Trabalho Individual INE5429 Segurança em Computação AES e S-AES

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1 **S-AES**

Tabela 1 – Dados de entrada

-	Decimal	Hex
Texto Claro	10022	0x2726
Chave secreta	8022	0x1F56

Tabela 2 — Representação de bits - Texto claro

0010	0010
0111	0110

Tabela 3 – Representação de nibbles - Texto claro

Tabela 4 – Representação de bits - Chave secreta

w_0	w_1
0001 1111	0101 0110

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1.1 Expansão da chave

 $= 0100 \ 0011$

 $= 0000 \ 1011$

$\mathbf{w_2} = w_0$ = w_0 = 0001 1111	$\oplus g(w_1)$ $\oplus Rcon(1)$ $\oplus Rcon(1)$ $\oplus Rcon(1)$ $\oplus Rcon(1)$ $\oplus 1000\ 0000$ $\oplus 0000\ 0001$ $\oplus 0000\ 0001$	$\oplus SubNib(RotNib(w_1))$ $\oplus SubNib(RotNib(0101\ 0110))$ $\oplus SubNib(0110\ 0101)$ $\oplus 1000\ 0001$ $\oplus 1000\ 0001$
$\mathbf{w_3} = w_2 \oplus w_1$ = 0001 1110 \oplus 0101 0110 = 0100 1000		
$\begin{aligned} \mathbf{w_4} &= w_2 \\ &= 0001 \ 1110 \\ &= 0100 \ 0011 \end{aligned}$	$⊕ g(w_3)$ $⊕ Rcon(2)$ $⊕ Rcon(2)$ $⊕ Rcon(2)$ $⊕ Rcon(2)$ $⊕ (10000 \mod 10011) 0000$ $⊕ 0011 0000$ $⊕ 0101 1101$ $⊕ 0101 1101$	$\oplus SubNib(RotNib(w_3))$ $\oplus SubNib(RotNib(0100\ 1000))$ $\oplus SubNib(1000\ 0100)$ $\oplus 0110\ 1101$ $\oplus 0110\ 1101$ $\oplus 0110\ 1101$
$\mathbf{w_5} = w_4$	$\oplus w_3$	

Tabela5 – Expansão - Chave secreta

 $\oplus\ 0100\ 1000$

w_0	w_1	w_2	w_3	w_4	w_5
0001 1111	0101 0110	0001 1110	0100 1000	0100 0011	0000 1011
1 F	5 6	1 E	4 8	4 3	0 B

1.2 Encriptação

1.2.1 Rodada 0

1.2.1.1 Incluir chave da rodada

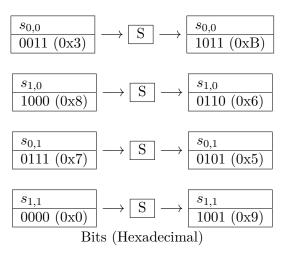
Tabela 6 – Inclusão da chave

0010 (0x2)	0010 (0x2)		0001 (0x1)	0101 (0x5)	_	0011 (0x3)	0111 (0x7)		
0111 (0x7)	0110 (0x6)	θ	1111 (0xF)	0110 (0x6)	_	1000 (0x8)	0000 (0x0)		
Bits (Hexadecimal)									

1.2.2 **Rodada 1**

1.2.2.1 Substituição de nibble

Tabela 7 – Substituição de nibble



1.2.2.2 **Deslocar linhas**

Tabela 8 – Deslocamento de linha

1011 (0xB)	0101 (0x5)		1011 (0xB)	0101 (0x5)					
0110 (0x6)	1001 (0x9)		1001 (0x9)	0110 (0x6)					
Bits (Hexadecimal)									

1.2.2.3 Misturar colunas

$$S'_{0,0} = S_{0,0} \oplus (4 \cdot S_{1,0})$$

$$= 1011 \oplus ((x^2 \cdot (x^3 + 1)) \mod (x^4 + x + 1))$$

$$= 1011 \oplus ((x^5 + x^2) \mod (x^4 + x + 1))$$

$$= 1011 \oplus (-x)$$

$$= 1011 \oplus (0)$$

$$= 1001$$

$$S'_{1,0} = (4 \cdot S_{0,0}) \oplus S_{1,0}$$

$$= ((x^2 \cdot (x^3 + x + 1)) \mod (x^4 + x + 1)) \oplus 1001$$

$$= ((x^5 + x^3 + x^2) \mod (x^4 + x + 1)) \oplus 1001$$

$$= (x^3 - x) \oplus 1001$$

$$= (x^3 + x) \oplus 1001$$

$$= (x^3 + x) \oplus 1001$$

$$= 1010 \oplus 1001$$

$$= 0011$$

$$S'_{0,1} = S_{0,1} \oplus (4 \cdot S_{1,1})$$

$$= 0101 \oplus (x^2 \cdot (x^2 + x)) \mod (x^4 + x + 1)$$

$$= 0101 \oplus (x^4 + x^3) \mod (x^4 + x + 1)$$

$$= 0101 \oplus (x^3 - x - 1)$$

$$= 0101 \oplus 1011$$

$$= 1110$$

$$S'_{1,1} = (4 \cdot S_{0,1}) \oplus S_{1,1}$$

$$= (x^2 \cdot (x^2 + 1) \mod (x^4 + x + 1) \oplus 0110)$$

$$= (x^4 + x^2) \mod (x^4 + x + 1) \oplus 0110$$

$$= (x^2 - x - 1) \oplus 0110$$

$$= 0111 \oplus 0110$$

Tabela 9 – Saída: Mistura de colunas

= 0001

1001 (0x9)	1110 (0xE)
0011 (0x3)	0001 (0x1)

Bits (Hexadecimal)

1.2.2.4 Incluir chave da rodada

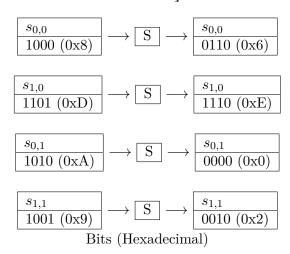
Tabela 10 – Inclusão da chave

1001 (0x9)	1110 (0xE)		0001 (0x1)	0100 (0x4)	_	1000 (0x8)	1010 (0xA)		
0011 (0x3)	0001 (0x1)	\oplus	1110 (0xE)	1000 (0x8)	_	1101 (0xD)	1001 (0x9)		
Bits (Hexadecimal)									

1.2.3 Rodada 2

1.2.3.1 Substituição de nibble

Tabela 11 – Substituição de nibble



1.2.3.2 **Deslocar linhas**

Tabela 12 – Deslocamento de linha

0110 (0x6)	0000 (0x0)		0110 (0x6)	0000 (0x0)				
1110 (0xE)	0010 (0x2)		0010 (0x2)	1110 (0xE)				
Bits (Hexadecimal)								

1.2.3.3 Incluir chave da rodada

Tabela 13 – Inclusão da chave

0110 (0x6)	0000 (0x0)		0100 (0x4)	0000 (0x0)	_	0010 (0x2)	0000 (0x0)	
0010 (0x2)	1110 (0xE)	\oplus	0011 (0x3)	1011 (0xB)	_	0001 (0x1)	0101 (0x5)	
Bits (Hexadecimal)								

1.2.3.4 Saída: Texto Cifrado

Binário: 0010 0001 0000 0101

Hexadecimal: 0x2105

1.3 Decriptação

1.3.1 Rodada 0

1.3.1.1 Incluir chave da rodada

Tabela 14 – Inclusão da chave

0010 (0x2)	0000 (0x0)		0100 (0x4)	0000 (0x0)	_	0110 (0x6)	0000 (0x0)		
0001 (0x1)	0101 (0x5)		0011 (0x3)	1011 (0xB)	_	0010 (0x2)	$1110 \; (0xE)$		
Bits (Hexadecimal)									

1.3.2 **Rodada 1**

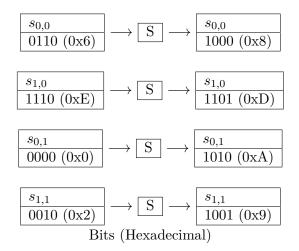
1.3.2.1 **Deslocar linhas invertidas**

Tabela 15 – Deslocamento de linha invetida

0110 (0x6)	0000 (0x0)		0110 (0x6)	0000 (0x0)								
0010 (0x2)	1110 (0xE)		1110 (0xE)	0010 (0x2)								
Bits (Hexadecimal)												

1.3.2.2 Substituição de nibble invertido

Tabela 16 – Substituição de nibble invertido



1.3.2.3 Incluir chave da rodada

Tabela 17 - Inclusão da chave

1000 (0x8)	1010 (0xA)		0001 (0x1)	0100 (0x4)	_	1001 (0x9)	1110 (0xE)
1101 (0xD)	1001 (0x9)		1110 (0xE)	1000 (0x8)	_	0011 (0x3)	0001 (0x1)
		_	Bits (Hexa	adecimal)			

1.3.2.4 Misturar colunas invertidas

$$\mathbf{S'_{0,0}} = (9 \cdot S_{0,0}) \oplus (2 \cdot S_{1,0})$$

$$= ((x^3 + 1) \cdot (x^3 + 1)) \mod (x^4 + x + 1) \oplus (x \cdot (x + 1))$$

$$= ((x^6 + 2x^3 + 1) \mod (x^4 + x + 1)) \oplus (x^2 + x)$$

$$= (x^3 - x^2 + 1) \oplus (x^2 + x)$$

$$= 1101 \oplus 0110$$

$$= 1011$$

$$\mathbf{S'_{1,0}} = (2 \cdot S_{0,0}) \oplus (9 \cdot S_{1,0})$$

$$= (x \cdot (x^3 + 1)) \oplus ((x^3 + 1) \cdot (x + 1))$$

$$= (x \cdot (x^3 + 1)) \oplus ((x^3 + 1) \cdot (x + 1))$$

$$= (x^4 + x) \mod (x^4 + x + 1) \oplus (x^4 + x^3 + x + 1) \mod (x^4 + x + 1)$$

$$= -1 \oplus x^3$$

$$= 0001 \oplus 1000$$

$$= 1001$$

$$\mathbf{S}'_{0,1} = (9 \cdot S_{0,1}) \oplus (2 \cdot S_{1,1})$$

$$= ((x^3 + 1) \cdot (x^3 + x^2 + x) \mod (x^4 + x + 1)) \oplus x \cdot 1$$

$$= (x^6 + x^5 + x^4 + x^3 + x^2 + x) \mod (x^4 + x + 1) \oplus x$$

$$= -x^2 - x - 1 \oplus x$$

$$= 0111 \oplus 0010$$

$$= 0101$$

$$\mathbf{S}'_{1,1} = (2 \cdot S_{0,1}) \oplus (9 \cdot S_{1,1})$$

$$= (x \cdot (x^3 + x^2 + x)) \mod (x^4 + x + 1) \oplus (x^3 + 1) \cdot 1$$

$$= (x^4 + x^3 + x^2) \mod (x^4 + x + 1) \oplus (x^3 + 1)$$

$$= (x^3 + x^2 + x + 1) \oplus (x^3 + 1)$$

$$= 1111 \oplus 1001$$

$$= 0110$$

Tabela 18 – Saída: Mistura de colunas invertidas

1011 (0xB)	0101 (0x5)
1001 (0x9)	0110 (0x6)

Bits (Hexadecimal)

1.3.3 Rodada 2

1.3.3.1 **Deslocar linhas invertidas**

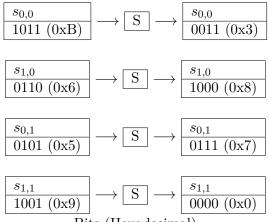
Tabela 19 – Deslocamento de linha invertida

1011 (0xB)	0101 (0x5)	 1011 (0xB)	(/
1001 (0x9)	0110 (0x6)	0110 (0x6)	1001 (0x9)

Bits (Hexadecimal)

1.3.3.2 Substituição de nibble invertido

Tabela 20 – Substituição de nibble invertido



Bits (Hexadecimal)

1.3.3.3 Incluir chave da rodada

Tabela 21 – Inclusão da chave

0011 (0x3)	0111 (0x7)		0001 (0x1)	0101 (0x5)	_	0010 (0x2)	0010 (0x2)
1000 (0x8)	0000 (0x0)	Ð	1111 (0xF)	0110 (0x6)	_	0111 (0x7)	0110 (0x6)
			Bits (Hexa	adecimal)			

1.3.3.4 Saída: Texto Decifrado

Binário: 0010 0111 0010 0110

Hexadecimal: 0x2726

Decimal: 10022

2 **AES**

Para a simulação do AES 128
bits foi utilizada como base a planilha disponibilizada por (NAYUKI, 2016).

Tabela 22 – Dados de entrada

Texto Claro	0x000000000000000000000000000000000000
Chave	0x0000000000000000000000000001F56
Texto Cifrado	0x26BC0CCCE44C0066DE46EF3C44EBC555

2.1 Encriptação

Tabela 23 – Execução AES
(18bits) - Encriptação

Rodadas	Iníc	io da	Rod	ada	Apó	s Sul	$_{ m Byt}$	es	Apó	s Shi	ftRov	vs	Apó	s Mi	xColu	ımns	Cha	ve de	Rod	ada
Rodada 0	00	00	00	00													00	00	00	00
	00	00	00	00													00	00	00	00
	00	00	00	27													00	00	00	1F
	00	00	00	26													00	00	00	56
Rodada 1	00	00	00	00	63	63	63	63	63	63	63	63	51	07	63	63	62	62	62	62
	00	00	00	00	63	63	63	63	63	63	63	63	51	$_{\mathrm{CF}}$	63	63	C0	C0	C0	C0
	00	00	00	38	63	63	63	07	63	07	63	63	35	AB	63	63	В1	B1	B1	AE
	00	00	00	70	63	63	63	51	51	63	63	63	07	07	63	63	63	63	63	35
Rodada 2	33	65	01	01	C3	4D	7C	7C	C3	4D	7C	7C	03	7A	FA	A1	DA	B8	DA	B8
	91	0F	A3	A3	81	76	0A	0A	76	0A	0A	81	5A	C6	CA	FB	24	E4	24	E4
	84	1A	D2	CD	5F	A2	B5	BD	B5	BD	5F	A2	0C	E3	0D	07	27	96	27	89
	64	64	00	56	43	43	63	B1	B1	43	43	63	E4	E6	57	61	C9	AA	C9	FC
Rodada 3	D9	C2	20	19	35	25	B7	D4	35	25	B7	D4	7F	F3	F6	2B	B7	0F	D5	6D
Ttoddad 0	7E	22	EE	1F	F3	93	28	C0	93	28	C0	F3	62	86	0D	9E	83	67	43	A7
	2B	75	2A	8E	F1	9D	E5	19	E5	19	F1	9D	95	4C	F5	1B	97	01	26	AF
	2D	4C	9E	9D	D8	29	0B	5E	5E	D8	29	0B	95	F5	A1	1F	A5	0F	C6	3A
Rodada 4	C8	FC	23	46	E8	B0	26	5A	E8	B0	26	5A	81	83	20	C1	E3	EC	39	54
22044444	E1	E1	4E	39	F8	F8	2F	12	F8	2F	12	F8	96	66	B6	0A	FA	9D	DE	79
	02	4D	D3	B4	77	E3	66	8D	66	8D	77	E3	9D	92	AD	EB	17	16	30	9F
	30	FA	67	25	04	2D	85	3F	3F	04	2D	85	C3	61	55	E4	99	96	50	6A
Rodada 5	62	6F	19	95	AA	A8	D4	2A	AA	A8	D4	2A	19	A8	2F	90	45	A9	90	C4
Teodada o	6C	FB	68	73	50	0F	45	8F	0F	45	8F	50	4F	31	3B	00	21	BC	62	1B
	8A	84	9D	74	7E	5F	5E	92	5E	92	7E	5F	32	0B	1F	79	15	03	33	AC
	5A	F7	05	8E	BE	68	6B	19	19	BE	68	6B	86	53	46	A7	B9	2F	7F	15
Rodada 6	5C	01	BF	54	4A	7C	08	20	4A	7C	08	20	35	C8	26	D8	CA	63	F3	37
Teodada o	6E	8D	59	1B	9F	5D	CB	AF	5D	CB	AF	9F	54	81	12	47	BO	0C	6E	75
	27	08	2C	D5	CC	30	71	03	71	03	CC	30	AC	$^{2}\mathrm{E}$	14	E9	4C	4F	7C	D0
	3F	7C	39	B2	75	10	12	37	37	75	10	12	9C	A6	5B	EB	A5	8A	F5	E0
Rodada 7	FF	AB	D5	EF	16	62	03	DF	16	62	03	DF	A5	F4	F3	15	17	74	87	В0
	E4	8D	7C	32	69	5D	10	23	5D	10	23	69	48	66	0C	СЗ	C0	CC	A2	D7
	E0	61	68	39	E1	$_{\mathrm{EF}}$	45	12	45	12	E1	$_{\mathrm{EF}}$	BC	60	6A	44	AD	E2	9E	4E
	39	2C	AE	0B	12	71	E4	2B	2B	12	71	E4	74	80	25	2F	3F	B5	40	A0
Rodada 8	B2	80	74	A5	37	CD	92	06	37	CD	92	06	4D	62	3E	05	99	ED	6A	DA
	88	AA	AE	14	C4	AC	E4	FA	AC	E4	FA	C4	DD	04	76	ED	EF	23	81	56
	11	82	F4	0A	82	13	$_{\mathrm{BF}}$	67	BF	67	82	13	6B	29	D6	33	4D	AF	31	7F
	4B	35	65	8F	ВЗ	96	4D	73	73	B3	96	4D	AC	B2	E2	47	D8	6D	2D	8D
Rodada 9	D4	8F	54	DF	48	73	20	9E	48	73	20	9E	3F	E5	0C	8C	33	DE	B4	6E
	32	27	F7	BB	23	CC	68	EA	CC	68	EA	23	18	4A	73	9E	3D	1E	9F	C9
	26	86	E7	4C	F7	44	94	29	94	29	F7	44	2B	E4	86	B0	10	$_{\mathrm{BF}}$	8E	F1
	74	$_{ m DF}$	$_{\mathrm{CF}}$	CA	92	9E	8A	74	74	92	9E	8A	68	EB	5A	D1	8F	E2	$_{\mathrm{CF}}$	42
Rodada 10	0C	3B	B8	E2	FE	E2	6C	98	FE	E2	6C	98					D8	06	B2	DC
	25	54	EC	57	3F	20	CE	5B	20	CE	5B	3F					9C	82	1D	D4
	3В	5B	08	41	E2	39	30	83	30	83	E2	39					3C	83	0D	FC
	E7	09	95	93	94	01	2A	DC	DC	94	01	2A					10	F2	3D	7F
Saída	26	E4	DE	44																
	BC	4C	46	EB																
	0C	00	EF	C5																
	CC	66	3C	55																

2.2 **Decriptação**

Tabela 24 – Execução AES(18bits) - Decriptação

Rodadas	Fir	nal da	roda	ada		Shift	Rows	i		Subl	Bytes		K	ley So	chedu	N	MixColumns			
Saída	00	00	00	00																
	00	00	00	00																
	00	00	00	27																
	00	00	00	26																
Rodada 10	51	07	63	63	63	63	63	63	00	00	00	00	00	00	00	00	63	63	63	63
	51	$_{\mathrm{CF}}$	63	63	63	63	63	63	00	00	00	00	00	00	00	00	63	63	63	63
	35	AB	63	63	63	63	63	07	00	00	00	38	00	00	00	1F	63	07	63	63
	07	07	63	63	63	63	63	51	00	00	00	70	00	00	00	56	51	63	63	63
Rodada 9	03	7A	FA	A1	C3	4D	7C	7C	33	65	01	01	62	62	62	62	С3	4D	7C	7C
	5A	C6	CA	FB	81	76	0A	0A	91	0F	A3	A3	C0	C0	C0	C0	76	0A	0A	81
	0C	E3	0D	07	5F	A2	B5	$^{\mathrm{BD}}$	84	1A	D2	$^{\rm CD}$	B1	B1	B1	AE	B5	BD	5F	A2
	E4	E6	57	61	43	43	63	B1	64	64	00	56	63	63	63	35	B1	43	43	63
Rodada 8	7F	F3	F6	$^{2\mathrm{B}}$	35	25	В7	D4	D9	C2	20	19	DA	В8	DA	B8	35	25	В7	D4
	62	86	0D	9E	F3	93	28	C0	7E	22	$_{\rm EE}$	1F	24	E4	24	E4	93	28	C0	F3
	95	4C	F5	1B	F1	9D	E_5	19	2B	75	2A	8E	27	96	27	89	E_5	19	F1	9D
	95	F5	A1	1F	D8	29	0B	5E	2D	4C	9E	9D	C9	AA	C9	FC	5E	D8	29	0B
Rodada 7	81	83	20	C1	E8	B0	26	5A	C8	FC	23	46	В7	0F	D5	6D	E8	B0	26	5A
	96	66	$_{\rm B6}$	0A	F8	F8	2F	12	E1	E1	4E	39	83	67	43	A7	F8	2F	12	F8
	9D	92	$^{\mathrm{AD}}$	$_{\rm EB}$	77	E3	66	8D	02	4D	D3	B4	97	01	26	AF	66	8D	77	E3
	СЗ	61	55	E4	04	2D	85	3F	30	FA	67	25	A5	0F	C6	3A	3F	04	2D	85
Rodada 6	19	A8	2F	90	AA	A8	D4	2A	62	6F	19	95	E3	EC	39	54	AA	A8	D4	2A
,	4F	31	3B	00	50	0F	45	8F	6C	FB	68	73	FA	9D	DE	79	0F	45	8F	50
	32	0B	1F	79	7E	5F	5E	92	8A	84	9D	74	17	16	30	9F	5E	92	7E	5F
	86	53	46	A7	BE	68	6B	19	5A	F7	05	8E	99	96	50	6A	19	BE	68	6B
Rodada 5	35	C8	26	D8	4A	7C	08	20	5C	01	BF	54	45	A9	90	C4	4A	7C	08	20
	54	81	12	47	9F	5D	СВ	AF	6E	8D	59	1B	21	BC	62	1B	5D	CB	AF	9F
	AC	2E	14	E9	CC	30	71	03	27	08	2C	D5	15	03	33	AC	71	03	CC	30
D 1 1 1	9C	A6	5B	EB	75	10	12	37	3F	7C	39	B2	B9	2F	7F	15	37	75	10	12
Rodada 4	A5	F4	F3	15 Co	16	62 5 D	03	DF	FF	AB	D5	EF	CA	63	F3	37	16 5D	62	03	DF
	48	66	OC	C3	69	5D	10	23	E4	8D	7C	32	B0	0C	6E	75 Do	5D	10	23	69 EE
	BC	60	6A	44	E1	EF	45	12	E0	61	68	39 op	4C	4F	7C	D0	45	12	E1	EF
D - J - J - 9	74	80	25	2F	12	71 CD	E4	2B	39	2C	AE 74	0B	A5	8A	F5	E0	2B	12	71	E4
Rodada 3	4D DD	62	3E	05 ED	37 C4	CD	92	06 EA	B2	80		A5	17	74 CC	87	B0	37	CD	92	06
	6B	04 29	76 D6	ED 33	82	AC 13	E4 BF	FA 67	88 11	AA 82	AE F4	14 0A	C0 AD	E2	A2 9E	D7 4E	AC BF	E4 67	FA 82	C4 13
	AC	B2	E2	47	B3	96	4D	73	4B	35	65	8F	3F	B5	40	A0	73	B3	96	4D
Rodada 2	3F	E5	0C	8C	48	73	20	9E	D4	8F	54	DF	99	ED	6A	DA	48	73	20	9E
1touaua 2	18	4A	73	9E	23	CC	68	EA	32	27	F7	BB	EF	23	81	56	CC	68	EA	23
	2B	E4	86	B0	F7	44	94	29	26	86	E7	4C	4D	AF	31	7F	94	29	F7	44
	68	EB	5A	D1	92	9E	8A	74	74	DF	CF	CA	D8	6D	2D	8D	74	92	9E	8A
Rodada 1	FE	E2	6C	98	FE	E2	6C	98	0C	3B	B8	E2	33	DE	B4	6E	1.4	32	JL	071
Loudda 1	20	CE	5B	3F	3F	20	CE	5B	25	54	EC	57	3D	1E	9F	C9				
	30	83	E2	39	E2	39	30	83	3B	5B	08	41	10	BF	8E	F1				
	DC	94	01	2A	94	01	2A	DC	E7	09	95	93	8F	E2	CF	42				
Rodada 0	26	E4	DE	44	0.1	01	211	20	L.	00	55	50	D8	06	B2	DC				
2.oudda 0	BC	4C	46	EB									9C	82	1D	D4				
	0C	00	EF	C5									3C	83	0D	FC				
	CC	66	3C	55									10	F2	3D	7F				
	LOO	00	30	55									10	1. 7	317	11				

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