

GTD in Haskell

for-loops

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Revisiting `for`-loops

Refresher

`for`-loops have been a staple of imperative languages, starting with FORTRAN's `DO` loop in 1957.

Syntax

```
        DO label counter = first, last, step
            statements
label    statement
```

Example

```
6  DO 9, COUNTER = 1, 5, 1
7      WRITE (6,8) COUNTER
8      FORMAT( I2 )
9  CONTINUE
```

History

- 1957: FORTRAN provided the **DO** loop
- 1960: COBOL provided **PERFORM** verb with many options akin to for-loop
- 1964: BASIC gives us FOR-NEXT loops and PL/I offers first **break** in **LEAVE**
- 1968: Algol 68 gives us the first universal **for** loop as we know it today.

...

Mainstream languages today offer various types of for-loops:

- numeric/traditional **for**-loops (most common, low-level)
- iterator-based **for**-loops (most common, high-level)
- vectorized **for**-loops (niche, e.g. FORTRAN 95)
- compound **for**-loops (niche, e.g. ALGOL 68)

Low-level for vs higher-level "foreach" (Java)

```
1 Pet pets[] = {  
2     new Pet("Garfield"),  
3     new Pet("Odie"),  
4     new Pet("Nermal")  
5 };
```

```
1 // low-level (numeric) for-loop  
2 for (int i = 0; i < pets.length; i++) {  
3     System.out.println(pets[i]);  
4 }
```

```
1 // slightly higher-level "foreach"-style loop  
2 for (int pet : pets) {  
3     System.out.println(pet);  
4 }
```

for examples (Ruby)

```
pets = [  
  Pet.new("Garfield"),  
  Pet.new("Odie"),  
  Pet.new("Nermal")  
]  
  
// language construct  
for pet in pets  
  puts pet  
end  
  
// non-language-based variations  
pets.size.times { |idx| puts pets[idx] }  
pets.each_index { |idx| puts pets[idx] }  
pets.each { |pet| puts pet }
```


Basic intuition

iteration (*noun*): the action or a process of iterating or repeating: such as

- a procedure in which repetition of a sequence of operations yields results successively closer to a desired result
- the repetition of a sequence of computer instructions a specified number of times or until a condition is met — compare to recursion

Related:

- recursion
- **while** loops

Levels of abstraction

Application developers

Build software for specialized domains not generalized computing, e.g.:

- publishing
- investment management
- cloud infrastructure management
- generating computer music
- financial planning
- dashboarding data
- pretty much every piece of software is "domain-specific"

Focus on the *what* not the *how*

When building programs with `for`, `while`, `do` loops, we are focusing on the *how*.

We need to remove:

- mutable state
- loops

But how?

Higher-level primitives

- `filter`
- `map`
- `reduce`
- `apply effects`
- `composition of the above`

filter in Ruby

```
1  pets = [  
2    { name: "garfield", type: :cat },  
3    { name: "odie", type: :dog },  
4    { name: "nermal", type: :cat }  
5  ]  
6  cats = pets.select { |p| p[:type] == :cat }  
7  # cats contains garfield and nermal  
8  
9  dogs = pets.select { |p| p[:type] == :dog }  
10 # dogs contains odie
```

map in Ruby

```
1 pet_names = pets.collect do |p|
2   p[:name]
3 end
4 # => ["garfield", "odie", "nermal"]
5
6 pet_types = pets.collect do |p|
7   p[:type]
8 end
```

In many mainstream languages like Ruby these data transformations are available on basic data structures like **Array**-s, **Hash**-es, etc.

reduce in Ruby

```
1 cat_names_string = cat_names.join(', ')
2 # => "garfield, nermal"
```

In Ruby (and most mainstream languages) these *reduction* methods are monomorphic. This means we can't **abstract** in the general case!

apply [side] effects in Ruby

```
1 # typically found inline in controller or  
  ↪ view code in Rails  
2 pets.each do |p|  
3   stream_via_websockets_to_client p.to_json  
4 end
```

In Ruby we have **side** effects not *explicit* effects. This means we compromise on:

- testability
- composition

`stream_via_websockets_to_client` isn't abstracted therefore can't test without mocking which isn't high assurance testing technique.

Composition of filter, map, reduce, apply, part 1

```
1  cats = pets.select { |p| p[:type] == :cat }
2  cat_names = cats.collect { |c| c[:name] }
3  # => ["garfield", "nermal"]
4
5  puts cat_names.join(', ')
6  # print to stdout "garfield, nermal"
```

In mainstream languages we **must**:

- use `#select`, `#collect`, `#join`, etc. - the *pure* versions
 - not `#select!`, `#collect!`, or manual mutation.
- not include side effects in the given blocks
- not use `#each`
- no language support to ensure expressions are used not statements with side effects

State of the `union(filter, map, reduce, apply)` in mainstream languages

- No general abstraction (specific to data structures like **Array**, **Hash**, etc.)
- Still hard to test [side] effects in isolation
- Some composition possible, but not general case
- Lack of reuse since **Enumerable** interface is implemented per data type wholesale

Can we do better?

Assume we have the following Haskell defined

```
1  import Relude
2
3  data Pet
4      = MkPet { petName :: Text
5                , petAnimal :: Animal }
6              deriving (Show)
7  data Animal = Cat | Dog deriving (Show)
8  mkCat name = MkPet name Cat
9  mkDog name = MkPet name Dog
10
11  pets =
12      [ mkCat "garfield"
13        , mkDog "odie"
14        , mkCat "nermal" ]
```

filter in Haskell

```
1  isCat (MkPet _ Cat) = True
2  isCat _              = False
3
4  isDog (MkPet _ Dog) = True
5  isDog _              = False
6
7  cats = filter isCat pets
8  dogs = filter isDog pets
```

map in Haskell

```
1  -- Using (<&>) which is flipped fmap aka (<$>)
2  -- >>> :t (<&>)
3  --(<&>) :: Functor f => f a -> (a -> b) -> f b
4  -- >>> catNames
5  -- ["garfield","nermal"]
6  catNames = cats <&> petName
7
8  -- >>> :t (<$>) -- alias for fmap
9  --(<$>) :: Functor f => (a -> b) -> f a -> f b
10 -- >>> dogNames
11 -- ["odie"]
12 dogNames = petName <$> dogs
```

reduce in Haskell

```
1  import Data.List.NonEmpty (NonEmpty)
2  import Data.Semigroup (Semigroup (sconcat))
3
4  -- >>> :i NonEmpty
5  -- data NonEmpty a = a :| [a]
6  --
7  -- >>> :i Semigroup
8  -- class Semigroup a where
9  --   (< >) :: a -> a -> a
10 --   sconcat :: NonEmpty a -> a
11 --   stimes :: Integral b => b -> a -> a
12 --   {-# MINIMAL (< >) #-}
13 reduceS :: Semigroup a => NonEmpty a -> a
14 reduceS = sconcat
```


reduce in Haskell (usage)

```
1  namesString = mconcat . intersperse ", "
2
3  -- >>> catNamesString
4  -- "garfield, nermal"
5  catNamesString = namesString catNames
6
7  -- >>> dogNamesString
8  -- "odie"
9  dogNamesString = namesString dogNames
```

ap-ply in Haskell

```
1 import Data.Traversable (traverse)
2 import System.IO (FilePath, IOMode (..),
   ↪ hPutStrLn, withFile)
3
4 writePet :: Handle -> Pet -> IO Pet
5 writePet h (MkPet name animal) = hPutStrLn h
   ↪ $ textRep
6   where textRep = show animal <> " { " <>
   ↪       show name <> "}"
7
8 writePets :: FilePath -> [Pet] -> IO ()
9 writePets path = withFile path WriteMode $
   ↪ void $ traverse (writePet h)
```

Composition in Haskell

- With Haskell's equational reasoning we get substitution for free
- Use general abstract notion of **Functor** to transform underlying data
- Implementing **fmap** for your datatype in the **Functor** instance will net you many **Functor** based functions for free
- Similarly for **Monoid**, **Traversable**, **Applicative**, etc.
- Use general abstract notion of **Monoid** to reduce a list of values to one (e.g. sum, concatenate)
- Compose effectful actions in the **same** effect context (in this case IO, below):

```
1  main = do
2      writePets "/tmp/cats.txt" cats
3      writePets "/tmp/dogs.txt" dogs
```

Results (Haskell)

```
$ cat /tmp/cats.txt  
Cat{ "garfield"}  
Cat{ "nermal"}  
  
$ cat /tmp/dogs.txt  
Dog{ "odie"}
```

What's Next?

Encoding **for** in Haskell using more generalized but expressive types:

```
1  -- In Data.Traversable
2  for
3      -- constraints
4      :: (Traversable t, Applicative f)
5      -- "traversable" collection of elements
6      => t a
7      -- given an element of the collection
8      --   ↪ produce a b in context f
9      --   ↪ (a -> f b)
10     -- produce in the effect context f the
11     --   ↪ resultant traversable of b's
12     -> f (t b)
```

Just when you thought you were safe...

The monadic form:

```
1  -- In Control.Monad
2  forM
3      -- constraints
4      :: (Traversable t, Monad m)
5      -- "traversable" collection of a's
6      => t a
7      -- monadic function to perform for each
8      --   ↪ element
9      -- produce new "traversable" in monadic
10     --   ↪ effect context m
10    -> m (t b)
```

Most specific forms of for

```
1  -- In import Data.Foldable
2  foldMap
3      -- constraints
4      :: (Foldable t, Monoid m)
5      -- function to convert a collection element
6      --   ↪ to monoidal value
7      => (a -> m)
8      -- collection of "foldable" of a elements
9      -> t a
10     -- produce reduced monoidal result
11     -> m
12
13 scorePet (MkPet _ Cat) = Sum 1
14 scorePet (MkPet _ Dog) = Sum 2
15
16 example = foldMap scorePet pets
```