Thesis Defense Presentation

by

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Department of Mathematics

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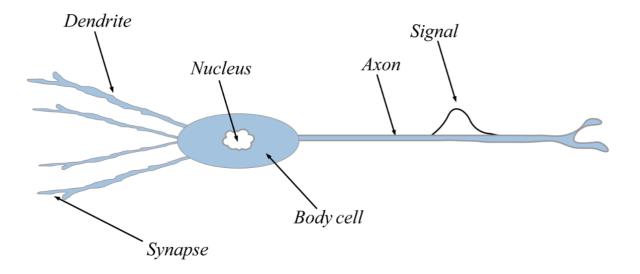
Outline

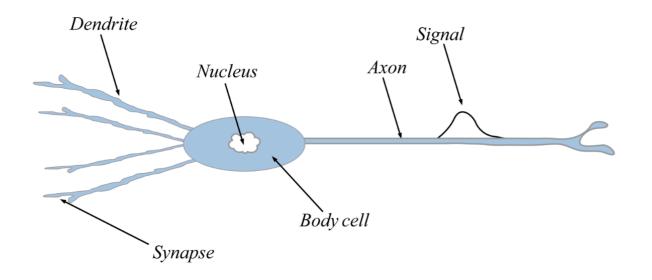
• Neural networks (particularly, *linear networks*)

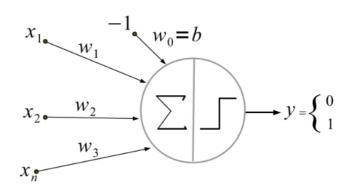
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- Neural network training as a minimization problem in real analysis

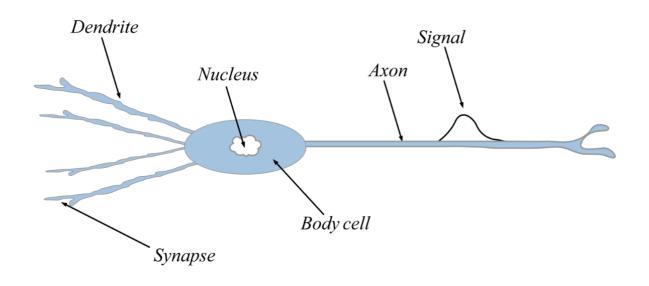
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- Neural network training as solving polynomial systems in algebraic geometry

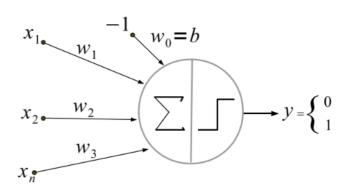
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- Neural network training as solving polynomial systems in algebraic geometry
- Results and contributions



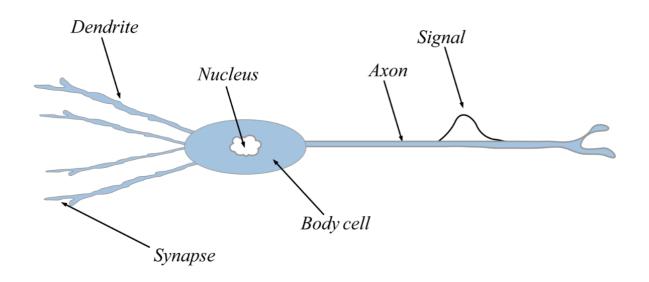


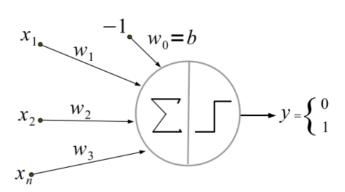






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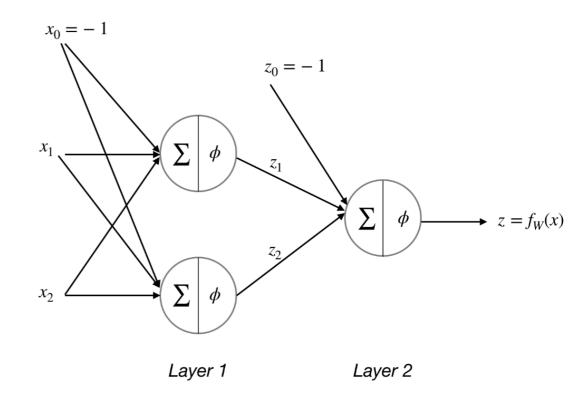




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 (output function)

Neural network

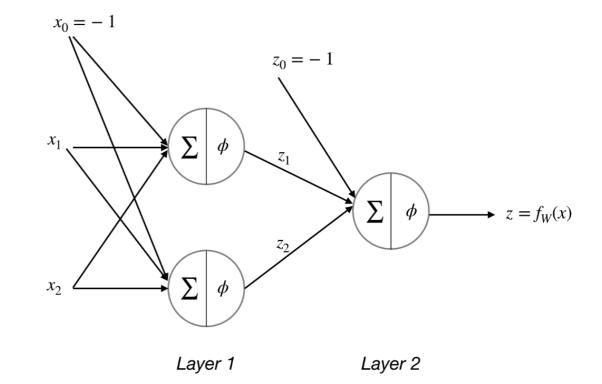
Neural network



Neural network

• The weight matrix of the *i*th layer

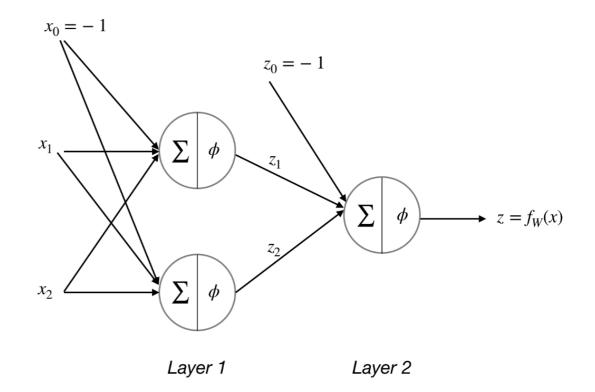
$$W_i = egin{pmatrix} w_{i11} & \cdots & w_{i1n} \ dots & & dots \ w_{ir1} & \cdots & w_{irn} \end{pmatrix}$$



Neural network

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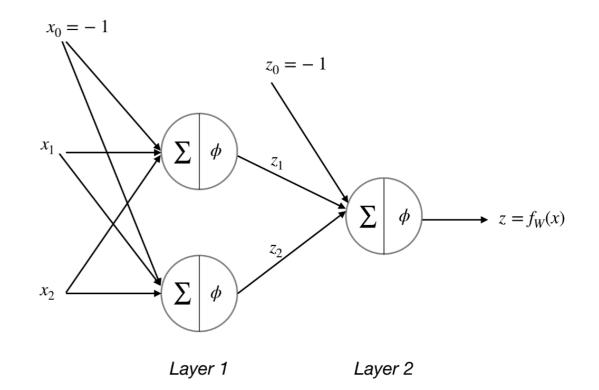
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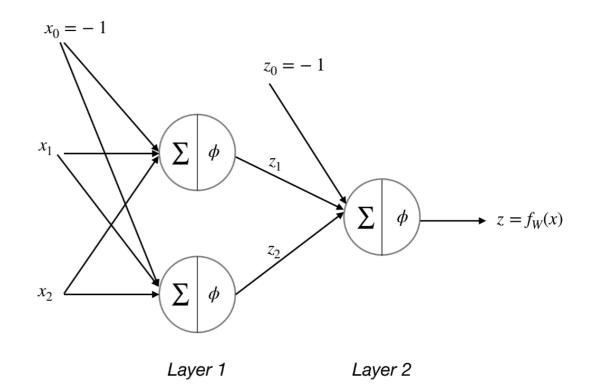
• Output of the *i*th layer

$$z^{(i)} = \phi \Big(W_i \begin{bmatrix} -1 \\ z^{(i-1)} \end{bmatrix} \Big)$$

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• Output of the network

$$f_W(x) = z^{(H+1)} = \phi\left(W_{H+1} \begin{bmatrix} -1\\ z^{(H)} \end{bmatrix}\right) \tag{2}$$

Linear network

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We want to use f_W to approximate some function $y: \mathbb{R}^n \longrightarrow \mathbb{R}^p$

Training a linear network

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- Error (or loss), $\mathcal{L}: \mathbb{R}^n \longrightarrow \mathbb{R}$ defined by:

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(3)

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Not convex

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Motivation for an algebraic geometry view

Training a linear network

From a problem in Real Analysis to a problem in Algebraic Geometry

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From a problem in Real Analysis to a problem in Algebraic Geometry

We introduce the following relaxations:

• $\min_{W \in \mathbb{R}^N} \mathcal{L}(W)$

Training a linear network

From a problem in Real Analysis to a problem in Algebraic Geometry

•
$$\min_{W \in \mathbb{R}^N} \mathcal{L}(W)$$
 \longrightarrow

Training a linear network

From a problem in Real Analysis to a problem in Algebraic Geometry

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$$\min_{W \in \mathbb{R}^N} \mathcal{L}(W) \longrightarrow \nabla \mathcal{L}(W) = 0$$

Training a linear network

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Training a linear network

From a problem in Real Analysis to a problem in Algebraic Geometry

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Polynomial system

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Training a linear network

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$$(1.13)$$

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Solutions are precisely the critical points

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- Square
- Sparse
- Deg 2H + 1

(1.13)

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Homotopy Continuation

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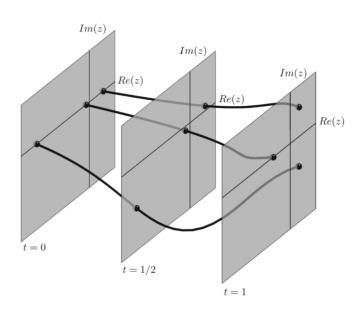
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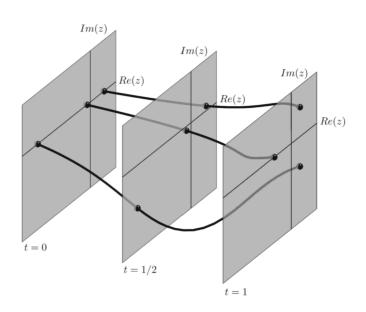
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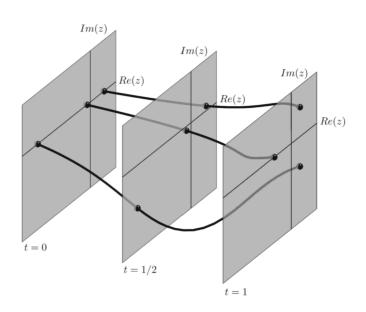
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Proposition 3.6 (upper bound on solutions in $(\mathbb{C}^*)^N$, $\mathcal{B}_{\mathbb{C}^*}$). Consider a linear network with

$$H=1, \ m=1, \ d_x=n, \ d_y=p \ \ and \ d_1=d. \ \ Let \ (W_1,W_2)=(\begin{pmatrix} a_{1,1} & \cdots & a_{1,n} \\ \vdots & \vdots & \\ a_{d,1} & \cdots & a_{d,n} \end{pmatrix}, \ \begin{pmatrix} b_{1,1} & \cdots & b_{1,d} \\ \vdots & \vdots & \\ b_{p,1} & \cdots & b_{p,d} \end{pmatrix})$$

denote a solution to the gradient polynomial system (3.1). Then, there are at most

$$\mathcal{B}_{\mathbb{C}^*} = (4p)^d$$

solutions for which $a_{1,1},...,a_{d,n} \in \mathbb{C}^*$ and $b_{1,1},...,b_{p,d} \in \mathbb{C}^*$.

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Theorem 3.9 (upper bound on complex critical points, $\mathcal{B}_{\mathbb{C}}$). Consider a linear network with

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No	d	d_x	d_y	N	CBB	BKK	$\mathcal{B}_{\mathbb{C}}$	$\mathcal{B}_{\mathbb{C}^*}$	$N_{\mathbb{C}}$	$N_{\mathbb{C}^*}$	$max\{N_{\mathbb{R}}\}$
1	1	1	1	2	9	5	5	4	5	4	3
2	1	2	1	3	27	9	5	4	5	4	3
3	1	3	1	4	81	13	5	4	5	4	3
4	1	1	2	3	27	9	9	8	9	8	3
5	1	2	2	4	81	33	9	8	9	8	3
6	1	3	2	5	243	73	9	8	9	8	3
7	1	1	3	4	81	13	13	12	13	12	3
8	1	2	3	5	243	73	13	12	13	12	3
9	1	3	3	6	729	245	13	12	13	12	3
10	2	1	1	4	81	25	25	16	9	0	4
11	2	2	1	6	729	81	25	16	9	0	5
12	2	3	1	8	6561	169	25	16	9	0	5
13	2	1	2	6	729	81	81	64	33	16	9
14	2	2	2	8	6561	1089	81	64	33	16	9
15	2	3	2	10	59049	5329	81	64	33	16	9
16	2	1	3	8	6561	169	169	144	73	48	9
17	2	2	3	10	59049	5329	169	144	73	48	9
18	2	3	3	12	531441	60025	169	144	73	48	9
19	3	1	1	6	729	125	125	64	13	0	7
20	3	2	1	9	19683	729	125	64	13	0	7
21	3	3	1	12	531441	2197	125	64	13	0	7
22	3	1	2	9	19683	729	729	512	73	0	19
23	3	2	2	12	531441	35937	729	512	73	0	19
24	3	3	2	15	14348907	389017	729	512	73	0	19
25	3	1	3	12	531441	2197	2197	1728	245	64	27
26	3	2	3	15	14348907	389017	2197	1728	245	64	27

Table 3.1: Case: H=1, m=1. Comparison of upper bounds on the number of complex critical points of a linear network. d= number of neurons in each layer, $d_x=$ input dimension and $d_y=$ output dimension. N= total number of weights in the network. CBB and BKK refer to the classical Bezout bound and the BKK bound respectively. $\mathcal{B}_{\mathbb{C}}$ and $\mathcal{B}_{\mathbb{C}^*}$ refer to the new bounds on the number of critical points in $(\mathbb{C})^N$ and $(\mathbb{C}^*)^N$ respectively. $N_{\mathbb{C}}$ and $N_{\mathbb{C}^*}$ refer to the actual number of critical points in $(\mathbb{C})^N$ and $(\mathbb{C}^*)^N$ respectively. $\max\{N_{\mathbb{R}}\}$ = maximum number of real solutions observed thin each sample.

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Critical points with some zero weights lie on particular coordinate subspaces

Proposition 3.1 (no stray zeros in W_1). Consider a linear network with $H = 1, m = 1, d_x = n, d_y = p$ and $d_1 = d$. Let $(W_1, W_2) = \begin{pmatrix} a_{1,1} & \cdots & a_{1,n} \\ \vdots & \vdots & \vdots \\ a_{d,1} & \cdots & a_{d,n} \end{pmatrix}, \begin{pmatrix} b_{1,1} & \cdots & b_{1,d} \\ \vdots & \vdots & \vdots \\ b_{p,1} & \cdots & b_{p,d} \end{pmatrix}$) denote a solution to the regularized gradient polynomial system (3.1). If $a_{i,j} = 0$, then $a_{i,s} = 0$ for all s = 1, ..., n.

Proposition 3.3 (no stray zeros in W_2). Consider a linear network with $H = 1, m = 1, d_x = n, d_y = p$ and $d_1 = d$. Let $(W_1, W_2) = \begin{pmatrix} a_{1,1} & \cdots & a_{1,n} \\ \vdots & \vdots & \vdots \\ a_{d,1} & \cdots & a_{d,n} \end{pmatrix}, \begin{pmatrix} b_{1,1} & \cdots & b_{1,d} \\ \vdots & \vdots & \vdots \\ b_{p,1} & \cdots & b_{p,d} \end{pmatrix}$) denote a solution to the regularized gradient polynomial system (3.1). Then,

$$b_{k,i} = 0 \implies b_{\cdot,i} = 0$$

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Proposition 3.2 (null rows of W_1 match null columns of W_2). Consider a linear network with

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```
W2 W1
01 00
0111
subspace dim: 4
solution count: 8
W2 W1
11 11
11 11
subspace dim: 8
solution count: 16
W2 W1
10 11
10 00
subspace dim: 4
solution count: 8
W2 W1
00 00
00 00
subspace dim: 0
solution count: 1
```

Further research

- 3. run experiments to check if increasing m causes BBK bound to be attained? for what value of m for a given architecture?
- 4. Conversely, can we show that the BKK bound is never reached? Even though our systems are non-generic, we might be able to solve the corresponding facial system may have no solutions, in which case, the BKK bound is a strict bound.
- 5. prove zero patterns for H > 1
- 6. extend table for H = 1, m=1 but higher values if d_x, d_y, d_i by solving reduced systems.



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