CSC 374/407: Computer Systems II

Lecture 5
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Reading

- Bryant & O'Hallaron "Computer Systems, 2nd Ed."
 - Chapter 13.1-13.5: Concurrent Programming
- Hoover "System Programming"
 - · (None)

Topics

- What threads are and why use them
- How to create threads
- Critical sections and unsafe thread programming
- Synchronization
- Application: Producer/Consumer

What threads are and why use them

Purpose: to have multiple "little processes" parallelize action of one program

Threads are like processes, but they share:

- Code segment
- Heap segment
- Data/BSS segments
- Stack segment

They differ in that:

- Have different registers (including condition codes)
- Start in different regions of stack
- Have unique thread id

How to compile for threads

At the beginning of your program:

```
#include <thisHeader.h>
#include <pthread.h>
#include <thatHeader.h>
```

When you link:

```
unix> gcc . . . -lpthread . . .
```

Why are they called "p-threads?"

- P = POSIX = Portable Operating Sys Interface for UniX
- Standardizes "Unix" across Linux, BSD, Solaris, etc.

How to create threads

```
int pthread create
    (pthread t* restrict
                                   thread,
      /* Pointer to a pthread t object to identify
       the child thread */
     const pthread attr t* restrict attr,
      /* Pointer to optional object for properties
       of child. You can just say NULL. */
     void *(*fncName)(void*),
      /* Name of function to run:
         void* fncName(void* ptr) */
     void *restrict
                                    arq
      /* Ptr to object that is arg to fncName() */
Return value is:
     - 0: success
     -anything else: ERROR!
```

How to wait for threads

Compare threads with fork:

```
/* I use fork() to create a process */
#include <stdlib.h>
#include <stdio.h>
#include <sys/types.h>
#include <sys/wait.h>
int main()
      childsStatus;
int
pid t childId
              = fork();
char* argsText = /* whatever */
if (childId == 0)
  execlp("doThisFile",..,argText,.. );
waitpid(childId,&childStatus,0);
return(0);
```

```
/* I use pthread_create() */
#include <stdlib.h>
#include <stdio.h>
#include <pthread.h>
int main()
pthread t
                 childId:
SendObject
                 arg = /*whatever*/
ReceiveObject*
                 childsStatusPtr;
pthread create(&childId,NULL,
                doThisFnc,&arg
pthread_join(childId,&childStatusPtr);
return(0);
```

Example program:

```
/* Compile with:
* gcc -lpthread thread ex1.c -o thread ex1
*/
#include <stdlib.h>
#include <stdio.h>
#include <string.h>
#include <pthread.h>
const int N = 2;
/* thread routine */
void *thread_routine(void *vargp)
  int id = *(int*)vargp;
  char* ptr;
  printf("Hello from child thread %d\n", id);
  switch (id)
    case 0: ptr = strdup("Hello "); break;
    case 1: ptr = strdup("there!"); break;
  return(ptr);
```

```
int main()
  int
  char*
             msqPtr;
  pthread t tid[N];
  for (i = 0; i < N; i++)
     pthread create
        (&tid[i], NULL,
        thread routine, (void *)&i);
  for (i = 0; i < N; i++)
     pthread_join(tid[i], (void**)&msgPtr);
     puts(msgPtr);
     free(msgPtr);
  return(0);
```

Time for you!

Can you write a threaded program that can:

- 1. **Prove** that all threads use the same **stack** (did we just do that?)
- 2. **Prove** that all threads use the same **global var data space** (for global vars and static vars inside functions)
- 3. **Prove** that all threads use the same **heap**.

What's wrong with this?

```
/* Compile with:
* gcc -lpthread badcnt.c -o badcnt
*/
#include <stdlib.h>
#include <stdio.h>
#include <pthread.h>
unsigned int cnt = 0; /* shared */
unsigned int NUM ITERS;
/* thread routine */
void *count(void *arg)
  int i:
  for (i=0; i<NUM ITERS; i++)
     cnt++;
  return NULL:
```

```
int main(int argc, char* argv[])
  pthread t tid1, tid2;
  if ((argc \ge 2) \&\& isdigit(*argv[1]))
     NUM_ITERS = atoi(argv[1]);
   else
     const int LINE SIZE = 16;
     char line[LINE SIZE];
     do {
       printf("How many iterations? ");
       fgets(line,LINE SIZE,stdin);
    while (!isdigit(line[0]));
     NUM ITERS = atoi(line);
  pthread create(&tid1, NULL, count, NULL);
  pthread create(&tid2, NULL, count, NULL);
  pthread_join(tid1, NULL);
  pthread_join(tid2, NULL);
  printf("Should be %d is %d\n",NUM ITERS*2,cnt);
  return(0);
```

Output:

[jphillips@localhost]\$ badcnt 10
Should be 20 is 20
[jphillips@localhost]\$ badcnt 100
Should be 200 is 200
[jphillips@localhost]\$ badcnt 1000
Should be 2000 is 2000
[jphillips@localhost]\$ badcnt 10000
Should be 20000 is 20000
[jphillips@localhost]\$ badcnt 100000
Should be 200000 is 200000
[jphillips@localhost]\$ badcnt 100000
Should be 20000000 is 1830352
[jphillips@localhost]\$ badcnt 1000000
Should be 200000000 is 15214384

Last two values are WRONG!

Critical sections

```
(gdb) disass count
0x400688:
          push
                %rbp
0x400689: mov
                %rsp,%rbp
                %rdi,0xffffffffffe8(%rbp)
0x40068c: mov
0x400690: movl
                $0x0,0xfffffffffff(%rbp)
0x400697: jmp
               0x4006ac <count+36>
0x400699: mov 2098629(%rip),%eax # 0x600c64 <cnt>
0x40069f: add
              $0x1,%eax
0x4006a2: mov %eax,2098620(%rip) # 0x600c64 <cnt
0x4006a8: addl $0x1,0xffffffffffff(%rbp)
               Oxfffffffffff(%rbp),%edx
0x4006ac: mov
              2098611(%rip),%eax #0x600c68 < NUM ITERS>
0x4006af: mov
0x4006b5: cmp
               %eax,%edx
0x4006b7: jb
              0x400699 < count+17>
0x4006b9: mov
               $0x0,%eax
0x4006be: leaveg
0x4006bf: retq
```

These three must be done by one thread at a time!

Critical sections (2)

Imagine some scenarios where they are **not** done atomically

Other ways threads can step on each other's toes

- 1. Working with the same global variable
 - (Just covered)
- 2. Working with the same static var within a fnc.
 - Can you think of an example?
- 3. A function returning a pointer to a static var.
 - Can you think of an example?
- 4. Calling a thread-unsafe function.
 - Can you think of an example?

Functions: safe and unsafe

Reentrant functions call no shared variables

Always thread safe

Most standard C Library functions are thread safe

- e.g. printf(), malloc()

Unsafe to safe

- asctime() (unsafe!)
- ctime() (unsafe!)
- gethostbyaddr() (unsafe!)
- inet_ntoa(unsafe!)
- localtime() (unsafe!)
- rand() (unsafe!)

```
asctime_r() (safe!)
ctime_r() (safe!)
aethostbyaddr_r()(safe!)
```

gethostbyaddr_r()(safe!)

NO SAFE VERSION!

localtime_r() (safe!)

rand_r() (**safe!**)

The semaphore solution

Classic solution: Dijkstra's P and V operations on semaphores.

- semaphore: non-negative integer synchronization variable.
 - sem_wait(s): [while(s==0) wait(); s--;]
 Originally named P(), Dutch for "Proberen" (test)
 - sem_post(s): [s++;]
 - Originally named V(), Dutch for "Verhogen" (increment)
- OS guarantees that operations between brackets
 [] are executed indivisibly.
 - Only one P or V operation at a time can modify s.
 - When while loop in P terminates, only that P can decrement s.

Semaphore invariant: $(s \ge 0)$

(1) POSIX Semaphores

What to include:

- #include <semaphore.h>

Types and functions:

- sem_t semaphore;
- sem_init(sem_t* semPtr, int flag, int value)
 - Initialize pointed-to semaphore, with value, if flag
 == 1 then semaphore can be forked
- sem_destroy(sem_t* semPtr)
 - Destroy pointed-to semaphore. If it's negative then block.

POSIX Semaphores, cont'd

- sem_wait(sem_t* semPtr)
 - Decrement pointed-to semaphore. If it's negative then block.
- sem_post(sem_t* semPtr)
 - Increment pointed-to semaphore. Wake one blocked process if any.
- sem_getvalue(sem_t* semPtr, int* valuePtr)
 - Get value of pointed to semaphore.

POSIX semaphore solution

```
sem t
             sem:
/* P operation on semaphore sem */
void P(sem t *semPtr) {
 if (sem_wait(semPtr)) {
  fputs("sem_wait error\n",stderr);
  exit(-1);
/* V operation on semaphore sem */
void V(sem t *semPtr) {
 if (sem_post(semPtr)) {
  fputs("sem post error\n",stderr);
  exit(-1);
```

```
/* thread routine */
void *count(void *arg) {
  int i;
  for (i=0; i<NUM ITERS; i++) {
     P(&sem);
     cnt++;
     V(&sem);
  return NULL;
/* in main */
/* initialize sem to 1 */
  if (sem_init(\&sem, 0, 1) < 0) {
     fputs("sem_init error\n",stderr);
     exit(1);
```

(2) pthread_mutex solution

- pthread_mutex_t lock
- pthread_mutex_init(addressOfLock,NULL)
 - Makes a lock and initializes it to default values.
- pthread_mutex_destroy(addressOfLock)
 - Makes a lock and initializes it to default values.
- pthread_mutex_lock(address of lock)
 - Blocks thread until lock obtained, then obtains lock and blocks other threads until lock released.
- pthread_mutex_unlock(address of lock)
 - Releases lock allowing other threads to obtain it.

pthread_mutex solution

```
pthread mutex t cnt lock;
/* thread routine */
void *count(void *arg) {
  pthread_mutex_t* mutexLockPtr;
  mutexLockPtr = (pthread mutex t*)arg;
  for (int i=0; i<NUM ITERS; i++) {
     pthread_mutex_lock(mutexLockPtr);
    cnt++:
     pthread mutex unlock(mutexLockPtr);
  return NULL;
/* in main() */
pthread_mutex init(&cnt lock,NULL);
```

[jphillips@localhost]\$ bettercnt 10 Should be 20 is 20 [jphillips@localhost]\$ bettercnt 100 Should be 200 is 200 [jphillips@localhost]\$ bettercnt 1000 Should be 2000 is 2000 [jphillips@localhost]\$ bettercnt 10000 Should be 20000 is 20000 [jphillips@localhost]\$ bettercnt 100000 Should be 200000 is 200000 [jphillips@localhost]\$ bettercnt 1000000 Should be 2000000 is 2000000 [jphillips@localhost]\$ bettercnt 10000000 Should be 20000000 is 20000000 [jphillips@localhost]\$ bettercnt 100000000 Should be 200000000 is 200000000

Producer-Consumer

One or more threads produce something

Place in buffer

One or more threads consume something

Retrieve from buffer

Critical section

Access of pointers/indices for buffer

Problem!

- A producer may gain access to a full buffer
- A consumer may gain access to an empty buffer

Oh no! An example in <u>C++</u>!

- A buffer is an <u>object</u>.
- Make the <u>object</u> thread-safe.
- Lessons carry over to Java, C#, etc.

Unsafe_Buffer.h

```
class Buffer
  enum \{ SIZE = 16 \};
  int array [SIZE];
  int inIndex ;
  int outIndex ;
  int numItems ;
public:
  Buffer
                ()
    inIndex = outIndex
      = num\overline{I}tems = 0;
  ~Buffer
                ()
```

```
int
     getNumItems () const
{ return(numItems ); }
void putIn (int i)
 while (getNumItems() >= SIZE)
    printf("Full! Waiting!\n");
   usleep(10);
  array [inIndex ] = i;
  countArray[array [inIndex ]]++;
  usleep(10 + rand() % 10);
  inIndex ++;
  numItems ++;
  if (inIndex >= SIZE)
    inIndex = 0;
```

Unsafe_Buffer.h

```
int pullOut ()
 while (getNumItems() <= 0)</pre>
   printf("Empty! Waiting!\n");
   usleep(10);
  }
  countArray[array_[outIndex_]]--;
 int toReturn = array_[outIndex_];
 usleep(10 + rand() % 10);
  outIndex ++;
 numItems --;
 if (outIndex >= SIZE)
   outIndex = 0;
  return(toReturn);
```

```
#include
                <stdlib.h>
#include
                <stdio.h>
#include
             <unistd.h>
#include
                <pthread.h>
int*
                countArray;
                "Unsafe Buffer.h"
#include
const int
                NUM INTEGERS TO BUFFER = 0 \times 1000;
void* stuffIntegersIn (void* vPtr)
  for (int i = 0; i < NUM INTEGERS TO BUFFER;
                                                  i++)
    ((Buffer*)vPtr)->putIn(i);
  return(NULL);
}
```

```
void* pullIntegersOut (void* vPtr)
{
  for (int i = 0; i < NUM_INTEGERS_TO_BUFFER; i++)
  {
    int j = ((Buffer*)vPtr)->pullOut();
    printf("Trial %d got %d.\n",i,j);
    fflush(stdout);
  }
  return(NULL);
}
```

```
int
       main
                ()
 pthread t
                producer0;
 pthread t
                producer1;
 pthread t
                consumer0;
 pthread t
                consumer1;
 Buffer
                buffer;
 countArray
  (int*)calloc(NUM INTEGERS TO BUFFER, sizeof(int));
 pthread create(&producer0, NULL, stuffIntegersIn, &buffer);
 pthread create(&producer1,NULL,stuffIntegersIn,&buffer);
 pthread create(&consumer0, NULL, pullIntegersOut, &buffer);
 pthread create(&consumer1, NULL, pullIntegersOut, &buffer);
```

```
pthread join(producer1,NULL);
 pthread join(producer0,NULL);
 pthread join(consumer1,NULL);
 pthread join(consumer0, NULL);
 for (int i = 0; i < NUM INTEGERS_TO_BUFFER; i++)</pre>
    if (countArray[i] < 0)</pre>
      printf("%d was gotten too many times!\n",i);
    else
    if (countArray[i] > 0)
      printf("%d was put too many times!\n",i);
 return(EXIT SUCCESS);
}
```

Your turn!

(1) Run it as-is. Any problems?

Your turn!

(2) Make it thread-safe.

```
void myMethod ()
{
   pthread_mutex_lock(&lock);
   doCriticalSection();
   pthread_mutex_unlock(&lock);
}
```

NOTE:

- 1) Each object gets its **own** lock.
- 2) All threads accessing the same object use same lock

Any problems still?

Solution: pthread_cond!

- pthread cond t cond
- pthread cond init(addressOfCondition, NULL)
 - Makes a condition and initializes it to default values.
- pthread cond destroy(addressOfCondition)
 - Destroys condition.
- pthread cond wait(addrOfCondition, addrOfLock)
 - Blocks thread until lock released and condition signaled. Then obtains lock again.
- pthread_cond_signal(addressOfCondition)
 - Signal one waiting thread that condition is met.
- pthread cond broadcast(addressOfCondition)
 - Signal all waiting threads that condition is met.

The <u>proper</u> solution

```
void myMethod ()
{
  pthread_mutex_lock(&lock);
  while ( !this->isReady() )
    pthread_cond_wait(&wCond,&lock);
  doCriticalSection();
  pthread_cond_signal(&sCond);
  pthread_mutex_unlock(&lock);
}
```

NOTE:

- (1) Method **isReady()** determines if this object is ready for thread to do **myMethod()**.
- (2) Conditions wcond and scond may or may not be same.

We have mutual exclusion, are guaranteed to be safe?

Sorry but No!

- Race conditions
 - When result is dependent on order of processes or threads
- Deadlock
 - Thread/process A has lock 1 and is waiting for lock 2
 - Thread/process B has lock 2 and is waiting for lock 1

Next time: Memory!