Grafuri (Graphs)

Reprezentări

```
Matrice de adiacență: /də:ə/
Lista de adiacență: wester/Nede> li
```

• Liste de adiacență: vector<Node> list[NodeCount];

• Lista muchiilor/arcelor: vector<pair<Node, Node>> list;

Grafuri neorientate (Undirected Graphs)

```
Noduri (Vertices); Muchii (Edges)
Graf complete (Complete Graph): toate muchiile adiacente
Graf partial (partial?): stergem muchii
Subgraf (Subgraph): ştergem noduri
Lant (Walk / Chain): adiacente 2 cate 2
Lant elementar (Trail): lant cu noduri distincte
Ciclu (Cycle): lant cu primu' = ultimu'
Ciclu elementar (elementary?): You get the point
Graf aciclic (Acyclic): natürlich
Conex (Connected): \forall v, w \exists lant
Componenta conexă (Connected component): subgraf maximal (aka tătî bucata)
Lanț eulerian (Eulerian Path/ just 'Path'): lanț care vizitează fiecare muchie o singura dată
Ciclu eulerian (Eulerian cycle/circuit): możesz to rozgryźć
Graf eulerian (Eulerian Graph): Graf cu ciclu eulerian
Bipartit (Bipartite): putem împarti nodurile în 2 partiții, în care 2 noduri din aceeași partiție nu sunt
    adiacente
Lant hamiltonian (Hamiltonian path): lant elementar cu toate nodurile
Ciclu hamiltonian (Hamiltonian cycle): duh
```

Grafuri orientate (Directed Graphs)

```
Vârfuri (Vertices); Arce (Edges/arc)
Grad interior: câte intră
Grad exterior: câte ies
Drum (Directed Walk): "lanț"
Graph Turneu (Tour): intre oricare 2 varfuri distincte exista un arc
Slab Conex (Weakly connected): daca-l transformam în neorientat, e conex
Tare Conex (Strongly connected): există drum intre oricare 2 varfuri, dus-intors
Modalitati de parcurgere: Breadth first, Depth first
```

Grafuri ponderate (Weighted Graphs)

Dijkstra

```
function Dijkstra(Graph, source):
    create vertex set Q
    for each vertex v in Graph:
        dist[v] = INFINITY
        prev[v] = UNDEFINED
        add v to Q
    dist[source] = 0
```

```
8
        while Q is not empty:
9
             u = vertex in Q with min dist[u]
10
             remove u from Q
11
             for each neighbor v of u:
                                                     # only v that are still in Q
12
                 alt = dist[u] + length(u, v)
13
                 if alt < dist[v]:
14
                     dist[v] = alt
15
                     prev[v] = u
16
17
        return dist[], prev[]
18
```

Roy-Floyd

```
let dist be a |V| \times |V| array of minimum distances initialized to \infty
    let next be a |V| \times |V| array of vertex indices initialized to null
3
    procedure FloydWarshallWithPathReconstruction():
        for each edge (u, v) do
5
            dist[u][v] = w(u, v) # The weight of the edge (u, v)
6
            next[u][v] = v
        for each vertex v do
            dist[v][v] = 0
9
            next[v][v] = v
10
        for k from 1 to |V| do # standard Floyd-Warshall implementation
11
            for i from 1 to |V|
12
                 for j from 1 to |V|
13
                     if dist[i][j] > dist[i][k] + dist[k][j] then
                         dist[i][j] = dist[i][k] + dist[k][j]
15
                         next[i][j] = next[i][k]
16
17
    procedure Path(u, v)
18
        if next[u][v] = null then
19
            return []
20
        path = [u]
21
        while u \neq v
22
            u = next[u][v]
23
            path.append(u)
24
        return path
25
```

Arbori (Trees)

Subarbori (Subtree): fii

Nivelul unui nod (Level): distanta pana la radacina

Inaltimea unui nod (Height): cel mai lung lant pana la un nod terminal

Arbore partial al unui graf: graf partial care este arbore

Arbore partial de cost minim (minimum-spanning-tree): avem o functie f care determina costul fiecarei muchii, gasim un arbore partial aî suma costurilor sale sa fie minima.

Reprezentări

```
like a graph: /də:ə/
Ref descendente: struct Node { int data; Vector<Node*> children; };
```

• Ref ascendente: int parinte[NodeCount] = {-1, ...};

Kruskal

```
algorithm Kruskal(G) is
1
   A := \emptyset
2
   for each v \in G.V do
3
        MAKE-SET(v)
4
   for each (u, v) in G.E ordered by weight(u, v), increasing do
5
        if FIND-SET(u) \neq FIND-SET(v) then
6
           A := A \cup \{(u, v)\}
           UNION(FIND-SET(u), FIND-SET(v))
   return A
9
```

Prim

```
void prim(double mat[sz][sz], ssize_t len) {
1
        int* s = new int[len] {-1}, j;
2
        for (int i = 1; i < len; ++i) s[i] = 0;
3
        for (int k = 1; k < len; ++k) {
            double min = inf;
5
            for (int i = 0; i < len; ++i)
6
                if (s[i] != -1 \&\& min > mat[i][s[i]]) {
                     min = mat[i][s[i]];
                     j = i;
9
10
            std::cout << "edge: " << j << "-" << s[j] << "\n";
11
            for (int i = 0; i < len; ++i)
12
                if (s[i] != -1 && mat[i][s[i]] > mat[i][j])
13
                     s[i] = j;
14
            s[j] = -1;
15
        }
16
    }
```

Arbori binari (Binary Tree)

Arbori binari completi (Perfect binary tree): pe fiecare nivel s are exact 2^s noduri

Aproape complet (Complete binary tree): pe ultimul nivel lipsesc doar primele din stanga/ ultimele din dreapta

Ansamblul Heap binar minim (Min-heap): Arbore aproape complet cu cheia oricarui parinte mai mica sau egala cu a fiului

Heap sort: stergem si punem intr-un vector

Coada de priorități (Priority queue): tl;dr a fancy heap

AVL Trees

```
struct Node {
    Key key;
    Node *left, *right;
```

```
int h = 1;
4
        static int height(const Node* p) { return (p == nullptr) ? 0 : p->h; }
5
        void updateHeight() { h = std::max(height(left), height(right)) + 1; }
6
        int balanceFactor() const { return height(left) - height(right); }
        static void rightRotate(Node*& n) {
            Node* t = n->left;
            n->left = t->right;
10
            t->right = n;
11
            n->updateHeight();
            t->updateHeight();
13
            n = t;
14
        }
        static void leftRotate(Node*& n) { /* like above but swap right and left*/ }
        static void balance(Node*& node) {
            node->updateHeight();
            int factor = node->balanceFactor();
            if (factor > 1) {
                if (node->left->balanceFactor() < 0) leftRotate(node->left);
21
                rightRotate(node);
            } else if (factor < -1) {</pre>
23
                if (node->right->balanceFactor() > 0) rightRotate(node->right);
                leftRotate(node);
25
            }
26
        }
27
        static void add(Node*& node, const Key& key) {
            if (node == nullptr) { node = new Node(key); return; }
29
            assert(key != node->key, "Value already in tree");
30
            if (key < node->key) add(node->left, key);
31
            else if (key > node->key) add(node->right, key);
32
            balance(node);
33
        }
34
35
        static void remove(Node*& n, const Key& key) {
36
            if (n == nullptr) { printf("not found?\n"); return; }
37
            if (key < n->key) remove(n->left, key);
38
            else if (key > n->key) remove(n->right, key);
39
            else { // key == n->key
40
                if (n->left == nullptr) {
41
                     if (n->right == nullptr) { n = nullptr; return; }
42
                    Node * tmp = n;
43
                    n = n->right;
44
                     tmp->right = nullptr;
45
                     delete tmp;
46
                } else if (n->right == nullptr) {
                    Node * tmp = n;
                    n = n->left;
49
                    tmp->left = nullptr;
                     delete tmp;
                } else {
                     constexpr auto findMax = [](Node*& r, auto& findMax) -> Node*& {
                             if (r->right == nullptr) return r;
54
                             return findMax(r->right, findMax);
55
```

```
};
56
                       Node*& tmp = findMax(n->left, findMax);
57
                       n->key = tmp->key;
58
                       remove(n->left, tmp->key);
59
                  }
60
             }
61
             if (n == nullptr) return;
62
             balance(n);
63
         }
64
    };
65
```

Binary heap

```
Vector<Val> data;
    constexpr static int parent(int i) { return (i - 1) / 2; }
2
    constexpr static int left(int i) { return (2 * i + 1); }
    constexpr static int right(int i) { return (2 * i + 2); }
    constexpr bool empty() const { return data.empty(); }
    void push(const Val& v) {
        data.push_back(v);
        for (size_t i = data.size()-1; ;) {
            if (i == 0) break;
            size_t j = parent(i);
10
            if (data[i] > data[j]) {
                 std::swap(data[i], data[j]);
12
                 i = j;
13
            }
        }
15
16
    Val pop() {
17
        Val res = std::move(data[0]);
18
        data[0] = data.back();
19
        data.pop_back();
20
        size_t i = 0;
21
        size_t sz = data.size();
22
        while (i < sz) {
23
            size_t j = left(i);
24
            if (j < sz) {
25
                 if (j+1 < sz \&\& data[j+1] > data[j]) ++j;
26
                 if (data[i] < data[j]) {</pre>
27
                     std::swap(data[i], data[j]);
28
                     i = j;
29
                 } else break;
30
            } else break;
31
32
        return res;
33
    }
34
```

Tablele de dispersie (Hash tables)

```
functii: h(k) = k \mod m; h(k) = \lfloor m(k\varphi - [k\varphi]) \rfloor, \varphi = \frac{\sqrt{5}-1}{2}
Or: do it again: h(k,i) = h_1(k) + i \mod m; h(k,i) = h_1(k) + c_1i + c_2i^2 \mod m
```

Binary search trees

```
void insert(Node*& root, int key, int value) {
1
      if (!root)
2
        root = new Node(key, value);
3
      else if (key == root->key)
        root->value = value;
      else if (key < root->key)
6
        insert(root->left, key, value);
      else // key > root->key
        insert(root->right, key, value);
10
    constexpr NodeRef findRef(const Key& key) {
11
        auto impl = [&key] (Node*& t, auto& impl) {
                         if (t == nullptr || key == t->key) return NodeRef { t};
13
                         if (key < t->key) return impl(t->left, impl);
                         return impl(t->right, impl);
15
                     };
        return impl(root, impl);
18
19
    constexpr void remove(NodeRef r) {
20
        assert(r.val != nullptr, "Node not found");
21
        Node*& n = r.val;
22
        if (n->left == nullptr) {
23
            if (n->right == nullptr) { n = nullptr; return; }
24
            Node* tmp = n;
25
            n = n->right;
26
            tmp->right = nullptr;
27
            delete tmp;
28
        } else if (n->right == nullptr) {
29
            Node* tmp = n;
30
            n = n->left;
31
            tmp->left = nullptr;
32
            delete tmp;
33
        } else {
34
            constexpr auto findMax = [] (Node*& r, auto& findMax) -> Node*& {
35
                                 if (r->right == nullptr) return r;
36
                                 return findMax(r->right, findMax);
37
38
            Node*& replacement = findMax(n->left, findMax);
39
            Node* tmp = replacement;
40
            std::swap(replacement->key, n->key);
41
            std::swap(replacement->val, n->val);
42
            replacement = replacement->left;
43
            tmp->right = tmp->left = nullptr;
44
            delete tmp;
        }
46
47
    constexpr void remove(const Key& key) { remove(findRef(key)); }
```