

# Deductive Program Verification with Vampire

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```
theory Scratch
```

```
imports Main
```

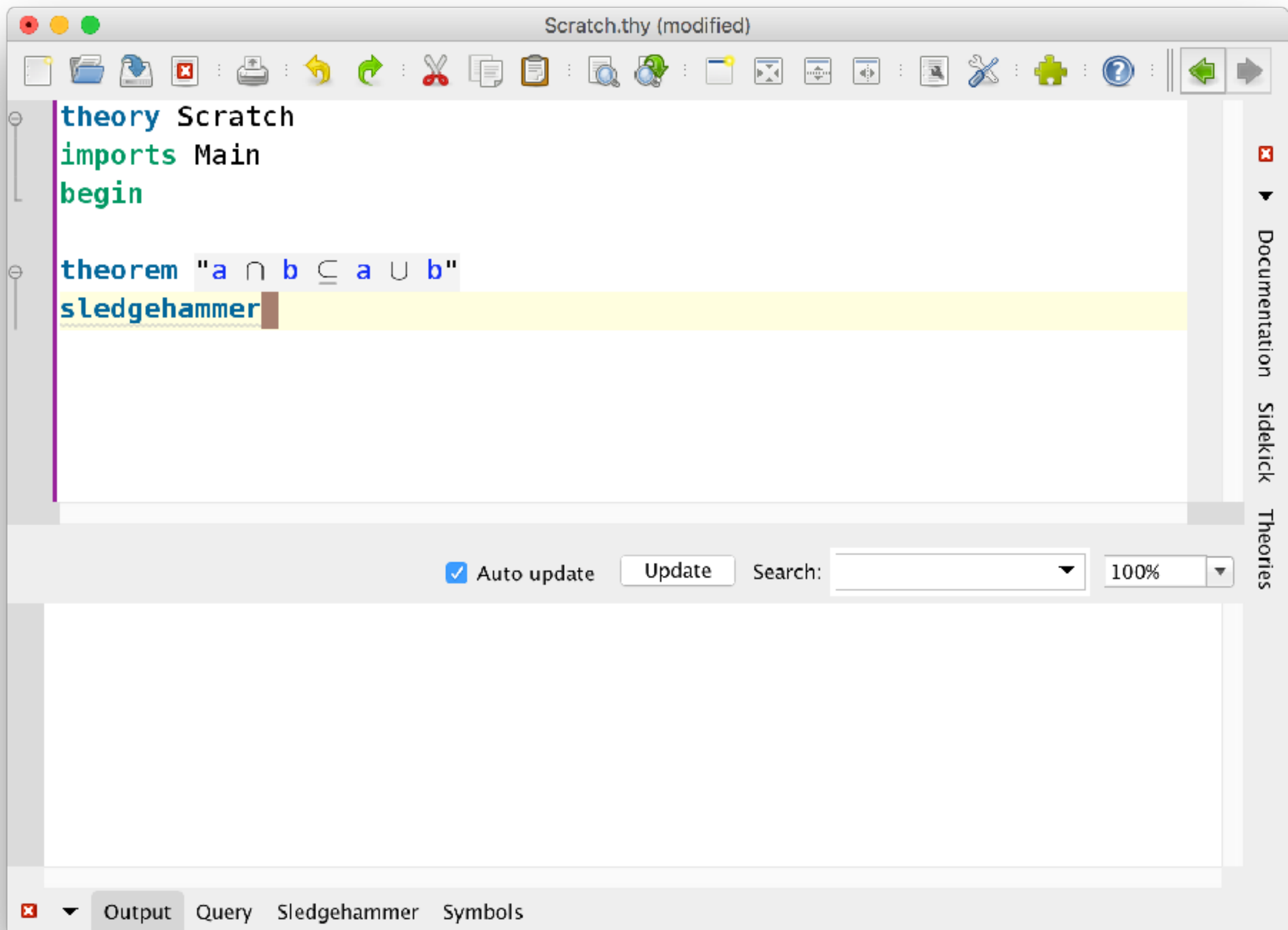
```
begin
```

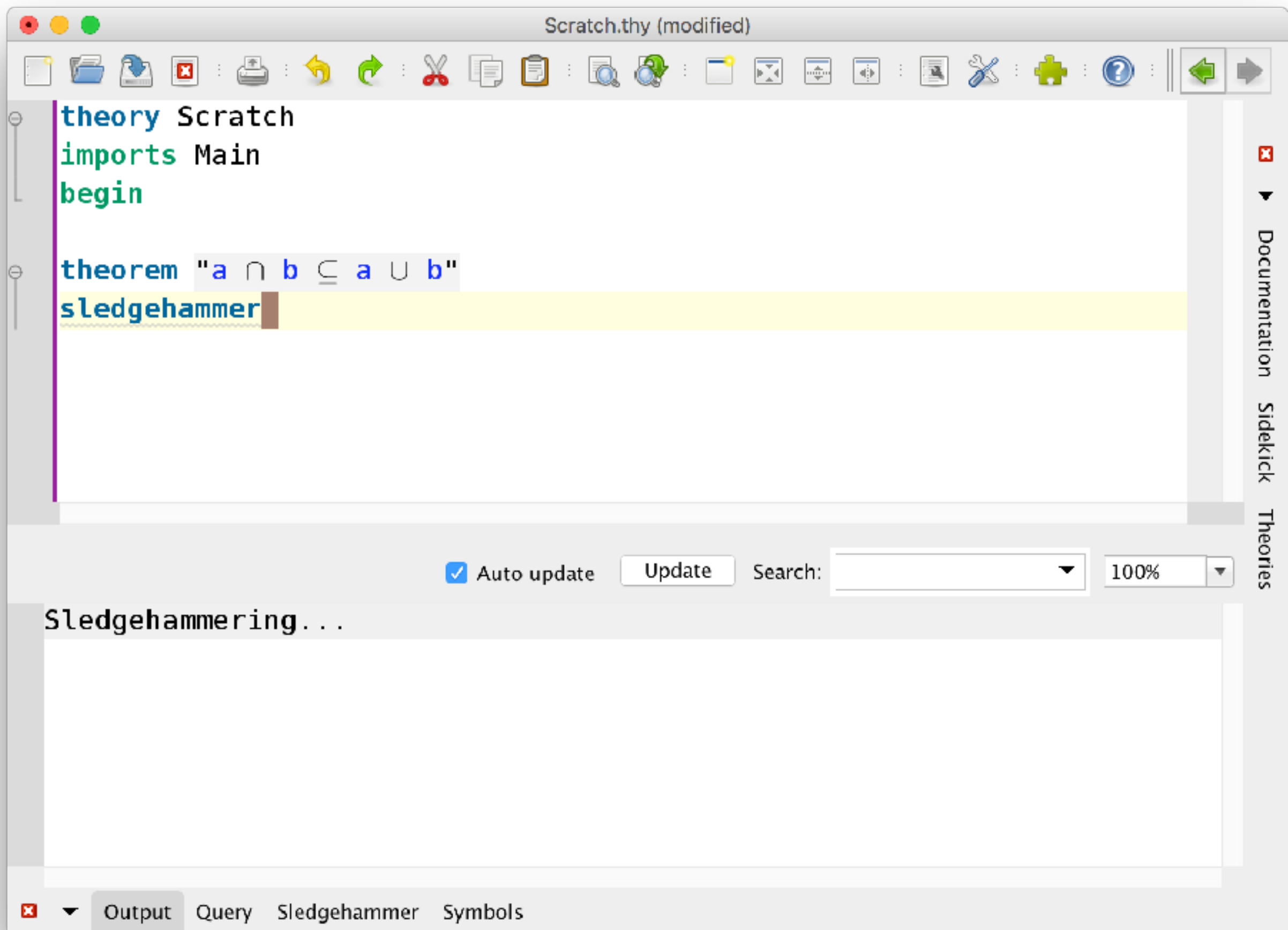
```
theorem " $a \cap b \subseteq a \cup b$ "
```

```
proof (prove): depth 0
```


```
goal (1 subgoal):
```

```
1.  $a \cap b \subseteq a \cup b$ 
```





Scratch.thy (modified)



```
theory Scratch
imports Main
begin

theorem "a ∩ b ⊆ a ∪ b"
sledgehammer
```

☒ Auto update

Update



Search:

100%

Sledgehammering...

"cvc4": Try this: by (simp add: le\_infI2) (0.9 ms).


"remote\_vampire": Try this: by (simp add: inf.coboundedI2) (0.4 ms).


  Output

Query

Sledgehammer

Symbols

 Documentation

 Sidekick

Theories



```
theory Scratch
```

```
imports Main
```

```
begin
```

```
theorem "a  $\cap$  b  $\subseteq$  a  $\cup$  b"
```

```
by (simp add: le_infI2)
```



Documentation

Sidekick

Theories

☒ Auto update

Update

Search:

100%

```
theorem ?a  $\cap$  ?b  $\subseteq$  ?a  $\cup$  ?b
```

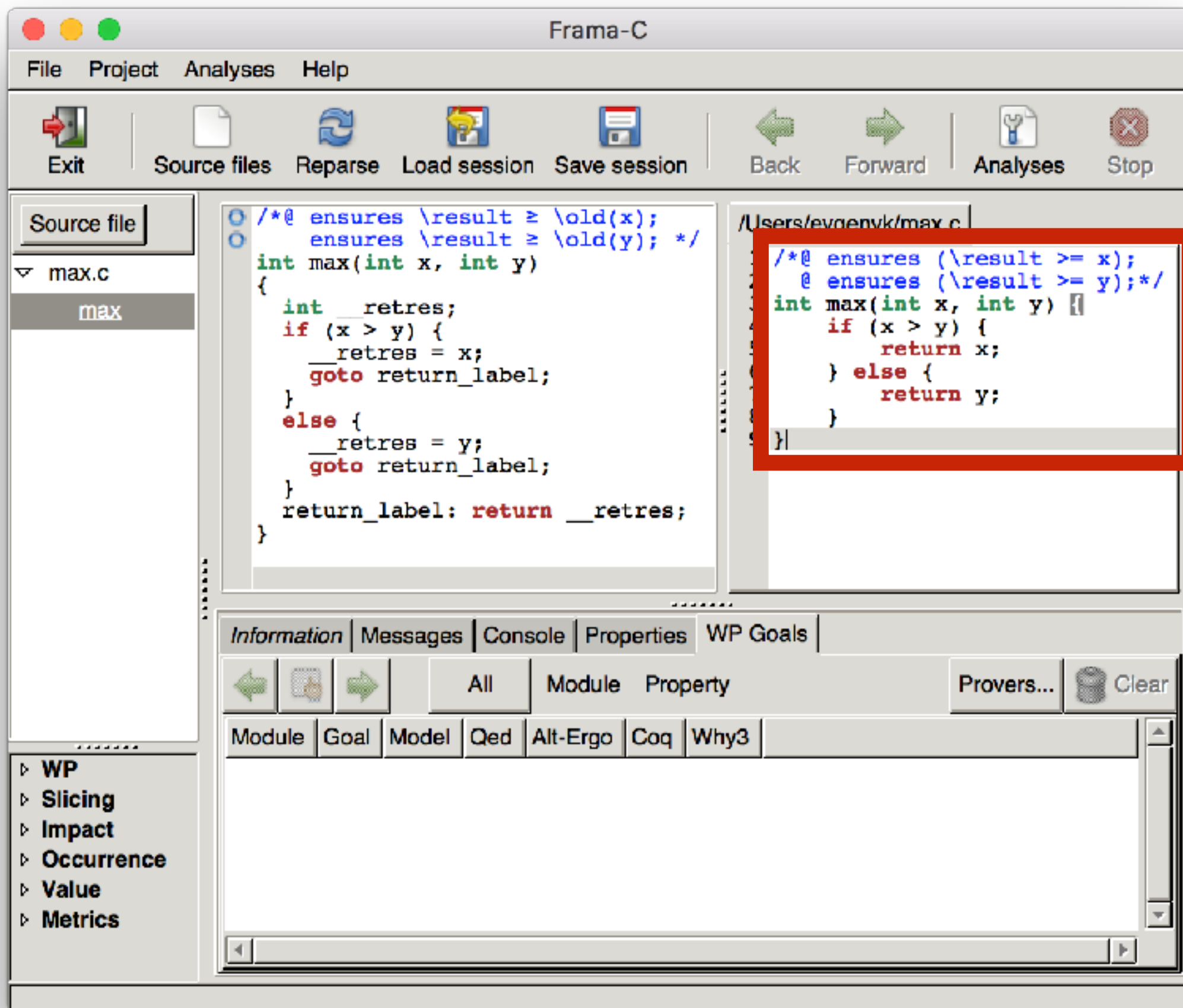


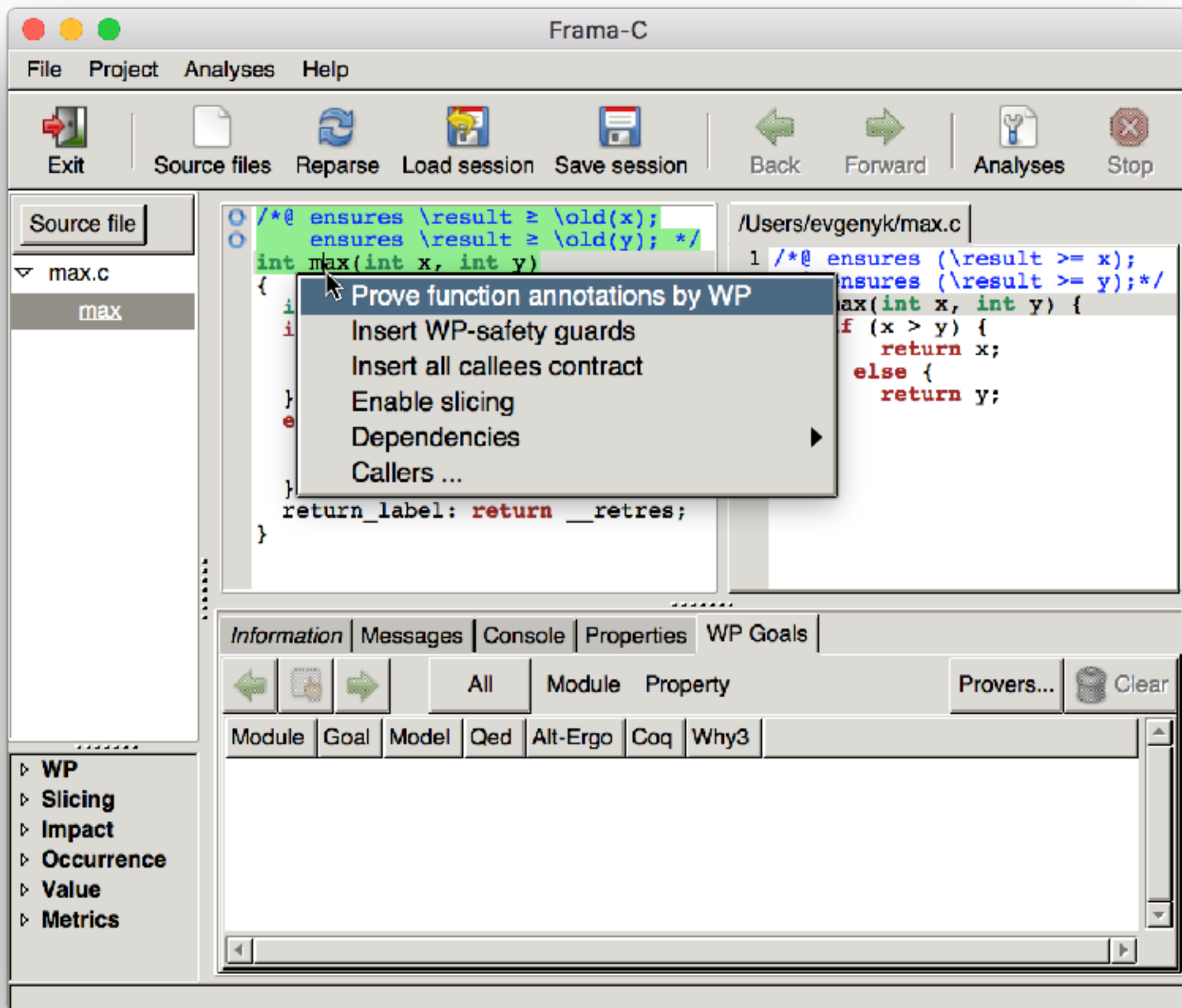
Output

Query

Sledgehammer

Symbols







Frama-C

File

Project

Analyses

Help

Exit

Source files

Reparse

Load session

Save session

Back

Forward

Analyses

Stop

Source file

max.c

max

```

/*@ ensures \result ≥ \old(x);
   ensures \result ≥ \old(y); */
int max(int x, int y)
{
    int __retres;
    if (x > y) {
        __retres = x;
        goto return_label;
    }
    else {
        __retres = y;
        goto return_label;
    }
    return_label: return __retres;
}

```

/Users/evgenyk/max.c

```

1 /*@ ensures (\result >= x);
2   @ ensures (\result >= y); */
3 int max(int x, int y) {
4     if (x > y) {
5         return x;
6     } else {
7         return y;
8     }
9 }

```

Information

Messages

Console

Properties

WP Goals

←

👉

→

All

Module

Property

Provers...

Clear

| Module | Goal           | Model | Qed | Alt-Ergo | Coq | Why3 |
|--------|----------------|-------|-----|----------|-----|------|
| max    | Post-condition | Typed | ●   |          |     |      |
| max    | Post-condition | Typed | ●   |          |     |      |

WP

Slicing

Impact

Occurrence

Value

Metrics

User's  
domain

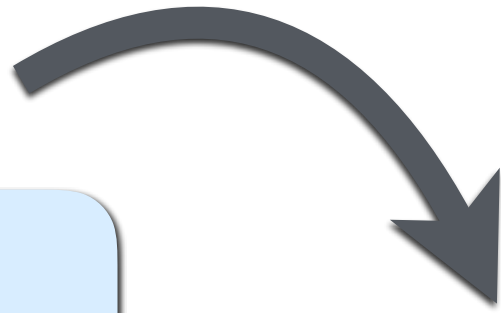
```
graph LR; A[User's domain] -- curved arrow --> B[FOL]; B -- straight arrow --> C[FOL theorem prover];
```

The diagram illustrates a workflow starting from a 'User's domain' (blue box) on the left. A curved arrow points from the top of this box to a central green box labeled 'FOL'. From the 'FOL' box, a straight arrow points to a yellow box on the right labeled 'FOL theorem prover'.

FOL

FOL  
theorem  
prover

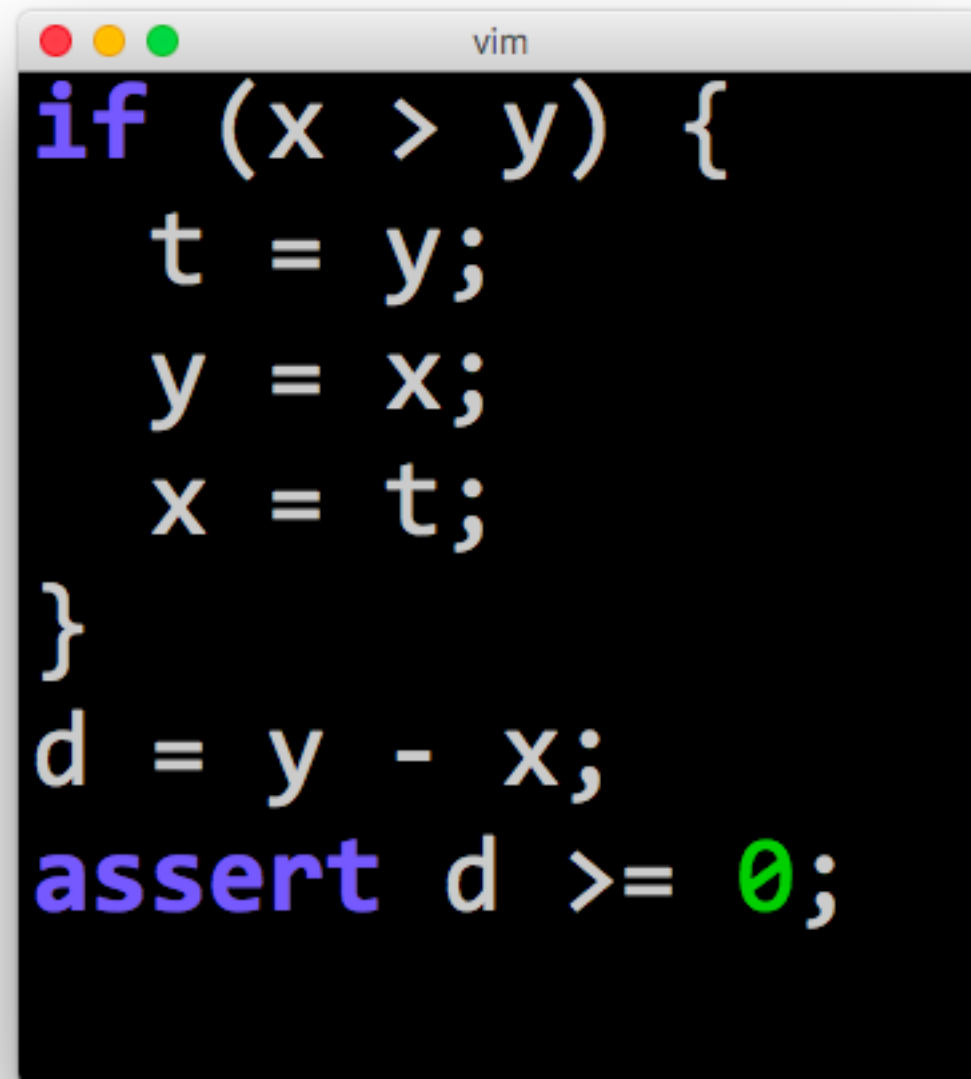
User's  
domain



FOL

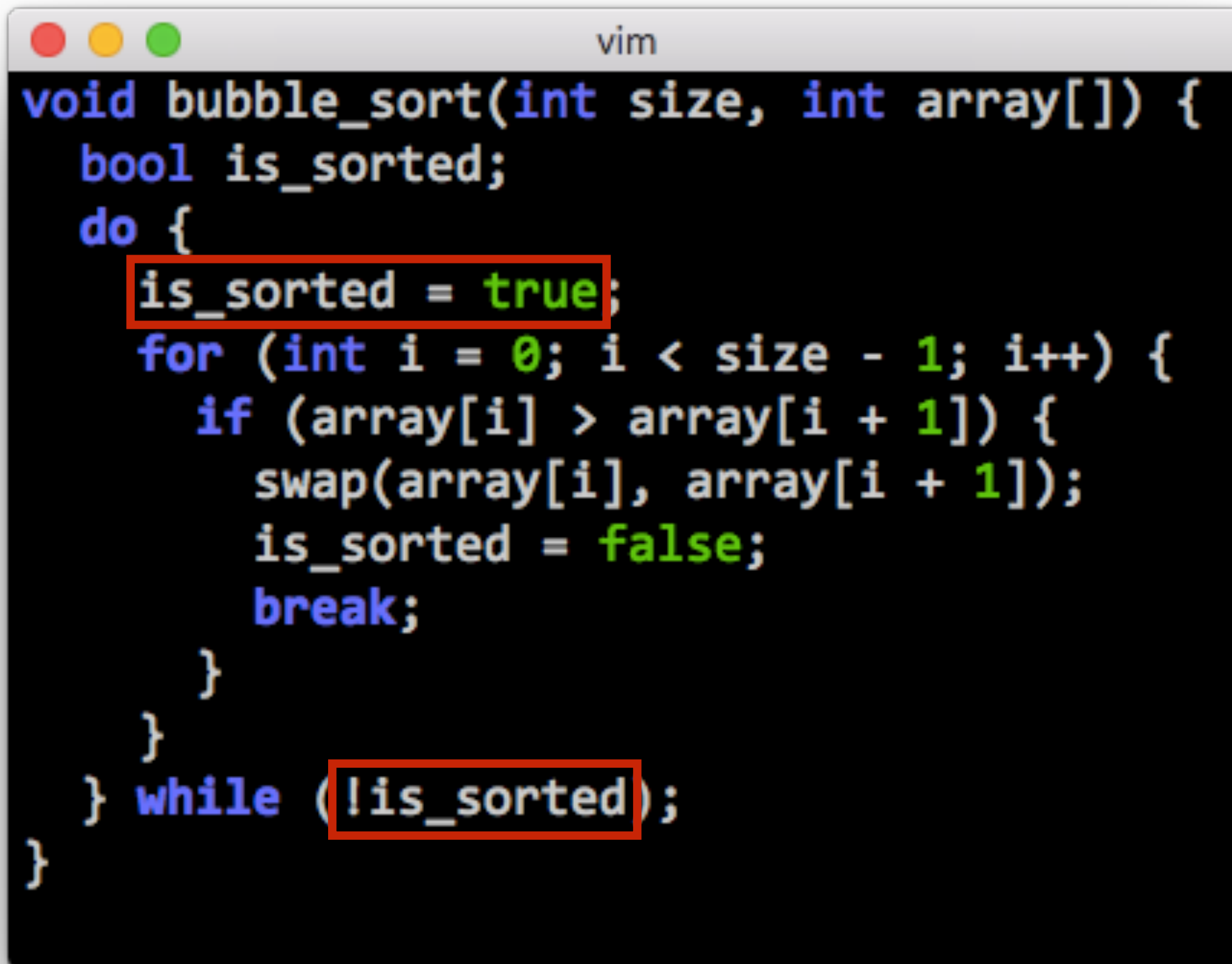


FOL  
theorem  
prover



```
vim
if (x > y) {
    t = y;
    y = x;
    x = t;
}
d = y - x;
assert d >= 0;
```





```
vim
void bubble_sort(int size, int array[]) {
    bool is_sorted;
    do {
        is_sorted = true;
        for (int i = 0; i < size - 1; i++) {
            if (array[i] > array[i + 1]) {
                swap(array[i], array[i + 1]);
                is_sorted = false;
                break;
            }
        }
    } while (!is_sorted);
}
```

$$true \neq false$$

$$(\forall x : bool)(x = true \vee x = false)$$

- First-class boolean sort

$(\forall x : \text{bool})(x \vee F(x))$

$Q((\forall x : \sigma)P(x))$

- Condition expressions

**if**  $\varphi$  **then**  $\psi_1$  **else**  $\psi_2$

$Q(\text{if } \varphi \text{ then } s \text{ else } t)$

- Local definitions

**let**  $f(x) = g(x, g(x))$  **in**  $Q(f(a), f(b))$

- Tuples

**let**  $(a, b) = (b, a)$  **in**  $f(a, b)$



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- Tuples

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- First-class boolean sort

$$(\forall x : \text{bool})(x = \text{true} \vee F(x))$$

$$Q((\forall x : \sigma)P(x))$$

- Condition expressions

$$(\varphi \Rightarrow \psi_1) \wedge (\neg \varphi \Rightarrow \psi_2)$$

$$Q(\text{if } \varphi \text{ then } s \text{ else } t)$$

- Local definitions

$$\text{let } f(x) = g(x, g(x)) \text{ in } Q(f(a), f(b))$$

- Tuples

$$\text{let } (a, b) = (b, a) \text{ in } f(a, b)$$

- First-class boolean sort

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$$\text{let } (a, b) = (b, a) \text{ in } f(a, b)$$

- First-class boolean sort

$$(\forall x : \textit{bool})(x = \textit{true} \vee F(x))$$

$$Q(g_1)$$

- Condition expressions

$$(\varphi \Rightarrow \psi_1) \wedge (\neg \varphi \Rightarrow \psi_2)$$

$$Q(g_2)$$

- Local definitions

$$Q(g_3(a), g_3(b))$$

- Tuples

$$f(g_4, g_5)$$

$$(\textit{Premise}_1 \wedge \dots \wedge \textit{Premise}_n) \Rightarrow \textit{Conjecture}$$

$$(\textit{Premise}_1 \wedge \dots \wedge \textit{Premise}_n \wedge \neg \textit{Conjecture}) \Rightarrow \perp$$



- CNF:  $C_1 \wedge \dots \wedge C_n$
- Clause:  $(\forall x_1 : \sigma_1) \dots (\forall x_n : \sigma_n)(L_1 \vee \dots \vee L_n)$
- Literal:  $A$  or  $\neg A$
- Atom:  $P(t_1, \dots, t_n)$  or  $t_1 = t_2$

$$(\forall x : \tau)(P(x) \wedge Q(x) \Leftrightarrow R(x, f(x)))$$



$$\left\{ \begin{array}{l} \neg P(x) \vee \neg Q(x) \vee R(x, f(x)) \\ P(x) \vee \neg R(x, f(x)) \\ Q(x) \vee \neg R(x, f(x)) \end{array} \right\}$$

$$(A_1 \wedge B_1) \vee \dots \vee (A_n \wedge B_n)$$

$$(R_1 \vee \dots \vee R_n) \wedge (R_1 \Leftrightarrow (A_1 \wedge B_1)) \wedge \dots \wedge (R_n \Leftrightarrow (A_1 \wedge B_n))$$

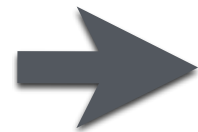
$$\frac{A \vee C_1 \quad \neg A' \vee C_2}{(C_1 \vee C_2)\theta} \quad \theta = mgu(A, A')$$

$$\frac{l = r \vee C_1 \quad L[s] \vee C_2}{(L[r] \vee C_1 \vee C_2)\theta} \quad \theta = mgu(l, s)$$

$$(\forall x : \textit{bool})(x = \textit{true} \vee x = \textit{false})$$

$$\frac{x = \textit{true} \vee x = \textit{false} \quad y = \textit{true} \vee y = \textit{false}}{x = y \vee x = \textit{false} \vee y = \textit{false}}$$

FOL



FO  
CNF



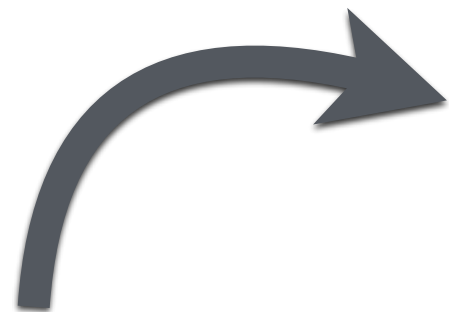
FOL  
theorem  
prover



if  $\varphi$  then  $\psi_1$  else  $\psi_2$

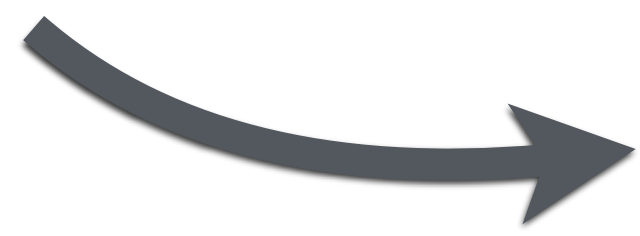


$$\left\{ \begin{array}{l} \neg\varphi, \psi_1 \\ \varphi, \psi_2 \end{array} \right\}$$



$$\left\{ \begin{array}{l} Q(x) \vee \neg G(x) \\ \neg \varphi \vee G(s) \\ \varphi \vee \neg G(t) \end{array} \right\}$$

$Q(\text{if } \varphi \text{ then } s \text{ else } t)$



$$\left\{ \begin{array}{l} \neg \varphi \vee Q(s) \\ \varphi \vee Q(t) \end{array} \right\}$$

$$Q((\forall x : \sigma)P(x))$$

$Q(\text{if } (\forall x : \sigma)P(x) \text{ then } \textit{true} \text{ else } \textit{false})$

$$(\forall x : \textit{bool})(x \vee F(x))$$

$$(\top \vee F(\textit{true})) \wedge (\perp \vee F(\textit{false}))$$

$F(false)$

$$\left\{ \begin{array}{l} g = f(b) \\ R(g, g) \end{array} \right\}$$

let  $a = f(b)$  in  $R(a, a)$

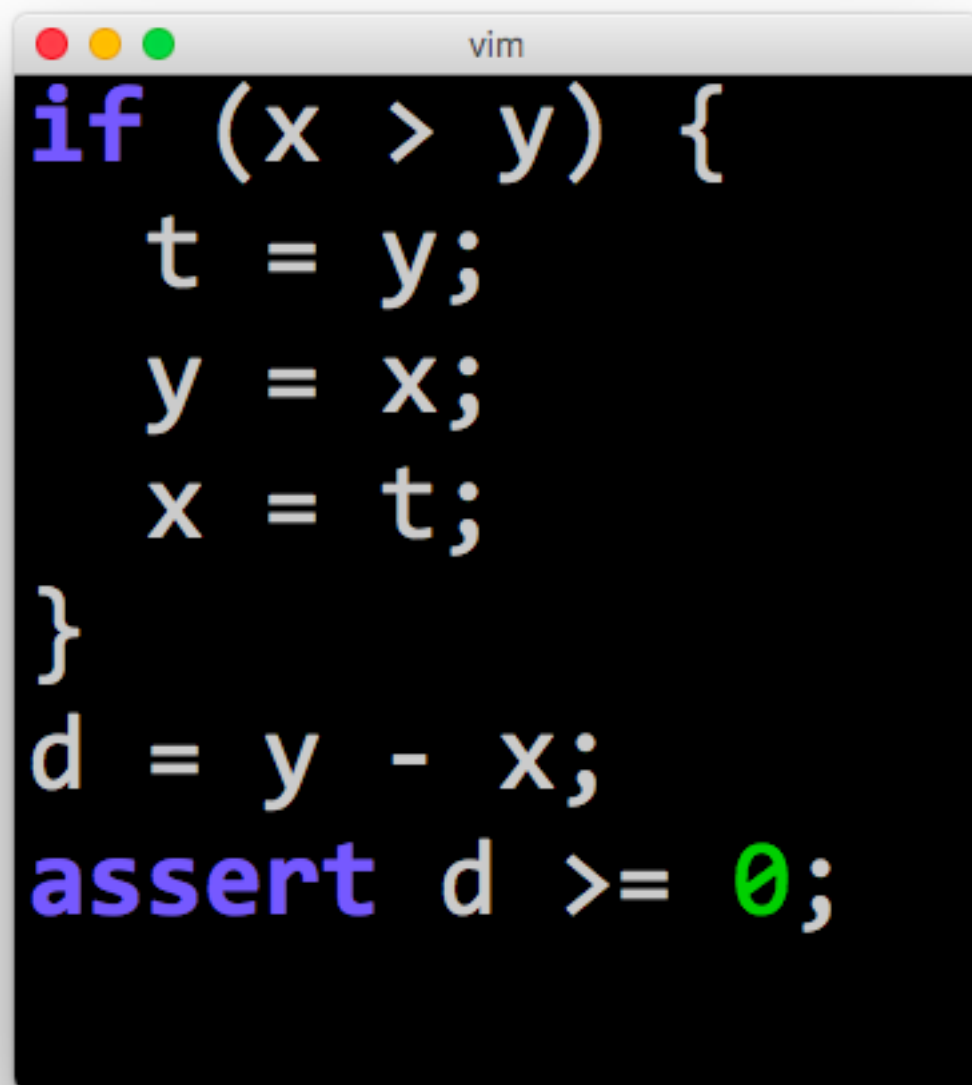
$$\{ R(f(b), f(b)) \}$$



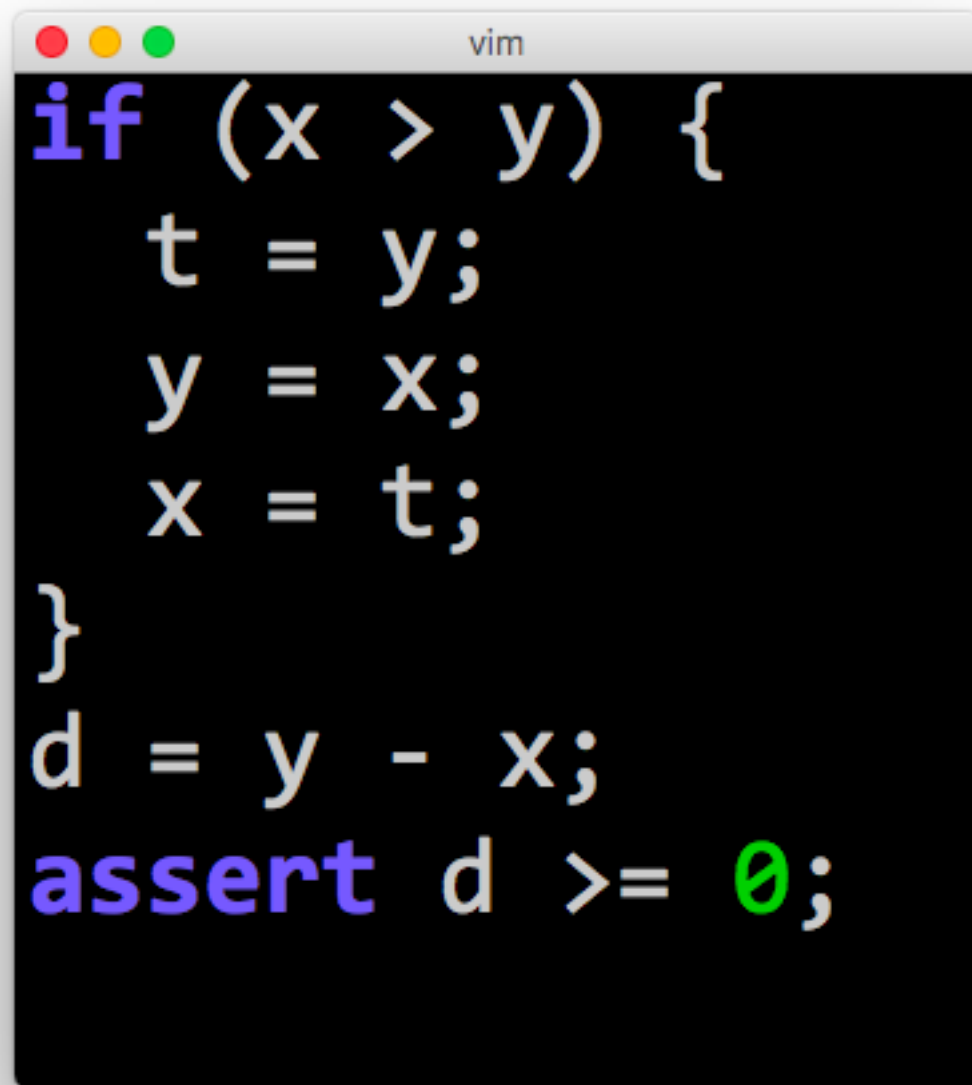
**let**  $P(x) = Q(x) \vee R(x)$  **in**  $P(a) \wedge P(b)$

**let**  $(a, b) = (b, a)$  **in**  $f(a, b)$


**let**  $t = (b, a)$  **in**  $f(\pi_1(t), \pi_2(t))$

A screenshot of a vim editor window. The window has a title bar with three colored buttons (red, yellow, green) and the text 'vim'. The code is displayed on a black background with syntax highlighting: 'if' and 'assert' are in blue, '0' is in green, and other tokens are in white. The code is as follows:

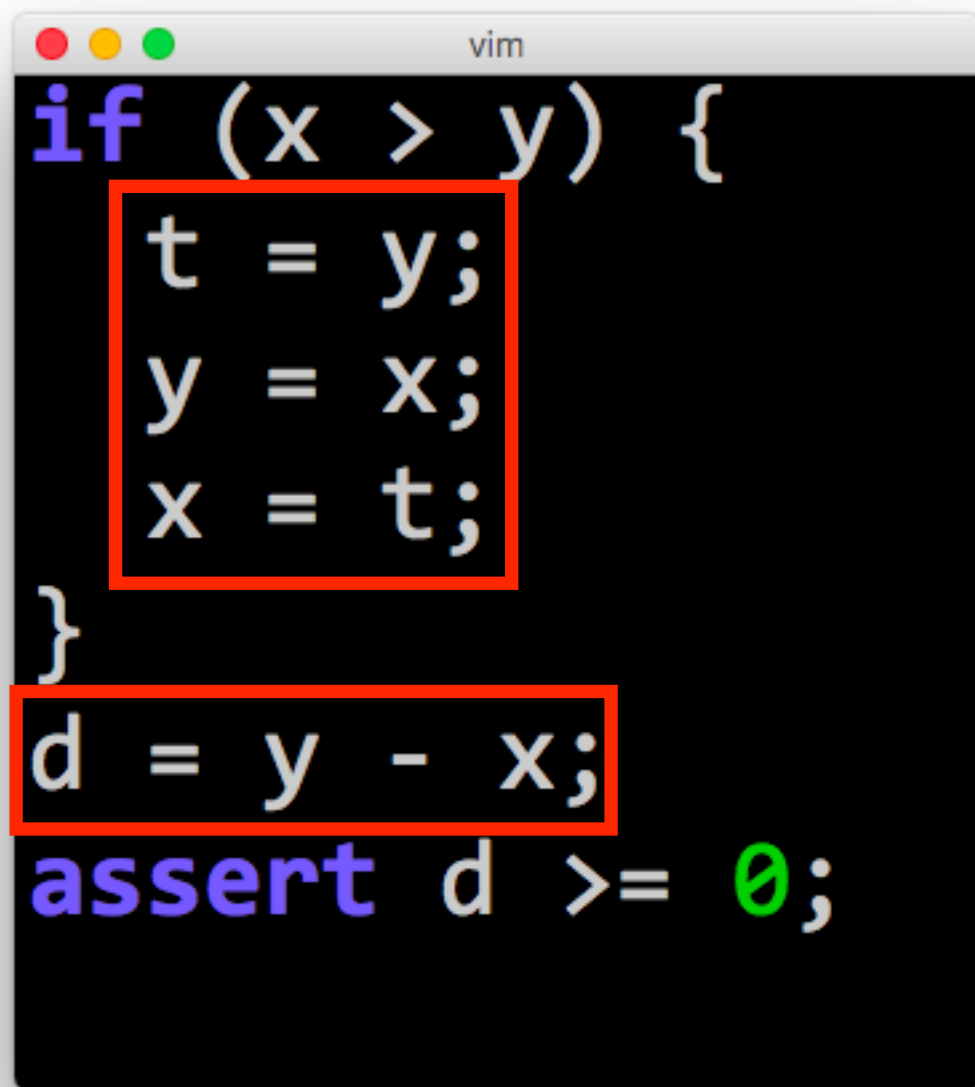
```
if (x > y) {  
    t = y;  
    y = x;  
    x = t;  
}  
d = y - x;  
assert d >= 0;
```



```
vim
if (x > y) {
    t = y;
    y = x;
    x = t;
}
d = y - x;
assert d >= 0;
```

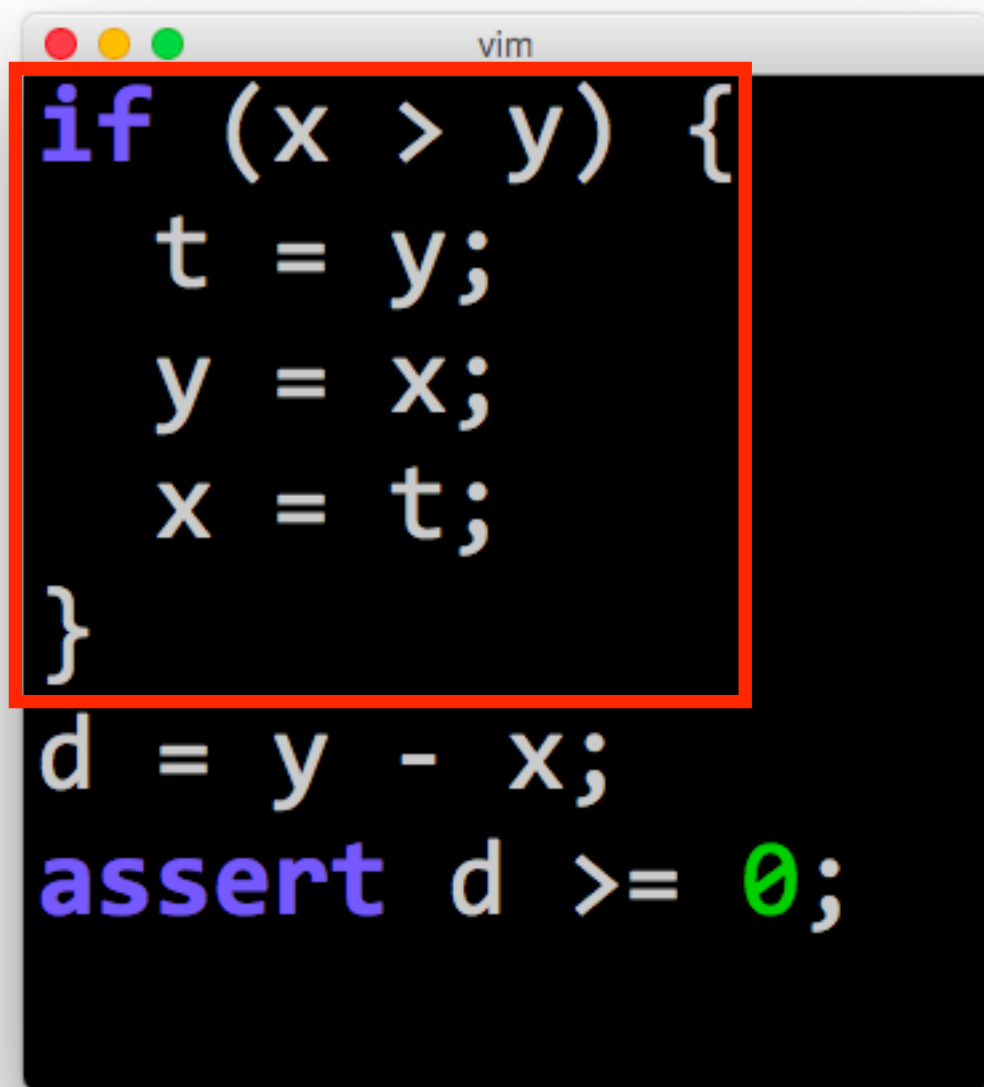


$\text{let } (x, y, t) =$   
    if  $x > y$  then  
        let  $t = y$  in  
        let  $y = x$  in  
        let  $x = t$  in  
         $(x, y, t)$   
    else  
         $(x, y, t)$  in  
let  $d = x - y$   
in  $d \geq 0$




```
vim
if (x > y) {
    t = y;
    y = x;
    x = t;
}
d = y - x;
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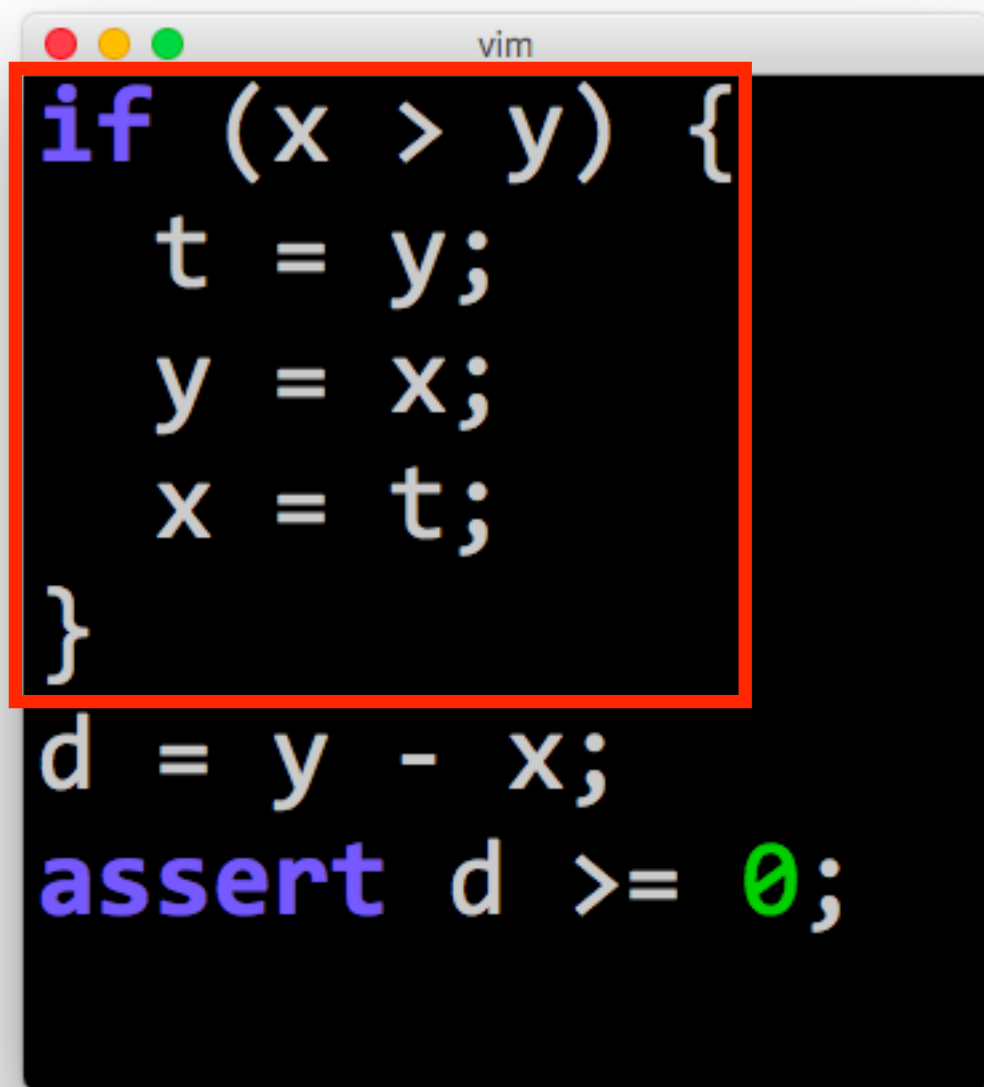
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        let  $x = t$  in  
         $(x, y, t)$   
    else  
         $(x, y, t)$  in  
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in  $d \geq 0$




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vim
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    x = t;
}
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```



$\text{let } (x, y, t) =$   
    if  $x > y$  then  
        let  $t = y$  in  
        let  $y = x$  in  
        let  $x = t$  in  
         $(x, y, t)$   
    else  
         $(x, y, t)$  in  
let  $d = x - y$   
    in  $d \geq 0$




```
vim
if (x > y) {
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$\text{let } (x, y, t) =$   
    if  $x > y$  then  
        let  $t = y$  in  
        let  $y = x$  in  
        let  $x = t$  in  
             $(x, y, t)$   
    else  
         $(x, y, t)$  in  
 $\text{let } d = x - y$   
    in  $d \geq 0$





```
vim
if (x > y) {
    t = y;
    y = x;
    x = t;
}
d = y - x;
assert d >= 0;
```

```
let  $(x, y, t) =$ 
    if  $x > y$  then
        let  $t = y$  in
        let  $y = x$  in
        let  $x = t$  in
         $(x, y, t)$ 
    else
         $(x, y, t)$  in
let  $d = x - y$ 
in  $d \geq 0$ 
```

```
vim
% Boolean variables
tff(1, axiom, ![X:$o]: (X | ~X)).

% Formulas as terms
tff(2, axiom, ![X:$o, Y:$o]: (impl(X, Y) <=> (~X | Y))).

% if-then-else
tff(3, axiom, ![X:$int, Y:$int]:
    (max(X, Y) = $ite($greatereq(X, Y), X, Y))).

tff(4, axiom, ![X:$int, Y:$int]: $ite(max(X, Y) = X,
    $greatereq(X, Y),
    $greatereq(Y, X))).

% let-in
tff(5, axiom, $let(array(I:$int) := $ite(I = 3, 5, array(I)),
    $sum(array(2), array(3)))).

tff(6, axiom, $let(a := b; b := a, f(a, b))).

~
~
~
~
```

