The First, let's show that given $\delta \in (0,1)$ and given Θ_{δ} $0 < \varepsilon_1 \le \varepsilon_2 < 1$, we have that $M_{\mathcal{H}}(\varepsilon_1, \delta) \ge M_{\mathcal{H}}(\varepsilon_2, \delta)$. Since \mathcal{H} is PAC-learnable with sample complexity $M_{\mathcal{H}}(v_1)$, we set that given a sample of size $M_{\widetilde{\rho}} = M_{\mathcal{H}}(\varepsilon_1, \delta)$, we have that with probability of at least $1-\delta$ over $S \sim D_{\mathcal{H}} = M_{\mathcal{H}}(\varepsilon_1, \delta)$, $S \sim D_{\mathcal{H}} = M_{\mathcal{H}}(\varepsilon_1, \delta)$ $S \sim D_{\mathcal{H}} = M_{\mathcal{H}}(\varepsilon_1, \delta)$ $S \sim D_{\mathcal{H}} = M_{\mathcal{H}}(\varepsilon_1, \delta)$ $S \sim D_{\mathcal{H}} = E_2$, for the sample of size $M_{\mathcal{H}} = M_{\mathcal{H}}(\varepsilon_1, \delta)$

we also have that w.p. of at least 1-5 over $SnD^{mexapper}M_1$, Lo, ϵ $(A(5)) \times E_2$. Hence, the sample size me M_2 required for Lo, ϵ $(A(5)) \epsilon$ be smaller than E_2 W.p. at least 1-5 over SnD^{m_2} is guaranteed to work, if $M_2 > M_1(E_1, \delta)$. Therefore, the sample complexity for E_2 is at most $M_1(E_1, \delta)$. Hence, we have shown that $M_1(E_1, \delta) \geq M_1(E_2, \delta)$

(2) Now, let's show that given $\xi \in (0,1)$ and given $0 < \delta_1 < \delta_2 < 1$, we have that $m_{\mathcal{H}}(\varepsilon, \delta_1) \ge m_{\mathcal{H}}(\varepsilon, \delta_2)$. Since \mathcal{H} is PAC-learnable, given any $m \ge m(\varepsilon, \delta_1)$, we have $P_{S\sim p^m}(A(s) < \varepsilon) \ge 1 - \delta_{1}$. Since $\delta_1 < \delta_2$, we have that for any $m \ge m(\varepsilon, \delta_1)$, $P_{S\sim p^m}(A(s) < \varepsilon) \ge 1 - \delta_1 \ge 1 - \delta_2$. Therefore, the sample complexity for δ_2 is at most $m_{\mathcal{H}}(\varepsilon, \delta_1)$ thence, $m_{\mathcal{H}}(\varepsilon, \delta_1) \ge m(\varepsilon, \delta_2)$

By D and (2) we set that mx is monotonically nonincreasing in each of its parameters.

2. e. d.

Exercise 2

1. Algorithm:

Go through every x in the sample and check whether it belongs to our discrete domain &XX.

If we find such x, then function hz, where z=x
is the ontcome of our algorithm. Due to realizability
assumption, there can only be one such x in the sample
and, therefore, L_s(hz) will be O.

If we don't find such x, then h is the outcome of our algorithm. Since in this case we don't have any positives in the sample, Ls(h) will be O.

2. If the algorithm described above outputs hz, then, due to the realizability assumption it will mean that our hypothesis hz correctly identifies the only positive in our domain, and hence, Lp, f (hz) with be 0, which is less than & for all & >0.

The only case in which our algorithm won't identify the correct labeling function is when the Only positive indistribution (let's call it xs) is not selected to our sample Sx and our algorithm incorrectly returns h. Since all items in the sample are i.i.d. selected, the probability of their happening is $(1-P(x))^m$, where m is the Size of our sample. We wont this probability to be \$2000 probability of at most p.

Since in this case, the population risk will 7 be equal to $P_D(x_3)$, we get the inequality for the size of the sumple m: $(1-\epsilon)^m \in \mathcal{O}$

 $m \geq \lceil \log_{(1-\epsilon)} S \rceil$

the required upper bound on the sample complexity

Exercise 3 Fix some distribution Dover X.

Assuming realizability, Addition Let R* with radius n* he the concentric circle that generates the labels and let f be the Corresponding hypothesis, Let r<r* be assessable to once radius of the concentric circle R, such that the probability mass of the area between R and R* is exactly E. (Let's denote that area R)

For perhas rar positive rexemple Allowalthe sugapher 8 Let A be the algorithm that returns the smallest contangles evolvsing all positive examples in the training set. From realizability assumption we get that A does not mislable any negative examples (if it does, there should be smaller circle enclosing all positive examples, which contradicts definition of A(SI), Hence, A is an ERM.

Since A(S) is the smallest circle enclosing all positive examples, A(S) = R*. IF R contains any positive example, then the boundary of A lies between R* and R. Since the probability mass of that region is E, empirial risk of A(5) in this case will be at most E. The probability of positive example being inside R is 1-E, hence the probability of all positive samples

the be inside R is (1-E) , where m is the Size of the sample. Therefore, the probability that at least

one positive is in R is 1-(1-E)m

We set that P(Lo, p(A(S)) < E) > 1-(1-E) = 1-ème

Rewriting this inequality, we get $m \not\in \frac{\log(1/\delta)}{\varepsilon}$.

Therefore, we have show that \exists algorithm f, s.b. \forall ε , $\delta > 0$, \forall distributions D, \forall labeling fous ε , s.t. D, ε ealizable by \mathcal{H} , given $m \geq \lceil \frac{\log(1/\delta)}{\varepsilon} \rceil$, then $W \cdot P$. And at least $1-\delta$ over $S \cap D^m$, we have L_{D}, ε (A(S)) $< \varepsilon$

tience, It is PAC-learnable and its sample complexity is bounded by $m_{\mathcal{H}}(\xi, \delta) \leq \Gamma \frac{\log(1/\sigma)}{\varepsilon}$

q. e. d.

Exercise 6 Let 21 be agnostic PAC-leurnable and let A be a successful agnostic PAC learner for H. Then, Almo consider my (0,12 -> IV, s. E. Ve, 5>0, \forall D on $\forall x Y$, given $m \ge m_{\mathcal{H}}(\mathcal{E}, \delta)$ w. p. $\ge (1-\delta)$ over $\delta = 0^m$ it is true that 1-p(A(S)) < inf Lp(h) + E Using definition (3.1 from the book) of the true error, le get that Lo(h)=P(x,v)~0 (x) 77] Therefore, given m= m= (E, S), we know that: $A = P_{s \sim 0m} \left[LP_0(A(s)) \leq \inf_{h \in \mathcal{H}} P_{(x,y) \sim 0} Lh(x) \neq y \right] + \varepsilon \right] \geq (1 - \delta)$ For the binary clussifier loss function is defined as following: Exist $\{(h,(x,s))=\begin{cases}0 & \text{if } h(x)=y\\1 & \text{if } h(x)\neq y\end{cases}$ Hence, [[(h,(+,4))] = 0 . P [h(x)=4] + 1. P [h(x) # 4] = (x,4)~b (x,4)~b $= P \left[h(x) \neq \gamma\right]$ Therefore, For our algorithm A and Function Mrc(·,·), we beared get that: $A2 \qquad P_{s \sim 0^m} \left[L_b(A(s)) \leq \inf_{h \in \mathcal{H}} \left[P(h,(x,y)) \right] + E \right] \geq (1-5)$

Hence, H is PAC-learnable and A is a successful PAC learner for H by definition.

Exercise 7 By definition 3.1 of the true risk, we get that the true risk for randomly closen classifier g is equal to: $L_{b}(g) = P L_{g(x)} \neq y J = \begin{cases} P(y=1|x), i \in g(x) = 0 \\ P(y=0|x), i \in g(x) = 1 \end{cases}$ tredictor is optimal, if it has the lowest possible value of Lo(g). As we can see from the expression above, Lo(g) va has the lowest possible vulue, when the predictor is such that Lo(g) = P(y=1/x), when $P(y=1|x) \leq P(y=o|x)$ and $L_{0}(y) = P(y=o|x)$ when $P(y=o|x) \leq P(y=l|x)$ Now, consider to (x) on (Bayes Optimal Predictor) $L_{D}(f_{O}(x)) \geq DM \left\{ P(y=1|x), i \in f_{O}(x) = 0 \right\} = \left(by \ desinition \right)$ $= \begin{cases} P(y=1|x), & \text{if } P(y=1|x) < \frac{\pi}{2} \\ P(y=0|x), & \text{if } P(y=1|x) > \frac{\pi}{2} \end{cases} = \begin{cases} P(y=1|x), & \text{if } P(y=1|x) < \frac{\pi}{2} \\ P(y=0|x), & \text{if } P(y=0|x) < \frac{\pi}{2} \end{cases}$ $= \begin{cases} P(y=1|x), & P(y=1|x) < P(y=0|x) \\ P(y=0|x), & P(y=0|x) < P(y=1|x) \end{cases}$ (since P(y=1/x)+P(y=0/x)=1) As we can see, LoCg) has the lowest possible Value Hence, & classifier g from X to Vo, 14, Lo (fo) \in Lo (g)

. e. a