

Exercise 1

Consider a system where a client can be randomly routed to one (and only one) among n servers, with probability p_i , $1 \leq i \leq n$. Each server's service time is an exponentially distributed RV, with a mean $\frac{1}{\mu_i}$. All servers are independent.

- 1) Find the CDF and PDF of the service time *of a client*
- 2) Find the mean and variance of the service time
- 3) Consider an alternative design of the same system, where a client request is sent to *all the servers simultaneously*, and is only considered served when
 - a. *at least one* server has processed it
 - b. *all* the servers have processed it.

Find the CDF of the service time of a client in these cases.

- 4) Assume that $\mu_i = \mu$. Is one design of the system (among the initial option and options *a* and *b* at point 3) faster? Why?

Exercise 2

A counseling practice offers individual advice to its clients. It admits both *singles* and *couples*, but counsels its clients *individually* (spouses are requested to wait outside the counseling room). The arrival rate at the counseling practice is λ . An arrival may be a *single*, with probability π , and a *couple*, with probability $1 - \pi$. Individual counseling takes an exponentially distributed time, with a rate μ .

- 1) Model the practice as a queueing system, draw the CTMC and write down the global equilibrium equations.
- 2) Find the stability condition, justify it, and compute the PGF of the SS probabilities.
- 3) Compute the mean number of jobs in the system. Verify in the limit cases.
- 4) Compute the probability p_0 that the practice is empty, and the probability that only one client is in, p_1

Solution of Exercise 1

1) Call S the service time RV. Due to the Law of Total Probability, we have:

$$\begin{aligned} F(s) &= P\{S \leq s\} \\ &= \sum_{i=1}^n P\{S \leq s \mid \text{server} = i\} \cdot P\{\text{server} = i\} \\ &= \sum_{i=1}^n (1 - e^{-\mu_i \cdot s}) \cdot p_i \\ &= \sum_{i=1}^n F_i(s) \cdot p_i \end{aligned}$$

Hence:

$$f(s) = \frac{d}{ds} F(s) = \frac{d}{ds} \sum_{i=1}^n F_i(s) \cdot p_i = \sum_{i=1}^n p_i \cdot \mu_i \cdot e^{-\mu_i \cdot s} = \sum_{i=1}^n f_i(s) \cdot p_i$$

2) The mean service time is $\sum_{i=1}^n \frac{1}{\mu_i} \cdot p_i$. The variance is $\sum_{i=1}^n \frac{1}{\mu_i^2} \cdot p_i$, due to independence.

3) Case a . is a textbook case of minimum of independent exponential RVs. The theory says that the answers are $F(s) = 1 - e^{-\tau \cdot s}$, where $\tau = \sum_{i=1}^n \mu_i$

For case b , we have:

$$\begin{aligned} F(s) &= P\{S \leq s\} = P\{\max\{S_i\} \leq s\} \\ &= P\{S_1 \leq s, S_2 \leq s, \dots, S_n \leq s\} \\ &= \prod_{i=1}^n (1 - e^{-\mu_i \cdot s}) \\ &= \prod_{i=1}^n F_i(s) \end{aligned}$$

4) When all the servers are indistinguishable, the CDF for the three systems is, respectively:

- $F_o(s) = 1 - e^{-\mu \cdot s}$ (original)
- $F_a(s) = 1 - e^{-n \cdot \mu \cdot s}$ (design a)
- $F_b(s) = (1 - e^{-\mu \cdot s})^n$ (design b)

It is easy to see that $\forall s > 0, F_a(s) > F_o(s) > F_b(s)$. In fact:

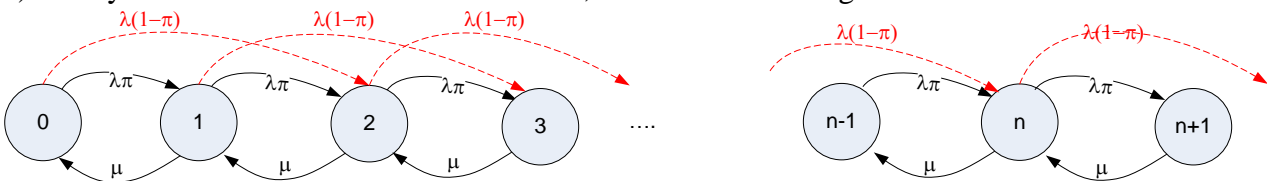
$$\begin{aligned} F_a(s) &> F_o(s) \\ 1 - e^{-n \cdot \mu \cdot s} &> 1 - e^{-\mu \cdot s} \\ e^{-n \cdot \mu \cdot s} &< e^{-\mu \cdot s} \\ -n\mu s &< -\mu s \\ n &> 1 \end{aligned}$$

Moreover, $F_o(s) > F_b(s)$ iff $1 - e^{-\mu \cdot s} > (1 - e^{-\mu \cdot s})^n$, which is obvious since the l.h.s. is between 0 and 1.

The above implies that design a is (probabilistically) faster than the original design, which is faster than b .

Exercise 2 – Solution

1) The system is an M/M/1 with bulk arrivals, and the CTMC diagram is below.



The global equilibrium equations are:

- State 0: $\lambda \cdot p_0 = \mu \cdot p_1$
- State 1: $(\lambda + \mu) \cdot p_1 = \mu \cdot p_2 + \lambda \cdot \pi \cdot p_0$
- State 2: $(\lambda + \mu) \cdot p_2 = \mu \cdot p_3 + \lambda \cdot \pi \cdot p_1 + \lambda \cdot (1 - \pi) \cdot p_0$
- State n : $(\lambda + \mu) \cdot p_n = \mu \cdot p_{n+1} + \lambda \cdot \pi \cdot p_{n-1} + \lambda \cdot (1 - \pi) \cdot p_{n-2}$

2) The RV that describes the size of the arrival is a Bernoullian g , such that $g_1 = P\{g = 1\} = \pi$, $g_2 = P\{g = 2\} = 1 - \pi$, hence $E[g] = 1 \cdot \pi + 2 \cdot (1 - \pi) = 2 - \pi$, $\mathbf{G}(z) = \pi \cdot z + (1 - \pi) \cdot z^2$. The computations for a generic $\mathbf{G}(z)$ can be found on the QT notes, and read:

- $\rho = \frac{\lambda}{\mu} E[g] = \frac{\lambda}{\mu} \cdot (2 - \pi)$. Note that, if $\pi = 1$, then the system is an M/M/1, and the stability condition is the usual one. Instead, if $\pi = 0$, the system is one with constant-batch bulk arrivals.

$$- \mathbf{P}(z) = \frac{\mu \cdot (1 - \rho) \cdot (1 - z)}{\mu \cdot (1 - z) - \lambda \cdot z \cdot [1 - \mathbf{G}(z)]}$$

By substituting the above $\mathbf{G}(z)$ into the above expression, after a few straightforward computations, we get:

$$\begin{aligned} \mathbf{P}(z) &= \frac{\mu \cdot (1 - \rho) \cdot (1 - z)}{\mu \cdot (1 - z) - \lambda \cdot z \cdot [1 - \pi \cdot z - (1 - \pi) \cdot z^2]} \\ &= \frac{\mu \cdot (1 - \rho) \cdot \cancel{(1 - z)}}{\mu \cdot \cancel{(1 - z)} - \lambda \cdot z \cdot \cancel{(1 - z)} [1 + z - \pi \cdot z]} \\ &= \frac{\mu - \lambda \cdot (2 - \pi)}{\mu - \lambda \cdot z - \lambda \cdot z^2 \cdot (1 - \pi)} \end{aligned}$$

3) Both expressions require computing the first derivative of the above expression, which is:

$$\frac{\partial}{\partial z} \mathbf{P}(z) = \frac{\partial}{\partial z} \left(\frac{\mu - \lambda \cdot (2 - \pi)}{\mu - \lambda \cdot z - \lambda \cdot z^2 \cdot (1 - \pi)} \right) = \frac{\lambda \cdot [1 + 2z \cdot (1 - \pi)] \cdot [\mu - \lambda \cdot (2 - \pi)]}{[\mu - \lambda \cdot z - \lambda \cdot z^2 \cdot (1 - \pi)]^2}$$

From the above, we obtain:

$$E[N] = \frac{\partial}{\partial z} \mathbf{P}(z) \Big|_{z=1} = \frac{\lambda \cdot [3 - 2\pi] \cdot [\mu - (2 - \pi)\lambda]}{[\mu - (2 - \pi)\lambda]^2} = \frac{(3 - 2\pi)\lambda}{\mu - (2 - \pi)\lambda}$$

When $\pi = 1$ the system is an M/M/1, and the above expression reads $E[N] = \frac{\lambda}{\mu - \lambda} = \frac{\rho}{1 - \rho}$.

When $\pi = 0$ the system is a constant-batch one, with $b = 2$, and the expression is $E[N] = \frac{3\lambda}{\mu - 2\lambda}$. The expression on the notes reads $E[N] = \frac{\rho \cdot (b+1)}{2 \cdot (1 - \rho)}$, which is equal to the former after some straightforward substitutions.

4) It is $p_0 = \lim_{z \rightarrow 0} \mathbf{P}(z) = 1 - \frac{\lambda}{\mu} \cdot (2 - \pi) = 1 - \rho$. This was expected, since ρ is the system utilization. Moreover, it is $p_1 = \frac{\partial}{\partial z} \mathbf{P}(z) \Big|_{z=0} = \frac{\lambda \cdot [\mu - \lambda \cdot (2 - \pi)]}{\mu^2} = \frac{\lambda}{\mu} \cdot \left[1 - \left(\frac{\lambda}{\mu} \right) \cdot (2 - \pi) \right] = \frac{\rho \cdot (1 - \rho)}{2 - \pi}$. If $\pi = 1$, the expression is the one of an M/M/1 system.