

综合练习

人员

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上周作业检查

上周作业链接: <https://cppoj.kids123code.com/contest/753>

2025-0907 周日8:30 (逆元)										
刷新										
#	用户名	姓名	编程分	时间	A	B	C	D	E	F
1	yangjunyan	杨俊彦	600	16999	100	100	100	100	100	100
2	yuanchenjun	袁晨峻	500	2297	100	100	100	100	100	
3	yangyongcheng	杨咏丞	500	2448	100	100	100	100	100	
4	liuyichen	刘奕辰	500	2597	100	100	100	100	100	
5	lijinshu	李锦澍	500	2653	100	100	100	100	100	
6	liyuanqian	李雨谦	400	1604	100	100	100	100		
7	wangluwenlong	王陆文龙	400	1691	100	100	100	100		
8	zhouzhirun	周治润	400	1716	100	100	100	100		
9	xuruiqian	许睿谦	300	509	100	100	100	0		
10	caoyuan	曹塬	300	570	100	100	100			
11	yuyue	于跃	233	49	100	33	100			
12	suitianyi	隋天乙	100	18	0		100	0		

本周作业

<https://cppoj.kids123code.com/contest/792> (课上讲了 A ~ E 题, 课后作业是 F 题)

课堂表现

今天上课先做了一道完全背包的题目, 能看出全班大部分同学把之前背包的内容已经遗忘很多了, 同学们课下要再去复习复习之前学的背包题目。

课堂内容

A Piece of Cake (上周作业)

map 套 pair 维护

判断一个草莓在哪一块蛋糕上, 这个过程可以用 二分查找 $\log n$ 来实现, 然后用 map 套 pair 记录这块蛋糕上草莓数量 +1

```
#include <bits/stdc++.h>
#define x first
#define y second
```

```

using namespace std;

typedef pair<int,int> PII;
const int maxn = 2e5 + 5;
int a[maxn], b[maxn];
struct node {
    int x, y;
} w[maxn];

int main()
{
    int W, H; cin >> W >> H;
    int T; cin >> T;
    for (int i = 1; i <= T; ++i) cin >> w[i].x >> w[i].y;

    int n; cin >> n;
    for (int i = 1; i <= n; ++i) cin >> a[i];
    int m; cin >> m;
    for (int i = 1; i <= m; ++i) cin >> b[i];

    map<PII, int> mp;
    for (int i = 1; i <= T; ++i) {
        int x = w[i].x, y = w[i].y;
        int tx = lower_bound(a+1, a+n+1, x) - a, ty = lower_bound(b+1, b+m+1, y) - b;
        mp[{tx,ty}]++;
    }

    int minn = 1e9, maxx = 0;
    if (mp.size() != (n+1)*(m+1)) minn = 0;
    for (auto it : mp) minn = min(minn, it.y), maxx = max(maxx, it.y);
    cout << minn << " " << maxx << endl;
    return 0;
}

```

疯狂的采药

完全背包 模板题

```

#include <bits/stdc++.h>

using namespace std;

typedef long long LL;
const int maxn = 1e7 + 5;
LL f[maxn];

int main()
{
    int T, n; cin >> T >> n;
    while (n -- ) {
        int a, b; cin >> a >> b;
    }
}

```

```
    for (int i = a; i <= T; ++i) f[i] = max(f[i], f[i-a]+b);
}
cout << f[T] << endl;
return 0;
}
```

纸币问题 3

完全背包 模板题

```
#include <bits/stdc++.h>

using namespace std;

const int maxn = 1e4 + 5;
const int mod = 1e9 + 7;
int f[maxn];

int main()
{
    int n, m; cin >> n >> m;
    f[0] = 1;
    for (int i = 1; i <= n; ++i) {
        int x; cin >> x;
        for (int j = x; j <= m; ++j) f[j] = (f[j] + f[j-x]) % mod;
    }
    cout << f[m] << endl;
    return 0;
}
```

纸币问题 2

先枚举背包体积, 在枚举纸币的类别即可

```
#include <bits/stdc++.h>

using namespace std;

const int maxn = 1e4 + 5;
const int mod = 1e9 + 7;
int w[maxn], f[maxn];

int main()
{
    int n, m; cin >> n >> m;
    for (int i = 1; i <= n; ++i) cin >> w[i];

    f[0] = 1;
    for (int i = 1; i <= m; ++i) {
```

```

        for (int j = 1; j <= n; ++j) {
            if (i >= w[j]) f[i] = (f[i] + f[i-w[j]]) % mod;
        }
    }
    cout << f[m] << endl;
    return 0;
}

```

Colorful Blocks

题目要求最多有 k 对相邻对的颜色相同, 因此可以从 $0 \sim n$ 枚举相邻对的数量

若有 i 对相邻, 那么有 $C(n-1, i) * m * (m-1)^{(n-i-1)}$ 种方案

```

#include <bits/stdc++.h>

using namespace std;

typedef long long LL;
const int maxn = 2e5 + 5;
const int mod = 998244353;
int fac[maxn], i_fac[maxn];

int qmod(int a, int k) {
    int res = 1;
    while (k) {
        if (k&1) res = (LL)res * a % mod;
        a = (LL)a * a % mod;
        k >>= 1;
    }
    return res;
}

void init() {
    fac[0] = i_fac[0] = 1;
    for (int i = 1; i < maxn; ++i) fac[i] = (LL)fac[i-1]*i % mod;
    for (int i = 1; i < maxn; ++i) i_fac[i] = qmod(fac[i], mod-2);
}

int C(int n, int m) {
    return (LL)fac[n] * i_fac[m] % mod * i_fac[n-m] % mod;
}

int main()
{
    init();

    int n, m, k; cin >> n >> m >> k;
    int res = 0;
    for (int i = 0; i <= k; ++i) {
        int t = (LL)C(n-1, i) * m % mod * qmod(m-1, n-i-1) % mod;
    }
}

```

```

    res = (res + t) % mod;
}
cout << res << endl;
return 0;
}

```

Count Bracket Sequences

dp, 设 $f[i][j]$ 代表 以第 i 位结尾, $1 \sim i$ 中 $(-)=j$ 时的方案数

答案是 $f[n][0]$, 初值是 $f[0][0] = 1$

转移分为 左括号转移 和 右括号转移 两种情况

```

#include <bits/stdc++.h>

using namespace std;

typedef long long LL;
const int maxn = 3000 + 5;
const int mod = 998244353;
char s[maxn];
int f[maxn][maxn]; // f[i][j]: 以第 i 位结尾, 1~i 中 ( - ) ==j 时的方案数

int main()
{
    cin >> (s+1);
    int n = strlen(s+1);

    f[0][0] = 1;
    for (int i = 1; i <= n; ++i) {
        for (int j = 0; j <= i; ++j) {
            if (s[i]=='(' || s[i]=='?') {
                if (j > 0) f[i][j] = f[i-1][j-1];
            }
            if (s[i]==')' || s[i]=='?') {
                f[i][j] = (f[i][j] + f[i-1][j+1]) % mod;
            }
        }
    }

    if (s[n] == '(') cout << 0 << endl;
    else cout << f[n][0] << endl;
    return 0;
}

```