

EE185523

# Consensus Protocol (Discrete)

Yurid Eka Nugraha  
TSP DTE FTEIC ITS

# Multiagent states and communication protocols

Consider  $x[k+1] = Ax[k]$ ,  $x[0] = x_0$ ,  $x \in \mathbb{R}^n$ ,  $A \in \mathbb{R}^{n \times n}$

or..

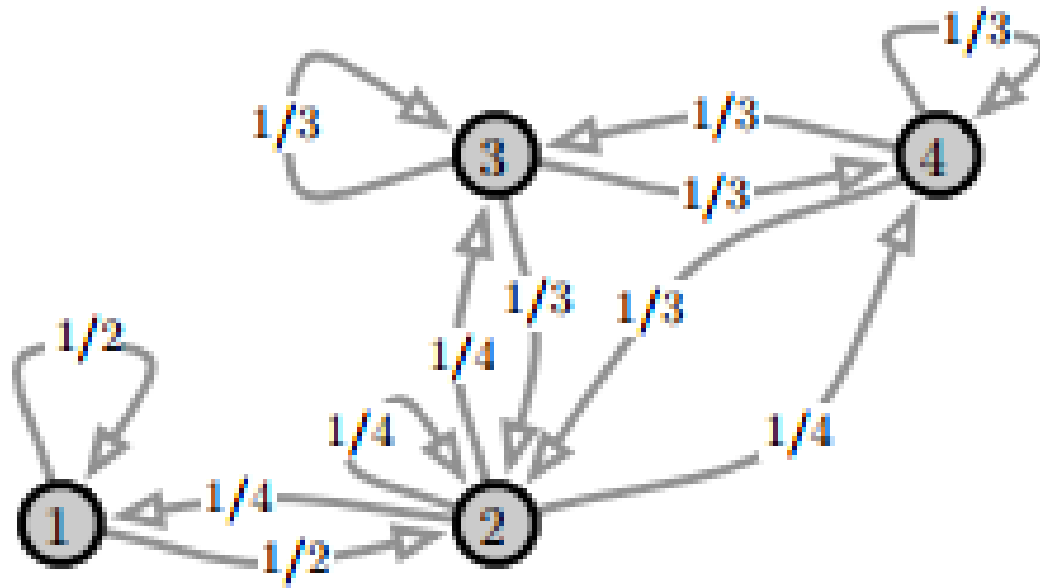
$$x[k] = A^k x_0$$

where..

$A$  row-stochastic.

what will  $x$  be?

# Graph with row-stochastic matrix



$$A = \begin{bmatrix} 1/2 & 1/2 & 0 & 0 \\ 1/4 & 1/4 & 1/4 & 1/4 \\ 0 & 1/3 & 1/3 & 1/3 \\ 0 & 1/3 & 1/3 & 1/3 \end{bmatrix}$$

Theorem:

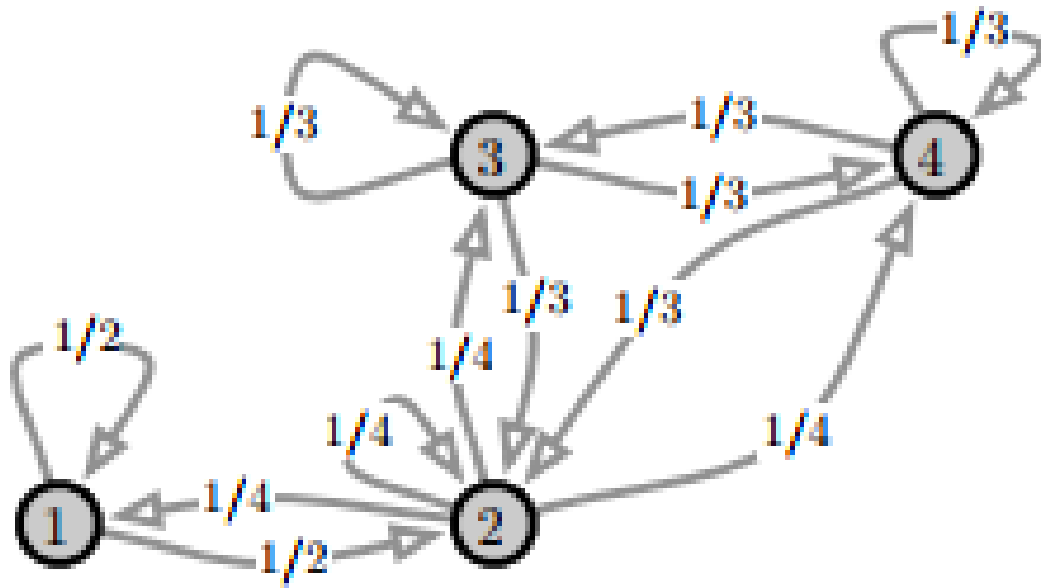
For non-negative matrix  $A$  :

$$\lim_{k \rightarrow \infty} \frac{A^k}{\lambda^k} = vw^\top$$

$w$  : left dominant eigenvector

$v$  : right dominant eigenvector

# Graph with row-stochastic matrix



$$A = \begin{bmatrix} 1/2 & 1/2 & 0 & 0 \\ 1/4 & 1/4 & 1/4 & 1/4 \\ 0 & 1/3 & 1/3 & 1/3 \\ 0 & 1/3 & 1/3 & 1/3 \end{bmatrix}$$

For row-stochastic matrix,

$$\lim_{k \rightarrow \infty} A^k = \mathbf{1}_n w^T$$

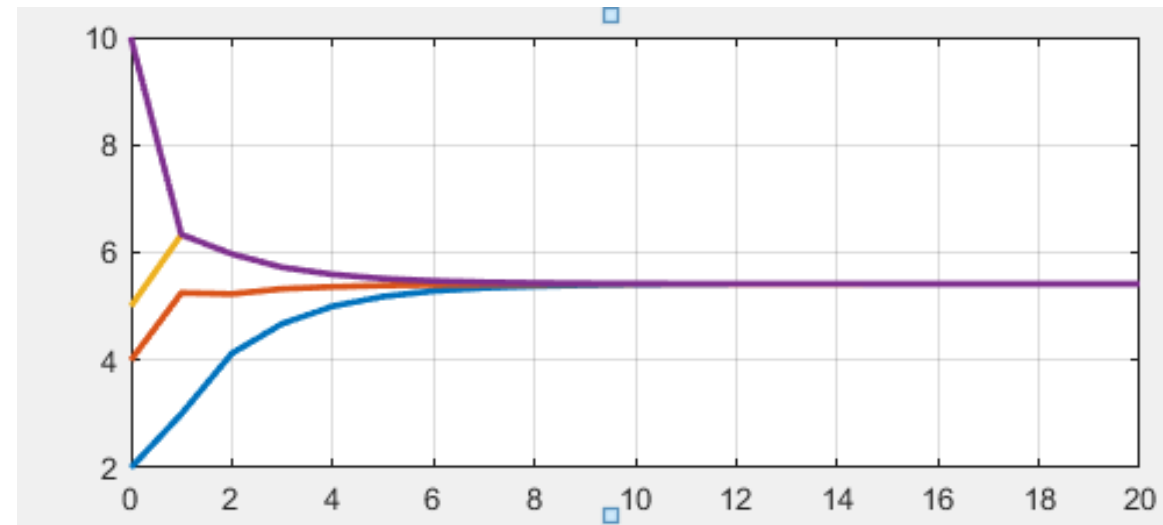
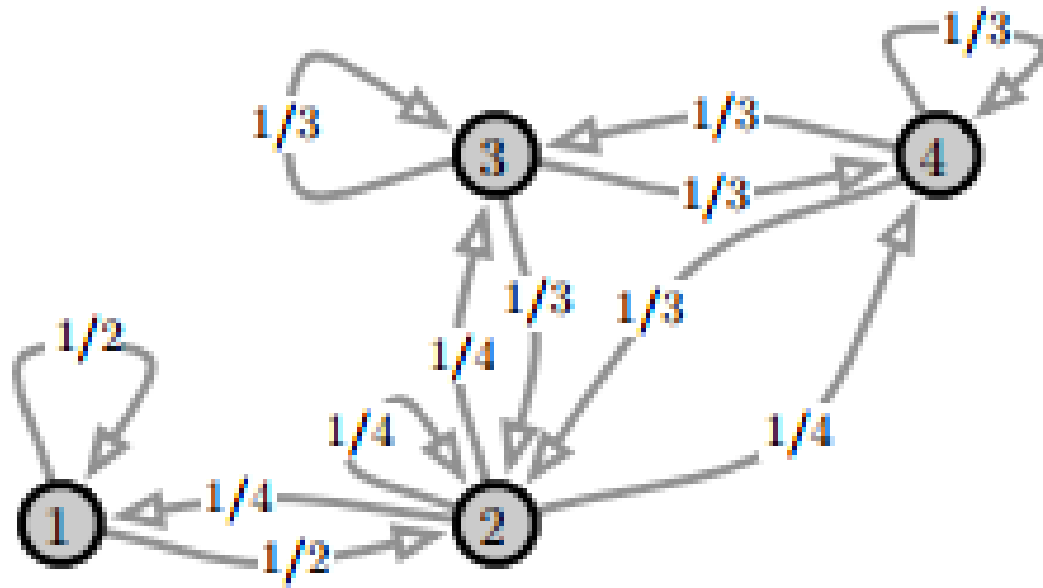
$w$  : left dominant eigenvector with sum 1

$$w : [1/6/1/3/1/4/1/4]^T$$

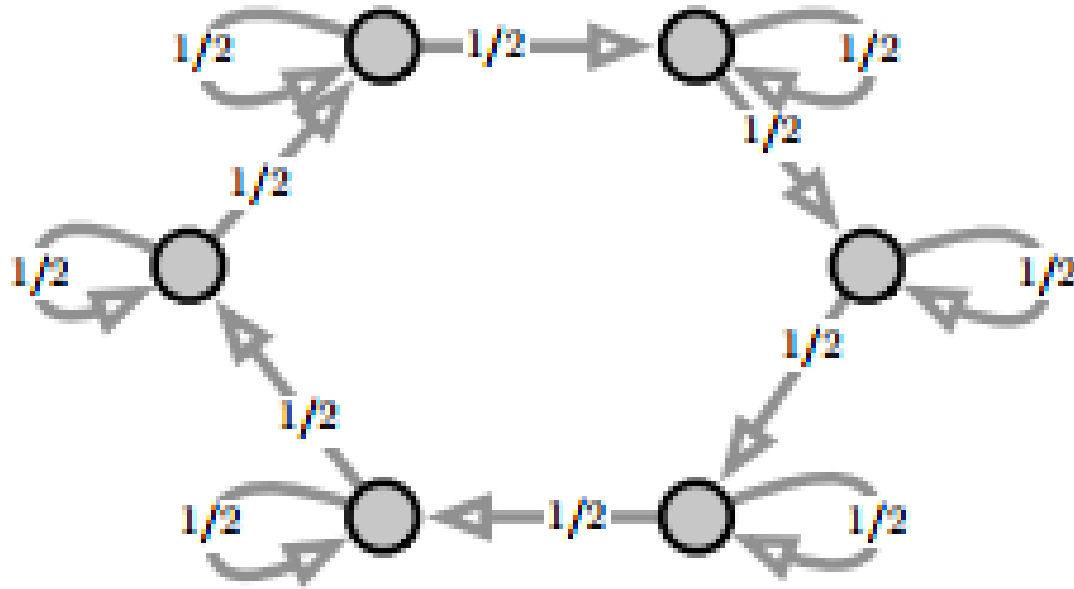
$$\lim_{k \rightarrow \infty} x(k) = (1/6)x_1(0) + (1/3)x_2(0) + (1/4)x_3(0) + (1/4)x_4(0)$$

Where do they converge?

# Graph with row-stochastic matrix



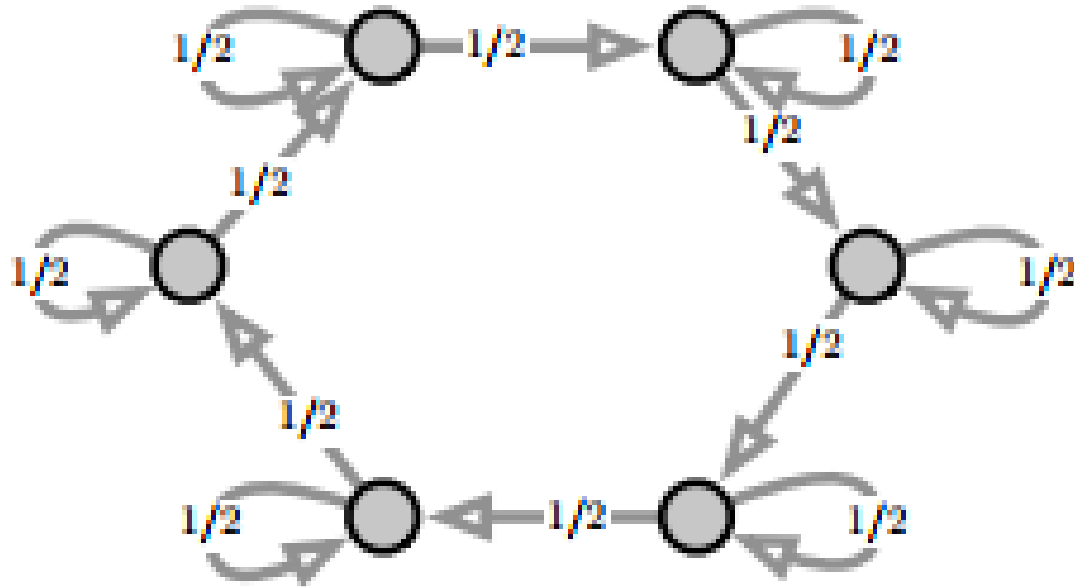
# Graph with row-stochastic matrix



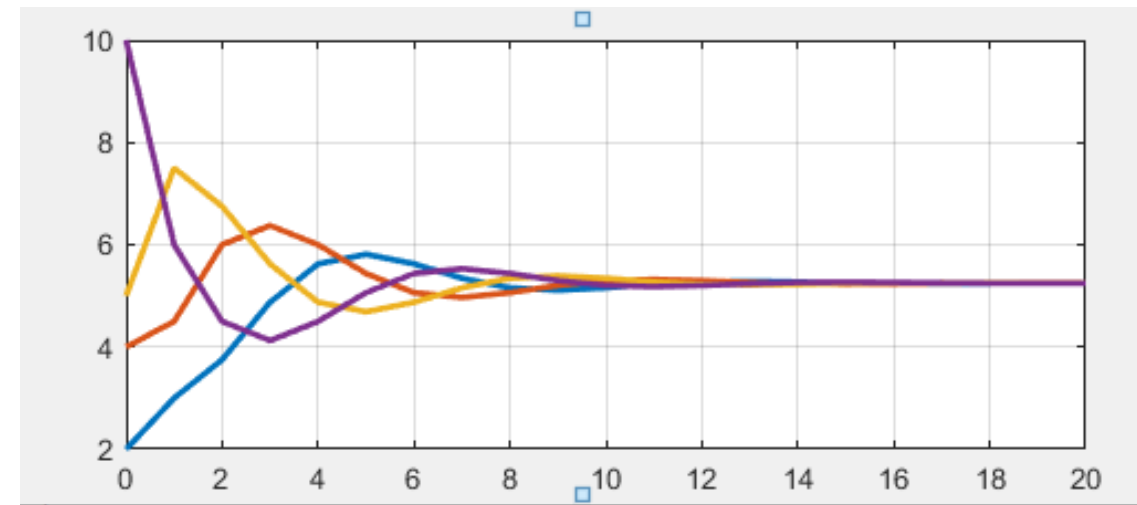
$$A = \begin{bmatrix} 1/2 & 1/2 & \dots & 0 \\ 0 & 1/2 & \dots & 0 \\ \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & 1/2 \end{bmatrix}$$

Difference from previous graph?  
Row-stochastic?

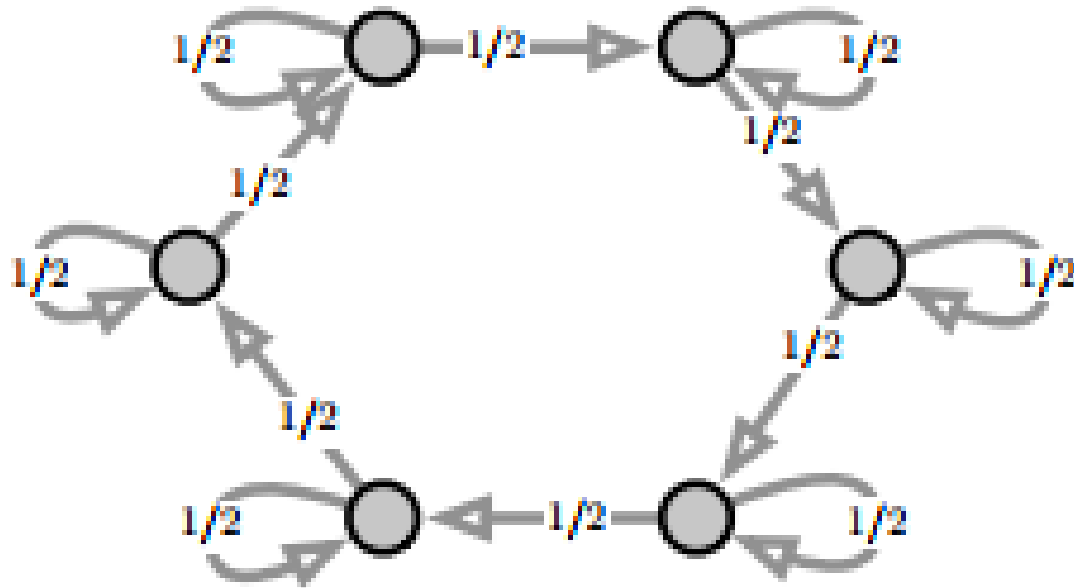
# Graph with row-stochastic matrix



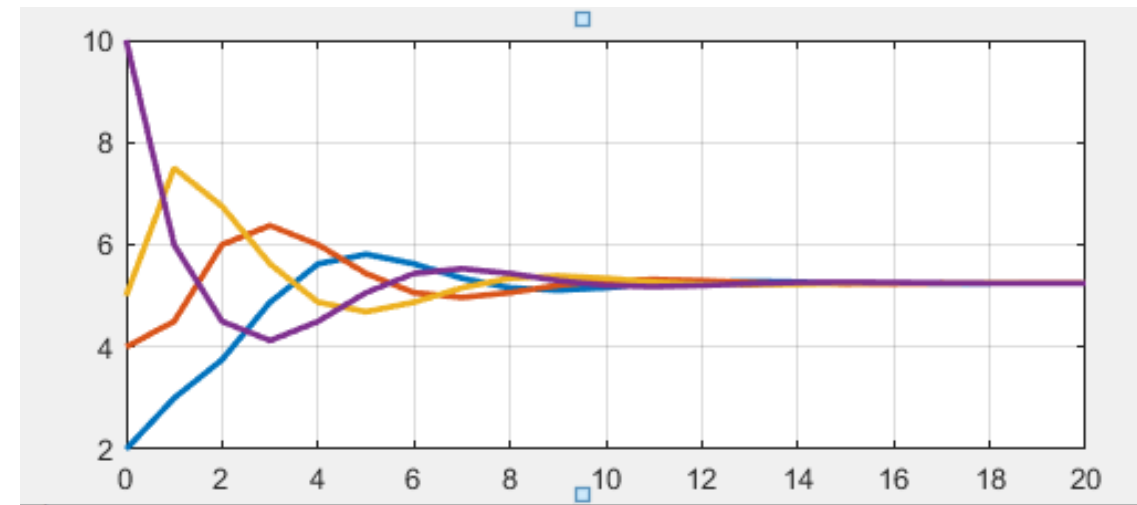
For  $n = 4$ :



# Graph with row-stochastic matrix



For  $n = 4$ :



For column-stochastic:

$$\lim_{k \rightarrow \infty} x(k) = \lim_{k \rightarrow \infty} A^k x(0) = \text{average}(x_0) \mathbf{1}_n$$



# Periodic and aperiodic digraphs

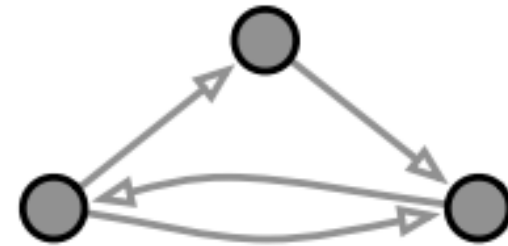
- A strongly-connected directed graph is periodic if there exists a  $k > 1$ , called the *period*, that divides the length of every simple cycle of the graph.
- In other words: a digraph is periodic if the greatest common divisor of the lengths of all its simple cycles is larger than one.
- A digraph is aperiodic if it is not periodic.



(a) A periodic digraph with period 2



(b) An aperiodic digraph with simple cycles of length 1 and 2.



(c) An aperiodic digraph with simple cycles of length 2 and 3.

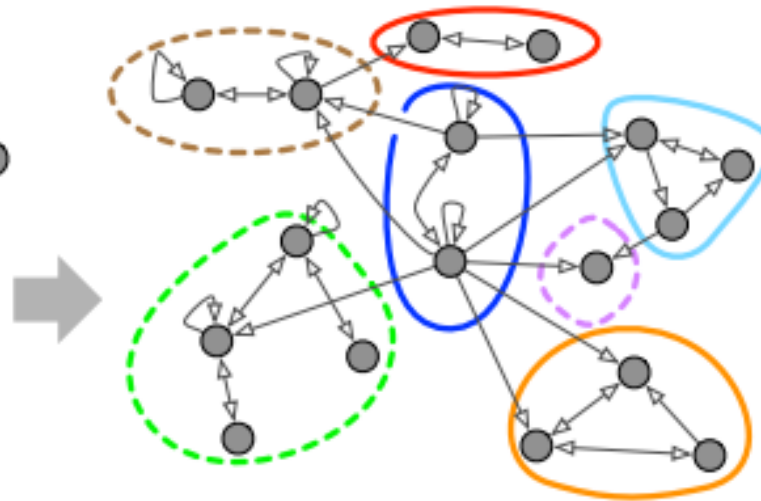
# Condensation digraphs

- A subgraph  $H$  is a strongly connected component of  $G$  if  $H$  is strongly connected and any other subgraph of  $G$  strictly containing  $H$  is not strongly connected.
- The condensation digraph of a digraph  $G$  denoted by  $C(G)$ , is defined as follows: the nodes of  $C(G)$  are the strongly connected components of  $G$ , and there exists a directed edge in  $C(G)$  from node  $H_1$  to node  $H_2$  if and only if there exists a directed edge in  $G$  from a node of  $H_1$  to a node of  $H_2$ .
- The condensation digraph has no self-loops.

# Condensation digraphs



Digraph



Connected components



Condensation digraph

# Consensus with a globally-reachable aperiodic strongly-connected component

Theorem:

Let  $A$  be row-stochastic matrix associated with a digraph  $G$ .

If  $A$  is semi-convergent and  $\lim_{k \rightarrow \infty} A^k = \mathbf{1}_n w^\top$  where  $\mathbf{1}_n^\top w = 1$ ,  $w^\top A = w^\top$ , then  $G$  contains a **globally reachable node**.

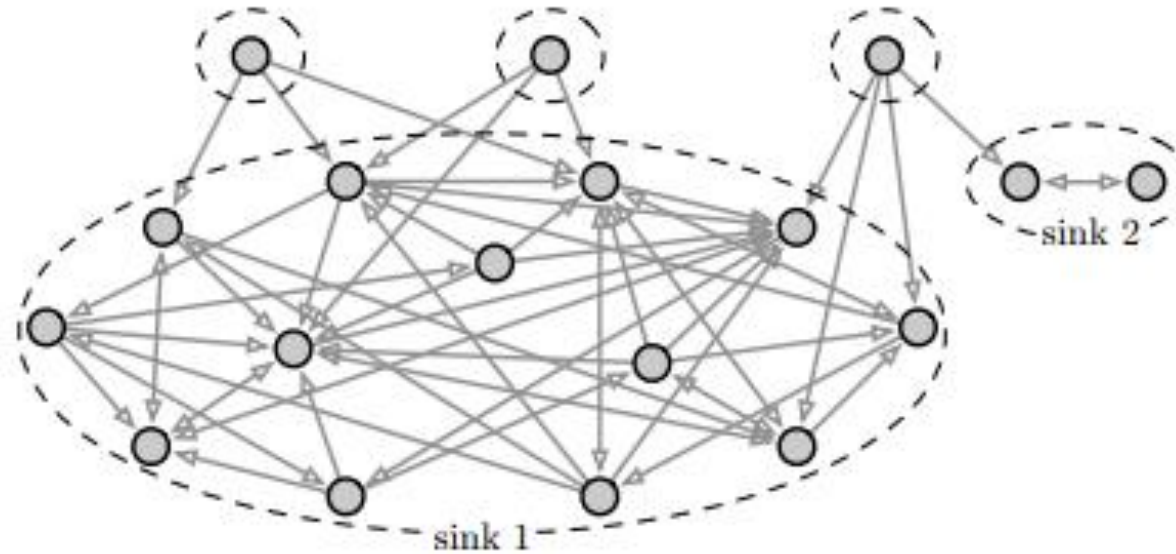
Furthermore, matrix  $A$  is said to be *indecomposable* and

- (i) The solution of the model  $x(k+1) = Ax(k)$  satisfies  $\lim_{k \rightarrow \infty} x(k) = (w^\top x(0)) \mathbf{1}_n$
- (ii) If additionally  $A$  is doubly-stochastic, then

$$\lim_{k \rightarrow \infty} x(k) = \frac{\mathbf{1}_n^\top x(0)}{n} = \text{average}(x(0)) \mathbf{1}_n$$

# Averaging system reaching averaging disagreement

Digraph without globally reachable nodes:



# Convergence for row-stochastic matrices with multiple aperiodic sinks

Theorem:

Let  $A$  be row-stochastic matrix associated with a digraph  $G$  and  $n_s$  be the number of sinks in the condensation digraph  $C(G)$ .

If  $A$  is semi-convergent, then the solution to  $x(k+1) = Ax(k)$  satisfies

$$\lim_{k \rightarrow \infty} x_i(k) = \begin{cases} (w^P)^T x(0), & \text{if node } i \text{ belongs to sink } p, \\ \sum_{p=1}^{n_s} z_{i,p} ((w^P)^T x(0)), & \text{otherwise} \end{cases}$$

Where  $z_{i,p}$ ,  $p \in \{1, \dots, n_s\}$  are convex combination coefficient and  $z_{i,p} > 0$  if and only if there exists a directed walk from node  $i$  to the sink  $p$

# Convergence for row-stochastic matrices with multiple aperiodic sinks

Message:

- Convergence does not occur to consensus (not all components of the state are equal)
- The final value of all nodes is independent of the initial values at nodes which are not in the sinks of the condensation digraph

# Convergence for row-stochastic matrices with multiple aperiodic sinks

