Large-Scale and Multi-Structured Databases Column Databases Design Tips Prof. Pietro Ducange







Users (and queries) Guide the DB Design

In the following, we show some examples of *typical queries* that may led to choose *column DB*:

- How many new orders were placed in the Northwest region yesterday?
- When did a particular customer last place an order?
- What orders are en route to customers in London, England?
- What products in the Ohio warehouse have fewer than the stock keeping minimum number of items?







Information Needed for the DB Design

Queries provide information needed to effectively design column family databases.

The information includes:

- Entities
- Attributes of entities
- Query criteria
- Derived values

Designers start with this information and then use **the features of column-based** databases management systems to **select** the most appropriate **implementation**.







How to Use the Extracted Information

Entities

A single row describe the instance of a single entity. Rows are uniquely identified by row keys.

Attributes of entities

Attributes of entities are modeled using columns.

Query criteria

The selection criteria should be used to determine optimal ways to organize data with tables and column families and how to build partitions

Derived values

it is an indication that additional attributes may be needed to store derived data (example "count of orders placed yesterday by a user").







Differences with Relational DBs

- Column databases are implemented as sparse and multidimensional maps
- Columns can vary between rows
- Columns can be added dynamically
- Joins are not used because data is denormalized
- It is suggested to use *separate keyspace* for each application.







Sparse and Multi-dimensional Maps

Street	City	State	Province	Zip	Postal Code	Country
178 Main St.	Boise	ID		83701		U.S.
89 Woodridge	Baltimore	MD		21218	8	U.S.
293 Archer St.	Ottawa		ON		K1A 2C5	Canada
8713 Alberta DR	Vancouver		ВС		VSK 0AI	Canada

Image extracted from: "Dan Sullivan, NoSQL For Mere Mortals, Addison-Wesley, 2015

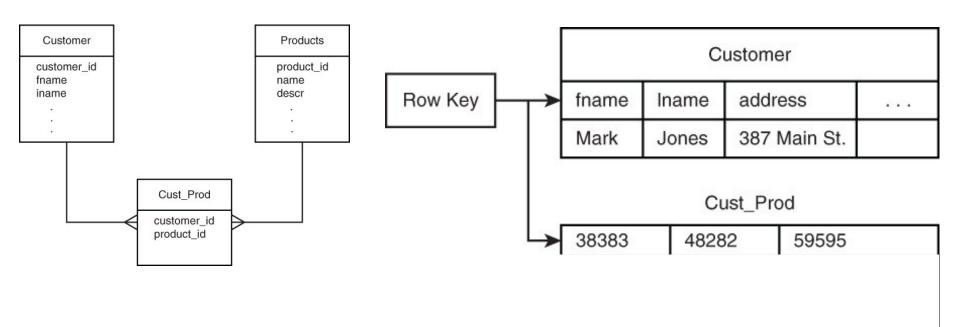






Denormalization and Valueless Columns

Example: Many to many relationships



Relational DB

We use 3 tables, including a join table.

Column DB

In order to avoid join operations, each customer includes a **set of column names** that correspond to purchased products (lds).

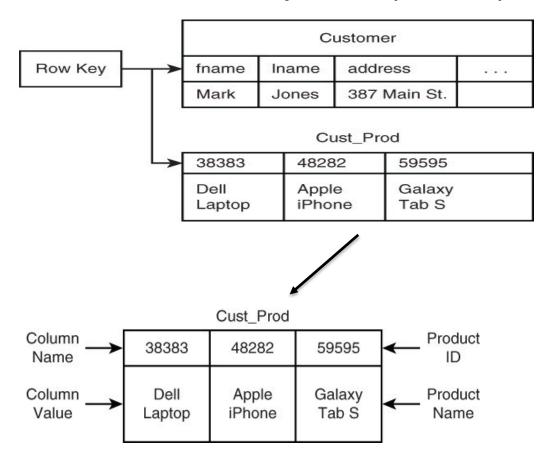






Use Both Column Names and Column Values to Store Data

Example: Many to many relationships



Because the column value is not used for anything else, we can store the product name there.

This can avoid the access to the product table, if we want to produce a report listing products bought by a customer.

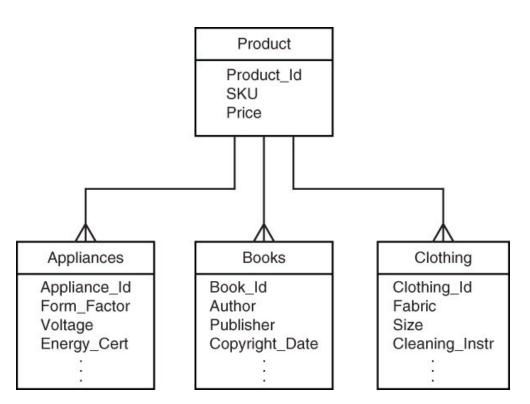
Keeping a copy of the product name in the customer table will *increase* the amount of *storage* used (*we improve the read performance*).







Model an Entity with a Single Row



A single instance of one entity, such as a particular customer or a specific product, should have *all its attributes* in a *single row*.

In this case, typical of column DBs, some rows store more column values than others.

The figure shows the classical organization in a relational DB of an entity that can assume different "shapes".

In order to *ensure* the *atomicity* of the operation in column DB, the same entity than in the figure must be organized in just *one table with 4 column families*.







Model an Entity with a Single Row

Row id

Product_columns

Applicance_columns

Row id

Product_columns

Clothing_columns

Row id

Product_columns

Book_columns

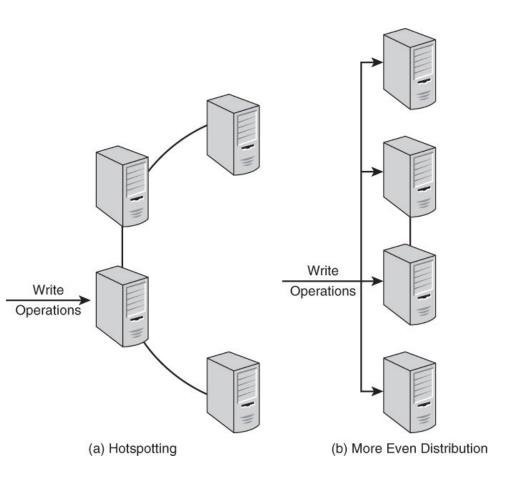






Avoid Hotspotting in Row Keys

Hotspotting occurs when many operations are performed on a small number of servers



We can prevent hotspotting by hashing sequential values generated by other systems and used as row key in column DBs.

Alternatively, we could **add a random string** as a prefix to the sequential value.

These strategies would *eliminate* the effects of the *lexicographic orde*r of the source file on the data load process.







Keep an Appropriate Number of Column Value Versions

- Some column DBMSs provides for column value versions
- The number of versions maintained is controlled by database parameters (min and max number of data versions)
- The versioning parameters depends on application requirements
- Versioning is useful if we need to roll back changes we did to column values.







Column Value Versions

Column Family					
Column Name ₁	Column Name ₂	Column Name ₃			
value _{1a} :timestamp _{1a}	value _{2a} : timestamp _{2a}	value _{3a} : timestamp _{3a}			
value _{1b} :timestamp _{1b}	value _{2b} :timestamp _{2b}	value _{3b} : timestamp _{3b}			
value _{1c} :timestamp _{1c}	value _{2c} :timestamp _{2c}	value _{3c} :timestamp _{3c}			







Avoid Complex Data Structures in Column Values

Any kind of data structure may be stored in a column value such a JSON file:

```
{
    "customer_id":187693,
    "name": "Kiera Brown",
    "address" : {
        "street" : "1232 Sandy Blvd.",
        "city" : "Vancouver",
        "state" : "Washington",
        "zip" : "99121"
        },
    "first_order" : "01/15/2013",
    "last_order" : "06/27/2014"
}
```

Notice that:

- Using *separate columns* for each attribute makes it easier to apply database features to the attributes (*indexing*).
- Separating attributes into individual columns allows you to use different column families if needed.







Indexing: Primary and Secondary Indexes

Index: a specific data structure in a DBMS that allows the database engine to *retrieve* data *faster* than it otherwise would.

In column databases, we can *look up* a *column value* to quickly *find rows* that reference that column value.

SELECT fname, Iname FROM customers WHERE state = 'OR';

Primary indexes are indexes on the **row keys** of a table (created and maintained by the DBMS).

Secondary indexes are indexes created on one or more column value

Either the database system or your application can create and manage secondary indexes.

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Secondary Indexes Managed by the Database Management System

General and common sense rule:

"if we need secondary indexes on column values and the column family database system provides automatically managed secondary indexes, then we should use them."

Example: we can speedup this query

SELECT fname, Iname FROM customers WHERE state = 'OR' AND Iname = 'Smith'

By setting an index on the *state* and on the *lname* attributes.

Typically, the DBMS will use the *most selective* index first.

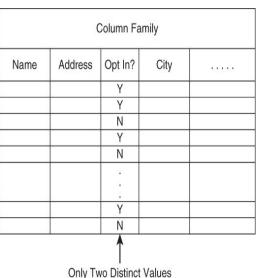
The automatic use of secondary indexes has the advantage the we do not have to change our code to use the indexes, if the application requirements will change.







When avoiding to use automatically managed indexes



Name	Address	City	State	Email
				ralken@gmail.com
				iman123@gmail.com
				dans37@yahoo.com
				marypdx@gmail.com
				gwashington@aol.com
				kcameron@future.com
				info@mybbiz.com

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						Sparsely Populated
						Populated

An index will probably not help much, especially if there are *roughly equal numbers* of each value.

Indexes may not help much here because the index will have to *maintain so much data* it could take *too much time* to search the index and retrieve the data.

Too few values associated to a specific column. It is not worth creating and managing an index data structure for this attribute.



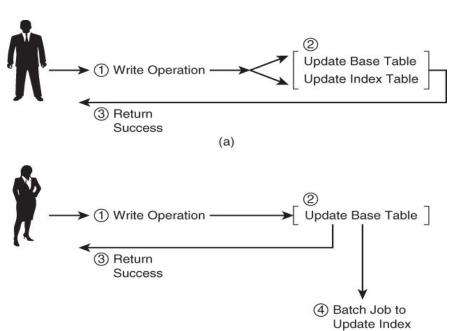




Create and Manage Secondary Indexes Using Tables

In this case, we have to *explicitly create* and manage tables to store data we would like to access via the index.

The *programmer* will be be *responsible* for maintaining the indexes and two main strategies may be adopted:



(b)

- a) *Updating an index table during write* operations keeps data synchronized but *increases the time* needed to complete a write operation.
- b) **Batch updates** introduce periods of time when the **data** is **not** synchronized, but this may be acceptable in some cases.





Table



Suggested Readings

Chapter 11 of the book "Dan Sullivan, NoSQL For Mere Mortals, Addison-Wesley, 2015."

All the images used in this presentation have been extracted from the aforementioned book.





