2);
}
//2 球体积并
ld sphere_union(ld x1,ld y1,ld z1,ld r1,ld x2,ld y2,ld z2,ld r2)
{
   ld d=dis(x1,y1,z1,x2,y2,z2);
   //相离
   if(d>=pow2(r1+r2))return (pow3(r1)+pow3(r2))*4*pi/3;
   //包含
   if(d<=pow2(r1-r2))return pow3(max(r1,r2))*4*pi/3;</pre>
   //相交
```

```
ld h1=r1+r1*cos(r2,r1,sqrt(d)),h2=r2+r2*cos(r1,r2,sqrt(d));
    return cap(r1,h1)+cap(r2,h2);
}
int main()
{
    double x1,y1,z1,r1,x2,y2,z2,r2;
    sf("%1f%1f%1f%1f%1f%1f%1f%1f",&x1,&y1,&z1,&r1,&x2,&y2,&z2,&r2);
    pf("%.12Lf\n",sphere_union(x1,y1,z1,r1,x2,y2,z2,r2));
    return 0;
}
数据结构
仙人掌
```

```
仙人掌:任意一条边至多只出现在一条简单回路的无向连通图称为仙人掌。
 转化为圆方树,然后根据树的算法来做一些问题,注意区分圆点和方点
这题: 求带环(环和环之间无公共边) 无向图两点间的最短路径
 */
#include <iostream>
#include <cstring>
#include <algorithm>
using namespace std;
const int N = 12010, M = N * 3;
int n, m, Q, new_n;
int h1[N], h2[N], e[M], w[M], ne[M], idx;
int dfn[N], low[N], cnt;
int s[N], stot[N], fu[N], fw[N];
int fa[N][14], depth[N], d[N];
int A, B;
void add(int h[], int a, int b, int c)
   e[idx] = b, w[idx] = c, ne[idx] = h[a], h[a] = idx ++ ;
}
void build_circle(int x, int y, int z)
   int sum = z;
   for (int k = y; k != x; k = fu[k])
      s[k] = sum;
      sum += fw[k];
   s[x] = stot[x] = sum;
   add(h2, x, ++ new_n, 0);
   for (int k = y; k != x; k = fu[k])
   {
       stot[k] = sum;
      add(h2, new_n, k, min(s[k], sum - s[k]));
}
void tarjan(int u, int from)
   dfn[u] = low[u] = ++ cnt;
   for (int i = h1[u]; ~i; i = ne[i])
   {
      int j = e[i];
```

```
if (!dfn[j])
           fu[j] = u, fw[j] = w[i];
           tarjan(j, i);
           low[u] = min(low[u], low[j]);
           if (dfn[u] < low[j]) add(h2, u, j, w[i]);</pre>
       }
       else if (i != (from ^ 1)) low[u] = min(low[u], dfn[j]);
   for (int i = h1[u]; \sim i; i = ne[i])
       int j = e[i];
       if (dfn[u] < dfn[j] && fu[j] != u)</pre>
           build_circle(u, j, w[i]);
}
void dfs lca(int u, int father)
   depth[u] = depth[father] + 1;
   fa[u][0] = father;
   for (int k = 1; k <= 13; k ++ )
       fa[u][k] = fa[fa[u][k - 1]][k - 1];
   for (int i = h2[u]; \sim i; i = ne[i])
       int j = e[i];
       d[j] = d[u] + w[i];
       dfs_lca(j, u);
   }
}
int lca(int a, int b)
   if (depth[a] < depth[b]) swap(a, b);</pre>
   for (int k = 13; k >= 0; k -- )
       if (depth[fa[a][k]] >= depth[b])
           a = fa[a][k];
   if (a == b) return a;
   for (int k = 13; k >= 0; k -- )
       if (fa[a][k] != fa[b][k])
       {
           a = fa[a][k];
           b = fa[b][k];
       }
   A = a, B = b;
   return fa[a][0];
}
int main()
```

```
{
   scanf("%d%d%d", &n, &m, &Q);
   new_n = n;
   memset(h1, -1, sizeof h1);
   memset(h2, -1, sizeof h2);
   while (m -- )
   {
       int a, b, c;
       scanf("%d%d%d", &a, &b, &c);
       add(h1, a, b, c), add(h1, b, a, c);
   tarjan(1, -1);
   dfs_lca(1, 0);
   while (Q -- )
       int a, b;
       scanf("%d%d", &a, &b);
       int p = 1ca(a, b);
       if (p <= n) printf("%d\n", d[a] + d[b] - d[p] * 2);</pre>
       else
       {
           int da = d[a] - d[A], db = d[b] - d[B];
           int 1 = abs(s[A] - s[B]);
           int dm = min(1, stot[A] - 1);
           printf("%d\n", da + dm + db);
   }
   return 0;
}
CDQ
```

```
/*
处理三维偏序问题,
每个node 的三维不能完全相等,完全相等的话加权做
#include <iostream>
#include <cstring>
#include <algorithm>
using namespace std;
const int N = 100010, M = 200010;
int n, m;
struct Data
{
   int a, b, c, s, res;
   bool operator< (const Data& t) const</pre>
       if (a != t.a) return a < t.a;</pre>
       if (b != t.b) return b < t.b;</pre>
       return c < t.c;
   bool operator== (const Data& t) const
       return a == t.a && b == t.b && c == t.c;
}q[N], w[N];
int tr[M], ans[N];
int lowbit(int x)
   return x & -x;
}
void add(int x, int v)
   for (int i = x; i < M; i += lowbit(i)) tr[i] += v;</pre>
int query(int x)
{
   int res = 0;
   for (int i = x; i; i -= lowbit(i)) res += tr[i];
   return res;
}
```

```
void merge sort(int 1, int r)
{
   if (1 >= r) return;
   int mid = 1 + r \gg 1;
   merge_sort(1, mid), merge_sort(mid + 1, r);
   int i = 1, j = mid + 1, k = 0;
   while (i <= mid && j <= r)
       if (q[i].b \le q[j].b) add(q[i].c, q[i].s), w[k ++] = q[i ++];
       else q[j].res += query(q[j].c), w[k ++ ] = q[j ++ ];
   while (i <= mid) add(q[i].c, q[i].s), w[k ++ ] = q[i ++ ];
   while (j <= r) q[j].res += query(q[j].c), w[k ++ ] = q[j ++ ];
   for (i = 1; i <= mid; i ++ ) add(q[i].c, -q[i].s);
   for (i = 1, j = 0; j < k; i ++, j ++) q[i] = w[j];
}
int main()
   scanf("%d%d", &n, &m);
   for (int i = 0; i < n; i ++ )
   {
       int a, b, c;
       scanf("%d%d%d", &a, &b, &c);
       q[i] = \{a, b, c, 1\};
   sort(q, q + n);
   int k = 1;
   for (int i = 1; i < n; i ++)
       if (q[i] == q[k - 1]) q[k - 1].s ++ ;
       else q[k ++] = q[i];
   merge_sort(0, k - 1);
   for (int i = 0; i < k; i ++)
       ans[q[i].res + q[i].s - 1] += q[i].s;
   for (int i = 0; i < n; i ++ ) printf("%d\n", ans[i]);</pre>
   return 0;
}
kruskal 重构树
int pa[N];
void init(int n) {
   for (int i = 0; i <= n; i++) {</pre>
       pa[i] = i;
   }
}
```

```
int find(int a) {
   return pa[a] == a ? a : pa[a] = find(pa[a]);
}
struct edge {
   int from, to, 1;
};
int w[N];
edge e[N];
vector<int> g[N];
int kruskal(int n, int m) {
   int kcnt = n;
   init(n);
   sort(e + 1, e + 1 + m, [](edge a, edge b) { return a.l < b.l; });</pre>
   for (int i = 1; i <= m; i++) {</pre>
       int u = find(e[i].from);
       int v = find(e[i].to);
       if (u == v) continue;
       w[++kcnt] = e[i].1;
       pa[kcnt] = pa[u] = pa[v] = kcnt;
       g[u].push_back(kcnt);
       g[v].push_back(kcnt);
       g[kcnt].push back(u);
       g[kcnt].push_back(v);
   return kcnt;
}
普通莫队
#include <bits/stdc++.h>
using namespace std;
const int N = 1e6 + 10, M = 1e6 + 10;
int a[N];
struct node {
      int id, l, r;
} mp[M];
int len;
int ans[M], cnt[1000010];
int getNum(int 1) {
```

return 1 / len;

}

```
//左指针的分块,右指针的大小
bool cmp (const node &a, const node & b) {
      if(getNum(a.l) == getNum(b.l)) return a.r < b.r;</pre>
      return a.l < b.l;</pre>
}
/* 奇偶优化
struct node {
 int l, r, id;
 bool operator<(const node &x) const {</pre>
   if (l / unit != x.l / unit) return l < x.l;</pre>
   if ((l / unit) & 1)
     return r < x.r; // 注意这里和下面一行不能写小于 (大于) 等于
   return r > x.r;
};
void add(int x, int& res) {
      if(cnt[x] == 0) res++;
      cnt[x] ++;
}
void del(int x, int& res) {
      cnt[x] --;
      if(cnt[x] == 0) res --;
}
int main() {
      ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
      int n;
      cin >> n;
      for(int i = 1; i <= n; ++ i) {
            cin >> a[i];
      }
      int m;
      cin >> m;
      len = sqrt((double)n * n / m);
      for(int i = 1; i <= m; ++ i) {</pre>
            mp[i].id = i;
            cin >> mp[i].1 >> mp[i].r;
      sort(mp + 1, mp + m + 1, cmp);
      //离线处理询问
      int res = 0, i = 0, j = 0;
      for(int k = 1; k <= m; ++ k) {</pre>
            int id = mp[k].id, 1 = mp[k].1, r = mp[k].r;
```

```
while(j < r) add(a[++j], res);</pre>
            while(j > r) del(a[j--], res);
            while(i < 1) del(a[i++], res);</pre>
            while(i > 1) add(a[--i], res);
            ans[id] = res;
      }
      for(int i = 1; i <= m; ++ i) {</pre>
            cout << ans[i] << endl;</pre>
      }
      return 0;
}
带修莫队
#include <bits/stdc++.h>
using namespace std;
const int N = 10010;
int a[N], cnt[1000010], ans[N];
int len, mq, mc;
struct Query {
      int id, 1, r, t;
} q[N];
struct Modify {
      int p, c;
} c[N];
int getNum(int x) {
      return x / len;
}
// L 所在块的编号, r 所在块的编号, t 升序
bool cmp(const Query& a, const Query& b) {
      if(getNum(a.1) == getNum(b.1) && getNum(a.r) == getNum(b.r)) {
            return a.t < b.t;</pre>
      if(getNum(a.1) == getNum(b.1)) return a.r < b.r;</pre>
      return a.l < b.l;</pre>
}
void add(int x, int& res) {
    if (!cnt[x]) res ++ ;
```

```
cnt[x] ++ ;
}
void del(int x, int& res) {
    cnt[x] --;
    if (!cnt[x]) res -- ;
}
int main() {
      ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
      int n, m;
      cin >> n >> m;
      char op;
      int x, y;
      for(int i = 1; i <= n; ++ i) {</pre>
             cin >> a[i];
      for(int i = 1; i <= m; ++ i) {</pre>
             cin >> op >> x >> y;
       if (op == 'Q') q[++ mq] = \{mq, x, y, mc\};
       else c[ ++ mc] = \{x, y\};
      }
 ///
      len = cbrt((double)n * mc) + 1;
  sort(q + 1, q + mq + 1, cmp);
      int i = 1, j = 0, t = 0, res = 0;
      for(int k = 1; k <= mq; ++ k) {</pre>
             int id = q[k].id, l = q[k].l, r = q[k].r, tm = q[k].t;
             while(j < r) add(a[++ j], res);</pre>
             while(j > r) del(a[j --], res);
             while(i < 1) del(a[i ++], res);</pre>
             while(i > 1) add(a[-- i], res);
             while(t < tm) {</pre>
                   ++ t;
                   if(c[t].p >= i && c[t].p <= j) {
                          del(a[c[t].p], res);
                          add(c[t].c, res);
                   swap(a[c[t].p], c[t].c);
             while(t > tm) {
                   if(c[t].p >= i && c[t].p <= j) {
                          del(a[c[t].p], res);
                          add(c[t].c, res);
                   }
```

回滚莫队

```
/*
离线,询问按左端点升序为第一关键字,右端点升序为第二关键字
对于都在块内的点直接暴力,否则跨块:
若当前左端点所属的块与上一个不同,则将左端点初始为当前块的右端点+1,右端点初始
为当前块的右端点
左端点每次暴力,右端点单调
*/
#include <iostream>
#include <cstring>
#include <cstdio>
#include <algorithm>
#include <cmath>
#include <vector>
using namespace std;
typedef long long LL;
const int N = 100010;
int n, m, len;
int w[N], cnt[N];
LL ans[N];
struct Query
   int id, 1, r;
}q[N];
vector<int> nums;
int get(int x)
   return x / len;
}
bool cmp(const Query& a, const Query& b)
   int i = get(a.1), j = get(b.1);
   if (i != j) return i < j;
   return a.r < b.r;</pre>
}
void add(int x, LL& res)
{
   cnt[x] ++ ;
   res = max(res, (LL)cnt[x] * nums[x]);
}
int main()
```

```
{
   scanf("%d%d", &n, &m);
   len = sqrt(n);
   for (int i = 1; i <= n; i ++ ) scanf("%d", &w[i]),
nums.push_back(w[i]);
   sort(nums.begin(), nums.end());
   nums.erase(unique(nums.begin(), nums.end()), nums.end());
   for (int i = 1; i <= n; i ++ )
       w[i] = lower bound(nums.begin(), nums.end(), w[i]) -
nums.begin();
   for (int i = 0; i < m; i ++)
       int 1, r;
       scanf("%d%d", &1, &r);
       q[i] = \{i, 1, r\};
   sort(q, q + m, cmp);
   for (int x = 0; x < m;)
       int y = x;
       while (y < m \&\& get(q[y].1) == get(q[x].1)) y ++ ;
       int right = get(q[x].1) * len + len - 1;
       // 暴力求块内的询问
       while (x < y && q[x].r <= right)
           LL res = 0;
           int id = q[x].id, 1 = q[x].1, r = q[x].r;
           for (int k = 1; k <= r; k ++ ) add(w[k], res);
           ans[id] = res;
           for (int k = 1; k <= r; k ++ ) cnt[w[k]] -- ;
           x ++ ;
       }
       // 求块外的询问
       LL res = 0;
       int i = right, j = right + 1;
       while (x < y)
           int id = q[x].id, 1 = q[x].1, r = q[x].r;
           while (i < r) add(w[ ++ i], res);</pre>
           LL backup = res;
           while (j > 1) add(w[ -- j], res);
           ans[id] = res;
          while (j < right + 1) cnt[w[j ++ ]] -- ;
          res = backup;
           x ++ ;
```

```
memset(cnt, 0, sizeof cnt);
    }
   for (int i = 0; i < m; i ++ ) printf("%lld\n", ans[i]);</pre>
    return 0;
}
线段树合并分裂
11 nodetot, recycnt, bac[N << 5], ch[N << 5][2], rt[N];</pre>
ll val[N << 5];
11 newnod() { return (recycnt ? bac[recycnt--] : ++nodetot); }
void recyc(ll p) {
    bac[++recycnt] = p, ch[p][0] = ch[p][1] = val[p] = 0;
    return;
}
void pushdown(ll p) {
}
void pushup(ll p) {
   val[p] = 0;
    if (ch[p][0]) val[p] += val[ch[p][0]];
   if (ch[p][1]) val[p] += val[ch[p][1]];
}
void modify(l1 &p, l1 l, l1 r, l1 pos, l1 v) {
    if (!p) { p = newnod(); }
   if (1 == r) {
       val[p] += v;
       return;
   11 \text{ mid} = (1 + r) >> 1;
   pushdown(p);
   if (pos <= mid) { modify(ch[p][0], 1, mid, pos, v); }</pre>
    else { modify(ch[p][1], mid + 1, r, pos, v); }
    pushup(p);
   return;
}
11 query(11 p, 11 l, 11 r, 11 x1, 11 xr) {
   if (xr < 1 | | r < x1) { return 0; }
    if (xl <= 1 && r <= xr) { return val[p]; }</pre>
   11 \text{ mid} = (1 + r) >> 1;
```

```
// pushdown(p);
   return query(ch[p][0], 1, mid, x1, xr) + query(ch[p][1], mid + 1, r,
x1, xr);
11 kth(ll p, ll l, ll r, ll k) {
   if (1 == r) { return 1; }
   11 \text{ mid} = (1 + r) >> 1;
   pushdown(p);
   if (val[ch[p][0]] >= k) { return kth(ch[p][0], 1, mid, k); }
   else { return kth(ch[p][1], mid + 1, r, k - val[ch[p][0]]); }
}
11 merge(ll x, ll y, ll l, ll r) {
   if (!x || !y) {
       return x + y;
      // 只有一边有点,不用合并
   ll p = newnod(); // 创建一个新结点 p
                                // 边界(某些时候可以省略,见下面一个代
   if (1 == r) {
石马)
       val[p] = val[x] + val[y];
       return p;
   }
// pushdown(x), pushdown(y);
   11 \text{ mid} = (1 + r) >> 1;
   ch[p][0] = merge(ch[x][0], ch[y][0], 1, mid);
   ch[p][1] = merge(ch[x][1], ch[y][1], mid + 1, r);
   recyc(x), recyc(y);
                               // 垃圾回收
                                 // pushup
   pushup(p);
   return p;
}
void split(ll x, ll &y, ll k) {
   if (x == 0) return;
   y = newnod();
   11 v = val[ch[x][0]];
// pushdown(x);
   if (k > v) { split(ch[x][1], ch[y][1], k - v); }
   else { swap(ch[x][1], ch[y][1]); }
   if (k < v) { split(ch[x][0], ch[y][0], k); }</pre>
   val[y] = val[x] - k;
   val[x] = k;
   return;
}
```

主席树

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 N = 1 << 20;
11 ch[N << 5][2], rt[N], tot;</pre>
11 val[N << 5];
11 update(11 a, 11 b) {
   return a + b;
}
ll build(ll l, ll r) { // 建树
   11 p = ++tot;
   if (1 == r) {
       //初始化
       val[p] = 0;
       return p;
   11 \text{ mid} = (1 + r) >> 1;
   ch[p][0] = build(1, mid);
   ch[p][1] = build(mid + 1, r);
   val[p] = update(val[ch[p][0]], val[ch[p][1]]);
   return p; // 返回该子树的根节点
}
ll modify(ll pre, ll l, ll r, ll pos, ll v) { // 插入操作
   11 \text{ now} = ++\text{tot};
   ch[now][0] = ch[pre][0], ch[now][1] = ch[pre][1];
   if (1 == r) {
       val[now] = val[pre] + v;
       return now;
   11 \text{ mid} = (1 + r) >> 1;
   if (pos <= mid)</pre>
       ch[now][0] = modify(ch[now][0], 1, mid, pos, v);
   else
       ch[now][1] = modify(ch[now][1], mid + 1, r, pos, v);
   val[now] = update(val[ch[now][0]], val[ch[now][1]]);
   return now;
}
ll kth(ll pre, ll now, ll l, ll r, ll k) { // 查询操作
   11 \text{ mid} = (1 + r) >> 1;
   ll x = val[ch[now][0]] - val[ch[pre][0]]; // 通过区间减法得到左儿子的
信息
   if (1 == r) return 1;
```

```
if (k <= x) // 说明在左儿子中
       return kth(ch[pre][0], ch[now][0], 1, mid, k);
   else // 说明在右儿子中
       return kth(ch[pre][1], ch[now][1], mid + 1, r, k - x);
}
ll query(ll pre, ll now, ll l, ll r, ll ql, ll qr) { // 查询操作
   if (q1 <= 1 && r <= qr) {
       return val[now] - val[pre];
   if (qr < 1 || r < q1) {
       return 0;
   11 \text{ mid} = (1 + r) >> 1;
   11 lv = query(ch[pre][0], ch[now][0], 1, mid, q1, qr);
   11 rv = query(ch[pre][1], ch[now][1], mid + 1, r, ql, qr);
   return update(lv, rv);
}
//修改查询记得用rt[]!!!
LCT
```

```
11 ch[N][2], f[N], sum[N], val[N], tag[N], siz[N], siz2[N];
inline void pushup(ll p) {
   sum[p] = sum[ch[p][0]] ^ sum[ch[p][1]] ^ val[p];
   siz[p] = siz[ch[p][0]] + siz[ch[p][1]] + 1 + siz2[p];
}
inline void pushdown(11 p) {
   if (tag[p]) {
       if (ch[p][0]) swap(ch[ch[p][0]][0], ch[ch[p][0]][1]),
tag[ch[p][0]] ^= 1;
       if (ch[p][1]) swap(ch[ch[p][1]][0], ch[ch[p][1]][1]),
tag[ch[p][1]] ^= 1;
      tag[p] = 0;
   }
}
11 getch(11 x) { return ch[f[x]][1] == x; }
bool isroot(11 x) { return ch[f[x]][0] != x && ch[f[x]][1] != x; }
inline void rotate(11 x) {
   11 y = f[x], z = f[y], k = getch(x);
   if (!isroot(y)) ch[z][ch[z][1] == y] = x;
   // 上面这句一定要写在前面,普通的 Splay 是不用的,因为 isRoot (后面会讲)
   ch[y][k] = ch[x][!k], f[ch[x][!k]] = y;
   ch[x][!k] = y, f[y] = x, f[x] = z;
   pushup(y), pushup(x);
}
// 从上到下一层一层 pushDown 即可
void update(11 p) {
   if (!isroot(p)) update(f[p]);
   pushdown(p);
}
inline void splay(11 x) {
   update(x); // 马上就能看到啦。 在
   // Splay 之前要把旋转会经过的路径上的点都 PushDown
   for (11 fa; fa = f[x], !isroot(x); rotate(x)) {
      if (!isroot(fa)) rotate(getch(fa) == getch(x) ? fa : x);
}
// 回顾一下代码
inline void access(11 x) {
   for (11 p = 0; x; p = x, x = f[x]) {
       splay(x), siz2[x] += siz[ch[x][1]] - siz[p], ch[x][1] = p,
```

```
pushup(x);
   }
inline void makeroot(11 p) {
   access(p);
   splay(p);
   swap(ch[p][0], ch[p][1]);
   tag[p] ^= 1;
}
inline void split(ll a, ll b) {
   makeroot(a);
   access(b);
   splay(b);
}
inline 11 find(11 p) {
   access(p), splay(p);
   while (ch[p][0]) pushdown(p), p = ch[p][0];
   splay(p);
   return p;
}
inline void link(11 x, 11 y) {
   makeroot(y);
   makeroot(x);
   if (find(y) != x) {
       f[x] = y;
       siz2[y] += siz[x];
   }
}
inline void cut(11 x, 11 y) {
   makeroot(x);
   if (find(y) == x \&\& f[y] == x) {
       ch[x][1] = f[y] = 0;
       pushup(x);
   }
}
void init(int n) {
   for (int i = 1; i <= n; i++) siz[i] = 1;
}
Splay1
#include <bits/stdc++.h>
using namespace std;
```

```
struct Splay {
   static const int N = 100005;
   int rt, tot, fa[N], ch[N][2], val[N], cnt[N], sz[N];
   // rt=根编号, tot=总节点, fa=父节点编号, ch=左/右儿子编号, val=节点的
值,cnt=权值出现次数,sz=子树大小
   void maintain(int x) { //更新x 节点字数大小
       sz[x] = sz[ch[x][0]] + sz[ch[x][1]] + cnt[x];
   }
   bool get(int x) {
       return x == ch[fa[x]][1];
   } //返回节点是父亲的0/1-左/右儿子
   void clear(int x) { //销毁节点x
       ch[x][0] = ch[x][1] = fa[x] = val[x] = sz[x] = cnt[x] = 0;
   }
   void rotate(int x) { //旋转
       int y = fa[x], z = fa[y], chk = get(x);
       ch[y][chk] = ch[x][chk ^ 1];
      fa[ch[x][chk ^ 1]] = y;
       ch[x][chk ^ 1] = y;
      fa[y] = x;
      fa[x] = z;
      if (z) ch[z][y == ch[z][1]] = x;
      maintain(x);
      maintain(y);
   }
   void splay(int x) { //将x 节点移动到根
      for (int f = fa[x]; f = fa[x], f; rotate(x))
          if (fa[f]) rotate(get(x) == get(f) ? f : x);
       rt = x;
   }
   void ins(int k) { //插入
       if (!rt) {
          val[++tot] = k;
          cnt[tot]++;
          rt = tot;
          maintain(rt);
          return;
       int cnr = rt, f = 0;
      while (1) {
          if (val[cnr] == k) {
              cnt[cnr]++;
```

```
maintain(cnr);
           maintain(f);
           splay(cnr);
           break;
       }
       f = cnr;
       cnr = ch[cnr][val[cnr] < k];</pre>
       if (!cnr) {
           val[++tot] = k;
           cnt[tot]++;
           fa[tot] = f;
           ch[f][val[f] < k] = tot;
           maintain(tot);
           maintain(f);
           splay(tot);
           break;
       }
   }
}
int rk(int k) { // k 权值的排名
   int res = 0, cnr = rt;
   while (1) {
       if (k < val[cnr]) {
           cnr = ch[cnr][0];
       } else {
           res += sz[ch[cnr][0]];
           if (k == val[cnr]) {
               splay(cnr);
               return res + 1;
           }
           res += cnt[cnr];
           cnr = ch[cnr][1];
       }
   }
}
int kth(int k) { //第 k 名的权值
   int cnr = rt;
   while (1) {
       if (ch[cnr][0] && k <= sz[ch[cnr][0]]) {</pre>
           cnr = ch[cnr][0];
       } else {
           k -= cnt[cnr] + sz[ch[cnr][0]];
           if (k <= 0) {
               splay(cnr);
               return val[cnr];
           cnr = ch[cnr][1];
```

```
}
      }
   }
   int pre() { //前驱节点编号
       int cnr = ch[rt][0];
       while (ch[cnr][1]) cnr = ch[cnr][1];
       splay(cnr);
       return cnr;
   } // 若需要得到前驱 tree.ins(x), printf("%d\n", tree.val[tree.pre
()]),
      // tree.del(x);
   int nxt() { //后驱节点编号
       int cnr = ch[rt][1];
       while (ch[cnr][0]) cnr = ch[cnr][0];
       splay(cnr);
       return cnr;
   } // 若需要得到后驱 tree.ins(x), printf("%d\n", tree.val[tree.pre
()]),
      // tree.del(x);
   void del(int k) { //删除k值
       rk(k);
       if (cnt[rt] > 1) {
          cnt[rt]--;
          maintain(rt);
          return;
       }
       if (!ch[rt][0] && !ch[rt][1]) {
          clear(rt);
          rt = 0;
          return;
       if (!ch[rt][0]) {
          int cnr = rt;
          rt = ch[rt][1];
          fa[rt] = 0;
          clear(cnr);
          return;
       if (!ch[rt][1]) {
          int cnr = rt;
          rt = ch[rt][0];
          fa[rt] = 0;
          clear(cnr);
          return;
       }
       int cnr = rt;
```

```
int x = pre();
       splay(x);
       fa[ch[cnr][1]] = x;
       ch[x][1] = ch[cnr][1];
       clear(cnr);
       maintain(rt);
   }
} tree;
splay2
11 ch[N][2], f[N], sum[N], val[N], tag[N], siz[N];
inline void pushup(ll p) {
   sum[p] = sum[ch[p][0]] ^ sum[ch[p][1]] ^ val[p];
   siz[p] = siz[ch[p][0]] + siz[ch[p][1]] + 1;
}
inline void pushdown(ll p) {
   if (tag[p]) {
       if (ch[p][0]) swap(ch[ch[p][0]][0], ch[ch[p][0]][1]),
tag[ch[p][0]] ^= 1;
       if (ch[p][1]) swap(ch[ch[p][1]][0], ch[ch[p][1]][1]),
tag[ch[p][1]] ^= 1;
       tag[p] = 0;
   }
}
11 getch(11 x) { return ch[f[x]][1] == x; }
bool isroot(ll x) { return ch[f[x]][0] != x && ch[f[x]][1] != x; }
inline void rotate(ll x) {
   11 y = f[x], z = f[y], k = getch(x);
   if (!isroot(y)) ch[z][ch[z][1] == y] = x;
   // 上面这句一定要写在前面,普通的 Splay 是不用的,因为 is Root (后面会讲)
   ch[y][k] = ch[x][!k], f[ch[x][!k]] = y;
   ch[x][!k] = y, f[y] = x, f[x] = z;
   pushup(y), pushup(x);
}
// 从上到下一层一层 pushDown 即可
void update(ll p) {
   if (!isroot(p)) update(f[p]);
   pushdown(p);
}
```

```
inline void splay(ll x) {
   update(x); // 马上就能看到啦。 在
   // Splay 之前要把旋转会经过的路径上的点都 PushDown
   for (ll fa; fa = f[x], !isroot(x); rotate(x)) {
       if (!isroot(fa)) rotate(getch(fa) == getch(x) ? fa : x);
   }
}
// 回顾一下代码
inline void access(ll x) {
   for (11 p = 0; x; p = x, x = f[x]) {
       splay(x), ch[x][1] = p, pushup(x);
   }
}
inline void makeroot(ll p) {
   access(p);
   splay(p);
   swap(ch[p][0], ch[p][1]);
   tag[p] ^= 1;
}
inline void split(ll a, ll b) {
   makeroot(a);
   access(b);
   splay(b);
}
inline 11 find(11 p) {
   access(p), splay(p);
   while (ch[p][0]) pushdown(p), p = ch[p][0];
   splay(p);
   return p;
}
inline void link(ll x, ll y) {
   makeroot(x);
   if (find(y) != x) f[x] = y;
}
inline void cut(ll x, ll y) {
   makeroot(x);
   if (find(y) == x \&\& f[y] == x) {
       ch[x][1] = f[y] = 0;
       pushup(x);
   }
}
```

```
Treap
#include <bits/stdc++.h>
using namespace std;
struct node {
    node* ch[2];
    int r;
    int v;
    int cmp(int const& a) const {
        if (v == a) return -a;
        return a > v ? 1 : 0;
    }
};
void rotate(node*& a, int d) {
    node* k = a \rightarrow ch[d ^ 1];
    a->ch[d ^ 1] = k->ch[d];
    k \rightarrow ch[d] = a;
    a = k;
}
void insert(node*& a, int x) {
    if (a == NULL) {
        a = new node;
         a \rightarrow ch[0] = a \rightarrow ch[1] = NULL;
         a \rightarrow v = x;
        a \rightarrow r = rand();
    } else {
        int d = a - > cmp(x);
        insert(a->ch[d], x);
        if (a->ch[d]->r > a->r) rotate(a, d ^ 1);
    }
void remove(node*& a, int x) {
    int d = a - > cmp(x);
    if (d == -1) {
        if (a->ch[0] == NULL)
             a = a \rightarrow ch[1];
        else if (a->ch[1] == NULL)
             a = a \rightarrow ch[0];
        else {
             int d2 = a \rightarrow ch[1] \rightarrow r \rightarrow a \rightarrow ch[0] \rightarrow r ? 0 : 1;
             rotate(a, d2);
             remove(a->ch[d2], x);
        }
    } else {
         remove(a->ch[d], x);
int find(node*& a, int x) {
    if (a == NULL)
```

```
return 0;
   else if (a->v == x)
        return 1;
    else {
        int d = a - > cmp(x);
        return find(a->ch[d], x);
    }
}
int main() {
   node* a = NULL;
   int k, 1;
   while (cin >> k >> 1) {
        if (k == 1)
            insert(a, 1);
       else if (k == 2)
            remove(a, 1);
       else {
            cout << find(a, 1) << endl;</pre>
    }
}
```

```
舞蹈链(多重覆盖)
#include <bits/stdc++.h>
using namespace std;
struct DLX {
   static const int maxn = 1000;
                                    //列的上限
   static const int maxr = 1000;
                                    //解的上限
   static const int maxnode = 5000; //总结点数上限
   static const int INF = 1000000000;
   int n, sz;
   int S[maxn];
   int row[maxnode], col[maxnode];
   int L[maxnode], R[maxnode], U[maxnode], D[maxnode];
   int ansd, ans[maxr];
   int vis[maxnode];
   void init(int n) {
       this->n = n;
       //虚拟节点
       for (int i = 0; i <= n; i++) {</pre>
          U[i] = i;
          D[i] = i;
```

```
L[i] = i - 1;
         R[i] = i + 1;
      R[n] = 0;
      L[0] = n;
      sz = n + 1;
      memset(S, 0, sizeof(S));
   }
   void addRow(int r, vector<int> columns) {
      int first = sz;
      for (int i = 0; i < columns.size(); i++) {</pre>
         int c = columns[i];
         L[sz] = sz - 1;
         R[sz] = sz + 1;
         D[sz] = c;
         U[sz] = U[c];
         D[U[c]] = sz;
         U[c] = sz;
         row[sz] = r;
         col[sz] = c;
         S[c]++;
         SZ++;
      R[sz - 1] = first;
      L[first] = sz - 1;
#define FOR(i, A, s) for (int i = A[s]; i != s; i = A[i])
   void remove(int c) {
      FOR(i, D, c) \{ L[R[i]] = L[i], R[L[i]] = R[i]; \}
   }
   void restore(int c) {
      FOR(i, U, c) { L[R[i]] = i, R[L[i]] = i; }
   int f check() //精确覆盖区估算剪枝
   {
      /*
      强剪枝。这个
      剪枝利用的思想是A*搜索中的估价函数。即,对于当前的递归深度K 下的矩
阵,估计其最好情况下(即最少还需要多少步)才能出解。也就是,如果将能够覆盖当
      前列的所有行全部选中,去掉这些行能够覆盖到的列,将这个操作作为一层深
度。重复此操作直到所有列全部出解的深度是多少。如果当前深度加上这个估价函数返
      回值,其和已然不能更优(也就是已经超过当前最优解),则直接返回,不必再
搜。
      int ret = 0;
```

```
FOR(c, R, 0) vis[c] = true;
      FOR(c, R, ∅)
      if (vis[c]) {
          ret++;
          vis[c] = false;
          FOR(i, D, c)
          FOR(j, R, i) vis[col[j]] = false;
       }
      return ret;
   // d 为递归深度
   void dfs(int d, vector<int>& v) {
      if (d + f check() >= ansd) return;
      if (R[0] == 0) {
          if (d < ansd) {</pre>
             ansd = d;
             v.clear();
             for (int i = 0; i < ansd; i++) {</pre>
                 v.push_back(ans[i]);
             }
                 //找到解
          }
          return; //记录解的长度
       }
      //找到S最小的列c
      int c = R[0];
      FOR(i, R, ∅)
      if (S[i] < S[c])
          c = i;
                    //第一个未删除的列
                    //删除第c列
      FOR(i, D, c) { //用结点i 所在的行能覆盖的所有其他列
          ans[d] = row[i];
          remove(i);
          FOR(j, R, i) remove(j); //删除结点i 所在的能覆的所有其他列
          dfs(d + 1, v);
          FOR(j, L, i) restore(j);
          restore(i); //恢复结点i所在的行能覆盖的所有其他列
                     //恢复第c列
      }
   }
   bool solve(vector<int>& v) {
      v.clear();
      ansd = INF;
      dfs(0, v);
      return !v.empty();
   }
};
//使用时 init 初始化,vector 中存入 r 行结点列表用 addRow 加行,solve(ans)后答
```

案按行的选择在 ans 中 DLX dlx; int main() { int n, m; cin >> n >> m; dlx.init(m); for (int i = 1; i <= n; i++) {</pre> vector<int> v; for (int j = 1; j <= m; j++) {</pre> int a; cin >> a; if (a == 1) v.push_back(j); dlx.addRow(i, v); } vector<int> ans; dlx.solve(ans); for (int i = 0; i < ans.size(); i++) cout << ans[i];</pre> }

舞蹈链 (精确覆盖)

```
#include <bits/stdc++.h>
using namespace std;
struct DLX {
                                    //列的上限
   static const int maxn = 1000;
   static const int maxr = 1000;
                                     //解的上限
   static const int maxnode = 5000; //总结点数上限
   int n, sz;
   int S[maxn];
   int row[maxnode], col[maxnode];
   int L[maxnode], R[maxnode], U[maxnode], D[maxnode];
   int ansd, ans[maxr];
   void init(int n) {
       this->n = n;
       //虚拟节点
       for (int i = 0; i <= n; i++) {
          U[i] = i;
          D[i] = i;
          L[i] = i - 1;
          R[i] = i + 1;
       }
       R[n] = 0;
       L[0] = n;
       sz = n + 1;
       memset(S, 0, sizeof(S));
   }
   void addRow(int r, vector<int> columns) {
       int first = sz;
       for (int i = 0; i < columns.size(); i++) {</pre>
           int c = columns[i];
           L[sz] = sz - 1;
           R[sz] = sz + 1;
          D[sz] = c;
          U[sz] = U[c];
          D[U[c]] = sz;
          U[c] = sz;
          row[sz] = r;
           col[sz] = c;
          S[c]++;
          SZ++;
       }
       R[sz - 1] = first;
       L[first] = sz - 1;
   }
```

```
#define FOR(i, A, s) for (int i = A[s]; i != s; i = A[i])
   void remove(int c) {
      L[R[c]] = L[c];
      R[L[c]] = R[c];
      FOR(i, D, c)
      FOR(j, R, i) {
         U[D[j]] = U[j];
         D[U[j]] = D[j];
         --S[col[j]];
   }
   void restore(int c) {
      FOR(i, U, c)
      FOR(j, L, i) {
         ++S[col[j]];
         U[D[j]] = j;
         D[U[j]] = j;
      L[R[c]] = c;
      R[L[c]] = c;
   }
   // d 为递归深度
   bool dfs(int d) {
      if (R[0] == 0) {
                     //找到解
         ansd = d;
         return true; //记录解的长度
      }
      //找到S最小的列c
      int c = R[0];
      FOR(i, R, 0) if (S[i] < S[c]) c = i; //第一个未删除的列
      remove(c);
                    //删除第c列
      FOR(i, D, c) { //用结点i 所在的行能覆盖的所有其他列
         ans[d] = row[i];
         FOR(j, R, i) remove(col[j]); //删除结点i所在的能覆的所有其他
列
         if (dfs(d + 1)) return true;
         FOR(j, L, i) restore(col[j]); //恢复结点i 所在的行能覆盖的所有
其他列
      restore(c); //恢复第c列
      return false;
   }
```

```
bool solve(vector<int>& v) {
       v.clear();
       if (!dfs(0)) return false;
       for (int i = 0; i < ansd; i++) v.push_back(ans[i]);</pre>
        return true;
};
//使用时 init 初始化,vector 中存入 r 行结点列表用 addRow 加行,solve(ans)后答
案按行的选择在 ans 中
数论
lucas 求组合数
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
11 p;
const int maxn = 1e5 + 10;
ll \ qpow(ll \ x, ll \ n)
      11 \text{ res} = 1;
      while(n){
            if(n & 1) res = (res * x) % p;
            x = (x * x) % p;
            n \gg 1;
      }
      return res;
}
11 C(11 up, 11 down){
      if(up > down) return 0;
      11 \text{ res} = 1;
//
      for(int i = up + 1; i <= down; ++ i){</pre>
//
            res = (res * i) % p;
//
//
      for(int i = 1; i <= down - up; ++ i){</pre>
//
            res = (res * qpow(i, p - 2)) % p;
//
      for(int i = 1, j = down; i <= up; ++ i, -- j){</pre>
            res = (res * j) % p;
            res = (res * qpow(i, p - 2)) % p;
      }
```

```
return res;
}
11 lucas(ll up, ll down){
     if(up 
     return C(up % p, down % p) * lucas(up / p, down / p) % p;
}
int main(){
     ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
     int T;
     cin >> T;
     while (T --){
           11 down, up;
           cin >> down >> up >> p;
           cout << lucas(up, down) % p << endl;</pre>
     }
     return 0;
}
扩展欧几里得求逆元
typedef long long 11;
void extgcd(ll a,ll b,ll& d,ll& x,ll& y){
   if(!b){ d=a; x=1; y=0;}
   else{ extgcd(b,a%b,d,y,x); y-=x*(a/b); }
}
ll inverse(ll a,ll n){
   11 d,x,y;
   extgcd(a,n,d,x,y);
   return d==1?(x+n)%n:-1;
}
```

逆元线性递推 inv 阶乘组合数

```
11 fac[N];// n!
ll invfac[N]; // n!的inv
11 invn[N]; //n 的inv
inline void init() {
   fac[0] = fac[1] = invfac[0] = invfac[1] = invn[0] = invn[1] = 1;
   for (int i = 2; i < N; ++i) {
      fac[i] = fac[i - 1] * i % mod;
      invn[i] = (mod - mod / i) * invn[mod % i] % mod;
      invfac[i] = invfac[i - 1] * invn[i] % mod;
   }
}
11 C(11 up, 11 down) {
   if (up > down) return 0;
   if (up < 0 || down < 0) return 0;
   11 res = fac[down];
   res = res * invfac[down - up] % mod;
   res = res * invfac[up] % mod;
   return res;
}
//先init
数学
一些范围
1~n 的质数个数
1~2e9 中拥有最多约数个数的数拥有的约数个数
约 1600
n 个不同的点可以构成n^{n-2}棵不同的树
```

勾股数

1.任何一个勾股数(a,b,c)内的三个数同时乘以一个正整数 n 得到的新数组(na, nb, nc) 仍然是勾股数,

于是找 abc 互质的勾股数

勾股数/圆上格点数

- 一,当 a 为大于 1 的奇数 2n+1 时, $b=2n^2+2n$, $c=2n^2+2n+1$ (把 a 拆成两个连续的自然数)
- 二, 当 a 为大于 4 的偶数 2n 时, $b = n^2 1$, $c = n^2 + 1$ (只想得到互质的数的话: a=4n, $b = 4n^2 1$, $c = 4n^2 + 1$

公式 1

a=2mnt

 $b = (m^2 - n^2) t$

 $c = (m^2 + n^2) t$

(t 是倍数)

完全公式

a=m, $b=(m^2/k-k)/2$, $c=(m^2/k+k)/2$

其中 m ≥3

- 1. 当 m 确定为任意一个 ≥3 的奇数时, k={1, m² 的所有小于 m 的因子}
- 2. 当 m 确定为任意一个 ≥4 的偶数时, k={m^2/2 的所有小于 m 的偶数因子}

高斯整数/高斯素数

二维平面转化为复数平面,

4n+1 的素数,都能分解成高斯素数,4n+3 的素数,他们本身就是高斯素数,2特殊

(乘以1, -1, i, -i 四个

半径为 \sqrt{n} 的圆上的格点数,先将 n 分解质因数,对每个不是高斯素数的数分解成共轭的高斯素数,分配数比指数多 1,指数是偶数的话,有一种方法分配,不然就没有格点

2 = (1+i)(1+i) ,但是这对数格点数没有影响,因为要乘-i。

它是一个周期函数,同时是一个积性函数,

再来看这个问题,

\$\$45 = 3^2 \times 5 \\ 半径为 \sqrt{45} 圆上格点数问题 = 4 \times (f(1)+f(3)+f(3^2)) \times(f(1)+f(5))\\ =4 \times (f(1)+f(3)+f(5)+f(9)+f(15)+f(45))\$\$

最后转化为 45 的所有约数

exgcd

```
11 ex gcd(ll a, ll b, ll &x, ll &y) {
   if (b == 0) {
       x = 1;
       y = 0;
       return a;
   11 d = ex gcd(b, a \% b, x, y);
   11 \text{ temp} = x;
   x = y;
   y = temp - a / b * y;
   return d;
}
Pollard Rho+Miller-Robin
typedef long long 11;
namespace Miller Rabin {
   const ll Pcnt = 12;
   const 11 p[Pcnt] = \{2, 3, 5, 7, 11, 13, 17, 19, 61, 2333, 4567,
24251};
   11 pow(ll a, ll b, ll p) {
       11 \text{ ans} = 1;
       for (; b; a = (_int128) a * a % p, b >>= 1)if (b & 1)ans =
( int128) ans * a % p;
       return ans;
   }
   bool check(ll x, ll p) {
       if (x \% p == 0 || pow(p \% x, x - 1, x) ^ 1)return true;
       11 t, k = x - 1;
       while ((k ^ 1) & 1) {
           t = pow(p \% x, k >>= 1, x);
           if (t ^ 1 && t ^ x - 1)return true;
           if (!(t ^ x - 1))return false;
       return false;
   }
   inline bool MR(ll x) { //用这个
```

```
if (x < 2)return false;</pre>
       for (int i = 0; i ^ Pcnt; ++i) {
           if (!(x ^ p[i]))return true;
           if (check(x, p[i]))return false;
       return true;
   }
}
namespace Pollard_Rho {
#define Rand(x) (111*rand()*rand()%(x)+1)
   11 gcd(const 11 a, const 11 b) { return b ? gcd(b, a % b) : a; }
   11 mul(const 11 x, const 11 y, const 11 X) {
       11 k = (1.0L * x * y) / (1.0L * X) - 1, t = (_int128) x * y -
(<u>__int128</u>) k * X;
       while (t < \emptyset)t += X;
       return t;
   }
   11 PR(const 11 x, const 11 y) {
       int t = 0, k = 1;
       11 \ v0 = Rand(x - 1), \ v = v0, \ d, \ s = 1;
       while (true) {
           v = (mul(v, v, x) + y) \% x, s = mul(s, abs(v - v0), x);
           if (!(v ^ v0) || !s)return x;
           if (++t == k) {
               if ((d = gcd(s, x)) ^ 1)return d;
               v0 = v, k <<= 1;
           }
       }
   }
   void Resolve(ll x, ll &ans) {
       if (!(x ^ 1) || x <= ans)return;</pre>
       if (Miller Rabin::MR(x)) {
           if (ans < x)ans = x;
           return;
       }
       11 y = x;
       while ((y = PR(x, Rand(x))) == x);
       while (!(x \% y))x /= y;
       Resolve(x, ans);
       Resolve(y, ans);
   }
   long long check(ll x) { //用这个,素数返回本身
       11 \text{ ans} = 0;
       Resolve(x, ans);
```

```
return ans;
   }
}
FFT
#include <iostream>
#include <cstring>
#include <algorithm>
#include <cmath>
using namespace std;
const int N = 300010;
const double PI = acos(-1);
int n, m;
struct Complex
{
   double x, y;
   Complex operator+ (const Complex& t) const
       return \{x + t.x, y + t.y\};
   Complex operator- (const Complex& t) const
       return {x - t.x, y - t.y};
   Complex operator* (const Complex& t) const
       return {x * t.x - y * t.y, x * t.y + y * t.x};
}a[N], b[N];
int rev[N], bit, tot;
void fft(Complex a[], int inv)
{
   for (int i = 0; i < tot; i ++ )</pre>
       if (i < rev[i])
           swap(a[i], a[rev[i]]);
   for (int mid = 1; mid < tot; mid <<= 1)</pre>
       auto w1 = Complex({cos(PI / mid), inv * sin(PI / mid)});
       for (int i = 0; i < tot; i += mid * 2)</pre>
           auto wk = Complex(\{1, \emptyset\});
           for (int j = 0; j < mid; j ++, wk = wk * w1)
```

```
auto x = a[i + j], y = wk * a[i + j + mid];
              a[i + j] = x + y, a[i + j + mid] = x - y;
           }
       }
   }
}
int main()
   scanf("%d%d", &n, &m);
   for (int i = 0; i <= n; i ++ ) scanf("%lf", &a[i].x);</pre>
   for (int i = 0; i <= m; i ++ ) scanf("%lf", &b[i].x);</pre>
   while ((1 << bit) < n + m + 1) bit ++;
   tot = 1 << bit;
   for (int i = 0; i < tot; i ++ )</pre>
       rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (bit - 1));
   fft(a, 1), fft(b, 1);
   for (int i = 0; i < tot; i ++ ) a[i] = a[i] * b[i];</pre>
   fft(a, -1);
   for (int i = 0; i <= n + m; i ++ )</pre>
       printf("%d ", (int)(a[i].x / tot + 0.5));
   return 0;
}
作者: yxc
链接: https://www.acwing.com/activity/content/code/content/664840/
来源: AcWing
著作权归作者所有。商业转载请联系作者获得授权,非商业转载请注明出处。
BSGS
求a^t \equiv b \pmod{p} (a,p) = 1 的最小的 t
t \in [1, k^2]
对 b \times a^y 建立 hash 表, 枚举 x 看是否有解
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
unordered map<int , int> mp;
int bsgs(int a, int p, int b) {
      if (1 % p == b % p) return 0; // 特判0 是不是解
```

```
mp.clear();
       int k = sqrt(p) + 1;
      for(int i = 0, j = b \% p; i < k; ++ i, j = (11)j * a % p) {
             mp[j] = i;
       }
       int ak = 1;
       for(int i = 0; i < k; ++i) {</pre>
             ak = (11)ak * a % p;
       }
      for(int i = 1, j = ak % p; i <= k; ++ i, j = (l1)j * ak % p) {
              if(mp.count(j)) return (ll)i * k - mp[j];
       }
      return -1;
}
int main() {
       ios::sync_with_stdio(∅);
      cin.tie(0); cout.tie(0);
       int a, p, b;
      while(cin >> a >> p >> b, a | p | b) {
              int res;
              res = bsgs(a, p, b);
              if(res == -1) {
                     cout << "No Solution\n";</pre>
              }
              else {
                     cout << res << endl;</pre>
              }
       }
       return 0;
}
扩展 BSGS
求a^t \equiv b(\text{mod}p) 的最小的 t
当(a, p)! = 1
(a,p) = d d \nmid b 无解
a^t \equiv b \pmod{p} , a^t + kp = b 两边同时除以 d, \frac{a}{d}a^{t-1} + k\frac{p}{d} = \frac{b}{d}
```

```
t' = t - 1, p' = \frac{p}{d}, b' = \frac{b}{a} (\frac{a}{d})^{-1}
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
unordered_map<11, 11> mp;
11 bsgs(ll a, ll p, ll b) {
       if(1 % p == b % p) return 0; // 特判0 是不是解
      mp.clear();
      11 k = sqrt(p) + 1;
      for(11 i = 0, j = b \% p; i < k; ++i, j = (11)j * a \% p) {
             mp[j] = i;
       }
      11 ak = 1;
      for(ll i = 0; i < k; ++i) {</pre>
             ak = (11) ak * a % p;
      for(ll i = 1, j = ak % p;i <= k; ++i, j = (ll)j * ak % p) {</pre>
             if(mp.count(j)) return (ll) i * k - mp[j];
       }
       return -1;
}
11 gcd(ll x, ll y) {
       return x % y == 0 ? y : gcd(y, x % y);
}
void extgcd(l1 a,l1 b,l1& d,l1& x,l1& y){
    if(!b){
        d = a; x = 1; y = 0;
    else{
        extgcd(b, a%b, d, y, x);
        y -= x * (a / b);
    }
}
11 inverse(ll a,ll n){
    11 d,x,y;
```

```
extgcd(a,n,d,x,y);
    return d == 1 ? (x + n) % n : -1;
}
int main() {
      11 a, p, b;
      while(cin >> a >> p >> b, a | p | b) {
             11 d = gcd(a, p);
             if(d == 1) {
                    11 res = bsgs(a, p, b);
                    if(res == -1) {
                          cout << "No Solution\n";</pre>
                    }
                    else {
                          cout << res << endl;</pre>
                    }
             }
             else {
                    if(b % d != 0) {
                          cout << "No Solution\n";</pre>
                          continue;
                    }
                    else {
                          p = p / d;
                          b = (b / d) * inverse(a / d, p);
                          11 res = bsgs(a, p, b);
                          if(res == -1) {
                                 cout << "No Solution\n";</pre>
                          }
                          else {
                                 cout << res + 1 << endl;</pre>
                          }
                   }
             }
      }
      return 0;
}
二次剩余
```

解的数量

对于 $x^2 \equiv n \pmod{p}$ 能满足 n 是 mod p 的二次剩余的 n 一共有 $\frac{p-1}{2}$ 个(不包括 0),非二次剩余为 $\frac{p-1}{2}$ 个

勒让德符号

欧拉判别准则

若 n 是二次剩余,当且仅当 $n^{\frac{p-1}{2}} \equiv 1 \pmod{p}$

若 n 是非二次剩余,当且仅当 $n^{\frac{p-1}{2}} \equiv -1 \pmod{p}$

Cipolla

找到一个数 a 满足 a^2-n 是 **非二次剩余** ,至于为什么要找满足非二次剩余的数,在下文会给出解释。 这里通过生成随机数再检验的方法来实现,由于非二次剩余的数量为 $\frac{p-1}{2}$,接近 $\frac{p}{2}$,所以期望约 2 次就可以找到这个数。

建立一个 "复数域",并不是实际意义上的复数域,而是根据复数域的概念建立的一个类似的域。 在复数中 $i^2=-1$,这里定义 $i^2=a^2-n$,于是就可以将所有的数表达为A+Bi 的形式,这里的 和 都是模意义下的数,类似复数中的实部和虚部。

在有了 i 和 a 后可以直接得到答案, $x^2 \equiv n \pmod{p}$ 的解为 $(a + i)^{\frac{p+1}{2}}$ 。

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
int t;
11 n, p;
11 w;
struct num { //建立一个复数域
     11 x, y;
};
num mul(num a, num b, ll p) { //复数乘法
     num ans = \{0, 0\};
     ans.x = ((a.x * b.x \% p + a.y * b.y \% p * w \% p) \% p + p) \% p;
     ans.y = ((a.x * b.y % p + a.y * b.x % p) % p + p) % p;
     return ans;
}
ll binpow real(ll a, ll b, ll p) { //实部快速幂
     11 \text{ ans} = 1;
     while (b) {
           if (b & 1) ans = ans * a % p;
           a = a * a % p;
           b >>= 1;
     return ans % p;
}
ll binpow_imag(num a, ll b, ll p) { //虚部快速幂
     num ans = \{1, 0\};
     while (b) {
           if (b & 1) ans = mul(ans, a, p);
           a = mul(a, a, p);
           b >>= 1;
     }
     return ans.x % p;
}
11 cipolla(11 n, 11 p) {
     n %= p;
     if (p == 2) return n;
     if (binpow_real(n, (p - 1) / 2, p) == p - 1) return -1;
     11 a;
     while (1) { //生成随机数再检验找到满足非二次剩余的 a
           a = rand() \% p;
           w = ((a * a % p - n) % p + p) % p;
```

```
if (binpow_real(w, (p - 1) / 2, p) == p - 1) break;
       }
      num x = \{a, 1\};
      return binpow_imag(x, (p + 1) / 2, p);
}
卡特兰数
卡特兰数 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012,...
C_n = \frac{1}{n+1} \sum_{i=0}^{n} (C_n^i)^2
C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i} (C_0 = 1)
超级卡特兰数 1, 1, 3, 11, 45, 197, 903, 4279, 20793, 103049,... (从第 0 项开始)
大施罗德数(OEIS A006318)1, 2, 6, 22, 90, 394, 1806, 8558, 41586, 206098,...
超级卡特兰数的两倍(除第一项)
快速幂
11 qpow(ll a, ll b) {
    11 \text{ ans} = 1;
    while (b) {
        if (b & 1) ans = (ans * a) % mod;
        a = (a * a) % mod;
        b >>= 1;
    }
    return ans;
}
```

龟速乘快速幂(快速幂爆 longlong

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
11 qmul(11 a, 11 b, 11 p) {
      11 \text{ res} = 0;
      while(b) {
            if(b & 1) res = (res + a) % p;
            a = (a + a) \% p;
            b >>= 1;
      return res;
}
11 qpow(11 x, 11 n, 11 p) {
      11 \text{ res} = 1;
      while(n) {
            if(n & 1) res = qmul(res, x, p);
            x = qmul(x, x, p);
            n \gg 1;
      return res % p; // 1 0 1
}
int main() {
      11 b, p, k;
      cin \gg b \gg p \gg k;
      11 ans = qpow(b, p, k);
      printf("%11d^%11d mod %11d=%11d", b, p, k, ans);
      return 0;
}
莫比乌斯反演
```

莫比乌斯函数

n的所有约数的莫比乌斯的和

反演

$$(\sqrt{2}). \, \overline{Z}F(n) = \sum_{n/d} f(d), \, \mathcal{D}f(n) = \sum_{n/d} \mu(\frac{d}{n})F(d)$$

```
构造F(n)和f(n) 使 f(n)为目标,F(n)好求
1
求满足a \le x \le b, c \le y \le d 且 gcd(x, y) = k 的 xy 的对数
f(n) = gcd(x, y) = n 的xy的对数
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const int N = 50010;
11 primes[N], mu[N], sum[N], cnt;
bool st[N];
void init() {
      mu[1] = 1;
      for(int i = 2; i < N; ++ i) {</pre>
            if(!st[i]) {
                   primes[cnt ++] = i;
                   mu[i] = -1;
            }
            for(int j = 0; primes[j] * i < N; ++ j) {</pre>
                   st[primes[j] * i] = 1;
                   if(i % primes[j] == 0) break;
                   mu[primes[j] * i] = -mu[i];
            }
      }
      for(int i = 1; i < N; ++ i) {</pre>
            sum[i] = sum[i - 1] + mu[i];
      }
}
ll g(ll n, ll x) {
      return n / (n / x);
}
11 f (int a, int b, int k) {
      a = a / k, b = b / k;
      11 \text{ res} = 0;
      11 n = min(a, b);
```

```
for(11 1 = 1, r; 1 \le n; 1 = r + 1) {
              r = min(n, min(g(a, 1), g(b, 1)));
              res += (sum[r] - sum[l - 1]) * (a / l) * (b / l);
       }
       return res;
}
int main() {
       ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
       init();
       int T;
       cin >> T;
       while(T --) {
              int a, b, c, d, k;
              cin >> a >> b >> c >> d >> k;
              cout << f(b, d, k) - f(a - 1, d, k) - f(b, c - 1, k)
                            + f(a - 1, c - 1, k) << endl;
       }
       return 0;
}
2
求\sum_{i=1}^{N}\sum_{j=1}^{M}d\left( ij\right)
// d(ij) = \sum_{x|i} \sum_{y|j} [(x, y) = 1]
f(n) = \sum_{i=1}^{N} \sum_{j=1}^{M} \sum_{x/i} \sum_{y/j} [(x, y) = n]
两次整数分块
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const int N = 50010;
int primes[N], cnt, mu[N], sum[N], h[N];
bool st[N];
inline int g(int n, int x) {
       return n / (n / x);
}
```

```
void init() {
      mu[1] = 1;
      for(int i = 2; i < N; ++i) {</pre>
            if(!st[i]){
                   primes[cnt++] = i;
                   mu[i] = -1;
            }
            for(int j = 0; primes[j] * i < N; ++j) {</pre>
                   st[primes[j] * i] = 1;
                   if(i % primes[j] == 0) break;
                   mu[primes[j] * i] = -mu[i];
            }
      }
      for(int i = 1; i < N; ++ i) {</pre>
            sum[i] = sum[i - 1] + mu[i];
      }
      for(int i = 1; i < N; ++i) {</pre>
            for(int l = 1, r; l \leftarrow i; l = r + 1) {
                   r = min(i, g(i, 1));
                   h[i] += (r - l + 1) * (i / l);
            }
      }
}
int main() {
      //ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
      init();
      int T;
      scanf("%d", &T);
      while(T--) {
            int n, m;
            scanf("%d %d", &n, &m);
            11 \text{ res} = 0;
            int k = min(n, m);
            for(int l = 1, r; l <= k; l = r + 1) {
                   r = min(k, min(g(n, 1), g(m, 1)));
                   res += (11)(sum[r] - sum[1 - 1]) * h[n / 1] * h[m /
1];
          printf("%lld\n", res);
      }
```

```
return 0;
}
```

博弈

SG 定理:

mex(minimal excludant)运算,表示最小的不属于这个集合的非负整数。例如 $mex\{0,1,2,4\}=3$ 、 $mex\{2,3,5\}=0$ 、 $mex\{\}=0$ 。

Sprague-Grundy 定理(SG 定理): 游戏和的 SG 函数等于各个游戏 SG 函数的 Nim 和。这样就可以将每一个子游戏分而治之,从而简化了问题。而 Bouton 定理就是 Sprague-Grundy 定理在 Nim 游戏中的直接应用,因为单堆的 Nim 游戏 SG 函数满足 SG(\mathbf{x}) = \mathbf{x} 。

Nimk:

普通的 NIM 游戏是在 n 堆石子中每次选一堆,取任意个石子,而 NIMK 游戏是在 n 堆石子中每次选择 k 堆, 1 <= k <= n,从这 k 堆中每堆里都取出任意数目的石子,取的石子数可以不同,其他规则相同。

对于普通的 NIM 游戏,我们采取的是对每堆的 SG 值进行异或,异或其实就是对每一个 SG 值二进制位上的数求和然后模 2,比如说 3^5 就是 011+101=112,然后对每一位都模 2 就变成了 110,所以 3^5=6。而 NIMK 游戏和 NIM 游戏的区别就在于模的不是 2,如果是取 k 堆,就模 k+1,所以取 1 堆的普通 NIM 游戏是模 2。当 k=2 时,3^5 \rightarrow 011+101=112,对每一位都模 3 之后三位二进制位上对应的数仍然是1,1,2。那么当且仅当每一位二进制位上的数都是 0 的时候,先手必败,否则先手必胜。

anti_nim

描述

和最普通的 Nim 游戏相同,不过是取走最后一个石子的人输。

先手必胜条件

以下两个条件满足其一即可:

- 1. 所有堆的石子个数=1,且异或和=0(其实这里就是有偶数堆的意思)。
- 2. 至少存在一堆石子个数>1,且异或和≠0。

高精度 GCD

```
#include <bits/stdc++.h>
using namespace std;
string add(string a, string b) {
```

```
const int L = 1e5;
   string ans;
   int na[L] = {0}, nb[L] = {0};
   int la = a.size(), lb = b.size();
   for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0';</pre>
   for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0';</pre>
   int lmax = la > lb ? la : lb;
   for (int i = 0; i < lmax; i++)</pre>
       na[i] += nb[i], na[i + 1] += na[i] / 10, na[i] %= 10;
   if (na[lmax]) lmax++;
   for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
   return ans;
}
string mul(string a, string b) {
   const int L = 1e5;
   string s;
   int na[L], nb[L], nc[L],
       La = a.size(), Lb = b.size(); // na 存储被乘数, nb 存储乘数, nc 存
储积
   fill(na, na + L, \emptyset);
   fill(nb, nb + L, \emptyset);
   fill(nc, nc + L, 0); //将na,nb,nc 都置为0
   for (int i = La - 1; i >= 0; i--)
       na[La - i] =
          a[i] - '0'; //将字符串表示的大整形数转成 i 整形数组表示的大整形数
   for (int i = Lb - 1; i >= 0; i--) nb[Lb - i] = b[i] - '0';
   for (int i = 1; i <= La; i++)</pre>
       for (int j = 1; j <= Lb; j++)</pre>
          nc[i + j - 1] +=
              na[i] *
              nb[j]; // a 的第 i 位乘以 b 的第 j 位为积的第 i+j-1 位 ( 先不考虑
讲位)
   for (int i = 1; i <= La + Lb; i++)
       nc[i + 1] += nc[i] / 10, nc[i] %= 10; //统一处理进位
   if (nc[La + Lb]) s += nc[La + Lb] + '0'; //判断第 i+j 位上的数字是不是
0
   for (int i = La + Lb - 1; i >= 1; i--)
       s += nc[i] + '0'; //将整形数组转成字符串
   return s;
int sub(int *a, int *b, int La, int Lb) {
   if (La < Lb) return -1; //如果 a 小于 b,则返回-1
   if (La == Lb) {
       for (int i = La - 1; i >= 0; i--)
           if (a[i] > b[i])
              break;
          else if (a[i] < b[i])
              return -1; //如果 a 小于 b , 则返回-1
   }
```

```
for (int i = 0; i < La; i++) //高精度减法
      a[i] -= b[i];
      if (a[i] < 0) a[i] += 10, a[i + 1] --;
   for (int i = La - 1; i >= 0; i--)
      if (a[i]) return i + 1; //返回差的位数
   return 0;
                             //返回差的位数
string div(string n1, string n2,
         int nn) // n1, n2 是字符串表示的被除数,除数, nn 是选择返回商还是余
数
{
   const int L = 1e5;
   string s, v; // s 存商, v 存余数
   int a[L], b[L], r[L],
      La = n1.size(), Lb = n2.size(), i,
      tp = La; // a, b 是整形数组表示被除数,除数,tp 保存被除数的长度
   fill(a, a + L, 0);
   fill(b, b + L, \emptyset);
   fill(r, r + L, 0); //数组元素都置为0
   for (i = La - 1; i >= 0; i--) a[La - 1 - i] = n1[i] - '0';
   for (i = Lb - 1; i >= 0; i--) b[Lb - 1 - i] = n2[i] - '0';
   if (La < Lb | (La == Lb && n1 < n2)) {
      // cout<<0<<endl;</pre>
      return n1;
                   //如果 a < b , 则商为 0 , 余数为被除数
   int t = La - Lb; //除被数和除数的位数之差
   for (int i = La - 1; i >= 0; i--) //将除数扩大10^t 倍
      if (i >= t)
          b[i] = b[i - t];
      else
          b[i] = 0;
   Lb = La;
   for (int j = 0; j <= t; j++) {
      int temp;
      while ((temp = sub(a, b + j, La, Lb - j)) >=
            0) //如果被除数比除数大继续减
      {
          La = temp;
          r[t - j]++;
       }
   for (i = 0; i < L - 10; i++)
      r[i + 1] += r[i] / 10, r[i] %= 10; //统一处理进位
   while (!r[i]) i--; //将整形数组表示的商转化成字符串表示的
   while (i >= 0) s += r[i--] + '0';
   // cout<<s<<endl;</pre>
```

```
i = tp;
   while (!a[i]) i--; //将整形数组表示的余数转化成字符串表示的</span>
   while (i >= 0) v += a[i--] + '0';
   if (v.empty()) v = "0";
   // cout<<v<<endl;</pre>
   if (nn == 1) return s;
   if (nn == 2) return v;
}
bool judge(string s) //判断 s 是否为全 0 串
   for (int i = 0; i < s.size(); i++)</pre>
       if (s[i] != '0') return false;
   return true;
string gcd(string a, string b) //求最大公约数
{
   string t;
   while (!judge(b)) //如果余数不为0,继续除
                        //保存被除数的值
       t = a;
                        //用除数替换被除数
       a = b;
       b = div(t, b, 2); //用余数替换除数
   return a;
}
//o(无法估计)
高精度乘法(FFT)
#include <bits/stdc++.h>
using namespace std;
#define L(x) (1 << (x))
const double PI = acos(-1.0);
const int Maxn = 133015;
double ax[Maxn], ay[Maxn], bx[Maxn], by[Maxn];
char sa[Maxn / 2], sb[Maxn / 2];
int sum[Maxn];
int x1[Maxn], x2[Maxn];
int revv(int x, int bits) {
   int ret = 0;
   for (int i = 0; i < bits; i++) {</pre>
       ret <<= 1;
       ret |= x \& 1;
       x >>= 1;
   }
   return ret;
}
void fft(double* a, double* b, int n, bool rev) {
   int bits = 0;
```

```
while (1 << bits < n) ++bits;
   for (int i = 0; i < n; i++) {
       int j = revv(i, bits);
       if (i < j) swap(a[i], a[j]), swap(b[i], b[j]);</pre>
    for (int len = 2; len <= n; len <<= 1) {</pre>
       int half = len >> 1;
       double wmx = cos(2 * PI / len), wmy = sin(2 * PI / len);
       if (rev) wmy = -wmy;
       for (int i = 0; i < n; i += len) {</pre>
           double wx = 1, wy = 0;
           for (int j = 0; j < half; j++) {</pre>
               double cx = a[i + j], cy = b[i + j];
               double dx = a[i + j + half], dy = b[i + j + half];
               double ex = dx * wx - dy * wy, ey = dx * wy + dy * wx;
               a[i + j] = cx + ex, b[i + j] = cy + ey;
               a[i + j + half] = cx - ex, b[i + j + half] = cy - ey;
               double wnx = wx * wmx - wy * wmy, wny = wx * wmy + wy * wm
х;
               wx = wnx, wy = wny;
           }
        }
    if (rev) {
       for (int i = 0; i < n; i++) a[i] /= n, b[i] /= n;
    }
int solve(int a[], int na, int b[], int nb, int ans[]) {
    int len = max(na, nb), ln;
    for (ln = 0; L(ln) < len; ++ln)
       ;
    len = L(++ln);
    for (int i = 0; i < len; ++i) {</pre>
       if (i >= na)
           ax[i] = 0, ay[i] = 0;
       else
           ax[i] = a[i], ay[i] = 0;
   fft(ax, ay, len, ∅);
    for (int i = 0; i < len; ++i) {</pre>
       if (i >= nb)
           bx[i] = 0, by[i] = 0;
       else
           bx[i] = b[i], by[i] = 0;
   fft(bx, by, len, ∅);
    for (int i = 0; i < len; ++i) {</pre>
       double cx = ax[i] * bx[i] - ay[i] * by[i];
       double cy = ax[i] * by[i] + ay[i] * bx[i];
        ax[i] = cx, ay[i] = cy;
```

```
fft(ax, ay, len, 1);
   for (int i = 0; i < len; ++i) ans[i] = (int)(ax[i] + 0.5);
   return len;
string mul(string sa, string sb) {
    int 11, 12, 1;
    int i;
    string ans;
   memset(sum, 0, sizeof(sum));
   11 = sa.size();
   12 = sb.size();
   for (i = 0; i < 11; i++) \times 1[i] = sa[11 - i - 1] - '0';
   for (i = 0; i < 12; i++) \times 2[i] = sb[12 - i - 1] - '0';
    1 = solve(x1, 11, x2, 12, sum);
   for (i = 0; i < l | sum[i] >= 10; i++) // 进位
       sum[i + 1] += sum[i] / 10;
       sum[i] %= 10;
    }
    l = i;
                                                 // 检索最高位
   while (sum[1] <= 0 && 1 > 0) 1--;
   for (i = 1; i >= 0; i--) ans += sum[i] + '0'; // 倒序输出
   return ans;
int main() {
   cin.sync_with_stdio(false);
    string a, b;
   while (cin >> a >> b) cout << mul(a, b) << endl;</pre>
   return 0;
}
//o(nlogn)
高精度乘法(乘单精度
#include <bits/stdc++.h>
using namespace std;
string mul(string a, int b) //高精度a 乘单精度b
{
    const int L = 100005;
    int na[L];
    string ans;
    int La = a.size();
   fill(na, na + L, \emptyset);
   for (int i = La - 1; i >= 0; i--) na[La - i - 1] = a[i] - '0';
    int w = 0;
   for (int i = 0; i < La; i++)</pre>
       na[i] = na[i] * b + w, w = na[i] / 10, na[i] = na[i] % 10;
   while (w) na[La++] = w % 10, w /= 10;
```

```
La--;
   while (La >= 0) ans += na[La--] + '0';
   return ans;
}
//o(n)
高精度乘法(朴素)
#include <bits/stdc++.h>
using namespace std;
string mul(string a, string b) //高精度乘法a,b,均为非负整数
   const int L = 1e5;
   string s;
   int na[L], nb[L], nc[L],
       La = a.size(), Lb = b.size(); // na 存储被乘数, nb 存储乘数, nc 存
储积
   fill(na, na + L, \emptyset);
   fill(nb, nb + L, \emptyset);
   fill(nc, nc + L, 0); //将na,nb,nc 都置为0
   for (int i = La - 1; i >= 0; i--)
       na[La - i] =
          a[i] - '0'; //将字符串表示的大整形数转成i整形数组表示的大整形数
   for (int i = Lb - 1; i >= 0; i--) nb[Lb - i] = b[i] - '0';
   for (int i = 1; i <= La; i++)</pre>
       for (int j = 1; j <= Lb; j++)</pre>
          nc[i + j - 1] +=
              na[i] *
              nb[j]; // a 的第 i 位乘以 b 的第 j 位为积的第 i+j-1 位 ( 先不考虑
讲位)
   for (int i = 1; i <= La + Lb; i++)</pre>
       nc[i + 1] += nc[i] / 10, nc[i] %= 10; //统一处理进位
   if (nc[La + Lb]) s += nc[La + Lb] + '0'; //判断第 i+j 位上的数字是不是
0
   for (int i = La + Lb - 1; i >= 1; i--)
       s += nc[i] + '0'; //将整形数组转成字符串
   return s;
}
//o(n^2)
高精度除法 (除单精度)
#include <bits/stdc++.h>
using namespace std;
string div(string a, int b) //高精度a 除以单精度b
   string r, ans;
   int d = 0;
```

```
if (a == "0") return a; //特判
for (int i = 0; i < a.size(); i++) {
    r += (d * 10 + a[i] - '0') / b + '0'; //求出商
    d = (d * 10 + (a[i] - '0')) % b; //求出余数
}
int p = 0;
for (int i = 0; i < r.size(); i++)
    if (r[i] != '0') {
        p = i;
        break;
    }
return r.substr(p);
}
```

高精度除法 (除高精度) #include <bits/stdc++.h> using namespace std; int sub(int *a, int *b, int La, int Lb) { if (La < Lb) return -1; //如果 a 小于 b,则返回-1 **if** (La == Lb) { for (int i = La - 1; i >= 0; i--) **if** (a[i] > b[i]) break; **else if** (a[i] < b[i]) return -1; //如果 a 小于 b ,则返回-1 for (int i = 0; i < La; i++) //高精度减法 { a[i] -= b[i];if (a[i] < 0) a[i] += 10, a[i + 1]--;</pre> for (int i = La - 1; i >= 0; i--) if (a[i]) return i + 1; //返回差的位数 //返回差的位数 return 0; } string div(string n1, string n2, int nn) // n1,n2 是字符串表示的被除数,除数,nn 是选择返回商还是余数 { const int L = 1e5; string s, v; // s 存商, v 存余数 int a[L], b[L], r[L], La = n1.size(), Lb = n2.size(), i, tp = La; // a,b 是整形数组表示被除数,除数,tp 保存被除数的长度 fill(a, a + L, 0);fill(b, b + L, \emptyset); fill(r, r + L, 0); //数组元素都置为0

```
for (i = La - 1; i >= 0; i--) a[La - 1 - i] = n1[i] - '0';
   for (i = Lb - 1; i >= 0; i--) b[Lb - 1 - i] = n2[i] - '0';
   if (La < Lb || (La == Lb && n1 < n2)) {</pre>
      // cout<<0<<endl;</pre>
       return n1;
                   //如果 a<b,则商为0,余数为被除数
   int t = La - Lb; //除被数和除数的位数之差
   for (int i = La - 1; i >= 0; i--) //将除数扩大10^t 倍
       if (i >= t)
          b[i] = b[i - t];
       else
          b[i] = 0;
   Lb = La;
   for (int j = 0; j <= t; j++) {</pre>
       int temp;
       while ((temp = sub(a, b + j, La, Lb - j)) >=
             0) //如果被除数比除数大继续减
       {
          La = temp;
          r[t - j]++;
       }
   for (i = 0; i < L - 10; i++)
       r[i + 1] += r[i] / 10, r[i] %= 10; //统一处理进位
   while (!r[i]) i--; //将整形数组表示的商转化成字符串表示的
   while (i >= 0) s += r[i--] + '0';
   // cout<<s<<endl;</pre>
   i = tp;
   while (!a[i]) i--; //将整形数组表示的余数转化成字符串表示的</span>
   while (i >= 0) v += a[i--] + '0';
   if (v.empty()) v = "0";
   // cout<<v<<endl;</pre>
   if (nn == 1) return s; //返回商
   if (nn == 2) return v; //返回余数
}
//o(n^2)
```

高精度加法

```
#include <bits/stdc++.h>
using namespace std;
string add(string a, string b) //只限两个非负整数相加
{
    const int L = 1e5;
    string ans;
    int na[L] = {0}, nb[L] = {0};
    int la = a.size(), lb = b.size();
```

```
for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0';</pre>
   for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0';
    int lmax = la > lb ? la : lb;
    for (int i = 0; i < lmax; i++)</pre>
       na[i] += nb[i], na[i + 1] += na[i] / 10, na[i] %= 10;
    if (na[lmax]) lmax++;
    for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
    return ans:
}
//o(n)
高精度减法
#include <bits/stdc++.h>
using namespace std;
string sub(string a, string b) // 只限大的非负整数减小的非负整数
{
    const int L = 1e5;
    string ans;
    int na[L] = \{0\}, nb[L] = \{0\};
    int la = a.size(), lb = b.size();
   for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0';
   for (int i = 0; i < 1b; i++) nb[1b - 1 - i] = b[i] - '0';
    int lmax = la > lb ? la : lb;
    for (int i = 0; i < lmax; i++) {</pre>
       na[i] -= nb[i];
       if (na[i] < 0) na[i] += 10, na[i + 1]--;</pre>
   while (!na[--lmax] && lmax > 0)
    lmax++;
    for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
    return ans;
}
//o(n)
高精度阶乘
#include <bits/stdc++.h>
using namespace std;
string fac(int n) {
   const int L = 100005;
   int a[L];
    string ans;
    if (n == 0) return "1";
   fill(a, a + L, \emptyset);
    int s = 0, m = n;
```

while (m) a[++s] = m % 10, m /= 10;

```
for (int i = n - 1; i >= 2; i--) {
    int w = 0;
    for (int j = 1; j <= s; j++)
        a[j] = a[j] * i + w, w = a[j] / 10, a[j] = a[j] % 10;
    while (w) a[++s] = w % 10, w /= 10;
}
while (!a[s]) s--;
while (s >= 1) ans += a[s--] + '0';
return ans;
}
//o(n^2)
```

高精度进制转换

```
#include <bits/stdc++.h>
using namespace std;
//将字符串表示的10 进制大整数转换为m 进制的大整数
// 并返回 m 进制大整数的字符串
bool judge(string s) //判断串是否为全零串
   for (int i = 0; i < s.size(); i++)</pre>
      if (s[i] != '0') return 1;
   return 0;
}
string solve(
   string s, int n,
   int m) // n 进制转 m 进制只限 0-9 进制, 若涉及带字母的进制, 稍作修改即可
{
   string r, ans;
   int d = 0;
   if (!judge(s)) return "0"; //特判
   while (judge(s))
                           //被除数不为 0 则继续
   {
      for (int i = 0; i < s.size(); i++) {</pre>
         r += (d * n + s[i] - '0') / m + '0'; //求出商
         d = (d * n + (s[i] - '0')) % m;
                                          //求出余数
      }
      s = r;
                    //把商赋给下一次的被除数
      r = "";
                    //把商清空
      ans += d + '0'; //加上进制转换后数字
                    //清空余数
      d = 0;
   reverse(ans.begin(), ans.end()); //倒置下
   return ans;
}
```

```
高精度幂
#include <bits/stdc++.h>
#define L(x) (1 << (x))
using namespace std;
const double PI = acos(-1.0);
const int Maxn = 133015;
double ax[Maxn], ay[Maxn], bx[Maxn], by[Maxn];
char sa[Maxn / 2], sb[Maxn / 2];
int sum[Maxn];
int x1[Maxn], x2[Maxn];
int revv(int x, int bits) {
    int ret = 0;
    for (int i = 0; i < bits; i++) {</pre>
        ret <<= 1;
       ret |= x \& 1;
       x >>= 1;
    }
    return ret;
void fft(double* a, double* b, int n, bool rev) {
    int bits = 0;
   while (1 << bits < n) ++bits;
    for (int i = 0; i < n; i++) {</pre>
        int j = revv(i, bits);
        if (i < j) swap(a[i], a[j]), swap(b[i], b[j]);</pre>
    for (int len = 2; len <= n; len <<= 1) {</pre>
        int half = len >> 1;
       double wmx = cos(2 * PI / len), wmy = sin(2 * PI / len);
        if (rev) wmy = -wmy;
        for (int i = 0; i < n; i += len) {</pre>
           double wx = 1, wy = 0;
           for (int j = 0; j < half; j++) {</pre>
               double cx = a[i + j], cy = b[i + j];
               double dx = a[i + j + half], dy = b[i + j + half];
               double ex = dx * wx - dy * wy, ey = dx * wy + dy * wx;
               a[i + j] = cx + ex, b[i + j] = cy + ey;
               a[i + j + half] = cx - ex, b[i + j + half] = cy - ey;
               double wnx = wx * wmx - wy * wmy, wny = wx * wmy + wy * wm
х;
               wx = wnx, wy = wny;
           }
        }
    if (rev) {
```

```
for (int i = 0; i < n; i++) a[i] /= n, b[i] /= n;
   }
int solve(int a[], int na, int b[], int nb, int ans[]) {
   int len = max(na, nb), ln;
   for (ln = 0; L(ln) < len; ++ln)</pre>
       ;
   len = L(++ln);
   for (int i = 0; i < len; ++i) {</pre>
       if (i >= na)
           ax[i] = 0, ay[i] = 0;
       else
           ax[i] = a[i], ay[i] = 0;
   fft(ax, ay, len, ∅);
   for (int i = 0; i < len; ++i) {</pre>
       if (i >= nb)
           bx[i] = 0, by[i] = 0;
       else
           bx[i] = b[i], by[i] = 0;
   fft(bx, by, len, ∅);
   for (int i = 0; i < len; ++i) {</pre>
       double cx = ax[i] * bx[i] - ay[i] * by[i];
       double cy = ax[i] * by[i] + ay[i] * bx[i];
       ax[i] = cx, ay[i] = cy;
   fft(ax, ay, len, 1);
   for (int i = 0; i < len; ++i) ans[i] = (int)(ax[i] + 0.5);
   return len;
}
string mul(string sa, string sb) {
   int 11, 12, 1;
   int i;
   string ans;
   memset(sum, 0, sizeof(sum));
   11 = sa.size();
   12 = sb.size();
   for (i = 0; i < 11; i++) \times 1[i] = sa[11 - i - 1] - '0';
   for (i = 0; i < 12; i++) \times 2[i] = sb[12 - i - 1] - '0';
   1 = solve(x1, 11, x2, 12, sum);
   for (i = 0; i < l || sum[i] >= 10; i++) // 进位
       sum[i + 1] += sum[i] / 10;
       sum[i] \% = 10;
   l = i;
                                                  // 检索最高位
   while (sum[1] \leftarrow 0 \&\& 1 > 0) 1--;
   for (i = 1; i >= 0; i--) ans += sum[i] + '0'; // 倒序输出
```

```
return ans;
}
string Pow(string a, int n) {
   if (n == 1) return a;
   if (n & 1) return mul(Pow(a, n - 1), a);
   string ans = Pow(a, n / 2);
   return mul(ans, ans);
}
//o(nlognlogm)
高精度平方根
#include <bits/stdc++.h>
using namespace std;
const int L = 2015;
string add(string a, string b) // 只限两个非负整数相加
   string ans;
   int na[L] = {0}, nb[L] = {0};
   int la = a.size(), lb = b.size();
   for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0';</pre>
   for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0';</pre>
   int lmax = la > lb ? la : lb;
   for (int i = 0; i < lmax; i++)</pre>
       na[i] += nb[i], na[i + 1] += na[i] / 10, na[i] %= 10;
   if (na[lmax]) lmax++;
   for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
   return ans;
}
string sub(string a, string b) // 只限大的非负整数减小的非负整数
   string ans;
   int na[L] = {0}, nb[L] = {0};
   int la = a.size(), lb = b.size();
   for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0';
   for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0';</pre>
   int lmax = la > lb ? la : lb;
   for (int i = 0; i < lmax; i++) {</pre>
       na[i] -= nb[i];
       if (na[i] < 0) na[i] += 10, na[i + 1]--;</pre>
   while (!na[--lmax] && lmax > 0)
   lmax++;
   for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
   return ans;
}
string mul(string a, string b) //高精度乘法a,b,均为非负整数
```

```
string s;
   int na[L], nb[L], nc[L],
       La = a.size(), Lb = b.size(); // na 存储被乘数, nb 存储乘数, nc 存
储积
   fill(na, na + L, \emptyset);
   fill(nb, nb + L, \emptyset);
   fill(nc, nc + L, 0); //将na,nb,nc 都置为0
   for (int i = La - 1; i >= 0; i--)
       na[La - i] =
          a[i] - '0'; //将字符串表示的大整形数转成 i 整形数组表示的大整形数
   for (int i = Lb - 1; i >= 0; i--) nb[Lb - i] = b[i] - '0';
   for (int i = 1; i <= La; i++)</pre>
       for (int j = 1; j <= Lb; j++)</pre>
          nc[i + j - 1] +=
              na[i] *
              nb[j]; // a 的第 i 位乘以 b 的第 j 位为积的第 i+j-1 位 ( 先不考虑
进位)
   for (int i = 1; i <= La + Lb; i++)</pre>
       nc[i + 1] += nc[i] / 10, nc[i] %= 10; //统一处理进位
   if (nc[La + Lb]) s += nc[La + Lb] + '0'; //判断第 i+j 位上的数字是不是
0
   for (int i = La + Lb - 1; i >= 1; i--)
       s += nc[i] + '0'; //将整形数组转成字符串
   return s;
int sub(int *a, int *b, int La, int Lb) {
   if (La < Lb) return -1; //如果 a 小于 b ,则返回-1
   if (La == Lb) {
       for (int i = La - 1; i >= 0; i--)
          if (a[i] > b[i])
              break;
          else if (a[i] < b[i])
              return -1; //如果 a 小于 b ,则返回-1
   for (int i = 0; i < La; i++) //高精度减法
   {
       a[i] -= b[i];
       if (a[i] < 0) a[i] += 10, a[i + 1]--;</pre>
   for (int i = La - 1; i >= 0; i--)
       if (a[i]) return i + 1; //返回差的位数
   return 0;
                             //返回差的位数
string div(string n1, string n2,
         int nn) // n1,n2 是字符串表示的被除数,除数,nn 是选择返回商还是余
数
{
   string s, v; // s 存商, v 存余数
```

```
int a[L], b[L], r[L],
       La = n1.size(), Lb = n2.size(), i,
       tp = La; // a, b 是整形数组表示被除数,除数,tp 保存被除数的长度
   fill(a, a + L, 0);
   fill(b, b + L, \emptyset);
   fill(r, r + L, 0); //数组元素都置为0
   for (i = La - 1; i >= 0; i--) a[La - 1 - i] = n1[i] - '0';
   for (i = Lb - 1; i >= 0; i--) b[Lb - 1 - i] = n2[i] - '0';
   if (La < Lb || (La == Lb && n1 < n2)) {</pre>
       // cout<<0<<endl;</pre>
       return n1;
                   //如果 a<b,则商为0,余数为被除数
   }
   int t = La - Lb; //除被数和除数的位数之差
   for (int i = La - 1; i >= 0; i--) //将除数扩大10^t 倍
       if (i >= t)
          b[i] = b[i - t];
       else
          b[i] = 0;
   Lb = La;
   for (int j = 0; j <= t; j++) {
       int temp;
       while ((temp = sub(a, b + j, La, Lb - j)) >=
             0) //如果被除数比除数大继续减
       {
          La = temp;
          r[t - j]++;
       }
   for (i = 0; i < L - 10; i++)
       r[i + 1] += r[i] / 10, r[i] %= 10; //统一处理进位
   while (!r[i]) i--; //将整形数组表示的商转化成字符串表示的
   while (i >= 0) s += r[i--] + '0';
   // cout<<s<<endl;</pre>
   i = tp;
   while (!a[i]) i--; //将整形数组表示的余数转化成字符串表示的</span>
   while (i >= 0) v += a[i--] + '0';
   if (v.empty()) v = "0";
   // cout<<v<<endl;</pre>
   if (nn == 1) return s;
   if (nn == 2) return v;
bool cmp(string a, string b) {
   if (a.size() < b.size()) return 1; // a 小于等于 b 返回真
   if (a.size() == b.size() && a <= b) return 1;</pre>
   return 0;
string DeletePreZero(string s) {
   int i;
   for (i = 0; i < s.size(); i++)</pre>
```

```
if (s[i] != '0') break;
   return s.substr(i);
}
string BigInterSqrt(string n) {
   n = DeletePreZero(n);
   string l = "1", r = n, mid, ans;
   while (cmp(1, r)) {
       mid = div(add(1, r), "2", 1);
       if (cmp(mul(mid, mid), n))
           ans = mid, l = add(mid, "1");
       else
           r = sub(mid, "1");
   return ans;
}
// o(n^3)
高精度取模(对单精度)
#include <bits/stdc++.h>
using namespace std;
int mod(string a, int b)//高精度 a 除以单精度 b
{
   int d=0;
   for(int i=0;i<a.size();i++) d=(d*10+(a[i]-'0'))%b;//求出余数
   return d;
}
//o(n)
欧拉筛
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const int N = 1000005;
int phi[N], prime[N], cnt;
bool st[N];
void get_eulers() {
   phi[1] = 1;
   for (int i = 2; i < N; i++) {</pre>
       if (!st[i]) {
           prime[cnt++] = i;
           phi[i] = i - 1;
       for (int j = 0; prime[j] * i < N; j++) {</pre>
           st[prime[j] * i] = 1;
           if (i % prime[j] == 0) {
```

组合数(逆元线性递推

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 \mod = 1e9 + 7;
const 11 \text{ maxn} = 3e4 + 5;
11 inv[maxn], fac[maxn];
11 qpow(11 a, 11 b) {
   11 \text{ ans} = 1;
   while (b) {
       if (b & 1) ans = (ans * a) % mod;
       a = (a * a) \% mod;
       b >>= 1;
   return ans;
}
11 c(11 n, 11 m) {
    if (n < 0 || m < 0 || n < m) return 0;
   return fac[n] * inv[n - m] % mod * inv[m] % mod;
}
void init() {
   fac[0] = 1;
   for (int i = 1; i < maxn; i++) {</pre>
       fac[i] = fac[i - 1] * i % mod;
   inv[maxn - 1] = qpow(fac[maxn - 1], mod - 2);
   for (11 i = maxn - 2; i >= 0; i--) {
       inv[i] = (inv[i + 1] * (i + 1)) % mod;
    }
}
```

中国剩余定理

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const int maxn = 20;
11 A[maxn], B[maxn];
11 exgcd(11 a, 11 b, 11 & x, 11 & y) {
      if(b == 0) {
            x = 1, y = 0;
            return a;
      }
      11 d = exgcd(b, a % b, y, x);
      y -= (a / b) * x;
      return d;
}
int main() {
      int n;
      cin >> n;
      11 M = 111;
      for(int i = 0; i < n; ++ i) {
            cin >> A[i] >> B[i];
            M = M * A[i];
      }
      11 \text{ ans} = 0;
      11 x, y;
      for(int i = 0; i < n; ++ i) {
            11 \text{ Mi} = M / A[i];
            exgcd(Mi, A[i], x, y);
            ans += B[i] * Mi * x;
      }
      cout << (ans % M + M) % M;</pre>
}
```

图论

```
有源汇上下界最大小流
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
struct Edge {
   11 from, to, cap, flow, mn;
   Edge(ll a, ll b, ll c, ll d, ll e) : from(a), to(b), cap(c), flow(d),
mn(e) {}
};
11 n, m;
struct Dinic {
   static const ll maxn = 50010; // 点的大小,记得改
   static const 11 inf = 0x3f3f3f3f3f3f3f3f3f;
   11 N, M, S, T;
   vector<Edge> edges;
   vector<ll> G[maxn];
   bool vis[maxn];
   11 d[maxn];
   11 cur[maxn];
   void AddEdge(ll from, ll to, ll cap, ll c) {
       edges.push back(Edge(from, to, cap, 0, c));
       edges.push_back(Edge(to, from, 0, 0, c));
       M = edges.size();
       G[from].push_back(M - 2);
       G[to].push_back(M - 1);
   }
   bool BFS() {
       memset(vis, 0, sizeof(vis));
       queue<11> Q;
       Q.push(S);
       d[S] = 0;
       vis[S] = 1;
       while (!Q.empty()) {
           11 x = Q.front();
           Q.pop();
           for (ll i = 0; i < G[x].size(); i++) {</pre>
               Edge& e = edges[G[x][i]];
               if (!vis[e.to] && e.cap > e.flow) {
                  vis[e.to] = 1;
                  d[e.to] = d[x] + 1;
                  Q.push(e.to);
               }
```

```
}
       }
       return vis[T];
   }
   11 DFS(11 x, 11 a) {
       if (x == T || a == 0) return a;
       11 flow = 0, f;
       for (11& i = cur[x]; i < G[x].size(); i++) {</pre>
           Edge& e = edges[G[x][i]];
           if (d[x] + 1 == d[e.to] &&
               (f = DFS(e.to, min(a, e.cap - e.flow))) > ∅) {
               e.flow += f;
               edges[G[x][i] ^ 1].flow -= f;
               flow += f;
               a -= f;
               if (a == 0) break;
           }
       return flow;
   }
   void deleteEdge(ll u, ll v) {
       11 siz = edges.size();
       for(ll i = 0; i < siz; ++ i) {</pre>
           if(edges[i].from == u && edges[i].to == v) {
               edges[i].cap = edges[i].flow = 0;
               edges[i ^1].cap = edges[i ^1].flow = 0;
               break;
           }
       }
   }
   11 getValue() {
       return edges[2 * m].flow;
   }
   11 Maxflow(11 S, 11 T) {
       this->S = S, this->T = T;
       11 flow = 0;
       while (BFS()) {
           memset(cur, 0, sizeof(cur));
           flow += DFS(S, inf);
       return flow;
} MF;
```

```
int main() {
   11 s, t;
   cin >> n >> m >> s >> t;
 // n 个点, m 条边, 给的源点汇点
   ll mp[50010] = {0}; // 点的大小,记得改
   for(ll i = 1; i <= m; ++ i) {</pre>
       11 a, b, c, d; // 从 a 到 b 有一条下界 c 上界 d 的边
       cin >> a >> b >> c >> d;
       mp[b] += c;
       mp[a] -= c;
       MF.AddEdge(a, b, d - c, c);
   MF.AddEdge(t, s, 1e18, 0); //
   11 tot = 0;
   for(ll i = 1; i <= n; ++ i) {</pre>
       if(mp[i] > 0) {
           tot += mp[i];
           MF.AddEdge(0, i , mp[i], 0);
       }
       else {
           MF.AddEdge(i, n + 1, -mp[i], 0);
       }
   }
   if( MF.Maxflow(0, n + 1) != tot) {
       cout << "No Solution" << endl;</pre>
   }
   else {
       ll res = MF.getValue(); // 从t到s边的流量
       MF.deleteEdge(t, s);
     //cout << res + MF.Maxflow(s, t) << endl; // 最大流
       cout << res - MF.Maxflow(t, s) << endl; // 最小流
   }
   return 0;
}
```

树链剖分

```
11 fa[N], son[N], dep[N], siz[N], dfn[N], rnk[N], top[N];
11 dfscnt;
vector<11> g[N];
11 tree[N << 1];</pre>
11 lazy[N << 1];</pre>
void dfs1(ll u, ll f, ll d) {
    son[u] = -1;
   siz[u] = 1;
   fa[u] = f;
   dep[u] = d;
   for (auto v:g[u]) {
       if (v == f) continue;
       dfs1(v, u, d + 1);
       siz[u] += siz[v];
       if (son[u] == -1 \mid | siz[v] > siz[son[u]]) son[u] = v;
   }
}
void dfs2(11 u, 11 t) {
    dfn[u] = ++dfscnt;
   rnk[dfscnt] = u;
   top[u] = t;
   if (son[u] == -1) return;
   dfs2(son[u], t);
   for (auto v:g[u]) {
       if (v == son[u] || v == fa[u]) continue;
       dfs2(v, v);
    }
}
11 lca(11 a, 11 b) {
   while (top[a] != top[b]) {
       if (dep[top[a]] < dep[top[b]]) swap(a, b);</pre>
       a = fa[top[a]];
   return dep[a] < dep[b] ? a : b;</pre>
}
void init() {
    for (11 i = 0; i < N; i++) g[i].clear();</pre>
   for (11 i = 0; i < (N << 1); i++) {
       tree[i] = 0;
       lazy[i] = 0;
   dfscnt = 0;
}
```

```
void pushdown(11 k, 11 1, 11 r) {
    if (k >= N \mid | lazy[k] == 0) return;
    11 len = (r - 1 + 1) / 2;
   tree[k << 1] = tree[k << 1] + len * lazy[k];
   tree[k << 1 | 1] = tree[k << 1 | 1] + len * lazy[k];
   lazy[k << 1] = lazy[k << 1] + lazy[k];
   lazy[k << 1 | 1] = lazy[k << 1 | 1] + lazy[k];
   lazy[k] = 0;
}
11 merge_range(11 a, 11 b) {
   11 \text{ ans} = a + b;
    return ans;
void change_range(ll k, ll l, ll r, ll ql, ll qr, ll x) {
    if (r < q1 \mid | qr < 1) return;
    if (q1 <= 1 && r <= qr) {
       tree[k] = tree[k] + x * (r - 1 + 1);
       lazy[k] = lazy[k] + x;
       return;
    pushdown(k, 1, r);
   11 \text{ mid} = (1 + r) >> 1;
    change range(k << 1, 1, mid, ql, qr, x);</pre>
    change_range(k << 1 | 1, mid + 1, r, ql, qr, x);
   tree[k] = merge_range(tree[k << 1], tree[k << 1 | 1]);</pre>
}
11 query range(11 k, 11 1, 11 r, 11 q1, 11 qr) {
   if (r < ql || qr < 1)return 0;
   if (ql <= 1 && r <= qr) {
       return tree[k];
    }
   pushdown(k, 1, r);
   11 \text{ mid} = (1 + r) >> 1;
   11 lq = query_range(k << 1, 1, mid, q1, qr);</pre>
   11 rq = query_range(k << 1 | 1, mid + 1, r, ql, qr);</pre>
   return merge range(lq, rq);
}
11 query_path(11 a, 11 b) {
   11 \text{ sum} = 0;
   while (top[a] != top[b]) {
       if (dep[top[a]] < dep[top[b]]) swap(a, b);</pre>
       sum = sum + query_range(1, 1, N, dfn[top[a]], dfn[a]);
       //dfn[top[a]]~dfn[a]
       a = fa[top[a]];
    }
```

```
if (dep[a] > dep[b]) swap(a, b);
   //点权
   sum = sum + query_range(1, 1, N, dfn[a], dfn[b]);
   //if (a != b) sum = sum + query_range(1, 1, N, dfn[a] + 1, dfn[b]);
   //dfn[a]\sim dfn[b],x
   return sum;
}
void change_path(ll a, ll b, ll x) {
   while (top[a] != top[b]) {
       if (dep[top[a]] < dep[top[b]]) swap(a, b);</pre>
       change_range(1, 1, N, dfn[top[a]], dfn[a], x);
       //dfn[top[a]]~dfn[a]
       a = fa[top[a]];
   if (dep[a] > dep[b]) swap(a, b);
   //点权
   change_range(1, 1, N, dfn[a], dfn[b], x);
   //if (a != b) change_range(1, 1, N, dfn[a] + 1, dfn[b], x);
   //dfn[a]\sim dfn[b],x
}
```

虚树

```
11 fa[N], son[N], dep[N], siz[N], dfn[N], rnk[N], top[N];
11 dfscnt;
vector<11> g[N];
11 mmin[N];
void dfs1(ll u, ll f, ll d) {
    son[u] = -1;
    siz[u] = 1;
   fa[u] = f;
   dep[u] = d;
   for (auto v:g[u]) {
       if (v == f) continue;
       dfs1(v, u, d + 1);
       siz[u] += siz[v];
       if (son[u] == -1 \mid | siz[v] > siz[son[u]]) son[u] = v;
   }
}
void dfs2(11 u, 11 t) {
    dfn[u] = ++dfscnt;
    rnk[dfscnt] = u;
   top[u] = t;
   if (son[u] == -1) return;
   dfs2(son[u], t);
   for (auto v:g[u]) {
       if (v == son[u] || v == fa[u]) continue;
       dfs2(v, v);
   }
}
11 lca(11 a, 11 b) {
   while (top[a] != top[b]) {
       if (dep[top[a]] < dep[top[b]]) swap(a, b);</pre>
       a = fa[top[a]];
   return dep[a] < dep[b] ? a : b;</pre>
}
struct edge {
   11 s, t, v;
edge e[N];
vector<int> vg[N];
int sta[N], tot;
int h[N];
void build(int *H, int num) {
    sort(H + 1, H + 1 + num, [](int a, int b) { return dfn[a] <</pre>
```

```
dfn[b]; });
   sta[tot = 1] = 1, vg[1].clear();// 1 号节点入栈, 清空 1 号节点对应的邻
接表,设置邻接表边数为1
   for (int i = 1, 1; i <= num; ++i) {
      if (H[i] == 1) continue; //如果 1 号节点是关键节点就不要重复添加
      1 = lca(H[i], sta[tot]); //计算当前节点与栈顶节点的 LCA
      if (1 != sta[tot]) { //如果 LCA 和栈顶元素不同,则说明当前节点不再
当前栈所存的链上
         while (dfn[1] < dfn[sta[tot - 1]]) {//当次大节点的 Dfs 序大于
LCA 的 Dfs 序
            vg[sta[tot - 1]].push back(sta[tot]);
            vg[sta[tot]].push_back(sta[tot - 1]);
            tot--:
         } //把与当前节点所在的链不重合的链连接掉并且弹出
         if (dfn[1] > dfn[sta[tot - 1]]) { //如果 LCA 不等于次大节点
(这里的大于其实和不等于没有区别)
            vg[1].clear();
            vg[1].push_back(sta[tot]);
            vg[sta[tot]].push_back(1);
            sta[tot] = 1;//说明 LCA 是第一次入栈,清空其邻接表,连边后弹
出栈顶元素,并将 LCA 入栈
         } else {
            vg[1].push_back(sta[tot]);
            vg[sta[tot]].push back(1);
            tot--; //说明 LCA 就是次大节点,直接弹出栈顶元素
         }
      }
      vg[H[i]].clear();
      sta[++tot] = H[i];
      //当前节点必然是第一次入栈,清空邻接表并入栈
   for (int i = 1; i < tot; ++i) {
      vg[sta[i]].push back(sta[i + 1]);
      vg[sta[i + 1]].push back(sta[i]);
   } //剩余的最后一条链连接一下
   return:
spfa 最短路及负环
```

```
#include<bits/stdc++.h>
using namespace std;
typedef long long 11;
const int N = 1 \ll 20;
struct edge {
   11 to, len;
};
vector<edge> g[N];
11 d[N], cnt[N], vis[N];
bool spfa(11 s, 11 n) {
   queue<int> que;
   for (int i = 1; i <= n; i++) { //防止不连通,全加进去
      que.push(i);
      vis[i] = 1;
   while (!que.empty()) {
      11 p = que.front();
      que.pop();
      vis[p] = 0;
      for (auto x:g[p]) {
          if (d[x.to] > d[p] + x.len) {
             d[x.to] = d[p] + x.len;
             cnt[x.to] = cnt[p] + 1;
             if (!vis[x.to]) {
                if (cnt[x.to] > n) return 0;
                vis[x.to] = 1;
                que.push(x.to);
             }
          }
      }
   return 1;
}
二分图匹配(匈牙利)
//大量使用了memset,但常数貌似很小?HDU6808 跑了998ms (限制5000ms),然而
这个代int main()不是HDU6808的
#include<bits/stdc++.h>
using namespace std;
const int maxn=505;// 最大点数
const int inf=0x3f3f3f3f;// 距离初始值
struct HK_Hungary{//这个板子从1开始,0点不能用,nx 为左边点数,ny 为右边点数
   int nx,ny;//左右顶点数量
   vector<int>bmap[maxn];
   int cx[maxn];//cx[i]表示左集合i 顶点所匹配的右集合的顶点序号
```

```
int cy[maxn]; //cy[i]表示右集合i 顶点所匹配的左集合的顶点序号
   int dx[maxn];
   int dy[maxn];
   int dis;
   bool bmask[maxn];
   void init(int a,int b){
      nx=a,ny=b;
      for(int i=0;i<=nx;i++){</pre>
          bmap[i].clear();
      }
   }
   void add_edge(int u,int v){
      bmap[u].push back(v);
   bool searchpath(){//寻找 增广路径
      queue<int>Q;
      dis=inf:
      memset(dx,-1,sizeof(dx));
      memset(dy,-1,sizeof(dy));
      for(int i=1;i<=nx;i++){//cx[i]表示左集合i 顶点所匹配的右集合的顶点
序号
          if(cx[i]==-1){//将未遍历的节点 入队 并初始化次节点距离为0
             Q.push(i);
             dx[i]=0;
      }//广度搜索增广路径
      while(!Q.empty()){
          int u=Q.front();
          Q.pop();
          if(dx[u]>dis) break;//取右侧节点
          for(int i=0;i<bmap[u].size();i++){</pre>
             int v=bmap[u][i];//右侧节点的增广路径的距离
             if(dy[v]==-1){
                 dy[v]=dx[u]+1;//v 对应的距离 为 u 对应距离加 1
                 if(cy[v]==-1)dis=dy[v];
                 else{
                    dx[cy[v]]=dy[v]+1;
                    Q.push(cy[v]);
                 }
             }
          }
      }
      return dis!=inf;
   int findpath(int u){//寻找路径 深度搜索
      for(int i=0;i<bmap[u].size();i++){</pre>
          int v=bmap[u][i];//如果该点没有被遍历过 并且距离为上一节点+1
          if(!bmask[v]&&dy[v]==dx[u]+1){//对该点染色
             bmask[v]=1;
```

```
if(cy[v]!=-1&&dy[v]==dis)continue;
               if(cy[v]==-1||findpath(cy[v])){
                  cy[v]=u;cx[u]=v;
                  return 1;
               }
           }
       }
       return 0;
   }
   int MaxMatch(){//得到最大匹配的数目
       int res=0;
       memset(cx,-1,sizeof(cx));
       memset(cy,-1,sizeof(cy));
       while(searchpath()){
           memset(bmask,0,sizeof(bmask));
           for(int i=1;i<=nx;i++){</pre>
               if(cx[i]==-1){
                  res+=findpath(i);
               }
           }
       return res;
}HK;
int main(){
   int nn,n,m;
   cin>>nn;
   while(nn--){
       scanf("%d%d",&n,&m);
       HK.init(n,m);//左端点和右端点数量
       for(int i=1;i<=n;i++){</pre>
           int snum;
           cin>>snum;
           int v;
           for(int j=1;j<=snum;j++){</pre>
               cin>>v;
               HK.add_edge(i,v);//连边
           }
       cout<<HK.MaxMatch()<<endl;//求最大匹配
   return 0;
}
```

```
强连通(kosaraju
#include <bits/stdc++.h>
using namespace std;
struct SCC {
   static const int MAXV = 100000;
   int V;
   vector<int> g[MAXV], rg[MAXV], vs;
   bool used[MAXV];
   int cmp[MAXV];
   void add_edge(int from, int to) {
       g[from].push_back(to);
       rg[to].push back(from);
   }
   void dfs(int v) {
       used[v] = 1;
       for (int i = 0; i < g[v].size(); i++) {</pre>
           if (!used[g[v][i]]) dfs(g[v][i]);
       vs.push_back(v);
   }
   void rdfs(int v, int k) {
       used[v] = 1;
       cmp[v] = k;
       for (int i = 0; i < rg[v].size(); i++) {</pre>
           if (!used[rg[v][i]]) rdfs(rg[v][i], k);
       }
   }
   int solve() {
       memset(used, 0, sizeof(used));
       vs.clear();
       for (int v = 1; v <= V; v++) {</pre>
           if (!used[v]) dfs(v);
       memset(used, 0, sizeof(used));
       int k = 0;
       for (int i = (int)vs.size() - 1; i >= 0; i--) {
           if (!used[vs[i]]) rdfs(vs[i], ++k);
       }
       return k;
   }
   void init(int n) {
       V = n;
       vs.clear();
       for (int i = 0; i < MAXV; i++) {</pre>
```

```
g[i].clear();
           rg[i].clear();
           used[i] = 0;
           cmp[i] = 0;
       }
    }
} scc;
//记得调用 init()
强连通(tarjan
#include <bits/stdc++.h>
using namespace std;
struct SCC {
    static const int MAXN = 100000;
   vector<int> g[MAXN];
    int dfn[MAXN], lowlink[MAXN], sccno[MAXN], dfs_clock, scc_cnt;
    stack<int> S;
   void dfs(int u) {
       dfn[u] = lowlink[u] = ++dfs clock;
       S.push(u);
       for (int i = 0; i < g[u].size(); i++) {</pre>
           int v = g[u][i];
           if (!dfn[v]) {
               dfs(v);
               lowlink[u] = min(lowlink[u], lowlink[v]);
           } else if (!sccno[v]) {
               lowlink[u] = min(lowlink[u], dfn[v]);
           }
       if (lowlink[u] == dfn[u]) {
           ++scc_cnt;
           for (;;) {
               int x = S.top();
               S.pop();
               sccno[x] = scc_cnt;
               if (x == u) break;
           }
       }
    }
```

void solve(int n) {

dfs_clock = scc_cnt = 0;

memset(sccno, 0, sizeof(sccno));

```
memset(dfn, 0, sizeof(dfn));
    memset(lowlink, 0, sizeof(lowlink));
    for (int i = 1; i <= n; i++) {
        if (!dfn[i]) dfs(i);
    }
    }
} scc;

// scc_cnt 为SCC 计数器,sccno[i]为i 所在SCC 的编号
// vector<int> g[MAXN]中加边
//之后再补充init()
```

```
强连通(tarjan 无 vector
#include <bits/stdc++.h>
using namespace std;
struct SCC {
   static const int MAXN = 5000;
   static const int MAXM = 2000000;
   int dfs clock, edge cnt = 1, scc cnt;
   int head[MAXN];
   int dfn[MAXN], lowlink[MAXN];
   int sccno[MAXN];
   stack<int> s;
   struct edge {
       int v, next;
   } e[MAXM];
   void add_edge(int u, int v) {
       e[edge_cnt].v = v;
       e[edge_cnt].next = head[u];
       head[u] = edge_cnt++;
   }
   void tarjan(int u) {
       int v;
       dfn[u] = lowlink[u] = ++dfs_clock; //每次dfs, u 的次序号增加1
                                        //将u入栈
       s.push(u);
       for (int i = head[u]; i != -1; i = e[i].next) //访问从 u 出发的边
       {
          v = e[i].v;
          if (!dfn[v]) //如果ν没被处理过
          {
              tarjan(v); // dfs(v)
              lowlink[u] = min(lowlink[u], lowlink[v]);
           } else if (!sccno[v])
              lowlink[u] = min(lowlink[u], dfn[v]);
```

```
if (dfn[u] == lowlink[u]) {
           scc_cnt++;
           do {
              v = s.top();
              s.pop();
              sccno[v] = scc_cnt;
           } while (u != v);
       }
   }
   int find scc(int n) {
       for (int i = 1; i <= n; i++)</pre>
           if (!dfn[i]) tarjan(i);
       return scc cnt;
   }
   void init() {
       scc_cnt = dfs_clock = 0;
       edge_cnt = 1; //不用初始化 e 数组,省时间
       while (!s.empty()) s.pop();
       memset(head, -1, sizeof(head));
       memset(sccno, 0, sizeof(sccno));
       memset(dfn, 0, sizeof(dfn));
       memset(lowlink, 0, sizeof(lowlink));
} scc;
```

最大流

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
struct Edge {
   11 from, to, cap, flow;
   Edge(ll a, ll b, ll c, ll d) : from(a), to(b), cap(c), flow(d) {}
};
struct Dinic {
   static const ll maxn = 10000;
   static const 11 inf = 0x3f3f3f3f3f3f3f3f3f;
   11 N, M, S, T;
   vector<Edge> edges;
   vector<11> G[maxn];
   bool vis[maxn];
   11 d[maxn];
   11 cur[maxn];
```

```
void AddEdge(ll from, ll to, ll cap) {
    edges.push_back(Edge(from, to, cap, ∅));
    edges.push_back(Edge(to, from, 0, 0));
   M = edges.size();
   G[from].push back(M - 2);
   G[to].push_back(M - 1);
}
bool BFS() {
   memset(vis, 0, sizeof(vis));
   queue<11> Q;
   Q.push(S);
   d[S] = 0;
   vis[S] = 1;
   while (!Q.empty()) {
       11 \times = Q.front();
       Q.pop();
       for (ll i = 0; i < G[x].size(); i++) {</pre>
           Edge& e = edges[G[x][i]];
           if (!vis[e.to] && e.cap > e.flow) {
               vis[e.to] = 1;
               d[e.to] = d[x] + 1;
               Q.push(e.to);
           }
       }
   return vis[T];
}
11 DFS(11 x, 11 a) {
   if (x == T || a == 0) return a;
    11 flow = 0, f;
   for (11& i = cur[x]; i < G[x].size(); i++) {</pre>
       Edge& e = edges[G[x][i]];
       if (d[x] + 1 == d[e.to] &&
           (f = DFS(e.to, min(a, e.cap - e.flow))) > 0) {
           e.flow += f;
           edges[G[x][i] ^ 1].flow -= f;
           flow += f;
           a -= f;
           if (a == 0) break;
       }
    }
   return flow;
}
11 Maxflow(11 S, 11 T) {
   this->S = S, this->T = T;
```

```
11 flow = 0;
     while (BFS()) {
        memset(cur, 0, sizeof(cur));
        flow += DFS(S, inf);
     return flow;
  }
} MF;
//有源汇上下界最大流, 跑完可行流后, s-t 的最大流即为答案
//有源汇上下届最小流,不连无穷边, s-t 跑最大流,再加上 t-s 无穷边,再跑最大流,
无穷边流量为答案
//最大权闭合子图
//构造一个新的流网络,建一个源点 s 和汇点 t,从 s 向原图中所有点权为正数的点建一
条容量为点权的边,
//从点权为负数的点向t建一条容量为点权绝对值的边,原图中各点建的边都建成容量为
正无穷的边。
//然后求从 s 到 t 的最小割,再用所有点权为正的权值之和减去最小割,就是我们要求的
最大权值和了。
//最大密度子图
//01 分数规划
//addedge(S, V, m), addedge(E, 1), addedge(V, T, 2*q-deg(v)+m)
//h(q)=n*m-maxflow(S,T)
```

最大流 (double)

```
#include <iostream>
#include <cstring>
#include <algorithm>

using namespace std;

struct Dinic {
    static constexpr int N = 10010, M = 100010, INF = 1e8;
    static constexpr double eps = 1e-8;

// int n, m, S, T;
    int S, T;
    int S, T;
    int h[N], e[M], ne[M], idx;
    double f[M];
    int q[N], d[N], cur[N]; // d 表示从源点开始走到该点的路径上所有边的容量的最小值
```

```
void AddEdge(int a, int b, double c)
{
   e[idx] = b, f[idx] = c, ne[idx] = h[a], h[a] = idx ++ ;
   e[idx] = a, f[idx] = 0, ne[idx] = h[b], h[b] = idx ++ ;
}
bool bfs()
{
   int hh = 0, tt = 0;
   memset(d, -1, sizeof d);
   q[0] = S, d[S] = 0, cur[S] = h[S];
   while (hh <= tt)</pre>
   {
       int t = q[hh ++ ];
       for (int i = h[t]; ~i; i = ne[i])
           int ver = e[i];
           if (d[ver] == -1 && f[i] > 0)
           {
               d[ver] = d[t] + 1;
               cur[ver] = h[ver];
               if (ver == T) return true;
               q[ ++ tt] = ver;
           }
       }
   return false;
}
double find(int u, double limit)
{
   if (u == T) return limit;
   double flow = 0;
   for (int i = cur[u]; ~i && flow < limit; i = ne[i])</pre>
       cur[u] = i;
       int ver = e[i];
       if (d[ver] == d[u] + 1 && f[i] > 0)
           double t = find(ver, min(f[i], limit - flow));
           if (t < eps) d[ver] = -1;
           f[i] -= t, f[i ^ 1] += t, flow += t;
       }
    }
   return flow;
}
double Maxflow(int S, int T)
      this->S = S, this->T = T;
```

```
double r = 0, flow;
    while (bfs()) while (flow = find(S, INF)) r += flow;
    return r;
}
void init() //////
{
    memset(h, -1, sizeof h);
    idx = 0;
}
MF;
// ?èinit
```

```
最小费用最大流
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
struct Edge {
   11 from, to, cap, flow, cost;
   Edge(ll u, ll v, ll c, ll f, ll w):from(u), to(v), cap(c), flow(f),
cost(w) {}
};
struct MCMF {
   static const 11 maxn = 6000;
   static const 11 INF = 0x3f3f3f3f3f3f3f3f;
   11 n, m;
   vector<Edge> edges;
   vector<ll> G[maxn];
   11 inq[maxn];
   11 d[maxn];
   11 p[maxn];
   11 a[maxn];
   void init(ll n) {
       this->n = n;
       for (ll i = 1; i <= n; i++) G[i].clear();</pre>
       edges.clear();
   }
   void add_edge(ll from, ll to, ll cap, ll cost) {
       from++,to++;//原板子无法使用 0 点,故修改
       edges.push_back(Edge(from, to, cap, 0, cost));
       edges.push_back(Edge(to, from, 0, 0, -cost));
```

```
m = edges.size();
       G[from].push back(m - 2);
       G[to].push_back(m - 1);
   }
   bool BellmanFord(ll s, ll t, ll& flow, ll& cost) {
       for (ll i = 1; i <= n; ++i) d[i] = INF;
       memset(inq, 0, sizeof(inq));
       d[s] = 0, inq[s] = 1, p[s] = 0, a[s] = INF;
       queue<11> Q;
       Q.push(s);
       while (!Q.empty()) {
           11 u = Q.front();
           Q.pop();
           inq[u] = 0;
           for (ll i = 0; i < G[u].size(); ++i) {</pre>
              Edge& e = edges[G[u][i]];
              if (e.cap > e.flow && d[e.to] > d[u] + e.cost) {
                  d[e.to] = d[u] + e.cost;
                  p[e.to] = G[u][i];
                  a[e.to] = min(a[u], e.cap - e.flow);
                  if (!inq[e.to]) {
                     Q.push(e.to);
                      inq[e.to] = 1;
                  }
              }
           }
       if (d[t] == INF) return false;
       flow += a[t];
       cost += (11)d[t] * (11)a[t];
       for (11 u = t; u != s; u = edges[p[u]].from) {
           edges[p[u]].flow += a[t];
           edges[p[u] ^ 1].flow -= a[t];
       return true;
   }
   //需要保证初始网络中没有负权圈
   11 MincostMaxflow(ll s, ll t, ll& cost) {
       S++, t++; // 原板子无法使用 Ø 点,故修改
       11 flow = 0;
       cost = 0;
       while (BellmanFord(s, t, flow, cost));
       return flow;
   }
\} mcmf; // 若固定流量 k,增广时在 fLow+a>=k 的时候只增广 k-fLow 单位的流量,
然后终止程序
```

```
树分治
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 10005;
const int INF = 1000000000;
struct edge {
   int to, length;
   edge() {}
   edge(int a, int b) : to(a), length(b) {}
};
vector<edge> g[MAXN];
bool centroid[MAXN];
int subtree_size[MAXN];
int ans;
//计算子树大小
int compute subtree size(int v, int p) {
   int c = 1;
   for (int i = 0; i < g[v].size(); i++) {</pre>
       int w = g[v][i].to;
       if (w == p || centroid[w]) continue;
       c += compute_subtree_size(w, v);
   subtree_size[v] = c;
   return c;
}
//查找重心, t 为连通分量大小
// pair (最大子树顶点数,顶点编号)
pair<int, int> search_centroid(int v, int p, int t) {
   pair<int, int> res = pair<int, int>(INF, -1);
   int s = 1, m = 0;
   for (int i = 0; i < g[v].size(); i++) {</pre>
       int w = g[v][i].to;
       if (w == p || centroid[w]) continue;
       res = min(res, search_centroid(w, v, t));
       m = max(m, subtree size[w]);
       s += subtree size[w];
   }
   m = max(m, t - s);
```

```
res = min(res, pair<int, int>(m, v));
   return res;
}
void init(int n) {
   memset(centroid, 0, sizeof(centroid));
   memset(subtree_size, 0, sizeof(subtree_size));
   for (int i = 0; i <= n; i++) g[i].clear();</pre>
   ans = 0;
}
int solve(int u) {
   compute_subtree_size(u, -1);
   int s = search_centroid(u, -1, subtree_size[u]).second;
   centroid[s] = 1;
   for (int i = 0; i < g[s].size(); i++) {</pre>
       int v = g[s][i].to;
       if (centroid[v]) continue;
       /*solve()*/
   }
   /*do something*/
   centroid[s] = 0;
   return ans;
}
```

```
拓扑排序
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 100000;
int c[MAXN];
int topo[MAXN], t, V;
vector<int> g[MAXN];
bool dfs(int u) {
   c[u] = -1;
   for (int i = 0; i < g[u].size(); i++) {</pre>
       int v = g[u][i];
       if (c[v] < 0)
           return false;
       else if (!c[v] && !dfs(v))
           return false;
   c[u] = 1;
   topo[t--] = u;
   return true;
}
```

```
bool toposort(int n) {
    V = n;
    t = n;
    memset(c, 0, sizeof(c));
    for (int u = 1; u <= V; u++)
        if (!c[u] && !dfs(u)) return false;
    return true;
}</pre>
```

```
最近公共祖先(倍增)
#include <algorithm>
#include <cstdio>
#include <cstring>
#include <iostream>
using namespace std;
const int MAX = 600000;
struct edge {
   int t, nex;
} e[MAX << 1];</pre>
int head[MAX], tot;
int depth[MAX], fa[MAX][22], lg[MAX];
void add_edge(int x, int y) {
   e[++tot].t = y;
   e[tot].nex = head[x];
   head[x] = tot;
   e[++tot].t = x;
   e[tot].nex = head[y];
   head[y] = tot;
}
void dfs(int now, int fath) {
   fa[now][0] = fath;
   depth[now] = depth[fath] + 1;
   for (int i = 1; i <= lg[depth[now]]; ++i)</pre>
       fa[now][i] = fa[fa[now][i - 1]][i - 1];
   for (int i = head[now]; i; i = e[i].nex)
       if (e[i].t != fath) dfs(e[i].t, now);
}
int lca(int x, int y) {
   if (depth[x] < depth[y]) swap(x, y);</pre>
   while (depth[x] > depth[y]) x = fa[x][lg[depth[x] - depth[y]] - 1];
```

```
if (x == y) return x;
   for (int k = \lg[depth[x]] - 1; k \ge 0; --k)
       if (fa[x][k] != fa[y][k]) x = fa[x][k], y = fa[y][k];
   return fa[x][0];
}
void init(int n, int root) {
   for (int i = 1; i \le n; ++i) lg[i] = lg[i - 1] + (1 << lg[i - 1] ==
i);
   dfs(root, ∅);
}
最近公共祖先(线段树)
#include <bits/stdc++.h>
using namespace std;
int n, m, root;
const int MAX_N = 500005;
const int MAX = 1 << 20;
vector<int> g[MAX_N];
vector<int> vs;
pair<int, int> tree[MAX * 2 + 10];
int fir[MAX N];
int fa[MAX_N];
int dep[MAX N];
void dfs(int k, int p, int d) {
   fa[k] = p;
   dep[k] = d;
   vs.push back(k);
   for (int i = 0; i < g[k].size(); i++) {</pre>
       if (g[k][i] != p) {
           dfs(g[k][i], k, d + 1);
           vs.push back(k);
       }
   }
}
void build(int k) {
   if (k >= MAX) return;
   build(k << 1);
   build(k \langle\langle 1 | 1 \rangle\rangle;
   tree[k] = min(tree[k << 1], tree[k << 1 | 1]);
pair<int, int> query(int k, int s, int e, int l, int r) {
   if (e < 1 | | r < s) return pair<int, int>(INT_MAX, 0);
   if (1 <= s && e <= r) return tree[k];
   return min(query(k << 1, s, (s + e) >> 1, l, r),
              query(k << 1 | 1, ((s + e) >> 1) + 1, e, l, r));
void init() {
   dfs(root, root, ∅);
   for (int i = 0; i < MAX * 2 + 10; i++) tree[i] = pair<int, int>(INT M
```

```
AX, 0);
    for (int i = MAX; i < MAX + vs.size(); i++)</pre>
       tree[i] = pair<int, int>(dep[vs[i - MAX]], vs[i - MAX]);
   for (int i = 0; i < vs.size(); i++) {</pre>
       if (fir[vs[i]] == 0) fir[vs[i]] = i + 1;
   build(1);
int lca(int a, int b) {
    return query(1, 1, MAX, min(fir[a], fir[b]), max(fir[a], fir[b])).se
cond;
int main() {
    scanf("%d%d%d", &n, &m, &root);
    for (int i = 1; i < n; i++) {</pre>
       int a, b;
       scanf("%d%d", &a, &b);
       g[a].push back(b);
       g[b].push_back(a);
    }
   init();
   for (int i = 1; i <= m; i++) {</pre>
       int a, b;
       scanf("%d%d", &a, &b);
       printf("%d\n", lca(a, b));
    }
}
KM
#include<bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 maxN = 310;
const ll INF = 1e16;
struct KM {
    11 mp[maxN][maxN], link_x[maxN], link_y[maxN], N;
   bool visx[maxN], visy[maxN];
    11 que[maxN << 1], top, fail, pre[maxN];</pre>
```

```
11 hx[maxN], hy[maxN], slk[maxN];
inline ll check(ll i) {
    visx[i] = true;
    if (link_x[i]) {
       que[fail++] = link_x[i];
        return visy[link_x[i]] = true;
    }
    while (i) {
       link_x[i] = pre[i];
        swap(i, link_y[pre[i]]);
    }
    return 0;
}
void bfs(ll S) {
    for (ll i = 1; i <= N; i++) {</pre>
       slk[i] = INF;
       visx[i] = visy[i] = false;
    }
    top = 0;
    fail = 1;
   que[0] = S;
   visy[S] = true;
    while (true) {
       11 d;
       while (top < fail) {</pre>
           for (ll i = 1, j = que[top++]; i <= N; i++) {</pre>
```

```
if (!visx[i] && slk[i] >= (d = hx[i] + hy[j] - mp[i]
[j])) {
                        pre[i] = j;
                        if (d) slk[i] = d;
                        else if (!check(i)) return;
                    }
                }
            }
            d = INF;
            for (ll i = 1; i <= N; i++) {</pre>
                if (!visx[i] && d > slk[i]) d = slk[i];
            }
            for (ll i = 1; i <= N; i++) {</pre>
                if (visx[i]) hx[i] += d;
                else slk[i] -= d;
                if (visy[i]) hy[i] -= d;
            }
            for (ll i = 1; i <= N; i++) {</pre>
                if (!visx[i] && !slk[i] && !check(i)) return;
            }
        }
    }
   void init() {
        for (ll i = 1; i <= N; i++) {</pre>
            link_x[i] = link_y[i] = 0;
            visy[i] = false;
       }
        for (ll i = 1; i <= N; i++) {</pre>
```

```
hx[i] = 0;
            for (ll j = 1; j <= N; j++) {</pre>
                if (hx[i] < mp[i][j]) hx[i] = mp[i][j];</pre>
            }
        }
    }
} km;
int main() {
    ios::sync_with_stdio(∅);
    11 n;
    cin >> n;
    11 \text{ ans} = 0;
    for (int i = 1; i <= n; i++) {</pre>
        11 a, b, c, d;
        cin >> a >> b >> c >> d;
        ans += a * a + b * b;
        for (int j = 1; j <= n; j++) {</pre>
            km.mp[i][j] = -(c + d * (j - 1)) * (c + d * (j - 1));
              cout << -km.mp[i][j] << ' ';</pre>
//
//
              cin >> km.mp[i][j];
//
              km.mp[i][j] = -km.mp[i][j];
        }
//
        cout << endl;</pre>
    }
    km.N = n;
    km.init();
```

```
for (int i = 1; i <= km.N; i++) km.bfs(i);</pre>
    for (int i = 1; i <= n; i++) ans -= km.mp[i][km.link_x[i]];</pre>
    cout << ans << endl;</pre>
}
prufer 序列
#include <iostream>
#include <cstdio>
#include <cstring>
#include <algorithm>
using namespace std;
const int N = 100010;
int n, m;
int f[N], d[N], p[N];
void tree2prufer()
{
    for (int i = 1; i < n; i ++ )</pre>
    {
       scanf("%d", &f[i]);
       d[f[i]] ++;
    }
    for (int i = 0, j = 1; i < n - 2; j ++ )
    {
        while (d[j]) j ++;
       p[i ++ ] = f[j];
```

```
while (i < n - 2 \&\& -- d[p[i - 1]] == 0 \&\& p[i - 1] < j) p[i ++ ]
= f[p[i - 1]];
   }
   for (int i = 0; i < n - 2; i ++ ) printf("%d ", p[i]);</pre>
}
void prufer2tree()
{
   for (int i = 1; i <= n - 2; i ++ )</pre>
    {
        scanf("%d", &p[i]);
       d[p[i]] ++;
    }
   p[n - 1] = n;
   for (int i = 1, j = 1; i < n; i ++, j ++)
   {
       while (d[j]) j ++;
       f[j] = p[i];
       while (i < n - 1 \&\& -- d[p[i]] == 0 \&\& p[i] < j) f[p[i]] = p[i +
1], i ++;
    }
   for (int i = 1; i <= n - 1; i ++ ) printf("%d ", f[i]);</pre>
}
int main()
{
```

```
scanf("%d%d", &n, &m);
   if (m == 1) tree2prufer();
   else prufer2tree();
   return 0;
}
朱刘算法
#include <iostream>
#include <cstring>
#include <cstdio>
#include <algorithm>
#include <cmath>
#define x first
#define y second
using namespace std;
typedef pair<double, double> PDD;
const int N = 110;
const double INF = 1e8;
int n, m;
PDD q[N];
bool g[N][N];
double d[N][N], bd[N][N];
int pre[N], bpre[N];
int dfn[N], low[N], ts, stk[N], top;
```

```
int id[N], cnt;
bool st[N], ins[N];
void dfs(int u) {
   st[u] = true;
   for (int i = 1; i <= n; i++)</pre>
       if (g[u][i] && !st[i])
           dfs(i);
}
bool check_con() {
   memset(st, 0, sizeof st);
   dfs(1);
   for (int i = 1; i <= n; i++)</pre>
       if (!st[i])
           return false;
   return true;
}
double get_dist(int a, int b) {
   double dx = q[a].x - q[b].x;
   double dy = q[a].y - q[b].y;
   return sqrt(dx * dx + dy * dy);
}
void tarjan(int u) {
   dfn[u] = low[u] = ++ts;
   stk[++top] = u, ins[u] = true;
```

```
int j = pre[u];
    if (!dfn[j]) {
       tarjan(j);
       low[u] = min(low[u], low[j]);
    } else if (ins[j]) low[u] = min(low[u], dfn[j]);
    if (low[u] == dfn[u]) {
       int y;
       ++cnt;
       do {
           y = stk[top--], ins[y] = false, id[y] = cnt;
        } while (y != u);
    }
}
double work() {
    double res = 0;
    for (int i = 1; i <= n; i++)</pre>
       for (int j = 1; j <= n; j++)</pre>
            if (g[i][j]) d[i][j] = get_dist(i, j);
           else d[i][j] = INF;
   while (true) {
       for (int i = 1; i <= n; i++) {</pre>
            pre[i] = i;
           for (int j = 1; j <= n; j++)</pre>
               if (d[pre[i]][i] > d[j][i])
```

```
}
        memset(dfn, ∅, sizeof dfn);
        ts = cnt = 0;
        for (int i = 1; i <= n; i++)</pre>
            if (!dfn[i])
                tarjan(i);
        if (cnt == n) {
            for (int i = 2; i <= n; i++) res += d[pre[i]][i];</pre>
            break;
       }
        for (int i = 2; i <= n; i++)</pre>
            if (id[pre[i]] == id[i])
                res += d[pre[i]][i];
        for (int i = 1; i <= cnt; i++)</pre>
            for (int j = 1; j <= cnt; j++)</pre>
                bd[i][j] = INF;
        for (int i = 1; i <= n; i++)</pre>
            for (int j = 1; j <= n; j++)</pre>
                if (d[i][j] < INF && id[i] != id[j]) {</pre>
                    int a = id[i], b = id[j];
                    if (id[pre[j]] == id[j]) bd[a][b] = min(bd[a][b], d[i]
[j] - d[pre[j]][j]);
                    else bd[a][b] = min(bd[a][b], d[i][j]);
```

pre[i] = j;

```
}
       n = cnt;
       memcpy(d, bd, sizeof d);
   }
   return res;
}
int main() {
   while (~scanf("%d%d", &n, &m)) {
       for (int i = 1; i <= n; i++) scanf("%lf%lf", &q[i].x, &q[i].y);</pre>
       memset(g, 0, sizeof g);
       while (m--) {
           int a, b;
           scanf("%d%d", &a, &b);
           if (a != b && b != 1) g[a][b] = true;
       }
       if (!check_con()) puts("poor snoopy");
       else printf("%.21f\n", work());
   }
   return 0;
}
欧拉回路
#include <bits/stdc++.h>
```

```
using namespace std;
typedef long long ll;
const int N = 1e6 + 10;
int stk[N], top;
struct edge {
   int to, idx;
};
vector<edge> g[N];
namespace Euler1 { //有向图欧拉回路
   bool vis[N];
   int cur[N];
   void dfs(int u, const int &w) {
       vis[abs(w)] = true;
       for (int &i = cur[u]; i < g[u].size();) {</pre>
           int idx = g[u][i].idx, v = g[u][i].to;
           i++;
           if (!vis[abs(idx)]) dfs(v, idx);
       }
       stk[++top] = w;
   }
   bool solve(int n) {
       // init();
```

```
for (int i = 0; i <= n; i++) cur[i] = 0;</pre>
       for (int i = 0; i <= n; i++) vis[i] = 0;</pre>
       // calculate degree
       for (int i = 1; i <= n; i++) {</pre>
           if (g[i].size() & 1) return false;
       }
       // Hierholzer
       for (int i = 1; i <= n; i++)</pre>
           if (!g[i].empty()) {
               dfs(i, ∅);
               break;
           }
       return true;
    }
} // namespace Euler1
namespace Euler2 { // 无向图欧拉回路
   int deg[N], cur[N];
   void dfs(int u, const int &w) {
       for (int &i = cur[u]; i < g[u].size();) {</pre>
           int idx = g[u][i].idx, v = g[u][i].to;
           i++;
           dfs(v, idx);
       }
       stk[++top] = w;
    }
```

```
bool solve(int n) {
       // init
        for (int i = 0; i <= n; i++) deg[i] = 0;</pre>
        for (int i = 0; i <= n; i++) cur[i] = 0;</pre>
        // calculate degree
        for (int i = 1; i <= n; ++i) {</pre>
            for (auto x: g[i]) deg[i]++, deg[x.to]--;
        }
        for (int i = 1; i <= n; ++i)</pre>
            if (deg[i]) return false;
        // Hierholzer
        for (int i = 1; i <= n; ++i)</pre>
            if (!g[i].empty()) {
                dfs(i, 0);
                break;
            }
        return true;
    }
} // namespace Euler2
int main() {
    int t, n, m;
   cin >> t >> n >> m;
   for (int u, v, i = 1; i <= m; i++) {</pre>
        cin >> u >> v;
        g[u].push_back({v, i});
       if (t == 1) g[v].push_back({u, -i});
    }
```

```
bool flag = t == 1 ? Euler1::solve(n) : Euler2::solve(n);

// output

if (!flag || (m > 0 && top - 1 < m))
    puts("NO");

else {
    puts("YES");
    for (int i = top - 1; i > 0; --i) printf("%d%c", stk[i], " \n"[i == 1]);
    }
    return 0;
}
```

点分树

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 N = 2e5 + 10;
11 age[N];
struct edge {
   ll to, val;
};
struct father {
   11 u, num;
   11 dist;
};
struct son {
   11 age, dist;
   bool operator<(const son &s) const {</pre>
       return age < s.age;</pre>
   }
};
vector<father> f[N];
vector<vector<son> > s[N];
vector<edge> g[N];
```

```
bool st[N];
11 siz[N];
11 getsiz(ll u, ll fa) {
   if (st[u]) return 0;
   siz[u] = 1;
   for (auto x:g[u]) {
       if (x.to == fa) continue;
       if (st[x.to]) continue;
       siz[u] += getsiz(x.to, u);
   }
   return siz[u];
}
void getwc(11 u, 11 fa, 11 tot, 11 &wc) {
   if (st[u]) return;
   11 \text{ mmax} = 0, \text{ sum} = 1;
   for (auto x:g[u]) {
       if (x.to == fa) continue;
       if (st[x.to]) continue;
       getwc(x.to, u, tot, wc);
       mmax = max(mmax, siz[x.to]);
       sum += siz[x.to];
   }
   mmax = max(mmax, tot - sum);
   if (2 * mmax <= tot) wc = u;
}
```

```
void getdist(ll u, ll fa, ll now, ll rt, ll kth, vector<son> &v) {
   if (st[u]) return;
   f[u].push_back({rt, kth, now});
   v.push_back({age[u], now});
   for (auto x:g[u]) {
       if (x.to == fa || st[x.to]) continue;
       getdist(x.to, u, now + x.val, rt, kth, v);
   }
}
void calc(ll u) {
   if (st[u]) return;
   getwc(u, -1, getsiz(u, -1), u);
   st[u] = 1;
   for (auto x: g[u]) {
       if (st[x.to]) continue;
       s[u].push back(vector<son>(0));
       auto &v = s[u].back();
       v.push back({-0x3f3f3f3f, 0});
       v.push_back({0x3f3f3f3f, 0});
       getdist(x.to, u, x.val, u, (11) s[u].size() - 1, v);
       sort(v.begin(), v.end(), [](son a, son b) { return a.age <</pre>
b.age; });
       for (11 i = 1; i < v.size(); i++) {
           v[i].dist += v[i - 1].dist;
   }
```

```
for (auto x:g[u]) {
       calc(x.to);
   }
}
11 query(11 u, 11 1, 11 r) {
   11 \text{ ans} = 0;
   for (auto x:f[u]) {
       if (1 \le age[x.u] \& age[x.u] \le r) ans += x.dist;
       for (ll i = 0; i < s[x.u].size(); i++) {</pre>
           if (i == x.num) continue;
           auto &v = s[x.u][i];
           11 btn = lower bound(v.begin(), v.end(), (son) \{1, 0\}) -
v.begin() - 1;
           11 top = upper_bound(v.begin(), v.end(), (son) {r, 0}) -
v.begin() - 1;
           ans += v[top].dist - v[btn].dist;
           ans += (top - btn) * x.dist;
       }
    }
   for (auto v:s[u]) {
       11 btn = lower_bound(v.begin(), v.end(), (son) {1, 0}) -
v.begin() - 1;
       11 top = upper_bound(v.begin(), v.end(), (son) {r, 0}) -
v.begin() - 1;
       ans += v[top].dist - v[btn].dist;
    }
   return ans;
}
```

```
signed main() {
   ios::sync_with_stdio(false);
   cin.tie(nullptr);
   cout.tie(nullptr);
   11 n, q, a;
   cin >> n >> q >> a;
   for (ll i = 1; i <= n; i++) cin >> age[i];
   for (11 i = 1; i < n; i++) {
       11 x, y, z;
       cin >> x >> y >> z;
       g[x].push_back({y, z});
       g[y].push_back({x, z});
   }
   calc(1);
   11 \text{ ans} = 0;
   while (q--) {
       11 u, 1, r;
       cin >> u >> 1 >> r;
       1 = (1 + ans) \% a;
       r = (r + ans) % a;
       if (1 > r) swap(1, r);
       ans = query(u, 1, r);
       cout << ans << endl;</pre>
   }
```

} 线性代数 高斯消元

```
#include <iostream>
#include <vector>
using namespace std;
const double eps = 1e-8;
void sway(vector<double>& a, vector<double>& b) {
   vector<double> s;
   for (int i = 0; i < a.size(); i++) {
       s.push back(a[i]);
   }
   a.clear();
   for (int i = 0; i < b.size(); i++) {
       a.push_back(b[i]);
   b.clear();
   for (int i = 0; i < s.size(); i++) {
       b.push_back(s[i]);
vector<double> gauss jordan(const vector<vector<double> >& A,
                          const vector<double>& b) {
   int n = A.size();
   vector<vector<double> > B(n, vector<double>(n + 1));
   for (int i = 0; i < n; i++)
       for (int j = 0; j < n; j++) B[i][j] = A[i][j];
   for (int i = 0; i < n; i++) B[i][n] = b[i];
   for (int i = 0; i < n; i++) {
       int pivot = i;
       for (int j = i; j < n; j++) {
           if (abs(B[j][i]) > abs(B[pivot][i])) pivot = j;
       swap(B[i], B[pivot]);
       if (abs(B[i][i]) < eps) return vector<double>();
       for (int j = i + 1; j \le n; j++) B[i][j] /= B[i][i];
       for (int j = 0; j < n; j++) {
           if (i != j) {
              for (int k = i + 1; k \le n; k++) B[j][k] -= B[j][i] *
B[i][k];
           }
       }
   vector<double> x(n);
   for (int i = 0; i < n; i++) x[i] = B[i][n];
   return x;
int main() {
   int n, m;
   cin >> n >> m;
   vector<vector<double> > mat(n, vector<double>(m));
   for (int i = 0; i < n; i++) {
```

```
for (int j = 0; j < m; j++) {
            cin >> mat[i][j];
        }
    }
   vector<double> val(n);
   for (int i = 0; i < n; i++) cin >> val[i];
   vector<double> ans = gauss jordan(mat, val);
   for (int i = 0; i < ans.size(); i++) cout << ans[i] << ' ';</pre>
}
矩阵行列式
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 \mod = 1e9 + 7;
struct Matrix {
    static const 11 MAXN = 300;
    11 a[MAXN][MAXN];
   void init() { memset(a, 0, sizeof(a)); }
    11 det(11 n) {
       for (int i = 0; i < n; i++)</pre>
           for (int j = 0; j < n; j++) a[i][j] = (a[i][j] + mod) % mod;</pre>
        11 \text{ res} = 1;
        for (int i = 0; i < n; i++) {</pre>
           if (!a[i][i]) {
               bool flag = false;
               for (int j = i + 1; j < n; j++) {
                   if (a[j][i]) {
                       flag = true;
                        for (int k = i; k < n; k++) {</pre>
                           swap(a[i][k], a[j][k]);
                        }
                       res = -res;
                        break;
                   }
               }
               if (!flag) return 0;
           }
           for (int j = i + 1; j < n; j++) {</pre>
               while (a[j][i]) {
                   ll t = a[i][i] / a[j][i];
                   for (int k = i; k < n; k++) {</pre>
                        a[i][k] = (a[i][k] - t * a[j][k]) % mod;
                       swap(a[i][k], a[j][k]);
                   res = -res;
               }
```

```
}
    res *= a[i][i];
    res %= mod;
}
return (res + mod) % mod;
}
mat;
```

线性基

//

};
//对原集合的每个数 val 转为 2 进制,从高位向低位扫,对于当前位为 1 的,若 lba[i]
不存在就令 lba[i]=x,否则令 val=val`xor`lba[i]
//使用: 直接 insert

// ------线性基模板

}

线性基 2

线性基 能表示的线性空间与原向量 能表示的线性空间等价 用高斯消元得到线性基

val ^= lba[i];

先输入数组 a[] 中

```
int n, k;
ll a[N];
```

```
void getVec() {
   k = 0;
   for(int i = 62; i >= 0; -- i) {
       for(int j = k; j < n; ++ j) {</pre>
            if(a[j] >> i & 1) {
               swap(a[j], a[k]);
               break;
            }
       if(!(a[k] >> i & 1)) continue;
       for(int j = 0; j < n; ++j) {</pre>
            if(j != k && (a[j] >> i & 1)) {
               a[j] ^= a[k];
            }
        }
       ++k;
       if(k == n) break;
}
```

这里注意最后的线性基是 a[]中从 0 到 k-1 个,在前的是高位

```
矩阵(加减乘快速幂
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 N = 305;
const 11 mod = 998244353;
//矩阵类模板
struct Matrix {
   11 n, m;
   ll a[N][N];
   void set(ll _a, ll _b) {
```

```
n = a, m = b;
}
Matrix() {
   clear();
}
void clear() {
   n = m = 0;
   memset(a, 0, sizeof(a));
}
Matrix operator+(const Matrix &b) const {
   Matrix tmp;
   tmp.n = n;
   tmp.m = m;
   for (ll i = 0; i < n; ++i)</pre>
       for (11 j = 0; j < m; ++j)
           tmp.a[i][j] = (a[i][j] + b.a[i][j]) % mod;
    return tmp;
}
Matrix operator-(const Matrix &b) const {
   Matrix tmp;
   tmp.n = n;
   tmp.m = m;
   for (ll i = 0; i < n; ++i) {</pre>
       for (11 j = 0; j < m; ++j)
```

```
tmp.a[i][j] = (a[i][j] - b.a[i][j] + mod) % mod;
   }
    return tmp;
}
Matrix operator*(const Matrix &b) const {
   Matrix tmp;
   tmp.clear();
   tmp.n = n;
   tmp.m = b.m;
   for (ll i = 0; i < n; ++i)</pre>
       for (11 j = 0; j < b.m; ++j)
           for (11 k = 0; k < m; ++k) {
               tmp.a[i][j] += a[i][k] * b.a[k][j];
               tmp.a[i][j] %= mod;
           }
    return tmp;
}
Matrix get(ll x) {//幂运算
   Matrix E;
   E.clear();
   E.set(n, m);
   for (ll i = 0; i < n; ++i)</pre>
       E.a[i][i] = 1;
   if (x == 0) return E;
   else if (x == 1) return *this;
```

```
Matrix tmp = get(x / 2);
   tmp = tmp * tmp;
   if (x % 2) tmp = tmp * (*this);
   return tmp;
}
void exgcd(ll _a, ll _b, ll &x, ll &y) {
   if (!_b)return x = 1, y = 0, void();
   exgcd(_b, _a % _b, y, x);
   y -= x * (_a / _b);
}
11 inv(ll p) {
   11 x, y;
   exgcd(p, mod, x, y);
   return (x + mod) % mod;
}
Matrix inv() {
   Matrix E = *this;
   11 is[N], js[N];
   for (11 k = 0; k < E.n; k++) {
       is[k] = js[k] = -1;
       for (ll i = k; i < E.n; i++) // 1</pre>
           for (ll j = k; j < E.n; j++)</pre>
               if (E.a[i][j]) {
                   is[k] = i, js[k] = j;
                   break;
```

```
}
           if (is[k] == -1) {
               E.clear();
               return E;
           }
           for (ll i = 0; i < E.n; i++) // 2
               swap(E.a[k][i], E.a[is[k]][i]);
           for (ll i = 0; i < E.n; i++)</pre>
               swap(E.a[i][k], E.a[i][js[k]]);
           if (!E.a[k][k]) {
               E.clear();
               return E;
           }
           E.a[k][k] = inv(E.a[k][k]); // 3
           for (11 j = 0; j < E.n; j++)</pre>
               if (j != k) // 4
                   (E.a[k][j] *= E.a[k][k]) %= mod;
           for (ll i = 0; i < E.n; i++)</pre>
               if (i != k) // 5
                   for (11 j = 0; j < E.n; j++)
                       if (j != k)
                           (E.a[i][j] += mod - E.a[i][k] * E.a[k][j] % mo
d) %= mod;
           for (ll i = 0; i < E.n; i++)</pre>
               if (i != k) // 就是这里不同
                   E.a[i][k] = (mod - E.a[i][k] * E.a[k][k] % mod) % mod;
       }
       for (ll k = E.n - 1; k >= 0; k--) { // 6
```

稀疏矩阵乘法

```
struct Matrix{
    int n,m;
   int a[maxn][maxn];////
   void clear(){
       n=m=0;
       memset(a,0,sizeof(a));
   Matrix operator * (const Matrix &b) const{
       Matrix tmp;
       tmp.clear();
       tmp.n=n;tmp.m=b.m;
       for (int k=0;k<m;++k){</pre>
           for (int i=0;i<n;++i){</pre>
            if(a[i][k]==0) continue;
            for(int j=0;j<b.m;++j){</pre>
                   if(b.a[k][j]==0) continue;
                   tmp.a[i][j]+=a[i][k]*b.a[k][j];
                   tmp.a[i][j]%=mod;
                   }
       }
       return tmp;
};
//稀疏矩阵乘法
```

杂项

```
mt19937
#include <random>
#include <iostream>
int main()
{
   std::random device rd; //获取随机数种子
   std::mt19937 gen(rd()); //Standard mersenne_twister_engine seeded wi
th rd()
   std::uniform int distribution<> dis(0, 9);
   for (int n = 0; n < 20; ++n)
       std::cout << dis(gen) << ' ';
   std::cout << '\n';</pre>
   system("pause");
   return 0;
}
//可能的结果: 72214140472109192351
doule: std::uniformrealdistribution<> dis(0, 9);
#include <iostream>
#include <chrono>
#include <random>
using namespace std;
int main()
{
     // 随机数种子
     unsigned seed = std::chrono::system_clock::now().time_since epoch
     mt19937 rand_num(seed); // 大随机数
     uniform_int_distribution<long long> dist(0, 1000000000); // 给定
范围
      cout << dist(rand_num) << endl;</pre>
      return 0;
}
注意: 代码中的 rand num 和 dist 都是自己定义的对象,不是系统的。
洗牌算法
#include <random>
#include <algorithm>
#include <iterator>
#include <iostream>
int main()
```

```
{
   std::vector<int> v = { 1, 2, 3, 4, 5, 6, 7, 8, 9, 10 };
   std::random device rd;
   std::mt19937 g(rd());
   std::shuffle(v.begin(), v.end(), g);
   std::copy(v.begin(), v.end(), std::ostream_iterator<int>(std::cout,
" "));
   std::cout << "\n";</pre>
   system("pause");
   return 0;
}
快读
inline int read(){
   int X=0,w=0;char ch=0;
   while(!isdigit(ch)){w|=ch=='-';ch=getchar();}
   while(isdigit(ch))X=(X<<3)+(X<<1)+(ch^48), ch=getchar();
   return w?-X:X;
}
fread 快读
#include <bits/stdc++.h>
using namespace std;
char next_char() {
      static char buf[1 << 20], *first, *last;</pre>
      if(first == last) {
            last = buf + fread(buf, 1, 1 << 20, stdin);</pre>
            first = buf;
      return first == last ? EOF : *first ++;
}
inline int read(){
      int x = 0, w = 0; char ch = 0;
     while(!isdigit(ch)) {w |= ch == '-'; ch = next_char(); }
      while(isdigit(ch)) \{x = (x << 3) + (x << 1) + (ch ^ 48), ch = next\}
_char(); }
      return w ? -x : x;
}
int main(){
      freopen("1.txt", "r", stdin); // 交代码的时候一定要去掉 aaa
```

```
int T;
cin >> T;
while(T --){
    int x = read();
    cout << x << endl;
}</pre>
```

朝鲜大哥快读

```
#define FI(n) FastIO::read(n)
#define FO(n) FastIO::write(n)
#define Flush FastIO::Fflush()
//程序末尾写上 Flush:
namespace FastIO {
   const int SIZE = 1 << 16;</pre>
   char buf[SIZE], obuf[SIZE], str[60];
   int bi = SIZE, bn = SIZE, opt;
   double D[] = {0.1, 0.01, 0.001, 0.0001, 0.00001, 0.000001,
0.0000001, 0.00000001, 0.000000001, 0.0000000001};
   int read(char *s) {
       while (bn) {
           for (; bi < bn && buf[bi] <= ' '; bi++);</pre>
           if (bi < bn)</pre>
               break;
           bn = fread(buf, 1, SIZE, stdin);
           bi = 0;
       int sn = 0;
       while (bn) {
           for (; bi < bn && buf[bi] > ' '; bi++)
               s[sn++] = buf[bi];
           if (bi < bn)</pre>
               break;
           bn = fread(buf, 1, SIZE, stdin);
           bi = 0;
       }
       s[sn] = 0;
       return sn;
   }
   bool read(int &x) {
       int n = read(str), bf = 0;
       if (!n)
           return 0;
       int i = 0;
       if (str[i] == '-')
           bf = 1, i++;
       else if (str[i] == '+')
           i++;
       for (x = 0; i < n; i++)
           x = x * 10 + str[i] - '0';
       if (bf)
           x = -x;
       return 1;
   }
```

```
bool read(long long &x) {
   int n = read(str), bf;
   if (!n)
       return 0;
   int i = 0;
   if (str[i] == '-')
       bf = -1, i++;
   else
       bf = 1;
   for (x = 0; i < n; i++)
       x = x * 10 + str[i] - '0';
   if (bf < 0)
       x = -x;
   return 1;
}
void write(int x) {
   if (x == 0)
       obuf[opt++] = '0';
   else {
       if (x < 0)
           obuf[opt++] = '-', x = -x;
       int sn = 0;
       while (x)
           str[sn++] = x % 10 + '0', x /= 10;
       for (int i = sn - 1; i >= 0; i--)
           obuf[opt++] = str[i];
   if (opt >= (SIZE >> 1)) {
       fwrite(obuf, 1, opt, stdout);
       opt = 0;
}
void write(long long x) {
   if (x == 0)
       obuf[opt++] = '0';
   else {
       if (x < 0)
           obuf[opt++] = '-', x = -x;
       int sn = 0;
       while (x)
           str[sn++] = x \% 10 + '0', x /= 10;
       for (int i = sn - 1; i >= 0; i--)
           obuf[opt++] = str[i];
   if (opt >= (SIZE >> 1)) {
       fwrite(obuf, 1, opt, stdout);
       opt = 0;
   }
```

```
}
   void write(unsigned long long x) {
      if (x == 0)
          obuf[opt++] = '0';
      else {
          int sn = 0;
          while (x)
             str[sn++] = x % 10 + '0', x /= 10;
          for (int i = sn - 1; i >= 0; i--)
             obuf[opt++] = str[i];
      if (opt >= (SIZE >> 1)) {
          fwrite(obuf, 1, opt, stdout);
          opt = 0;
      }
   }
   void write(char x) {
      obuf[opt++] = x;
      if (opt >= (SIZE >> 1)) {
          fwrite(obuf, 1, opt, stdout);
          opt = 0;
      }
   }
   void Fflush() {
      if (opt)
          fwrite(obuf, 1, opt, stdout);
      opt = 0;
}; // namespace FastIO
模拟退火
"优化的随机算法"
连续函数找区间最优
// 找一个点,与平面中的 n 个点的距离和最近
//进行多次模拟退火避免局部最大值
#include <bits/stdc++.h>
#include <ctime>
using namespace std;
const int maxn = 110;
int n;
```

```
#define x first
#define y second
typedef pair<double, double> PDD;
PDD q[maxn];
double ans = 1e8;
double rand(double 1, double r) {
   return (double) rand() / RAND_MAX * (r - 1) + 1;
}
double getDist(PDD a, PDD b) {
   double dx = a.x - b.x;
   double dy = a.y - b.y;
   return sqrt(dx * dx + dy * dy);
}
double calc(PDD p) {
   double res = 0;
   for(int i = 0; i < n; ++ i) {</pre>
       res += getDist(q[i], p);
   }
   ans = min(ans, res);
   return res;
}
double simulate anneal() {
   PDD cur(rand(0, 10000), rand(0, 10000)); // 随机一个起点
   for(double T = 1e4; T > 1e-4; T = T * 0.99) { // 初始温度,末态温度,
衰减系数,一般调整衰减系数 0.999 0.95
       PDD np(rand(cur.x - T, cur.x + T), rand(cur.y - T, cur.y + T));
// 随机新点
       double delta = calc(np) - calc(cur);
       if(exp(-delta / T) > rand(0, 1)) cur = np; //如果新点比现在的点更
优, 必过去, 不然有一定概率过去
}
int main() {
   cin >> n;
   for(int i = 0; i < n; ++ i) {
       cin >> q[i].x >> q[i].y;
   }
   while((double) clock() / CLOCKS_PER_SEC < 0.8) { // 卡时 // 或 for
```

```
(100)
       simulate_anneal();
   cout << (int)(ans + 0.5) << endl;</pre>
   return 0;
}
//n个点带权费马点 // 平衡点||吊打 XXX
//n 个二维坐标点, 带重物重量, 找平衡点
//进行一次模拟退火,但是在局部最大值周围多次跳动(以提高精度
#include <cmath>
#include <cstdio>
#include <cstdlib>
#include <ctime>
const int N = 10005;
int n, x[N], y[N], w[N];
double ansx, ansy, dis;
double Rand() { return (double)rand() / RAND_MAX; }
double calc(double xx, double yy) {
 double res = 0;
 for (int i = 1; i <= n; ++i) {</pre>
   double dx = x[i] - xx, dy = y[i] - yy;
   res += sqrt(dx * dx + dy * dy) * w[i];
 if (res < dis) dis = res, ansx = xx, ansy = yy;</pre>
 return res;
void simulateAnneal() {
 double t = 100000;
 double nowx = ansx, nowy = ansy;
 while (t > 0.001) {
   double nxtx = nowx + t * (Rand() * 2 - 1);
   double nxty = nowy + t * (Rand() * 2 - 1);
   double delta = calc(nxtx, nxty) - calc(nowx, nowy);
   if (exp(-delta / t) > Rand()) nowx = nxtx, nowy = nxty;
   t *= 0.97;
 for (int i = 1; i <= 1000; ++i) {
   double nxtx = ansx + t * (Rand() * 2 - 1);
   double nxty = ansy + t * (Rand() * 2 - 1);
   calc(nxtx, nxty);
 }
}
```

```
int main() {
  srand(time(∅));
 scanf("%d", &n);
 for (int i = 1; i <= n; ++i) {</pre>
    scanf("%d%d%d", &x[i], &y[i], &w[i]);
    ansx += x[i], ansy += y[i];
  }
 ansx /= n, ansy /= n, dis = calc(ansx, ansy);
  simulateAnneal();
 printf("%.31f %.31f\n", ansx, ansy);
 return 0;
}
整体二分
11 bit[N];
void add_bit(ll k, ll a) {
   while (k < N) {
       bit[k] = bit[k] + a;
       k += k \& -k;
   }
}
11 query_bit(ll k) {
   11 \text{ ans} = 0;
   while (k) {
       ans = ans + bit[k];
       k -= k \& -k;
   return ans;
}
struct node {
   ll x, y, k, id, type;
};
node q[N], q1[N], q2[N];
11 ans[N], now[N], tot, totx;
void solve(11 1, 11 r, 11 q1, 11 qr) {
   if (ql > qr) return;
   if (1 == r) {
       for (11 i = q1; i <= qr; i++) {</pre>
           if (q[i].type == 2) {
               ans[q[i].id] = 1;
           }
       }
       return;
   11 \text{ mid} = (1 + r) >> 1;
```

```
11 cq1 = 0, cq2 = 0;
   for (ll i = ql; i <= qr; i++) {</pre>
       if (q[i].type == 1) {
           if (q[i].y <= mid) {
               add_bit(q[i].x, q[i].k);
               q1[++cq1] = q[i];
           } else {
               q2[++cq2] = q[i];
           }
       } else {
           ll sum = query_bit(q[i].y) - query_bit(q[i].x - 1);
           if (sum >= q[i].k) {
               q1[++cq1] = q[i];
           } else {
               q2[++cq2] = q[i];
               q2[cq2].k -= sum;
           }
       }
   for (ll i = 1; i <= cq1; i++) if (q1[i].type == 1) add_bit(q1[i].x,</pre>
-q1[i].k);
   for (ll i = 1; i <= cq1; i++) q[ql + i - 1] = q1[i];</pre>
   for (ll i = 1; i \le cq2; i++) q[ql + cq1 + i - 1] = q2[i];
   solve(1, mid, ql, ql + cq1 - 1);
   solve(mid + 1, r, ql + cq1, qr);
}
void init() {
   totx = 0;
   tot = 0;
   memset(bit, 0, sizeof bit);
}
```

字符串

```
#include <bits/stdc++.h>
using namespace std;
const int maxn = 100005;
char s[maxn];
char s_new[maxn * 2];
int p[maxn * 2];
int Manacher(char* a, int 1) {
    s new[0] = '$';
}
```

```
s_new[1] = '#';
   int len = 2;
   for (int i = 0; i < 1; i++) {</pre>
       s_{new}[len++] = a[i];
       s_new[len++] = '#';
   s new[len] = '\0';
   int id;
   int mx = 0;
   int mmax = 0;
   for (int i = 1; i < len; i++) {</pre>
       p[i] = i < mx ? min(p[2 * id - i], mx - i) : 1;
       while (s_new[i + p[i]] == s_new[i - p[i]]) p[i]++;
       if (mx < i + p[i]) {
           id = i;
           mx = i + p[i];
       mmax = max(mmax, p[i] - 1);
   }
   return mmax;
}
int main() {
   cin >> s;
   cout << Manacher(s, strlen(s));</pre>
}
AC 自动机
#include <bits/stdc++.h>
using namespace std;
struct AC {
   static const int maxnode = 200005;
   static const int sigma_size = 26;
   char T[maxnode];
   int ch[maxnode][sigma_size];
   int val[maxnode], fail[maxnode], last[maxnode];
   int sz;
   vector<pair<int, int> > ans;
   void init() {
       sz = 1;
       memset(ch[0], 0, sizeof(ch[0]));
       ans.clear();
   }
   int idx(const char &c) { return c - 'a'; }
   void insert(string s, int v) {
```

```
int u = 0, n = s.length();
    for (int i = 0; i < n; i++) {</pre>
        int c = idx(s[i]);
        if (!ch[u][c]) {
           memset(ch[sz], 0, sizeof(ch[sz]));
           val[sz] = 0;
           ch[u][c] = sz++;
        }
        u = ch[u][c];
    val[u] = v;
}
void get_fail() {
    queue<int> que;
    fail[0] = 0;
    for (int c = 0; c < sigma_size; c++) {</pre>
        int u = ch[0][c];
        if (u) {
           fail[u] = 0;
           que.push(u);
           last[u] = 0;
        }
    while (!que.empty()) {
        int r = que.front();
        que.pop();
        for (int c = 0; c < sigma_size; c++) {</pre>
           int u = ch[r][c];
           if (!u) continue;
           que.push(u);
           int v = fail[r];
           while (v && !ch[v][c]) v = fail[v];
           fail[u] = ch[v][c];
           last[u] = val[fail[u]] ? fail[u] : last[fail[u]];
        }
    }
}
void print(int j) {
    if (j) {
        ans.push_back(pair<int, int>(j, val[j]));
        print(last[j]);
    }
}
void find() {
    int n = strlen(T);
    int j = 0;
    for (int i = 0; i < n; i++) {</pre>
```

```
int c = idx(T[i]);
    while (j && !ch[j][c]) j = fail[j];
    j = ch[j][c];
    if (val[j])
        print(j);
    else if (last[j])
        print(last[j]);
    }
}
ac; //字符串下标从 0 开始
```

KMP //next 数组等价于前缀函数 #include<bits/stdc++.h> using namespace std; typedef long long 11; int kmp(char *s1,int *p1,char *s2=0,int *p2=0){//必须先求 s1 的 next 数组, 即 kmp(s1,p1); 再 kmp(s1,p1,s2,p2); int n=strlen(s1); **if**(p2==0){ p1[0]=0; for(int i=1;s1[i]!='\0';i++){ int j=p1[i-1]; while(j>0&&s1[i]!=s1[j])j=p1[j-1]; **if**(s1[i]==s1[j])j++; p1[i]=j; } } else{ for(int i=0;s2[i]!='\0';i++){ int j=i==0?0:p2[i-1]; while(j>0&&s2[i]!=s1[j])j=p1[j-1]; **if**(s2[i]==s1[j])j++; p2[i]=j; if(j==n)return i-n+2;//返回位置 } return 0; } int main(){ char s1[15],s2[105]; int p1[15],p2[105]; cin>>s1>>s2; kmp(s1,p1); cout<<kmp(s1,p1,s2,p2)<<endl;</pre>

```
return 0;
}
KMP 2
#include <bits/stdc++.h>
using namespace std;
struct KMP {
   static const int MAXN = 1000010;
   char T[MAXN], P[MAXN];
   int fail[MAXN];
   vector<int> ans;
   void init() { ans.clear(); }
   void get_fail() {
       int m = strlen(P);
       fail[0] = fail[1] = 0;
       for (int i = 1; i < m; i++) {</pre>
           int j = fail[i];
           while (j && P[i] != P[j]) j = fail[j];
           fail[i + 1] = (P[i] == P[j] ? j + 1 : 0);
       }
   }
   void find() {
       int n = strlen(T), m = strlen(P);
       get_fail();
       int j = 0;
       for (int i = 0; i < n; i++) {</pre>
           while (j && P[j] != T[i]) j = fail[j];
           if (P[j] == T[i]) j++;
           if (j == m) ans.push_back(i - m + 1);
       }
   }
} kmp; //P 为模式串,下标从 0 开始,输入后直接调用 find()
```

```
Tire
#include <bits/stdc++.h>
using namespace std;
struct Trie {
    static const int maxnode = 200005;
    static const int sigma_size = 26;
    int ch[maxnode][sigma_size];
    int val[maxnode];
    int sz;
```

```
Trie() {
       sz = 1;
       memset(ch[0], 0, sizeof(ch[0]));
   }
   int idx(const char &c) { return c - 'a'; }
   void insert(string s, int v) {
       int u = 0, n = s.length();
       for (int i = 0; i < n; i++) {</pre>
           int c = idx(s[i]);
           if (!ch[u][c]) {
               memset(ch[sz], 0, sizeof(ch[sz]));
               val[sz] = 0;
               ch[u][c] = sz++;
           u = ch[u][c];
       val[u] = v;
   }
   int find(string s) {
       int u = 0, n = s.length();
       for (int i = 0; i < n; i++) {</pre>
           int c = idx(s[i]);
           if (!ch[u][c]) return 0;
           u = ch[u][c];
       }
       return val[u];
} trie;
```

后缀数组

```
#include <bits/stdc++.h>
using namespace std;
struct SuffixArray {
    static const int MAXN = 1100000;
    char s[MAXN];
    int sa[MAXN], t[MAXN], t1[MAXN], c[MAXN], ra[MAXN], height[MAXN], m;
    inline void init() { memset(this, 0, sizeof(SuffixArray)); }

    inline void get_sa(int n) {
        m = 256;
        int *x = t, *y = t1;
        for (int i = 1; i <= m; i++) c[i] = 0;
        for (int i = 1; i <= n; i++) c[x[i] = s[i]]++;
        for (int i = 1; i <= m; i++) c[i] += c[i - 1];
        for (int i = n; i >= 1; i--) sa[c[x[i]]--] = i;
```

```
for (int k = 1; k <= n; k <<= 1) {
           int p = 0;
           for (int i = n - k + 1; i <= n; i++) y[++p] = i;
           for (int i = 1; i <= n; i++)</pre>
               if (sa[i] > k) y[++p] = sa[i] - k;
           for (int i = 1; i <= m; i++) c[i] = 0;</pre>
           for (int i = 1; i <= n; i++) c[x[y[i]]]++;
           for (int i = 1; i <= m; i++) c[i] += c[i - 1];</pre>
           for (int i = n; i >= 1; i--) sa[c[x[y[i]]]--] = y[i];
           std::swap(x, y);
           p = x[sa[1]] = 1;
           for (int i = 2; i <= n; i++) {</pre>
               x[sa[i]] = (y[sa[i - 1]] == y[sa[i]] &&
                           y[sa[i - 1] + k] == y[sa[i] + k])
                              ? p
                              : ++p;
           if (p >= n) break;
           m = p;
       }
   }
   inline void get_height(int n) {
       int i, j, k = 0;
       for (int i = 1; i <= n; i++) ra[sa[i]] = i;</pre>
       for (int i = 1; i <= n; i++) {</pre>
           if (k) k--;
           int j = sa[ra[i] - 1];
           while (s[i + k] == s[j + k]) k++;
           height[ra[i]] = k;
       }
   }
} SA;
       //字符串下标从一开始
```

可持久化字典树

```
struct Trie01 {
    static const int maxnode = 2000005;
    static const int sigma_size = 2;
    int ch[maxnode << 5][sigma_size], val[maxnode << 5];
    int rt[maxnode];
    int sz;

Trie01() {
        sz = 0;
        memset(ch[0], 0, sizeof(ch[0]));
    }</pre>
```

```
void insert(int &now, int pre, int v) {
    now = ++sz;
    for (int i = 30; i >= 0; i--) {
        int k = ((v >> i) & 1);
        ch[now][k] = ++sz;
        ch[now][k ^ 1] = ch[pre][k ^ 1];
        val[ch[now][k]] = val[ch[pre][k]] + 1;
        now = ch[now][k];
        pre = ch[pre][k];
    }
}
trie;
```

对拍

```
windows 环境下 bat 对拍
@echo off
:loop
      dataa.exe > data.txt
      biaocheng.exe < data.txt > ac.txt
      A.exe < data.txt > test.txt
      fc ac.txt test.txt
      if not errorlevel 1 goto loop
pause
goto loop
其中要改的部分(标红辽):
@echo off
:loop
 dataa.exe > data.txt
 $\color{red}{biaocheng.exe}$ < data.txt > ac.txt
 $\color{red}{A.exe}$ < data.txt > test.txt
 fc ac.txt test.txt
if not errorlevel 1 goto loop
pause
goto loop
文件以.bat 作为后缀
```

将三个程序(数据生成文件(dataa),标程或暴力代码(biaocheng), 要看的代码(A))放在同一目录下,

记得加 freopen

随机数生成 code

```
#include <iostream>
#include <cstdlib>
#include <ctime>
using namespace std;
int main(){
      freopen("data.txt", "w", stdout);
   srand((int)time(0));
   int T = rand() % 100000;
   cout << T << endl;</pre>
   for (int i = 0; i < T; i++){
      cout << rand() % 100;</pre>
}
```

rand()似乎只有三万多,需要更大的数的话要乘一下