Introduction to Real-World Graphs

Maximilien Danisch

LIP6 - CNRS and Sorbonne Université

first_name.last_name@lip6.fr

Outline

- Graphs and networks
- 2 Time and memory complexity
- 3 Algorithms

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Graph

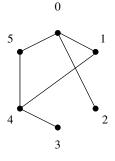
A graph G = (V, E) is a couple of sets.

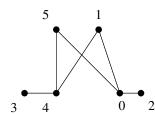
- V is a set of vertices (or nodes) fr: sommet ou nœud
- $E \subseteq (V \times V)$ is a set of edges (or links) fr: lien, arête

Example

•
$$V = \{0, 1, 2, 3, 4, 5\}$$

•
$$E = \{(0,1), (0,2), (3,4), (4,5), (5,0), (1,4)\}$$





Graph theory

Graph theory is a very important and well studied field:

https://en.wikipedia.org/wiki/Graph_theory

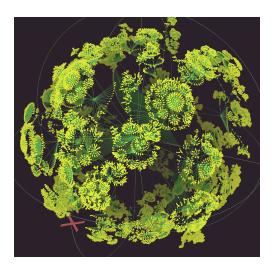
Question: Why is this useful to us?

Graph theory

Graph theory is a very important and well studied field: https://en.wikipedia.org/wiki/Graph_theory

Question: Why is this useful to us? Many systems can be modeled using graphs!

Internet: computers connected by internet connections



Twitter: profiles connected by follow-links



Many more examples

Graphs extracted from the "real-world" are called networks, complex networks or real-world graphs.

Question: Give some more examples of real-world graphs?

Common properties of real-world graphs

Very large:

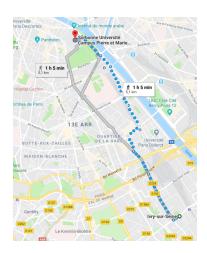
Several G edges

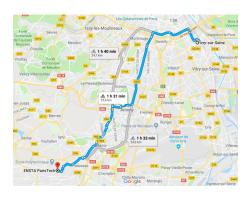
Sparse:

Not many edges compared to what could be

- Short distances and diameter:
 - $\sim \log(n)$; Six degrees of separation; Kevin Bacon game
- Heterogeneous degree distribution:
 - \sim power law; contains hubs (very high degree nodes)
- Dense locally (many triangles):
 High clustering coefficient & transitivity ratio
- **3** ...

A real-world problem: GPS





Many more examples

Question: Give some more real-world problems related to graphs?

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Time and memory complexity

Let *G* be a graph of *n* nodes and *m* edges.

Question: What do you think about the following complexities? $\mathcal{O}(m)$ and $\mathcal{O}(n^2)$

Storing a graph in main memory

Three main formats:

- List of edges
- Adjacency matrix
- Adjacency array

Question: Implementation? Pros and cons? Complexities?

https://github.com/maxdan94/LoadGraph

Adjacency Array in C

```
//compact adjacency array datastructure
typedef struct {
   unsigned n; //number of nodes
   unsigned m; //number of edges
   unsigned *cd; //cumulative degree cd[0]=0 length=n+1
   unsigned *adj; //concatenated lists of neighbors
} adjarray;
```

Question: Remarks? How to list the neighbors of node u?

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Breadth First Search algorithm (BFS)

Algorithm 1 BFS from node s

```
function BFS(G,s)

FIFO \leftarrow CreateFIFO()

FIFO.add(s)

Mark(s)

while FIFO not empty do

u \leftarrow FIFO.pop()

output u

for each v neighbor of u in G do

if v is not marked then

FIFO.add(v)

Mark(v)
```

Question: Implementation? Complexity?

Breadth First Search algorithm (BFS)

- Question: How to compute the diameter of a graph? Complexity?
- Question: How to compute a good lower bound to the diameter using only a few BFS? Complexity?
- Question: How to generalize BFS if edges have weights?

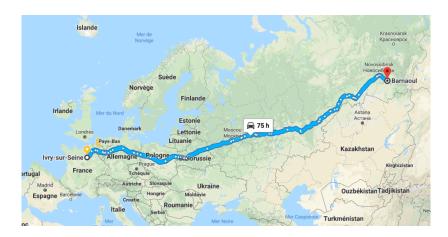
Dijkstra algorithm

Algorithm 2 Dijkstra on a weighted graph G_w from node s

```
function DIJKSTRA(G_w,s)
for each node u in G_w do
D[u] \leftarrow \mathsf{INFINITY}
D[s] = 0
while G_w not empty do
u \leftarrow \mathsf{node} \text{ in } G_w \text{ with minimum } D[u]
output u
for each v neighbor of u in G do
tmp \leftarrow D[v] + w_{uv}
if tmp < D[v] \text{ then}
D[v] = tmp
remove u from G_w
```

Question: Remarks? Complexity?

Dijkstra algorithm



Question: Is Dijkstra fast enough to answer that query in 1sec?

Listing triangles

Question: Suggest an algorithm to list all triangles?

(Fast) computation of the number of triangles

Algorithm: for each edge (u, v) in the graph, compute the intersection of the neighborhoods of u and v.

Remarks:

- Algorithm slow if u or v have high degree
- Triangle (u, v, w) is found 3 times

Idea: reduce as much as possible the size of the lists involved in the computation of the intersection:

- sort the nodes in non-increasing order of degree and re-index the graph (such that $u < v \Rightarrow d(u) \geq d(v)$)
- given an edge (u, v), we consider only the neighbors w of u (resp. v) such that u < w (resp. v < w)

(Fast) computation of the number of triangles

Algorithm 3 List all triangles in *G*

```
function TRLIST(G)
for each node u in G do
tsl[u] \leftarrow truncated and sorted list of neighbors of u
for each edge (u, v) in G do
W \leftarrow Intersect(tsl[u], tsl[v])
for each node w in W do
output triangle \{u, v, w\}
```

Question: Remarks? Complexity?

Listing k-cliques

Question: Given an integer *k*, suggest an algorithm to list all k-cliques?

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https://papers-gamma.link/paper/32