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Preface

This is our Team Notebook for ACM ICPC and other Competitive Programming contests. Notable sources are:

- Introduction to Algorithm 3rd edition
- Competitive Programming 2 by Felix and Steven Halim
- Topcoder Algorithm Tutorials
- https://sites.google.com/site/indy256/
- http://stanford.edu/~liszt90/acm/notebook.html
- Dongskar Pedongi and DELAPAN.3gp Team Notebook
- Google, Wikipedia

Regards,

hehehe

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Template

C++

```
#include <bits/stdc++.h> <vector> <map> <set> <queue> <deque> <stack>
<algorithm> <sstream> <iostream> <iomanip> <fstream> <cstring> <cmath>
<cstdlib> <ctime> <cassert> <limits> <numeric> <utility>
using namespace std;
#ifdef DEBUG
 #define debug(...) printf( VA ARGS )
 #define GetTime() fprintf(stderr, "Running time: %.31f
second\n", ((double)clock())/CLOCKS PER SEC)
#else
 #define debug(...)
 #define GetTime()
#endif
//type definitions
typedef long long 11;
typedef double db;
typedef pair<int,int> pii;
typedef vector<int> vint;
//abbreviations
#define A first
#define B second
#define F first
#define S second
#define MP make pair
```

```
#define PB push back
//macros
\#define REP(i,n) for (int i = 0; i < (n); ++i)
#define REPD(i,n) for (int i = (n)-1; 0 \le i; --i)
#define FOR(i,a,b) for (int i = (a); i \le (b); ++i)
#define FORD(i,a,b) for (int i = (a); (b) \leq i; --i)
#define FORIT(it,c) for ( typeof ((c).begin()) it = (c).begin(); it !=
(c).end(); it++)
#define ALL(a) (a).begin(),(a).end()
#define SZ(a) ((int)(a).size())
#define RESET(a,x) memset(a,x,sizeof(a))
\#define EXIST(a,s) ((s).find(a) != (s).end())
\#define MX(a,b) a = max((a),(b));
\#define MN(a,b) a = min((a),(b));
inline void OPEN (const string &s) {
 freopen((s + ".in").c str(), "r", stdin);
 freopen((s + ".out").\overline{c} str(), "w", stdout);
/* ----- end of template ----- */
```

Graph Theory

Articulation Point

```
/** Articulation Point **/
/* complexity : O(|V| + |E|) */
#define MAXN 100100
int n, m, low[MAXN], num[MAXN], parent[MAXN], art[MAXN], root,
rootChildren, counter;
vector<int> adj[MAXN];
void dfs(int u) {
 low[u] = num[u] = counter++;
 FORIT(it, adj[u]) {
   int v = *it;
    if (num[v] == -1) {
      parent[v] = u;
      if (u == root) rootChildren++;
      dfs(v);
      if (low[v] >= num[u]) art[u] = 1;
      MN(low[u], low[v]);
    else if (v != parent[u]) {
      MN(low[u], num[v]);
```

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```
int main() {
// read the graph here. It should be 0-indexed
// initialization
 counter = 0;
 REP(i, n) {
  num[i] = -1;
  low[i] = parent[i] = art[i] = 0;
// perform the dfs
 REP(i, n) {
   if (num[i] == -1) {
     root = i, rootChildren = 0;
     dfs(i);
     art[root] = (rootChildren > 1);
   }
 }
// now the articulation points are stored in art[]
 return 0;
```

Articulation Bridge

```
/** Bridge **/
/* complexity : O(|V| + |E| + |E| log |E|) */
#define MAXN 100100
int n, low[MAXN], num[MAXN], parent[MAXN], bridge[MAXN], counter;
vector<pii> adj[MAXN]; // adj[u].PB(MP(v, idx of edge));
void dfs(int u) {
 low[u] = num[u] = counter++;
 FORIT(it, adj[u]) {
   int v = it-A;
   if (num[v] == -1) {
     parent[v] = u;
     dfs(v);
     if (low[v] > num[u]) bridge[it->B] = 1;
     MN(low[u], low[v]);
   else if (v != parent[u]) {
     MN(low[u], num[v]);
 }
int main() {
 // read the graph here. it should be 0-indexed
 // should not work if multiple edges exist
 // initialization
  counter = 0;
 REP(i, n) {
```

```
num[i] = -1;
  low[i] = parent[i] = 0;
}
REP(i, m) {
  bridge[i] = 0;
}
// perform the dfs
REP(i, n) {
  if (num[i] == -1) {
    dfs(i);
  }
}
// the bridges are stored in bridge[]
return 0;
}
```

Tarjan's Directed SCC

```
/** Tarjan's Directed Strongly Connected Component **/
/* complexity : O(|V| + |E|) */
#define MAXN 100100
int n, low[MAXN], num[MAXN], visited[MAXN], counter;
vector<int> adj[MAXN], s;
vector<vector<int> > scc;
void dfs(int u) {
 low[u] = num[u] = counter++;
 s.PB(u);
 visited[u] = 1;
 FORIT(it, adj[u]) {
   int v = *it;
    if (num[v] == -1) dfs(v);
    if (visited[v]) {
     MN(low[u], low[v]);
 if (low[u] == num[u]) {
   vector<int> temp;
   int v = -1;
   while (u != v) {
     v = s.back(); s.pop back(); visited[v] = 0;
     temp.PB(v);
    scc.PB(temp);
int main() {
 // read the graph here. it should be 0-indexed
 // initialization
```

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```
counter = 0;
scc.clear();
REP(i, n) {
    num[i] = -1;
    low[i] = visited[i] = 0;
}
// perform the dfs
REP(i, n) {
    if (num[i] == -1) {
        dfs(i);
    }
}
// the components are stored in scc
return 0;
}
```

Max Flow

```
#define MAXN 1100
#define INF 0x3FFFFFFF
int res[MAXN][MAXN], vis[MAXN];
/** Maximum Flow **/
/* Edmond Karp | complexity : O(|V|*(|V|+|E|)) */
void augment (int v, int minEdge, int &s, int &f, vector <int > &p) {
 if (v == s) { f = minEdge; return; }
 else if (p[v] != -1) {
   augment(p[v], min(minEdge, res[p[v]][v]), s,f,p); res[p[v]][v]-= f;
res[v][p[v]] += f;
 - }
int maxFlowEdmondKarp(int n, int source, int target) {
 int mf = 0;
 while (1) {
  int f = 0;
   vector<int> dist(n+5,INF);
   dist[source] = 0;
   queue<int> q; q.push(source);
   vector<int> p; p.assign(n+5,-1);
   while (!a.emptv()) {
     int u = q.front(); q.pop();
     if (u == target) break;
     for (int v = 0; v < n; v++)
       if (res[u][v] > 0 && dist[v] == INF)
          dist[v] = dist[u] + 1, q.push(v), p[v] = u;
     augment(target, INF, source, f, p);
     if (f == 0) break;
     mf += f;
   return mf;
```

```
/* Ford Fulkerson | complexity : O(|V|^2 F) */
 int findPath(int n, int u, int t, int f){
   if (u == t) return f;
   vis[u] = 1;
   for (int v = 0; v < n; ++v){
      if (!vis[v] && res[u][v] > 0){
        int df = findPath(n, v, t, min(f,res[u][v]));
       if (df > 0) {
         res[u][v] -= df;
         res[v][u] += df;
         return df:
   return 0;
 int maxFlowFordFulkerson(int n, int source, int target) {
   for (int flow = 0;;) {
      for (int i = 0; i < n; ++i) vis[i] = 0;
       int df = findPath (n, source, target, INF);
      if (df == 0) return flow;
     flow += df;
   }
/* WARNING: res will be modified during the process */
```

Max Flow Min Cost

```
/** Max Flow Min Cost **/
/* complexity: O(min(E^2 V log V, E log V F)) */
const int max nodes = 200000;
int nodes = maxnodes;
int prio[maxnodes], curflow[maxnodes], prevedge[maxnodes],
prevnode[maxnodes], q[maxnodes], pot[maxnodes];
bool inqueue[maxnodes];
struct Edge {
 int to, f, cap, cost, rev;
vector<Edge> graph[maxnodes];
void addEdge(int s, int t, int cap, int cost) {
  Edge a = \{t, 0, cap, cost, graph[t].size()\};
  Edge b = \{s, 0, 0, -\cos t, \operatorname{graph}[s].\operatorname{size}()\};
  graph[s].push back(a);
  graph[t].push back(b);
void bellmanFord(int s, int dist[]) {
 fill(dist, dist + nodes, 1000000000);
```

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```
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```

```
dist[s] = 0;
 int at = 0:
 q[qt++] = s;
 for (int qh = 0; (qh - qt) % nodes != 0; qh++) {
   int u = q[qh % nodes];
   inqueue[u] = false;
   for (int i = 0; i < (int) graph[u].size(); <math>i++) {
     Edge &e = graph[u][i];
     if (e.cap <= e.f) continue;</pre>
     int v = e.to;
     int ndist = dist[u] + e.cost;
     if (dist[v] > ndist) {
       dist[v] = ndist;
       if (!inqueue[v]) {
         inqueue[v] = true;
          q[qt++ % nodes] = v;
       - }-
   - }
 }
pii minCostFlow(int s, int t, int maxf) {
 // bellmanFord can be safely commented if edges costs are non-negative
 bellmanFord(s, pot);
 int flow = 0;
 int flowCost = 0;
 while (flow < maxf) {</pre>
 priority queue<11, vector<11>, greater<11> > q;
   q.push(s);
   fill(prio, prio + nodes, 1000000000);
   prio[s] = 0;
   curflow[s] = 10000000000;
   while (!q.empty()) {
     11 \text{ cur} = q.top();
     int d = cur \gg 32;
     int u = cur;
     q.pop();
      if (d != prio[u]) continue;
      for (int i = 0; i < (int) graph[u].size(); <math>i++) {
       Edge &e = graph[u][i];
       int v = e.to;
       if (e.cap <= e.f) continue;</pre>
       int nprio = prio[u] + e.cost + pot[u] - pot[v];
       if (prio[v] > nprio) {
         prio[v] = nprio;
         q.push(((ll) nprio \ll 32) + v);
         prevnode[v] = u;
         prevedge[v] = i;
          curflow[v] = min(curflow[u], e.cap - e.f);
```

```
if (prio[t] == 10000000000) break;
for (int i = 0; i < nodes; i++) pot[i] += prio[i];
   int df = min(curflow[t], maxf - flow);
flow += df;
for (int v = t; v != s; v = prevnode[v]) {
   Edge &e = graph[prevnode[v]][prevedge[v]];
   e.f += df;
   graph[v][e.rev].f -= df;
   flowCost += df * e.cost;
   }
}
return make_pair(flow, flowCost);

/* usage example:
   * addEdge (source, target, capacity, cost)
   * minCostFlow(source, target, INF) -> <flow, flowCost>
   */
```

Lowest Common Ancestor

```
/** Lowest Common Ancestor **/
/* complexity : LCApre : O(N log N), LCAquery : O(log N) */
/* legend:
 * N : number of vertices. WARNING: zero based
 * T : direct parent. T[v] is parent of v
 * L : L[v] is the level of v. zero/one based is okay
 * P : dp table of size [MAXN] [LOGMAXN]. P[v][i] is the 2^i-th parent of v
#define MAXN 100100
#define LOGMAXN 18
int L[MAXN], P[MAXN][LOGMAXN], T[MAXN], N;
void pre(){
 int i, j;
  //we initialize every element in P with -1
  for (i = 0; i < N; i++) {
   for (j = 0; 1 << j < N; j++) {
      P[i][j] = -1;
   - }
  1
  //the first ancestor of every node i is T[i]
  for (i = 0; i < N; i++) {
   P[i][0] = T[i];
  //bottom up dynamic programing
  for (j = 1; 1 << j < N; j++) {
    for (i = 0; i < N; i++) {
```

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```
if (P[i][j-1] != -1) {
       P[i][j] = P[P[i][j-1]][j-1];
   }
}
int query(int p, int q){
int log, i;
 //if p is situated on a higher level than q then we swap them
 if (L[p] < L[q]) {</pre>
  swap(p,q);
 //we compute the value of [log(L[p)]
 for (log = 1; 1 << log <= L[p]; log++);</pre>
   log--;
 //we find the ancestor of node p situated on the same level
 //with q using the values in P
 for (i = log; i >= 0; i--) {
   if (L[p] - (1 << i) >= L[q]) {
     p = P[p][i];
 }
 if (p == q) return p;
     //we compute LCA(p, q) using the values in P
 for (i = log; i >= 0; i--) {
   if (P[p][i] != -1 && P[p][i] != P[q][i]) {
     p = P[p][i];
     q = P[q][i];
 }
 return T[p];
```

Blossom

```
/** Maximum Matching on General Graph **/
/* Blossom | O(V^3) */
int lca(vector<int> &match, vector<int> &base, vector<int> &p, int a, int
b) {
  vector<bool> used(SZ(match));
  while (true) {
    a = base[a];
    used[a] = true;
    if (match[a] == -1) break;
    a = p[match[a]];
```

```
while (true) {
   b = base[b];
   if (used[b]) return b;
   b = p[match[b]];
 return -1;
void markPath(vector<int> &match, vector<int> &base, vector<bool> &blossom,
vector<int> &p, int v, int b, int children) {
 for (; base[v] != b; v = p[match[v]]) {
   blossom[base[v]] = blossom[base[match[v]]] = true;
    p[v] = children;
   children = match[v];
int findPath(vector<vector<int> > &graph, vector<int> &match, vector<int>
&p, int root) {
 int n = SZ(graph);
 vector<bool> used(n);
 FORIT(it, p) *it = -1;
 vector<int> base(n);
 for (int i = 0; i < n; ++i) base[i] = i;
 used[root] = true;
 int qh = 0;
 int qt = 0;
 vector<int> q(n);
 q[qt++] = root;
  while (qh < qt) {</pre>
   int v = q[qh++];
   FORIT(it, graph[v]) {
      int to = *it;
      if (base[v] == base[to] || match[v] == to) continue;
      if (to == root || match[to] != -1 && p[match[to]] != -1) {
       int curbase = lca(match, base, p, v, to);
        vector<bool> blossom(n);
        markPath(match, base, blossom, p, v, curbase, to);
        markPath(match, base, blossom, p, to, curbase, v);
        for (int i = 0; i < n; ++i) {
         if (blossom[base[i]]) {
           base[i] = curbase;
           if (!used[i]) {
             used[i] = true;
              q[qt++] = i;
      } else if (p[to] == -1) {
       p[to] = v;
        if (match[to] == -1) return to;
```

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```
to = match[to];
       used[to] = true;
       q[qt++] = to;
   }
 return -1;
int maxMatching(vector<vector<int> > graph) {
 int n = SZ(graph);
 vector<int> match(n, -1);
 vector<int> p(n);
 for (int i = 0; i < n; ++i) {
  if (match[i] == -1) {
     int v = findPath(graph, match, p, i);
     while (v != -1) {
       int pv = p[v];
       int ppv = match[pv];
       match[v] = pv;
       match[pv] = v;
       v = ppv;
 int matches = 0;
 for (int i = 0; i < n; ++i) {
  if (match[i] != -1) {
     ++matches;
   }
 return matches / 2;
```

Minimum Cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
//
// Running time:
// O(|V|^3)
//
// INPUT:
// - graph, constructed using AddEdge()
//
// OUTPUT:
// - (min cut value, nodes in half of min cut)
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
```

```
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
 int N = weights.size();
 VI used(N), cut, best cut;
 int best weight = -1;
  for (int phase = N-1; phase >= 0; phase--) {
   VI w = weights[0];
    VI added = used;
    int prev, last = 0;
    for (int i = 0; i < phase; <math>i++) {
      prev = last;
      last = -1:
      for (int j = 1; j < N; j++)
        if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
      if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j];</pre>
          for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];</pre>
            used[last] = true;
        cut.push back(last);
        if (best weight == -1 || w[last] < best weight) {</pre>
          best cut = cut;
          best weight = w[last];
      } else {
        for (int j = 0; j < N; j++)
          w[j] += weights[last][j];
        added[last] = true;
 return make pair(best weight, best cut);
```

String Processing

Knuth-Morris-Pratt

```
/** Knuth-Morris-Pratt **/
/* Complexity: O(N) */
void buildFailTable(char *pattern, int *t) {
  int i = 0, j = -1, m = strlen(pattern);
  t[0] = -1;
  while (i < m) {
    while (j >= 0 && pattern[i] != pattern[j]) j = t[j];
    i++; j++;
  t[i] = j;
```

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```
}

vector<int> kmpSearch(char *pattern, char *text) {
  vector<int> res;
  int i = 0, j = 0, n = strlen(text), m = strlen(pattern);
  int t[m+5];
  buildFailTable(pattern,t);
  while (i < n) {
    while (j >= 0 && text[i] != pattern[j]) j = t[j];
    i++; j++;
    if (j == m) {
        res.push_back(i-j);
        j = t[j];
    }
    return res;
}
```

Z-Algorithm

```
/* Z-Algorithm */
// Z[i] is the longest substring starting from i which is also a prefix of
// Z[0] is not set
int L = 0, R = 0;
for (int i = 1; i < n; ++i) {
 if (i > R) {
   L = R = i;
    while (R < n && s[R] == s[R-L]) ++R;
   Z[i] = R-L; --R;
 1
  else {
   int k = i-L;
   if (Z[k] < R-i+1) Z[i] = Z[k];
    else {
     L = i;
      while (R < n && s[R] \Longrightarrow s[R-L]) ++R;
      Z[i] = R-L; --R;
 - }-
```

Suffix Array

```
/** Suffix Array **/
/* complexity: O(N log N) */

#define MAXN 200000

char T[MAXN+5]; // input
int n; // length
int RA[MAXN+5], tempRA[MAXN+5]; // rank array
int SA[MAXN+5], tempSA[MAXN+5]; // suffix array
```

```
int c[MAXN+5]; //for counting/radix sort
void countingSort(int k) {
 int sum, maxi = max(300,n);
 memset(c,0,sizeof(c));
 for (int i = 0; i < n; i++)
   c[i+k < n ? RA[i+k] : 0]++;
 for (int i = sum = 0; i < maxi; i++) {</pre>
   int t = c[i]; c[i] = sum;
    sum += t;
 for (int i = 0; i < n; i++)
    tempSA[c[SA[i]+k<n?RA[SA[i]+k]:0]++] = SA[i];
 for (int i = 0; i < n; i++) SA[i] = tempSA[i];
void SuffixArray Construct() {
 int r:
  for (int i = 0; i < n; i++) RA[i] = T[i]-'.';</pre>
 for (int i = 0; i < n; i++) SA[i] = i;
 for (int k = 1; k < n; k <<= 1) {
   countingSort(k);
   countingSort(0);
   tempRA[SA[0]] = r = 0;
    for (int i = 1; i < n; i++)
    tempRA[SA[i]] =
    (RA[SA[i]] == RA[SA[i-1]] \& RA[SA[i]+k] == RA[SA[i-1]+k]) ? r : ++r;
    for (int i = 0; i < n; i++) RA[i] = tempRA[i];
```

Suffix Tree

```
/** Suffix Tree Ukkonen's algorithm **/
/* Complexity: O(N) (Warning: large multiplier) */
const string ALPHABET = "abcdefghijklmnopgrstuvwxyz0123456789\1\2";
const int NALPHABET = 38;
struct Node {
 int begin, end, depth;
 Node* parent;
 Node** children:
 Node* suffixLink:
 Node (int begin, int end, int depth, Node* parent) {
    this->begin = begin;
    this->end = end;
    this->depth = depth;
    this->parent = parent;
    this->children = new Node*[NALPHABET];
    for (int i = 0; i < NALPHABET; ++i) {</pre>
      this->children[i] = NULL:
```

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```
~Node() {
   delete[] children;
};
Node* buildSuffixTree(string s) {
 int n = s.length();
 char* a = new char[n];
 for (int i = 0; i < n; ++i) {
  a[i] = (char) ALPHABET.find(s[i]);
 Node* root = new Node(0, 0, 0, NULL);
 Node* node = root;
 for (int i = 0, tail = 0; i < n; ++i, ++tail) {
   Node* last = NULL;
   while (tail >= 0) {
     Node* ch = node->children[a[i - tail]];
     while (ch != NULL && tail >= ch->end - ch->begin) {
      tail -= (ch->end - ch->begin);
       node = ch;
       ch = ch->children[a[i - tail]];
     if (ch == NULL) {
       node->children[a[i]] = new Node(i, n, node->depth + node->end -
node->begin, node);
       if (last != NULL) {
         last->suffixLink = node;
       last = NULL:
     } else {
       char t = a[ch->begin + tail];
       if (t == a[i]) {
         if (last != NULL) {
           last->suffixLink = node;
         break;
       } else {
         Node* splitNode = new Node(ch->begin, ch->begin + tail, node-
>depth + node->end - node->begin, node);
         splitNode->children[a[i]] = new Node(i, n, ch->depth + tail,
splitNode):
         splitNode->children[t] = ch;
         ch->begin += tail;
         ch->depth += tail;
         ch->parent = splitNode;
         node->children[a[i - tail]] = splitNode;
         if (last != NULL) {
           last->suffixLink = splitNode;
```

```
last = splitNode;
      if (node == root) {
        --tail;
      } else {
        node = node->suffixLink;
    }
  delete[] a;
  return root:
/* Example: longest common substring */
int lcsLength;
Node* lcsNode;
int traverseLCS(Node* node, const vector<int>& stops, const int target) {
 for (int i = 0; i < stops.size(); ++i) {</pre>
    if (node->begin <= stops[i] && stops[i] < node->end) {
      return 1 << i:
  int mask = 0;
  for (int f = 0; f < ALPHABET.length(); ++f) {</pre>
    if (node->children[f] != NULL) {
      mask |= traverseLCS(node->children[f], stops, target);
  if (mask == target) {
   int curLength = node->depth + node->end - node->begin;
   if (lcsLength < curLength) {</pre>
      lcsLength = curLength;
      lcsNode = node;
  return mask;
int longestCommonSubstring(const vector<string> &ss) {
 int totalN = 0;
 int n = ss.size();
  for (int i = 0; i < n; ++i) {
   totalN += ss[i].length() + 1;
```

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```
string s:
s.resize(totalN);
int offset = 0:
vector<int> stops:
for (int i = 0; i < n; ++i) {
 for (int j = 0; j < ss[i].length(); ++j) {
    s[offset + j] = ss[i][j];
 offset += ss[i].length() + 1;
 s[offset - 1] = '0' + i;
 stops.push back(offset - 1);
Node* tree = buildSuffixTree(s);
lcsLength = 0;
lcsNode = NULL;
traverseLCS(tree, stops, (1 << n) - 1);
delete tree;
return lcsLength;
```

Aho-Corasick

```
/** Aho-Corasick Dictionary Matching **/
const int NALPHABET = 26;
struct Node {
 Node** children, go;
 bool leaf:
 char charToParent;
 Node* parent, suffLink, dictSuffLink;
 int count, value;
 Node() {
   children = new Node*[NALPHABET];
   go = new Node*[NALPHABET];
   for (int i = 0; i < NALPHABET; ++i) {
     children[i] = go[i] = NULL;
   parent = suffLink = dictSuffLink = NULL;
   leaf = false;
   count = 0;
 }
};
Node* createRoot() {
 Node* node = new Node();
 node->suffLink = node;
 return node;
void addString(Node* node, const string& s, int value = -1) {
for (int i = 0; i < s.length(); ++i) {
```

```
int c = s[i] - 'a';
    if (node->children[c] == NULL) {
      Node* n = new Node();
      n->parent = node;
      n->charToParent = s[i];
      node \rightarrow children[c] = n;
   node = node->children[c];
  node->leaf = true;
  node->count++;
  node->value = value;
Node* suffLink(Node* node);
Node* dictSuffLink(Node* node);
Node* go (Node* node, char ch);
int calc (Node* node);
Node* suffLink(Node* node) {
  if (node->suffLink == NULL) {
    if (node->parent->parent == NULL) {
      node->suffLink = node->parent;
      node->suffLink = go(suffLink(node->parent), node->charToParent);
  return node->suffLink;
Node* dictSuffLink(Node* node) {
  if (node->dictSuffLink == NULL) {
    Node* n = suffLink(node);
    if (node == n) {
      node->dictSuffLink = node;
    } else {
      while (!n->leaf && n->parent != NULL) {
        n = dictSuffLink(n);
      node->dictSuffLink = n;
   }
  return node->dictSuffLink;
Node* go (Node* node, char ch) {
 int c = ch - 'a';
 if (node->go[c] == NULL) {
    if (node->children[c] != NULL) {
      node->go[c] = node->children[c];
      node->go[c] = node->parent == NULL ? node : go(suffLink(node), ch);
```

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```
return node->go[c]:
int calc(Node* node) {
 if (node->parent == NULL) {
 return 0;
 } else {
   return node->count + calc(dictSuffLink(node));
 }
int main() {
 Node* root = createRoot();
 addString(root, "a", 0);
 addString(root, "aa", 1);
 addString(root, "abc", 2);
 string s("abcaadc");
 Node* node = root;
 for (int i = 0; i < s.length(); ++i) {
  node = qo(node, s[i]);
  Node* temp = node;
   while (temp != root) {
    if (temp->leaf) {
       printf("string (%d) occurs at position %d\n", temp->value, i);
     temp = dictSuffLink(temp);
   - }
 }
 return 0;
```

Mathematics

Extended Euclid

```
/** Extended Euclid | returns <x,y> where ax + by = gcd(a,b) **/
/* complexity: O(min(log(a),log(b))) */
pair<11,ll> extendedEuclid(ll a, ll b) {
    ll x = 0, y = 1, lastx = 1, lasty = 0;
    while (b != 0) {
        ll quotient = a / b;
        /* (a, b) = (b, a mod b) */
        ll temp = a;
        a = b;
        b = temp % b;
        /* (x, lastx) = (lastx - quotient*x, x) */
        temp = x;
        x = lastx - quotient * x;
        lastx = temp;
```

```
/* (y, lasty) = (lasty - quotient*y, y) */
temp = y;
y = lasty - quotient * y;
lasty = temp;
}
return make_pair(lastx, lasty);
}
```

Diophantine

```
// computes x and y such that ax + by = c; on failure, x = y =-1
void linear_diophantine(int a, int b, int c, int &x, int &y) {
  int d = gcd(a,b);
  if (c%d) {
    x = y = -1;
  } else {
    x = c/d * mod_inverse(a/d, b/d);
    y = (c-a*x)/b;
  }
}
```

Chinese Reminder Theorem

```
// Chinese remainder theorem (special case): find z such that
// z % x = a, z % y = b. Here, z is unique modulo M = lcm(x,y).
// Return (z, M). On failure, M = -1.
PII chinese remainder theorem (int x, int a, int y, int b) {
 int s, t;
 int d = extended euclid(x, y, s, t);
 if (a%d != b%d) return make pair(0, -1);
 return make pair (mod(s*b*x+t*a*y,x*y)/d, x*y/d);
// Chinese remainder theorem: find z such that
// z % x[i] = a[i] for all i. Note that the solution is
// unique modulo M = lcm i (x[i]). Return (z,M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese remainder theorem (const VI &x, const VI &a) {
 PII ret = make pair(a[0], x[0]);
 for (int i = 1; i < x.size(); i++) {
    ret = chinese remainder theorem(ret.second, ret.first, x[i], a[i]);
    if (ret.second == -1) break;
  return ret:
```

Rabin Miller Primality Test

```
/** Works for all 64-bit integers **/
bool rabinMillerPrimalityTest(long long n) {
   if ((n & 1) == 0) return n == 2;
   if (n == 1) return false;

long long a[] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};
```

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```
int s = 0;
long long d = n - 1:
while ((d \& 1) == 0) {
 d /= 2LL:
 ++s;
for (int i = 0; i < 12; ++i) {
 if (a[i] >= n) break;
 long long ad = powerMod(a[i], d, n);
 if (ad != 1) {
   bool composite = true;
   for (int j = 0; j < s; ++j) {
     if (ad == n - 1) {
        composite = false;
       break:
      ad = (ad * ad) % n;
    if (composite) return false;
return true;
```

Lagrange Interpolation

```
/** Lagrange Polynomial Interpolation **/
/* complexity: O(n^2) */
class lagrangeInterpolation {
public:
 lagrangeInterpolation (): x (0), y (0) {}
 void addCoef (db x, db y){
   x .push back(x);
   y .push back(y);
 db interpolate (db x) {
   db value = 0;
   for (int i = 0; i < (int)x .size(); ++i){
     db \ addum = y [i];
     for (int j = 0; j < (int)x .size(); ++j) if (i != j){
      addum *= (x - x [j]);
       addum /= (x [i] - x [j]);
     value += addum;
   return value;
 vector<db> x , y ;
};
```

```
class modularInterpolation {
public:
  modularInterpolation (const 11 &modu) : modu (modu), x (0), y (0) {}
  void addCoef (ll x, ll y) {
   x %= modu ;
   if (x < OLL) x += modu ;
   x .push back(x);
    y %= modu ;
    if (y < 0LL) y += modu ;</pre>
    y .push back (y);
  ll interpolate (ll x) {
   x %= modu ;
    if (x < OLL) x += modu ;
    for (int i = 0; i < (int)x . size(); ++i) if (x [i] == x) return y [i];
      11 value = OLL;
    for (int i = 0; i < (int)x .size(); ++i){
      ll addum = y [i];
      for (int j = 0; j < (int)x .size(); ++j) if (<math>j != i){
        ll delta1 = (x - x [j] + modu) % modu;
        ll delta2 = (x [i] - x [j] + modu ) % modu ;
        addum = (addum * delta1) % modu ;
        addum = (addum * multInverse(delta2, modu )) % modu ;
      value += addum;
      value %= modu ;
   return value:
  const 11 modu ;
  vector<ll> x , y ;
/* WARNING: no two x [i] should be the same */
```

Fast Fourier Transform

```
/** Fast Fourier Transform **/
/* complexity: 0(N log N) */
vector< complex<db> > iterativeDFT (const vector< complex<db> > &seq, int
direction) {
  int n = SZ(seq);
  int bits = 0;
  int tmp_n = n;
  complex<db> *placeholder = new complex<db>[n];
  complex<db> *tmp = new complex<db>[n];
```

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```
while (tmp n > 1) {
   ++bits;
   tmp n /= 2;
 REP(i,n){
   int res = 0;
   int tmp i = i;
 REP(j,bits){
     if (tmp i % 2) res += (1 << (bits-j-1));
     tmp i /= 2;
   placeholder[i] = seq[res];
 for (int comp size = 2; comp size <= n; comp size *= 2){</pre>
   for (int j = 0; j < n; j += comp size) {
     int n mem = comp size / 2;
     db w mult exp i = 2. * acos(-1.) / (db)comp size;
     if (!direction) w mult exp i *= -1.;
     complex<db> w mult (cos(w mult exp_i),sin(w_mult_exp_i));
     complex<db> w (1., 0.);
     for (int k = 0; k < comp size; ++k){
       int idx = k % n mem;
       tmp[k] = placeholder[j+idx] + w * placeholder[j+n mem+idx];
       w = w * w mult;
     for (int k = 0; k < comp size; ++k){
       placeholder[j+k] = tmp[k];
 vector< complex<db> > result;
 for (int i = 0; i < n; ++i) result.PB(placeholder[i]);</pre>
 delete[] placeholder;
 delete[] tmp;
 return result;
vector<db> FFT (vector<db> a, vector<db> b) {
 if (SZ(a) == 0) a.PB(0.);
 if (SZ(b) == 0) b.PB(0.);
 int n final elements = SZ(a) + SZ(b) - 1;
 int actual size = 1;
 while (actual size < max(SZ(a), SZ(b))) {
  actual size *= 2;
 actual size *= 2;
 while (SZ(a) < actual size) a.PB(0.);</pre>
 while (SZ(b) < actual size) b.PB(0.);</pre>
```

```
vector< complex<db> > dft_input_a, dft_input_b;
REP(i,actual_size) {
    dft_input_a.PB(complex<db> (a[i], 0.));
    dft_input_b.PB(complex<db> (b[i], 0.));
}

dft_input_a = iterativeDFT (dft_input_a, 1);
dft_input_b = iterativeDFT (dft_input_b, 1);
REP(i,actual_size) {
    dft_input_a[i] = dft_input_a[i] * dft_input_b[i];
}
dft_input_a = iterativeDFT (dft_input_a, 0);

vector<db> res;
REP(i,n_final_elements) {
    res.PB(dft_input_a[i].real() / (db) actual_size);
}
return res;
}
```

Karatsuba

```
typedef vector<long long> vll;
vll karatsubaMultiply(const vll &a, const vll &b) {
 int n = a.size();
 vll res(n + n);
 if (n <= 32) {
   for (int i = 0; i < n; i++)
     for (int j = 0; j < n; j++)
       res[i + j] += a[i] * b[j];
    return res;
 int k = n \gg 1;
 vll al(a.begin(), a.begin() + k);
 vll a2(a.begin() + k, a.end());
 vll b1(b.begin(), b.begin() + k);
 vll b2(b.begin() + k, b.end());
 vll a1b1 = karatsubaMultiply(a1, b1);
 vll a2b2 = karatsubaMultiply(a2, b2);
 for (int i = 0; i < k; i++)
   a2[i] += a1[i];
 for (int i = 0; i < k; i++)
   b2[i] += b1[i];
 vll r = karatsubaMultiply(a2, b2);
 for (int i = 0; i < (int) alb1.size(); i++)</pre>
   r[i] -= a1b1[i];
  for (int i = 0; i < (int) a2b2.size(); i++)
```

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```
r[i] -= a2b2[i];

for (int i = 0; i < (int) r.size(); i++)
    res[i + k] += r[i];
for (int i = 0; i < (int) alb1.size(); i++)
    res[i] += alb1[i];
for (int i = 0; i < (int) a2b2.size(); i++)
    res[i + n] += a2b2[i];
return res;
}</pre>
```

Simplex

```
// Two-phase simplex algorithm for solving linear programs of the form
// maximize c^T x
// subject to Ax <= b
                 x >= 0
11
// INPUT: A -- an m x n matrix
    b -- an m-dimensional vector
      c -- an n-dimensional vector
11
      x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
       above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD:
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
 int m, n;
 VI B, N;
 VVD D:
 LPSolver (const VVD &A, const VD &b, const VD &c) :
     m(b.size()), n(c.size()), N(n+1), B(m), D(m+2, VD(n+2)) {
   for (int i = 0; i < m; i++)
```

```
for (int j = 0; j < n; j++) D[i][j] = A[i][j];
    for (int i = 0; i < m; i++) {
      B[i] = n+i; D[i][n] = -1; D[i][n+1] = b[i];
    for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
      N[n] = -1; D[m+1][n] = 1;
  void Pivot(int r, int s) {
    for (int i = 0; i < m+2; i++) if (i != r)
      for (int j = 0; j < n+2; j++) if (j != s)
        D[i][j] = D[r][j] * D[i][s] / D[r][s];
    for (int j = 0; j < n+2; j++) if (j != s) D[r][j] /= D[r][s];
      for (int i = 0; i < m+2; i++) if (i != r) D[i][s] /= -D[r][s];
        D[r][s] = 1.0 / D[r][s];
    swap(B[r], N[s]);
  bool Simplex(int phase) {
    int x = phase == 1 ? m+1 : m;
    while (true) {
      int s = -1;
      for (int j = 0; j \le n; j++) {
        if (phase == 2 && N[j] == -1) continue;
        if (s == -1 \mid | D[x][\dot{\gamma}] < D[x][s] \mid | D[x][\dot{\gamma}] == D[x][s] && N[\dot{\gamma}] <
N[s]) s = j;
      if (D[x][s] >= -EPS) return true;
      int r = -1;
      for (int i = 0; i < m; i++) {
        if (D[i][s] <= 0) continue;</pre>
        if (r == -1 || D[i][n+1] / D[i][s] < D[r][n+1] / D[r][s] ||
          D[i][n+1] / D[i][s] == D[r][n+1] / D[r][s] && B[i] < B[r]) r = i;
      if (r == -1) return false;
      Pivot(r, s);
  }
  DOUBLE Solve (VD &x) {
    int r = 0:
    for (int i = 1; i < m; i++) if (D[i][n+1] < D[r][n+1]) r = i;
    if (D[r][n+1] \leftarrow -EPS) {
      Pivot(r, n);
      if (!Simplex(1) || D[m+1][n+1] < -EPS) return -</pre>
numeric limits < DOUBLE > :: infinity();
      for (int i = 0; i < m; i++) if (B[i] == -1) {
        int s = -1:
        for (int j = 0; j \le n; j++)
          if (s == -1 \mid | D[i][j] < D[i][s] \mid | D[i][j] == D[i][s] && N[j] <
N[s]) s = i;
        Pivot(i, s);
```

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```
if (!Simplex(2)) return numeric limits<DOUBLE>::infinity();
   for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n+1];
     return D[m][n+1];
 }
};
int main() {
 const int m = 4;
 const int n = 3;
 DOUBLE A[m][n] = {
   \{ 6, -1, 0 \},
   \{-1, -5, 0\},\
   { 1, 5, 1 },
  \{-1, -5, -1\}
 DOUBLE b[m] = \{ 10, -4, 5, -5 \};
 DOUBLE c[n] = \{ 1, -1, 0 \};
 VVD A(m);
 VD b(b, b+m);
 VD c(c, c+n);
 for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);</pre>
 LPSolver solver (A, b, c);
 VD x;
 DOUBLE value = solver.Solve(x);
 cerr << "VALUE: "<< value << endl;</pre>
 cerr << "SOLUTION:";</pre>
 for (size t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
   cerr << endl:
 return 0:
```

Gauss Jordan Flimination

```
// Gauss-Jordan elimination with full pivoting.
//
// Uses:
// (1) solving systems of linear equations (AX=B)
// (2) inverting matrices (AX=I)
// (3) computing determinants of square matrices
//
// Running time: O(n^3)
//
// INPUT: a[][] = an nxn matrix
// b[][] = an nxm matrix
//
// OUTPUT: X = an nxm matrix (stored in b[][])
// A^{-1} = an nxn matrix (stored in a[][])
// returns determinant of a[][]
```

```
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a, VVT &b) {
 const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T \det = 1;
  for (int i = 0; i < n; i++) {
   int pj = -1, pk = -1;
    for (int j = 0; j < n; j++) if (!ipiv[j])
      for (int k = 0; k < n; k++) if (!ipiv[k])
        if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k;}
    if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl;</pre>
exit(0);}
    ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
    irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
    a[pk][pk] = 1.0;
    for (int p = 0; p < n; p++) a[pk][p] *= c;
    for (int p = 0; p < m; p++) b[pk][p] *= c;
    for (int p = 0; p < n; p++) if (p != pk) {
      c = a[p][pk];
      a[p][pk] = 0;
      for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
      for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  }
  for (int p = n-1; p \ge 0; p--) if (irow[p] != icol[p]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);
  return det;
```

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```
int main() {
 const int n = 4;
 const int m = 2;
 double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \};
 double B[n][m] = \{\{1,2\},\{4,3\},\{5,6\},\{8,7\}\};
 VVT a(n), b(n);
 for (int i = 0; i < n; i++) {
   a[i] = VT(A[i], A[i] + n);
   b[i] = VT(B[i], B[i] + m);
 double det = GaussJordan(a, b);
 // expected: 60
 cout << "Determinant: " << det << endl;</pre>
 // expected: -0.233333 0.166667 0.133333 0.0666667
 //
            0.166667 0.166667 0.333333 -0.333333
 11
            0.233333 0.833333 -0.133333 -0.0666667
 //
           0.05 -0.75 -0.1 0.2
 cout << "Inverse: " << endl;</pre>
 for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++)
     cout << a[i][j] << ' ';
   cout << endl;
 // expected: 1.63333 1.3
           -0.166667 0.5
 //
            2.36667 1.7
 //
 //
           -1.85 -1.35
 cout << "Solution: " << endl;</pre>
 for (int i = 0; i < n; i++) {
   for (int j = 0; j < m; j++)
     cout << b[i][j] << ' ';
   cout << endl;</pre>
 }
```

Reduced Row Echelon Form

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
//
// Running time: O(n^3)
//
// INPUT: a[][] = an nxm matrix
//
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
// returns rank of a[][]
```

```
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
 int n = a.size();
 int m = a[0].size();
 int r = 0;
 for (int c = 0; c < m && r < n; c++) {
    int j = r;
    for (int i = r+1; i < n; i++)
      if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
    if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[j], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s;
    for (int i = 0; i < n; i++) if (i != r) {
     T t = a[i][c];
      for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];
    r++;
 return r;
int main(){
 const int n = 5;
 const int m = 4;
 double A[n][m] = {
{16,2,3,13},{5,11,10,8},{9,7,6,12},{4,14,15,1},{13,21,21,13} };
 VVT a(n);
 for (int i = 0; i < n; i++)
   a[i] = VT(A[i], A[i] + n);
  int rank = rref (a);
 // expected: 4
  cout << "Rank: " << rank << endl;</pre>
  // expected: 1 0 0 1
 //
            0 1 0 3
 //
            0 0 1 -3
 //
            0 0 0 2.78206e-15
           0 0 0 3.22398e-15
  cout << "rref: " << endl;</pre>
  for (int i = 0; i < 5; i++) {
```

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```
for (int j = 0; j < 4; j++)
    cout << a[i][j] << ' ';
    cout << endl;
}</pre>
```

Data Structures

K-d Tree

```
// A straightforward, but probably sub-optimal KD-tree implmentation that's
// probably good enough for most things (current it's a 2D-tree)
// - constructs from n points in O(n lg^2 n) time
// - handles nearest-neighbor query in O(lg n) if points are well
// - worst case for nearest-neighbor may be linear in pathological case
// Sonny Chan, Stanford University, April 2009
#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>
using namespace std;
// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point {
ntype x, y;
point(ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
1:
bool operator == (const point &a, const point &b) {return a.x == b.x && a.y ==
b.v;}
// sorts points on x-coordinate
bool on x(const point &a, const point &b) {return a.x < b.x;}
// sorts points on v-coordinate
bool on y(const point &a, const point &b) {return a.y < b.y;}
// squared distance between points
ntype pdist2(const point &a, const point &b) {
ntype dx = a.x-b.x, dy = a.y-b.y;
```

```
return dx*dx + dv*dv;
// bounding box for a set of points
struct bbox {
 ntype x0, x1, y0, y1;
 bbox() : x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}
 // computes bounding box from a bunch of points
 void compute(const vector<point> &v) {
    for (int i = 0; i < v.size(); ++i) {
     x0 = min(x0, v[i].x); x1 = max(x1, v[i].x);
      v0 = min(y0, v[i].y); y1 = max(y1, v[i].y);
  // squared distance between a point and this bbox, 0 if inside
  ntype distance(const point &p) {
    if (p.x < x0) {
      if (p.v < v0)
                         return pdist2(point(x0, y0), p);
      else if (p.y > y1) return pdist2(point(x0, y1), p);
                         return pdist2(point(x0, p.y), p);
      else
    else if (p.x > x1) {
                         return pdist2(point(x1, y0), p);
      if (p.v < v0)
      else if (p.y > y1) return pdist2(point(x1, y1), p);
                         return pdist2(point(x1, p.y), p);
    else {
      if (p.y < y0)
                         return pdist2(point(p.x, y0), p);
      else if (p.y > y1) return pdist2(point(p.x, y1), p);
                         return 0:
1:
// stores a single node of the kd-tree, either internal or leaf
struct kdnode {
 bool leaf; // true if this is a leaf node (has one point)
 point pt;
               // the single point of this is a leaf
 bbox bound; // bounding box for set of points in children
  kdnode *first. *second: // two children of this kd-node
  kdnode() : leaf(false), first(0), second(0) {}
  ~kdnode() { if (first) delete first; if (second) delete second; }
  // intersect a point with this node (returns squared distance)
 ntype intersect(const point &p) {
   return bound.distance(p);
  // recursively builds a kd-tree from a given cloud of points
 void construct(vector<point> &vp) {
   // compute bounding box for points at this node
```

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```
bound.compute(vp);
    // if we're down to one point, then we're a leaf node
   if (vp.size() == 1) {
     leaf = true;
     pt = vp[0];
   }
   else {
     // split on x if the bbox is wider than high (not best heuristic...)
     if (bound.x1-bound.x0 >= bound.y1-bound.y0)
       sort(vp.begin(), vp.end(), on x);
     // otherwise split on v-coordinate
     else
       sort(vp.begin(), vp.end(), on y);
     // divide by taking half the array for each child
     // (not best performance if many duplicates in the middle)
     int half = vp.size()/2;
     vector<point> vl(vp.begin(), vp.begin()+half);
     vector<point> vr(vp.begin()+half, vp.end());
      first = new kdnode(); first->construct(vl);
      second = new kdnode(); second->construct(vr);
 - }
1:
// simple kd-tree class to hold the tree and handle queries
struct kdtree (
 kdnode *root;
 // constructs a kd-tree from a points (copied here, as it sorts them)
 kdtree(const vector<point> &vp) {
   vector<point> v(vp.begin(), vp.end());
   root = new kdnode();
   root->construct(v);
 ~kdtree() { delete root; }
 // recursive search method returns squared distance to nearest point
 ntype search (kdnode *node, const point &p)
   if (node->leaf) {
     // commented special case tells a point not to find itself
     // if (p == node->pt) return sentry;
     // else
     return pdist2(p, node->pt);
   ntype bfirst = node->first->intersect(p);
   ntype bsecond = node->second->intersect(p);
   // choose the side with the closest bounding box to search first
   // (note that the other side is also searched if needed)
   if (bfirst < bsecond) {</pre>
     ntype best = search(node->first, p);
     if (bsecond < best)</pre>
```

```
best = min(best, search(node->second, p));
      return best;
    else {
      ntype best = search(node->second, p);
      if (bfirst < best)</pre>
        best = min(best, search(node->first, p));
      return best:
  // squared distance to the nearest
  ntype nearest(const point &p) {
    return search(root, p);
1:
int main() {
  // generate some random points for a kd-tree
 vector<point> vp;
  for (int i = 0; i < 100000; ++i) {
    vp.push back(point(rand()%100000, rand()%100000));
  kdtree tree (vp);
  // query some points
  for (int i = 0; i < 10; ++i) {
    point q(rand()%100000, rand()%100000);
    cout << "Closest squared distance to (" << q.x << ", " << q.y << ")"</pre>
    << " is " << tree.nearest(g) << endl;</pre>
  return 0;
```

Fenwick Tree

```
/** Fenwick Tree with Range Update **/
#define MAXN 100005

int n, bitMul[MAXN], bitAdd[MAXN];

void internalUpdate(int k, int mul, int add) {
   for (int x = k; x <= n; x += (x & -x)) {
     bitMul[x] += mul;
     bitAdd[x] += add;
   }
}

void update(int l, int r, int value) {
   internalUpdate(l, value, -value * (l - 1));
   internalUpdate(r, -value, value * r);
}</pre>
```

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```
int query(int k) {
  int mul = 0, add = 0;
  for (int x = k; x > 0; x -= (x & -x)) {
    mul += bitMul[x];
    add += bitAdd[x];
  }
  return mul * k + add;
}
```

Splay Tree

```
#include <cstdio>
#include <algorithm>
using namespace std;
const int N MAX = 130010;
const int oo = 0x3f3f3f3f3f;
struct Node {
 Node *ch[2], *pre;
int val, size;
 bool isTurned;
} nodePool[N MAX], *null, *root;
Node *allocNode(int val) {
 static int freePos = 0;
 Node *x = & nodePool[freePos ++];
 x->val = val, x->isTurned = false;
 x \rightarrow ch[0] = x \rightarrow ch[1] = x \rightarrow pre = null;
 x->size = 1;
 return x;
inline void update(Node *x) {
 x-size = x-ch[0]-size + x-ch[1]-size + 1;
inline void makeTurned(Node *x) {
if(x == null)
   return;
 swap (x->ch[0], x->ch[1]);
 x->isTurned ^= 1;
inline void pushDown(Node *x) {
if(x->isTurned) {
   makeTurned(x->ch[0]);
   makeTurned(x->ch[1]);
   x->isTurned ^= 1;
 }
inline void rotate(Node *x, int c) {
 Node *y = x-pre;
```

```
x->pre = y->pre;
  if(v->pre != null)
    y \rightarrow pre \rightarrow ch[y == y \rightarrow pre \rightarrow ch[1]] = x;
  y - ch[!c] = x - ch[c];
  if(x->ch[c] != null)
    x->ch[c]->pre = y;
  x->ch[c] = y, y->pre = x;
  update(y);
  if(y == root)
    root = x;
void splay(Node *x, Node *p) {
  while(x->pre != p) {
    if(x-pre-pre == p)
       rotate(x, x == x-pre-ch[0]);
    else {
      Node *y = x \rightarrow pre, *z = y \rightarrow pre;
       if(y == z->ch[0]) {
        if(x == y->ch[0])
           rotate(y, 1), rotate(x, 1);
           rotate(x, 0), rotate(x, 1);
      } else {
        if(x == y->ch[1])
           rotate(y, 0), rotate(x, 0);
           rotate(x, 1), rotate(x, 0);
    }
  update(x);
void select(int k, Node *fa) {
  Node *now = root;
  while(1) {
    pushDown (now);
    int tmp = now->ch[0]->size + 1;
    if (tmp == k)
      break;
    else if(tmp < k)</pre>
      now = now->ch[1], k -= tmp;
       now = now->ch[0];
  splay(now, fa);
Node *makeTree(Node *p, int 1, int r) {
 if(1 > r)
    return null:
  int mid = (l + r) / 2;
```

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```
Node *x = allocNode(mid):
 x->pre = p;
 x \rightarrow ch[0] = makeTree(x, 1, mid - 1);
 x \rightarrow ch[1] = makeTree(x, mid + 1, r);
 update(x);
 return x;
int main() {
 int n, m;
 null = allocNode(0);
 null->size = 0;
 root = allocNode(0);
 root->ch[1] = allocNode(oo);
 root->ch[1]->pre = root;
 update (root);
 scanf("%d%d", &n, &m);
 root->ch[1]->ch[0] = makeTree(root->ch[1], 1, n);
 splay(root->ch[1]->ch[0], null);
 while (m --) {
   int a, b;
   scanf("%d%d", &a, &b);
   a ++, b ++;
   select(a - 1, null);
   select(b + 1, root);
   makeTurned(root->ch[1]->ch[0]);
 for(int i = 1; i <= n; i ++) {</pre>
   select(i + 1, null);
   printf("%d ", root->val);
 }
```

DP Convex Hull Optimization

```
public class ConvexHullOptimization {
  long[] A = new long[1000000];
  long[] B = new long[1000000];
  int len;
  int ptr;

  // a descends
  public void addLine(long a, long b) {
    // intersection of (A[len-2],B[len-2]) with (A[len-1],B[len-1]) must
  lie to the left of intersection of (A[len-1],B[len-1]) with (a,b)
    while (len >= 2 && (B[len - 2] - B[len - 1]) * (a - A[len - 1]) >=
  (B[len - 1] - b) * (A[len - 1] - A[len - 2])) {
        --len;
    }
}
```

```
A[len] = a;
    B[len] = b;
    ++len;
 // x ascends
 public long minValue(long x) {
   ptr = Math.min(ptr, len - 1);
    while (ptr + 1 < len && A[ptr + 1] * x + B[ptr + 1] <= A[ptr] * x +
B[ptr]) {
      ++ptr;
    return A[ptr] * x + B[ptr];
 // Usage example
 public static void main(String[] args) {
    ConvexHullOptimization h = new ConvexHullOptimization();
    h.addLine(3, 0);
    h.addLine(2, 1);
    h.addLine(3, 2);
   h.addLine(0, 6);
    System.out.println(h.minValue(0));
    System.out.println(h.minValue(1));
    System.out.println(h.minValue(2));
    System.out.println(h.minValue(3));
```

Geometry

Point, Segment, Line, Circle

```
double acos(double x) {
  double ret = acos(x);
  if (ret == ret) return ret;
  if (x < 0) return acos(-1.0);
  return acos(1.0);
#define acos acos
\#define sqr(x) ((x)*(x))
const double PI = acos(-1);
const double EPS = 1e-9;
const double INF = 1e300;
struct point{
  double x, y;
  point() { x = y = 0; }
  point(double x, double y) : x(x), y(y) {}
};
struct segment {
```

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```
point p1, p2;
 segment() \{p1 = p2 = point(0,0);\}
 segment(point p1, point p2) : p1(p1), p2(p2) {}
/** basic operators and functions of point and segment **/
/* complexity: constant */
double cross(const point &p1, const point &p2) {
/* returns z-component of cross product of two points (vectors) */
 return p1.x * p2.y - p1.y * p2.x;
double dot(const point &p1, const point &p2) {
 /* returns dot product of two points (vectors) */
return p1.x * p2.x + p1.y * p2.y;
double getAngle(const point &p1, const point &p2) {
 /* returns angle formed by two vectors. WARNING: undirected angle */
 return fabs(acos(dot(p1,p2) / dist(p1,point(0,0)) /
dist(p2,point(0,0)));
double getAngle(const point &p1, const point &center, const point &p2) {
 /* returns angle formed by three points. WARNING: undirected angle */
 return getAngle(p1 - center, p2 - center);
double distToSegment(const point &p, const segment &s) {
 /* returns distance of a point to a segment */
 if (getAngle(s.p2, s.p1, p) > PI/2 + EPS || getAngle(s.p1, s.p2, p) >
PI/2 + EPS) return min(dist(p,s.p1), dist(p,s.p2));
 return fabs(cross(s.p1 - p, s.p2 - p)) / dist(s.p1, s.p2);
double distToLine(const point &p, const segment &s){
 /* returns distance of a point to a line (its orthogonal projection) */
 return fabs(cross(s.p1 - p, s.p2 - p)) / dist(s.p1, s.p2);
point rotate (const point &p, const double &alpha) {
 /* rotates a point with respect to the origin. alpha in radians */
 return point (p.x * cos(alpha) - p.y * sin(alpha), p.x * sin(alpha) + p.y
* cos(alpha));
point rotate (const point &p, const point &center, const double &alpha) {
 /* rotates a point with respect to point center. alpha in radians */
 return center + rotate(p - center, alpha);
point rescale (const point &p, const double s) {
 return point(p.x * s, p.y * s);
point dilate (const point &p, const double Factor) {
 return rescale(p, Factor);
point dilate (const point &p, const point &center, double factor) {
 return dilate(p- center, factor) + center;
bool isRightTurn(const point &p1, const point &p2, const point &p3) {
```

```
return cross(p2 - p1, p3 - p2) <= 0;
  /* straight returns true */
bool isOnSameSide (const point &p1, const point &p2, const segment &s) {
 double z1 = cross(s.p2 - s.p1, p1 - s.p1);
  double z2 = cross(s.p2 - s.p1, p2 - s.p1);
  return (z1 + EPS < 0 && z2 + EPS < 0) || (0 < z1 - EPS && 0 < z2 - EPS)
|| fabs(z1) < EPS || fabs(z2) < EPS;
 /* on segment returns true */
bool isOnLine (const point &p, const segment &1) {
 return fabs((1.p1.y - p.y) * (1.p2.x - p.x) - (1.p2.y - p.y) * (1.p1.x -
p.x)) < EPS;
bool isOnSegment(const point &p, const segment &s){
  return fabs(dist(p, s.p1) + dist(p, s.p2) - dist(s.p1, s.p2)) < EPS;
bool isIntersecting(const segment &s1, const segment &s2){
  return !(isOnSameSide(s1.p1,s1.p2,s2) || isOnSameSide(s2.p1,s2.p2,s1)) ||
isOnSegment(s1.p1,s2) || isOnSegment(s1.p2,s2) || isOnSegment(s2.p1,s1) ||
isOnSegment(s2.p2,s1);
bool isParallel(const segment &s1, const segment &s2) {
 return fabs((s1.p1.y-s1.p2.y)*(s2.p1.x-s2.p2.x)-(s2.p1.y-
s2.p2.v)*(s1.p1.x-s1.p2.x)) < EPS;
point intersection (const segment &s1, const segment &s2) {
  /* assumes !isParallel(s1,s2) */
  double x1 = s1.p1.x - s1.p2.x;
  double x2 = s2.p1.x - s2.p2.x;
  double y1 = s1.p1.y - s1.p2.y;
  double y2 = s2.p1.y - s2.p2.y;
  double cross1 = cross(s1.p1, s1.p2);
  double cross2 = cross(s2.p1, s2.p2);
  return point ((cross1 * x2 - cross2 * x1) / (x1 * y2 - x2 * y1), (cross1
* y2 - cross2 * y1) / (x1 * y2 - x2 * y1));
point projection (const point &p, const segment &s) {
  /* projects p onto line s */
  return rescale(s.p2 - s.p1, dot(p - s.p1, s.p2 - s.p1) / <math>sqr(length(s)))
+ s.p1;
/** introducing circle **/
struct circle {
  point center:
  double r:
  circle() { center = point(0, 0); r = 0; }
 circle(point p, double r) : center(p), r(r) {}
vector<point> intersectionLineCircle(const segment &1, const circle &c){
  vector<point> res;
  double dx = 1.p2.x - 1.p1.x;
```

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```
double dy = 1.p2.y - 1.p1.y;
 double dr = length(1);
 double d = cross(1.p1 - c.center,1.p2 - c.center);
 if (sqr(c.r) * sqr(dr) - sqr(d) + EPS < 0) return res;</pre>
 double det = sqrt(fabs(sqr(c.r) * sqr(dr) - sqr(d)));
 double sdx = dy < 0 ? -dx : dx;
 double sdv = fabs(dv);
 res.push back(c.center + point((d*dy + sdx * det)/sqr(dr), (-d*dx + sdy *
det)/sqr(dr)));
 if (det > EPS) res.push back(c.center + point((d*dy - sdx * det)/sqr(dr),
(-d*dx - sdy * det)/sqr(dr)));
 return res;
vector<point> intersectionSegmentCircle(const segment &s, const circle &c){
 vector<point> res, res = intersectionLineCircle(s,c);
 for (vector<point>::iterator it = res.begin(); it != res.end(); ++it){
   if (isOnSegment(*it,s)) res.push back(*it);
 return res;
```

Polygons (Area, Orientation)

```
/** introducing polygon **/
typedef vector<point> polygon;
/** Check position of a point with respect to a polygon **/
/* complexity : O(N) */
bool isPointInsidePolygon(point p, polygon poly){
 /* ray casting to the right */
 segment ray (p,p+point(1,0));
 int n = (int)poly.size();
 /* counts the number of intersections */
 int nIntersection = 0;
 for (int i = 0; i < n; ++i){
   segment side(poly[i],poly[(i+1)%n]);
   if (isOnSegment(p,side)) return false;
   if (isParallel(ray,side)) continue;
   point x = intersection(ray, side);
   if (isOnSegment(x,side) && dot(x-p,ray.p2-p) > 0) {
      /* special case: x is one of vertices of sides */
     if (x == side.p1){
       if (isRightTurn(p,x,side.p2)) nIntersection ++;
      else if (x == side.p2){
       if (isRightTurn(p,x,side.pl)) nIntersection ++;
      else nIntersection ++:
 return nIntersection % 2 == 1;
```

Convex Hull

```
/** Convex Hull | monotone chain algorithm **/
/* complexity : O(N log N) */
polygon convexHull(polygon p){
   int m = 0, n = p.size();
   polygon hull(2*n);
   sort(p.begin(),p.end());
   for (int i = 0; i < n; ++i){
      while (m >= 2 && isRightTurn(hull[m-2],hull[m-1],p[i])) --m;
      hull[m++] = p[i];
   }
   for (int i = n-1, t = m+1; i >= 0; --i){
      while (m >= t && isRightTurn(hull[m-2],hull[m-1],p[i])) --m;
      hull[m++] = p[i];
   }
   hull.resize(m);
   return hull;
}
```

Dealunay Triangulation

```
// Slow but simple Delaunay triangulation. Does not handle
// degenerate cases (from O'Rourke, Computational Geometry in C)
// Running time: O(n^4)
// INPUT: x[] = x-coordinates
11
       v[] = v-coordinates
// OUTPUT: triples = a vector containing m triples of indices
                    corresponding to triangle vertices
#include <vector>
using namespace std;
typedef double T:
struct triple {
 int i, j, k;
 triple() {}
 triple(int i, int j, int k) : i(i), j(j), k(k) {}
vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y) {
 int n = x.size():
 vector<T> z(n);
 vector<triple> ret;
  for (int i = 0; i < n; i++)
   z[i] = x[i] * x[i] + v[i] * v[i];
  for (int i = 0; i < n-2; i++) {
```

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```
for (int j = i+1; j < n; j++) {
     for (int k = i+1; k < n; k++) {
       if (j == k) continue;
       double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])*(z[j]-z[i]);
       double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])*(z[k]-z[i]);
       double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i])*(y[j]-y[i]);
       bool flag = zn < 0;
       for (int m = 0; flag && m < n; m++)</pre>
         flag = flag && ((x[m]-x[i])*xn +
            (y[m]-y[i])*yn +
            (z[m]-z[i])*zn <= 0);
       if (flag) ret.push back(triple(i, j, k));
 return ret;
int main() {
 T \times S[] = \{0, 0, 1, 0.9\};
 T ys[]=\{0, 1, 0, 0.9\};
 vectorT x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
 vector<triple> tri = delaunayTriangulation(x, y);
 //expected: 0 1 3
           0 3 2
 int i;
 for(i = 0; i < tri.size(); i++)</pre>
   printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
 return 0;
```

Miscellaneous

Graph Theorems

Erdos-Gallai. A sequence of nonnegative integers $d_1 \ge \cdots \ge d_n$ is a sequence of degree of an undirected graph iff $\sum d_i$ is even and $\sum_{i=1}^k d_i \le k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$

Fulkerson-Chen-Anstee. A sequence $\left((a_1,b_1),\ldots,(a_n,b_n)\right)$ of nonnegative integer pairs with $a_1\geq \cdots \geq a_n$ is a sequence of (in, outdeg) of a directed graph iff $\sum a_i=\sum b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k)$

Lindstrom-Gessel-Viennot. The number of non-intersecting path from A to B in a directed acyclic graph is equal to the determinant of ... (elements (i,j) of matrix denotes the number of ways to go from A_i to B_j).

Koenig's. In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover.

Brook's. For any connected undirected graph G with maximum degree Δ , the chromatic number of G is at most Δ unless G is a complete graph or an odd cycle, in which case the chromatic number is $\Delta + 1$.

Combinatorics

Lucas Theorem. $\binom{n}{m} = \prod_{i=0}^k \binom{n_i}{m_i} \pmod{p}$ where $n = \overline{n_k n_{k-1} \dots n_0}$ and $m = \overline{m_k m_{k-1} \dots m_0}$ in base p.

Stirling Number of the First Kind. s(n, k) denotes the number of n-permutation with k cycles. s(n + 1, k) = ns(n, k) + s(n, k - 1).

Stirling Number of the Second Kind. S(n,k) denotes the number of partition a set of n into k non-empty subsets. S(n+1,k)=kS(n,k)+S(n,k-1). $S(n,k)=\frac{1}{k!}\sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^n$.

Gambler's Ruin. Two players with n_1 and n_2 points each are playing, each turn P1 has probability of winning p and P2 has probability q=1-p. The probability of P1 losing all his points is $\left(1-\left(\frac{p}{q}\right)^{n_2}\right)/\left(1-\left(\frac{p}{q}\right)^{n_1+n_2}\right)$.

Notes

std::lower_bound. Returns an iterator pointing to the first element in the range [first,last) which **does not compare less than** val.

std::upper_bound. Returns an iterator pointing to the first element in the range [first,last) which **compares greater** than val.