

INSTITUTE OF TECHNOLOGY

DHULE (M.S.)

DEPARMENT OF COMPUTER ENGINEERING

Subject: Data Structure Lab **Subject Code:** 24AFAIPCL309

Class: SY AIML & Data Science **Expt. No.**: 03(B)

Title: Write a progra to implement a Doubly linked list

Problem

To implement insertion, deletion, and traversal operations on a Doubly Linked List.

Staement:

CodeBlock

Software Required

Doubly Linked List

Theory:

A Doubly Linked List in C is a versatile data structure that connects nodes in a sequential chain using pointers. Each node in this list consists of three essential components:

- **Data**: Stores the actual value or information of the node.
- **Previous Pointer**: Contains the address of the preceding node in the sequence, allowing backward traversal.
- **Next Pointer**: Holds the address of the next node, enabling forward traversal.



Unlike a Singly Linked List, which can only be traversed in the forward direction, a doubly linked list allows traversal in both directions (forward and backward). This makes certain operations, like deletion, easier and more efficient.

Representation of Doubly Linked List:

Insertion Example:

Suppose we insert: 10, 20, 30

Step 1: Insert 10 (Head = 10)

NULL <- [10] -> NULL

Step 2: Insert 20

NULL <- [10] <-> [20] -> NULL

Step 3: Insert 30

NULL <- [10] <-> [20] <-> [30] -> NULL

Deletion Example

Delete node at position 2 (value = 20):

Before: NULL <- [10] <-> [20] <-> [30] -> NULL

After: NULL <- [10] <-> [30] -> NULL

Advantages of Doubly Linked List

- Traversal is possible in both directions.
- Deletion is more efficient since we have access to the previous node.
- Insertions and deletions in the middle are faster compared to singly linked lists.

Disadvantages

Requires extra memory for the prev pointer.

More complex implementation compared to singly linked list.

Applications

- Used in implementing navigation (forward & backward in browsers).
- Undo/redo functionality in text editors.
- · Memory management systems.
- Degues and advanced data structures.

Operations in Doubly Linked List:

1. Insertion at End

- A new node is created and inserted at the end of the list.
- If the list is empty, the new node becomes the **head**.
- Otherwise, traversal is done to the last node, and the new node is linked after it.

Algorithm:

- 1. Create a new node with given data.
- 2. If the list is empty:
 - Set head = newNode.
- 3. Else:
 - Traverse to the last node.
 - Update last node's next pointer to newNode.
 - Set newNode's prev to last node.

2. Deletion of a Node (by Position)

- The node at a given position is removed.
- If it is the first node, the head pointer is updated.
- If it is a middle or last node, pointers of adjacent nodes are updated.

Algorithm:

- 1. If the list is empty, display "List is empty".
- 2. If position = 1:
 - Update head to head->next.
 - If new head exists, set head->prev = NULL.
 - Free the first node.
- 3. Else:

- Traverse to the node at given position.
- Update its previous node's next to its next.
- Update its next node's prev to its prev.
- Free the node.

3. Display

- Traverse from head to the end, printing each node's data.
- Connections are shown using <->.

```
#include <stdio.h>
#include <stdlib.h>
// Structure of a node
struct Node {
  int data;
  struct Node* next;
  struct Node* prev;
struct Node* head = NULL; // Global head pointer
// Insert element at end
void insertElement(int value) {
  struct Node* newNode = (struct Node*)malloc(sizeof(struct Node));
  newNode->data = value;
  newNode->next = NULL;
  if (head == NULL) {
    newNode->prev = NULL;
    head = newNode;
  } else {
    struct Node* temp = head;
    while (temp->next != NULL) {
      temp = temp->next;
    temp->next = newNode;
    newNode->prev = temp;
  }
// Delete node by position
void deleteNode(int pos) {
  if (head == NULL) {
    printf("List is empty\n");
    return;
  struct Node* temp = head;
  if (pos == 1) {
    head = temp->next;
```

```
if (head != NULL) {
      head->prev = NULL;
    free(temp);
    return;
  }
  for (int i = 1; i < pos && temp != NULL; i++) {
    temp = temp->next;
  }
  if (temp == NULL) {
    printf("Position out of range!\n");
    return;
  }
  if (temp->prev != NULL) {
    temp->prev->next = temp->next;
  }
  if (temp->next != NULL) {
    temp->next->prev = temp->prev;
  }
  free(temp);
// Display the list
void display() {
  struct Node* temp = head;
  if (temp == NULL) {
    printf("List is empty\n");
    return;
  printf("Doubly Linked List: ");
Output:
--- Doubly Linked List Menu ---
1. Insert Element
2. Delete Node (by position)
3. Display List
4. Exit
Enter choice: 1
Enter value: 10
Enter choice: 1
Enter value: 20
Enter choice: 1
Enter value: 30
```

Enter choice: 3

Doubly Linked List: 10 <-> 20 <-> 30 <-> NULL

Enter choice: 2

Enter position to delete: 2

Enter choice: 3

Doubly Linked List: 10 <-> 30 <-> NULL



INSTITUTE OF TECHNOLOGY

DHULE (M.S.) DEPARMENT OF COMPUTER ENGINEERING

Subject: Data Structure Lab Subject Code: 24AFAIPCL309

Class: SY AIML & Data Science Expt. No.: 03(C)

Title: Write a program to implement a Circular Linked List

Problem Staement To implement insertion, deletion, and traversal operations on a Circular Linked

List.

Software Required

CodeBlock

Theory:

Circular Linked List

A **linked list** is a linear data structure where elements (called **nodes**) are connected using pointers. Unlike arrays, linked lists do not require contiguous memory allocation, which makes insertion and deletion operations more efficient.

A Circular Singly Linked List (CSLL) is a special type of linked list where:

- Each node contains data and a pointer to the next node.
- The **last node** does not point to NULL, but instead points back to the **head node**, forming a **circle**.

This circular connection allows traversal to continue from the last node back to the first node, making the list suitable for applications where a circular traversal is needed (e.g., round-robin scheduling, circular queues).

1. Circular Singly Linked List

Here, the address of the last node consists of the address of the first node.



Advantages of Circular Singly Linked List

- 1. Efficient traversal: You can start at any node and reach all others in a circular manner.
- 2. Saves space compared to linear linked lists where last node points to NULL.
- 3. Useful in **applications requiring cyclic iteration** (e.g., multiplayer games, round-robin scheduling, buffers).

Disadvantages

- 1. Traversal is more complex compared to linear linked lists.
- 2. Extra care is needed while performing insertion and deletion to maintain circular structure.
- 3. No direct access to elements (sequential access only).

Algorithm:

Insertion at End

```
Step 1: Create newNode and assign value to data.

Step 2: If head == NULL:
    head = newNode
    newNode->next = head

Else:
    temp = head
    While (temp->next != head):
    temp = temp->next
    temp->next = newNode
    newNode->next = head
```

```
2. Deletion by Position
Step 1: If head == NULL \rightarrow print "List is empty" and return.
Step 2: If pos == 1:
       If head->next == head:
         free(head)
         head = NULL
       Else:
         Find last node
         last->next = head->next
         temp = head
         head = head->next
         free(temp)
Step 3: Else:
       Traverse list until position
       If position invalid → print "Out of range"
       Else adjust prev->next = temp->next
          free(temp)
3. Display List
Step 1: If head == NULL \rightarrow print "List is empty" and return.
Step 2: Start with temp = head
Step 3: Do
       Print temp->data
       temp = temp->next
    While (temp != head)
Step 4: Print (head) to indicate circular nature.
#include <stdio.h>
#include <stdlib.h>
// Structure of a node
struct Node {
 int data;
 struct Node* next;
};
```

struct Node* head = NULL; // Global head pointer

struct Node* newNode = (struct Node*)malloc(sizeof(struct Node));

// Function to insert element at the end

void insertElement(int value) {

newNode->data = value;

```
if (head == NULL) { // First node
    head = newNode;
    newNode->next = head; // Circular link
  } else {
    struct Node* temp = head;
    while (temp->next != head) { // Traverse till last node
      temp = temp->next;
    temp->next = newNode;
    newNode->next = head; // Point back to head
  }
}
// Function to delete node by position
void deleteNode(int pos) {
  if (head == NULL) {
    printf("List is empty\n");
    return;
  }
  struct Node* temp = head;
  // Case 1: Delete first node
  if (pos == 1) {
    if (head->next == head) { // Only one node
      free(head);
      head = NULL;
      return;
    }
    struct Node* last = head;
    while (last->next != head) { // Find last node
      last = last->next;
    last->next = head->next;
    head = head->next;
    free(temp);
    return;
  }
  // Case 2: Delete at other position
  struct Node* prev = NULL;
  for (int i = 1; i < pos && temp->next != head; <math>i++) {
    prev = temp;
    temp = temp->next;
  }
  if (temp == head) { // Position out of range
    printf("Position out of range!\n");
```

```
return;
  }
  prev->next = temp->next;
  free(temp);
// Function to display the list
void display() {
  if (head == NULL) {
    printf("List is empty\n");
    return;
  }
  struct Node* temp = head;
  printf("Circular Singly Linked List: ");
  do {
    printf("%d -> ", temp->data);
    temp = temp->next;
  } while (temp != head);
  printf("(head)\n");
}
// Main function
int main() {
  int choice, value, pos;
  while (1) {
    printf("\n--- Circular Singly Linked List Menu ---\n");
    printf("1. Insert Element\n");
    printf("2. Delete Node (by position)\n");
    printf("3. Display List\n");
    printf("4. Exit\n");
    printf("Enter choice: ");
    scanf("%d", &choice);
    switch (choice) {
       case 1:
         printf("Enter value: ");
         scanf("%d", &value);
         insertElement(value);
         break;
       case 2:
         printf("Enter position to delete: ");
         scanf("%d", &pos);
         deleteNode(pos);
         break;
       case 3:
         display();
         break;
```

```
case 4:
         exit(0);
       default:
         printf("Invalid choice!\n");
    }
  }
  return 0;
Output
--- Circular Singly Linked List Menu ---
1. Insert Element
2. Delete Node (by position)
3. Display List
4. Exit
Enter choice: 1
Enter value: 10
--- Circular Singly Linked List Menu ---
1. Insert Element
2. Delete Node (by position)
3. Display List
4. Exit
Enter choice: 1
Enter value: 20
--- Circular Singly Linked List Menu ---
1. Insert Element
2. Delete Node (by position)
3. Display List
4. Exit
Enter choice: 1
Enter value: 30
--- Circular Singly Linked List Menu ---
1. Insert Element
2. Delete Node (by position)
3. Display List
4. Exit
Enter choice: 3
Circular Singly Linked List: 10 -> 20 -> 30 -> (head)
--- Circular Singly Linked List Menu ---
1. Insert Element
2. Delete Node (by position)
3. Display List
4. Exit
Enter choice: 2
```

Enter position to delete: 2 --- Circular Singly Linked List Menu ---1. Insert Element 2. Delete Node (by position) 3. Display List 4. Exit Enter choice: 3 Circular Singly Linked List: 10 -> 30 -> (head) --- Circular Singly Linked List Menu ---1. Insert Element 2. Delete Node (by position) 3. Display List 4. Exit Enter choice: 1 Enter value: 40 --- Circular Singly Linked List Menu ---1. Insert Element 2. Delete Node (by position) 3. Display List 4. Exit Enter choice: 3 Circular Singly Linked List: 10 -> 30 -> 40 -> (head) --- Circular Singly Linked List Menu ---1. Insert Element 2. Delete Node (by position) 3. Display List 4. Exit Enter choice: 4



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DHULE (M.S.) DEPARMENT OF COMPUTER ENGINEERING

Subject : Data Structure Lab **Subject Code :** 24AFAIPCL309

Class: SY AIML & Data Science Expt. No.: 04

Title: Write a program to implement a Stack Using linked list such that the push and pop operation of the stack still take O(1) time.

Problem Staement To implement a Stack Using linked list

Software Required CodeBlock

Theory:

Stack Using Linked List

What is a Stack?

A stack is a linear data structure that follows the Last In First Out (LIFO) principle. This means the element added last is removed first.

Stack Operations:

- **Push** Insert an element onto the stack
- **Pop** Remove the top element
- **Peek/Top** View the top element without removing it
- **isEmpty** Check if the stack is empty

Why Use Linked List for Stack?

- Dynamic size: No need to predefine size as in array-based stack
- Efficient: Memory is allocated only when needed

Linked List Stack Structure:

Each node contains:

- data value of the element
- next pointer to the next node

The **top** of the stack is represented by the head of the linked list.

Algorithm

```
PUSH(x)
```

```
1. Create a new node
```

- 2. Assign data = x
- 3. Set new_node->next = top
- 4. Update top = new_node

POP()

- 1. If top is NULL
 - Stack Underflow (Nothing to pop)
- 2. Else
 - Temp = top
 - top = top->next
 - Delete temp

PEEK()

- 1. If top is NULL
 - Stack is empty
- 2. Else
 - Return top->data

isEmpty()

#include <stdio.h>
#include <stdlib.h>

Return true if top == NULL, else false

```
// Define structure for a stack node
struct Node {
   int data;
   struct Node* next;
};

// Initialize stack as empty
struct Node* top = NULL;

// Function to push an element onto the stack
void push(int value) {
   struct Node* newNode = (struct Node*)malloc(sizeof(struct Node));
   if (newNode == NULL) {
      printf("Stack Overflow! Unable to allocate memory.\n");
      return;
   }
   newNode->data = value;
```

```
newNode->next = top; // Link new node to previous top
  top = newNode;
                      // Make new node the top
  printf("%d pushed to stack.\n", value);
// Function to pop an element from the stack
void pop() {
  if (top == NULL) {
    printf("Stack Underflow! No elements to pop.\n");
  struct Node* temp = top;
  printf("%d popped from stack.\n", top->data);
  top = top->next;
  free(temp);
// Function to peek at the top element
void peek() {
  if (top == NULL) {
    printf("Stack is empty! Nothing to peek.\n");
  } else {
    printf("Top element is: %d\n", top->data);
  }
}
// Function to display elements of the stack
void display() {
  if (top == NULL) {
    printf("Stack is empty.\n");
    return;
  }
  struct Node* temp = top;
  printf("Stack elements (top to bottom): ");
  while (temp != NULL) {
    printf("%d -> ", temp->data);
    temp = temp->next;
  }
  printf("NULL\n");
// Main function to perform stack operations
int main() {
  int choice, value;
  while (1) {
    printf("\n--- Stack using Linked List ---\n");
    printf("1. Push\n");
    printf("2. Pop\n");
    printf("3. Peek\n");
    printf("4. Display\n");
    printf("5. Exit\n");
    printf("Enter your choice: ");
    scanf("%d", &choice);
    switch (choice) {
      case 1:
```

```
printf("Enter value to push: ");
         scanf("%d", &value);
         push(value);
         break;
       case 2:
         pop();
         break;
       case 3:
         peek();
         break;
       case 4:
         display();
         break;
       case 5:
         printf("Exiting...\n");
         exit(0);
       default:
         printf("Invalid choice! Please try again.\n");
    }
  }
 return 0;
Output:
--- Stack using Linked List ---
1. Push
2. Pop
3. Peek
4. Display
5. Exit
Enter your choice: 1
Enter value to push: 10
10 pushed to stack.
--- Stack using Linked List ---
1. Push
2. Pop
3. Peek
4. Display
5. Exit
Enter your choice: 1
Enter value to push: 20
20 pushed to stack.
--- Stack using Linked List ---
Enter your choice: 4
Stack elements (top to bottom): 20 -> 10 -> NULL
Enter your choice: 3
Top element is: 20
Enter your choice: 2
20 popped from stack.
Enter your choice: 4
```

Stack elements (top to bottom): 10 -> NULL



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DHULE (M.S.) DEPARMENT OF COMPUTER ENGINEERING

Subject: Data Structure Lab Subject Code: 24AFAIPCL309

Class: SY AIML & Data Science Expt. No.: 7

Title: Write a program to implement a Insertion Sorting

Problem Staement

To implement insertion sort.

Software Required

CodeBlock

Theory:

Insertion Sort is one of the simplest and most intuitive sorting algorithms. It is based on the idea of **building a sorted list one element at a time** by comparing and inserting each new element into its correct position in an already sorted part of the list.

Concept and Working Principle

- The array is divided into two parts:
 - Left side (sorted part) elements that are already arranged in order.
 - Right side (unsorted part) elements that are yet to be arranged.
- The algorithm starts from the **second element** (index 1), assuming that the first element is already sorted.
- It picks the current element (called **key**) and compares it with the elements on the left side.
- All elements that are greater than the key are shifted one position to the right to make space for the key.
- The key is then placed in its correct sorted position.

• This process is repeated for all elements in the array until the entire list becomes sorted.

Advantages:

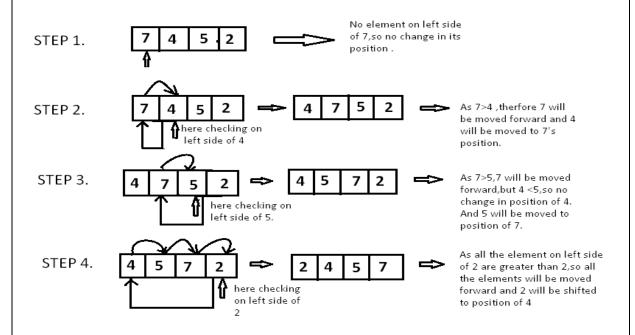
- 1. Simple and easy to understand.
- 2. Efficient for small datasets.
- 3. Performs well for nearly sorted data.
- 4. Requires no additional memory (in-place algorithm).
- 5. Stable sort preserves relative order of equal elements.

5. Disadvantages:

- 1. Inefficient for large datasets $(O(n^2)$ time complexity).
- 2. More comparisons and shifts as the number of elements increases.
- 3. Not suitable for large lists compared to algorithms like Merge Sort or Ouick Sort.

Example:

Take Array[]=[7,4,5,2]



ALGORITHM:

- 1. Start
- 2. Read the number of elements and the array elements.
- 3. Repeat for i = 1 to n-1:
 - o Set temp = arr[i]
 - \circ Set j = i 1
 - While j >= 0 and arr[j] > temp:
 - Move arr[j] to position j + 1
 - Decrement j by 1
 - Insert key at position j + 1
- 4. Display the sorted array.
- 5. Stop

```
#include <stdio.h>
int main() {
  int a[100], n, i, j, temp;
  printf("Enter number of elements: ");
  scanf("%d", &n);
  printf("Enter %d elements:\n", n);
  for (i = 0; i < n; i++)
    scanf("%d", &a[i]);
  // Insertion Sort logic
  for (i = 1; i < n; i++) {
    temp = a[i];
    j = i - 1;
    while (j \ge 0 \&\& a[j] > temp) {
       a[j + 1] = a[j];
       j--;
    a[j + 1] = temp;
```

```
printf("Sorted array:\n");
for (i = 0; i < n; i++)
printf("%d ", a[i]);

return 0;
}

Output:
Enter number of elements: 5
Enter 5 elements:
5 2 4 6 1
Sorted array:
1 2 4 5 6
```



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DHULE (M.S.) DEPARMENT OF COMPUTER ENGINEERING

Subject: Data Structure Lab Subject Code: 24AFAIPCL309

Class: SY AIML & Data Science Expt. No.: 8

Title: Write a program to implement a Merge Sort

Problem Staement

To implement merge sort

Software Required CodeBlock

Theory:

Theory:

Merge Sort is a divide and conquer sorting algorithm.

It divides the unsorted array into smaller subarrays until each subarray contains a single element, then merges these subarrays to produce new sorted subarrays until there is only one sorted array remaining.

Working Steps:

- 1. **Divide:** Split the array into two halves.
- 2. **Conquer:** Recursively sort the two halves.
- 3. **Combine:** Merge the two sorted halves into a single sorted array.

Concept of Divide and Conquer:

The **Divide and Conquer** approach breaks a problem into smaller subproblems, solves each subproblem independently, and then combines the results to obtain the final solution.

Merge Sort applies this idea in three main steps:

1. Divide:

The array is divided into two halves (approximately equal parts).

This division continues recursively until each subarray contains a single element.

2. Conquer (Sort):

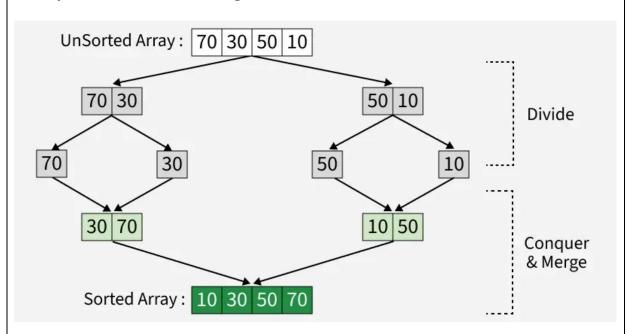
Each of these smaller subarrays is considered sorted because a single element is trivially sorted.

3. Combine (Merge):

The sorted subarrays are merged together in a way that the resulting array is also sorted.

This merging process continues until all subarrays are combined into one fully sorted array.

Example to Understand Merge Sort:



Let's consider an example array:

[70,30,50,10]

Step 1 – Divide:

Divide the array into two halves recursively:

[70,30] and [50,10]

And again:

[70] [30] [50] [10]

Step 2 - Conquer (Merge):

Now merge pairs of subarrays while sorting them:

[30,70] [10,50]

Step 3 – Combine:

Continue merging the sorted subarrays:

[10,30,50,70]

Finally:

[10,30,50,70]

Thus, the array is completely sorted.

Advantages:

- 1. Very efficient for large datasets.
- 2. Guarantees **O(n log n)** time complexity for all cases.
- 3. It is a **stable** sorting algorithm.
- 4. Performs well on linked lists (does not require random access).
- 5. Can be easily implemented using recursion.

Disadvantages:

- 1. Requires **extra memory** (O(n)) for temporary arrays during the merge process.
- 2. Not suitable for small datasets compared to simpler algorithms like **Insertion Sort** or **Bubble Sort**.
- 3. Recursive calls increase overhead and may cause stack overflow for very large arrays.

Algorithm: MergeSort(A, low, high)

- 1. If low < high
 - 1. Find mid = (low + high) / 2
 - 2. Call MergeSort(A, low, mid)
 - 3. Call MergeSort(A, mid + 1, high)
 - 4. Call Merge(A, low, mid, high)

Algorithm: Merge(A, low, mid, high)

- 1. Create temporary arrays for left and right subarrays.
- 2. Compare elements of both subarrays and copy the smaller one to the main array.
- 3. Copy the remaining elements of left or right subarray (if any).

```
#include <stdio.h>
void merge(int a[], int left, int mid, int right) {
  int temp[100];
  int i = left: // start of first half
  int j = mid + 1; // start of second half
  int k = left; // start of temp array
  // Compare and copy smaller element into temp
  while (i <= mid && j <= right) {
     if (a[i] < a[i])
       temp[k++] = a[i++];
     else
       temp[k++] = a[j++];
  }
  // Copy remaining elements (if any) from first half
  while (i \le mid)
     temp[k++] = a[i++];
  // Copy remaining elements (if any) from second half
  while (j <= right)
     temp[k++] = a[j++];
  // Copy back to original array
  for (i = left; i \le right; i++)
     a[i] = temp[i];
}
void mergeSort(int a[], int left, int right) {
  if (left < right) {
    int mid = (left + right) / 2;
    // Divide the array into two halves
```

```
mergeSort(a, left, mid);
    mergeSort(a, mid + 1, right);
    // Merge the two sorted halves
    merge(a, left, mid, right);
}
int main() {
  int a[100], n, i;
  printf("Enter number of elements: ");
  scanf("%d", &n);
  printf("Enter %d elements:\n", n);
  for (i = 0; i < n; i++)
    scanf("%d", &a[i]);
  // Call merge sort
  mergeSort(a, 0, n - 1);
  printf("Sorted array:\n");
  for (i = 0; i < n; i++)
    printf("%d ", a[i]);
  return 0;
}
Output:
Enter number of elements: 6
Enter 6 elements:
1035281
Sorted array:
1235810
```