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COMPUTER NETWORKS AND INTERNET PROTOCOLS

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**Line Rate
1Gbps**

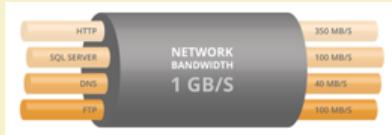


Internet QoS

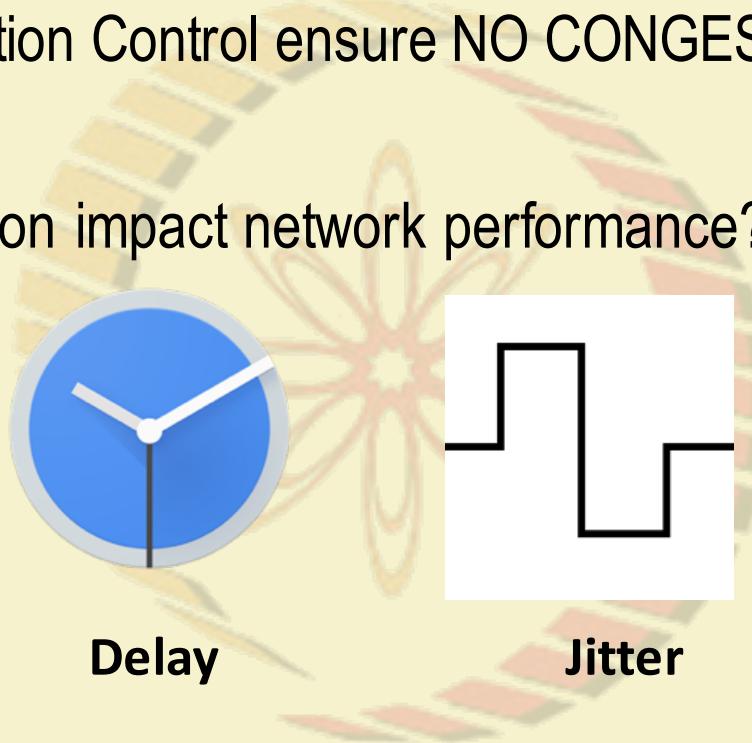
What is QoS

Revisiting Congestion

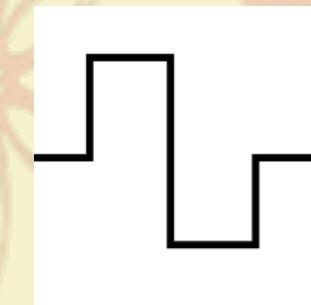
- Does TCP Congestion Control ensure NO CONGESTION in the network?
- How does congestion impact network performance?



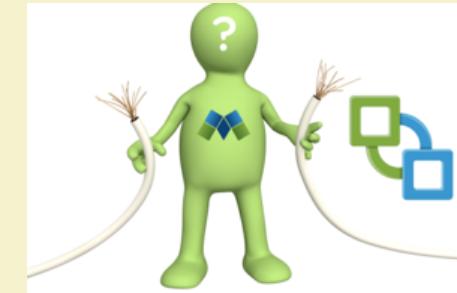
Bandwidth



Delay



Jitter



Loss

Network Bandwidth

- Amount of data that can be transmitted over a link within a fixed amount of time.
- *“When a drain chronically runs slow even though it isn’t plugged, it’s time to get a bigger pipe.”* – QoS vs more bandwidth, by Tim Greene
- Some applications in the network are bandwidth hungry – video applications; congestion limits per user bandwidth.
- Improve bandwidth – design networks with high capacity

Delay

- Three components – (a) **transmission delay**, (b) **propagation delay**, (c) **queueing delay**
- **Transmission Delay:** Amount of time to push all the packet bits in the network. Bandwidth 8 Mbps, Packet size (including headers) 1 MB, Transmission delay?
- **Propagation Delay:** Time to transfer one bit from one end of the link to another end of the link; usually depends on the underlying communication media
- **Queuing Delay:** Delay at the interface buffer; the major delay component

Delay

- In general,
Queuing delay >> Transmission delay + Propagation delay
- Packet multiplexing in the network devices (like routers, switches) impact the queuing delay
- Impact of congestion – increase in queuing delay



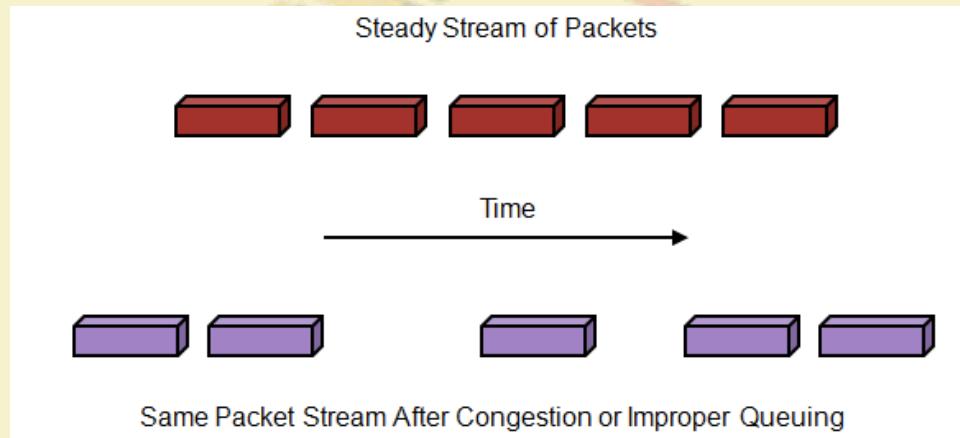
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Jitter

- Variation in End to End delay



- Why jitter impacts application performance – video streaming; consider delay variation among different video frames

Loss

- A relative measure of the number of packets (or segments or bits) that were not received compared to the total number of packets (or segments or bits) transmitted.
- Loss is a function of availability
 - If the network is available (capacity more than the demand) then loss would be generally zero (**Note: This assumption is not true for wireless networks**)
- Congestion increases data loss from the intermediate network devices

Application QoS

- Different application requires different level of QoS – delay, jitter and bandwidth

	Loss	Delay (One-way)	Jitter	Bandwidth
Voice	$\leq 1\%$	$\leq 150\text{ms}$	$\leq 30\text{ms}$	21 Kbps - 320 Kbps
Interactive Video	$\leq 1\%$	$\leq 150\text{ms}$	$\leq 30\text{ms}$	On demand
Streaming Video	$\leq 5\%$	$\leq \text{Buffer time}$	On buffer time	On demand
Data	-	-	-	Best Effort

Source: <http://www.ciscopress.com/articles/article.asp?p=357102>

A Formal Definition of QoS

“Quality of Service (QoS) refers to the **capability of a network to provide better service to selected network traffic over various technologies**, including Frame Relay, Asynchronous Transfer Mode (ATM), Ethernet and 802.1 networks, SONET, and IP-routed networks that **may use any or all of these underlying technologies**. The primary goal of QoS is to **provide priority including dedicated bandwidth**, controlled jitter and latency (required by some real-time and interactive traffic), and **improved loss characteristics**. “

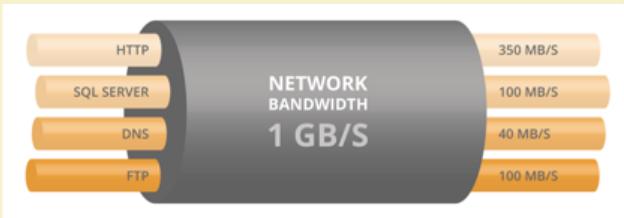
Source: Cisco - http://docwiki.cisco.com/wiki/Quality_of_Service_Networking

Ensure QoS over a Packet Switching Network

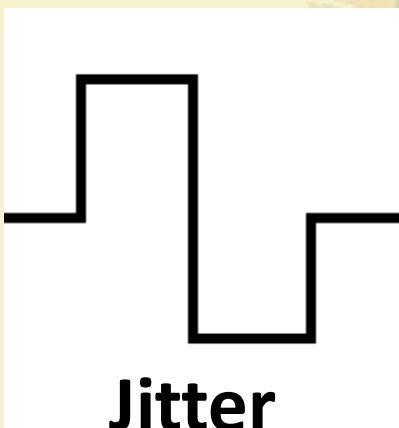
- What applications need from the network
- How to regulate the traffic that enters the network
- How to reserve resources at router to guarantee performance
- Whether the network can safely accept more traffic

Section 4, Chapter – THE NETWORK LAYER, Computer Network by Tanenbaum

Primary QoS Parameters



Bandwidth



Jitter



Loss



Delay



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Flow

- A stream of packets from a source to a destination is called a **flow**.
- Source to Destination:
 - Machine to Machine
 - Process to Process
 - Application to Application
 - Socket to Socket
- Different flows require different levels of QoS

Why QoS is Considered at the Network Layer

- Maintaining QoS requires both per-hop and end-to-end behavior
- End-to-end performance needs to be monitored
 - End-to-end delay
 - End-to-end bandwidth
 - End-to-end jitter
 - Total end-to-end data loss



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Why QoS is Considered at the Network Layer

- However, resource reservation needs to be on per-hop basis – otherwise end-to-end requirements can not be guaranteed
- Network layer bridges end-to-end (Transport) and Per-hop (Data Link Layer)



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Application Classes based on QoS

- Constant bit rate (e.g. telephone applications – VoIP)
- Real time variable bit rate (e.g. videoconferencing)
- Non real-time variable bit rate (e.g. on demand video streaming – IPTV)
- Available bit rate or Best effort (e.g. File transfer)



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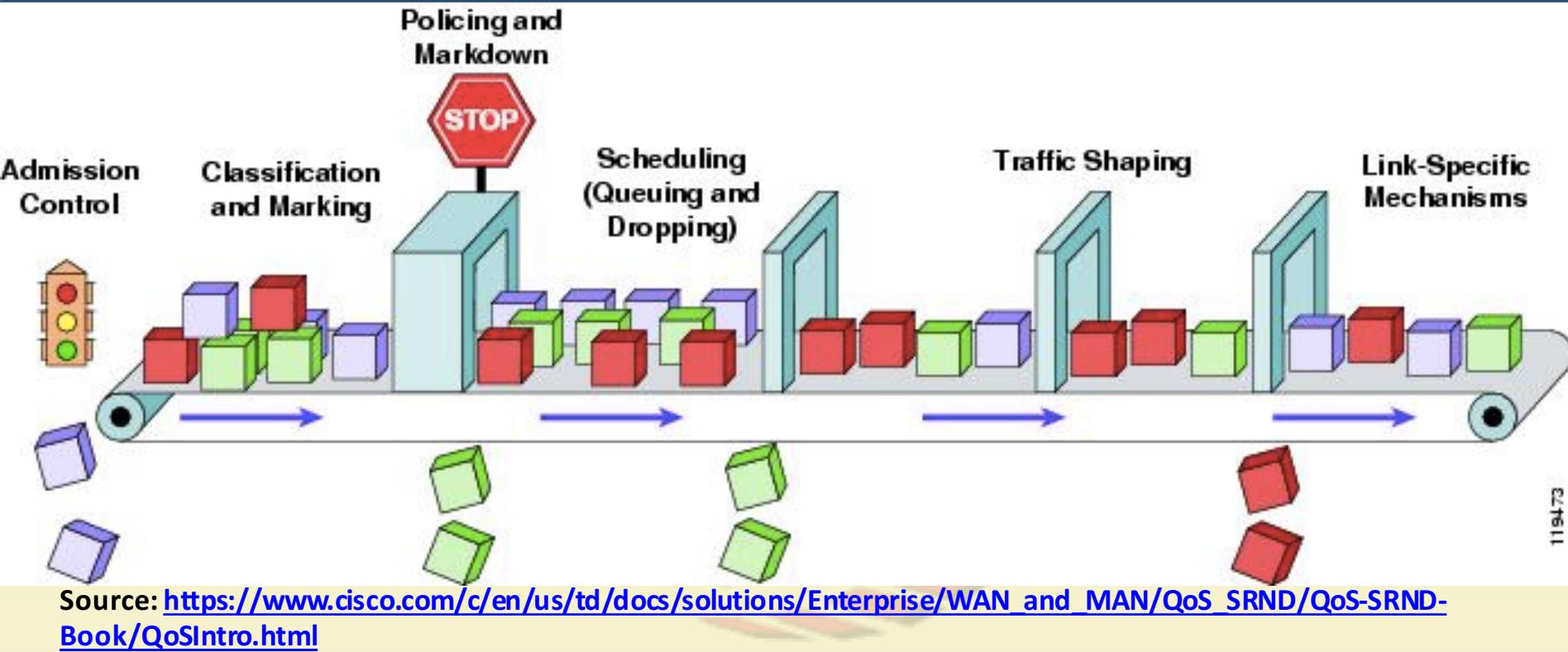
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Internet QoS Basic QoS Architecture

Basic QoS Architecture

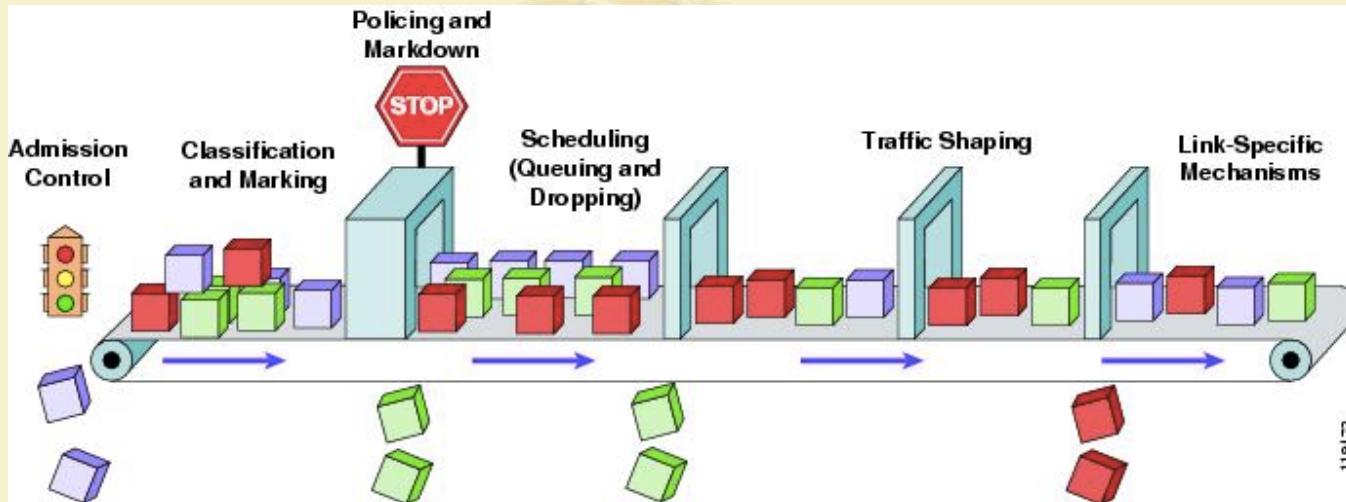


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Basic QoS Architecture



- **Admission Control:** Ensures that new flows are entered in the network only if the QoS of all the existing flows along with the new flows can be satisfied

Admission Control

- Remember the phase from cellular networks – “*all lines are busy; please dial after sometimes*”
- The network does not allow new flows if all the network resources are blocked in servicing the existing flows based on their QoS requirements



YOU SHALL NOT PASS!

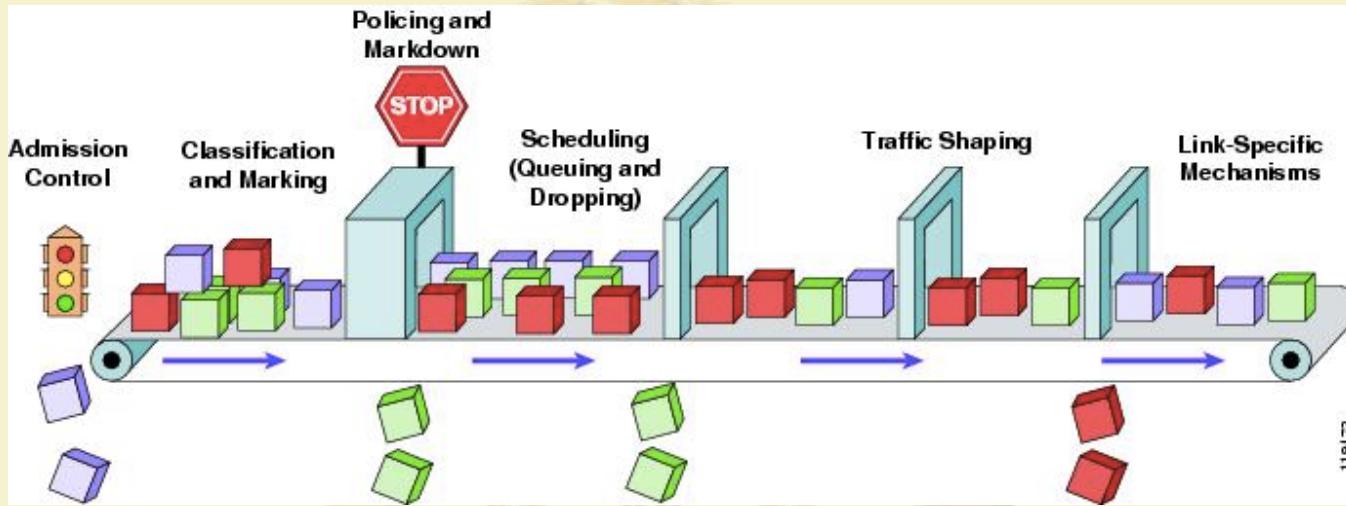


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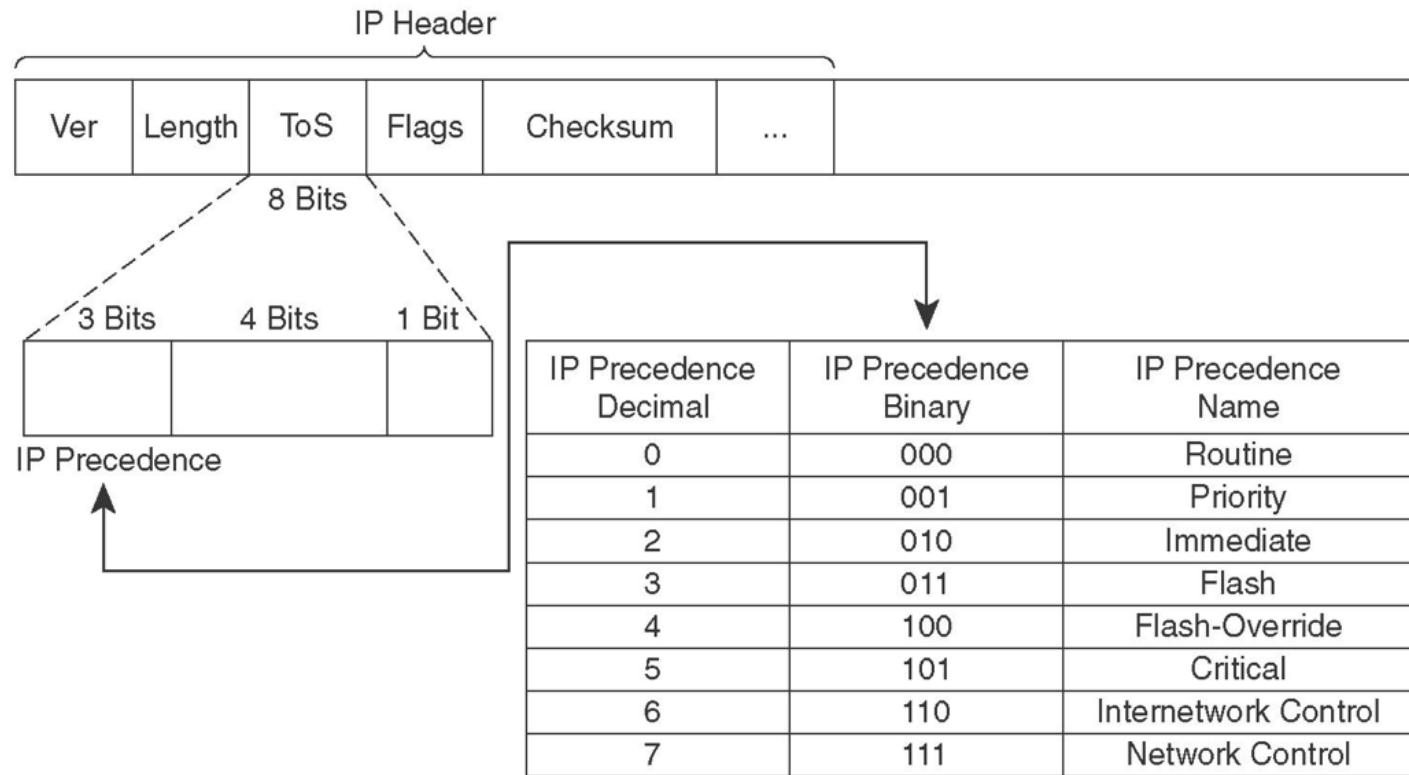
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Basic QoS Architecture



- **Classification and Marking:** Classifies the packets based on their application QoS requirements, and mark the packets accordingly

IP Type of Service (ToS) Field

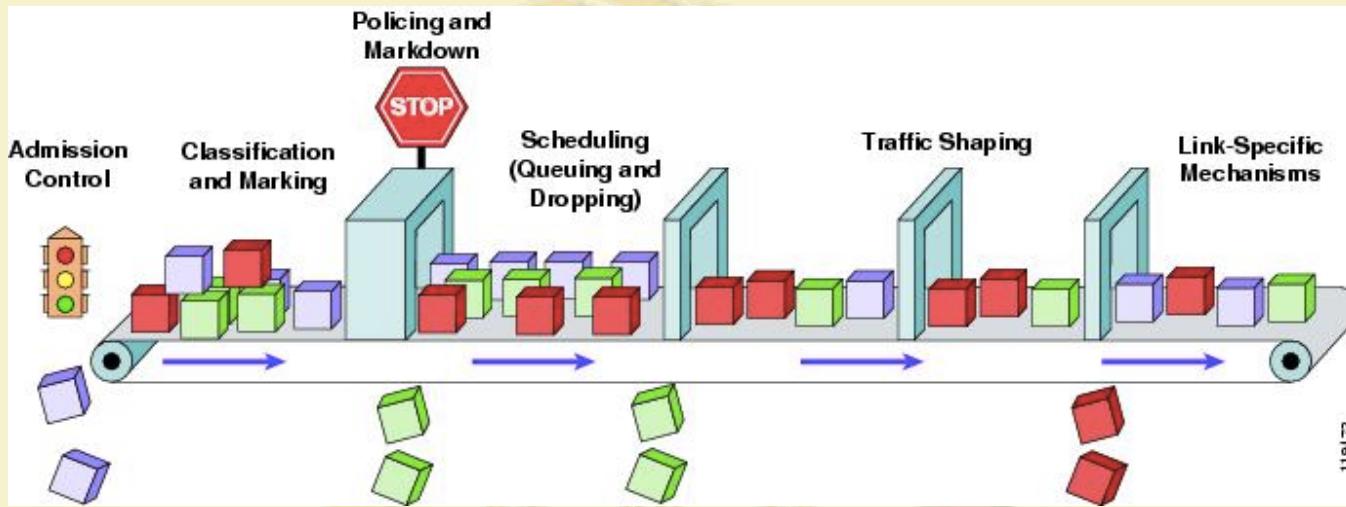


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Basic QoS Architecture



- **Policing and Markdown:** Monitor the flow characteristics and take appropriate actions based on the flow QoS

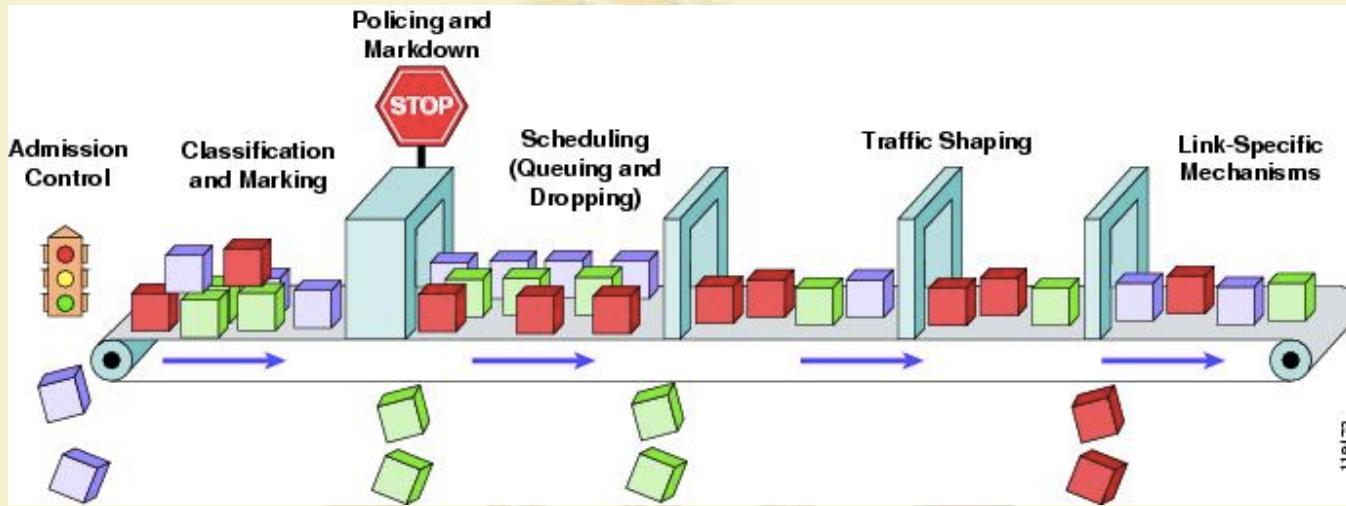
Traffic Policing

- **Service Level Agreements (SLA)**: An agreement or a contract between the customer and the service provider to maintain QoS of an application

```
1 ip sla 11
2   icmp-echo 10.0.10.11 source-ip 68.68.1.2
3   frequency 5
4 !
5 track 10 ip sla 11 reachability
6   delay down 15 up 15
7 !
8 ip sla schedule 11 life forever start-time now
9 ip sla enable reaction-alerts
```

- **Traffic policing** monitors the flow of traffic and mark them to take appropriate actions (reduce priority, drop etc.)

Basic QoS Architecture



- **Scheduling:** Based on markdown by the traffic polices, schedule the traffic into output buffers of an interface



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Traffic Scheduling based on Queuing

- Maintain multiple Queues at the interface, the scheduling mechanism services the queues based on the scheduling policy
- Example: Priority Queuing

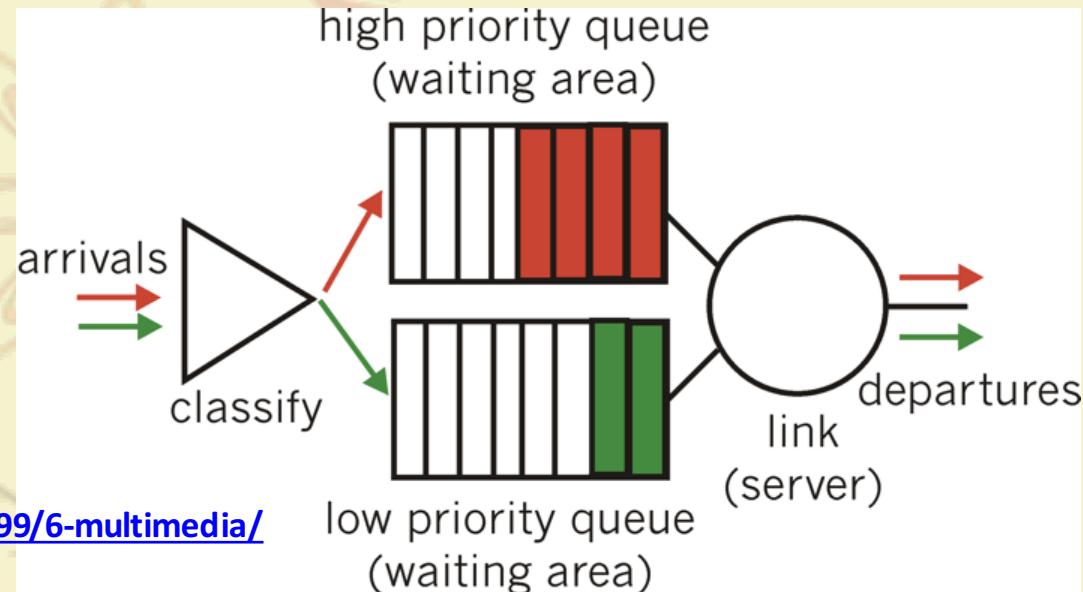
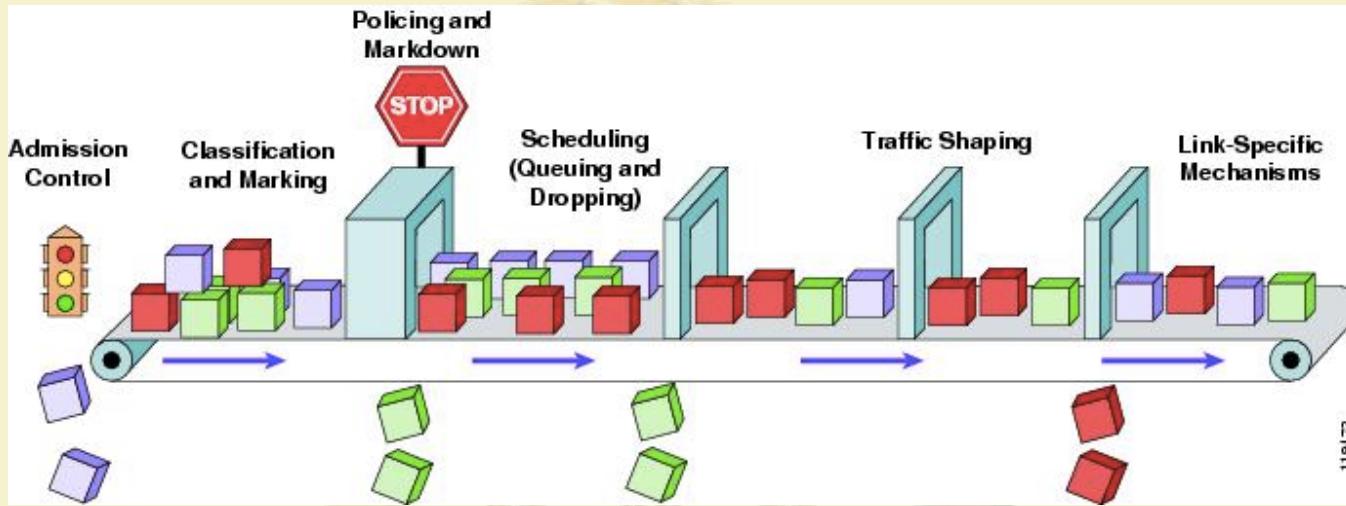


Image source: <http://www2.ic.uff.br/~michael/kr1999/6-multimedia/>

Basic QoS Architecture



- **Traffic Shaping:** Control the outgoing traffic rate irrespective of the incoming traffic rate (e.g. constant bit rate output from the interface buffer)

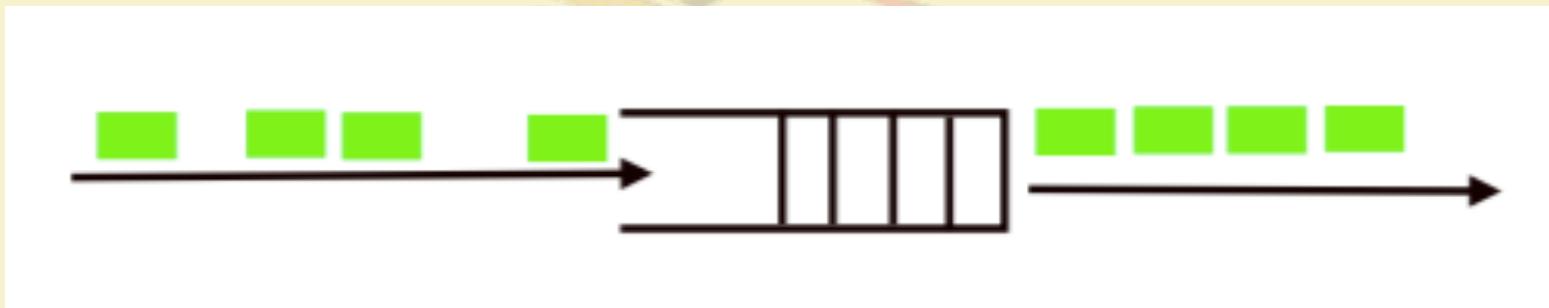


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Traffic Shaping

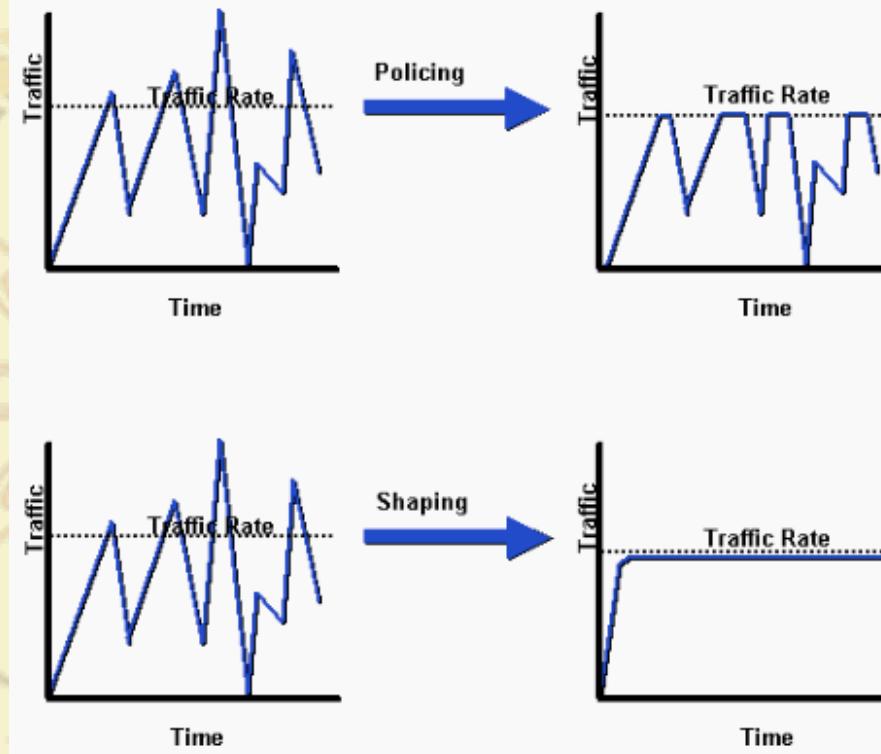


- Input traffic is bursty
- Output traffic has constant packet rate – reduces the jitter

Traffic Policing versus Traffic Shaping

Source:

<https://www.cisco.com/c/en/us/support/docs/quality-of-service-qos/qos-policing/19645-policevsshape.html>





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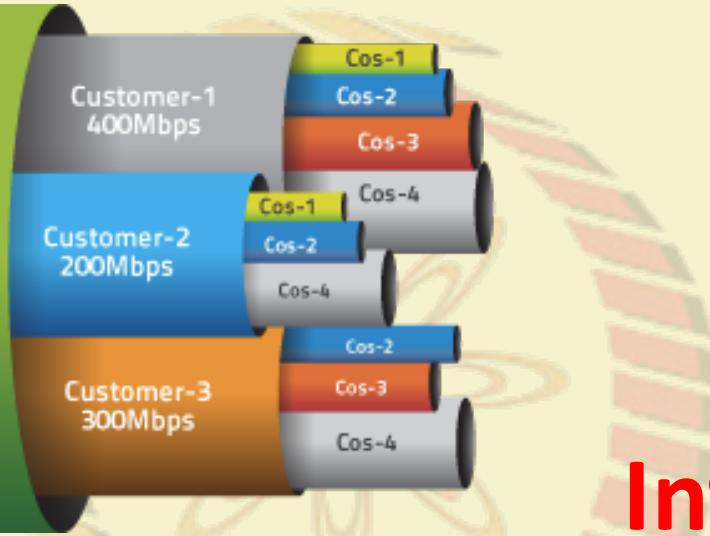
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Internet QoS Digging Further - Policing, Shaping and Scheduling



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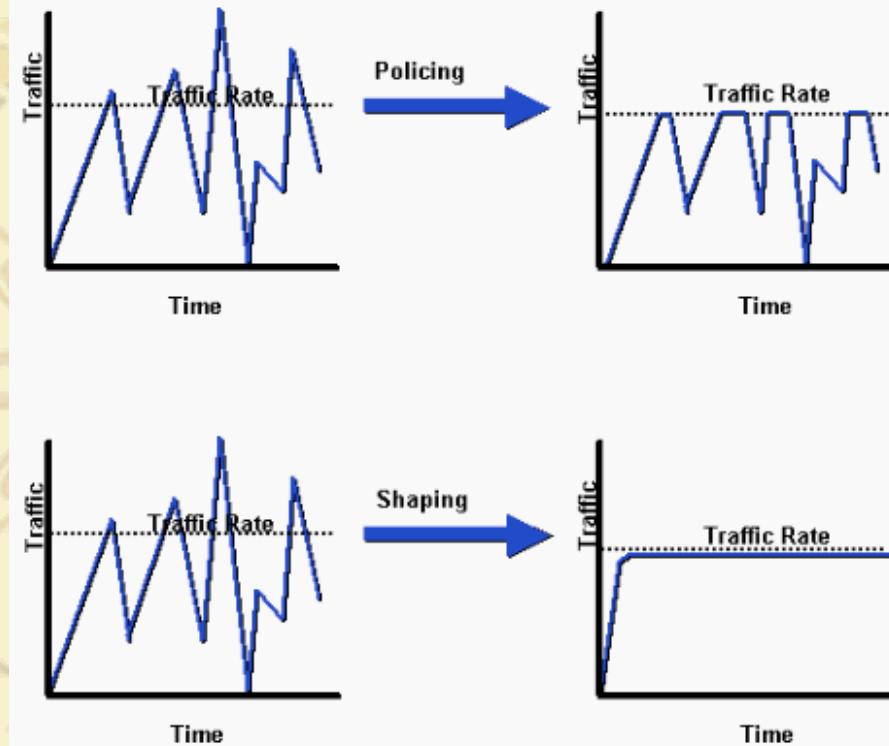


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Traffic Policing versus Traffic Shaping

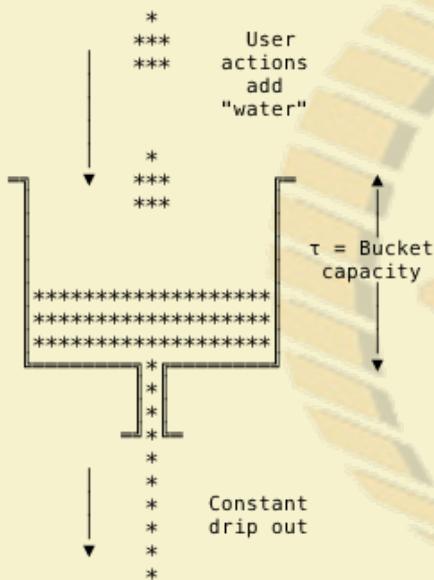
Source:

<https://www.cisco.com/c/en/us/support/docs/quality-of-service-qos/qos-policing/19645-policevsshape.html>



Leaky Bucket for Traffic Policing

- Incoming packets are put in the packet queue – the packet queue works as a bucket
- A single server queue with constant service time
- If the bucket overflows that the packets are discarded
- Input rate can vary but the output remains constant

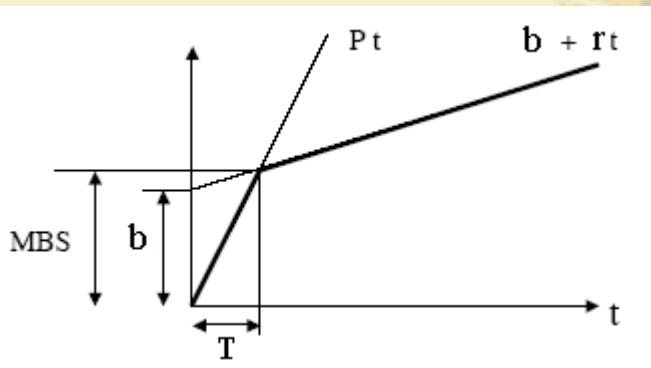
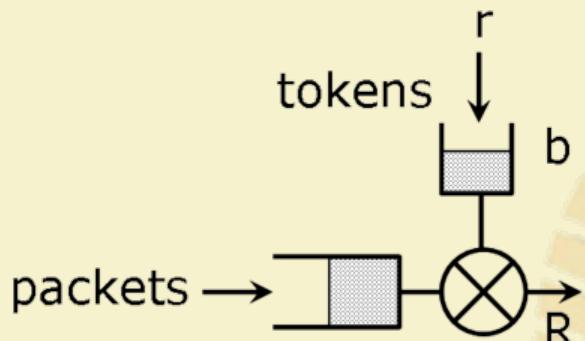


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Token Bucket for Traffic Policing



- Incoming packets are put in the packet queue.
- Token generation rate = r tokens/second, bucket size = b
- The rate of output traffic is bounded by the token generation rate
- Output rate $R(t) = \min(Pt, b + rt)$



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Leaky Bucket versus Token Bucket

- **Leaky Bucket:** Smooth out traffic but does not permit burstiness
- **Token Bucket:** Smooth out traffic and also permits burstiness – if there is no incoming packet, tokens are get added in the token bucket, and the burst traffic is permitted up to the amount of token accumulated
- Both the leaky bucket and the token bucket algorithms can be used for **traffic shaping**, with the addition of a **playout buffer**
 - A buffer to add some additional delay to the packets that arrived too early.



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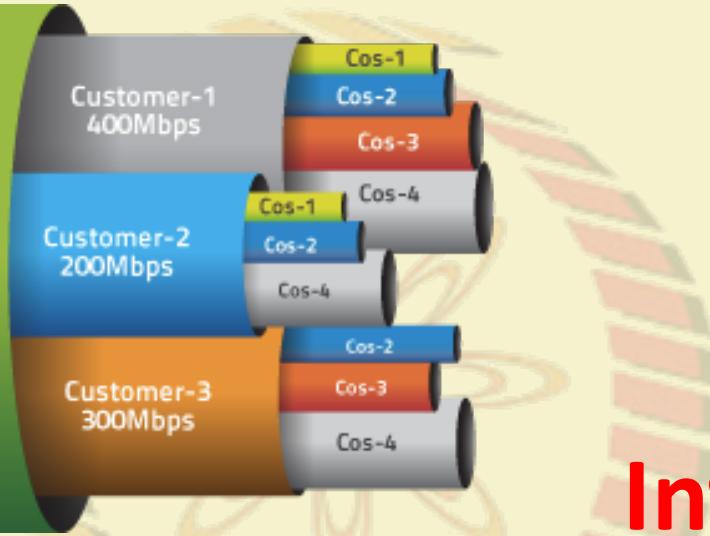
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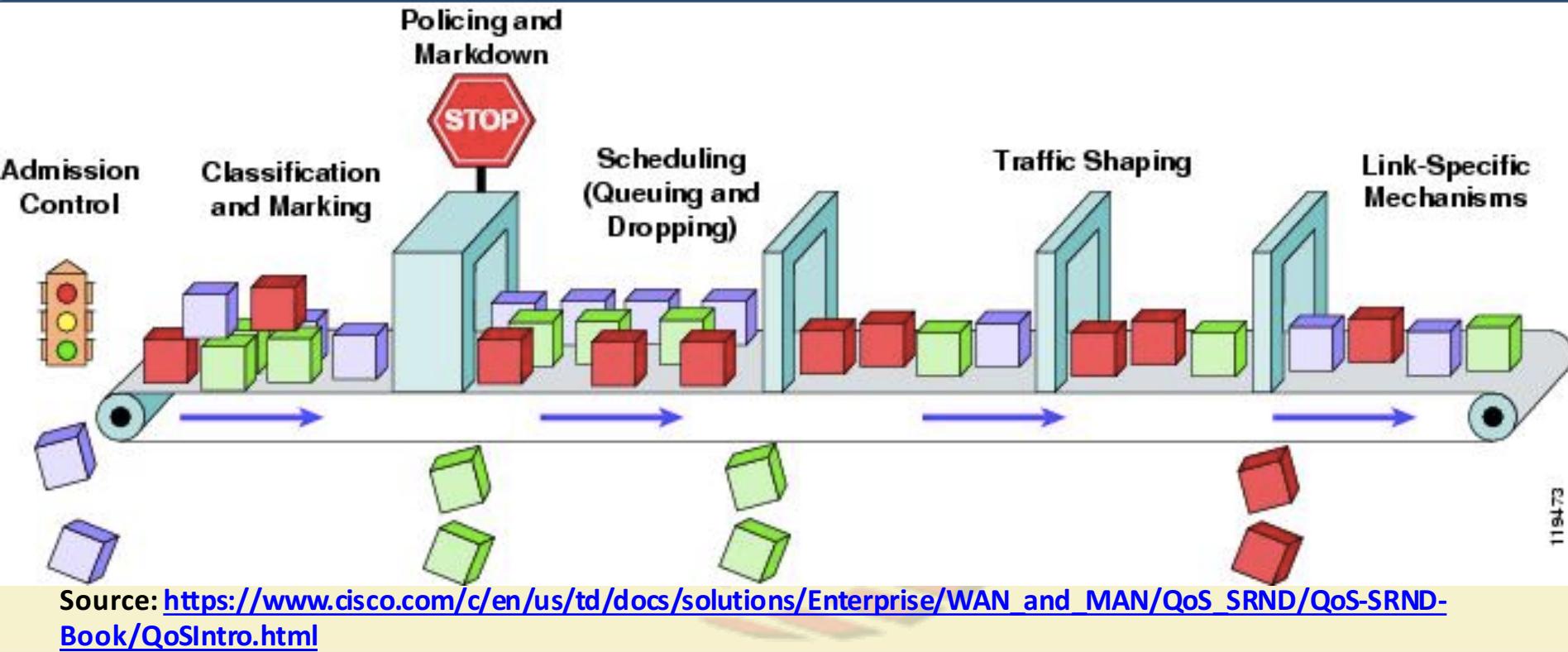
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Internet QoS

Digging Further - Queuing and Congestion Avoidance

Basic QoS Architecture



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Scheduling

- Classification and Marking - marked the data packers into different traffic classes.
 - The marked traffics are of different priority classes and require different level of QoS based on their SLA.
 - High priority delay-sensitive traffic (VoIP)
 - High bandwidth requirements (VoD, IPTV)
 - Best Effort services (HTTP, FTP)



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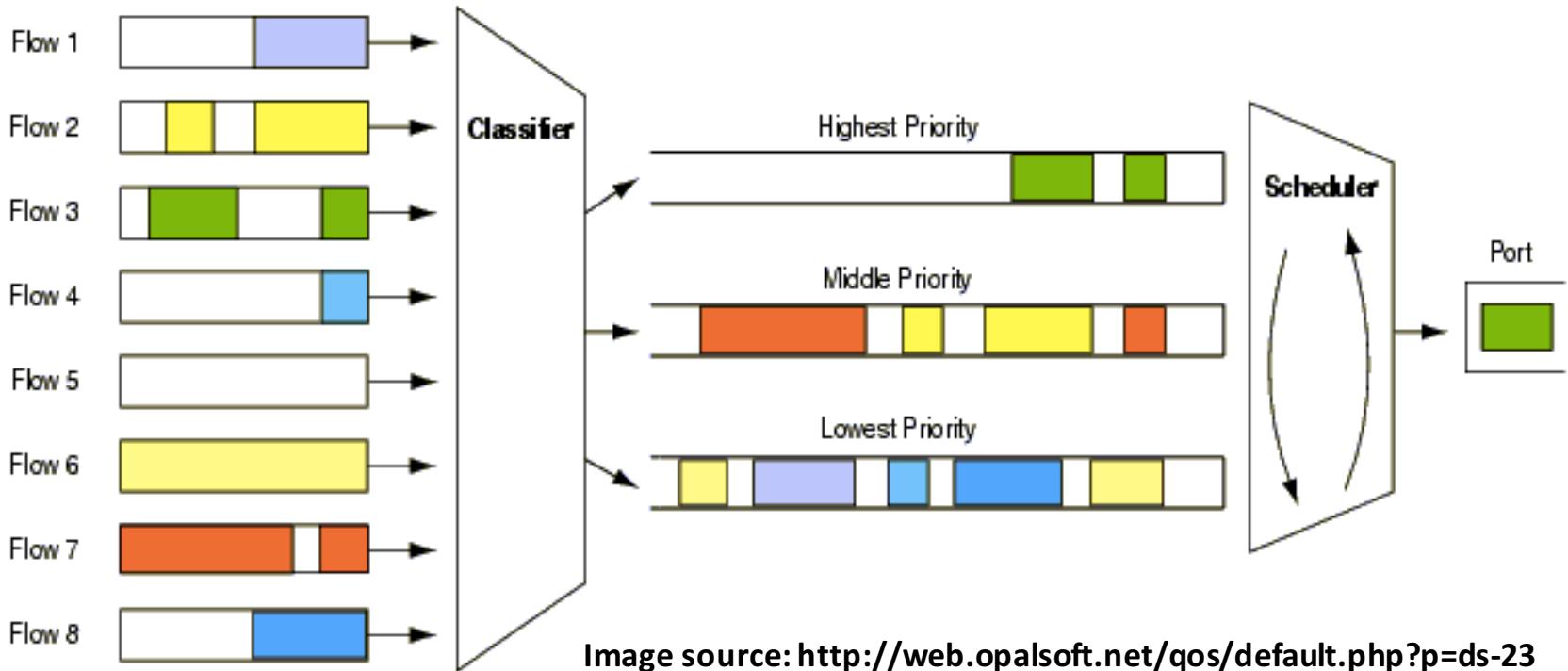
Multi-Class Scheduling

- We call them as different “class”es of traffic
 - Traffic class 1: High priority delay sensitive traffic
 - Traffic class 2: Medium priority bandwidth hungry traffic
 - Traffic class 3: Low priority best effort traffic
- Different classes of traffic requires different treatments

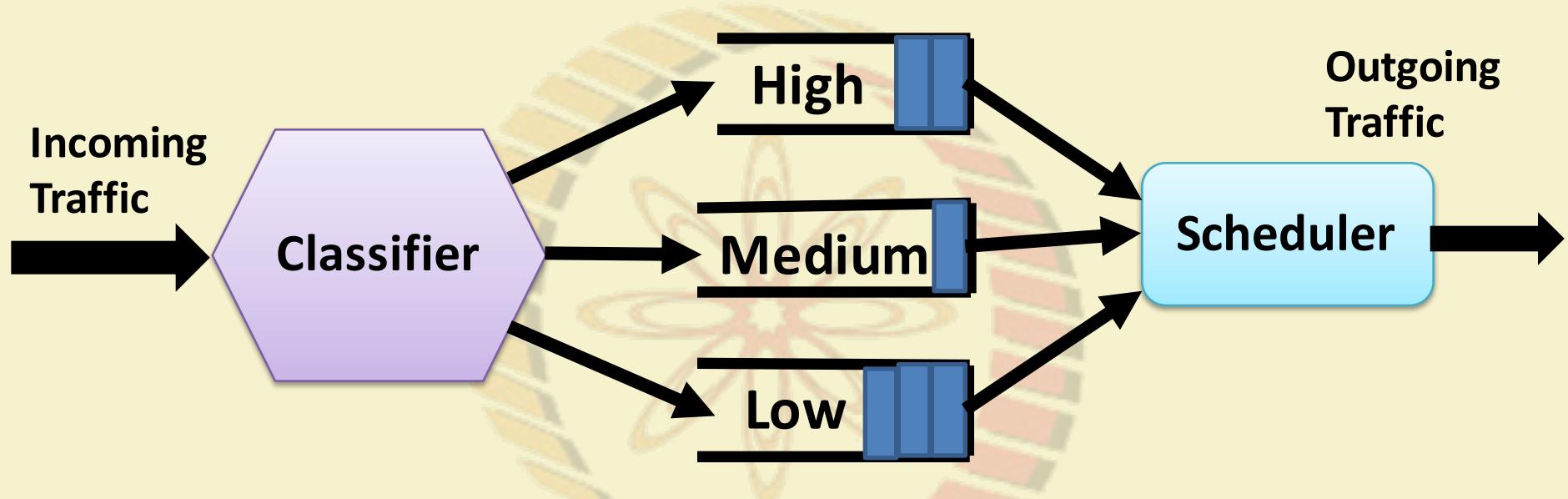
Multi-Class Scheduling

- Different classes of traffic requires different treatments
 - Traffic class 1: Ensure minimum queuing delay for the packets
 - Traffic class 2: Ensure sufficient bandwidth for the packets
 - Traffic class 3: No specific requirements, serve using best-effort services
- Use different queuing strategies

Queuing for Multi-Class Scheduling

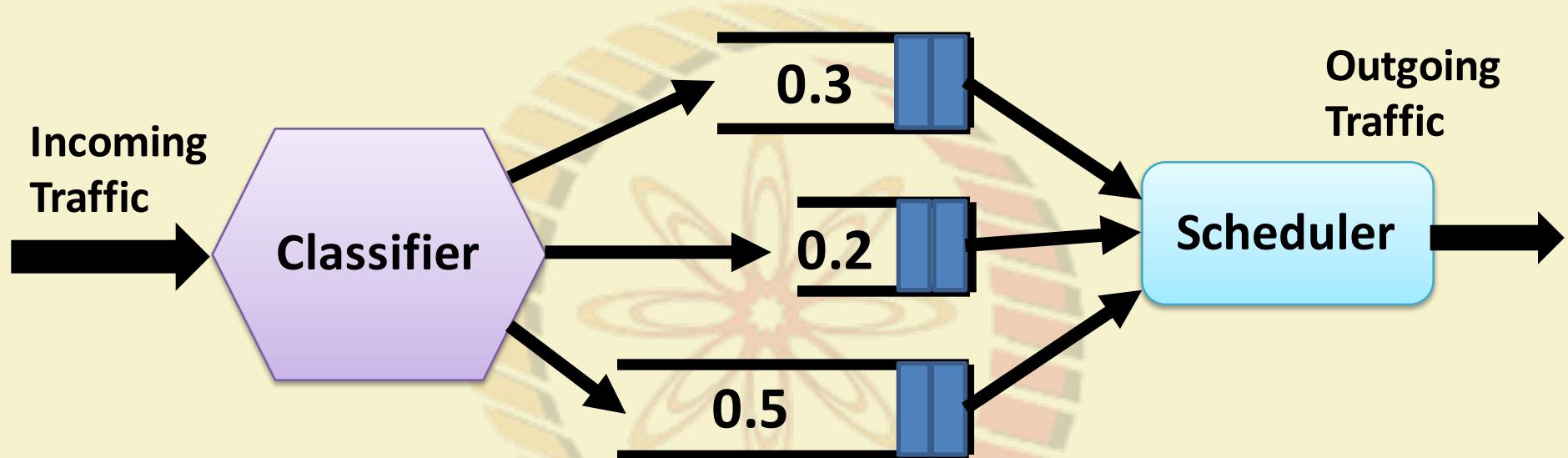


Scheduling – Priority Queuing



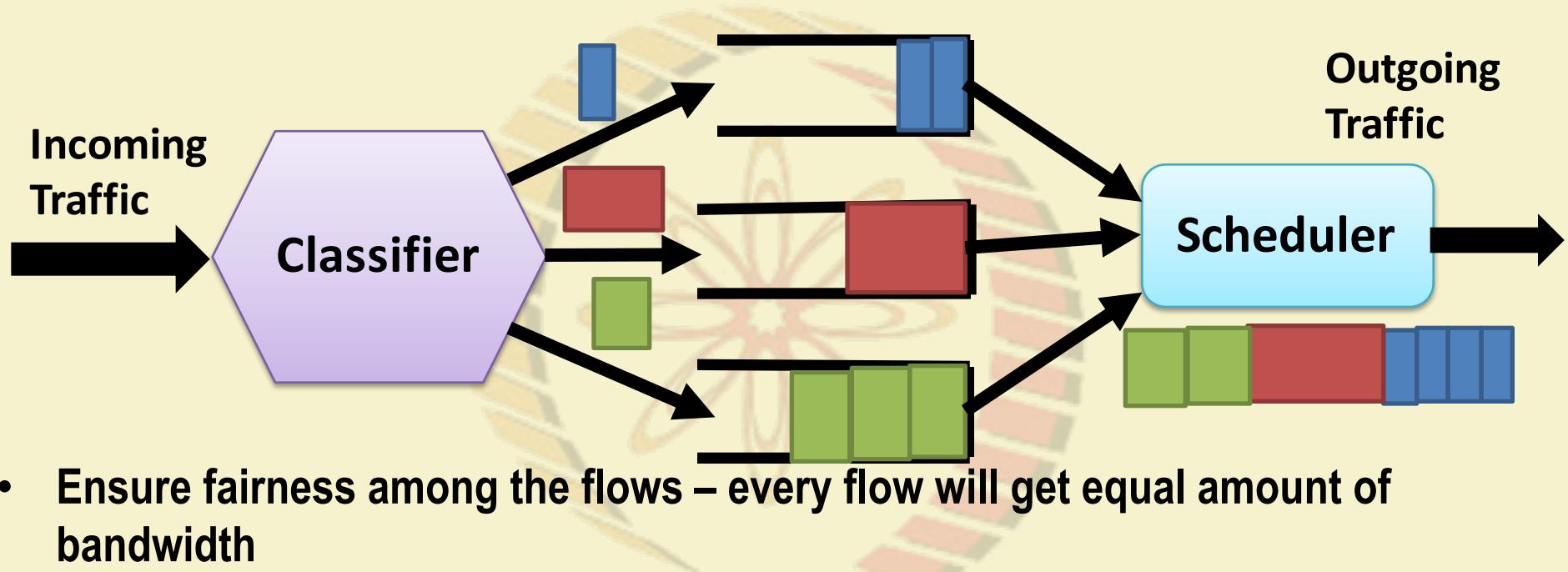
- Transfer traffic from lower priority queues once the higher priority queues are empty

Scheduling – Custom Queueing



- Different queues are of different size – works like transmission window
- Supports guaranteed bandwidth

Scheduling – Weighted Fair Queuing (WFQ)



Congestion Avoidance in the Internet

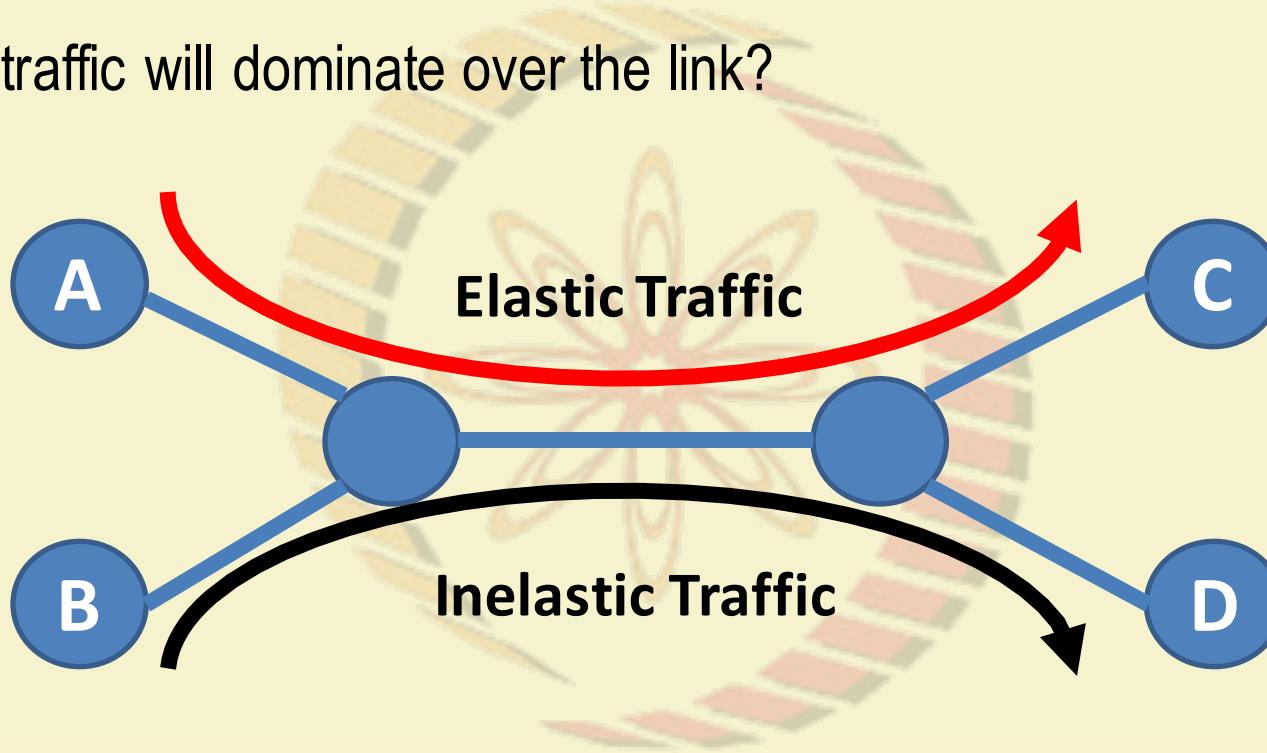
- Avoid congestion by controlling the queue parameters
- **Different from TCP congestion control**
 - TCP response on detection of the congestion
 - **Congestion avoidance** – ensure that a flow performance does not get affected by the congestion
- If congestion avoidance is there at the network layer, do we still need congestion control at transport layer?

Why Congestion Avoidance is Necessary for QoS

- Internet carries multiple data packets from different applications having different QoS requirements
- **Elastic traffic** – TCP like traffic, elastic nature of flow control based on AIMD
- **Inelastic Traffic** – UDP traffic, smoothed or controlled or constant bit rate traffic
- **Inelastic traffic is preferred for real time applications** – Can you say why?
(NB: Don't confuse with YouTube!)

Why Congestion Avoidance is Necessary for QoS

- Which traffic will dominate over the link?



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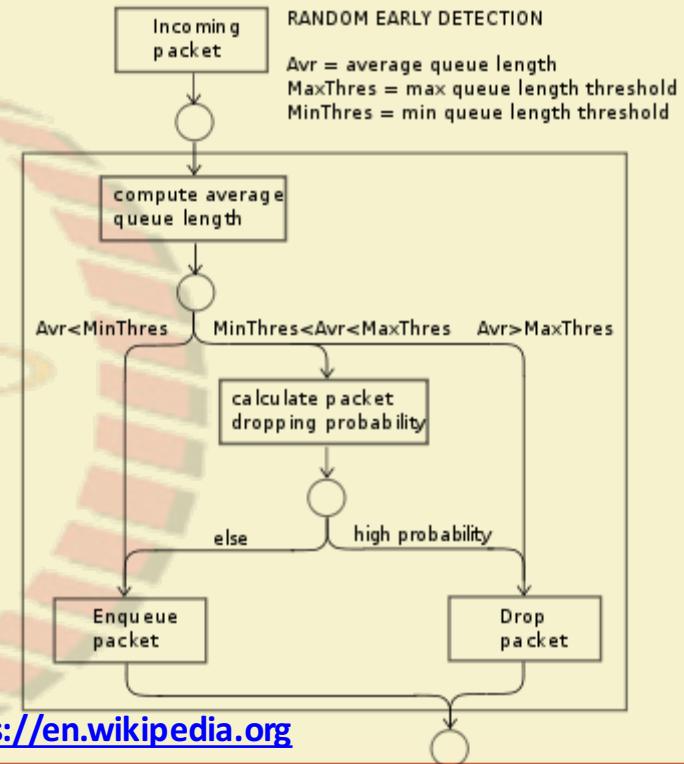
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Random Early Detection (RED)

- Drop probability is different for different traffics, depending on the nature – elastic or inelastic
- RED smooths out the drop probability across all the flows depending on the congestion probability
- Detect the possibility of congestion in the Internet. If congestion possibility is high, **RANDOMLY** drop packets before the enqueue

Random Early Detection (RED)

- Determine the possibility of packet drop by observing the average queue length

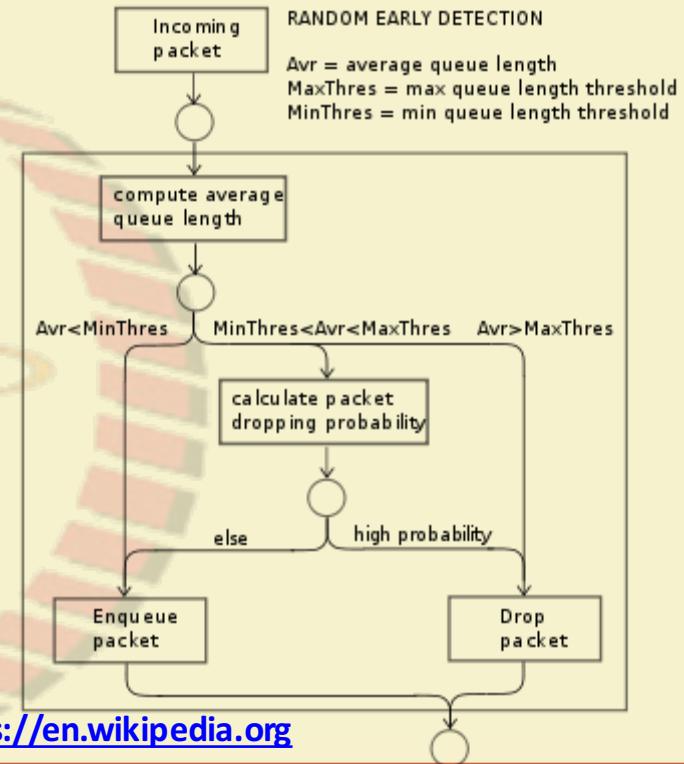


Source: <https://en.wikipedia.org>

Random Early Detection (RED)

- Let $d(k)$ denotes the packet drop probability and Max_p be the maximum packet drop probability

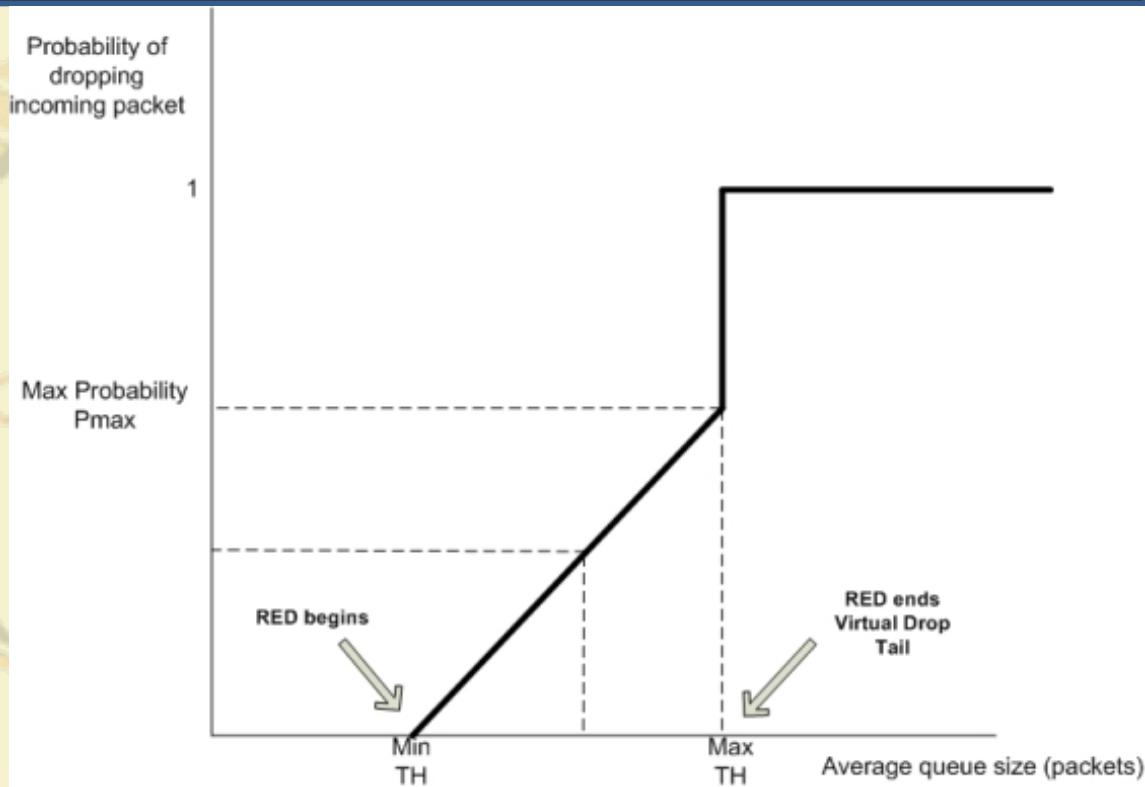
$$d(k) = Max_p \times \left(\frac{k - MinThresh}{MaxThresh - MinThresh} \right)$$



Source: <https://en.wikipedia.org>

Random Early Detection (RED)

Source: <http://www.eng.uwi.tt/>



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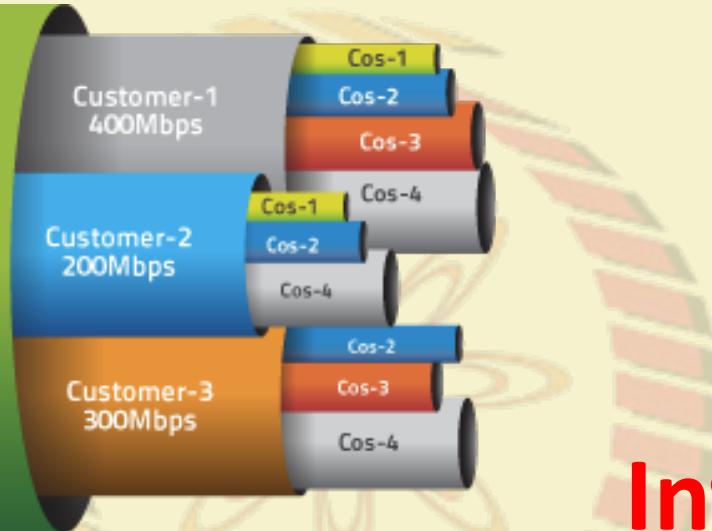
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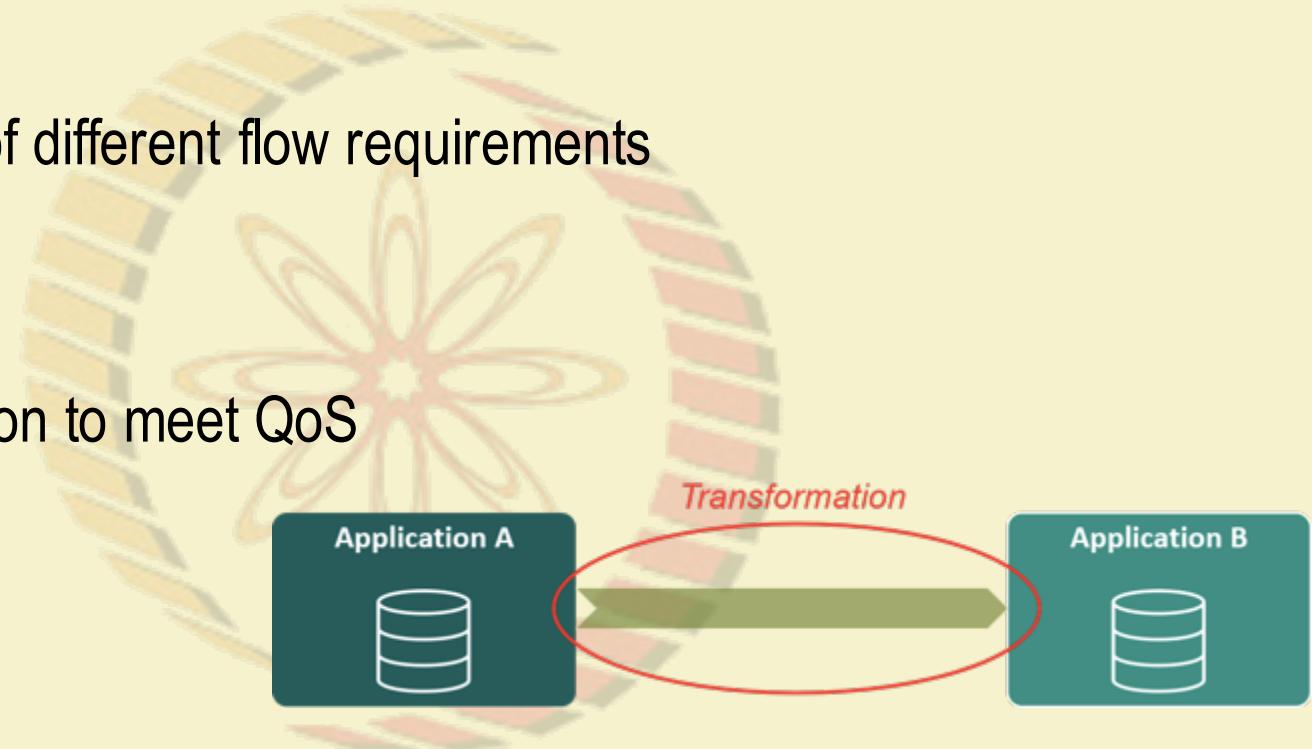
Internet QoS IntServ and DiffServ Architectures

Internet Service Architecture (ISA)

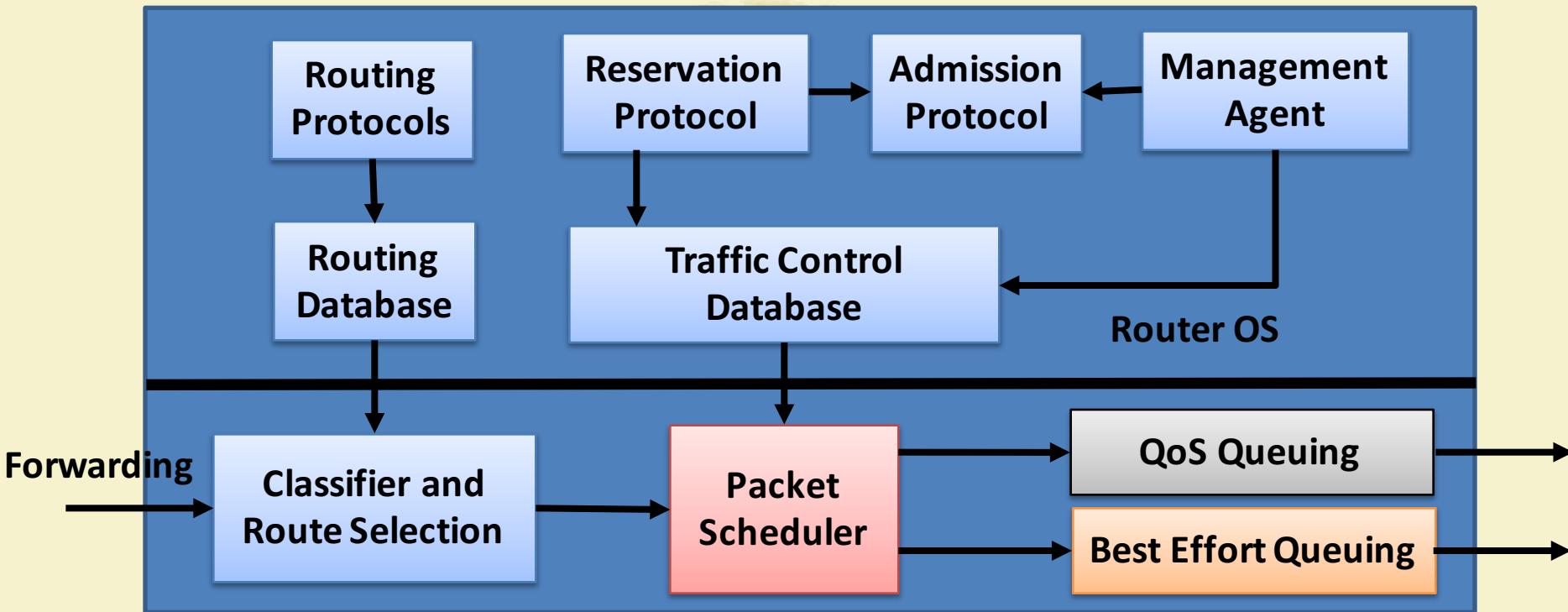
- Integrated Service (IntServ) QoS architecture over the Internet
- **Admission Control**
 - For QoS, reservation required for new flow
 - Resource Reservation Protocol (RSVP)
- **Routing Algorithm** - routing decision based on QoS parameters

Internet Service Architecture (ISA)

- **Queuing**
 - Take account of different flow requirements
- **Discard Policy**
 - Avoid congestion to meet QoS



ISA in a Router



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Resource Reservation Protocol (RSVP)

- A network control protocol that allows data receiver to request a special end to end quality of service for its data flows
- A network control protocol, not a routing protocol - although works over IP
- Designed to operate with current and future unicast and multicast routing protocols

ISA and RSVP

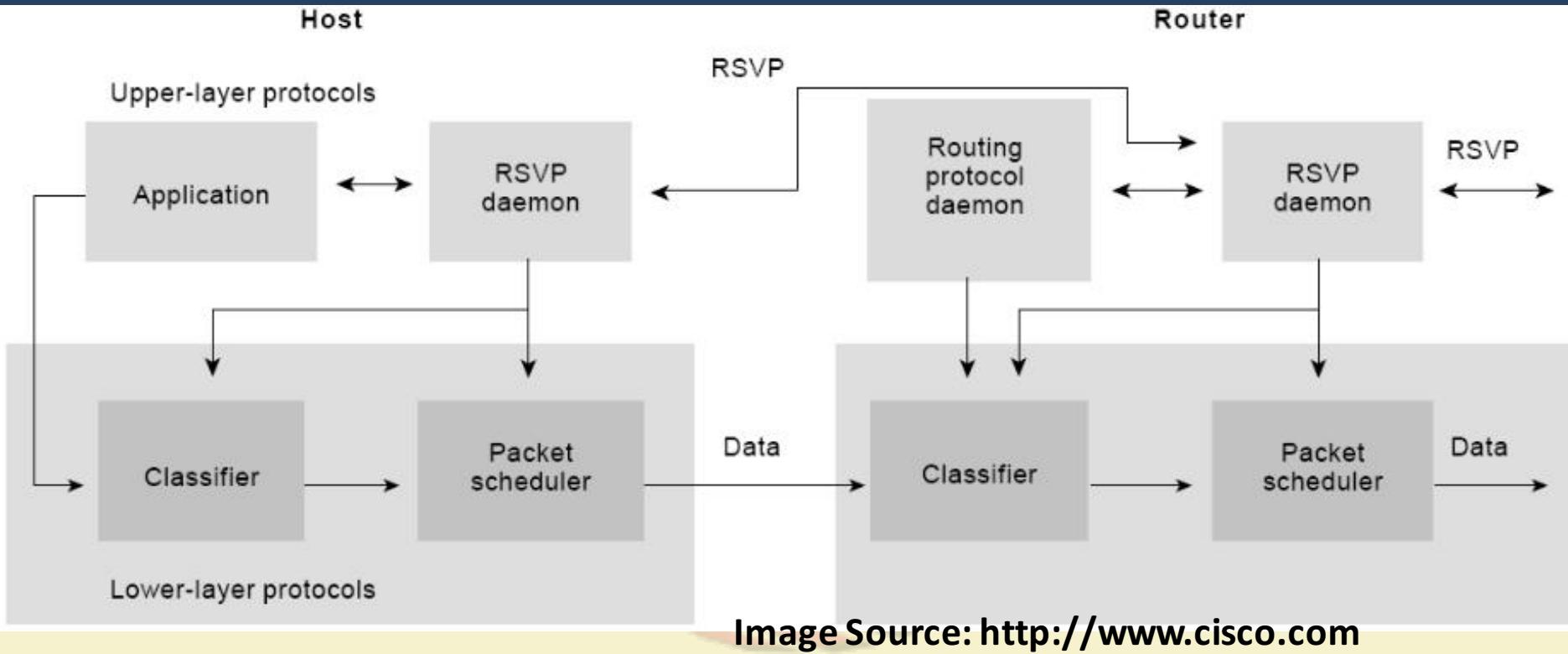


Image Source: <http://www.cisco.com>



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RSVP Terminologies

- Quality of Service is implemented for a particular data flow by a mechanism called “traffic control”
- **Packet Classifier:** Determines the QoS
- **Packet Scheduler:** Link layer dependent mechanism to determine which particular packets are forwarded
 - For each outgoing interface, the scheduler achieves the desired QoS.

RSVP Reservation Procedure

- During reservation setup, an RSVP QoS request is passed to two local decision modules: 1. Admission control 2. Policy control.
- **Admission Control:** Determines whether the node has sufficient available resources to supply the requested resources.
- **Policy Control:** Determines whether the user has administrative permission to make the reservation.

RSVP Reservation Procedure

- If both checks succeed, parameters are set in the packet classifier and in the link layer interface to obtain the desired QoS.
- If either checks fails, the RSVP program returns an error notification to the application process that generated the request.

RSVP Reservation Model

- An RSVP request consists of : **flowspec** together with a **filterspec**. This pair is called the “flow descriptor”.
- The **flowspec** specifies a desired QoS.
- The **filterspec** together with the session specification defines the set of data packets.

RSVP Reservation Model

- The **flowspec** is used to set parameters in the packet scheduler, while the **filterspec** is used to set parameters in the packet classifier
- The **flowspec** in a reservation request will generally include a service class and two sets of numeric parameters:
 - Rspec - defines the desired QoS,
 - Tspec - describes the data flow.

flowspec Structure

```
typedef struct flowspec {  
    ULONG TokenRate;  
    ULONG TokenBucketSize;  
    ULONG PeakBandwidth;  
    ULONG Latency;  
    ULONG DelayVariation;  
    SERVICETYPE ServiceType;  
    ULONG MaxSduSize;  
    ULONG MinimumPolicedSize;  
} FLOWSPEC, *PFLOWSPEC, *LPFLOWSPEC;
```



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Problems with RSVP

- The RSVP daemon needs to maintain per-flow states at intermediate routers.
- Use of per-flow state and per-flow processing raises scalability concerns for large network.



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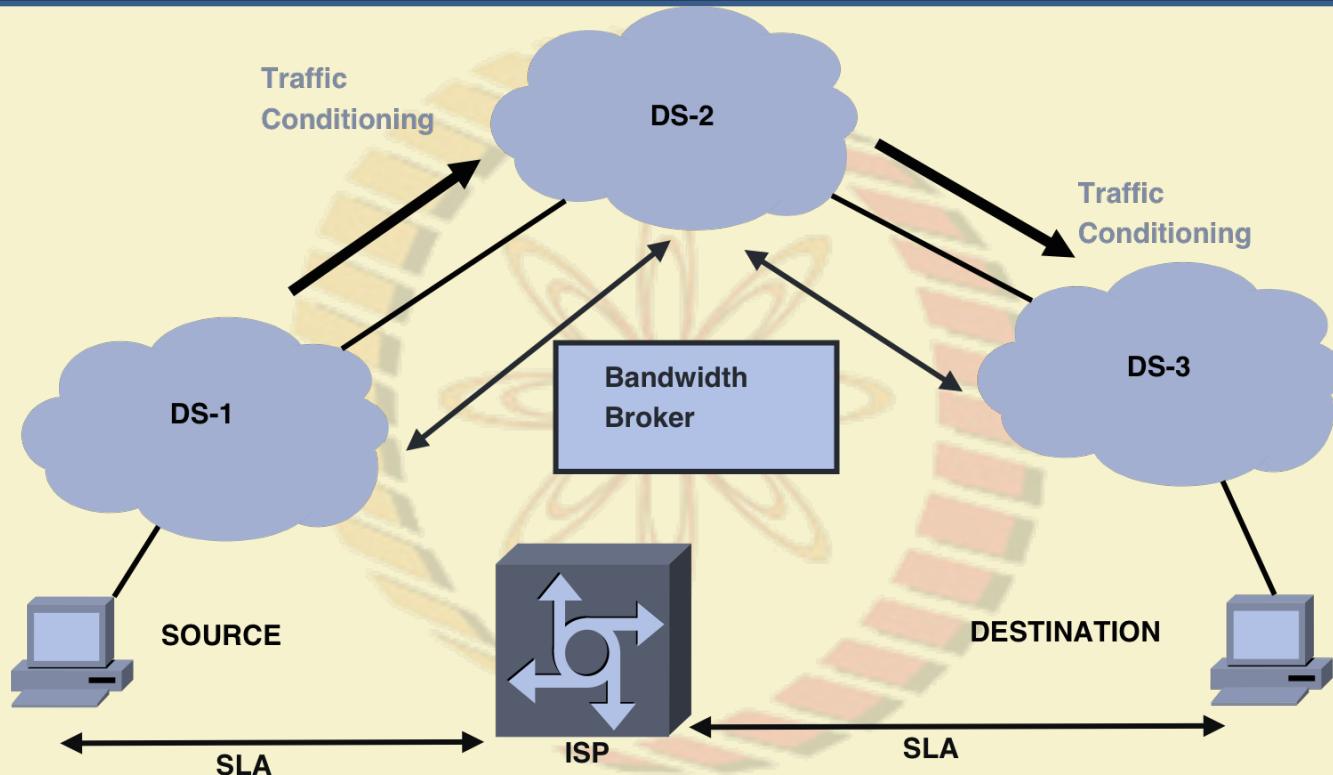


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Differentiated Service Architecture (DiffServ)

- Coarse-grained, class-based mechanism for traffic management
- **Packet classifier:** uses a 6-bit differentiated services code point (DSCP) in the 8-bit Differentiated Services field (DS field) in the IP header
- DiffServ-aware routers implement **per-hop behaviors** (PHBs) - defines the packet-forwarding properties associated with a class of traffic
- DiffServ recommends a standardized set of traffic classes.
- A group of routers that implement common, administratively defined DiffServ policies are referred to as a **DiffServ (DS) domain**.

DiffServ Architecture



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Bandwidth Broker

- A **Bandwidth Broker** is an agent that has some knowledge of an organization's priorities and policies and allocates QoS resources with respect to those policies. (RFC 2638)
- In order to achieve an end-to-end allocation of resources across separate domains, the Bandwidth Broker managing a domain will have to communicate with its adjacent peers, which allows end-to-end services to be constructed out of purely bilateral agreements.

SLA and TCA

- **Service Level Agreement (SLA):** A set of parameters and their values which together define the service offered to a traffic stream by a differentiated service (DS) domain.
- **Traffic Conditioning Agreement (TCA):** A set of parameters and their values which together specify a set of classifier rules and traffic profile.

Traffic Classification and Conditioning

- In a DS domain, the boundary node interconnects the current DS domain to other DS or non DS capable domains.
- The **classification and conditioning process** of a boundary node in a DS domain is responsible for mapping packets to a forwarding class supported in the network and ensuring that the traffic from a customer conforms to their SLAs.

Traffic Conditioning

- Traffic conditioning is a *set of control functions* that is applied to classified packets streams in order to enforce traffic conditioning agreements; which are made between customers and service providers.
 - have four components - **meter, marker, shapers and droppers**
 - a meter is used to measure the classified traffic stream against a traffic profile.
 - the state of the meter may then be used to enable a marking, shaping or dropping action.

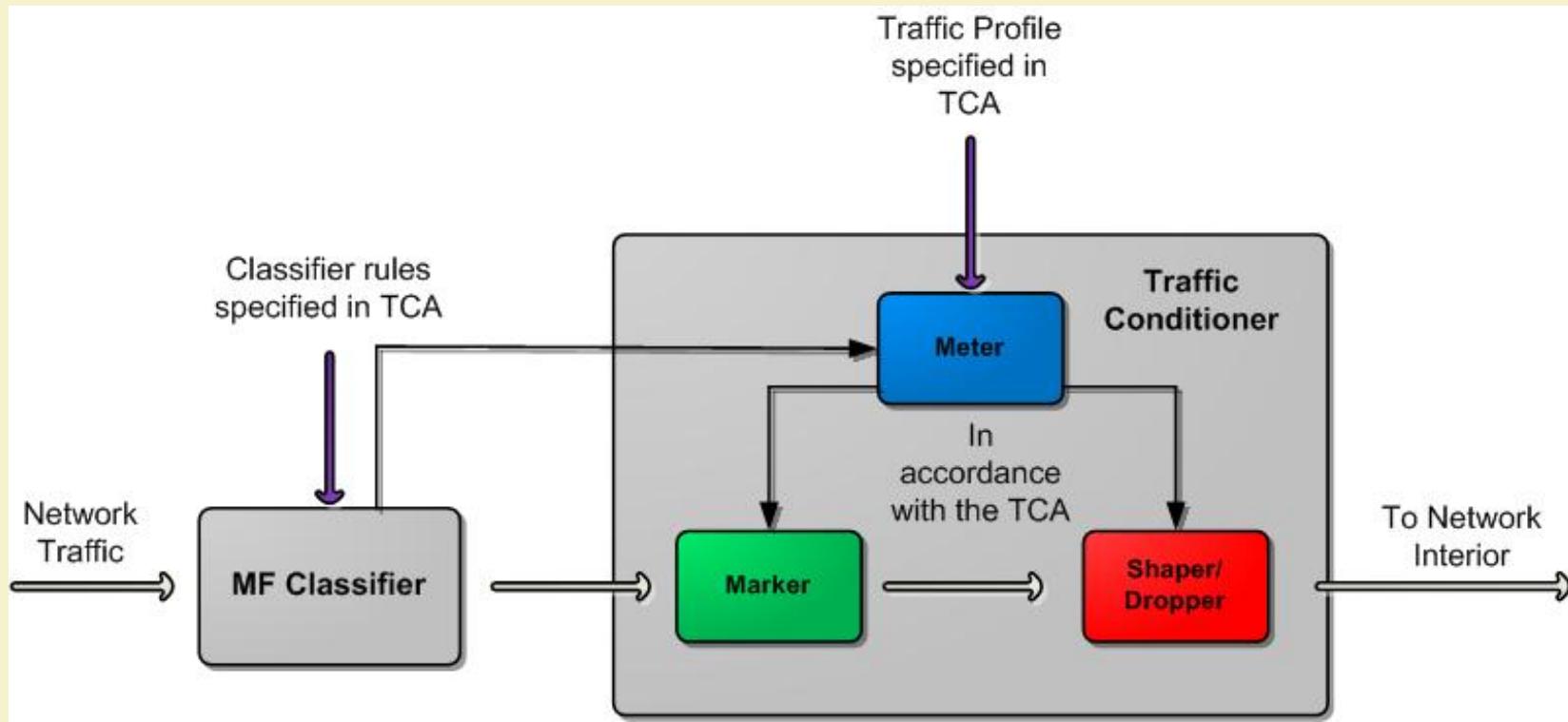


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Traffic Conditioner



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Classification and Marking - Per Hop Behaviors (PHB)

- **Default PHB** - for best-effort services
- **Expedited Forwarding (EF) PHB** - dedicated to low-loss, low-latency traffic
- **Assured Forwarding (AF) PHB** - gives assurance of delivery under prescribed conditions (different queuing behaviors)
- **Class Selector PHBs** - which maintain backward compatibility with the IP Precedence field.

Working Steps of a DS Domain

- The source (users) make a contract with the ISP for a specific SLA. Source sends request message to first hop router
- First hop router sends request to BB, which sends back either a accept or reject
- If accept, either source or first hop router will mark DSCP and start sending packets
- Edge router checks compliance with SLA and does policing. Excess packets are either discarded or marked as low priority to comply with the SLA
- Core routers will just look at DSCP and decide PHB

Further Readings

- Cisco White Papers on QoS -
 - <https://www.cisco.com/c/en/us/products/ios-nx-os-software/quality-of-service-qos/index.html>
 - http://docwiki.cisco.com/wiki/Quality_of_Service_Networking



thank you!



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