

## Problem 1:

Let's compare some basic math functions to refresh our memory. For each of the following, just write which function is *asymptotically greater* (So, you should be thinking about your asymptotic notations!). Show your reasoning for the same.

1.  $10000000000n^2$  vs  $n^3$ :  $n^3$  is asymptotically greater than  $n^2$ . No matter how large the coefficient for  $n^2$ , as long as it is a constant  $n^3$  will be asymptotically greater.
2.  $n^2 \log(n)$  vs  $n(\log(n))^{10}$ :  $n(\log n)^{10}$  will be asymptotically greater because it grows faster
3.  $n \log n$  vs  $2n$ :  $n^{\log n}$  is asymptotically greater because between base  $n$  and base 2, base  $n$  will grow faster than the constant 2.
4.  $2n$  vs  $2^{2n}$ : asymptotically these two should be the same because they have the same base of 2 and they are to the power of  $n$ . The only difference is the constant in front of the second  $n$ , however asymptotically this would not make a difference.

## Problem 2:

Now let's examine some [pseudo]code and apply asymptotic notation to it.

```
isPrime(n) :  
  
for(i = 2, i*i <= n; i++) { //i^2 = n-> i = sqrt(n)  
  
    if(n % i == 0) { //O(1)  
  
        return false //O(1)  
  
    }  
  
    return true // O(1)
```

What is the

1. Best Case:  $O(1)$  \*\*if  $n$  is 2 then  $n \bmod 2$  would be 0 and return false in the first iteration.
2. Worst Case:  $O(\sqrt{n})$
3. Average Case:  $O(\sqrt{n})$

Time complexity for the above function?  $O(\sqrt{n})$  because on average the function has to loop through the entire for loop which increases by one and stops after  $i$  reaches a value of square root of  $n$ .