## ESO207A: Data Structures and Algorithms Theoretical Assignment 4

Deadline: April 18, 2018

April 7, 2018

## Instructions

- You are not allowed (explicitly) to use the internet for this assignment unless otherwise specified.
- All submissions must be in the pdf format.
- Please write precise answers.
- Marks of subparts of questions may not be evenly distributed.
- This assignment may take time and hence you are advised to start early.

**Problem 1.** (40 marks) The DFS that you learnt in class is not just a tool to traverse a graph but a very powerful tool. You have seen its application in computing strongly connected components of a directed Graph. We will see one more application of DFS in this assignment. Note that we will be dealing only with **directed** graphs unless otherwise specified.

Consider a directed Graph G = (V,E). Let |V| = m and |E| = n. A graph G is said to be a **Unique Path Graph** iff for every  $u \in V$ , if  $v \in V$  is reachable from u then there exists a unique path from u to v.

Given a graph G = (V, E) the task is to determine whether G is a unique path graph.

- 1 [20 marks] Consider the DFS tree  $T_v$  of a vertex v. If G is a unique path graph can there exist (1) forward edges, (2) cross edges and (3) backward edges in  $T_v$ ? Extend your argument to all vertices of G to give a necessary and sufficient condition (iff) for G to be a unique path graph. Prove it.
- 2 [20 marks] Convert and maybe modify the above into an  $\mathcal{O}(m^2)$  time algorithm and give a pseudo-code for your algorithm. (A C code will not fetch any credit. An  $\mathcal{O}(mn)$  time algorithm will fetch 25% credit).

**Problem 2.** (60 marks) Consider a set of N Boolean variables  $X = \{x_1, \ldots, x_N\}$ . Consider an expression E in the Conjunctive Normal Form (CNF) with K conjunctions such that  $E = a_1 \wedge a_2 \wedge \ldots a_K$  where each  $a_k$  has the form  $a_k = y_i \vee y_j$  where  $y_i, y_j$  are either variables in X or their negations. For example, E can be the expression  $E = (x_1 \vee x_3) \wedge (\neg x_1 \vee x_4) \wedge (x_2 \vee x_3) \wedge (x_2 \vee \neg x_4)$ .

Your task is to determine whether a given expression E is satisfiable, that is, does there exist an assignment of truth values of variables in X such that E becomes true. If yes, then you must report the corresponding assignment of truth values.

The following sub-parts will guide you towards the solution

- 1 [15 marks] Observe that  $a \lor b \equiv \neg a \implies b$ . Convert E into this new form and use it to construct a graph. Clearly specify your Vertex set and Edge set.
- 2 [25 Marks] Relate the problem to a problem covered in class to determine whether E is satisfiable. The following observations may help
  - (a) a and  $\neg a$  cannot have the same truth value
  - (b) If  $a \implies b \implies c \implies a$ , then a, b and c must have the same truth value.
- 3 [20 Marks] Finally if E is satisfiable, give a method to assign truth values to the variables in X.