

# Critical Points of Deep Linear Networks in $\mathbb{C}^N$

## Thesis Defense Presentation

by

Ayush Bharadwaj

Department of Mathematics

San Francisco State University

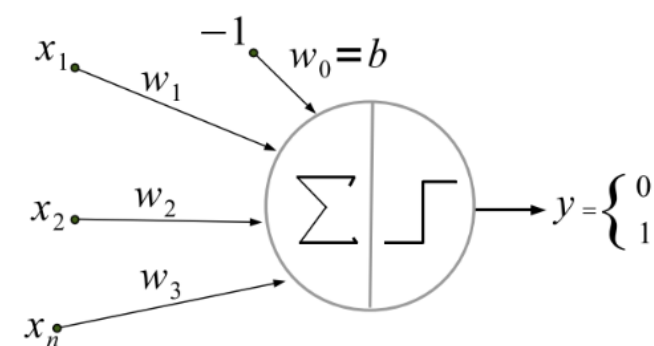
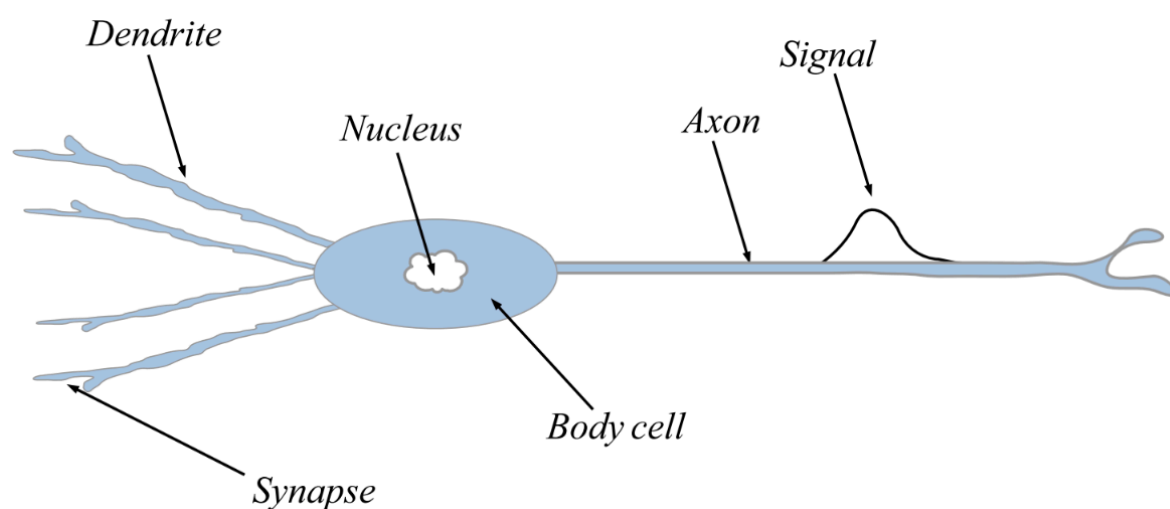
# Critical Points of Deep Linear Networks in $\mathbb{C}^N$

## Outline

- Neural networks (particularly, *linear networks*)
- Neural network training as a minimization problem in real analysis
- Neural network training as solving polynomial systems in algebraic geometry
- Results

# Neural Networks

## Neuron



- $f_w(x_1, \dots, x_n) = \phi \left( \sum_{j=1}^n w_j x_j - b \right)$  **(output function)**

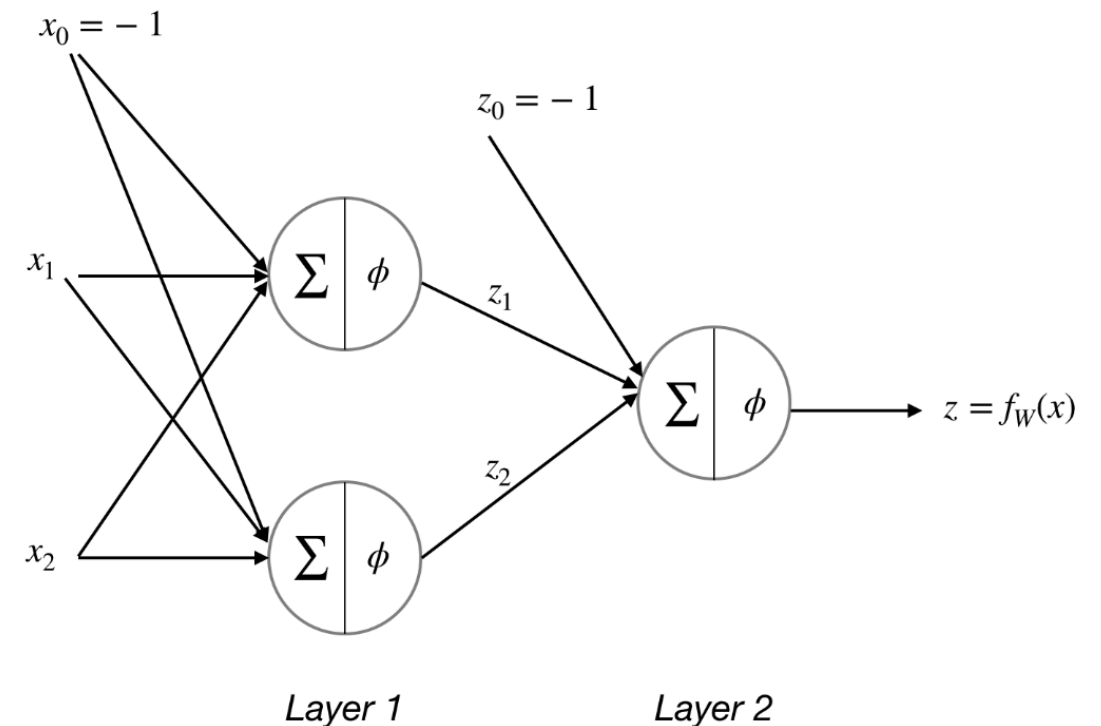
Image source: Ovidiu Calin. *Deep Learning Architectures - A Mathematical Approach*. Springer Series in the Data Sciences. Springer Cham, 2020. isbn: 978-3-030-36720-6.

# Neural Networks

## Neural network

- The weight matrix of the  $i$ th layer

$$W_i = \begin{pmatrix} w_{i11} & \cdots & w_{i1n} \\ \vdots & & \vdots \\ w_{ir1} & \cdots & w_{irn} \end{pmatrix}$$



We drop the bias term  $b$  and choose  $\phi = \mathbf{I}$     Neuron output  $f_w(x_1, \dots, x_n) = \sum_{j=1}^n w_j x_j$

This simplifies the network output function:

## Linear Network

$$f_W(x) = W_{H+1} W_H \cdots W_2 W_1 x \quad (2)$$

We want to use  $f_W$  to approximate some function  $y : \mathbb{R}^n \longrightarrow \mathbb{R}^p$

# Neural Networks

## Training a linear network

**Minimization problem in  
real analysis**

- Error (or *loss*),  $\mathcal{L} : \mathbb{R}^n \longrightarrow \mathbb{R}$  defined by:

$$\mathcal{L}(W) = \frac{1}{2} \sum_{i=1}^m \|W_{H+1}W_H \dots W_2W_1x^{(i)} - y^{(i)}\|^2 \quad (3)$$

- We want to solve the minimization problem:

$$\min_{W \in \mathbb{R}^N} \mathcal{L}(W) \quad (4)$$

**Not convex**

# Neural Networks

## Training a linear network

Non-convexity of  $\mathcal{L}(W)$  presents challenges

1. Hard to guarantee that all local minima have been found by the algorithm.
2. Hard to make assertions about the number and location of minima before hand.

**Motivation for an  
algebraic geometry view**

# Neural Networks

## Training a linear network

From a problem in Real Analysis to a problem in Algebraic Geometry

We introduce the following relaxations:

- $\min_{W \in \mathbb{R}^N} \mathcal{L}(W) \longrightarrow \nabla \mathcal{L}(W) = 0$
- $W \in \mathbb{R}^N \longrightarrow W \in \mathbb{C}^N$

We solve:

$$\boxed{\nabla \mathcal{L}(W) = 0, \quad W \in \mathbb{C}^N} \quad (5)$$

**Polynomial system**

# Neural Networks

## Training a linear network

From a problem in Real Analysis to a problem in Algebraic Geometry

*Example 1.2* (A 4-weight network). Let us consider a 2-layer network with  $W_1 = [\alpha_1, \alpha_2]$   $W_2 = \begin{bmatrix} \beta_1 \\ \beta_2 \end{bmatrix}$

- $\mathcal{L}(W) = \mathcal{L}(\alpha_1, \alpha_2, \beta_1, \beta_2) = \frac{1}{2} \sum_{i=1}^2 ||W_2 W_1 x^{(i)} - y^{(i)}||^2$
- $\nabla \mathcal{L}(W) = \nabla \mathcal{L}(\alpha_1, \alpha_2, \beta_1, \beta_2) = 0$  gives

**Nice polynomial system**

**Solutions are  
precisely  
the critical points**

$$5\alpha_1\beta_1^2 + 5\alpha_1\beta_2^2 + 11\alpha_2\beta_1^2 + 11\alpha_2\beta_2^2 - 7\beta_1 - 10\beta_2 = 0$$

$$11\alpha_1\beta_1^2 + 11\alpha_1\beta_2^2 + 25\alpha_2\beta_1^2 + 25\alpha_2\beta_2^2 - 15\beta_1 - 22\beta_2 = 0$$

$$5\alpha_1^2\beta_1 + 22\alpha_1\alpha_2\beta_1 + 25\alpha_2^2\beta_1 - 7\alpha_1 - 15\alpha_2 = 0$$

$$5\alpha_1^2\beta_2 + 22\alpha_1\alpha_2\beta_2 + 25\alpha_2^2\beta_2 - 10\alpha_1 - 22\alpha_2 = 0$$

(1.13)

- **Square**
- **Sparse**
- **Deg  $2H + 1$**



# Algebraic Geometry

## Main Ideas

Algebraic geometry provides results and methods to analyze and solve systems of polynomial equations. In particular, it provides us ways to:

1. exploit the monomial structure of a polynomial system to place upper bounds on the number of complex solutions **beforehand**

**addresses  
challenge 2**

2. use these upper bounds to algorithmically find **all** complex solutions

**addresses  
challenge 1**



**Homotopy continuation**

Tighter upper bounds  
lead to faster solutions

**Question:** How to find a good upper bound for our system

# Algebraic Geometry

## Main Ideas

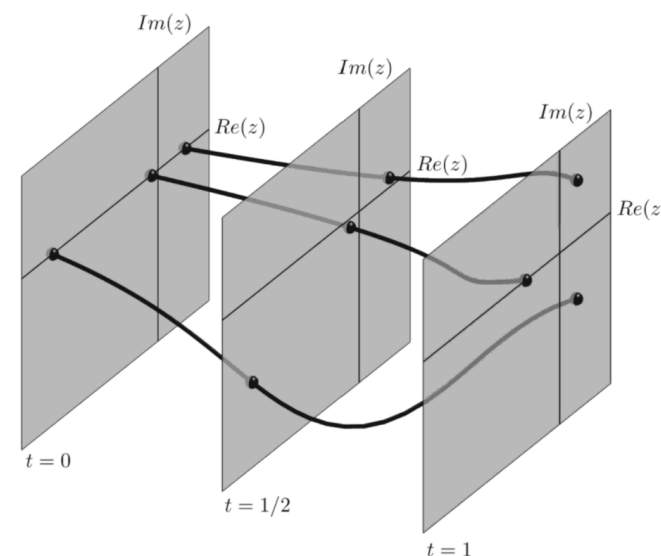
### Homotopy Continuation

- Suppose we want to solve the system  $F(x) = [f_1(x), \dots, f_N(x)] = 0$  **Target system**
- We generate another polynomial system  $G(x) = [g_1(x), \dots, g_N(x)] = 0$  **Start system**
- $G(x) = 0$  is guaranteed to have at least as many isolated solutions as  $F(x) = 0$  and these solutions are known beforehand
- We define a parameterized family of systems:

$$H(x, t) = tG(x) + (1 - t)F(x), \quad t \in [0, 1]$$

A tighter upper bound means fewer paths to track and therefore provides a more efficient way to solve the target system

### Homotopy



# Algebraic Geometry

## Main Ideas

### Well known upper bounds

**Theorem 2.13 (Classical Bezout Bound).** *Let  $f_1, \dots, f_N$  be polynomials in  $\mathbb{C}[x_1, \dots, x_N]$ .*

*Then the number of isolated solutions of the system  $f_1(x) = \dots = f_N(x) = 0$  is bounded above by the product  $\deg(f_1) \cdots \deg(f_N)$ .*

**Theorem 2.22 (Bernstein's Theorem).** *Let  $f_1, \dots, f_N \in \mathbb{C}[x_1, \dots, x_N]$  be Laurent polynomials with Newton polytopes  $Q_1, \dots, Q_N$ . The number of isolated solutions of the system*

*$f_1(x) = \dots = f_N(x) = 0$  in  $(\mathbb{C}^*)^N$  is bounded above by the mixed volume  $\mathcal{M}(Q_1, \dots, Q_N)$ .*

<b>Question:</b> Can we do better?    Yes!
--

# Results

## Number of complex critical points

Critical points with all non-zero weights

**Proposition 3.6** (upper bound on solutions in  $(\mathbb{C}^*)^N, \mathcal{B}_{\mathbb{C}^*}$ ). *Consider a linear network with  $H = 1, m = 1, d_x = n, d_y = p$  and  $d_1 = d$ . Let  $(W_1, W_2) = (\begin{pmatrix} a_{1,1} & \cdots & a_{1,n} \\ \vdots & & \vdots \\ a_{d,1} & \cdots & a_{d,n} \end{pmatrix}, \begin{pmatrix} b_{1,1} & \cdots & b_{1,d} \\ \vdots & & \vdots \\ b_{p,1} & \cdots & b_{p,d} \end{pmatrix})$  denote a solution to the gradient polynomial system (3.1). Then, there are at most*

$$\mathcal{B}_{\mathbb{C}^*} = (4p)^d$$

*solutions for which  $a_{1,1}, \dots, a_{d,n} \in \mathbb{C}^*$  and  $b_{1,1}, \dots, b_{p,d} \in \mathbb{C}^*$ .*

# Results

## Number of complex critical points

All critical points

**Theorem 3.9** (upper bound on complex critical points,  $\mathcal{B}_{\mathbb{C}}$ ). *Consider a linear network with*

*$H = 1$ ,  $m = 1$ ,  $d_x = n$ ,  $d_y = p$  and  $d_1 = d$ . This network has at most*

$$\mathcal{B}_{\mathbb{C}} = (1 + 4p)^d$$

*complex critical points.*

No	$d$	$d_x$	$d_y$	$N$	CBB	BKK	$\mathcal{B}_{\mathbb{C}}$	$\mathcal{B}_{\mathbb{C}^*}$	$N_{\mathbb{C}}$	$N_{\mathbb{C}^*}$	$\max\{N_{\mathbb{R}}\}$
1	1	1	1	2	9	5	5	4	5	4	3
2	1	2	1	3	27	9	5	4	5	4	3
3	1	3	1	4	81	13	5	4	5	4	3
4	1	1	2	3	27	9	9	8	9	8	3
5	1	2	2	4	81	33	9	8	9	8	3
6	1	3	2	5	243	73	9	8	9	8	3
7	1	1	3	4	81	13	13	12	13	12	3
8	1	2	3	5	243	73	13	12	13	12	3
9	1	3	3	6	729	245	13	12	13	12	3
10	2	1	1	4	81	25	25	16	9	0	4
11	2	2	1	6	729	81	25	16	9	0	5
12	2	3	1	8	6561	169	25	16	9	0	5
13	2	1	2	6	729	81	81	64	33	16	9
14	2	2	2	8	6561	1089	81	64	33	16	9
15	2	3	2	10	59049	5329	81	64	33	16	9
16	2	1	3	8	6561	169	169	144	73	48	9
17	2	2	3	10	59049	5329	169	144	73	48	9
18	2	3	3	12	531441	60025	169	144	73	48	9
19	3	1	1	6	729	125	125	64	13	0	7
20	3	2	1	9	19683	729	125	64	13	0	7
21	3	3	1	12	531441	2197	125	64	13	0	7
22	3	1	2	9	19683	729	729	512	73	0	19
23	3	2	2	12	531441	35937	729	512	73	0	19
24	3	3	2	15	14348907	389017	729	512	73	0	19
25	3	1	3	12	531441	2197	2197	1728	245	64	27
26	3	2	3	15	14348907	389017	2197	1728	245	64	27

Table 3.1: Case:  $H = 1, m = 1$ . Comparison of upper bounds on the number of complex critical points of a linear network.  $d$  = number of neurons in each layer,  $d_x$  = input dimension and  $d_y$  = output dimension.  $N$  = total number of weights in the network. CBB and BKK refer to the classical Bezout bound and the BKK bound respectively.  $\mathcal{B}_{\mathbb{C}}$  and  $\mathcal{B}_{\mathbb{C}^*}$  refer to the new bounds on the number of critical points in  $(\mathbb{C})^N$  and  $(\mathbb{C}^*)^N$  respectively.  $N_{\mathbb{C}}$  and  $N_{\mathbb{C}^*}$  refer to the actual number of critical points in  $(\mathbb{C})^N$  and  $(\mathbb{C}^*)^N$  respectively.  $\max\{N_{\mathbb{R}}\}$  = maximum number of real solutions observed<sup>14</sup> within each sample.

# Results

## Location of complex critical points

Critical points with some zero weights lie on particular coordinate subspaces

**Proposition 3.1** (no stray zeros in  $W_1$ ). *Consider a linear network with  $H = 1, m = 1, d_x = n, d_y = p$  and  $d_1 = d$ . Let  $(W_1, W_2) = \left( \begin{pmatrix} a_{1,1} & \cdots & a_{1,n} \\ \vdots & & \vdots \\ a_{d,1} & \cdots & a_{d,n} \end{pmatrix}, \begin{pmatrix} b_{1,1} & \cdots & b_{1,d} \\ \vdots & & \vdots \\ b_{p,1} & \cdots & b_{p,d} \end{pmatrix} \right)$  denote a solution to the regularized gradient polynomial system (3.1). If  $a_{i,j} = 0$ , then  $a_{i,s} = 0$  for all  $s = 1, \dots, n$ .*

**Proposition 3.3** (no stray zeros in  $W_2$ ). *Consider a linear network with  $H = 1, m = 1, d_x = n, d_y = p$  and  $d_1 = d$ . Let  $(W_1, W_2) = \left( \begin{pmatrix} a_{1,1} & \cdots & a_{1,n} \\ \vdots & & \vdots \\ a_{d,1} & \cdots & a_{d,n} \end{pmatrix}, \begin{pmatrix} b_{1,1} & \cdots & b_{1,d} \\ \vdots & & \vdots \\ b_{p,1} & \cdots & b_{p,d} \end{pmatrix} \right)$  denote a solution to the regularized gradient polynomial system (3.1). Then,*

$$b_{k,i} = 0 \implies b_{\cdot,i} = 0$$



# Results

## Location of complex critical points

Critical points with some zero weights lie on particular coordinate subspaces

**Proposition 3.2** (null rows of  $W_1$  match null columns of  $W_2$ ). *Consider a linear network with*

*$H = 1, m = 1, d_x = n, d_y = p$  and  $d_1 = d$ . Let  $(W_1, W_2) = \left( \begin{pmatrix} a_{1,1} & \cdots & a_{1,n} \\ \vdots & & \vdots \\ a_{d,1} & \cdots & a_{d,n} \end{pmatrix}, \begin{pmatrix} b_{1,1} & \cdots & b_{1,d} \\ \vdots & & \vdots \\ b_{p,1} & \cdots & b_{p,d} \end{pmatrix} \right)$*

*denote a solution to the regularized gradient polynomial system (3.1). Then,*

$$a_{i,\cdot} = 0 \iff b_{\cdot,i} = 0$$



# Results

## Location of complex critical points

Critical points with some zero weights lie on particular coordinate subspaces

W1		W2		
0	0	0	*	8 solutions
*	*	0	*	
*	*	*	0	8 solutions
0	0	*	0	
*	*	*	*	16 solutions
*	*	*	*	
0	0	0	0	1 solution
0	0	0	0	

Figure 3.1: Case:  $H = 1$ ,  $m = 1$ ,  $d = 2$ ,  $d_x = 2$ ,  $d_y = 2$  (corresponds to line no. 14 in Table 3.1. Number of solutions corresponding to each zero-pattern occurring in the weight matrices  $W_1$  and  $W_2$ . \* represents a non-zero entry.

# Recap

- Neural networks (particularly, *linear networks*)
- Neural network training as a minimization problem in real analysis
- Neural network training as solving polynomial systems in algebraic geometry
- Results:
  - New upper bounds  $\mathcal{B}_{\mathbb{C}}, \mathcal{B}_{\mathbb{C}^*}$  on the number of complex critical points of 1-hidden layer networks
  - Structure in the location of complex critical points with some zero weights

# Further research

- Are  $\mathcal{B}_{\mathbb{C}^*}$  and  $\mathcal{B}_{\mathbb{C}}$  ever attained? Maybe for large  $m$ ?
- Conversely, can we show that  $\mathcal{B}_{\mathbb{C}^*}$  and  $\mathcal{B}_{\mathbb{C}}$  are never attained?
- Prove zero patterns hold for  $H > 1$ .