

Introduction to Mathematical Logic, Spring 2018

Homework 1

B04902012 Han-Sheng Liu, CSIE, NTU

E-mail: b04902012@ntu.edu.tw

(1) (a)

p	q	$q \rightarrow p$	$p \rightarrow (q \rightarrow p)$
T	T	T	T
T	F	T	T
F	T	F	T
F	F	T	T

(b)

p	q	r	$p \rightarrow (q \rightarrow r)$	$(p \rightarrow q) \rightarrow (p \rightarrow r)$	$(p \rightarrow (q \rightarrow r)) \rightarrow ((p \rightarrow q) \rightarrow (p \rightarrow r))$
T	T	T	T	T	T
T	T	F	F	F	T
T	F	T	T	T	T
T	F	F	T	T	T
F	T	T	T	T	T
F	T	F	T	T	T
F	F	T	T	T	T
F	F	F	T	T	T

(c)

p	q	$p \vee q$	$p \rightarrow (p \vee q)$
T	T	T	T
T	F	T	T
F	T	T	T
F	F	F	T

(d)

p	q	$\neg p \rightarrow q$	$p \rightarrow (\neg p \rightarrow q)$
T	T	T	T
T	F	T	T
F	T	T	T
F	F	F	T

All formulas are tautology.

(2) The formulas are constructed as following:

$$\begin{aligned}\alpha_k &= \bigvee_{1 \leq i \leq n-1} p_{k,i} \\ &\quad \bigwedge_{1 \leq i < j \leq n-1} \neg(p_{k,i} \wedge p_{k,j}) \\ &\quad \bigwedge_{1 \leq l \leq n} \bigwedge_{1 \leq i \leq n-1} \neg(p_{k,i} \wedge p_{l,i}) \\ X &= \{\alpha_k | 1 \leq k \leq n\}\end{aligned}$$

X corresponds to $(n-1)$ -colorability of an n -clique G . The construction is a generalized version of the formula set in the proof of *Lemma 3.6*.

X itself is trivially unsatisfiable, since an n -clique is not $(n-1)$ -colorable. However, every proper subset of X is satisfiable. Intuitively, a proper subset corresponds to a proper induced subgraph of G , which is an m -clique with $m < n$ and is thus $(n-1)$ -colorable.

Here gives an truth assignment for a proper set Y of X . Let $|Y| = m$, $m < n$. W.L.O.G, assume that $Y = \{\alpha_k | 1 \leq k \leq m\}$. A satisfying assignment is

$$p_{k,i} = \begin{cases} \text{T} & , \text{ if } k \leq m \text{ and } k = i. \\ \text{F} & , \text{ otherwise.} \end{cases}$$

That is, color the vertex i with color i .

(3)

$$\begin{aligned}X \models \alpha &\Rightarrow X \cup \{\neg\alpha\} \text{ is not satisfiable} \\ &\Rightarrow X \cup \{\neg\alpha\} \text{ is not finitely satisfiable} \\ &\Rightarrow \exists X' \subseteq X \cup \{\neg\alpha\} \text{ s.t. } X' \text{ is finite and non-satisfiable} \\ &\Rightarrow X' \cup \{\neg\alpha\} \text{ is a finite subset of } X \text{ and non-satisfiable} \\ &\Rightarrow X' \setminus \{\neg\alpha\} \models \alpha\end{aligned}$$

i.e. the set $X' \setminus \{\neg\alpha\}$ is the desired X_0 .