

Iterative DFS

- ▶ Use a stack to allow for backtracking during DFS.
- ▶ Initialize stack by placing a start vertex v .
- ▶ As long as stack is nonempty pop the last vertex and mark it visited if it is not visited.
- ▶ Then push all the other end vertices of the edges incident on the current vertex.

Iterative DFS

```
IterativeDFS( $G, v$ ) {  
  // Initialization  
  index = 0;  
   $T = \Phi$ ;  
  makeNull( $S$ ); // Define an empty stack  
  for all ( $v \in V$ )  
    mark[ $v$ ] = unvisited;  
  choose( $s$ ); // Start vertex  
   $S.push(s)$ ;  
  // Remaining part in next slide  
}
```

Iterative DFS

```
while (!isEmpty(S)) {  
    v = S.pop();  
    if (marked[v] == "unvisited") {  
        mark[v] = "visited";  
        dfn[v] = ++index;  
        for all  $w \in \text{ADJ}_G(v)$  {  
            if (marked[w] == "unvisited") {  
                S.push(w);  
            }  
        }  
    }  
}
```

Iterative DFS

- ▶ Can have multiple copies vertices on the stack.
- ▶ But total number of iterations of stack loop cannot exceed number edges.
- ▶ Thus size of stack cannot exceed $|E|$.
- ▶ Try out how you can avoid having multiple copies a vertex in the stack.

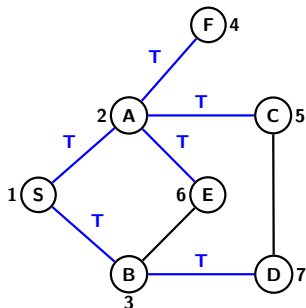
Breadth First Search

```
// Initialization  
index = 0;  
T =  $\Phi$ ;  
Q = NULL;  
for all ( $v \in V$ )  
    mark[v] = "unvisited";  
choose( $s$ ); // Start vertex  
bfn[ $s$ ] = ++index;  
ENQUEUE(Q,  $s$ );  
// Remaining part in next slide
```

Breadth First Search

```
while (!isEmpty(Q)) {  
    v = DEQUEUE(Q);  
    mark[v] = "visited";  
    for all (w ∈ ADJG(v)) {  
        if (mark[w] == "unvisited") {  
            mark[w] = "visited";  
            T = T ∪ {(v, w)};  
            bfn[v] = ++index;  
            ENQUEUE(Q, w);  
        }  
    }  
}
```

Breadth First Search



v	w	Action	Queue
-	-	bf _n (S) = 1	{S}
S	A	bf _n (A) = 1	{A}
	B	bf _n (B) = 2	{A,B}
A	F	bf _n (F) = 4	{B,F}
	C	bf _n (C) = 5	{B,F,C}
	E	bf _n (E) = 6	{B,F,C,E}
B	D	BF _N (D) = 7	{F,C,E,D}
	E	None	{F,C,E,D}
	S	None	{F,C,E,D}
F	A	None	{C, E, D}
C	A	None	{E, D}
	D	None	{E, D}
E	A	None	{D}
	B	None	{D}
D	C	None	{}
	B	None	{}

Classification of Edges by BFS

- ▶ There can be no back edges or forward edges in BFS of undirected graphs.
- ▶ For each tree edge (u, v) , $\text{dist}[v] = \text{dist}[u] + 1$
- ▶ For each cross edge (u, v) $\text{dist}[u] = \text{dist}[v]$ or $\text{dist}[v] = \text{dist}[u] + 1$