### **Contents**

1	Basic         1           1.1 .vimrc         1
2	Combinatorics         1           2.1 FFT         1           2.2 FWT         1           2.3 NTT         2           2.4 permanent         2
3	Data Structure23.1 Heavy Light Decomposition
4	Flow         3           4.1 CostFlow         3           4.2 Dinic         3
5	Geometry         3           5.1 Line and points
6	Graph       5         6.1 2-SAT       5         6.2 maximal cliques       5         6.3 Tarjan SCC       5
7	Number Theory         6           7.1 basic         6           7.2 Chinese Remainder Theorem         6           7.3 Discrete Log         6           7.4 Lucas         6           7.5 Meissel-Lehmer Pl         6           7.6 Miller Rabin with Pollard rho         7           7.7 Primitive Root         7
8	String         7           8.1 AC automaton         7           8.2 Gusfield         8           8.3 KMP         8           8.4 Manacher         8           8.5 Suffix Array         8           8.6 Suffix Automaton         9
9	Formulas       10         9.0.1 Pick's theorem       10         9.0.2 Graph Properties       10         9.0.3 Number Theory       10         9.0.4 Combinatorics       10         9.0.5 巨次, 巨次和       10         9.0.6 Burnside's lemma       10         9.0.7 Count on a tree       11

# 1 Basic

#### 1.1 .vimrc

```
syntax on
set nu ai bs=2 sw=2 et ve=all cb=unnamed mouse=a ruler
incsearch hlsearch
```

# 2 Combinatorics

### 2.1 FFT

```
typedef complex<double> cpx;
const double PI = acos(-1);
vector<cpx> FFT(vector<cpx> &P, bool inv = 0) {
  assert(__builtin_popcount(P.size()) == 1);
  int lg = 31 - __builtin_clz(P.size()), n = 1 << lg;
// == P.size();</pre>
  for (int j = 1, i = 0; j < n - 1; ++j) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1);
    if (j < i) swap(P[i], P[j]);</pre>
  } //bit reverse
  auto w1 = \exp((2 - 4 * inv) * PI / n * cpx(0, 1)); //
        order is 1<<lg
  for (int i = 1; i <= lg; ++i) {
    auto wn = pow(w1, 1<<(lg - i)); // order is 1<<i</pre>
    for (int k = 0; k < (1 << lg); k += 1 << i) {
      cpx base = 1;
      for (int j = 0; j < (1 << i - 1); ++j, base = base * wn) {
         auto t = base * P[k + j + (1 << i - 1)];
         auto u = P[k + j];
         P[k + j] = u + t;
         P[k + j + (1 \ll i - 1)] = u - t;
    }
  if(inv)
    for (int i = 0; i < n; ++i) P[i] /= n;
  return P;
} //faster performance with calling by reference
```

### 2.2 FWT

```
vector<int> fast_OR_transform(vector<int> f, bool
     inverse) {
   for (int i = 0; (2 << i) <= f.size(); ++i)
     for (int j = 0; j < f.size(); j += 2 << i)
  for (int k = 0; k < (1 << i); ++k)
    f[j + k + (1 << i)] += f[j + k] * (inverse? -1</pre>
                : 1);
  return f;
vector<int> rev(vector<int> A) {
   for (int i = 0; i < A.size(); i += 2) swap(A[i], A[i
       ^ (A.size() - 1)]);
   return A;
vector<int> fast_AND_transform(vector<int> f, bool
     inverse) {
   return rev(fast_OR_transform(rev(f), inverse));
vector<int> fast_XOR_transform(vector<int> f, bool
     inverse) {
   for (int i = 0; (2 << i) <= f.size(); ++i)
     for (int j = 0; j < f.size(); j += 2 << i)
for (int k = 0; k < (1 << i); ++k) {</pre>
          int u = f[j + k], v = f[j + k + (1 << i)];

f[j + k + (1 << i)] = u - v, f[j + k] = u + v;
  if (inverse) for (auto &a : f) a /= f.size();
   return f;
}
```

#### 2.3 NTT

```
/* p == (a << n) + 1
         1 \ll n
                                          root
   n
                       97
   5
         32
                                    3
                       193
   6
         64
         128
                       257
                                          3
   8
         256
                       257
                                    1
                                          3
   9
         512
                       7681
                                    15
                                          17
   10
                       12289
                                    12
         1024
                                          11
   11
         2048
                       12289
                                          11
   12
         4096
                       12289
                                          11
   13
         8192
                       40961
                                          3
                       65537
   14
         16384
   15
         32768
                       65537
                                          3
   16
         65536
                       65537
                                    1
                                          3
   17
         131072
                       786433
                                    6
                                          10
                                          10 (605028353,
   18
         262144
                       786433
                                    3
        2308, 3)
   19
         524288
                       5767169
                                    11
                                          3
   20
         1048576
                       7340033
                                          3
         2097152
   21
                       23068673
                       104857601
   22
         4194304
                                          3
                                    25
                                          3
   23
         8388608
                       167772161
                                    20
   24
         16777216
                       167772161
                                    10
   25
         33554432
                       167772161
                                          3 (1107296257, 33,
        10)
   26
         67108864
                       469762049 7
                                          31 */
   27
         134217728
                       2013265921 15
LL root = 10, p = 786433, a = 3;
LL powM(LL x, LL b) {
  LL s = 1, m = x \% p;
  for (; b; m = m * m % p, b >>= 1)
    if (b\&1) s = s * m % p;
  return s;
vector<LL> NTT(vector<LL> P, bool inv = 0) {
  assert(__builtin_popcount(P.size()) == 1);
  int lg = 31 - __builtin_clz(P.size()), n = 1 << lg;</pre>
       // == P.size();
  for (int j = 1, i = 0; j < n - 1; ++j) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1);
     if (j < i) swap(P[i], P[j]);</pre>
  } //bit reverse
  LL w1 = powM(root, a * (inv? p - 2: 1)); // order is
       1<<lg
  for (LL i = 1; i \le lg; ++i)
    LL wn = powM(w1, 1 << (lg - i)); // order is 1 << i
    for (int k = 0; k < (1 << lg); k += 1 << i) {
       LL base = 1;
       for (int j = 0; j < (1 << i - 1); ++j, base = base * wn % p) {
LL t = base * P[k + j + (1 << i - 1)] % p;
         LL u = P[k + j] \% p;
         P[k + j] = (u + t) \% p
         P[k + j + (1 \ll i - 1)] = (u - t + p) \% p;
       }
    }
  if(inv){
    LL invN = powM(n, p - 2);
    transform(P.begin(), P.end(), P.begin(), [&](LL a)
         {return a * invN % p;});
  return P:
} //faster performance with calling by reference
```

## 2.4 permanent

```
typedef vector<vector<LL> > mat;
LL permanent(mat A) {
   LL n = A.size(), ans = 0, *tmp = new LL[n], add;
   for (int pgray = 0, s = 1, gray, i; s < 1 << n; ++s)
      {
       gray = s ^ s >> 1, add = 1;
       i = __builtin_ctz(pgray ^ gray);
       for (int j = 0; j < n; ++j)
            add *= tmp[j] += A[i][j] * (gray>>i&1 ? 1 : -1);
       ans += add * (s&1^n&1? -1 : 1), pgray = gray;
```

```
return ans;
}
// how many ways to put rooks on a matrix with 0,1 as
    constrain
// 1 - ok to put
// 0 - not ok to put
```

## 3 Data Structure

## 3.1 Heavy Light Decomposition

```
struct HLD {
  using Tree = vector<vector<int>>>;
   vector<int> par, head, vid, len, inv;
  HLD(const Tree &g) : par(g.size()), head(g.size()),
    vid(g.size()), len(g.size()), inv(g.size()) {
     int k = 0;
     vector<int> size(g.size(), 1);
     function<void(int, int)> dfs_size = [&](int u, int
       for (int v : g[u]) {
         if (v != p) {
           dfs_size(v, u);
            size[u] += size[v];
       }
     };
     function<void(int, int, int)> dfs_dcmp = [&](int u,
           int p, int h) {
       par[u] = p;
       head[u] = h;
       vid[u] = k++;
       inv[vid[u]] = u;
       for (int v : g[u]) {
         if (v != p && size[u] < size[v] * 2) {</pre>
            dfs_dcmp(v, u, h);
         }
       for (int v : g[u]) {
         if (v != p && size[u] >= size[v] * 2) {
           dfs_dcmp(v, u, v);
       }
     dfs_size(0, -1);
     dfs_dcmp(0, -1, 0);
for (int i = 0; i < g.size(); ++i) {</pre>
       ++len[head[i]];
   template<typename T>
   void foreach(int u, int v, T f) {
     while (true) {
       if (vid[u] > vid[v])
         if (head[u] == head[v])
            f(vid[v] + 1, vid[u], 0);
           break
            f(vid[head[u]], vid[u], 1);
            u = par[head[u]];
       } else {
         if (head[u] == head[v]) +
            f(vid[u] + 1, vid[v], 0);
           break;
         } else
            f(vid[head[v]], vid[v], 0);
            v = par[head[v]];
       }
    }
  }
|};
```

### 4 Flow

## 4.1 CostFlow

```
template <class TF, class TC>
struct CostFlow {
   static const int MAXV = 205;
   static const TC INF = 0x3f3f3f3f3f;
   struct Edge {
     int v, r;
     TF f;
     TC c;
     Edge(int _v, int _r, TF _f, TC _c) : v(_v), r(_r), f(_f), c(_c) {}
   int n, s, t, pre[MAXV], pre_E[MAXV], inq[MAXV];
  TF fl;
  TC dis[MAXV], cost;
  vector<Edge> E[MAXV];
  CostFlow(int _n, int _s, int _t) : n(_n), s(_s), t(_t
     ), fl(0), cost(0) {}
  void add_edge(int u, int v, TF f, TC c) {
    E[u].emplace_back(v, E[v].size(), f, c);
    E[v].emplace_back(u, E[u].size() - 1, 0, -c);
  pair<TF, TC> flow() {
  while (true) {
    for (int i = 0; i < n; ++i) {</pre>
          dis[i] = INF;
          inq[i] = 0;
       dis[s] = 0;
       queue<int> que;
       que.emplace(s);
       while (not que.empty()) {
          int u = que.front();
          que.pop();
          inq[u] = 0;
          for (int i = 0; i < E[u].size(); ++i) {
  int v = E[u][i].v;</pre>
             TC w = E[\bar{u}][\bar{i}].c;
             if (E[u][i].f > 0 and dis[v] > dis[u] + w) {
               pre[v] = u;
               pre_E[v] = i;
               dis[v] = dis[u] + w;
if (not inq[v]) {
                  inq[v] = 1;
                  que.emplace(v);
            }
          }
        if (dis[t] == INF) break;
       TF tf = INF;
        for (int v = t, u, l; v != s; v = u) {
          u = pre[v];
          l = pre_E[v];
          tf = min(tf, E[u][l].f);
        for (int v = t, u, l; v != s; v = u) {
          u = pre[v];
          l = pre_E[v];
          E[u][l].f -= tf;
          E[v][E[u][l].r].f += tf;
        cost += tf * dis[t];
       fl += tf;
     return {fl, cost};
};
```

#### 4.2 Dinic

```
template <class T>
struct Dinic {
  static const int MAXV = 10000;
  static const T INF = 0x3f3f3f3f3f;
```

```
struct Edge {
     int v;
     Tf;
     int re;
     Edge(int _v, T _f, int _re) : v(_v), f(_f), re(_re)
   int n, s, t, level[MAXV];
   vector<Edge> E[MAXV];
   int now[MAXV];
   Dinic(int _n, int _s, int _t) : n(_n), s(_s), t(_t)
       {}
   void add_edge(int u, int v, T f, bool bidirectional =
        false) {
     E[u].emplace_back(v, f, E[v].size());
     E[v].emplace_back(u, 0, E[u].size() - 1);
     if (bidirectional) {
       E[v].emplace_back(u, f, E[u].size() - 1);
   bool BFS() {
     memset(level, -1, sizeof(level));
     queue<int> que;
     que.emplace(s);
     level[s] = 0;
     while (not que.empty()) {
       int u = que.front();
       que.pop();
       for (auto it : E[u]) {
   if (it.f > 0 and level[it.v] == -1) {
     level[it.v] = level[u] + 1;
           que.emplace(it.v);
       }
     return level[t] != -1;
   T DFS(int u, T nf) {
     if (u == t) return nf;
     T res = 0;
     while (now[u] < E[u].size()) {</pre>
       Edge &it = E[u][now[u]];
if (it.f > 0 and level[it.v] == level[u] + 1) {
         T tf = DFS(it.v, min(nf, it.f));
         res += tf;
         nf -= tf;
         it.f -= tf;
         E[it.v][it.re].f += tf;
         if (nf == 0) return res;
       } else
         ++now[u];
     if (not res) level[u] = -1;
     return res;
   T flow(T res = 0) {
     while (BFS()) {
       T temp;
       memset(now, 0, sizeof(now))
       while (temp = DFS(s, INF)) {
         res += temp;
         res = min(res, INF);
       }
     }
     return res;
|};
      Geometry
```

# 5.1 2D Geometry

```
namespace geo {
  using pt = complex<double>;
  using cir = pair<pt, double>;
  const double EPS = 1e-4;
  const double PI = acos(-1);
  pt cent(cir C) { return C.first; }
```

```
p.erase(unique(p.begin(), p.end()), p.end());
double radi(cir C) { return C.second; }
int dcmp(double x) {
                                                                int n = p.size(), m = 0;
  if (abs(x) < EPS) return 0;
                                                               vector<pt> ch(n + 1);
  return x > 0 ? 1 : -1;
                                                                    is cleared
bool less(pt a, pt b) {
  return real(a) < real(b) || real(a) == real(b) &&</pre>
                                                                      <= 0) --m;
      imag(a) < imag(b);
                                                                 ch[m++] = p[i];
bool more(pt a, pt b) {
  return real(a) > real(b) || real(a) == real(b) &&
      imag(a) > imag(b);
                                                                      <= 0) --m;
                                                                  ch[m++] = p[i];
double dot(pt a, pt b) { return real(conj(a) * b); }
double cross(pt a, pt b) { return imag(conj(a) * b);
                                                               return ch;
double sarea(pt a, pt b, pt c) { return cross(b - a,
    c - a); }
int ori(pt a, pt b, pt c) { return dcmp(sarea(a, b, c
                                                               pt B = tri[1] - tri[0];
                                                               pt C = tri[2] - tri[0];
    )); }
                                                               double det = 2 * cross(B, C);
double angle(pt a, pt b) { return acos(dot(a, b) /
    abs(a) / abs(b)); }
pt rotate(pt a, double rad) { return a * pt(cos(rad),
     sin(rad));
                                                                              det:
pt normal(pt a) { return pt(-imag(a) / abs(a), real(a
                                                               return {r + tri[0], abs(r)};
    ) / abs(a)); }
pt normalized(pt a) { return a / abs(a); }
pt get_line_intersection(pt p, pt v, pt q, pt w) {
   // L = p + t * v, J = q + s * w
                                                               assert(tri.size() == 3);
                                                               pt ans = 0;
  assert(dcmp(cross(v, w)));
                                                                double div = 0;
  return p + v * cross(w, p - q) / cross(v, w);
                                                                for (int i = 0; i < 3; ++i) {
double distance_to_line(pt p, pt a, pt b) {
                                                                 3]);
ans += l * tri[i], div += l;
  // the line contains two distinct points a,
  return abs(cross(b - a, p - a) / abs(b - a));
                                                               ans /= div;
double distance_to_segment(pt p, pt a, pt b) {
  // similar to previous function
  if (a == b) return abs(p - a);
  pt v1(b - a), v2(p - a), v3(p - b);
  if (dcmp(dot(v1, v2)) < 0) return abs(v2);
else if (dcmp(dot(v1, v3)) > 0) return abs(v3);
  return abs(cross(v1, v2)) / abs(v1);
                                                                    return {p};
pt get_line_projection(pt p, pt a, pt b) {
  pt v = b - a;
  return a + dot(v, p - a) / dot(v, v) * v;
bool segment_proper_intersection(pt a1, pt a2, pt b1,
     pt b2) {
  double det1 = ori(a1, a2, b1) * ori(a1, a2, b2);
  double det2 = ori(b1, b2, a1) * ori(b1, b2, a2);
                                                                  pt ed) {
  return det1 < 0 && det2 < 0;
double polygon_area(vector<pt> p) {
  double area = 0;
  for (int i = 1; i < p.size() - 1; ++i) {</pre>
    area += sarea(p[0], p[i], p[i + 1]);
                                                                    * r = 0
  return area / 2;
                                                               if (dcmp(det) < 0) return {};</pre>
bool is_point_on_segment(pt p, pt a1, pt a2) {
  return dcmp(sarea(p, a1, a2)) == 0 && dcmp(dot(a1 -
       p, a2 - p) < 0;
int is_point_in_polygon(pt p, vector<pt> gon) {
  int wn = 0;
                                                                  d) {
  int n = gon.size();
                                                               auto a = cent(c), b = cent(d);
  for (int i = 0; i < n; ++i) {
    if (is_point_on_segment(p, gon[i], gon[(i + 1) %
        n])) return -1;
    int k = dcmp(cross(gon[(i + 1) % n] - gon[i], p -
                                                                    0) return {};
         gon[i]))
    int d1 = dcmp(imag(gon[i]) - imag(p));
    int d2 = dcmp(imag(gon[(i + 1) % n] - imag(p)));
    wn += k > 0 \&\& d1 <= 0 \&\& d2 > 0;
    wn -= k < 0 \&\& d2 <= 0 \&\& d1 > 0;
  return wn != 0;
                                                          };
vector<pt> convex_hull(vector<pt> p) {
  sort(p.begin(), p.end(), less);
```

```
for (int i = 0; i < n; ++i) { // note that border
    while (m > 1 \&\& ori(ch[m - 2], ch[m - 1], p[i])
  for (int i = n - 2, k = m; i >= 0; --i) {
   while (m > k \& ori(ch[m - 2], ch[m - 1], p[i])
 ch.erase(ch.begin() + m - (n > 1), ch.end());
cir circumscribed_circle(vector<pt> tri) {
  pt r = pt(imag(C) * norm(B) - imag(B) * norm(C)
            real(B) * norm(C) - real(C) * norm(B)) /
cir inscribed_circle(vector<pt> tri) {
    double l = abs(tri[(i + 1)\% 3] - tri[(i + 2)\%]
  return {ans, distance_to_line(ans, tri[0], tri[1])
vector<pt> tangent_line_through_point(cir c, pt p) {
  if (dcmp(abs(cent(c) - p) - radi(c)) < 0) return</pre>
  {};
else if (dcmp(abs(cent(c) - p) - radi(c)) == 0)
  auto theta = acos(radi(c) / abs(cent(c) - p));
 auto norm_v = normalized(p - cent(c));
return {cent(c) + radi(c) * rotate(norm_v, +theta),
          cent(c) + radi(c) * rotate(norm_v, -theta)
vector<pt> get_line_circle_intersection(cir d, pt st,
  pt v = ed - st, p = st - cent(d);
 // t^2 * norm(v) + 2 * t * dot(p, v) + norm(p) - r
  auto get_point = [=](double t) { return st + t * v;
  if (dcmp(det) == 0) return {get_point(-b / 2 / a)};
  return {get_point((-b + sqrt(det)) / 2 / a)
          get_point((-b - sqrt(det)) / 2 / a);
vector<pt> get_circle_circle_intersection(cir c, cir
  auto r = radi(c), s = radi(d), g = abs(a - b);
  if (dcmp(g) == 0) return \{\}; // may be C == D
  if (dcmp(r + s - g) < 0 \text{ or } dcmp(abs(r - s) - g) >
  auto C_to_D = normalized(b - a);
  auto theta = acos((r * r + g * g - s * s) / (2 * r)
  if (dcmp(theta) == 0) return {a + r * C_to_D};
 else return {a + rotate(r * C_to_D, theta), a +
    rotate(r * C_to_D, -theta)};
```

# 6 Graph

# 6.1 2-SAT

```
#include <cstdio>
#include <vector>
#include <stack>
#include <cstring>
using namespace std;
const int N = 2010;
struct two_SAT {
  int n:
  vector<int> G[N], revG[N];
  stack<int> finish;
  bool sol[N], visit[N];
  int cmp[N];
  void init(int _n) {
    n = _n;
     for (int i = 0; i < N; i++) {
       G[i].clear();
       revG[i].clear();
    }
  }
  void add_edge(int u, int v) {
    // 2 * i -> i is True, 2 * i + 1 -> i is False
     G[u].push_back(v)
     G[v^1].push_back(u^1);
     revG[v].push_back(u);
     revG[u^1].push_back(v^1);
  void dfs(int v) {
     visit[v] = true;
     for ( auto i:G[v] ) {
       if ( !visit[i] ) dfs(i);
     finish.push(v);
  }
  void revdfs(int v, int id) {
    visit[v] = true;
for ( auto i:revG[v] ) {
       if ( !visit[i] ) revdfs(i,id);
     cmp[v] = id;
  int scc() {
    memset( visit, 0, sizeof(visit) );
for (int i = 0; i < 2 * n; i++) {</pre>
       if ( !visit[i] ) dfs(i);
     int id = 0:
    memset( visit, 0, sizeof(visit) );
while ( !finish.empty() ) {
  int v = finish.top(); finish.pop();
       if ( visit[v] ) continue;
       revdfs(v,++id);
     }
     return id;
  bool solve() {
     scc();
     for (int i = 0; i < n; i++) {
  if ( cmp[2*i] == cmp[2*i+1] ) return 0;</pre>
       sol[i] = (cmp[2*i] > cmp[2*i+1]);
     return 1;
  }
} sat;
int main() {
  // ( a or not b ) and ( b or c ) and ( not c or not a
  sat.init(3);
  sat.add_edge( 2*0+1, 2*1+1 );
sat.add_edge( 2*1+1, 2*2+0 );
sat.add_edge( 2*2+0, 2*0+1 );
printf("%d\n", sat.solve() );
  return 0;
```

# 6.2 maximal cliques

```
#include <bits/stdc++.h>
using namespace std;
const int N = 60;
typedef long long LL;
struct Bron_Kerbosch {
  int n, res;
  LL edge[N];
  void init(int _n) {
    n = _n;
    for (int i = 0; i \le n; i++) edge[i] = 0;
  void add_edge(int u, int v) {
    if ( u == v ) return;
    edge[u] |= 1LL << v;
    edge[v] l= 1LL \ll u;
  void go(LL R, LL P, LL X) {
    if (P == 0 & X == 0)
      res = max( res, __builtin_popcountll(R) ); //
          notice LL
      return;
    if (
           _builtin_popcountll(R) + __builtin_popcountll
        (P) <= res ) return;
    for (int i = 0; i <= n; i++) {
      LL v = 1LL << i;
      if (P&v) {
        go( R | v, P & edge[i], X & edge[i] );
        P &= ~v;
        X \mid = v;
      }
    }
  int solve() {
    res = 0;
    go( 0LL, ( 1LL << (n+1) ) - 1, 0LL );
    return res;
   BronKerbosch1(R, P, X):
      if P and X are both empty:
        report R as a maximal clique
      for each vertex v in P:
        BronKerbosch1(R \square {v}, P \square N(v), X \square N(v))
        P := P \setminus \{v\}
        X := X \square \{v\}
} MaxClique;
int main() {
  MaxClique.init(6);
  MaxClique.add_edge(1,2);
  MaxClique.add_edge(1,5);
  MaxClique.add_edge(2,5);
  MaxClique.add_edge(4,5);
  MaxClique.add_edge(3,2);
  MaxClique.add_edge(4,6);
  MaxClique.add\_edge(3,4);
  cout << MaxClique.solve() << "\n";</pre>
  return 0;
```

# 6.3 Tarjan SCC

```
#include <cstdio>
#include <vector>
#include <stack>
#include <cstring>
using namespace std;

const int N = 10010;
struct Tarjan {
  int n;
  vector<int> G[N], revG[N];
  stack<int> finish;
  bool visit[N];
```

```
int cmp[N]:
  void init(int _n) {
    n = _n;
    for (int i = 0; i <= n; i++) {
      G[i].clear();
      revG[i].clear();
  }
  void add_edge(int u, int v) {
    G[u].push_back(v)
    revG[v].push_back(u);
  void dfs(int v) {
    visit[v] = true;
    for ( auto i:G[v] ) {
      if ( !visit[i] ) dfs(i);
    finish.push(v);
  void revdfs(int v, int id) {
    visit[v] = true;
    for ( auto i:revG[v] ) {
      if ( !visit[i] ) revdfs(i,id);
    cmp[v] = id;
  int solve() {
    memset( visit, 0, sizeof(visit) );
    for (int i = 0; i < n; i++) {
      if ( !visit[i] ) dfs(i);
    int id = 0;
    memset( visit, 0, sizeof(visit) );
while ( !finish.empty() ) {
      int v = finish.top(); finish.pop();
      if ( visit[v] ) continue;
      revdfs(v,++id);
    return id;
  }
} scc;
int main() {
  int V, É;
scanf("%d %d", &V, &E);
  scc.init(V);
  for (int i = 0; i < E; i++) {
    int u, v;
scanf("%d %d", &u, &v);
    scc.add_edge(u-1,v-1);
  printf("%d\n", scc.solve() );
  return 0;
```

# 7 Number Theory

#### 7.1 basic

```
PLL exd_gcd(LL a, LL b) {
 if (a\% b == 0) return \{0, 1\};
 PLL T = exd\_gcd(b, a \% b);
 return {T.second, T.first - a / b * T.second};
LL mul(LL x, LL y, LL mod) {
 LL ans = 0, m = x, s = 0, sgn = (x > 0) xor (y > 0)?
      -1: 1;
  for (x = abs(x), y = abs(y); y; y >>= 1, m <<= 1, m =
       m >= mod? m - mod: m)
    if (y&1) s += m, s = s >= mod? s - mod: s;
 return s * sgn;
LL dangerous_mul(LL a, LL b, LL mod){ // 10 times
    faster than the above in average, but could be
    prone to wrong answer (extreme low prob?)
  return (a * b - (LL)((long double)a * b / mod) * mod)
       % mod;
```

```
LL powmod(LL x, LL p, LL mod) {
  LL s = 1, m = x % mod;
  for (; p; m = mul(m, m, mod), p >>= 1)
    if (p&1) s = mul(s, m, mod);
  return s;
}
```

### 7.2 Chinese Remainder Theorem

```
typedef long long LL;
typedef pair<LL, LL> PLL;
PLL exd_gcd(LL a, LL b);
LL CRT(vector<PLL> &eqs) {
    LL prod = 1, ans = 0, ni, ns;
    for (auto P: eqs) prod *= P.second;
    for (auto P: eqs) {
        ni = P.second, ns = prod / ni;
        (ans += ns * P.first % prod * exd_gcd(ni, ns).
            second) %= prod;
    }
    return (ans + prod) % prod;
}
```

# 7.3 Discrete Log

#### 7.4 Lucas

# 7.5 Meissel-Lehmer PI

```
LL PI(LL m);
const int MAXM = 1000, MAXN = 650, UPBD = 1000000;
// 650 ~ PI(cbrt(1e11))
LL pi[UPBD] = {0}, phi[MAXM][MAXN];
vector<LL> primes;
void init() {
  fill(pi + 2, pi + UPBD, 1);
  for (LL p = 2; p < UPBD; ++p)
    if (pi[p]) {
     for (LL N = p * p; N < UPBD; N += p)
        pi[N] = 0;
     primes.push_back(p);</pre>
```

```
for (int i = 1; i < UPBD; ++i) pi[i] += pi[i - 1];
for (int i = 0; i < MAXM; ++i)</pre>
    phi[i][0] = i;
  for (int i = 1; i < MAXM; ++i)
    LL P_2(LL m, LL n) {
  LL ans = 0;
  for (LL i = n; primes[i] * primes[i] <= m and i <</pre>
      primes.size(); ++i)
    ans += PI(m / primes[i]) - i;
  return ans;
LL PHI(LL m, LL n) {
  if (m < MAXM and n < MAXN) return phi[m][n];</pre>
  if (n == 0) return m;
  LL p = primes[n - 1];
  if (m < UPBD) {
    if (m <= p) return 1;
if (m <= p * p * p) return pi[m] - n + 1 + P_2(m, n)</pre>
  return PHI(m, n - 1) - PHI(m / p, n - 1);
LL PI(LL m) {
  if (m < UPBD) return pi[m];</pre>
  LL y = cbrt(m) + 10, n = pi[y]
  return PHI(m, n) + n - 1 - P_2(m, n);
```

## 7.6 Miller Rabin with Pollard rho

```
// Miller_Rabin
LL abs(LL a) {return a > 0? a: -a;}
bool witness(LL a, LL n, LL u, int t) {
  LL x = modpow(a, u, n), nx;
  for (int i = 0; i < t; ++i, x = nx){
    nx = mul(x, x, n);
     if (nx == 1 \text{ and } x != 1 \text{ and } x != n - 1) \text{ return 1};
  return x != 1;
const LL wits[7] = \{2, 325, 9375, 28178, 450775,
9780504, 1795265022};
bool miller_rabin(LL n, int s = 7) {
  if (n < 2) return 0;
  if (n\&1^1) return n == 2;
  LL u = n - 1, t = 0, a; // n == (u << t) + 1 while <math>(u\&1^1) u >>= 1, ++t;
  while (s--)
    if (a = wits[s] % n and witness(a, n, u, t)) return
  return 1;
// Pollard_rho
LL f(LL x, LL n) {
  return mul(x, x, n) + 1;
LL pollard_rho(LL n) {
  if (n&1^1) return 2;
  while (true) {
    LL x = rand() \% (n - 1) + 1, y = 2, d = 1;
     for (int sz = 2; d == 1; y = x, sz <<= 1)
       for (int i = 0; i < sz and d <= 1; ++i)
     x = f(x, n), d = \_gcd(abs(x - y), n);
if (d and n - d) return d;
  }
}
```

#### 7.7 Primitive Root

```
vector<LL> factor(LL N) {
  vector<LL> ans;
  for (LL p = 2, n = N; p * p <= n; ++p)
   if (N % p == 0) {
     ans.push_back(p);
}</pre>
```

```
while (N \% p == 0) N /= p;
  if (N != 1) ans.push_back(N);
  return ans;
LL find_root(LL p) {
  LL ans = 1;
  for (auto q: factor(p - 1)) {
    LL a = rand() \% (p - 1) + 1, b = (p - 1) / q;
    while (powmod(a, b, p) == 1) a = rand() % (p - 1) +
    while (b % q == 0) b /= q;
    ans = mul(ans, powmod(a, b, p), p);
  return ans;
}
bool is_root(LL a, LL p) {
  for (auto q: factor(p - 1))
    if (powmod(a, (p - 1) / q, p) == 1)
      return false;
  return true;
```

# 8 String

# 8.1 AC automaton

```
// SIGMA[0] will not be considered
const string SIGMA =
     _0123456789ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrs
vector<int> INV_SIGMA;
const int SGSZ = 63;
struct PMA {
  PMA *next[SGSZ]; // next[0] is for fail
  vector<int> ac;
  PMA *last; // state of longest accepted string that
       is pre of this
  PMA(): last(nullptr) { fill(next, next + SGSZ,
       nullptr); }
};
template<typename T>
PMA *buildPMA(const vector<T> &p) {
  PMA *root = new PMA;
  for (int i = 0; i < p.size(); ++i) { // make trie</pre>
    PMA *t = root;
    for (int j = 0; j < p[i].size(); ++j) {
  int c = INV_SIGMA[p[i][j]];
   ___</pre>
      if (t->next[c] == nullptr) t->next[c] = new PMA;
      t = t->next[c];
    t->ac.push_back(i);
  queue<PMA *> que; // make failure link using bfs
  for (int c = 1; c < SGSZ; ++c) {</pre>
    if (root->next[c]) {
      root->next[c]->next[0] = root;
       que.push(root->next[c]);
    } else root->next[c] = root;
  while (!que.empty()) {
    PMA *t = que.front();
    que.pop();
    for (int c = 1; c < SGSZ; ++c) {
  if (t->next[c]) {
         que.push(t->next[c]);
         PMA *r = t->next[0];
         while (!r->next[c]) r = r->next[0];
         t->next[c]->next[0] = r->next[c];
t->next[c]->last = r->next[c]->ac.size() ? r->
             next[c] : r->next[c]->last;
    }
  return root;
```

```
void destructPMA(PMA *root) {
  queue<PMA *> que;
  que.emplace(root)
  while (!que.empty()) {
    PMA *t = que.front();
    que.pop();
    for (int c = 1; c < SGSZ; ++c) {
      if (t->next[c] && t->next[c] != root) que.emplace
           (t->next[c]);
    delete t:
  }
}
template<typename T>
map<int, int> match(const T &t, PMA *v) {
  map<int, int> res;
  for (int i = 0; i < t.size(); ++i) {</pre>
    int c = INV_ŚIGMA[t[i]];
    while (!v-\text{next}[c]) v = v-\text{next}[0];
    v = v->next[c];
    for (int j = 0; j < v -> ac.size(); ++j) ++res[v -> ac[
    for (PMA *q = v->last; q; q = q->last) {
      for (int j = 0; j < q->ac.size(); ++j) ++res[q->
    }
  }
  return res;
}
```

## 8.2 Gusfield

```
template<typename T>
vector<int> gusfield(const T &s) {
 vector<int> z(s.size(), s.size()); // z[i] := max k
    for z[0, k) = z[i, i + k)
  for (int i = 1, L = 0, R = 0; i < s.size(); ++i) {
    if (R < i) {
      L = R = i;
      while (R < s.size() \&\& s[R] == s[R - L]) ++R;
      z[i] = R - L;
      --R;
    } else {
      int k = i - L;
      if(z[k] < R - i + 1) {
        z[i] = z[k];
      } else {
        while (R < s.size() \&\& s[R] == s[R - L]) ++R;
        z[i] = R - L;
         --R;
      }
    }
  return z;
```

#### 8.3 KMP

```
template<typename T>
vector<int> build_kmp(const T &s) {
  vector<int> f(s.size());
  int fp = f[0] = -1;
  for (int i = 1; i < s.size(); ++i) {
    while (~fp && s[fp + 1] != s[i]) fp = f[fp];
    if (s[fp + 1] == s[i]) ++fp;
    f[i] = fp;
  }
  return f;
}</pre>
```

## 8.4 Manacher

# 8.5 Suffix Array

```
// -----O(NlgNlgN)-----
vector<int> sa_db(const string &s) {
  int n = s.size();
  vector<int> sa(n), r(n), t(n);
  for (int i = 0; i < n; ++i) r[sa[i] = i] = s[i];
for (int h = 1; t[n - 1] != n - 1; h *= 2) {
    auto cmp = [&](int i, int j) {
  if (r[i] != r[j]) return r[i] < r[j];</pre>
      return i + h < n \& j + h < n ? r[i + h] < r[j +
           h] : i > j;
    sort(sa.begin(), sa.end(), cmp);
    for (int i = 0; i + 1 < n; ++i) t[i + 1] = t[i] +
         cmp(sa[i], sa[i + 1]);
    for (int i = 0; i < n; ++i) r[sa[i]] = t[i];
  return sa;
}
// O(N) -- CF: 1e6->31ms,18MB;1e7->296ms;158MB;3e7->856
    ms,471MB
bool is_lms(const string &t, int i) {
  return i > 0 && t[i - 1] == 'L' && t[i] == 'S';
template<typename T>
vector<int> induced_sort(const T &s, const string &t,
    const vector<int> &lmss, int sigma = 256) {
  vector<int> sa(s.size(), -1);
  vector<int> bin(sigma + 1);
  for (auto it = s.begin(); it != s.end(); ++it) {
    ++bin[*it + 1];
  int sum = 0;
  for (int i = 0; i < bin.size(); ++i) {</pre>
    sum += bin[i];
    bin[i] = sum;
  vector<int> cnt(sigma);
  for (auto it = lmss.rbegin(); it != lmss.rend(); ++it
    int ch = s[*it];
sa[bin[ch + 1] - 1 - cnt[ch]] = *it;
    ++cnt[ch];
  cnt = vector<int>(sigma);
  for (auto it = sa.begin(); it != sa.end(); ++it) {
    if (*it <= 0 || t[*it - 1] == 'S') continue;
int ch = s[*it - 1];
    sa[bin[ch] + cnt[ch]] = *it - 1;
    ++cnt[ch];
  cnt = vector<int>(sigma);
  for (auto it = sa.rbegin(); it != sa.rend(); ++it) {
```

```
if (*it <= 0 || t[*it - 1] == 'L') continue;
         int ch = s[*it - \bar{1}];
         sa[bin[ch + 1] - 1 - cnt[ch]] = *it - 1;
         ++cnt[ch];
    return sa;
}
template<typename T>
vector<int> sa_is(const T &s, int sigma = 256) {
    string t(s.size(), 0);
     t[s.size() - 1] = 'S'
    for (int i = int(s.size()) - 2; i >= 0; --i) {
  if (s[i] < s[i + 1]) t[i] = 'S';
}</pre>
         else if (s[i] > s[i + 1]) t[i] = 'L';
         else t[i] = t[i + 1];
    vector<int> lmss;
    for (int i = 0; i < s.size(); ++i) {
        if (is_lms(t, i)) {
             lmss.emplace_back(i);
    vector<int> sa = induced_sort(s, t, lmss, sigma);
    vector<int> sa_lms;
    for (int i = 0; i < sa.size(); ++i) {
  if (is_lms(t, sa[i])) {</pre>
              sa_lms.emplace_back(sa[i]);
    int lmp_ctr = 0;
    vector<int> lmp(s.size(), -1);
    lmp[sa_lms[0]] = lmp_ctr;
     for (int i = 0; i + 1 < sa_lms.size(); ++i) {</pre>
         int diff = 0;
         for (int d = 0; d < sa.size(); ++d) {</pre>
             if (s[sa_lms[i] + d] != s[sa_lms[i + 1] + d] ||
                       is_{ms}(t, sa_{ms}[i] + d) != is_{
                                i + 1] + d)) {
                  diff = 1; // something different in range of
                            lms
             } else if (d > 0 && is_lms(t, sa_lms[i] + d) &&
                       is_{ms}(t, sa_{ms}[i + 1] + d)) {
                  break; // exactly the same
             }
         if (diff) ++lmp_ctr;
         lmp[sa_lms[i + 1]] = lmp_ctr;
    vector<int> lmp_compact;
     for (int i = 0; i < lmp.size(); ++i) {</pre>
         if (~lmp[i]) {
             lmp_compact.emplace_back(lmp[i]);
    if (lmp_ctr + 1 < lmp_compact.size()) {</pre>
         sa_lms = sa_is(lmp_compact, lmp_ctr + 1);
    } else {
         for (int i = 0; i < lmp_compact.size(); ++i) {</pre>
              sa_lms[lmp_compact[i]] = i;
    vector<int> seed;
    for (int i = 0; i < sa_lms.size(); ++i) {</pre>
        seed.emplace_back(lmss[sa_lms[i]]);
    return induced_sort(s, t, seed, sigma);
} // s must end in char(0)
// O(N) lcp, note that s must end in '\0'
vector<int> build_lcp(const string &s, const vector<int</pre>
         > &sa, const vector<int> &rank) {
     int n = s.size();
```

```
vector<int> lcp(n);
for (int i = 0, h = 0; i < n; ++i) {
   if (rank[i] == 0) continue;</pre>
     int j = sa[rank[i] - 1];
     if (h > 0) --h;
    for ( ; j + h < n && i + h < n; ++h) {
  if (s[j + h] != s[i + h]) break;</pre>
     lcp[rank[i] - 1] = h;
  return lcp; // lcp[i] := lcp(s[sa[i]..-1], s[sa[i +
       17...-17)
}
// O(N) build segment tree for lcp
vector<int> build_lcp_rmq(const vector<int> &lcp) {
  vector<int> sgt(lcp.size() << 2);</pre>
  function<void(int, int, int)> build = [&](int t, int
    lb, int rb) {
     if (rb - lb == 1) return sgt[t] = lcp[lb], void();
     int mb = lb + rb \gg 1;
    build(t << 1, lb, mb);
build(t << 1 | 1, mb, rb);
sgt[t] = min(sgt[t << 1], sgt[t << 1 | 1]);</pre>
  build(1, 0, lcp.size());
  return sgt;
// O(IPI + lg ITI) pattern searching, returns last
     index in sa
int match(const string &p, const string &s, const
     vector<int> &sa, const vector<int> &rmq) { // rmq
     is segtree on lcp
  int t = 1, lb = 0, rb = s.size(); // answer in [lb,
       rb)
  int lcplp = 0; // lcp(char(0), p) = 0
  while (rb - lb > 1) {
     int mb = 1b + rb >> 1
     int lcplm = rmq[t << 1];</pre>
     if (lcplp < lcplm) t = t << 1 | 1, lb = mb;
     else if (lcplp > lcplm) t = t << 1, rb = mb;</pre>
     else {
       int lcpmp = lcplp;
       while (lcpmp < p.size() && p[lcpmp] == s[sa[mb] +
             lcpmp]) ++lcpmp;
       if (lcpmp == p.size() || p[lcpmp] > s[sa[mb] +
            lcpmp]) t = t << 1 | 1, lb = mb, lcplp =
            lcomp:
       else t = t << 1, rb = mb;
    }
  if (lcplp < p.size()) return -1;</pre>
  return sa[lb];
```

#### 8.6 Suffix Automaton

```
template<typename T>
struct SuffixAutomaton {
  vector<map<int, int>> edges;// edges[i] : the
      labeled edges from node i
  vector<int> link;
                              // link[i]
                                           : the parent
       of i
  vector<int> length;
                              // length[i] : the length
       of the longest string in the ith class
  int last;
                              // the index of the
      equivalence class of the whole string
  vector<bool> is_terminal;
                              // is_terminal[i] : some
      suffix ends in node i (unnecessary)
  vector<int> occ;
                              // occ[i] : number of
      matches of maximum string of node i (unnecessary)
  SuffixAutomaton(const T &s) : edges({map<int, int>()
      }), link({-1}), length({0}), last(0), occ({0}) {
    for (int i = 0; i < s.size(); ++i) {
      edges.push_back(map<int, int>());
      length.push_back(i + 1);
      link.push_back(0);
      occ.push_back(1);
      int r = edges.size() - 1;
```

```
int p = last; // add edges to r and find p with
           link to q
       while (p \ge 0 \& edges[p].find(s[i]) == edges[p].
           end()) {
         edges[p][s[i]] = r;
         p = link[p];
       if (~p) {
         int q = edges[p][s[i]];
         if (length[p] + 1 == length[q]) { // no need to}
              split q
           link[r] = q;
         } else { // split q, add qq
           edges.push_back(edges[q]); // copy edges of
           length.push_back(length[p] + 1);
           link.push_back(link[q]); // copy parent of q
           occ.push_back(0);
           int qq = edges.size() - 1; // qq is new
                parent of q and r
           link[q] = qq;
           link[r] = qq;
           while (p >= 0 \&\& edges[p][s[i]] == q) { // }
               what points to q points to qq
             edges[p][s[i]] = qq;
             p = link[p];
        }
       last = r;
    } // below unnecessary
    is_terminal = vector<bool>(edges.size());
    for (int p = last; p > 0; p = link[p]) is_terminal[
    p] = 1; // is_terminal calculated
    vector<int> cnt(link.size()), states(link.size());
         // sorted states by length
    for (int i = 0; i < link.size(); ++i) ++cnt[length[</pre>
         i]];
    for (int i = 0; i < s.size(); ++i) cnt[i + 1] +=</pre>
         cnt[i];
    for (int i = link.size() - 1; i >= 0; --i) states
    [--cnt[length[i]]] = i;
for (int i = link.size() - 1; i >= 1; --i) occ[link
         [states[i]]] += occ[states[i]]; // occ
         calculated
};
```

## **Formulas**

#### 9.0.1 Pick's theorem

Pick's theorem provides a simple formula for calculating the area A of this polygon in terms of the number i of lattice points in the interior located in the polygon and the number b of lattice points on the boundary placed on the polygon's

$$A = i + \frac{b}{2} - 1$$

#### 9.0.2 Graph Properties

- Euler's Formula V-E+F=2 For a planar graph, F=E-V+n+1, n is the numbers of components For a planar graph,  $E\leq 3V-6$  For a connected graph G, let I(G) be the size of maximum independent set,  $M(\mathcal{G})$  be the size of maximum matching, Cv(G) be the size of minimum vertex cover, Ce(G) be the size of minimum edge cover.
- 4. For any connected graph:

$$\begin{array}{ll} \text{(a)} & I(G)+Cv(G)=|V| \\ \text{(b)} & M(G)+Ce(G)=|V| \end{array}$$

5. For any bipartite:

$$\begin{array}{ll} \text{(a)} & I(G) = Cv(G) \\ \text{(b)} & M(G) = Ce(G) \end{array}$$

```
double l=0,=m,stop=1.0/n/n;
while(r-l>=stop){
  double(mid);
  if((n*m-sol.maxFlow(s,t))/2>eps)l=mid;
  else r=mid;
```

```
build(l);
sol.maxFlow(s,t);
vector<int> ans;
for(int i=1;i<=n;++i)</pre>
  if(sol.vis[i])ans.push_back(i);
```

#### 9.0.3 Number Theory

```
1. g(m) = \sum_{d|m} f(d) \Leftrightarrow f(m) = \sum_{d|m} \mu(d) \times g(m/d)
```

2.  $\phi(x), \mu(x)$  are Möbius inverse

3.  $\sum_{i=1}^{m}\sum_{j=1}^{m}[\gcd(i,j)=1]=\sum \mu(d)\left\lfloor\frac{n}{d}\right\rfloor\left\lfloor\frac{m}{d}\right\rfloor$ 4.  $\sum_{i=1}^{n}\sum_{j=1}^{n}lcm(i,j)=n\sum_{d\mid n}d\times\phi(d)$ 

#### 9.0.4 Combinatorics

```
1. Harmonic series H_n = \ln(n) + \gamma + 1/(2n) - 1/(12n^2) + 1/(120n^4)
```

2. 
$$\gamma = 0.57721566490153286060651209008240243104215$$

3. Gray Code:  $= n \oplus (n >> 1)$ 

4. Catalan Number: 
$$\frac{C_n^{kn}}{n(k-1)+1}$$
,  $C_m^n = \frac{n!}{m!(n-m)!}$ 

6. 
$$H(n,m) \cong x_1 + x_2 \dots + x_n = k, num = C_k^{n+k-1}$$

7. Stirling number of  $2^{nd}$  kind: n 人分 k 組方法數目

(a) S(0,0) = S(n,n) = 1

(b) S(n,0) = 0

(c) 
$$S(n,k) = kS(n-1,k) + S(n-1,k-1)$$

8. Bell number, n 人分任意多組方法數目

(a)  $B_0 = 1$ 

(a)  $B_0=1$ (b)  $B_n=\sum_{i=0}^n S(n,i)$ (c)  $B_{n+1}=\sum_{k=0}^n C_k^n B_k$ (d)  $B_{p+n}\equiv B_n+B_{n+1} modp$ , p is prime (e)  $B_p{m+n}\equiv mB_n+B_{n+1} modp$ , p is prime (f) From  $B_0:1,1,2,5,15,52$ , 203,877,4140,21147,115975

9. Derangement, 錯排, 下有人在自己位置上

(a)  $D_n=n!(1-\frac{1}{1!}+\frac{1}{2!}-\frac{1}{3!}\ldots+(-1)^n\frac{1}{n!})$ (b)  $D_n=(n-1)(D_{n-1}+D_{n-2}), D_0=1, D_1=0$ (c) From  $D_0:1,0,1,2,9,44$ ,

265, 1854, 14833, 133496

10. Binomial Equality

(a)  $\sum_{k} \binom{r}{m+k} \binom{s}{n-k} = \binom{r+s}{m+n}$ (b)  $\sum_{k} \binom{l}{m+k} \binom{s}{n+k} = \binom{l+s}{l-m+n}$ (c)  $\sum_{k} \binom{l}{m+k} \binom{s+k}{n} (-1)^k = (-1)^{l+m} \binom{s-m}{n-l}$ (d)  $\sum_{k \le l} \binom{l-k}{m} \binom{s}{k-n} (-1)^k = (-1)^{l+m} \binom{s-m-1}{l-n-m}$ 

(d)  $\sum_{k \le l} {m \choose k-n} {(-1) \choose n} = {-1 \choose l}$ (e)  $\sum_{0 \le k \le l} {l-k \choose m} {q+k \choose n} = {l+q+1 \choose m+n+1}$ (f)  ${r \choose k} = {(-1)^k} {k-r-1 \choose k}$ (g)  ${r \choose m} {m \choose k} = {r \choose k} {r-k \choose m-k}$ (h)  $\sum_{k \le n} {r+k \choose k} = {r+n+1 \choose n}$ (i)  $\sum_{0 \le k \le n} {m \choose m} = {n+1 \choose m+1}$ 

(j)  $\sum_{k \le m} {m+r \choose k} x^k y^k = \sum_{k \le m} {r \choose k} (-x)^k (x+y)^{m-k}$ 

## 9.0.5 [上次, 上次和

1.  $a^b \% P = a^{b\%\varphi(p) + \varphi(p)}, b \ge \varphi(p)$ 

2.  $1^3 + 2^3 + 3^3 + \ldots + n^3 = \frac{n^4}{4} + \frac{n^3}{2} + \frac{n^2}{4}$ 

3.  $1^4 + 2^4 + 3^4 + \dots + n^4 = \frac{n^5}{5} + \frac{n^4}{2} + \frac{n^3}{3} - \frac{n}{30}$ 4.  $1^5 + 2^5 + 3^5 + \dots + n^5 = \frac{n^6}{6} + \frac{n^5}{2} + \frac{5n^4}{12} - \frac{n^2}{12}$ 

5.  $0^k + 1^k + 2^k + \ldots + n^k = P(k), P(k) = \frac{(n+1)^{k+1} - \sum_{i=0}^{k-1} C_i^{k+1} P(i)}{k+1}, P(0) = \frac{(n+1)^{k+1} - \sum_{i=0}^{k-1} C_i^{k+1} P(i)}{k+1}$ 

6.  $\sum_{k=0}^{m-1} k^n = \frac{1}{n+1} \sum_{k=0}^n C_k^{n+1} B_k m^{n+1-k}$ 

7.  $\sum_{j=0}^{m} C_j^{m+1} B_j = 0, B_0 = 1$ 

8. 除了  $B_1 = -1/2$ ,剩下的奇數項都是 0

 $B_2 = 1/6, B_4 = -1/30, B_6 = 1/42, B_8 = -1/30, B_{10} = 5/66, B_{12} = -691/2730, B_{14} = 7/6, B_{16} = -3617/510, B_{18} = 43867/798, B_{20} = -3617/510, B_{18} = 43867/798, B_{20} = -3617/510, B_{20}$ -174611/330,

#### 9.0.6 Burnside's lemma

- 1.  $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$
- $2. \ X^g = t^{c(g)}$
- 3. G 表示有幾種轉法, $X^g$  表示在那種轉法下,有幾種是會保持對稱的,t 是 $\mathbb{E}$ 數, c(g) 是循環節不動的面數。
- 4. 正立方體塗三匠色,轉 0 有  $3^6$  個元素不變,轉 90 有 6 種,每種有  $3^3$  不變,180 有  $3 \times 3^4$ ,120(角) 有  $8 \times 3^2$ ,180(邊) 有  $6 \times 3^3$ ,全部  $\frac{1}{24}$  ( $3^6 + 6 \times 3^3 + 3 \times 3^4 + 8 \times 3^2 + 6 \times 3^3$ ) = 57

# 9.0.7 Count on a tree

- 1. Rooted tree:  $s_{n+1} = \frac{1}{n} \sum_{i=1}^n (i \times a_i \times \sum_{j=1}^{\lfloor n/i \rfloor} a_{n+1-i \times j})$
- 2. Unrooted tree:
  - (a)  $\operatorname{Odd}: a_n \sum_{i=1}^{n/2} a_i a_{n-i}$  (b)  $\operatorname{Even}: Odd + \frac{1}{2} a_{n/2} (a_{n/2} + 1)$
- 3. Spanning Tree

  - (a) 完全圖  $n^n-2$  (b) 般圖 (Kirchhoff's theorem) $M[i][i]=degree(V_i), M[i][j]=-1,$ if have E(i,j),0 if no edge. delete any one row and col in A, ans=det(A)