# 综合混练

## 人员

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## 上周作业检查

上周作业链接: https://vjudge.net/contest/747412



# 本周作业

https://vjudge.net/contest/751260, (课上讲了上周比赛的 B C D E F G 题, 课后作业是本周比赛的 A B C D E 题)

### 课堂表现

今天的 C 题和 G 题要稍微复杂一些, 但是这两个题知识点都不难, 一个是搜索, 一个是二分+前缀和, 大家要认真补一补这两道题, 不要觉得麻烦就不做了

# 课堂内容

#### CF1245D Shichikuji and Power Grid

建一个虚拟源点 n+1 号点, 他与 i 号点的边权是 c[i], i 与 j 的边权是 dis(i,j)\*(k[i]+k[j])

此时, 原问题被转化为一个最小生成树问题

这个图是一个完全图, 因此建议用 prim 算法实现

最后, 需要输出选择哪些点或哪些路径时, 可以用一个数组记录一下每个点的 d 值是由谁转移而来的

如果是由 n+1 号点转移而来的, 那么就是单独点这一个点; 如果不是 n+1 号点转移而来, 就相当于连接了上一个点和这一个点

```
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
const int maxn = 2000 + 5;
const LL inf = 0x3f3f3f3f3f3f3f3f3f;
int x[maxn], y[maxn], c[maxn], k[maxn], pre[maxn];
LL w[maxn][maxn], d[maxn];
bool st[maxn];
int dis(int a, int b) { return abs(x[a]-x[b]) + abs(y[a]-y[b]); }
int main()
  int n; cin >> n;
  for (int i = 1; i <= n; ++i) cin >> x[i] >> y[i];
  for (int i = 1; i <= n; ++i) cin >> c[i];
  for (int i = 1; i <= n; ++i) cin >> k[i];
  for (int i = 1; i <= n+1; ++i) {
   for (int j = 1; j <= n+1; ++j) {
      if (i == j) continue;
      if (i == n+1) w[i][j] = c[j];
      else if (j == n+1) w[i][j] = c[i];
      else w[i][j] = (LL)dis(i,j) * (k[i]+k[j]);
    }
  }
  LL res = 0;
  memset(d, 0x3f, sizeof(d));
  vector<int> vec1;
```

```
vector<pair<int,int>> vec2;
  for (int i = 0; i <= n; ++i) {
    int id = n+1;
    if (i) {
      for (int j = 1; j <= n+1; ++j) {
       if (st[j]) continue;
        if (id==n+1 || d[j] < d[id]) id = j;
      }
      res += d[id];
      if (pre[id] == n+1) vec1.push_back(id);
      else vec2.push_back({pre[id], id});
    }
    st[id] = true;
    for (int j = 1; j <= n+1; ++j) {
     if (id == j) continue;
     if (w[id][j] < d[j]) d[j] = w[id][j], pre[j] = id;</pre>
    }
  }
 cout << res << endl;</pre>
 cout << vec1.size() << endl;</pre>
 for (int i : vec1) cout << i << " ";
 cout << endl;</pre>
 cout << vec2.size() << endl;</pre>
 for (pair<int,int> it : vec2) cout << it.first << " " << it.second << endl;</pre>
  return 0;
}
```

#### UVA1599 理想路径 Ideal Path

先从 n 号点做一遍 bfs, 求出来每个点到 n 号点的最短距离 dis[i]

然后从 1 号点往后搜索遍历, 在保证最短路径的前提下, 找到颜色字典序最小的情况

```
#include <bits/stdc++.h>

using namespace std;

const int maxn = 1e5 + 5;
struct node {
   int to, c;
};
vector<node> vec[maxn];
int n, m;
int dis[maxn];
vector<int> ans;
bool st[maxn];

void bfs() {
   queue<int> q; q.push(n); dis[n] = 0;
```

```
while (!q.empty()) {
    int u = q.front(); q.pop();
    for (node it : vec[u]) {
      if (dis[it.to] == -1) q.push(it.to), dis[it.to] = dis[u]+1;
  }
}
void solve() {
  queue<int> q; q.push(1); st[1] = true;
  while (!st[n]) {
    vector<int> points;
    while (!q.empty()) points.push_back(q.front()), q.pop();
    int min_color = 1e9+10;
    for (int i : points) {
     for (node it : vec[i]) {
        if (dis[it.to] == dis[i]-1) min_color = min(min_color, it.c);
      }
    }
    ans.push_back(min_color);
    for (int i : points) {
      for (node it : vec[i]) {
        if (dis[it.to]==dis[i]-1 && !st[it.to] && it.c==min_color) q.push(it.to),
st[it.to] = true;
      }
    }
  }
}
int main()
  ios::sync_with_stdio(false);
  cin.tie(0);
  while (cin >> n >> m) {
    for (int i = 1; i <= n; ++i) {
      vec[i].clear(); dis[i] = -1; st[i] = false;
    }
    ans.clear();
    for (int i = 1; i <= m; ++i) {
      int a, b, c; cin >> a >> b >> c;
      vec[a].push_back({b,c}), vec[b].push_back({a,c});
    bfs(); solve();
    cout << dis[1] << "\n";</pre>
    for (int i = 0; i < (int)ans.size(); ++i) {
     if (i) cout << " ";
      cout << ans[i];</pre>
    }
```

```
cout << "\n";
}
return 0;
}</pre>
```

#### **CF148E Porcelain**

可以先预处理, 求出来每一行分别选 1个, 2个, 3个, ... 时能获得的最大价值

接下来, 就是在每一行挑几个, 问正好选够 m 个时, 能获得的最大价值是多少, 问题转化为了一个分组背包的板子题

```
#include <bits/stdc++.h>
using namespace std;
const int N = 100 + 5, M = 10000 + 5;
int len[N], w[N], pre[N], c[N][N], f[M];
int get_sum(int 1, int r) { return pre[r] - pre[1-1]; }
void solve(int id) {
 cin >> len[id];
 for (int i = 1; i \leftarrow len[id]; ++i) cin >> w[i], pre[i] = pre[i-1]+w[i];
  for (int i = 1; i <= len[id]; ++i) {
   for (int j = 0, k = i; j <= i; ++j, --k) {
      c[id][i] = max(c[id][i], pre[j] + get_sum(len[id]-k+1,len[id]));
    }
}
int main()
  int n, m; cin >> n >> m;
 for (int i = 1; i \le n; ++i) solve(i);
  for (int i = 1; i <= n; ++i) {
   for (int j = m; j \ge 0; --j) {
      for (int k = 1; k \le len[i]; ++k) {
        if (j-k \ge 0) f[j] = max(f[j], f[j-k]+c[i][k]);
      }
    }
  cout << f[m] << endl;</pre>
  return 0;
}
```

Gym - 103428G Shinyruo and KFC

#### 最多只有 sqrt(n) 种不同的餐品数量, 每种餐品可以 O(k) 来做

```
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
const int maxn = 1e5 + 5;
const int mod = 998244353;
int w[maxn], fac[maxn], inv_fac[maxn], c[maxn];
int qmod(int a, int k) {
 int res = 1;
 while (k) {
   if (k&1) res = (LL)res*a % mod;
   a = (LL)a*a \% mod;
    k >>= 1;
  }
 return res;
}
int inv(int x) { return qmod(x, mod-2); }
int C(int n, int m) {
 if (n < m) return 0;
 return (LL)fac[n]*inv_fac[m]%mod*inv_fac[n-m]%mod;
}
struct node {
 int val, cnt;
};
int main()
{
  fac[0] = 1;
 for (int i = 1; i < maxn; ++i) fac[i] = (LL)fac[i-1]*i % mod;
 inv_fac[maxn-1] = inv(fac[maxn-1]);
 for (int i = maxn-2; i \ge 0; --i) inv_fac[i] = (LL)inv_fac[i+1]*(i+1) % mod;
  int n, m; cin >> n >> m;
  for (int i = 1; i <= n; ++i) cin >> w[i], c[w[i]]++;
 vector<node> vec;
  for (int i = 1; i < maxn; ++i) {
   if (c[i]) vec.push_back({i, c[i]});
  for (int k = 1; k <= m; ++k) {
   int res = 1;
   for (auto it : vec) {
     int val = it.val, cnt = it.cnt;
      res = (LL)res * qmod(C(k,val), cnt) % mod;
```

```
cout << res << "\n";
}
return 0;
}</pre>
```

#### **CF2026D Sums of Segments**

二分 + 前缀和, 先预处理出每一段的和以及每一段的长度, 后面用 二分查找 定位到第 x 个数在哪一段, 然后可以 O(1) 求前面的和

```
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
const int maxn = 3e5 + 5;
int w[maxn], len[maxn];
LL pre_w[maxn], pre_len[maxn];
LL cw[maxn], pre_cw[maxn];
LL seg[maxn], pre_seg[maxn];
int n;
LL get_sum(int 1, int r, LL a[]) { return (l<=r ? a[r]-a[l-1] : 0); }
LL calc(LL x) {
 if (!x) return ∅;
 int k = lower_bound(pre_len+1, pre_len+n+1, x) - pre_len;
  LL res = pre_seg[k-1];
 int nums = x - pre_len[k-1];
 res += get_sum(k, k+nums-1, pre_cw);
  res -= get_sum(k, k+nums-1, pre_w)*(n-k+1-nums);
 return res;
}
int main()
{
 cin >> n;
  for (int i = 1; i <= n; ++i) {
   cin >> w[i], pre_w[i] = pre_w[i-1] + w[i];
    cw[i] = (LL)w[i]*(n-i+1), pre_cw[i] = pre_cw[i-1] + cw[i];
    len[i] = n-i+1, pre_len[i] = pre_len[i-1] + len[i];
  for (int i = 1; i <= n; ++i) {
    seg[i] = get_sum(i, n, pre_cw), pre_seg[i] = pre_seg[i-1] + seg[i];
  }
  int m; cin >> m;
```

```
while (m -- ) {
   LL l, r; cin >> l >> r;
   cout << calc(r) - calc(l-1) << "\n";
}
return 0;
}</pre>
```