

# Compiler Technology of Programming Languages

## **Code Generation**

### **Control Structures and Loops**

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# Simple If

```
if ( a )
```

```
{
```

```
    b = 1;
```

```
}
```

```
ldr w9, [x29, #-4]
```

```
cmp w9, #0
```

```
beq  Lexit1
```

```
mov  w10,  #1
```

```
str   w10,  b
```

```
Lexit1:
```

```
....
```

What if a is an expression? e.g. if (a&b), (a+b)

```
cmp w9, #0    where w9 = expr.place
```

```
beq Lexit1
```

```
cmp w9, w10
```

```
ble Lexit1
```

How to generate code for (a>b) ?

# If-else

```
if ( a )  
{  
    b = 1;  
}  
  
else  
{  
    b = 0;  
}
```

```
ldr w9, [x29, #-4]  
cmp w9, #0  
beq  Lelse1  
mov  w10, #1  
str  w10, b  
b    Lexit1
```

Lelse1:

```
mov  w11, #0  
str  w11, b
```

Lexit1:

....

# Nested If-else

```

if ( a )
{
    b = 1;
}
else
{
    if ( j )
    {
        b = 0;
    }
    else
    {
        c = 0;
    }
    a = 0;
}

```

```

ldr    w9,    a
cmp    w9, #0
beq    Lelse1
mov    w10,   #1
str    w10,   b
b      Lexit1

Lelse1:
ldr    w11,   j
cmp    w11, #0
beq    Lelse2
mov    w12,   #0
str    w12,   b
b      Lexit2

Lelse2:
mov    w13,   #0
str    w13,   c

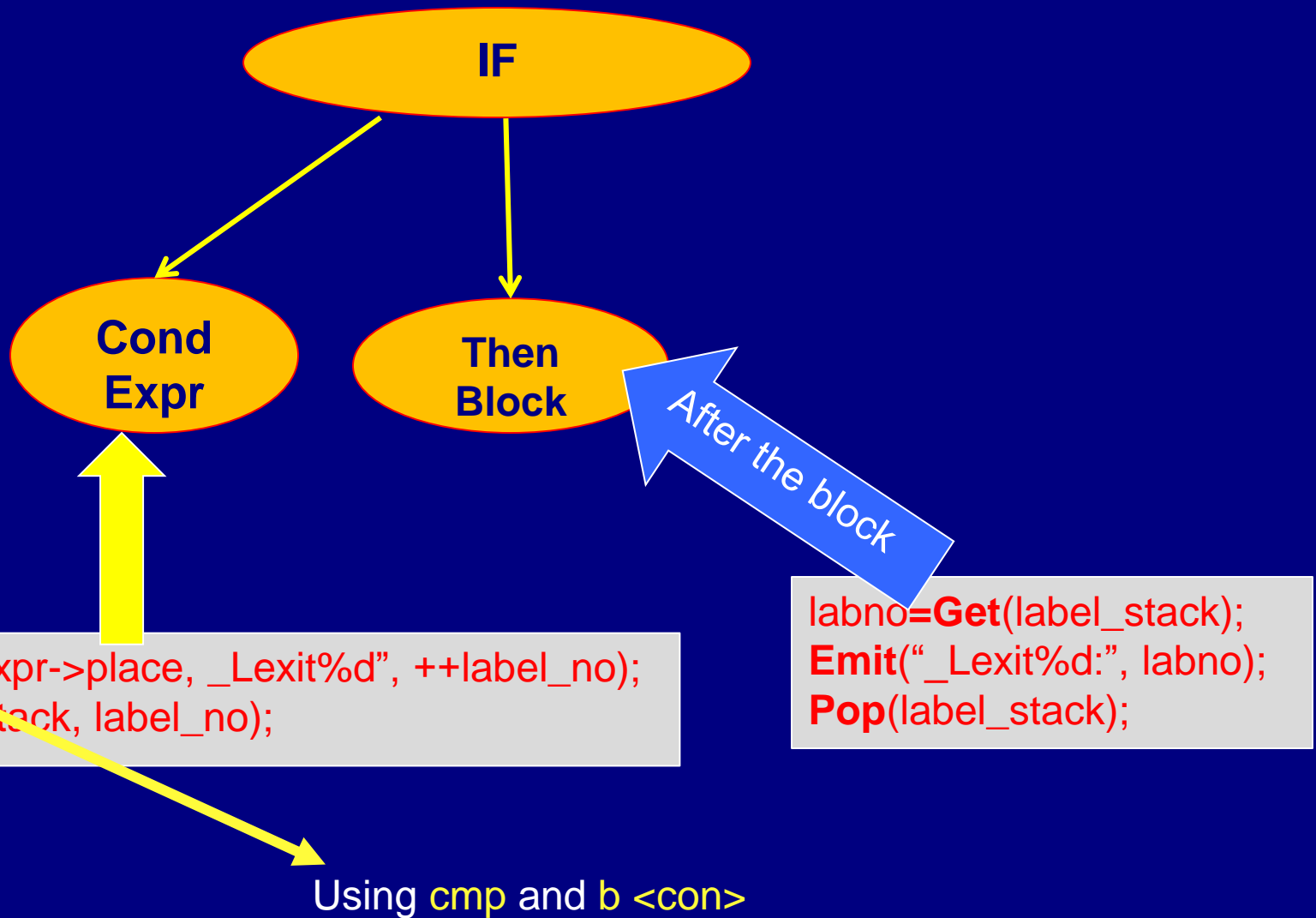
Lexit2:
mov    w9, #0
str    w9,   a

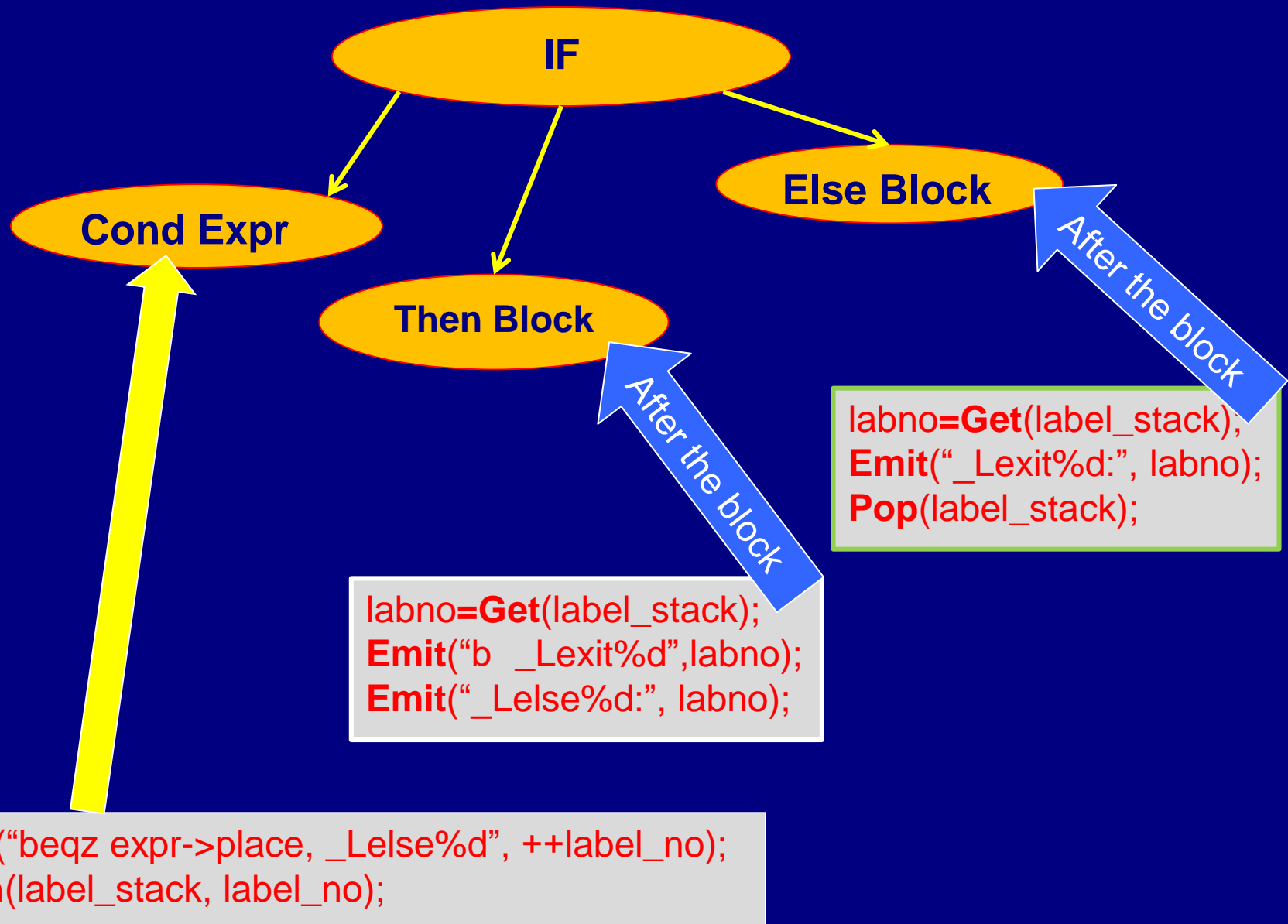
Lexit1:

```

# Code Generation for If Structures

- A test that branches to the else part (if-else) or the exit part (simple if) if the condition is false.
- Unique labels such as Lelse1, Lexit2
- Need to handle nested structures, so label number should be maintained on a stack.
- Since the code gen walks AST via recursive calls, label number can be maintained in a local variable.
- Jump to Lexit (skip over the else part) at the end of the *then* block.





# Avoid using the label stack

```
Procedure visit(AST IF n) {
```

```
    int local_label_number;
```

```
    local_label_number= ++label_no; /* keeps a local  
    copy of the current label_no. If there are nested  
    structures, this local label will be saved on the calling  
    stack automatically. */
```

```
    .....
```

```
    CodeGen(...);
```

```
    ...
```

```
    gen_labels();
```

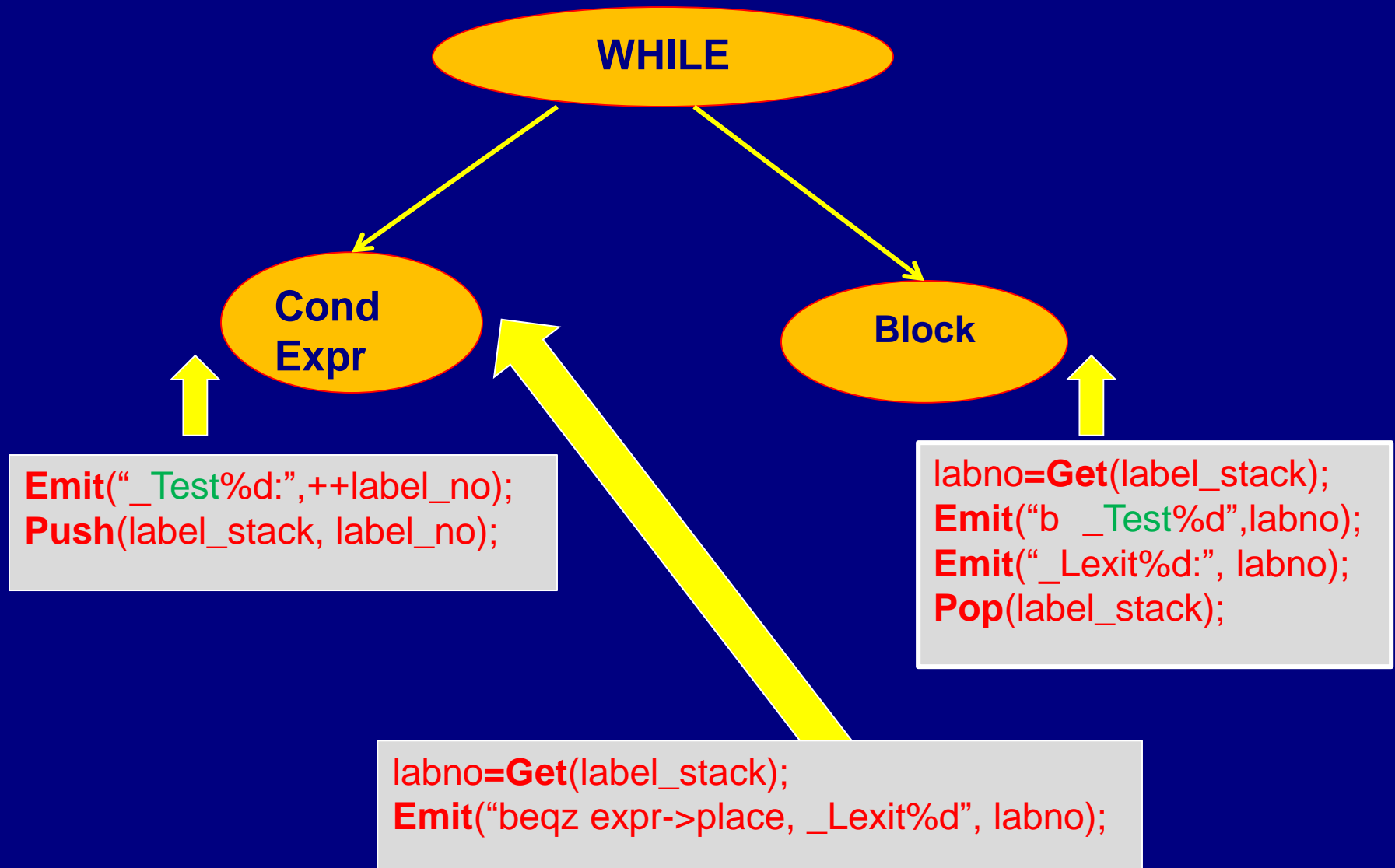
```
}
```



# While Loop

```
while ( a )  
{  
    stmt;  
}
```

```
_Test1:  
    ldr    w9,    a  
    cmp    w9,    #0  
    beq    _Lexit1  
    .... Code for stmt ...  
    b      _Test1  
_Lexit1:  
    ....
```

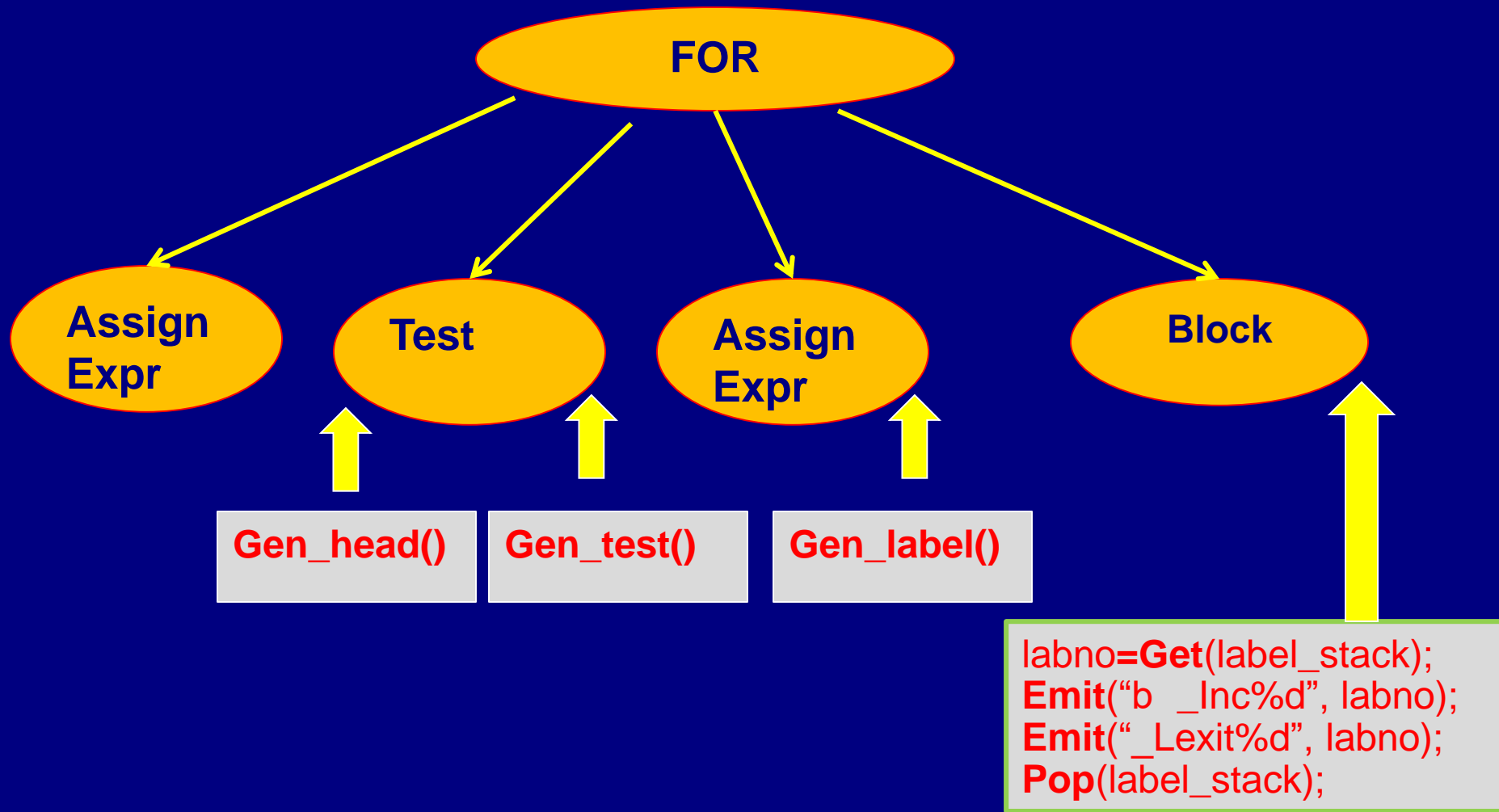


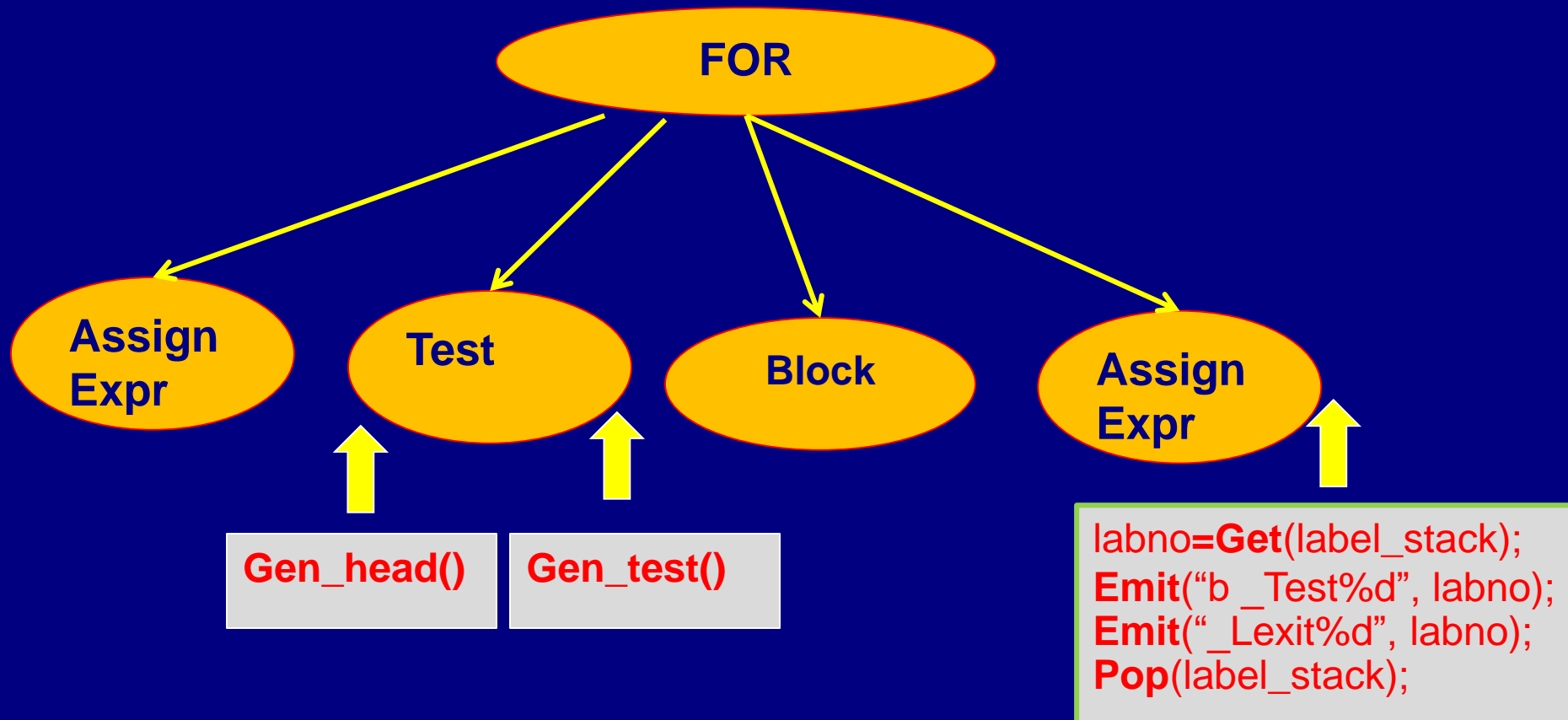
# For loop

```
for (j=0; j < n; j=j+1)
    stmt;
```

```

mov     w9,#0
str w9, j
_Test1:
    eval expr j<n in w10
    cmp w10,#0
    beq  _Lexit1
    b    _Body1
_Inc1:
    ldr w11, j
    add w11,w11,#1
    str w11, j
    b    _Test1
_Body1:
    code for stmt;
    b    _Inc1
_Lexit1:
```





# CodeGen Routines

## gen\_head

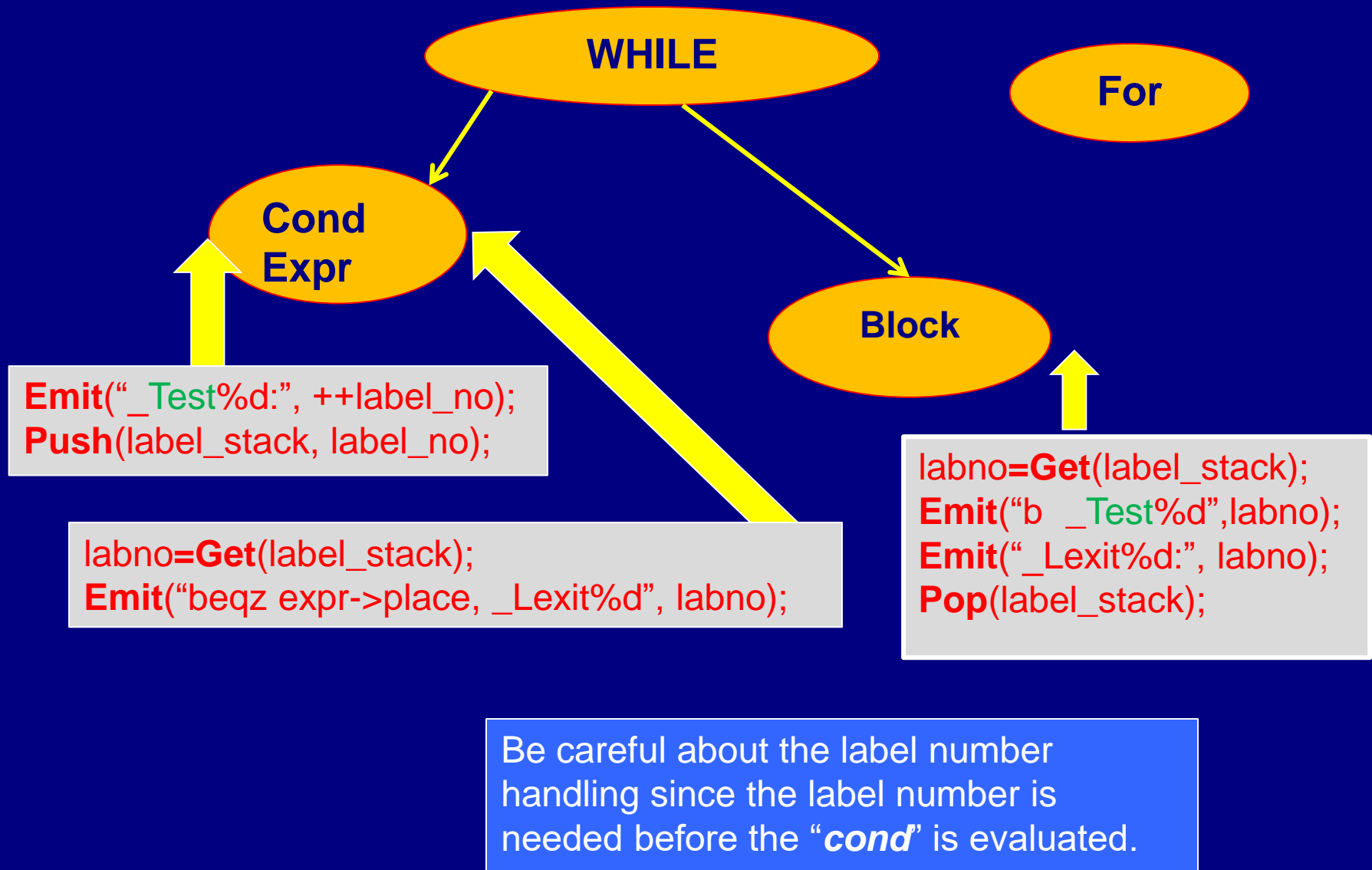
```
emit("_Test%d:", ++label_no);  
push(label_stack, label_no);
```

## gen\_test

```
labno=get(label_stack);  
if expr->type is integer but not integer const  
    emit("beqz    expr->place, _Lexit%d", labno);  
emit("b        _Body%d", label_no);  
emit("_Inc%d:", label_no);
```

## gen\_label

```
labno=get(label_stack);  
emit("b        _Test%d:", labno);  
emit("_Body%d:", labno);
```



# Boolean Expressions

■ (I && J && K)

assuming I, J, and K are in register w9, w10, w11  
the generated code may look like:

```
beqz    w9, F1
beqz    w10, F1
beqz    w11, F1
```

T1:

```
mov     w12, #1
b       E1
```

F1:

```
mov     r7, #0
```

E1:



# Boolean Expressions in Control Structures

## ■ Using Jump Code Sequence

If (**I** && **J** && **K**) Then... Else ...

    beqz        w9, Else

    beqz        w10, Else

    beqz        w11, Else

Then ....

....

Else

....

# Boolean Expressions in Control Structures

- The expression could be an assignment expression

If (a = (I && J && K)) Then... Else ...

For assignment expressions, we do need to generate a boolean value.

- && could mixed with ||

(a < b) && (a < 10) || (b > 5) || (c == 0)

# Jump Code for relop\_expr

if ((a < b) && (a < 10) || (b > 5)) then X else Y

assuming a, and b are in register w9, w10

```
bge      w9, w10, F1 // cmp and bge are folded
bge      w9, 10, F1
b         Then // jump directly instead of mov w11 #0
F1:      // (a<b) && (a<10) is False, need to eval (b>5)
ble      w10,5      ELSE
Then:
X
ELSE:
Y
```

# Function Call

- `save_registers();`
- `passing_param();`
- `gen_proc_call(name);`
- `gen_after_return();`

# Function Call

- `save_registers();`

Optional – depends on whether caller save registers are used

- `passing_param();`

Not required in Part I

- `gen_proc_call(name);`

generate “bl name” instruction

- `gen_after_return();`

generate a copy inst that moves x0 to x7 to a target reg

remove stack space allocated for arguments

restore saved registers if any