





21)  $\{w \mid w \in (a+b)^*, |w| = \text{even}\}$  


22)  $\{w \mid |w| = \text{odd}\}$  


23)  $\{w \mid |w| \text{ is divisible by } 3\}$  


24)  $\{w \mid |w| = 2 \pmod 3\}$  


H.W.

25)  $\{w \mid \#a(w) = 2\}$  

26)  $\{w \mid \#a(w) \leq 2\}$  

27)  $\{w \mid \#a(w) \geq 2\}$  

28)  $\{w \mid \#a(w) = \text{even}\}$  

29)  $\{w \mid \#a(w) = \text{odd}\}$  

$$30) \{w \mid w \in (a+b)^*, \text{ 2}^{\text{nd}} \text{ symbol of } w \text{ is 'a'}\} = (a+b)a(a+b)^* \\ = (aa+ba)(a+b)^*$$

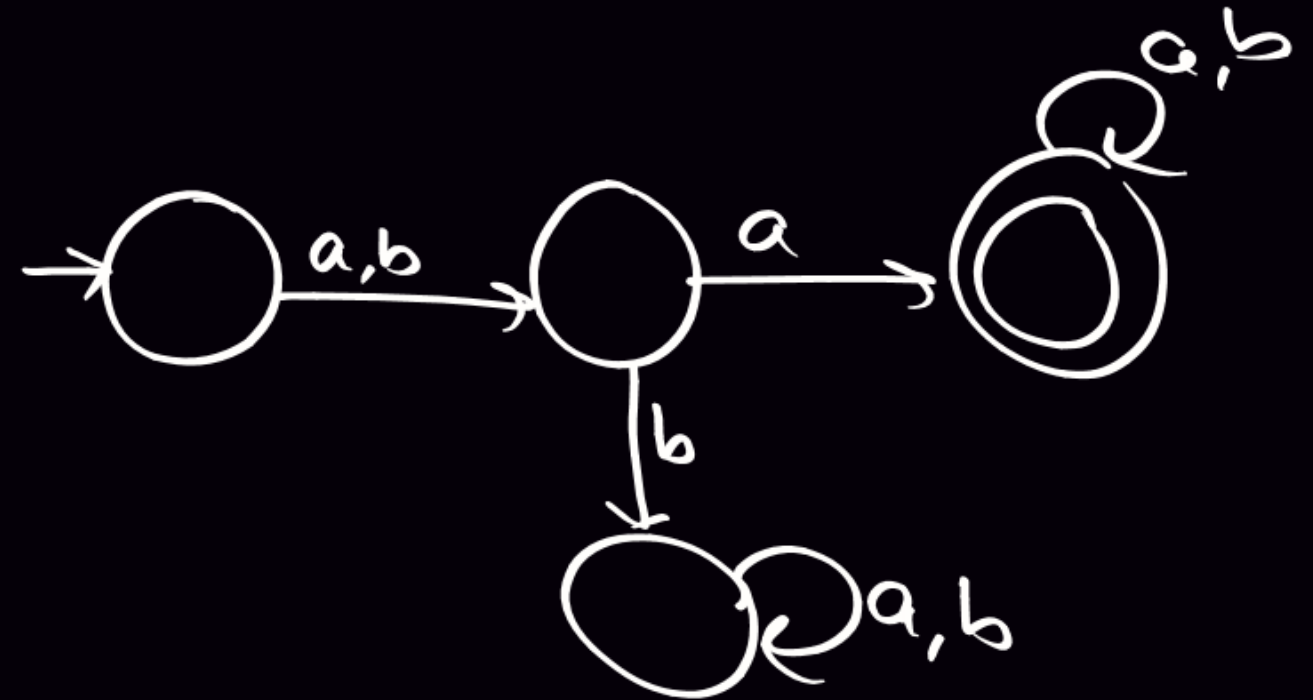
$K^{\text{th}}$  Symbol is 'a'

$\Downarrow$

$$= (K+1)+1$$

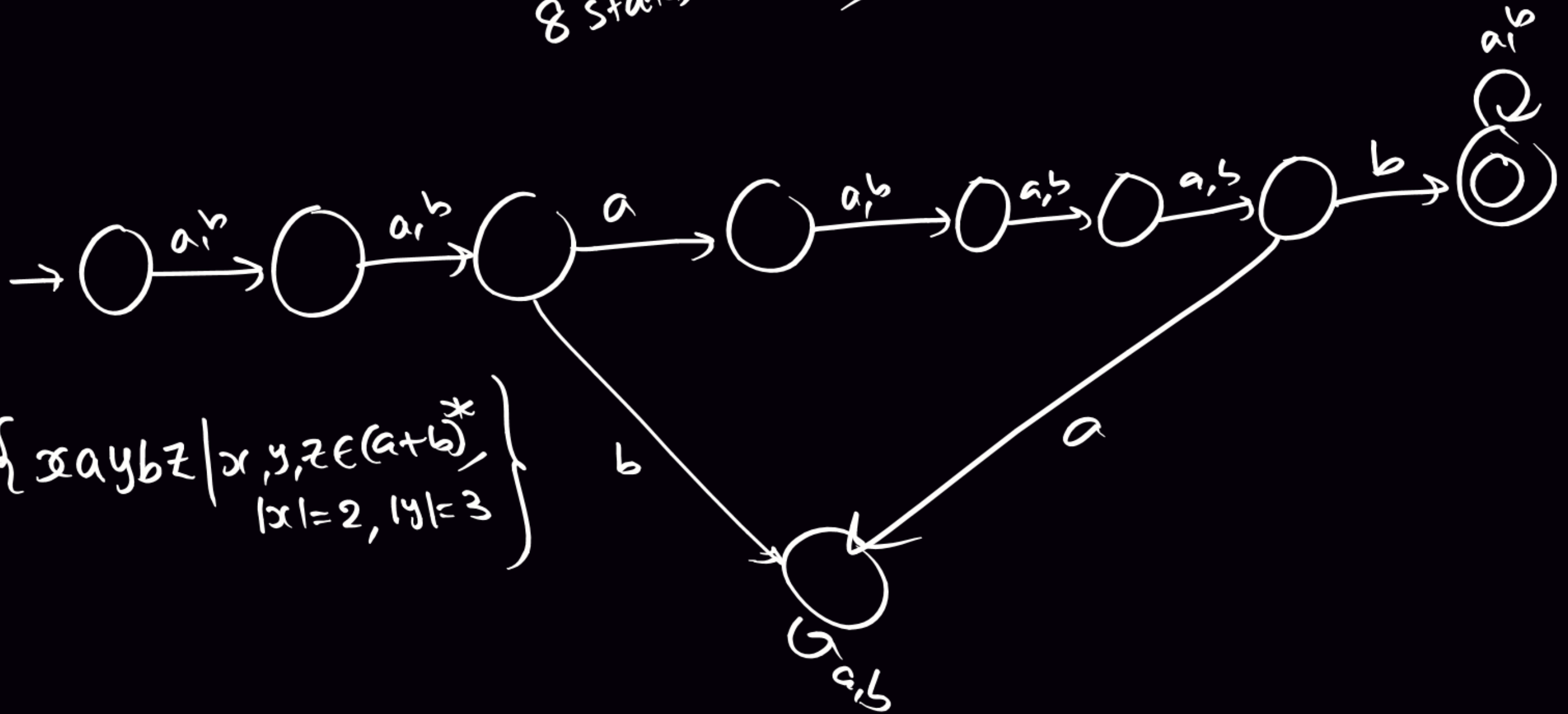
$= K+2$  states  
in min DFA

Min = aa or ba



31)  $L = (a+b)^2 a (a+b)^3 \underline{b} (a+b)^*$

$\downarrow$   
 $8 \text{ states} + 1 = \underline{\underline{9 \text{ states}}}$



$$L = \left\{ xaybz \mid x,y,z \in (a+b)^*, \right. \\ \left. |x|=2, |y|=3 \right\}$$

32)  $L = (a+b)^* a (a+b)^2$

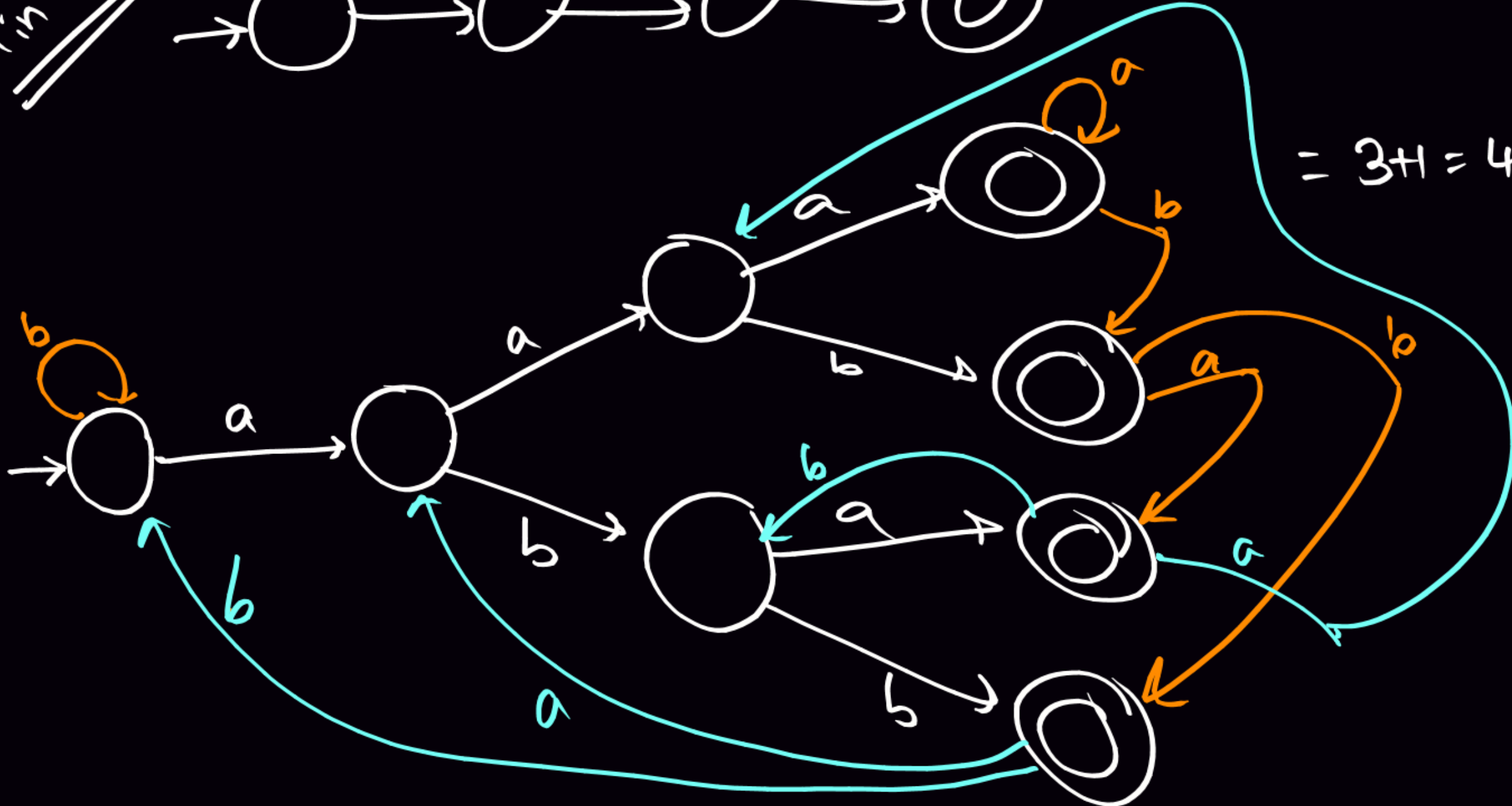
$3^{\text{rd}}$  last is 'a'

Min NFA:



$= 2^3 = 8$  states  
in min DFA

$= 3+1 = 4$  states  
in min NFA



33)  $L = \{a^{3n+2} \mid n \geq 0\} = \{w \mid w \in a^*, |w| \equiv 2 \pmod{3}\} \Rightarrow 3 \text{ states}$

34)  $L = \{a^{3n+7} \mid n \geq 0\} = \{a^7, a^{10}, a^{13}, a^{16}, \dots\} \Rightarrow 8 \text{ states}$

35)  $L = \{a^{100n+23} \mid n \geq 0\} \Rightarrow 100 \text{ states}$

36)  $L = \{a^{100n + \underline{1234}} \mid n \geq 0\} \Rightarrow 1235 \text{ states}$

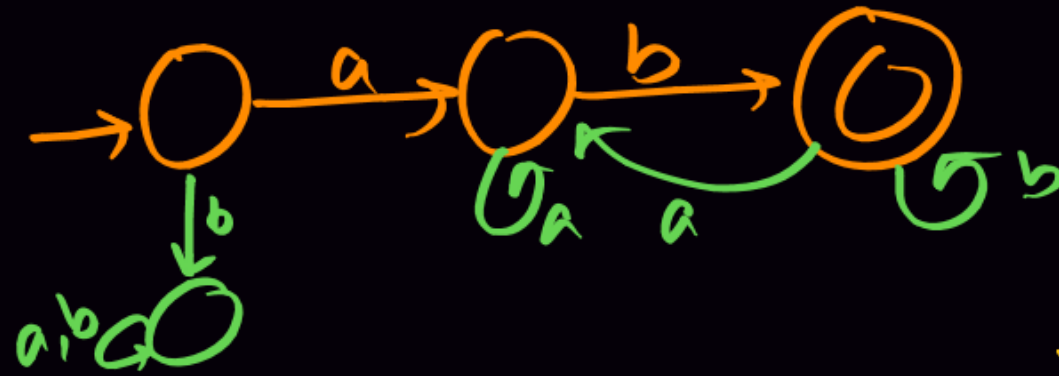
$$\{a^{K_1 n + K_2} \mid n \geq 0\}$$

If  $K_1 > K_2 \Rightarrow K_1 \text{ states}$   
 otherwise  $K_2 + 1 \text{ states}$



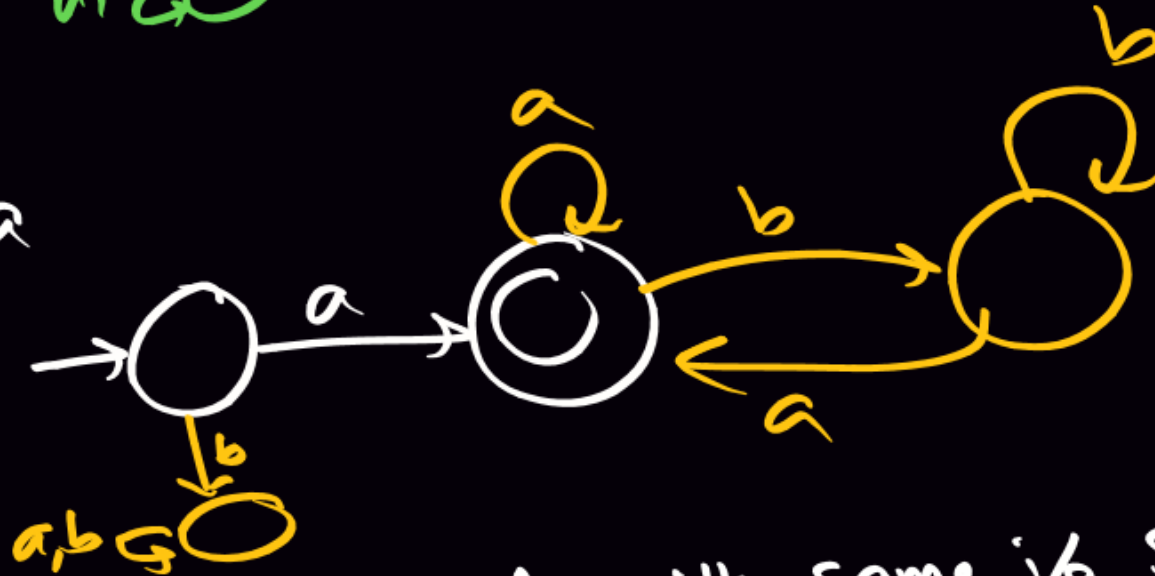
37)  $L = a(a+b)^*b$

→ 4 states in min DFA  
3 " " min NFA



38)  $L = a + a(a+b)^*a$

Min = a  
→ 4 states in min DFA

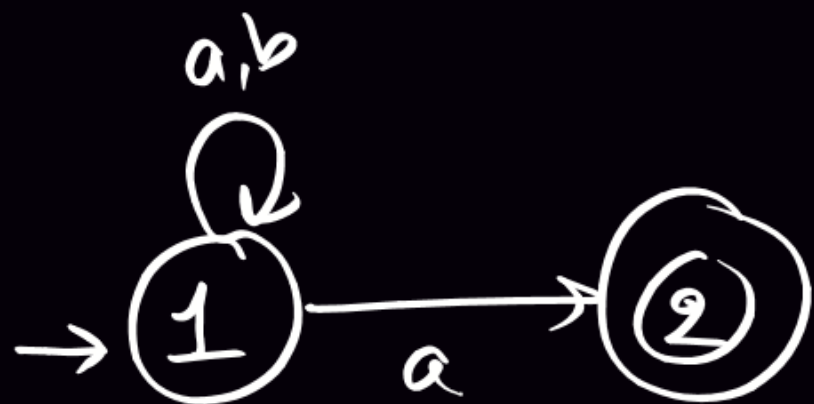


39)  $L = \{w \mid w \in (a+b)^*, w \text{ starts and ends with same symbol}\}$   
→ 5 states  
 $= a + a\Sigma^*a + b + b\Sigma^*b$



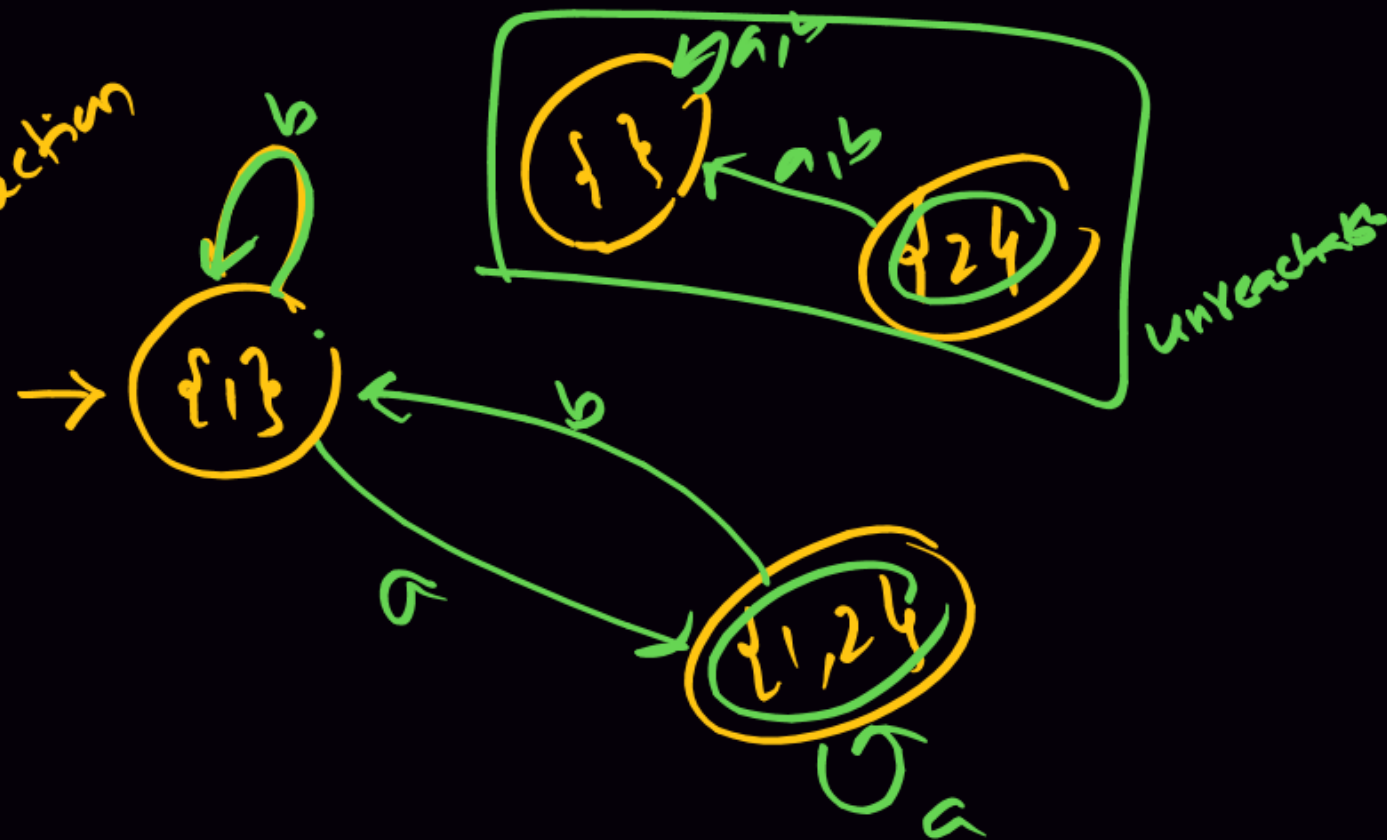
40)  $L = \{w \mid w \in (a+b)^*, w \text{ starts and ends with different symbols}\}$   
→ 5 states  
 $= a\Sigma^*b + b\Sigma^*a$

# Conversion From NFA to DFA:



NFA

Subset construction



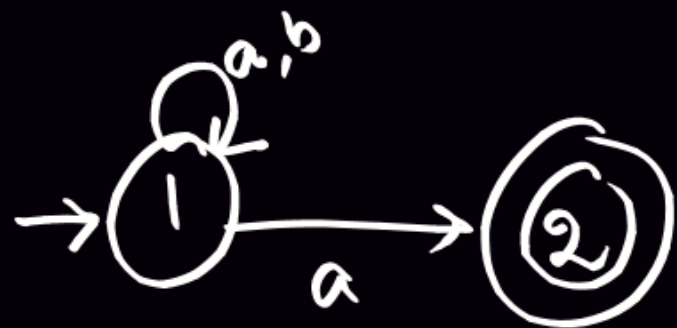
DFA

$Q_{NFA} = \{1, 2\}$

$\delta$	a	b
$\rightarrow 1$	$\{1, 2\}$	$\{1\}$
$*2$	$\{1\}$	$\{1\}$

$$Q_{DFA} = P(Q_{NFA})$$

$= \{ \emptyset, \{1\}, \{2\}, \{1, 2\} \}$   
 State of DFA  
 Subset of  $Q_{NFA}$



	a	b
→ 1	{1, 2}	{1}
* 2	{}	{}

⇒

	a	b
→ 1	{1, 2} ✓	{1} ✓
* 2	{}	{}

$$\delta_{\cup}(\{1, 2\}, a) = \delta_N(1, a) \cup \delta_N(2, a)$$



NFA  $\Rightarrow$  DFA  $\Rightarrow$  Min DFA

$n$  states       $2^n$  states      almost  $2^n$  states  
( $\leq 2^n$  states)

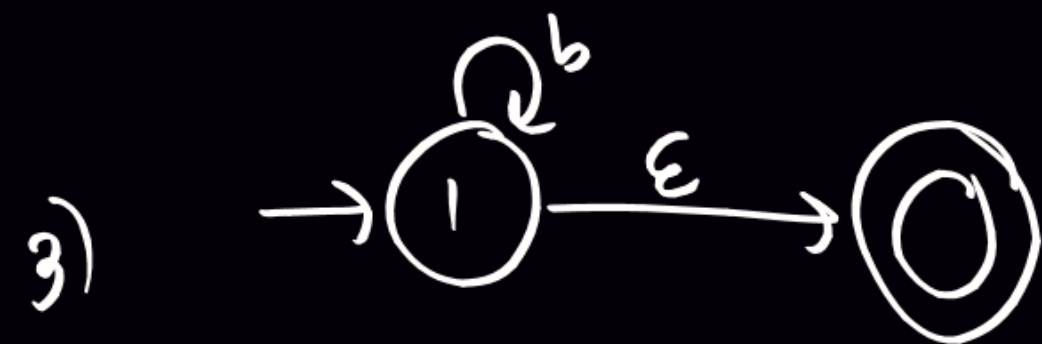


# NFA with $\epsilon$ moves:

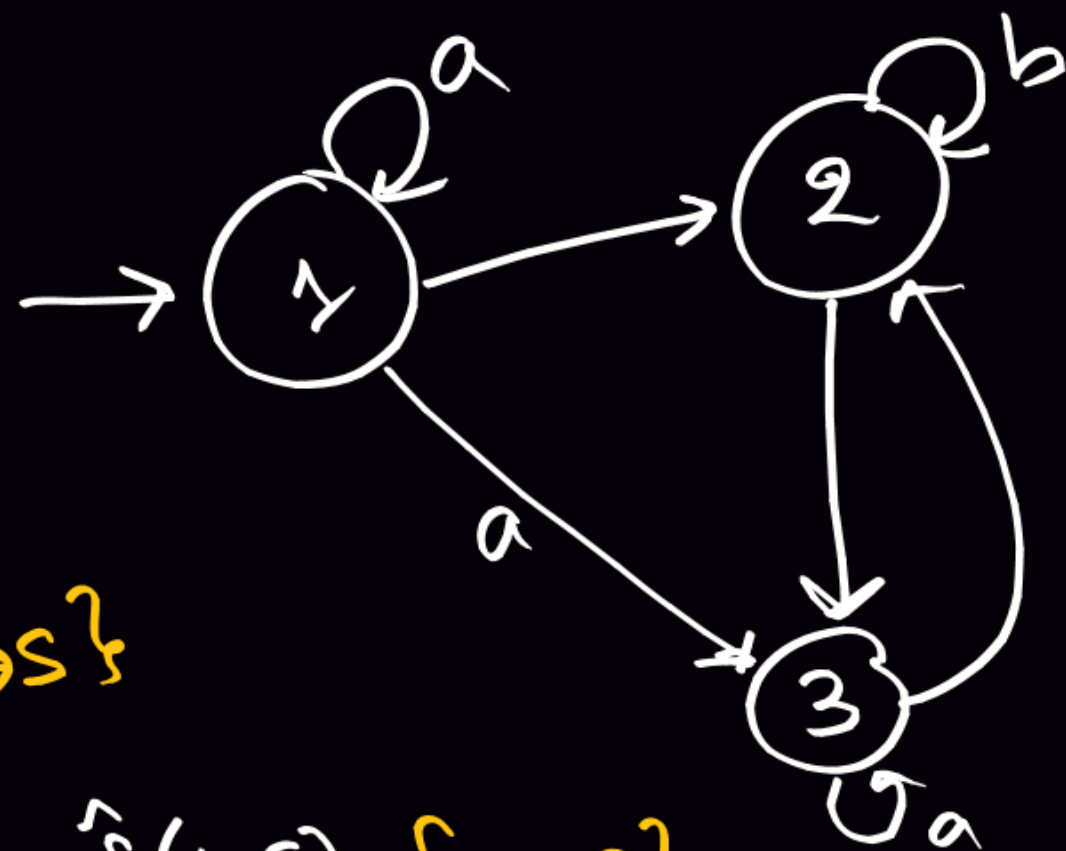


$\delta(1, \epsilon) = \{1\}$

$Q \times \Sigma \rightarrow 2^Q$



	DFA	NFA without $\epsilon$ moves	NFA with $\epsilon$ moves
1)	✓	✓	✓
2)	X	✓	✓
3)	X	X	✓
$\delta:$	$Q \times \Sigma \rightarrow Q$	$Q \times \Sigma \rightarrow 2^Q$	$Q \times \Sigma_{\epsilon} \rightarrow 2^Q$



$$\begin{aligned} \epsilon\text{-closure}(q) &= \hat{\delta}(q, \epsilon) \\ &= \{s \mid q \xRightarrow{\epsilon} s\} \end{aligned}$$

$$\epsilon\text{-closure}(1) = \hat{\delta}(1, \epsilon) = \{1, 2, 3\}$$

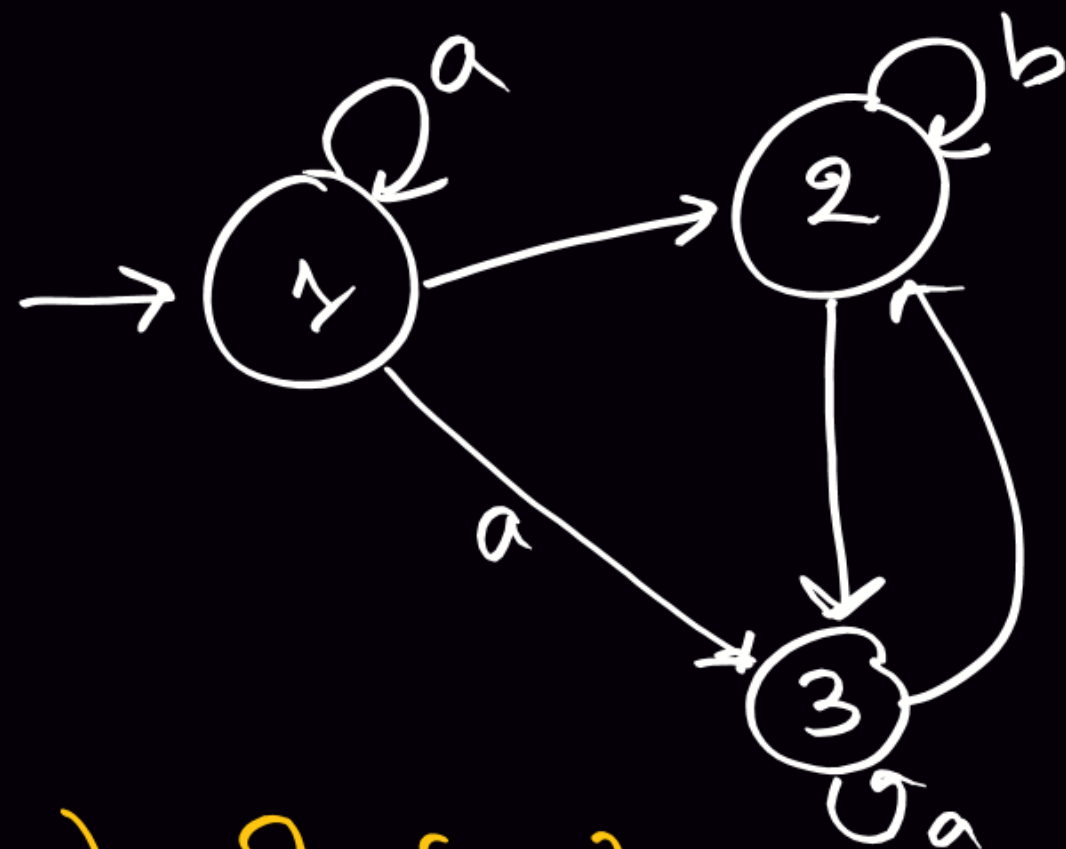
$$\epsilon\text{-closure}(2) = \hat{\delta}(2, \epsilon) = \{2, 3\}$$

$$\epsilon\text{-closure}(3) = \hat{\delta}(3, \epsilon) = \{3, 2\}$$

$$1 \xRightarrow{\epsilon} 1$$

$$1 \xRightarrow{\epsilon} 2$$

$$1 \xRightarrow{\epsilon} 3$$



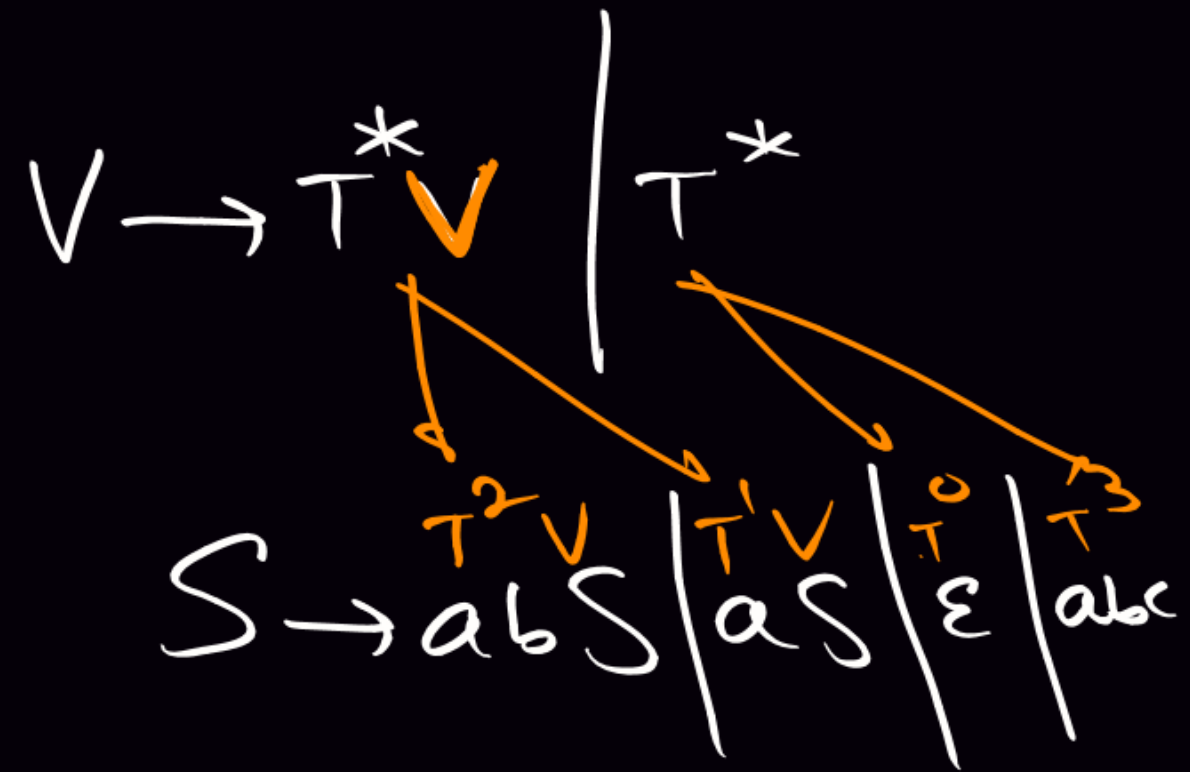
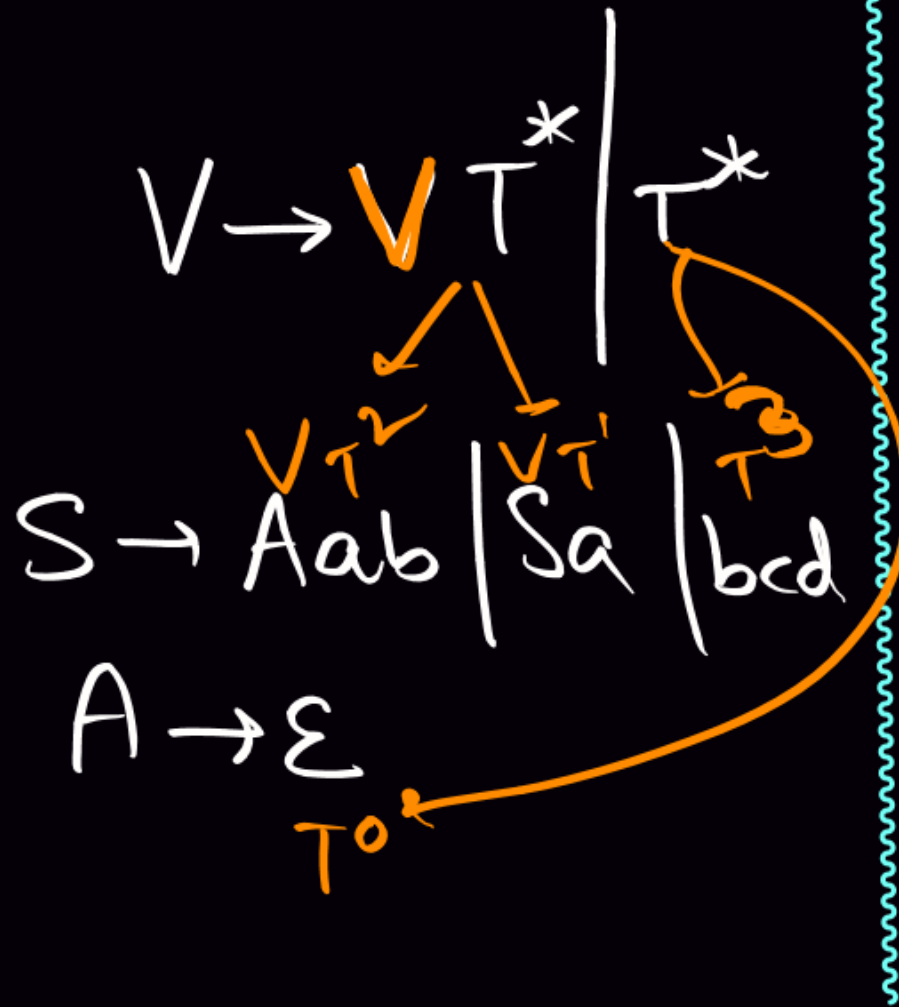
$$\hat{\delta}(2, bba) = ? = \{3, 2\}$$



# Regular Grammar (RG)

→ It is Left Linear Grammar (LLG) OR Right Linear Grammar (RLG)

$G = (V, T, P, S)$   
 → set of non-terminals





	LLG	RLG	RG
1) $S \rightarrow \varepsilon$	✓	✓	✓
2) $S \rightarrow a bc \varepsilon$	✓	✓	✓
3) $S \rightarrow Aa Bb$ $A \rightarrow \varepsilon, B \rightarrow ab$	✓	✗	✓
4) $S \rightarrow Sa bS \varepsilon$	✗	✗	✗
5) $S \rightarrow Sa Sb \varepsilon$	✓	✗	✓
6) $S \rightarrow A, A \rightarrow \varepsilon$	✓	✓	✓
7) $S \rightarrow AB, A \rightarrow a, B \rightarrow b$	✗	✗	✗

Find Language generated by following RGs.

1)  $S \rightarrow \epsilon$        $L = L(S) = \{\epsilon\}$

2)  $S \rightarrow a | \epsilon$        $L = \{a, \epsilon\}$

3)  $S \rightarrow a | b | \epsilon$        $L = \{a, b, \epsilon\}$        $\sum_{a,b,\epsilon}$        $\sum_{a,b}$        $\sum_{\epsilon}$

4)  $S \rightarrow ab | \epsilon$        $L = \{ab, \epsilon\}$

5)  $S \rightarrow ab | c | \epsilon$        $L = \{\epsilon, c, ab\}$

$$6) \begin{array}{l} S \rightarrow Aa \\ A \rightarrow b \end{array} \uparrow A=b \quad L = \{ba\}$$

$$7) \begin{array}{l} S \rightarrow Aa \mid Bb \\ A \rightarrow c \\ B \rightarrow d \end{array} \begin{array}{l} \uparrow A=c \\ \uparrow B=d \end{array} \quad L = \{ca, db\}$$

$$8) \begin{array}{l} S \rightarrow Aa \\ A \rightarrow \epsilon \mid b \end{array} \uparrow L(A) = \{\epsilon, b\} \Rightarrow L = \{a, ba\}$$

$$9) \begin{array}{l} S \rightarrow aB \\ B \rightarrow \epsilon \mid a \end{array} \uparrow B = \epsilon + a \Rightarrow L = aB = \{a, aa\}$$

$$10) S \rightarrow Sa | b \quad \boxed{L = ba^*}$$

$\left. \begin{array}{l} b \checkmark \\ ba \checkmark \\ baa \checkmark \\ baab \checkmark \\ \vdots \end{array} \right\} \Rightarrow ba^*$

$$11) S \rightarrow S \textcircled{a} | \epsilon \quad L = \epsilon \cdot a^* = a^*$$

$$12) S \rightarrow S \textcircled{a} | a \quad L = a a^* = a^+ = a^* a$$

$$13) S \rightarrow S \textcircled{aa} | \epsilon \quad L = (aa)^*$$

$$14) S \rightarrow \underbrace{Sa | Sb}_{S(a+b)} | \epsilon \quad L = (a+b)^*$$

$$15) S \rightarrow \underbrace{Sa | Sb}_{S(a+b)} | \underbrace{a | b}_{a+b} \quad L = \cancel{a^+} (a+b)^* = (a+b)^+$$

$$16) S \rightarrow aS | \epsilon \quad L = a^*$$

$$17) S \rightarrow aS | bS | \epsilon \quad L = (a+b)^*$$

$$18) S \rightarrow aS | bS | a | b \quad L = (a+b)^+$$

$$19) S \rightarrow \underbrace{aaS | bS}_{(aa+b)S} | c \quad L = (aa+b)^* c$$

$$20) S \rightarrow aA \quad A \rightarrow aA | bA | \epsilon \Rightarrow A = (a+b)^* \Rightarrow L = a(a+b)^*$$

$$21) S \rightarrow aS | bS | A$$

$$A \rightarrow aB$$

$$B \rightarrow aB | bB | \epsilon$$

H.P.

$$22) S \rightarrow Sa | Sb | A$$

$$A \rightarrow Ba$$

$$B \rightarrow Ba | Bb | \epsilon$$

$$23) S \rightarrow aS | bS | a$$

$$24) S \rightarrow Sa | Sb | a$$

$$25) S \rightarrow aA | bB$$

$$A \rightarrow aA | \epsilon$$

$$B \rightarrow bB | \epsilon$$

$$26) S \rightarrow Aa | Bb$$

$$A \rightarrow Aa | \epsilon$$

$$B \rightarrow Bb | \epsilon$$

$$27) S \rightarrow abS | \epsilon$$

$$28) S \rightarrow aA | bA$$

$$A \rightarrow a | b$$