MLDS-413 Introduction to Databases and Information Retrieval

Lecture 6 Extended ER Diagrams, SELECT Query Steps

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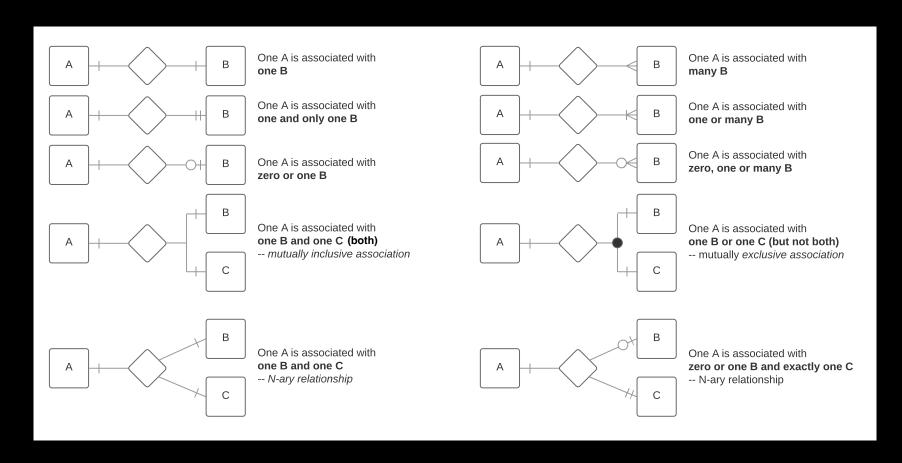
Last Lecture

- Three policies for handling deletion of foreign keys
 - Restrict, Cascade, set NULL
- Table relationships
 - One to Many: most foreign keys
 - One to One: primary key is foreign key. Used for subset tables
 - Many to Many: implemented with a linking table
- Introduced ER diagrams
 - Cardinality & participation constraints
 - Crow's foot notation
 - Schema vs. instance
- Introduced SQL
 - Syntax diagrams and grammars
 - SELECT queries with filtering, sorting, limiting, arithmetic, and grouping

Overview of Database Design

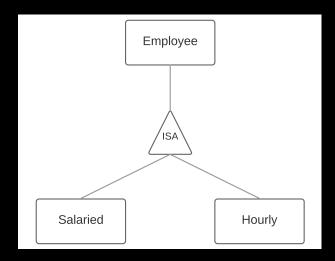
- Requirements analysis
- Conceptual design: (ER Model is used at this stage)
 - What are the entities and relationships (and events) in the system we are describing?
 - It is typical to define only entities and relationships; events can be considered as entities
 - What information about these entities and relationships should we store in the database?
 - These are the "attributes" of the entity or relationship
 - Which of the attributes are primary keys?
 - · Uniquely distinguish items that belong to the same entity
 - Primary keys are indicated in ER diagrams by underlying them
 - What are the *integrity constraints* or rules of engagement that hold?
 - · Cardinality? Participation?
 - A database schema in the ER Model can be represented pictorially (ER diagrams)
- Logical design
 - Can map an ER diagram into a relational schema (DBMS data model)
- Physical design
 - File types, indexes, disk layout

ER Diagrams cheat sheet

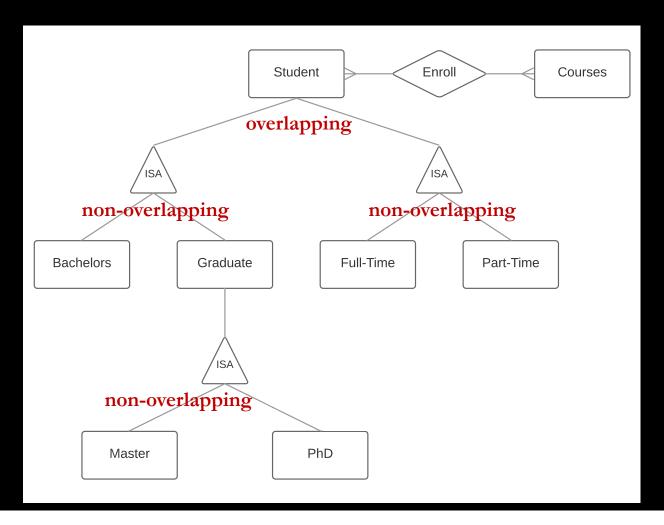


Extended ERDs: ISA ("is a") Hierarchies

- As in C++, or other programming languages, attributes are inherited
- If we declare A ISA B, every A entity is also considered to be a B entity
- Covering vs. overlapping constraints
 - Covering constraints: does every employee have to be either salaried or hourly?
 - Overlapping constraints: can Joe be a salaried employee as well as an hourly employee?

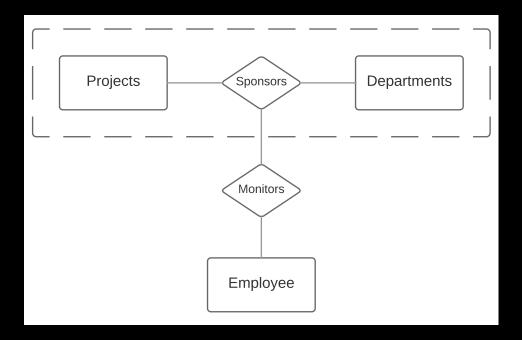


Complex ISA hierarchies



Aggregation

- Used to model a relationship involving another relationship
- Allows us to *treat a relationship as an entity* for purposes of participation in (other) relationships

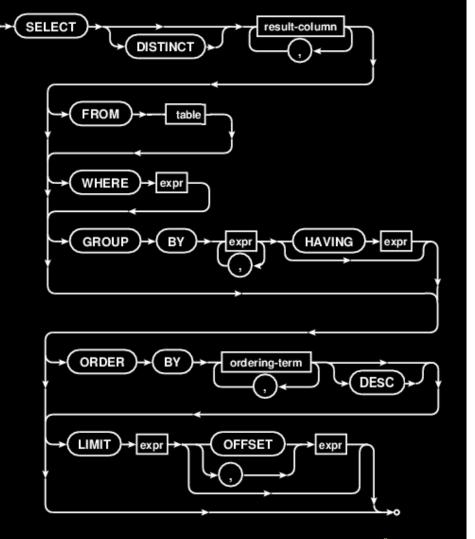


Part 2: SELECT Query Steps

SQLite SELECT Syntax

For example:

SELECT FirstName, LastName
FROM customers
WHERE City = "Paris";



SELECT queries are series of filtering and manipulation steps

- 1. The FROM expression gives the starting point a full table The final result will be a subset or aggregation of this
- 2. The WHERE expression keeps only those rows passing some test
 This expression can be very complex, but it must be something than can be
 evaluated on each row, one at a time
- 3. GROUP BY combines rows if something about them is the same
- 4. The SELECT result-columns are computed, including aggregation At this point we have thrown out the columns we don't need
- 5. HAVING expression keeps only the aggregated rows passing a test
- 6. ORDER BY sorts what's left
- 7. LIMIT truncates the results to just a certain number of rows

SELECT steps (abbreviated)

- 1. FROM chooses the table of interest
- 2. WHERE throws out irrelevant rows
- 3. GROUP BY identifies rows to combine
- 4. SELECT tells what values to return (allowing math and aggregation)
- 5. HAVING throws out irrelevant rows (after aggregation)
- 6. ORDER BY sorts
- 7. LIMIT throws out rows based on their position in the results

Each step gets closer to the specific result you want

- 1. FROM chooses the table of interest
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Products table has the price info, so we start there:



ProductNumber	ProductName	ProductDescription	RetailPrice	QuantityOnHand	CategoryID
Filter	Filter	Filter	Filter	Filter	Filter
1	Trek 9000 Mountain Bike	NULL	1200	6	2
2	Eagle FS-3 Mountain Bike	NULL	1800	8	2
3	Dog Ear Cyclecomputer	NULL	75	20	1
4	Victoria Pro All Weather Tires	NULL	54.95	20	4
5	Dog Ear Helmet Mount Mirrors	NULL	7.45	12	1
6	Viscount Mountain Bike	NULL	635	5	2
7	Viscount C-500 Wireless Bike Computer	NULL	49	30	1
8	Kryptonite Advanced 2000 U-Lock	NULL	50	20	1
9	Nikoma Lok-Tight U-Lock	NULL	33	12	1
10	Viscount Microshell Halmet	KII II I	36	20	1

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We only need the bike rack products, so we filter on CategoryID = 5

SELECT *
 FROM Products
 WHERE CategoryID = 5;

ProductNumber	ProductName	ProductDescription	RetailPrice	QuantityOnHand	CategoryID
39	Road Warrior Hitch Pack	NULL	175	6	5
40	Ultimate Export 2G Car Rack	NULL	180	8	5

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A GROUP BY statement is not needed because we will group all of the rows together

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Grouping

The GROUP BY clause combines multiple rows and lets you perform aggregation math functions

```
SELECT AlbumId,
SUM(Milliseconds/1000/60) AS AlbumMinutes
FROM tracks GROUP BY AlbumId ORDER BY AlbumMinutes;
```

Result:

AlbumId	AlbumMinutes
340	0.86300000
345	1.11065000
318	1.68821667

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We want the RetailPrice column, and we want to aggregate all the rows with the average function

SELECT AVG(RetailPrice)
 FROM Products
 WHERE CategoryID = 5;

AVG(RetailPrice)

177.5

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We want all products, so no need to filter

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 FROM Products;

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- SELECT tells what values to return (allowing math and aggregation)
- HAVING throws out irrelevant rows (after aggregation)
- ORDER BY sorts
- LIMIT throws out rows based on

The GROUP BY statement groups together all rows of the same product category
Note: syntax below only for illustration

SELECT FROM Products GROUP BY CategoryID;

ProductNumber	ProductName	ProductName F		ProductDescription Re		ce Qu	antityOnHand	CategoryID
3	Dog Ear Cyclecomputer	mputer		NULL		20		1
5	Dog Ear Helmet Mount Mirrors		NULL		7.45	12		1
• • •								
1	Trek 9000 Mountain Bike	NULL		1200)	6		2
2	Eagle FS-3 Mountain Bike	NULL		1800	0	8		2
•••								
23	Ultra-Pro Rain Jacket	NU	'LL	85	5	30		3
24	StaDry Cycling Pants	NU	'LL	69)	22	19	3

- 1. FROM chooses the table of interest
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We want the CategoryID and RetailPrice columns, and we want to aggregate the rows in each group with the average function

SELECT CategoryID, AVG(RetailPrice) FROM Products GROUP BY CategoryID;

CategoryID	AVG(RetailPrice)
1	66.1916666666667
2	1321.25
3	51.25
4	79.76555555556
5	177.5
6	29.0