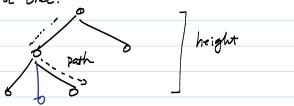
## <Lecture 20> 2.3 Non-CFL

Thm 2.34. [Pumping lenna for CFL] For any CFL A, there is a non-negative integer p ("pumping (ength") s.t. every string s can be divided into uvxyz satisfying where |5| > p

- (1). for each izo, uvigyiz e A
- (2) /y1 >0,
- (3 ( vxy 1 < p

Preparation for proving the lemna parse tree:



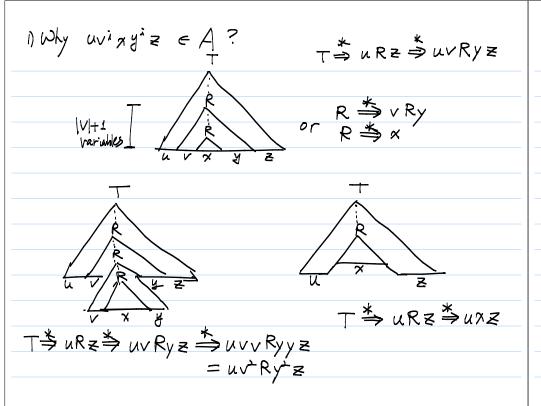
branching factor: largest number of children of any node, i.e. the maximum length of string in rhs of any rule Lemma1: The maximum length of a string that can be generated from a parse tree with height h and branching factor b is

Proof of the pumping lemma.

- o If L is context-free, there is a GFG G. Let b be the branching factor. Let  $P=b^{|V|+1}$
- Let s be any string st  $|s| \ge p$ . Since  $b^{|v|+1} > b^{lv}$ Any parse tree for s has height  $\ge |v|+1$
- o Consider a parse tree Twith the smallest number of nodes. T contains a path that contains IVI+1 variables and 1 terminal.

by pigeon hole prinicple

- O Since there are IVI variables, there must be a variable repeated in the path.
- o Let R be such a variable that repeats among the lowest |v|+1 variables on the path.



0 Why |vy >0 ?

o why Ivxy [≤p?

Example  $B = 30^n l^n 2^n ln \ge 0$  is not CF

- P Assume B is CF, There is a pumping length p. Choose  $s = a^p b^p c^p$ . Consider any division of s into uvxyz satisfying Cond 2,3.
- By condition 2, v or y is nonempty.
   By condition 3, (luxy|≤p), vxy contains at most
   2 different symbols.

Therefore uvyy'z contains at least one symbol that is repeated more than p times, and at least one symbol that is repeated p times thence it doesn't belong to B, contradicting Condition 1