

# MA 214 - Introduction to Numerical Analysis

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## 1 Interpolation Theory

### 1.1 Introduction

- Given finite set of points, reconstructing the original curve is interpolation.
- There will be obviously infinitely many curve.

- Interpolation problem

Given  $n+1$  real distinct points:  $x_0, x_1, \dots, x_n$  and real numbers:  $y_0, y_1, \dots, y_n$

Find a function  $f : \mathbb{R} \rightarrow \mathbb{R}$  such that

$$f(x_i) = y_i \quad \text{for } i = 0, 1, \dots, n$$

Such a function is called **interpolant** and points  $x_i$  are called **interpolation points**.

We attempt to rebuild original function using polynomial functions. This is called polynomial interpolation and function is polynomial interpolant.

- A polynomial is function of the form

$$p(x) = a_0 + a_1x + \dots + a_nx^n$$

- $\mathbb{P}_n$  is the set of polynomials consisting of all polynomials of degree  $\leq n$

### 1.2 Polynomial Interpolation

**Theorem** (Joseph-Louis Lagrange Theorem). *Given  $n + 1$  data points with unique  $x_i$ s, then there exists a unique polynomial  $p_n \in \mathbb{P}_n$  such that*

$$p(x_i) = y_i \quad \text{for } i = 0, 1, \dots, n$$

*Proof.* (1) This can be shown by linear algebra. In a  $n$  degree polynomial, we substitute the points and get  $n + 1$  equations in  $n+1$  variables (coeff) and all the rows are unique (since  $x_0, x_1, \dots, x_n$  are unique), hence in  $AX = b$ ,  $|A| \neq 0$ . □

*Proof.* (2) Part 1: Uniqueness : If there is an interpolant, then the interpolant is unique  
Let there be 2 interpolants,  $p_n$  and  $q_n$  and let  $r(x) = p(x) - q(x)$ ,

$$r(x) = 0 \quad \text{for } i = 0, 1, \dots, n$$

This contradicts the fundamental theorem of Algebra. (A polynomial of degree  $n$  can have at most  $n$  real roots). Therefore

$$\begin{aligned} r(x) &= 0 \quad \forall x \in \mathbb{R} \\ p(x) &= q(x) \quad \forall x \in \mathbb{R} \end{aligned}$$

Part 2: Existence (construction):

Given  $n+1$  data points, build  $n+1$  Lagrange polynomials

$$L_k^n(x_i) = \begin{cases} 0 & \text{for } i \neq k \\ 1 & \text{for } i = k \end{cases}$$

$$L_k^n(x) = \frac{(x - x_0) \dots (x - x_{k-1})(x - x_{k+1}) \dots (x - x_n)}{(x_k - x_0) \dots (x_k - x_{k-1})(x_k - x_{k+1}) \dots (x_k - x_n)}$$

$$p(x) = \sum_{k=0}^n y_k L_k^n(x)$$

□

### 1.3 Closeness between functions

Given two continuous functions  $f, g : [a, b] \rightarrow \mathbb{R}$ , to evaluate how close the functions are consider the following

$$\max_{x \in [a, b]} |f(x) - g(x)|$$

### 1.4 Set of continuous Functions

$C[a, b]$  is the set of all continuous functions on  $[a, b]$

$C[a, b]$  is a infinite dimensional vector space

$$f, g \in C \implies f + g \in C \text{ and } \lambda f \in C$$

We define norm on  $C[a, b]$  as

$$\|f\| = \max_{x \in [a, b]} |f(x)|$$

$C^k[a, b]$  denotes the set of all functions which are continuously k-times differentiable

### 1.5 Polynomial Approximation and Error

**Theorem** (Weierstrass approximation Theorem). *Given a function  $f \in C[a, b]$  and given  $\epsilon > 0$ , there exists a polynomial  $p(x)$  such that,*

$$\|f(x) - p\| < \epsilon$$

**Using Langrange's recipe to approximate**

Take  $n + 1$  interpolation points in the  $[a, b]$  and collect the function values at all the points. We have  $n + 1$  data points. Using Lagrange polynomials, find the interpolant

**Theorem** (Error equation). *Let  $f \in C^k[a, b]$ ,  $x_0, x_1, \dots, x_n \in [a, b]$  and  $p \in \mathbb{P}_n$  be the interpolant using these points, then for all  $x$ , there exists a  $\zeta = \zeta(x) \in (a, b)$  such that*

$$f(x) - p(x) = \frac{1}{(n+1)!} f^{(n+1)}(\zeta) \prod_{k=0}^n (x - x_k)$$

**Note:** Here  $\zeta$  is dependent on the  $x$ , i.e, for every  $x$  you choose,  $\zeta$  generally changes.

*Proof.* Consider the function,

$$\psi(t) = (f(t) - p(t)) \prod_{k=0}^n (x - x_k) - (f(x) - p(x)) \prod_{k=0}^n (t - x_k)$$

This  $n + 2$  roots ( $n+1$  data points and  $x$ ), applying rolle theorem's gives us that  $\phi^{(1)}(t)$  has at least  $n+1$  roots. Applying like this repeatedly on its derivatives, we get that  $f^{(n+1)}$  has at least 1 root in  $[a, b]$ . Assuming the root to be  $\zeta$ . We have,

$$f(x) - p(x) = \frac{1}{(n+1)!} f^{(n+1)}(\zeta) \prod_{k=0}^n (x - x_k)$$

□

**Approximating the error:**

Taking norm on both the sides of error equation, we have,

$$\max_{x \in [a, b]} |f(x) - p(x)| = \frac{1}{(n+1)!} \|f^{(n+1)}(\zeta)\| \prod_{k=0}^n (x - x_k) \quad (1)$$

$$\max_{x \in [a, b]} |f(x) - p(x)| \leq \frac{1}{(n+1)!} \|f^{(n+1)}\| \max_{x \in [a, b]} \prod_{k=0}^n (x - x_k) \quad (2)$$

**Chebyshev interpolation points:**

$$x_k = \frac{a+b}{2} + \frac{b-a}{2} \cos\left(\frac{j\pi}{n}\right)$$

These points minimise  $\max_{x \in [a, b]} \prod_{k=0}^n (x - x_k)$  and therefore preferred over equally spaced points on real line. These points can be visualised as projections of equally spaced points on the arc of the semicircle with  $\frac{a+b}{2}$  as center and  $\frac{b-a}{2}$  as radius.

## 1.6 Another method for calculating interpolant

This is similar to the linear algebra method (given as proof(1) to Joseph-Louis Lagrange Theorem) for finding the interpolant.

Consider the polynomial

$$p(x) = a_0 + a_1(x - x_0) + a_2(x - x_0)(x - x_1) + \dots + a_n(x - x_0)\dots(x - x_{n-1})$$

Find the coefficients  $a_0, a_1, \dots, a_n$  by substituting the data points. On substituting  $x_0$ , we get  $a_0$ , again on substituting  $x_1$  and using  $a_0$ , we get  $a_1$  and so on.

## 1.7 Divided difference - recursion relation

- **Divided difference** : It is the coefficient of  $x_n$  in the interpolant  $p \in \mathbb{P}_n$  and denoted by  $f[x_0, x_1, \dots, x_n]$ .

Using Lagrange polynomials, we have

$$p(x) = \sum_{k=0}^n f(x_k) \prod_{j=0, j \neq k}^n \frac{x - x_j}{x_k - x_j}$$

So the divided difference is

$$f[x_0, x_1, \dots, x_n] = \sum_{k=0}^n f(x_k) \prod_{j=0, j \neq k}^n \frac{1}{x_k - x_j}$$

- interpolant  $p(x)$  of  $x_0, x_1, \dots, x_n$  can be calculated using interpolants  $a(x)$  and  $b(x)$  of  $x_1, x_2, \dots, x_n$  and  $x_0, x_1, \dots, x_{n-1}$  respectively as

$$p(x) = \frac{(x - x_0)a(x) - (x - x_n)b(x)}{x_n - x_0}$$

**Theorem** (Divided difference recursion theorem).

$$f[x_0, x_1, \dots, x_{m+1}] = \frac{[f[x_1, x_2, \dots, x_{m+1}] - f[x_0, x_1, \dots, x_m]]}{x_{m+1} - x_0}$$

*Proof.* Let  $p(x)$  be the interpolant for  $x_0, x_1, \dots, x_m$  and  $q(x)$  be the interpolant for  $x_1, x_2, \dots, x_{m+1}$ . Then,

$$L(x) = \frac{(x - x_0)q(x) + (x_{m+1} - x)p(x)}{x_{m+1} - x_0}$$

is an interpolant. Since, interpolant is unique, considering coeff of  $x_m$  we have,

$$f[x_0, x_1, \dots, x_m] = \frac{f[x_1, x_2, \dots, x_{m+1}] - f[x_0, x_1, \dots, x_m]}{x_{m+1} - x_0}$$

□

**Theorem** (Interpolant using divided differences). Suppose  $x_0, x_1, \dots, x_n$  be the data points. Then interpolant  $p \in \mathbb{P}_n$  is

$$p(x) = f[x_0] + f[x_0, x_1](x - x_0) + \dots + f[x_0, x_1, \dots, x_n] \prod_{j=0}^{n-1} (x - x_j)$$

*Proof.* We prove this by induction. Base case  $n = 0$  is trivially satisfied. Assume that this is satisfied for  $p_k$ ,

$$p_k(x) = f[x_0] + f[x_0, x_1](x - x_0) + \dots + f[x_0, x_1, \dots, x_k] \prod_{j=0}^{k-1} (x - x_j)$$

Consider the polynomial  $p_{k+1}(x) - p_k(x) \in \mathbb{P}_{k+1}$  which has  $x_0, x_1, \dots, x_k$  as roots. Hence,

$$p_{k+1}(x) - p_k(x) = c \prod_{j=0}^k (x - x_j)$$

Comparing leading coefficient on both sides, we have  $c = f[x_0, x_1, \dots, x_k]$ . Hence,

$$p_{k+1}(x) = f[x_0] + f[x_0, x_1](x - x_0) + \dots + f[x_0, x_1, \dots, x_{k+1}] \prod_{j=0}^k (x - x_j)$$

By PMI,

$$p(x) = f[x_0] + f[x_0, x_1](x - x_0) + \dots + f[x_0, x_1, \dots, x_n] \prod_{j=0}^{n-1} (x - x_j)$$

□

## 1.8 Time complexity of the algorithms

(Slide is useless)

- Lagrange's method -  $O(n^2)$  : Computing each Lagrange polynomial can be done in  $O(n)$  (Finding the coefficients given roots can be done in  $O(n \log^2(n))$ ). We need to do this  $n$  times. So,  $O(n^2)$ .
- Divided differences -  $O(n^2)$  : Summing operations in each stage -  $n + (n-1) + (n-2) + \dots + 1$ . Hence,  $O(n^2)$ .
- Divided difference can be considered better because we can extend from  $n$  to  $n+1$  and so on without discarding previous computation.

## 1.9 Weierstrass theorem consequences

- For Weierstrass theorem, take  $\epsilon_n = \frac{1}{n}$ . Then weierstrass theorem proves the existence of sequence of polynomials  $p^{(1)}, p^{(2)}, \dots$  such that

$$\lim_{n \rightarrow \infty} \|f - p^{(n)}\| = 0$$

- If  $f$  is not a polynomial, then

$$\lim_{n \rightarrow \infty} \text{degree of } p(n) = \infty$$

To prove (crude) this, assume that there exist a non-polynomial function  $f$  such that  $\|f - p^{(n)}\| = 0$ , then  $f = p^{(n)}$  which contradicts the assumption that  $f$  is not a polynomial.

## 1.10 Spline Interpolation

- **Piece wise polynomial:**  $\phi \in C[a, b]$  is a piecewise polynomial function, if there exists  $a = x_0 < x_1 < \dots < x_n = b$  such that  $\phi \in \mathbb{P}_m$  when  $x \in [x_i, x_{i+1}]$  for all  $i = 0, 1, \dots, n$  and some  $m > 0$ .
- Piece wise polynomial  $\phi$  need not be polynomial in whole domain.
- Splines interpolation for  $f \in C[a, b]$ 
  - Pick some data points  $x_0, x_1, \dots, x_n$  such that  $a = x_0 < x_1 < \dots < x_n = b$
  - Fix  $m \leq n$
  - Build  $\phi$  in each subinterval  $[x_i, x_{i+1}]$  using the following conditions:

$$\phi(x_i) = f_i \quad \text{for } i = 0, 1, \dots, n$$

$$\lim_{h \rightarrow 0+} \phi(x_i - h) = \lim_{h \rightarrow 0+} \phi(x_i + h) \quad \text{for } i = 1, 2, \dots, n-1$$

$$\lim_{h \rightarrow 0+} \frac{d\phi(x_i - h)}{dx} = \lim_{h \rightarrow 0+} \frac{d\phi(x_i + h)}{dx} \quad \text{for } i = 1, 2, \dots, n-1$$

$$\lim_{h \rightarrow 0+} \frac{d^2\phi(x_i - h)}{dx^2} = \lim_{h \rightarrow 0+} \frac{d^2\phi(x_i + h)}{dx^2} \quad \text{for } i = 1, 2, \dots, n-1$$

$\vdots$

$$\lim_{h \rightarrow 0+} \frac{d^{m-1}\phi(x_i - h)}{dx^{m-1}} = \lim_{h \rightarrow 0+} \frac{d^{m-1}\phi(x_i + h)}{dx^{m-1}} \quad \text{for } i = 1, 2, \dots, n-1$$

- We have  $(n+1) + m(n-1) = n(m+1) - (m-1)$  conditions. We need  $m-1$  more conditions.

**Theorem** (Error - Linear Splines). Let  $f \in C^2[a, b]$  and  $s_L(x)$  be the interpolating **linear spline** at  $(n+1)$  knots  $a = x_0 < x_1 < \dots < x_n = b$  and let  $h$  be the maximum subinterval length, then

$$\|f - s_L\| \leq \frac{h^2}{8} \|f''\|$$

*Proof.* Consider the interval  $[x_i, x_{i+1}]$ ,  $\phi(x)$  is the interpolating polynomial in this interval, using the error equation for interpolating polynomials,

$$f(x) - s_L(x) = \frac{1}{2}f''(\zeta)(x - x_i)(x - x_{i+1})$$

Taking absolute value on both the sides,

$$\begin{aligned} |f(x) - s_L(x)| &= \frac{1}{2}|f''(\zeta)|(x - x_i)(x - x_{i+1}) \\ &\leq \frac{1}{2}\|f''\|\frac{h_i^2}{4} \quad \text{where } h_i = \frac{x_{i+1} - x_i}{2} \\ &\leq \frac{h_i^2}{8}\|f''\| \end{aligned}$$

Considering  $h = \max(h_i)$ , then for  $x \in [a, b]$

$$\max_{x \in [a, b]} |f(x) - s_L(x)| \leq \frac{1}{8}h^2\|f''\|$$

□