MLDS-413 Introduction to Databases and Information Retrieval

Lecture 15
Indexing Databases

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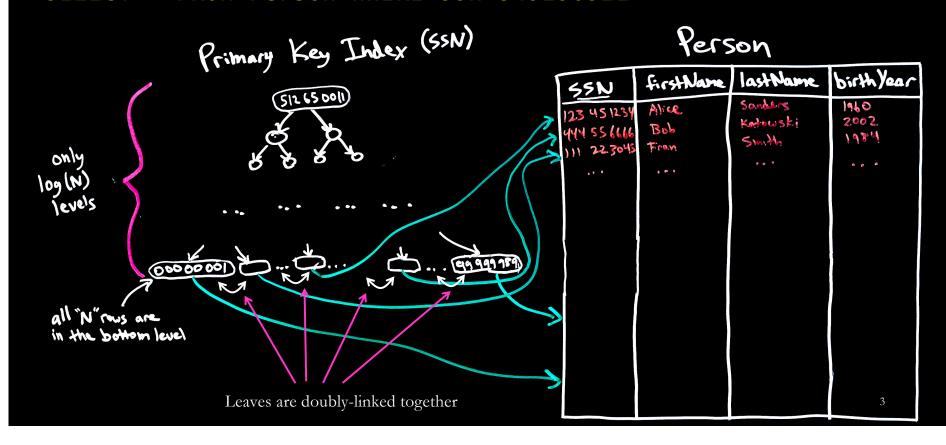
Slides adapted from Steve Tarzia

Last Lecture

- Computer memory can be thought of as one big array
 - Expensive to move large chunks of data
 - Let data remain permanently in a memory address, refer to it with the address number
- A Tree (or graph in general) can be stored as a list of nodes referring to other nodes by memory address
- Trees can serve as an *index* to find data quickly in a database
 - Logarithmic #steps, tree can be memory-resident > fast access, no need to sort data
- Insertions and deletions in a tree are fast as well
 - Most of the data remain in place; we just change a few references
- Trees need to implement *balancing rules* during insertions and deletions
 - e.g., B+-trees are self-balancing (will not cover self-balancing trees in this class)

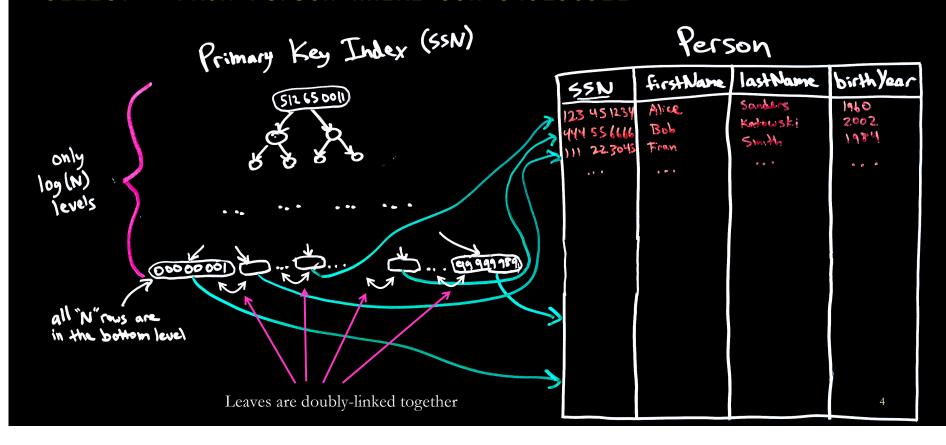
Review of how an index finds rows (equality, aka probe)

SELECT * FROM Person WHERE SSN=543230921



Review of how an index finds rows (range query)

SELECT * FROM Person WHERE SSN>543230921



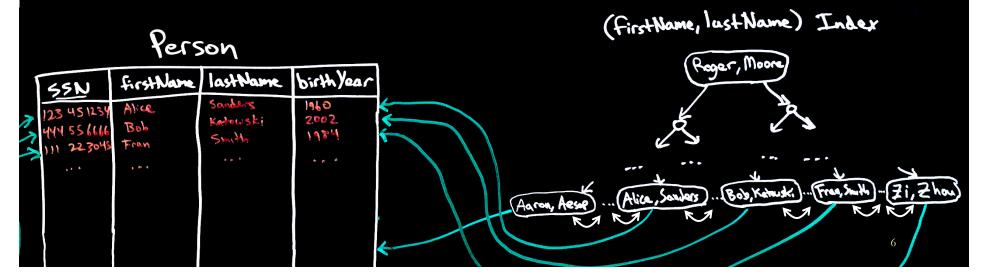
Multiple indexes allow finding rows quickly based on multiple criteria

- Need two indexes to quickly get results for both:
 - SELECT * FROM Person WHERE SSN=543230921
 - SELECT * FROM Person WHERE birthYear BETWEEN 1979 AND 1983



Composite indexes

- Useful when WHERE clauses involve pairs of column values
- SELECT * FROM Person WHERE firstName="Roger" and lastName="Moore"
- Index sorted in lexicographic order of composite keys, e.g., index on <firstName, lastName>
- Requires less space than two separate indexes
- Can find the matching pair of values with one lookup
- The index above is fast if you search for firstName, or "firstName, lastName"
- ...but slow for lastName alone



B+ Trees in practice (cool facts!)

- A typical B+ tree node has on average 134 children
 - Typical order: 100 (i.e., $100 \le \text{keys per node} \le 200$)
 - Typical fill-factor: 67%
 - $2 \times 100 \times 0.67 = 134$
- Index trees rarely have more than 4 or 5 levels:
 - Height 3: $134^3 = 2,406,104$ entries
 - Height 4: $134^4 = 322,417,936$ entries
 - Height 5: $134^5 = 44,840,334,375$ entries
- Top levels can always be in memory:
 - Level 1 = 1 page = 8 KB
 - Level 2 = 134 pages = 1 MB
 - Level 3 = 17,956 pages = 140 MB
 - Level 4 = 2,406,104 pages = 18 GB (OK for a server)
 - Level 5 = 322,417,936 pages = 2.4 TB (probably not in memory)

Key and Index terminology in SQL

- Plain "key" is just a field (or combination of) we can use to find rows quickly
- An "index" on that "key" is just a data structure we can use to find rows quickly
 - Just create a search tree (B+ tree) on that key
- "Unique key" is a key that prevents duplicates
 - Bottom level of index (B+ tree) on the unique key has no repeated values
 - DBMS can use the tree to quickly search for existing rows with that value before allowing a row insertion (or column update) to proceed
- "Primary key" is just a unique key, but there can only be one per table
 - We think of the primary key as the most important unique key in the table
- "Foreign key" makes a column's values match a column in another table
 - Usually it's the primary key in the other table

When to index columns?

- Generally, add an index if the column is:
 - Used in WHERE conditions, or
 - Used in JOIN ... ON conditions, or
 - A foreign key refers to it
- Also helpful if the column is:
 - In a MIN or MAX aggregation function

(examples follow)

```
SELECT *
FROM employees AS E
WHERE (E.salary > 150000 and E.age=45) OR E.age=30
```

- 1. Index on <age>
- 2. Index on <salary>
- 3. Index on <age, salary>
- 4. Index on <salary, age>
- 5. None
- 6. Something else

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FROM employees AS E
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- 1. Index on <age>
- 2. Index on <salary>
- 3. Index on <age, salary>
- 4. Index on <salary, age>
- 5. None
- 6. Something else

```
SELECT *
FROM employees AS E
        JOIN departments AS D
        ON E.deptID = D.deptID
WHERE (E.salary > 150000 and E.age=45) OR E.age=30
```

- 1. Index E on <age, salary>
- 2. Index E on <age, salary> and on <deptID> (i.e., 2 indexes on table E)
- 3. Index E on <age, salary, deptID>
- 4. Index D on <deptID>
- 5. No index on D

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SELECT *
FROM employees AS E
     JOIN departments AS D
     ON E.deptID = D.deptID
WHERE (E.salary > 150000 and E.age=45) OR E.age=30
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- 1. Index E on <age, salary>
- 2. Index E on <age, salary> and on <deptID> (i.e., 2 indexes on table E)
- 3. Index E on <age, salary, deptID>
- 4. Index D on <deptID>
- 5. No index on D

```
SELECT MAX(E.salary)
FROM employees AS E
WHERE E.age=30
```

- 1. Index on <age>
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- 4. Index on <salary, age>
- 5. Index on <age> and on <salary>
- 6. None
- 7. Something else

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SELECT MAX(E.salary)
FROM employees AS E
WHERE E.age=30
```

- 1. Index on <age>
- 2. Index on <salary>
- 3. Index on <age, salary>
- 4. Index on <salary, age>
- 5. Index on <age> and on <salary>
- 6. None
- 7. Something else

Creating indexes

- Indexes are usually defined when the table is created
 - Usually define at least a primary key
- But you may later realize that certain queries are too slow
 - Without proper indexes, DBMS will have to examine every row in the table to find the relevant rows
 - Adding one or more indexes may dramatically speed up a query

Basic syntax:

CREATE INDEX index_name ON table_name (column_name)