

## Tutorial - 3

Q1: write linear search pseudocode to search an element in a sorted array with minimum comparisons.

```
→ while (low <= high)
{
    mid = (low + high) / 2;
    if (arr[mid] == key)
        return True;
    else if (arr[mid] > key)
        high = mid - 1;
    else
        low = mid + 1;
}
return False;
```

Q2: write pseudo code for iterative and recursive insertion sort. Insertion sort is called online sorting, why? what about other sorting algorithms that has been discussed.

```
→ Iterative insertion sort - for (int i = 1; i < n; i++)
{
    j = i - 1;
    x = A[j];
    while (j > -1 && A[j] > x)
    {
        A[j+1] = A[j];
        j--;
    }
    A[j+1] = x;
}
```

```

Recursive insertion sort - void insertionSort(int arr[], int n)
{
    if (n <= 1)
        return;
    insertionSort(arr, n-1);
    int last = arr[n-1];
    i = n-2;
    while (i >= 0 && arr[i] > last)
    {
        arr[i+1] = arr[i];
        i--;
    }
    arr[i+1] = last;
}

```

Insertion sort is online sorting because whenever a new element come, insertion sort define its right place.

Q3: complexity of all sorting algorithms.

→ Bubble sort -  $O(n^2)$

Insertion sort -  $O(n^2)$

Selection sort -  $O(n^2)$

Merge sort -  $O(n \cdot \log n)$

Quick sort -  $O(n \log n)$

Count sort -  $O(n)$

Bucket sort -  $O(n)$

Q4: Divide all sorting algorithm into inplace / stable / online sorting.

→ Online sorting → Insertion sort

Stable sorting → Merge sort, Insertion sort, Bubble sort.



Inplace sorting  $\rightarrow$  Bubble sort, Insertion sort, Selection sort.

Q.5. Write recursive / iterative pseudo code for binary search. What is time & space complexity of linear & binary search?

$\rightarrow$  Iterative binary search - complexity  $- O(\log n)$

```
while (low <= high)
{
    int mid = (low + high) / 2;
    if (a[mid] == key)
        return true;
    else if (a[mid] > key)
        high = mid - 1;
    else
        low = mid + 1;
}
```

Recursive binary search - complexity  $- O(\log n)$

```
while (low <= high)
{
    int mid = (low + high) / 2;
    if (a[mid] == key)
        return true;
    else if (a[mid] > key)
        binarysearch(a, low, mid - 1);
    else
        binarysearch(a, mid + 1, high);
}
return false;
```

Q.6. Write ~~recursive~~ recurrence relation for binary recursive search.

$$\rightarrow T(n) = T(n/2) + T(n/2) + C,$$

Q7. Find two indexes such that  $A[i] + A[j] = K$  in minimum time complexity.

```
map<int, int> m;
```

```
for (int i = 0; i < arr.size(); i++)
```

```
{ if (m.find(target - arr[i]) != m.end())
```

```
{ m[arr[i]] = i;
```

```
else
```

```
{ cout << i << " " << m[arr[i]] << j << endl; }
```

```
}
```

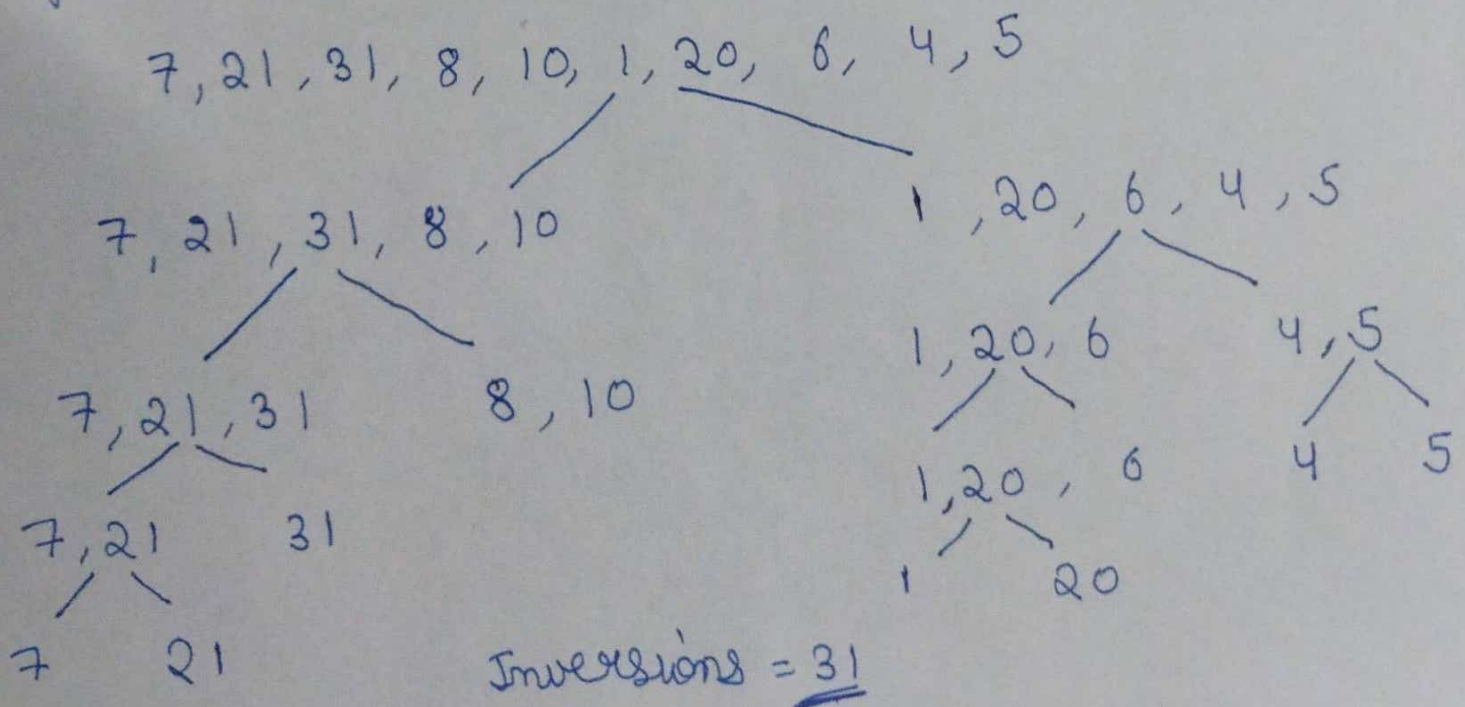
Q8. which sorting is best for practical use?

Quicksort is fastest general purpose sort, in most practical situation quicksort is method of choice. If stability is important & space is available, merge sort will be best.

Q9. what do you mean by number of inversion in an array? Count the no. of inversion in arr[] =

{ 7, 21, 31, 8, 10, 1, 20, 6, 4, 5 } using merge sort.

Inversion indicates how far as close the array is from being sorted.





Q10. In which cases Quick sort will give the best & worst case time complexity?

Worst case → when the pivot is always an extreme (smallest or largest). This happens when input array is sorted in reverse order.

Best case → when pivot element is as near to middle element.  $O(n \log n)$ .

Q11. Write recurrence relation of merge & quick sort in best & worst case? what are similarities & differences between complexities.

→ Merge sort →  $T(n) = 2T(\frac{n}{2}) + O(n)$

Quick sort →  $T(n) = 2T(n/2) + n + 1$

### Quick sort

- splitting is done in any ratio.
- smaller array is suitable.
- Inefficient for large array.
- Internal sorting
- Not stable

### Merge sort

- Array is just parted into 2 halves.
- suitable for any size of array.
- more efficient.
- external sorting.
- stable