

# Planar Graph

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# We are going to see

- 1 Introduction
- 2 Faces of a planar graph
- 3 Euler's Theorem
- 4 Closing Remarks



# Motivation

- PCB (Printed Circuit Board) is designed by planar graph.

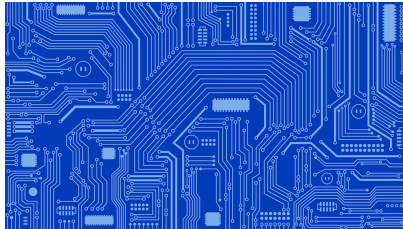


Figure: PCB Circuit



# Planar Graph

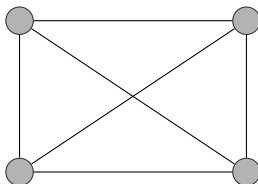
## Definition

A graph is called **planar** whether it can be drawn in the plane in such a way that no two edges cross.



# Planar Graph

A clique on 4 nodes

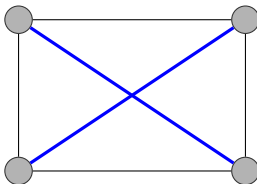


- Let's see whether it's a planar graph or not
- It has 6 edges, two of which crossing each other
- So, it is not a Planar Graph



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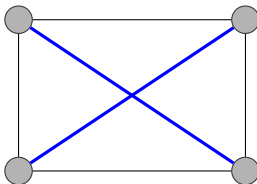


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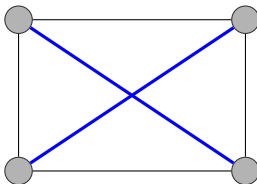


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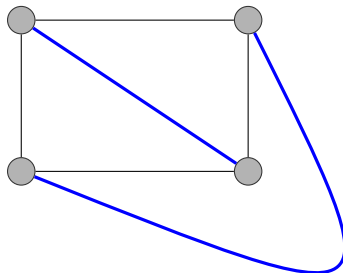


- Let's see whether it's a planar graph or not
- It has 6 edges, two of which crossing each other
- So, it is not a Planar Graph





# Planar Graph



Another clique on 4 nodes with six edges  
It's a planar Graph because no edges are crossing each other.



# Faces on a Planar Graph

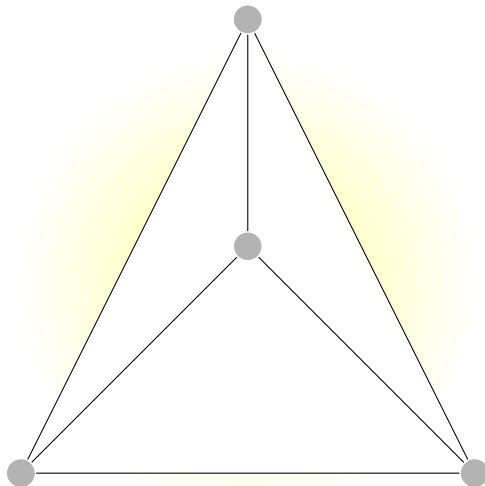
## Definition

Consider a planar graph  $G = (V, E)$ . A face is defined to be an area of the plane that is bounded by edges. A planar graph divides the planes into one or more **faces**. One of these faces always will be *infinite*.



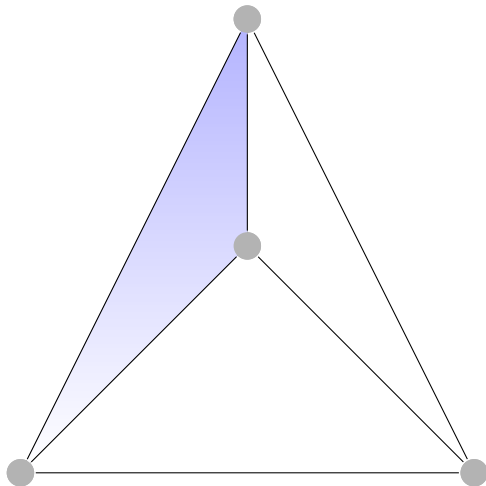
# Finding Faces of a Planar Graph

**Infinite Face**



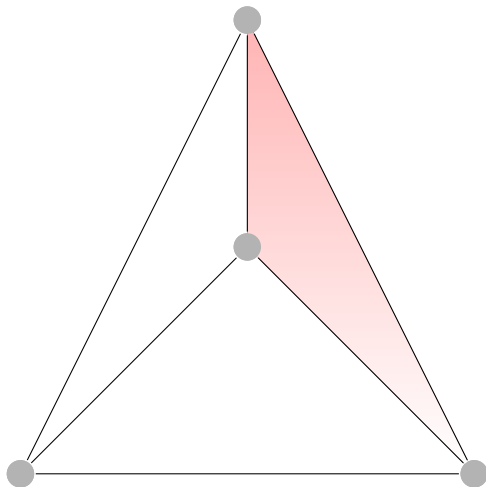
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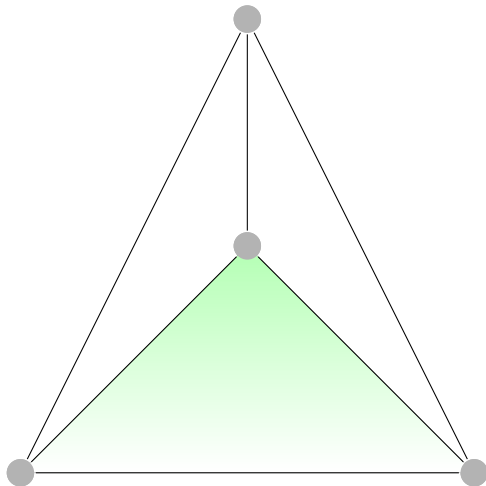
# Finding Faces of a Planar Graph

**Finite Face**



# Finding Faces of a Planar Graph

**Finite Face**



# Euler's Theorem on Planar Graph

- Let  $G$  be a connected planar graph (drawn without crossing edges).
- Define
  - $V$  = number of vertices
  - $E$  = number of edges
  - $F$  = number of faces, including the “infinite” face
- Then  $V - E + F = 2$ .



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# Proof

## PROOF IDEA

Proof by induction by the number of cycles

## BASE CASE

- $G$  has no cycles
- Since  $G$  is connected, it must be a tree. So,  $e = v - 1$  and  $f = 1$ .

$$\begin{aligned}v - e + f &= v - (v - 1) + 1 \\&= 1 + 1 \\&= 2\end{aligned}$$



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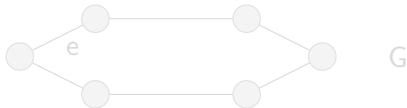
$$\begin{aligned}v - e + f &= v - (v - 1) + 1 \\&= 1 + 1 \\&= 2\end{aligned}$$



# Proof – Continued

## INDUCTION

- Let  $G$  has at least one cycle containing edge  $e$
- let  $G' = G - e$



$v'$  = number of vertices

$e'$  = number of edges

$f'$  = number of faces

- Now, in  $G'$ ,  $f' = f - 1$  and  $v' = v$  and  $e' = e - 1$
- By induction hypothesis:

$$v' - e' + f' = 2$$

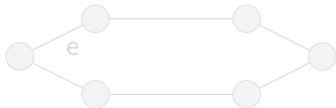
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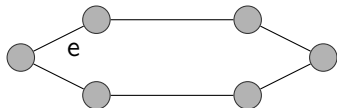
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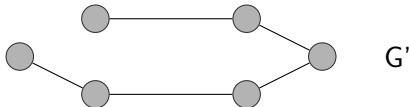
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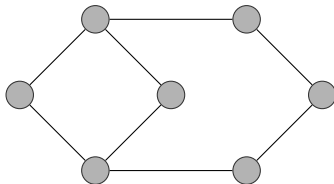
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$$v - e + f = 2$$



# Question

**Can you redraw this graph to alter the number of faces?**





# Corollary of Euler's Theorem

## An interesting fact

No matter how we redraw a planar graph it will always have the same number of faces

## Proof

$$f = 2 - v + e$$



# Acknowledgement

- [https://mathweb.ucsd.edu/~gptesler/154/slides/154\\_planar\\_20-handout.pdf](https://mathweb.ucsd.edu/~gptesler/154/slides/154_planar_20-handout.pdf)
- <https://www.slideserve.com/thor/planar-graphs>
- [https://en.wikipedia.org/wiki/Planar\\_graph](https://en.wikipedia.org/wiki/Planar_graph)



## Further Reads

<https://ieeexplore.ieee.org/document/5206863>

A very interesting application of planar graph is in Computer Vision  
Efficient planar graph cuts with applications in Computer Vision  
by Frank R. Schmidt et al.

