testFile

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STL

bitset

C++ bitset 用法

bitset

C++的 bitset 在 bitset 头文件中，它是一种类似数组的结构，它的每一个元素只能是 0 或 1 ，每个元素仅用 1 bit 空间。

bitset 数组与 vector 数组区别

bitset 声明数组:bitset<100> number[10]

vector 声明数组:vector number[10];

bitset<每个 bitset 元素的长度(没有占满前面全部自动补 0)> 元素

bitset 内置转化函数：可将 bitset 转化为 string,unsigned long,unsigned long long。

构造

```
        bitset<4> bitset1;    //无参构造，长度为 4，默认每一位为 0

    bitset<8> bitset2(12);    //长度为 8，二进制保存，前面用 0 补充

    string s = "100101";
    bitset<10> bitset3(s);    //长度为 10，前面用 0 补充

    char s2[] = "10101";
    bitset<13> bitset4(s2);    //长度为 13，前面用 0 补充

    cout << bitset1 << endl;    //0000
    cout << bitset2 << endl;    //00001100
    cout << bitset3 << endl;    //0000100101
    cout << bitset4 << endl;    //0000000010101
```

函数

```
        bitset<8> foo ("10011011");

    cout << foo.count() << endl;    //5    （count 函数用来求 bitset 中 1 的位数，foo 中共有 5 个 1
    cout << foo.size() << endl;    //8    （size 函数用来求 bitset 的大小，一共有 8 位

    cout << foo.test(0) << endl;    //true    （test 函数用来查下标处的元素是 0 还是 1，并返回 false 或 true，此处 foo[0] 为 1，返回 true
    cout << foo.test(2) << endl;    //false    （同理，foo[2] 为 0，返回 false

    cout << foo.any() << endl;    //true    （any 函数检查 bitset 中是否有 1
    cout << foo.none() << endl;    //false    （none 函数检查 bitset 中是否没有 1
    cout << foo.all() << endl;    //false    （all 函数检查 bitset 中是全部为 1
```

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```
#include <bits/stdc++.h>
#define ll long long
using namespace std;
int t,n,m;
char str[1010];
bitset<500> number[30];
int main() {
        ios::sync_with_stdio(false); cin.tie(0); cout.tie(0);
    //freopen("test.in","r",stdin);
    //freopen("test.out","w",stdout);
        scanf("%d",&t);
        while(t--)
        {
                scanf("%d %d",&n,&m);
                for(int i=0;i<m;i++)
                {
                        scanf("%s",str);
                        number[i]=bitset<500>(str);
                }
                int len=1<<m,ans=m+1;
                for(int i=1;i<len;i++)
                {
                        int t=i,s=0;
                        bitset<500> num(0);
                        for(int j=0;j<m&&t>0;j++)
                        {
                                if(t&1)
                                {
                                        num=num|number[j];
                                        s++;
                                }
                                t>>=1;
                        }
                        if(num.count()==n) ans=min(ans,s);
                }
                if(ans==m+1) printf("-1\n");
                else printf("%d\n",ans);
        }
        return 0;
}
```

windows 环境下的对拍

```
@echo off
:loop
        dataa.exe > data.txt
        biaocheng.exe < data.txt > ac.txt
        A.exe < data.txt > test.txt
        fc ac.txt test.txt
        if not errorlevel 1 goto loop
pause
goto loop
```

其中要改的部分（标红辽）：

@echo off
:loop
 dataa.exe > data.txt
 $\color{red}{biaocheng.exe}$ < data.txt > ac.txt
 $\color{red}{A.exe}$ < data.txt > test.txt
 fc ac.txt test.txt
 if not errorlevel 1 goto loop
pause
goto loop

文件以.bat 作为后缀

将三个程序（数据生成文件（dataa），标程或暴力代码（biaocheng），要看的代码（A））放在同一目录下，

记得加 freopen

随机数记得加 srand((int)time(0));

随机数生成 code

```
#include <iostream>
#include <cstdlib>
#include <ctime>
using namespace std;

int main(){
        freopen("data.txt", "w", stdout);
```

```
    srand((int)time(0));
    int T = rand() % 100000;
    cout << T << endl;

    for (int i = 0; i < T; i++){
          cout << rand() % 100;
    }
}
```

rand() 似乎只有三万多，需要更大的数的话要乘一下

动态规划

占位

图论

KM

```
#include<bits/stdc++.h>

using namespace std;
typedef long long ll;
const ll maxN = 310;
const ll INF = 1e16;

struct KM {
    ll mp[maxN][maxN], link_x[maxN], link_y[maxN], N;
    bool visx[maxN], visy[maxN];
    ll que[maxN << 1], top, fail, pre[maxN];
    ll hx[maxN], hy[maxN], slk[maxN];

    inline ll check(ll i) {
        visx[i] = true;
        if (link_x[i]) {
            que[fail++] = link_x[i];
            return visy[link_x[i]] = true;
        }
        while (i) {
            link_x[i] = pre[i];
            swap(i, link_y[pre[i]]);
        }
        return 0;
    }

    void bfs(ll S) {
        for (ll i = 1; i <= N; i++) {
            slk[i] = INF;
            visx[i] = visy[i] = false;
        }
        top = 0;
        fail = 1;
        que[0] = S;
        visy[S] = true;
        while (true) {
            ll d;
            while (top < fail) {
                for (ll i = 1, j = que[top++]; i <= N; i++)
 {
                    if (!visx[i] && slk[i] >= (d = hx[i] + h
y[j] - mp[i][j])) {
                        pre[i] = j;
                        if (d) slk[i] = d;
                        else if (!check(i)) return;
                    }
                }
            }
            d = INF;
            for (ll i = 1; i <= N; i++) {
                if (!visx[i] && d > slk[i]) d = slk[i];
            }
            for (ll i = 1; i <= N; i++) {
                if (visx[i]) hx[i] += d;
                else slk[i] -= d;
                if (visy[i]) hy[i] -= d;
            }
            for (ll i = 1; i <= N; i++) {
                if (!visx[i] && !slk[i] && !check(i)) retur
n;
            }
        }
    }

    void init() {
        for (ll i = 1; i <= N; i++) {
            link_x[i] = link_y[i] = 0;
            visy[i] = false;
        }
        for (ll i = 1; i <= N; i++) {
            hx[i] = 0;
            for (ll j = 1; j <= N; j++) {
                if (hx[i] < mp[i][j]) hx[i] = mp[i][j];
            }
        }
    }
} km;

int main() {
    ios::sync_with_stdio(0);

    ll n;
    cin >> n;
    ll ans = 0;
    for (int i = 1; i <= n; i++) {
        ll a, b, c, d;
        cin >> a >> b >> c >> d;
        ans += a * a + b * b;
        for (int j = 1; j <= n; j++) {
            km.mp[i][j] = -(c + d * (j - 1)) * (c + d * (j
- 1));
//            cout << -km.mp[i][j] << ' ';
//            cin >> km.mp[i][j];
//            km.mp[i][j] = -km.mp[i][j];
        }
//        cout << endl;
    }
    km.N = n;
    km.init();
    for (int i = 1; i <= km.N; i++) km.bfs(i);
    for (int i = 1; i <= n; i++) ans -= km.mp[i][km.link_x
[i]];
    cout << ans << endl;
}
```

prufer 序列

```
#include <iostream>
#include <cstdio>
#include <cstring>
#include <algorithm>

using namespace std;

const int N = 100010;

int n, m;
int f[N], d[N], p[N];

void tree2prufer()
{
    for (int i = 1; i < n; i ++ )
    {
        scanf("%d", &f[i]);
        d[f[i]] ++ ;
    }

    for (int i = 0, j = 1; i < n - 2; j ++ )
    {
        while (d[j]) j ++ ;
        p[i ++ ] = f[j];
        while (i < n - 2 && -- d[p[i - 1]] == 0 && p[i - 1]
 < j) p[i ++ ] = f[p[i - 1]];
    }

    for (int i = 0; i < n - 2; i ++ ) printf("%d ", p[i]);
}

void prufer2tree()
{
```

```
    for (int i = 1; i <= n - 2; i ++ )
    {
        scanf("%d", &p[i]);
        d[p[i]] ++ ;
    }
    p[n - 1] = n;

    for (int i = 1, j = 1; i < n; i ++, j ++ )
    {
        while (d[j]) j ++ ;
        f[j] = p[i];
        while (i < n - 1 && -- d[p[i]] == 0 && p[i] < j) f
[p[i]] = p[i + 1], i ++ ;
    }

    for (int i = 1; i <= n - 1; i ++ ) printf("%d ", f[i]);
}

int main()
{
    scanf("%d%d", &n, &m);
    if (m == 1) tree2prufer();
    else prufer2tree();

    return 0;
}
```

spfa 最短路及负环

```
#include<bits/stdc++.h>
using namespace std;
typedef long long ll;
const int N = 1 << 20;
struct edge {
    ll to, len;
};

vector<edge> g[N];
ll d[N], cnt[N], vis[N];

bool spfa(ll s, ll n) {
    queue<int> que;
    for (int i = 1; i <= n; i++) {  //防止不连通，全加进去
        que.push(i);
        vis[i] = 1;
    }
    while (!que.empty()) {
        ll p = que.front();
        que.pop();
        vis[p] = 0;
        for (auto x:g[p]) {
            if (d[x.to] > d[p] + x.len) {
                d[x.to] = d[p] + x.len;
                cnt[x.to] = cnt[p] + 1;
                if (!vis[x.to]) {
                    if (cnt[x.to] > n) return 0;
                    vis[x.to] = 1;
                    que.push(x.to);
                }
            }
        }
    }
    return 1;
}
```

一些定理

Hall 定理：若二分图存在完美匹配，且大小为 n，那么取任意 1≤k≤n，均满足 X 集选出 k 个不同的点，它们连向 Y 集的点的个数不小于 k。

二分图匹配（HK 匈牙利匹配）

```
//大量使用了memset，但常数貌似很小？HDU6808 跑了998ms（限制5000ms），然而这个代int main()不是HDU6808的
#include<bits/stdc++.h>
using namespace std;

const int maxn=505;// 最大点数
const int inf=0x3f3f3f3f;// 距离初始值
struct HK_Hungary{//这个板子从1开始，0点不能用,nx为左边点数，ny为右边点数
    int nx,ny;//左右顶点数量
    vector<int>bmap[maxn];
    int cx[maxn];//cx[i]表示左集合i顶点所匹配的右集合的顶点序号
    int cy[maxn]; //cy[i]表示右集合i顶点所匹配的左集合的顶点序号
    int dx[maxn];
    int dy[maxn];
    int dis;
    bool bmask[maxn];
    void init(int a,int b){
        nx=a,ny=b;
        for(int i=0;i<=nx;i++){
            bmap[i].clear();
        }
    }
    void add_edge(int u,int v){
        bmap[u].push_back(v);
    }
    bool searchpath(){//寻找 增广路径
        queue<int>Q;
        dis=inf;
        memset(dx,-1,sizeof(dx));
        memset(dy,-1,sizeof(dy));
        for(int i=1;i<=nx;i++){//cx[i]表示左集合i顶点所匹配的右集合的顶点序号
            if(cx[i]==-1){//将未遍历的节点 入队 并初始化次节点距离为0
                Q.push(i);
                dx[i]=0;
            }
        }//广度搜索增广路径
        while(!Q.empty()){
            int u=Q.front();
            Q.pop();
            if(dx[u]>dis) break;//取右侧节点
            for(int i=0;i<bmap[u].size();i++){
                int v=bmap[u][i];//右侧节点的增广路径的距离
                if(dy[v]==-1){
                    dy[v]=dx[u]+1;//v对应的距离 为u对应距离加1
                    if(cy[v]==-1)dis=dy[v];
                    else{
                        dx[cy[v]]=dy[v]+1;
                        Q.push(cy[v]);
                    }
                }
            }
        }
        return dis!=inf;
    }
    int findpath(int u){//寻找路径 深度搜索
        for(int i=0;i<bmap[u].size();i++){
            int v=bmap[u][i];//如果该点没有被遍历过 并且距离为上一节点+1
            if(!bmask[v]&&dy[v]==dx[u]+1){//对该点染色
                bmask[v]=1;
                if(cy[v]!=-1&&dy[v]==dis)continue;
                if(cy[v]==-1||findpath(cy[v])){
                    cy[v]=u;cx[u]=v;
                    return 1;
                }
            }
        }
        return 0;
    }
    int MaxMatch(){//得到最大匹配的数目
        int res=0;
        memset(cx,-1,sizeof(cx));
        memset(cy,-1,sizeof(cy));
        while(searchpath()){
            memset(bmask,0,sizeof(bmask));
            for(int i=1;i<=nx;i++){
                if(cx[i]==-1){
                    res+=findpath(i);
                }
            }
        }
        return res;
    }
}HK;
```

```
int main(){
    int nn,n,m;
    cin>>nn;
    while(nn--){
        scanf("%d%d",&n,&m);
        HK.init(n,m);//左端点和右端点数量
        for(int i=1;i<=n;i++){
            int snum;
            cin>>snum;
            int v;
            for(int j=1;j<=snum;j++){
                cin>>v;
                HK.add_edge(i,v);//连边
            }
        }
        cout<<HK.MaxMatch()<<endl;//求最大匹配
    }
    return 0;
}
```

带花树

```
/*
����h�� n ���� m ��� ������c������c���
��p�轉
������'��
��h��h�������������'���p�������
�ʒ��� n ������������ i ������'���� i p
�L���t����ý����p�������� 0��
*/

#include<bits/stdc++.h>
using namespace std;
#define I inline int
#define V inline void
#define FOR(i,a,b) for(int i=a;i<=b;i++)
#define REP(u) for(int i=h[u],v;v=e[i].t,i;i=e[i].n)
const int N=1e3+1,M=1e5+1;
queue<int>q;
int n,m,tot,qwq,ans;
int h[N],lk[N],tag[N],fa[N],pre[N],dfn[N];
struct edge{int t,n;}e[M];
V link(int x,int y){lk[x]=y,lk[y]=x;}
V add_edge(int x,int y){
          if(!lk[x]&&!lk[y])link(x,y),ans++;
          e[++tot]=(edge){y,h[x]},h[x]=tot;
          e[++tot]=(edge){x,h[y]},h[y]=tot;
}
V rev(int x){if(x)rev(x[pre][lk]),link(x,pre[x]);}
I find(int x){return fa[x]==x?x:fa[x]=find(fa[x]);}
I lca(int x,int y){
          for(qwq++;;x=x[lk][pre],swap(x,y))
                    if(dfn[x=find(x)]==qwq)return x;
                    else if(x)dfn[x]=qwq;
}
V shrink(int x,int y,int p){
          for(;find(x)!=p;x=pre[y]){
                    pre[x]=y,y=lk[x],fa[x]=fa[y]=p;
                    if(tag[y]==2)tag[y]=1,q.push(y);
          }
}
I blossom(int u){
          FOR(i,1,n)tag[i]=pre[i]=0,fa[i]=i;
          tag[u]=1,q=queue<int>(),q.push(u);
          for(int p;!q.empty();q.pop())REP(u=q.front())
                    if(tag[v]==1)
                              p=lca(u,v),shrink(u,v,p),sh
rink(v,u,p);
                    else if(!tag[v]){
                              pre[v]=u,tag[v]=2;
                              if(!lk[v])return rev(v),1;
                              else tag[lk[v]]=1,q.push(lk
[v]);
                    }
          return 0;
}
int main(){
          scanf("%d%d",&n,&m);
          for(int x,y;m--;add_edge(x,y))scanf("%d%d",&x,&
y);
          FOR(i,1,n)ans+=!lk[i]&&blossom(i);
          cout<<ans<<'\n';
          FOR(i,1,n)cout<<lk[i]<<' ';
          return 0;
}
```

带花树 2

```
// graph
template <typename T>
class graph {
 public:
  struct edge {
    int from;
    int to;
    T cost;
  };
  vector<edge> edges;
  vector<vector<int> > g;
  int n;
  graph(int _n) : n(_n) { g.resize(n); }
  virtual int add(int from, int to, T cost) = 0;
};

// undirectedgraph
template <typename T>
class undirectedgraph : public graph<T> {
 public:
  using graph<T>::edges;
  using graph<T>::g;
  using graph<T>::n;

  undirectedgraph(int _n) : graph<T>(_n) {}
  int add(int from, int to, T cost = 1) {
    assert(0 <= from && from < n && 0 <= to && to < n);
    int id = (int)edges.size();
    g[from].push_back(id);
    g[to].push_back(id);
    edges.push_back({from, to, cost});
    return id;
  }
};

// blossom / find_max_unweighted_matching
template <typename T>
vector<int> find_max_unweighted_matching(const undirected
graph<T> &g) {
  std::mt19937 rng(chrono::steady_clock::now().time_since
_epoch().count());
  vector<int> match(g.n, -1);   // p��
  vector<int> aux(g.n, -1);     // ʰ�����
  vector<int> label(g.n);       // "o" or "i"
  vector<int> orig(g.n);        // ����
  vector<int> parent(g.n, -1);  // ���x̄�
  queue<int> q;
  int aux_time = -1;

  auto lca = [&](int v, int u) {
    aux_time++;
    while (true) {
      if (v != -1) {
        if (aux[v] == aux_time) {  // ū��ù��j̀ X̀
��LCA
          return v;
        }
        aux[v] = aux_time;
        if (match[v] == -1) {
          v = -1;
        } else {
          v = orig[parent[match[v]]];  // ��p̀���κ��x̄
����ψ̀
        }
      }
      swap(v, u);
    }
  };  // lca

  auto blossom = [&](int v, int u, int a) {
    while (orig[v] != a) {
      parent[v] = u;
      u = match[v];
      if (label[u] == 1) {  // ��'����ĭ"o" ̀̀̀̀
̀·
        label[u] = 0;
        q.push(u);
```

```
    }
    orig[v] = orig[u] = a;  // ����
    v = parent[u];
  }
};  // blossom

auto augment = [&](int v) {
  while (v != -1) {
    int pv = parent[v];
    int next_v = match[pv];
    match[v] = pv;
    match[pv] = v;
    v = next_v;
  }
};  // augment

auto bfs = [&](int root) {
  fill(label.begin(), label.end(), -1);
  iota(orig.begin(), orig.end(), 0);
  while (!q.empty()) {
    q.pop();
  }
  q.push(root);
  // ��'����Ï "o", ������"0"����"o", "1"�
���"i"
  label[root] = 0;
  while (!q.empty()) {
    int v = q.front();
    q.pop();
    for (int id : g.g[v]) {
      auto &e = g.edges[id];
      int u = e.from ^ e.to ^ v;
      if (label[u] == -1) {  // ū�ó�õ
        label[u] = 1;        // ��� "i"
        parent[u] = v;
        if (match[u] == -1) {  // ū�ôþ���
          augment(u);          // Ψ������·��
          return true;
        }
        // ū���þ���  ������þ��j灭��queue ��
��翋����
        label[match[u]] = 0;
        q.push(match[u]);
        continue;
      } else if (label[u] == 0 && orig[v] != orig[u]) {
        // ū��Ψõ  ¥ȳ"ĭ"o"  ����ū"��"
        int a = lca(orig[v], orig[u]);
        // ��LCA ∠������
        blossom(u, v, a);
        blossom(v, u, a);
      }
    }
  }
  return false;
};  // bfs

auto greedy = [&]() {
  vector<int> order(g.n);
  // ������ order
  iota(order.begin(), order.end(), 0);
  shuffle(order.begin(), order.end(), rng);

  // ������þ��j�þ��
  for (int i : order) {
    if (match[i] == -1) {
      for (auto id : g.g[i]) {
        auto &e = g.edges[id];
        int to = e.from ^ e.to ^ i;
        if (match[to] == -1) {
          match[i] = to;
          match[to] = i;
          break;
        }
      }
    }
  }
};  // greedy

// ʰ��'����þ��
greedy();
// ��δþ��������·
for (int i = 0; i < g.n; i++) {
  if (match[i] == -1) {
    bfs(i);
  }
  }
  return match;
}
```

强连通（kosaraju）

```
#include <bits/stdc++.h>
using namespace std;
struct SCC {
    static const int MAXV = 100000;
    int V;
    vector<int> g[MAXV], rg[MAXV], vs;
    bool used[MAXV];
    int cmp[MAXV];

    void add_edge(int from, int to) {
        g[from].push_back(to);
        rg[to].push_back(from);
    }

    void dfs(int v) {
        used[v] = 1;
        for (int i = 0; i < g[v].size(); i++) {
            if (!used[g[v][i]]) dfs(g[v][i]);
        }
        vs.push_back(v);
    }

    void rdfs(int v, int k) {
        used[v] = 1;
        cmp[v] = k;
        for (int i = 0; i < rg[v].size(); i++) {
            if (!used[rg[v][i]]) rdfs(rg[v][i], k);
        }
    }

    int solve() {
        memset(used, 0, sizeof(used));
        vs.clear();
        for (int v = 1; v <= V; v++) {
            if (!used[v]) dfs(v);
        }
        memset(used, 0, sizeof(used));
        int k = 0;
        for (int i = (int)vs.size() - 1; i >= 0; i--) {
            if (!used[vs[i]]) rdfs(vs[i], ++k);
        }
        return k;
    }

    void init(int n) {
        V = n;
        vs.clear();
        for (int i = 0; i < MAXV; i++) {
            g[i].clear();
            rg[i].clear();
            used[i] = 0;
            cmp[i] = 0;
        }
    }

} scc;

//记得调用 init()
```

强连通（tarjan 无 vector）

```
#include <bits/stdc++.h>
using namespace std;
struct SCC {
    static const int MAXN = 5000;
    static const int MAXM = 2000000;
    int dfs_clock, edge_cnt = 1, scc_cnt;
    int head[MAXN];
    int dfn[MAXN], lowlink[MAXN];
    int sccno[MAXN];
    stack<int> s;

    struct edge {
        int v, next;
    } e[MAXM];
```

```
    void add_edge(int u, int v) {
        e[edge_cnt].v = v;
        e[edge_cnt].next = head[u];
        head[u] = edge_cnt++;
    }

    void tarjan(int u) {
        int v;
        dfn[u] = lowlink[u] = ++dfs_clock;  //每次 dfs，u 的次序号增加 1
        s.push(u);                          //将 u 入栈
        for (int i = head[u]; i != -1; i = e[i].next)  //访问从 u 出发的边
        {
            v = e[i].v;
            if (!dfn[v])  //如果 v 没被处理过
            {
                tarjan(v);  // dfs(v)
                lowlink[u] = min(lowlink[u], lowlink[v]);
            } else if (!sccno[v])
                lowlink[u] = min(lowlink[u], dfn[v]);
        }
        if (dfn[u] == lowlink[u]) {
            scc_cnt++;
            do {
                v = s.top();
                s.pop();
                sccno[v] = scc_cnt;
            } while (u != v);
        }
    }

    int find_scc(int n) {
        for (int i = 1; i <= n; i++)
            if (!dfn[i]) tarjan(i);
        return scc_cnt;
    }

    void init() {
        scc_cnt = dfs_clock = 0;
        edge_cnt = 1;  //不用初始化 e 数组，省时间
        while (!s.empty()) s.pop();
        memset(head, -1, sizeof(head));
        memset(sccno, 0, sizeof(sccno));
        memset(dfn, 0, sizeof(dfn));
        memset(lowlink, 0, sizeof(lowlink));
    }
} scc;
```

强连通（tarjan）

```
#include <bits/stdc++.h>
using namespace std;

struct SCC {
    static const int MAXN = 100000;
    vector<int> g[MAXN];
    int dfn[MAXN], lowlink[MAXN], sccno[MAXN], dfs_clock, scc_cnt;
    stack<int> S;

    void dfs(int u) {
        dfn[u] = lowlink[u] = ++dfs_clock;
        S.push(u);
        for (int i = 0; i < g[u].size(); i++) {
            int v = g[u][i];
            if (!dfn[v]) {
                dfs(v);
                lowlink[u] = min(lowlink[u], lowlink[v]);
            } else if (!sccno[v]) {
                lowlink[u] = min(lowlink[u], dfn[v]);
            }
        }
        if (lowlink[u] == dfn[u]) {
            ++scc_cnt;
            for (;;) {
                int x = S.top();
                S.pop();
                sccno[x] = scc_cnt;
                if (x == u) break;
            }
        }
    }

    void solve(int n) {
        dfs_clock = scc_cnt = 0;
        memset(sccno, 0, sizeof(sccno));
        memset(dfn, 0, sizeof(dfn));
        memset(lowlink, 0, sizeof(lowlink));
        for (int i = 1; i <= n; i++) {
            if (!dfn[i]) dfs(i);
        }
    }
} scc;

// scc_cnt 为 SCC 计数器，sccno[i]为 i 所在 SCC 的编号
// vector<int> g[MAXN]中加边
//之后再补充 init()
```

拓扑排序

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 100000;

int c[MAXN];
int topo[MAXN], t, V;
vector<int> g[MAXN];

bool dfs(int u) {
    c[u] = -1;
    for (int i = 0; i < g[u].size(); i++) {
        int v = g[u][i];
        if (c[v] < 0)
            return false;
        else if (!c[v] && !dfs(v))
            return false;
    }
    c[u] = 1;
    topo[t--] = u;
    return true;
}

bool toposort(int n) {
    V = n;
    t = n;
    memset(c, 0, sizeof(c));
    for (int u = 1; u <= V; u++)
        if (!c[u] && !dfs(u)) return false;
    return true;
}
```

数链剖分

```
ll fa[N], son[N], dep[N], siz[N], dfn[N], rnk[N], top[N];
ll dfscnt;
vector<ll> g[N];
ll tree[N << 1];
ll lazy[N << 1];

void dfs1(ll u, ll f, ll d) {
    son[u] = -1;
    siz[u] = 1;
    fa[u] = f;
    dep[u] = d;
    for (auto v:g[u]) {
        if (v == f) continue;
        dfs1(v, u, d + 1);
        siz[u] += siz[v];
        if (son[u] == -1 || siz[v] > siz[son[u]]) son[u] = v;
    }
}

void dfs2(ll u, ll t) {
    dfn[u] = ++dfscnt;
    rnk[dfscnt] = u;
    top[u] = t;
    if (son[u] == -1) return;
    dfs2(son[u], t);
    for (auto v:g[u]) {
        if (v == son[u] || v == fa[u]) continue;
        dfs2(v, v);
```

```
    }
}

ll lca(ll a, ll b) {
    while (top[a] != top[b]) {
        if (dep[top[a]] < dep[top[b]]) swap(a, b);
        a = fa[top[a]];
    }
    return dep[a] < dep[b] ? a : b;
}

void init() {
    for (ll i = 0; i < N; i++) g[i].clear();
    for (ll i = 0; i < (N << 1); i++) {
        tree[i] = 0;
        lazy[i] = 0;
    }
    dfscnt = 0;
}


void pushdown(ll k, ll l, ll r) {
    if (k >= N || lazy[k] == 0) return;
    ll len = (r - l + 1) / 2;
    tree[k << 1] = tree[k << 1] + len * lazy[k];
    tree[k << 1 | 1] = tree[k << 1 | 1] + len * lazy[k];
    lazy[k << 1] = lazy[k << 1] + lazy[k];
    lazy[k << 1 | 1] = lazy[k << 1 | 1] + lazy[k];
    lazy[k] = 0;
}

ll merge_range(ll a, ll b) {
    ll ans = a + b;
    return ans;
}

void change_range(ll k, ll l, ll r, ll ql, ll qr, ll x) {
    if (r < ql || qr < l)return;
    if (ql <= l && r <= qr) {
        tree[k] = tree[k] + x * (r - l + 1);
        lazy[k] = lazy[k] + x;
        return;
    }
    pushdown(k, l, r);
    ll mid = (l + r) >> 1;
    change_range(k << 1, l, mid, ql, qr, x);
    change_range(k << 1 | 1, mid + 1, r, ql, qr, x);
    tree[k] = merge_range(tree[k << 1], tree[k << 1 | 1]);
}

ll query_range(ll k, ll l, ll r, ll ql, ll qr) {
    if (r < ql || qr < l)return 0;
    if (ql <= l && r <= qr) {
        return tree[k];
    }
    pushdown(k, l, r);
    ll mid = (l + r) >> 1;
    ll lq = query_range(k << 1, l, mid, ql, qr);
    ll rq = query_range(k << 1 | 1, mid + 1, r, ql, qr);
    return merge_range(lq, rq);
}

ll query_path(ll a, ll b) {
    ll sum = 0;
    while (top[a] != top[b]) {
        if (dep[top[a]] < dep[top[b]]) swap(a, b);
        sum = sum + query_range(1, 1, N, dfn[top[a]], dfn
[a]);
        //dfn[top[a]]~dfn[a]
        a = fa[top[a]];
    }
    if (dep[a] > dep[b]) swap(a, b);
    //点权
    sum = sum + query_range(1, 1, N, dfn[a], dfn[b]);
    //边权
    //if (a != b) sum = sum + query_range(1, 1, N, dfn[a]
+ 1, dfn[b]);
    //dfn[a]~dfn[b],x
    return sum;
}

void change_path(ll a, ll b, ll x) {
    while (top[a] != top[b]) {
        if (dep[top[a]] < dep[top[b]]) swap(a, b);
        change_range(1, 1, N, dfn[top[a]], dfn[a], x);
        //dfn[top[a]]~dfn[a]
        a = fa[top[a]];
    }
    if (dep[a] > dep[b]) swap(a, b);
    //点权
    change_range(1, 1, N, dfn[a], dfn[b], x);
    //边权
    //if (a != b) change_range(1, 1, N, dfn[a] + 1, dfn[b],
 x);
    //dfn[a]~dfn[b],x
}
```

最大流

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;

struct Edge {
    ll from, to, cap, flow;
    Edge(ll a, ll b, ll c, ll d) : from(a), to(b), cap(c),
 flow(d) {}
};

struct Dinic {
    static const ll maxn = 10000;
    static const ll inf = 0x3f3f3f3f3f3f3f3f;
    ll N, M, S, T;
    vector<Edge> edges;
    vector<ll> G[maxn];
    bool vis[maxn];
    ll d[maxn];
    ll cur[maxn];

    void AddEdge(ll from, ll to, ll cap) {
        edges.push_back(Edge(from, to, cap, 0));
        edges.push_back(Edge(to, from, 0, 0));
        M = edges.size();
        G[from].push_back(M - 2);
        G[to].push_back(M - 1);
    }

    bool BFS() {
        memset(vis, 0, sizeof(vis));
        queue<ll> Q;
        Q.push(S);
        d[S] = 0;
        vis[S] = 1;
        while (!Q.empty()) {
            ll x = Q.front();
            Q.pop();
            for (ll i = 0; i < G[x].size(); i++) {
                Edge& e = edges[G[x][i]];
                if (!vis[e.to] && e.cap > e.flow) {
                    vis[e.to] = 1;
                    d[e.to] = d[x] + 1;
                    Q.push(e.to);
                }
            }
        }
        return vis[T];
    }

    ll DFS(ll x, ll a) {
        if (x == T || a == 0) return a;
        ll flow = 0, f;
        for (ll& i = cur[x]; i < G[x].size(); i++) {
            Edge& e = edges[G[x][i]];
            if (d[x] + 1 == d[e.to] &&
                (f = DFS(e.to, min(a, e.cap - e.flow))) > 0)
 {
                e.flow += f;
                edges[G[x][i] ^ 1].flow -= f;
                flow += f;
                a -= f;
                if (a == 0) break;
            }
        }
        return flow;
    }
```

```
    ll Maxflow(ll S, ll T) {
        this->S = S, this->T = T;
        ll flow = 0;
        while (BFS()) {
            memset(cur, 0, sizeof(cur));
            flow += DFS(S, inf);
        }
        return flow;
    }
} MF;

//有源汇上下界最大流，跑完可行流后，s-t 的最大流即为答案

//有源汇上下届最小流，不连无穷边，s-t 跑最大流，再加上 t-s 无穷边，再跑最大流，无穷边流量为答案

//最大权闭合子图
//构造一个新的流网络，建一个源点 s 和汇点 t，从 s 向原图中所有点权为正数的点建一条容量为点权的边，
//从点权为负数的点向 t 建一条容量为点权绝对值的边，原图中各点建的边都建成容量为正无穷的边。
//然后求从 s 到 t 的最小割，再用所有点权为正的权值之和减去最小割，就是我们要求的最大权值和了。

//最大密度子图
//01 分数规划
//addedge(S,V,m),addedge(E,1),addedge(V,T,2*g-deg(v)+m)
//h(g)=n*m-maxflow(S,T)
```

最大流（double）

```
#include <iostream>
#include <cstring>
#include <algorithm>

using namespace std;

struct Dinic {
        static constexpr int N = 10010, M = 100010, INF = 1e8;
        static constexpr double eps = 1e-8;
//      int n, m, S, T;
        int S, T;
        int h[N], e[M], ne[M], idx;
        double f[M];
        int q[N], d[N], cur[N]; // d ??'??1?妓'? ?õ??·?????????6 ??????????C

        void AddEdge(int a, int b, double c)
        {
            e[idx] = b, f[idx] = c, ne[idx] = h[a], h[a] = idx ++ ;
            e[idx] = a, f[idx] = 0, ne[idx] = h[b], h[b] = idx ++ ;
        }

        bool bfs()
        {
            int hh = 0, tt = 0;
            memset(d, -1, sizeof d);
            q[0] = S, d[S] = 0, cur[S] = h[S];
            while (hh <= tt)
            {
                int t = q[hh ++ ];
                for (int i = h[t]; ~i; i = ne[i])
                {
                    int ver = e[i];
                    if (d[ver] == -1 && f[i] > 0)
                    {
                        d[ver] = d[t] + 1;
                        cur[ver] = h[ver];
                        if (ver == T) return true;
                        q[ ++ tt] = ver;
                    }
                }
            }
            return false;
        }

        double find(int u, double limit)
        {
            if (u == T) return limit;
            double flow = 0;
            for (int i = cur[u]; ~i && flow < limit; i = ne[i])
            {
                cur[u] = i;
                int ver = e[i];
                if (d[ver] == d[u] + 1 && f[i] > 0)
                {
                    double t = find(ver, min(f[i], limit - flow));
                    if (t < eps) d[ver] = -1;
                    f[i] -= t, f[i ^ 1] += t, flow += t;
                }
            }
            return flow;
        }

        double Maxflow(int S, int T)
        {
                this->S = S, this->T = T;
            double r = 0, flow;
            while (bfs()) while (flow = find(S, INF)) r += flow;
            return r;
        }
        void init() ////////
        {
                memset(h, -1, sizeof h);
                idx = 0;
        }
} MF;

// ?��init
```

最小费用最大流

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;

struct Edge {
    ll from, to, cap, flow, cost;
    Edge(ll u, ll v, ll c, ll f, ll w):from(u), to(v), cap(c), flow(f), cost(w) {}
};

struct MCMF {
    static const ll maxn = 6000;
    static const ll INF = 0x3f3f3f3f3f3f3f3f;
    ll n, m;
    vector<Edge> edges;
    vector<ll> G[maxn];
    ll inq[maxn];
    ll d[maxn];
    ll p[maxn];
    ll a[maxn];

    void init(ll n) {
        this->n = n;
        for (ll i = 1; i <= n; i++) G[i].clear();
        edges.clear();
    }

    void add_edge(ll from, ll to, ll cap, ll cost) {
        from++,to++;//原板子无法使用 0 点，故修改
        edges.push_back(Edge(from, to, cap, 0, cost));
        edges.push_back(Edge(to, from, 0, 0, -cost));
        m = edges.size();
        G[from].push_back(m - 2);
        G[to].push_back(m - 1);
    }

    bool BellmanFord(ll s, ll t, ll& flow, ll& cost) {
        for (ll i = 1; i <= n; ++i) d[i] = INF;
        memset(inq, 0, sizeof(inq));
        d[s] = 0, inq[s] = 1, p[s] = 0, a[s] = INF;
        queue<ll> Q;
        Q.push(s);
        while (!Q.empty()) {
            ll u = Q.front();
```

```
            Q.pop();
            inq[u] = 0;
            for (ll i = 0; i < G[u].size(); ++i) {
                Edge& e = edges[G[u][i]];
                if (e.cap > e.flow && d[e.to] > d[u] + e.co
st) {
                    d[e.to] = d[u] + e.cost;
                    p[e.to] = G[u][i];
                    a[e.to] = min(a[u], e.cap - e.flow);
                    if (!inq[e.to]) {
                        Q.push(e.to);
                        inq[e.to] = 1;
                    }
                }
            }
        }
        if (d[t] == INF) return false;
        flow += a[t];
        cost += (ll)d[t] * (ll)a[t];
        for (ll u = t; u != s; u = edges[p[u]].from) {
            edges[p[u]].flow += a[t];
            edges[p[u] ^ 1].flow -= a[t];
        }
        return true;
    }

    //需要保证初始网络中没有负权圈
    ll MincostMaxflow(ll s, ll t, ll& cost) {
        s++,t++;//原板子无法使用0点，故修改
        ll flow = 0;
        cost = 0;
        while (BellmanFord(s, t, flow, cost));
        return flow;
    }
} mcmf;  //  若固定流量k，增广时在flow+a>=k 的时候只增广k-f
low 单位的流量，然后终止程序
//下标从0开始
```

最小路径覆盖

对于有向无环图（DAG）

定义：在一个有向图中，找出最少的路径，使得这些路径经过了所有的点。

最小路径覆盖分为**最小不相交路径覆盖**和**最小可相交路径覆盖**。

最小不相交路径覆盖：每一条路径经过的顶点各不相同。

最小可相交路径覆盖：每一条路径经过的顶点可以相同。

DAG 的最小不相交路径覆盖：

把原图的每个点 v 拆成v_x和v_y两个点，如果有一条有向边$A \to B$，就加边$A_x \to B_y$，这样就得到一个二分图，最小路径覆盖=原图的节点数-新图的最大匹配数

DAG 的最小可相交路径覆盖：

先用 floyd 求出原图的传递闭包，即若 a 到 b 有路径，则加边$a \to b$，转化为最小不相交路径覆盖问题

最近公共祖先（倍增）

```
#include <algorithm>
#include <cstdio>
#include <cstring>
#include <iostream>
using namespace std;
const int MAX = 600000;

struct edge {
    int t, nex;
} e[MAX << 1];
int head[MAX], tot;

int depth[MAX], fa[MAX][22], lg[MAX];

void add_edge(int x, int y) {
    e[++tot].t = y;
    e[tot].nex = head[x];
    head[x] = tot;

    e[++tot].t = x;
    e[tot].nex = head[y];
    head[y] = tot;
}

void dfs(int now, int fath) {
    fa[now][0] = fath;
    depth[now] = depth[fath] + 1;
    for (int i = 1; i <= lg[depth[now]]; ++i)
        fa[now][i] = fa[fa[now][i - 1]][i - 1];
    for (int i = head[now]; i; i = e[i].nex)
        if (e[i].t != fath) dfs(e[i].t, now);
}

int lca(int x, int y) {
    if (depth[x] < depth[y]) swap(x, y);
    while (depth[x] > depth[y]) x = fa[x][lg[depth[x] - de
pth[y]] - 1];
    if (x == y) return x;
    for (int k = lg[depth[x]] - 1; k >= 0; --k)
        if (fa[x][k] != fa[y][k]) x = fa[x][k], y = fa[y]
[k];
    return fa[x][0];
}

void init(int n, int root) {
    for (int i = 1; i <= n; ++i) lg[i] = lg[i - 1] + (1 <<
lg[i - 1] == i);
    dfs(root, 0);
}
```

最近公共祖先（线段树）

```
#include <bits/stdc++.h>
using namespace std;
int n, m, root;
const int MAX_N = 500005;
const int MAX = 1 << 20;
vector<int> g[MAX_N];
vector<int> vs;
pair<int, int> tree[MAX * 2 + 10];
int fir[MAX_N];
int fa[MAX_N];
int dep[MAX_N];
void dfs(int k, int p, int d) {
    fa[k] = p;
    dep[k] = d;
    vs.push_back(k);
    for (int i = 0; i < g[k].size(); i++) {
        if (g[k][i] != p) {
            dfs(g[k][i], k, d + 1);
            vs.push_back(k);
        }
    }
}
void build(int k) {
    if (k >= MAX) return;
    build(k << 1);
    build(k << 1 | 1);
    tree[k] = min(tree[k << 1], tree[k << 1 | 1]);
}
pair<int, int> query(int k, int s, int e, int l, int r) {
    if (e < l || r < s) return pair<int, int>(INT_MAX, 0);
    if (l <= s && e <= r) return tree[k];
    return min(query(k << 1, s, (s + e) >> 1, l, r),
               query(k << 1 | 1, ((s + e) >> 1) + 1, e, l,
r));
}
void init() {
    dfs(root, root, 0);
    for (int i = 0; i < MAX * 2 + 10; i++) tree[i] = pair<i
nt, int>(INT_MAX, 0);
    for (int i = MAX; i < MAX + vs.size(); i++)
        tree[i] = pair<int, int>(dep[vs[i - MAX]], vs[i -
MAX]);
    for (int i = 0; i < vs.size(); i++) {
        if (fir[vs[i]] == 0) fir[vs[i]] = i + 1;
    }
    build(1);
```

```
}
int lca(int a, int b) {
    return query(1, 1, MAX, min(fir[a], fir[b]), max(fir
[a], fir[b])).second;
}
int main() {
    scanf("%d%d%d", &n, &m, &root);
    for (int i = 1; i < n; i++) {
        int a, b;
        scanf("%d%d", &a, &b);
        g[a].push_back(b);
        g[b].push_back(a);
    }
    init();
    for (int i = 1; i <= m; i++) {
        int a, b;
        scanf("%d%d", &a, &b);
        printf("%d\n", lca(a, b));
    }
}
```

有源汇上下界最大小流

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;

struct Edge {
    ll from, to, cap, flow, mn;
    Edge(ll a, ll b, ll c, ll d, ll e) : from(a), to(b), ca
p(c), flow(d), mn(e) {}
};

ll n, m;

struct Dinic {
    static const ll maxn = 50010; // 点的大小，记得改
    static const ll inf = 0x3f3f3f3f3f3f3f3f;
    ll N, M, S, T;
    vector<Edge> edges;
    vector<ll> G[maxn];
    bool vis[maxn];
    ll d[maxn];
    ll cur[maxn];

    void AddEdge(ll from, ll to, ll cap, ll c) {
        edges.push_back(Edge(from, to, cap, 0, c));
        edges.push_back(Edge(to, from, 0, 0, c));
        M = edges.size();
        G[from].push_back(M - 2);
        G[to].push_back(M - 1);
    }

    bool BFS() {
        memset(vis, 0, sizeof(vis));
        queue<ll> Q;
        Q.push(S);
        d[S] = 0;
        vis[S] = 1;
        while (!Q.empty()) {
            ll x = Q.front();
            Q.pop();
            for (ll i = 0; i < G[x].size(); i++) {
                Edge& e = edges[G[x][i]];
                if (!vis[e.to] && e.cap > e.flow) {
                    vis[e.to] = 1;
                    d[e.to] = d[x] + 1;
                    Q.push(e.to);
                }
            }
        }
        return vis[T];
    }

    ll DFS(ll x, ll a) {
        if (x == T || a == 0) return a;
        ll flow = 0, f;
        for (ll& i = cur[x]; i < G[x].size(); i++) {
            Edge& e = edges[G[x][i]];
            if (d[x] + 1 == d[e.to] &&
                (f = DFS(e.to, min(a, e.cap - e.flow))) > 0)
{
                e.flow += f;
                edges[G[x][i] ^ 1].flow -= f;
                flow += f;
                a -= f;
                if (a == 0) break;
            }
        }
        return flow;
    }

    void deleteEdge(ll u, ll v) {
        ll siz = edges.size();
        for(ll i = 0; i < siz; ++ i) {
            if(edges[i].from == u && edges[i].to == v) {
                edges[i].cap = edges[i].flow = 0;
                edges[i ^ 1].cap = edges[i ^ 1].flow = 0;

                break;
            }

        }

    }

    ll getValue() {
        return edges[2 * m].flow;
    }

    ll Maxflow(ll S, ll T) {
        this->S = S, this->T = T;
        ll flow = 0;
        while (BFS()) {
            memset(cur, 0, sizeof(cur));
            flow += DFS(S, inf);
        }
        return flow;
    }
} MF;

int main() {
    ll s, t;
    cin >> n >> m >> s >> t;
  // n 个点，m 条边，给的源点汇点

    ll mp[50010] = {0}; // 点的大小，记得改
    for(ll i = 1; i <= m; ++ i) {
        ll a, b, c, d; // 从 a 到 b 有一条下界 c 上界 d 的边
        cin >> a >> b >> c >> d;
        mp[b] += c;
        mp[a] -= c;
        MF.AddEdge(a, b, d - c, c);
    }
    MF.AddEdge(t, s, 1e18, 0); //
    ll tot = 0;
    for(ll i = 1; i <= n; ++ i) {
        if(mp[i] > 0) {
            tot += mp[i];
            MF.AddEdge(0, i , mp[i], 0);
        }
        else {
            MF.AddEdge(i, n + 1, -mp[i], 0);
        }
    }

    if( MF.Maxflow(0, n + 1) != tot) {
        cout << "No Solution" << endl;
    }
    else {
        ll res = MF.getValue(); // 从 t 到 s 边的流量
        MF.deleteEdge(t, s);
      //cout << res + MF.Maxflow(s, t) << endl; // 最大流
        cout << res - MF.Maxflow(t, s) << endl; // 最小流
    }

    return 0;
}
```

朱刘算法

```
#include <iostream>
#include <cstring>
#include <cstdio>
#include <algorithm>
#include <cmath>
```

```
#define x first
#define y second

using namespace std;

typedef pair<double, double> PDD;

const int N = 110;
const double INF = 1e8;

int n, m;
PDD q[N];
bool g[N][N];
double d[N][N], bd[N][N];
int pre[N], bpre[N];
int dfn[N], low[N], ts, stk[N], top;
int id[N], cnt;
bool st[N], ins[N];

void dfs(int u) {
    st[u] = true;
    for (int i = 1; i <= n; i++)
        if (g[u][i] && !st[i])
            dfs(i);
}

bool check_con() {
    memset(st, 0, sizeof st);
    dfs(1);
    for (int i = 1; i <= n; i++)
        if (!st[i])
            return false;
    return true;
}

double get_dist(int a, int b) {
    double dx = q[a].x - q[b].x;
    double dy = q[a].y - q[b].y;
    return sqrt(dx * dx + dy * dy);
}

void tarjan(int u) {
    dfn[u] = low[u] = ++ts;
    stk[++top] = u, ins[u] = true;

    int j = pre[u];
    if (!dfn[j]) {
        tarjan(j);
        low[u] = min(low[u], low[j]);
    } else if (ins[j]) low[u] = min(low[u], dfn[j]);

    if (low[u] == dfn[u]) {
        int y;
        ++cnt;
        do {
            y = stk[top--], ins[y] = false, id[y] = cnt;
        } while (y != u);
    }
}

double work() {
    double res = 0;
    for (int i = 1; i <= n; i++)
        for (int j = 1; j <= n; j++)
            if (g[i][j]) d[i][j] = get_dist(i, j);
            else d[i][j] = INF;

    while (true) {
        for (int i = 1; i <= n; i++) {
            pre[i] = i;
            for (int j = 1; j <= n; j++)
                if (d[pre[i]][i] > d[j][i])
                    pre[i] = j;
        }

        memset(dfn, 0, sizeof dfn);
        ts = cnt = 0;
        for (int i = 1; i <= n; i++)
            if (!dfn[i])
                tarjan(i);

        if (cnt == n) {
            for (int i = 2; i <= n; i++) res += d[pre[i]]
[i];
            break;
        }

        for (int i = 2; i <= n; i++)
            if (id[pre[i]] == id[i])
                res += d[pre[i]][i];

        for (int i = 1; i <= cnt; i++)
            for (int j = 1; j <= cnt; j++)
                bd[i][j] = INF;

        for (int i = 1; i <= n; i++)
            for (int j = 1; j <= n; j++)
                if (d[i][j] < INF && id[i] != id[j]) {
                    int a = id[i], b = id[j];
                    if (id[pre[j]] == id[j]) bd[a][b] = min
(bd[a][b], d[i][j] - d[pre[j]][j]);
                    else bd[a][b] = min(bd[a][b], d[i][j]);
                }

        n = cnt;
        memcpy(d, bd, sizeof d);
    }

    return res;
}

int main() {
    while (~scanf("%d%d", &n, &m)) {
        for (int i = 1; i <= n; i++) scanf("%lf%lf", &q[i].
x, &q[i].y);

        memset(g, 0, sizeof g);
        while (m--) {
            int a, b;
            scanf("%d%d", &a, &b);
            if (a != b && b != 1) g[a][b] = true;
        }

        if (!check_con()) puts("poor snoopy");
        else printf("%.2lf\n", work());
    }

    return 0;
}
```

树上启发式合并

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const int N = 2e5 + 10;

int vis[N], now;

vector<int> g[N];
int fa[N], son[N], siz[N], ans[N];

void insert(int pos) {
    vis[pos] = 1;
    now = now + 1 - vis[pos - 1] - vis[pos + 1];
}

void remove(int pos) {
    vis[pos] = 0;
    now = now - 1 + vis[pos - 1] + vis[pos + 1];
}

void dfs1(ll u, ll f) {
    siz[u] = 1;
    fa[u] = f;
    son[u] = -1;
    for (auto v:g[u]) {
        if (v == f) continue;
        dfs1(v, u);
        siz[u] += siz[v];
        if (son[u] == -1 || siz[v] > siz[son[u]]) son[u] =
v;
    }
}

void add(int u, int exc, int op) {
```

```
    if (op) insert(u);
    else remove(u);
    for (auto x:g[u]) {
        if (x == fa[u] || x == exc) continue;
        add(x, exc, op);
    }
}

void dfs(ll u, ll opt) {
    for (auto x:g[u]) {
        if (x == fa[u] || x == son[u]) continue;
        dfs(x, 0);
    }
    if (son[u] != -1) dfs(son[u], 1);
    add(u, son[u], 1);
    ans[u] = now;
    if (!opt) {
        add(u, 0, 0);
    }
}


int main() {
    ios::sync_with_stdio(false),
            cin.tie(nullptr),
            cout.tie(nullptr);
    int t;
    cin >> t;
    int test = 0;
    while (t--) {
        int n;
        cin >> n;
        for (int i = 1; i < n; i++) {
            int a, b;
            cin >> a >> b;
            g[a].push_back(b);
            g[b].push_back(a);
        }
        cout << "Case #" << ++test << ": ";
        dfs1(1, -1);
        dfs(1, 0);
        for (int i = 1; i <= n; i++) {
            if (i != 1) cout << ' ';
            cout << ans[i];
        }
        cout << endl;
        for (int i = 1; i <= n; i++) g[i].clear();
    }
}
```

树分治

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 10005;
const int INF = 1000000000;
struct edge {
    int to, length;
    edge() {}
    edge(int a, int b) : to(a), length(b) {}
};


vector<edge> g[MAXN];

bool centroid[MAXN];
int subtree_size[MAXN];

int ans;

//计算子树大小
int compute_subtree_size(int v, int p) {
    int c = 1;
    for (int i = 0; i < g[v].size(); i++) {
        int w = g[v][i].to;
        if (w == p || centroid[w]) continue;
        c += compute_subtree_size(w, v);
    }
    subtree_size[v] = c;
    return c;
}

//查找重心，t 为连通分量大小
// pair（最大子树顶点数，顶点编号）
pair<int, int> search_centroid(int v, int p, int t) {
    pair<int, int> res = pair<int, int>(INF, -1);
    int s = 1, m = 0;
    for (int i = 0; i < g[v].size(); i++) {
        int w = g[v][i].to;
        if (w == p || centroid[w]) continue;
        res = min(res, search_centroid(w, v, t));
        m = max(m, subtree_size[w]);
        s += subtree_size[w];
    }
    m = max(m, t - s);
    res = min(res, pair<int, int>(m, v));
    return res;
}

void init(int n) {
    memset(centroid, 0, sizeof(centroid));
    memset(subtree_size, 0, sizeof(subtree_size));
    for (int i = 0; i <= n; i++) g[i].clear();
    ans = 0;
}

int solve(int u) {
    compute_subtree_size(u, -1);
    int s = search_centroid(u, -1, subtree_size[u]).second;
    centroid[s] = 1;
    for (int i = 0; i < g[s].size(); i++) {
        int v = g[s][i].to;
        if (centroid[v]) continue;
        /*solve()*/
    }
    /*do something*/
    centroid[s] = 0;
    return ans;
}
```

欧拉回路

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const int N = 1e6 + 10;


int stk[N], top;
struct edge {
    int to, idx;
};

vector<edge> g[N];

namespace Euler1 {  //有向图欧拉回路
    bool vis[N];
    int cur[N];

    void dfs(int u, const int &w) {
        vis[abs(w)] = true;
        for (int &i = cur[u]; i < g[u].size();) {
            int idx = g[u][i].idx, v = g[u][i].to;
            i++;
            if (!vis[abs(idx)]) dfs(v, idx);
        }
        stk[++top] = w;
    }

    bool solve(int n) {
        // init();
        for (int i = 0; i <= n; i++) cur[i] = 0;
        for (int i = 0; i <= n; i++) vis[i] = 0;
        // calculate degree
        for (int i = 1; i <= n; i++) {
            if (g[i].size() & 1) return false;
        }
        // Hierholzer
        for (int i = 1; i <= n; i++)
            if (!g[i].empty()) {
                dfs(i, 0);
                break;
            }
        return true;
```

```
    }
} // namespace Euler1

namespace Euler2 {  // 无向图欧拉回路
    int deg[N], cur[N];

    void dfs(int u, const int &w) {
        for (int &i = cur[u]; i < g[u].size();) {
            int idx = g[u][i].idx, v = g[u][i].to;
            i++;
            dfs(v, idx);
        }
        stk[++top] = w;
    }

    bool solve(int n) {
        // init
        for (int i = 0; i <= n; i++) deg[i] = 0;
        for (int i = 0; i <= n; i++) cur[i] = 0;
        // calculate degree
        for (int i = 1; i <= n; ++i) {
            for (auto x: g[i]) deg[i]++, deg[x.to]--;
        }
        for (int i = 1; i <= n; ++i)
            if (deg[i]) return false;
        // Hierholzer
        for (int i = 1; i <= n; ++i)
            if (!g[i].empty()) {
                dfs(i, 0);
                break;
            }
        return true;
    }
} // namespace Euler2

int main() {
    int t, n, m;
    cin >> t >> n >> m;
    for (int u, v, i = 1; i <= m; i++) {
        cin >> u >> v;
        g[u].push_back({v, i});
        if (t == 1) g[v].push_back({u, -i});
    }
    // solve
    bool flag = t == 1 ? Euler1::solve(n) : Euler2::solve
(n);
    // output
    if (!flag || (m > 0 && top - 1 < m))
        puts("NO");
    else {
        puts("YES");
        for (int i = top - 1; i > 0; --i) printf("%d%c", st
k[i], " \n"[i == 1]);
    }
    return 0;
}
```

点分树

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const ll N = 2e5 + 10;

ll age[N];
struct edge {
    ll to, val;
};

struct father {
    ll u, num;
    ll dist;
};

struct son {
    ll age, dist;

    bool operator<(const son &s) const {
        return age < s.age;
    }
};

vector<father> f[N];
vector<vector<son> > s[N];
vector<edge> g[N];
bool st[N];
ll siz[N];

ll getsiz(ll u, ll fa) {
    if (st[u]) return 0;
    siz[u] = 1;
    for (auto x:g[u]) {
        if (x.to == fa) continue;
        if (st[x.to]) continue;
        siz[u] += getsiz(x.to, u);
    }
    return siz[u];
}

void getwc(ll u, ll fa, ll tot, ll &wc) {
    if (st[u]) return;
    ll mmax = 0, sum = 1;
    for (auto x:g[u]) {
        if (x.to == fa) continue;
        if (st[x.to]) continue;
        getwc(x.to, u, tot, wc);
        mmax = max(mmax, siz[x.to]);
        sum += siz[x.to];
    }
    mmax = max(mmax, tot - sum);
    if (2 * mmax <= tot) wc = u;
}

void getdist(ll u, ll fa, ll now, ll rt, ll kth, vector<so
n> &v) {
    if (st[u]) return;
    f[u].push_back({rt, kth, now});
    v.push_back({age[u], now});
    for (auto x:g[u]) {
        if (x.to == fa || st[x.to]) continue;
        getdist(x.to, u, now + x.val, rt, kth, v);
    }
}

void calc(ll u) {
    if (st[u]) return;
    getwc(u, -1, getsiz(u, -1), u);

    st[u] = 1;

    for (auto x: g[u]) {
        if (st[x.to]) continue;
        s[u].push_back(vector<son>(0));
        auto &v = s[u].back();
        v.push_back({-0x3f3f3f3f, 0});
        v.push_back({0x3f3f3f3f, 0});
        getdist(x.to, u, x.val, u, (ll) s[u].size() - 1,
v);
        sort(v.begin(), v.end(), [](son a, son b) { return
 a.age < b.age; });
        for (ll i = 1; i < v.size(); i++) {
            v[i].dist += v[i - 1].dist;
        }
    }
    for (auto x:g[u]) {
        calc(x.to);
    }
}

ll query(ll u, ll l, ll r) {
    ll ans = 0;
    for (auto x:f[u]) {
        if (l <= age[x.u] && age[x.u] <= r) ans += x.dist;
        for (ll i = 0; i < s[x.u].size(); i++) {
            if (i == x.num) continue;
            auto &v = s[x.u][i];
            ll btn = lower_bound(v.begin(), v.end(), (son)
{l, 0}) - v.begin() - 1;
            ll top = upper_bound(v.begin(), v.end(), (son)
{r, 0}) - v.begin() - 1;
            ans += v[top].dist - v[btn].dist;
            ans += (top - btn) * x.dist;
        }
    }
    for (auto v:s[u]) {
        ll btn = lower_bound(v.begin(), v.end(), (son) {l,
 0}) - v.begin() - 1;
```

```
        ll top = upper_bound(v.begin(), v.end(), (son) {r,
 0}) - v.begin() - 1;
        ans += v[top].dist - v[btn].dist;

    }
    return ans;
}

signed main() {
    ios::sync_with_stdio(false);
    cin.tie(nullptr);
    cout.tie(nullptr);

    ll n, q, a;
    cin >> n >> q >> a;
    for (ll i = 1; i <= n; i++) cin >> age[i];
    for (ll i = 1; i < n; i++) {
        ll x, y, z;
        cin >> x >> y >> z;
        g[x].push_back({y, z});
        g[y].push_back({x, z});
    }

    calc(1);

    ll ans = 0;
    while (q--) {
        ll u, l, r;
        cin >> u >> l >> r;
        l = (l + ans) % a;
        r = (r + ans) % a;
        if (l > r) swap(l, r);
        ans = query(u, l, r);
        cout << ans << endl;
    }
}
```

虚树

```
ll fa[N], son[N], dep[N], siz[N], dfn[N], rnk[N], top[N];
ll dfscnt;
vector<ll> g[N];
ll mmin[N];

void dfs1(ll u, ll f, ll d) {
    son[u] = -1;
    siz[u] = 1;
    fa[u] = f;
    dep[u] = d;
    for (auto v:g[u]) {
        if (v == f) continue;
        dfs1(v, u, d + 1);
        siz[u] += siz[v];
        if (son[u] == -1 || siz[v] > siz[son[u]]) son[u] =
v;
    }
}

void dfs2(ll u, ll t) {
    dfn[u] = ++dfscnt;
    rnk[dfscnt] = u;
    top[u] = t;
    if (son[u] == -1) return;
    dfs2(son[u], t);
    for (auto v:g[u]) {
        if (v == son[u] || v == fa[u]) continue;
        dfs2(v, v);
    }
}

ll lca(ll a, ll b) {
    while (top[a] != top[b]) {
        if (dep[top[a]] < dep[top[b]]) swap(a, b);
        a = fa[top[a]];
    }
    return dep[a] < dep[b] ? a : b;
}

struct edge {
    ll s, t, v;
};
edge e[N];
```

```
vector<int> vg[N];
int sta[N], tot;
int h[N];

void build(int *H, int num) {
    sort(H + 1, H + 1 + num, [](int a, int b) { return dfn
[a] < dfn[b]; });
    sta[tot = 1] = 1, vg[1].clear();// 1 号节点入栈，清空 1
号节点对应的邻接表，设置邻接表边数为 1
    for (int i = 1, l; i <= num; ++i) {
        if (H[i] == 1) continue; //如果 1 号节点是关键节点就
不要重复添加
        l = lca(H[i], sta[tot]); //计算当前节点与栈顶节点的
LCA
        if (l != sta[tot]) { //如果 LCA 和栈顶元素不同，则说
明当前节点不再当前栈所存的链上
            while (dfn[l] < dfn[sta[tot - 1]]) {//当次大节
点的 Dfs 序大于 LCA 的 Dfs 序
                vg[sta[tot - 1]].push_back(sta[tot]);
                vg[sta[tot]].push_back(sta[tot - 1]);
                tot--;
            } //把与当前节点所在的链不重合的链连接掉并且弹出
            if (dfn[l] > dfn[sta[tot - 1]]) { //如果 LCA 不
等于次大节点（这里的大于其实和不等于没有区别）
                vg[l].clear();
                vg[l].push_back(sta[tot]);
                vg[sta[tot]].push_back(l);
                sta[tot] = l;//说明 LCA 是第一次入栈，清空其
邻接表，连边后弹出栈顶元素，并将 LCA 入栈
            } else {
                vg[l].push_back(sta[tot]);
                vg[sta[tot]].push_back(l);
                tot--; //说明 LCA 就是次大节点，直接弹出栈顶元
素
            }
        }
        vg[H[i]].clear();
        sta[++tot] = H[i];
        //当前节点必然是第一次入栈，清空邻接表并入栈
    }
    for (int i = 1; i < tot; ++i) {
        vg[sta[i]].push_back(sta[i + 1]);
        vg[sta[i + 1]].push_back(sta[i]);
    } //剩余的最后一条链连接一下
    return;
}
```

多项式

多项式全家桶

```
//#pragma GCC optimize(2)
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;

const int N = 3000007;
const int p = 998244353, gg = 3, ig = 332738118, img = 865
83718;
const int mod = 998244353;

int qpow(int a, int b) {
    int res = 1;
    while (b) {
        if (b & 1) res = 1ll * res * a % mod;
        a = 1ll * a * a % mod;
        b >>= 1;
    }
    return res;
}

namespace Poly {
#define mul(x, y) (1ll * x * y >= mod ? 1ll * x * y % mod :
 1ll * x * y)
#define minus(x, y) (1ll * x - y < 0 ? 1ll * x - y + mod :
1ll * x - y)
#define plus(x, y) (1ll * x + y >= mod ? 1ll * x + y - mod :
 1ll * x + y)
```

```
#define ck(x) (x >= mod ? x - mod : x)//取模运算太慢了

    typedef vector<int> poly;
    const int G = 3;//根据具体的模数而定，原根可不一定不一样！！！
    //一般模数的原根为 2 3 5 7 10 6
    const int inv_G = qpow(G, mod - 2);
    int RR[N], inv[N];
//    int deer[2][22][N];

    void init(const int t) {//预处理出来NTT里需要的w和wn，砍掉了一个log的时间
//        for (int p = 1; p <= t; ++p) {
//            int buf1 = qpow(G, (mod - 1) / (1 << p));
//            int buf0 = qpow(inv_G, (mod - 1) / (1 << p));
//            deer[0][p][0] = deer[1][p][0] = 1;
//            for (int i = 1; i < (1 << p); ++i) {
//                deer[0][p][i] = 1ll * deer[0][p][i - 1] * buf0 % mod;//逆
//                deer[1][p][i] = 1ll * deer[1][p][i - 1] * buf1 % mod;
//            }
//        }
        inv[1] = 1;
        for (int i = 2; i <= (1 << t); ++i)
            inv[i] = 1ll * inv[mod % i] * (mod - mod / i) % mod;
    }

    int NTT_init(int n) {//快速数论变换预处理
        int limit = 1, L = 0;
        while (limit <= n) limit <<= 1, L++;
        for (int i = 0; i < limit; ++i)
            RR[i] = (RR[i >> 1] >> 1) | ((i & 1) << (L - 1));
        return limit;
    }

//  省空间用
    int deer[2][N];

    void NTT(poly &A, int type, int limit) {//快速数论变换
        A.resize(limit);
        for (int i = 0; i < limit; ++i)
            if (i < RR[i])
                swap(A[i], A[RR[i]]);
        for (int mid = 2, j = 1; mid <= limit; mid <<= 1, ++j) {
            int len = mid >> 1;

//          省空间用
            int buf1 = qpow(G, (mod - 1) / (1 << j));
            int buf0 = qpow(inv_G, (mod - 1) / (1 << j));
            deer[0][0] = deer[1][0] = 1;
            for (int i = 1; i < (1 << j); ++i) {
                deer[0][i] = 1ll * deer[0][i - 1] * buf0 % mod;//逆
                deer[1][i] = 1ll * deer[1][i - 1] * buf1 % mod;
            }

            for (int pos = 0; pos < limit; pos += mid) {
//                int *wn = deer[type][j];
//              省空间用
                int *wn = deer[type];
                for (int i = pos; i < pos + len; ++i, ++wn) {
                    int tmp = 1ll * (*wn) * A[i + len] % mod;
                    A[i + len] = ck(A[i] - tmp + mod);
                    A[i] = ck(A[i] + tmp);
                }
            }
        }
        if (type == 0) {
            for (int i = 0; i < limit; ++i)
                A[i] = 1ll * A[i] * inv[limit] % mod;
        }
    }

    poly poly_mul(poly A, poly B) {//多项式乘法
        int deg = A.size() + B.size() - 1;
        int limit = NTT_init(deg);
        poly C(limit);
        NTT(A, 1, limit);
        NTT(B, 1, limit);
        for (int i = 0; i < limit; ++i)
            C[i] = 1ll * A[i] * B[i] % mod;
        NTT(C, 0, limit);
        C.resize(deg);
        return C;
    }

    poly poly_inv(poly &f, int deg) {//多项式求逆
        if (deg == 1)
            return poly(1, qpow(f[0], mod - 2));

        poly A(f.begin(), f.begin() + deg);
        poly B = poly_inv(f, (deg + 1) >> 1);
        int limit = NTT_init(deg << 1);
        NTT(A, 1, limit), NTT(B, 1, limit);
        for (int i = 0; i < limit; ++i)
            A[i] = B[i] * (2 - 1ll * A[i] * B[i] % mod + mod) % mod;
        NTT(A, 0, limit);
        A.resize(deg);
        return A;
    }

    poly poly_dev(poly f) {//多项式求导
        int n = f.size();
        for (int i = 1; i < n; ++i) f[i - 1] = 1ll * f[i] * i % mod;
        return f.resize(n - 1), f;//f[0] = 0，这里直接扔了，从1开始
    }

    poly poly_idev(poly f) {//多项式求积分
        int n = f.size();
        for (int i = n - 1; i; --i) f[i] = 1ll * f[i - 1] * inv[i] % mod;
        return f[0] = 0, f;
    }

    poly poly_ln(poly f, int deg) {//多项式求对数
        poly A = poly_idev(poly_mul(poly_dev(f), poly_inv(f, deg)));
        return A.resize(deg), A;
    }

    poly poly_exp(poly &f, int deg) {//多项式求指数
        if (deg == 1)
            return poly(1, 1);

        poly B = poly_exp(f, (deg + 1) >> 1);
        B.resize(deg);
        poly lnB = poly_ln(B, deg);
        for (int i = 0; i < deg; ++i)
            lnB[i] = ck(f[i] - lnB[i] + mod);

        int limit = NTT_init(deg << 1);//n -> n^2
        NTT(B, 1, limit), NTT(lnB, 1, limit);
        for (int i = 0; i < limit; ++i)
            B[i] = 1ll * B[i] * (1 + lnB[i]) % mod;
        NTT(B, 0, limit);
        B.resize(deg);
        return B;
    }

    poly poly_sqrt(poly &f, int deg) {//多项式开方
        if (deg == 1) return poly(1, 1);
        poly A(f.begin(), f.begin() + deg);
        poly B = poly_sqrt(f, (deg + 1) >> 1);
        poly IB = poly_inv(B, deg);
        int limit = NTT_init(deg << 1);
        NTT(A, 1, limit), NTT(IB, 1, limit);
        for (int i = 0; i < limit; ++i)
            A[i] = 1ll * A[i] * IB[i] % mod;
        NTT(A, 0, limit);
        for (int i = 0; i < deg; ++i)
            A[i] = 1ll * (A[i] + B[i]) * inv[2] % mod;
        A.resize(deg);
        return A;
    }

    poly poly_pow(poly f, int k) {//多项式快速幂
        if(f.size()==1){
            f[0] = qpow(f[0],k);
```

```
            return f;
        }
        f = poly_ln(f, f.size());
        for (auto &x: f) x = 1ll * x * k % mod;
        return poly_exp(f, f.size());
    }

    poly poly_cos(poly f, int deg) {//多项式三角函数（cos）
        poly A(f.begin(), f.begin() + deg);
        poly B(deg), C(deg);
        for (int i = 0; i < deg; ++i)
            A[i] = 1ll * A[i] * img % mod;

        B = poly_exp(A, deg);
        C = poly_inv(B, deg);
        int inv2 = qpow(2, mod - 2);
        for (int i = 0; i < deg; ++i)
            A[i] = 1ll * (1ll * B[i] + C[i]) % mod * inv2 %
 mod;
        return A;
    }

    poly poly_sin(poly f, int deg) {//多项式三角函数（sin）
        poly A(f.begin(), f.begin() + deg);
        poly B(deg), C(deg);
        for (int i = 0; i < deg; ++i)
            A[i] = 1ll * A[i] * img % mod;

        B = poly_exp(A, deg);
        C = poly_inv(B, deg);
        int inv2i = qpow(img << 1, mod - 2);
        for (int i = 0; i < deg; ++i)
            A[i] = 1ll * (1ll * B[i] - C[i] + mod) % mod *
inv2i % mod;
        return A;
    }

    poly poly_arcsin(poly f, int deg) {
        poly A(f.size()), B(f.size()), C(f.size());
        A = poly_dev(f);
        B = poly_mul(f, f);
        for (int i = 0; i < deg; ++i)
            B[i] = minus(mod, B[i]);
        B[0] = plus(B[0], 1);
        C = poly_sqrt(B, deg);
        C = poly_inv(C, deg);
        C = poly_mul(A, C);
        C = poly_idev(C);
        return C;
    }

    poly poly_arctan(poly f, int deg) {
        poly A(f.size()), B(f.size()), C(f.size());
        A = poly_dev(f);
        B = poly_mul(f, f);
        B[0] = plus(B[0], 1);
        C = poly_inv(B, deg);
        C = poly_mul(A, C);
        C = poly_idev(C);
        return C;
    }
}

using namespace Poly;
```

字符串

AC 自动机

```
#include <bits/stdc++.h>
using namespace std;
struct AC {
    static const int maxnode = 200005;
    static const int sigma_size = 26;
    char T[maxnode];
    int ch[maxnode][sigma_size];
    int val[maxnode], fail[maxnode], last[maxnode];
    int sz;
    vector<pair<int, int> > ans;

    void init() {
        sz = 1;
        memset(ch[0], 0, sizeof(ch[0]));
        ans.clear();
    }

    int idx(const char &c) { return c - 'a'; }

    void insert(string s, int v) {
        int u = 0, n = s.length();
        for (int i = 0; i < n; i++) {
            int c = idx(s[i]);
            if (!ch[u][c]) {
                memset(ch[sz], 0, sizeof(ch[sz]));
                val[sz] = 0;
                ch[u][c] = sz++;
            }
            u = ch[u][c];
        }
        val[u] = v;
    }

    void get_fail() {
        queue<int> que;
        fail[0] = 0;
        for (int c = 0; c < sigma_size; c++) {
            int u = ch[0][c];
            if (u) {
                fail[u] = 0;
                que.push(u);
                last[u] = 0;
            }
        }
        while (!que.empty()) {
            int r = que.front();
            que.pop();
            for (int c = 0; c < sigma_size; c++) {
                int u = ch[r][c];
                if (!u) continue;
                que.push(u);
                int v = fail[r];
                while (v && !ch[v][c]) v = fail[v];
                fail[u] = ch[v][c];
                last[u] = val[fail[u]] ? fail[u] : last[fai
l[u]];
            }
        }
    }

    void print(int j) {
        if (j) {
            ans.push_back(pair<int, int>(j, val[j]));
            print(last[j]);
        }
    }

    void find() {
        int n = strlen(T);
        int j = 0;
        for (int i = 0; i < n; i++) {
            int c = idx(T[i]);
            while (j && !ch[j][c]) j = fail[j];
            j = ch[j][c];
            if (val[j])
                print(j);
            else if (last[j])
                print(last[j]);
        }
    }
} ac;   //字符串下标从 0 开始
```

KMP 2

```
#include <bits/stdc++.h>
using namespace std;
struct KMP {
    static const int MAXN = 1000010;
    char T[MAXN], P[MAXN];
    int fail[MAXN];
    vector<int> ans;

    void init() { ans.clear(); }

    void get_fail() {
```

```
        int m = strlen(P);
        fail[0] = fail[1] = 0;
        for (int i = 1; i < m; i++) {
            int j = fail[i];
            while (j && P[i] != P[j]) j = fail[j];
            fail[i + 1] = (P[i] == P[j] ? j + 1 : 0);
        }
    }

    void find() {
        int n = strlen(T), m = strlen(P);
        get_fail();
        int j = 0;
        for (int i = 0; i < n; i++) {
            while (j && P[j] != T[i]) j = fail[j];
            if (P[j] == T[i]) j++;
            if (j == m) ans.push_back(i - m + 1);
        }
    }
} kmp;  //P 为模式串，下标从 0 开始，输入后直接调用 find()
```

kmp

```
//next 数组等价于前缀函数
#include<bits/stdc++.h>
using namespace std;
typedef long long ll;

int kmp(char *s1,int *p1,char *s2=0,int *p2=0){//必须先求 s1 的 next 数组，即 kmp(s1,p1);再 kmp(s1,p1,s2,p2);
    int n=strlen(s1);
    if(p2==0){
        p1[0]=0;
        for(int i=1;s1[i]!='\0';i++){
            int j=p1[i-1];
            while(j>0&&s1[i]!=s1[j])j=p1[j-1];
            if(s1[i]==s1[j])j++;
            p1[i]=j;
        }
    }
    else{
        for(int i=0;s2[i]!='\0';i++){
            int j=i==0?0:p2[i-1];
            while(j>0&&s2[i]!=s1[j])j=p1[j-1];
            if(s2[i]==s1[j])j++;
            p2[i]=j;
            if(j==n)return i-n+2;//返回位置
        }
    }
    return 0;
}
int main(){
    char s1[15],s2[105];
    int p1[15],p2[105];
    cin>>s1>>s2;
    kmp(s1,p1);
    cout<<kmp(s1,p1,s2,p2)<<endl;
    return 0;
}
```

regex

元字符	描述
\	将下一个字符标记符、或一个向后引用、或一个八进制转义符。例如，“\n”匹配\n。“\n”匹配换行符。序列“\”匹配“\”而“(”则匹配“(”。即相当于多种编程语言中都有的“转义字符”的概念。
^	匹配输入字行首。如果设置了 RegExp 对象的 Multiline 属性，^也匹配“\n”或“\r”之后的位置。
$	匹配输入行尾。如果设置了 RegExp 对象的 Multiline 属性，$也匹配“\n”或“\r”之前的位置。
*	匹配前面的子表达式任意次。例如，zo *能匹配“z”，也能匹配“zo”以及“zoo”*。等价于{0,}。
+	匹配前面的子表达式一次或多次(大于等于 1 次）。例如，“zo+”能匹配“zo”以及“zoo”，但不能匹配“z”。+等价于{1,}。
?	匹配前面的子表达式零次或一次。例如，“do(es)?”可以匹配“do”或“does”。?等价于{0,1}。
{*n*}	*n* 是一个非负整数。匹配确定的 *n* 次。例如，“o{2}”不能匹配“Bob”中的“o”，但是能匹配“food”中的两个 o。
{*n*,}	*n* 是一个非负整数。至少匹配 *n* 次。例如，“o{2,}”不能匹配“Bob”中的“o”，但能匹配“fooooood”中的所有 o。“o{1,}”等价于“o+”。“o{0,}”则等价于“o*”。
{*n*,*m*}	*m* 和 *n* 均为非负整数，其中 *n*<=*m*。最少匹配 *n* 次且最多匹配 *m* 次。例如，“o{1,3}”将匹配“fooooood”中的前三个 o 为一组，后三个 o 为一组。“o{0,1}”等价于“o?”。请注意在逗号和两个数之间不能有空格。
?	当该字符紧跟在任何一个其他限制符（*,+,?, {n}, {n,}, {n,m}*）后面时，匹配模式是非贪婪的。非贪婪模式尽可能少地匹配所搜索的字符串，而默认的贪婪模式则尽可能多地匹配所搜索的字符串。例如，对于字符串“oooo”，“o+”将尽可能多地匹配“o”，得到结果[“oooo”]，而“o+?”将尽可能少地匹配“o”，得到结果 ['o', 'o', 'o', 'o']
.点	匹配除“\n”和“\r”之外的任何单个字符。要匹配包括“\n”和“\r”在内的任何字符，请使用像“[\s\S]”的模式。
(pattern)	匹配 pattern 并获取这一匹配。所获取的匹配可以从产生的 Matches 集合得到，在 VBScript 中使用 SubMatches 集合，在 JScript 中则使用 *0*...9 属性。要匹配圆括号字符，请使用“(”或“)”。
(?:pattern)	非获取匹配，匹配 pattern 但不获取匹配结果，不进行存储供以后使用。这在使用或字符“(|)”来组合一个模式的各个部分时很有用。例如“industr(?:y|ies)”就是一个比“industry|industries”更简略的表达式。
(?=pattern)	非获取匹配，正向肯定预查，在任何匹配 pattern 的字符串开始处匹配查找字符串，该匹配不需要获取供以后使用。例如，“Windows(?=95|98|NT|2000)”能匹配“Windows2000”中的“Windows”，但不能匹配“Windows3.1”中的“Windows”。预查不消耗字符，也就是说，在一个匹配发生后，在最后一次匹配之后立即开始下一次匹配的搜索，而不是从包含预查的字符之后开始。
(?!pattern)	非获取匹配，正向否定预查，在任何不匹配 pattern 的字符串开始处匹配查找字符串，该匹配不需要获取供以后使用。例如“Windows(?!95|98|NT|2000)”能匹配“Windows3.1”中的“Windows”，但不能匹配“Windows2000”中的“Windows”。
(?<=pattern)	非获取匹配，反向肯定预查，与正向肯定预查类似，只是方向相反。例如，“(?<=95|98|NT|2000)Windows”能匹配“2000Windows”中的“Windows”，但不能匹配“3.1Windows”中的“Windows”。*python 的正则表达式没有完全按照正则表达式规范实现，所以一些高级特性建议使用其他语言如 java、scala 等
(?<!pattern)	非获取匹配，反向否定预查，与正向否定预查类似，只是方向相反。例如“(?<!95|98|NT|2000)Windows”能匹配“3.1Windows”中的“Windows”，但不能匹配“2000Windows”中的“Windows”。*python 的正则表达式没有完全按照正则表达式规范实现，所以一些高级特性建议使用其他语言如 java、scala 等
x|y	匹配 x 或 y。例如，“z|food”能匹配“z”或“food”(此处请谨慎)。“[z|f]ood”则匹配“zood”或“food”。
[xyz]	字符集合。匹配所包含的任意一个字符。例如，“[abc]”可以匹配“plain”中的“a”。
[^xyz]	负值字符集合。匹配未包含的任意字符。例如，“abc”可以匹配“plain”中的“plin”任一字符。
[a-z]	字符范围。匹配指定范围内的任意字符。例如，“[a-z]”可以匹配“a”到“z”范围内的任意小写字母字符。注意:只有连字符在字符组内部时,并且出现在两个字符之间时,才能表示字符的范围; 如果出字符组的开头,则只能表示连字符本身.
[^a-z]	负值字符范围。匹配任何不在指定范围内的任意字符。例如，“a-z”可以匹配任何不在“a”到“z”范围内的任意字符。
\b	匹配一个单词的边界，也就是指单词和空格间的位置（即正则表达式的“匹配”有两种概念，一种是匹配字符，一种是匹配位置，这里的\b 就是匹配位置的）。例如，“er\b”可以匹配“never”中的“er”，但不能匹配“verb”中的“er”；

元字符	描述
	“\b1”*可以匹配*“123”中的“1”，*但不能匹配*“213”中的“1_”。
\B	匹配非单词边界。“er\B”能匹配“verb”中的“er”，但不能匹配“never”中的“er”。
\cx	匹配由 x 指明的控制字符。例如，\cM 匹配一个 Control-M 或回车符。x 的值必须为 A-Z 或 a-z 之一。否则，将 c 视为一个原义的“c”字符。
\d	匹配一个数字字符。等价于[0-9]。grep 要加上-P，perl 正则支持
\D	匹配一个非数字字符。等价于 0-9。grep 要加上-P，perl 正则支持
\f	匹配一个换页符。等价于\x0c 和\cL。
\n	匹配一个换行符。等价于\x0a 和\cJ。
\r	匹配一个回车符。等价于\x0d 和\cM。
\s	匹配任何不可见字符，包括空格、制表符、换页符等等。等价于[\f\n\r\t\v]。
\S	匹配任何可见字符。等价于 \f\n\r\t\v。
\t	匹配一个制表符。等价于\x09 和\cI。
\v	匹配一个垂直制表符。等价于\x0b 和\cK。
\w	匹配包括下划线的任何单词字符。类似但不等价于“[A-Za-z0-9_]”，这里的"单词"字符使用 Unicode 字符集。
\W	匹配任何非单词字符。等价于“A-Za-z0-9_”。
\x*n*	匹配 *n*，其中 *n* 为十六进制转义值。十六进制转义值必须为确定的两个数字长。例如，“\x41”匹配“A”。“\x041”则等价于“\x04&1”。正则表达式中可以使用 ASCII 编码。
num	匹配 *num*，其中 *num* 是一个正整数。对所获取的匹配的引用。例如，“(.)\1”匹配两个连续的相同字符。
n	标识一个八进制转义值或一个向后引用。如果*n* *之前至少* n *个获取的子表达式，则* n *为向后引用。否则，如果* n *为八进制数字（0-7），则* n*为一个八进制转义值。
nm	标识一个八进制转义值或一个向后引用。如果*nm *之前至少有* nm *个获得子表达式，则* nm *为向后引用。如果*\nm *之前至少有* n *个获取，则* n *为一个后跟文字* m *的向后引用。如果前面的条件都不满足，若* n *和* m *均为八进制数字（0-7），则*\nm *将匹配八进制转义值* nm*。
nml	如果 *n* 为八进制数字（0-7），且 *m* 和 *l* 均为八进制数字（0-7），则匹配八进制转义值 *nml*。
\u*n*	匹配 *n*，其中 *n* 是一个用四个十六进制数字表示的 Unicode 字符。例如，\u00A9 匹配版权符号（©）。
\p{P}	小写 p 是 property 的意思，表示 Unicode 属性，用于 Unicode 正表达式的前缀。中括号内的“P”表示 Unicode 字符集七个字符属性之一：标点字符。其他六个属性：L：字母；M：标记符号（一般不会单独出现）；Z：分隔符（比如空格、换行等）；S：符号（比如数学符号、货币符号等）；N：数字（比如阿拉伯数字、罗马数字等）；C：其他字符。**注：此语法部分语言不支持，例：javascript*。
<>	匹配词（word）的开始（<）和结束（>）。例如正则表达式<the>能够匹配字符串"for the wise"中的"the"，但是不能匹配字符串"otherwise"中的"the"。注意：这个元字符不是所有的软件都支持的。
( )	将(和) 之间的表达式定义为“组”（group），并且将匹配这个表达式的字符保存到一个临时区域（一个正则表达式中最多可以保存 9 个），它们可以用 \1 到\9 的符号来引用。
|	将两个匹配条件进行逻辑“或”（or）运算。例如正则表达式(him|her) 匹配"it belongs to him"和"it belongs to her"，但是不能匹配"it belongs to them."。注意：这个元字符不是所有的软件都支持的。

Trie

```
#include <bits/stdc++.h>
using namespace std;
struct Trie {
    static const int maxnode = 200005;
    static const int sigma_size = 26;
    int ch[maxnode][sigma_size];
    int val[maxnode];
    int sz;

    Trie() {
        sz = 1;
        memset(ch[0], 0, sizeof(ch[0]));
    }

    int idx(const char &c) { return c - 'a'; }

    void insert(string s, int v) {
        int u = 0, n = s.length();
        for (int i = 0; i < n; i++) {
            int c = idx(s[i]);
            if (!ch[u][c]) {
                memset(ch[sz], 0, sizeof(ch[sz]));
                val[sz] = 0;
                ch[u][c] = sz++;
            }
            u = ch[u][c];
        }
        val[u] = v;
    }

    int find(string s) {
        int u = 0, n = s.length();
        for (int i = 0; i < n; i++) {
            int c = idx(s[i]);
            if (!ch[u][c]) return 0;
            u = ch[u][c];
        }
        return val[u];
    }
} trie;
```

可持久化字典树

```
struct Trie01 {
    static const int maxnode = 2000005;
    static const int sigma_size = 2;
    int ch[maxnode << 5][sigma_size], val[maxnode << 5];
    int rt[maxnode];
    int sz;

    Trie01() {
        sz = 0;
        memset(ch[0], 0, sizeof(ch[0]));
    }

    void insert(int &now, int pre, int v) {
        now = ++sz;
        for (int i = 30; i >= 0; i--) {
            int k = ((v >> i) & 1);
            ch[now][k] = ++sz;
            ch[now][k ^ 1] = ch[pre][k ^ 1];
            val[ch[now][k]] = val[ch[pre][k]] + 1;
            now = ch[now][k];
            pre = ch[pre][k];
        }
    }
} trie;
```

后缀数组

```
#include <bits/stdc++.h>
using namespace std;
struct SuffixArray {
    static const int MAXN = 1100000;
    char s[MAXN];
    int sa[MAXN], t[MAXN], t1[MAXN], c[MAXN], ra[MAXN], he
ight[MAXN], m;
    inline void init() { memset(this, 0, sizeof(SuffixArra
y)); }

    inline void get_sa(int n) {
        m = 256;
        int *x = t, *y = t1;
        for (int i = 1; i <= m; i++) c[i] = 0;
        for (int i = 1; i <= n; i++) c[x[i] = s[i]]++;
        for (int i = 1; i <= m; i++) c[i] += c[i - 1];
```

```
        for (int i = n; i >= 1; i--) sa[c[x[i]]--] = i;
        for (int k = 1; k <= n; k <<= 1) {
            int p = 0;
            for (int i = n - k + 1; i <= n; i++) y[++p] = i;
            for (int i = 1; i <= n; i++)
                if (sa[i] > k) y[++p] = sa[i] - k;
            for (int i = 1; i <= m; i++) c[i] = 0;
            for (int i = 1; i <= n; i++) c[x[y[i]]]++;
            for (int i = 1; i <= m; i++) c[i] += c[i - 1];
            for (int i = n; i >= 1; i--) sa[c[x[y[i]]]--] =
y[i];
            std::swap(x, y);
            p = x[sa[1]] = 1;
            for (int i = 2; i <= n; i++) {
                x[sa[i]] = (y[sa[i - 1]] == y[sa[i]] &&
                            y[sa[i - 1] + k] == y[sa[i] + k])
                               ? p
                               : ++p;
            }
            if (p >= n) break;
            m = p;
        }
    }

    inline void get_height(int n) {
        int i, j, k = 0;
        for (int i = 1; i <= n; i++) ra[sa[i]] = i;
        for (int i = 1; i <= n; i++) {
            if (k) k--;
            int j = sa[ra[i] - 1];
            while (s[i + k] == s[j + k]) k++;
            height[ra[i]] = k;
        }
    }

} SA;   //字符串下标从一开始
```

后缀自动机

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const int N = 2e6 + 10;

int tot = 1, last = 1;
struct Node {
    int len, fa;
    int ch[26];
} node[N];
char str[N];
ll f[N], ans;
int h[N], e[N], ne[N], idx;

void extend(int c) {
    int p = last, np = last = ++tot;
    f[tot] = 1;
    node[np].len = node[p].len + 1;
    for (; p && !node[p].ch[c]; p = node[p].fa) node[p].ch
[c] = np;
    if (!p) node[np].fa = 1;
    else {
        int q = node[p].ch[c];
        if (node[q].len == node[p].len + 1) node[np].fa =
q;
        else {
            int nq = ++tot;
            node[nq] = node[q], node[nq].len = node[p].len
+ 1;
            node[q].fa = node[np].fa = nq;
            for (; p && node[p].ch[c] == q; p = node[p].fa)
 node[p].ch[c] = nq;
        }
    }
}

void add(int a, int b) {
    e[idx] = b, ne[idx] = h[a], h[a] = idx++;
}

void dfs(int u) {
    for (int i = h[u]; ~i; i = ne[i]) {
        dfs(e[i]);
        f[u] += f[e[i]];
    }
    if (f[u] > 1) ans = max(ans, f[u] * node[u].len);
}

int main() {
    scanf("%s", str);
    for (int i = 0; str[i]; i++) extend(str[i] - 'a');
    memset(h, -1, sizeof h);
    for (int i = 2; i <= tot; i++) add(node[i].fa, i);
    dfs(1);
    printf("%lld\n", ans);

    return 0;
}
```

马拉车

```
#include <bits/stdc++.h>
using namespace std;
const int maxn = 100005;
char s[maxn];
char s_new[maxn * 2];
int p[maxn * 2];

int Manacher(char* a, int l) {
    s_new[0] = '$';
    s_new[1] = '#';
    int len = 2;
    for (int i = 0; i < l; i++) {
        s_new[len++] = a[i];
        s_new[len++] = '#';
    }
    s_new[len] = '\0';
    int id;
    int mx = 0;
    int mmax = 0;

    for (int i = 1; i < len; i++) {
        p[i] = i < mx ? min(p[2 * id - i], mx - i) : 1;
        while (s_new[i + p[i]] == s_new[i - p[i]]) p[i]++;
        if (mx < i + p[i]) {
            id = i;
            mx = i + p[i];
        }
        mmax = max(mmax, p[i] - 1);
    }
    return mmax;
}

int main() {
    cin >> s;
    cout << Manacher(s, strlen(s));
}
```

搜索

占位

数据结构

CDQ 分治

```
/*
处理三维偏序问题，
每个node 的三维不能完全相等，完全相等的话加权做
*/

#include <iostream>
#include <cstring>
#include <algorithm>

using namespace std;

const int N = 100010, M = 200010;

int n, m;
```

```
struct Data
{
    int a, b, c, s, res;

    bool operator< (const Data& t) const
    {
        if (a != t.a) return a < t.a;
        if (b != t.b) return b < t.b;
        return c < t.c;
    }
    bool operator== (const Data& t) const
    {
        return a == t.a && b == t.b && c == t.c;
    }
}q[N], w[N];
int tr[M], ans[N];

int lowbit(int x)
{
    return x & -x;
}

void add(int x, int v)
{
    for (int i = x; i < M; i += lowbit(i)) tr[i] += v;
}

int query(int x)
{
    int res = 0;
    for (int i = x; i; i -= lowbit(i)) res += tr[i];
    return res;
}

void merge_sort(int l, int r)
{
    if (l >= r) return;
    int mid = l + r >> 1;
    merge_sort(l, mid), merge_sort(mid + 1, r);
    int i = l, j = mid + 1, k = 0;
    while (i <= mid && j <= r)
        if (q[i].b <= q[j].b) add(q[i].c, q[i].s), w[k ++ ]
 = q[i ++ ];
        else q[j].res += query(q[j].c), w[k ++ ] = q[j ++ ];
    while (i <= mid) add(q[i].c, q[i].s), w[k ++ ] = q[i +
+ ];
    while (j <= r) q[j].res += query(q[j].c), w[k ++ ] = q
[j ++ ];
    for (i = l; i <= mid; i ++ ) add(q[i].c, -q[i].s);
    for (i = l, j = 0; j < k; i ++, j ++ ) q[i] = w[j];
}

int main()
{
    scanf("%d%d", &n, &m);
    for (int i = 0; i < n; i ++ )
    {
        int a, b, c;
        scanf("%d%d%d", &a, &b, &c);
        q[i] = {a, b, c, 1};
    }
    sort(q, q + n);

    int k = 1;
    for (int i = 1; i < n; i ++ )
        if (q[i] == q[k - 1]) q[k - 1].s ++ ;
        else q[k ++ ] = q[i];

    merge_sort(0, k - 1);
    for (int i = 0; i < k; i ++ )
        ans[q[i].res + q[i].s - 1] += q[i].s;

    for (int i = 0; i < n; i ++ ) printf("%d\n", ans[i]);

    return 0;
}
```

kruskal 重构树

```
int pa[N];

void init(int n) {
    for (int i = 0; i <= n; i++) {
        pa[i] = i;
    }
}

int find(int a) {
    return pa[a] == a ? a : pa[a] = find(pa[a]);
}

struct edge {
    int from, to, l;
};

int w[N];
edge e[N];
vector<int> g[N];

int kruskal(int n, int m) {
    int kcnt = n;
    init(n);
    sort(e + 1, e + 1 + m, [](edge a, edge b) { return a.l
< b.l; });
    for (int i = 1; i <= m; i++) {
        int u = find(e[i].from);
        int v = find(e[i].to);
        if (u == v) continue;
        w[++kcnt] = e[i].l;
        pa[kcnt] = pa[u] = pa[v] = kcnt;
        g[u].push_back(kcnt);
        g[v].push_back(kcnt);
        g[kcnt].push_back(u);
        g[kcnt].push_back(v);
    }
    return kcnt;
}
```

LCT

```
ll ch[N][2], f[N], sum[N], val[N], tag[N], siz[N], siz2[N];

inline void pushup(ll p) {
    sum[p] = sum[ch[p][0]] ^ sum[ch[p][1]] ^ val[p];
    siz[p] = siz[ch[p][0]] + siz[ch[p][1]] + 1 + siz2[p];
}

inline void pushdown(ll p) {
    if (tag[p]) {
        if (ch[p][0]) swap(ch[ch[p][0]][0], ch[ch[p][0]]
[1]), tag[ch[p][0]] ^= 1;
        if (ch[p][1]) swap(ch[ch[p][1]][0], ch[ch[p][1]]
[1]), tag[ch[p][1]] ^= 1;
        tag[p] = 0;
    }
}

ll getch(ll x) { return ch[f[x]][1] == x; }

bool isroot(ll x) { return ch[f[x]][0] != x && ch[f[x]][1]
 != x; }

inline void rotate(ll x) {
    ll y = f[x], z = f[y], k = getch(x);
    if (!isroot(y)) ch[z][ch[z][1] == y] = x;
    // 上面这句一定要写在前面，普通的 Splay 是不用的，因为 isRo
ot （后面会讲）
    ch[y][k] = ch[x][!k], f[ch[x][!k]] = y;
    ch[x][!k] = y, f[y] = x, f[x] = z;
    pushup(y), pushup(x);
}

// 从上到下一层一层 pushDown 即可
void update(ll p) {
    if (!isroot(p)) update(f[p]);
    pushdown(p);
}

inline void splay(ll x) {
    update(x);  // 马上就能看到啦。 在
    // Splay 之前要把旋转会经过的路径上的点都 PushDown
    for (ll fa; fa = f[x], !isroot(x); rotate(x)) {
        if (!isroot(fa)) rotate(getch(fa) == getch(x) ? fa
 : x);
    }
}
```

```
// 回顾一下代码
inline void access(ll x) {
    for (ll p = 0; x; p = x, x = f[x]) {
        splay(x), siz2[x] += siz[ch[x][1]] - siz[p], ch[x]
[1] = p, pushup(x);
    }
}

inline void makeroot(ll p) {
    access(p);
    splay(p);
    swap(ch[p][0], ch[p][1]);
    tag[p] ^= 1;
}

inline void split(ll a, ll b) {
    makeroot(a);
    access(b);
    splay(b);
}


inline ll find(ll p) {
    access(p), splay(p);
    while (ch[p][0]) pushdown(p), p = ch[p][0];
    splay(p);
    return p;
}

inline void link(ll x, ll y) {
    makeroot(y);
    makeroot(x);
    if (find(y) != x) {
        f[x] = y;
        siz2[y] += siz[x];
    }
}

inline void cut(ll x, ll y) {
    makeroot(x);
    if (find(y) == x && f[y] == x) {
        ch[x][1] = f[y] = 0;
        pushup(x);
    }
}

void init(int n) {
    for (int i = 1; i <= n; i++) siz[i] = 1;
}
```

Splay

```
ll ch[N][2], f[N], sum[N], val[N], tag[N], siz[N];

inline void pushup(ll p) {
    sum[p] = sum[ch[p][0]] ^ sum[ch[p][1]] ^ val[p];
    siz[p] = siz[ch[p][0]] + siz[ch[p][1]] + 1;
}

inline void pushdown(ll p) {
    if (tag[p]) {
        if (ch[p][0]) swap(ch[ch[p][0]][0], ch[ch[p][0]]
[1]), tag[ch[p][0]] ^= 1;
        if (ch[p][1]) swap(ch[ch[p][1]][0], ch[ch[p][1]]
[1]), tag[ch[p][1]] ^= 1;
        tag[p] = 0;
    }
}

ll getch(ll x) { return ch[f[x]][1] == x; }

bool isroot(ll x) { return ch[f[x]][0] != x && ch[f[x]][1]
 != x; }

inline void rotate(ll x) {
    ll y = f[x], z = f[y], k = getch(x);
    if (!isroot(y)) ch[z][ch[z][1] == y] = x;
    // 上面这句一定要写在前面，普通的Splay是不用的，因为 isRo
ot  (后面会讲)
    ch[y][k] = ch[x][!k], f[ch[x][!k]] = y;
    ch[x][!k] = y, f[y] = x, f[x] = z;
    pushup(y), pushup(x);
}

// 从上到下一层一层 pushDown 即可
void update(ll p) {
    if (!isroot(p)) update(f[p]);
    pushdown(p);
}

inline void splay(ll x) {
    update(x);  // 马上就能看到啦。 在
    // Splay 之前要把旋转会经过的路径上的点都 PushDown
    for (ll fa; fa = f[x], !isroot(x); rotate(x)) {
        if (!isroot(fa)) rotate(getch(fa) == getch(x) ? fa
 : x);
    }
}

// 回顾一下代码
inline void access(ll x) {
    for (ll p = 0; x; p = x, x = f[x]) {
        splay(x), ch[x][1] = p, pushup(x);
    }
}

inline void makeroot(ll p) {
    access(p);
    splay(p);
    swap(ch[p][0], ch[p][1]);
    tag[p] ^= 1;
}

inline void split(ll a, ll b) {
    makeroot(a);
    access(b);
    splay(b);
}


inline ll find(ll p) {
    access(p), splay(p);
    while (ch[p][0]) pushdown(p), p = ch[p][0];
    splay(p);
    return p;
}

inline void link(ll x, ll y) {
    makeroot(x);
    if (find(y) != x) f[x] = y;
}

inline void cut(ll x, ll y) {
    makeroot(x);
    if (find(y) == x && f[y] == x) {
        ch[x][1] = f[y] = 0;
        pushup(x);
    }
}
```

ST 表

```
#include <bits/stdc++.h>

using namespace std;
const int logn = 21;
const int N = 2000001;
int f[N][logn + 1], lg[N + 1];

void pre() {
    lg[1] = 0;
    for (int i = 2; i < N; i++) {
        lg[i] = lg[i / 2] + 1;
    }
}

int main() {
    ios::sync_with_stdio(false);
    int n, m;
    cin >> n >> m;
    for (int i = 1; i <= n; i++) cin >> f[i][0];
    pre();
    for (int j = 1; j <= logn; j++)
        for (int i = 1; i + (1 << j) - 1 <= n; i++)
            f[i][j] = max(f[i][j - 1], f[i + (1 << (j - 1))]
```

```
[j - 1]);
    for (int i = 1; i <= m; i++) {
        int x, y;
        cin >> x >> y;
        int s = lg[y - x + 1];
        printf("%d\n", max(f[x][s], f[y - (1 << s) + 1]
[s]));
    }
    return 0;
}
```

Treap

```
#include <bits/stdc++.h>
using namespace std;
struct node {
    node* ch[2];
    int r;
    int v;
    int cmp(int const& a) const {
        if (v == a) return -a;
        return a > v ? 1 : 0;
    }
};
void rotate(node*& a, int d) {
    node* k = a->ch[d ^ 1];
    a->ch[d ^ 1] = k->ch[d];
    k->ch[d] = a;
    a = k;
}
void insert(node*& a, int x) {
    if (a == NULL) {
        a = new node;
        a->ch[0] = a->ch[1] = NULL;
        a->v = x;
        a->r = rand();
    } else {
        int d = a->cmp(x);
        insert(a->ch[d], x);
        if (a->ch[d]->r > a->r) rotate(a, d ^ 1);
    }
}
void remove(node*& a, int x) {
    int d = a->cmp(x);
    if (d == -1) {
        if (a->ch[0] == NULL)
            a = a->ch[1];
        else if (a->ch[1] == NULL)
            a = a->ch[0];
        else {
            int d2 = a->ch[1]->r > a->ch[0]->r ? 0 : 1;
            rotate(a, d2);
            remove(a->ch[d2], x);
        }
    } else {
        remove(a->ch[d], x);
    }
}
int find(node*& a, int x) {
    if (a == NULL)
        return 0;
    else if (a->v == x)
        return 1;
    else {
        int d = a->cmp(x);
        return find(a->ch[d], x);
    }
}
int main() {
    node* a = NULL;
    int k, l;
    while (cin >> k >> l) {
        if (k == 1)
            insert(a, l);
        else if (k == 2)
            remove(a, l);
        else {
            cout << find(a, l) << endl;
        }
    }
}
```

y 🤖 Splay Plus

```
#include <iostream>
#include <cstdio>
#include <cstring>
#include <algorithm>

using namespace std;

const int N = 500010, INF = 1e9;

int n, m;
struct Node
{
    int s[2], p, v;
    int rev, same;
    int size, sum, ms, ls, rs;

    void init(int _v, int _p)
    {
        s[0] = s[1] = 0, p = _p, v = _v;
        rev = same = 0;
        size = 1, sum = ms = v;
        ls = rs = max(v, 0);
    }
}tr[N];
int root, nodes[N], tt;
int w[N];

void pushup(int x)
{
    auto &u = tr[x], &l = tr[u.s[0]], &r = tr[u.s[1]];
    u.size = l.size + r.size + 1;
    u.sum = l.sum + r.sum + u.v;
    u.ls = max(l.ls, l.sum + u.v + r.ls);
    u.rs = max(r.rs, r.sum + u.v + l.rs);
    u.ms = max(max(l.ms, r.ms), l.rs + u.v + r.ls);
}

void pushdown(int x)
{
    auto &u = tr[x], &l = tr[u.s[0]], &r = tr[u.s[1]];
    if (u.same)
    {
        u.same = u.rev = 0;
        if (u.s[0]) l.same = 1, l.v = u.v, l.sum = l.v * l.
size;
        if (u.s[1]) r.same = 1, r.v = u.v, r.sum = r.v * r.
size;
        if (u.v > 0)
        {
            if (u.s[0]) l.ms = l.ls = l.rs = l.sum;
            if (u.s[1]) r.ms = r.ls = r.rs = r.sum;
        }
        else
        {
            if (u.s[0]) l.ms = l.v, l.ls = l.rs = 0;
            if (u.s[1]) r.ms = r.v, r.ls = r.rs = 0;
        }
    }
    else if (u.rev)
    {
        u.rev = 0, l.rev ^= 1, r.rev ^= 1;
        swap(l.ls, l.rs), swap(r.ls, r.rs);
        swap(l.s[0], l.s[1]), swap(r.s[0], r.s[1]);
    }
}

void rotate(int x)
{
    int y = tr[x].p, z = tr[y].p;
    int k = tr[y].s[1] == x;
    tr[z].s[tr[z].s[1] == y] = x, tr[x].p = z;
    tr[y].s[k] = tr[x].s[k ^ 1], tr[tr[x].s[k ^ 1]].p = y;
    tr[x].s[k ^ 1] = y, tr[y].p = x;
    pushup(y), pushup(x);
}

void splay(int x, int k)
{
    while (tr[x].p != k)
    {
        int y = tr[x].p, z = tr[y].p;
        if (z != k)
            if ((tr[y].s[1] == x) ^ (tr[z].s[1] == y)) rota
```

```
te(x);
            else rotate(y);
        rotate(x);
    }
    if (!k) root = x;
}

int get_k(int k)
{
    int u = root;
    while (u)
    {
        pushdown(u);
        if (tr[tr[u].s[0]].size >= k) u = tr[u].s[0];
        else if (tr[tr[u].s[0]].size + 1 == k) return u;
        else k -= tr[tr[u].s[0]].size + 1, u = tr[u].s[1];
    }
}

int build(int l, int r, int p)
{
    int mid = l + r >> 1;
    int u = nodes[tt -- ];
    tr[u].init(w[mid], p);
    if (l < mid) tr[u].s[0] = build(l, mid - 1, u);
    if (mid < r) tr[u].s[1] = build(mid + 1, r, u);
    pushup(u);
    return u;
}

void dfs(int u)
{
    if (tr[u].s[0]) dfs(tr[u].s[0]);
    if (tr[u].s[1]) dfs(tr[u].s[1]);
    nodes[ ++ tt] = u;
}

int main()
{
    for (int i = 1; i < N; i ++ ) nodes[ ++ tt] = i;
    scanf("%d%d", &n, &m);
    tr[0].ms = w[0] = w[n + 1] = -INF;
    for (int i = 1; i <= n; i ++ ) scanf("%d", &w[i]);
    root = build(0, n + 1, 0);

    char op[20];
    while (m -- )
    {
        scanf("%s", op);
        if (!strcmp(op, "INSERT"))
        {
            int posi, tot;
            scanf("%d%d", &posi, &tot);
            for (int i = 0; i < tot; i ++ ) scanf("%d", &w
[i]);
            int l = get_k(posi + 1), r = get_k(posi + 2);
            splay(l, 0), splay(r, l);
            int u = build(0, tot - 1, r);
            tr[r].s[0] = u;
            pushup(r), pushup(l);
        }
        else if (!strcmp(op, "DELETE"))
        {
            int posi, tot;
            scanf("%d%d", &posi, &tot);
            int l = get_k(posi), r = get_k(posi + tot + 1);
            splay(l, 0), splay(r, l);
            dfs(tr[r].s[0]);
            tr[r].s[0] = 0;
            pushup(r), pushup(l);
        }
        else if (!strcmp(op, "MAKE-SAME"))
        {
            int posi, tot, c;
            scanf("%d%d%d", &posi, &tot, &c);
            int l = get_k(posi), r = get_k(posi + tot + 1);
            splay(l, 0), splay(r, l);
            auto& son = tr[tr[r].s[0]];
            son.same = 1, son.v = c, son.sum = c * son.size;
            if (c > 0) son.ms = son.ls = son.rs = son.sum;
            else son.ms = c, son.ls = son.rs = 0;
            pushup(r), pushup(l);
        }
        else if (!strcmp(op, "REVERSE"))
        {
            int posi, tot;
            scanf("%d%d", &posi, &tot);
            int l = get_k(posi), r = get_k(posi + tot + 1);
            splay(l, 0), splay(r, l);
            auto& son = tr[tr[r].s[0]];
            son.rev ^= 1;
            swap(son.ls, son.rs);
            swap(son.s[0], son.s[1]);
            pushup(r), pushup(l);
        }
        else if (!strcmp(op, "GET-SUM"))
        {
            int posi, tot;
            scanf("%d%d", &posi, &tot);
            int l = get_k(posi), r = get_k(posi + tot + 1);
            splay(l, 0), splay(r, l);
            printf("%d\n", tr[tr[r].s[0]].sum);
        }
        else printf("%d\n", tr[root].ms);
    }

    return 0;
}
```

y 🤖 Splay

```
#include <bits/stdc++.h>

using namespace std;
const int N = 1e6 + 10;
struct node {
    int p, v, s[2];
    int siz, tag;
    void init(int _v, int _p) {
        v = _v, p = _p;
        siz = 1;
    }
};
node tr[N];
int root, idx;

void pushup(int x) { tr[x].siz = tr[tr[x].s[0]].siz + tr[t
r[x].s[1]].siz + 1; }

void pushdown(int x) {
    if (tr[x].tag) {
        swap(tr[x].s[0], tr[x].s[1]);
        tr[tr[x].s[0]].tag ^= 1;
        tr[tr[x].s[1]].tag ^= 1;
        tr[x].tag = 0;
    }
}
void rotate(int x) {
    int y = tr[x].p, z = tr[y].p;
    int k = tr[y].s[1] == x;
    tr[y].s[k] = tr[x].s[k ^ 1], tr[tr[y].s[k]].p = y;
    tr[x].s[k ^ 1] = y, tr[y].p = x;
    tr[z].s[tr[z].s[1] == y] = x, tr[x].p = z;
    pushup(y), pushup(x);
}

void splay(int x, int k) {
    while (tr[x].p != k) {
        int y = tr[x].p, z = tr[y].p;
        if (z != k) {
            if ((tr[z].s[1] == y) ^ (tr[y].s[1] == x)) {
                rotate(x);
            } else {
                rotate(y);
            }
        }
        rotate(x);
    }
    if (!k) root = x;
}

void insert(int v) {
    int u = root, p = 0;
    while (u) p = u, u = tr[u].s[v > tr[u].v];
    u = ++idx;
    if (p) tr[p].s[v > tr[p].v] = u;
    tr[u].init(v, p);
    splay(u, 0);
```

```
}

int getk(int k) {
    int u = root;
    while (1) {
        pushdown(u);
        if (k <= tr[tr[u].s[0]].siz) {
            u = tr[u].s[0];
        } else if (k == tr[tr[u].s[0]].siz + 1) {
            splay(u, 0);
            return u;
        } else {
            k -= tr[tr[u].s[0]].siz + 1, u = tr[u].s[1];
        }
    }
}

int n, m;
void output(int u) {
    if (u == 0) return;
    pushdown(u);
    output(tr[u].s[0]);
    if (1 <= tr[u].v && tr[u].v <= n) cout << tr[u].v << '
';
    output(tr[u].s[1]);
}

int main() {
    ios::sync_with_stdio(0), cin.tie(0), cout.tie(0);
    cin >> n >> m;
    for (int i = 0; i <= n + 1; i++) insert(i);
    while (m--) {
        int a, b;
        cin >> a >> b;
        int id1 = getk(a), id2 = getk(b + 2);
        splay(id1, 0), splay(id2, id1);
        tr[tr[id2].s[0]].tag ^= 1;
    }
    output(root);
}
```

主席树

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const ll N = 1 << 20;

ll ch[N << 5][2], rt[N], tot;
ll val[N << 5];

ll update(ll a, ll b) {
    return a + b;
}

ll build(ll l, ll r) {  // 建树
    ll p = ++tot;
    if (l == r) {
        //初始化
        val[p] = 0;
        return p;
    }
    ll mid = (l + r) >> 1;
    ch[p][0] = build(l, mid);
    ch[p][1] = build(mid + 1, r);
    val[p] = update(val[ch[p][0]], val[ch[p][1]]);
    return p;  // 返回该子树的根节点
}

ll modify(ll pre, ll l, ll r, ll pos, ll v) {  // 插入操作
    ll now = ++tot;
    ch[now][0] = ch[pre][0], ch[now][1] = ch[pre][1];
    if (l == r) {
        val[now] = val[pre] + v;
        return now;
    }
    ll mid = (l + r) >> 1;
    if (pos <= mid)
        ch[now][0] = modify(ch[now][0], l, mid, pos, v);
    else
        ch[now][1] = modify(ch[now][1], mid + 1, r, pos,
v);
    val[now] = update(val[ch[now][0]], val[ch[now][1]]);
    return now;
}

ll kth(ll pre, ll now, ll l, ll r, ll k) {  // 查询操作
    ll mid = (l + r) >> 1;
    ll x = val[ch[now][0]] - val[ch[pre][0]];  // 通过区间
减法得到左儿子的信息
    if (l == r) return l;
    if (k <= x)  // 说明在左儿子中
        return kth(ch[pre][0], ch[now][0], l, mid, k);
    else  // 说明在右儿子中
        return kth(ch[pre][1], ch[now][1], mid + 1, r, k -
x);
}

ll query(ll pre, ll now, ll l, ll r, ll ql, ll qr) {  // 查
询操作
    if (ql <= l && r <= qr) {
        return val[now] - val[pre];
    }
    if (qr < l || r < ql) {
        return 0;
    }
    ll mid = (l + r) >> 1;
    ll lv = query(ch[pre][0], ch[now][0], l, mid, ql, qr);
    ll rv = query(ch[pre][1], ch[now][1], mid + 1, r, ql,
qr);
    return update(lv, rv);
}
//修改查询记得用rt[]!!!
```

仙人掌

```
/*
 仙人掌:任意一条边至多只出现在一条简单回路的无向连通图称为仙人
掌。
 转化为圆方树，然后根据树的算法来做一些问题，注意区分圆点和方点
 这题:求带环（环和环之间无公共边）无向图两点间的最短路径
 */

#include <iostream>
#include <cstring>
#include <algorithm>

using namespace std;

const int N = 12010, M = N * 3;

int n, m, Q, new_n;
int h1[N], h2[N], e[M], w[M], ne[M], idx;
int dfn[N], low[N], cnt;
int s[N], stot[N], fu[N], fw[N];
int fa[N][14], depth[N], d[N];
int A, B;

void add(int h[], int a, int b, int c)
{
    e[idx] = b, w[idx] = c, ne[idx] = h[a], h[a] = idx ++ ;
}

void build_circle(int x, int y, int z)
{
    int sum = z;
    for (int k = y; k != x; k = fu[k])
    {
        s[k] = sum;
        sum += fw[k];
    }
    s[x] = stot[x] = sum;
    add(h2, x, ++ new_n, 0);
    for (int k = y; k != x; k = fu[k])
    {
        stot[k] = sum;
        add(h2, new_n, k, min(s[k], sum - s[k]));
    }
}

void tarjan(int u, int from)
{
    dfn[u] = low[u] = ++ cnt;
    for (int i = h1[u]; ~i; i = ne[i])
```

```
    {
        int j = e[i];
        if (!dfn[j])
        {
            fu[j] = u, fw[j] = w[i];
            tarjan(j, i);
            low[u] = min(low[u], low[j]);
            if (dfn[u] < low[j]) add(h2, u, j, w[i]);
        }
        else if (i != (from ^ 1)) low[u] = min(low[u], dfn
[j]);
    }
    for (int i = h1[u]; ~i; i = ne[i])
    {
        int j = e[i];
        if (dfn[u] < dfn[j] && fu[j] != u)
            build_circle(u, j, w[i]);
    }
}

void dfs_lca(int u, int father)
{
    depth[u] = depth[father] + 1;
    fa[u][0] = father;
    for (int k = 1; k <= 13; k ++ )
        fa[u][k] = fa[fa[u][k - 1]][k - 1];
    for (int i = h2[u]; ~i; i = ne[i])
    {
        int j = e[i];
        d[j] = d[u] + w[i];
        dfs_lca(j, u);
    }
}

int lca(int a, int b)
{
    if (depth[a] < depth[b]) swap(a, b);
    for (int k = 13; k >= 0; k -- )
        if (depth[fa[a][k]] >= depth[b])
            a = fa[a][k];
    if (a == b) return a;
    for (int k = 13; k >= 0; k -- )
        if (fa[a][k] != fa[b][k])
        {
            a = fa[a][k];
            b = fa[b][k];
        }
    A = a, B = b;
    return fa[a][0];
}

int main()
{
    scanf("%d%d%d", &n, &m, &Q);
    new_n = n;
    memset(h1, -1, sizeof h1);
    memset(h2, -1, sizeof h2);
    while (m -- )
    {
        int a, b, c;
        scanf("%d%d%d", &a, &b, &c);
        add(h1, a, b, c), add(h1, b, a, c);
    }
    tarjan(1, -1);
    dfs_lca(1, 0);

    while (Q -- )
    {
        int a, b;
        scanf("%d%d", &a, &b);
        int p = lca(a, b);
        if (p <= n) printf("%d\n", d[a] + d[b] - d[p] * 2);
        else
        {
            int da = d[a] - d[A], db = d[b] - d[B];
            int l = abs(s[A] - s[B]);
            int dm = min(l, stot[A] - l);
            printf("%d\n", da + dm + db);
        }
    }

    return 0;
}
```

区间 max

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const int N = 1 << 20;

struct node {
    int mmax, semax, cnt;
    ll sum;
};

node tree[N << 1];
int init[N << 1];

node merge_range(node a, node b) {
    node ans;
    ans.sum = a.sum + b.sum;
    if (a.mmax == b.mmax) {
        ans.mmax = a.mmax;
        ans.cnt = a.cnt + b.cnt;
        ans.semax = max(a.semax, b.semax);
    } else {
        if (a.mmax < b.mmax) swap(a, b);
        ans.mmax = a.mmax;
        ans.cnt = a.cnt;
        ans.semax = max(a.semax, b.mmax);
    }
    return ans;
}

void build(int k, int l, int r) {
    if (l == r) {
        tree[k] = {init[l], -1, 1, init[l]};
        return;
    }
    int mid = (l + r) >> 1;
    build(k << 1, l, mid);
    build(k << 1 | 1, mid + 1, r);
    tree[k] = merge_range(tree[k << 1], tree[k << 1 | 1]);
}


void pushdown(int k, int l, int r) {
    if (l == r) return;
    if (tree[k].mmax < tree[k << 1].mmax) {
        tree[k << 1].sum -= 1LL * (tree[k << 1].mmax - tree
[k].mmax) * tree[k << 1].cnt;
        tree[k << 1].mmax = tree[k].mmax;
    }
    if (tree[k].mmax < tree[k << 1 | 1].mmax) {
        tree[k << 1 | 1].sum -= 1LL * (tree[k << 1 | 1].mma
x - tree[k].mmax) * tree[k << 1 | 1].cnt;
        tree[k << 1 | 1].mmax = tree[k].mmax;
    }
}


node query(int k, int l, int r, int ql, int qr) {
    if (qr < l || r < ql) return {0, -1, 1, 0};
    if (ql <= l && r <= qr) {
        return tree[k];
    }
    pushdown(k, l, r);
    int mid = (l + r) >> 1;
    node lq = query(k << 1, l, mid, ql, qr);
    node rq = query(k << 1 | 1, mid + 1, r, ql, qr);
    return merge_range(lq, rq);
}

void modify(int k, int l, int r, int ql, int qr, int x) {
    if (qr < l || r < ql) return;
    if (ql <= l && r <= qr && tree[k].semax < x) {
        if (x < tree[k].mmax) {
            tree[k].sum -= 1LL * (tree[k].mmax - x) * tree
[k].cnt;
            tree[k].mmax = x;
        }
        return;
    }
    pushdown(k, l, r);
    int mid = (l + r) >> 1;
    modify(k << 1, l, mid, ql, qr, x);
    modify(k << 1 | 1, mid + 1, r, ql, qr, x);
```

```
    tree[k] = merge_range(tree[k << 1], tree[k << 1 | 1]);
}


signed main() {
//    freopen("data.txt", "r", stdin);
//    freopen("test1.txt", "w", stdout);
    int t;
    scanf("%d", &t);
    while (t--) {
        int n, q;
        scanf("%d%d", &n, &q);
        for (int i = 1; i <= n; i++) scanf("%d", &init[i]);
        build(1, 1, n);
        while (q--) {
            int x, y, op, val;
            scanf("%d%d%d", &op, &x, &y);
            if (op == 0) {
                scanf("%d", &val);
                modify(1, 1, n, x, y, val);
            } else if (op == 1) {
                node ans = query(1, 1, n, x, y);
                printf("%d\n", ans.mmax);
            } else {
                node ans = query(1, 1, n, x, y);
                printf("%lld\n", ans.sum);
            }
        }
    }
}
```

回滚莫队

```
/*
离线，询问按左端点升序为第一关键字，右端点升序为第二关键字
对于都在块内的点直接暴力，否则跨块：
若当前左端点所属的块与上一个不同，则将左端点初始为当前块的右端点+1，右端点初始为当前块的右端点
左端点每次暴力，右端点单调
*/

#include <iostream>
#include <cstring>
#include <cstdio>
#include <algorithm>
#include <cmath>
#include <vector>

using namespace std;

typedef long long LL;
const int N = 100010;

int n, m, len;
int w[N], cnt[N];
LL ans[N];
struct Query
{
    int id, l, r;
}q[N];
vector<int> nums;

int get(int x)
{
    return x / len;
}

bool cmp(const Query& a, const Query& b)
{
    int i = get(a.l), j = get(b.l);
    if (i != j) return i < j;
    return a.r < b.r;
}

void add(int x, LL& res)
{
    cnt[x] ++ ;
    res = max(res, (LL)cnt[x] * nums[x]);
}

int main()
{
    scanf("%d%d", &n, &m);
    len = sqrt(n);
    for (int i = 1; i <= n; i ++ ) scanf("%d", &w[i]), nums.
push_back(w[i]);
    sort(nums.begin(), nums.end());
    nums.erase(unique(nums.begin(), nums.end()), nums.end
());
    for (int i = 1; i <= n; i ++ )
        w[i] = lower_bound(nums.begin(), nums.end(), w[i])
 - nums.begin();

    for (int i = 0; i < m; i ++ )
    {
        int l, r;
        scanf("%d%d", &l, &r);
        q[i] = {i, l, r};
    }
    sort(q, q + m, cmp);

    for (int x = 0; x < m;)
    {
        int y = x;
        while (y < m && get(q[y].l) == get(q[x].l)) y ++ ;
        int right = get(q[x].l) * len + len - 1;

        // 暴力求块内的询问
        while (x < y && q[x].r <= right)
        {
            LL res = 0;
            int id = q[x].id, l = q[x].l, r = q[x].r;
            for (int k = l; k <= r; k ++ ) add(w[k], res);
            ans[id] = res;
            for (int k = l; k <= r; k ++ ) cnt[w[k]] -- ;
            x ++ ;
        }

        // 求块外的询问
        LL res = 0;
        int i = right, j = right + 1;
        while (x < y)
        {
            int id = q[x].id, l = q[x].l, r = q[x].r;
            while (i < r) add(w[ ++ i], res);
            LL backup = res;
            while (j > l) add(w[ -- j], res);
            ans[id] = res;
            while (j < right + 1) cnt[w[j ++ ]] -- ;
            res = backup;
            x ++ ;
        }
        memset(cnt, 0, sizeof cnt);
    }

    for (int i = 0; i < m; i ++ ) printf("%lld\n", ans[i]);
    return 0;
}
```

带修莫队

```
#include <bits/stdc++.h>
using namespace std;

const int N = 10010;

int a[N], cnt[1000010], ans[N];

int len, mq, mc;

struct Query {
        int id, l, r, t;
} q[N];

struct Modify {
        int p, c;
} c[N];

int getNum(int x) {
        return x / len;
}

// l 所在块的编号，r 所在块的编号，t 升序

bool cmp(const Query& a, const Query& b) {
        if(getNum(a.l) == getNum(b.l) && getNum(a.r) ==
```

```
 getNum(b.r)) {
                    return a.t < b.t;
          }
          if(getNum(a.l) == getNum(b.l)) return a.r < b.r;
          return a.l < b.l;
}

void add(int x, int& res) {
    if (!cnt[x]) res ++ ;
    cnt[x] ++ ;
}

void del(int x, int& res) {
    cnt[x] -- ;
    if (!cnt[x]) res -- ;
}


int main() {
          ios::sync_with_stdio(0); cin.tie(0); cout.tie
(0);

          int n, m;
          cin >> n >> m;
          char op;
          int x, y;
          for(int i = 1; i <= n; ++ i) {
                    cin >> a[i];
          }
          for(int i = 1; i <= m; ++ i) {
                    cin >> op >> x >> y;
       if (op == 'Q') q[++ mq] = {mq, x, y, mc};
       else c[ ++ mc] = {x, y};
          }

  ///
          len = cbrt((double)n * mc) + 1;
  sort(q + 1, q + mq + 1, cmp);

          int i = 1, j = 0, t = 0, res = 0;
          for(int k = 1; k <= mq; ++ k) {
                    int id = q[k].id, l = q[k].l, r = q[k].
r, tm = q[k].t;
                    while(j < r) add(a[++ j], res);
                    while(j > r) del(a[j --], res);
                    while(i < l) del(a[i ++], res);
                    while(i > l) add(a[-- i], res);
                    while(t < tm) {
                              ++ t;
                              if(c[t].p >= i && c[t].p <=
 j) {
                                        del(a[c[t].p], re
s);
                                        add(c[t].c, res);
                              }
                              swap(a[c[t].p], c[t].c);
                    }
                    while(t > tm) {
                              if(c[t].p >= i && c[t].p <=
 j) {
                                        del(a[c[t].p], re
s);
                                        add(c[t].c, res);
                              }
                              swap(a[c[t].p], c[t].c);
                              -- t;
                    }
                    ans[id] = res;
          }

          for(int i = 1; i <= mq; ++ i) {
                    cout << ans[i] << endl;
          }
}
```

普通莫队

```
#include <bits/stdc++.h>
using namespace std;

const int N = 1e6 + 10, M = 1e6 + 10;
int a[N];

struct node {
          int id, l, r;
} mp[M];

int len;
int ans[M], cnt[1000010];

int getNum(int l) {
          return l / len;
}

//左指针的分块，右指针的大小
bool cmp (const node &a, const node & b) {
          if(getNum(a.l) == getNum(b.l)) return a.r < b.r;
          return a.l < b.l;
}
/* 奇偶优化
struct node {
  int l, r, id;
  bool operator<(const node &x) const {
    if (l / unit != x.l / unit) return l < x.l;
    if ((l / unit) & 1)
      return r <  x.r;  // 注意这里和下面一行不能写小于（大
于）等于
    return r > x.r;
  }
};
*/

void add(int x, int& res) {
          if(cnt[x] == 0) res++;
          cnt[x] ++;
}

void del(int x, int& res) {
          cnt[x] --;
          if(cnt[x] == 0) res --;
}

int main() {
          ios::sync_with_stdio(0); cin.tie(0); cout.tie
(0);

          int n;
          cin >> n;
          for(int i = 1; i <= n; ++ i) {
                    cin >> a[i];
          }
          int m;
          cin >> m;
          len = sqrt((double)n * n / m);
          for(int i = 1; i <= m; ++ i) {
                    mp[i].id = i;
                    cin >> mp[i].l >> mp[i].r;
          }
          sort(mp + 1, mp + m + 1, cmp);

          //离线处理询问
          int res = 0, i = 0, j = 0;
          for(int k = 1; k <= m; ++ k) {
                    int id = mp[k].id, l = mp[k].l, r = mp
[k].r;
                    while(j < r) add(a[++j], res);
                    while(j > r) del(a[j--], res);
                    while(i < l) del(a[i++], res);
                    while(i > l) add(a[--i], res);
                    ans[id] = res;
          }

          for(int i = 1; i <= m; ++ i) {
                    cout << ans[i] << endl;
          }
          return 0;
}
```

树状数组（fenwick）

```
template <typename T>
struct fenwick {
    vector<T> fenw;
    int n;

    fenwick(int _n) : n(_n) {
```

```
        fenw.resize(n);
    }

    void clear(){
        fenw.clear();
        fenw.resize(n);
    }

    void modify(int x, T v) {
        while (x < n) {
            fenw[x] += v;
            //if(fenw[x]>=mod)fenw[x]-=mod;
            x |= (x + 1);
        }
    }

    T get(int x) {
        T v{};
        while (x >= 0) {
            v += fenw[x];
            //if(v>=mod)v-=mod;
            x = (x & (x + 1)) - 1;
        }
        return v;
    }

    T gets(int l,int r){
        T res=get(r)-get(l-1);
        //if(res<0)res+=mod;
        return res;
    }
};
```

线段树合并分裂

```
ll nodetot, recycnt, bac[N << 5], ch[N << 5][2], rt[N];
ll val[N << 5];

ll newnod() { return (recycnt ? bac[recycnt--] : ++nodeto
t); }

void recyc(ll p) {
    bac[++recycnt] = p, ch[p][0] = ch[p][1] = val[p] = 0;
    return;
}

void pushdown(ll p) {

}

void pushup(ll p) {
    val[p] = 0;
    if (ch[p][0]) val[p] += val[ch[p][0]];
    if (ch[p][1]) val[p] += val[ch[p][1]];
}

void modify(ll &p, ll l, ll r, ll pos, ll v) {
    if (!p) { p = newnod(); }
    if (l == r) {
        val[p] += v;
        return;
    }
    ll mid = (l + r) >> 1;
//    pushdown(p);
    if (pos <= mid) { modify(ch[p][0], l, mid, pos, v); }
    else { modify(ch[p][1], mid + 1, r, pos, v); }
    pushup(p);
    return;
}

ll query(ll p, ll l, ll r, ll xl, ll xr) {
    if (xr < l || r < xl) { return 0; }
    if (xl <= l && r <= xr) { return val[p]; }
    ll mid = (l + r) >> 1;
//    pushdown(p);
    return query(ch[p][0], l, mid, xl, xr) + query(ch[p]
[1], mid + 1, r, xl, xr);
}

ll kth(ll p, ll l, ll r, ll k) {
    if (l == r) { return l; }
    ll mid = (l + r) >> 1;
//    pushdown(p);
    if (val[ch[p][0]] >= k) { return kth(ch[p][0], l, mid,
 k); }
    else { return kth(ch[p][1], mid + 1, r, k - val[ch[p]
[0]]); }
}

ll merge(ll x, ll y, ll l, ll r) {
    if (!x || !y) {
        return x + y;
    }   // 只有一边有点，不用合并
    ll p = newnod(); // 创建一个新结点 p
    if (l == r) {                 // 边界（某些时候可以省略，
见下面一个代码）
        val[p] = val[x] + val[y];
        return p;
    }
//    pushdown(x), pushdown(y);
    ll mid = (l + r) >> 1;
    ch[p][0] = merge(ch[x][0], ch[y][0], l, mid);
    ch[p][1] = merge(ch[x][1], ch[y][1], mid + 1, r);
    recyc(x), recyc(y);           // 垃圾回收
    pushup(p);                     // pushup
    return p;
}

void split(ll x, ll &y, ll k) {
    if (x == 0) return;
    y = newnod();
    ll v = val[ch[x][0]];
//    pushdown(x);
    if (k > v) { split(ch[x][1], ch[y][1], k - v); }
    else { swap(ch[x][1], ch[y][1]); }
    if (k < v) { split(ch[x][0], ch[y][0], k); }
    val[y] = val[x] - k;
    val[x] = k;
    return;
}
```

舞蹈链（多重覆盖）

```
#include <bits/stdc++.h>
using namespace std;
struct DLX {
    static const int maxn = 1000;     //列的上限
    static const int maxr = 1000;     //解的上限
    static const int maxnode = 5000;  //总结点数上限
    static const int INF = 1000000000;
    int n, sz;
    int S[maxn];

    int row[maxnode], col[maxnode];
    int L[maxnode], R[maxnode], U[maxnode], D[maxnode];

    int ansd, ans[maxr];

    int vis[maxnode];

    void init(int n) {
        this->n = n;

        //虚拟节点
        for (int i = 0; i <= n; i++) {
            U[i] = i;
            D[i] = i;
            L[i] = i - 1;
            R[i] = i + 1;
        }
        R[n] = 0;
        L[0] = n;

        sz = n + 1;
        memset(S, 0, sizeof(S));
    }

    void addRow(int r, vector<int> columns) {
        int first = sz;
        for (int i = 0; i < columns.size(); i++) {
            int c = columns[i];
            L[sz] = sz - 1;
            R[sz] = sz + 1;
            D[sz] = c;
            U[sz] = U[c];
```

```
            D[U[c]] = sz;
            U[c] = sz;
            row[sz] = r;
            col[sz] = c;
            S[c]++;
            sz++;
        }
        R[sz - 1] = first;
        L[first] = sz - 1;
    }
#define FOR(i, A, s) for (int i = A[s]; i != s; i = A[i])
    void remove(int c) {
        FOR(i, D, c) { L[R[i]] = L[i], R[L[i]] = R[i]; }
    }

    void restore(int c) {
        FOR(i, U, c) { L[R[i]] = i, R[L[i]] = i; }
    }
    int f_check()  //精确覆盖区估算剪枝
    {
        /*
        强剪枝。这个
        剪枝利用的思想是A*搜索中的估价函数。即，对于当前的递归深度K下的矩阵，估计其最好情况下（即最少还需要多少步）才能出解。也就是，如果将能够覆盖当
        前列的所有行全部选中，去掉这些行能够覆盖到的列，将这个操作作为一层深度。重复此操作直到所有列全部出解的深度是多少。如果当前深度加上这个估价函数返
        回值，其和已然不能更优（也就是已经超过当前最优解），则直接返回，不必再搜。
        */

        int ret = 0;
        FOR(c, R, 0) vis[c] = true;
        FOR(c, R, 0)
        if (vis[c]) {
            ret++;
            vis[c] = false;
            FOR(i, D, c)
            FOR(j, R, i) vis[col[j]] = false;
        }
        return ret;
    }
    // d 为递归深度
    void dfs(int d, vector<int>& v) {
        if (d + f_check() >= ansd) return;
        if (R[0] == 0) {
            if (d < ansd) {
                ansd = d;
                v.clear();
                for (int i = 0; i < ansd; i++) {
                    v.push_back(ans[i]);
                }
            }       //找到解
            return;  //记录解的长度
        }

        //找到S最小的列c
        int c = R[0];
        FOR(i, R, 0)
        if (S[i] < S[c])
            c = i;      //第一个未删除的列
                        //删除第c列
        FOR(i, D, c) {  //用结点i所在的行能覆盖的所有其他列
            ans[d] = row[i];
            remove(i);
            FOR(j, R, i) remove(j);  //删除结点i所在的能覆的所有其他列
            dfs(d + 1, v);
            FOR(j, L, i) restore(j);
            restore(i);  //恢复结点i所在的行能覆盖的所有其他列
        }               //恢复第c列
    }

    bool solve(vector<int>& v) {
        v.clear();
        ansd = INF;
        dfs(0, v);
        return !v.empty();
    }
};
```

```
//使用时 init 初始化，vector 中存入 r 行结点列表用 addRow 加行，solve(ans)后答案按行的选择在 ans 中
DLX dlx;
int main() {
    int n, m;
    cin >> n >> m;
    dlx.init(m);
    for (int i = 1; i <= n; i++) {
        vector<int> v;
        for (int j = 1; j <= m; j++) {
            int a;
            cin >> a;
            if (a == 1) v.push_back(j);
        }
        dlx.addRow(i, v);
    }
    vector<int> ans;
    dlx.solve(ans);
    for (int i = 0; i < ans.size(); i++) cout << ans[i];
}
```

舞蹈链（精确覆盖）

```
#include <bits/stdc++.h>
using namespace std;
struct DLX {
    static const int maxn = 1000;     //列的上限
    static const int maxr = 1000;     //解的上限
    static const int maxnode = 5000;  //总结点数上限
    int n, sz;
    int S[maxn];

    int row[maxnode], col[maxnode];
    int L[maxnode], R[maxnode], U[maxnode], D[maxnode];

    int ansd, ans[maxr];

    void init(int n) {
        this->n = n;

        //虚拟节点
        for (int i = 0; i <= n; i++) {
            U[i] = i;
            D[i] = i;
            L[i] = i - 1;
            R[i] = i + 1;
        }
        R[n] = 0;
        L[0] = n;

        sz = n + 1;
        memset(S, 0, sizeof(S));
    }

    void addRow(int r, vector<int> columns) {
        int first = sz;
        for (int i = 0; i < columns.size(); i++) {
            int c = columns[i];
            L[sz] = sz - 1;
            R[sz] = sz + 1;
            D[sz] = c;
            U[sz] = U[c];
            D[U[c]] = sz;
            U[c] = sz;
            row[sz] = r;
            col[sz] = c;
            S[c]++;
            sz++;
        }
        R[sz - 1] = first;
        L[first] = sz - 1;
    }
#define FOR(i, A, s) for (int i = A[s]; i != s; i = A[i])
    void remove(int c) {
        L[R[c]] = L[c];
        R[L[c]] = R[c];
        FOR(i, D, c)
        FOR(j, R, i) {
            U[D[j]] = U[j];
            D[U[j]] = D[j];
            --S[col[j]];
        }
    }
```

```
    void restore(int c) {
        FOR(i, U, c)
        FOR(j, L, i) {
            ++S[col[j]];
            U[D[j]] = j;
            D[U[j]] = j;
        }
        L[R[c]] = c;
        R[L[c]] = c;
    }

    // d 为递归深度
    bool dfs(int d) {
        if (R[0] == 0) {
            ansd = d;     //找到解
            return true;  //记录解的长度
        }

        //找到S 最小的列c
        int c = R[0];
        FOR(i, R, 0) if (S[i] < S[c]) c = i;  //第一个未删除的列

        remove(c);      //删除第 c 列
        FOR(i, D, c) {  //用结点 i 所在的行能覆盖的所有其他列
            ans[d] = row[i];
            FOR(j, R, i) remove(col[j]);  //删除结点 i 所在的能覆的所有其他列
            if (dfs(d + 1)) return true;
            FOR(j, L, i) restore(col[j]);  //恢复结点 i 所在的行能覆盖的所有其他列
        }
        restore(c);  //恢复第 c 列

        return false;
    }

    bool solve(vector<int>& v) {
        v.clear();
        if (!dfs(0)) return false;
        for (int i = 0; i < ansd; i++) v.push_back(ans[i]);
        return true;
    }
};
//使用时 init 初始化，vector 中存入 r 行结点列表用 addRow 加行，solve(ans)后答案按行的选择在 ans 中
```

数论

BSGS 扩展 BSGS

BSGS

求$a^t \equiv b \pmod p$ (a,p) = 1 的最小的 t

$$t = x \times k - y, x \in [1,k], y \in [0,k-1]$$
$$t \in [1,k^2]$$

$$a^k x \equiv b \times a^y \pmod p$$

对 $b \times a^y$ 建立 hash 表，枚举 x 看是否有解

```
#include <bits/stdc++.h>
using namespace std;

typedef long long ll;

unordered_map<int , int> mp;

int bsgs(int a, int p, int b) {

        if (1 % p == b % p) return 0; // 特判0 是不是解
        mp.clear();

        int k = sqrt(p) + 1;

        for(int i = 0, j = b % p; i < k; ++ i, j = (ll)j * a % p) {
                mp[j] = i;
        }

        int ak = 1;
        for(int i = 0; i < k; ++i) {
                ak = (ll)ak * a % p;
        }

        for(int i = 1, j = ak % p; i <= k; ++ i, j = (ll)j * ak % p) {
                if(mp.count(j)) return (ll)i * k - mp[j];
        }

        return -1;
}

int main() {
        ios::sync_with_stdio(0);
        cin.tie(0); cout.tie(0);

        int a, p, b;
        while(cin >> a >> p >> b, a | p | b) {
                int res;
                res = bsgs(a, p, b);
                if(res == -1) {
                        cout << "No Solution\n";
                }
                else {
                        cout << res << endl;
                }
        }

        return 0;
}
```

扩展 BSGS

求$a^t \equiv b \pmod p$ 的最小的 t

当$(a,p) != 1$

$(a,p) = d$ $d \nmid b$ 无解

$a^t \equiv b \pmod p$ ， $a^t + kp = b$ 两边同时除以 d， $\frac{a}{d}a^{t-1} + k\frac{p}{d} = \frac{b}{d}$

$$a^{t-1} \equiv \frac{b}{d}\left(\frac{a}{d}\right)^{-1}$$

$$t' = t-1, p' = \frac{p}{d}, b' = \frac{b}{a}\left(\frac{a}{d}\right)^{-1}$$

```
#include <bits/stdc++.h>
using namespace std;

typedef long long ll;

unordered_map<ll, ll> mp;

ll bsgs(ll a, ll p, ll b) {

        if(1 % p == b % p) return 0; // 特判0 是不是解
        mp.clear();

        ll k = sqrt(p) + 1;

        for(ll i = 0, j = b % p; i < k; ++i, j = (ll)j * a % p) {
                mp[j] = i;
        }

        ll ak = 1;
        for(ll i = 0; i < k; ++i) {
                ak = (ll) ak * a % p;
        }

        for(ll i = 1, j = ak % p;i <= k; ++i, j = (ll)j * ak % p) {
                if(mp.count(j)) return (ll) i * k - mp[j];
```

```
        }

        return -1;
}

ll gcd(ll x, ll y) {
        return x % y == 0 ? y : gcd(y, x % y);
}

void extgcd(ll a,ll b,ll& d,ll& x,ll& y){
    if(!b){
        d = a; x = 1; y = 0;
    }
    else{
        extgcd(b, a%b, d, y, x);
        y -= x * (a / b);
    }
}

ll inverse(ll a,ll n){
    ll d,x,y;
    extgcd(a,n,d,x,y);
    return d == 1 ? (x + n) % n : -1;

}

int main() {
        ll a, p, b;

        while(cin >> a >> p >> b, a | p | b) {
                ll d = gcd(a, p);
                if(d == 1) {
                        ll res = bsgs(a, p, b);
                        if(res == -1) {
                                cout << "No Solut
ion\n";
                        }
                        else {
                                cout << res << en
dl;
                        }
                }
                else {
                        if(b % d != 0) {
                                cout << "No Solut
ion\n";
                                continue;
                        }
                        else {
                                p = p / d;
                                b = (b / d) * inv
erse(a / d, p);
                                ll res = bsgs(a,
p, b);
                                if(res == -1) {
                                        cout <<
 "No Solution\n";
                                }
                                else {
                                        cout <<
 res + 1 << endl;
                                }
                        }
                }
        }

        return 0;
}
```

burnside&polya

burnside 引理

burnside ：用$D(a_j)$ 表示在置换a_j下不变的元素的个数，L 表示本质不同的方案数（等价类）：

$$L = \frac{1}{|G|}\sum_{j=1}^{s} D(a_j)$$

定理：$|E_k| \times |Z_k| = |G|, k = 1,2,\ldots,n,$ 该定理的一个重要研究对象是群的元素个数，其中Z_k 是 K 不动置换类，设 G 是 1,2，..n 的置换群，若 k 是 1 到 n 中某个元素，则 G 中使 K 保持不变的置换的全体，为Z_k. E_k是等价类，设 G 是 1,2，..n 的置换群，若 k 是 1 到 n 中某个元素，k 在 G 的作用下的轨迹，为E_k，即 k 在 G 的作用下能变化成的所有元素的集合

$$\sum_{k=1}^{n}|Z_k| = \sum_{i=1}^{L}\sum_{k\in E_i}|Z_k| = \sum_{i=1}^{L}|E_i| \times |Z_i| = L \times |G|$$

每个置换的不动点的平均值就是不同方案数

任何一个置换都可以拆解成若干个循环置换

polya 定理

polya: 设 G 是 p 个对象的一个置换群，用 m 种颜色涂染 p 个对象，则不同染色的方案为：

$$L = \frac{1}{|G|}\left(m^{c(g_1)} + m^{c(g_2)} + \ldots + m^{c(g_s)}\right)$$

其中$G = \{g_1, g_2 \ldots g_s\}, c(g_i)$ 为置换g_i为置换的循环节数

浅证：$D(a_j) = m^{c(g_i)}$

每个置换的不动点有公式可以求 $每个循环的方案数^{循环数}$

(不同循环直接完全独立)

Cipolla

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;

ll mod;
ll I_mul_I; // 虚数单位的平方

struct Complex {
    ll real, imag;

    Complex(ll real = 0, ll imag = 0) : real(real), imag(i
mag) {}
};

inline bool operator==(Complex x, Complex y) {
    return x.real == y.real and x.imag == y.imag;
}

inline Complex operator*(Complex x, Complex y) {
    return Complex((x.real * y.real + I_mul_I * x.imag % m
od * y.imag) % mod,
                   (x.imag * y.real + x.real * y.imag) % mo
d);
}

Complex power(Complex x, ll k) {
    Complex res = 1;
    while (k) {
        if (k & 1) res = res * x;
        x = x * x;
        k >>= 1;
    }
    return res;
}

bool check_if_residue(ll x) {
    return power(x, (mod - 1) >> 1) == 1;
}

void solve(ll n, ll &x0, ll &x1) {

    ll a = rand() % mod;
    while (!a or check_if_residue((a * a + mod - n) % mod))
        a = rand() % mod;
    I_mul_I = (a * a + mod - n) % mod;
    x0 = ll(power(Complex(a, 1), (mod + 1) >> 1).real);
    x1 = mod - x0;
}
```

```
signed main() {
    ios::sync_with_stdio(false);
    cin.tie(nullptr);
    cout.tie(nullptr);

    ll t;
    cin >> t;
    while (t--) {
        ll n;
        cin >> n >> mod;
        if (n == 0) {
            cout << 0 << endl;
            continue;
        }
        if (!check_if_residue(n)) {
            cout << "Hola!" << endl;
            continue;
        }
        ll x0, x1;
        solve(n, x0, x1);
        if (x0 > x1) swap(x0, x1);
        cout << x0 << ' ' << x1 << endl;
    }
}
```

exgcd

```
ll ex_gcd(ll a, ll b, ll &x, ll &y) {
    if (b == 0) {
        x = 1;
        y = 0;
        return a;
    }
    ll d = ex_gcd(b, a % b, x, y);
    ll temp = x;
    x = y;
    y = temp - a / b * y;
    return d;
}
```

FFT

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const int N = 1e7 + 10;


const double Pi = acos(-1.0);

struct Complex {
    double x, y;

    Complex(double xx = 0, double yy = 0) { x = xx, y = yy;
 }
} a[N], b[N];

Complex operator+(Complex _a, Complex _b) { return Complex
(_a.x + _b.x, _a.y + _b.y); }

Complex operator-(Complex _a, Complex _b) { return Complex
(_a.x - _b.x, _a.y - _b.y); }

Complex operator*(Complex _a, Complex _b) {
    return Complex(_a.x * _b.x - _a.y * _b.y, _a.x * _b.y
+ _a.y * _b.x);
} //不懂的看复数的运算那部分

int L, r[N];
int limit = 1;

void fft(Complex *A, int type) {
    for (int i = 0; i < limit; i++)
        if (i < r[i]) swap(A[i], A[r[i]]); //求出要迭代的序
列
    for (int mid = 1; mid < limit; mid <<= 1) { //待合并区
间的长度的一半
        Complex Wn(cos(Pi / mid), type * sin(Pi / mid)); //
单位根
        for (int R = mid << 1, j = 0; j < limit; j += R) {
//R 是区间的长度，j 表示前已经到哪个位置了
            Complex w(1, 0); //幂
            for (int k = 0; k < mid; k++, w = w * Wn) { //枚
举左半部分
                Complex x = A[j + k], y = w * A[j + mid + k];
 //蝴蝶效应
                A[j + k] = x + y;
                A[j + mid + k] = x - y;

            }
        }
    }
}

void FFT(int n, int m) {
    limit = 1;
    L = 0;
    while (limit <= n + m) limit <<= 1, L++;
    for (int i = 0; i < limit; i++) r[i] = (r[i >> 1] >> 1)
 | ((i & 1) << (L - 1));
    // 在原序列中 i 与 i/2 的关系是 ： i 可以看做是i/2 的二进
制上的每一位左移一位得来
    // 那么在反转后的数组中就需要右移一位，同时特殊处理一下奇
数
    fft(a, 1), fft(b, 1);
    for (int i = 0; i <= limit; i++) a[i] = a[i] * b[i];
    fft(a, -1);
    for (int i = 0; i <= n + m; i++) a[i].x /= limit;
}

int main() {
    int n, m;
    cin >> n >> m;
    for (int i = 0; i <= n; i++) cin >> a[i].x;
    for (int i = 0; i <= m; i++) cin >> b[i].x;
    FFT(n, m);
    for (int i = 0; i <= n + m; i++) cout << (int) (a[i].x
+ 0.5) << ' ';
    return 0;
}
```

FWT

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const int mod = 998244353;

void add(int &x, int y) {
    (x += y) >= mod && (x -= mod);
}

void sub(int &x, int y) {
    (x -= y) < 0 && (x += mod);
}

namespace FWT {
    int extend(int n) {
        int N = 1;
        for (; N < n; N <<= 1);
        return N;
    }

    void FWTor(std::vector<int> &a, bool rev) {
        int n = a.size();
        for (int l = 2, m = 1; l <= n; l <<= 1, m <<= 1) {
            for (int j = 0; j < n; j += l)
                for (int i = 0; i < m; i++) {
                    if (!rev) add(a[i + j + m], a[i + j]);
                    else sub(a[i + j + m], a[i + j]);
                }
        }
    }

    void FWTand(std::vector<int> &a, bool rev) {
        int n = a.size();
        for (int l = 2, m = 1; l <= n; l <<= 1, m <<= 1) {
            for (int j = 0; j < n; j += l)
                for (int i = 0; i < m; i++) {
                    if (!rev) add(a[i + j], a[i + j + m]);
```

```
                else sub(a[i + j], a[i + j + m]);
            }
        }
    }

    void FWTxor(std::vector<int> &a, bool rev) {
        int n = a.size(), inv2 = (mod + 1) >> 1;
        for (int l = 2, m = 1; l <= n; l <<= 1, m <<= 1) {
            for (int j = 0; j < n; j += l)
                for (int i = 0; i < m; i++) {
                    int x = a[i + j], y = a[i + j + m];
                    if (!rev) {
                        a[i + j] = (x + y) % mod;
                        a[i + j + m] = (x - y + mod) % mod;
                    } else {
                        a[i + j] = 1LL * (x + y) * inv2 % mo
d;
                        a[i + j + m] = 1LL * (x - y + mod) *
 inv2 % mod;
                    }
                }
        }
    }

    std::vector<int> Or(std::vector<int> a1, std::vector<i
nt> a2) {
        int n = std::max(a1.size(), a2.size()), N = extend
(n);
        a1.resize(N), FWTor(a1, false);
        a2.resize(N), FWTor(a2, false);
        std::vector<int> A(N);
        for (int i = 0; i < N; i++) A[i] = 1LL * a1[i] * a2
[i] % mod;
        FWTor(A, true);
        return A;
    }

    std::vector<int> And(std::vector<int> a1, std::vector<
int> a2) {
        int n = std::max(a1.size(), a2.size()), N = extend
(n);
        a1.resize(N), FWTand(a1, false);
        a2.resize(N), FWTand(a2, false);
        std::vector<int> A(N);
        for (int i = 0; i < N; i++) A[i] = 1LL * a1[i] * a2
[i] % mod;
        FWTand(A, true);
        return A;
    }

    std::vector<int> Xor(std::vector<int> a1, std::vector<
int> a2) {
        int n = std::max(a1.size(), a2.size()), N = extend
(n);
        a1.resize(N), FWTxor(a1, false);
        a2.resize(N), FWTxor(a2, false);
        std::vector<int> A(N);
        for (int i = 0; i < N; i++) A[i] = 1LL * a1[i] * a2
[i] % mod;
        FWTxor(A, true);
        return A;
    }
};

int main() {
    int n;
    scanf("%d", &n);
    n = (1 << n);
    std::vector<int> a1(n), a2(n);
    for (int i = 0; i < n; i++) scanf("%d", &a1[i]);
    for (int i = 0; i < n; i++) scanf("%d", &a2[i]);
    std::vector<int> A;
    A = FWT::Or(a1, a2);
    for (int i = 0; i < n; i++) {
        printf("%d%c", A[i], " \n"[i == n - 1]);
    }
    A = FWT::And(a1, a2);
    for (int i = 0; i < n; i++) {
        printf("%d%c", A[i], " \n"[i == n - 1]);
    }
    A = FWT::Xor(a1, a2);
    for (int i = 0; i < n; i++) {
        printf("%d%c", A[i], " \n"[i == n - 1]);
    }
    return 0;
}
```

lucas 求组合数

```
#include <bits/stdc++.h>
using namespace std;

typedef long long ll;

ll p;

const int maxn = 1e5 + 10;

ll qpow(ll x, ll n){
        ll res = 1;
        while(n){
                if(n & 1) res = (res * x) % p;
                x = (x * x) % p;
                n >>= 1;
        }

        return res;
}

ll C(ll up, ll down){
        if(up > down) return 0;
        ll res = 1;

//      for(int i = up + 1; i <= down; ++ i){
//              res = (res * i) % p;
//      }
//      for(int i = 1; i <= down - up; ++ i){
//              res = (res * qpow(i, p - 2)) % p;
//      }

        for(int i = 1, j = down; i <= up; ++ i, -- j){
                res = (res * j) % p;
                res = (res * qpow(i, p - 2)) % p;
        }

        return res;
}


ll lucas(ll up, ll down){
        if(up < p && down < p) return C(up, down);
        return C(up % p, down % p) * lucas(up / p, down
/ p) % p;
}

int main(){
        ios::sync_with_stdio(0); cin.tie(0); cout.tie
(0);

        int T;
        cin >> T;
        while (T --){
                ll down, up;
                cin >> down >> up >> p;

                cout << lucas(up, down) % p << endl;
        }

        return 0;
}
```

min_25 筛

```
/*
https://loj.ac/p/6053
筛积性函数f的前缀和
f(1)=1
f(p^e)=f xor e
n<=1e10，LOJ 347ms 本地1100ms
*/
#include<bits/stdc++.h>
using namespace std;
typedef long long ll;
const ll mod=1e9+7,inv3=333333336;
const int N=1e5+5;//开到sqrt(n)即可
```

```
ll prime[N],sp0[N],sp1[N],sp2[N],g0[N<<1],g1[N<<1],g2[N<<1];
ll pnum,min25n,sqrn,w[N<<1],ind1[N],ind2[N];
bool notp[N];

void pre() { //预处理，线性筛
    notp[1]=1;
    for(int i=1; i<N; i++) {
        if(!notp[i]) {
            prime[++pnum]=i;
            sp0[pnum]=(sp0[pnum-1]+1)%mod;//p^0 前缀和（p 指质数），可以按需增删，下标意义为第 pnum 个质数的前缀和，而 g 的实际下标意义为 w 之前的前缀和，两者有所区别
            sp1[pnum]=(sp1[pnum-1]+i)%mod;//p^1 前缀和
            sp2[pnum]=(sp2[pnum-1]+1ll*i*i)%mod;//p^2 前缀和
        }
        for(int j=1; j<=pnum&&prime[j]*i<N; j++) {
            notp[i*prime[j]]=1;
            if(i%prime[j]==0)break;
        }
    }
}

void min25(ll n) {
    ll tot=0;
    min25n=n;
    sqrn=sqrt(n);
    for(ll i=1; i<=n; i=n/(n/i)+1) {
        w[++tot]=n/i;//实际下标
        ll x=w[tot]%mod;
        g0[tot]=x-1;//x^0 前缀和
        g1[tot]=x*(x+1)/2%mod-1;//x^1 前缀和
        g2[tot]=x*(x+1)/2%mod*(2*x+1)%mod*inv3%mod-1;//x^2 前缀和
        if(n/i<=sqrn)ind1[n/i]=tot;//离散下标
        else ind2[n/(n/i)]=tot;//离散下标
    }
    for(int i=1; i<=pnum; i++) {//扩展埃氏筛，筛质数部分前缀和
        for(int j=1; j<=tot&&prime[i]*prime[i]<=w[j]; j++) {
            int id=w[j]/prime[i]<=sqrn?ind1[w[j]/prime[i]]:ind2[n/(w[j]/prime[i])];
            g0[j]-=(g0[id]-sp0[i-1]+mod)%mod;
            g1[j]-=prime[i]*(g1[id]-sp1[i-1]+mod)%mod;
            g2[j]-=prime[i]*prime[i]%mod*(g2[id]-sp2[i-1]+mod)%mod;
            g0[j]%=mod,g1[j]%=mod,g2[j]%=mod;
            if(g0[j]<0)g0[j]+=mod;
            if(g1[j]<0)g1[j]+=mod;
            if(g2[j]<0)g2[j]+=mod;
        }
    }
}

//该前缀和不计算 f(1)，需要自行加上
ll S(ll x,int y) {//x 以内最小质因子大于第 y 个因子的前缀和
    if(prime[y]>=x)return 0;
    int id=x<=sqrn?ind1[x]:ind2[min25n/x];
    ll ans=(((g1[id]-g0[id])-(sp1[y]-sp0[y]))%mod+mod)%mod;//x 以内大于第 y 个因子的质数部分前缀和
    if(x>=2&&y<1)ans=(ans+2)%mod;//特判包含 f(2)的情况
    for(int i=y+1; i<=pnum&&prime[i]*prime[i]<=x; i++) {//筛合数部分前缀和
        ll pe=prime[i];
        for(int e=1; pe<=x; e++,pe=pe*prime[i]) {
            ll fpe=prime[i]^e;//f(p^e)
            ans=(ans+fpe%mod*(S(x/pe,i)+(e!=1)))%mod;
        }
    }
    return ans%mod;
}

int main() {
    pre();//预处理一次即可
    ll n;
    scanf("%lld",&n);
    min25(n);//每个不同的 n 都要调用一次该函数，再调用 S(n,0)
    printf("%lld\n",S(n,0)+1);//加上 f(1)
    return 0;
}
```

NTT

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;

const int N = 4e6 + 10;
const ll mod = 998244353, G = 3, Gi = 332748118;

int limit = 1, L, r[N];
ll a[N], b[N];

ll qpow(ll _a, ll _b) {
    ll ans = 1;
    while (_b) {
        if (_b & 1) ans = (ans * _a) % mod;
        _b >>= 1;
        _a = (_a * _a) % mod;
    }
    return ans;
}

void ntt(ll *A, int type) {
    auto swap = [](ll &_a, ll &_b) {
        _a ^= _b, _b ^= _a, _a ^= _b;
    };
    for (int i = 0; i < limit; i++)
        if (i < r[i]) swap(A[i], A[r[i]]);
    for (int mid = 1; mid < limit; mid <<= 1) {
        ll Wn = qpow(type == 1 ? G : Gi, (mod - 1) / (mid << 1));
        for (int j = 0; j < limit; j += (mid << 1)) {
            ll w = 1;
            for (int k = 0; k < mid; k++, w = (w * Wn) % mod) {
                int x = A[j + k], y = w * A[j + k + mid] % mod;
                A[j + k] = (x + y) % mod,
                        A[j + k + mid] = (x - y + mod) % mod;
            }
        }
    }
}

void NTT(int n, int m) {
    limit = 1;
    L = 0;
    while (limit <= n + m) limit <<= 1, L++;
    for (int i = 0; i < limit; i++) r[i] = (r[i >> 1] >> 1) | ((i & 1) << (L - 1));
    ntt(a, 1), ntt(b, 1);
    for (int i = 0; i < limit; i++) a[i] = (a[i] * b[i]) % mod;
    ntt(a, -1);
    ll inv = qpow(limit, mod - 2);
    for (int i = 0; i <= n + m; i++) a[i] = a[i] * inv % mod;
}

int main() {
    int n, m;
    cin >> n >> m;
    for (int i = 0; i <= n; i++) {
        cin >> a[i];
        a[i] = (a[i] + mod) % mod;
    }
    for (int i = 0; i <= m; i++) {
        cin >> b[i];
        b[i] = (b[i] + mod) % mod;
    }
    NTT(n, m);
    for (int i = 0; i <= n + m; i++) cout << a[i] << ' ';
}
```

Pollard_Rho+Miller-Robin

```
typedef long long ll;
namespace Miller_Rabin {
    const ll Pcnt = 12;
```

```
    const ll p[Pcnt] = {2, 3, 5, 7, 11, 13, 17, 19, 61, 233
3, 4567, 24251};

    ll pow(ll a, ll b, ll p) {
        ll ans = 1;
        for (; b; a = (__int128) a * a % p, b >>= 1)if (b &
 1)ans = (__int128) ans * a % p;
        return ans;
    }

    bool check(ll x, ll p) {
        if (x % p == 0 || pow(p % x, x - 1, x) ^ 1)return t
rue;
        ll t, k = x - 1;
        while ((k ^ 1) & 1) {
            t = pow(p % x, k >>= 1, x);
            if (t ^ 1 && t ^ x - 1)return true;
            if (!(t ^ x - 1))return false;
        }
        return false;
    }

    inline bool MR(ll x) {  //用这个
        if (x < 2)return false;
        for (int i = 0; i ^ Pcnt; ++i) {
            if (!(x ^ p[i]))return true;
            if (check(x, p[i]))return false;
        }
        return true;
    }
}
namespace Pollard_Rho {
#define Rand(x) (lll*rand()*rand()%(x)+1)

    ll gcd(const ll a, const ll b) { return b ? gcd(b, a %
b) : a; }

    ll mul(const ll x, const ll y, const ll X) {
        ll k = (1.0L * x * y) / (1.0L * X) - 1, t = (__int1
28) x * y - (__int128) k * X;
        while (t < 0)t += X;
        return t;
    }

    ll PR(const ll x, const ll y) {
        int t = 0, k = 1;
        ll v0 = Rand(x - 1), v = v0, d, s = 1;
        while (true) {
            v = (mul(v, v, x) + y) % x, s = mul(s, abs(v -
v0), x);
            if (!(v ^ v0) || !s)return x;
            if (++t == k) {
                if ((d = gcd(s, x)) ^ 1)return d;
                v0 = v, k <<= 1;
            }
        }
    }

    void Resolve(ll x, ll &ans) {
        if (!(x ^ 1) || x <= ans)return;
        if (Miller_Rabin::MR(x)) {
            if (ans < x)ans = x;
            return;
        }
        ll y = x;
        while ((y = PR(x, Rand(x))) == x);
        while (!(x % y))x /= y;
        Resolve(x, ans);
        Resolve(y, ans);
    }

    long long check(ll x) { //用这个，素数返回本身
        ll ans = 0;
        Resolve(x, ans);
        return ans;
    }
}
```

prufer

Prufer 序列 (Prufer code)，这是一种将带标号的树用一个唯一的整数序列表示的方法。

Prufer 序列可以将一个带标号 n 个结点的树用$[1,n]$中的$n-2$个整数表示。你也可以把它理解为完全图的生成树与数列之间的双射。

显然你不会想不开拿这玩意儿去维护树结构。这玩意儿常用组合计数问题上。

线性建立 prufer

Prufer 是这样建立的：每次选择一个编号最小的叶结点并删掉它，然后在序列中记录下它连接到的那个结点。重复 n - 2 次后就只剩下两个结点，算法结束。

线性构造的本质就是维护一个指针指向我们将要删除的结点。首先发现，叶结点数是非严格单调递减的。要么删一个，要么删一个得一个。

于是我们考虑这样一个过程：维护一个指针 p 。初始时 p 指向编号最小的叶结点。同时我们维护每个结点的度数，方便我们知道在删除结点的时候是否产生新的叶结点。操作如下：

1. 删除 指向的结点，并检查是否产生新的叶结点。
2. 如果产生新的叶结点，假设编号为 x ，我们比较 p,x 的大小关系。如果 x>p，那么不做其他操作；否则就立刻删除 x，然后检查删除 x 后是否产生新的叶结点，重复 2 步骤，直到未产生新节点或者新节点的编号>p 。
3. 让指针 p 自增直到遇到一个未被删除叶结点为止；

循环上述操作 n - 2 次，就完成了序列的构造。

```
// 从原文摘的代码，同样以 0 为起点
vector<vector<int>> adj;
vector<int> parent;

void dfs(int v) {
  for (int u : adj[v]) {
    if (u != parent[v]) parent[u] = v, dfs(u);
  }
}

vector<int> pruefer_code() {
  int n = adj.size();
  parent.resize(n), parent[n - 1] = -1;
  dfs(n - 1);

  int ptr = -1;
  vector<int> degree(n);
  for (int i = 0; i < n; i++) {
    degree[i] = adj[i].size();
    if (degree[i] == 1 && ptr == -1) ptr = i;
  }

  vector<int> code(n - 2);
  int leaf = ptr;
  for (int i = 0; i < n - 2; i++) {
    int next = parent[leaf];
    code[i] = next;
    if (--degree[next] == 1 && next < ptr) {
      leaf = next;
    } else {
      ptr++;
      while (degree[ptr] != 1) ptr++;
      leaf = ptr;
    }
  }
  return code;
}
```

性质

1. 在构造完 Prufer 序列后原树中会剩下两个结点，其中一个一定是编号最大的点 。
2. 每个结点在序列中出现的次数是其度数减 1 。（没有出现的就是叶结点）

线性 prufer 转化成树

同线性构造 Prufer 序列的方法。在删度数的时候会产生新的叶结点，于是判断这个叶结点与指针 p 的大小关系，如果更小就优先考虑它

```
// 原文摘代码
vector<pair<int, int>> pruefer_decode(vector<int> const& code) {
  int n = code.size() + 2;
  vector<int> degree(n, 1);
  for (int i : code) degree[i]++;

  int ptr = 0;
  while (degree[ptr] != 1) ptr++;
  int leaf = ptr;

  vector<pair<int, int>> edges;
  for (int v : code) {
    edges.emplace_back(leaf, v);
    if (--degree[v] == 1 && v < ptr) {
      leaf = v;
    } else {
      ptr++;
      while (degree[ptr] != 1) ptr++;
      leaf = ptr;
    }
  }
  edges.emplace_back(leaf, n - 1);
  return edges;
}
```

cayley 公式

完全图K_n 有n^{n-2} 棵生成树。

用 Prufer 序列证:任意一个长度为 n - 2 的值域 [1, n] 的整数序列都可以通过 Prufer 序列双射对应一个生成树，于是方案数就是n^{n-2} 。

图连通方案数

一个 n 个点 m 条边的带标号无向图有 k 个连通块。我们希望添加 k - 1 条边使得整个图连通。求方案数。

设s_i表示每个连通块的数量。我们对 k 个连通块构造 Prufer 序列，然后你发现这并不是普通的 Prufer 序列。因为每个连通块的连接方法很多。不能直接淦就设啊。于是设d_i为第 i 个连通块的度数。由于度数之和是边数的两倍，于是$\sum_{i=1}^{k} d_i = 2k - 2$ 。则对于给定的 d 序列构造 Prufer 序列的方案数是

$$\tbinom{k - 2}{d_1 - 1, d_2 - 1, \dots, d_k - 1} = \frac{(k - 2)!}{(d_1 - 1)!(d_2 - 1)! \dots (d_k - 1)!}$$

对于第 i 个连通块，它的连接方式有$s_i^{d_i}$种，因此对于给定 d 序列使图连通的方案数是

$$\tbinom{k - 2}{d_1 - 1, d_2 - 1, \dots, d_k - 1} \prod_{i = 1}^{k}s_i^{d_i}$$

现在我们要枚举 d 序列，式子变成

$$\sum_{d_i\geq 1. \sum_{i = 1}^{k} d_i = 2k - 2} \tbinom{k - 2}{d_1 - 1, d2 - 1, \dots ,d_k - 1} \prod_{i = 1}^{k}s_i^{d_i}$$

根据多元二项式定理

$$(x_1+\dots+x_m)^{p}= \sum_{c_i \geq 0, \sum_{i = 1}^{m} c_i = p} \tbinom{p}{C_1, C_2, \dots , C_m} \prod_{i= 1}^{m}x_i^{C_i}$$

对原式换元，设$e_i = d_i - 1$ ，显然有$\sum_{i=1}^{k} e_i = k - 2$

$$\Rightarrow \sum_{e_i\ge 0, \sum_{i= 1}^{k} e_i = k - 2} \tbinom{k - 2}{e_1, e_2, \dots, e_k} \prod_{i = 1}^{k}s_i ^{e_i+ 1} \\ 化简 \Rightarrow (s_1 + s_2 + \dots + s_k)^{k - 2} \prod_{i = 1}^{k}s_i \\ \Rightarrow n^{k - 2} \prod_{i = 1}^{k} s_i$$

中国剩余定理

```
#include <bits/stdc++.h>

using namespace std;

typedef long long ll;

const int maxn = 20;

ll A[maxn], B[maxn];

ll exgcd(ll a, ll b, ll & x, ll & y) {
          if(b == 0) {
                    x = 1, y = 0;
                    return a;
          }

          ll d = exgcd(b, a % b, y, x);

          y -= (a / b) * x;

          return d;
}

int main() {
          int n;
          cin >> n;
          ll M = 1ll;
          for(int i = 0; i < n; ++ i) {
                    cin >> A[i] >> B[i];
                    M = M * A[i];
          }

          ll ans = 0;

          ll x, y;

          for(int i = 0; i < n; ++ i) {
                    ll Mi = M / A[i];
                    exgcd(Mi, A[i], x, y);
                    ans += B[i] * Mi * x;
          }

          cout << (ans % M + M) % M;

}
```

二次剩余

解的数量

对于$x^2 \equiv n(\mod p)$ 能满足 n 是 mod p 的二次剩余的 n 一共有$\frac{p-1}{2}$个（不包括 0），非二次剩余为$\frac{p-1}{2}$个

勒让德符号

$$\left(\frac{n}{p}\right) = \begin{cases} 1, p \nmid n, n是p的二次剩余 \\ -1, p \nmid n, n不是p的二次剩余 \\ 0, p|n \end{cases}$$

欧拉判别准则

$$\left(\frac{n}{p}\right) \equiv n^{\frac{p-1}{2}} (\mod p)$$

若 n 是二次剩余，当且仅当$n^{\frac{p-1}{2}} \equiv 1(\mod p)$

若 n 是非二次剩余，当且仅当$n^{\frac{p-1}{2}} \equiv -1(\mod p)$

Cipolla

找到一个数 a 满足$a^2 - n$是 **非二次剩余** ，至于为什么要找满足非二次剩余的数，在下文会给出解释。 这里通过生成随机数再检验的方法来实现，由于非二次剩余的数量为 $\frac{p-1}{2}$，接近$\frac{p}{2}$，所以期望约 2 次就可以找到这个数。

建立一个 " 复数域 " ，并不是实际意义上的复数域，而是根据复数域的概念建立的一个类似的域。 在复数中$i^2 = -1$ ，这里定义$i^2 = a^2 - n$ ，

于是就可以将所有的数表达为$A + Bi$ 的形式，这里的 和 都是模意义下的数，类似复数中的实部和虚部。

在有了 i 和 a 后可以直接得到答案，$x^2 \equiv n(\mod p)$ 的解为$(a + i)^{\frac{p+1}{2}}$。

```
#include <bits/stdc++.h>
using namespace std;

typedef long long ll;
int t;
ll n, p;
ll w;

struct num {          //建立一个复数域

          ll x, y;
};

num mul(num a, num b, ll p) {  //复数乘法
          num ans = {0, 0};
          ans.x = ((a.x * b.x % p + a.y * b.y % p * w % p)
 % p + p) % p;
          ans.y = ((a.x * b.y % p + a.y * b.x % p) % p + p)
 % p;
          return ans;
}

ll binpow_real(ll a, ll b, ll p) {        //实部快速幂
          ll ans = 1;
          while (b) {
                    if (b & 1) ans = ans * a % p;
                    a = a * a % p;
                    b >>= 1;
          }
          return ans % p;
}

ll binpow_imag(num a, ll b, ll p) {       //虚部快速幂
          num ans = {1, 0};
          while (b) {
                    if (b & 1) ans = mul(ans, a, p);
                    a = mul(a, a, p);
                    b >>= 1;
          }
          return ans.x % p;
}

ll cipolla(ll n, ll p) {
          n %= p;
          if (p == 2) return n;
          if (binpow_real(n, (p - 1) / 2, p) == p - 1) ret
urn -1;
          ll a;
          while (1) {          //生成随机数再检验找到满足非
二次剩余的a
                    a = rand() % p;
                    w = ((a * a % p - n) % p + p) % p;
                    if (binpow_real(w, (p - 1) / 2, p) ==
p - 1) break;
          }
          num x = {a, 1};
          return binpow_imag(x, (p + 1) / 2, p);
}
```

勾股数圆上格点数

勾股数

$$a^2 + b^2 = c^2$$

1.任何一个勾股数(a,b,c)内的三个数同时乘以一个正整数 n 得到的新数组(na, nb, nc)仍然是勾股数，

于是找 abc 互质的勾股数

一，当 a 为大于 1 的奇数 2n+1 时，$b = 2n^2 + 2n$，$c = 2n^2 + 2n + 1$

（把 a 拆成两个连续的自然数）

二，当 a 为大于 4 的偶数 2n 时，$b = n^2 - 1, c = n^2 + 1$

（只想得到互质的数的话：a=4n，$b = 4n^2 - 1, c = 4n^2 + 1$

公式 1

a=2mnt

b=（m^2-n^2）t

c=（m^2+n^2）t

（t 是倍数）

完全公式

a=m，b=(m^2 / k - k) / 2，c=(m^2 / k + k) / 2 ①

其中 m ≥3

1. 当 m 确定为任意一个 ≥3 的奇数时，k={1，m^2 的所有小于 m 的因子}

2. 当 m 确定为任意一个 ≥4 的偶数时，k={m^2 / 2 的所有小于 m 的偶数因子}

高斯整数/高斯素数

3B1B 的视频

洛谷某题

二维平面转化为复数平面，

4n+1 的素数，都能分解成高斯素数，4n+3 的素数，他们本身就是高斯素数，2 特殊

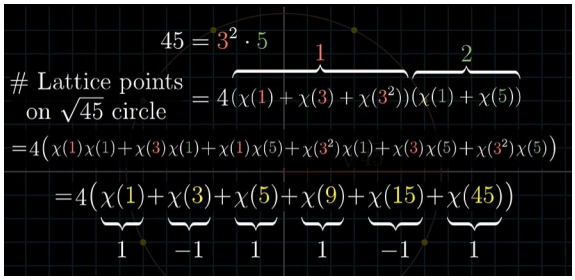
（乘以 1， -1， i，-i 四个

半径为 $\sqrt{n}$ 的圆上的格点数，先将 n 分解质因数，对每个不是高斯素数的数分解成共轭的高斯素数，分配数比指数多 1，指数是偶数的话，有一种方法分配，不然就没有格点

2 = (1+ i)(1 + i) ，但是这对数格点数没有影响，因为要乘-i。

$$引入f(x) = \begin{cases} 1, x为素数x = 4n + 1 \\ -1, x为素数x = 4n + 3 \\ 0, x为偶数 \end{cases}$$

它是一个周期函数，同时是一个积性函数，



再来看这个问题，

$$45 = 3^2 \times 5 \\ 半径为 \sqrt{45} 圆上格点数问题 = 4 \times (f(1)+f(3)+f(3^2)) \times(f(1)+f(5))\\ =4 \times (f(1)+f(3)+f(5)+f(9)+f(15)+f(45))$$

最后转化为 45 的所有约数

$$f(x) = \begin{cases} 1 ,x 为素数 x = 4n+1 \\ -1, x 为素数 x = 4n+3 \\ 0, x 为偶数\\ \end{cases}\\ 半径为\sqrt { n}的圆上的格点数（二维坐标轴 xy 都为整数的点）是 4 \times \sum_{d|n}f(d)$$

博弈拾遗

SG 定理：

mex(minimal excludant)运算，表示最小的不属于这个集合的非负整数。例如 mex{0,1,2,4}=3、mex{2,3,5}=0、mex{}=0。
Sprague-Grundy 定理（SG 定理）：游戏和的 SG 函数等于各个游戏 SG 函数的 Nim 和。这样就可以将每一个子游戏分而治之，从而简化了问题。而 Bouton 定理就是 Sprague-Grundy 定理在 Nim 游戏中的直接应用，因为单堆的 Nim 游戏 SG 函数满足 SG(x) = x。

Nimk：

普通的 NIM 游戏是在 n 堆石子中每次选一堆，取任意个石子，而 NIMK 游戏是在 n 堆石子中每次选择 k 堆，1<=k<=n，从这 k 堆中每堆里都取出任意数目的石子，取的石子数可以不同，其他规则相同。
对于普通的 NIM 游戏，我们采取的是对每堆的 SG 值进行异或，异或其实就是对每一个 SG 值二进制位上的数求和然后模 2，比如说 3^5 就是 011+101=112，然后对每一位都模 2 就变成了 110，所以 3^5=6。而 NIMK 游戏和 NIM 游戏的区别就在于模的不是 2，如果是取 k 堆，就模 k+1，所以取 1 堆的普通 NIM 游戏是模 2。当 k=2 时,3^5→011+101=112，对每一位都模 3 之后三位二进制位上对应的数仍然是 1，1，2。那么当且仅当每一位二进制位上的数都是 0 的时候，先手必败，否则先手必胜。

anti_nim

描述

和最普通的 Nim 游戏相同，不过是取走最后一个石子的人输。

先手必胜条件

以下两个条件满足其一即可：

1. 所有堆的石子个数=1，且异或和=0（其实这里就是有偶数堆的意思）。
2. 至少存在一堆石子个数>1，且异或和≠0。

卡特兰

卡特兰数 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012,...

$$C_n = \frac{1}{n+1}C_{2n}^n = C_{2n}^n - C_{2n}^{n-1}$$

$$C_n = \frac{1}{n+1}\sum_{i=0}^{n}(C_n^i)^2$$

$$C_n = \frac{4n-2}{n+1}C_{n-1}(C_0 = 1)$$

$$C_{n+1} = \sum_{i=0}^{n} C_i\, C_{n-i}(C_0 = 1)$$

超级卡特兰数 1, 1, 3, 11, 45, 197, 903, 4279, 20793, 103049,...（从第 0 项开始）

$$F_n * (n+1) = (6 * n - 3) * F_{n-1} - (n-2) * F_{n-2}$$

大施罗德数(OEIS A006318)1, 2, 6, 22, 90, 394, 1806, 8558, 41586, 206098,...

超级卡特兰数的两倍（除第一项）

卡特兰三角

卡特兰三角

卡特兰数：由 n 个 X 和 n 个 Y 组成的一个序列中，满足**所有前缀中 Y 出现的次数不超过 X 出现的次数**的序列的个数

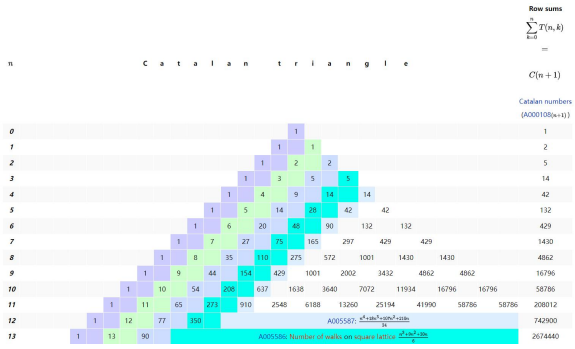
$$C_n = \frac{1}{n+1}\binom{2n}{n}$$

卡特兰三角：由 n 个 X 和 k 个 Y 组成的一个序列，满足**所有前缀中 Y 出现的次数-X 出现的次数小于 m** 的序列的个数

$$C_m(n,k) = \begin{cases} \binom{n+k}{k}, 0 \le k < m \\ \binom{n+k}{k} - \binom{n+k}{k-m}, m \le k \le n+m-1 \\ 0, k > n+m-1 \end{cases}$$

卡特兰三角（OEIS）：$T(n,k) = T(n-1,k) + T(n,k-1)$

$$T(n+1,n+1) = \sum_{k=0}^{n} T(n,k)$$



原根

```
#include<bits/stdc++.h>

using namespace std;

ll qpow(ll n, ll m, ll p) {
    ll ans = 1;
    n %= p;
    for (; m; m >>= 1, n = n * n % p)
        if (m & 1)ans = ans * n % p;
    return ans;
}

ll pri[N], tot;

ll getRoot(ll p) //求质数 p 的最小原根
{
    tot = 0;
    ll n = p - 1, sq = sqrt(p + 0.5);
    for (ll i = 2; i <= sq; i++)
        if (n % i == 0) {
            pri[tot++] = i;
            while (n % i == 0)n /= i;
        }
    if (n > 1)pri[tot++] = n;
    for (ll g = 2; g <= p - 1; g++) //试探每一个 g 是否原根
    {
        ll flag = 1;
        for (ll i = 0; i < tot; i++)
            if (qpow(g, (p - 1) / pri[i], p) == 1) {
                flag = 0;
                break;
            }
```

```
        if (flag)return g;
    }
    return -1; //没有原根
}
```

快速幂

```
ll qpow(ll a, ll b) {
    ll ans = 1;
    while (b) {
        if (b & 1) ans = (ans * a) % mod;
        a = (a * a) % mod;
        b >>= 1;
    }
    return ans;
}
```

扩展欧拉定理

用于在底数与模数不互质的情况下将质数降将至与模数同阶大小，从而使用快速幂

$$a^c = \begin{cases} a^{c \mod \phi(m)} , \gcd(a, m) = 1 \\ a^c, gcd(a, m)\not = 1 \and c < \phi(m) \\ a^{c \mod \phi(m) + \phi(m)}, gcd(a, m)\not = 1 \and c \ge \phi(m) \end{cases}$$

证明以及引理：

欧拉定理： $a^{\phi(m)} \equiv 1 \pmod m$

证明欧拉：记 x_i 为第 i 个与 m 互质的数，则小于 m 的范围内共有$\phi(m)$个这样的数

$p_i = a \times x_i$

$\triangle : \{p_i\}$ 两两不同余且与 m 互质，$\{x_i\}$ 两两不同余

所有p_i 的模 m 的集合与$\{x_i\}$ 相等 $\Rightarrow$ 他们的积模 m 相等

$$\Rightarrow \prod_{i=1}^{\phi m} p_i = a^{\phi(m)} \prod_{i=1}^{\phi(m)} x_i = \prod_{i=1}^{\phi(m)} x_i \pmod m$$

扩展欧拉：

$$a^c = \begin{cases} a^{c \mod \phi(m)} , \gcd(a, m) = 1 \\ a^c, gcd(a, m)\not = 1 \and c < \phi(m) \\ a^{c \mod \phi(m) + \phi(m)}, gcd(a, m)\not = 1 \and c \ge \phi(m) \end{cases}$$

证明扩展欧拉(3)：

1. $\phi(p^r) = (p-1) \times p^r$ ，P 为质数
2. $\exist a, b, x, y , s.t. x^a \times y^b = k, 都有 a， b\le \phi(k) $
3. $\exist r \le c ,s.t. a^{\phi(m)+r} \equiv a ^r (\mod m)$

证明其中 3：$m = t \times a^r$，其中 $gcd(a, t) = 1$

又 ϕ 是一个积性函数，故 $\phi(t) \mid \phi(m)$

$$a^{\phi(t)} \equiv 1 \pmod t \Rightarrow a^{\phi(m)} \equiv 1 \pmod t$$

两边同乘以$a^r \Rightarrow a^{\phi(m)+r} \equiv a^r \pmod m$

根据 2，$r \le \phi(m)$ 又 $c \ge \phi(m)$ ，得证

$$a^c \equiv a^{c-r+r} \equiv a^{c-r+\phi(m)+r} \equiv a^{c+\phi(m)} \pmod m$$

扩欧求逆元

```
#include <bits/stdc++.h>
using namespace std;

typedef  long long ll;

void extgcd(ll a,ll b,ll& d,ll& x,ll& y){
    if(!b){ d=a; x=1; y=0;}
    else{ extgcd(b,a%b,d,y,x); y-=x*(a/b); }
}

ll inverse(ll a,ll n){
    ll d,x,y;
    extgcd(a,n,d,x,y);
    return d==1?(x+n)%n:-1;
}

int main(){
          int x, y;
          //cin >> x >> y;
          while(1){
                    cin >> x >> y;
                    cout << inverse(x, y) << endl;
          }
          //cout << inverse(x, y) << endl;
}
```

数学知识

数学知识的一些范围（?

1～n 的质数个数

$\frac{n}{l_n n}$

1～2e9 中拥有最多约数个数的数拥有的约数个数

约 1600

n 个不同的点可以构成 n^{n-2} 棵不同的树

判断一个数是否为 11 的倍数

奇偶位置上的数位和的差是否为 11 的倍数

平方前缀和

$$\frac{n \times (n+1) \times (2 \times n+1)}{6}$$

立方前缀和

$$\left(\frac{n \times (n+1)}{2}\right)^2$$

库默尔定理

设 m,n 为正整数，p 为素数，则C_{m+n}^m 含 p 的幂次等于 m+n 在 p 进制下的进位次数

原根存在定理

一个数 m 存在原根当且仅当$m = 2,4,p^\alpha,2p^\alpha$，其中 p 为奇素数，$\alpha \in N^*$

整除分块（向上向下取整）

```
int x;
scanf("%d",&x);
int ans1=0,ans2=0;
//向下取整
for(int l=1,r;l<=x;l=r+1){
    int m=x/l;
    r=x/m;
    ans1+=(r-l+1)*m;
}
//向上取整
```

```
int R=1e5;
for(int l=1,r;l<=R;l=r+1){
    int m=(x+l-1)/l;
    r=m!=1?(x-1)/(m-1):R;
    ans2+=(r-l+1)*m;
}
```

格雷码

```
int gray_encode(int num) {
    return num ^ (num >> 1);
}

int gray_decode(int num) {
    int head;
    if (!num) return 0;
    head = 1 << int(log(num) / log(2));
    return head + gray_decode((num ^ head) ^ (head >> 1));
}
```

欧拉筛（素数）

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
const int N = 1000005;
int phi[N], prime[N], cnt;
bool st[N];

void get_eulers() {
    phi[1] = 1;
    for (int i = 2; i < N; i++) {
        if (!st[i]) {
            prime[cnt++] = i;
            phi[i] = i - 1;
        }
        for (int j = 0; prime[j] * i < N; j++) {
            st[prime[j] * i] = 1;
            if (i % prime[j] == 0) {
                phi[prime[j] * i] = phi[i] * prime[j];
                break;
            }
            phi[prime[j] * i] = phi[i] * (prime[j] - 1);
        }
    }
}

int main() {
    get_eulers();
    ll n;
    cin >> n;
    ll ans = 0;
    for (int i = 1; i <= n; i++) ans += phi[i];
    cout << ans;
}
```

欧拉筛（莫比乌斯）

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const int N = 1e5 + 10;

bool vis[N];
ll prime[N], mu[N];

void init_mu() {
    ll cnt = 0;
    mu[1] = 1;
    for (ll i = 2; i < N; i++) {
        if (!vis[i]) {
            prime[cnt++] = i;
            mu[i] = -1;
        }
        for (ll j = 0; j < cnt && i * prime[j] < N; j++) {
            vis[i * prime[j]] = 1;
            if (i % prime[j] == 0) {
                mu[i * prime[j]] = 0;
                break;
            } else { mu[i * prime[j]] = -mu[i]; }
        }
    }
}

int main() {
    init_mu();
}
```

欧拉降幂

~~不知道它有什么用毕竟已经有快速幂了~~

这里有一张图可以很好的说明欧拉降幂是什么

$$a^b \equiv \begin{cases} a^{b\%\varphi(n)} (\bmod\ n) & n,a\text{互质} \\ a^b (\bmod\ n) & b < \varphi(n) \\ a^{b\%\varphi(n)+\varphi(n)} (\bmod\ n) & b \geq \varphi(n) \end{cases}$$

```
//其实只是想试一下 markdown 怎么用
//假装这里有代码
```

然后下面这个是用 $\LaTeX$公式写的

$$a^b \equiv \begin{cases} a^{b\%\varphi(n)} (\mathrm{mod} n) & n,a\text{互质} \\ a^b (\mathrm{mod} n) & b < \varphi(n) \\ a^{b\%\varphi(n)+\varphi(n)} (\mathrm{mod} n) & b \geq \varphi(n) \end{cases}$$

组合数

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const ll mod = 1e9 + 7;
const ll maxn = 3e4 + 5;
ll inv[maxn], fac[maxn];

ll qpow(ll a, ll b) {
    ll ans = 1;
    while (b) {
        if (b & 1) ans = (ans * a) % mod;
        a = (a * a) % mod;
        b >>= 1;
    }
    return ans;
}

ll c(ll n, ll m) {
    if (n < 0 || m < 0 || n < m) return 0;
    return fac[n] * inv[n - m] % mod * inv[m] % mod;
}

void init() {
    fac[0] = 1;
    for (int i = 1; i < maxn; i++) {
        fac[i] = fac[i - 1] * i % mod;
    }
    inv[maxn - 1] = qpow(fac[maxn - 1], mod - 2);
    for (ll i = maxn - 2; i >= 0; i--) {
        inv[i] = (inv[i + 1] * (i + 1)) % mod;
    }
}
```

莫比乌斯反演

莫比乌斯函数

$$对n进行因数分解：n = P_1^{\alpha_1}P_2^{\alpha_2}...P_k^{\alpha_k}，则\mu(n) = \begin{cases} 1, n = 1 \\ 0, \forall \alpha_i \geq 2 \\ \pm 1, (-1)^k \end{cases}$$

n 的所有约数的莫比乌斯的和

$$S(n) = \sum_{d|n} \mu(d) = \begin{cases} 1, n = 1 \\ 0, else \end{cases}$$

反演

$$(一般不用)1.若F(n) = \sum_{d|n} f(d)，则f(n) = \sum_{d|n} \mu(d)F\left(\frac{n}{d}\right)$$

$$(\surd)2.若F(n) = \sum_{n|d} f(d)，则f(n) = \sum_{n|d} \mu\left(\frac{d}{n}\right)F(d)$$

构造[$MathProcessingError$]$F(n)和f(n)$ 使 f(n)为目标，F(n)好求

1

求满足$a \leq x \leq b, c \leq y \leq d$ 且 gcd(x, y) = k 的 xy 的对数

$$F(n) = gcd(x, y) = n的倍数的xy的对数$$

$$f(n) = gcd(x, y) = n的xy的对数$$

```
#include <bits/stdc++.h>
using namespace std;

typedef long long ll;

const int N = 50010;

ll primes[N], mu[N], sum[N], cnt;
bool st[N];

void init() {
          mu[1] = 1;

          for(int i = 2; i < N; ++ i) {
                    if(!st[i]) {
                              primes[cnt ++] = i;
                              mu[i] = -1;
                    }

                    for(int j = 0; primes[j] * i < N; ++ j)
 {
                              st[primes[j] * i] = 1;
                              if(i % primes[j] == 0) brea
k;
                              mu[primes[j] * i] = -mu[i];

                    }
          }

          for(int i = 1; i < N; ++ i) {
                    sum[i] = sum[i - 1] + mu[i];
          }
}

ll g(ll n, ll x) {
          return n / (n / x);
}

ll f (int a, int b, int k) {
          a = a / k, b = b / k;
          ll res = 0;

          ll n = min(a, b);

          for(ll l = 1, r; l <= n; l = r + 1) {
                    r = min(n, min(g(a, l), g(b, l)));
                    res += (sum[r] - sum[l - 1]) * (a / l)
 * (b / l);
          }

          return res;
}

int main() {
          ios::sync_with_stdio(0); cin.tie(0); cout.tie
(0);

          init();

          int T;
          cin >> T;
          while(T --) {
                    int a, b, c, d, k;
                    cin >> a >> b >> c >> d >> k;
                    cout << f(b, d, k) - f(a - 1, d, k) -
f(b, c - 1, k)
                                                  + f(a - 1, c - 1,
 k) << endl;
          }

          return 0;
}
```

2

求$\sum_{i=1}^{N}\sum_{j=1}^{M} d(ij)$

// $d(ij) = \sum_{x|i}\sum_{y|j}[(x, y) = 1]$

$$F(n) = \sum_{i=1}^{N}\sum_{j=1}^{M}\sum_{x|i}\sum_{y|j}[n|(x, y)]$$

$$f(n) = \sum_{i=1}^{N}\sum_{j=1}^{M}\sum_{x|i}\sum_{y|j}[(x, y) = n]$$

$$F(n) = \sum_{i=1}^{N}\sum_{j=1}^{M}\sum_{x|i}\sum_{y|j}[n|(x, y)] = \sum_{x=1}^{N}\sum_{y=1}^{M}\lfloor\frac{N}{x}\rfloor\lfloor\frac{M}{y}\rfloor[n|(x, y)]$$

$$= \sum_{x'}^{\frac{N}{n}}\sum_{y'}^{\frac{M}{n}}\lfloor\frac{N}{x'n}\rfloor\lfloor\frac{M}{y'n}\rfloor$$

两次整数分块

```
#include <bits/stdc++.h>
using namespace std;

typedef long long ll;
const int N = 50010;

int primes[N], cnt, mu[N], sum[N], h[N];
bool st[N];

inline int g(int n, int x) {
          return n / (n / x);
}

void init() {
          mu[1] = 1;
          for(int i = 2; i < N; ++i) {
                    if(!st[i]){
                              primes[cnt++] = i;
                              mu[i] = -1;
                    }
                    for(int j = 0; primes[j] * i < N; ++j)
 {
```

```
                                st[primes[j] * i] = 1;
                                if(i % primes[j] == 0) brea
k;
                                mu[primes[j] * i] = -mu[i];
                        }


                }

                for(int i = 1; i < N; ++ i) {
                        sum[i] = sum[i - 1] + mu[i];
                }

                for(int i = 1; i < N; ++i) {
                        for(int l = 1, r; l <= i; l = r + 1) {
                                r = min(i, g(i, l));
                                h[i] += (r - l + 1) * (i /
l);
                        }
                }
}

int main() {
        //ios::sync_with_stdio(0); cin.tie(0); cout.tie
(0);
        init();

        int T;
        scanf("%d", &T);
        while(T--) {
                int n, m;
                scanf("%d %d", &n, &m);
                ll res = 0;
                int k = min(n, m);
                for(int l = 1, r; l <= k; l = r + 1) {
                        r = min(k, min(g(n, l), g(m,
 l)));
                        res += (ll)(sum[r] - sum[l
- 1]) * h[n / l] * h[m / l];
                }
            printf("%lld\n", res);
        }

        return 0;
}
```

逆元线性递推 inv 阶乘逆元组合数

```
ll fac[N];// n!
ll invfac[N]; // n!的inv
ll invn[N]; //n 的inv

inline void init() {
    fac[0] = fac[1] = invfac[0] = invfac[1] = invn[0] = in
vn[1] = 1;
    for (int i = 2; i < N; ++i) {
        fac[i] = fac[i - 1] * i % mod;
        invn[i] = (mod - mod / i) * invn[mod % i] % mod;
        invfac[i] = invfac[i - 1] * invn[i] % mod;
    }
}

ll C(ll up, ll down) {
    if (up > down) return 0;
    if (up < 0 || down < 0) return 0;
    ll res = fac[down];
    res = res * invfac[down - up] % mod;
    res = res * invfac[up] % mod;
    return res;
}

//先init
```

杂项

fread 快读

```
#include <bits/stdc++.h>
using namespace std;

char next_char() {
        static char buf[1 << 20], *first, *last;
        if(first == last) {
                last = buf + fread(buf, 1, 1 << 20, st
din);
                first = buf;
        }
        return first == last ? EOF : *first ++;
}

inline int read(){
        int x = 0, w = 0; char ch = 0;
        while(!isdigit(ch)) {w |= ch == '-'; ch = next_c
har(); }
        while(isdigit(ch)) {x = (x << 3) + (x << 1) + (c
h ^ 48), ch = next_char(); }
        return w ? -x : x;
}

int main(){
        freopen("1.txt", "r", stdin); // ▯▯▯▯▯▯▯ʰ
▯▯ʰ▯▯Çᵹ▯▯aaa
        int T;
        cin >> T;
        while(T --){
                int x = read();
                cout << x << endl;
        }
}
```

int128 输出

```
inline void print(__int128 x) {
    if (x < 0) {
        putchar('-');
        x = -x;
    }
    if (x > 9)
        print(x / 10);
    putchar(x % 10 + '0');
}
```

mt19937

```
mt19937
#include <random>
#include <iostream>

int main()
{
    std::random_device rd;  //获取随机数种子
    std::mt19937 gen(rd()); //Standard mersenne_twister_en
gine seeded with rd()
    std::uniform_int_distribution<> dis(0, 9);

    for (int n = 0; n<20; ++n)
        std::cout << dis(gen) << ' ';
    std::cout << '\n';
    system("pause");
    return 0;
}

//可能的结果：7 2 2 1 4 1 4 0 4 7 2 1 0 9 1 9 2 3 5 1
```

doule ： std::uniform*real*distribution<> dis(0, 9);

```
#include <iostream>
#include <chrono>
#include <random>
using namespace std;
int main()
{
        // 随机数种子
        unsigned seed = std::chrono::system_clock::now
().time_since_epoch().count();
        mt19937 rand_num(seed);  // 大随机数
        uniform_int_distribution<long long> dist(0, 100
0000000);  // 给定范围
        cout << dist(rand_num) << endl;
```

```
        return 0;
}
```

注意： 代码中的 rand_num 和 dist 都是自己定义的对象，不是系统的。

洗牌算法

```
#include <random>
#include <algorithm>
#include <iterator>
#include <iostream>

int main()
{
    std::vector<int> v = { 1, 2, 3, 4, 5, 6, 7, 8, 9, 10 };

    std::random_device rd;
    std::mt19937 g(rd());

    std::shuffle(v.begin(), v.end(), g);

    std::copy(v.begin(), v.end(), std::ostream_iterator<in
t>(std::cout, " "));
    std::cout << "\n";

    system("pause");
    return 0;
}
```

快读 read

```
inline int read(){
    int X=0,w=0;char ch=0;
    while(!isdigit(ch)){w|=ch=='-';ch=getchar();}
    while(isdigit(ch))X=(X<<3)+(X<<1)+(ch^48),ch=getchar
();
    return w?-X:X;
}
```

整体二分

```
ll bit[N];

void add_bit(ll k, ll a) {
    while (k < N) {
        bit[k] = bit[k] + a;
        k += k & -k;
    }
}

ll query_bit(ll k) {
    ll ans = 0;
    while (k) {
        ans = ans + bit[k];
        k -= k & -k;
    }
    return ans;
}

struct node {
    ll x, y, k, id, type;
};
node q[N], q1[N], q2[N];
ll ans[N], now[N], tot, totx;

void solve(ll l, ll r, ll ql, ll qr) {
    if (ql > qr) return;
    if (l == r) {
        for (ll i = ql; i <= qr; i++) {
            if (q[i].type == 2) {
                ans[q[i].id] = l;
            }
        }
        return;
    }
    ll mid = (l + r) >> 1;
    ll cq1 = 0, cq2 = 0;
    for (ll i = ql; i <= qr; i++) {
        if (q[i].type == 1) {
            if (q[i].y <= mid) {
                add_bit(q[i].x, q[i].k);
                q1[++cq1] = q[i];
            } else {
                q2[++cq2] = q[i];
            }
        } else {
            ll sum = query_bit(q[i].y) - query_bit(q[i].x -
1);
            if (sum >= q[i].k) {
                q1[++cq1] = q[i];
            } else {
                q2[++cq2] = q[i];
                q2[cq2].k -= sum;
            }
        }
    }
    for (ll i = 1; i <= cq1; i++) if (q1[i].type == 1) add_
bit(q1[i].x, -q1[i].k);
    for (ll i = 1; i <= cq1; i++) q[ql + i - 1] = q1[i];
    for (ll i = 1; i <= cq2; i++) q[ql + cq1 + i - 1] = q2
[i];
    solve(l, mid, ql, ql + cq1 - 1);
    solve(mid + 1, r, ql + cq1, qr);

}

void init() {
    totx = 0;
    tot = 0;
    memset(bit, 0, sizeof bit);
}
```

朝鲜大哥快读

```
#define FI(n) FastIO::read(n)
#define FO(n) FastIO::write(n)
#define Flush FastIO::Fflush()
//程序末尾写上   Flush;

namespace FastIO {
    const int SIZE = 1 << 16;
    char buf[SIZE], obuf[SIZE], str[60];
    int bi = SIZE, bn = SIZE, opt;
    double D[] = {0.1, 0.01, 0.001, 0.0001, 0.00001, 0.000
001, 0.0000001, 0.00000001, 0.000000001, 0.0000000001};

    int read(char *s) {
        while (bn) {
            for (; bi < bn && buf[bi] <= ' '; bi++);
            if (bi < bn)
                break;
            bn = fread(buf, 1, SIZE, stdin);
            bi = 0;
        }
        int sn = 0;
        while (bn) {
            for (; bi < bn && buf[bi] > ' '; bi++)
                s[sn++] = buf[bi];
            if (bi < bn)
                break;
            bn = fread(buf, 1, SIZE, stdin);
            bi = 0;
        }
        s[sn] = 0;
        return sn;
    }

    bool read(int &x) {
        int n = read(str), bf = 0;
        if (!n)
            return 0;
        int i = 0;
        if (str[i] == '-')
            bf = 1, i++;
        else if (str[i] == '+')
            i++;
        for (x = 0; i < n; i++)
            x = x * 10 + str[i] - '0';
        if (bf)
            x = -x;
        return 1;
    }

    bool read(long long &x) {
        int n = read(str), bf;
```

```
    if (!n)
        return 0;
    int i = 0;
    if (str[i] == '-')
        bf = -1, i++;
    else
        bf = 1;
    for (x = 0; i < n; i++)
        x = x * 10 + str[i] - '0';
    if (bf < 0)
        x = -x;
    return 1;
}

void write(int x) {
    if (x == 0)
        obuf[opt++] = '0';
    else {
        if (x < 0)
            obuf[opt++] = '-', x = -x;
        int sn = 0;
        while (x)
            str[sn++] = x % 10 + '0', x /= 10;
        for (int i = sn - 1; i >= 0; i--)
            obuf[opt++] = str[i];
    }
    if (opt >= (SIZE >> 1)) {
        fwrite(obuf, 1, opt, stdout);
        opt = 0;
    }
}

void write(long long x) {
    if (x == 0)
        obuf[opt++] = '0';
    else {
        if (x < 0)
            obuf[opt++] = '-', x = -x;
        int sn = 0;
        while (x)
            str[sn++] = x % 10 + '0', x /= 10;
        for (int i = sn - 1; i >= 0; i--)
            obuf[opt++] = str[i];
    }
    if (opt >= (SIZE >> 1)) {
        fwrite(obuf, 1, opt, stdout);
        opt = 0;
    }
}

void write(unsigned long long x) {
    if (x == 0)
        obuf[opt++] = '0';
    else {
        int sn = 0;
        while (x)
            str[sn++] = x % 10 + '0', x /= 10;
        for (int i = sn - 1; i >= 0; i--)
            obuf[opt++] = str[i];
    }
    if (opt >= (SIZE >> 1)) {
        fwrite(obuf, 1, opt, stdout);
        opt = 0;
    }
}

void write(char x) {
    obuf[opt++] = x;
    if (opt >= (SIZE >> 1)) {
        fwrite(obuf, 1, opt, stdout);
        opt = 0;
    }
}

void Fflush() {
    if (opt)
        fwrite(obuf, 1, opt, stdout);
    opt = 0;
}
}; // namespace FastIO
```

枚举子集

```
cin >> n;
for (int s = n; s; s = (s - 1) & n) {
    cout << bitset<8>(s) << endl;
}
```

模拟退火

“优化的随机算法”

连续函数找区间最优

// 找一个点，与平面中的 n 个点的距离和最近

//进行多次模拟退火避免局部最大值

```
#include <bits/stdc++.h>
#include <ctime>
using namespace std;

const int maxn = 110;

int n;

#define x first
#define y second

typedef pair<double, double> PDD;

PDD q[maxn];
double ans = 1e8;

double rand(double l, double r) {
    return (double) rand() / RAND_MAX * (r - l) + l;
}

double getDist(PDD a, PDD b) {
    double dx = a.x - b.x;
    double dy = a.y - b.y;
    return sqrt(dx * dx + dy * dy) ;
}

double calc(PDD p) {
    double res = 0;
    for(int i = 0; i < n; ++ i) {
        res += getDist(q[i], p);
    }
    ans = min(ans, res);
    return res;
}

double simulate_anneal() {
    PDD cur(rand(0, 10000), rand(0, 10000)); // 随机一个起点
    for(double T = 1e4; T > 1e-4; T = T * 0.99) { // 初始温度，末态温度，衰减系数，一般调整衰减系数0.999 0.95
        PDD np(rand(cur.x - T, cur.x + T), rand(cur.y - T, cur.y + T)); // 随机新点
        double delta = calc(np) - calc(cur);
        if(exp(-delta / T) > rand(0, 1)) cur = np; //如果新点比现在的点更优，必过去，不然有一定概率过去
    }

}

int main() {
    cin >> n;
    for(int i = 0; i < n; ++ i) {
        cin >> q[i].x >> q[i].y;
    }

    while((double) clock() / CLOCKS_PER_SEC < 0.8) { // 卡时 // 或for（100）
        simulate_anneal();
    }

    cout << (int)(ans + 0.5) << endl;

    return 0;
}
```

// n 个点带权费马点 // 平衡点||吊打 XXX

//n 个二维坐标点，带重物重量，找平衡点

//进行一次模拟退火，但是在局部最大值周围多次跳动（以提高精度

```
#include <cmath>
#include <cstdio>
#include <cstdlib>
#include <ctime>

const int N = 10005;
int n, x[N], y[N], w[N];
double ansx, ansy, dis;

double Rand() { return (double)rand() / RAND_MAX; }
double calc(double xx, double yy) {
  double res = 0;
  for (int i = 1; i <= n; ++i) {
    double dx = x[i] - xx, dy = y[i] - yy;
    res += sqrt(dx * dx + dy * dy) * w[i];
  }
  if (res < dis) dis = res, ansx = xx, ansy = yy;
  return res;
}
void simulateAnneal() {
  double t = 100000;
  double nowx = ansx, nowy = ansy;
  while (t > 0.001) {
    double nxtx = nowx + t * (Rand() * 2 - 1);
    double nxty = nowy + t * (Rand() * 2 - 1);
    double delta = calc(nxtx, nxty) - calc(nowx, nowy);
    if (exp(-delta / t) > Rand()) nowx = nxtx, nowy = nxty;
    t *= 0.97;
  }
  for (int i = 1; i <= 1000; ++i) {
    double nxtx = ansx + t * (Rand() * 2 - 1);
    double nxty = ansy + t * (Rand() * 2 - 1);
    calc(nxtx, nxty);
  }
}
int main() {
  srand(time(0));
  scanf("%d", &n);
  for (int i = 1; i <= n; ++i) {
    scanf("%d%d%d", &x[i], &y[i], &w[i]);
    ansx += x[i], ansy += y[i];
  }
  ansx /= n, ansy /= n, dis = calc(ansx, ansy);
  simulateAnneal();
  printf("%.3lf %.3lf\n", ansx, ansy);
  return 0;
}
```

测试时常用的代码

```
//�ļ���������ʱ�ļ�

#ifdef ONLINE_JUDGE
#else
    freopen("in.txt","r",stdin);
    //freopen("out.txt","w",stdout);
#endif
//���ٴ���stdin/stdout��in.txt/out.txt��"r"��"w"
Ĭ����ð


#include<ctime>
    clock_t ST,ED;
    ST=clock();
    //�������ֵij���
    ED=clock();
    cout<<ED-ST<<"ms"<<endl;


#include<ctime>
#include<cstdlib>
    srand(time(0));//��'��
    rand();//����[0,RAND_MAX]ֵ(int)
```

```
RAND_MAX��cstdlib�e(ɽ)壬h��Ĭ0x7fff(32767)
```

算法基础

占位

线性代数

矩阵类模板_加减乘快速幂

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const ll N = 305;
const ll mod = 998244353;

//矩阵类模板
struct Matrix {
    ll n, m;
    ll a[N][N];

    void set(ll _a, ll _b) {
        n = _a, m = _b;
    }

    Matrix() {
        clear();
    }

    void clear() {
        n = m = 0;
        memset(a, 0, sizeof(a));
    }

    Matrix operator+(const Matrix &b) const {
        Matrix tmp;
        tmp.n = n;
        tmp.m = m;
        for (ll i = 0; i < n; ++i)
            for (ll j = 0; j < m; ++j)
                tmp.a[i][j] = (a[i][j] + b.a[i][j]) % mod;
        return tmp;
    }

    Matrix operator-(const Matrix &b) const {
        Matrix tmp;
        tmp.n = n;
        tmp.m = m;
        for (ll i = 0; i < n; ++i) {
            for (ll j = 0; j < m; ++j)
                tmp.a[i][j] = (a[i][j] - b.a[i][j] + mod) %
 mod;
        }

        return tmp;
    }

    Matrix operator*(const Matrix &b) const {
        Matrix tmp;
        tmp.clear();
        tmp.n = n;
        tmp.m = b.m;
        for (ll i = 0; i < n; ++i)
            for (ll j = 0; j < b.m; ++j)
                for (ll k = 0; k < m; ++k) {
                    tmp.a[i][j] += a[i][k] * b.a[k][j];
                    tmp.a[i][j] %= mod;
                }
        return tmp;
    }

    Matrix get(ll x) {//幂运算
        Matrix E;
        E.clear();
        E.set(n, m);
        for (ll i = 0; i < n; ++i)
            E.a[i][i] = 1;
```

```
        if (x == 0) return E;
        else if (x == 1) return *this;
        Matrix tmp = get(x / 2);
        tmp = tmp * tmp;
        if (x % 2) tmp = tmp * (*this);
        return tmp;
    }

    void exgcd(ll _a, ll _b, ll &x, ll &y) {
        if (!_b)return x = 1, y = 0, void();
        exgcd(_b, _a % _b, y, x);
        y -= x * (_a / _b);
    }

    ll inv(ll p) {
        ll x, y;
        exgcd(p, mod, x, y);
        return (x + mod) % mod;
    }

    Matrix inv() {
        Matrix E = *this;
        ll is[N], js[N];
        for (ll k = 0; k < E.n; k++) {
            is[k] = js[k] = -1;
            for (ll i = k; i < E.n; i++) // 1
                for (ll j = k; j < E.n; j++)
                    if (E.a[i][j]) {
                        is[k] = i, js[k] = j;
                        break;
                    }
            if (is[k] == -1) {
                E.clear();
                return E;
            }
            for (ll i = 0; i < E.n; i++) // 2
                swap(E.a[k][i], E.a[is[k]][i]);
            for (ll i = 0; i < E.n; i++)
                swap(E.a[i][k], E.a[i][js[k]]);
            if (!E.a[k][k]) {
                E.clear();
                return E;
            }
            E.a[k][k] = inv(E.a[k][k]); // 3
            for (ll j = 0; j < E.n; j++)
                if (j != k) // 4
                    (E.a[k][j] *= E.a[k][k]) %= mod;
            for (ll i = 0; i < E.n; i++)
                if (i != k) // 5
                    for (ll j = 0; j < E.n; j++)
                        if (j != k)
                            (E.a[i][j] += mod - E.a[i][k] *
E.a[k][j] % mod) %= mod;
            for (ll i = 0; i < E.n; i++)
                if (i != k) // 就是这里不同
                    E.a[i][k] = (mod - E.a[i][k] * E.a[k][k]
 % mod) % mod;
        }
        for (ll k = E.n - 1; k >= 0; k--) { // 6
            for (ll i = 0; i < E.n; i++)
                swap(E.a[js[k]][i], E.a[k][i]);
            for (ll i = 0; i < E.n; i++)
                swap(E.a[i][is[k]], E.a[i][k]);
        }
        return E;
    }
};
//矩阵模板结束
```

矩阵类模板_稀疏矩阵乘法

```
struct Matrix{
    int n,m;
    int a[maxn][maxn];////
    void clear(){
        n=m=0;
        memset(a,0,sizeof(a));
    }
    Matrix operator * (const Matrix &b) const{
        Matrix tmp;
        tmp.clear();
        tmp.n=n;tmp.m=b.m;
        for (int k=0;k<m;++k){
            for (int i=0;i<n;++i){
                    if(a[i][k]==0) continue;
                    for(int j=0;j<b.m;++j){
                            if(b.a[k][j]==0) continue;
                            tmp.a[i][j]+=a[i][k]*b.a[k]
[j];
                tmp.a[i][j]%=mod;
                                    }
                            }
        }
        return tmp;
    }
};
//稀疏矩阵乘法
```

矩阵行列式

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
const ll mod = 1e9 + 7;
struct Matrix {
    static const ll MAXN = 300;
    ll a[MAXN][MAXN];

    void init() { memset(a, 0, sizeof(a)); }

    ll det(ll n) {
        for (int i = 0; i < n; i++)
            for (int j = 0; j < n; j++) a[i][j] = (a[i][j]
+ mod) % mod;
        ll res = 1;
        for (int i = 0; i < n; i++) {
            if (!a[i][i]) {
                bool flag = false;
                for (int j = i + 1; j < n; j++) {
                    if (a[j][i]) {
                        flag = true;
                        for (int k = i; k < n; k++) {
                            swap(a[i][k], a[j][k]);
                        }
                        res = -res;
                        break;
                    }
                }
                if (!flag) return 0;
            }

            for (int j = i + 1; j < n; j++) {
                while (a[j][i]) {
                    ll t = a[i][i] / a[j][i];
                    for (int k = i; k < n; k++) {
                        a[i][k] = (a[i][k] - t * a[j][k]) %
mod;
                        swap(a[i][k], a[j][k]);
                    }
                    res = -res;
                }
            }
            res *= a[i][i];
            res %= mod;
        }
        return (res + mod) % mod;
    }
} mat;
```

线性基 2

线性基 能表示的线性空间与原向量 能表示的线性空间等价

用高斯消元得到线性基

先输入数组 a[] 中

```
int n, k;
ll a[N];

void getVec() {
```

```
    k = 0;

    for(int i = 62; i >= 0; -- i) {
        for(int j = k; j < n; ++ j) {
            if(a[j] >> i & 1) {
                swap(a[j], a[k]);
                break;
            }
        }
        if(!(a[k] >> i & 1)) continue;
        for(int j = 0; j < n; ++j) {
            if(j != k && (a[j] >> i & 1)) {
                a[j] ^= a[k];
            }
        }
        ++k;
        if(k == n) break;
    }

}
```

这里注意最后的线性基是 a[]中从 0 到 k-1 个，在前的是**高位**

线性基模板

```
//

const int maxbit = 62;                    //maxbit���_�
�

struct L_B{
          ll lba[maxbit];
          L_B(){
        memset(lba, 0, sizeof(lba));
    }

          void Insert(ll val){           //����
        for(int i = maxbit - 1; i >= 0; -- i) // ��ì�λ��
�λ¡
            if(val & (1ll << i)){ //
                if(!lba[i]){
                    lba[i] = val;
                    break;
                }
                val ^= lba[i];
            }
    }
};
//��ц����є�ÿ����valлЇ2���ш��ì�λ���λ¡����
��Ќ�ĵлЇ1ó¿���lba[i]�����ì���lba[i]=x�����
���val=val`xor`lba[i]
//'ðã„��insert
// --------------���h�ǵ��
```

高斯消元

```
#include <iostream>
#include <vector>
using namespace std;
const double eps = 1e-8;
void sway(vector<double>& a, vector<double>& b) {
    vector<double> s;
    for (int i = 0; i < a.size(); i++) {
        s.push_back(a[i]);
    }
    a.clear();
    for (int i = 0; i < b.size(); i++) {
        a.push_back(b[i]);
    }
    b.clear();
    for (int i = 0; i < s.size(); i++) {
        b.push_back(s[i]);
    }
}
vector<double> gauss_jordan(const vector<vector<double> >
& A,
                            const vector<double>& b) {
    int n = A.size();
    vector<vector<double> > B(n, vector<double>(n + 1));
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++) B[i][j] = A[i][j];
    for (int i = 0; i < n; i++) B[i][n] = b[i];
    for (int i = 0; i < n; i++) {
        int pivot = i;
        for (int j = i; j < n; j++) {
            if (abs(B[j][i]) > abs(B[pivot][i])) pivot = j;
        }
        swap(B[i], B[pivot]);
        if (abs(B[i][i]) < eps) return vector<double>();
        for (int j = i + 1; j <= n; j++) B[i][j] /= B[i][i];
        for (int j = 0; j < n; j++) {
            if (i != j) {
                for (int k = i + 1; k <= n; k++) B[j][k] -=
B[j][i] * B[i][k];
            }
        }
    }
    vector<double> x(n);
    for (int i = 0; i < n; i++) x[i] = B[i][n];
    return x;
}
int main() {
    int n, m;
    cin >> n >> m;
    vector<vector<double> > mat(n, vector<double>(m));
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < m; j++) {
            cin >> mat[i][j];
        }
    }
    vector<double> val(n);
    for (int i = 0; i < n; i++) cin >> val[i];
    vector<double> ans = gauss_jordan(mat, val);
    for (int i = 0; i < ans.size(); i++) cout << ans[i] <<
' ';
}
```

组合数学

斯特林数

百度百科讲的超好

第一类斯特林数（无符号第一类）

定义：$\begin{bmatrix} n \\ k \end{bmatrix}$ 表示将 n 个两两不同的元素，划分为 k 个非空圆排列的方案数。

递推式 $\begin{bmatrix} n \\ n \end{bmatrix} = \begin{bmatrix} n-1 \\ k-1 \end{bmatrix} + (n-1)\begin{bmatrix} n-1 \\ k \end{bmatrix}$

升阶函数

$$x^{n\uparrow} = x(x+1)(x+2)\cdots(x+n-1) = \sum_{k=0}^{n} s_u(n,k)\,x^k$$

（每一项系数则为无符号第一类斯特林数，求前 n 项和则为取 x=1）

性质

无符号Stirling数有如下性质：

① $s_u(0,0)=1$

② $s_u(n,0)=0$

③ $s_u(n,n)=1$

④ $s_u(n,1)=(n-1)!$

⑤ $s_u(n,n-1)=C(n,2)$

⑥ $s_u(n,2)=(n-1)!\cdot\sum_{i=1}^{n-1}\frac{1}{i}$

⑦ $s_u(n,n-2)=2\cdot C(n,3)+3\cdot C(n,4)$

⑧ $\sum_{k=0}^{n}s_u(n,k)=n!$ 证明可令升阶函数中的x=1，比较两边系数。

第二类斯特林数

定义：$\{{n\atop k}\}$ 表示将 n 个两两不同的元素，划分为 k 个非空子集的方案数。

递推式 $\{{n\atop k}\}=\{{n-1\atop k-1}\}+k\{{n-1\atop k}\}$

性质

① $S(n,0)=0^n$

② $S(n,1)=1$

③ $S(n,n)=1$

④ $S(n,2)=2^{n-1}-1$

⑤ $S(n,n-1)=C(n,2)$

⑥ $S(n,n-2)=C(n,3)+3\cdot C(n,4)$

⑦ $S(n,3)=\frac{1}{2}(3^{n-1}+1)-2^{n-1}$

⑧ $S(n,n-3)=C(n,4)+10\cdot C(n,5)+15\cdot C(n$

⑨ $\sum_{k=0}^{n}S(n,k)=B_n$ ， B_n 是贝尔数。

通项公式：

$$S(n,m)=\frac{1}{m!}\sum_{k=0}^{m}(-1)^k C(m,k)(m-$$

两类Stirling数之间的关系

两类Stirling数之间的递推式和实际含义很类似，事实上他们之间存在一个互为转置的转化关系：

$$\sum_{k=0}^{n}S(n,k)s(k,m)=\sum_{k=0}^{n}s(n,k)S(k,m)$$

计算几何

zyx 的计算几何

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const int N = 1e6 + 10;

const double eps = 1e-9;
const double PI = acos(-1.0);
const double dinf = 1e99;
const ll inf = 0x3f3f3f3f3f3f3f3f;
struct Line;

struct Point {
    double x, y;

    Point() { x = y = 0; }

    Point(const Line &a);

    Point(const double &a, const double &b) : x(a), y(b) {}

    Point operator+(const Point &a) const {
        return {x + a.x, y + a.y};
    }

    Point operator-(const Point &a) const {
        return {x - a.x, y - a.y};
    }

    Point operator*(const double &a) const {
        return {x * a, y * a};
    }

    Point operator/(const double &d) const {
        return {x / d, y / d};
    }

    bool operator==(const Point &a) const {
        return abs(x - a.x) + abs(y - a.y) < eps;
    }

    // 标准化，转化为膜长为1
    void standardize() {
        *this = *this / sqrt(x * x + y * y);
    }
};


double norm(const Point &p) { return p.x * p.x + p.y * p.y;
 }

//逆时针转 90 度
Point orth(const Point &a) { return Point(-a.y, a.x); }

//两点间距离
double dist(const Point &a, const Point &b) {
    return sqrt((a.x - b.x) * (a.x - b.x) + (a.y - b.y) *
(a.y - b.y));
}

//两点间距离的平方
double dist2(const Point &a, const Point &b) {
    return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y -
 b.y);
}

struct Line {
```

```
    Point s, t;

    Line() {}

    Line(const Point &a, const Point &b) : s(a), t(b) {}

};


struct Circle {
    Point o;
    double r;

    Circle() {}

    Circle(Point P, double R = 0) { o = P, r = R; }
};

//向量的模长
double length(const Point &p) {
    return sqrt(p.x * p.x + p.y * p.y);
}

//线段的长度
double length(const Line &l) {
    Point p(l);
    return length(p);
}

Point::Point(const Line &a) { *this = a.t - a.s; }

istream &operator>>(istream &in, Point &a) {
    in >> a.x >> a.y;
    return in;
}

ostream &operator<<(ostream &out, Point &a) {
    out << fixed << setprecision(10) << a.x << ' ' << a.y;
    return out;
}

//点积
double dot(const Point &a, const Point &b) { return a.x * b.x + a.y * b.y; }

//叉积
double det(const Point &a, const Point &b) { return a.x * b.y - a.y * b.x; }

//正负判断
int sgn(const double &x) { return fabs(x) < eps ? 0 : (x > 0 ? 1 : -1); }

//平方
double sqr(const double &x) { return x * x; }

//将向量 a 逆时针旋转 ang（弧度制）
Point rotate(const Point &a, const double &ang) {
    double x = cos(ang) * a.x - sin(ang) * a.y;
    double y = sin(ang) * a.x + cos(ang) * a.y;
    return {x, y};
}

//点 p 在线段 seg 上，<=0 则包含端点
bool sp_on(const Line &seg, const Point &p) {
    Point a = seg.s, b = seg.t;
    return !sgn(det(p - a, b - a)) && sgn(dot(p - a, p - b)) <= 0;
}

//点 p 在直线 line 上
bool lp_on(const Line &line, const Point &p) {
    Point a = line.s, b = line.t;
    return !sgn(det(p - a, b - a));
}

//凸包，下标从 0 开始，<=0 则凸包中不包含共线点
int andrew(Point *point, Point *convex, int n) {
    sort(point, point + n, [](Point a, Point b) {
        if (a.x != b.x) return a.x < b.x;
        return a.y < b.y;
    });
    int top = 0;
    for (int i = 0; i < n; i++) {
        while ((top > 1) && det(convex[top - 1] - convex[top - 2], point[i] - convex[top - 1]) <= 0)
            top--;
        convex[top++] = point[i];
    }
    int tmp = top;
    for (int i = n - 2; i >= 0; i--) {
        while ((top > tmp) && det(convex[top - 1] - convex[top - 2], point[i] - convex[top - 1]) <= 0)
            top--;
        convex[top++] = point[i];
    }
    if (n > 1) top--;
    return top;
}

//斜率
double slope(const Point &a, const Point &b) { return (a.y - b.y) / (a.x - b.x); }

//斜率
double slope(const Line &a) { return slope(a.s, a.t); }

//两直线的焦点
Point ll_intersection(const Line &a, const Line &b) {
    double s1 = det(Point(a), b.s - a.s), s2 = det(Point(a), b.t - a.s);
    if (sgn(s1) == 0 && sgn(s2) == 0) return a.s;
    return (b.s * s2 - b.t * s1) / (s2 - s1);
}

//两线段交点 p，返回 0 为无交点，2 为交点为端点，1 为相交
int ss_cross(const Line &a, const Line &b, Point &p) {
    int d1 = sgn(det(a.t - a.s, b.s - a.s));
    int d2 = sgn(det(a.t - a.s, b.t - a.s));
    int d3 = sgn(det(b.t - b.s, a.s - b.s));
    int d4 = sgn(det(b.t - b.s, a.t - b.s));
    if ((d1 ^ d2) == -2 && (d3 ^ d4) == -2) {
        p = ll_intersection(a, b);
        return 1;
    }
    if (!d1 && sp_on(a, b.s)) {
        p = b.s;
        return 2;
    }
    if (!d2 && sp_on(a, b.t)) {
        p = b.t;
        return 2;
    }
    if (!d3 && sp_on(b, a.s)) {
        p = a.s;
        return 2;
    }
    if (!d4 && sp_on(b, a.t)) {
        p = a.t;
        return 2;
    }
    return 0;
}

//两向量直接的相对位置关系，含义见英文注释
int ccw(const Point &a, Point b, Point c) {
    b = b - a, c = c - a;
    if (sgn(det(b, c)) > 0) return +1;  // "COUNTER_CLOCKWISE"
    if (sgn(det(b, c)) < 0) return -1; // "CLOCKWISE"
    if (sgn(dot(b, c)) < 0) return +2;      // "ONLINE_BACK"
    if (sgn(norm(b) - norm(c)) < 0) return -2;  // "ONLINE_FRONT"
    return 0;                           // "ON_SEGMENT"
}


//点 p 在线 l 上的投影位置
Point project(const Line &l, const Point &p) {
    Point base(l);
    double r = dot(base, p - l.s) / sqr(length(base));
    return l.s + (base * r);
}

//线段 l 和点 p 的距离
double sp_dist(const Line &l, const Point &p) {
    if (l.s == l.t) return dist(l.s, p);
```

```
    Point x = p - l.s, y = p - l.t, z = l.t - l.s;
    if (sgn(dot(x, z)) < 0)return length(x);//P 距离 A 更近
    if (sgn(dot(y, z)) > 0)return length(y);//P 距离 B 更近
    return abs(det(x, z) / length(z));//面积除以底边长
}

//直线 l 和点 p 的距离
double lp_dist(const Line &l, const Point &p) {
    Point x = p - l.s, y = p - l.t, z = l.t - l.s;
    return abs(det(x, z) / length(z));//面积除以底边长
}

//圆 c 和直线 l 的交点，返回值为交点的数量，ans 为交点位置
int cl_cross(const Circle &c, const Line &l, pair<Point, P
oint> &ans) {
    Point a = c.o;
    double r = c.r;
    Point pr = project(l, a);
    double dis = dist(pr, a);
    double tmp = r * r - dis * dis;
    if (sgn(tmp) == 1) {
        double base = sqrt(max(0.0, r * r - dis * dis));
        Point e(l);
        e.standardize();
        e = e * base;
        ans = make_pair(pr + e, pr - e);
        return 2;
    } else if (sgn(tmp) == 0) {
        ans = make_pair(pr, pr);
        return 1;
    } else return 0;
}

//圆 c 和线段 l 交点个数，下面 cs_cross 用到
int intersectCS(Circle c, Line l) {
    if (sgn(norm(project(l, c.o) - c.o) - c.r * c.r) > 0)
return 0;
    double d1 = length(c.o - l.s), d2 = length(c.o - l.t);
    if (sgn(d1 - c.r) <= 0 && sgn(d2 - c.r) <= 0) return 0;
    if ((sgn(d1 - c.r) < 0 && sgn(d2 - c.r) > 0) || (sgn(d1
 - c.r) > 0 && sgn(d2 - c.r) < 0)) return 1;
    Point h = project(l, c.o);
    if (dot(l.s - h, l.t - h) < 0) return 2;
    return 0;
}

//圆和线段交点，返回交点数量
int cs_cross(Circle c, Line s, pair<Point, Point> &ans) {
    Line l(s);
    int num = cl_cross(c, l, ans);
    int res = intersectCS(c, s);
    if (res == 2) return 2;
    if (num > 1) {
        if (dot(l.s - ans.first, l.t - ans.first) > 0) swap
(ans.first, ans.second);
        ans.second = ans.first;
    }
    return res;
}

//两圆交点，位置关系见注释
int cc_cross(const Circle &cir1, const Circle &cir2, pair<
Point, Point> &ans) {
    const Point &c1 = cir1.o, &c2 = cir2.o;
    const double &r1 = cir1.r, &r2 = cir2.r;
    double x1 = c1.x, x2 = c2.x, y1 = c1.y, y2 = c2.y;
    double d = length(c1 - c2);
    if (sgn(fabs(r1 - r2) - d) > 0) return 0;  //内含
    if (sgn(r1 + r2 - d) < 0) return 4; //相离
    double a = r1 * (x1 - x2) * 2, b = r1 * (y1 - y2) * 2,
c = r2 * r2 - r1 * r1 - d * d;
    double p = a * a + b * b, q = -a * c * 2, r = c * c - b
 * b;

    double cosa, sina, cosb, sinb;
    //One Intersection
    if (sgn(d - (r1 + r2)) == 0 || sgn(d - fabs(r1 - r2)) =
= 0) {
        cosa = -q / p / 2;
        sina = sqrt(1 - sqr(cosa));
        Point p0(x1 + r1 * cosa, y1 + r1 * sina);
        if (sgn(dist(p0, c2) - r2)) p0.y = y1 - r1 * sina;
        ans = pair<Point, Point>(p0, p0);
        if (sgn(r1 + r2 - d) == 0) return 3;    //外切
        else return 1;  //内切
    }
    //Two Intersections
    double delta = sqrt(q * q - p * r * 4);
    cosa = (delta - q) / p / 2;
    cosb = (-delta - q) / p / 2;
    sina = sqrt(1 - sqr(cosa));
    sinb = sqrt(1 - sqr(cosb));
    Point p1(x1 + r1 * cosa, y1 + r1 * sina);
    Point p2(x1 + r1 * cosb, y1 + r1 * sinb);
    if (sgn(dist(p1, c2) - r2)) p1.y = y1 - r1 * sina;
    if (sgn(dist(p2, c2) - r2)) p2.y = y1 - r1 * sinb;
    if (p1 == p2) p1.y = y1 - r1 * sina;
    ans = pair<Point, Point>(p1, p2);
    return 2;   //  相交
}

//点 p 关于直线 l 的对称点
Point lp_sym(const Line &l, const Point &p) {
    return p + (project(l, p) - p) * 2;
}

//返回两向量的夹角
double alpha(const Point &t1, const Point &t2) {
    double theta;
    theta = atan2((double) t2.y, (double) t2.x) - atan2((d
ouble) t1.y, (double) t1.x);
    if (sgn(theta) < 0)
        theta += 2.0 * PI;
    return theta;
}

//【射线法】判断点 A 是否在任意多边形 Poly 以内，下标从 1 开始
（为保险起见，可以在判断前将所有点随机旋转一个角度防止被卡）
int pip(const Point *P, const int &n, const Point &a) {
    int cnt = 0;
    double tmp;
    for (int i = 1; i <= n; ++i) {
        int j = i < n ? i + 1 : 1;
        if (sp_on(Line(P[i], P[j]), a))return 2;//点在多边
形上
        if (a.y >= min(P[i].y, P[j].y) && a.y < max(P[i].y,
 P[j].y))//纵坐标在该线段两端点之间
            tmp = P[i].x + (a.y - P[i].y) / (P[j].y - P[i].
y) * (P[j].x - P[i].x), cnt += sgn(tmp - a.x) > 0;//交点在
A 右方
    }
    return cnt & 1;//穿过奇数次则在多边形以内
}

//判断 AL 是否在 AR 右边
bool pip_convex_jud(const Point &a, const Point &L, const
Point &R) {
    return sgn(det(L - a, R - a)) > 0;//必须严格以内
}

//【二分法】判断点 A 是否在凸多边形 Poly 以内，下标从 0 开始
bool pip_convex(const Point *P, const int &n, const Point
&a) {
    //点按逆时针给出
    if (pip_convex_jud(P[0], a, P[1]) || pip_convex_jud(P
[0], P[n - 1], a)) return 0;//在 P[0_1]或 P[0_n-1]外
    if (sp_on(Line(P[0], P[1]), a) || sp_on(Line(P[0], P[n
 - 1]), a)) return 2;//在 P[0_1]或 P[0_n-1]上
    int l = 1, r = n - 2;
    while (l < r) {//二分找到一个位置 pos 使得 P[0]_A 在 P[0_p
os],P[0_(pos+1)]之间
        int mid = (l + r + 1) >> 1;
        if (pip_convex_jud(P[0], P[mid], a))l = mid;
        else r = mid - 1;
    }
    if (pip_convex_jud(P[l], a, P[l + 1]))return 0;//在 P[p
os_(pos+1)]外
    if (sp_on(Line(P[l], P[l + 1]), a))return 2;//在 P[pos_
(pos+1)]上
    return 1;
}
// 多边形是否包含线段
// 因此我们可以先求出所有和线段相交的多边形的顶点，然后按照 X-Y
坐标排序(X 坐标小的排在前面，对于 X 坐标相同的点，Y 坐标小的排在
前面，
```

```
// 这种排序准则也是为了保证水平和垂直情况的判断正确)，这样相邻的两个点就是在线段上相邻的两交点，如果任意相邻两点的中点也在多边形内，
// 则该线段一定在多边形内。

//【判断多边形 A 与多边形 B 是否相离】
int pp_judge(Point *A, int n, Point *B, int m) {
    for (int i1 = 1; i1 <= n; ++i1) {
        int j1 = i1 < n ? i1 + 1 : 1;
        for (int i2 = 1; i2 <= m; ++i2) {
            int j2 = i2 < m ? i2 + 1 : 1;
            Point tmp;
            if (ss_cross(Line(A[i1], A[j1]), Line(B[i2], B[j2]), tmp)) return 0;//两线段相交
            if (pip(B, m, A[i1]) || pip(A, n, B[i2]))return 0;//点包含在内
        }
    }
    return 1;
}

//【任意多边形 P 的面积】，下标从 0 开始
double area(Point *P, int n) {
    double S = 0;
    for (int i = 0; i < n; i++) S += det(P[i], P[(i + 1) % n]);
    return S * 0.5;
}

//多边形和圆的面积交 ， 下表从 0 开始
double pc_area(Point *p, int n, const Circle &c) {
    if (n < 3) return 0;
    function<double(Circle, Point, Point)> dfs = [&](Circle c, Point a, Point b) {
        Point va = c.o - a, vb = c.o - b;
        double f = det(va, vb), res = 0;
        if (sgn(f) == 0) return res;
        if (sgn(max(length(va), length(vb)) - c.r) <= 0) return f;
        Point d(dot(va, vb), det(va, vb));
        if (sgn(sp_dist(Line(a, b), c.o) - c.r) >= 0) return c.r * c.r * atan2(d.y, d.x);
        pair<Point, Point> u;
        int cnt = cs_cross(c, Line(a, b), u);
        if (cnt == 0) return res;
        if (cnt > 1 && sgn(dot(u.second - u.first, a - u.first)) > 0) swap(u.first, u.second);
        res += dfs(c, a, u.first);
        if (cnt == 2) res += dfs(c, u.first, u.second) + dfs(c, u.second, b);
        else if (cnt == 1) res += dfs(c, u.first, b);
        return res;
    };
    double res = 0;
    for (int i = 0; i < n; i++) {
        res += dfs(c, p[i], p[(i + 1) % n]);
    }
    return res * 0.5;
}

Line Q[N];

//【半平面交】
int judge(Line L, Point a) { return sgn(det(a - L.s, L.t - L.s)) > 0; }//判断点 a 是否在直线 L 的右边
int halfcut(Line *L, int n, Point *P) {
    sort(L, L + n, [](const Line &a, const Line &b) {
        double d = atan2((a.t - a.s).y, (a.t - a.s).x) - atan2((b.t - b.s).y, (b.t - b.s).x);
        return sgn(d) ? sgn(d) < 0 : judge(a, b.s);
    });

    int m = n;
    n = 0;
    for (int i = 0; i < m; ++i)
        if (i == 0 || sgn(atan2(Point(L[i]).y, Point(L[i]).x) - atan2(Point(L[i - 1]).y, Point(L[i - 1]).x)))
            L[n++] = L[i];
    int h = 1, t = 0;
    for (int i = 0; i < n; ++i) {
        while (h < t && judge(L[i], ll_intersection(Q[t], Q[t - 1]))) --t;//当队尾两个直线交点不是在直线 L[i] 上或者左边时就出队
        while (h < t && judge(L[i], ll_intersection(Q[h], Q[h + 1]))) ++h;//当队头两个直线交点不是在直线 L[i] 上或者左边时就出队
        Q[++t] = L[i];

    }
    while (h < t && judge(Q[h], ll_intersection(Q[t], Q[t - 1]))) --t;
    while (h < t && judge(Q[t], ll_intersection(Q[h], Q[h + 1]))) ++h;
    n = 0;
    for (int i = h; i <= t; ++i) {
        P[n++] = ll_intersection(Q[i], Q[i < t ? i + 1 : h]);
    }
    return n;
}

Point V1[N], V2[N];

//【闵可夫斯基和】求两个凸包{P1},{P2}的向量集合{V}={P1+P2}构成的凸包
int mincowski(Point *P1, int n, Point *P2, int m, Point *V) {
    for (int i = 0; i < n; ++i) V1[i] = P1[(i + 1) % n] - P1[i];
    for (int i = 0; i < m; ++i) V2[i] = P2[(i + 1) % m] - P2[i];
    int t = 0, i = 0, j = 0;
    V[t++] = P1[0] + P2[0];
    while (i < n && j < m) V[t] = V[t - 1] + (sgn(det(V1[i], V2[j])) > 0 ? V1[i++] : V2[j++]), t++;
    while (i < n) V[t] = V[t - 1] + V1[i++], t++;
    while (j < m) V[t] = V[t - 1] + V2[j++], t++;
    return t;
}

//【三点确定一圆】向量垂心法
Circle external_circle(const Point &A, const Point &B, const Point &C) {
    Point P1 = (A + B) * 0.5, P2 = (A + C) * 0.5;
    Line R1 = Line(P1, P1 + orth(B - A));
    Line R2 = Line(P2, P2 + orth(C - A));
    Circle O;
    O.o = ll_intersection(R1, R2);
    O.r = length(A - O.o);
    return O;
}

//三角形内接圆
Circle internal_circle(const Point &A, const Point &B, const Point &C) {
    double a = dist(B, C), b = dist(A, C), c = dist(A, B);
    double s = (a + b + c) / 2;
    double S = sqrt(max(0.0, s * (s - a) * (s - b) * (s - c)));
    double r = S / s;

    return Circle((A * a + B * b + C * c) / (a + b + c), r);
}

//动态凸包
struct ConvexHull {

    int op;

    struct cmp {
        bool operator()(const Point &a, const Point &b) const {
            return sgn(a.x - b.x) < 0 || sgn(a.x - b.x) == 0 && sgn(a.y - b.y) < 0;
        }
    };

    set<Point, cmp> s;

    ConvexHull(int o) {
        op = o;
        s.clear();
    }

    inline int PIP(Point P) {
        set<Point>::iterator it = s.lower_bound(Point(P.x, -dinf));//找到第一个横坐标大于 P 的点
```

```
        if (it == s.end())return 0;
        if (sgn(it->x - P.x) == 0) return sgn((P.y - it->y)
 * op) <= 0;//比较纵坐标大小
        if (it == s.begin())return 0;
        set<Point>::iterator j = it, k = it;
        --j;
        return sgn(det(P - *j, *k - *j) * op) >= 0;//看叉姬
1
    }

    inline int judge(set<Point>::iterator it) {
        set<Point>::iterator j = it, k = it;
        if (j == s.begin())return 0;
        --j;
        if (++k == s.end())return 0;
        return sgn(det(*it - *j, *k - *j) * op) >= 0;//看叉
姬
    }

    inline void insert(Point P) {
        if (PIP(P))return;//如果点P已经在凸壳上或凸包里就不
插入了
        set<Point>::iterator tmp = s.lower_bound(Point(P.x,
 -dinf));
        if (tmp != s.end() && sgn(tmp->x - P.x) == 0)s.eras
e(tmp);//特判横坐标相等的点要去掉
        s.insert(P);
        set<Point>::iterator it = s.find(P), p = it;
        if (p != s.begin()) {
            --p;
            while (judge(p)) {
                set<Point>::iterator temp = p--;
                s.erase(temp);
            }
        }
        if ((p = ++it) != s.end()) {
            while (judge(p)) {
                set<Point>::iterator temp = p++;
                s.erase(temp);
            }
        }
    }
} up(1), down(-1);

int PIC(Circle C, Point a) { return sgn(length(a - C.o) -
C.r) <= 0; }//判断点A是否在圆C内
void Random(Point *P, int n) { for (int i = 0; i < n; ++i)
swap(P[i], P[(rand() + 1) % n]); }//随机一个排列
//【求点集P的最小覆盖圆】 O(n)
Circle min_circle(Point *P, int n) {
//  random_shuffle(P,P+n);
    Random(P, n);
    Circle C = Circle(P[0], 0);
    for (int i = 1; i < n; ++i)
        if (!PIC(C, P[i])) {
            C = Circle(P[i], 0);
            for (int j = 0; j < i; ++j)
                if (!PIC(C, P[j])) {
                    C.o = (P[i] + P[j]) * 0.5, C.r = length
(P[j] - C.o);
                    for (int k = 0; k < j; ++k) if (!PIC(C,
P[k])) C = external_circle(P[i], P[j], P[k]);
                }
        }
    return C;
}


int temp[N];

//最近点对
double closest_point(Point *p, int n) {
    function<double(int, int)> merge = [&](int l, int r) {
        double d = dinf;
        if (l == r) return d;
        if (l + 1 == r) return dist(p[l], p[r]);
        int mid = (l + r) >> 1;
        double d1 = merge(l, mid);
        double d2 = merge(mid + 1, r);
        d = min(d1, d2);
        int i, j, k = 0;
        for (i = l; i <= r; i++) {
            if (sgn(abs(p[mid].x - p[i].x) - d) <= 0)
                temp[k++] = i;
        }
        sort(temp, temp + k, [&](const int &a, const int &b)
 {
            return sgn(p[a].y - p[b].y) < 0;
        });
        for (i = 0; i < k; i++) {
            for (j = i + 1; j < k && sgn(p[temp[j]].y - p[t
emp[i]].y - d) <= 0; j++) {
                double d3 = dist(p[temp[i]], p[temp[j]]);
                d = min(d, d3);
            }
        }
        return d;
    };
    sort(p, p + n, [&](const Point &a, const Point &b) {
        if (sgn(a.x - b.x) == 0) return sgn(a.y - b.y) < 0;
        else return sgn(a.x - b.x) < 0;
    });
    return merge(0, n - 1);
}

//圆和点的切线
int tangent(const Circle &c1, const Point &p2, pair<Point,
 Point> &ans) {
    Point tmp = c1.o - p2;
    int sta;
    if (sgn(norm(tmp) - c1.r * c1.r) < 0) return 0;
    else if (sgn(norm(tmp) - c1.r * c1.r) == 0) sta = 1;
    else sta = 2;
    Circle c2 = Circle(p2, sqrt(max(0.0, norm(tmp) - c1.r
* c1.r)));
    cc_cross(c1, c2, ans);
    return sta;
}

//圆和圆的切线
int tangent(Circle c1, Circle c2, vector<Line> &ans) {
    ans.clear();
    if (sgn(c1.r - c2.r) < 0) swap(c1, c2);
    double g = norm(c1.o - c2.o);
    if (sgn(g) == 0) return 0;
    Point u = (c2.o - c1.o) / sqrt(g);
    Point v = orth(u);
    for (int s = 1; s >= -1; s -= 2) {
        double h = (c1.r + s * c2.r) / sqrt(g);
        if (sgn(1 - h * h) == 0) {
            ans.push_back(Line(c1.o + u * c1.r, c1.o + (u +
 v) * c1.r));
        } else if (sgn(1 - h * h) >= 0) {
            Point uu = u * h, vv = v * sqrt(1 - h * h);
            ans.push_back(Line(c1.o + (uu + vv) * c1.r, c2.
o - (uu + vv) * c2.r * s));
            ans.push_back(Line(c1.o + (uu - vv) * c1.r, c2.
o - (uu - vv) * c2.r * s));
        }
    }

    return ans.size();
}

//两圆面积交
double areaofCC(Circle c1, Circle c2) {
    if (c1.r > c2.r) swap(c1, c2);
    double nor = norm(c1.o - c2.o);
    double dist = sqrt(max(0.0, nor));

    if (sgn(c1.r + c2.r - dist) <= 0) return 0;

    if (sgn(dist + c1.r - c2.r) <= 0) return c1.r * c1.r *
PI;

    double val;
    val = (nor + c1.r * c1.r - c2.r * c2.r) / (2 * c1.r * d
ist);
    val = max(val, -1.0), val = min(val, 1.0);
    double theta1 = acos(val);
    val = (nor + c2.r * c2.r - c1.r * c1.r) / (2 * c2.r * d
ist);
    val = max(val, -1.0), val = min(val, 1.0);
    double theta2 = acos(val);
    return (theta1 - sin(theta1 + theta1) * 0.5) * c1.r *
c1.r + (theta2 - sin(theta2 + theta2) * 0.5) * c2.r * c2.r;
```

```
}

//https://onlinejudge.u-aizu.ac.jp/courses/library/4/CGL/all/CGL_4_C
//把凸包切一刀
int convexCut(Point *p, Point *ans, int n, Line l) {
    int top = 0;
    for (int i = 0; i < n; i++) {
        Point a = p[i], b = p[(i + 1) % n];
        if (ccw(l.s, l.t, a) != -1) ans[top++] = a;
        if (ccw(l.s, l.t, a) * ccw(l.s, l.t, b) < 0)
            ans[top++] = ll_intersection(Line(a, b), l);
    }
    return top;
}

//两球体积交
double SphereCross(double d, double r1, double r2) {
    if (r1 < r2) swap(r1, r2);
    if (sgn(d - r1 - r2) >= 0) return 0;
    if (sgn(d + r2 - r1) <= 0) return 4.0 / 3 * PI * r2 * r2 * r2;
    double co = (r1 * r1 + d * d - r2 * r2) / (2.0 * d * r1);
    double h = r1 * (1 - co);
    double ans = (1.0 / 3) * PI * (3.0 * r1 - h) * h * h;
    co = (r2 * r2 + d * d - r1 * r1) / (2.0 * d * r2);
    h = r2 * (1 - co);
    ans += (1.0 / 3) * PI * (3.0 * r2 - h) * h * h;
    return ans;
}
```

几何一些定理（或知识点？

多面体欧拉定理

多面体欧拉定理是指对于简单多面体，其各维对象数总满足一定的数学关系，在三维空间中多面体欧拉定理可表示为：
“顶点数-棱长数+表面数=2”。
简单多面体即表面经过连续变形可以变为球面的多面体。

球体积交和并

```
#include<bits/stdc++.h>
#define fi first
#define sf scanf
#define se second
#define pf printf
#define pb push_back
#define mp make_pair
#define sz(x) ((int)(x).size())
#define all(x) (x).begin(),(x).end()
#define mem(x,y) memset((x),(y),sizeof(x))
#define fup(i,x,y) for(int i=(x);i<=(y);++i)
#define fdn(i,x,y) for(int i=(x);i>=(y);--i)
typedef long long ll;
typedef long double ld;
typedef unsigned long long ull;
typedef std::pair<int,int> pii;
using namespace std;

const ld pi=acos(-1);

ld pow2(ld x){return x*x;}

ld pow3(ld x){return x*x*x;}

ld dis(ld x1,ld y1,ld z1,ld x2,ld y2,ld z2)
{
    return pow2(x1-x2)+pow2(y1-y2)+pow2(z1-z2);
}

ld cos(ld a,ld b,ld c){return (b*b+c*c-a*a)/(2*b*c);}

ld cap(ld r,ld h){return pi*(r*3-h)*h*h/3;} // 球缺体积公式，h 为球缺的高

//2 球体积交
ld sphere_intersect(ld x1,ld y1,ld z1,ld r1,ld x2,ld y2,ld z2,ld r2)
{
    ld d=dis(x1,y1,z1,x2,y2,z2);
    //相离
    if(d>=pow2(r1+r2))return 0;
    //包含
    if(d<=pow2(r1-r2))return pow3(min(r1,r2))*4*pi/3;
    //相交
    ld h1=r1-r1*cos(r2,r1,sqrt(d)),h2=r2-r2*cos(r1,r2,sqrt(d));
    return cap(r1,h1)+cap(r2,h2);
}

//2 球体积并
ld sphere_union(ld x1,ld y1,ld z1,ld r1,ld x2,ld y2,ld z2,ld r2)
{
    ld d=dis(x1,y1,z1,x2,y2,z2);
    //相离
    if(d>=pow2(r1+r2))return (pow3(r1)+pow3(r2))*4*pi/3;
    //包含
    if(d<=pow2(r1-r2))return pow3(max(r1,r2))*4*pi/3;
    //相交
    ld h1=r1+r1*cos(r2,r1,sqrt(d)),h2=r2+r2*cos(r1,r2,sqrt(d));
    return cap(r1,h1)+cap(r2,h2);
}

int main()
{
    double x1,y1,z1,r1,x2,y2,z2,r2;
    sf("%lf%lf%lf%lf%lf%lf%lf%lf",&x1,&y1,&z1,&r1,&x2,&y2,&z2,&r2);
    pf("%.12Lf\n",sphere_union(x1,y1,z1,r1,x2,y2,z2,r2));
    return 0;
}
```

自适应辛普森

```
double f(double x) {
}

double simpson(double l, double r) {
    double mid = (l + r) / 2;
    return (r - l) * (f(l) + 4 * f(mid) + f(r)) / 6;  // 辛普森公式
}

double asr(double l, double r, double EPS, double ans) {
    double mid = (l + r) / 2;
    double fl = simpson(l, mid), fr = simpson(mid, r);
    if (abs(fl + fr - ans) <= 15 * EPS)
        return fl + fr + (fl + fr - ans) / 15;  // 足够相似的话就直接返回
    return asr(l, mid, EPS / 2, fl) +
           asr(mid, r, EPS / 2, fr);  // 否则分割成两段递归求解
}
```

计算几何全家桶

```
#include <bits/stdc++.h>

using namespace std;
typedef long long ll;
const ll N = 1 << 20;
const ll mod = 1e9 + 7;
const double dinf = 1e99;
const int inf = 0x3f3f3f3f;
const ll linf = 0x3f3f3f3f3f3f3f3f;

const double eps = 1e-9;
const double PI = acos(-1.0);

struct Line;

struct Point {
    double x, y;

    Point() { x = y = 0; }

    Point(const Line &a);
```

```
    Point(const double &a, const double &b) : x(a), y(b) {}

    Point operator+(const Point &a) const {
        return {x + a.x, y + a.y};
    }

    Point operator-(const Point &a) const {
        return {x - a.x, y - a.y};
    }

    Point operator*(const double &a) const {
        return {x * a, y * a};
    }

    Point operator/(const double &d) const {
        return {x / d, y / d};
    }

    bool operator==(const Point &a) const {
        return abs(x - a.x) + abs(y - a.y) < eps;
    }

    void standardize() {
        *this = *this / sqrt(x * x + y * y);
    }
};

Point normal(const Point &a) { return Point(-a.y, a.x); }

double dist(const Point &a, const Point &b) {
    return sqrt((a.x - b.x) * (a.x - b.x) + (a.y - b.y) *
(a.y - b.y));
}

double dist2(const Point &a, const Point &b) {
    return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y -
 b.y);
}

struct Line {
    Point s, t;

    Line() {}

    Line(const Point &a, const Point &b) : s(a), t(b) {}

};

struct circle {
    Point o;
    double r;

    circle() {}

    circle(Point P, double R = 0) { o = P, r = R; }
};

double length(const Point &p) {
    return sqrt(p.x * p.x + p.y * p.y);
}

double length(const Line &l) {
    Point p(l);
    return length(p);
}

Point::Point(const Line &a) { *this = a.t - a.s; }

istream &operator>>(istream &in, Point &a) {
    in >> a.x >> a.y;
    return in;
}

double dot(const Point &a, const Point &b) {
    return a.x * b.x + a.y * b.y;
}

double det(const Point &a, const Point &b) {
    return a.x * b.y - a.y * b.x;
}

int sgn(const double &x) { return fabs(x) < eps ? 0 : (x >
 0 ? 1 : -1); }

double sqr(const double &x) { return x * x; }

Point rotate(const Point &a, const double &ang) {
    double x = cos(ang) * a.x - sin(ang) * a.y;
    double y = sin(ang) * a.x + cos(ang) * a.y;
    return {x, y};
}

//点在线段上 <=0 包含端点
bool sp_on(const Line &seg, const Point &p) {
    Point a = seg.s, b = seg.t;
    return !sgn(det(p - a, b - a)) && sgn(dot(p - a, p - b))
 <= 0;
}

bool lp_on(const Line &line, const Point &p) {
    Point a = line.s, b = line.t;
    return !sgn(det(p - a, b - a));
}

//等于不包含共线
int andrew(Point *point, Point *convex, int n) {
    sort(point, point + n, [](Point a, Point b) {
        if (a.x != b.x) return a.x < b.x;
        return a.y < b.y;
    });
    int top = 0;
    for (int i = 0; i < n; i++) {
        while ((top > 1) && det(convex[top - 1] - convex[to
p - 2], point[i] - convex[top - 1]) <= 0)
            top--;
        convex[top++] = point[i];
    }
    int tmp = top;
    for (int i = n - 2; i >= 0; i--) {
        while ((top > tmp) && det(convex[top - 1] - convex
[top - 2], point[i] - convex[top - 1]) <= 0)
            top--;
        convex[top++] = point[i];
    }
    if (n > 1) top--;
    return top;
}

double slope(const Point &a, const Point &b) {
    return (a.y - b.y) / (a.x - b.x);
}

double slope(const Line &a) {
    return slope(a.s, a.t);
}

Point ll_intersection(const Line &a, const Line &b) {
    double s1 = det(Point(a), b.s - a.s), s2 = det(Point
(a), b.t - a.s);
    return (b.s * s2 - b.t * s1) / (s2 - s1);
}

int ss_cross(const Line &a, const Line &b, Point &p) {
    int d1 = sgn(det(a.t - a.s, b.s - a.s));
    int d2 = sgn(det(a.t - a.s, b.t - a.s));
    int d3 = sgn(det(b.t - b.s, a.s - b.s));
    int d4 = sgn(det(b.t - b.s, a.t - b.s));
    if ((d1 ^ d2) == -2 && (d3 ^ d4) == -2) {
        p = ll_intersection(a, b);
        return 1;
    }
    if (!d1 && sp_on(a, b.s)) {
        p = b.s;
        return 2;
    }
    if (!d2 && sp_on(a, b.t)) {
        p = b.t;
        return 2;
    }
    if (!d3 && sp_on(b, a.s)) {
        p = a.s;
        return 2;
    }
    if (!d4 && sp_on(b, a.t)) {
        p = a.t;
        return 2;
    }
    return 0;
```

```
}

Point project(const Line &l, const Point &p) {
    Point base(l);
    double r = dot(base, p - l.s) / sqr(length(base));
    return l.s + (base * r);
}

double sp_dist(const Line &l, const Point &p) {
    if (l.s == l.t) return dist(l.s, p);
    Point x = p - l.s, y = p - l.t, z = l.t - l.s;
    if (sgn(dot(x, z)) < 0)return length(x);//P 距离 A 更近
    if (sgn(dot(y, z)) > 0)return length(y);//P 距离 B 更近
    return abs(det(x, z) / length(z));//面积除以底边长
}

double lp_dist(const Line &l, const Point &p) {
    Point x = p - l.s, y = p - l.t, z = l.t - l.s;
    return abs(det(x, z) / length(z));//面积除以底边长
}

int lc_cross(const Line &l, const Point &a, const double &
r, pair<Point, Point> &ans) {
    int num = 0;
    Point pr = project(l, a);
    double dis = dist(pr, a);
    double tmp = r * r - dis * dis;
    if (sgn(tmp) == 1) num = 2;
    else if (sgn(tmp) == 0) num = 1;
    else return 0;
    double base = sqrt(r * r - dis * dis);
    Point e(l);
    e.standardize();
    e = e * base;
    ans = make_pair(pr + e, pr - e);
    return num;
}

int cc_cross(const Point &c1, const double &r1, const Poin
t &c2, const double &r2, pair<Point, Point> &ans) {
    double x1 = c1.x, x2 = c2.x, y1 = c1.y, y2 = c2.y;
    double d = length(c1 - c2);
    if (sgn(fabs(r1 - r2) - d) > 0) return -1;  //内含
    if (sgn(r1 + r2 - d) < 0) return 0; //相离
    double a = r1 * (x1 - x2) * 2, b = r1 * (y1 - y2) * 2,
c = r2 * r2 - r1 * r1 - d * d;
    double p = a * a + b * b, q = -a * c * 2, r = c * c - b
 * b;

    double cosa, sina, cosb, sinb;
    //One Intersection
    if (sgn(d - (r1 + r2)) == 0 || sgn(d - fabs(r1 - r2)) =
= 0) {
        cosa = -q / p / 2;
        sina = sqrt(1 - sqr(cosa));
        Point p0(x1 + r1 * cosa, y1 + r1 * sina);
        if (sgn(dist(p0, c2) - r2)) p0.y = y1 - r1 * sina;
        ans = pair<Point, Point>(p0, p0);
        return 1;
    }
    //Two Intersections
    double delta = sqrt(q * q - p * r * 4);
    cosa = (delta - q) / p / 2;
    cosb = (-delta - q) / p / 2;
    sina = sqrt(1 - sqr(cosa));
    sinb = sqrt(1 - sqr(cosb));
    Point p1(x1 + r1 * cosa, y1 + r1 * sina);
    Point p2(x1 + r1 * cosb, y1 + r1 * sinb);
    if (sgn(dist(p1, c2) - r2)) p1.y = y1 - r1 * sina;
    if (sgn(dist(p2, c2) - r2)) p2.y = y1 - r1 * sinb;
    if (p1 == p2) p1.y = y1 - r1 * sina;
    ans = pair<Point, Point>(p1, p2);
    return 2;
}

Point lp_sym(const Line &l, const Point &p) {
    return p + (project(l, p) - p) * 2;
}

double alpha(const Point &t1, const Point &t2) {
    double theta;
    theta = atan2((double) t2.y, (double) t2.x) - atan2((d
ouble) t1.y, (double) t1.x);
    if (sgn(theta) < 0)
        theta += 2.0 * PI;
    return theta;
}

int pip(const Point *P, const int &n, const Point &a) {//
【射线法】判断点 A 是否在任意多边形 Poly 以内
    int cnt = 0;
    int tmp;
    for (int i = 1; i <= n; ++i) {
        int j = i < n ? i + 1 : 1;
        if (sp_on(Line(P[i], P[j]), a))return 2;//点在多边
形上
        if (a.y >= min(P[i].y, P[j].y) && a.y < max(P[i].y,
 P[j].y))//纵坐标在该线段两端点之间
            tmp = P[i].x + (a.y - P[i].y) / (P[j].y - P[i].
y) * (P[j].x - P[i].x), cnt += sgn(tmp - a.x) > 0;//交点在
A 右方
    }
    return cnt & 1;//穿过奇数次则在多边形以内
}

bool pip_convex_jud(const Point &a, const Point &L, const
Point &R) {//判断 AL 是否在 AR 右边
    return sgn(det(L - a, R - a)) > 0;//必须严格以内
}

bool pip_convex(const Point *P, const int &n, const Point
&a) {//【二分法】判断点 A 是否在凸多边形 Poly 以内
    //点按逆时针给出
    if (pip_convex_jud(P[0], a, P[1]) || pip_convex_jud(P
[0], P[n - 1], a)) return 0;//在 P[0_1]或 P[0_n-1]外
    if (sp_on(Line(P[0], P[1]), a) || sp_on(Line(P[0], P[n
 - 1]), a)) return 2;//在 P[0_1]或 P[0_n-1]上
    int l = 1, r = n - 2;
    while (l < r) {//二分找到一个位置 pos 使得 P[0]_A 在 P[0_p
os],P[0_(pos+1)]之间
        int mid = (l + r + 1) >> 1;
        if (pip_convex_jud(P[0], P[mid], a))l = mid;
        else r = mid - 1;
    }
    if (pip_convex_jud(P[l], a, P[l + 1]))return 0;//在 P[p
os_(pos+1)]外
    if (sp_on(Line(P[l], P[l + 1]), a))return 2;//在 P[pos_
(pos+1)]上
    return 1;
}
// 多边形是否包含线段
// 因此我们可以先求出所有和线段相交的多边形的顶点，然后按照 X-Y
坐标排序(X 坐标小的排在前面，对于 X 坐标相同的点，Y 坐标小的排在
前面，
// 这种排序准则也是为了保证水平和垂直情况的判断正确)，这样相邻
的两个点就是在线段上相邻的两交点，如果任意相邻两点的中点也在多
边形内，
// 则该线段一定在多边形内。

int pp_judge(Point *A, int n, Point *B, int m) {//【判断多
边形 A 与多边形 B 是否相离】
    for (int i1 = 1; i1 <= n; ++i1) {
        int j1 = i1 < n ? i1 + 1 : 1;
        for (int i2 = 1; i2 <= m; ++i2) {
            int j2 = i2 < m ? i2 + 1 : 1;
            Point tmp;
            if (ss_cross(Line(A[i1], A[j1]), Line(B[i2], B
[j2]), tmp)) return 0;//两线段相交
            if (pip(B, m, A[i1]) || pip(A, n, B[i2]))return
 0;//点包含在内
        }
    }
    return 1;
}

double area(Point *P, int n) {//【任意多边形 P 的面积】
    double S = 0;
    for (int i = 1; i <= n; i++) S += det(P[i], P[i < n ? i
 + 1 : 1]);
    return S / 2.0;
}

Line Q[N];

int judge(Line L, Point a) { return sgn(det(a - L.s, L.t -
```

```
 L.s)) > 0; }//判断点 a 是否在直线 L 的右边
int halfcut(Line *L, int n, Point *P) {//【半平面交】
    sort(L, L + n, [](const Line &a, const Line &b) {
        double d = atan2((a.t - a.s).y, (a.t - a.s).x) - at
an2((b.t - b.s).y, (b.t - b.s).x);
        return sgn(d) ? sgn(d) < 0 : judge(a, b.s);
    });

    int m = n;
    n = 0;
    for (int i = 0; i < m; ++i)
        if (i == 0 || sgn(atan2(Point(L[i]).y, Point(L[i]).
x) - atan2(Point(L[i - 1]).y, Point(L[i - 1]).x)))
            L[n++] = L[i];
    int h = 1, t = 0;
    for (int i = 0; i < n; ++i) {
        while (h < t && judge(L[i], ll_intersection(Q[t],
Q[t - 1]))) --t;//当队尾两个直线交点不是在直线 L[i] 上或者左边
时就出队
        while (h < t && judge(L[i], ll_intersection(Q[h],
Q[h + 1]))) ++h;//当队头两个直线交点不是在直线 L[i] 上或者左边
时就出队
        Q[++t] = L[i];

    }
    while (h < t && judge(Q[h], ll_intersection(Q[t], Q[t
- 1]))) --t;
    while (h < t && judge(Q[t], ll_intersection(Q[h], Q[h
+ 1]))) ++h;
    n = 0;
    for (int i = h; i <= t; ++i) {
        P[n++] = ll_intersection(Q[i], Q[i < t ? i + 1 :
h]);
    }
    return n;
}

Point V1[N], V2[N];

int mincowski(Point *P1, int n, Point *P2, int m, Point *V)
 {//【闵可夫斯基和】求两个凸包{P1},{P2}的向量集合{V}={P1+P2}
构成的凸包
    for (int i = 0; i < n; ++i) V1[i] = P1[(i + 1) % n] - P
1[i];
    for (int i = 0; i < m; ++i) V2[i] = P2[(i + 1) % m] - P
2[i];
    int t = 0, i = 0, j = 0;
    V[t++] = P1[0] + P2[0];
    while (i < n && j < m) V[t] = V[t - 1] + (sgn(det(V1[i],
 V2[j])) > 0 ? V1[i++] : V2[j++]), t++;
    while (i < n) V[t] = V[t - 1] + V1[i++], t++;
    while (j < m) V[t] = V[t - 1] + V2[j++], t++;
    return t;
}

circle getcircle(const Point &A, const Point &B, const Poi
nt &C) {//【三点确定一圆】向量垂心法
    Point P1 = (A + B) * 0.5, P2 = (A + C) * 0.5;
    Line R1 = Line(P1, P1 + normal(B - A));
    Line R2 = Line(P2, P2 + normal(C - A));
    circle O;
    O.o = ll_intersection(R1, R2);
    O.r = length(A - O.o);
    return O;
}

struct ConvexHull {

    int op;

    struct cmp {
        bool operator()(const Point &a, const Point &b) co
nst {
            return sgn(a.x - b.x) < 0 || sgn(a.x - b.x) ==
0 && sgn(a.y - b.y) < 0;
        }
    };

    set<Point, cmp> s;

    ConvexHull(int o) {
        op = o;
        s.clear();
    }

    inline int PIP(Point P) {
        set<Point>::iterator it = s.lower_bound(Point(P.x,
 -dinf));//找到第一个横坐标大于 P 的点
        if (it == s.end())return 0;
        if (sgn(it->x - P.x) == 0) return sgn((P.y - it->y)
 * op) <= 0;//比较纵坐标大小
        if (it == s.begin())return 0;
        set<Point>::iterator j = it, k = it;
        --j;
        return sgn(det(P - *j, *k - *j) * op) >= 0;//看叉姬
1
    }

    inline int judge(set<Point>::iterator it) {
        set<Point>::iterator j = it, k = it;
        if (j == s.begin())return 0;
        --j;
        if (++k == s.end())return 0;
        return sgn(det(*it - *j, *k - *j) * op) >= 0;//看叉
姬
    }

    inline void insert(Point P) {
        if (PIP(P))return;//如果点 P 已经在凸壳上或凸包里就不
插入了
        set<Point>::iterator tmp = s.lower_bound(Point(P.x,
 -inf));
        if (tmp != s.end() && sgn(tmp->x - P.x) == 0)s.eras
e(tmp);//特判横坐标相等的点要去掉
        s.insert(P);
        set<Point>::iterator it = s.find(P), p = it;
        if (p != s.begin()) {
            --p;
            while (judge(p)) {
                set<Point>::iterator temp = p--;
                s.erase(temp);
            }
        }
        if ((p = ++it) != s.end()) {
            while (judge(p)) {
                set<Point>::iterator temp = p++;
                s.erase(temp);
            }
        }
    }
} up(1), down(-1);

int PIC(circle C, Point a) { return sgn(length(a - C.o) -
C.r) <= 0; }//判断点 A 是否在圆 C 内
void Random(Point *P, int n) { for (int i = 0; i < n; ++i)
swap(P[i], P[(rand() + 1) % n]); }//随机一个排列
circle min_circle(Point *P, int n) {//【求点集 P 的最小覆盖
圆】 O(n)
//  random_shuffle(P,P+n);
    Random(P, n);
    circle C = circle(P[0], 0);
    for (int i = 1; i < n; ++i)
        if (!PIC(C, P[i])) {
            C = circle(P[i], 0);
            for (int j = 0; j < i; ++j)
                if (!PIC(C, P[j])) {
                    C.o = (P[i] + P[j]) * 0.5, C.r = length
(P[j] - C.o);
                    for (int k = 0; k < j; ++k) if (!PIC(C,
P[k])) C = getcircle(P[i], P[j], P[k]);
                }
        }
    return C;
}
```

高精度

高精度 GCD

```
#include <bits/stdc++.h>
using namespace std;
string add(string a, string b) {
    const int L = 1e5;
    string ans;
```

```
    int na[L] = {0}, nb[L] = {0};
    int la = a.size(), lb = b.size();
    for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0
';
    for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0
';
    int lmax = la > lb ? la : lb;
    for (int i = 0; i < lmax; i++)
        na[i] += nb[i], na[i + 1] += na[i] / 10, na[i] %= 1
0;
    if (na[lmax]) lmax++;
    for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
    return ans;
}
string mul(string a, string b) {
    const int L = 1e5;
    string s;
    int na[L], nb[L], nc[L],
        La = a.size(), Lb = b.size();  // na 存储被乘数，nb
存储乘数，nc 存储积
    fill(na, na + L, 0);
    fill(nb, nb + L, 0);
    fill(nc, nc + L, 0);  //将na,nb,nc 都置为0
    for (int i = La - 1; i >= 0; i--)
        na[La - i] =
            a[i] - '0';  //将字符串表示的大整形数转成i 整形数
组表示的大整形数
    for (int i = Lb - 1; i >= 0; i--) nb[Lb - i] = b[i] - '
0';
    for (int i = 1; i <= La; i++)
        for (int j = 1; j <= Lb; j++)
            nc[i + j - 1] +=
                na[i] *
                nb[j];  // a 的第i 位乘以b 的第j 位为积的第i+
j-1 位（先不考虑进位）
    for (int i = 1; i <= La + Lb; i++)
        nc[i + 1] += nc[i] / 10, nc[i] %= 10;  //统一处理进
位
    if (nc[La + Lb]) s += nc[La + Lb] + '0';  //判断第i+j
位上的数字是不是0
    for (int i = La + Lb - 1; i >= 1; i--)
        s += nc[i] + '0';  //将整形数组转成字符串
    return s;
}
int sub(int *a, int *b, int La, int Lb) {
    if (La < Lb) return -1;  //如果a 小于b，则返回-1
    if (La == Lb) {
        for (int i = La - 1; i >= 0; i--)
            if (a[i] > b[i])
                break;
            else if (a[i] < b[i])
                return -1;  //如果a 小于b，则返回-1
    }
    for (int i = 0; i < La; i++)  //高精度减法
    {
        a[i] -= b[i];
        if (a[i] < 0) a[i] += 10, a[i + 1]--;
    }
    for (int i = La - 1; i >= 0; i--)
        if (a[i]) return i + 1;  //返回差的位数
    return 0;                    //返回差的位数
}
string div(string n1, string n2,
           int nn)  // n1,n2 是字符串表示的被除数，除数,nn 是
选择返回商还是余数
{
    const int L = 1e5;
    string s, v;  // s 存商,v 存余数
    int a[L], b[L], r[L],
        La = n1.size(), Lb = n2.size(), i,
        tp = La;  // a，b 是整形数组表示被除数，除数，tp 保存
被除数的长度
    fill(a, a + L, 0);
    fill(b, b + L, 0);
    fill(r, r + L, 0);  //数组元素都置为0
    for (i = La - 1; i >= 0; i--) a[La - 1 - i] = n1[i] - '
0';
    for (i = Lb - 1; i >= 0; i--) b[Lb - 1 - i] = n2[i] - '
0';
    if (La < Lb || (La == Lb && n1 < n2)) {
        // cout<<0<<endl;
        return n1;
    }               //如果a<b,则商为0，余数为被除数
    int t = La - Lb;  //除被数和除数的位数之差
    for (int i = La - 1; i >= 0; i--)  //将除数扩大10^t 倍
        if (i >= t)
            b[i] = b[i - t];
        else
            b[i] = 0;
    Lb = La;
    for (int j = 0; j <= t; j++) {
        int temp;
        while ((temp = sub(a, b + j, La, Lb - j)) >=
               0)  //如果被除数比除数大继续减
        {
            La = temp;
            r[t - j]++;
        }
    }
    for (i = 0; i < L - 10; i++)
        r[i + 1] += r[i] / 10, r[i] %= 10;  //统一处理进位
    while (!r[i]) i--;  //将整形数组表示的商转化成字符串表示
的
    while (i >= 0) s += r[i--] + '0';
    // cout<<s<<endl;
    i = tp;
    while (!a[i]) i--;  //将整形数组表示的余数转化成字符串表
示的</span>
    while (i >= 0) v += a[i--] + '0';
    if (v.empty()) v = "0";
    // cout<<v<<endl;
    if (nn == 1) return s;
    if (nn == 2) return v;
}
bool judge(string s)  //判断s 是否为全0 串
{
    for (int i = 0; i < s.size(); i++)
        if (s[i] != '0') return false;
    return true;
}
string gcd(string a, string b)  //求最大公约数
{
    string t;
    while (!judge(b))  //如果余数不为0，继续除
    {
        t = a;            //保存被除数的值
        a = b;            //用除数替换被除数
        b = div(t, b, 2);  //用余数替换除数
    }
    return a;
}

//o(无法估计)
```

高精度乘法（FFT）

```
#include <bits/stdc++.h>
using namespace std;
#define L(x) (1 << (x))
const double PI = acos(-1.0);
const int Maxn = 133015;
double ax[Maxn], ay[Maxn], bx[Maxn], by[Maxn];
char sa[Maxn / 2], sb[Maxn / 2];
int sum[Maxn];
int x1[Maxn], x2[Maxn];
int revv(int x, int bits) {
    int ret = 0;
    for (int i = 0; i < bits; i++) {
        ret <<= 1;
        ret |= x & 1;
        x >>= 1;
    }
    return ret;
}
void fft(double* a, double* b, int n, bool rev) {
    int bits = 0;
    while (1 << bits < n) ++bits;
    for (int i = 0; i < n; i++) {
        int j = revv(i, bits);
        if (i < j) swap(a[i], a[j]), swap(b[i], b[j]);
    }
    for (int len = 2; len <= n; len <<= 1) {
        int half = len >> 1;
```

```
        double wmx = cos(2 * PI / len), wmy = sin(2 * PI /
len);
        if (rev) wmy = -wmy;
        for (int i = 0; i < n; i += len) {
            double wx = 1, wy = 0;
            for (int j = 0; j < half; j++) {
                double cx = a[i + j], cy = b[i + j];
                double dx = a[i + j + half], dy = b[i + j +
half];
                double ex = dx * wx - dy * wy, ey = dx * wy
+ dy * wx;
                a[i + j] = cx + ex, b[i + j] = cy + ey;
                a[i + j + half] = cx - ex, b[i + j + half] =
 cy - ey;
                double wnx = wx * wmx - wy * wmy, wny = wx *
 wmy + wy * wmx;
                wx = wnx, wy = wny;
            }
        }
    }
    if (rev) {
        for (int i = 0; i < n; i++) a[i] /= n, b[i] /= n;
    }
}
int solve(int a[], int na, int b[], int nb, int ans[]) {
    int len = max(na, nb), ln;
    for (ln = 0; L(ln) < len; ++ln)
        ;
    len = L(++ln);
    for (int i = 0; i < len; ++i) {
        if (i >= na)
            ax[i] = 0, ay[i] = 0;
        else
            ax[i] = a[i], ay[i] = 0;
    }
    fft(ax, ay, len, 0);
    for (int i = 0; i < len; ++i) {
        if (i >= nb)
            bx[i] = 0, by[i] = 0;
        else
            bx[i] = b[i], by[i] = 0;
    }
    fft(bx, by, len, 0);
    for (int i = 0; i < len; ++i) {
        double cx = ax[i] * bx[i] - ay[i] * by[i];
        double cy = ax[i] * by[i] + ay[i] * bx[i];
        ax[i] = cx, ay[i] = cy;
    }
    fft(ax, ay, len, 1);
    for (int i = 0; i < len; ++i) ans[i] = (int)(ax[i] + 0.
5);
    return len;
}
string mul(string sa, string sb) {
    int l1, l2, l;
    int i;
    string ans;
    memset(sum, 0, sizeof(sum));
    l1 = sa.size();
    l2 = sb.size();
    for (i = 0; i < l1; i++) x1[i] = sa[l1 - i - 1] - '0';
    for (i = 0; i < l2; i++) x2[i] = sb[l2 - i - 1] - '0';
    l = solve(x1, l1, x2, l2, sum);
    for (i = 0; i < l || sum[i] >= 10; i++)  // 进位
    {
        sum[i + 1] += sum[i] / 10;
        sum[i] %= 10;
    }
    l = i;
    while (sum[l] <= 0 && l > 0) l--;              // 检索最
高位
    for (i = l; i >= 0; i--) ans += sum[i] + '0';  // 倒序
输出
    return ans;
}
int main() {
    cin.sync_with_stdio(false);
    string a, b;
    while (cin >> a >> b) cout << mul(a, b) << endl;
    return 0;
}

//o(nLogn)
```

高精度乘法（乘单精）

```
#include <bits/stdc++.h>
using namespace std;
string mul(string a, int b)  //高精度a 乘单精度b
{
    const int L = 100005;
    int na[L];
    string ans;
    int La = a.size();
    fill(na, na + L, 0);
    for (int i = La - 1; i >= 0; i--) na[La - i - 1] = a[i]
 - '0';
    int w = 0;
    for (int i = 0; i < La; i++)
        na[i] = na[i] * b + w, w = na[i] / 10, na[i] = na[i]
 % 10;
    while (w) na[La++] = w % 10, w /= 10;
    La--;
    while (La >= 0) ans += na[La--] + '0';
    return ans;
}

//o(n)
```

高精度乘法（朴素）

```
#include <bits/stdc++.h>
using namespace std;
string mul(string a, string b)  //高精度乘法a,b,均为非负整
数
{
    const int L = 1e5;
    string s;
    int na[L], nb[L], nc[L],
        La = a.size(), Lb = b.size();  // na 存储被乘数，nb
存储乘数，nc 存储积
    fill(na, na + L, 0);
    fill(nb, nb + L, 0);
    fill(nc, nc + L, 0);  //将na,nb,nc 都置为0
    for (int i = La - 1; i >= 0; i--)
        na[La - i] =
            a[i] - '0';  //将字符串表示的大整形数转成i 整形数
组表示的大整形数
    for (int i = Lb - 1; i >= 0; i--) nb[Lb - i] = b[i] - '
0';
    for (int i = 1; i <= La; i++)
        for (int j = 1; j <= Lb; j++)
            nc[i + j - 1] +=
                na[i] *
                nb[j];  // a 的第i 位乘以b 的第j 位为积的第i+
j-1 位（先不考虑进位）
    for (int i = 1; i <= La + Lb; i++)
        nc[i + 1] += nc[i] / 10, nc[i] %= 10;  //统一处理进
位
    if (nc[La + Lb]) s += nc[La + Lb] + '0';  //判断第i+j
位上的数字是不是0
    for (int i = La + Lb - 1; i >= 1; i--)
        s += nc[i] + '0';  //将整形数组转成字符串
    return s;
}

//o(n^2)
```

高精度减法

```
#include <bits/stdc++.h>
using namespace std;
string sub(string a, string b)  //只限大的非负整数减小的非负
整数
{
    const int L = 1e5;
    string ans;
    int na[L] = {0}, nb[L] = {0};
    int la = a.size(), lb = b.size();
    for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0
';
    for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0
';
    int lmax = la > lb ? la : lb;
```

```
    for (int i = 0; i < lmax; i++) {
        na[i] -= nb[i];
        if (na[i] < 0) na[i] += 10, na[i + 1]--;
    }
    while (!na[--lmax] && lmax > 0)
        ;
    lmax++;
    for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
    return ans;
}

//o(n)
```

高精度加法

```
#include <bits/stdc++.h>
using namespace std;
string add(string a, string b)  //只限两个非负整数相加
{
    const int L = 1e5;
    string ans;
    int na[L] = {0}, nb[L] = {0};
    int la = a.size(), lb = b.size();
    for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0
';
    for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0
';
    int lmax = la > lb ? la : lb;
    for (int i = 0; i < lmax; i++)
        na[i] += nb[i], na[i + 1] += na[i] / 10, na[i] %= 1
0;
    if (na[lmax]) lmax++;
    for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
    return ans;
}

//o(n)
```

高精度取模（对单精）

```
#include <bits/stdc++.h>
using namespace std;
int mod(string a,int b)//高精度a除以单精度b
{
    int d=0;
    for(int i=0;i<a.size();i++) d=(d*10+(a[i]-'0'))%b;//求
出余数
    return d;
}

//o(n)
```

高精度幂

```
#include <bits/stdc++.h>
#define L(x) (1 << (x))
using namespace std;
const double PI = acos(-1.0);
const int Maxn = 133015;
double ax[Maxn], ay[Maxn], bx[Maxn], by[Maxn];
char sa[Maxn / 2], sb[Maxn / 2];
int sum[Maxn];
int x1[Maxn], x2[Maxn];
int revv(int x, int bits) {
    int ret = 0;
    for (int i = 0; i < bits; i++) {
        ret <<= 1;
        ret |= x & 1;
        x >>= 1;
    }
    return ret;
}
void fft(double* a, double* b, int n, bool rev) {
    int bits = 0;
    while (1 << bits < n) ++bits;
    for (int i = 0; i < n; i++) {
        int j = revv(i, bits);
        if (i < j) swap(a[i], a[j]), swap(b[i], b[j]);
    }
    for (int len = 2; len <= n; len <<= 1) {
        int half = len >> 1;
        double wmx = cos(2 * PI / len), wmy = sin(2 * PI /
len);
        if (rev) wmy = -wmy;
        for (int i = 0; i < n; i += len) {
            double wx = 1, wy = 0;
            for (int j = 0; j < half; j++) {
                double cx = a[i + j], cy = b[i + j];
                double dx = a[i + j + half], dy = b[i + j +
half];
                double ex = dx * wx - dy * wy, ey = dx * wy
+ dy * wx;
                a[i + j] = cx + ex, b[i + j] = cy + ey;
                a[i + j + half] = cx - ex, b[i + j + half] =
 cy - ey;
                double wnx = wx * wmx - wy * wmy, wny = wx *
 wmy + wy * wmx;
                wx = wnx, wy = wny;
            }
        }
    }
    if (rev) {
        for (int i = 0; i < n; i++) a[i] /= n, b[i] /= n;
    }
}
int solve(int a[], int na, int b[], int nb, int ans[]) {
    int len = max(na, nb), ln;
    for (ln = 0; L(ln) < len; ++ln)
        ;
    len = L(++ln);
    for (int i = 0; i < len; ++i) {
        if (i >= na)
            ax[i] = 0, ay[i] = 0;
        else
            ax[i] = a[i], ay[i] = 0;
    }
    fft(ax, ay, len, 0);
    for (int i = 0; i < len; ++i) {
        if (i >= nb)
            bx[i] = 0, by[i] = 0;
        else
            bx[i] = b[i], by[i] = 0;
    }
    fft(bx, by, len, 0);
    for (int i = 0; i < len; ++i) {
        double cx = ax[i] * bx[i] - ay[i] * by[i];
        double cy = ax[i] * by[i] + ay[i] * bx[i];
        ax[i] = cx, ay[i] = cy;
    }
    fft(ax, ay, len, 1);
    for (int i = 0; i < len; ++i) ans[i] = (int)(ax[i] + 0.
5);
    return len;
}
string mul(string sa, string sb) {
    int l1, l2, l;
    int i;
    string ans;
    memset(sum, 0, sizeof(sum));
    l1 = sa.size();
    l2 = sb.size();
    for (i = 0; i < l1; i++) x1[i] = sa[l1 - i - 1] - '0';
    for (i = 0; i < l2; i++) x2[i] = sb[l2 - i - 1] - '0';
    l = solve(x1, l1, x2, l2, sum);
    for (i = 0; i < l || sum[i] >= 10; i++)  // 进位
    {
        sum[i + 1] += sum[i] / 10;
        sum[i] %= 10;
    }
    l = i;
    while (sum[l] <= 0 && l > 0) l--;              // 检索最
高位
    for (i = l; i >= 0; i--) ans += sum[i] + '0';  // 倒序
输出
    return ans;
}
string Pow(string a, int n) {
    if (n == 1) return a;
    if (n & 1) return mul(Pow(a, n - 1), a);
    string ans = Pow(a, n / 2);
    return mul(ans, ans);
}
```

```
//o(nLognLogm)
```

高精度平方根

```
#include <bits/stdc++.h>
using namespace std;
const int L = 2015;
string add(string a, string b)  //只限两个非负整数相加
{
    string ans;
    int na[L] = {0}, nb[L] = {0};
    int la = a.size(), lb = b.size();
    for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0
';
    for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0
';
    int lmax = la > lb ? la : lb;
    for (int i = 0; i < lmax; i++)
        na[i] += nb[i], na[i + 1] += na[i] / 10, na[i] %= 1
0;
    if (na[lmax]) lmax++;
    for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
    return ans;
}
string sub(string a, string b)  //只限大的非负整数减小的非负
整数
{
    string ans;
    int na[L] = {0}, nb[L] = {0};
    int la = a.size(), lb = b.size();
    for (int i = 0; i < la; i++) na[la - 1 - i] = a[i] - '0
';
    for (int i = 0; i < lb; i++) nb[lb - 1 - i] = b[i] - '0
';
    int lmax = la > lb ? la : lb;
    for (int i = 0; i < lmax; i++) {
        na[i] -= nb[i];
        if (na[i] < 0) na[i] += 10, na[i + 1]--;
    }
    while (!na[--lmax] && lmax > 0)
        ;
    lmax++;
    for (int i = lmax - 1; i >= 0; i--) ans += na[i] + '0';
    return ans;
}
string mul(string a, string b)  //高精度乘法a,b,均为非负整
数
{
    string s;
    int na[L], nb[L], nc[L],
        La = a.size(), Lb = b.size();  // na 存储被乘数，nb
存储乘数，nc 存储积
    fill(na, na + L, 0);
    fill(nb, nb + L, 0);
    fill(nc, nc + L, 0);  //将na,nb,nc 都置为0
    for (int i = La - 1; i >= 0; i--)
        na[La - i] =
            a[i] - '0';  //将字符串表示的大整形数转成i 整形数
组表示的大整形数
    for (int i = Lb - 1; i >= 0; i--) nb[Lb - i] = b[i] - '
0';
    for (int i = 1; i <= La; i++)
        for (int j = 1; j <= Lb; j++)
            nc[i + j - 1] +=
                na[i] *
                nb[j];  // a 的第i 位乘以b 的第j 位为积的第i+
j-1 位（先不考虑进位）
    for (int i = 1; i <= La + Lb; i++)
        nc[i + 1] += nc[i] / 10, nc[i] %= 10;  //统一处理进
位
    if (nc[La + Lb]) s += nc[La + Lb] + '0';  //判断第i+j
位上的数字是不是0
    for (int i = La + Lb - 1; i >= 1; i--)
        s += nc[i] + '0';  //将整形数组转成字符串
    return s;
}
int sub(int *a, int *b, int La, int Lb) {
    if (La < Lb) return -1;  //如果a 小于b，则返回-1
    if (La == Lb) {
        for (int i = La - 1; i >= 0; i--)
            if (a[i] > b[i])
                break;
            else if (a[i] < b[i])
                return -1;  //如果a 小于b，则返回-1
    }
    for (int i = 0; i < La; i++)  //高精度减法
    {
        a[i] -= b[i];
        if (a[i] < 0) a[i] += 10, a[i + 1]--;
    }
    for (int i = La - 1; i >= 0; i--)
        if (a[i]) return i + 1;  //返回差的位数
    return 0;                    //返回差的位数
}
string div(string n1, string n2,
           int nn)  // n1,n2 是字符串表示的被除数，除数,nn 是
选择返回商还是余数
{
    string s, v;  // s 存商,v 存余数
    int a[L], b[L], r[L],
        La = n1.size(), Lb = n2.size(), i,
        tp = La;  // a，b 是整形数组表示被除数，除数，tp 保存
被除数的长度
    fill(a, a + L, 0);
    fill(b, b + L, 0);
    fill(r, r + L, 0);  //数组元素都置为0
    for (i = La - 1; i >= 0; i--) a[La - 1 - i] = n1[i] - '
0';
    for (i = Lb - 1; i >= 0; i--) b[Lb - 1 - i] = n2[i] - '
0';
    if (La < Lb || (La == Lb && n1 < n2)) {
        // cout<<0<<endl;
        return n1;
    }               //如果a<b,则商为0，余数为被除数
    int t = La - Lb;  //除被数和除数的位数之差
    for (int i = La - 1; i >= 0; i--)  //将除数扩大10^t 倍
        if (i >= t)
            b[i] = b[i - t];
        else
            b[i] = 0;
    Lb = La;
    for (int j = 0; j <= t; j++) {
        int temp;
        while ((temp = sub(a, b + j, La, Lb - j)) >=
               0)  //如果被除数比除数大继续减
        {
            La = temp;
            r[t - j]++;
        }
    }
    for (i = 0; i < L - 10; i++)
        r[i + 1] += r[i] / 10, r[i] %= 10;  //统一处理进位
    while (!r[i]) i--;  //将整形数组表示的商转化成字符串表示
的
    while (i >= 0) s += r[i--] + '0';
    // cout<<s<<endl;
    i = tp;
    while (!a[i]) i--;  //将整形数组表示的余数转化成字符串表
示的</span>
    while (i >= 0) v += a[i--] + '0';
    if (v.empty()) v = "0";
    // cout<<v<<endl;
    if (nn == 1) return s;
    if (nn == 2) return v;
}
bool cmp(string a, string b) {
    if (a.size() < b.size()) return 1;  // a 小于等于b 返回
真
    if (a.size() == b.size() && a <= b) return 1;
    return 0;
}
string DeletePreZero(string s) {
    int i;
    for (i = 0; i < s.size(); i++)
        if (s[i] != '0') break;
    return s.substr(i);
}

string BigInterSqrt(string n) {
    n = DeletePreZero(n);
    string l = "1", r = n, mid, ans;
    while (cmp(l, r)) {
        mid = div(add(l, r), "2", 1);
```

```
        if (cmp(mul(mid, mid), n))
            ans = mid, l = add(mid, "1");
        else
            r = sub(mid, "1");
    }
    return ans;
}

// o(n^3)
```

高精度进制转换

```
#include <bits/stdc++.h>
using namespace std;
//将字符串表示的10进制大整数转换为m进制的大整数
//并返回m进制大整数的字符串
bool judge(string s)  //判断串是否为全零串
{
    for (int i = 0; i < s.size(); i++)
        if (s[i] != '0') return 1;
    return 0;
}
string solve(
    string s, int n,
    int m)  // n进制转m进制只限0-9进制，若涉及带字母的进制，稍作修改即可
{
    string r, ans;
    int d = 0;
    if (!judge(s)) return "0";  //特判
    while (judge(s))            //被除数不为0则继续
    {
        for (int i = 0; i < s.size(); i++) {
            r += (d * n + s[i] - '0') / m + '0';  //求出商
            d = (d * n + (s[i] - '0')) % m;       //求出余数
        }
        s = r;           //把商赋给下一次的被除数
        r = "";          //把商清空
        ans += d + '0';  //加上进制转换后数字
        d = 0;           //清空余数
    }
    reverse(ans.begin(), ans.end());  //倒置下
    return ans;
}

//o(n^2)
```

高精度阶乘

```
#include <bits/stdc++.h>
using namespace std;
string fac(int n) {
    const int L = 100005;
    int a[L];
    string ans;
    if (n == 0) return "1";
    fill(a, a + L, 0);
    int s = 0, m = n;
    while (m) a[++s] = m % 10, m /= 10;
    for (int i = n - 1; i >= 2; i--) {
        int w = 0;
        for (int j = 1; j <= s; j++)
            a[j] = a[j] * i + w, w = a[j] / 10, a[j] = a[j]
 % 10;
        while (w) a[++s] = w % 10, w /= 10;
    }
    while (!a[s]) s--;
    while (s >= 1) ans += a[s--] + '0';
    return ans;
}

//o(n^2)
```

高精度除法（除单精）

```
#include <bits/stdc++.h>
using namespace std;
string div(string a, int b)  //高精度a除以单精度b
{
    string r, ans;
    int d = 0;
    if (a == "0") return a;  //特判
    for (int i = 0; i < a.size(); i++) {
        r += (d * 10 + a[i] - '0') / b + '0';  //求出商
        d = (d * 10 + (a[i] - '0')) % b;       //求出余数
    }
    int p = 0;
    for (int i = 0; i < r.size(); i++)
        if (r[i] != '0') {
            p = i;
            break;
        }
    return r.substr(p);
}

//o(n)
```

高精度除法（除高精）

```
#include <bits/stdc++.h>
using namespace std;
int sub(int *a, int *b, int La, int Lb) {
    if (La < Lb) return -1;  //如果a小于b，则返回-1
    if (La == Lb) {
        for (int i = La - 1; i >= 0; i--)
            if (a[i] > b[i])
                break;
            else if (a[i] < b[i])
                return -1;  //如果a小于b，则返回-1
    }
    for (int i = 0; i < La; i++)  //高精度减法
    {
        a[i] -= b[i];
        if (a[i] < 0) a[i] += 10, a[i + 1]--;
    }
    for (int i = La - 1; i >= 0; i--)
        if (a[i]) return i + 1;  //返回差的位数
    return 0;                    //返回差的位数
}
string div(string n1, string n2, int nn)
// n1,n2是字符串表示的被除数，除数,nn是选择返回商还是余数
{
    const int L = 1e5;
    string s, v;  // s存商,v存余数
    int a[L], b[L], r[L], La = n1.size(), Lb = n2.size(),
i, tp = La;
    // a，b是整形数组表示被除数，除数，tp保存被除数的长度
    fill(a, a + L, 0);
    fill(b, b + L, 0);
    fill(r, r + L, 0);  //数组元素都置为0
    for (i = La - 1; i >= 0; i--) a[La - 1 - i] = n1[i] - '
0';
    for (i = Lb - 1; i >= 0; i--) b[Lb - 1 - i] = n2[i] - '
0';
    if (La < Lb || (La == Lb && n1 < n2)) {
        // cout<<0<<endl;
        return n1;
    }               //如果a<b,则商为0，余数为被除数
    int t = La - Lb;  //除被数和除数的位数之差
    for (int i = La - 1; i >= 0; i--)  //将除数扩大10^t倍
        if (i >= t)
            b[i] = b[i - t];
        else
            b[i] = 0;
    Lb = La;
    for (int j = 0; j <= t; j++) {
        int temp;
        while ((temp = sub(a, b + j, La, Lb - j)) >=
               0)  //如果被除数比除数大继续减
        {
            La = temp;
            r[t - j]++;
        }
    }
    for (i = 0; i < L - 10; i++)
        r[i + 1] += r[i] / 10, r[i] %= 10;  //统一处理进位
    while (!r[i]) i--;  //将整形数组表示的商转化成字符串表示的
    while (i >= 0) s += r[i--] + '0';
    // cout<<s<<endl;
```

```
    i = tp;
    while (!a[i]) i--;  //将整形数组表示的余数转化成字符串表
示的</span>
    while (i >= 0) v += a[i--] + '0';
    if (v.empty()) v = "0";
    // cout<<v<<endl;
    if (nn == 1) return s;  //返回商
    if (nn == 2) return v;  //返回余数
}

//o(n^2)
```

龟速乘快速幂（快速幂爆 longlong

```
#include <bits/stdc++.h>
using namespace std;

typedef long long ll;

ll qmul(ll a, ll b, ll p) {
          ll res = 0;
          while(b) {
                    if(b & 1) res = (res + a) % p;
                    a = (a + a) % p;
                    b >>= 1;
          }
          return res;
}

ll qpow(ll x, ll n, ll p) {
          ll res = 1;
          while(n) {
                    if(n & 1) res = qmul(res, x, p);
                    x = qmul(x, x, p);
                    n >>= 1;
          }
          return res % p; // 1 0 1
}

int main() {
          ll b, p, k;
          cin >> b >> p >> k;
          ll ans = qpow(b, p, k);
          printf("%lld^%lld mod %lld=%lld", b, p, k, ans);

          return 0;
}
```