

Robust and Data-Driven Markov Decision Processes

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Markov decision process

Tuple $(\mathcal{S}, \mathcal{A}, q, p, r, \lambda)$ where

- $\mathcal{S} = \{1, \dots, S\}$ is the (finite) **state space**;
- $\mathcal{A} = \{1, \dots, A\}$ is the (finite) **action space**;
- $q = (q_1, \dots, q_S) \in \Delta(\mathcal{S})$ is the **initial state distribution**;
- $p : \mathcal{S} \times \mathcal{A} \rightarrow \Delta(\mathcal{S})$ is the **transition kernel** with elements $p(s' | s, a)$;
- $r : \mathcal{S} \times \mathcal{A} \rightarrow \mathbb{R}$ are the **expected one-step rewards**;
- $\lambda \in (0, 1)$ is the **discount factor**.

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- $r : \mathcal{S} \times \mathcal{A} \rightarrow \mathbb{R}$ is the reward function;
- $\lambda \in (0, 1]$ is the discount factor.

Objective

find policy π that maximizes the expected total discounted rewards:

$$\underset{\pi \in \Pi}{\text{maximize}} \quad \mathbb{E}_p \left[\sum_{t=1}^{\infty} \lambda^{t-1} \cdot r(s_t, \pi_t[s_t]) \right]$$

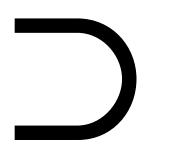
MDPs: Key Results

- Stationary deterministic policies are **optimal**:

$$\pi = \{\pi_t\}_{t=1}^{\infty}$$

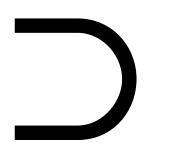
$$\pi_t : (\mathcal{S} \times \mathcal{A})^{t-1} \times \mathcal{S} \rightarrow \Delta(\mathcal{A})$$

history-dependent, randomized



$$\pi : \mathcal{S} \rightarrow \Delta(\mathcal{A})$$

*stationary,
randomized*



$$\pi : \mathcal{S} \rightarrow \mathcal{A}$$

*stationary,
deterministic*

- Stationary deterministic policies are **optimal**.
- Discounted rewards of a **fixed policy**

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- Discounted rewards of a **fixed policy**

$$v^\pi(s) = \mathbb{E}_p \left[\sum_{t=1}^{\infty} \lambda^{t-1} \cdot r(s_t, \pi[s_t]) \mid s_1 = s \right]$$

- Stationary deterministic policies are optimal.
- Discounted rewards of a fixed policy satisfy linear equations:

$$v^\pi(s) = r(s, \pi[s]) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, \pi[s]) \cdot v^\pi(s')$$

- Stationary deterministic policies are optimal.
- Discounted rewards of a fixed policy satisfy linear equations.
- Discounted rewards of an optimal policy satisfy nonlinear equations:

$$v^*(s) = \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^*(s') \right\}$$

- Stationary deterministic policies are **optimal**.
- Discounted rewards of a **fixed policy** satisfy **linear equations**.
- Discounted rewards of an **optimal policy** satisfy **nonlinear equations**:

$$v^*(s) = \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^*(s') \right\}$$

- v^* -greedy policy is **optimal**:

$$\pi^*(s) \in \arg \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^*(s') \right\}$$

MDPs: Solution Methods

- **Value iteration:**

$$v^\star(s) = \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^\star(s') \right\}$$

- **Value iteration:**

$$\cancel{v^*(s)} = \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot \cancel{v^*(s')} \right\}$$

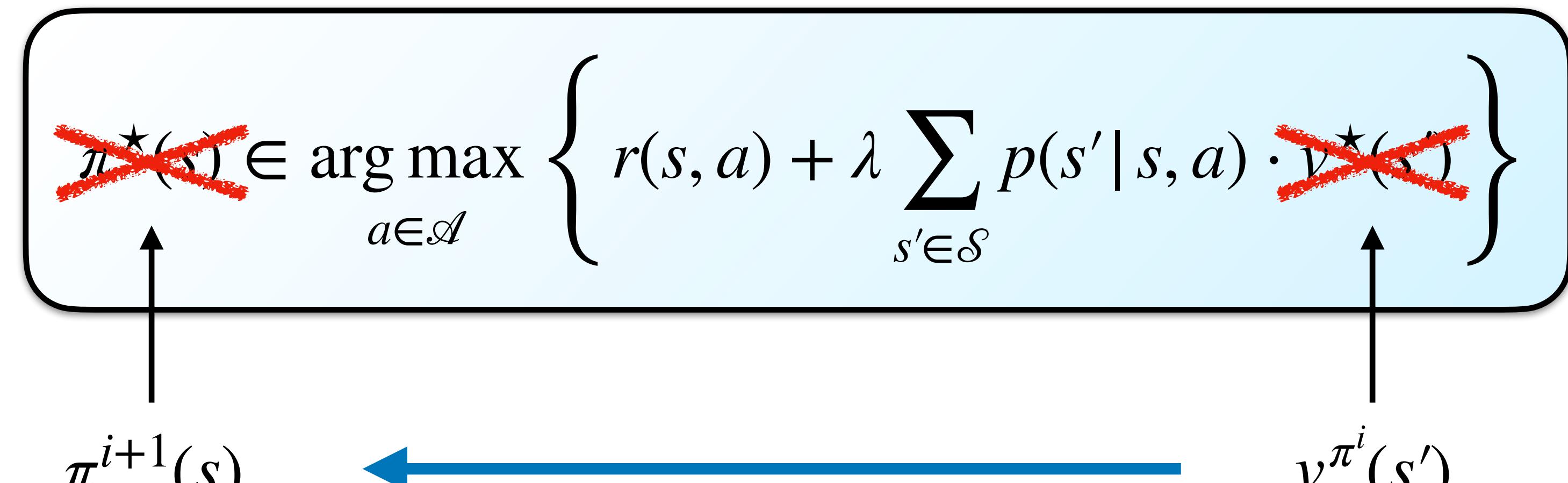
$v^{i+1}(s)$ $\xleftarrow{\hspace{10em}}$ $v^i(s')$

*Starting from any $v^0 \in \mathbb{R}^S$, converges at **linear rate** to v^* .*

- **Value iteration.**
- **(Modified) Policy iteration:**

$$\pi^\star(s) \in \arg \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^\star(s') \right\}$$

- **Value iteration.**
- **(Modified) Policy iteration:**



Policy improvement

Policy evaluation

Under suitable conditions, converges at superlinear rate to v^ .
Converges to (ϵ -)optimal policy in finitely many iterations.*

- **Value iteration.**
- **(Modified) Policy iteration.**
- **Linear programming:**

minimize
 $v \in \mathbb{R}^S$

$$\sum_{s \in \mathcal{S}} q(s) \cdot v(s)$$

subject to

$$v(s) = \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \quad \forall s \in \mathcal{S}$$

- **Value iteration.**
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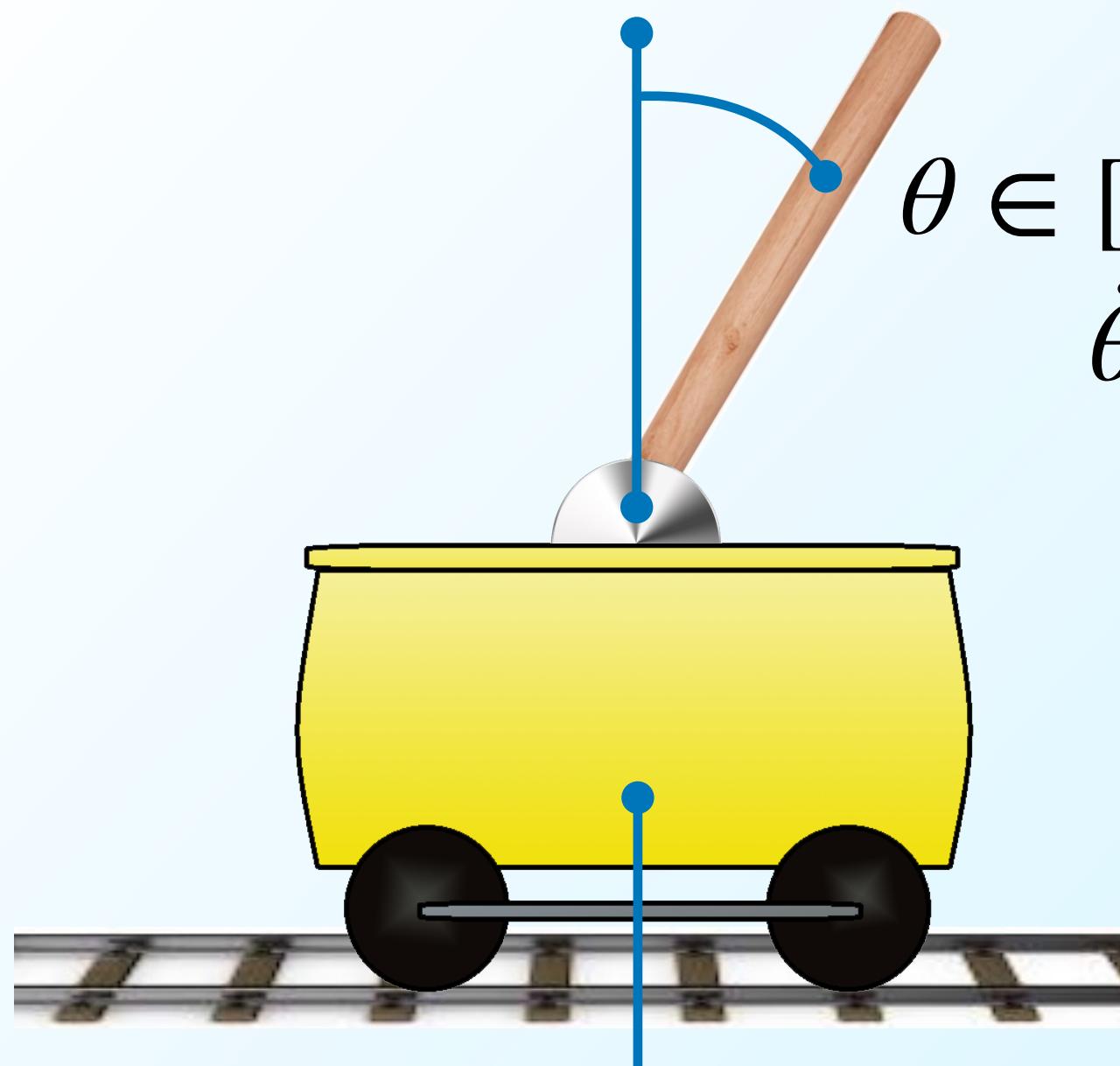
$$v(s) \geq \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \quad \forall s \in \mathcal{S}$$

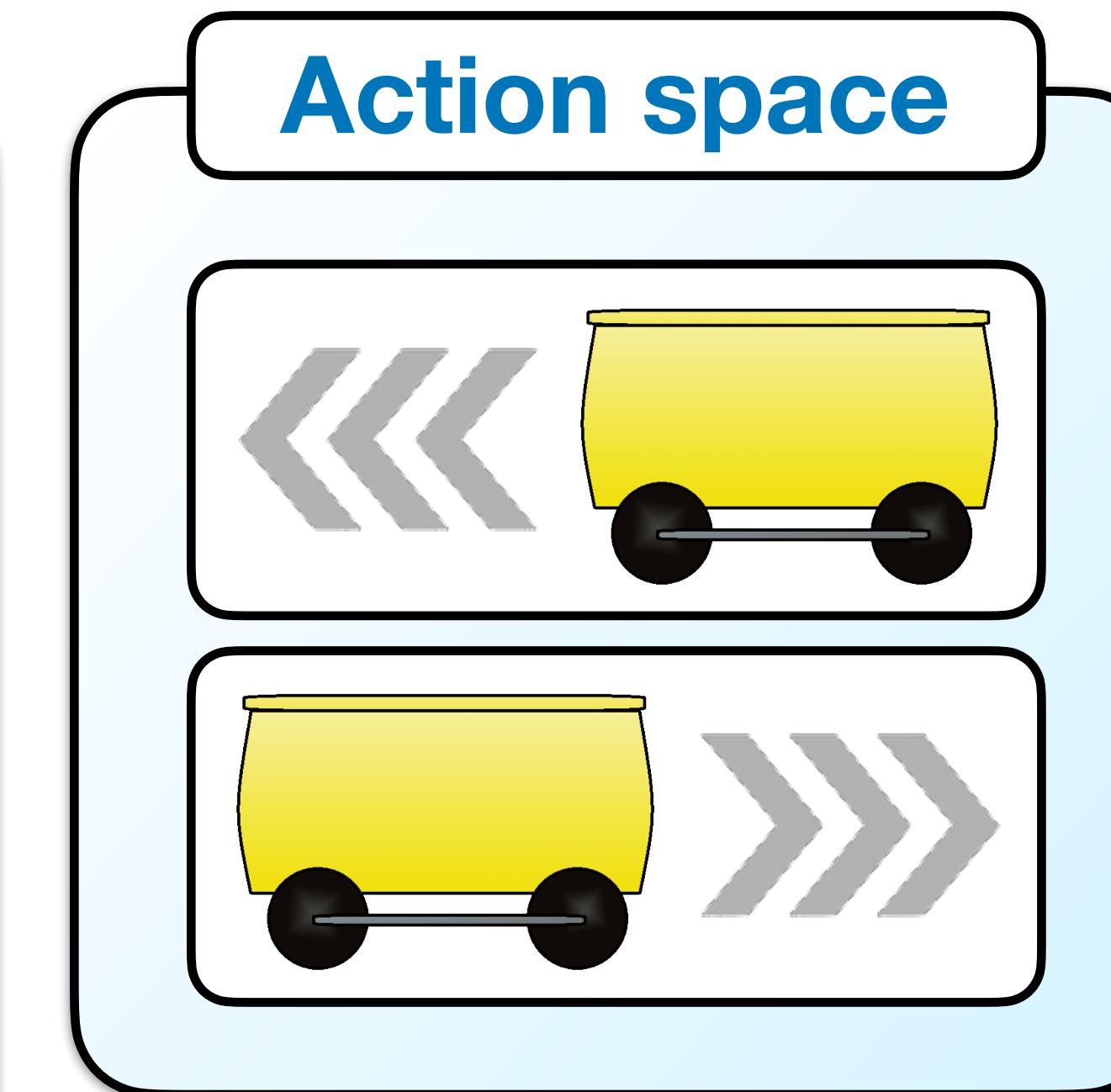
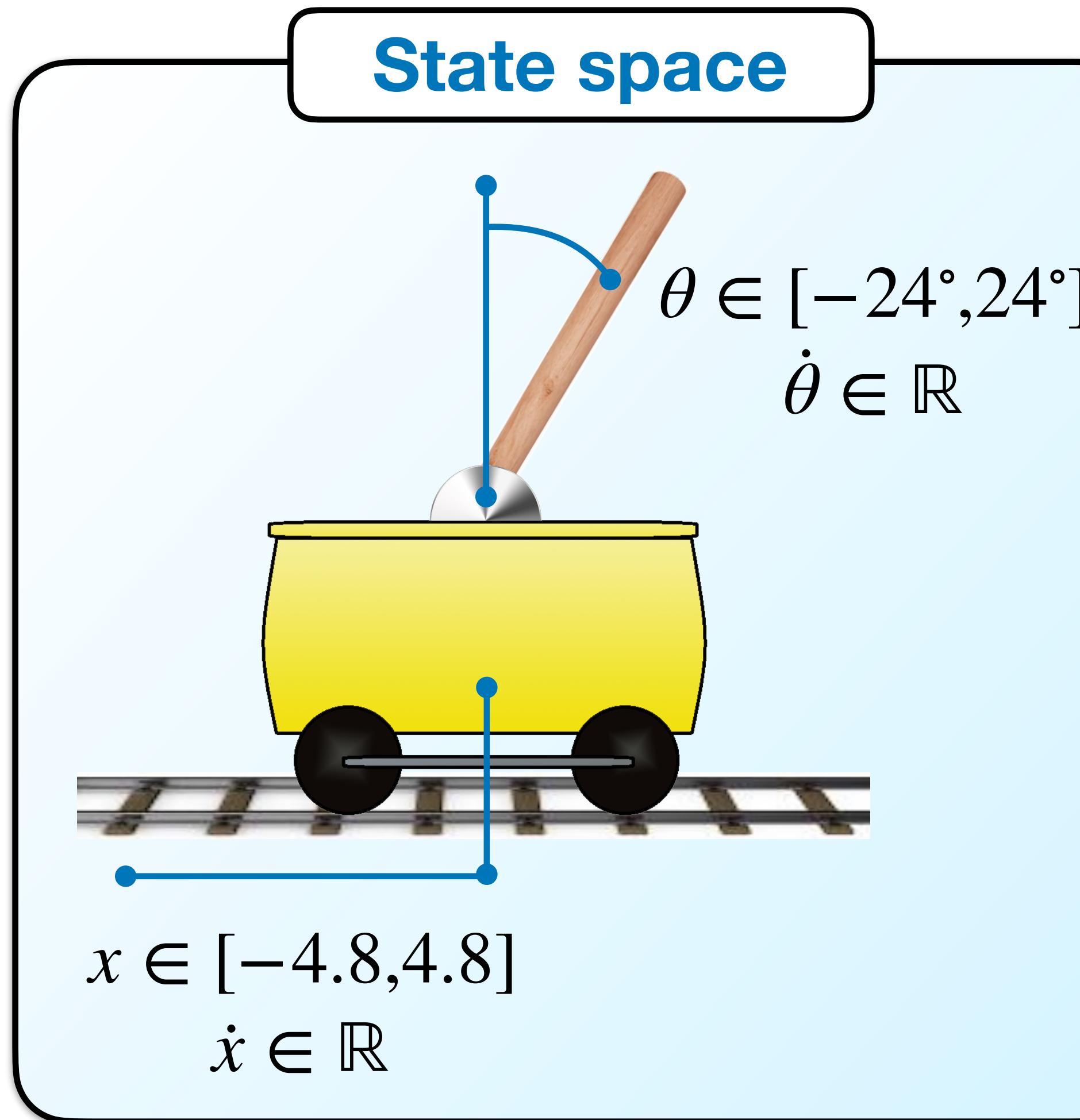
- **Value iteration.**
- **(Modified) Policy iteration.**
- **Linear programming:**

$$\begin{aligned} & \underset{v \in \mathbb{R}^S}{\text{minimize}} && \sum_{s \in \mathcal{S}} q(s) \cdot v(s) \\ & \text{subject to} && v(s) \geq r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \quad \forall s \in \mathcal{S}, \forall a \in \mathcal{A} \end{aligned}$$

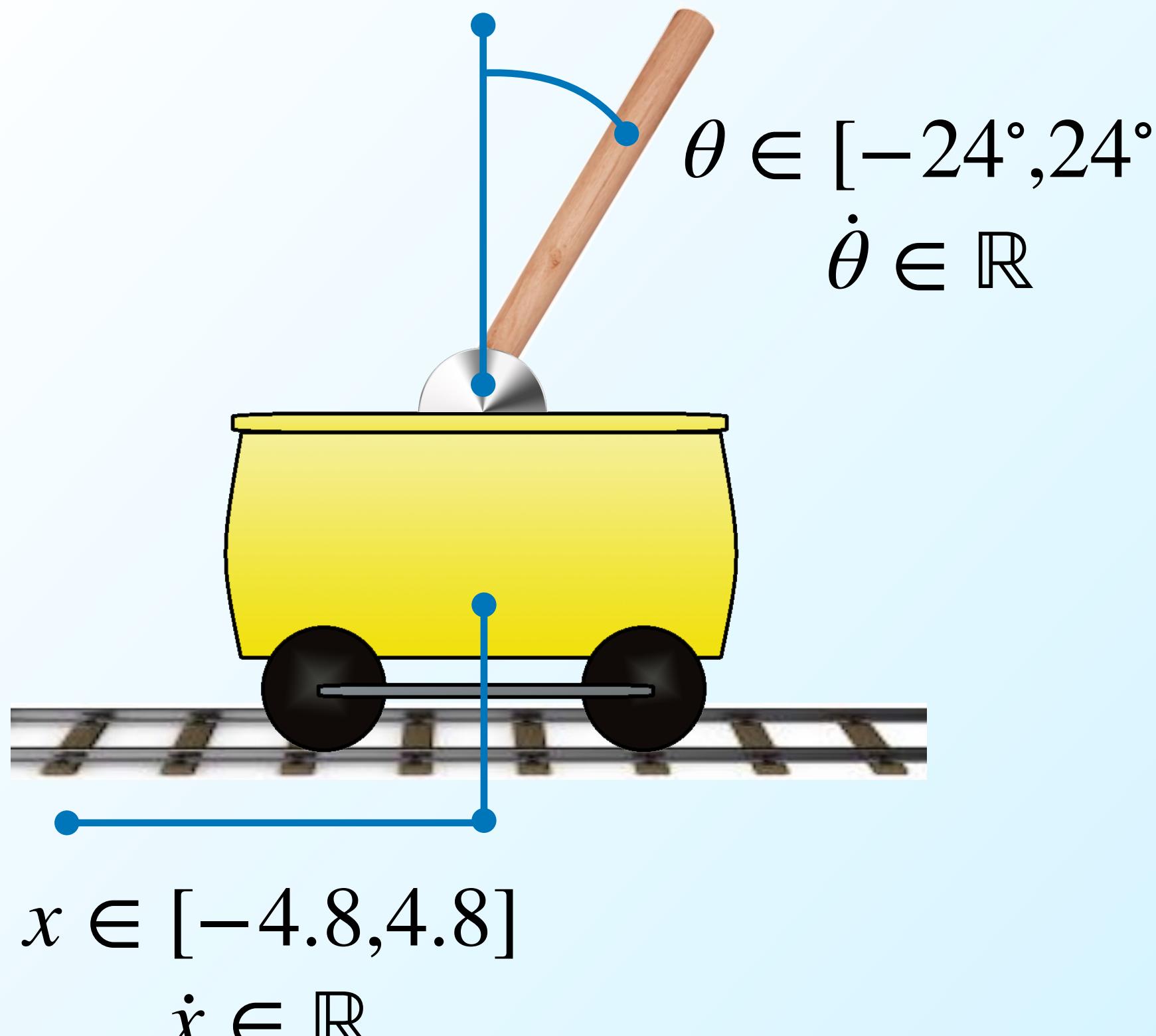
Optimally solved in polynomial time with standard solvers.

State space

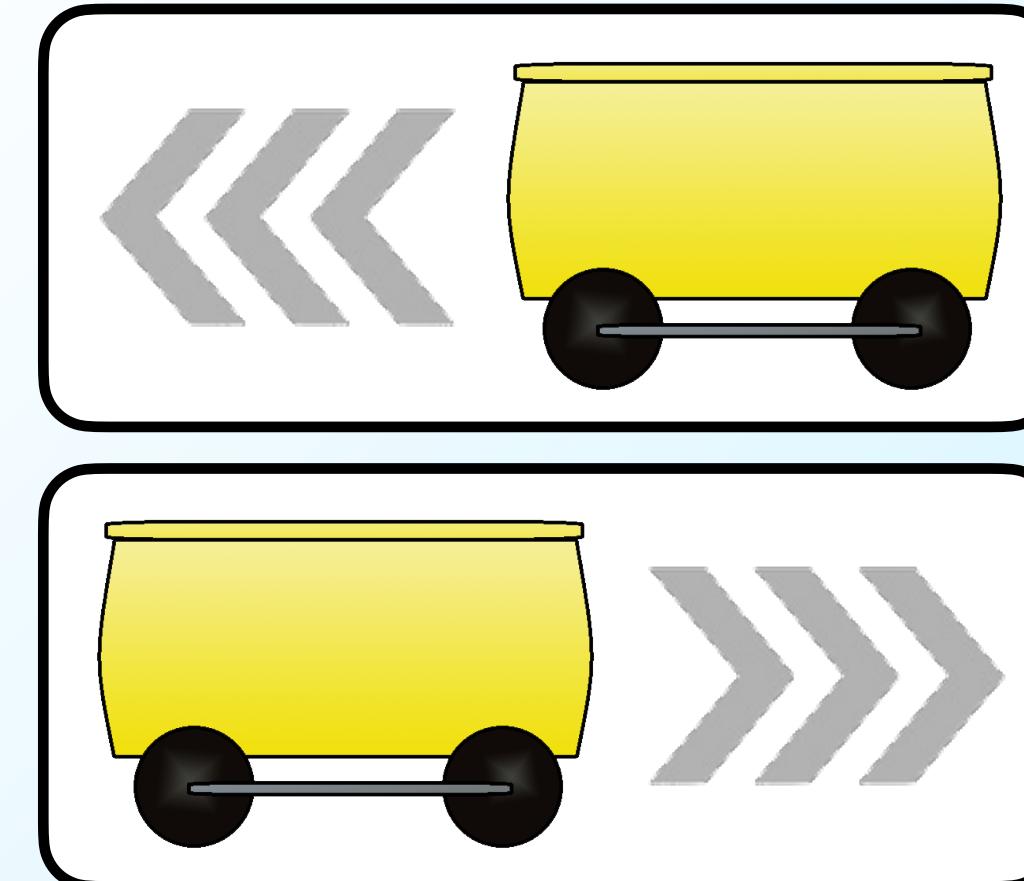




State space



Action space

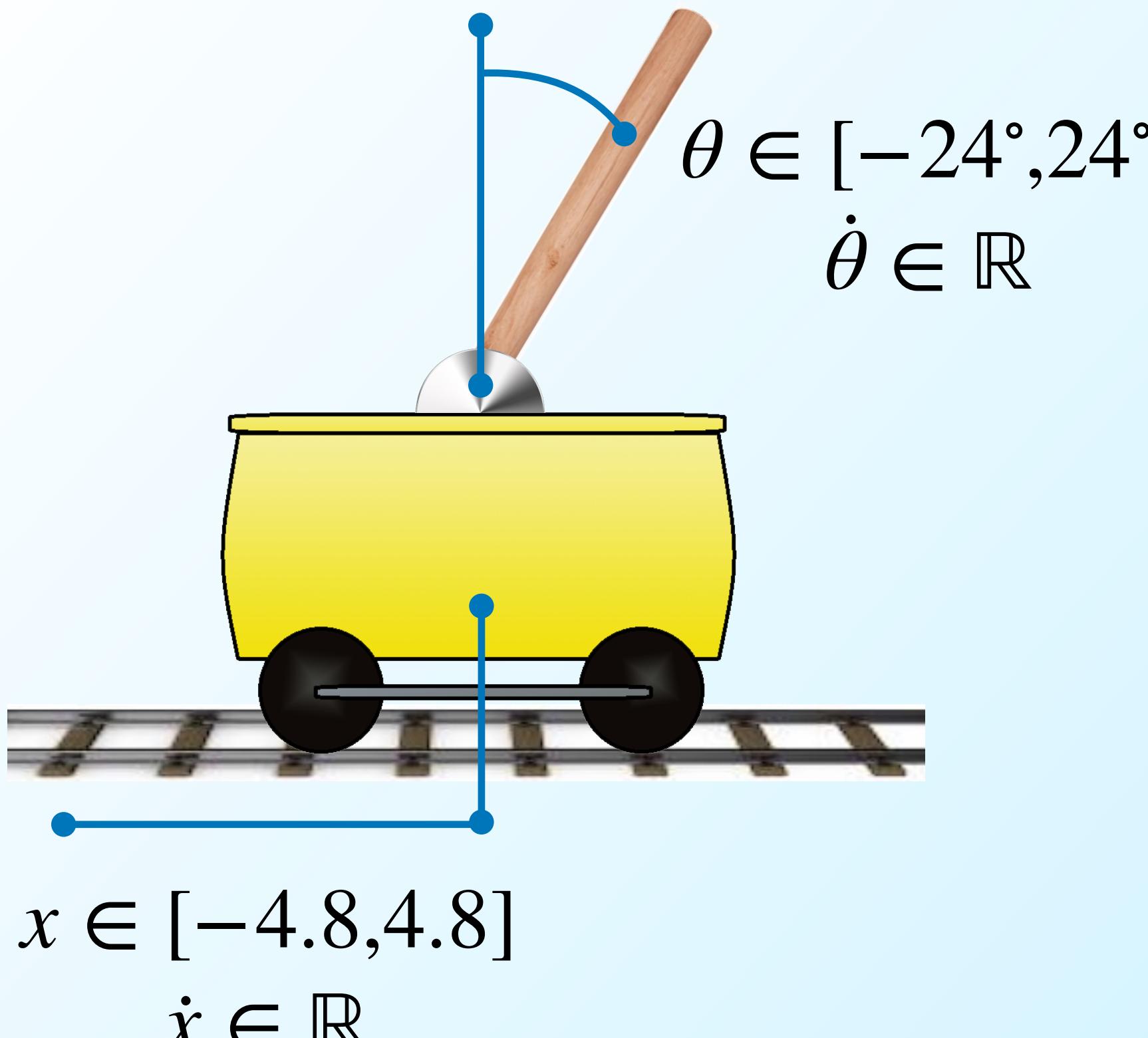


Initial state

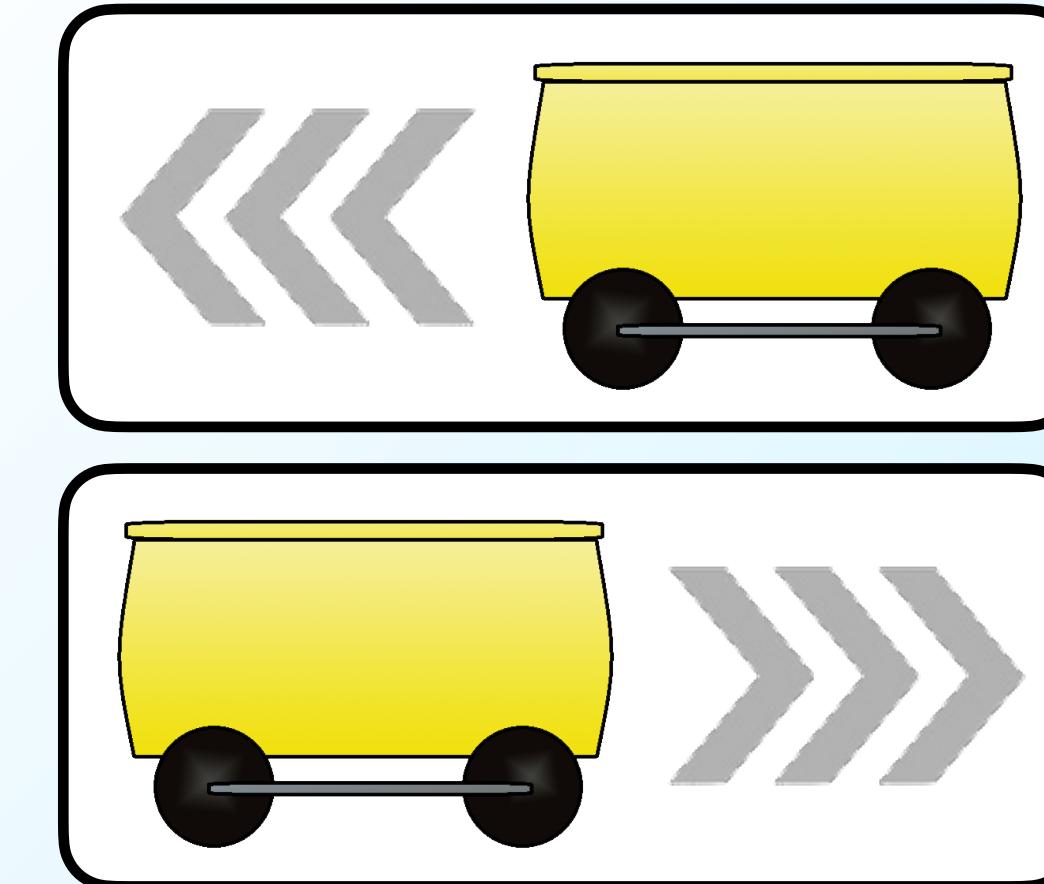
$$x, \dot{x}, \theta, \dot{\theta} \sim \mathcal{U}[-0.05, 0.05]$$

Cart Pole Example

State space



Action space



Transitions

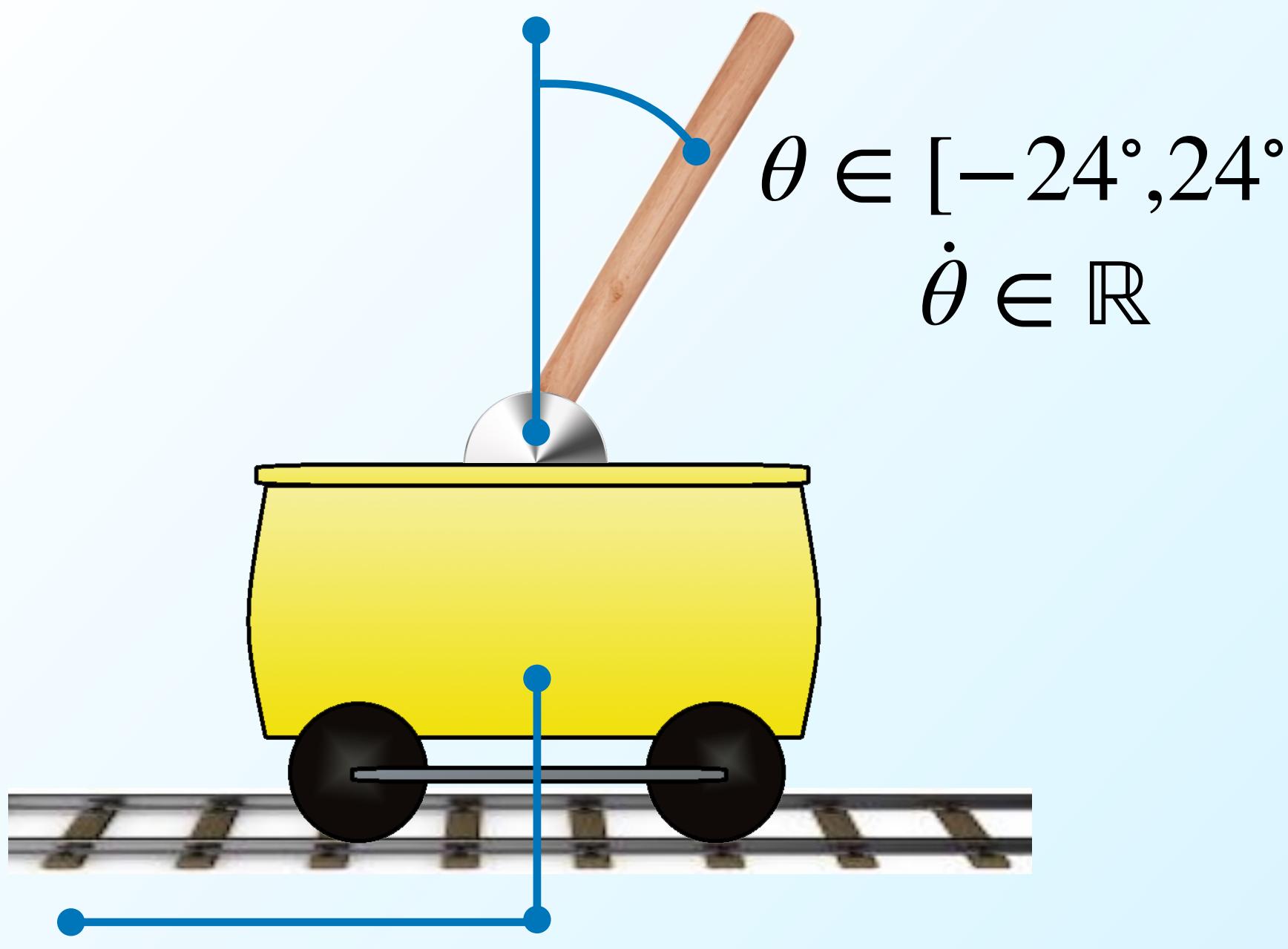
- deterministic via laws of mechanics
- terminate if
 - $x \notin [-2.4, 2.4]$
 - or $\theta \notin [-12^\circ, 12^\circ]$

Initial state

$$x, \dot{x}, \theta, \dot{\theta} \sim \mathcal{U}[-0.05, 0.05]$$

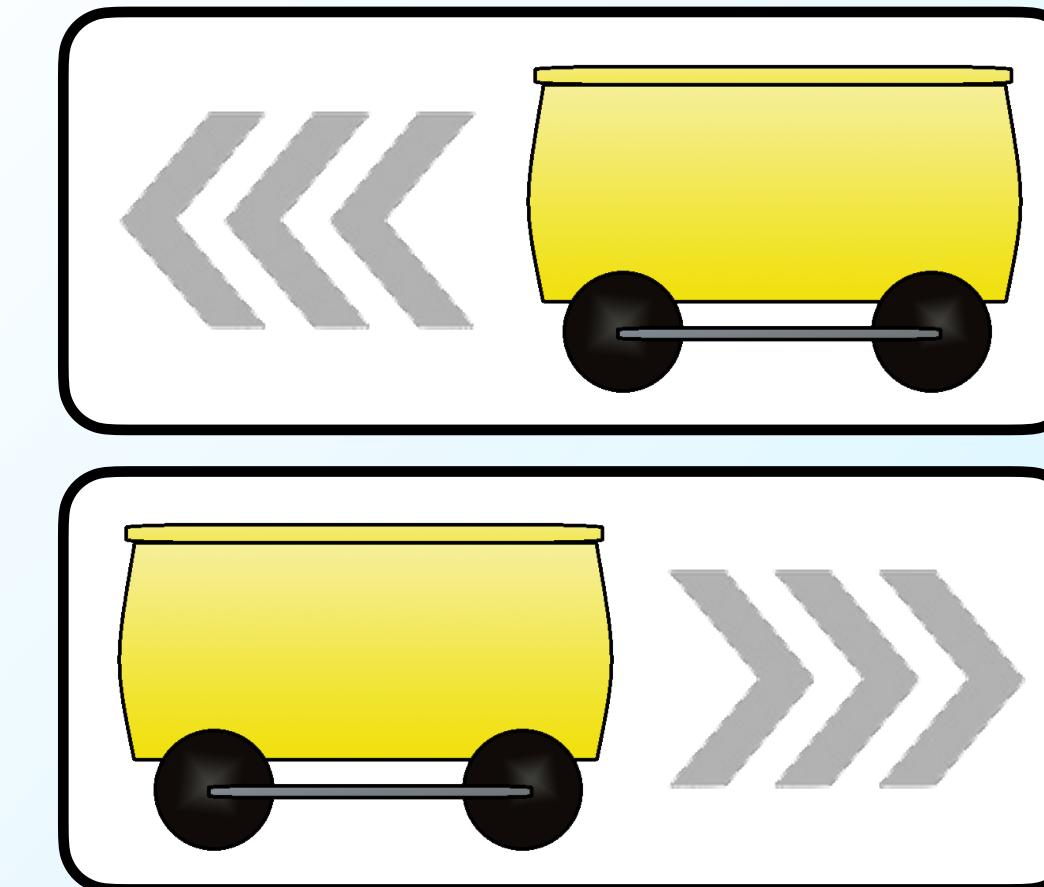
Cart Pole Example

State space



$$x \in [-4.8, 4.8]$$
$$\dot{x} \in \mathbb{R}$$

Action space



Initial state

$$x, \dot{x}, \theta, \dot{\theta} \sim \mathcal{U}[-0.05, 0.05]$$

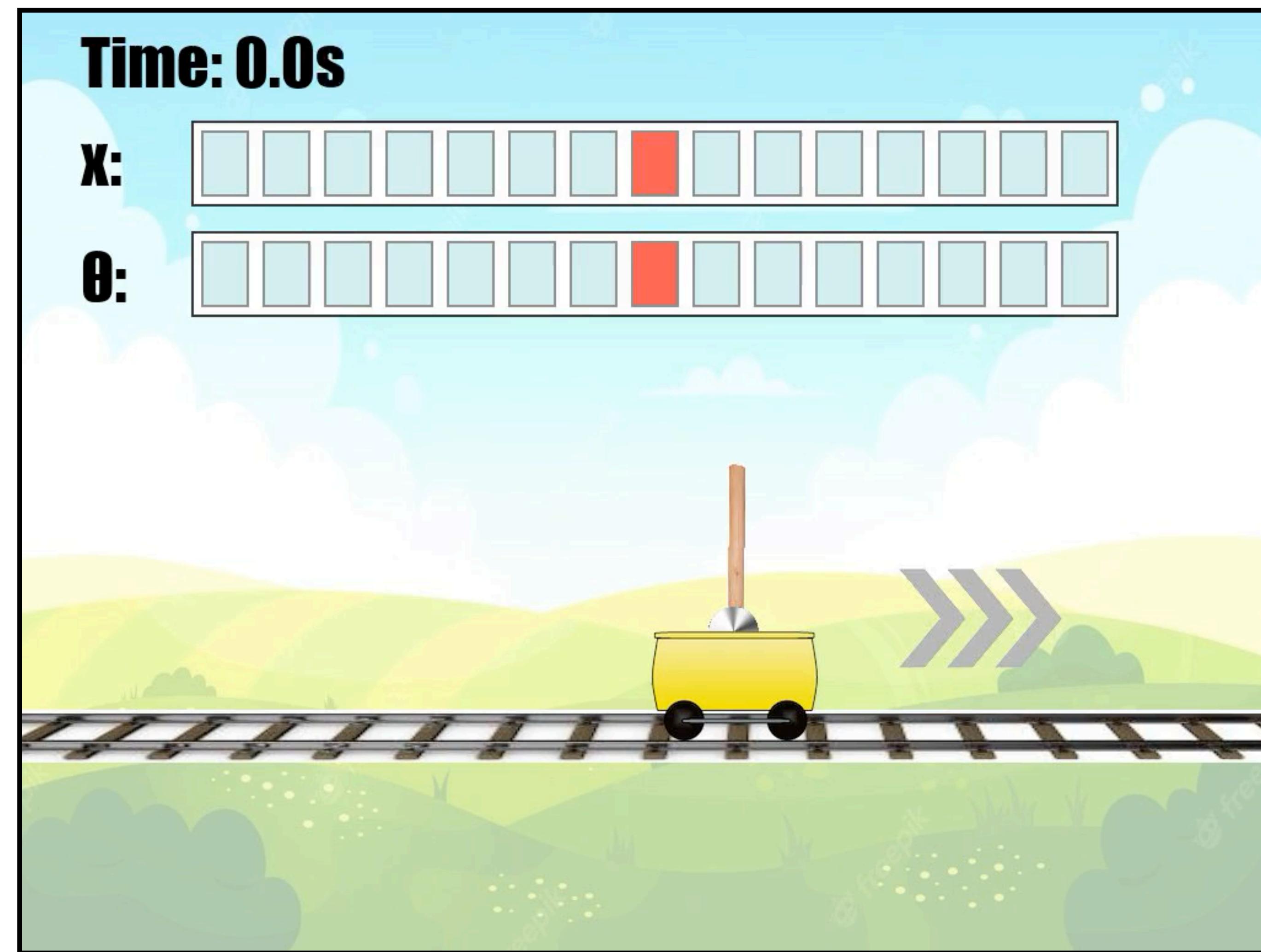
Transitions

- deterministic via laws of mechanics
- terminate if
$$x \notin [-2.4, 2.4]$$
or
$$\theta \notin [-12^\circ, 12^\circ]$$

Rewards

+1/non-terminated time step

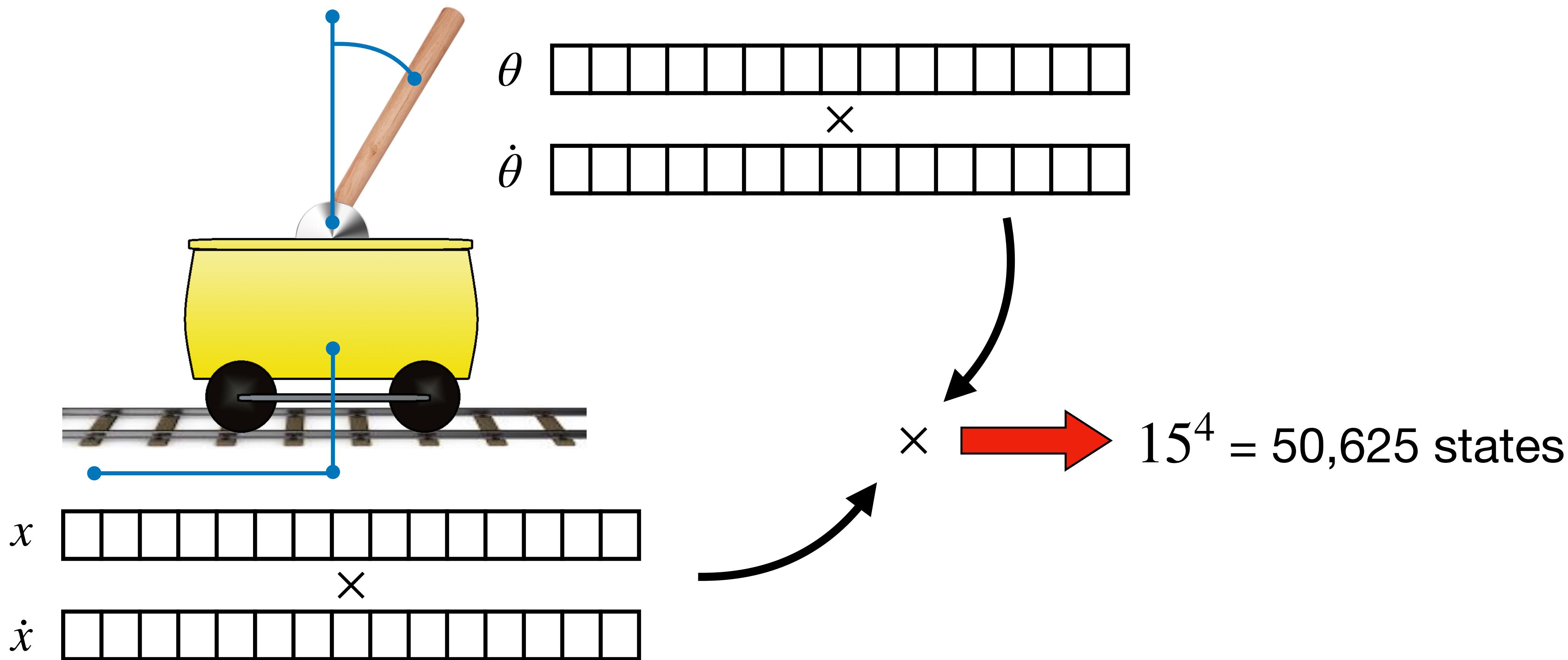
Cart Pole Example



Ambiguity and Robust MDPs

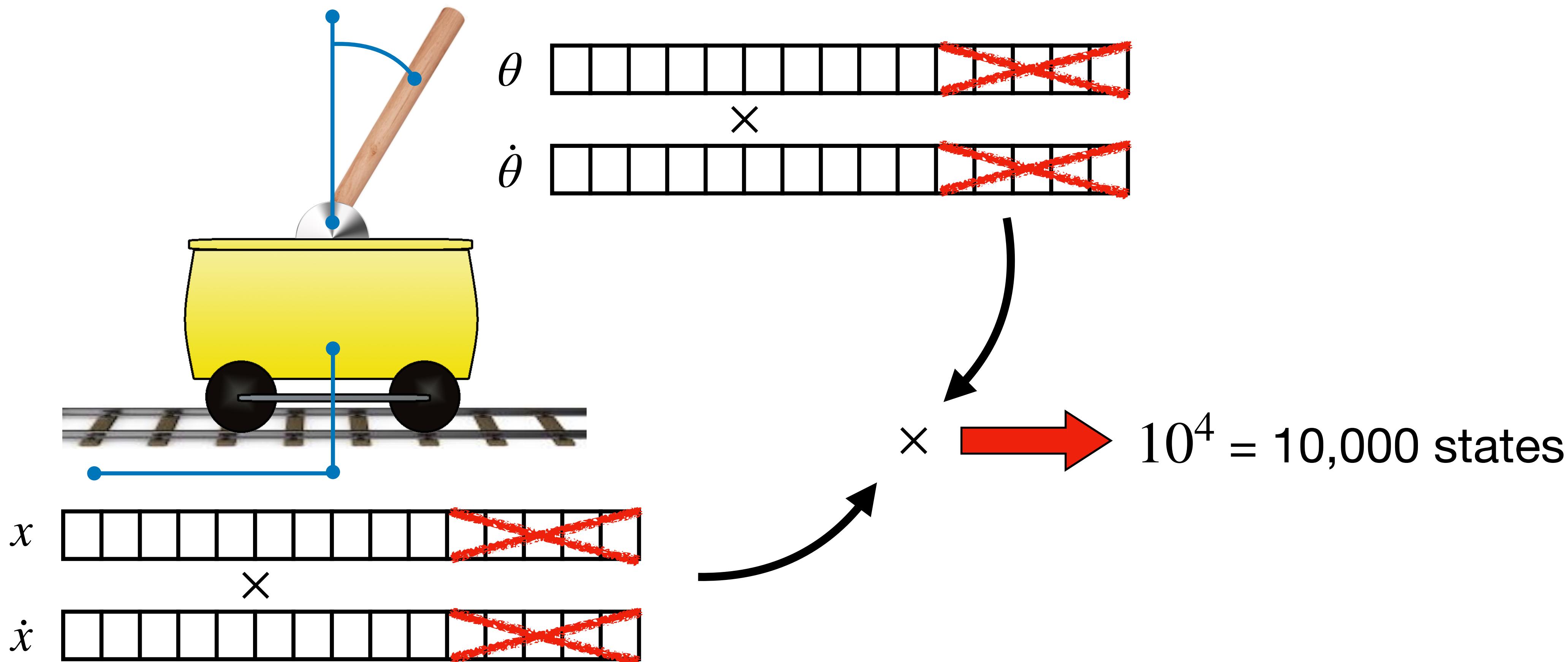
Two common sources of **ambiguity**:

- **Modelling errors:** 32.67 secs/run



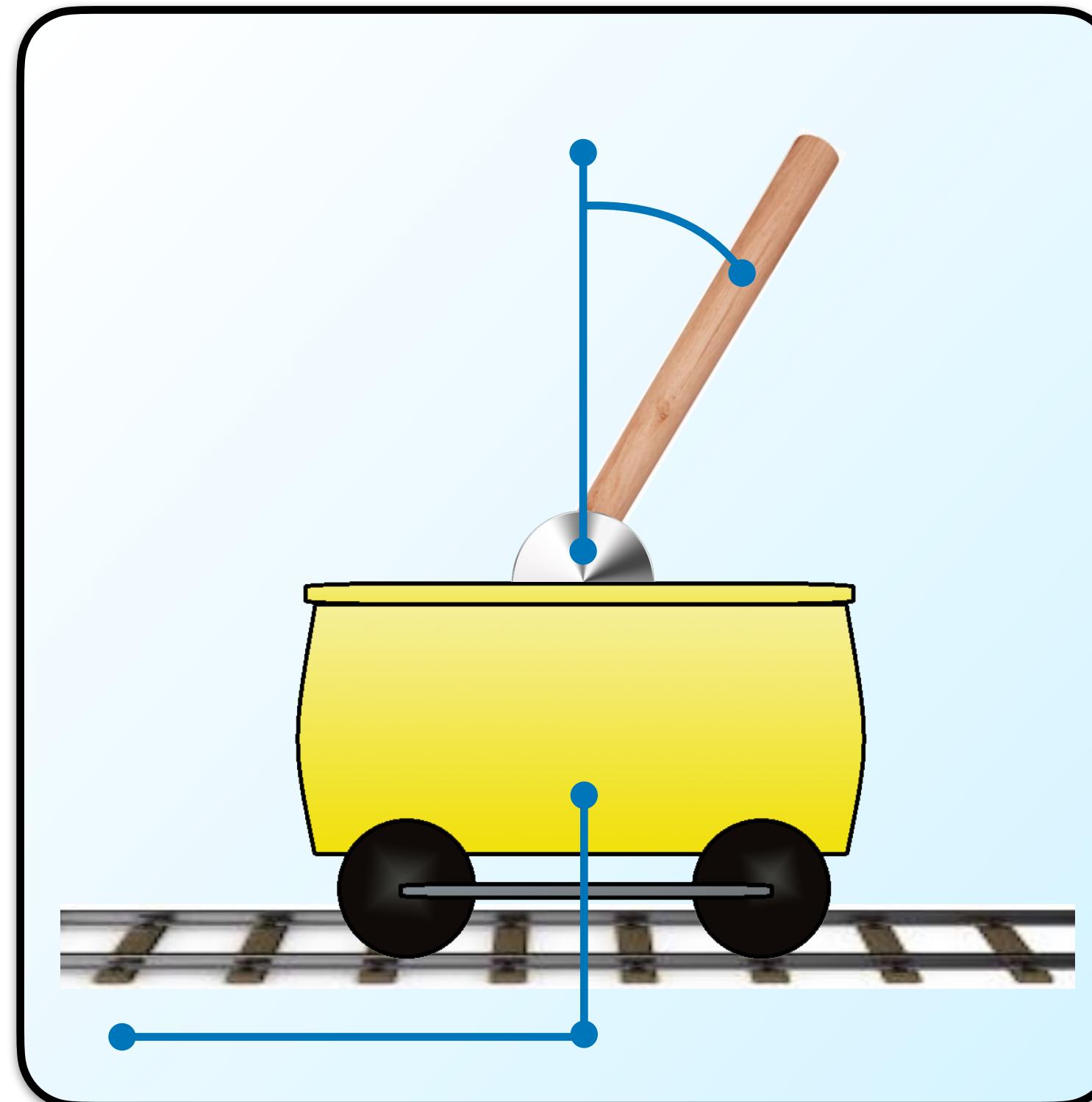
Two common sources of **ambiguity**:

- **Modelling errors:** 32.67 secs/run → 2.45 secs/run



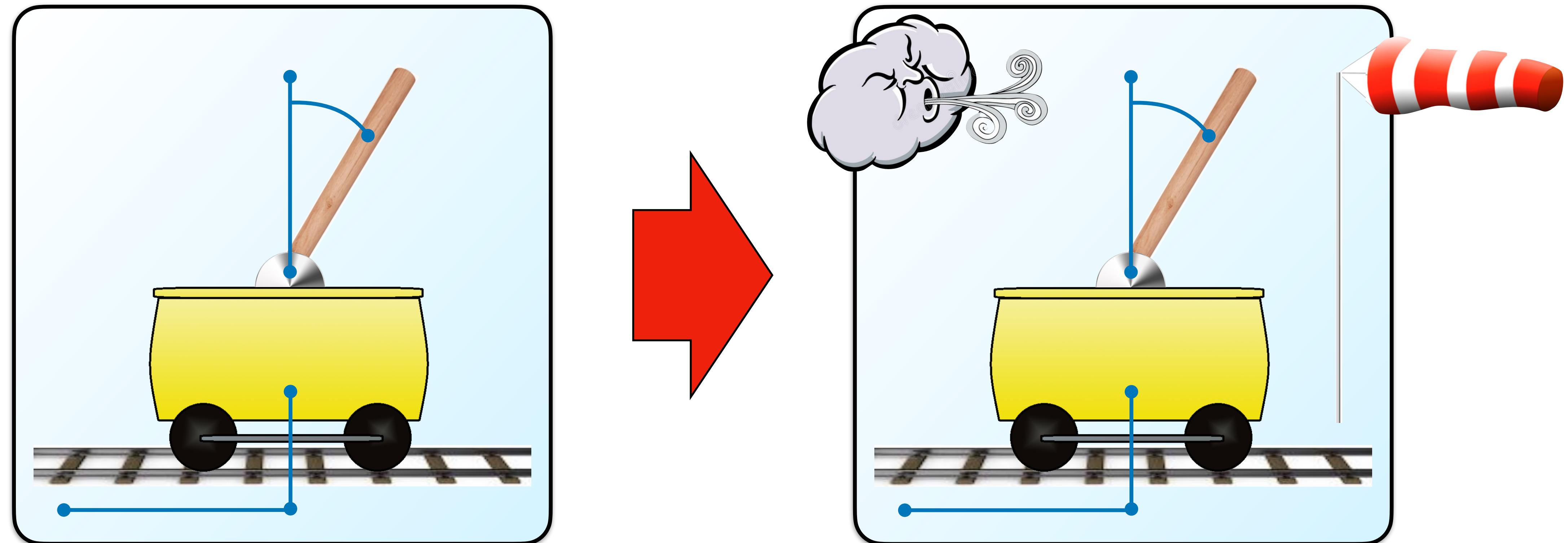
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Two common sources of **ambiguity**:

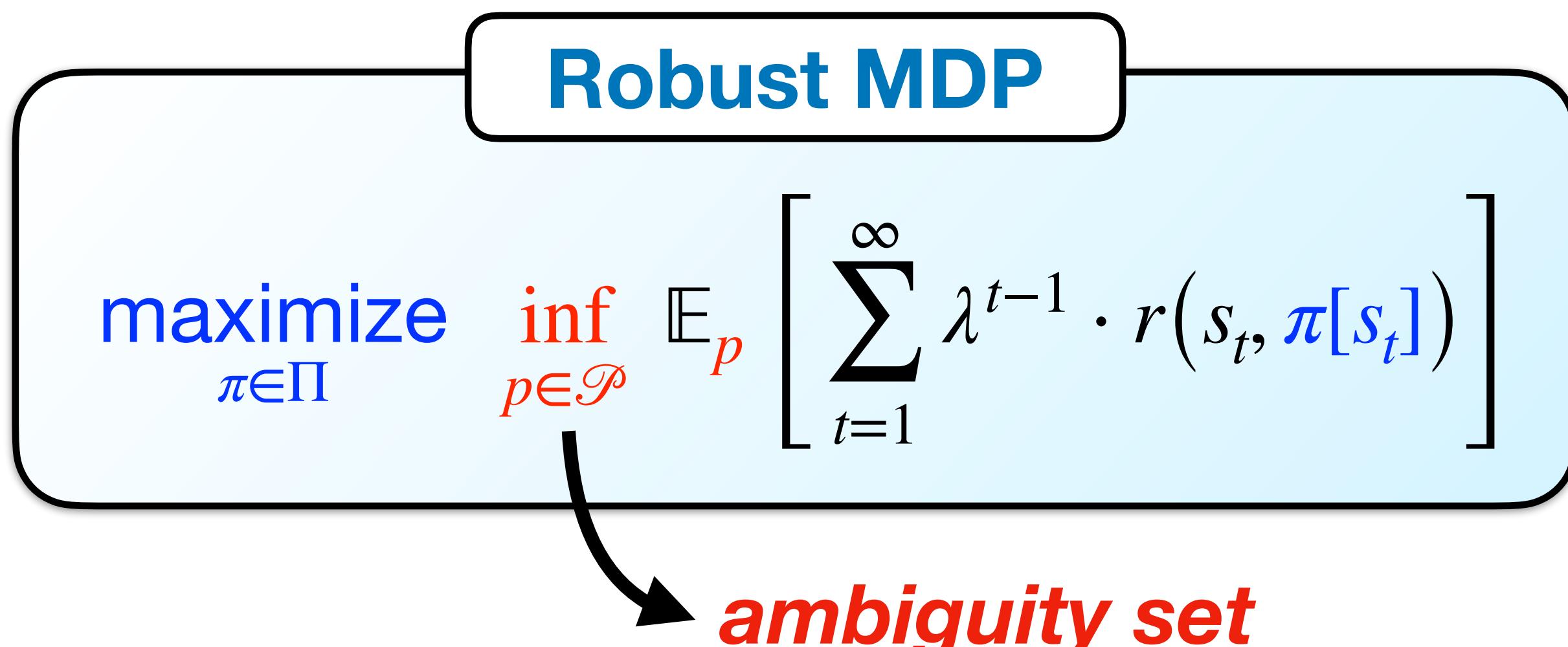
- **Modelling errors:** 32.67 secs/run → 2.45 secs/run
- **Estimation errors:** 32.67 secs/run → 4.68 secs/run



Two common sources of ambiguity:

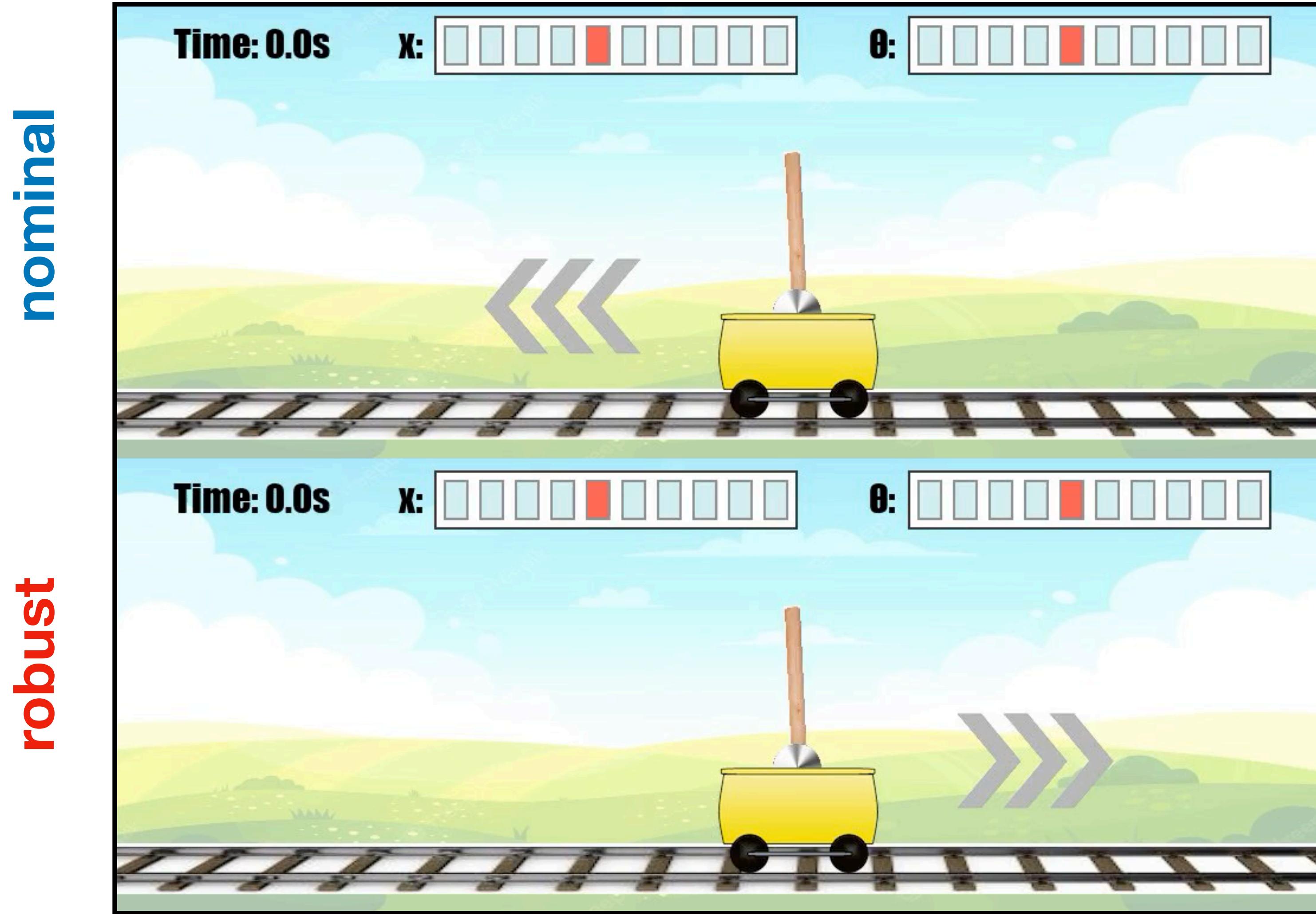
- **Modelling errors:** 32.67 secs/run → 2.45 secs/run
- **Estimation errors:** 32.67 secs/run → 4.68 secs/run

Impact of ambiguity can be alleviated via robust optimization:



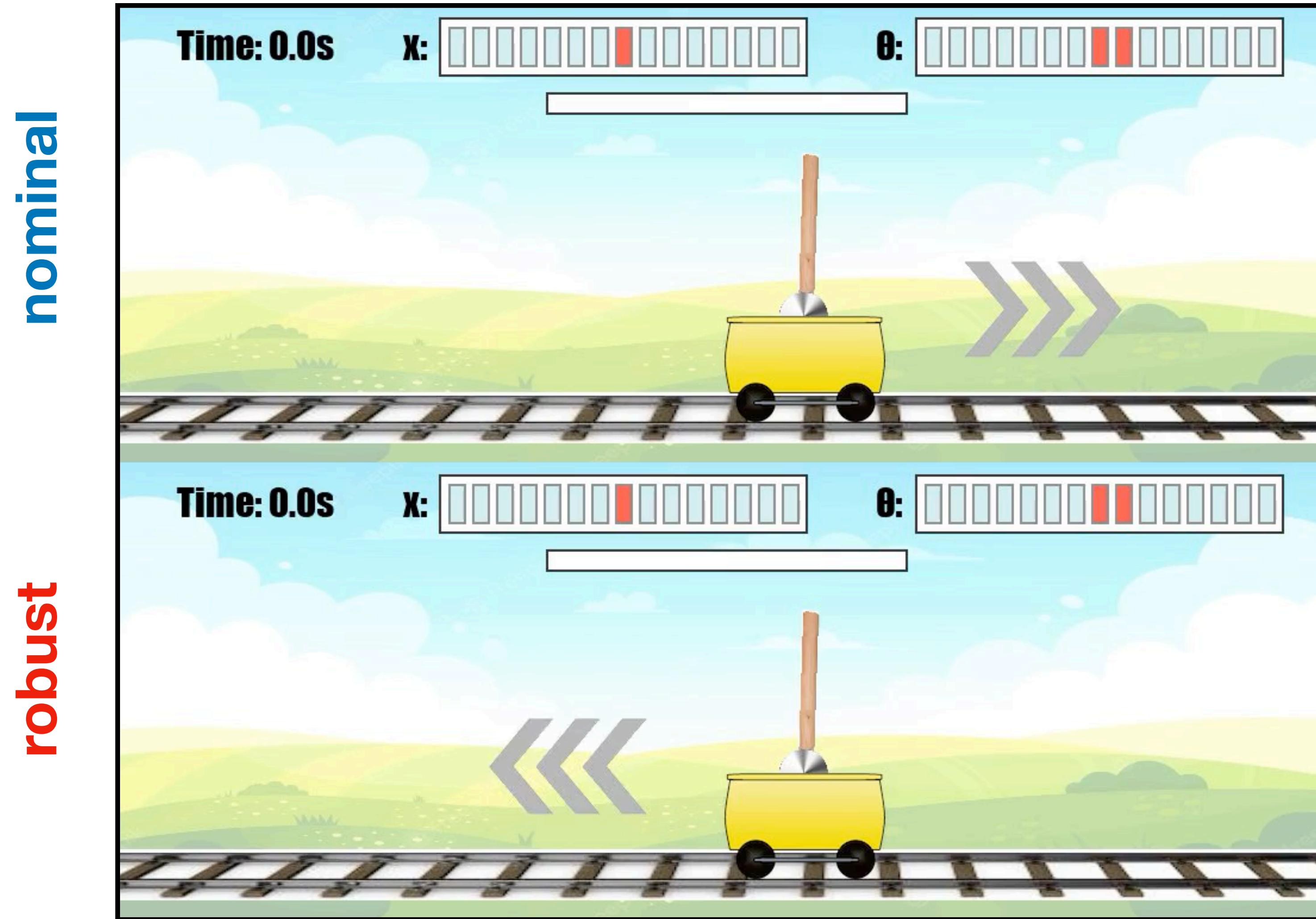
Robust MDPs admit interpretation as regularized MDPs!

Ambiguity: Modelling Errors



Modelling errors: 32.67 secs/run → 2.45 secs/run → 15.77 secs/run

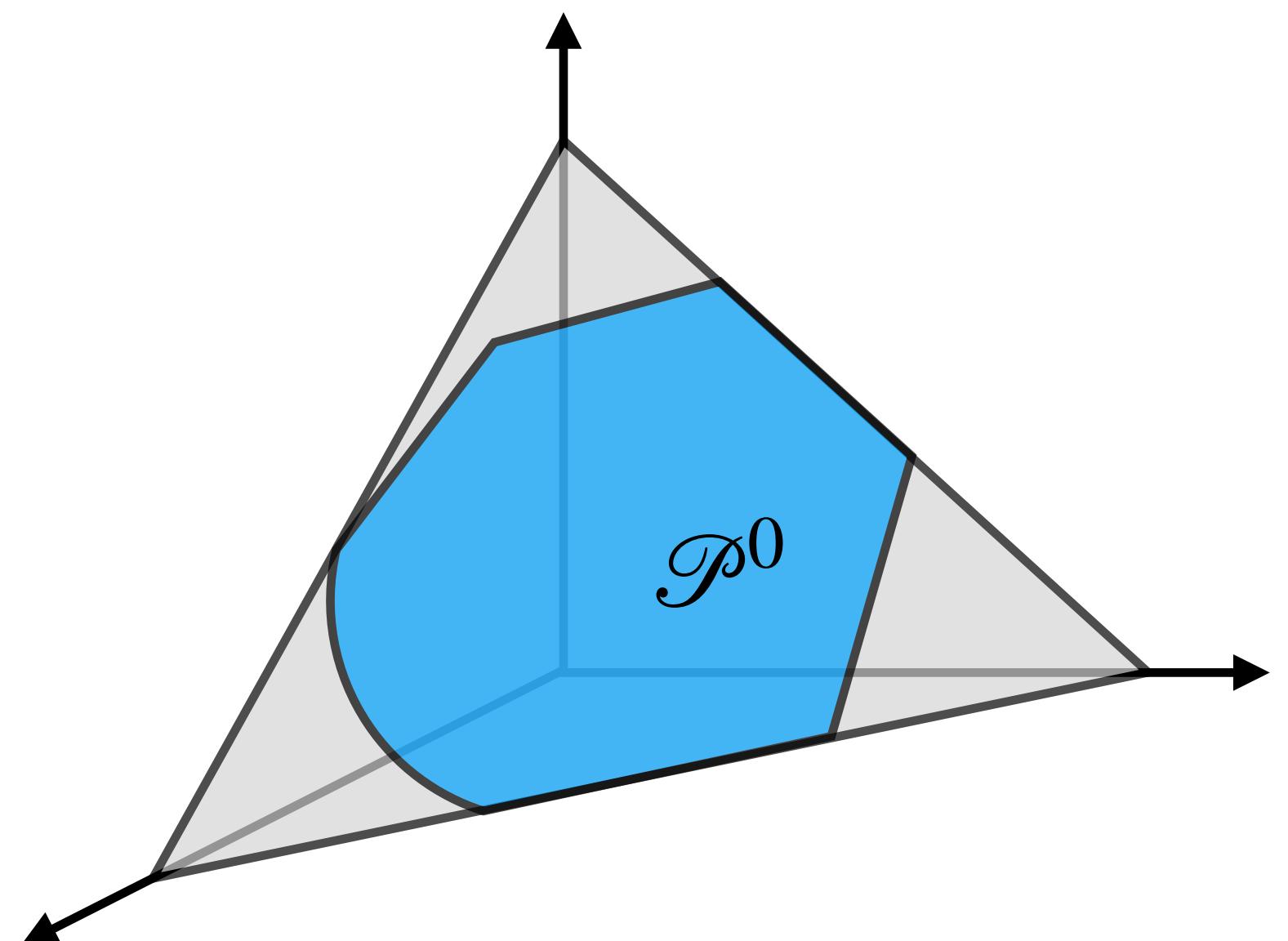
Ambiguity: Estimation Errors



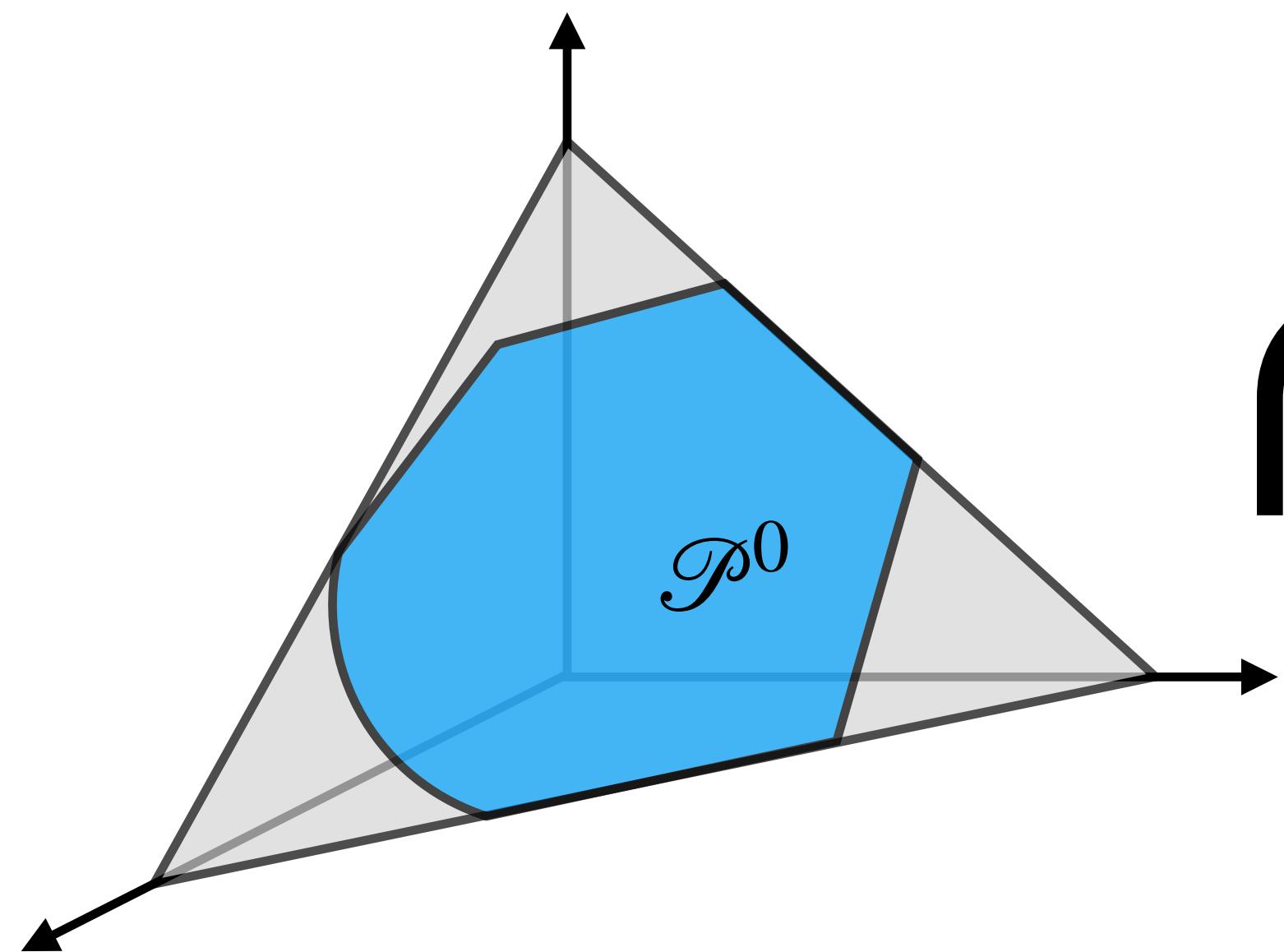
Estimation errors: 32.67 secs/run → 4.68 secs/run → 15.76 secs/run

Ambiguity Sets

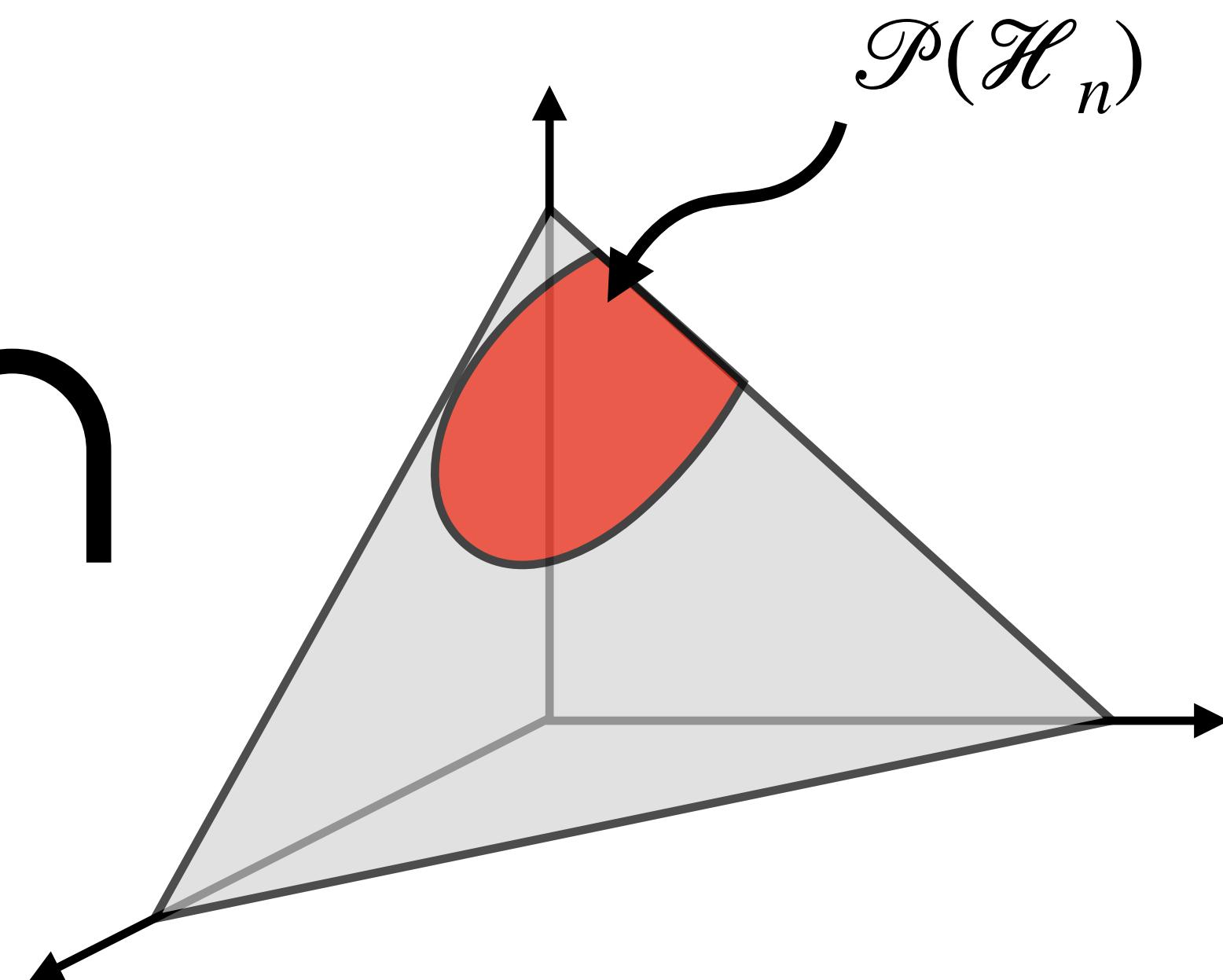
Structural ambiguity set



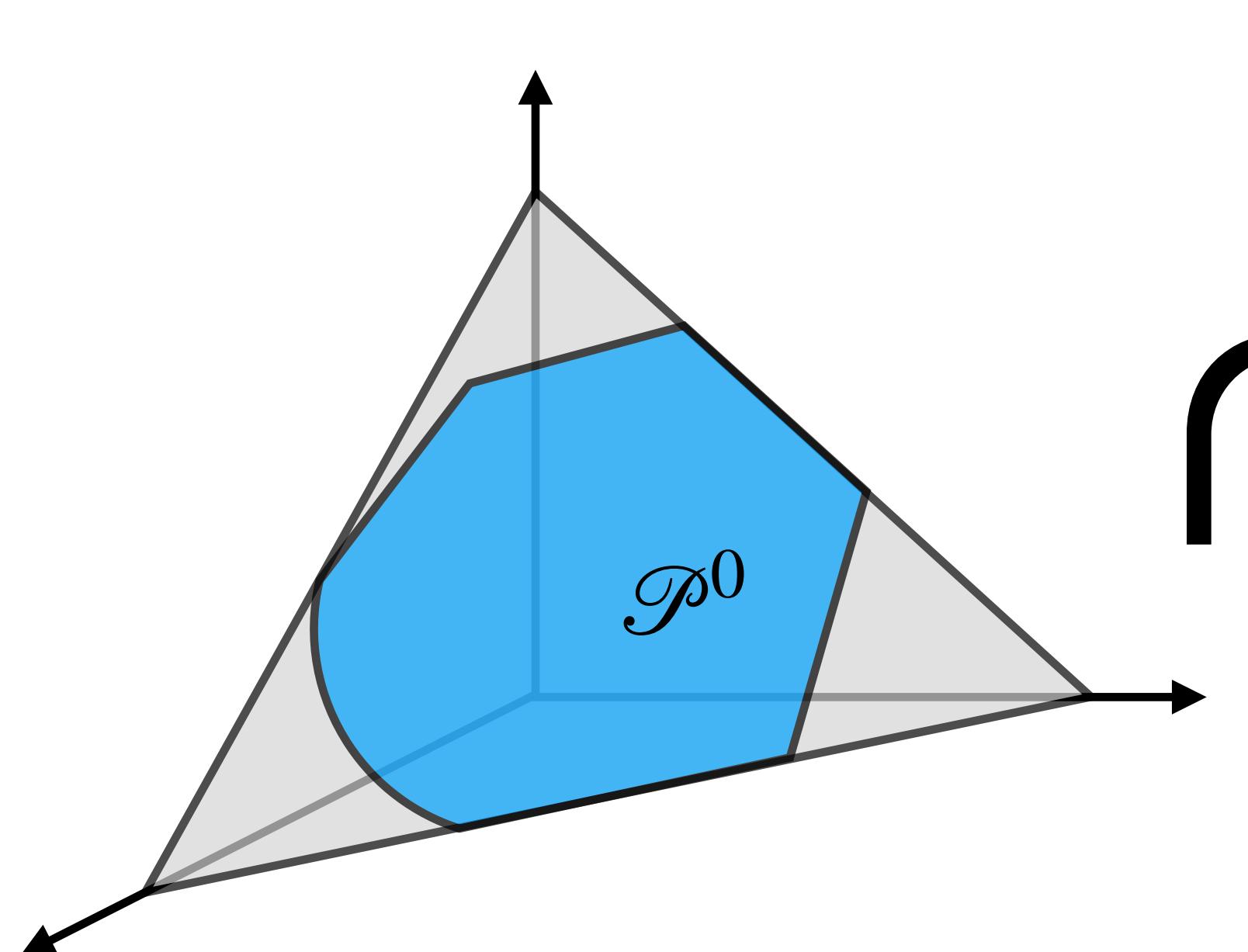
Structural ambiguity set



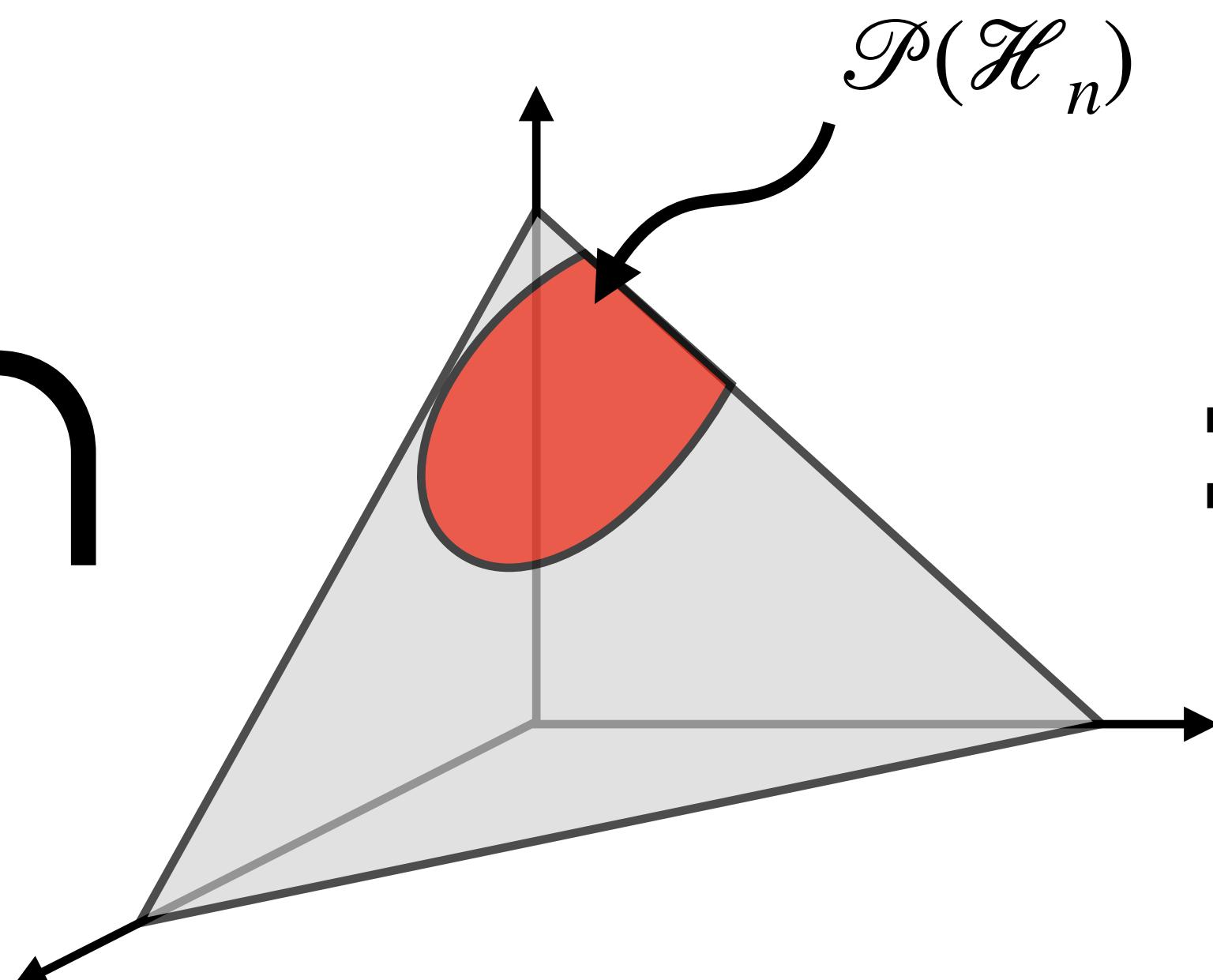
Historical sample



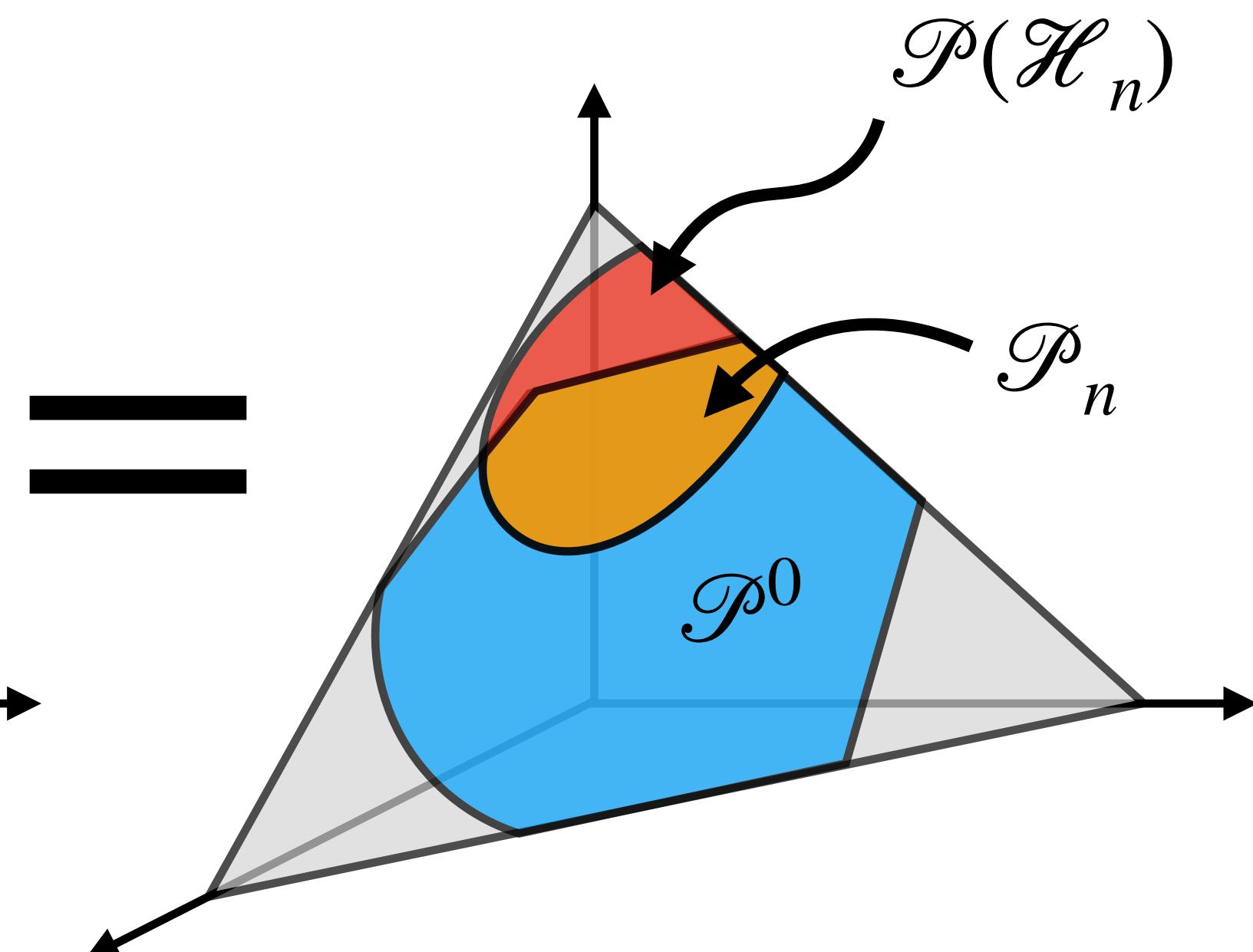
Structural ambiguity set



Historical sample

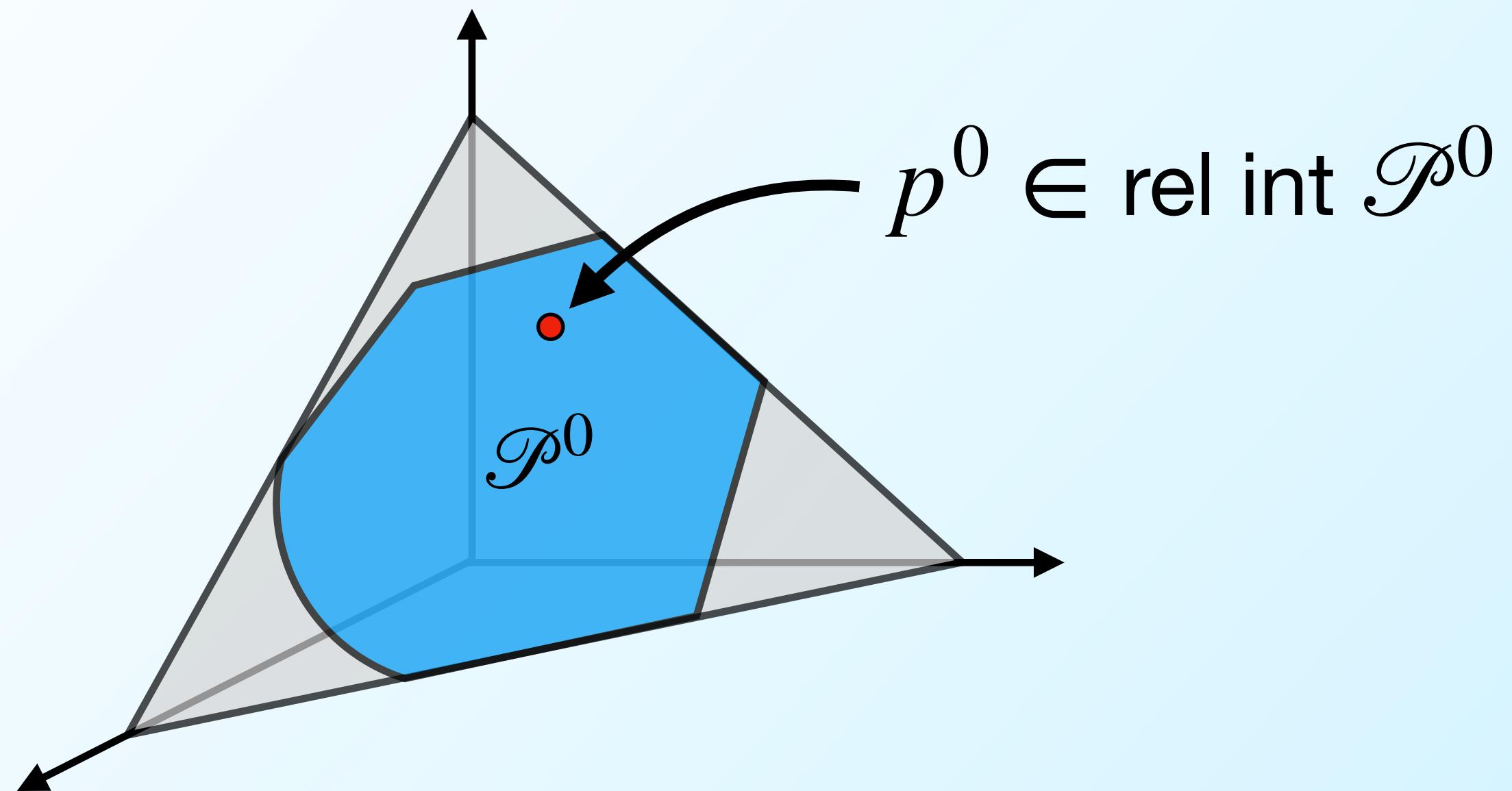


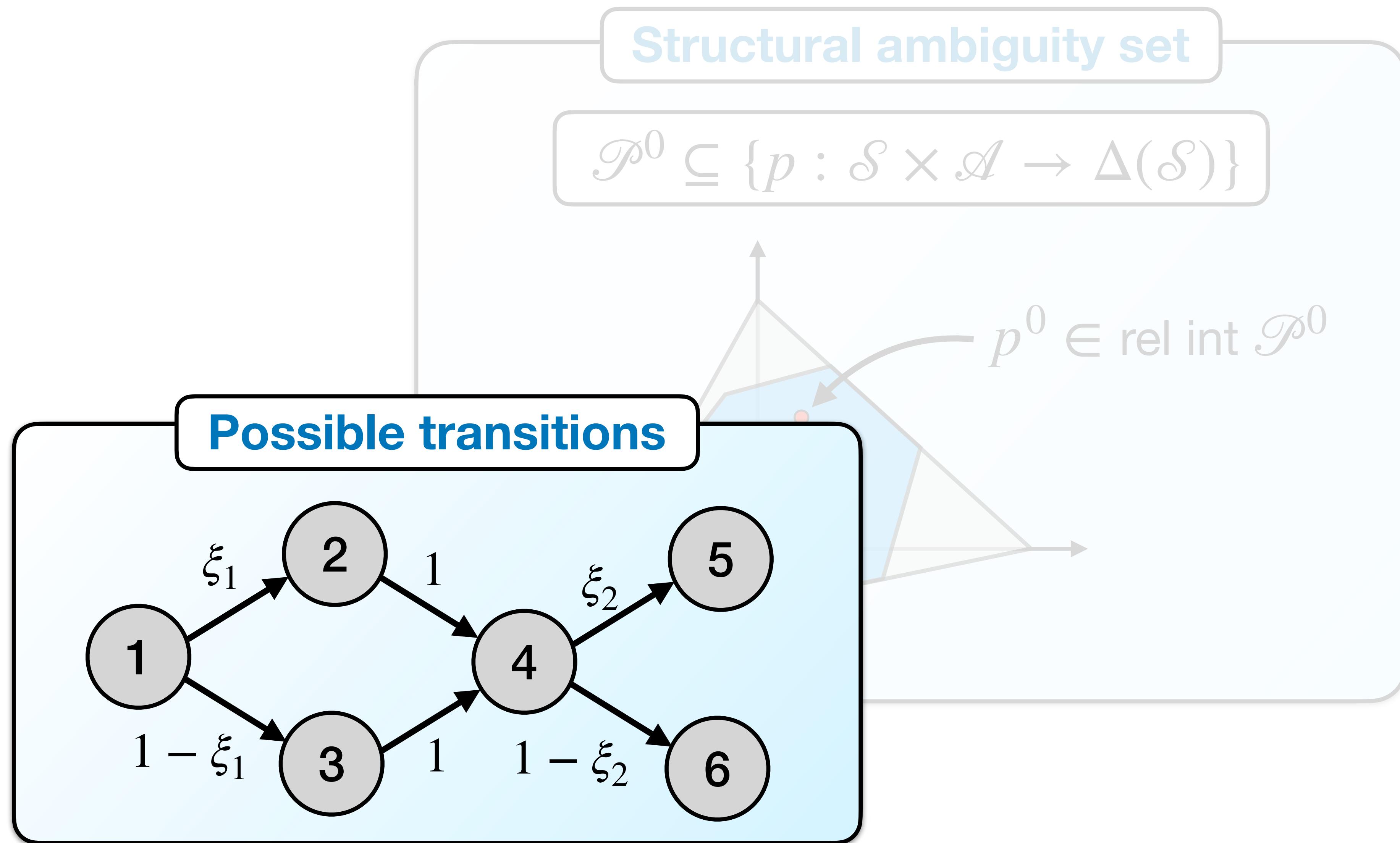
Out-of-sample guarantee



Structural ambiguity set

$$\mathcal{P}^0 \subseteq \{p : \mathcal{S} \times \mathcal{A} \rightarrow \Delta(\mathcal{S})\}$$



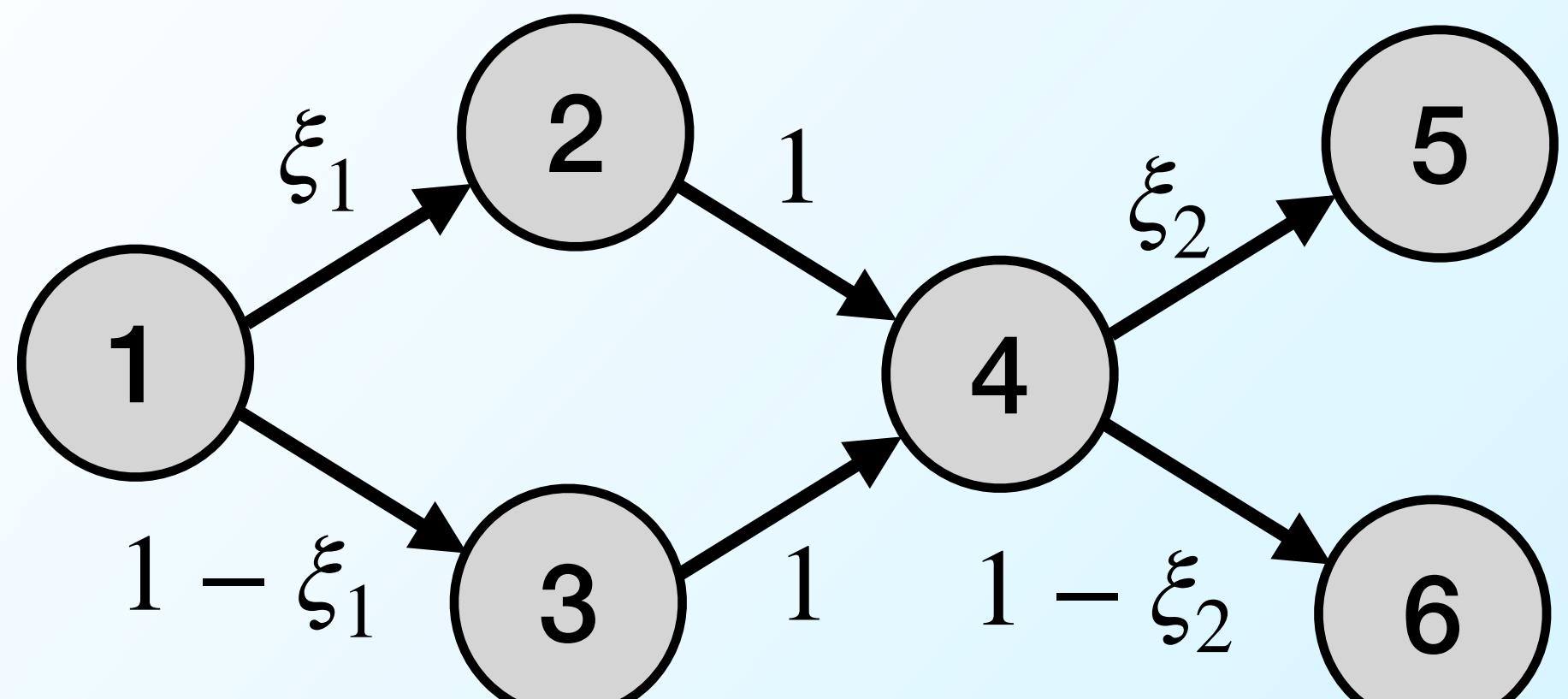


Structural ambiguity set

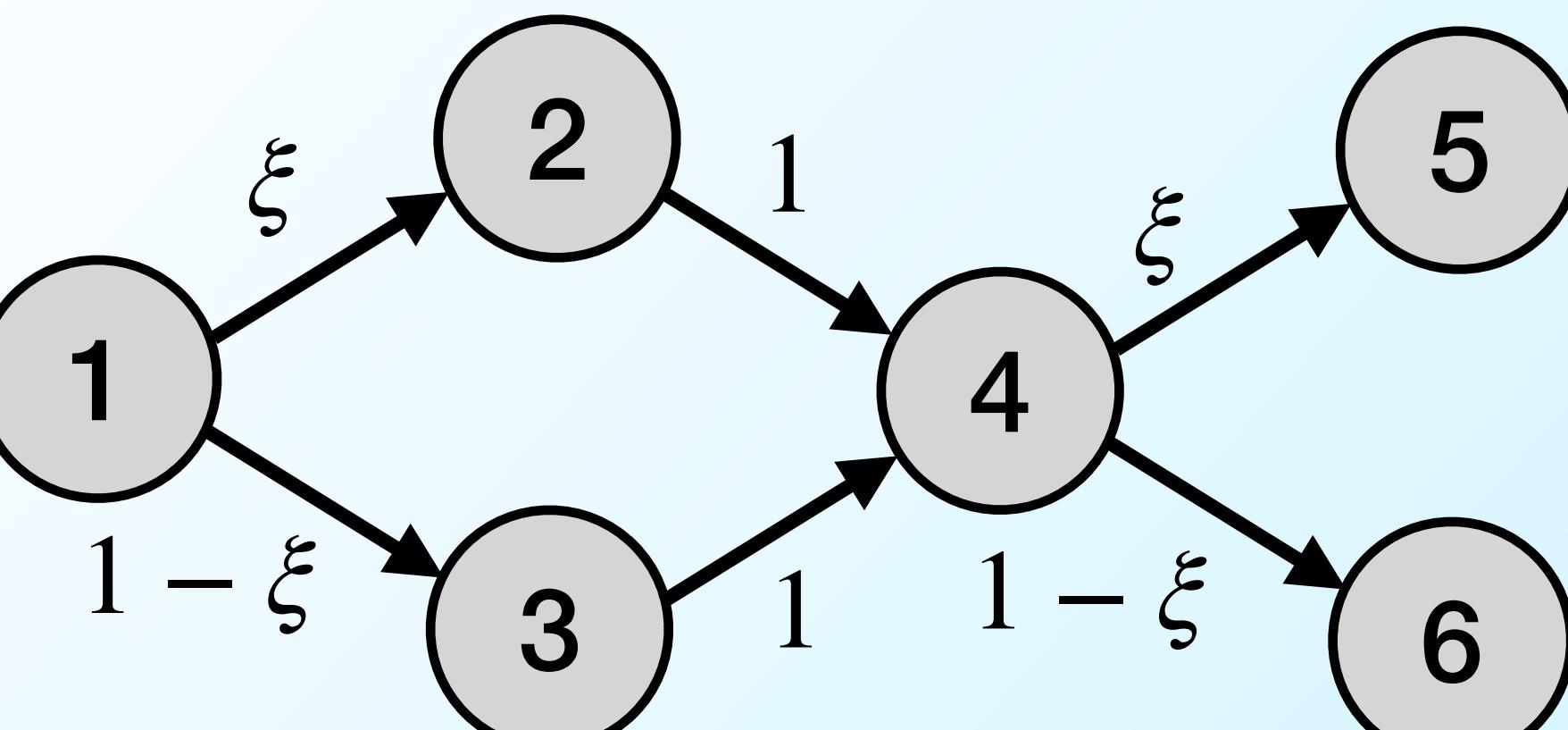
$$\mathcal{P}^0 \subseteq \{p : \mathcal{S} \times \mathcal{A} \rightarrow \Delta(\mathcal{S})\}$$

$$p^0 \in \text{rel int } \mathcal{P}^0$$

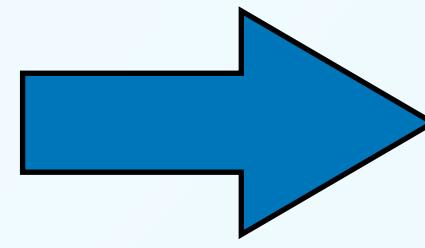
Possible transitions



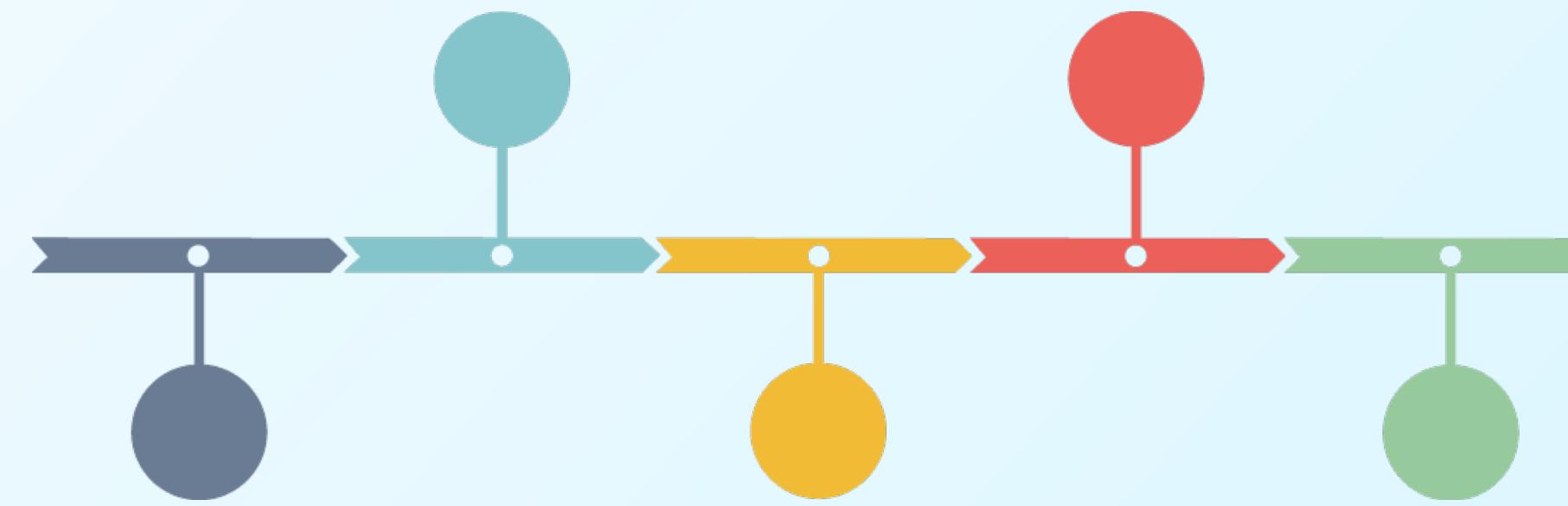
Equal probabilities



Historical sample



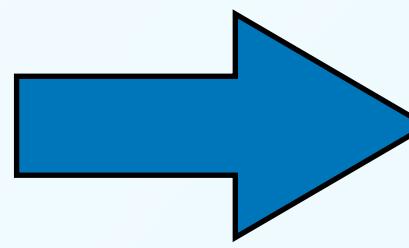
historical policy π^0
(stationary, randomized)



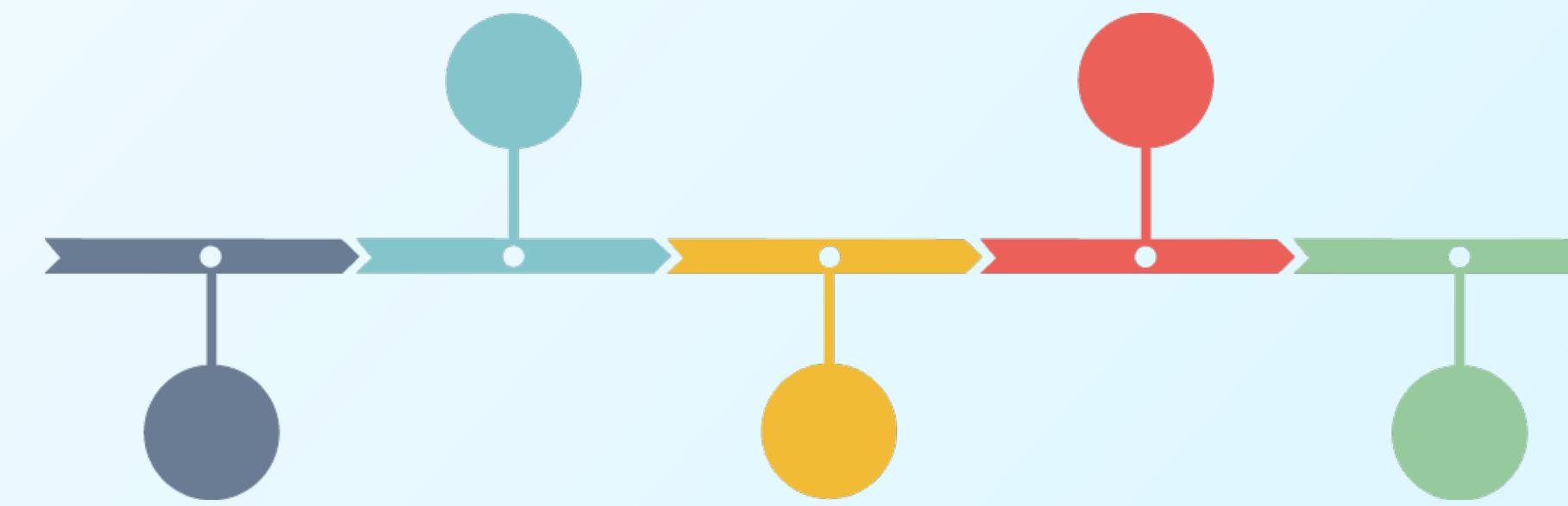
state-action history

$$\mathcal{H}_n = (s_1, a_1, \dots, s_n, a_n) \in (\mathcal{S} \times \mathcal{A})^n$$

Historical sample



historical policy π^0
(stationary, randomized)



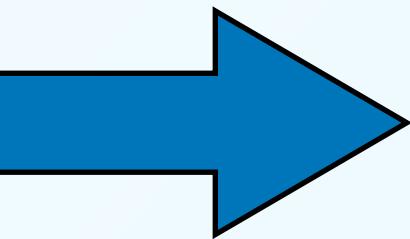
state-action history
 $\mathcal{H}_n = (s_1, a_1, \dots, s_n, a_n) \in (\mathcal{S} \times \mathcal{A})^n$

Likelihood, given history

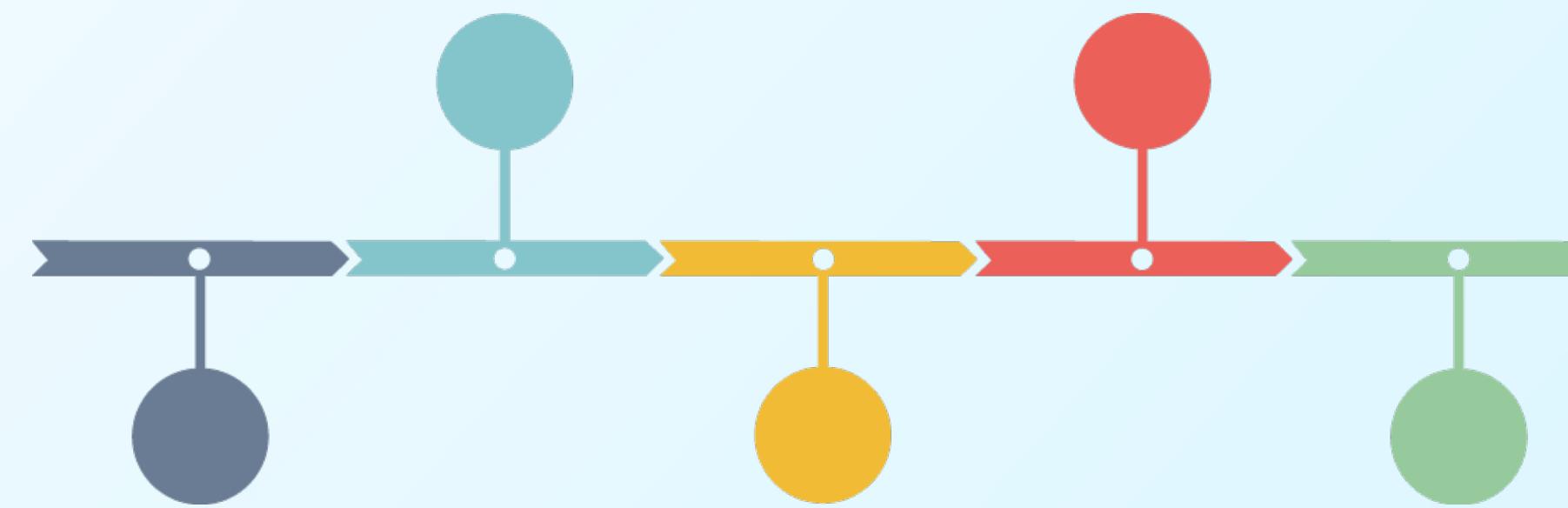
$$\mathcal{L}_n(p) = q(s_1) \cdot \pi^0(a_n | s_n) \cdot \prod_{t=1}^{n-1} [\pi^0(a_t | s_t) \cdot p(s_{t+1} | s_t, a_t)]$$

Historical sample

$$\mathcal{P}(\mathcal{H}_n) = \{p : \log \mathcal{L}_n(p) \geq \log \mathcal{L}_n(p^*) - \delta\}$$



historical policy π^0
(stationary, randomized)



state-action history
 $\mathcal{H}_n = (s_1, a_1, \dots, s_n, a_n) \in (\mathcal{S} \times \mathcal{A})^n$

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Theorem

Assumption: Historical policy π^0 visits every $s \in \mathcal{S}$ infinitely often as $n \rightarrow \infty$

Theorem

$\mathcal{P}_n = \mathcal{P}^0 \cap \mathcal{P}(\mathcal{H}_n)$ with $\delta = (1 - \beta)$ -quantile
of χ^2 -distribution with κ degrees of freedom

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$$\lim_{n \rightarrow \infty} \mathbb{P} [p^0 \in \mathcal{P}_n] = 1 - \beta$$

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2

$$\operatorname{plim}_{n \rightarrow \infty} [\sqrt{n} \cdot d^H(\mathcal{P}_n, \{p^0\})] = 0$$

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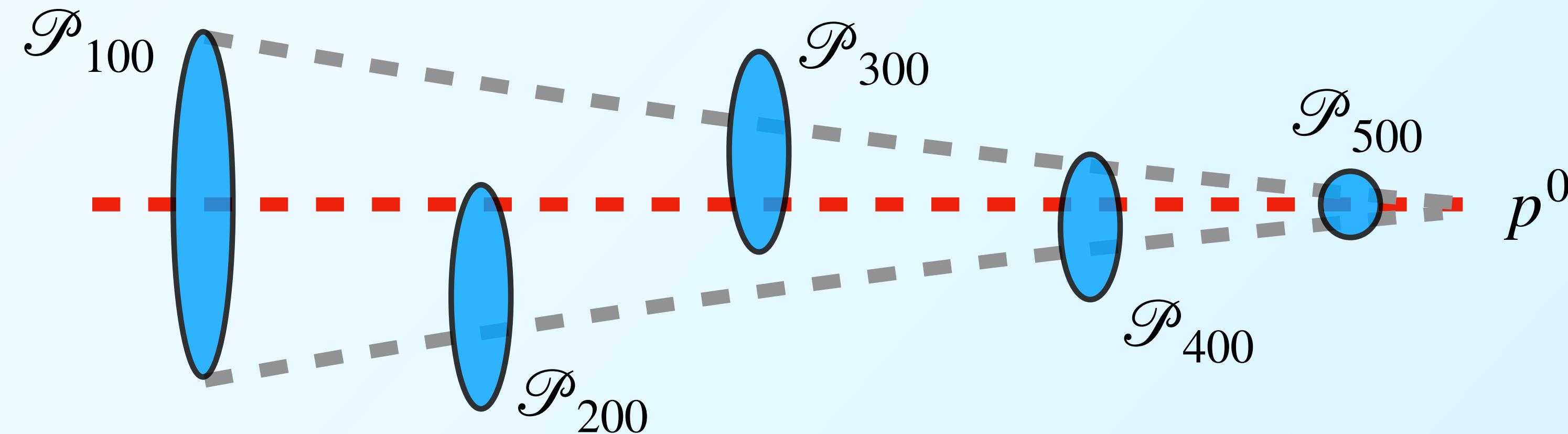
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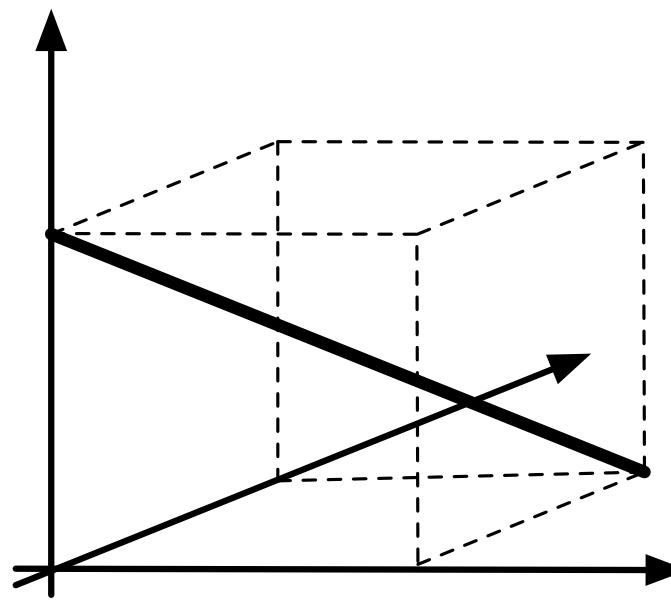
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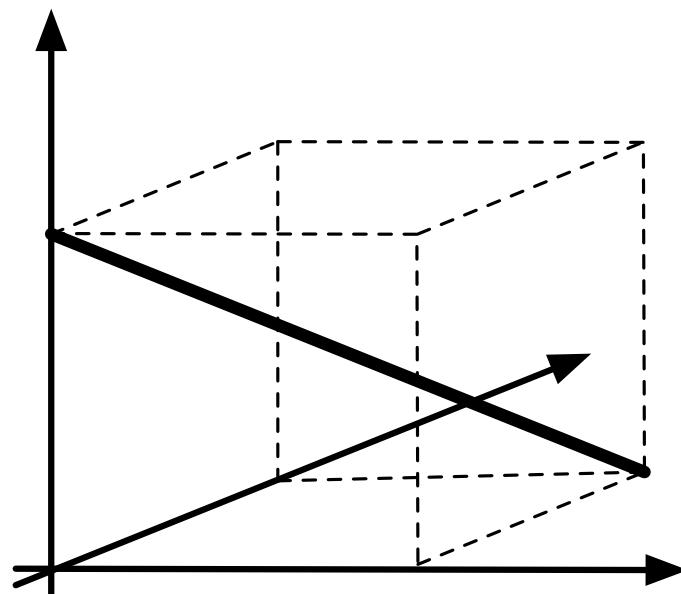
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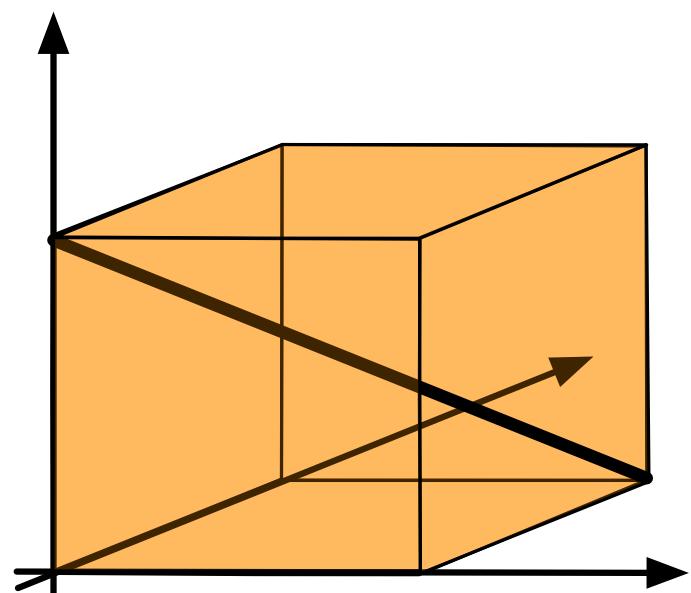
General (non-rectangular) ambiguity sets

- Optimal policy can be randomized & history-dependent
- Bellman optimality principle violated; NP-hard



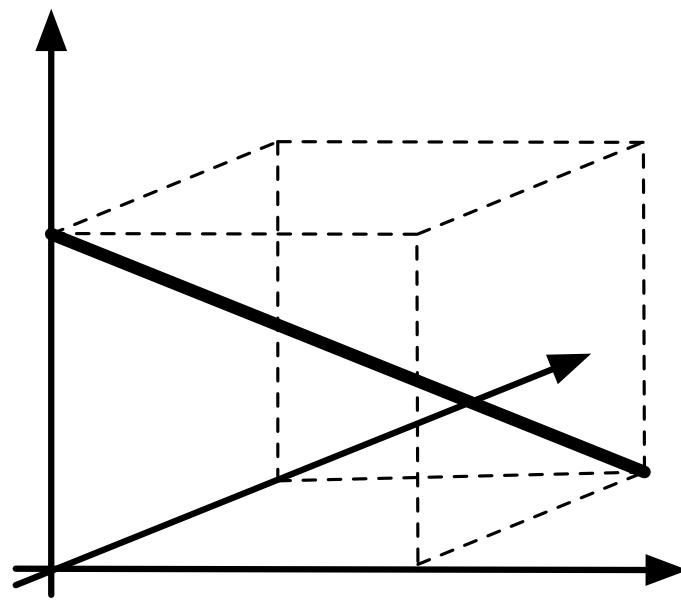
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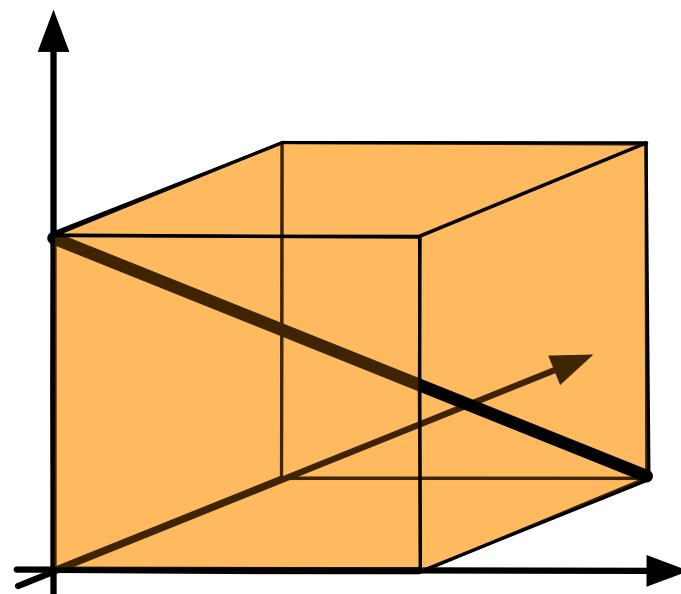
(s,a)-rectangular ambiguity sets

$$\mathcal{P} = \prod_{(s,a) \in \mathcal{S} \times \mathcal{A}} \mathcal{P}_{s,a} \quad \text{with} \quad \mathcal{P}_{s,a} \subseteq \Delta(\mathcal{S})$$



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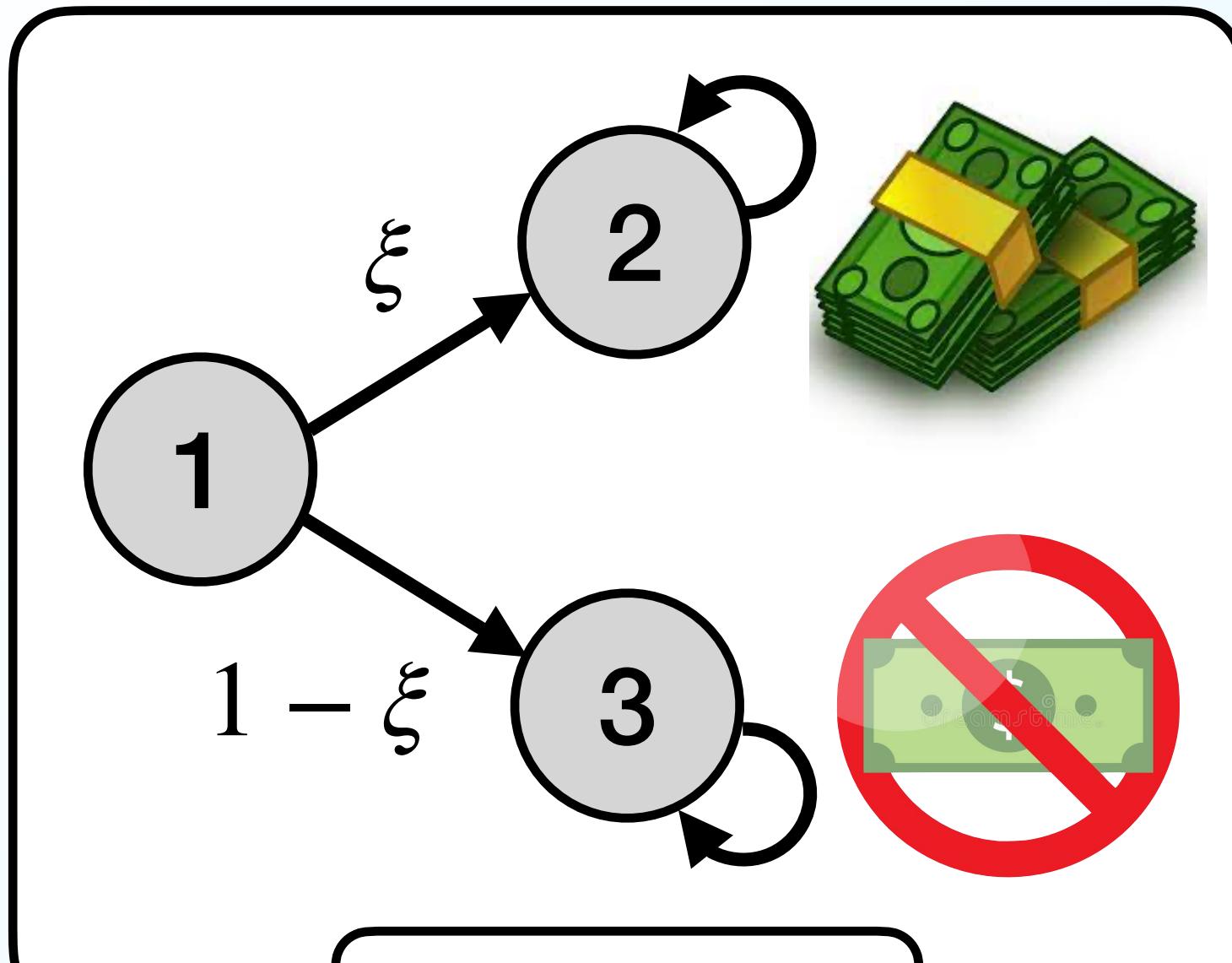


(s,a)-rectangular ambiguity sets

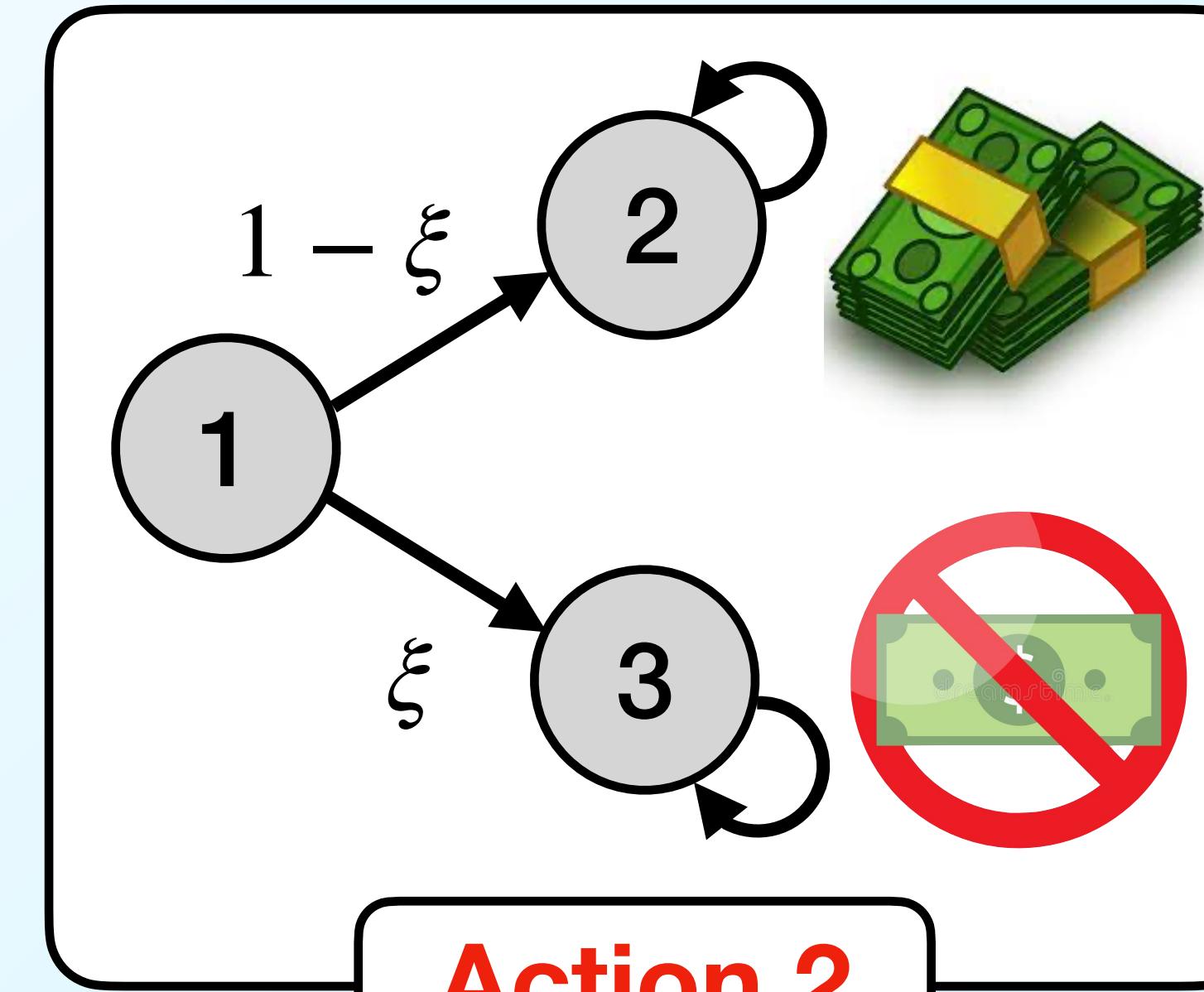
- Optimal policy stationary and deterministic
- Bellman optimality principle holds

General (non-rectangular) ambiguity sets

Example



Action 1



Action 2

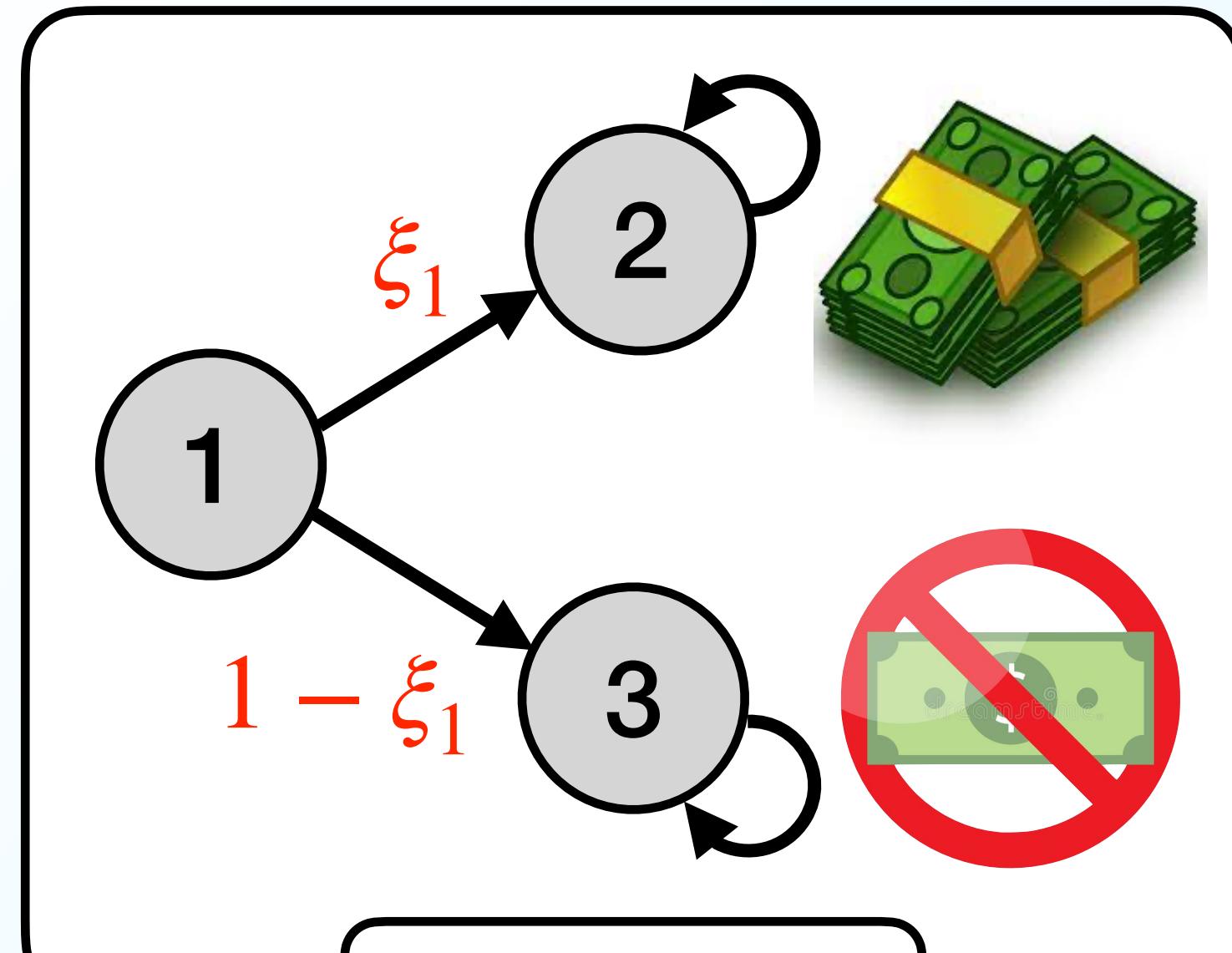
for some unknown $\xi \in [0,1]$



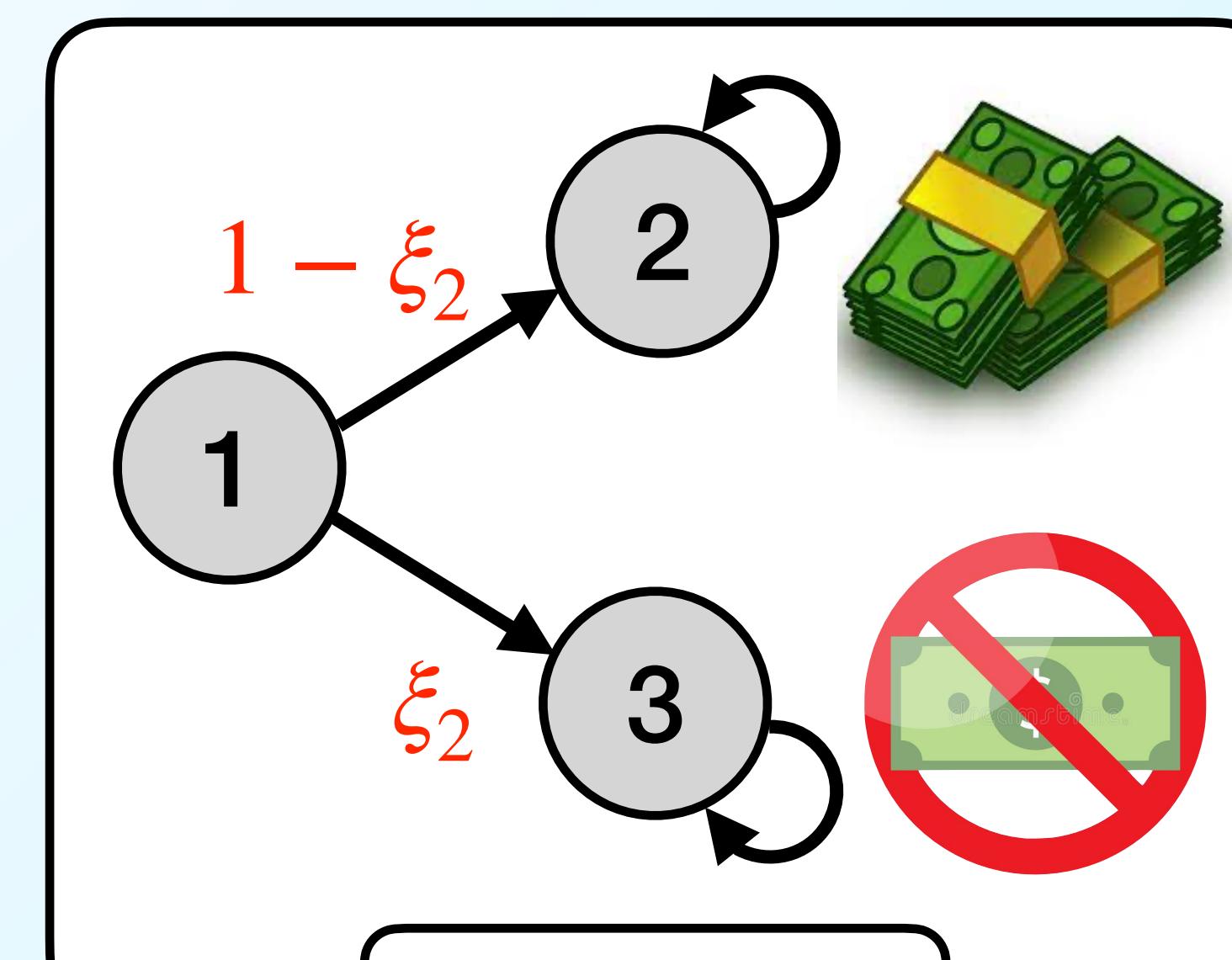
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Example



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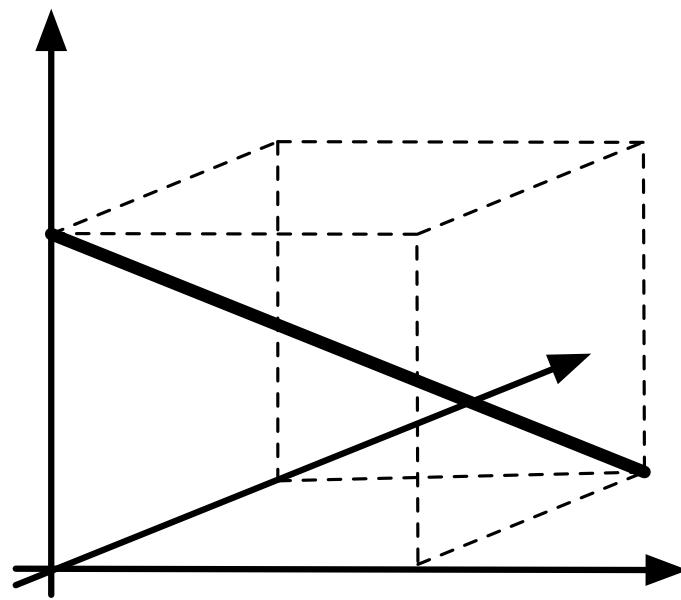


Action 2

for some unknown $\xi_1, \xi_2 \in [0,1]$

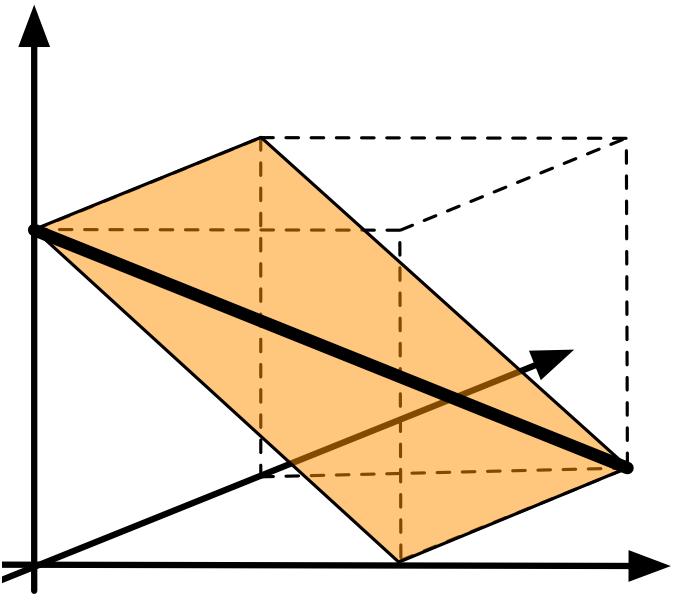


Bellman optimality principle holds



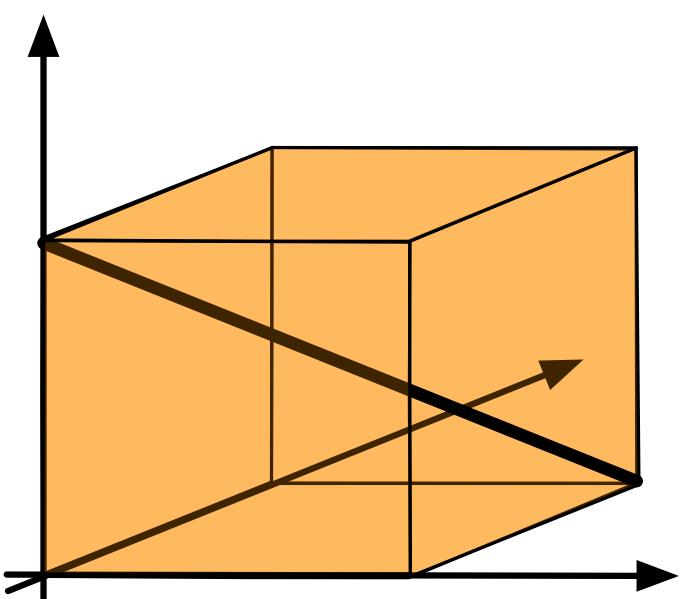
General (non-rectangular) ambiguity sets

- Optimal policy can be randomized & history-dependent
- Bellman optimality principle violated; NP-hard



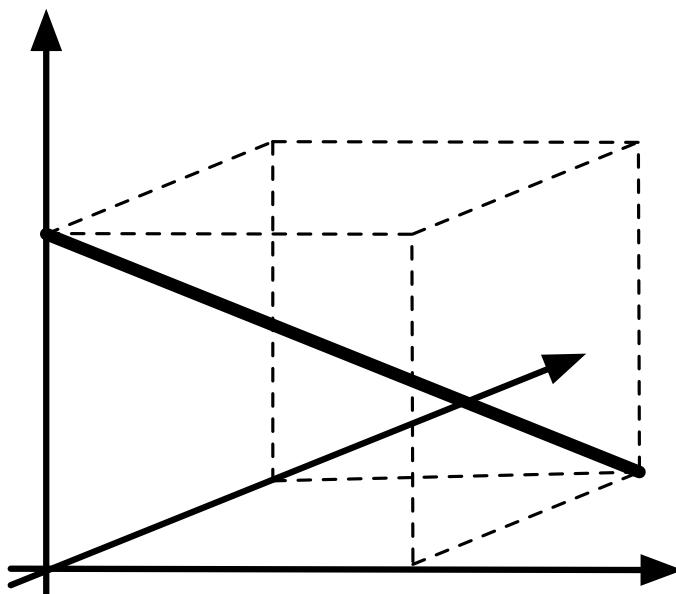
s -rectangular ambiguity sets

$$\mathcal{P} = \prod_{s \in \mathcal{S}} \mathcal{P}_s \quad \text{with} \quad \mathcal{P}_s \subseteq [\Delta(\mathcal{A})]^A$$



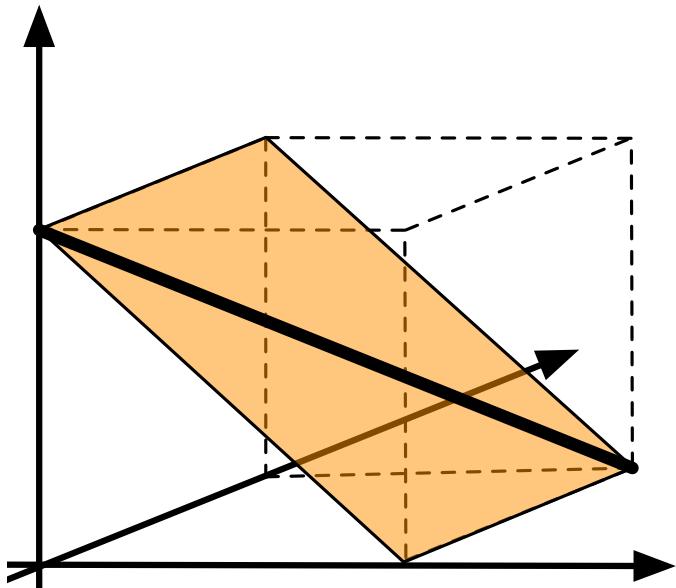
(s,a) -rectangular ambiguity sets

- Optimal policy stationary and deterministic
- Bellman optimality principle holds



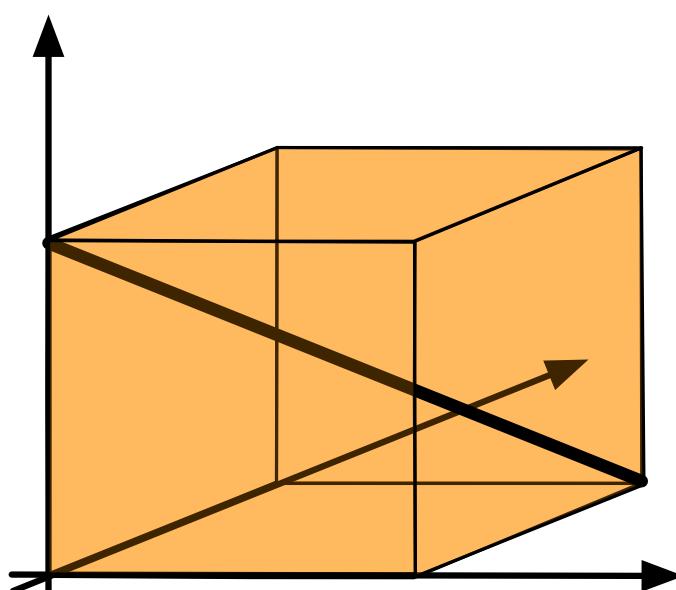
General (non-rectangular) ambiguity sets

- Optimal policy can be **randomized** & history-dependent
- Bellman optimality principle violated; NP-hard



s-rectangular ambiguity sets

- Optimal policy **stationary** but can be **randomized**
- Bellman optimality principle **holds**

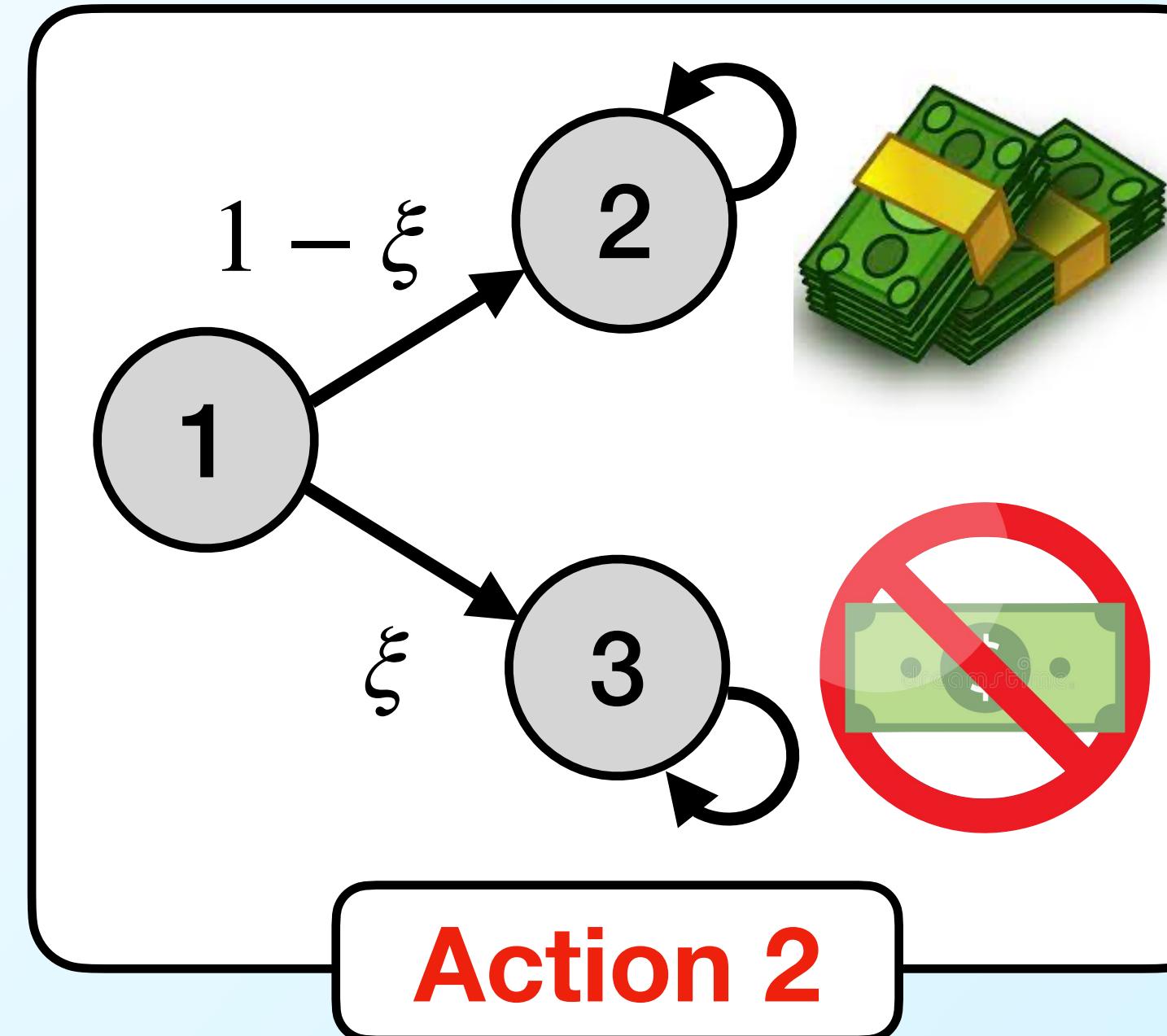
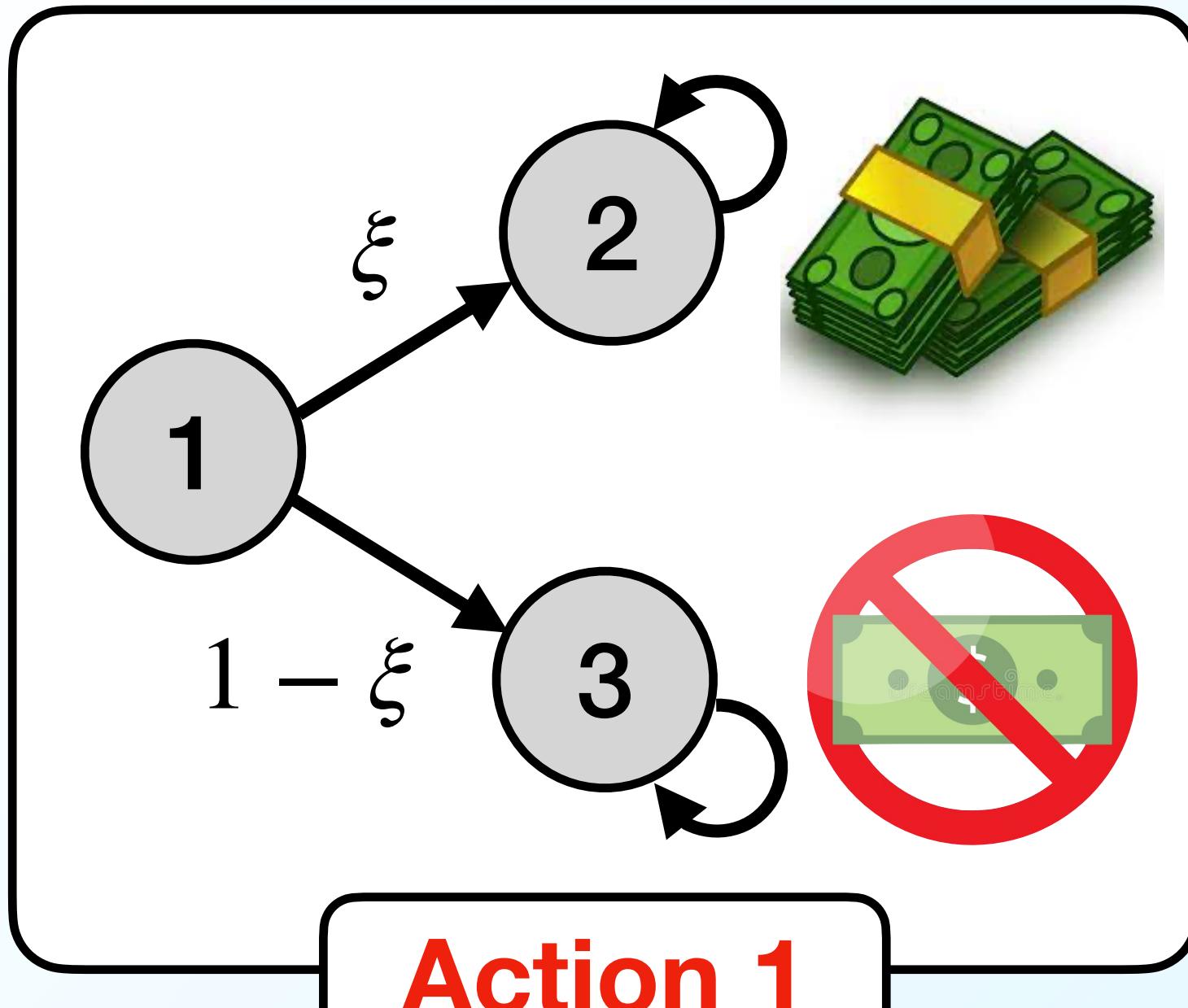


(s,a)-rectangular ambiguity sets

- Optimal policy **stationary** and **deterministic**
- Bellman optimality principle **holds**

General (non-rectangular) ambiguity sets

Example



for some unknown $\xi \in [0,1]$



Bellman optimality principle holds

s-Rectangular Ambiguity Sets: Bellman Operator

s-Rectangular Ambiguity Sets: Bellman Operator

Classical (non-robust) Bellman equations

$$v^*(s) = \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^*(s') \right\}$$

s-Rectangular Ambiguity Sets: Bellman Operator

Robust Bellman equations

$$v^\star(s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^\star(s') \right] \right\}$$

s-Rectangular Ambiguity Sets: Bellman Operator

Robust Bellman operator

$$[\mathfrak{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$

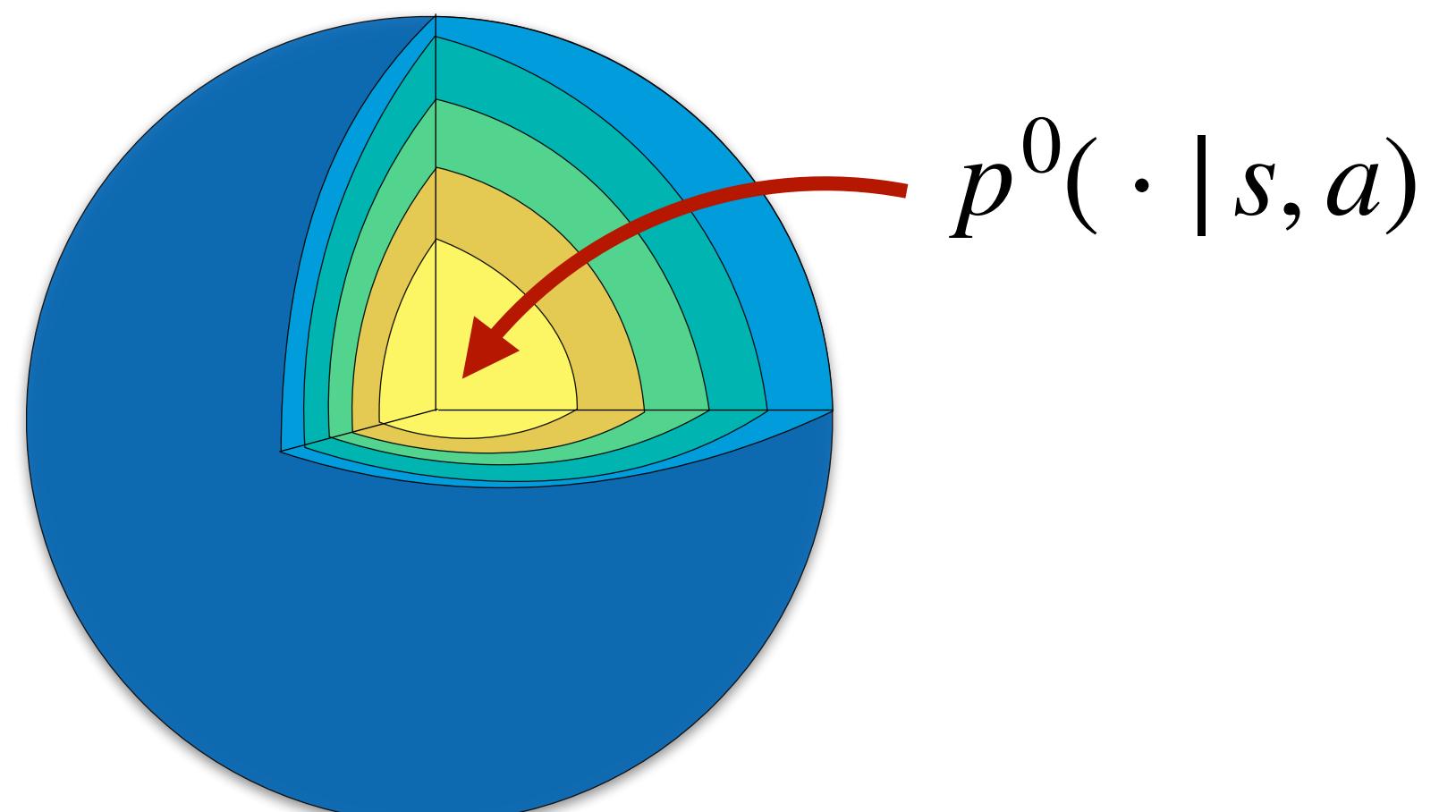
s-Rectangular Ambiguity Sets: Bellman Operator

Robust Bellman operator

$$[\mathfrak{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$

Distance-constrained s-rectangular ambiguity set

$$\mathcal{P} = \prod_{s \in \mathcal{S}} \mathcal{P}_s \quad \text{with} \quad \mathcal{P}_s = \left\{ p(\cdot | s, \cdot) : \sum_{a \in \mathcal{A}} d[p(\cdot | s, a), p^0(\cdot | s, a)] \leq \kappa \right\}$$



s-Rectangular Ambiguity Sets: Bellman Operator

$$[\mathcal{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$

s-Rectangular Ambiguity Sets: Bellman Operator

$$\begin{aligned} [\mathfrak{B}v](s) &= \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \\ &\downarrow \\ &= \min_{p \in \mathcal{P}_s} \max_{\pi \in \Delta(\mathcal{A})} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \end{aligned}$$

Minimax theorem: exchange order of min and max

s-Rectangular Ambiguity Sets: Bellman Operator

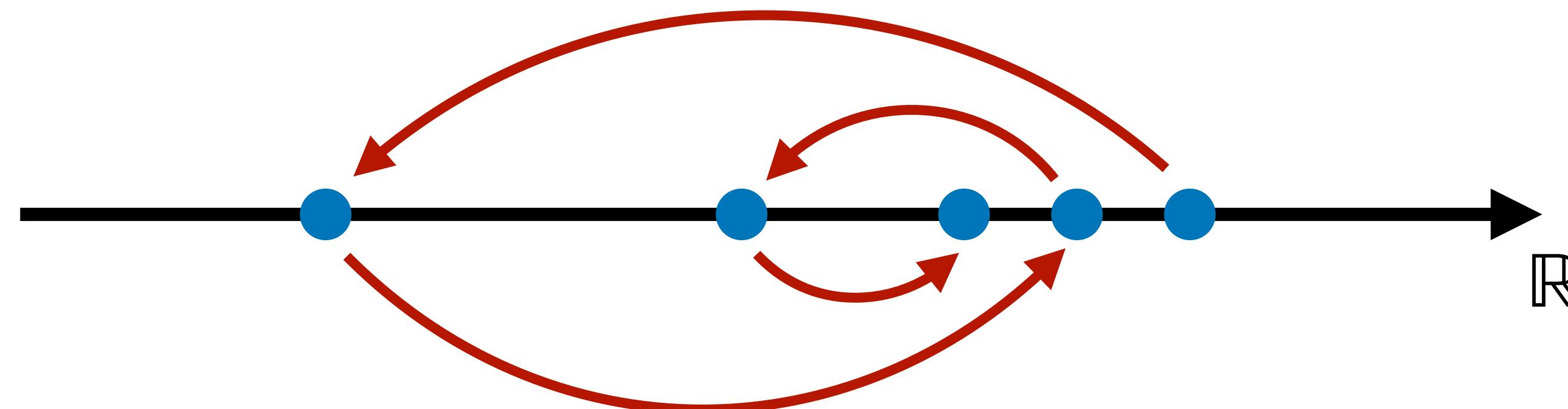
$$\begin{aligned}
 [\mathcal{B}v](s) &= \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \\
 &\downarrow \\
 &= \min_{p \in \mathcal{P}_s} \max_{\pi \in \Delta(\mathcal{A})} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \\
 &\downarrow \\
 &= \min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\}
 \end{aligned}$$

Linearity: we only need to consider $\text{ext } \Delta(\mathcal{A}) = \mathcal{A}$

s-Rectangular Ambiguity Sets: Bellman Operator

$$\begin{aligned}
 [\mathcal{B}v](s) &= \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \\
 &\downarrow \\
 &= \min_{p \in \mathcal{P}_s} \max_{\pi \in \Delta(\mathcal{A})} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \\
 &\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?
 \end{aligned}$$

Bisection search:



s-Rectangular Ambiguity Sets: Bellman Operator

$$\min_{p \in \mathcal{P}_s} \quad \max_{a \in \mathcal{A}} \quad \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

s-Rectangular Ambiguity Sets: Bellman Operator

$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$



$$\min_{p \in [\Delta(\mathcal{S})]^A} \left\{ \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} : \sum_{a \in \mathcal{A}} d[p(\cdot | s, a), p^0(\cdot | s, a)] \leq \kappa \right\} \leq \theta$$

s-Rectangular Ambiguity Sets: Bellman Operator

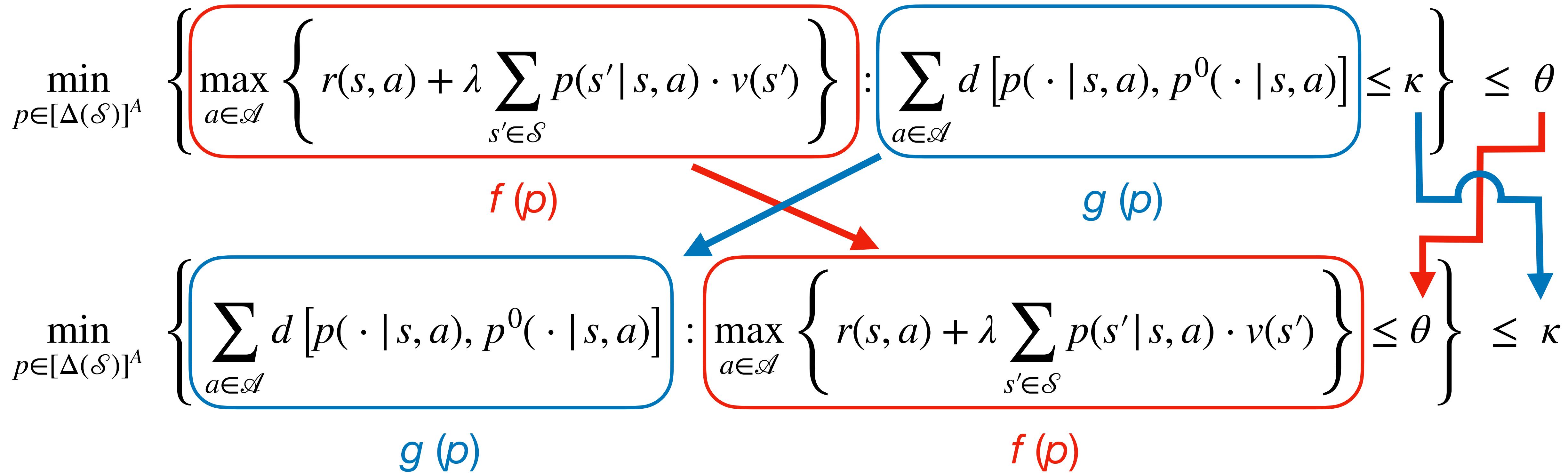
$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

$$\min_{p \in [\Delta(\mathcal{S})]^A} \left\{ \boxed{\max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\}} : \boxed{\sum_{a \in \mathcal{A}} d[p(\cdot | s, a), p^0(\cdot | s, a)] \leq \kappa} \leq \theta \right\}$$

f (p) *g (p)*

s-Rectangular Ambiguity Sets: Bellman Operator

$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$



s-Rectangular Ambiguity Sets: Bellman Operator

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g (p) *f (p)*

s-Rectangular Ambiguity Sets: Bellman Operator

$$\min_{p \in \mathcal{P}_s} \quad \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

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$$\Downarrow \iff \sum_{a \in \mathcal{A}} \min_{p_a \in \Delta(\mathcal{S})} \left\{ d \left[p(\cdot | s, a), p^0(\cdot | s, a) \right] : r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \leq \theta \right\} \leq \kappa$$

Separability: of both objective and constraints in $a \in \mathcal{A}$

s-Rectangular Ambiguity Sets: Bellman Operator

$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

$$\sum_{a \in \mathcal{A}} \min_{p_a \in \Delta(\mathcal{S})} \left\{ d[p(\cdot | s, a), p^0(\cdot | s, a)] : r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \leq \theta \right\} \leq \kappa$$

s-Rectangular Ambiguity Sets: Bellman Operator

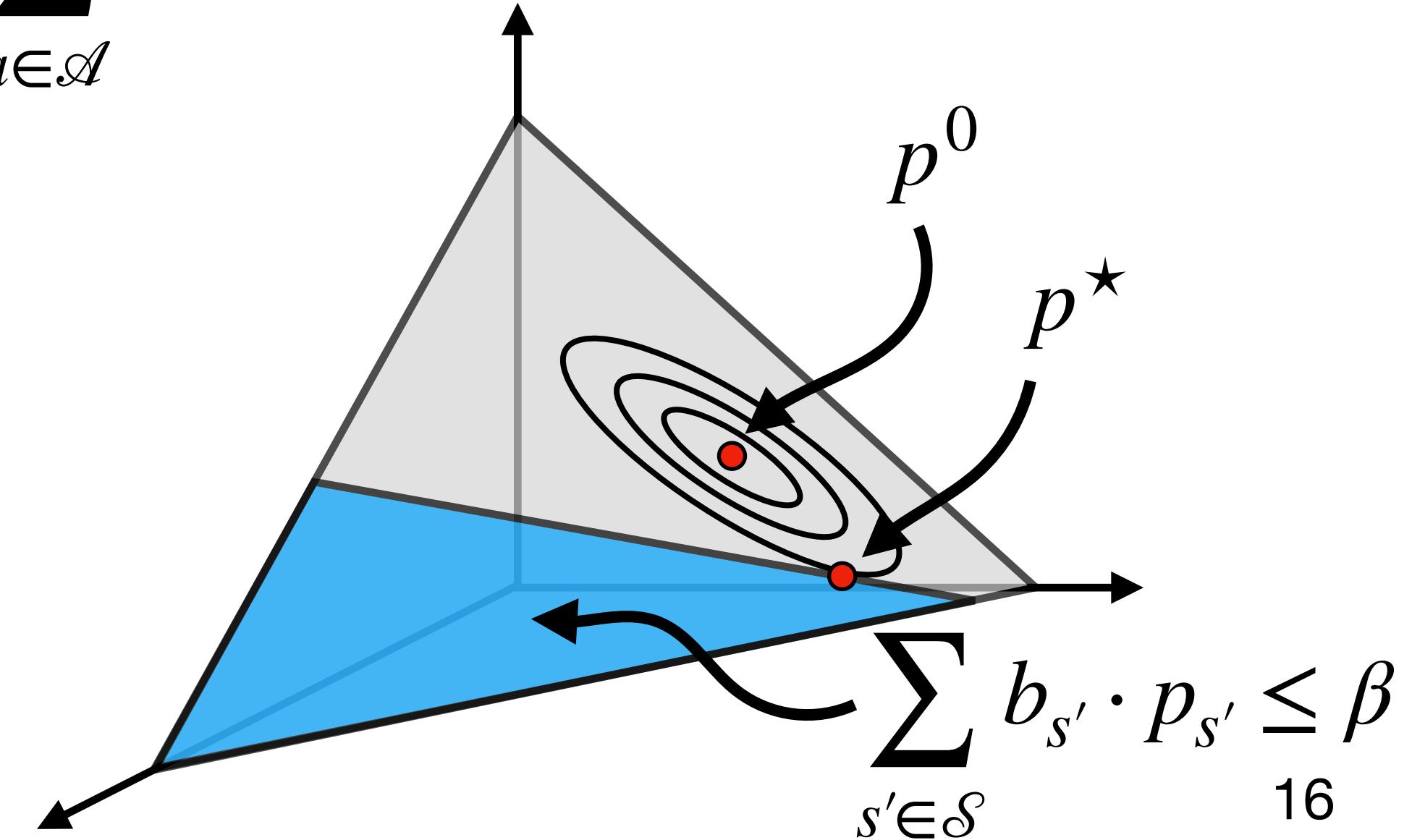
$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

$$\sum_{a \in \mathcal{A}} \min_{p_a \in \Delta(\mathcal{S})} \left\{ d[p(\cdot | s, a), p^0(\cdot | s, a)] : r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \leq \theta \right\} \leq \kappa$$

L \iff $\sum_{a \in \mathcal{A}} \mathfrak{P}(p^0; \lambda v, \theta - r(s | a)) \leq \kappa$

with $\mathfrak{P}(p^0; b, \beta) = \begin{bmatrix} \text{minimize}_{p} \\ \text{subject to} \end{bmatrix}$

$$d[p, p^0] \\ \sum_{s' \in \mathcal{S}} b_{s'} \cdot p_{s'} \leq \beta \\ p \in \Delta(\mathcal{S})$$



s-Rectangular Ambiguity Sets: Bellman Operator

Distance-constrained s-rectangular ambiguity set

$$\mathcal{P} = \prod_{s \in \mathcal{S}} \mathcal{P}_s \quad \text{with} \quad \mathcal{P}_s = \left\{ p(\cdot | s, \cdot) : \sum_{a \in \mathcal{A}} d[p(\cdot | s, a), p^0(\cdot | s, a)] \leq \kappa \right\}$$

s-Rectangular Ambiguity Sets: Bellman Operator

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Robust Bellman operator

$$[\mathfrak{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$

s-Rectangular Ambiguity Sets: Bellman Operator

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Projection problem

$$\mathfrak{P}(p^0; b, \beta) = \begin{bmatrix} \underset{p}{\text{minimize}} & d[p, p^0] \\ \text{subject to} & \sum_{s' \in \mathcal{S}} b_{s'} \cdot p_{s'} \leq \beta \\ & p \in \Delta(\mathcal{S}) \end{bmatrix}$$

s-Rectangular Ambiguity Sets: Bellman Operator

Distance-constrained s-rectangular ambiguity set

$$\mathcal{P} = \prod_{s \in \mathcal{S}} \mathcal{P}_s \quad \text{with} \quad \mathcal{P}_s = \left\{ p(\cdot | s, \cdot) : \sum_{a \in \mathcal{A}} d[p(\cdot | s, a), p^0(\cdot | s, a)] \leq \kappa \right\}$$

Theorem

Assume \mathfrak{P} can be computed **exactly** in time $\mathcal{O}(h(S))$.

Then \mathfrak{B} can be **computed to accuracy** $\epsilon > 0$ in time $\mathcal{O}(AS \cdot h(S) \cdot \log[\bar{R}/\epsilon])$.

s-Rectangular Ambiguity Sets: Bellman Operator

Distance-constrained s-rectangular ambiguity set

$$\mathcal{P} = \prod_{s \in \mathcal{S}} \mathcal{P}_s \quad \text{with} \quad \mathcal{P}_s = \left\{ p(\cdot | s, \cdot) : \sum_{a \in \mathcal{A}} d[p(\cdot | s, a), p^0(\cdot | s, a)] \leq \kappa \right\}$$

Theorem

Assume \mathfrak{P} can be computed **exactly** in time $\mathcal{O}(h(S))$.

Then \mathfrak{B} can be **computed to accuracy** $\epsilon > 0$ in time $\mathcal{O}(AS \cdot h(S) \cdot \log[\bar{R}/\epsilon])$.

Assume \mathfrak{P} can be **computed to any accuracy** $\delta > 0$ in time $\mathcal{O}(h(\delta))$. Then \mathfrak{B} can be **computed to accuracy** $\epsilon > 0$ in time $\mathcal{O}(AS \cdot h(\epsilon\kappa/[2A\bar{R} + A\epsilon]) \cdot \log[\bar{R}/\epsilon])$.

s-Rectangular Ambiguity Sets: Bellman Operator

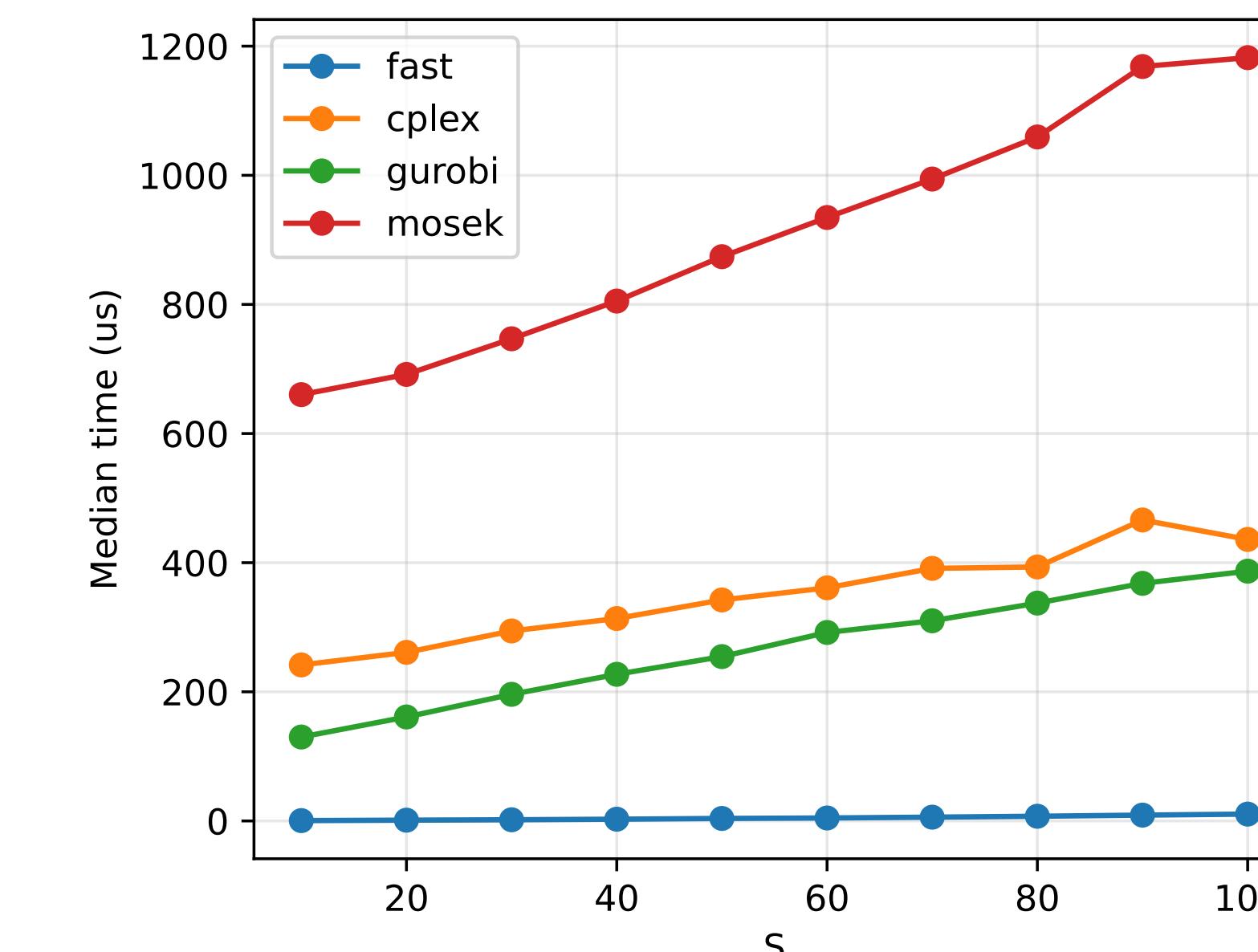
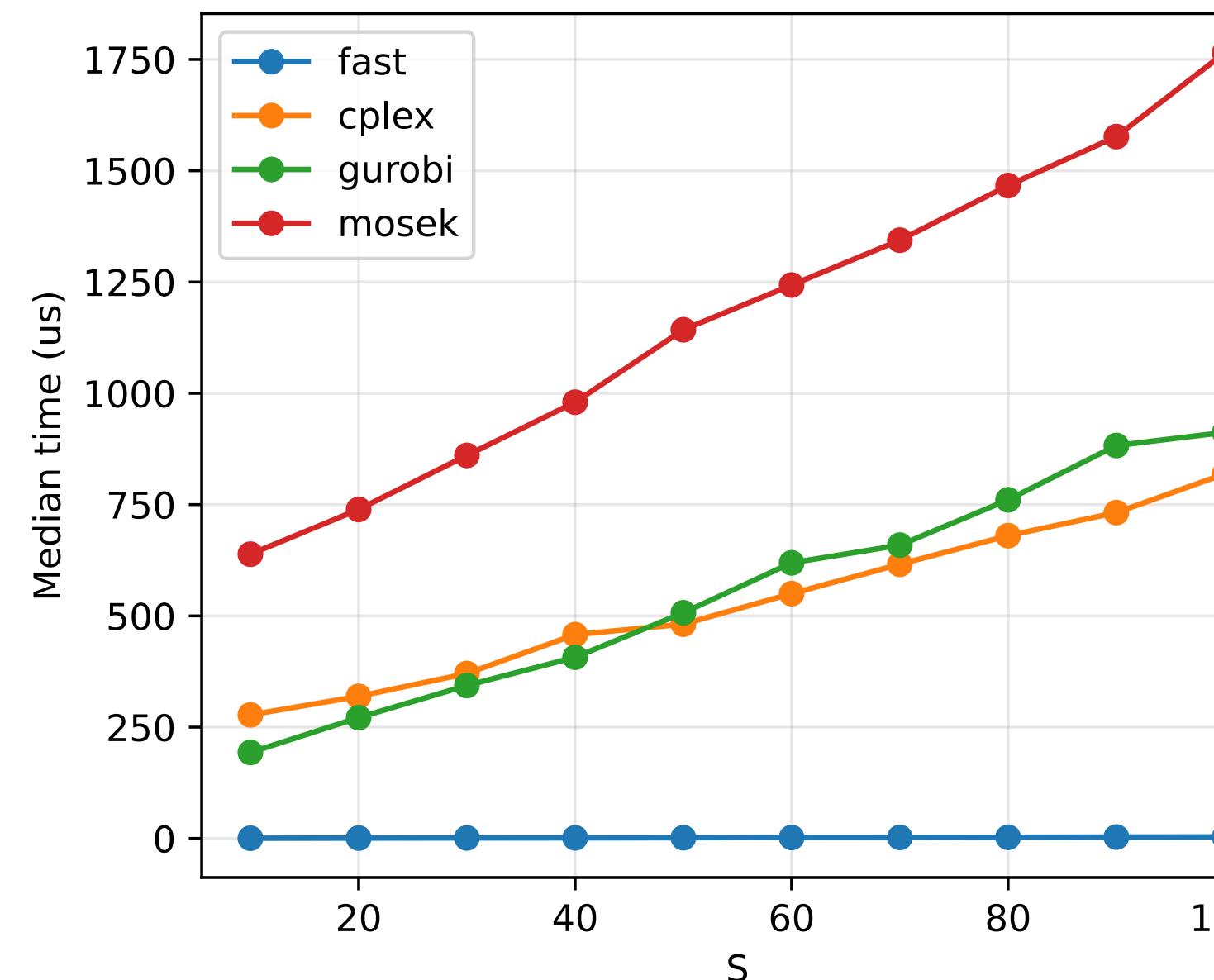
Divergence	$d_a(\cdot, p^0)$	Ours	Previous
KL-Divergence	$\sum_{s' \in \mathcal{S}} p(s' s, a) \cdot \log \left(\frac{p(s' s, a)}{p^0(s' s, a)} \right)$	$\mathcal{O}(S^2 A \cdot \log A)$	$\mathcal{O}(\ell^2 \cdot S^2 \cdot A)$
Burg Entropy	$\sum_{s' \in \mathcal{S}} p^0(s' s, a) \cdot \log \left(\frac{p^0(s' s, a)}{p(s' s, a)} \right)$	$\mathcal{O}(S^2 A \cdot \log A)$	(none)
Variation Distance	$\sum_{s' \in \mathcal{S}} p(s' s, a) - p^0(s' s, a) $	$\mathcal{O}(S^2 A \cdot \log S)$	$\mathcal{O}(S^2 A \cdot \log S)$
χ^2 -Distance	$\sum_{s' \in \mathcal{S}} \frac{[p(s' s, a) - p^0(s' s, a)]^2}{p^0(s' s, a)}$	$\mathcal{O}(S^2 A \cdot \log S)$	$\mathcal{O}(S^{4.5} \cdot A)$

s-Rectangular Ambiguity Sets: Bellman Operator

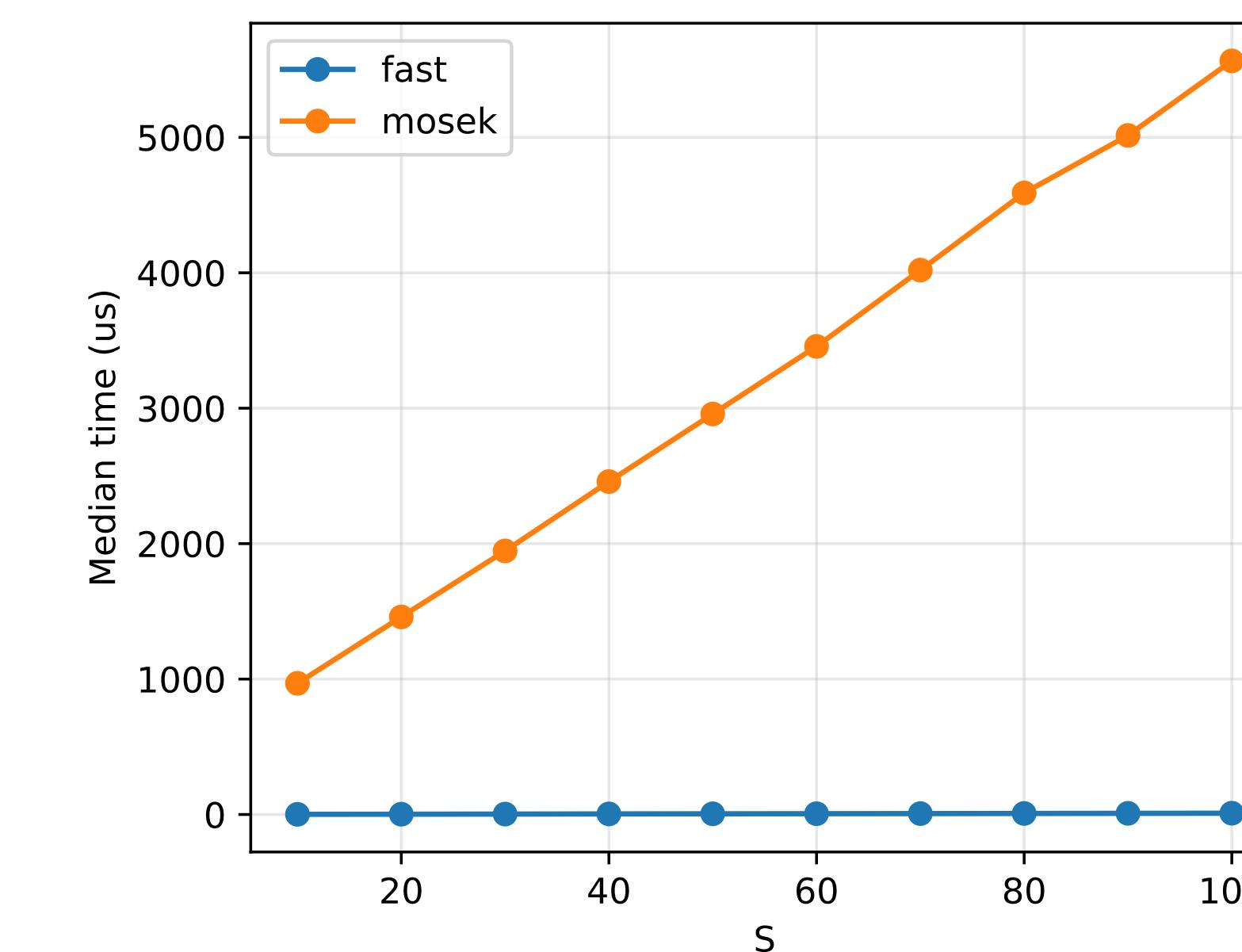
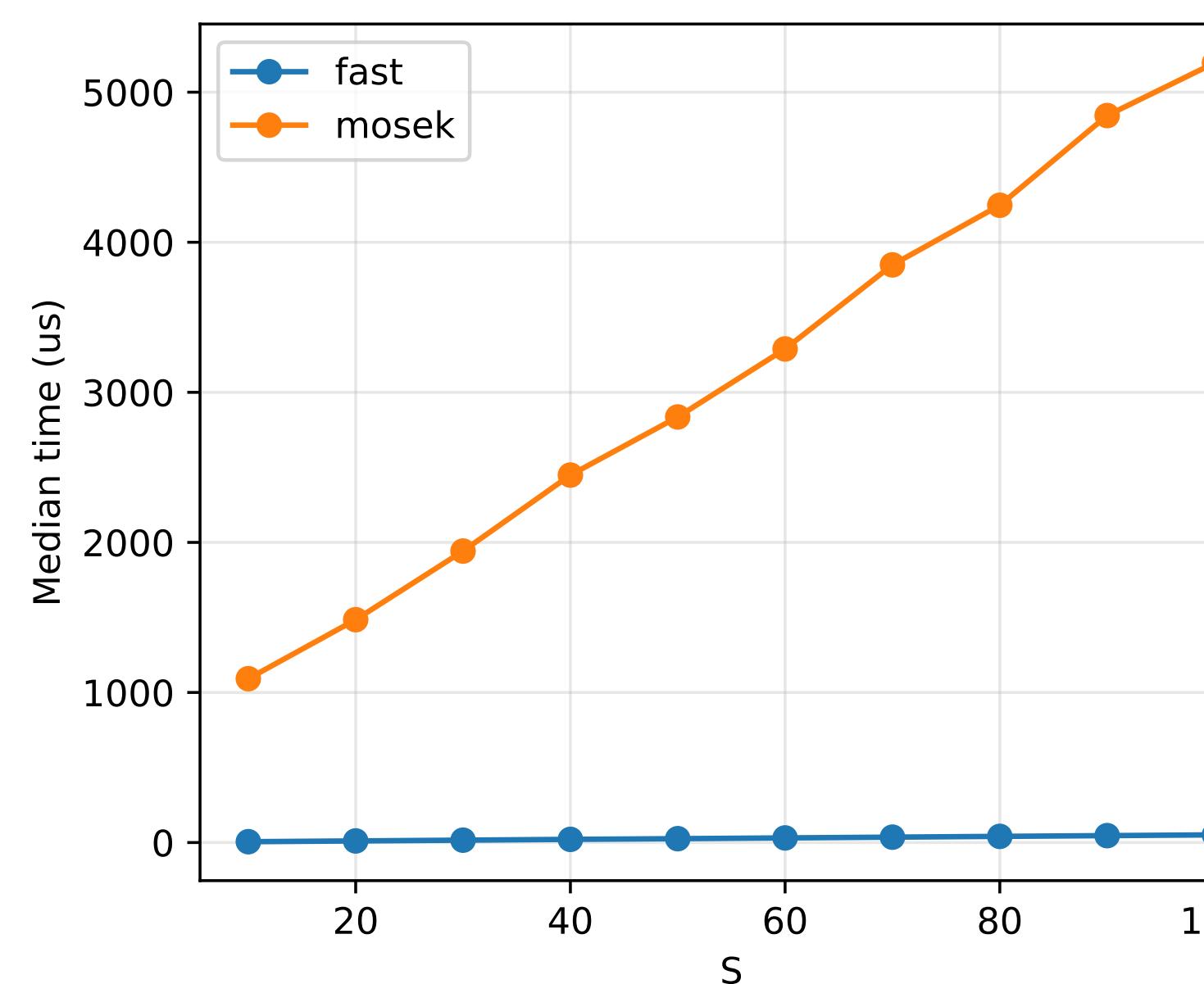
Divergence	$d_a(\cdot, p^0)$	Ours	Previous
KL-Divergence	$\sum_{s' \in \mathcal{S}} p(s' s, a) \cdot \log \left(\frac{p(s' s, a)}{p^0(s' s, a)} \right)$	$\mathcal{O}(S^2 A \cdot \log A)$	$\mathcal{O}(\ell^2 \cdot S^2 \cdot A)$
Burg Entropy	$\sum_{s' \in \mathcal{S}} p^0(s' s, a) \cdot \log \left(\frac{p^0(s' s, a)}{p(s' s, a)} \right)$	$\mathcal{O}(S^2 A \cdot \log A)$	(none)
Variation Distance	$\sum_{s' \in \mathcal{S}} p(s' s, a) - p^0(s' s, a) $	$\mathcal{O}(S^2 A \cdot \log S)$	$\mathcal{O}(S^2 A \cdot \log S)$
χ^2 -Distance	$\sum_{s' \in \mathcal{S}} \frac{[p(s' s, a) - p^0(s' s, a)]^2}{p^0(s' s, a)}$	$\mathcal{O}(S^2 A \cdot \log S)$	$\mathcal{O}(S^{4.5} \cdot A)$

s-Rectangular Ambiguity Sets: Projection Problem

1-norm



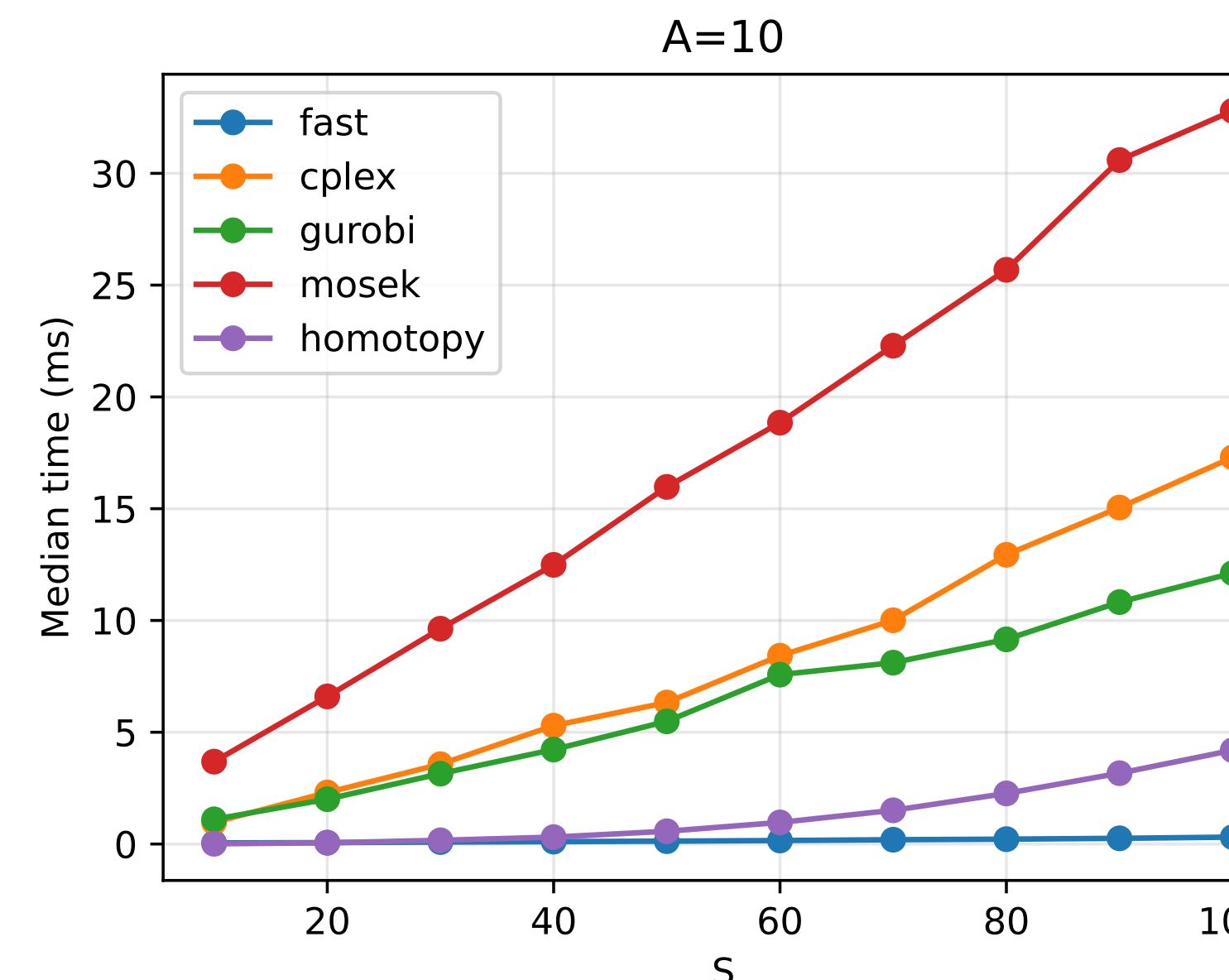
2-norm



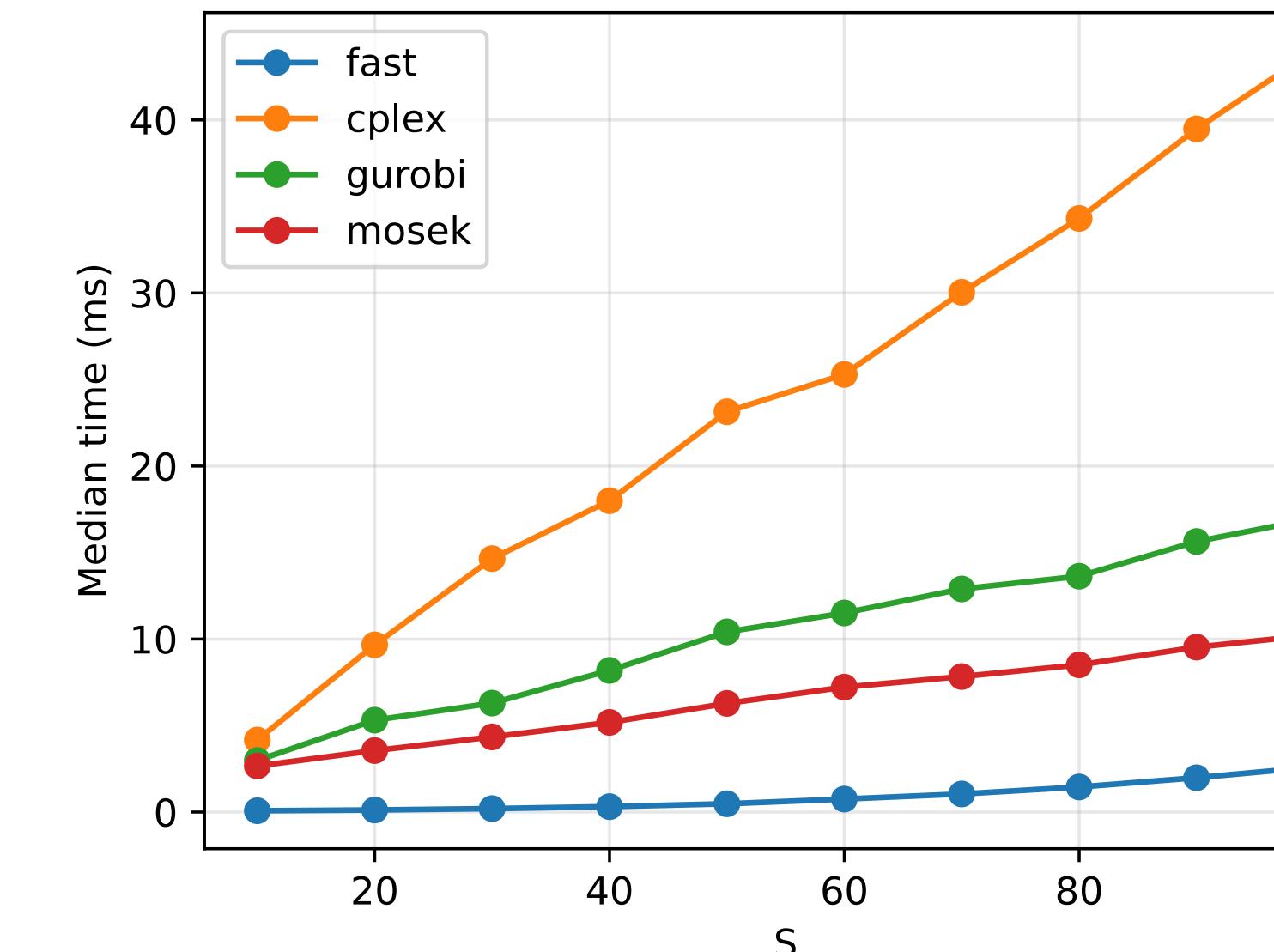
Breg

s-Rectangular Ambiguity Sets: Bellman Operator

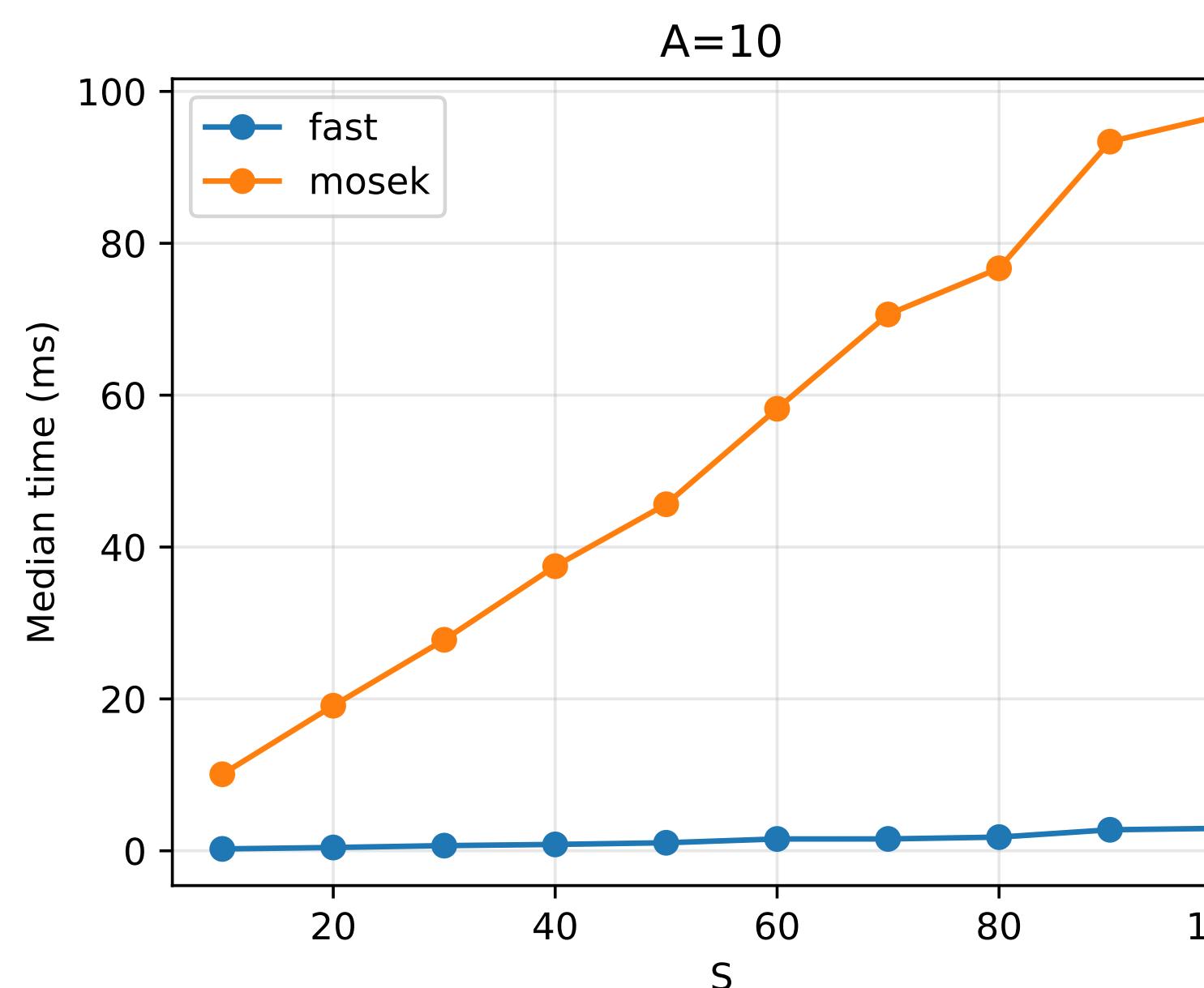
1-norm



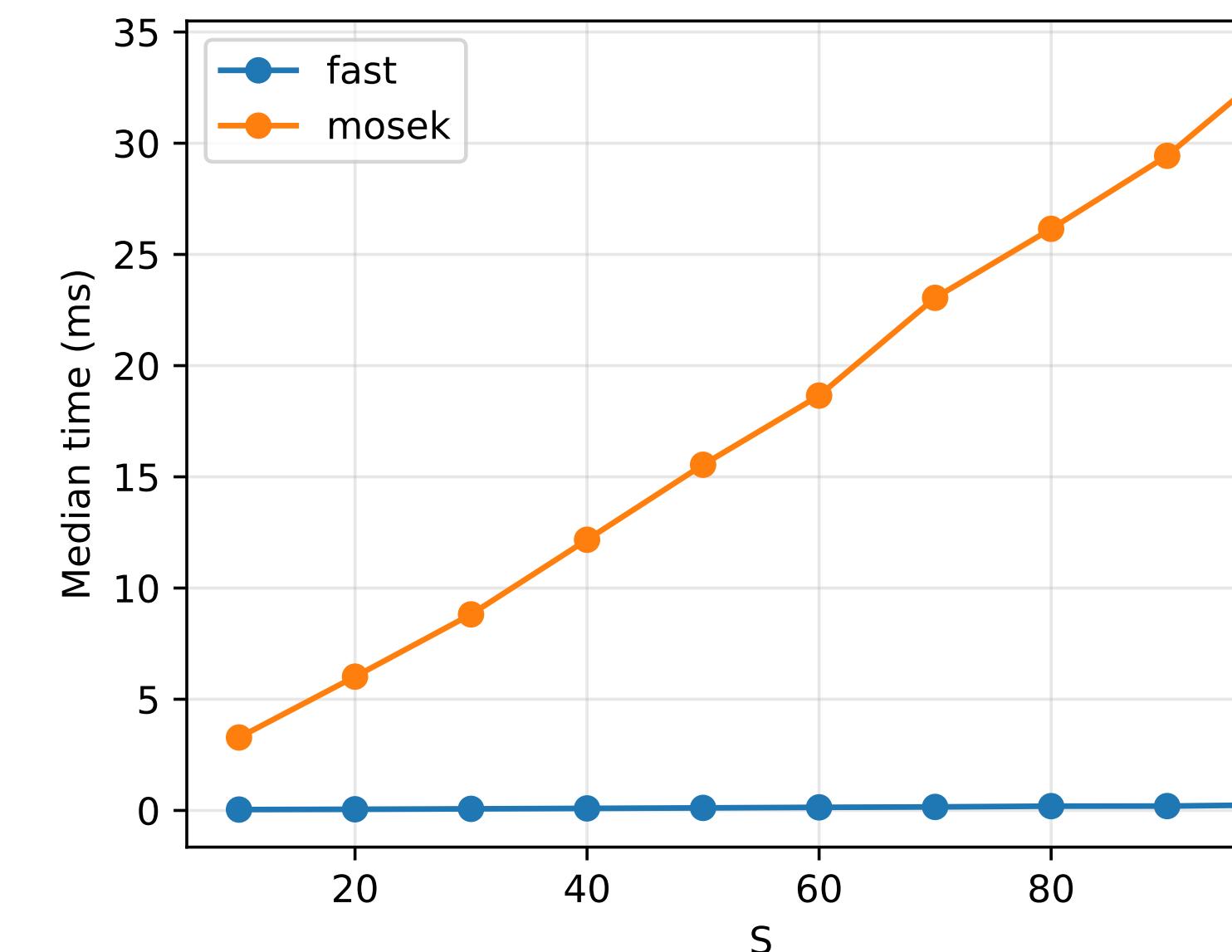
A=10



2-norm



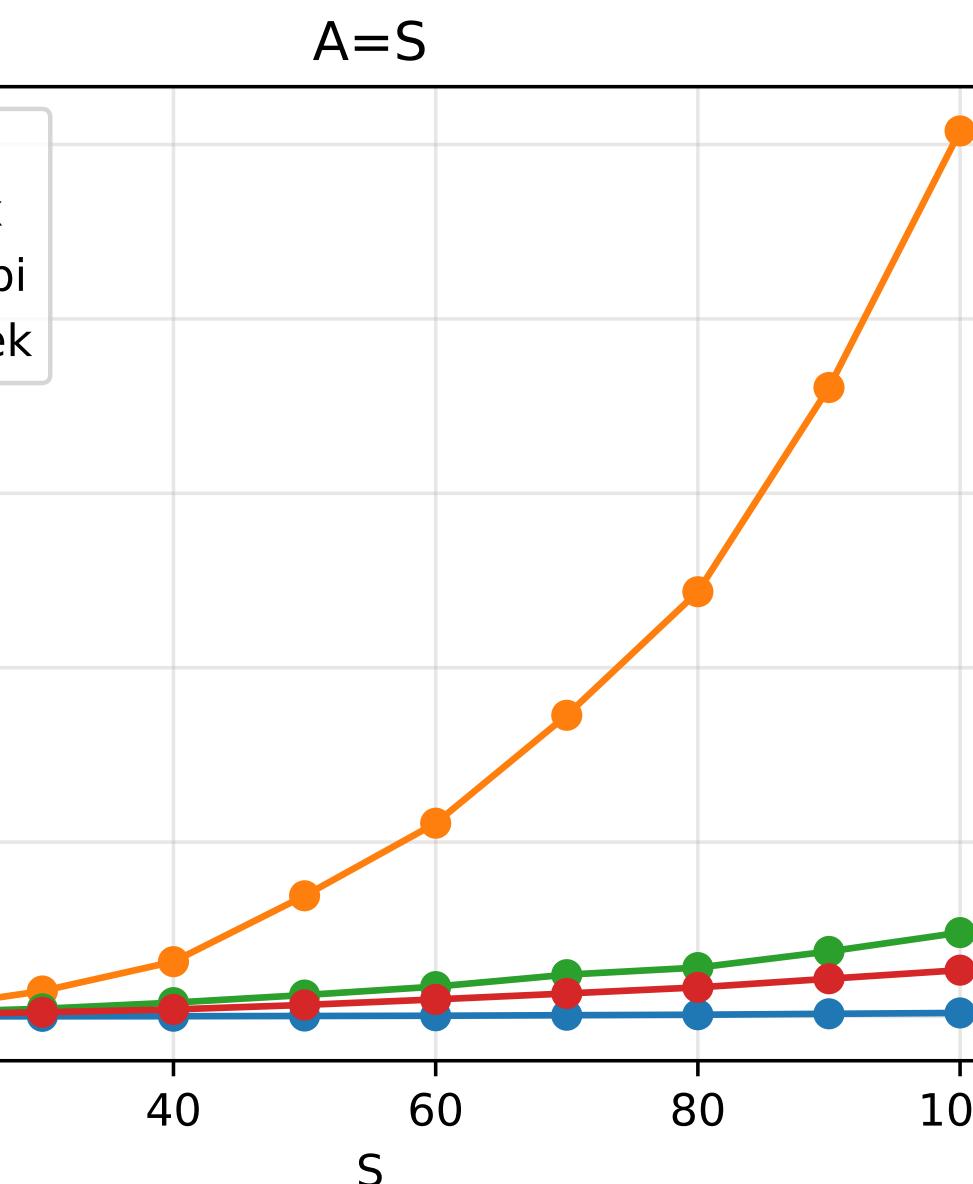
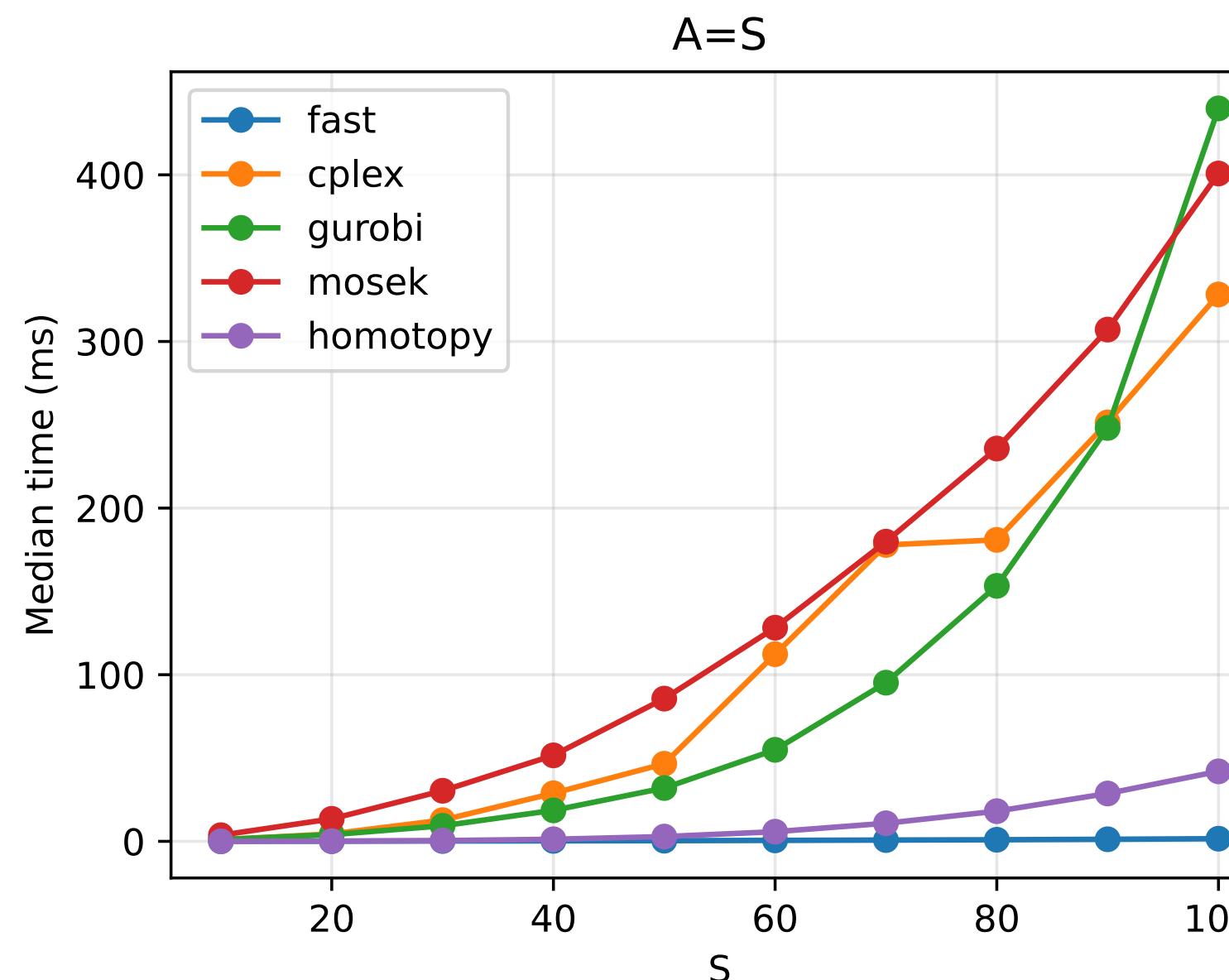
A=10



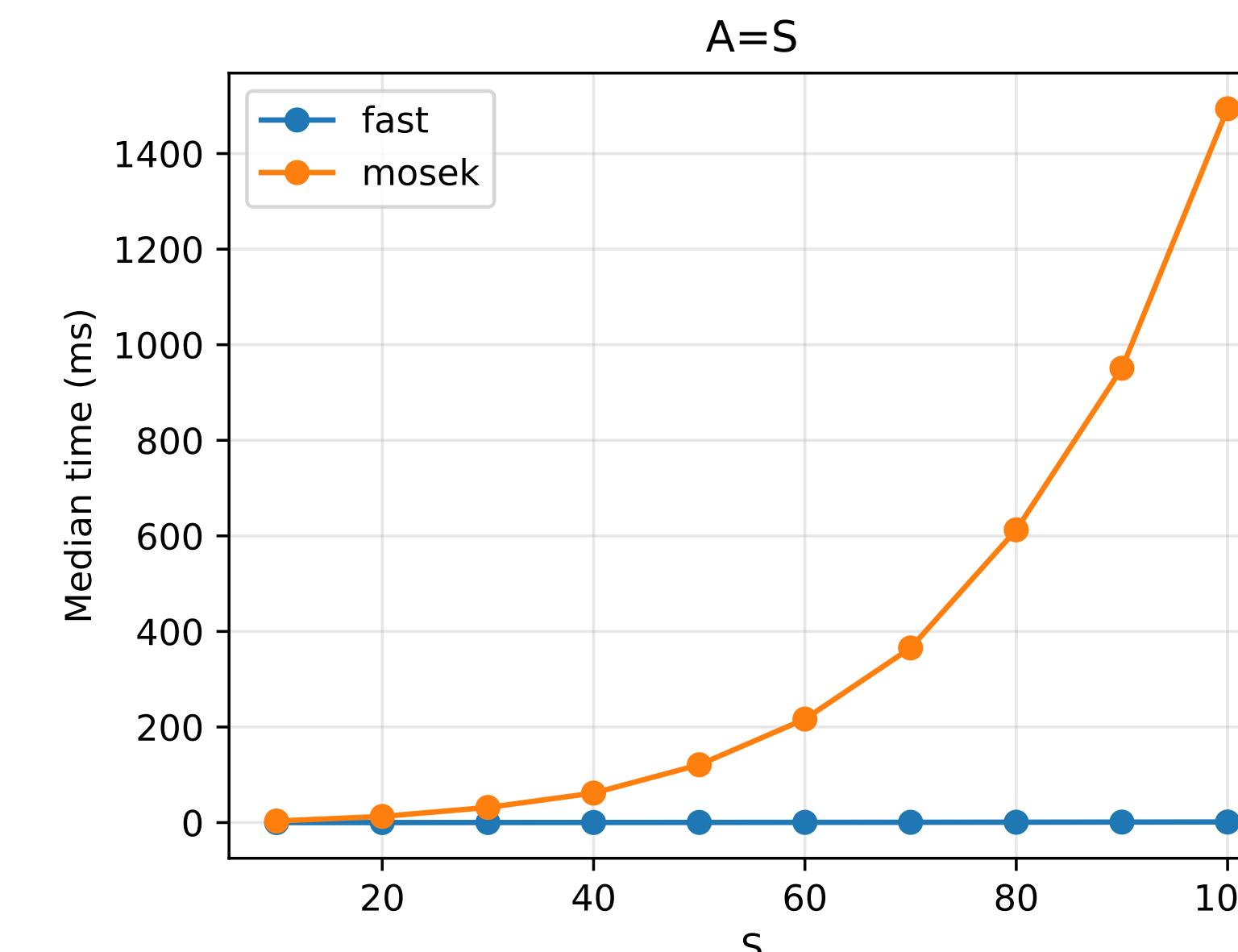
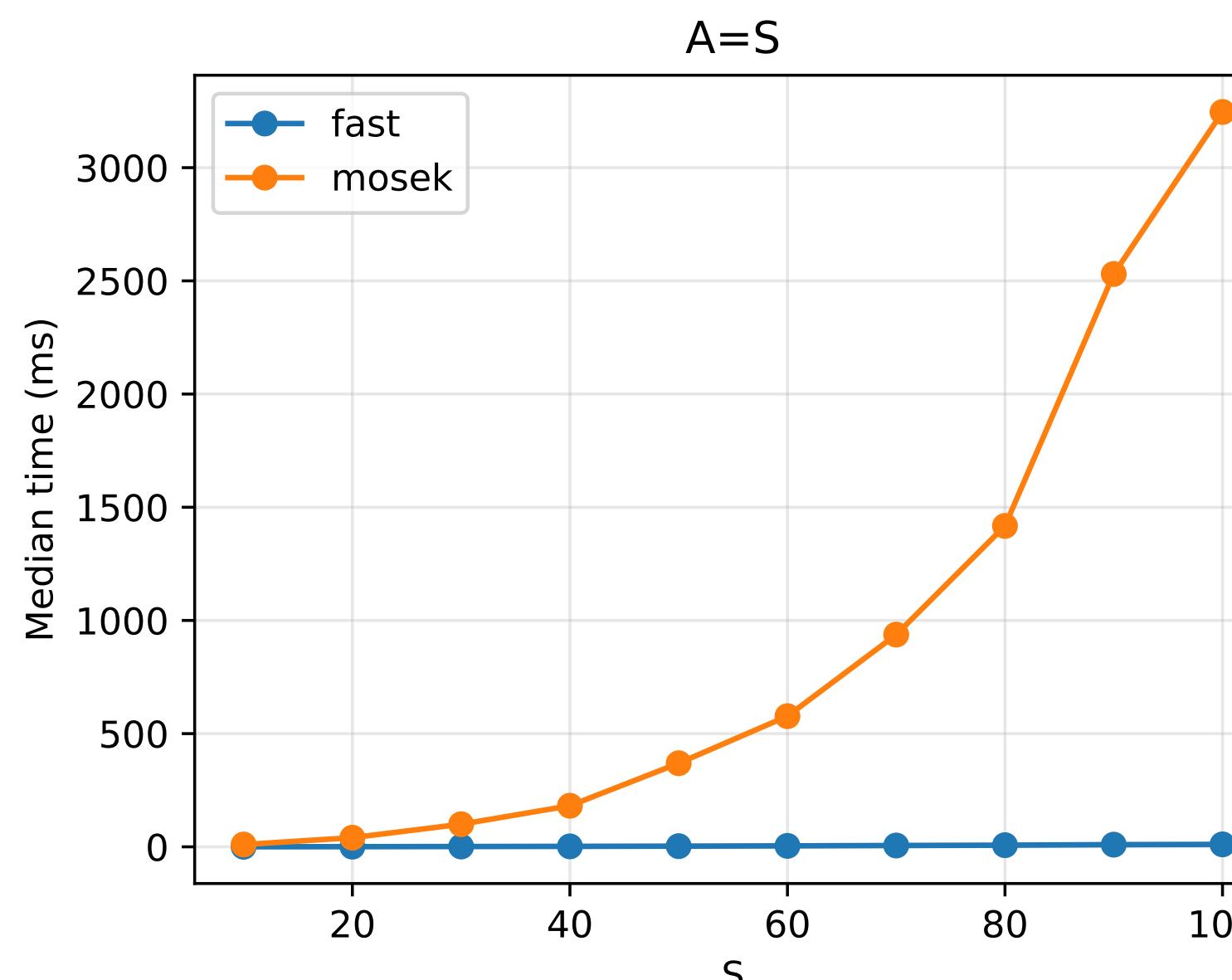
Burg

s-Rectangular Ambiguity Sets: Bellman Operator

1-norm



KL-Div



s-Rectangular Ambiguity Sets: Partial Policy Iteration

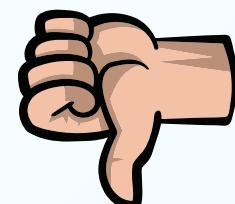
Reconsider idea of (modified) policy iteration:

s -Rectangular Ambiguity Sets: Partial Policy Iteration

Reconsider idea of (modified) policy iteration:

Policy improvement

$$[\mathfrak{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$



expensive operator: requires robust Bellman operator

s-Rectangular Ambiguity Sets: Partial Policy Iteration

Reconsider idea of (modified) policy iteration:

Policy improvement

$$[\mathfrak{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$



expensive operator: requires robust Bellman operator

Policy evaluation

$$[\mathfrak{B}(\pi)v](s) = \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$



cheaper operator: can be recast as Bellman operator
of a nominal MDP

s-Rectangular Ambiguity Sets: Robust Modified Policy Iteration

s-Rectangular Ambiguity Sets: Robust Modified Policy Iteration

Repeat: starting with $i = 0$, ν^0 arbitrary

Policy improvement

Compute $w^{i+1} = \mathcal{B}\nu^i$ and let π^{i+1} be the corresponding greedy policy



expensive operator:
robust value iteration

$$i = i + 1$$

Until $\|w^{i+1} - \vartheta^{i+1,N}\|_\infty < \frac{1-\lambda}{2} \cdot \delta$

s-Rectangular Ambiguity Sets: Robust Modified Policy Iteration

Repeat: starting with $i = 0$, ν^0 arbitrary

Policy improvement

Compute $w^{i+1} = \mathfrak{B}\nu^i$ and let π^{i+1} be the corresponding greedy policy

Policy evaluation

Compute sequence $\vartheta^{i+1,j+1} = \mathfrak{B}(\pi^{i+1})\vartheta^{i+1,j}$ with $\vartheta^{i+1,0} = w^{i+1}$ until
 $\|\vartheta^{i+1,j+1} - \vartheta^{i+1,j}\|_\infty \leq (1 - \lambda)\epsilon_{i+1}$

$i = i + 1$



cheaper operator:
no maximum involved

Until $\|w^{i+1} - \vartheta^{i+1,N}\|_\infty < \frac{1 - \lambda}{2} \cdot \delta$

s-Rectangular Ambiguity Sets: Partial Policy Iteration

Repeat: starting with $i = 0$, ν^0 arbitrary

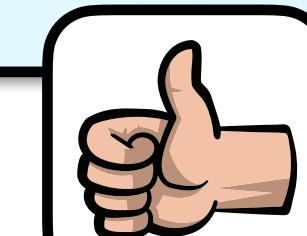
Policy improvement

Compute $w^{i+1} = \mathfrak{B}\nu^i$ and let π^{i+1} be the corresponding greedy policy

Policy evaluation

Alternate between **single robust Bellman evaluation** $\mathfrak{B}(\pi^{i+1})$ and
multiple nominal Bellman evaluations under **worst-case p** .

$$i = i + 1$$



cheap (!) operator:
nominal evaluations

Until $\|w^{i+1} - \vartheta^{i+1,N}\|_\infty < \frac{1-\lambda}{2} \cdot \delta$

s-Rectangular Ambiguity Sets: Partial Policy Iteration

Repeat: starting with $i = 0$, v^0 arbitrary

Policy improvement

Compute $w^{i+1} = \mathcal{B}v^i$ and π^{i+1} the corresponding greedy policy

Theorem

Assume $\epsilon_{i+1} < \lambda^c \cdot \epsilon_i$ for some $c > 1$. Then the optimality gap of partial policy iteration satisfies:

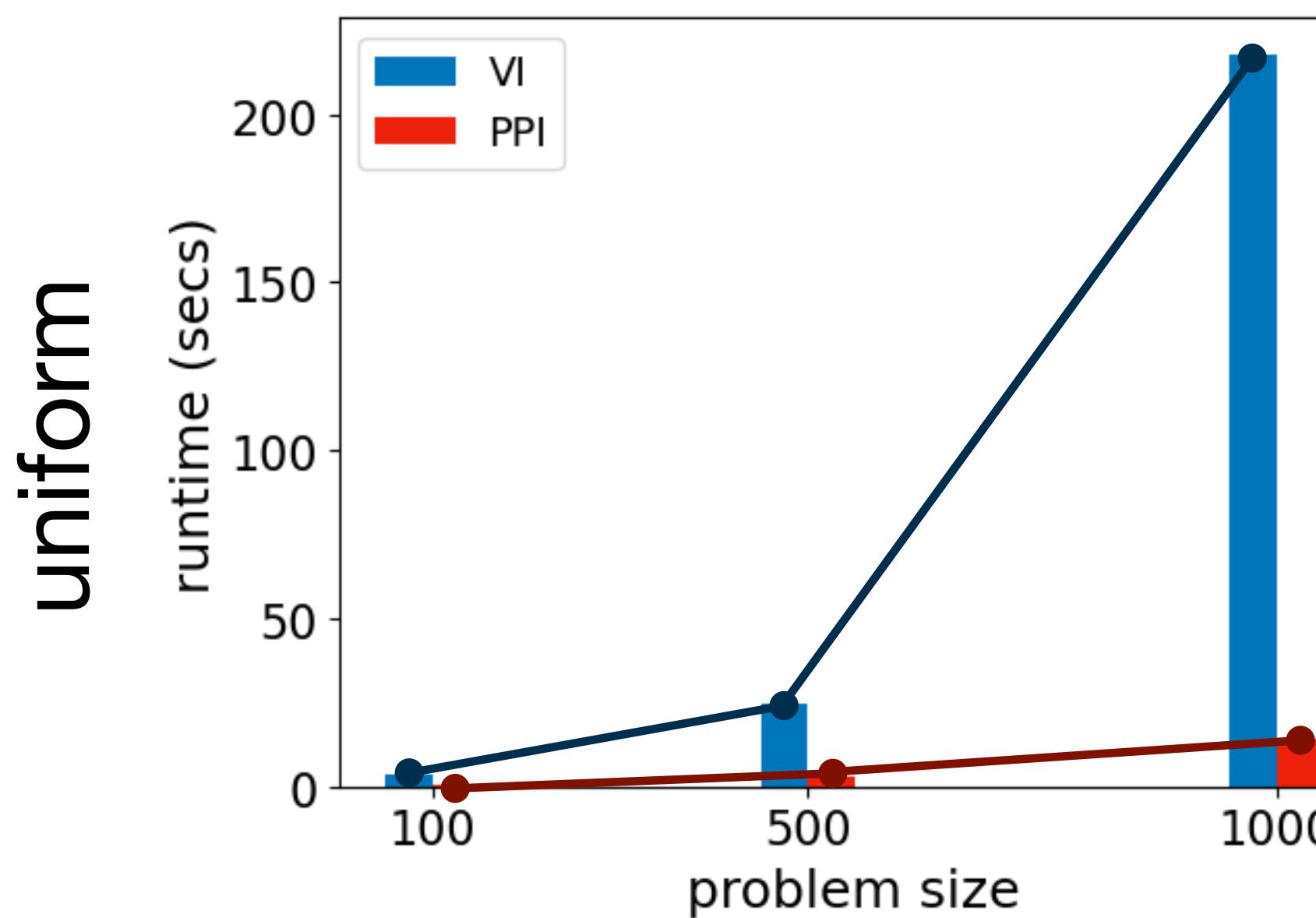
$$\|v(\pi^{i+1}) - v^\star\|_\infty \leq \lambda^i \left(\|v(\pi^1) - v^\star\|_\infty + \frac{2\epsilon_1}{(1 - \lambda^{c-1})(1 - \lambda)} \right)$$

$i = i + 1$

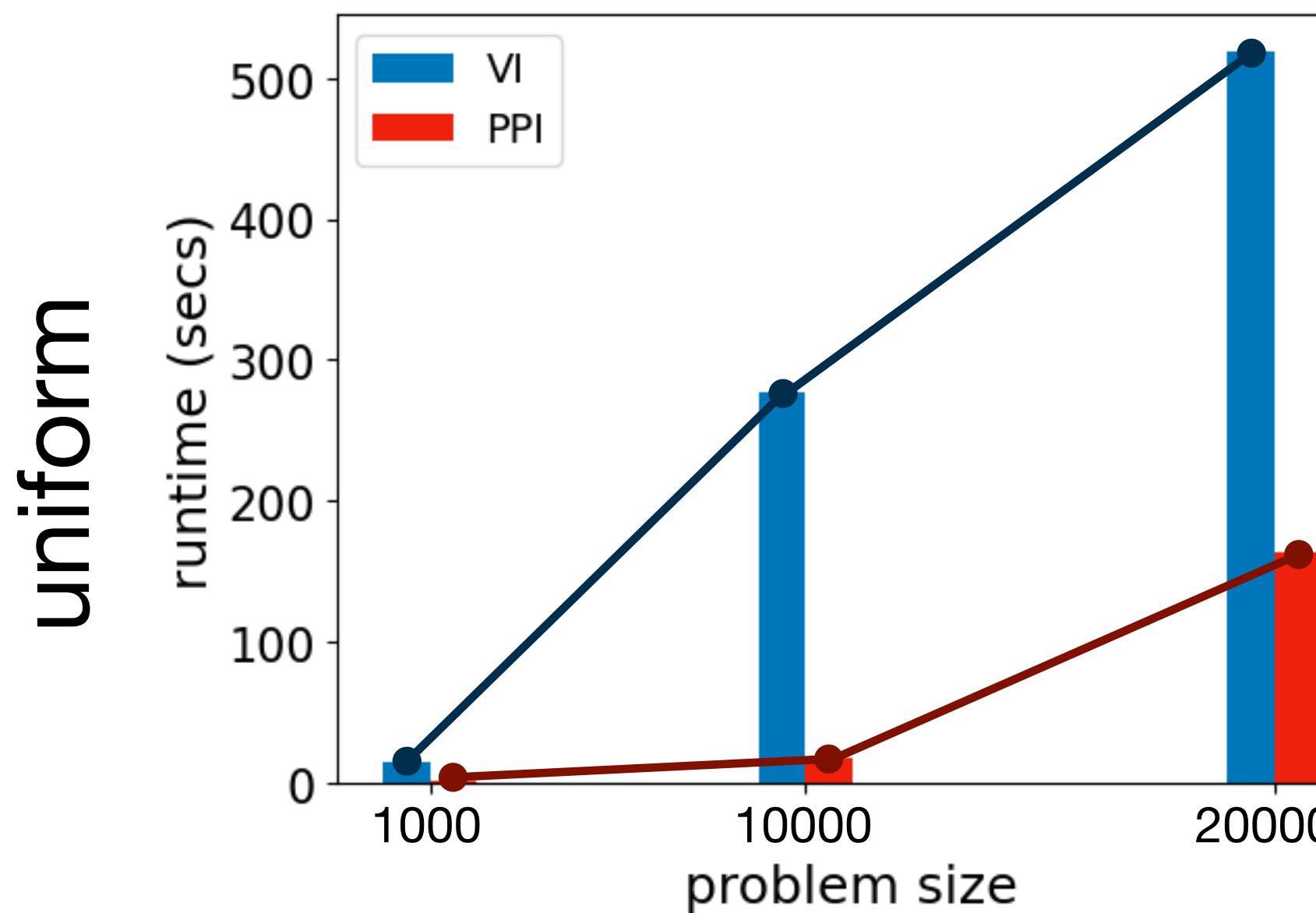
Until $\|w^{i+1} - \vartheta^{i+1,N}\|_\infty < \frac{1 - \lambda}{2} \cdot \delta$

s-Rectangular Ambiguity Sets: Partial Policy Iteration

Inventory



Cart Pole



Conclusions: MDPs – Now More Important Than Ever!

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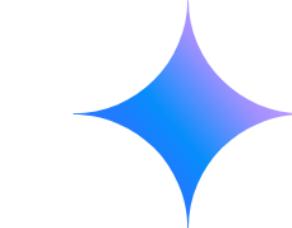
- Reinforcement learning is changing the world:

Conclusions: MDPs – Now More Important Than Ever!

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ChatGPT



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Claude

Reinforcement learning with human feedback

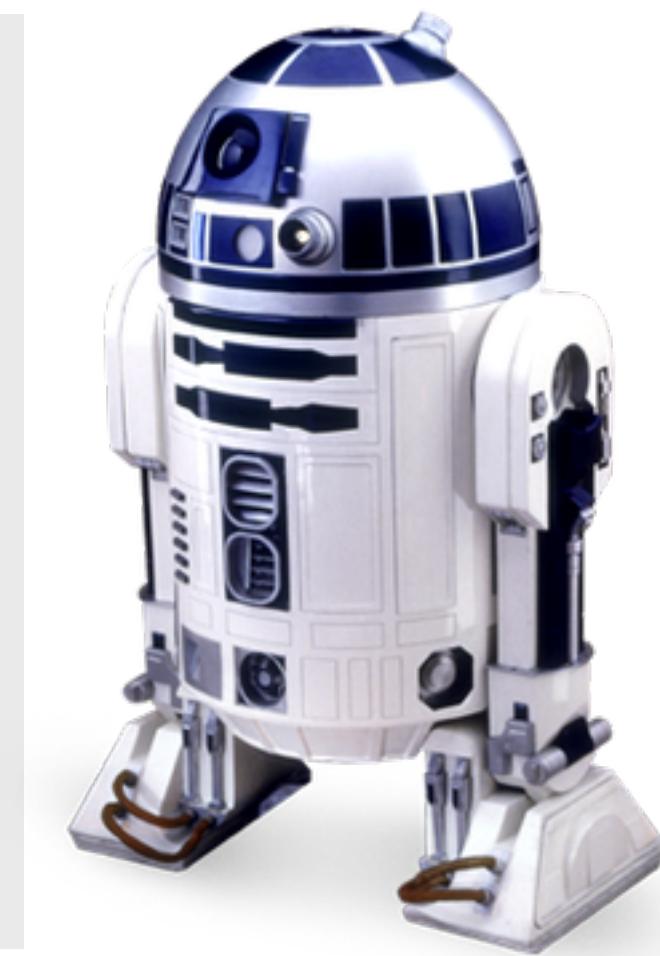
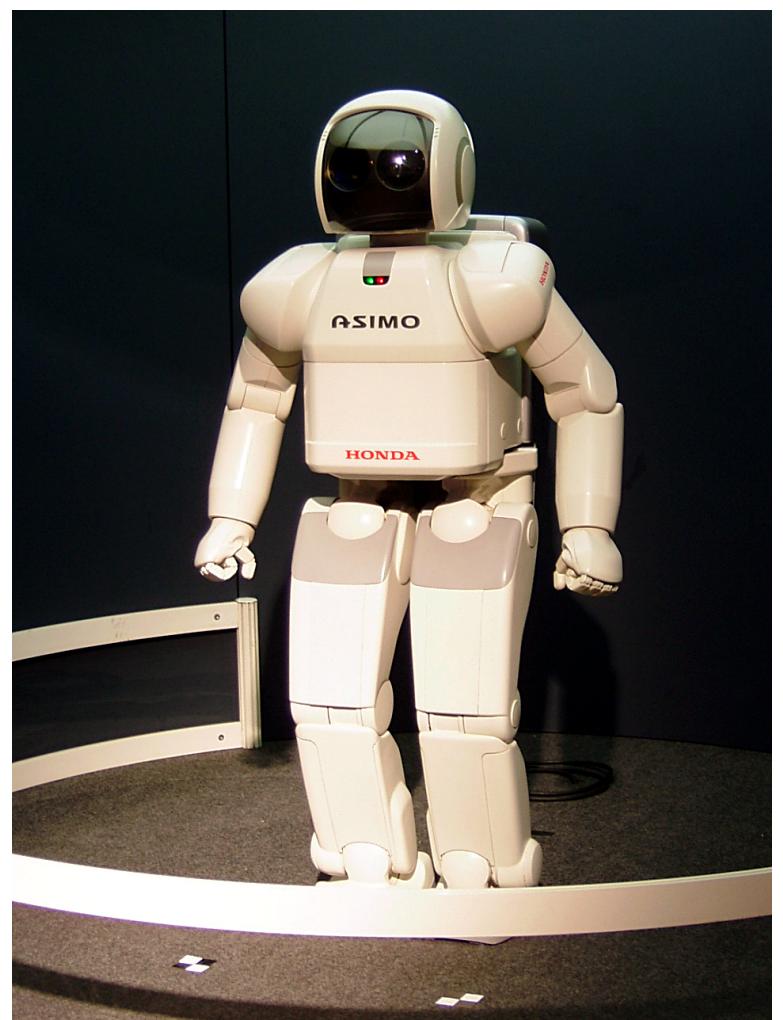
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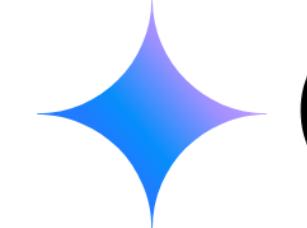
Robotics

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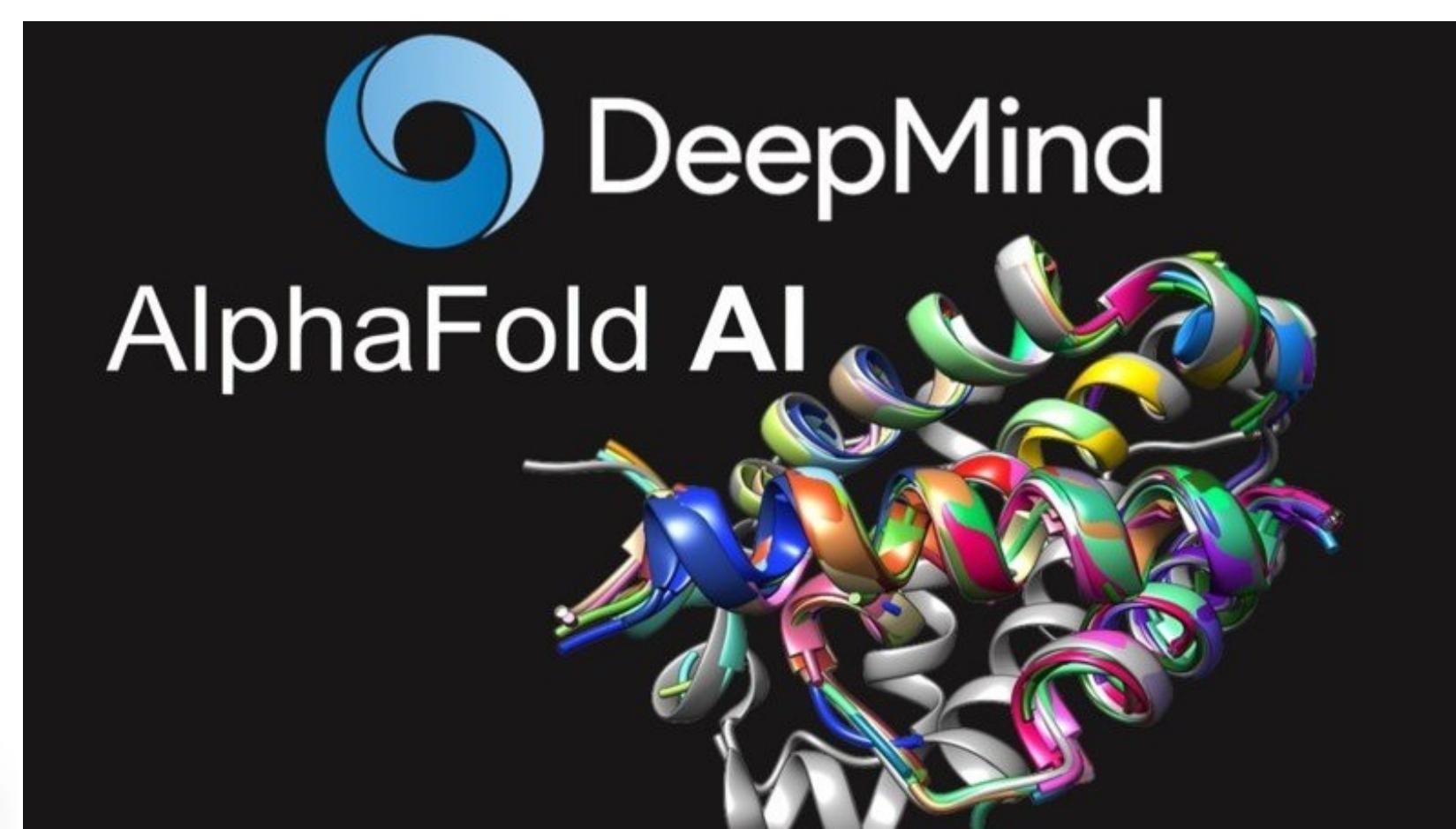
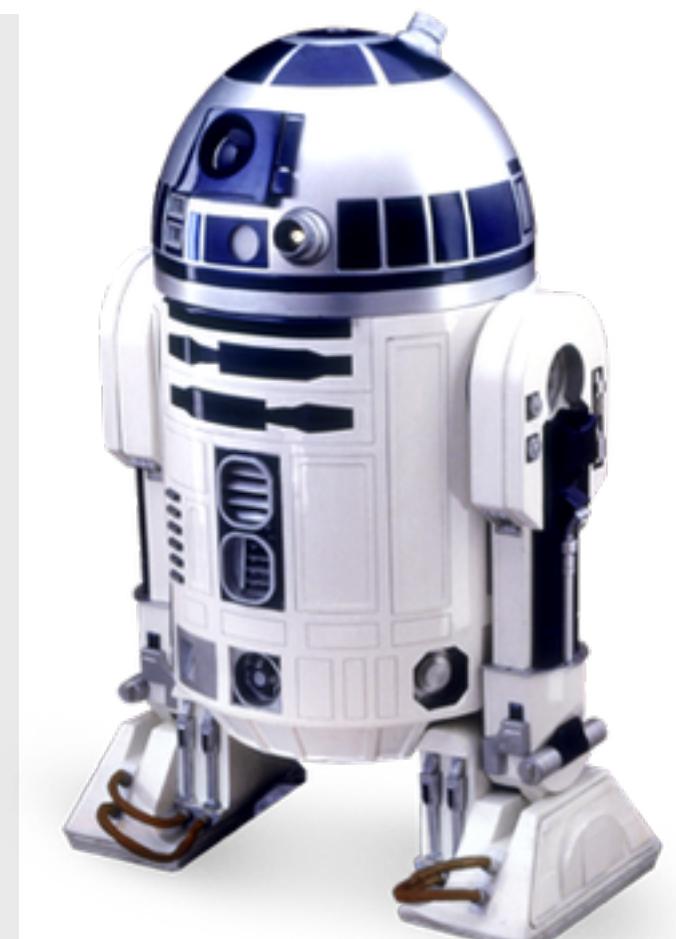
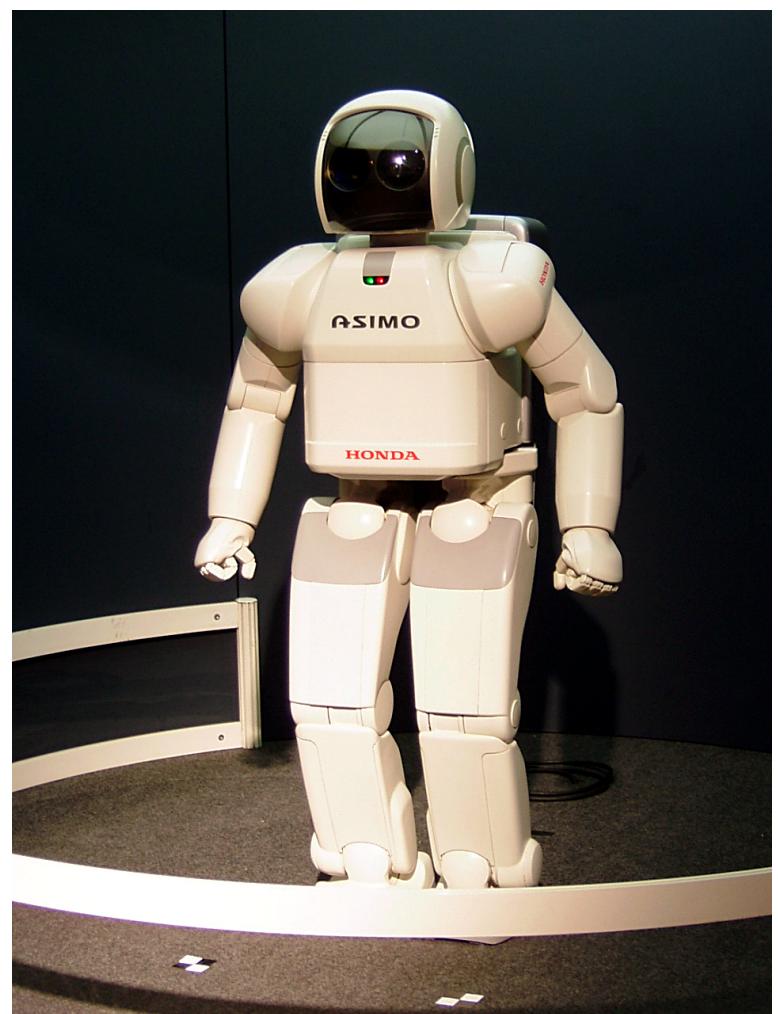


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Reinforcement learning with human feedback



Robotics

Medical discovery

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- MDPs give us the mathematical understanding that underpins all this.
- There are exciting new developments that bring both fields closer:
 - Linear MDPs: Reinforcement learning with linear dynamics
 - Weakly Coupled MDPs, Factored MDPs, Constrained MDPs, Safe RL, ...
 - covered extensively in top AI/ML conferences in recent years!

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