# Critical Points of Deep Linear Networks in $\mathbb{C}^N$

### Thesis Defense Presentation

by

Ayush Bharadwaj

Department of Mathematics

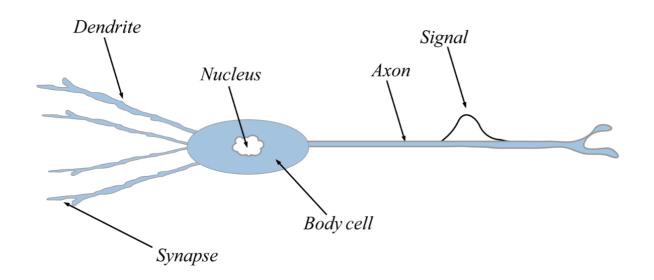
San Francisco State University

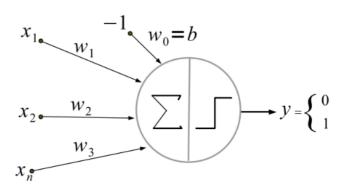
# Critical Points of Deep Linear Networks in $\mathbb{C}^N$

### Outline

- Neural networks (particularly, linear networks)
- Neural network training as a minimization problem in real analysis
- Neural network training as solving polynomial systems in algebraic geometry
- Results

### Neuron



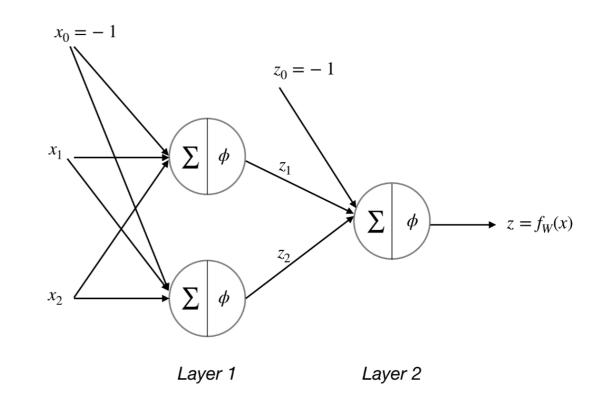


• 
$$f_w(x_1,...,x_n) = \phi\left(\sum_{j=1}^n w_j x_j - b\right)$$
 (output function)

### Neural network

• The weight matrix of the *i*th layer

$$W_i = egin{pmatrix} w_{i11} & \cdots & w_{i1n} \ dots & & dots \ w_{ir1} & \cdots & w_{irn} \end{pmatrix}$$



**Linear Network** 

We drop the bias term b and choose  $\phi = \mathbf{I}$  Neuron output  $f_w(x_1, ..., x_n) = \sum_{j=1}^n w_j x_j$ 

This simplifies the network output function:

$$f_W(x) = W_{H+1}W_H \cdots W_2W_1x \tag{2}$$

We want to use  $f_W$  to approximate some function  $y: \mathbb{R}^n \longrightarrow \mathbb{R}^p$ 

# Training a linear network

# Minimization problem in real analysis

• Error (or loss),  $\mathcal{L}: \mathbb{R}^n \longrightarrow \mathbb{R}$  defined by:

$$\mathcal{L}(W) = \frac{1}{2} \sum_{i=1}^{m} ||W_{H+1}W_{H}...W_{2}W_{1}x^{(i)} - y^{(i)}||^{2}$$
(3)

Not convex

• We want to solve the minimization problem:

$$\min_{W \in \mathbb{R}^N} \ \mathcal{L}(W) \tag{4}$$

# Training a linear network

Non-convexity of  $\mathcal{L}(W)$  presents challenges

- 1. Hard to guarantee that all local minima have been found by the algorithm.
- 2. Hard to make assertions about the number and location of minima before hand.

Motivation for an algebraic geometry view

# Training a linear network

From a problem in Real Analysis to a problem in Algebraic Geometry

We introduce the following relaxations:

• 
$$\min_{W \in \mathbb{R}^N} \mathcal{L}(W) \longrightarrow \nabla \mathcal{L}(W) = 0$$

$$\bullet \ \ W \in \mathbb{R}^N$$
 
$$\longrightarrow W \in \mathbb{C}^N$$

We solve:

$$\nabla \mathcal{L}(W) = 0, \quad W \in \mathbb{C}^N$$
 (5)

Polynomial system

# Training a linear network

From a problem in Real Analysis to a problem in Algebraic Geometry

Example 1.2 (A 4-weight network). Let us consider a 2-layer network with  $W_1 = [\alpha_1, \alpha_2]$   $W_2 = \begin{bmatrix} \beta_1 \\ \beta_2 \end{bmatrix}$ 

• 
$$\mathcal{L}(W) = \mathcal{L}(\alpha_1, \alpha_2, \beta_1, \beta_2) = \frac{1}{2} \sum_{i=1}^{2} ||W_2 W_1 x^{(i)} - y^{(i)}||^2$$

• 
$$\nabla \mathcal{L}(W) = \nabla \mathcal{L}(\alpha_1, \alpha_2, \beta_1, \beta_2) = 0$$
 gives

#### Nice polynomial system

Solutions are precisely the critical points

$$5\alpha_1\beta_1^2 + 5\alpha_1\beta_2^2 + 11\alpha_2\beta_1^2 + 11\alpha_2\beta_2^2 - 7\beta_1 - 10\beta_2 = 0$$

$$11\alpha_1\beta_1^2 + 11\alpha_1\beta_2^2 + 25\alpha_2\beta_1^2 + 25\alpha_2\beta_2^2 - 15\beta_1 - 22\beta_2 = 0$$

$$5\alpha_1^2\beta_1 + 22\alpha_1\alpha_2\beta_1 + 25\alpha_2^2\beta_1 - 7\alpha_1 - 15\alpha_2 = 0$$

$$5\alpha_1^2\beta_2 + 22\alpha_1\alpha_2\beta_2 + 25\alpha_2^2\beta_2 - 10\alpha_1 - 22\alpha_2 = 0$$

- Square
- Sparse
- Deg 2H + 1

(1.13)

# Algebraic Geometry

### Main Ideas

Algebraic geometry provides results and methods to analyze and solve systems of polynomial equations. In particular, it provides us ways to:

1. exploit the monomial structure of a polynomial system to place upper bounds on the number of complex solutions **beforehand** 

adresses challenge 2

2. use these upper bounds to algorithmically find all complex solutions

adresses challenge 1

#### **Homotopy continuation**

Tighter upper bounds lead to faster solutions

Question: How to find a good upper bound for our system

# Algebraic Geometry

### Main Ideas

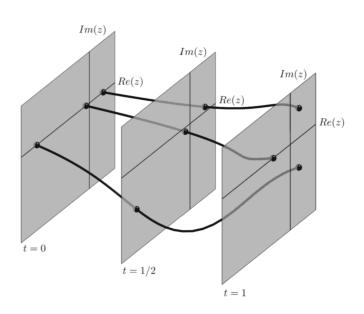
### **Homotopy Continuation**

- Suppose we want to solve the system  $F(x) = [f_1(x), ..., f_N(x)] = 0$  Target system
- We generate another polynomial system  $G(x) = [g_1(x), ..., g_N(x)] = 0$  Start system
- G(x) = 0 is guaranteed to have at least as many isolated solutions as F(x) = 0 and these solutions are known beforehand
- We define a parameterized family of systems:

$$H(x,t) = tG(x) + (1-t)F(x), t \in [0,1]$$

A tighter upper bound means fewer paths to track and therefore provides a more efficient way to solve the target system

#### **Homotopy**



# Algebraic Geometry

### Main Ideas

### Well known upper bounds

Theorem 2.13 (Classical Bezout Bound). Let  $f_1, ..., f_N$  be polynomials in  $\mathbb{C}[x_1, ..., x_N]$ .

Then the number of isolated solutions of the system  $f_1(x) = ... = f_N(x) = 0$  is bounded above by the product  $deg(f_1) \cdots deg(f_N)$ .

**Theorem 2.22** (Bernstein's Theorem). Let  $f_1, ..., f_N \in \mathbb{C}[x_1, ..., x_N]$  be Laurent polynomials with Newton polytopes  $Q_1, ..., Q_N$ . The number of isolated solutions of the system  $f_1(x) = ... = f_N(x) = 0$  in  $(\mathbb{C}^*)^N$  is bounded above by the mixed volume  $\mathcal{M}(Q_1, ..., Q_N)$ .

Question: Can we do better? Yes!

# Number of complex critical points

Critical points with all non-zero weights

**Proposition 3.6** (upper bound on solutions in  $(\mathbb{C}^*)^N$ ,  $\mathcal{B}_{\mathbb{C}^*}$ ). Consider a linear network with

$$H=1,\ m=1,\ d_x=n,\ d_y=p\ and\ d_1=d.$$
 Let  $(W_1,W_2)=(\begin{pmatrix} a_{1,1}&\cdots&a_{1,n}\\ \vdots\\ a_{d,1}&\cdots&a_{d,n} \end{pmatrix}, \begin{pmatrix} b_{1,1}&\cdots&b_{1,d}\\ \vdots\\ b_{p,1}&\cdots&b_{p,d} \end{pmatrix})$  denote a solution to the gradient polynomial system (3.1). Then, there are at most

$$\mathcal{B}_{\mathbb{C}^*} = (4p)^d$$

solutions for which  $a_{1,1},...,a_{d,n}\in\mathbb{C}^*$  and  $b_{1,1},...,b_{p,d}\in\mathbb{C}^*$ .

# Number of complex critical points All critical points

**Theorem 3.9** (upper bound on complex critical points,  $\mathcal{B}_{\mathbb{C}}$ ). Consider a linear network with

 $H=1,\ m=1,\ d_x=n,\ d_y=p\ \ and\ d_1=d.$  This network has at most

$$\mathcal{B}_{\mathbb{C}} = (1+4p)^d$$

 $complex\ critical\ points.$ 

No	d	$d_x$	$d_y$	N	CBB	BKK	$\mathcal{B}_{\mathbb{C}}$	$\mathcal{B}_{\mathbb{C}^*}$	$N_{\mathbb{C}}$	$N_{\mathbb{C}^*}$	$max\{N_{\mathbb{R}}\}$
1	1	1	1	2	9	5	5	4	5	4	3
2	1	2	1	3	27	9	5	4	5	4	3
3	1	3	1	4	81	13	5	4	5	4	3
4	1	1	2	3	27	9	9	8	9	8	3
5	1	2	2	4	81	33	9	8	9	8	3
6	1	3	2	5	243	73	9	8	9	8	3
7	1	1	3	4	81	13	13	12	13	12	3
8	1	2	3	5	243	73	13	12	13	12	3
9	1	3	3	6	729	245	13	12	13	12	3
10	2	1	1	4	81	25	25	16	9	0	4
11	2	2	1	6	729	81	25	16	9	0	5
12	2	3	1	8	6561	169	25	16	9	0	5
13	2	1	2	6	729	81	81	64	33	16	9
14	2	2	2	8	6561	1089	81	64	33	16	9
15	2	3	2	10	59049	5329	81	64	33	16	9
16	2	1	3	8	6561	169	169	144	73	48	9
17	2	2	3	10	59049	5329	169	144	73	48	9
18	2	3	3	12	531441	60025	169	144	73	48	9
19	3	1	1	6	729	125	125	64	13	0	7
20	3	2	1	9	19683	729	125	64	13	0	7
21	3	3	1	12	531441	2197	125	64	13	0	7
22	3	1	2	9	19683	729	729	512	73	0	19
23	3	2	2	12	531441	35937	729	512	73	0	19
24	3	3	2	15	14348907	389017	729	512	73	0	19
25	3	1	3	12	531441	2197	2197	1728	245	64	27
26	3	2	3	15	14348907	389017	2197	1728	245	64	27

Table 3.1: Case: H=1, m=1. Comparison of upper bounds on the number of complex critical points of a linear network. d= number of neurons in each layer,  $d_x=$  input dimension and  $d_y=$  output dimension. N= total number of weights in the network. CBB and BKK refer to the classical Bezout bound and the BKK bound respectively.  $\mathcal{B}_{\mathbb{C}}$  and  $\mathcal{B}_{\mathbb{C}^*}$  refer to the new bounds on the number of critical points in  $(\mathbb{C})^N$  and  $(\mathbb{C}^*)^N$  respectively.  $N_{\mathbb{C}}$  and  $N_{\mathbb{C}^*}$  refer to the actual number of critical points in  $(\mathbb{C})^N$  and  $(\mathbb{C}^*)^N$  respectively.  $\max\{N_{\mathbb{R}}\}$  = maximum number of real solutions observed thin each sample.

# Location of complex critical points

Critical points with some zero weights lie on particular coordinate subspaces

**Proposition 3.1** (no stray zeros in  $W_1$ ). Consider a linear network with  $H=1, m=1, d_x=1$ 

$$n, d_y = p \ and \ d_1 = d. \ Let (W_1, W_2) = \left( \begin{pmatrix} a_{1,1} & \cdots & a_{1,n} \\ \vdots & & & \\ a_{d,1} & \cdots & a_{d,n} \end{pmatrix}, \begin{pmatrix} b_{1,1} & \cdots & b_{1,d} \\ \vdots & & & \\ b_{p,1} & \cdots & b_{p,d} \end{pmatrix} \right) \ denote \ a \ solution \ to \ the$$

regularized gradient polynomial system (3.1). If  $a_{i,j} = 0$ , then  $a_{i,s} = 0$  for all s = 1, ..., n.

**Proposition 3.3** (no stray zeros in  $W_2$ ). Consider a linear network with  $H=1, m=1, d_x=1$ 

$$n, d_y = p \ and \ d_1 = d. \ Let (W_1, W_2) = (\begin{pmatrix} a_{1,1} & \cdots & a_{1,n} \\ \vdots & & & \\ a_{d,1} & \cdots & a_{d,n} \end{pmatrix}, \begin{pmatrix} b_{1,1} & \cdots & b_{1,d} \\ \vdots & & & \\ b_{p,1} & \cdots & b_{p,d} \end{pmatrix}) \ denote \ a \ solution \ to \ the$$

regularized gradient polynomial system (3.1). Then,

$$b_{k,i} = 0 \implies b_{\cdot,i} = 0$$

# Location of complex critical points

Critical points with some zero weights lie on particular coordinate subspaces

**Proposition 3.2** (null rows of  $W_1$  match null columns of  $W_2$ ). Consider a linear network with

$$H = 1, \ m = 1, \ d_x = n, \ d_y = p \ \ and \ d_1 = d. \ \ Let \ (W_1, W_2) = \left( \begin{pmatrix} a_{1,1} & \cdots & a_{1,n} \\ \vdots & \vdots & \vdots \\ a_{d,1} & \cdots & a_{d,n} \end{pmatrix}, \ \begin{pmatrix} b_{1,1} & \cdots & b_{1,d} \\ \vdots & \vdots & \vdots \\ b_{p,1} & \cdots & b_{p,d} \end{pmatrix} \right)$$

denote a solution to the regularized gradient polynomial system (3.1). Then,

$$a_{i,\cdot} = 0 \iff b_{\cdot,i} = 0$$

# Location of complex critical points

Critical points with some zero weights lie on particular coordinate subspaces

```
W2 W1
01 00
0111
subspace dim: 4
solution count: 8
W2 W1
11 11
11 11
subspace dim: 8
solution count: 16
W2 W1
10 11
10 00
subspace dim: 4
solution count: 8
W2 W1
00 00
00 00
subspace dim: 0
solution count: 1
```

# Recap

- Neural networks (particularly, linear networks)
- Neural network training as a minimization problem in real analysis
- Neural network training as solving polynomial systems in algebraic geometry
- Results:
  - New upper bounds  $\mathscr{B}_{\mathbb{C}}, \mathscr{B}_{\mathbb{C}^*}$  on the number of complex critical points of 1-hidden layer networks
  - Structure in the location of complex critical points with some zero weights

# Further research

- Are  $\mathcal{B}_{\mathbb{C}^*}$  and  $\mathcal{B}_{\mathbb{C}}$  ever attained? Maybe for large m?
- Conversely, can we show that  $\mathcal{B}_{\mathbb{C}^*}$  and  $\mathcal{B}_{\mathbb{C}}$  are never attained?
- Prove zero patterns hold for H > 1.