

## Problem 5

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1.  $P(\text{Yes}|\text{Senioritis}) = 4/5$   
 $P(\text{No}|\text{Senioritis}) = 1/5$   
 $P(\text{Yes}|\text{No Senioritis}) = 0$   
 $P(\text{No}|\text{No Senioritis}) = 1$

Given  $p = 1 - \frac{1}{\text{poly}(n)}$

So, we see that algorithm A always provides correct answer if the person does not have senioritis. However, if the person has senioritis, then the algorithm detects it with probability  $4/5$ .

In this part, we need to increase this probability to  $p$ .

For this problem, we output YES if the algorithm gives yes in any of the  $T$  trials, otherwise, we output NO.

The probability of not detecting the disease, given that the person has disease =  $1/5$

If we run this algorithm multiple times, let say for  $T = \log_5(\text{poly}(n))$ :

$$\begin{aligned} \text{probability of failure for all the } T \text{ iterations} &= \left(\frac{1}{5}\right)^T \\ &= \left(\frac{1}{5}\right)^{\log_5(\text{poly}(n))} \\ &= \left(\frac{1}{\text{poly}(n)}\right)^{\log_5 5} \\ &\leq \frac{1}{\text{poly}(n)} \end{aligned}$$

Probability that algorithm predicts correctly when the person has disease  $\geq 1 - \frac{1}{\text{poly}(n)}$

Probability that algorithm predicts correctly when the person does not has disease =  $1 \geq 1 - \frac{1}{\text{poly}(n)}$  as the algorithm A always gives the correct results for NO instances.

Analysing TC:

We know that algorithm A runs in  $O(n)$ . As in the algorithm discussed above, we run the algorithm A for  $T = (\log(\text{poly}(n)))$  trials, so we can say that the overall complexity of the proposed algorithm is  $O(n \log(\text{poly}(n)))$

2.  $P(\text{Yes}|\text{Senioritis}) = 4/5$   
 $P(\text{No}|\text{Senioritis}) = 1/5$

$$P(\text{Yes}|\text{No Senioritis}) = 1/5$$

$$P(\text{No}|\text{No Senioritis}) = 4/5$$

$$\text{Given } p = 1 - \frac{1}{\text{poly}(n)}$$

So, we see that algorithm A provides correct answer if the person does not have senioritis with a probability of 4/5. Also, if the person has senioritis, then the algorithm detects it with probability 4/5.

In this part, we need to increase both these probabilities to  $p$ .

For this problem, we output the majority answer which we get after  $T$  trials.

The probability of not detecting the disease, given that the person has disease = 1/5. The probability of detecting the disease, given that the person does not has disease = 1/5

If we run this algorithm multiple times, let say for  $T = \frac{50}{9} \log(\text{poly}(n))$ :

$$\begin{aligned} \text{probability of failure in majority of the } T \text{ iterations} &= \sum_{i=0}^{T/2} \binom{T}{i} \left(\frac{4}{5}\right)^i \left(\frac{1}{5}\right)^{T-i} \\ &\leq \sum_{i=0}^{T/2} \binom{T}{i} \left(\frac{4}{5}\right)^{T/2} \left(\frac{1}{5}\right)^{T/2} \\ &= \left(\frac{4}{25}\right)^{T/2} \sum_{i=0}^{T/2} \binom{T}{i} \\ &\leq \left(\frac{4}{25}\right)^{T/2} 2^{T/2} \\ &= \left(1 - \frac{9}{25}\right)^{T/2} \quad \text{by using the inequality } 1 - x \leq e^{-x} \\ &\leq e^{-\frac{9}{50}T} \\ &\leq \frac{1}{\text{poly}(n)} \end{aligned}$$

$$\text{Probability that algorithm predicts correctly} \geq 1 - \frac{1}{\text{poly}(n)}$$

Hence, the algorithm would now give the correct result whether the person has disease or not with a probability atleast  $p$ .

Analysing TC:

We know that algorithm A runs in  $O(n)$ . As in the algorithm discussed above, we run the algorithm A for  $T = (\log(\text{poly}(n)))$  trials, so we can say that the overall complexity of the proposed algorithm is  $O(n \log(\text{poly}(n)))$

3. For this part, consider  $q = 1/2 + \epsilon$

Now, we need to show that there exists an algorithm which runs in polynomial time which gives correct answers with probability atleast 3/4.

Using the same algorithm as used in part (b).

We run this algorithm multiple times, let say for  $T = \frac{1}{2\epsilon^2} \log(4)$ :

$$\begin{aligned} \text{probability of failure in majority of the } T \text{ iterations} &= \sum_{i=0}^{T/2} \binom{T}{i} \left(\frac{1}{2} + \epsilon\right)^i \left(\frac{1}{2} - \epsilon\right)^{T-i} \\ &\leq \sum_{i=0}^{T/2} \binom{T}{i} \left(\frac{1}{2} + \epsilon\right)^{T/2} \left(\frac{1}{2} - \epsilon\right)^{T/2} \\ &= \left(\frac{1 - 4\epsilon^2}{4}\right)^{T/2} \sum_{i=0}^{T/2} \binom{T}{i} \\ &\leq \left(\frac{1 - 4\epsilon^2}{4}\right)^{T/2} 2^{T/2} \\ &= (1 - 4\epsilon^2)^{T/2} \quad \text{by using the inequality } 1 - x \leq e^{-x} \\ &\leq e^{-2\epsilon^2 T} \\ &\leq \frac{1}{4} \end{aligned}$$

Probability that algorithm predicts correctly  $\geq 1 - \frac{1}{4}$

Hence, the algorithm would now give the correct result whether the person has disease or not with a probability atleast  $3/4$ .