Fourth Problem Set

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Problem 1

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\begin{split} B &\models C \text{ just in case } \models (B \to C) \\ \text{just in case } \bar{v}(B \to C) = T \\ \text{just in case} \\ \bar{v}(B) &= T \text{ and } \bar{v}(C) = T \\ \text{or} \\ \bar{v}(B) &= F \text{ and } \bar{v}(C) = T \\ \text{or} \\ \bar{v}(B) &= F \text{ and } \bar{v}(C) = F \\ \text{In the second two cases, } \bar{v}(B) &= F \text{ which means } \bar{v}(\neg B) = T \\ \text{Hence } &\models \neg B \\ \text{In the first two cases } \bar{v}(C) &= T \text{ which means } \models C \\ \text{So in all three cases } &\models \neg b \text{ or } \models C. \end{split}
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Problem 2

(a)

$$(x \wedge y \wedge \neg z) \vee (x \wedge \neg y \wedge z) \vee (\neg x \wedge y \wedge z) \vee (\neg x \wedge y \wedge \neg z) \vee (\neg x \wedge \neg y \wedge \neg z)$$

(b)

$$\neg(x \land y \land z) \land \neg(x \land \neg y \land \neg z) \land \neg(\neg x \land \neg y \land z)$$

$$\Rightarrow (\neg x \lor \neg y \lor \neg z) \land (\neg x \lor \neg y \lor z) \land (x \lor y \lor \neg z)$$

Problem 3

(a)

Let
$$C = \{\text{customers}\}\ J = \{\text{Jeans}\}\ J_x = \{\text{Jeans that fit } x \in C\}$$
. We know that $\forall x_1, \ldots, x_m \ \exists j_1, \ldots, j_m \ \text{such that } j_i \in J_{x_i} \ \text{where } i = 1, \ldots, m$

Assume J_x is finite $\forall x \in C$

Need to show that there is a way of assigning jeans to customers so that every $x \in C$ buys a $j \in J$ that fits and no two customers but the same pair of jeans.

(i.e. $\exists p: C \to J \text{ such that } p(x) \in J_x \text{ and } p \text{ is injective})$

Let $\Sigma \subseteq \overline{\mathbb{S}}$ such that $\exists v : \mathbb{S} \to \{T, F\}$ satisfying Σ iff There exists such a function p

We will use sentence letters: b_{xj} to build the wffs in Σ

Let
$$\Sigma = A \cup B \cup C$$
 where $A = \{ \bigvee_{j \in J_x} b_{xj}, \forall x \in C \}$
 $B = \{ b_{xj} \rightarrow \neg b_{xl}, \forall x \in C, \forall j, l \in J, j \neq l \}$
 $C = \{ b_{xj} \rightarrow \neg b_{yj}, \forall x, y \in C, \forall j \in J \}$

It is useful to note that by A every customer $x \in C$ buys a pair of jeans (i.e. $p(x) \in J_x$), by B every customer only buys one pair of jeans (i.e. p is well defined) and by C every pair of jeans is only bought by one customer (i.e. p is injective).

Now we can show that Σ is satisfiable using the compactness theorem.

Let $\Delta \subseteq \Sigma$ be finite.

Lemma: $\Delta \subseteq \Delta'$ and Δ' is satisfiable then Δ is satisfiable.

Let $C_0 = \{x \in C | b_{xj} \text{ shows up in some wff in } \Delta\}$ and $J_0 = \{j \in J | b_{xj} \text{ shows up in some wff in } \Delta\}$, these sets are both finite because Δ is finite. Now let $J_1 = J_0 \cup \bigcup_{x \in C_0} J_x$ and take Δ' to be the set:

$$\Delta' = \{ \text{all wffs in } \Sigma \text{ built from } b_{xj} \text{ with } x \in C_0 \text{ and } j \in J_1 \}$$

Now to show that Δ' is satisfiable, we have a finite set of customers C_0 and we can define

$$p':C_0\to J_1$$

We can show that $p'(x) \in J_{1_x}$ because the definition of J_1 says that it contains all of the sets of fitting jeans for each customer. We know that p' is also well defined and injective. Therefore Δ' is satisfiable. Meaning Δ is satisfiable and by the compactness theorem Σ is satisfiable.

(b)

We can show that the conclusion will not hold if we drop the extra hypothesis by providing a coutnerexample.

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Let C = \{c_0, c_1, c_2, \ldots\} and J = \{j_0, j_1, j_2, \ldots\}
There is at least one customer that fits infinitely many jeans because we dropped the extra hypothesis (i.e. say c_0 fits j_0, j_1, j_2, \ldots)
Say c_1 only fits j_0 c_2 only fits j_1
\vdots
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This satisfies the conditions of the problem because given a finite set $D = \{x_0, x_1, \ldots, x_m\} \subseteq C$ we can say c_i fits j_{i-1} for $i = 1, \ldots, m$ then c_0 fits j_m . Now if we set up a function $p: C \to J$ as:

$$p(c_i) = j_{i-1}, \forall i > 0$$

Then there is no pair of jeans for c_0 to buy.