Lesson 3

Local Aggregation

Capture the fort

Wholeness of the Lesson

The communication cost is the single most factor that determines the efficiency of a MapReduce job. The best way to achieve better efficiency is by carefully planning and organizing the data so that only the least amount of data is shuffled across the network.

"Capturing the fort" is the very first course of action.

LOCAL AGGREGATION

An improvement on the basic wordcount algorithm is shown next. Only mapper is modified in this example.

An associative array is introduced inside the mapper to tally up term counts within a single input-split. Instead of emitting a key-value pair for each term in the input-split, this version emits a key-value pair for each unique term in the input-split.

Given the fact that some words appear frequently within an input-split, this can yield substantial savings in the number of intermediate key-value pairs emitted.

Example

Assume that an input-split contains the word "science" 10 times.

The wordcount algorithm will emit the pair ("science", 1) ten times.

The "modified" algorithm shown next will emit only one pair: ("science", 10).

LOCAL AGGREGATION

```
class Mapper
  method initialize()
      H = new AssociativeArray()
  method map(docid a; doc d)
      for all term t in record r do
         H\{t\} = H\{t\} + 1.
  method close()
      for all term t in H do
         Emit(t, H{t})
```

LOCAL AGGREGATION

Prior to processing any record, the mapper's Initialize method is called, which initialize an associative array for holding term counts (in this example).

We can continue to accumulate partial term counts in the associative array for all records in this input-split through map method. Recall that map method is called for each record in the input-split.

Once all records in the input-split are processed, the close method is invoked and it emits all the key-value pairs.

In-mapper combining

With this technique, we are in essence incorporating the combiner functionality directly inside the mapper. There is no need to run a separate combiner,.

This design pattern in MapReduce is called, "in-mapper combining".

There are two main advantages to using this design pattern:

Advantages of the in-mapper combining pattern

First, it provides control over when local aggregation occurs and how it exactly takes place.

In contrast, the semantics of the combiner is under speciffied in MapReduce. For example, Hadoop makes no guarantees on how many times the combiner is applied, or that it is even applied at all. The combiner is provided as a semantics-preserving optimization to the execution framework, which has the option of using it, perhaps multiple times, or not at all.

Advantages of the in-mapper combining pattern

Second, in-mapper combining will typically be more efficient than using actual combiners.

One reason is the additional overhead associated with actually creating the key-value pairs. Combiners reduce the amount of intermediate data that is shuffled across the network, but don't actually reduce the number of key-value pairs that are emitted by the mappers in the first place.

LOCAL AGGREGATION

Note that there is a reduction in the number of intermediate key-value pairs that need to be shuffled across the network.

As far as the key-value pairs created, it changes from the number of terms in the input-split to the number of unique terms in the input-split.

Two important design principles to follow:

- 1. One key, One emit principle. A Mapper object will emit a key-value pair with a specific "key" only once.
- 2. Order 1 space complexity principle. The size of the "value" is independent of the number of records in the input-split.

First, it breaks the functional programming underpinnings of MapReduce, since state is being preserved across multiple input key-value pairs (records). Preserving state across multiple input instances means that algorithmic behavior may depend on the order in which input key-value pairs are encountered. This creates the potential for orderingdependent bugs, which are difficult to debug on large datasets in the general case (although the correctness of in-mapper combining for word count is easy to demonstrate).

Second, there is a fundamental scalability bottleneck associated with the in-mapper combining pattern. It critically depends on having sufficient memory to store intermediate results until the mapper has completely processed all key-value pairs in an input split.

One common solution to limiting memory usage when using the in-mapper combining technique is to "block" input key-value pairs and "flush" in-memory data structures periodically. Instead of emitting intermediate data only after every key-value pair has been processed, emit partial results after processing every n key-value pairs. This is straightforwardly implemented with a counter variable that keeps track of the number of input key-value pairs that have been processed.

As an alternative, the mapper could keep track of its own memory footprint and flush intermediate key-value pairs once memory usage has crossed a certain threshold. In both approaches, either the block size or the memory usage threshold needs to be determined empirically: with too large a value, the mapper may run out of memory, but with too small a value, opportunities for local aggregation may be lost.

Main Point 1

The local aggregation is accomplished through inmapper combining technique. In our daily life, we first organize and combine various ideas before talking to others in a spontaneous way. The most fundamental combiner is the unification of the Self with itself, which gives rise to knower, process of knowing, and known.

Average Problem

Consider a simple example: we have a large dataset where input keys are strings and input values are integers, and we wish to compute the mean of all integers associated with the same key. A real-world example might be a large user log from a popular website, where keys represent user ids and values represent some measure of activity such as elapsed time for a particular session - the task would correspond to computing the mean session length on a per-user basis, which would be useful understanding user demographics.

Class Activity

Write a MapReduce algorithm to solve the average problem

Assume that record contains (..., userID,..., time, ...)
There are access methods available for all "fields" in the record.
Thus you can assume the following:

```
u = r.getUserId()
t = r.getTime()
```

Average Problem

```
class Mapper
  method map(docid a, doc d)
      //Note: each record r has one user Id and
      // one time. So no loop is needed.
      u = r.getUserId()
      t = r.getTime()
      Emit(u, t)
```

Average Problem

```
class Reducer
  method Reduce(string u, integers [t1, t2, ...])
    sum = 0; cnt = 0;
    for all integer t in integers [t1, t2, ...] do
       sum = sum + t; cnt = cnt + 1;
    avg = sum / cnt
    Emit(u, avg)
```

In class exercise.

Create the in-mapper combining algorithm to solve the average problem.

If each mapper emits the average, the algorithm will produce the wrong result

```
(Mapper Input)
(u1, 1)
                   (u1, 4)
(u1, 2)
                   (u1, 5)
(u1, 3)
(Mapper output)
                   (u1, 4.5)
(u1, 2)
(Reducer output)
(u1, (2+4.5)/2)
The correct answer is (u1, 3).
```

```
(Mapper Input)
(u1, 1)
                  (u1, 4)
                  (u1, 5)
(u1, 2)
(u1, 3)
Mapper emits the pair (sum, count) as value.
(u1, (6, 3)) (u1, (9, 2))
Reducer emits
(u1, (6+9)/5)
That is (u1, 3). That is correct!
```

O(1) Space Complexity

The size of the "value" in the in-mapper Emit should not be affected by the number of values associated with the key in the Input-split.

For example, Average problem.

u1 10

u1 40

u1 60

u1 20

Emit is (u1, (130, 4))

O(1) Space Complexity

Now consider

u1 100

u1 80

Emit is (u1, (180, 2))

Note that size of (130, 4) is the same as (180, 2) (both are objects with two integers). Hence is a good design. Note that in the first case, there were FOUR values associated with u1. In the second case, there were only two values associated with u1.

However, the size remained the same. That is O(1) space complexity.

O(1) Space Complexity

On the other hand, if the in-Mapper has emitted (u1, List(10, 40, 60, 20)) and (u1, List(100, 80)) you see the size of "value" in the in-mapper combining Emit depends on the number of values associated with the key u1. In this example, first 4. Next 2. Thus the space complexity of the value in the in-Mapper combining version is O(n) or linear.

That is a bad design. We do not want that.

In the case of wordcount, the operation was addition. It is both associative and commutative.

The "average" is not associative.

You can modify an algorithm to "in-mapper combining" version without any change in the "emit" statements if the operations are associative and commutative.

If the operation is associative and commutative, you can use the same reducer since the mapper class's Emit statement's signature will not change.

Example: wordcount and In-Mapper Combining wordcount emit statements have the same signature. Here the operation on values is +.

We know a + b = b + a (commutative) and a + (b + c) = (a+b) + c (associative)

If the operation is not associative or not commutative, you can not use the same reducer since the mapper class's Emit statement's signature has changed.

Example: Average problem.

Signature of Emit in the Mapper of average problem : (key, value)

Signature of Emit in the In-Mapper Combiner version is: (key, Pair of values). We are forced to sent both sum and count as a pair. Otherwise program will not be correct. Thus mapper class's Emit statement's signature has changed.

In-Mapper Combining version of the average problem

```
class Mapper
  method initialize()
  H = new AssociativeArray()
```

```
method map(docid a, doc d)
      //Note: each record r has one user Id and
      // one time. So no loop is needed.
      u = r.getUserId()
      t = r.getTime()
      p = new Pair(t, 1)
      if (H{u} is null)
         H\{u\} = p
      else
         H\{u\} = H\{u\} + p //elementwise addition
```

```
method close()

for all term t in H do

Emit(t, H{t})
```

Average Problem

```
class Reducer
  method Reduce(string u, pair [(t1, c1), (t2, c2), ...])
  sum = 0; cnt = 0;
  for all pair (t, c) in [(t1, c1), (t2, c2), ...] do
     sum = sum + t; cnt = cnt + c;
  avg = sum / cnt
  Emit(u, avg)
```

Note: Since signature of Emit changed in the Mapper, we need to modify the Reducer.

Main Point 2

Using the In-mapper combining technique, each mapper is guaranteed to produce only one key-value pair for each key by combining the values. Bringing together all the relevant ideas, it is easy to convey the central idea in a concise and precise manner. Through the regular practice of the TM technique the functioning of the mind and body are brought together to operate in accord with all the laws of Nature.

CONNECTING THE PARTS OF KNOWLEDGE WITH THE WHOLENESS OF KNOWLEDGE

- 1. It is important that safeguards be taken to maintain the consistency and correctness.
- 2. In order to be an accurate representation of the part of the real world that it is modeling, a framework must guarantee certain features.

- 3. Transcendental consciousness: is the source of all perfection in life.
- **4.** Impulses within the Transcendental Field: It is these impulses that structure and are within, everything in the universe.
- **5.** Wholeness moving within itself: In unity consciousness the individual functions in perfect harmony with all the laws of nature.