



Backtracking



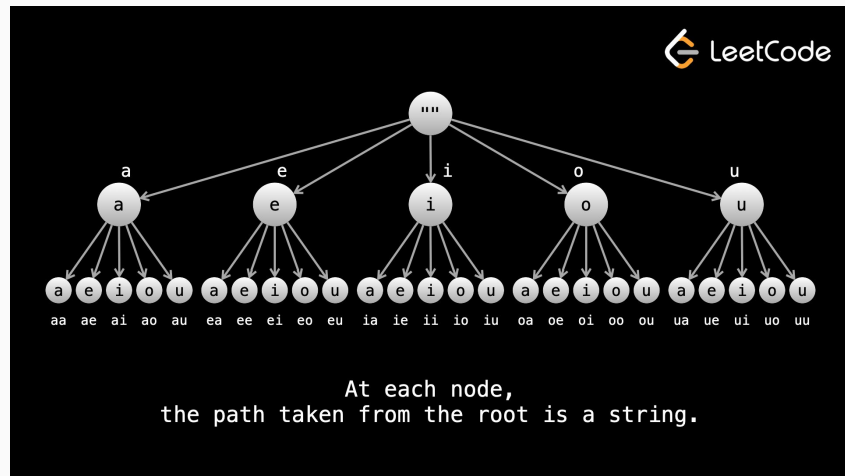
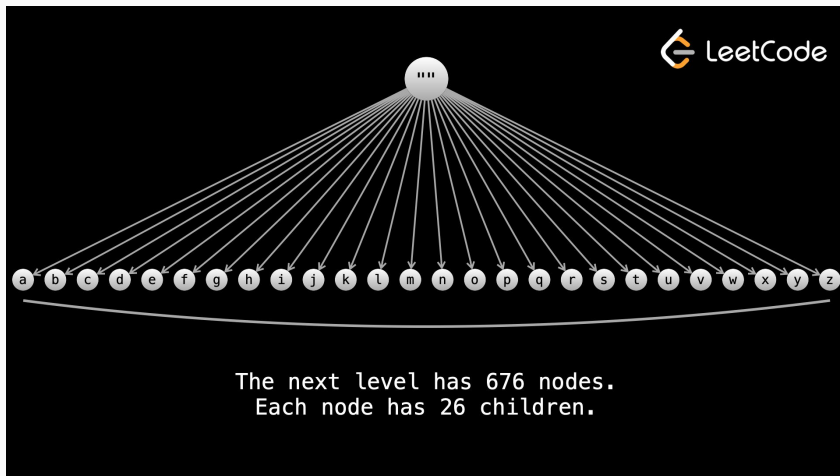
Backtracking

- Backtracking is a algorithmic technique where we solve problems recursively by “building candidates” for solutions and then abandon those candidates (“backtrack”) once we determine that candidate can no longer yield a valid solution.
- Optimization of “exhaustive search”, where we search through all possible candidates.
- This is an important concept because some problems can only be (reasonably) solved via backtracking.
- Most common type of problem that can be solved with backtracking is “find all possible ways to do something”
- Backtracking problems run in exponential runtime.



Example

Let's take a simple example where we're building all possible strings with a length of n . To check all possible strings, we'd be running a $O(26^n)$ runtime. But what if vowels-only was a constraint? We can discard all non-vowel subtrees and improve runtime to $O(5^n)$.



Backtracking skeleton

```
// let curr represent the thing you are building  
// it could be an array or a combination of variables
```

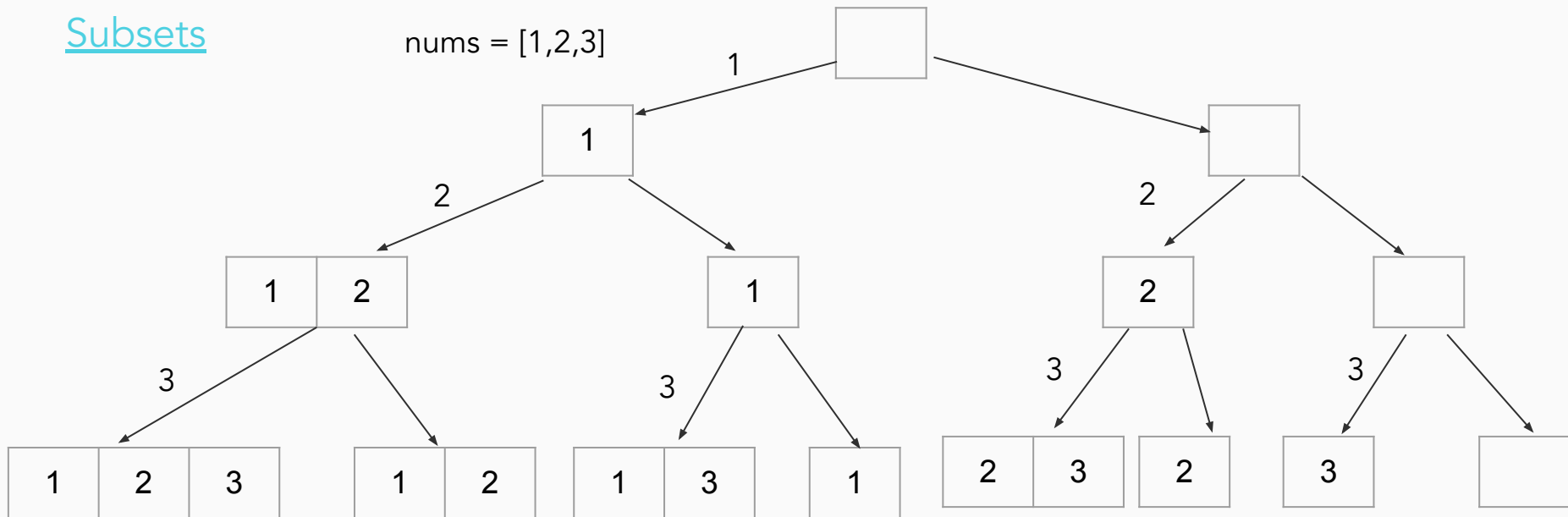
```
function backtrack(curr) {  
  if (base case) {  
    Increment or add to answer  
    return  
  }  
  
  for (iterate over input) {  
    Modify curr  
    backtrack(curr)  
    Undo whatever modification was done to curr  
  }  
}
```



Demo

Subsets

nums = [1,2,3]



For each node, there are two different paths we can take.
Either add `nums[i]` or skip it



Questions?



Let's practice!

- Review
 - [Combination Sum](#)
- Bonus
 - [N-Queens](#)

