ADVANCE ALGORITHM Experiment 1

Ayush Jain 60004200132 B3

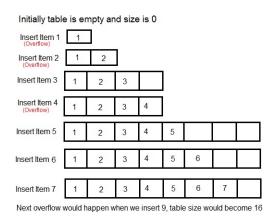
Aim: To implement Amortized Analysis.

Theory:

Amortized Analysis.

Amortized Analysis is used for algorithms where an occasional operation is very slow, but most of the other operations are faster. In Amortized Analysis, we analyze a sequence of operations and guarantee a worst-case average time which is lower than the worst-case time of a particular expensive operation. The example data structures whose operations are analyzed using Amortized Analysis are Hash Tables, Disjoint Sets and Splay Trees.

Let us consider an example of a simple hash table insertions. How do we decide table size? There is a trade-off between space and time, if we make hash-table size big, search time becomes low, but space required becomes high.



The solution to this trade-off problem is to use <u>Dynamic Table (or Arrays)</u>. The idea is to increase size of table whenever it becomes full. Following are the steps to follow when table becomes full.

- 1. Allocate memory for a larger table of size, typically twice the old table.
- 2. Copy the contents of old table to new table.
- 3. Free the old table.

Aggregate Method

The aggregate method is used to find the total cost. If we want to add a bunch of data, then we need to find the amortized cost by this formula.

For a sequence of n operations, the cost is –

$$\frac{Cost(n \ operations)}{n} = \frac{Cost(normal \ operations) + Cost(Expensive \ operations)}{n}$$

The Accounting Method

The accounting method is aptly named because it borrows ideas and terms from accounting. Here, each operation is assigned a charge, called the amortized cost. Some operations can be charged more or less than they actually cost. If an operation's amortized cost exceeds its actual cost, we assign the difference, called a credit, to specific objects in the data structure. Credit can be used later to help pay for other operations whose amortized cost is less than their actual cost. Credit can never be negative in any sequence of operations.

The Potential Method

The potential method is similar to the accounting method. However, instead of thinking about the analysis in terms of cost and credit, the potential method thinks of work already done as potential energy that can pay for later operations. This is similar to how rolling a rock up a hill creates potential energy that then can bring it back down the hill with no effort. Unlike the accounting method, however, potential energy is associated with the data structure as a whole, not with individual operations.

Code:

Aggregate Table:

```
class DynamicTable:
    def __init__(self, capacity=1):
       self.table = [0] * capacity
       self.size = 0
       self.capacity = capacity
    def add(self, element):
       if self.size == self.capacity:
           new_table = [0] * (self.capacity * 2)
           for i in range(self.size):
               new_table[i] = self.table[i]
           self.table = new_table
           self.capacity *= 2
        self.table.append(element)
        self.size += 1
    def size(self):
       return self.size
    def capacity(self):
        return self.capacity
    def numDoublings(self):
       return int(math.log2(self.capacity))
    def numCopyings(self):
       num_copyings = 0
       for i in range(1, self.numDoublings() + 1):
           num_copyings += 2**(i-1)
       return num copyings
table = DynamicTable()
cost = 0
f = 0
operation_cost = 1
print("Item No\tTable Size\tTable Cost\tCost of Operation")
print("============")
for i in range(1, 18):
   table.add(i)
    if f == 1:
```

```
cost = table.size
    f = 0
else:
    cost = 1
if table.size == table.capacity:
    f = 1
print(f"{i}\t{table.capacity}\t\t{cost}\t\t\t{operation_cost}")
```

Output:

Accounting using Multipop Code:

```
class Stack:
   def init (self):
       self.items = []
        self.cost = 0
        self.balance = 0
    def push(self, item):
        self.items.append(item)
        self.cost += 1
        self.balance += 1
    def pop(self):
        if not self.is empty():
            self.cost += 1
            self.balance -= 1
            return self.items.pop()
    def multi_pop(self, k):
        if k > len(self.items):
            k = len(self.items)
        for i in range(k):
            self.pop()
        self.cost += k
        self.balance -= k
    def is empty(self):
        return len(self.items) == 0
    def get_cost(self):
        return self.cost
   def get balance(self):
        return self.balance
    def display_table(self, n):
        print("{:<15}{:<15}{:<15}{:<15}".format("Operation", "Total Cost",</pre>
"Amortized Cost", "Balance"))
        print("-"*60)
        for i in range(1, n+1):
            if i % 3 == 1:
                self.push(i)
                print("{:<15}{:<15}{:<15.2f}{:<15}".format(f"push({i})",</pre>
self.get_cost(), self.get_cost()/i, self.get_balance()))
           elif i % 3 == 2:
                self.pop()
```

Output:

```
Enter the number of operations to perform on the stack: 5
Operation Total Cost Amortized Cost Balance
oush(1)
                            1.00
                                           1
              1
                                           0
oop()
              2
                            1.00
                                           0
multi pop(1) 2
                            0.67
oush(4)
              3
                            0.75
                                           1
                                           0
oop()
              4
                            0.80
...Program finished with exit code 0
Press ENTER to exit console.
```

Potential Method:

```
class DynamicTable:
    def __init__(self, capacity=1):
        self.table = [∅] * capacity
        self.size = 0
        self.capacity = capacity
    def add(self, element):
        if self.size == self.capacity:
            new_table = [0] * (self.capacity * 2)
            for i in range(self.size):
                new_table[i] = self.table[i]
            self.table = new table
            self.capacity *= 2
        self.table[self.size] = element
        self.size += 1
    def size(self):
        return self.size
    def capacity(self):
        return self.capacity
def potential(table):
    return 2*table.size - table.capacity
def cost(table, operation):
    if operation == "add":
        if table.size == table.capacity:
            return potential(table) + 1
        else:
            return 1
        raise Exception("Invalid operation")
table = DynamicTable()
total cost = 0
operation_cost = 1
operation_cos = 1
prev = 0
print("Item No.\tTable Size\tPotential\tOperation Cost\tTotal Cost\tAmortized
Cost")
for i in range(1, 18):
   operation cost = cost(table, "add")
```

```
total_cost += operation_cost
table.add(i)
# print("="*8)
pot = potential(table)
amortized_cost = operation_cost + (pot - prev)
prev = pot
# print(amortized_cost)
print(f"{i}\t\t {table.capacity}\t\t {potential(table)}\t\t {operation_cos}\t\t {operation_cost}\t\t {operation_cost}\t\t {amortized_cost}")
```

Output:

Item No.	Table Size	Potential	Operation Cost	Total Cost	Amortized Cost
1	1	1	1	1	2
2	2	2	1	2	3
3	4	2	1	3	3
4	4	4	1	1	3
5	8	2	1	5	3
6	8	4	1	1	3
7	8	6	1	1	3
8	8	8	1	1	3
9	16	2	1	9	3
10	16	4	1	1	3
11	16	6	1	1	3
12	16	8	1	1	3
13	16	10	1	1	3
14	16	12	1	1	3
15	16	14	1	1	3
16	16	16	1	1	3
17	32	2	1	17	3
Program finished with exit code 0					
Press ENTER to exit console.					

Conclusion: In conclusion, we learned the Amortized Analysis of the Algorithm where an occasional operation is very slow, but most of the other operations are faster.