Large-Scale and Multi-Structured Databases

Graph Databases Recap of Graph Theory

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Vertex

In the figure, we show *instances* of an *entity (city)* represented by a set of vertices. We are in the context of *highways networks*.





The *properties* that the vertices, namely the cities, may have are:

- Population
- Coordinates
- Name
- Geographic Region

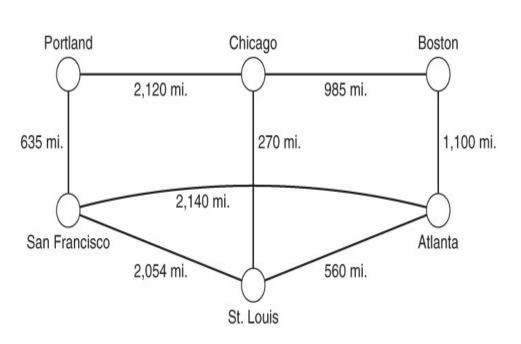




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Edges: Relationships between Vertices



The *properties* that the edge in the highways network, may have are:

- Distance
- Speed limit
- Number of lanes

A Weight is a commonly used property of edge.

It represents some value about the **strength** of the relationship (cost, distance, etc.)



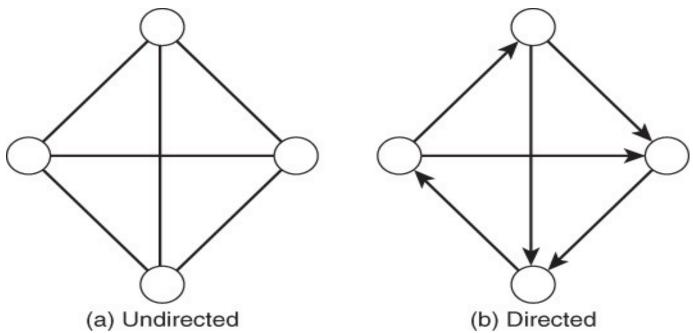




Undirected and Directed Graphs

Directed edges have a *direction*, whereas undirected edge does not have a direction.

The *presence* or not of a *direction* in edges depends on *specific applications*.



Example: Highways Network Graphs

Example: Family Relation Graphs

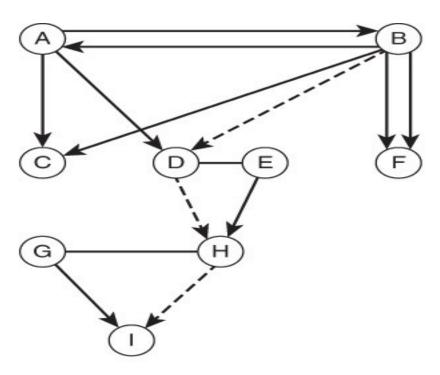






Path in a Graph

A *path* through a graph is a *set of vertices* along with the *edges* between those vertices.



Paths allow us to *capture* information about the *relationships* of vertices in a graph.

There may be *multiple paths* between two vertices. In this case (highways networks, for instance), often the problem *of finding the best* (the shortest, the cheapest, etc.) path is taken into consideration.

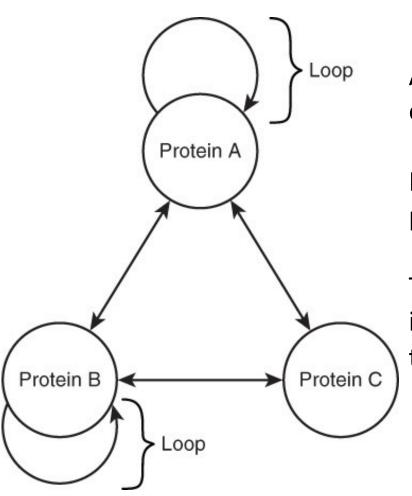








Loops in Graphs



A loops represents a *relation* that an entities has with *itself*.

Proteins, for example, may interacts with proteins of the same type.

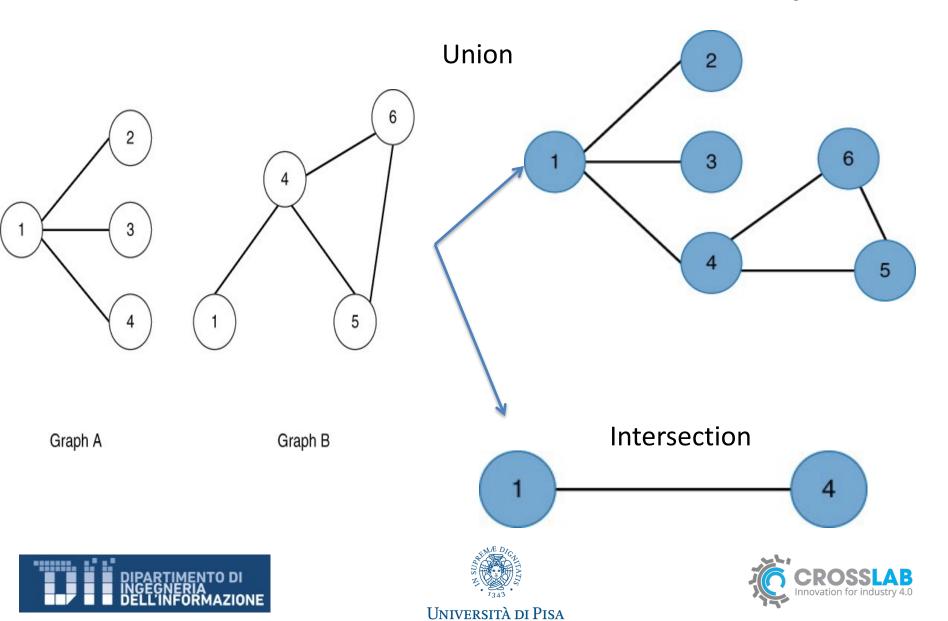
The presence of loop may *not have sense* in some specific domains, such as family tree graphs.



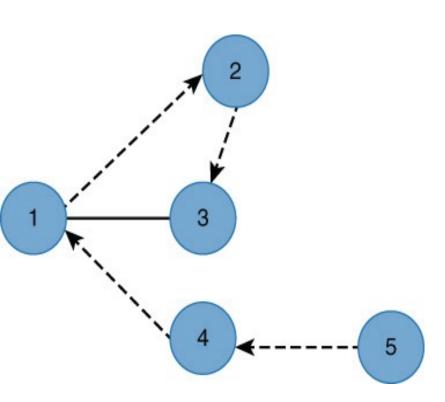




Union and Intersection of Graphs



Graph Traversal



Graph traversal is the process of *visiting* all *vertices* in a graph.

The purpose of this is usually to either **set** or **read** some property value in a graph.

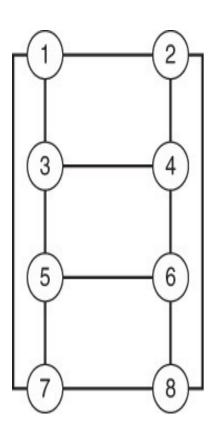
The traversal of the graph is often carried out in a particular way, such us *minimizing the cost* of the path.

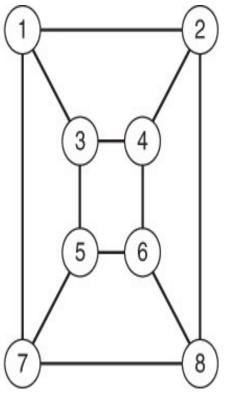






Isomorphism





Two graphs are *isomorphic* if:

- for each vertex in the first graph, there is a corresponding vertex in the other graph
- for each edge between a pair of vertices in the first graph, there is a corresponding edge between the corresponding vertices of the other graph.

Detecting isomorphism in graphs or sub graphs allows us to identify *useful patterns*.







Order and Size of Graphs

Order: the number of vertices

Size: the number of edges

These two measures are very important to evaluate time and memory occupancy required to handle specific operations.

Example: to find the largest subset of people in a social network that know each other.

The higher the order and the size of the graph, the harder will be to solve the query!







Degree and Closeness of a Vertex

Degree: is the number of edges linked to a vertex. It measures the **importance** of any given vertex in a graph.

Example: Consider a person with many family members and friends that he/she sees regularly (high degree vertex). What if that person contracts a contagious disease?

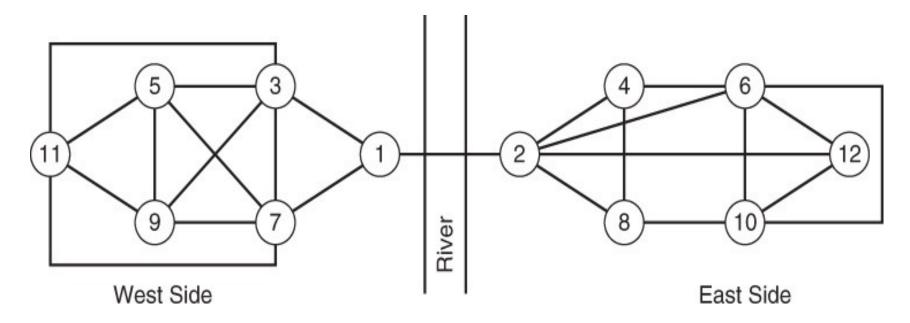
Closeness: is the measure of how far the vertex is from all others in the graph. It can be calculated, in a fully connected graph, as the **reciprocal of the sum** of the length of the **shortest paths** between the node and all other nodes in the graph.

It is useful to evaluate *the speed of information spreading in a network,* starting from a specific node.



Betweenness of a Vertex

Betweenness is a measure of how much of a bottleneck a given vertex is.



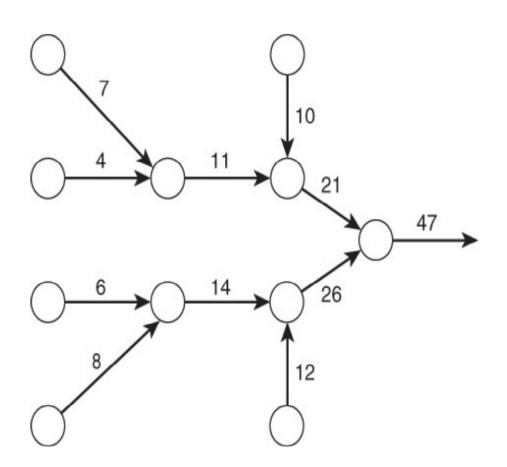
In the example above, both vertices 1 and 2 will have *high betweenness* levels. Indeed, they form a bottleneck in the graph and if one of them were removed, it would leave the graph *disconnected*.







Flow Networks



Applications: road systems, transportation networks, networks of storms drains

A flow network is a *direct graph*.

It adopts *capacitated* edges.

Each vertex has a set of *incoming* and *outcoming* edges.

In each vertex the sum of the capacities of incoming edges *cannot be greater* than the sum of the capacities of outcoming edges.

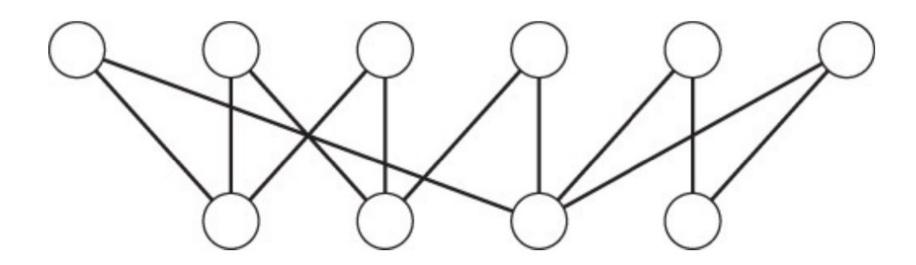
The unique exception to the previous rule regards the *source* and the *sink* nodes.







Bipartite Graph



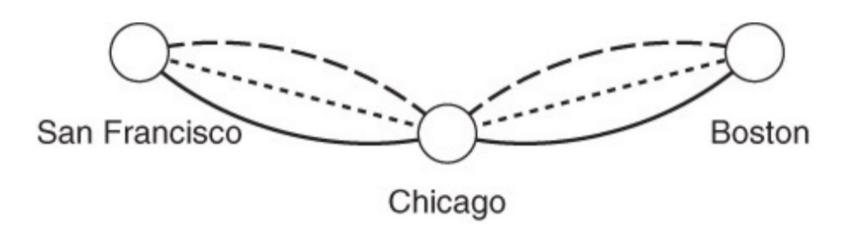
- There are two distinct sets of vertices.
- Each vertex is connected only with vertices of *the other set*.
- It can be used for modeling relationships between students and teachers.







Multigraph



- It is characterized by the presence of multiple edges between vertices.
- In the example above, each edge may represents a *different way* (by plane, by truck, etc.) to shipping item between the cities.
- Each edge can be equipped with different sets of properties.







Suggested Readings

Chapter 13 of the book "Dan Sullivan, NoSQL For Mere Mortals, Addison-Wesley, 2015".







Images

If not specified, the images shown in this lecture have been extracted from:

"Dan Sullivan, NoSQL For Mere Mortals, Addison-Wesley, 2015"





