

Data Structures & Algorithms (PCC-CS 301)

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Topics Covered

- 1. Binary tree formation from traversal sequences
- 2. BST searching
- 3. Data deletion from BST



Binary Tree

- Binary tree formation from the traversal
 - ☐ A binary tree can be formed if two traversal sequences are known
 - Pre-order and in-order can form the tree
 - Post-order and in-order can form the tree
 - But pre-order and post-order cannot form the tree

If there is only pre-order or post-order traversal is provided, you can form a BST as in-order is known as the sorted sequence of the BST



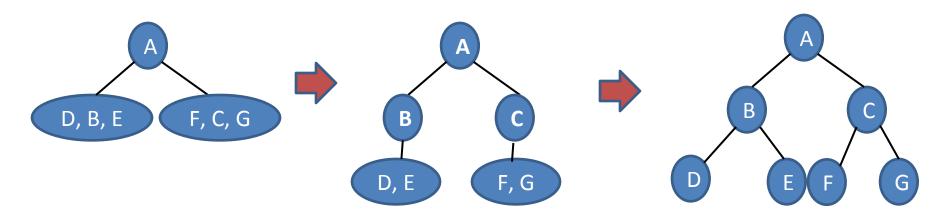
Binary Tree

Binary tree formation from the traversal

Pre-order: A -> B -> D -> E -> C -> F -> G

In-order: D -> B -> E -> A -> F -> C -> G

- Left most node in the pre-order will be parent
- Follow the in-order traversal
- The set of all nodes that reside in the left of root, are left sub-tree
- The set of all nodes that reside in the right of root, are right sub-tree.





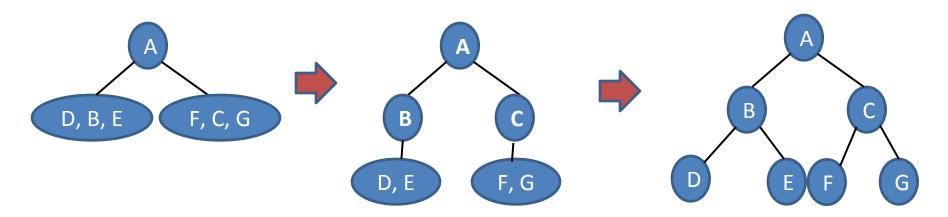
Binary Tree

Binary tree formation from the traversal

Post-order: D -> E -> B -> F -> G -> C -> A

In-order: D -> B -> E -> A -> F -> C -> G

- Right most node in the post-order will be parent
- Follow the in-order traversal
- The set of all nodes that reside in the left of root, are left sub-tree
- The set of all nodes that reside in the right of root, are right sub-tree.





- BST searching
 - ☐ Searching of a user given data into a BST
 - Start the searching from the root node
 - If the searching element is less than the current node, search in left sub-tree
 - If it is larger than the current searching node, search the right sub-tree next

Search for 8 =>

8 < 10 :: search the left sub-tree, next searching node is 5

8 > 5 :: search the right sub-tree, next searching node is 8

8 = 8 :: data found



BST searching (algorithm)

```
BST search(root, key) // key is the searching element
set ptr := root
while ptr != NULL
  if ptr -> data = key
     Print "search successful" and return
  else
     if ptr -> data < key
        ptr := ptr -> right
     else
        ptr := ptr -> left
 Print "Search unsuccessful"
 return
```

Time Complexity:

As we are searching a binary tree and data searching process will maximally be executed as the same number of iteration as the height of the tree.

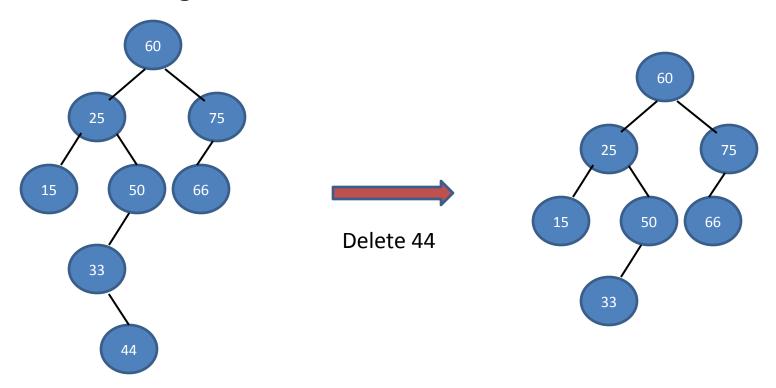
Hence, the time complexity of BST search is O(log₂n).



- Data deletion from BST
 - Data deletion from a BST depends on the position of deleting element
 - Deleting data has no children, it is a leaf node
 - ✓ Simply remove the node
 - Deleting data has only one child
 - ✔ Replace the deleting node with its only child node
 - Deleting data has two children
 - ✔ Replace the deleting node with its in-order successor
 - Delete the in-order successor by following case1 or case2 (in-order successor will not have any left child)
 - Replace it with actual deleting node

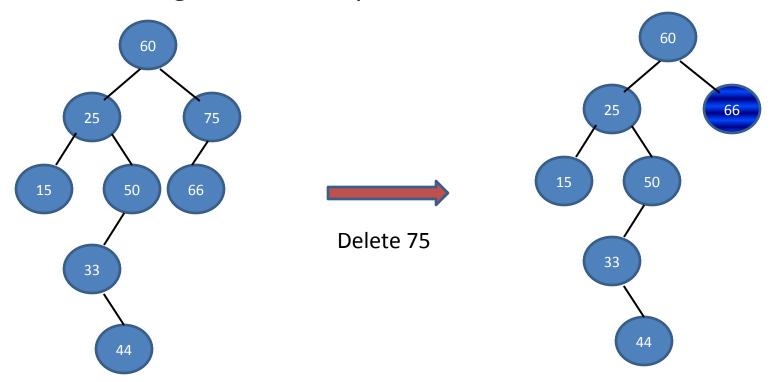


- Data deletion from BST
 - ☐ Deleting data has no children



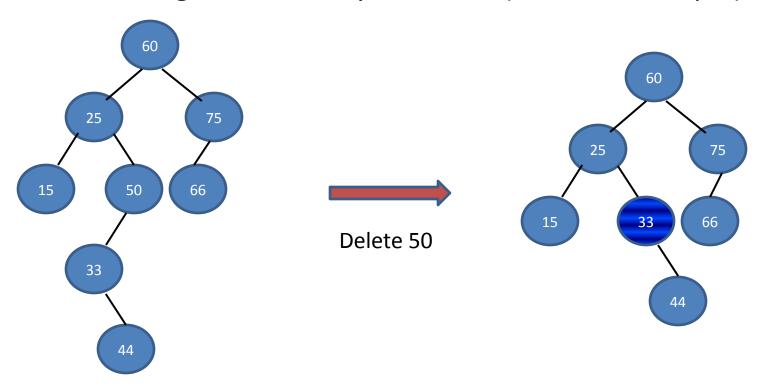


- Data deletion from BST
 - Deleting data has only one child



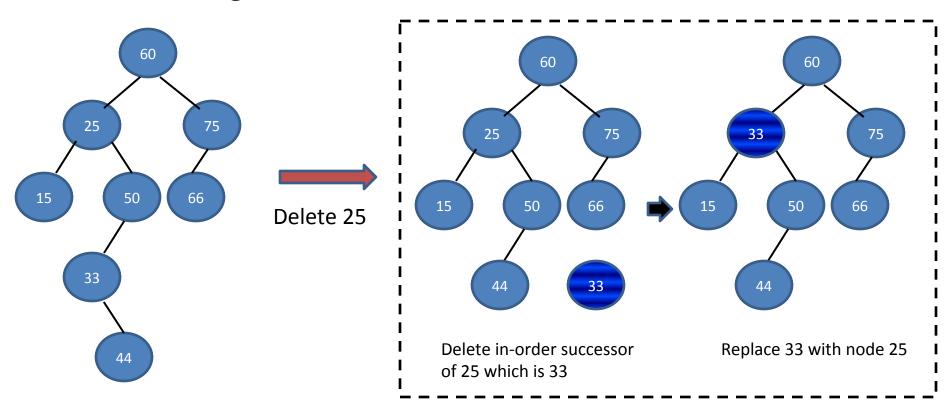


- Data deletion from BST
 - ☐ Deleting data has only one child (another example)





- Data deletion from BST
 - Deleting data has two children





- Data deletion from BST (complexity)
 - ☐ Time complexity
 - Deletion of any data from BST executes in two steps
 - searching the deleting data position which requires O(log₂n) time
 - ✓ Removing of selected data requires simple changing of few links, requires O(1) time
 - Total time complexity O(log₂n)



Queries?



Problem

Q1. Form the BST from the provided traversal

Post-order: 2, 8, 5, 3, 14, 12, 18, 25, 30, 20, 15, 10

- Q2. Delete following data from the BST shown below
 - a) delete 10
 - b) delete 60
 - c) delete 40

