

OS

Page
2/18/25

Assignment - 04.

- * Problem Statement : Write a program to stimulate the shortest remaining time first (SRTF) CPU scheduling algorithm in C.

- * Theory :

- I] Write SRTF Characteristics

→ It is a preemptive CPU scheduling algorithm where the process with the least remaining burst time is executed first if a new process arrives with a shorter remaining time, it preempts the current one. It minimizes average waiting time, favour short process, but may cause starvation for longer ones. Suitable for time - Critical tasks.

eg ⇒ Process	Arrival Time (AT)	Burst Time (BT)
P1	0	8
P2	1	4
P3	2	9
P4	3	5

t = 0 P1 runs (remaining = 7)

t = 1 P2 arrived (BT = 4 < P1's 7)

t = 1 → 5 P2 finishes at t = 5.

t = 5 P4 (BT = 5) vs P1(7) vs P3(9) → choose P4

t = 5 → 10 P4 finishes

t = 10 → 17 P1 finishes

t = 17 → 26 P3 finishes

Process	Completion time (CT)	TAT = CT - AT
P ₁	17	17
P ₂	5	4
P ₃	26	24
P ₄	10	7
		<u>52</u>

P ₁	P ₂	P ₄	P ₁	P ₃
0	1	5	10	17

$$\text{Average TAT} = \frac{52}{4} = 13$$

$$\text{Average CT} = \frac{26}{4} = 6.5$$

Eg:

Process	AT	BT	CT	TAT = CT - AT	WT = CT - AT
P ₁	0	6	9	9	9
P ₂	2	2	4	2	0
P ₃	4	1	5	1	0
P ₄	5	7	12	11	7
				<u>23</u>	

$$t = 0 \rightarrow 2$$

P₁ runs (remaining 4)

$$t = 2$$

P₂ arrives ($BT = 2 < P_1's\ 4$) \Rightarrow preempt

$$t = 2 \rightarrow 4$$

P₂ finished

$$t = 4$$

P₃ arrives

$$t = 4 \rightarrow 5$$

P₃ finished

$$t = 5$$

P₁ (4) vs P₄(7) \rightarrow Run P₁

$$t = 5 \rightarrow 9$$

P₁ finished

$$t = 9 \rightarrow 16$$

P₄ finished

P ₁	P ₂	P ₃	P ₁	P ₄
0	2	4	5	9

$$\text{Avg TAT} = \frac{23}{4} = 5.75$$

$$\text{Avg WT} = \frac{7}{4} = 1.75$$

- * Conclusion \Rightarrow SRTF minimizes average waiting time by prioritizing shorter jobs but risking risks starving longer process.

* FAQ's:

1]. How is SRTF different from shortest job first.

\rightarrow SRTF is a preemptive version of shortest job first (SJF). In SRTF a running process can be interrupted if a new process with a shorter remaining time arrives, while SJF is non-preemptive and runs chosen process till complete completion.

2]. What are the major drawbacks of SRTF scheduling.

\rightarrow

- i) Starvation: longer process may never execute if shorter job keep arriving.
- ii) Frequent context switches: Increases CPU overhead.
- iii) Complex to implement: Needs accurate burst time estimation.
- iv) Overhead in arrival tracking: (constantly checking for the process).

3] what is starvation in SRTF?

\rightarrow In SRTF starvation occurs when long process wait indefinitely because shorter process keep arriving and preempt them. This continuous delay prevents the longer job from getting CPU time, especially in system with frequently short-task arrivals.

CODE:

```
#include <stdio.h>

struct Process {
    int id, at, bt, rt, ct, tat, wt;
};

int main() {
    int n;
    printf("Enter number of processes: ");
    scanf("%d", &n);

    struct Process p[n];

    for (int i = 0; i < n; i++) {
        p[i].id = i + 1;
        printf("Enter arrival time and burst time for process %d: ", p[i].id);
        scanf("%d %d", &p[i].at, &p[i].bt);
        p[i].rt = p[i].bt;
    }

    int complete = 0, time = 0, min_rt, shortest;
    int found;
    float totalWT = 0, totalTAT = 0;

    // for Gantt chart
    int ganttOrder[1000], ganttTime[1000], gcIndex = 0;
    int prev = -1;

    while (complete < n) {
        found = 0;
        min_rt = 1000000000;
        shortest = -1;

        for (int i = 0; i < n; i++) {
            if (p[i].at <= time && p[i].rt > 0 && p[i].rt < min_rt) {
                min_rt = p[i].rt;
                shortest = i;
                found = 1;
            }
        }
```

```

    }

    if (!found) {
        time++;
        continue;
    }

    // record Gantt chart only when process changes
    if (shortest != prev) {
        ganttOrder[gclIndex] = p[shortest].id;
        ganttTime[gclIndex] = time;
        gclIndex++;
        prev = shortest;
    }

    p[shortest].rt--;
    time++;

    if (p[shortest].rt == 0) {
        complete++;
        p[shortest].ct = time;
        p[shortest].tat = p[shortest].ct - p[shortest].at;
        p[shortest].wt = p[shortest].tat - p[shortest].bt;

        totalWT += p[shortest].wt;
        totalTAT += p[shortest].tat;
    }
}

printf("\nID\tAT\tBT\tCT\tTAT\tWT\n");
for (int i = 0; i < n; i++) {
    printf("%d\t%d\t%d\t%d\t%d\t%d\n",
        p[i].id, p[i].at, p[i].bt,
        p[i].ct, p[i].tat, p[i].wt);
}

printf("\nAverage Turnaround Time = %.2f", totalTAT / n);
printf("\nAverage Waiting Time = %.2f\n", totalWT / n);

// Gantt chart
printf("\nGantt Chart:\n");
for (int i = 0; i < gclIndex; i++) {
    printf("P%d ", ganttOrder[i]);
}

```

```

printf("|\n0");
for (int i = 0; i < gclIndex; i++) {
    printf(" %d", ganttTime[i]);
}
printf(" %d\n", time);

return 0;
}

```

OUTPUT:

```

computer@computerVV:~$ gcc assign4os.c
computer@computerVV:~$ ./a.out
Enter number of processes: 4
Enter arrival time and burst time for process 1: 0 8
Enter arrival time and burst time for process 2: 1 4
Enter arrival time and burst time for process 3: 2 9
Enter arrival time and burst time for process 4: 3 5

ID      AT      BT      CT      TAT      WT
1       0       8       17      17      9
2       1       4       5       4       0
3       2       9       26      24      15
4       3       5       10      7       2

Average Turnaround Time = 13.00
Average Waiting Time   = 6.50

Gantt Chart:
| P1 | P2 | P4 | P1 | P3 |
0   0   1   5   10  17  26

```