Intro to Relational Database Design

This PDF is completed and recording is watched

- The 3 rules of database design
- Example: MovieLens
- Indexing to improve performance
- Stored procedures
- Transactions



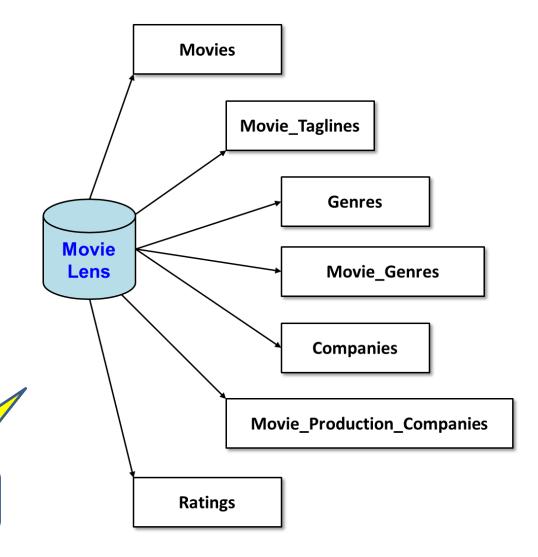
Database design

- Database design is an art...
- But there are 3 fundamental rules to always consider

MovieLens database

MovieLens

- <u>https://movielens.org/</u>
- -45K movies
- 26M reviews



Three fundamental rules of DB design

Here, the table breaks rule 1, as it is storing genres, movie multiple times. We need to store data once. Therefore, break down the table into multiple table like movies, genres which will have them once.

1. Store data exactly once

Assign a unique identifier called primary ID to each row in a table

2. Every row has a unique key to identify that row

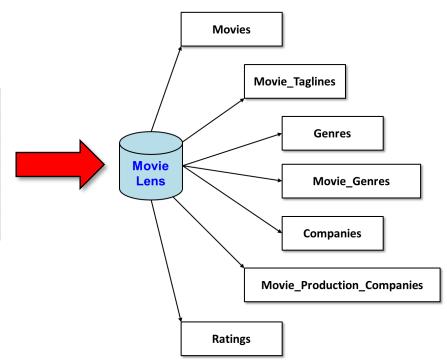
3. Store one value per cell

In Genres column, more than 1 value is stored. Its better to have 1 in each row. Eg Movie Genres -> This Table holds Genres for each movie.

- Id Movie Genre
- 1 Matrix Action
- 2. Matrix Science
- 3. Matrix Comedy

Original data (csv file)

Movie	Genres	Year	Revenue	Rating
The Matrix	Action; Science Fiction	1999	463517383	9
Toy Story	Animation; Comedy	1995	373554033	8
The Matrix	Action; Science Fiction	1999	463517383	10
The Matrix	Action; Science Fiction	1999	463517383	6



Movie_ID	Title	Release_Date	Runtime	Original_L anguage	Budget	Revenue
603	The Matrix	1999-03-30 00:00:00.000	136	en	63000000	463517383
862	Toy Story	1995-10-30 00:00:00.000	81	en	30000000	373554033
	•••					
<i>†</i>						

<u>Movies</u>	
	`
Movie	,
Lens	
\	\

Genres

Movie_ID	Rating
605	8
603	6
605	10
605	6
•••	

Movie General Movie_ID 862

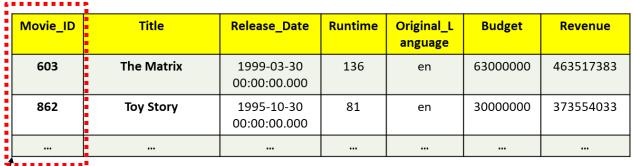
Genre_ID	Genre_Name		
28 Action			
878 Science Fictio			
16 Animation			
35	Comedy		
•••			

Ratings

Movie_ID	Genre_ID	
862	16	
862	35	
603	28	
603	878	
•••		

Primary key --- unique row identifier. Every table should have a PK.

Primary Key



Movie Genres

Genres

Genre_ID	Genre_Name			
28	Action			
878	Science Fiction			
16	Animation			
35	Comedy			

 Movie_ID
 Genre_ID

 862
 16

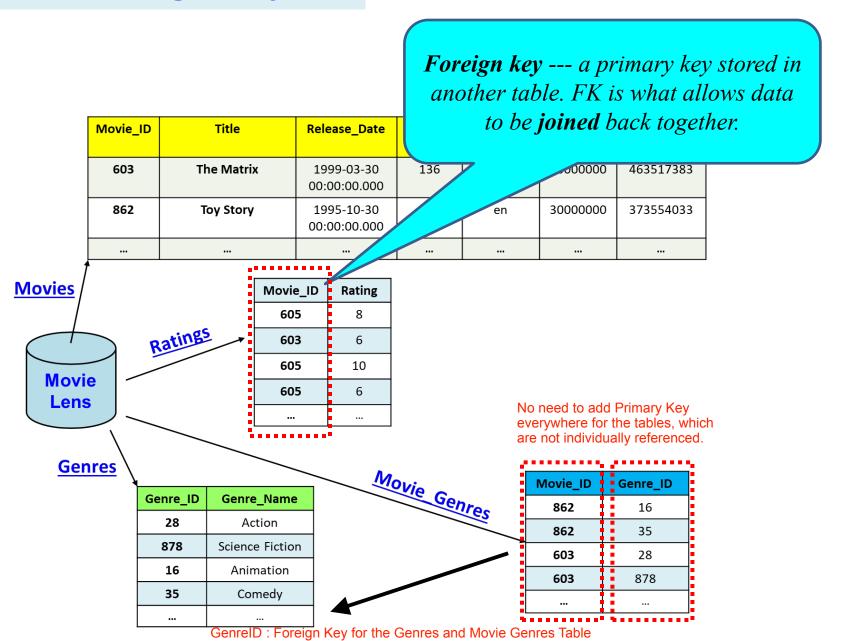
 862
 35

 603
 28

 603
 878

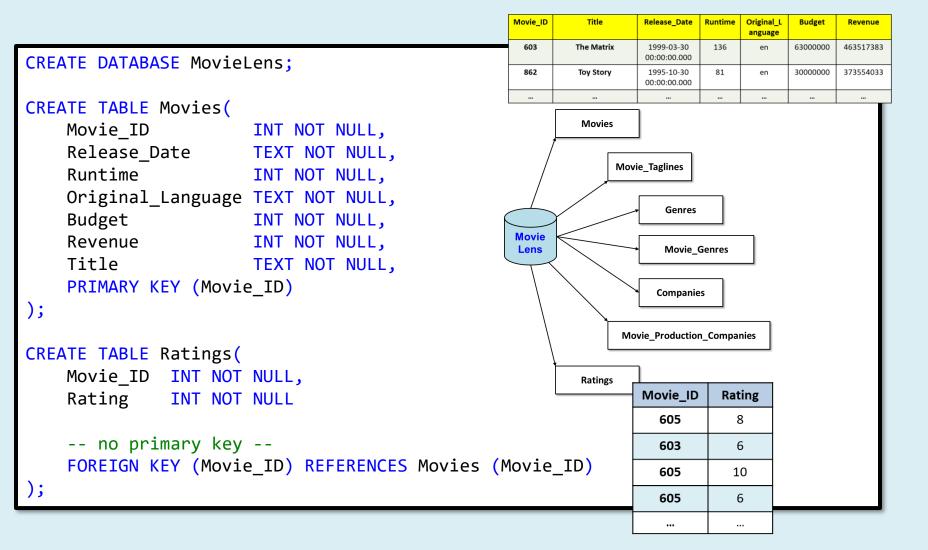
 ...
 ...

Foreign Key



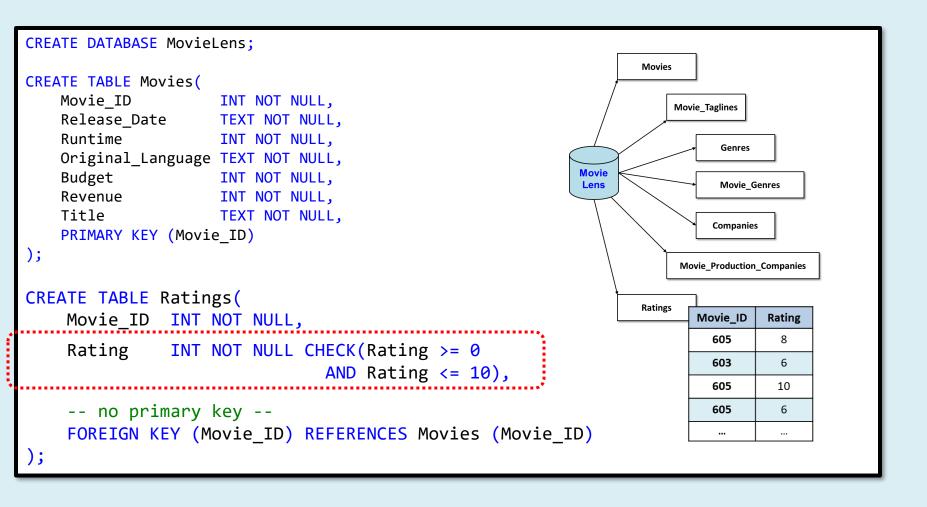
Creating MovieLens using DDL

Databases are created using SQL...



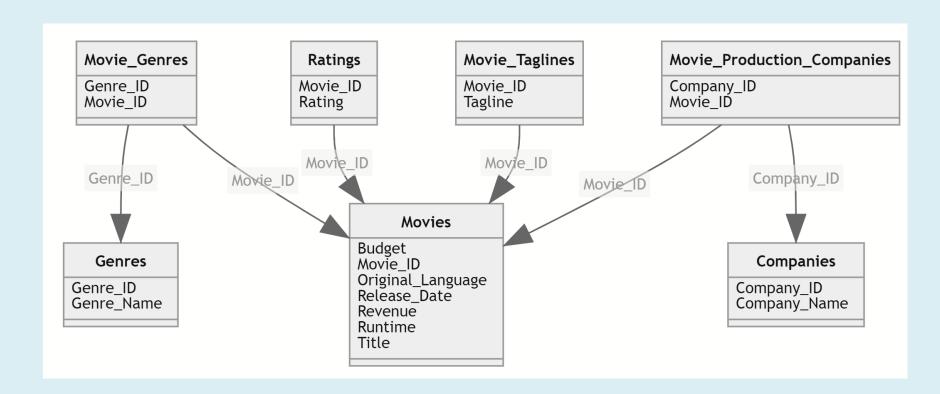
Constraints

Databases allow constraints on data – use them!



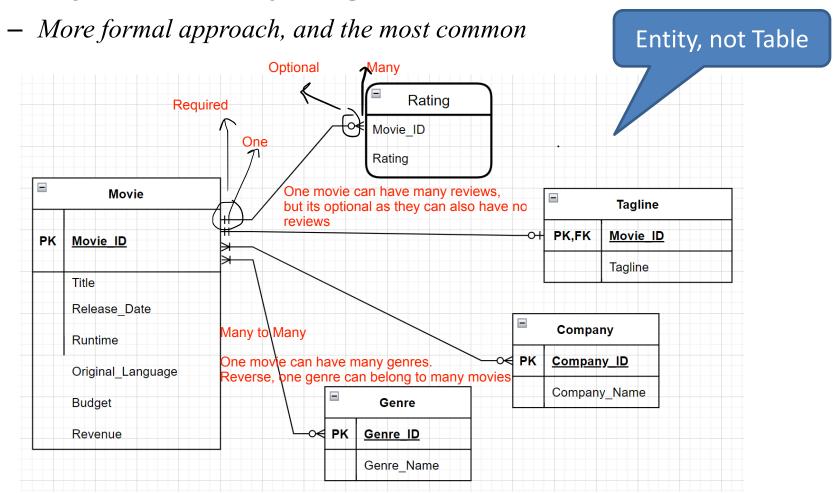
Database schema (design)

- One view is the SQL used to create the database
- Another view is some sort of visual diagram...



ER Diagram

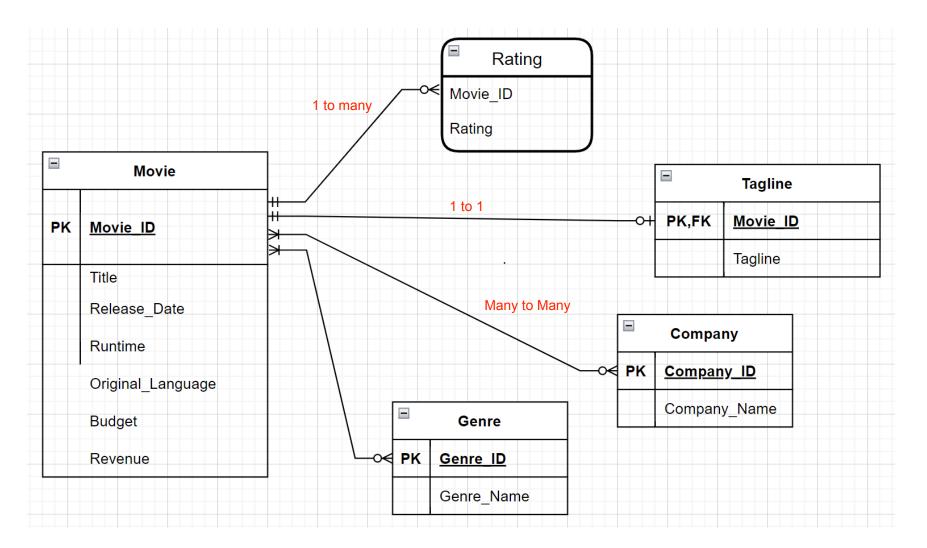
Entity Relationship diagram...



Drawn manually using free tool: https://app.diagrams.net/

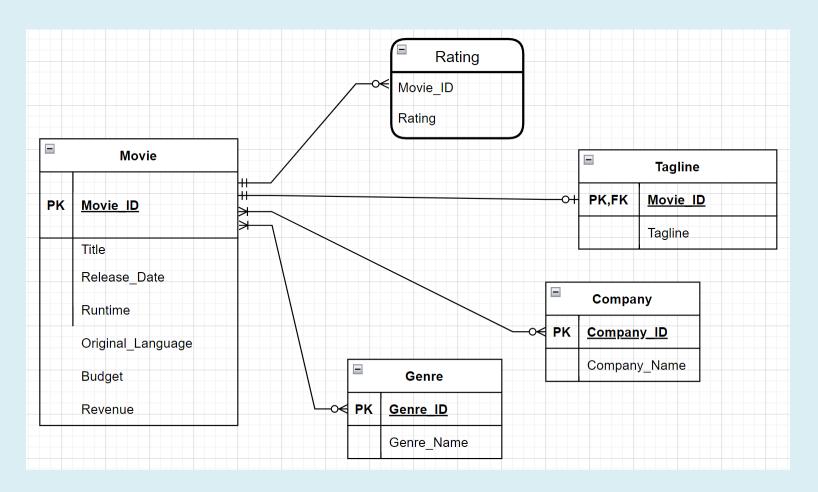
Types of relationships

- **–** 1-to-1
- *1-to-many*
- many-to-many



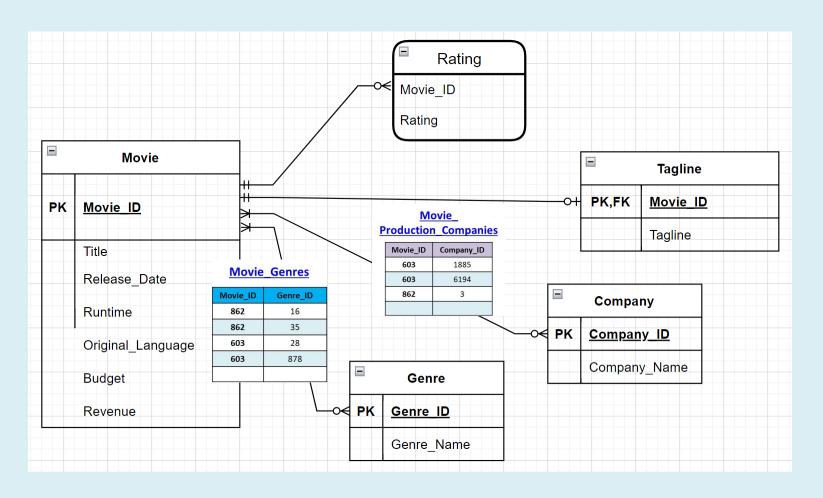
Question

Interestingly, what low-level detail is *missing*
 from the ER diagram?



Answer

 Many-to-many relationships <u>imply</u> a <u>linking</u> or bridge table to implement the connection...



Performance: Scan vs. Seek

- By default, DBMS has to "scan" when searching
 - -Row by row search of all movies in <u>persistent</u> store (file sys)

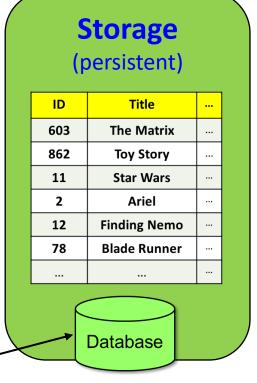
This is usually slow as it scans all rows of the database to fine the answer. To improve the searching, we can use seek, which works indexing.

• Example:

- Searching for a movie by title



DBMS



SELECT *
FROM Movies
WHERE Title = 'Blade Runner';

Query plan

- Query plan is how DBMS plans to execute query
- Most database tools can display plan...

```
ndexing-example....RFACE\joe-h (74))* 😐 🗙
 SELECT avg(cast(rating as float)) As AvgRating
          Movies
  FROM
  INNER JOIN Ratings
           ON Movies.Movie_ID = Ratings.Movie ID
  WHERE Title = 'Toy Story';
■ Results ■ Messages ♣ Execution plan
    AvgRating
     7.19786096256685
```

Indexing

With an index, DBMS can "seek"

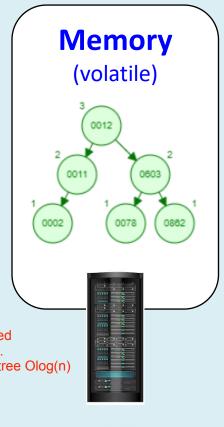
- Search of <u>volatile</u> memory + 1
 access to persistent store
- Primary key always indexed

DBMS builds @ startup

Indexes data residing in persistent storage

Whenever a database is indexed, a binary tree like structure is created in the RAM, and the indexes point to the actual data in the Hard Disk. Whenever we query something, we need to do a search on indexed tree Olog(n) and then 1 access to HDD.

SELECT Title
FROM Movies
WHERE Movie_ID = 78;



DBMS

Storage

(persistent)

row

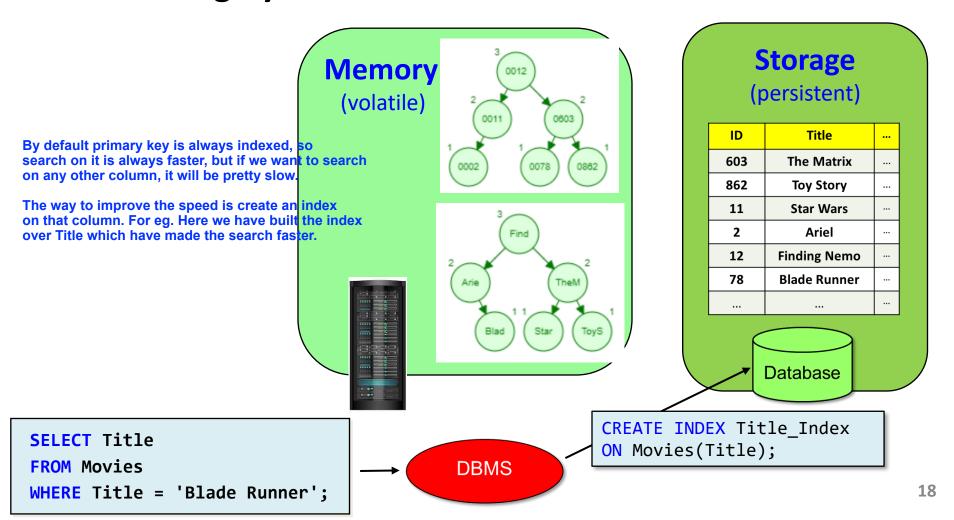
row
row
row
row

File

Database

Example #1

- Searching by movie id is fast (primary key is indexed)
- Searching by movie title is slow --- add index!



Example #2

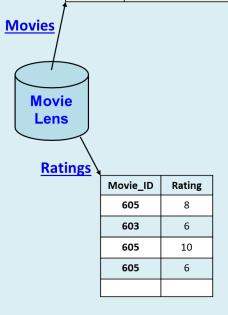
There are 26MB reviews

The problem here is Movie ID from Movies table is indexed so, we are able to find the movies pretty quickly but Ratings table is not indexed, so at the time of joining Movies and Rating, it has to go through each entry in Ratings (scan) and hence is slow

The solution is to index the Movie_ID from Ratings, this will speed up the query.

Computation of avg rating....

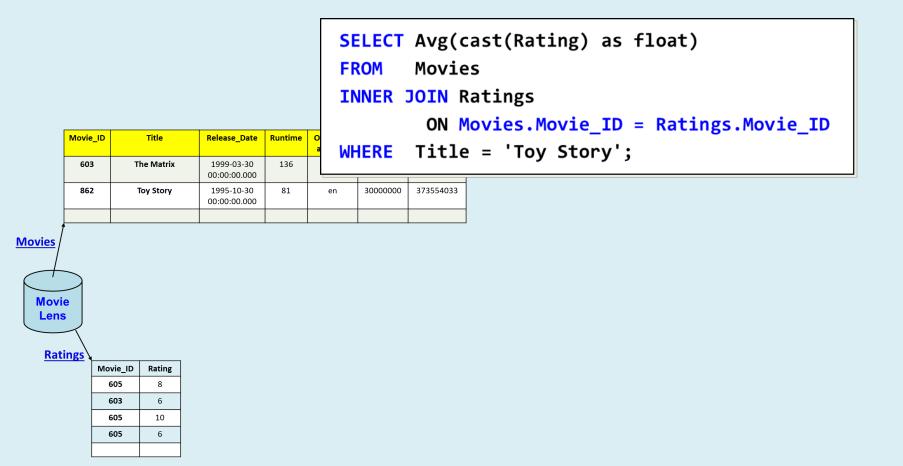
Movie_ID	Title	Release_Date	Runtime	Original_L anguage	Budget	Revenue
603	The Matrix	1999-03-30 00:00:00.000	136	en	63000000	463517383
862	Toy Story	1995-10-30 00:00:00.000	81	en	30000000	373554033



```
SELECT Avg(cast(Rating as float))
FROM Movies
INNER JOIN Ratings
          ON Movies.Movie_ID = Ratings.Movie_ID
WHERE Title = 'Toy Story';
```

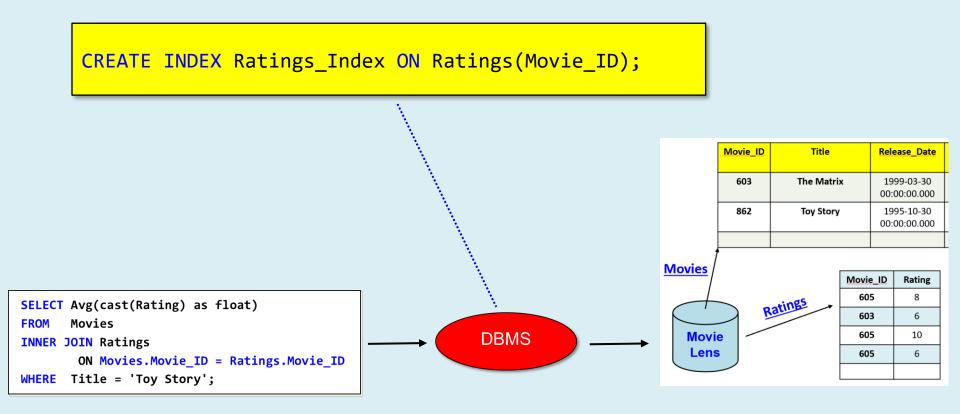
Demo

- Without index, query takes 4-5 seconds...
 - -Not bad, but we can do much better...



Solution: index Ratings by Movie_ID

With correct index, query takes 1/10th of a second

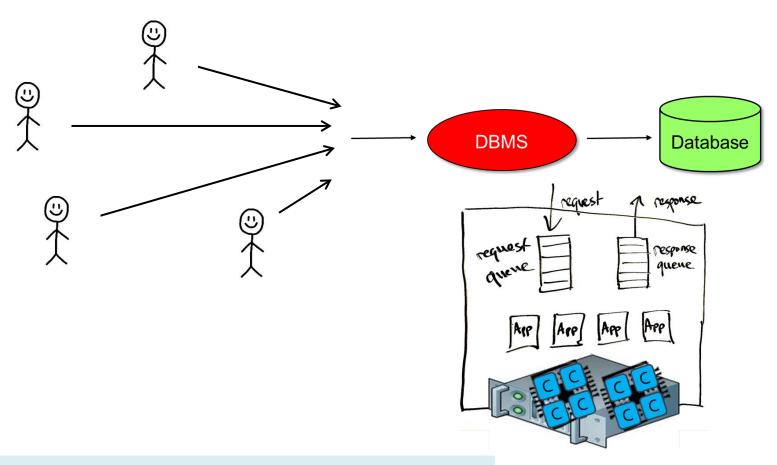


Trade-offs

- Why not create an index on every column?
 - Takes more server memory for underlying data structures
 - Slows down inserts / updates / deletes since have to maintain underyling data structures because these structures are needed to built again with each modification.

- General approach is to create indexes on the obvious columns, then deploy and monitor
 - See how DB is used, and add more indexes as needed

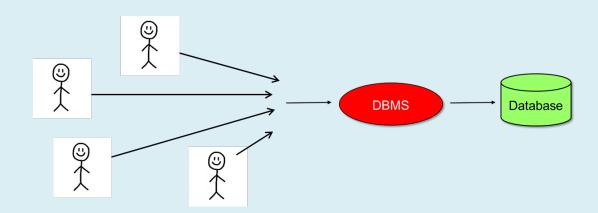
Databases are powerful systems



- Can handle thousands of concurrent users
- Handle simultaneous insert / update / delete
- Take advantage of multi-core hardware

Question

- We are building a ticket selling app
 - everyone is trying to buy 2 tickets



Here, since all users will access the Database at the same time and try to perfrom the same action, it can be unsafe and cause unwanted results. The query will make the seats go from 2 -> 0 -> -2 -> -4 etc...

This is unsafe in a multi-user scenario. How to fix?

```
sql = 'Select num_avail From Tickets Where eventid=1234;'
dbCursor.execute(sql);
row = dbCursor.fetchone()
numAvail = int(row[0])

if numAvail >= 2:
    success = True
    sql = "Update Tickets Set num_avail = num_avail-2 Where eventid=1234;'
    dbCursor.execute(sql)
else:
    success = False
```

Solution: Transactions

- Transaction isolate operations from other users
- Transactions make a series of operation atomic

```
BEGIN TRANSACTION;
  -- withdraw from checking
 UPDATE Accounts
         Balance = Balance - 100.00
  SET
  WHERE Account = 22197;
  -- deposit into savings
 UPDATE Accounts
         Balance = Balance + 100.00
  SET
 WHERE Account = 43992;
COMMIT;
```

Take this example, if power goes off after withdraw, then the data is lost, no way to resume and perform deposit.

Therefore, wrap the whole thing under one transaction object. Until the transaction is committed the database will not be permanently updated.

If anything goes wrong during the transaction, any temporary changes will be reverted back.

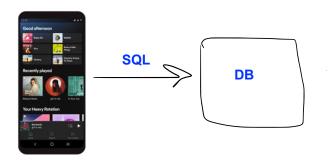
Beware

- For better performance, transactions default to "not-quite-perfect" isolation between users
- Solutions
 - (1) isolation level = Serializable
 - (2) default isolation level + "Select ... For Update;"

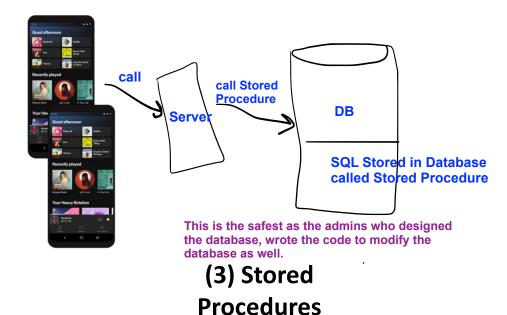
```
BEGIN TRANSACTION
sql = 'Select num_avail From Tickets Where eventid=1234 For Update;'
dbCursor.execute(sql);
row = dbCursor.fetchone()
numAvail = int(row[0])

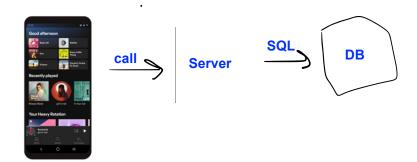
if numAvail >= 2:
    success = True
    sql = "Update Tickets Set num_avail = num_avail-2 Where eventid=1234;'
    dbCursor.execute(sql)
    COMMIT
else:
    success = False
    ROLLBACK
    This level of understanding is why you want
    DBAs writing stored procedures
```

Design: who executes the SQL?



(1) Client (project 01)





(2) Server (project 02)

Stored Procedures

Functions stored in the database

- Support parameters and nearly all features of SQL
- Can retrieve, insert, update, and delete data

```
EXEC AddMovieReview 603, 10;
  inserts a review into the database
Create Procedure AddMovieReview
 @id int,
 @rating tinyint
Begin
  BEGIN TRANSACTION;
  Insert Into Ratings(Movie_ID, Rating) Values(@id, @rating);
  COMMIT;
  RETURN @@ROWCOUNT; -- # of rows modified
End;
```

Top N movies with Min ratings

EXEC TopN 3, 500;



Movie_ID 🗸	Title	NumRatings 🗸	AvgRating 🗸
318	The Million Dollar Hotel	768	8.87890625
527	Once Were Warriors	559	8.554561717352415
296	Terminator 3: Rise of the Ma	755	8.512582781456954

```
-- retrieves top N movies with at least Min ratings
Create Procedure TopN
 @N int, -- retrieve N movies
 @Min int -- must have >= Min ratings
As
Begin
  Select Top (@N)
         Movies.Movie ID,
         Title,
         Count(Rating) As NumRatings,
         Avg(Cast(Rating as real)) As AvgRating
  From Movies
  Inner Join Ratings on Movies.Movie_ID = Ratings.Movie_ID
  Group By Movies. Movie ID, Movies. Title
  Having Count(Rating) >= @Min
 Order By AvgRating DESC, Title ASC;
End;
```

Why stored procedures?

- Faster
- More trustworthy!

- Typically written by database designer(s)
 - Experts in SQL
 - Understand details of database design

That's it, thank you!