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Preface

This is our Team Notebook for ACM ICPC and other Competitive Programming contests. Notable sources are:

- Introduction to Algorithm 3rd edition
- Competitive Programming 2 by Felix and Steven Halim
- Topcoder Algorithm Tutorials
- <https://sites.google.com/site/indy256/>
- <http://stanford.edu/~liszt90/acm/notebook.html>
- Dongskar Pedongi's Team Notebook
- Google, Wikipedia

Regards,

DELAPAN.3gp

Aufar Gilbran, Ahmad Zaky, Muntaha Ilmi

Institut Teknologi Bandung, Indonesia

Template

C++

```
#include <bits/stdc++.h> <vector> <map> <set> <queue> <deque> <stack>
<algorithm> <sstream> <iostream> <iomanip> <fstream> <cstring> <cmath>
<cstdlib> <ctime> <cassert> <limits> <numeric> <utility>
using namespace std;

#ifdef DEBUG
    #define debug(...) printf(__VA_ARGS__)
    #define GetTime() fprintf(stderr, "Running time: %.3lf
second\n", ((double)clock())/CLOCKS_PER_SEC)
#else
    #define debug(...)
    #define GetTime()
#endif

//type definitions
typedef long long ll;
typedef double db;
typedef pair<int,int> pii;
typedef vector<int> vint;

//abbreviations
#define A first
#define B second
#define F first
#define S second
```

```
#define MP make pair
#define PB push_back

//macros
#define REP(i,n) for (int i = 0; i < (n); ++i)
#define REPD(i,n) for (int i = (n)-1; 0 <= i; --i)
#define FOR(i,a,b) for (int i = (a); i <= (b); ++i)
#define FORD(i,a,b) for (int i = (a); (b) <= i; --i)
#define FORIT(it,c) for (__typeof ((c).begin()) it = (c).begin(); it !=
(c).end(); it++)
#define ALL(a) (a).begin(), (a).end()
#define SZ(a) ((int)(a).size())
#define RESET(a,x) memset(a,x,sizeof(a))
#define EXIST(a,s) ((s).find(a) != (s).end())
#define MX(a,b) a = max((a),(b));
#define MN(a,b) a = min((a),(b));

inline void OPEN(const string &s) {
    freopen((s + ".in").c_str(), "r", stdin);
    freopen((s + ".out").c_str(), "w", stdout);
}

/* ----- end of template ----- */
```

Graph Theory

Articulation Point

```
/** Articulation Point */
/* complexity : O(|V| + |E|) */

#define MAXN 100100

int n, m, low[MAXN], num[MAXN], parent[MAXN], art[MAXN], root,
rootChildren, counter;
vector<int> adj[MAXN];

void dfs(int u) {
    low[u] = num[u] = counter++;
    FORIT(it, adj[u]) {
        int v = *it;
        if (num[v] == -1) {
            parent[v] = u;
            if (u == root) rootChildren++;
            dfs(v);
            if (low[v] >= num[u]) art[u] = 1;
            MN(low[u], low[v]);
        }
        else if (v != parent[u]) {
            MN(low[u], num[v]);
        }
    }
}
```

```

    }
}

int main() {
    // read the graph here. It should be 0-indexed
    // initialization
    counter = 0;
    REP(i, n) {
        num[i] = -1;
        low[i] = parent[i] = art[i] = 0;
    }
    // perform the dfs
    REP(i, n) {
        if (num[i] == -1) {
            root = i, rootChildren = 0;
            dfs(i);
            art[root] = (rootChildren > 1);
        }
    }
    // now the articulation points are stored in art[]
    return 0;
}

```

Articulation Bridge

```

/** Bridge **/
/* complexity : O(|V| + |E| + |E| log |E|) */

#define MAXN 100100

int n, low[MAXN], num[MAXN], parent[MAXN], bridge[MAXN], counter;
vector<pii> adj[MAXN]; // adj[u].PB(MP(v, idx_of_edge));

void dfs(int u) {
    low[u] = num[u] = counter++;
    FORIT(it, adj[u]) {
        int v = it->A;
        if (num[v] == -1) {
            parent[v] = u;
            dfs(v);
            if (low[v] > num[u]) bridge[it->B] = 1;
            MN(low[u], low[v]);
        }
        else if (v != parent[u]) {
            MN(low[u], num[v]);
        }
    }
}

int main() {
    // read the graph here. it should be 0-indexed
    // should not work if multiple edges exist
    // initialization

```

```

    counter = 0;
    REP(i, n) {
        num[i] = -1;
        low[i] = parent[i] = 0;
    }
    REP(i, m) {
        bridge[i] = 0;
    }
    // perform the dfs
    REP(i, n) {
        if (num[i] == -1) {
            dfs(i);
        }
    }
    // the bridges are stored in bridge[]

    return 0;
}

```

Tarjan's Directed SCC

```

/** Tarjan's Directed Strongly Connected Component **/
/* complexity : O(|V| + |E|) */

#define MAXN 100100

int n, low[MAXN], num[MAXN], visited[MAXN], counter;
vector<int> adj[MAXN], s;
vector<vector<int>> scc;

void dfs(int u) {
    low[u] = num[u] = counter++;
    s.PB(u);
    visited[u] = 1;
    FORIT(it, adj[u]) {
        int v = *it;
        if (num[v] == -1) dfs(v);
        if (visited[v]) {
            MN(low[u], low[v]);
        }
    }
    if (low[u] == num[u]) {
        vector<int> temp;
        int v = -1;
        while (u != v) {
            v = s.back(); s.pop_back(); visited[v] = 0;
            temp.PB(v);
        }
        scc.PB(temp);
    }
}

int main() {

```

```

// read the graph here. it should be 0-indexed
// initialization
counter = 0;
scc.clear();
REP(i, n) {
    num[i] = -1;
    low[i] = visited[i] = 0;
}
// perform the dfs
REP(i, n) {
    if (num[i] == -1) {
        dfs(i);
    }
}
// the components are stored in scc
return 0;
}

```

Max Flow

```

#define MAXN 1100
#define INF 0x3FFFFFFF

int res[MAXN][MAXN], vis[MAXN];

/** Maximum Flow **/
/* Edmond Karp | complexity : O(|V|*(|V|+|E|)) */
void augment(int v, int minEdge, int &s, int &f, vector<int> &p){
    if (v == s) { f = minEdge; return; }
    else if (p[v] != -1) {
        augment(p[v], min(minEdge, res[p[v]][v]), s, f, p); res[p[v]][v] -= f;
        res[v][p[v]] += f;
    }
}

int maxFlowEdmondKarp(int n, int source, int target) {
    int mf = 0;
    while (1) {
        int f = 0;
        vector<int> dist(n+5, INF);
        dist[source] = 0;
        queue<int> q; q.push(source);
        vector<int> p; p.assign(n+5, -1);
        while (!q.empty()) {
            int u = q.front(); q.pop();
            if (u == target) break;
            for (int v = 0; v < n; v++)
                if (res[u][v] > 0 && dist[v] == INF)
                    dist[v] = dist[u] + 1, q.push(v), p[v] = u;
        }
        augment(target, INF, source, f, p);
        if (f == 0) break;
        mf += f;
    }
}

```

```

return mf;
}

/* Ford Fulkerson | complexity : O(|V|^2 F) */
int findPath(int n, int u, int t, int f){
    if (u == t) return f;
    vis[u] = 1;
    for (int v = 0; v < n; ++v){
        if (!vis[v] && res[u][v] > 0){
            int df = findPath(n, v, t, min(f, res[u][v]));
            if (df > 0){
                res[u][v] -= df;
                res[v][u] += df;
                return df;
            }
        }
    }
    return 0;
}

int maxFlowFordFulkerson(int n, int source, int target) {
    for (int flow = 0;;){
        for (int i = 0; i < n; ++i) vis[i] = 0;
        int df = findPath(n, source, target, INF);
        if (df == 0) return flow;
        flow += df;
    }
}

/* WARNING: res will be modified during the process */

```

Max Flow Min Cost

```

/** Max Flow Min Cost **/
/* complexity: O(min(E^2 V log V, E log V F)) */
const int maxnodes = 200000;

int nodes = maxnodes;
int prio[maxnodes], curflow[maxnodes], prevedge[maxnodes],
prevnode[maxnodes], q[maxnodes], pot[maxnodes];
bool inqueue[maxnodes];

struct Edge {
    int to, f, cap, cost, rev;
};

vector<Edge> graph[maxnodes];

void addEdge(int s, int t, int cap, int cost) {
    Edge a = {t, 0, cap, cost, graph[t].size()};
    Edge b = {s, 0, 0, -cost, graph[s].size()};
    graph[s].push_back(a);
    graph[t].push_back(b);
}

```

```

void bellmanFord(int s, int dist[]) {
    fill(dist, dist + nodes, 1000000000);
    dist[s] = 0;
    int qt = 0;
    q[qt++] = s;
    for (int qh = 0; (qh - qt) % nodes != 0; qh++) {
        int u = q[qh % nodes];
        inqueue[u] = false;
        for (int i = 0; i < (int) graph[u].size(); i++) {
            Edge &e = graph[u][i];
            if (e.cap <= e.f) continue;
            int v = e.to;
            int ndist = dist[u] + e.cost;
            if (dist[v] > ndist) {
                dist[v] = ndist;
                if (!inqueue[v]) {
                    inqueue[v] = true;
                    q[qt++ % nodes] = v;
                }
            }
        }
    }
}

pii minCostFlow(int s, int t, int maxf) {
    // bellmanFord can be safely commented if edges costs are non-negative
    bellmanFord(s, pot);
    int flow = 0;
    int flowCost = 0;
    while (flow < maxf) {
        priority_queue<ll, vector<ll>, greater<ll>> > q;
        q.push(s);
        fill(prio, prio + nodes, 1000000000);
        prio[s] = 0;
        curflow[s] = 1000000000;
        while (!q.empty()) {
            ll cur = q.top();
            int d = cur >> 32;
            int u = cur;
            q.pop();
            if (d != prio[u]) continue;
            for (int i = 0; i < (int) graph[u].size(); i++) {
                Edge &e = graph[u][i];
                int v = e.to;
                if (e.cap <= e.f) continue;
                int nprio = prio[u] + e.cost + pot[u] - pot[v];
                if (prio[v] > nprio) {
                    prio[v] = nprio;
                    q.push(((ll) nprio << 32) + v);
                    prevnode[v] = u;
                    prevedge[v] = i;
                    curflow[v] = min(curflow[u], e.cap - e.f);
                }
            }
        }
    }
}

```

```

    }
    if (prio[t] == 1000000000) break;
    for (int i = 0; i < nodes; i++) pot[i] += prio[i];
    int df = min(curflow[t], maxf - flow);
    flow += df;
    for (int v = t; v != s; v = prevnode[v]) {
        Edge &e = graph[prevnode[v]][prevedge[v]];
        e.f += df;
        graph[v][e.rev].f -= df;
        flowCost += df * e.cost;
    }
}

return make_pair(flow, flowCost);
}

/* usage example:
 * addEdge (source, target, capacity, cost)
 * minCostFlow(source, target, INF) -> <flow, flowCost>
 */

```

Lowest Common Ancestor

```

/** Lowest Common Ancestor */
/* complexity : LCApre : O(N log N), LCAquery : O(log N) */
/* legend:
 * N : number of vertices. WARNING: zero based
 * T : direct parent. T[v] is parent of v
 * L : L[v] is the level of v. zero/one based is okay
 * P : dp table of size [MAXN][LOGMAXN]. P[v][i] is the 2^i-th parent of v
 */

#define MAXN 100100
#define LOGMAXN 18

int L[MAXN], P[MAXN][LOGMAXN], T[MAXN], N;

void pre() {
    int i, j;

    //we initialize every element in P with -1
    for (i = 0; i < N; i++)
        for (j = 0; 1 << j < N; j++)
            P[i][j] = -1;

    //the first ancestor of every node i is T[i]
    for (i = 0; i < N; i++)
        P[i][0] = T[i];

    //bottom up dynamic programming
    for (j = 1; 1 << j < N; j++)
        for (i = 0; i < N; i++)
            if (P[i][j - 1] != -1)

```

```

        P[i][j] = P[P[i][j] - 1][j - 1];
    }
    int query(int p, int q){
        int log, i;

        //if p is situated on a higher level than q then we swap them
        if (L[p] < L[q]) swap(p,q);

        //we compute the value of [log(L[p])]
        for (log = 1; 1 << log <= L[p]; log++);
        log--;

        //we find the ancestor of node p situated on the same level
        //with q using the values in P
        for (i = log; i >= 0; i--)
            if (L[p] - (1 << i) >= L[q])
                p = P[p][i];

        if (p == q) return p;

        //we compute LCA(p, q) using the values in P
        for (i = log; i >= 0; i--)
            if (P[p][i] != -1 && P[q][i] != P[p][i])
                p = P[p][i], q = P[q][i];

        return T[p];
    }

```

Blossom

```

/** Maximum Matching on General Graph */
/* Blossom | O(V^3) */

int lca(vector<int> &match, vector<int> &base, vector<int> &p, int a, int
b) {
    vector<bool> used(SZ(match));
    while (true) {
        a = base[a];
        used[a] = true;
        if (match[a] == -1) break;
        a = p[match[a]];
    }
    while (true) {
        b = base[b];
        if (used[b]) return b;
        b = p[match[b]];
    }
    return -1;
}

void markPath(vector<int> &match, vector<int> &base, vector<bool> &blossom,
vector<int> &p, int v, int b, int children) {
    for (; base[v] != b; v = p[match[v]]) {

```

```

        blossom[base[v]] = blossom[base[match[v]]] = true;
        p[v] = children;
        children = match[v];
    }
}

int findPath(vector<vector<int> > &graph, vector<int> &match, vector<int>
&p, int root) {
    int n = SZ(graph);
    vector<bool> used(n);
    FORIT(it, p) *it = -1;
    vector<int> base(n);
    for (int i = 0; i < n; ++i) base[i] = i;

    used[root] = true;
    int qh = 0;
    int qt = 0;
    vector<int> q(n);
    q[qt++] = root;
    while (qh < qt) {
        int v = q[qh++];
        FORIT(it, graph[v]) {
            int to = *it;
            if (base[v] == base[to] || match[v] == to) continue;
            if (to == root || match[to] != -1 && p[match[to]] != -1) {
                int curbase = lca(match, base, p, v, to);
                vector<bool> blossom(n);
                markPath(match, base, blossom, p, v, curbase, to);
                markPath(match, base, blossom, p, to, curbase, v);
                for (int i = 0; i < n; ++i) {
                    if (blossom[base[i]]) {
                        base[i] = curbase;
                        if (!used[i]) {
                            used[i] = true;
                            q[qt++] = i;
                        }
                    }
                }
            }
            else if (p[to] == -1) {
                p[to] = v;
                if (match[to] == -1) return to;
                to = match[to];
                used[to] = true;
                q[qt++] = to;
            }
        }
    }
    return -1;
}

int maxMatching(vector<vector<int> > graph) {
    int n = SZ(graph);
    vector<int> match(n, -1);

```

```

vector<int> p(n);
for (int i = 0; i < n; ++i) {
    if (match[i] == -1) {
        int v = findPath(graph, match, p, i);
        while (v != -1) {
            int pv = p[v];
            int ppv = match[pv];
            match[v] = pv;
            match[pv] = v;
            v = ppv;
        }
    }
}

int matches = 0;
for (int i = 0; i < n; ++i) {
    if (match[i] != -1) {
        ++matches;
    }
}
return matches / 2;
}

```

Minimum Cut

```

// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
//
// Running time:
//   O(|V|^3)
//
// INPUT:
//   - graph, constructed using AddEdge()
//
// OUTPUT:
//   - (min cut value, nodes in half of min cut)

#include <cmath>
#include <vector>
#include <iostream>

using namespace std;

typedef vector<int> VI;
typedef vector<VI> VVI;

const int INF = 1000000000;

pair<int, VI> GetMinCut(VVI &weights) {
    int N = weights.size();
    VI used(N), cut, best_cut;
    int best_weight = -1;

    for (int phase = N-1; phase >= 0; phase--) {

```

```

VI w = weights[0];
VI added = used;
int prev, last = 0;
for (int i = 0; i < phase; i++) {
    prev = last;
    last = -1;
    for (int j = 1; j < N; j++)
        if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
    if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j];
        for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];
        used[last] = true;
        cut.push_back(last);
        if (best_weight == -1 || w[last] < best_weight) {
            best_cut = cut;
            best_weight = w[last];
        }
    } else {
        for (int j = 0; j < N; j++)
            w[j] += weights[last][j];
        added[last] = true;
    }
}
return make_pair(best_weight, best_cut);
}

```

String Processing

Knuth-Morris-Pratt

```

/** Knuth-Morris-Pratt */
/* Complexity: O(N) */
void buildFailTable(char *pattern, int *t){
    int i = 0, j = -1, m = strlen(pattern);
    t[0] = -1;
    while (i < m){
        while (j >= 0 && pattern[i] != pattern[j]) j = t[j];
        i++; j++;
        t[i] = j;
    }
}

vector<int> kmpSearch(char *pattern, char *text){
    vector<int> res;
    int i = 0, j = 0, n = strlen(text), m = strlen(pattern);
    int t[m+5];
    buildFailTable(pattern, t);
    while (i < n){
        while (j >= 0 && text[i] != pattern[j]) j = t[j];
        i++; j++;
        if (j == m){

```

```

        res.push_back(i-j);
        j = t[j];
    }
    return res;
}

```

Z-Algorithm

```

// z[i] is the longest substring starting from i which is also a prefix of
// s
// z[0] is not set
vector<int> z_function(string s) {
    int n = (int) s.length();
    vector<int> z(n);
    for (int i = 1, l = 0, r = 0; i < n; ++i) {
        if (i <= r)
            z[i] = min(r - i + 1, z[i - l]);
        while (i + z[i] < n && s[z[i]] == s[i + z[i]])
            ++z[i];
        if (i + z[i] - 1 > r)
            l = i, r = i + z[i] - 1;
    }
    return z;
}

```

Suffix Array

```

/** Suffix Array */
/* complexity: O(N log N) */

#define MAXN 200000

char T[MAXN+5]; // input
int n; // length
int RA[MAXN+5], tempRA[MAXN+5]; // rank array
int SA[MAXN+5], tempSA[MAXN+5]; // suffix array
int c[MAXN+5]; //for counting/radix sort

void countingSort(int k) {
    int sum, maxi = max(300, n);
    memset(c, 0, sizeof(c));
    for (int i = 0; i < n; i++)
        c[i+k < n ? RA[i+k] : 0]++;
    for (int i = sum = 0; i < maxi; i++) {
        int t = c[i]; c[i] = sum;
        sum += t;
    }
    for (int i = 0; i < n; i++)
        tempSA[c[SA[i]+k < n ? RA[SA[i]+k] : 0]++] = SA[i];
    for (int i = 0; i < n; i++) SA[i] = tempSA[i];
}

void SuffixArray_Construct() {

```

```

int r;
for (int i = 0; i < n; i++) RA[i] = T[i] - '.';
for (int i = 0; i < n; i++) SA[i] = i;
for (int k = 1; k < n; k <= 1) {
    countingSort(k);
    countingSort(0);
    tempRA[SA[0]] = r = 0;
    for (int i = 1; i < n; i++)
        tempRA[SA[i]] =
            (RA[SA[i]] == RA[SA[i-1]] && RA[SA[i]+k] == RA[SA[i-1]+k])
? r : ++r;
    for (int i = 0; i < n; i++) RA[i] = tempRA[i];
}

```

Suffix Tree

```

/** SUFFIX TREE UKKONEN */
class SuffixTree {
    static String alphabet = "abcdefghijklmnopqrstuvwxyz1234567890\1\2";
    static int alphabetSize = alphabet.length();

    static class Node {
        int depth; // from start of suffix
        int begin;
        int end;
        Node[] children;
        Node parent;
        Node suffixLink;

        Node(int begin, int end, int depth, Node parent) {
            children = new Node[alphabetSize];
            this.begin = begin;
            this.end = end;
            this.parent = parent;
            this.depth = depth;
        }

        boolean contains(int d) {
            return depth <= d && d < depth + (end - begin);
        }
    }

    public static Node buildSuffixTree(String s) {
        int n = s.length();
        byte[] a = new byte[n];
        for (int i = 0; i < n; i++) {
            a[i] = (byte) alphabet.indexOf(s.charAt(i));
        }
        Node root = new Node(0, 0, 0, null);
        Node cn = root;
        // root.suffixLink must be null, but that way it gets more convenient
        // processing
    }
}

```



```

root.suffixLink = root;
Node needsSuffixLink = null;
int lastRule = 0;
int j = 0;
for (int i = -1; i < n - 1; i++) { // strings s[j..i] already in tree,
    // add s[i+1] to it.
    int cur = a[i + 1]; // last char of current string
    for (; j <= i + 1; j++) {
        int curDepth = i + 1 - j;
        if (lastRule != 3) {
            cn = cn.suffixLink != null ? cn.suffixLink :
cn.parent.suffixLink;
            int k = j + cn.depth;
            while (curDepth > 0 && !cn.contains(curDepth - 1)) {
                k += cn.end - cn.begin;
                cn = cn.children[a[k]];
            }
            if (!cn.contains(curDepth)) { // explicit node
                if (needsSuffixLink != null) {
                    needsSuffixLink.suffixLink = cn;
                    needsSuffixLink = null;
                }
                if (cn.children[cur] == null) {
                    // no extension - add leaf
                    cn.children[cur] = new Node(i + 1, n, curDepth, cn);
                    lastRule = 2;
                } else {
                    cn = cn.children[cur];
                    lastRule = 3; // already exists
                    break;
                }
            } else { // implicit node
                int end = cn.begin + curDepth - cn.depth;
                if (a[end] != cur) { // split implicit node here
                    Node newn = new Node(cn.begin, end, cn.depth, cn.parent);
                    newn.children[cur] = new Node(i + 1, n, curDepth, newn);
                    newn.children[a[end]] = cn;
                    cn.parent.children[a[cn.begin]] = newn;
                    if (needsSuffixLink != null) {
                        needsSuffixLink.suffixLink = newn;
                    }
                    cn.begin = end;
                    cn.depth = curDepth;
                    cn.parent = newn;
                    cn = needsSuffixLink = newn;
                    lastRule = 2;
                } else if (cn.end != n || cn.begin - cn.depth < j) {
                    lastRule = 3;
                    break;
                } else {
                    lastRule = 1;
                }
            }
        }
    }
}

```

```

    }
}
root.suffixLink = null;
return root;
}

// usage example
static int lcsLength;
static int lcsBeginIndex;

// traverse suffix tree to find longest common substring
public static int findLCS(Node node, int i1, int i2) {
    if (node.begin <= i1 && i1 < node.end) {
        return 1;
    }
    if (node.begin <= i2 && i2 < node.end) {
        return 2;
    }
    int mask = 0;
    for (char f = 0; f < alphabetSize; f++) {
        if (node.children[f] != null) {
            mask |= findLCS(node.children[f], i1, i2);
        }
    }
    if (mask == 3) {
        int curLength = node.depth + node.end - node.begin;
        if (lcsLength < curLength) {
            lcsLength = curLength;
            lcsBeginIndex = node.begin;
        }
    }
    return mask;
}

// Usage example
public static void main(String[] args) {
    String s1 = "12345";
    String s2 = "124234";
    // build generalized suffix tree (see Gusfield, p.125)
    String s = s1 + '\1' + s2 + '\2';
    Node root = buildSuffixTree(s);
    lcsLength = 0;
    lcsBeginIndex = 0;
    // find longest common substring
    findLCS(root, s1.length(), s1.length() + s2.length() + 1);
    System.out.println(3 == lcsLength);
    System.out.println(s.substring(lcsBeginIndex - 1, lcsBeginIndex +
lcsLength - 1));
}
}

```

Aho-Corasick

```

/** DICTONARY MATCHING AHO-CORASICK */
public class AhoCorasick {

    static final int ALPHABET_SIZE = 26;

    static class Node {

        Node[] children = new Node[ALPHABET_SIZE];
        boolean leaf;
        Node parent;
        char charToParent;
        Node suffLink;
        Node[] go = new Node[ALPHABET_SIZE];
    }

    public static Node createRoot() {
        Node node = new Node();
        node.suffLink = node;
        return node;
    }

    public static void addString(Node node, String s) {
        for (char ch : s.toCharArray()) {
            int c = ch - 'a';
            if (node.children[c] == null) {
                Node n = new Node();
                n.parent = node;
                n.charToParent = ch;
                node.children[c] = n;
            }
            node = node.children[c];
        }
        node.leaf = true;
    }

    public static Node go(Node node, char ch) {
        int c = ch - 'a';
        if (node.go[c] == null) {
            if (node.children[c] != null) {
                node.go[c] = node.children[c];
            } else {
                node.go[c] = node.parent == null ? node : go(suffLink(node), ch);
            }
        }
        return node.go[c];
    }

    public static Node suffLink(Node node) {
        if (node.suffLink == null) {
            if (node.parent.parent == null) {
                node.suffLink = node.parent;
            }
        }
    }
}

```

```

    } else {
        node.suffLink = go(suffLink(node.parent), node.charToParent);
    }
}
return node.suffLink;
}

// Usage example
public static void main(String[] args) {
    Node tree = createRoot();
    addString(tree, "bc");
    addString(tree, "abc");

    String s = "tabc";
    Node node = tree;
    for (char ch : s.toCharArray()) {
        node = go(node, ch);
    }
    System.out.println(node.leaf);
}
}

```

Mathematics

Extended Euclid

```

/** Extended Euclid | returns <x,y> where ax + by = gcd(a,b) */
/* complexity: O(min(log(a),log(b))) */
pair<ll,ll> extendedEuclid(ll a, ll b){
    ll x = 0, y = 1, lastx = 1, lasty = 0;
    while (b != 0){
        ll quotient = a / b;
        /* (a, b) = (b, a mod b) */
        ll temp = a;
        a = b;
        b = temp % b;
        /* (x, lastx) = (lastx - quotient*x, x) */
        temp = x;
        x = lastx - quotient * x;
        lastx = temp;
        /* (y, lasty) = (lasty - quotient*y, y) */
        temp = y;
        y = lasty - quotient * y;
        lasty = temp;
    }
    return make_pair(lastx, lasty);
}

```

Diophantine

```

// computes x and y such that ax + by = c; on failure, x = y = -1
void linear_diophantine(int a, int b, int c, int &x, int &y) {
    int d = gcd(a,b);
}

```

```

if (c%d) {
    x = y = -1;
} else {
    x = c/d * mod_inverse(a/d, b/d);
    y = (c-a*x)/b;
}
}

```

Chinese Remainder Theorem

```

// Chinese remainder theorem (special case): find z such that
// z % x = a, z % y = b. Here, z is unique modulo M = lcm(x,y).
// Return (z,M). On failure, M = -1.
PII chinese_remainder_theorem(int x, int a, int y, int b) {
    int s, t;
    int d = extended_euclid(x, y, s, t);
    if (a%d != b%d) return make_pair(0, -1);
    return make_pair(mod(s*b*x+t*a*y,x*y)/d, x*y/d);
}

// Chinese remainder theorem: find z such that
// z % x[i] = a[i] for all i. Note that the solution is
// unique modulo M = lcm_i (x[i]). Return (z,M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &x, const VI &a) {
    PII ret = make_pair(a[0], x[0]);
    for (int i = 1; i < x.size(); i++) {
        ret = chinese_remainder_theorem(ret.second, ret.first, x[i], a[i]);
        if (ret.second == -1) break;
    }
    return ret;
}

```

Lagrange Interpolation

```

/** Lagrange Polynomial Interpolation */
/* complexity: O(n^2) */
class lagrangeInterpolation {
public:
    lagrangeInterpolation () : x_(0), y_(0) {}

    void addCoef (db x, db y){
        x_.push_back(x);
        y_.push_back(y);
    }

    db interpolate (db x){
        db value = 0;
        for (int i = 0; i < (int)x_.size(); ++i){
            db addum = y_[i];
            for (int j = 0; j < (int)x_.size(); ++j) if (i != j){
                addum *= (x - x_[j]);
            }
        }
    }
}

```

```

        addum /= (x_[i] - x_[j]);
    }
    value += addum;
}
return value;
}

vector<db> x_, y_;
};

class modularInterpolation {
public:
    modularInterpolation (const ll &modu) : modu_(modu), x_(0), y_(0) {}

    void addCoef (ll x, ll y){
        x %= modu_;
        if (x < 0LL) x += modu_;
        x_.push_back(x);

        y %= modu_;
        if (y < 0LL) y += modu_;
        y_.push_back(y);
    }

    ll interpolate (ll x){
        x %= modu_;
        if (x < 0LL) x += modu_;

        for (int i = 0; i < (int)x_.size(); ++i) if (x_[i] == x) return
y_[i];

        ll value = 0LL;
        for (int i = 0; i < (int)x_.size(); ++i){
            ll addum = y_[i];
            for (int j = 0; j < (int)x_.size(); ++j) if (j != i){
                ll delta1 = (x - x_[j] + modu_) % modu_;
                ll delta2 = (x_[i] - x_[j] + modu_) % modu_;
                addum = (addum * delta1) % modu_;
                addum = (addum * multInverse(delta2, modu_)) % modu_;
            }
            value += addum;
            value %= modu_;
        }

        return value;
    }

    const ll modu_;
    vector<ll> x_, y_;
};

/* WARNING: no two x_[i] should be the same */

```

Fast Fourier Transform

```

/** Fast Fourier Transform */
/* complexity: O(N log N) */
vector< complex<db> > iterativeDFT (const vector< complex<db> > &seq, int
direction) {
    int n = SZ(seq);
    int bits = 0;
    int tmp_n = n;
    complex<db> *placeholder = new complex<db>[n];
    complex<db> *tmp = new complex<db>[n];

    while (tmp_n > 1){
        ++bits;
        tmp_n /= 2;
    }

    REP(i,n){
        int res = 0;
        int tmp_i = i;
        REP(j,bits){
            if (tmp_i % 2) res += (1 << (bits-j-1));
            tmp_i /= 2;
        }
        placeholder[i] = seq[res];
    }

    for (int comp_size = 2; comp_size <= n; comp_size *= 2){
        for (int j = 0; j < n; j += comp_size){
            int n_mem = comp_size / 2;
            db w_mult_exp_i = 2. * acos(-1.) / (db)comp_size;
            if (!direction) w_mult_exp_i *= -1.;
            complex<db> w_mult (cos(w_mult_exp_i),sin(w_mult_exp_i));
            complex<db> w (1., 0.);
            for (int k = 0; k < comp_size; ++k){
                int idx = k % n_mem;
                tmp[k] = placeholder[j+idx] + w * placeholder[j+n_mem+idx];
                w = w * w_mult;
            }
            for (int k = 0; k < comp_size; ++k){
                placeholder[j+k] = tmp[k];
            }
        }
    }

    vector< complex<db> > result;
    for (int i = 0; i < n; ++i) result.PB(placeholder[i]);

    delete[] placeholder;
    delete[] tmp;
    return result;
}

```

```

vector<db> FFT(vector<db> a, vector<db> b) {
    if (SZ(a) == 0) a.PB(0.);
    if (SZ(b) == 0) b.PB(0.);
    int n_final_elements = SZ(a) + SZ(b) - 1;

    int actual_size = 1;
    while (actual_size < max(SZ(a), SZ(b))){
        actual_size *= 2;
    }
    actual_size *= 2;
    while (SZ(a) < actual_size) a.PB(0.);
    while (SZ(b) < actual_size) b.PB(0.);

    vector< complex<db> > dft_input_a, dft_input_b;
    REP(i,actual_size) {
        dft_input_a.PB(complex<db> (a[i], 0.));
        dft_input_b.PB(complex<db> (b[i], 0.));
    }

    dft_input_a = iterativeDFT (dft_input_a, 1);
    dft_input_b = iterativeDFT (dft_input_b, 1);
    REP(i,actual_size) {
        dft_input_a[i] = dft_input_a[i] * dft_input_b[i];
    }
    dft_input_a = iterativeDFT (dft_input_a, 0);

    vector<db> res;
    REP(i,n_final_elements) {
        res.PB(dft_input_a[i].real() / (db) actual_size);
    }
    return res;
}

```

Simplex

```

// Two-phase simplex algorithm for solving linear programs of the form
//
//      maximize      c^T x
//      subject to    Ax <= b
//                  x >= 0
//
// INPUT: A -- an m x n matrix
//        b -- an m-dimensional vector
//        c -- an n-dimensional vector
//        x -- a vector where the optimal solution will be stored
//
// OUTPUT: value of the optimal solution (infinity if unbounded
//        above, nan if infeasible)
//
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).

#include <iostream>

```

```

#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>

using namespace std;

typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;

const DOUBLE EPS = 1e-9;

struct LPSolver {
    int m, n;
    VI B, N;
    VVD D;

    LPSolver(const VVD &A, const VD &b, const VD &c) :
        m(b.size()), n(c.size()), N(n+1), B(m), D(m+2, VD(n+2)) {
        for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] =
A[i][j];
        for (int i = 0; i < m; i++) { B[i] = n+i; D[i][n] = -1; D[i][n+1] =
b[i]; }
        for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
        N[n] = -1; D[m+1][n] = 1;
    }

    void Pivot(int r, int s) {
        for (int i = 0; i < m+2; i++) if (i != r)
            for (int j = 0; j < n+2; j++) if (j != s)
                D[i][j] -= D[r][j] * D[i][s] / D[r][s];
        for (int j = 0; j < n+2; j++) if (j != s) D[r][j] /= D[r][s];
        for (int i = 0; i < m+2; i++) if (i != r) D[i][s] /= -D[r][s];
        D[r][s] = 1.0 / D[r][s];
        swap(B[r], N[s]);
    }

    bool Simplex(int phase) {
        int x = phase == 1 ? m+1 : m;
        while (true) {
            int s = -1;
            for (int j = 0; j <= n; j++) {
                if (phase == 2 && N[j] == -1) continue;
                if (s == -1 || D[x][j] < D[x][s] || D[x][j] == D[x][s] && N[j] < N[s]) s
= j;
            }
            if (D[x][s] >= -EPS) return true;
            int r = -1;
            for (int i = 0; i < m; i++) {
                if (D[i][s] <= 0) continue;
                if (r == -1 || D[i][n+1] / D[i][s] < D[r][n+1] / D[r][s] ||

```

```

                D[i][n+1] / D[i][s] == D[r][n+1] / D[r][s] && B[i] < B[r]) r = i;
            }
            if (r == -1) return false;
            Pivot(r, s);
        }

        DOUBLE Solve(VD &x) {
            int r = 0;
            for (int i = 1; i < m; i++) if (D[i][n+1] < D[r][n+1]) r = i;
            if (D[r][n+1] <= -EPS) {
                Pivot(r, n);
                if (!Simplex(1) || D[m+1][n+1] < -EPS) return -
numeric_limits<DOUBLE>::infinity();
                for (int i = 0; i < m; i++) if (B[i] == -1) {
                    int s = -1;
                    for (int j = 0; j <= n; j++)
                        if (s == -1 || D[i][j] < D[i][s] || D[i][j] == D[i][s] && N[j] < N[s])
s = j;
                    Pivot(i, s);
                }
                if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
                x = VD(n);
                for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n+1];
                return D[m][n+1];
            }
        };

        int main() {
            const int m = 4;
            const int n = 3;
            DOUBLE _A[m][n] = {
                { 6, -1, 0 },
                { -1, -5, 0 },
                { 1, 5, 1 },
                { -1, -5, -1 }
            };
            DOUBLE _b[m] = { 10, -4, 5, -5 };
            DOUBLE _c[n] = { 1, -1, 0 };

            VVD A(m);
            VD b(_b, _b + m);
            VD c(_c, _c + n);
            for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);

            LPSolver solver(A, b, c);
            VD x;
            DOUBLE value = solver.Solve(x);

            cerr << "VALUE: " << value << endl;
            cerr << "SOLUTION:";
            for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];

```

```

cerr << endl;
return 0;
}

```

Gauss Jordan Elimination

```

// Gauss-Jordan elimination with full pivoting.
//
// Uses:
// (1) solving systems of linear equations (AX=B)
// (2) inverting matrices (AX=I)
// (3) computing determinants of square matrices
//
// Running time: O(n^3)
//
// INPUT:   a[][] = an nxn matrix
//          b[][] = an nxm matrix
//
// OUTPUT:  X      = an nxm matrix (stored in b[][])
//          A^{-1} = an nxn matrix (stored in a[][])
//          returns determinant of a[][]

#include <iostream>
#include <vector>
#include <cmath>

using namespace std;

const double EPS = 1e-10;

typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;

T GaussJordan(VVT &a, VVT &b) {
    const int n = a.size();
    const int m = b[0].size();
    VI irow(n), icol(n), ipiv(n);
    T det = 1;

    for (int i = 0; i < n; i++) {
        int pj = -1, pk = -1;
        for (int j = 0; j < n; j++) if (!ipiv[j])
            for (int k = 0; k < n; k++) if (!ipiv[k])
                if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
        if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl;
            exit(0); }
        ipiv[pk]++;
        swap(a[pj], a[pk]);
        swap(b[pj], b[pk]);
        if (pj != pk) det *= -1;
        irow[i] = pj;
    }
}

```

```

        icol[i] = pk;

        T c = 1.0 / a[pk][pk];
        det *= a[pk][pk];
        a[pk][pk] = 1.0;
        for (int p = 0; p < n; p++) a[pk][p] *= c;
        for (int p = 0; p < m; p++) b[pk][p] *= c;
        for (int p = 0; p < n; p++) if (p != pk) {
            c = a[p][pk];
            a[p][pk] = 0;
            for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
            for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
        }
    }

    for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
        for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);
    }

    return det;
}

int main() {
    const int n = 4;
    const int m = 2;
    double A[n][n] = { {1,2,3,4},{1,0,1,0},{5,3,2,4},{6,1,4,6} };
    double B[n][m] = { {1,2},{4,3},{5,6},{8,7} };
    VVT a(n), b(n);
    for (int i = 0; i < n; i++) {
        a[i] = VT(A[i], A[i] + n);
        b[i] = VT(B[i], B[i] + m);
    }

    double det = GaussJordan(a, b);

    // expected: 60
    cout << "Determinant: " << det << endl;

    // expected: -0.233333 0.166667 0.133333 0.0666667
    //              0.166667 0.166667 0.333333 -0.333333
    //              0.233333 0.833333 -0.133333 -0.0666667
    //              0.05 -0.75 -0.1 0.2
    cout << "Inverse: " << endl;
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < n; j++)
            cout << a[i][j] << ' ';
        cout << endl;
    }

    // expected: 1.63333 1.3
    //              -0.166667 0.5
    //              2.36667 1.7
    //              -1.85 -1.35
}

```

```

cout << "Solution: " << endl;
for (int i = 0; i < n; i++) {
    for (int j = 0; j < m; j++)
        cout << b[i][j] << ' ';
    cout << endl;
}
}

```

Reduced Row Echelon Form

```

// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
//
// Running time: O(n^3)
//
// INPUT:    a[][] = an nxm matrix
//
// OUTPUT:   rref[][] = an nxm matrix (stored in a[][])
//            returns rank of a[][]

#include <iostream>
#include <vector>
#include <cmath>

using namespace std;

const double EPSILON = 1e-10;

typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;

int rref(VVT &a) {
    int n = a.size();
    int m = a[0].size();
    int r = 0;
    for (int c = 0; c < m && r < n; c++) {
        int j = r;
        for (int i = r+1; i < n; i++)
            if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
        if (fabs(a[j][c]) < EPSILON) continue;
        swap(a[j], a[r]);

        T s = 1.0 / a[r][c];
        for (int j = 0; j < m; j++) a[r][j] *= s;
        for (int i = 0; i < n; i++) if (i != r) {
            T t = a[i][c];
            for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];
        }
        r++;
    }
    return r;
}

```

```

}

int main(){
    const int n = 5;
    const int m = 4;
    double A[n][m] = {
        {16,2,3,13},{5,11,10,8},{9,7,6,12},{4,14,15,1},{13,21,21,13} };
    VVT a(n);
    for (int i = 0; i < n; i++)
        a[i] = VT(A[i], A[i] + n);
    int rank = rref(a);
    // expected: 4
    cout << "Rank: " << rank << endl;
    // expected: 1 0 0 1
    //           0 1 0 3
    //           0 0 1 -3
    //           0 0 0 2.78206e-15
    //           0 0 0 3.22398e-15
    cout << "rref: " << endl;
    for (int i = 0; i < 5; i++){
        for (int j = 0; j < 4; j++)
            cout << a[i][j] << ' ';
        cout << endl;
    }
}

```

Data Structures

K-d Tree

```

// -----
//
// A straightforward, but probably sub-optimal KD-tree implmentation that's
// probably good enough for most things (current it's a 2D-tree)
//
// - constructs from n points in O(n lg^2 n) time
// - handles nearest-neighbor query in O(lg n) if points are well
//   distributed
// - worst case for nearest-neighbor may be linear in pathological case
//
// Sonny Chan, Stanford University, April 2009
// -----
//
#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>

using namespace std;

// number type for coordinates, and its maximum value

```

```

typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();

// point structure for 2D-tree, can be extended to 3D
struct point {
    ntype x, y;
    point(ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
};

bool operator==(const point &a, const point &b){return a.x == b.x && a.y == b.y;}

// sorts points on x-coordinate
bool on_x(const point &a, const point &b){return a.x < b.x;}

// sorts points on y-coordinate
bool on_y(const point &a, const point &b){return a.y < b.y;}

// squared distance between points
ntype pdist2(const point &a, const point &b)
{
    ntype dx = a.x-b.x, dy = a.y-b.y;
    return dx*dx + dy*dy;
}

// bounding box for a set of points
struct bbox
{
    ntype x0, x1, y0, y1;
    bbox() : x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}
    // computes bounding box from a bunch of points
    void compute(const vector<point> &v) {
        for (int i = 0; i < v.size(); ++i) {
            x0 = min(x0, v[i].x);    x1 = max(x1, v[i].x);
            y0 = min(y0, v[i].y);    y1 = max(y1, v[i].y);
        }
    }
    // squared distance between a point and this bbox, 0 if inside
    ntype distance(const point &p) {
        if (p.x < x0) {
            if (p.y < y0) return pdist2(point(x0, y0), p);
            else if (p.y > y1) return pdist2(point(x0, y1), p);
            else return pdist2(point(x0, p.y), p);
        }
        else if (p.x > x1) {
            if (p.y < y0) return pdist2(point(x1, y0), p);
            else if (p.y > y1) return pdist2(point(x1, y1), p);
            else return pdist2(point(x1, p.y), p);
        }
        else {
            if (p.y < y0) return pdist2(point(p.x, y0), p);
            else if (p.y > y1) return pdist2(point(p.x, y1), p);
            else return 0;
        }
    }
};

```

```

    }
};

// stores a single node of the kd-tree, either internal or leaf
struct kdnnode
{
    bool leaf; // true if this is a leaf node (has one point)
    point pt; // the single point of this is a leaf
    bbox bound; // bounding box for set of points in children

    kdnnode *first, *second; // two children of this kd-node

    kdnnode() : leaf(false), first(0), second(0) {}
    ~kdnnode() { if (first) delete first; if (second) delete second; }

    // intersect a point with this node (returns squared distance)
    ntype intersect(const point &p) {
        return bound.distance(p);
    }

    // recursively builds a kd-tree from a given cloud of points
    void construct(vector<point> &vp)
    {
        // compute bounding box for points at this node
        bound.compute(vp);
        // if we're down to one point, then we're a leaf node
        if (vp.size() == 1) {
            leaf = true;
            pt = vp[0];
        }
        else {
            // heuristic...
            // split on x if the bbox is wider than high (not best)
            if (bound.x1-bound.x0 >= bound.y1-bound.y0)
                sort(vp.begin(), vp.end(), on_x);
            // otherwise split on y-coordinate
            else
                sort(vp.begin(), vp.end(), on_y);

            // divide by taking half the array for each child
            // (not best performance if many duplicates in the middle)
            int half = vp.size()/2;
            vector<point> vl(vp.begin(), vp.begin()+half);
            vector<point> vr(vp.begin()+half, vp.end());
            first = new kdnnode(); first->construct(vl);
            second = new kdnnode(); second->construct(vr);
        }
    }
};

// simple kd-tree class to hold the tree and handle queries
struct kdtree

```



```

{
    kdnode *root;

    // constructs a kd-tree from a points (copied here, as it sorts them)
    kdtree(const vector<point> &vp) {
        vector<point> v(vp.begin(), vp.end());
        root = new kdnode();
        root->construct(v);
    }
    ~kdtree() { delete root; }

    // recursive search method returns squared distance to nearest point
    ntype search(kdnode *node, const point &p)
    {
        if (node->leaf) {
            // commented special case tells a point not to find itself
            if (p == node->pt) return sentry;
            // else
            return pdist2(p, node->pt);
        }
        ntype bfirst = node->first->intersect(p);
        ntype bsecond = node->second->intersect(p);
        // choose the side with the closest bounding box to search first
        // (note that the other side is also searched if needed)
        if (bfirst < bsecond) {
            ntype best = search(node->first, p);
            if (bsecond < best)
                best = min(best, search(node->second, p));
            return best;
        }
        else {
            ntype best = search(node->second, p);
            if (bfirst < best)
                best = min(best, search(node->first, p));
            return best;
        }
    }
    // squared distance to the nearest
    ntype nearest(const point &p) {
        return search(root, p);
    }
};

int main()
{
    // generate some random points for a kd-tree
    vector<point> vp;
    for (int i = 0; i < 100000; ++i) {
        vp.push_back(point(rand()%100000, rand()%100000));
    }
    kdtree tree(vp);

    // query some points

```

```

        for (int i = 0; i < 10; ++i) {
            point q(rand()%100000, rand()%100000);
            cout << "Closest squared distance to (" << q.x << ", " << q.y <<
            ") "
                << " is " << tree.nearest(q) << endl;
        }

        return 0;
    }
}

```

Fenwick Tree

```

/** Fenwick Tree with Range Update */
#define MAXN 100005

int n, bitMul[MAXN], bitAdd[MAXN];

void internalUpdate(int k, int mul, int add) {
    for (int x = k; x <= n; x += (x & -x)) {
        bitMul[x] += mul;
        bitAdd[x] += add;
    }
}

void update(int l, int r, int value) {
    internalUpdate(l, value, -value * (l - 1));
    internalUpdate(r, -value, value * r);
}

int query(int k) {
    int mul = 0, add = 0;
    for (int x = k; x > 0; x -= (x & -x)) {
        mul += bitMul[x];
        add += bitAdd[x];
    }
    return mul * k + add;
}

```

Splay Tree

```

#include <cstdio>
#include <algorithm>
using namespace std;

const int N_MAX = 130010;
const int oo = 0x3f3f3f3f;
struct Node
{
    Node *ch[2], *pre;
    int val, size;
    bool isTurned;
} nodePool[N_MAX], *null, *root;

Node *allocNode(int val)

```

```

{
    static int freePos = 0;
    Node *x = &nodePool[freePos++];
    x->val = val, x->isTurned = false;
    x->ch[0] = x->ch[1] = x->pre = null;
    x->size = 1;
    return x;
}

inline void update(Node *x)
{
    x->size = x->ch[0]->size + x->ch[1]->size + 1;
}

inline void makeTurned(Node *x)
{
    if(x == null)
        return;
    swap(x->ch[0], x->ch[1]);
    x->isTurned ^= 1;
}

inline void pushDown(Node *x)
{
    if(x->isTurned)
    {
        makeTurned(x->ch[0]);
        makeTurned(x->ch[1]);
        x->isTurned ^= 1;
    }
}

inline void rotate(Node *x, int c)
{
    Node *y = x->pre;
    x->pre = y->pre;
    if(y->pre != null)
        y->pre->ch[y == y->pre->ch[1]] = x;
    y->ch[!c] = x->ch[c];
    if(x->ch[c] != null)
        x->ch[c]->pre = y;
    x->ch[c] = y, y->pre = x;
    update(y);
    if(y == root)
        root = x;
}

void splay(Node *x, Node *p)
{
    while(x->pre != p)
    {
        if(x->pre->pre == p)
            rotate(x, x == x->pre->ch[0]);
    }
}

```

```

else
{
    Node *y = x->pre, *z = y->pre;
    if(y == z->ch[0])
    {
        if(x == y->ch[0])
            rotate(y, 1), rotate(x, 1);
        else
            rotate(x, 0), rotate(x, 1);
    }
    else
    {
        if(x == y->ch[1])
            rotate(y, 0), rotate(x, 0);
        else
            rotate(x, 1), rotate(x, 0);
    }
}
update(x);
}

void select(int k, Node *fa)
{
    Node *now = root;
    while(1)
    {
        pushDown(now);
        int tmp = now->ch[0]->size + 1;
        if(tmp == k)
            break;
        else if(tmp < k)
            now = now->ch[1], k -= tmp;
        else
            now = now->ch[0];
    }
    splay(now, fa);
}

Node *makeTree(Node *p, int l, int r)
{
    if(l > r)
        return null;
    int mid = (l + r) / 2;
    Node *x = allocNode(mid);
    x->pre = p;
    x->ch[0] = makeTree(x, l, mid - 1);
    x->ch[1] = makeTree(x, mid + 1, r);
    update(x);
    return x;
}

int main()

```

```

{
    int n, m;
    null = allocNode(0);
    null->size = 0;
    root = allocNode(0);
    root->ch[1] = allocNode(0);
    root->ch[1]->pre = root;
    update(root);

    scanf("%d%d", &n, &m);
    root->ch[1]->ch[0] = makeTree(root->ch[1], 1, n);
    splay(root->ch[1]->ch[0], null);

    while (m --)
    {
        int a, b;
        scanf("%d%d", &a, &b);
        a ++, b ++;
        select(a - 1, null);
        select(b + 1, root);
        makeTurned(root->ch[1]->ch[0]);
    }

    for(int i = 1; i <= n; i ++)
    {
        select(i + 1, null);
        printf("%d ", root->val);
    }
}

```

DP Convex Hull Optimization

```

public class ConvexHullOptimization {

    long[] A = new long[1000000];
    long[] B = new long[1000000];
    int len;
    int ptr;

    // a descends
    public void addLine(long a, long b) {
        // intersection of (A[len-2],B[len-2]) with (A[len-1],B[len-1]) must
        lie to the left of intersection of (A[len-1],B[len-1]) with (a,b)
        while (len >= 2 && (B[len - 2] - B[len - 1]) * (a - A[len - 1]) >=
        (B[len - 1] - b) * (A[len - 1] - A[len - 2])) {
            --len;
        }
        A[len] = a;
        B[len] = b;
        ++len;
    }

    // x ascends

```

```

public long minValue(long x) {
    ptr = Math.min(ptr, len - 1);
    while (ptr + 1 < len && A[ptr + 1] * x + B[ptr + 1] <= A[ptr] * x +
    B[ptr]) {
        ++ptr;
    }
    return A[ptr] * x + B[ptr];
}

// Usage example
public static void main(String[] args) {
    ConvexHullOptimization h = new ConvexHullOptimization();
    h.addLine(3, 0);
    h.addLine(2, 1);
    h.addLine(3, 2);
    h.addLine(0, 6);
    System.out.println(h.minValue(0));
    System.out.println(h.minValue(1));
    System.out.println(h.minValue(2));
    System.out.println(h.minValue(3));
}
}

```

Geometry

Point, Segment, Line, Circle

```

double _acos(double x) {
    double ret = acos(x);
    if (ret == ret) return ret;
    if (x < 0) return acos(-1.0);
    return acos(1.0);
}

#define acos _acos
#define sqr(x) ((x)*(x))

const double PI = acos(-1);
const double EPS = 1e-9;
const double INF = 1e300;

struct point{
    double x, y;
    point() { x = y = 0; }
    point(double x, double y) : x(x), y(y) {}
};

struct segment {
    point p1, p2;
    segment() {p1 = p2 = point(0,0);}
    segment(point p1, point p2) : p1(p1), p2(p2) {}
};

/** basic operators and functions of point and segment */

```

```

/* complexity: constant */
double cross(const point &p1, const point &p2) {
    /* returns z-component of cross product of two points (vectors) */
    return p1.x * p2.y - p1.y * p2.x;
}
double dot(const point &p1, const point &p2) {
    /* returns dot product of two points (vectors) */
    return p1.x * p2.x + p1.y * p2.y;
}
double getAngle(const point &p1, const point &p2) {
    /* returns angle formed by two vectors. WARNING: undirected angle */
    return fabs(acos(dot(p1,p2) / dist(p1,point(0,0)) /
dist(p2,point(0,0))));
}
double getAngle(const point &p1, const point &center, const point &p2) {
    /* returns angle formed by three points. WARNING: undirected angle */
    return getAngle(p1 - center, p2 - center);
}
double distToSegment(const point &p, const segment &s) {
    /* returns distance of a point to a segment */
    if (getAngle(s.p2, s.p1, p) > PI/2 + EPS || getAngle(s.p1, s.p2, p) >
PI/2 + EPS) return min(dist(p,s.p1), dist(p,s.p2));
    return fabs(cross(s.p1 - p, s.p2 - p)) / dist(s.p1, s.p2);
}
double distToLine(const point &p, const segment &s){
    /* returns distance of a point to a line (its orthogonal projection) */
    return fabs(cross(s.p1 - p, s.p2 - p)) / dist(s.p1, s.p2);
}
point rotate(const point &p, const double &alpha) {
    /* rotates a point with respect to the origin. alpha in radians */
    return point(p.x * cos(alpha) - p.y * sin(alpha), p.x * sin(alpha) + p.y
* cos(alpha));
}
point rotate(const point &p, const point &center, const double &alpha){
    /* rotates a point with respect to point center. alpha in radians */
    return center + rotate(p - center, alpha);
}
point rescale(const point &p, const double s) {
    return point(p.x * s, p.y * s);
}
point dilate(const point &p, const double Factor){
    return rescale(p, Factor);
}
point dilate(const point &p, const point &center, double factor){
    return dilate(p- center, factor) + center;
}
bool isRightTurn(const point &p1, const point &p2, const point &p3){
    return cross(p2 - p1, p3 - p2) <= 0;
    /* straight returns true */
}
bool isOnSameSide(const point &p1, const point &p2, const segment &s){
    double z1 = cross(s.p2 - s.p1, p1 - s.p1);
    double z2 = cross(s.p2 - s.p1, p2 - s.p1);

```

```

    return (z1 + EPS < 0 && z2 + EPS < 0) || (0 < z1 - EPS && 0 < z2 - EPS)
|| fabs(z1) < EPS || fabs(z2) < EPS;
    /* on segment returns true */
}
bool isOnLine(const point &p, const segment &l){
    return fabs((l.p1.y - p.y) * (l.p2.x - p.x) - (l.p2.y - p.y) * (l.p1.x -
p.x)) < EPS;
}
bool isOnSegment(const point &p, const segment &s){
    return fabs(dist(p, s.p1) + dist(p, s.p2) - dist(s.p1, s.p2)) < EPS;
}
bool isIntersecting(const segment &s1, const segment &s2){
    return !(isOnSameSide(s1.p1,s1.p2,s2) || isOnSameSide(s2.p1,s2.p2,s1))
|| isOnSegment(s1.p1,s2) || isOnSegment(s1.p2,s2) || isOnSegment(s2.p1,s1)
|| isOnSegment(s2.p2,s1);
}
bool isParallel(const segment &s1, const segment &s2){
    return fabs((s1.p1.y-s1.p2.y)*(s2.p1.x-s2.p2.x)-(s2.p1.y-
s2.p2.y)*(s1.p1.x-s1.p2.x)) < EPS;
}
point intersection(const segment &s1, const segment &s2){
    /* assumes !isParallel(s1,s2) */
    double x1 = s1.p1.x - s1.p2.x;
    double x2 = s2.p1.x - s2.p2.x;
    double y1 = s1.p1.y - s1.p2.y;
    double y2 = s2.p1.y - s2.p2.y;
    double cross1 = cross(s1.p1, s1.p2);
    double cross2 = cross(s2.p1, s2.p2);
    return point ((cross1 * x2 - cross2 * x1) / (x1 * y2 - x2 * y1), (cross1
* y2 - cross2 * y1) / (x1 * y2 - x2 * y1));
}
point projection(const point &p, const segment &s){
    /* projects p onto line s */
    return rescale(s.p2 - s.p1, dot(p - s.p1, s.p2 - s.p1) / sqr(length(s)))
+ s.p1;
}

/** introducing circle */
struct circle {
    point center;
    double r;
    circle() { center = point(0, 0); r = 0; }
    circle(point p, double r) : center(p), r(r) {}
};
vector<point> intersectionLineCircle(const segment &l, const circle &c){
    vector<point> res;
    double dx = l.p2.x - l.p1.x;
    double dy = l.p2.y - l.p1.y;
    double dr = length(l);
    double d = cross(l.p1 - c.center,l.p2 - c.center);
    if (sqr(c.r) * sqr(dr) - sqr(d) + EPS < 0) return res;
    double det = sqrt(fabs(sqr(c.r) * sqr(dr) - sqr(d)));
    double sdx = dy < 0 ? -dx : dx;

```

```

double sdy = fabs(dy);
res.push_back(c.center + point((d*dy + sdx * det)/sqr(dr), (-d*dx + sdy
* det)/sqr(dr)));
if (det > EPS) res.push_back(c.center + point((d*dy - sdx *
det)/sqr(dr), (-d*dx - sdy * det)/sqr(dr)));
return res;
}
vector<point> intersectionSegmentCircle(const segment &s, const circle &c){
vector<point> res, _res = intersectionLineCircle(s,c);
for (vector<point>::iterator it = _res.begin(); it != _res.end(); ++it){
if (isOnSegment(*it,s)) res.push_back(*it);
}
return res;
}

```

Polygons (Area, Orientation)

```

/** introducing polygon */
typedef vector<point> polygon;

/** Check position of a point with respect to a polygon */
/* complexity : O(N) */
bool isPointInsidePolygon(point p, polygon poly){
/* ray casting to the right */
segment ray (p,p+point(1,0));
int n = (int)poly.size();
/* counts the number of intersections */
int nIntersection = 0;
for (int i = 0; i < n; ++i){
segment side(poly[i],poly[(i+1)%n]);
if (isOnSegment(p,side)) return false;
if (isParallel(ray,side)) continue;
point x = intersection(ray,side);
if (isOnSegment(x,side) && dot(x-p,ray.p2-p) > 0){
/* special case: x is one of vertices of sides */
if (x == side.p1){
if (isRightTurn(p,x,side.p2)) nIntersection ++;
}
else if (x == side.p2){
if (isRightTurn(p,x,side.p1)) nIntersection ++;
}
else nIntersection ++;
}
}
return nIntersection % 2 == 1;
}

```

Convex Hull

```

/** Convex Hull | monotone chain algorithm */
/* complexity : O(N log N) */
polygon convexHull(polygon p){
int m = 0, n = p.size();
polygon hull(2*n);

```

```

sort(p.begin(),p.end());
for (int i = 0; i < n; ++i){
while (m >= 2 && isRightTurn(hull[m-2],hull[m-1],p[i])) --m;
hull[m++] = p[i];
}
for (int i = n-1, t = m+1; i >= 0; --i){
while (m >= t && isRightTurn(hull[m-2],hull[m-1],p[i])) --m;
hull[m++] = p[i];
}
hull.resize(m);
return hull;
}

```

Delaunay Triangulation

```

// Slow but simple Delaunay triangulation. Does not handle
// degenerate cases (from O'Rourke, Computational Geometry in C)
//
// Running time: O(n^4)
//
// INPUT:    x[] = x-coordinates
//           y[] = y-coordinates
//
// OUTPUT:    triples = a vector containing m triples of indices
//               corresponding to triangle vertices

#include<vector>
using namespace std;

typedef double T;

struct triple {
int i, j, k;
triple() {}
triple(int i, int j, int k) : i(i), j(j), k(k) {}
};

vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y) {
int n = x.size();
vector<T> z(n);
vector<triple> ret;

for (int i = 0; i < n; i++)
z[i] = x[i] * x[i] + y[i] * y[i];

for (int i = 0; i < n-2; i++) {
for (int j = i+1; j < n; j++) {
for (int k = i+1; k < n; k++) {
if (j == k) continue;
double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-
y[i])*(z[j]-z[i]);
double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-
x[i])*(z[k]-z[i]);

```

```

double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-
x[i])*(y[j]-y[i]);
bool flag = zn < 0;
for (int m = 0; flag && m < n; m++)
    flag = flag && ((x[m]-x[i])*zn +
                    (y[m]-y[i])*yn +
                    (z[m]-z[i])*zn <= 0);
if (flag) ret.push_back(triple(i, j, k));
    }
}
return ret;
}

int main()
{
    T xs[]={0, 0, 1, 0.9};
    T ys[]={0, 1, 0, 0.9};
    vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
    vector<triple> tri = delaunayTriangulation(x, y);

    //expected: 0 1 3
    //           0 3 2

    int i;
    for (i = 0; i < tri.size(); i++)
        printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
    return 0;
}

```

Miscellaneous

Graph Theorems

Erdos-Gallai. A sequence of nonnegative integers $d_1 \geq \dots \geq d_n$ is a sequence of degree of an undirected graph iff $\sum d_i$ is even and $\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$

Fulkerson-Chen-Anstee. A sequence $((a_1, b_1), \dots, (a_n, b_n))$ of nonnegative integer pairs with $a_1 \geq \dots \geq a_n$ is a sequence of (in, outdeg) of a directed graph iff $\sum a_i = \sum b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i, k-1) + \sum_{i=k+1}^n \min(b_i, k)$

Lindstrom-Gessel-Viennot. The number of non-intersecting path from A to B in a directed acyclic graph is equal to the determinant of ... (elements (i, j) of matrix denotes the number of ways to go from A_i to B_j).

Koenig's. In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover.

Brook's. For any connected undirected graph G with maximum degree Δ , the chromatic number of G is at most Δ unless G is a complete graph or an odd cycle, in which case the chromatic number is $\Delta + 1$.

Combinatorics

Lucas Theorem. $\binom{n}{m} = \prod_{i=0}^k \binom{n_i}{m_i} \pmod{p}$ where $n = \overline{n_k n_{k-1} \dots n_0}$ and $m = \overline{m_k m_{k-1} \dots m_0}$ in base p .

Stirling Number of the First Kind. $s(n, k)$ denotes the number of n -permutation with k cycles. $s(n+1, k) = ns(n, k) + s(n, k-1)$.

Stirling Number of the Second Kind. $S(n, k)$ denotes the number of partition a set of n into k non-empty subsets. $S(n+1, k) = kS(n, k) + S(n, k-1)$. $S(n, k) = \frac{1}{k!} \sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^n$.

Gambler's Ruin. Two players with n_1 and n_2 points each are playing, each turn P1 has probability of winning p and P2 has probability $q = 1 - p$. The probability of P1 losing all his points is $\left(1 - \left(\frac{p}{q}\right)^{n_2}\right) / \left(1 - \left(\frac{p}{q}\right)^{n_1+n_2}\right)$.

Notes

std::lower_bound. Returns an iterator pointing to the first element in the range [first,last) which **does not compare less than** val.

std::upper_bound. Returns an iterator pointing to the first element in the range [first,last) which **compares greater** than val.



