- 1. Consider the schedules S1, S2, S3, and S4 below. State whether each schedule is strict, cascadeless, and recoverable. Explain your reasons by referring to the lecture notes. (15*3 points)
 - a. S1: R3(D); R3(B); R1(C); R2(C); R2(A); R4(D); R1(C); W1(C); W4(C); R3(D); C2; A4; C3; A1;

T1	Т2	Т3	Т4
		R3(D)	
		R3(B)	
R1(C)			
	R2(C)		
	R2(A)		
			R4(D)
R1(C)			
W1(C)			
			W4(C)
		R3(D)	
	C2		
			A4
		C3	
A1			

It is recoverable since every commit that has a read operation is committed after the item that is read was written and committed before. It is cascadeless since every transaction reads items that are written in committed transactions. It is not strict since T4 writes on C after T1 writes on C, but T4 commits before T1.

b. S2: R1(A); W3(C); R4(A); R4(C); W1(B); R4(C); W2(B); R2(A); W1(B); W4(B); R4(C); R1(C); C3; C1; C4; C2;

|--|

R1(A)			
		W3(C)	
			R4(A)
			R4(C)
W1(B)			
			R4(C)
	W2(B)		
	R2(A)		
W1(B)			
			W4(B)
			R4(C)
R1(C)			
		C3	
C1			
			C4
	C2		

It is recoverable since every commit that has a read operation is committed after the item that is read was written and committed before. It is not cascadeless since T1 reads C which was updated by T3 before T3 commits. Since it is not cascadeless, it is not strict as well.

c. S3: R2(C); R3(B); R3(C); R2(A); W4(A); R3(D); W4(C); R2(D); R4(B); R3(B); R4(C); R4(D); W1(D); C2; R4(B); R3(B); R1(B); W1(D); C1; W3(D); R4(B); C4; R3(C); C3;

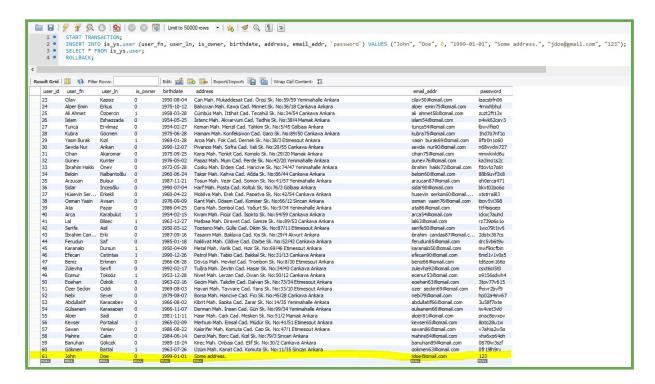
Т1	Т2	Т3	T4
	R2(C)		
		R3(B)	
		R3(C)	
	R2(A)		
			W4(A)
		R3(D)	

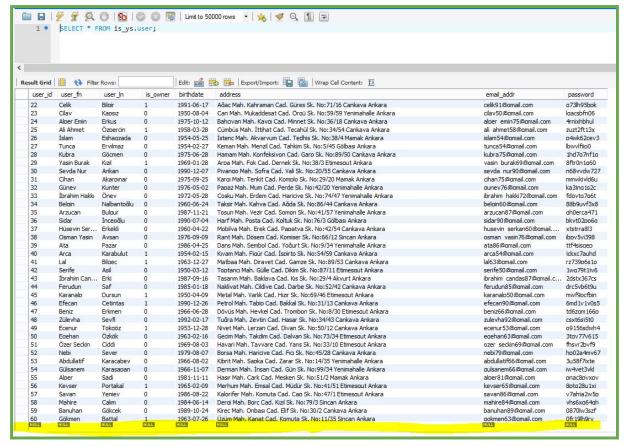
			W4(C)
	R2(D)		
			R4(B)
		R3(B)	
			R4(C)
			R4(D)
W1(D)			
	C2		
			R4(B)
		R3(B)	
R1(B)			
W1(D)			
C1			
		W3(D)	
			R4(B)
			C4
		R3(C)	
		C3	

It is recoverable since every commit that has a read operation is committed after the item that is read was written and committed before. It is cascadeless since every transaction reads items that are written in committed transactions. It is strict since no transaction reads or updates something that was updated by another transaction before that transaction commits.

- 2. While database management systems automatically handle transactions for us, we can manually create transactions and observe these concurrency problems. Follow the instructions below and observe how a phantom read occurs. You are supposed to use the database you created in your second assignment. (15 points)
 - Open MySQL Workbench, if you had not imported the previously provided schema (is_ys), import it.
 - Open a new SQL tab.
 - Start your transaction with "START TRANSACTION;" since we want to do it manually. You do not need to set the autocommit feature off since this statement temporarily turns it off (after a commit or rollback, it turns back to the automatic mode).
 - Write an insert query with the user table to create add a new user.

- Write a query to select all users.
- Finally, end your transaction with a rollback ("ROLLBACK;").
- Notice that since we rollback our transaction, nothing should change in the user table. Run this transaction. It should retrieve the users. Here, check the row you inserted. Although you did not commit, you should still see that your table has a new row. You are looking at a version of your table that never existed. Take a screenshot that features your query tab and the retrieved table. Highlight the row that looks inserted.
- Open another SQL tab, select the user table here as well. Notice that this table does not feature the new row since you did not commit. Take a screenshot of this version as well, and highlight how it differs from your previously retrieved table.
- Briefly explain what do transactions, commit, and rollback do.





- A transaction is a group of operations that has to be treated as a whole. If one operation fails, the other operations should fail as well, so that the data integrity is protected.
- Commit confirms your operations and saves the changes while rollback reverts the operations, achieving the last committed state of the database.
- 3. Consider the schedules S4, S5, and S6 below. State whether each schedule is conflict serializable and view serializable. Draw the precedence graph and show the related data items on arrows. If it is view serializable, show the equivalent serial schedule(s). Explain your reasons by referring to the lecture notes.
 - a. S4: W2(A); W3(C); W2(D); R2(C); R1(B); R2(A); R1(D); R3(B); R1(C); R2(B); C3; C2; C1; (10 points)

T1	Т2	Т3
	W2(A)	
		W3(C)
	W2(D)	
	R2(C)	
R1(B)		
	R2(A)	
R1(D)		

		R3(B)
R1(C)		
	R2(B)	
		C3
	C2	
C1		

Candidate equivalent serial schedules:

T1→T2→T3

 $T1 \rightarrow T3 \rightarrow T2$

T2→T3→T1

T2→T1→T3

T3→T1→T2

T3→T2→T1

Initial read:

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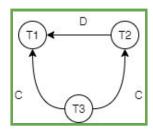
Final write:

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Update read:

T1 and T2 must read C after T3 writes on it, so T3 must come first. T1 must read D after T2 writes on it.

It is view serializable. Its equivalent serial schedule is $T3 \rightarrow T2 \rightarrow T1$.



Since there is no cycle, it is conflict serializable as well, which confirms its view serializability.

b. S5: W1(C); R3(C); W1(A); R1(C); R3(B); W2(C); W2(B); R1(C); R2(B); R1(B); R3(C); C1; C3; C2; (12 points)

Т1	Т2	Т3
W1(C)		

		R3(C)
W1(A)		
R1(C)		
		R3(B)
	W2(C)	
	W2(B)	
R1(C)		
	R2(B)	
R1(B)		
		R3(C)
C1		
		C3
	C2	

Candidate equivalent serial schedules:

 $\begin{array}{c} T1 \rightarrow T2 \rightarrow T3 \\ \hline T1 \rightarrow T3 \rightarrow T2 \end{array}$

T2→T3→T1

T2→T1→T3

12 11 13

 $\begin{array}{c} T3 \rightarrow T1 \rightarrow T2 \\ \hline T3 \rightarrow T2 \rightarrow T1 \end{array}$

Initial read:

T3 must read B before T2 writes on it, so T2 cannot come before T3.

Final write:

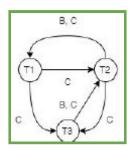
T2 must write on C after T1, so T1 cannot come before T2.

Update read:

T3 must read C after T1 writes on it, so T1 must come before T3.

T3 must read C after T2 writes on it, so T1 must come before T3.

These conditions eliminate all of our candidate equivalent schedules. Therefore, this schedule is not view serializable.



Since it is not view serializable and there are cycles in the graph, it is not conflict serializable either.

c. S6: R1(C); R2(D); R4(A); W1(C); R2(A); R3(B); R3(D); W1(C); W2(D); R1(A); R4(C); R3(A); W2(A); R4(A); R3(C); R2(A); W4(C); R4(A); C4; C2; C1; C3; (18 points)

T1	Т2	Т3	Т4
R1(C)			
	R2(D)		
			R4(A)
W1(C)			
	R2(A)		
		R3(B)	
		R3(D)	
W1(C)			
	W2(D)		
R1(A)			
			R4(C)
		R3(A)	
	W2(A)		
			R4(A)
		R3(C)	
	R2(A)		
			W4(C)
			R4(A)
			C4

	C2		
C1			
		C3	

Candidate equivalent serial schedules:

 $T1 \rightarrow T2 \rightarrow T3 \rightarrow T4$

 $T1 \rightarrow T2 \rightarrow T4 \rightarrow T3$

 $T1 \rightarrow T3 \rightarrow T2 \rightarrow T4$

 $T1 \rightarrow T3 \rightarrow T4 \rightarrow T2$

 $T1 \rightarrow T4 \rightarrow T2 \rightarrow T3$

 $T1 \rightarrow T4 \rightarrow T3 \rightarrow T2$

 $T2 \rightarrow T1 \rightarrow T3 \rightarrow T4$

 $T2 \rightarrow T1 \rightarrow T4 \rightarrow T3$

 $T2 \rightarrow T3 \rightarrow T1 \rightarrow T4$

 $T2 \rightarrow T3 \rightarrow T4 \rightarrow T1$

 $T2 \rightarrow T4 \rightarrow T1 \rightarrow T3$

 $T2 \rightarrow T4 \rightarrow T3 \rightarrow T1$

 $T3 \rightarrow T1 \rightarrow T2 \rightarrow T4$

 $T3 \rightarrow T1 \rightarrow T4 \rightarrow T2$

 $T3 \rightarrow T4 \rightarrow T1 \rightarrow T2$

 $T3 \rightarrow T4 \rightarrow T2 \rightarrow T1$

 $T3 \rightarrow T2 \rightarrow T1 \rightarrow T4$

 $\overrightarrow{\text{T3} \rightarrow \text{T2} \rightarrow \text{T4} \rightarrow \text{T1}}$

 $T4 \rightarrow T1 \rightarrow T2 \rightarrow T3$

 $T4 \rightarrow T1 \rightarrow T3 \rightarrow T2$

 $T4 \rightarrow T2 \rightarrow T3 \rightarrow T1$

 $T4 \rightarrow T2 \rightarrow T1 \rightarrow T3$

 $T4 \rightarrow T3 \rightarrow T1 \rightarrow T2$

 $T4 \rightarrow T3 \rightarrow T2 \rightarrow T1$

Initial read:

T1 must read C before T4 writes on it, so T4 cannot come before T1.

T4 must read D before T2 writes on it, so T2 cannot come before T4.

T4 must read A before T2 writes on it, so T2 cannot come before T4.

Final write:

T4 must write on C after T1, so T4 cannot come before T1.

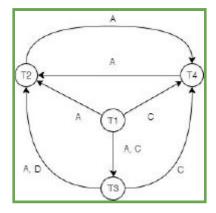
Update read:

T4 must read C after T1 writes on it, so T4 cannot come before T1.

T3 must read C after T1 writes on it, so T3 cannot come before T1.

T4 must read A after T2 writes on it, so T4 cannot come before T2.

These conditions eliminate all of our candidate equivalent schedules. Therefore, this schedule is not view serializable.



Since it is not view serializable and there is a cycle, it is not conflict serializable either.