



App Academy

Linked Lists



Linked Lists

- Linked Lists are made up of “list nodes”
- Each of these nodes encapsulate at minimum 2 components
 - Value
 - Pointers to other nodes
 - Next pointer
 - Prev pointer (for doubly linked list)
- We can represent these nodes using objects with key-value pairs
 - Example:

```
{  
    value: 0,  
    next: NODE,  
    prev: NODE  
}
```



Singly Linked Lists

ListNode	
value	next

ListNode 1	
red	

ListNode 2	
blue	

ListNode 3	
green	

RAM

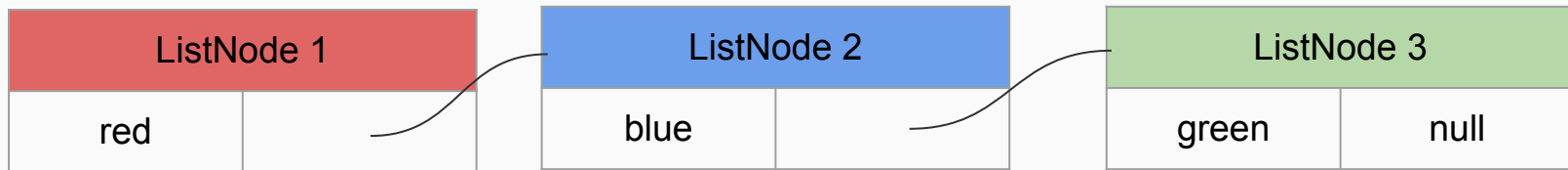
Value

Address



Singly Linked Lists

ListNode	
value	next



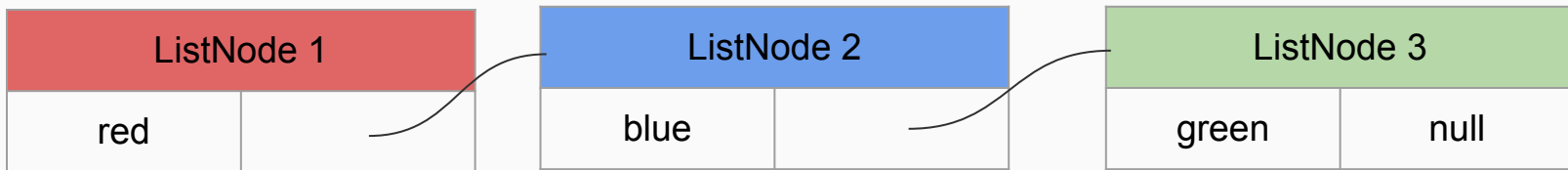
RAM

Value									
Address	\$0	\$4	\$8	\$12	\$16	\$20	\$24	\$28	\$32



Chaining .next

ListNode	
value	next



- From ListNode1, we can access ListNode2 by using `ListNode1.next`
- We can access ListNode3 from both ListNode1 or ListNode2 with the following:
 - `ListNode1.next.next`
 - `ListNode2.next`

Doubly Linked Lists

ListNode		
prev	value	next

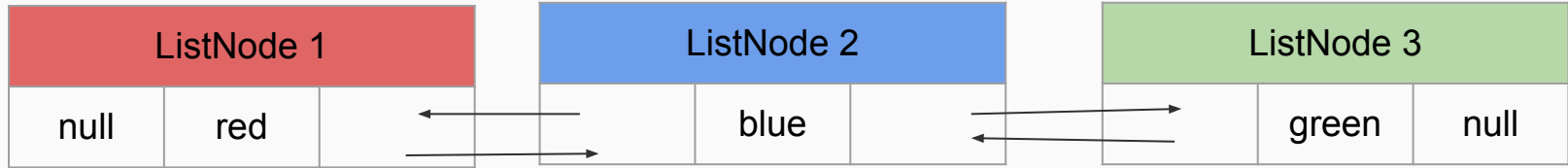
ListNode 1		
	red	

ListNode 2		
	blue	

ListNode 3		
	green	

Doubly Linked Lists

ListNode		
prev	value	next



How do we traverse a linked list?

Iterative:

```
def traverse(head):  
    current = head  
    while (current) :  
        current = current.next;  
    return None
```

Recursive:

```
def traverse(head):  
    if not head:  
        return None  
    return traverse(head.next)
```



Linked List Runtimes

Operations	Big-O Time
Read/ Write ith element	
Insert / Remove End	
Insert Middle or Beginning***	
Remove Middle or Beginning	



Comparing with Array Runtimes

Arrays

Operations	Big-O Time
Read/ Write ith element	$O(1)$
Insert / Remove End	$O(1)$
Insert Middle or Beginning	$O(n)$
Remove Middle or Beginning	$O(n)$

Linked Lists

Operations	Big-O Time
Read/ Write ith element	$O(n)$
Insert / Remove End	$O(1)$
Insert Middle or Beginning***	$O(1)$
Remove Middle or Beginning	$O(1)$

***Note that the act of inserting or removing a node from the middle of a linked list by itself is $O(1)$, but this assumes you already found the location for insertion / removal.



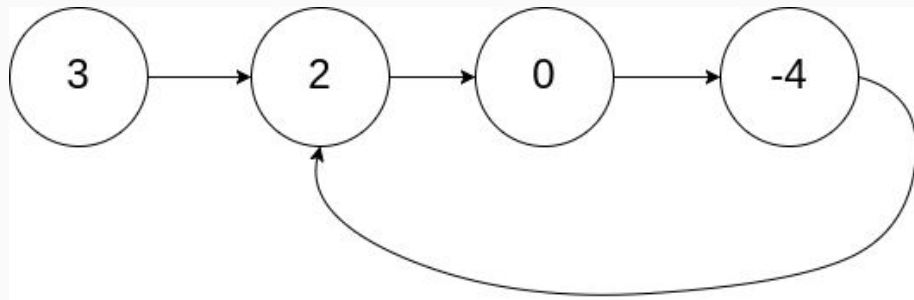
Demos

- [Middle of the Linked List](#)
- [Reverse Linked List](#)



Cycle Detection

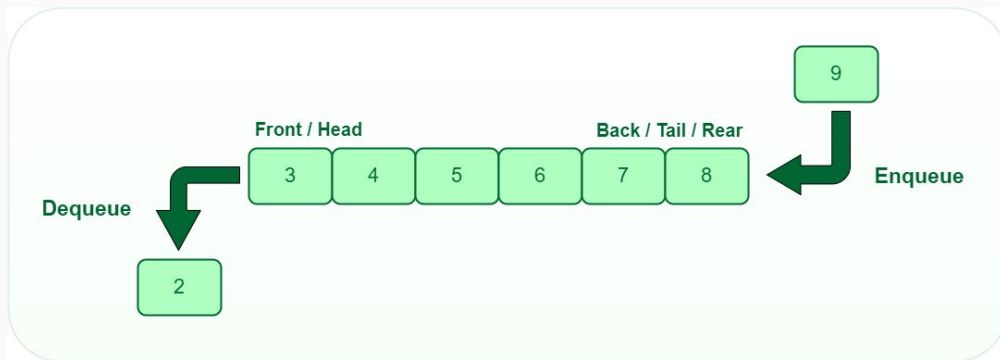
- Fast and slow pointers can also be used with Linked Lists to detect loops within the list.
- Demo:
 - [Linked List Cycle](#)



Queues

- FIFO (first in, first out)
- Because we can add/remove nodes at the beginning of a Linked List in $O(1)$ time, it is actually more efficient to create a queue this way versus an array.
- In python, we can use [deque from collections library](#)

Operations	Big-O Time
Enqueue	$O(1)$
Dequeue	$O(1)$



***There aren't a lot of problems where a queue shines on its own, which is why this slide is lackluster. Don't worry though, in trees & graphs, you'll see how queues are used to implement BFS - an extremely powerful algorithm.



Questions?



Let's practice!

- Review - try doing these both iteratively and recursively!
 - [Merge Two Sorted Lists](#)
 - [Remove Nth Node From End of List](#)
- Bonus
 - [Design Browser History](#)
 - [LRU Cache](#)

