

## Tutorial - 3

Q1: Write linear search pseudocode to search an element in a sorted array with minimum comparisons.

→ while ( $\text{low} \leq \text{high}$ )

{  $\text{mid} = (\text{low} + \text{high}) / 2;$

  'if ( $\text{arr}[\text{mid}] == \text{key}$ )

    return True;

  else if ( $\text{arr}[\text{mid}] > \text{key}$ )

$\text{high} = \text{mid} - 1;$

  else

$\text{low} = \text{mid} + 1;$

}

return False;

Q2: Write pseudo code for iterative and recursive insertion sort. Insertion sort is called online sorting. Why? What about other sorting algorithms that has been discussed.

→ Iterative insertion sort → for ( $\text{int } i=1; i < n; i++$ )

{  $j = i - 1;$

$x = A[j];$

  while ( $j > -1 \text{ and } A[j] > n$ )

$A[j+1] = A[j];$

$j--;$

$A[j+1] = n;$

3

Recursive insertion sort → void insertion sort (int arr), int n)

```

2 if (n <= 1)
    return;
insertionsort (arr, n-1);
int last = arr[n-1];
i = n-2;
while (j >= 0 && arr[j] > last)
    arr[j+1] = arr[j];
    i--;
3 arr[i+1] = last;

```

Insertion sort is online sorting because whenever a new element come, insertion sort define its right place.

Q3: complexity of all sorting algorithms .

- Bubblesort -  $O(n^2)$
- Insertionsort -  $O(n^2)$
- Selection sort -  $O(n^2)$
- Merge sort -  $O(n * \log n)$
- Quick sort -  $O(n \log n)$
- Count sort -  $O(n)$
- Bucket sort -  $O(n)$ .

Q4: Divide all sorting algorithm into in-place / stable / online sorting .

- Online sorting → Insertion sort
- stable sorting → Merge sort , insertion sort , Bubble sort .

Inplace sorting  $\rightarrow$  Bubble sort, Insertion sort, Selection sort.

Q. S. write recursive / iterative pseudo code for binary search. what is time & space complexity of linear & binary search.

$\rightarrow$  Iterative binary search - complexity -  $O(\log n)$

```
while (low <= high)
    int mid = (low + high) / 2;
    if (a[mid] == key)
        return true;
    else if (a[mid] > key)
        high = mid - 1;
    else
        low = mid + 1;
```

Recursive binary search - complexity -  $O(\log n)$

```
while (low <= high)
    int mid = (low + high) / 2;
    if (a[mid] == key)
        return true;
    else if (a[mid] > key)
        Binarysearch(a, low, mid - 1);
    else
        Binarysearch(a, mid + 1, high);
3
return false;
```

Q. b. Write recursive recurrence relation for binary recursive search.

$$\rightarrow T(n) = T(n/2) + T(n/2) + C,$$

Q7: Find two indexes such that  $A[i] + A[j] = K$  in minimum time complexity.

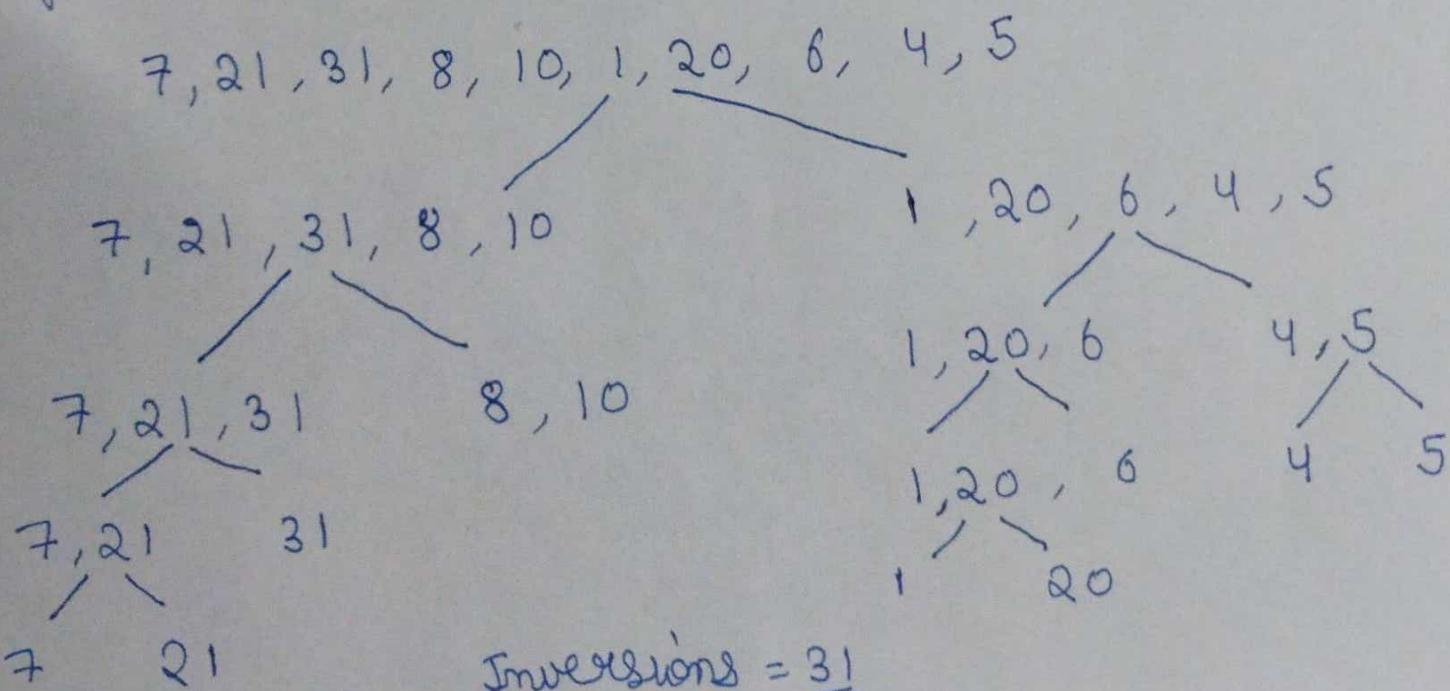
```
map<int, int> m;
for (int i = 0; i < arr.size(); i++) {
    if (m.find(target - arr[i]) != m.end())
        m[arr[i]] = i;
    else
        cout << i << " " << m[target - arr[i]];
}
```

Q8: Which sorting is best for practical use?

Quicksort is fastest general purpose sort. In most practical situation quicksort is method of choice. If stability is important & space is available, merge sort will be best.

Q9: What do you mean by number of inversion in an array? Count the no. of inversion in  $arr[] = \{7, 21, 31, 8, 10, 1, 20, 6, 4, 5\}$  using merge sort.

Inversion indicates how far as close the array is from being sorted.



Q10: In which cases Quick sort will give the best & worst case time complexity?

Worst case - when the pivot is always an extreme (smallest or largest). This happens when input array is sorted in reverse order.

Best case - when pivot element is as near to middle element.  $\Theta(n \log n)$ .

Q11: Write recurrence relation of merge & quick sort in best & worst case? What are similarities & differences between complexities.

$$\rightarrow \text{Merge sort} \rightarrow T(n) = 2T\left(\frac{n}{2}\right) + \Theta(n)$$

$$\text{Quick sort} \rightarrow T(n) = 2T(n/2) + n + 1$$

### Quick sort

- splitting is done in any ratio.
- smaller array is suitable.
- inefficient for large array.
- internal sorting
- Not stable

### Merge sort

- array is just partitioned into 2 halves.
- suitable for any size of array.
- more efficient.
- external sorting
- stable