# $Part\ II-Graph\ Theory$

Based on lectures by Prof. P. Russell Notes taken by Bhavik Mehta

Michaelmas 2017

# 0 Introduction

- 0.1 Preliminary
- 0.2 Informal definitions
- 0.3 Where do such structures arise?

**Theorem** (Schur's Theorem). Let n be a positive integer. Then if p is a sufficiently large positive integer, whenever  $\{1, 2, \ldots, p\}$  is partitioned into n parts, we can solve a+b=c with a,b,c all in some part.

# 1 Ramsey Theory

**Theorem** (Schur's Theorem reformulated). Let k be a positive integer. Then there is a positive integer n such that if the set  $[n] = \{1, 2, ..., n\}$  is coloured with k colours, we can find a, b, c with a + b = c and a, b, c the same colour.

**Proposition 1.** Let k be a positive integer. Then there is a positive integer n such that whenever the edges of  $K_n$  are coloured with k colours we can find a monochromatic triangle.

**Theorem 2** (Ramsey's Theorem). R(s,t) exists for all  $s,t \geq 2$ . Moreover, if s,t > 2 then  $R(s,t) \leq R(s-1,t) + R(s,t-1)$ .

Corollary 3. For all  $s, t \geq 2$ ,  $R(s, t) \leq 2^{s+t}$ , so  $R(s) \leq 4^s$ .

**Theorem 4** (Multicolour Ramsey Theorem). Let  $k \geq 1$  and  $s \geq 2$ . Then there exists some n such that whenever the edges of  $K_n$  are coloured with k colours, we can find a monochromatic  $K_s$ .

**Theorem 5** (Infinite Ramsey Theorem). Let  $k \geq 1$ . Whenever the edges of  $K_{\infty}$  are k-coloured, we have a monochromatic  $K_{\infty}$  subgraph.

Corollary 6. Any bounded sequence has a convergent subsequence.

#### 1.1 Basic Terminology

# 2 Extremal Graph Theory

### 2.1 Forbidden Subgraph Problem

**Theorem 7.** A graph is bipartite iff it contains no odd cycles.

**Theorem 8** (Mantel's Theorem). Let  $n \geq 3$ . Suppose |G| = n,  $e(g) \geq \lfloor \frac{n^2}{4} \rfloor$  and  $\triangle \not\subset G$ . Then  $G \cong K_{\lceil \frac{n}{2} \rceil, \lfloor \frac{n}{2} \rfloor}$ .

**Theorem 9** (Turán's Theorem). Let  $r \geq 2$  and  $|G| = n \geq r + 1$ . If  $e(G) \geq t_r(n)$  and  $K_{r+1} \not\subset G$  then  $G \cong T_r(n)$ .

Corollary 10. Let  $r \geq 2$ . As  $n \to \infty$ ,  $\operatorname{ex}(n; K_{r+1}) \sim (1 - \frac{1}{r})\binom{r}{2}$ .

**Theorem 11.** Let  $t \geq 2$ . Then  $\operatorname{ex}(n; K_{t,t}) = \mathcal{O}\left(n^{2-\frac{1}{t}}\right)$ .

**Theorem 12.** Let  $t \geq 2$ . Then  $z(n,t) = \mathcal{O}(n^{2-\frac{1}{t}})$ .

**Proposition 13.** Let H be a graph with at least one edge, and for  $n \ge |H|$ , let  $x_n = \frac{\operatorname{ex}(n;H)}{\binom{n}{2}}$ . Then  $(x_n)$  converges.

**Theorem 14** (Erdős-Stone Theorem). Let  $r, t \ge 1$  be integers, and let  $\epsilon > 0$  be real. Then  $\exists n_0$  such that  $\forall n \ge n_0$ ,

$$|G| = n, \ e(G) \ge \left(1 - \frac{1}{r} + \epsilon\right) \binom{n}{2} \implies K_{r+1}(t) \subset G.$$

Corollary 15. Let H be a graph with at least one edge. Then

$$ex(H) = 1 - \frac{1}{\chi(H) - 1}.$$

Corollary 16. For any infinite graph G,

$$ud(G) \in \left\{0, 1, \frac{1}{2}, \frac{2}{3}, \frac{3}{4}, \dots\right\}.$$

#### 2.2 Hamiltonian graphs

**Theorem 17** (Dirac's Theorem). Let  $|G| = n \ge 3$  and  $\delta(G) \ge \frac{n}{2}$ . Then G is Hamiltonian.

**Proposition 18.** Let G be a connected graph. Then

G Eulerian if and only if  $\forall v \in G$ , d(v) is even.

# 3 Graph Colouring

### 3.1 Planar Graphs

**Theorem 19** (Kuratowski's Theorem). Let G be a graph. Then G planar iff G contains no subdivision of  $K_5$  or  $K_{3,3}$ .

**Proposition 20.** Every tree of order at least 2 has a leaf.

**Proposition 21.** Let T be a tree,  $|T| = n \ge 1$ . Then e(T) = n - 1.

Proposition 22. Every tree is planar.

**Theorem 23** (Euler's Formula). Take G connected and planar. Take  $|G| = n \ge 1$ , e(G) = m with l faces. Then n - m + l = 2.

Corollary 24. Let G be planar,  $|G| = n \ge 3$ . Then  $e(G) \le 3n - 6$ .

**Proposition 25** (Six colour theorem). Any planar graph is 6-colourable.

**Theorem 26** (Five colour theorem). Any planar graph is 5-colourable.

**Theorem 27** (Four colour theorem). Any planar graph is 4-colourable.

#### 3.2 General Graphs

**Theorem 28** (Brooks' theorem). Let G be a connected graph that is neither complete nor an odd cycle. Then  $\chi(G) \leq \Delta(G)$ .

### 3.3 Graphs on surfaces

**Theorem 29** (Heawood's Theorem). Let S be a closed boundaryless surface of Euler characteristic  $E \leq 1$ . Then

$$\chi(S) \le \left\lfloor \frac{7 + \sqrt{49 - 24E}}{2} \right\rfloor.$$

#### 3.4 Edge Colouring

**Theorem 30** (Vizing's theorem). Let G be a graph. Then

$$\chi'(G) \le \Delta(G) + 1$$

# 4 Connectivity

#### 4.1 The Marriage Problem

**Theorem 31** (Hall's Marriage Theorem). Let G be a bipartite graph with bipartition (X, Y). Then G has a matching from X to Y iff G satisfies **Hall's condition**:

$$\forall A \subset X, |\Gamma(A)| \ge |A|$$

**Corollary 32** (Defect Hall). Let G be a bipartite graph with bipartition (X,Y) and let  $d \ge 1$ . Then G contains |X| - d independent edges if and only if  $\forall A \subset X$ ,  $|\Gamma(A)| \ge |A| - d$ .

Corollary 33 (Polyandrous Hall). Let G be a bipartite graph, bipartition (X,Y),  $d \geq 2$ . Then G contains a set of d|X| edges, each vertex in X in precisely d of them, each vertex in Y in at most one  $\iff \forall A \subset X, |\Gamma(A)| \geq d|A|$ .

### 4.2 Connectivity

**Theorem 34.** Let G be a graph and  $A, B \subset V(G)$ . Let

$$k = \min\{ |W| \mid W \text{ is an } AB\text{-separator} \}.$$

Then G contains k vertex-disjoint AB-paths.

**Corollary 35** (Menger's Theorem). Let G be an incomplete k-connected graph and let  $a, b \in V(G), a \neq b$ . Then G contains k independent ab-paths.

#### 4.3 Edge connectivity

Corollary 36 (Edge Menger). Let G be l-edge connected and  $a, b \in V(G)$  be distinct. Then G has l edge-disjoint ab-paths.

# 5 Probabilistic Techniques

#### 5.1 The Probabilistic Method

Theorem 37 (Erdős).

$$R(s) = \Omega(\sqrt{2}^s)$$

#### 5.2 Modifying a Random Graph

**Theorem 38.** If  $t \ge 2$  then  $z(n,t) = \Omega(n^{2-\frac{2}{t+1}})$ .

**Theorem 39.** Let  $g \geq 3$ ,  $k \geq 2$ . Then there is a graph G with no cycles of length  $\leq g$  and  $\chi(G) \geq k$ .

#### 5.3 The Structure of Random Graphs

**Proposition 40.**  $p = \frac{1}{n}$  is a sharp threshold for  $G \in \mathcal{G}(n,p)$  to contain a  $\triangle$ , in the sense that:

- if  $p = o(\frac{1}{n})$  then almost every  $G \in \mathcal{G}(n,p)$  has no  $\triangle$ , whereas
- if  $p = \omega(\frac{1}{n})$  then almost every  $G \in \mathcal{G}(n,p)$  has a  $\triangle$ .

**Theorem 41.** There exists a function  $d: \mathbb{N} \to \mathbb{N}$  such that a.e.  $G \in \mathcal{G}(n,p)$  has  $\omega(G) \in \{d-1,d,d+1\}$  (where d=d(n)).

Corollary 42. Almost every  $G \in \mathcal{G}(n,p)$  has

$$\chi(G) \ge (1 + o(1)) \frac{n \log \frac{1}{q}}{2 \log n}$$

where q = 1 - p.

# 6 Algebraic Methods

### 6.1 The Chromatic Polynomial

**Theorem 43** (Cut-fuse relation). Let G be a graph,  $e \in E(G)$ ,  $k \geq 1$ . Then  $f_G(k) = f_{G-e}(k) - f_{G/e}(k)$ .

Corollary 44. Let G be a graph. Then  $f_G$  is a polynomial.

Corollary 45. If |G| = n, e(G) = m then

$$f_G(X) = X^n - mX^{n-1} + \dots$$

# 6.2 Eigenvalues

**Theorem 46.** Let G be a graph,  $\Delta(G) = \Delta$ ,  $\lambda$  an eigenvalue of G. Then  $|\lambda| \leq \Delta$ . Moreover if G is connected then  $\Delta$  is an eigenvalue  $\iff G$  is  $\Delta$ -regular; in this case  $\Delta$  has multiplicity 1 and eigenvector  $\begin{pmatrix} 1\\1\\\vdots\\1 \end{pmatrix}$ .

### 6.3 Strongly Regular Graphs

**Theorem 47** (Rationality condition). Let G be (k, a, b)-strongly regular. Then

$$\frac{1}{2} \left\{ (n-1) \pm \frac{(a-b)(n-1) + 2k}{\sqrt{(b-a)^2 - 4(b-k)}} \right\} \in \mathbb{Z}.$$

8