



Sorting Algorithms



# Bubble Sort

```
let bubbleSort = (nums) => {  
  let len = nums.length;  
  for (let i = 0; i < len; i++) {  
    for (let j = 0; j < len; j++) {  
      if (nums[j] > nums[j + 1]) {  
        let tmp = nums[j];  
        nums[j] = nums[j + 1];  
        nums[j + 1] = tmp;  
      }  
    }  
  }  
  return nums;  
};
```

Runtime:  $O(n^2)$

visualization: <https://visualgo.net/en/sorting>

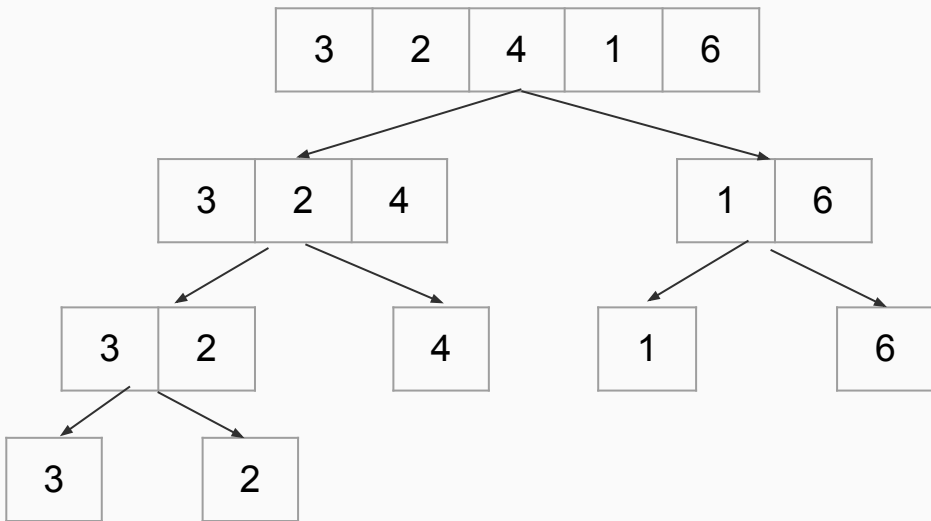
- This sorting algorithm scans through the array several times, taking the largest number and placing it onto the end.
- This algorithm is highly inefficient, and you would never want to sort this way unless specifically asked to.



# Merge Sort

Demo:

<https://leetcode.com/problems/sort-an-array/>



- Runtime:  $O(n \cdot \log(n))$
- visualization: <https://visualgo.net/en/sorting>
- This algorithm splits an array in half continuously and then returns the sorted halves.
- This sounds a lot like recursion!
  - Our base case is when we can no longer split any further; i.e. when `array.length = 1`



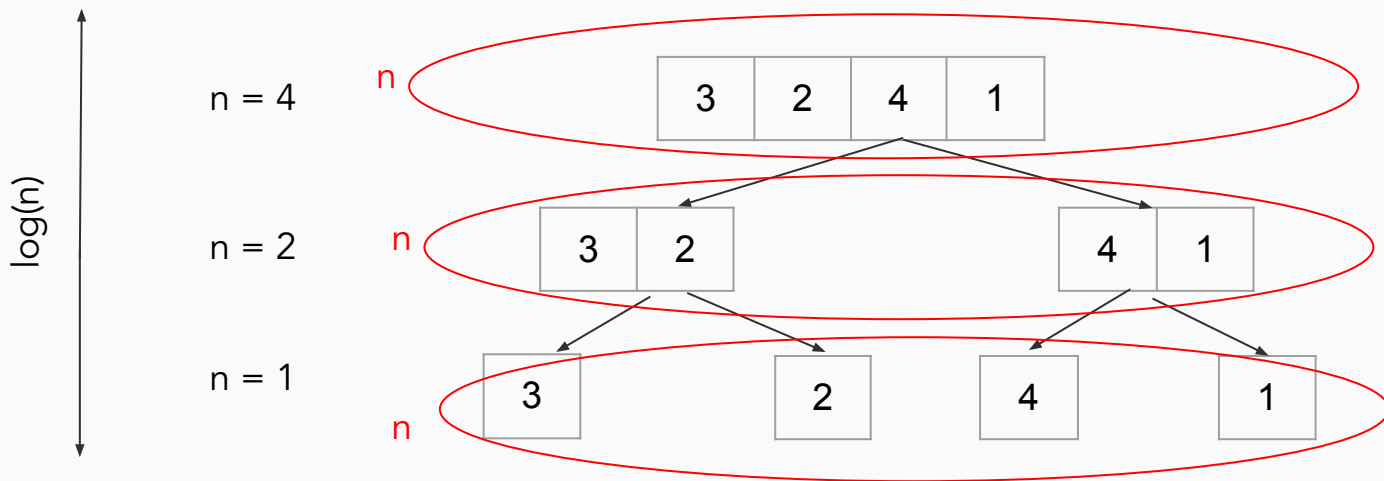
# Merge Sort (cont.)

```
const sortArray = function(nums) {  
  if (nums.length < 2) return nums;  
  
  const mid = Math.floor((nums.length) / 2)  
  
  const left = sortArray(nums.slice(0, mid));  
  const right = sortArray(nums.slice(mid));  
  
  return merge(left, right);  
};
```

```
const merge = function(left, right) {  
  let merged = [];  
  
  while (left.length && right.length) {  
    if (left[0] < right[0]) {  
      merged.push(left.shift());  
    } else {  
      merged.push(right.shift());  
    }  
  }  
  
  return merged.concat(...left, ...right);  
}
```

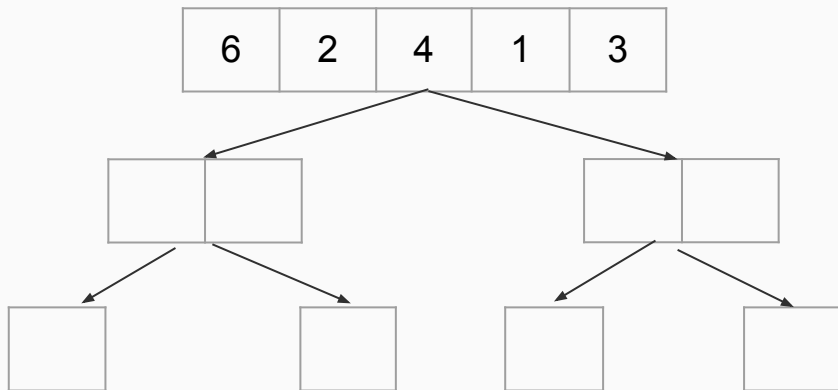


# Merge Sort (cont.)



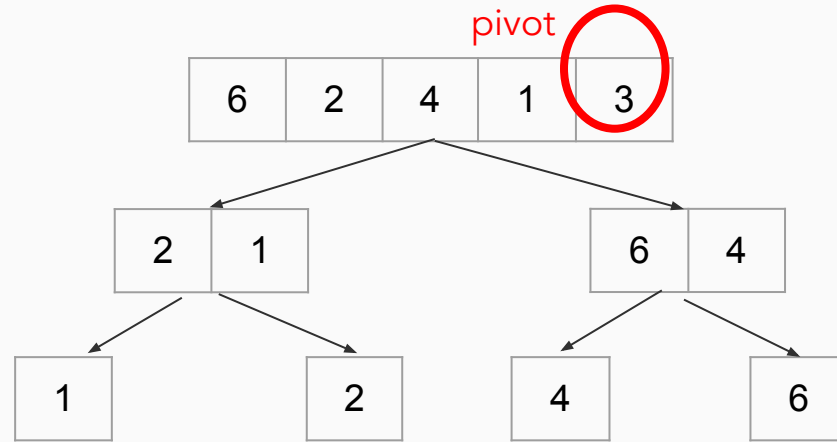
- Notice how the size of  $n$  is divided by 2 at each stage of the stack? That means the stack is  $\log(n)$  deep!
- Also notice that how wide this tree is the length of  $n$
- So we are doing  $n$  work  $\log(n)$  times. That's how we get  $n \cdot \log(n)$ !!!

# Quick Sort

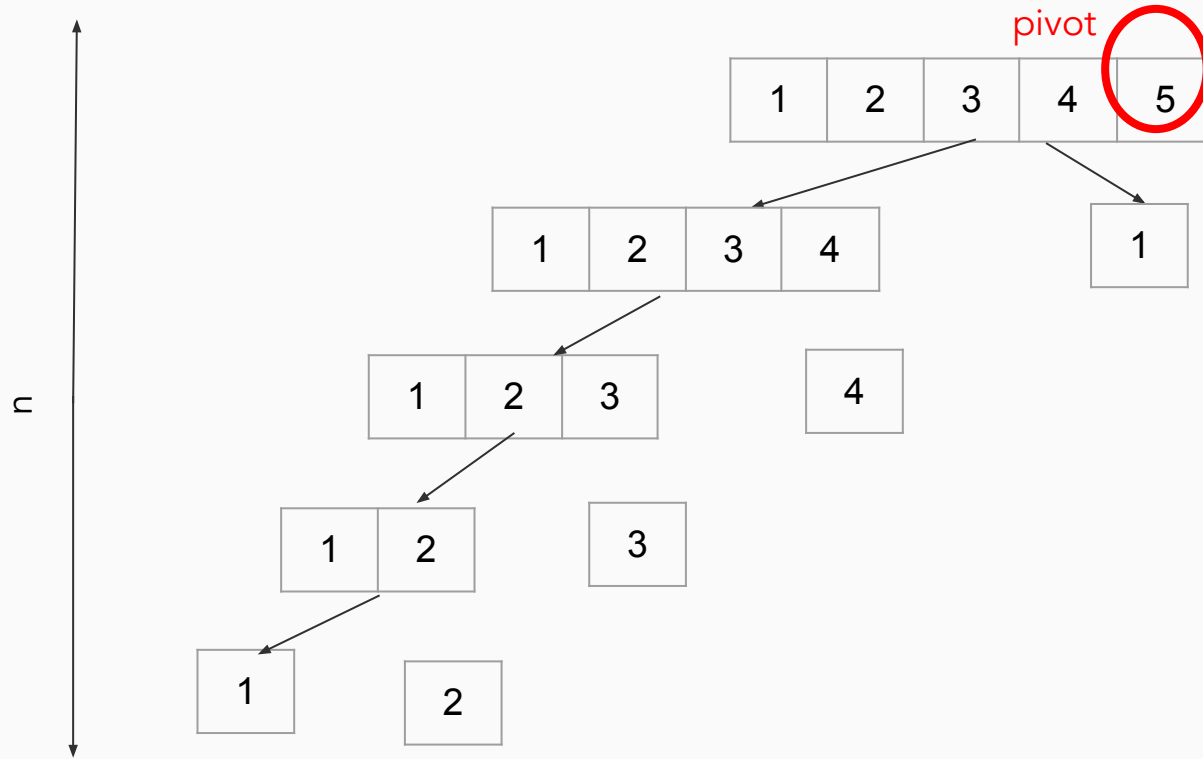


- Runtime:
  - WORST CASE:  $O(n^2)$
  - Amortized:  $(n \cdot \log(n))$
- Unlike merge sort which uses a temp array to merge, quicksort can sort IN PLACE with no extra memory.
- This algorithm selects a pivot and then partitions all values into “lesser” and “greater” buckets. It does this until there is only one element at most in either the lesser or greater sides and then concatenates them with the pivot.

# Quick Sort (cont.)



# Quick Sort - worst case





# Quick Sort (cont.)

```
const sortArray = function(nums, start = 0, end = nums.length-1) {  
  if (end - start + 1 <= 1) return nums;  
  
  let pivot = nums[end];  
  let left = start;  
  
  for (let i=start; i<=end; i++) {  
    if (nums[i] < pivot) {  
      tmp = nums[left];  
      nums[left] = nums[i];  
      nums[i] = tmp;  
      left++;  
    }  
  }  
  
  nums[end] = nums[left];  
  nums[left] = pivot;  
  
  sortArray(nums, start, left-1);  
  sortArray(nums, left+1, end);  
  
  return nums;  
}
```



# Bucket Sort

2	1	2	0	0	2
---	---	---	---	---	---

--	--	--

- Runtime:  $O(n)$
- ONLY ALLOWED UNDER CONSTRAINTS:
  - We must guarantee that all values fit within a finite range.
- As we iterate through the array, we can simply increment a value in an array.

# Bucket Sort

2	1	2	0	0	2
---	---	---	---	---	---

2	1	3
---	---	---

- Runtime:  $O(n)$
- ONLY ALLOWED UNDER CONSTRAINTS:
  - We must guarantee that all values fit within a finite range.
- As we iterate through the array, we can simply increment a value in an array.
- Demo: [Sort Colors](#)



# Questions?



# Let's practice!

- Review
  - [Sort Characters by Frequency](#)
  - [Top K Frequent Elements](#)

