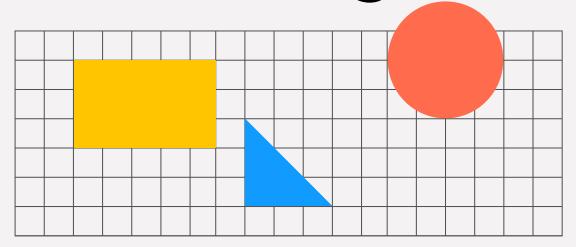
# Extended SQL Processing Tool



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### Standard SQL Setbacks

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- → Grouping and aggregate queries can be difficult to express in standard SQL
  - Can cause conceptual and implementation problems to database system
  - Complex SQL expressions with multiple joins, group by clauses, and subqueries can be difficult to comprehend

- → Why is this important?
  - Aggregation is intensively used in a wide range of applications
  - Crucial for query
    languages to be able to
    succinctly and efficiently
    express these queries

## Solution - MF Queries

#### Standard SQL

SELECT x.product, sum(x.quantity), sum(y.quantity), sum(z.quantity)
FROM Sales x, Sales y, Sales z
WHERE x.product=y.product AND
x.month=1 AND x.year=1997 AND
y.product=z.product AND y.month=2 AND
y.year=1997 AND z.month=3 AND
z.year=1997
GROUP BY x.product



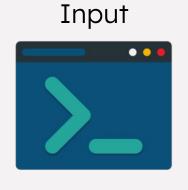
#### **Multiple-Feature Query**

SELECT product,
sum(X.quantity),
sum(Y.quantity),
sum(Z.quantity)
FROM Sales
WHERE year=''1997''
GROUP BY product: X, Y, Z
SUCH THAT X.month = 1,
Y.month = 2,
Z.month = 3

The multiple-feature query extends the GROUP BY clause to include a SUCH THAT clause which is used to define the grouping variables.

## **Project Overview**







### Python Generator



### Output

cust	prod	avg(quant)	max(quant)
Dan	   Ham	571.667	986
Claire	Fish	618.364	1000 j
Chae	Jelly	449.36	964 İ
Claire	Jellý	489.545	975 j
Boo	Ham	589.609	983 j
Emily	Ice	449.406	984 j
Emily	Ham	600.923	978 j
Boo	Dates	485.25	990 j
Mia	Jelly	431.409	940 j
Mia	Fish	457.583	950 j
Helen	Dates	448.955	949
Wally	Dates	438.167	978 j
Claire	Cherry	612.684	947 j
Sam	Ice	513.259	913 j
Chae	Apple	485.333	981
Sam	Eggs	533.357	963

#### **Phi Operators**

```
s = prod, x_sum_quant, y_sum_quant,
z_sum_quant
n = 3
v = prod
              O
                                 OR
F = [x_sum_quant, y_sum_quant,
z_sum_quant]
sigma = [X.month=1, Y.month=2,
Z.month=31
G = null
```

#### MF Query

```
SELECT prod,
sum(X.quant),
sum(Y.quant),
sum(Z.quant)
FROM sales
WHERE year=2017
GROUP BY prod : X, Y, Z
SUCH THAT X.month = 1,
Y.month = 2, Z.month = 3
```

process info() read\_file() schema info() H\_table()



mf\_struct

H table

## **SQL Limitations**

Upon execution of the below SQL query upon the Sales table, the system is unable to process the query structure. It conducts a 3-way self-join and computes the aggregate functions over the entire table. This method of evaluation can be very expensive and can produce inaccurate results, as shown below.

SELECT x.prod, sum(x.quant),
sum(y.quant), sum(z.quant)
FROM Sales x, Sales y, Sales z
WHERe x.prod=y.prod AND x.month=1
AND
 x.year=2017 AND y.prod=z.prod AND

y.month=2 AND y.year=2017 AND z.month=3 AND z.year=2017

GROUP BY x.prod

	prod character varying (20)	sum bigint	sum bigint	sum bigint
1	Apple	2099526	2261644	2119390
2	Butter	2055420	2459985	2545240
3	Cherry	2498560	2106560	2622080
4	Dates	662700	819450	570900
5	Eggs	3625622	1875236	4060121
6	Fish	2148800	3378880	3266560
7	Grapes	1295232	1355264	1308864
8	Ham	3938060	3444120	3187386
9	Ice	1535688	1868832	1177092
10	Jelly	1953567	1303731	1430247

## To rub it in some more...

This is the equivalent standard SQL query that generates a table with the correct results.

	prod character varying (20)	jan_sum numeric	feb_sum numeric <b>•</b>	mar_sum numeric
- 1	Apple	7341	7343	11645
2	Butter	7212	9647	7880
3	Cherry	9760	6583	8194
4	Dates	4418	5463	5709
5	Eggs	12677	3706	13579
6	Fish	5372	10559	10208
7	Grapes	6746	5294	6817
8	Ham	8561	9567	7699
9	Ice	9141	8652	4671
10	Jelly	8457	9117	5239

SELECT x.prod, SUM(x.sum\_quant) AS jan\_sum, SUM(y.sum\_quant) AS feb\_sum, SUM(z.sum\_quant) AS mar\_sum FROM (SELECT prod, SUM(quant) AS sum\_quant FROM sales WHERE month = 1 AND year = 2017GROUP BY prod) x JOIN (SELECT prod, SUM(quant) AS sum\_quant FROM sales WHERE month = 2 AND year = 2017GROUP BY prod) y ON x.prod = y.prodJOIN (SELECT prod, SUM(quant) AS sum\_quant FROM sales WHERE month = 3 AND year = 2017GROUP BY prod) z ON x.prod = z.prodGROUP BY x.prod;

(generated by ChatGPT)

I can't

even

read it

# Code Implementation

However, implementing an equivalent MF Query, using our generator, results in the correct implementation of the aggregate functions.

SELECT prod, sum(X.quant),
sum(Y.quant), sum(Z.quant)
FROM sales
WHERE year=2017
GROUP BY prod : X, Y, Z
SUCH THAT X.month = 1, Y.month = 2,
Z.month = 3



prod	sum(x.quant)	sum(y.quant)	+   sum(z.quant)
prou			Sum(2.quanc)   +
Apple	7341	7343	11645
Ham	8561	9567	7699
Fish	5372	10559	10208
Cherry	9760	6583	8194
Grapes	6746	5294	6817
Ice	9141	8652	4671
Jelly	8457	9117	5239
Butter	7212	9647	7880
Dates	4418	5463	5709
Eggs	12677	3706	13579
++			++

### Input Parsing: read\_file()

Reads an MF query file and each line is checked for specific keywords, and it extracts its components Inputs: filename

Outputs select, From, where, group\_by, such\_that, having:

```
ef read_file(filename):
  Reads an Extended SQL query and extracts its operands
  with open(filename, 'r') as file:
      lines = file.readlines()
      for line in lines:
          line = line.lower().strip() # make everything lowercase
      select = ""
      group by = ""
      such that = ""
      having = ""
      # Extract info
      if len(lines) > 0:
          select = lines[0][7:].strip().lower()
      if len(lines) > 1:
          From = lines[1][5:].strip().lower()
       if len(lines) > 2:
          if "where" in lines[2].lower():
              where = lines[2][6:].strip().lower()
```

### Schema Information: schema\_info()

connects to PostgreSQL database and queries the INFORMATION\_SCHEMA.COLUMNS table, retrieving column names and their data types and max lengths for the sales table

```
def schema info():
   This retrieves schema infomation from the database using 'information schema.columns'
   # hard coded query to send to database
   SELECT COLUMN NAME, DATA TYPE, CHARACTER MAXIMUM LENGTH
   FROM INFORMATION SCHEMA.COLUMNS
   WHERE TABLE SCHEMA = 'public' and TABLE NAME = 'sales'
   load dotenv()
   user = os.getenv('USER')
   password = os.getenv('PASSWORD')
   dbname = os.getenv('DBNAME')
   conn = psycopg2.connect("dbname="+dbname+" user="+user+" password="+password,
                               cursor factory=psycopg2.extras.DictCursor)
   cur = conn.cursor()
   cur.execute(query)
   data = cur.fetchall()
   conn.close()
   return data # return schema data
```

## Code Breakdown

#### Query Components: process\_info()

processes SELECT, GROUP BY, SUCH THAT, and aggregates populates the mf\_struct object

Outputs: group by vars and F vect: list of dictionaries for each aggregate

```
def process_info(select, group_by, such_that, having, mf_struct, schemaData):
   This function processes an MF query and extracts its information,
   populating mf_struct with 6 operands of Phi
   V = [] # list of grouping attributes
   F_VECT = [] # list of aggregate functions
   # potential aggregates
   aggregates = ["sum", "count", "avg", "min", "max"]
   # potential group by attributes
   groupAttrs = ["cust", "prod", "day", "month", "year", "state", "quant", "date"]
   # add projected vals to struct.s
   mf_struct.s = select
   # add group_by vars to mf_struct.v
   group by vars = []
   if len(such that) == 0:
       group_by = group_by.split(",") # normal SQL query (no such that conditions, group by seperated by comma)
   else:
```

#### SUCH THAT Conditions: process\_conditions()

Takes in mf\_struct and group by variables to separate them based on the such that clause. Returns such that conditions

## Code Breakdown

### H-table Construction: H table

Loops through database rows, groups rows based on attributes in GROUP BY, checks if SUCH THAT conditions are satisfied using eval\_conditions, and calculates aggregate values for SUM, AVG, MAX, etc

```
# only include unique rows based on group vars
if (len(mf_struct.v) != 0):
    row combo = tuple(row[col name.index(qv)] for qv in mf struct.v) # find index of group var based
else: # if no group by then use entire row as tuple
    row_combo = row
# if not in unique then add it (group by function)
if row_combo not in unique:
    unique.append(row combo)
# if there are aggregates, keep track of their values
if len(F_VECT) != 0:
    for f in F VECT:
        # if tuple of group by vals not already in agg_values, add new row
        if row_combo not in agg_values[f['agg']]:
            agg_values[f['agg']][row_combo] = []
        # filter based on such that conditions
        if (len(such that) != 0):
            conditions = process_conditions(mf_struct, group_by_vars)
            if eval_conditions(row, conditions, f) == False:
                continue # if it does not meet such that conditions, skip that row for the aggregate
        # get rid of period and only have column name to find index
        if '.' in f['arg']:
            name = f['arg'].split('.')[1]
            agg values[f['agg']][row combo].append(row[col name.index(name)])
```

#### **HAVING** clause

Filters the H-table based on the HAVING clause and normalizes column names for compatibility

```
def preprocess_having_clause(having, H):
    """
    Translates HAVING clause into a pandas-compatible query string.
    """
    having = having.replace("=", "==").replace("<>", "!=")
    for col in H.columns:
        if col in having:
              having = having.replace(col, f"`{col}`")
    return having
```

### **Technical Stack & Limitations**

Python for the main program logic, PostgreSQL for database and querying. Python Libraries:

psycopg2 for interacting with PostgreSQL pandas for data computation and dataframe tabulate for formatting output tables dotenv. for loading database credentials

Tools:

Query input is read from .txt files

#### **Limitations:**

- Our program performs multiple full scans of the database table to compute aggregates, applies SUCH THAT conditions, and builds the H-table, resulting in performance inefficiencies for large datasets, since minimal scanning techniques (e.g., combining scans or reducing redundant operations) are not implemented
- Additional limitations include hard coded columns (cust, prod, etc.) making our schema dependency static, possible performance overhead for large datasets due to pandas, operators sometimes ambiguous (i.e. use of minus in splitting/dates), and lack of query optimization as our aggregates are calculated
- In the future, these limitations, implementing minimal scanning, making the columns dynamics, and accounting for scalability can be addressed and improved upon.