

## LAB SESSION 10: Hangman game

Date of the Session: \_\_\_\_\_

Time of the Session: \_\_\_\_\_ to \_\_\_\_\_

Pre-Lab: What is the search where we examine the problem which arises when we try to plan ahead of the world and other agents are planning against us and write its algorithm and provide a code as an example.

Adversal search is a search, where we examine the problem which arises when we try to plan ahead of the world and other agents are planning against. We have studied the search strategies which are only associated with a single agent that to find the solution which is then expressed in the form of a sequence of actions

Lab:  
 Hangman is a classic word-guessing game. The user should guess the word correctly by entering alphabets of the user choice. The Program will get input as single alphabet from the user and it will matchmaking with the alphabets in the original word. If the user given alphabet matches with the alphabet from the original word then user must guess the remaining alphabets. Once the user successfully found the word he will be declared as winner otherwise user lose the game.

```

import random
words = ("which", "their", "about", "would", "these",
          "Others", "words", "could", "write")
def isWord(user_word, wordly_word):
    for x in user_word:
        print(x, end=" ")
    print()
    for i in range(len(user_word)):
        if (user_word[i]) == wordly_word[i]:
            print("?????", end=" ")
        else if user_word[i] in wordly_word:
            print("?????", end=" ")
        else:
            print(" ", end=" ")
    if user_word == wordly_word:
        return
    else:
        return 0
  
```

Post-Lab:  
The following problems that can be solved using Constraint Satisfaction Problems (CSP):

Test Case 1: Magic Square ( $\{[1,2,3],[4,5,6],[7,8,9]\}$ )

False

This is not a magic square. The numbers in the rows/columns/diagonals do not add up to the same value.

Let's try another list of lists.

Test Case 2: Magic Square ( $\{[2,7,6],[9,5,1],[4,3,8]\}$ )

True

Develop a python program that satisfies below operations.

2	7	6	→ 15
9	5	1	→ 15
4	3	8	→ 15
15	↓	15	15
15	15	15	15

import csp

x = csp variable (range (1,2,1))

y = csp variable (range (1,2,1))

z = csp variable (range (1,2,1))

Pythagorean triple-constraint = csp constraint (x,y,z).

and values = values (0) ↓↓ 2 + values (c)

\*\* 2 - = values (2) + + 2

if len(values) == 3:

else:

true

total-order-constraint = csp constraint

(x,y,z)<sup>86</sup>

value[i]

if lens(values) == 3:

else

True

Pythagorean-triples-problem > clp. constraint problems

for solution-assignments in esp. heuristic bacteria;

print ("x", x.values)

print ("y", y.values)

print ("z", z.values)

print()

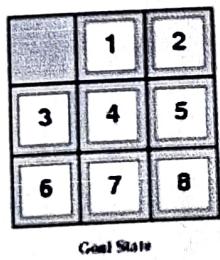
## LAB SESSION 11: Optimal Decisions in Games

Date of the Session: \_\_\_ / \_\_\_ / \_\_\_

Time of the Session: \_\_\_ to \_\_\_

Pre Lab:

Define Cost Function and idea cost function in 8-puzzle problem.  
Take the below picture as an example and calculate the heuristic functions for it (using the most two commonly used techniques).



cost function = returns the actual cost from the start node to the current node

-  $h_1(n)$  = number of misplaced tiles

-  $h_2(n)$  = meh hatten distance

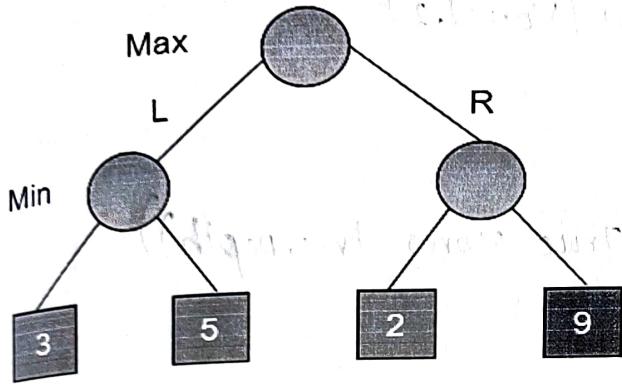
-  $h_3(n)$  = Gaschnig's

$$h_1(s) = ?$$

$$h_2(s) = ? \quad 3 + 1 + 2 + 2 + 3 + 3 + 2 = 18$$

$$h_3(s) = 8$$

Consider a game which has 4 final states and paths to reach the final state are from root to 4 leaves of a perfect binary tree as shown below. Assume you are the maximizing player and you get the first chance to move, i.e., you are at the root and your opponent at the next level. Which move you would make as a maximizing player considering that your opponent also plays optimally?



import math

def minimax (curDepth, nodeIndex, maxTurn, scores,  
+ targetDepth):

if (curDepth == budgetDepth):  
return scores [nodeIndex]

if (maxTurn):

return max (minimax (curDepth+1, node  
Index\*2, False, scores, targetDepth),

minimax (curDepth+1, nodeIndex\*2+1, False,  
scores, targetDepth))

else:

return min (minimax (curDepth+1, nodeIndex\*2,

True, scores, targetDepth),

```
minimax(asDepth + 1, nodeIndex * 2 + 1, True, scores,  
targetDepth)
```

```
scores = (3, 5, 2, 5, 12, 5, 23, 25)
```

```
Depth = math.log(len(scores), 2)
```

```
print(end="")
```

```
print(minimax(0, 0, True, scores, tree.Depth))
```

```
class Tic Tac Toe:
```

```
    def __init__(self):
```

```
        self.board = []
```

```
    def create_board(self):
```

```
        for i in range(3):
```

```
            row = []
```

```
            for j in range(3):
```

```
                row.append('..')
```

```
            self.board.append(row)
```

```
    def get_random_first_player(self):
```

```
        return random.randint(0, 0.5)
```

```
    def first_spot(self, row, col, player)
```

```
    def self.board[row][col] = player
```

```
    def is_player_win(self, player):
```

```
        win = None
```

```
        n = len(self.board)
```

```
        for j in range(n):
```

```
            win = True
```

```
            for j in range(n):
```

```
                if self.board[i][j] != player
```

```
                    if self.is_board_filled():
```

print ("match Draw!")

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break

player = self. swapPlayer. turn(player)  
print()

self. showBoard()

tic-tac-toe = TicTacToe()

tic-tac-toe. start()

return False

break

if win:

return win

for i in range(n):

win = True

for j in range(n):

if self. board[i][j] != player:  
win = False

break

if win:

return win

win = True

for i in range(n):

if self. board[i][j] != player:

win = False

n: len(self. board)

```
for i in range(n):
    win = True
    for j in range(n):
        if self.board[i][j] != player:
            win = False
            break
    if win:
        return win
for i in range(n):
    win = True
    for j in range(n):
        if self.board[j][i] != player:
            win = False
            break
    if win:
        return win
win = True
for i in range(n):
    if self.board[i][j] != player:
        win = False
        break
if win:
    return win
if self.player-win(player):
    print("I")
    break
```

$N = 9$

```
def printing (arr):  
    for i in range(N):  
        for j in range(N):  
            print (arr[i][j], end = " ")
```

```
print()  
def isSafe (grid, row, col, num):  
    for x in range(9):  
        if grid [x][col] == num:  
            return False
```

startRow = row - row % 3

```
startCol = col - col % 3  
for i in range(3):  
    for j in range(3):  
        if grid [i+startRow][j+startCol] == num:  
            return False  
return True
```

```
solveSudoku (grid, row, col):
```

Writing space for the Problem:(For Student's use only)

```

if (row == P-1 and col == N)
    return true

if col == N:
    row += 1
    col = 0
    if grid[row][col] == 0
        return solveSolution(grid, row, col+1)

for num in range(1, w+1, 1):
    if issafe(grid, row, col, num):
        grid[row][col] = num
        if solvesudoku(Grid, row, col+1):
            return True
        grid[row][col] = 0
return False.

```

```
grid = [[3, 0, 6, 5, 0, 8, 4, 0, 0],  
        [5, 2, 0, 0, 0, 0, 0, 0, 0],  
        [0, 8, 7, 0, 0, 0, 0, 3, 1],  
        [0, 0, 3, 0, 1, 0, 0, 8, 0],  
        [9, 0, 0, 8, 6, 3, 0, 0, 5],  
        [0, 5, 0, 0, 9, 0, 6, 0, 0],  
        [1, 3, 0, 0, 0, 0, 2, 5, 0],  
        [0, 0, 0, 0, 0, 0, 0, 7, 4],  
        [0, 0, 5, 2, 0, 6, 3, 0, 0]]
```

```
if (solve_sudoku(grid 0, 0)):
```

```
    print(grid)
```

```
else:
```

```
    print("no solution exists")
```

## LAB SESSION 12: Constraint Satisfaction problems

Date of the Session: 1/1/2023

Time of the Session: \_\_\_\_\_ to \_\_\_\_\_

Pre-Lab:

1. Solve the CSP problem:

TWO+TWO=FOUR

A constraint satisfaction problem (CSP) is defined by  $x$  is a set of  $n$  variables  $x_1, x_2, x_3, \dots, x_n$  each defined by a finite domain  $D_1, D_2, \dots, D_n$  of possible values. A solution is an assignment of values to be variables that satisfies all constraints.

$$\text{Constraints} \\ \Rightarrow \text{problem} = \text{TWO} + \text{TWO} = \text{FOUR}$$

variables =  $T \in \{0, \dots, 9\}; W \in \{0, \dots, 9\}; D \in \{0, \dots, 9\}$

$T \in \{0, \dots, 9\}; W \in \{0, \dots, 9\}, R \in \{0, \dots, 9\}$

$F \in \{0, \dots, 9\}; U \in \{0, \dots, 9\}, R \in \{0, \dots, 9\}$

$x_1 \in \{0, \dots, 9\}; x_2 \in \{0, \dots, 9\}; x_3 \in \{0, \dots, 9\}$

Constraints:

$$0 + 0 = R + 10^k x_1$$

$$x_1 + W + W = U + 10^k x_2$$

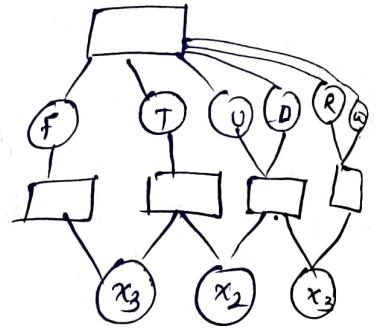
$$x_2 + T + T = 0 + 10^k x_3$$

$$x_3 = F$$

each letter has a different digit

- By cryptarithmetic

$$\begin{array}{r} \text{TWO} \\ + \text{TW O} \\ \hline \text{FOUR} \end{array}$$



Variables:  $f, T, U, W, R, O, x_1, x_2, x_3$

Domains:  $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$  (same domain parallel)

Example Constraints:

all dif ( $f, T, U, W, R, O$ )

A unary constraint:  $f \neq 0$

An n-ary constraint:  $O + O = R + 10 \times x_1$

Can add constraints to restrict the  $x_i$ 's to 0

$$\begin{array}{r} \text{TWO} \\ + \text{TW O} \\ \hline \text{FOUR} \end{array} \quad \begin{array}{r} \text{#65} \\ + \text{#65} \\ \hline \text{1530} \end{array}$$

A solution is an assignment to all variables from their domain so that all constraints are satisfied.

for any CSP, there might be a single, multiple or no solutions at all.

2. Discuss Zero-sum game rule with an example.

zero-sum games are found in many contexts. poker and gambling are popular examples of zero-sum games since the sum of amounts won by some players equals the combined losers of others. Games like chess and tennis, where there is one winner and one loser, are also zero-sum games.

Example of a zero-sum game: The game of matching pennies is often cited as an example of zero-sum game, according to game theory. The game involves two players A and B simultaneously placing a penny on the table. The payoff depends on whether the pennies match or not. If both pennies are heads or tails, player A wins and keeps player B's penny, if they do not match

then player B win and keeps player A's penny

A\B	Heads	Tails
Heads	(a) +1, -1	(b) -1, +1
Tails	(c) -1, +1	(d) +1, -1

As can be seen, the combined payoff for A and B in all four cell is zero.